

## A NOTE ON NON-BINARY ORTHOGONAL CODES

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## ABSTRACT

This paper presents three methods of constructing orthogonal signals whose amplitude levels are discrete, but not limited to binary: (1) method using M-sequence, (2) method by inspection, and (3) recursive method.

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## A NOTE ON NON-BINARY ORTHOGONAL CODES\*

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An effective set of signals for use in a channel with additive white Gaussian noise is the orthogonal set. Methods of constructing orthogonal continuous waveforms are widely studied. The construction of binary orthogonal codes is based primarily on Hadamard matrices. A Hadamard matrix is an orthogonal matrix whose elements are the integers +1 and -1. Hadamard matrices of various orders have been constructed<sup>1,2,3</sup> through the generation of pseudo-random sequences of the types (1) maximum length sequences (m-sequences), (2) quadratic residue sequence (or Legendre sequence), (3) twin prime sequence, and (4) Hall sequence. It seems that no such study has been made for the construction of orthogonal matrices using integers (or rational numbers) as elements, although their uses in non-binary coding can be anticipated. Furthermore, it is felt that such study may bring the two areas of endeavor, discrete coding and waveform design, closer to each other.

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Three methods are explored. They are summarized as follows.

(1) M-Sequences Over GF(p), p = 3, 5, 7, 11

To illustrate this method by an example, consider  $p = 5$  and an irreducible primitive polynomial of degree  $m = 2$  over  $GF(5)$

$$f(x) = x^2 + 3x + 3.$$

With the aid of the shift register circuit shown in Fig. 1, it is easy to see that a typical sequence generated by the polynomial is

$$* 1 0 2 -1 2 2 -2 0 1 2 1 1 -1 0 -2 1 -2 -2 2 0 -1 -2 -1 -1,$$

with period  $r = 5^2 - 1 = 24$ . By listing the above sequence and 11 successive cyclic shifts of the sequence in 12 rows, and retaining only the first 12 columns, then a  $12 \times 12$  orthogonal matrix using elements  $0, \pm 1, \pm 2$  is obtained.

1	0	2	-1	2	2	-2	0	1	2	1	1
0	2	-1	2	2	-2	0	1	2	1	1	-1
2	-1	2	2	-2	0	1	2	1	1	-1	0
-1	2	2	-2	0	1	2	1	1	-1	0	-2
2	2	-2	0	1	2	1	1	-1	0	-2	1
2	-2	0	1	2	1	1	-1	0	-2	1	-2
-2	0	1	2	1	1	-1	0	-2	1	-2	-2
0	1	2	1	1	-1	0	-2	1	-2	-2	2
1	2	1	1	-1	0	-2	1	-2	-2	2	0
2	1	1	-1	0	-2	1	-2	-2	2	0	-1
1	1	-1	0	-2	1	-2	-2	2	0	-1	-2
1	-1	0	-2	1	-2	-2	2	0	-1	-2	-1

This method is easily extended to generate  $n \times n$  orthogonal matrices,  $n = \frac{r}{2} = \frac{2^m - 1}{2}$ . Similar procedures, with proper mapping of the elements in the field of  $GF(p)$  onto elements of integers, or rational numbers (see table) can

be used to obtain  $n \times n$  orthogonal matrices with 3, 7 and 11 elements with  $n = \frac{r}{2} = \frac{p^m - 1}{2}$ ,  $p > 2$ .

Table of Mapping Elements of GF(p) to Integers or Rationals for the Construction of Orthogonal Matrices from the M-Sequences

p	Elem of GF(p)	Integ	p	Elem of GF(p)	Ratnl	Integ	p	Elem of GF(p)	Ratnl	Integ
2	0	1	7	0	0	0	11	0	0	0
	1	-1		1	1	3		1	1	2
3	0	0		2	2	6		2	2	4
	1	1		3	2/3	2		3	3	6
	2	-1	4	-2/3	-2	4		4	8	
5	0	0	5	-2	-6	5		-1/2	-1	
	1	1	6	-1	-3	6		1/2	1	
	2	-1				7		-4	-8	
	3	-2				8		-3	-6	
	4	-1				9		-2	-4	
						10	-1	-2		

The construction is based upon the properties<sup>1</sup> of the autocorrelation function  $\phi(\tau)$  of the m-sequences of p elements ( $p = 3, 5, 7, 11$ ) relative to certain mapping  $\eta$ . The autocorrelation function has the same period as the m-sequence, namely,  $r = p^m - 1$ . Under symmetrical mapping, as adopted here, the values of  $\phi(\tau)$  at  $\tau = 0$  and  $\tau = r/2$  differ in signs but equal in magnitude. Therefore, unlike the case for  $p = 2$ , a segment equal to the half period of the sequence is used for construction of the orthogonal matrices. Furthermore, for cases  $p = 7$  and 11,  $\phi(\tau)$  assumes non-zero values under ordinary mapping for  $\tau$  smaller than a half period. These are restored to zero by suitable remapping the elements of GF(7) and GF(11) onto specially chosen sets of rationals or integers. Similar procedures applied to the cases for  $p > 11$

result in solutions in mapping of elements of GF(13), etc onto elements of irrational or complex field.

(2) Construction by Inspection

The following orthogonal matrices are obtained by inspection.

(1)  $2 \times 2$  (3 level or less)

$$\begin{bmatrix} a & b \\ -b & a \end{bmatrix}$$

(2)  $4 \times 4$  (7 level or less)

$$\begin{bmatrix} a & b & c & d \\ -b & a & -d & c \\ -c & d & a & -b \\ -d & -c & b & a \end{bmatrix}$$

(3)  $8 \times 8$  (15 level or less)

$$\begin{bmatrix} a & b & -c & d & e & f & g & h \\ -b & a & d & c & -f & e & -h & g \\ c & -d & a & b & -g & h & e & -f \\ -d & -c & -b & a & -h & -g & f & e \\ -e & f & g & h & a & -b & c & -d \\ -f & -e & -h & g & b & a & -d & -c \\ -g & h & -e & -f & -c & d & a & -b \\ -h & -g & f & -c & d & c & b & a \end{bmatrix}.$$

By assigning suitable values to the letters, some of which may have the same value, orthogonal matrices of various elements can be constructed.

(3) Recursive Methods

Let A and B be two orthogonal matrices of size  $n \times n$ . Then the following recursive methods may be used to obtain new orthogonal matrices.

(a)  $C = A \cdot B$  size  $n \times n$ , different elements from A or B.

(b)  $C = \begin{bmatrix} 1 & 1 \\ -1 & 1 \end{bmatrix} \otimes A$  size  $2n \times 2n$ , same elements as A.  
 where  $\otimes$  denotes the Kronecker or tensor product.

$$(c) \quad C = \begin{bmatrix} a & b \\ -b & a \end{bmatrix} \otimes A \quad \text{size } 2n \times 2n, \text{ different elements as } A.$$

$$(d) \quad C = \begin{bmatrix} A & B \\ -B^T & D \end{bmatrix} \quad \text{size } 2n \times 2n, \text{ same elements as } A \text{ and } B \\ \text{combined provided } AB = BA.$$

In the last method, the matrix D is computed from:

$$D^T = B^{-1} AB.$$

However, if  $AB = BA$ , then

$$D^T = B^{-1} BA = A$$

and

$$D = A^T.$$

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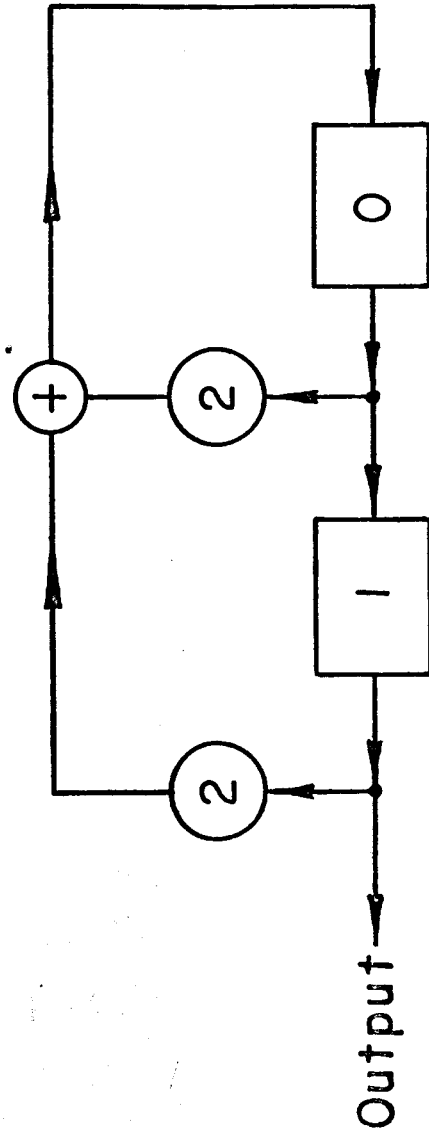


Fig. 1. Shift Register Circuit used to Generate M-Sequence Over GF(5).

Recursion Polynomial:  $f(x) = x^2 + 3x + 3$ .