A SOLUTION TO A COUNTABLE SYSTEM OF EQUATIONS ARISING IN MARKOVIAN DECISION PROCESSES

by

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Nontechnical Summary

Let X_0, X_1, \ldots be a sequence of non-negative integer valued random variables with the property that

$$Pr(X_{n+1} = j | X_0 = x_0, ..., X_{n-1} = x_{n-1}, X_n = i) = p_{ij}$$

for all i, j, x_0, \ldots, x_n , n. The collection of random variables $\{X_n^{+}\}$ is called a Markov chain and the p_{ij} are called transition probabilities. We refer to X_n as the state of the process at time n. Let w_i be the cost incurred at time n if the process is in state i at that time. Consider the system of equations

(1)
$$g + v_i = w_i + \sum_{j=0}^{\infty} p_{ij}v_j, i = 0, l, ...$$

in the unknown variables g, v_0 , v_1 , \cdots . Such a system arises in connection with constructing optimal rules for controlling Markovian decision processes. Also the numbers g, v_0 , v_1 , \cdots are of interest in their own right. Often g is the long run expected average cost and $v_i - v_j$ is the limit, as $n \to \infty$, of the difference between expected total cost during times 0, 1, \cdots , n given that the process starts in states i and j respectively.

We show in this paper that one solution to the system (1) is given by

(2)
$$g = \frac{c_{oo}}{m_{oo}}$$
 and $v_i = c_i - gm_{io}$, $i = 0, 1, ...$

Acquisitioned Document SQT provided that the expected time m_{io} required to go from state i to state 0 is finite and that the expected cost c_{io} incurred during that time is also finite, i = 0, 1, ... Notice that $v_0 = 0$.

As an illustration of the above ideas, consider a single item inventory model in which the demands in periods 1, 2, ... are independent. A demand of size one occurs with probability p, 0 , and $a demand of size zero occurs with probability 1 - p. Let <math>X_n$ denote the stock on hand at the beginning of period n. An order for one unit is placed in period n with immediate delivery if $X_n = 0$; otherwise, no order is placed in period n. There is a unit cost h for each unit of stock on hand after ordering in a period. There is a cost K for placing an order in a period. Under these assumptions the nonzero transition probabilities are $p_{oo} = p$, $p_{ol} = 1 - p$, $p_{ii} = 1 - p$, and $p_{i,i-1} = p$, i = 1, 2, ... Also $w_o = K + h$ and $w_i = hi$, i = 1, 2, Thus the system (1) becomes

 $g + v_0 = K + h + pv_0 + (l - p)v_1$

$$g + v_{i} = ih + pv_{i-1} + (1 - p)v_{i}, i = 1, 2, ...$$

The solution given in (2) is

 $v_{i} = \frac{hi(i-1)}{2p} - Ki, i = 0, 1, \dots$

g = pK + h,

ii

Thus g is here the long run expected average cost under the indicated ordering policy. Also v_i is the limit, as $n \to \infty$, of the amount by which the expected cost in periods 0, 1, ..., n starting with i units of stock on hand exceeds that starting with no stock on hand.

A SOLUTION TO A COUNTABLE SYSTEM OF EQUATIONS

ARISING IN MARKOVIAN DECISION PROCESSES

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Let $\{X_n\}$, n = 0, 1, ..., be a Markov chain having a state space consisting of the non-negative integers and having stationary transition probabilities $\{p_{ij}\}$. Let $\{w_l\}$, i = 0, 1, ..., be a sequence of real numbers. Consider the system of equations

(1)
$$g + v_i = w_i + \sum_{j=0}^{\infty} p_{ij}v_j, i = 0, 1, ...,$$

in the unknown variables $\{g, v_0, v_1, \dots\}$. In [2], the system (1) arises in connection with conditions for the existence and construction of optimal rules for controlling a Markovian decision process. For a finite state space existence of solutions to (1) is guaranteed by the condition that the Markov chain have at most one ergodic class of states. (See [3].) In this note we give conditions ensuring the existence (Theorem 1) and uniqueness (Theorem 2) of solutions to (1).

Let

$$Z_{n}(j) = \begin{cases} 1, & \text{if } X_{n} = j \text{ and if } X_{m} \neq 0 \text{ for } 0 < m \leq n \\ 0, & \text{otherwise} \end{cases}$$

$$j, n = 0, l, \dots,$$

$$o^{p} \underset{ij}{*} = E\left(\sum_{n=0}^{\infty} Z_{n}(j) | X_{o} = i\right), \quad i, j = 0, l, \dots,$$

 $m_{io} = \sum_{j=0}^{\infty} o_{ij}^{*}, i = 0, 1, \dots$

If the last series converges absolutely, then m_{io} is the mean first passage time from i to 0 and we say m_{io} is finite. If the m_{io} are all finite, as we assume throughout, then state 0 is positive recurrent and there is only one recurrent class.

Let
$$Y_n = \sum_{j=0}^{\infty} w_j Z_n(j)$$
 and $c_{i0} = E \left(\sum_{n=0}^{\infty} Y_n | X_0 = i \right)$

By an obvious generalization of Theorem 5 in [1, p. 81] we get $c_{i0} = \sum_{j=0}^{\infty} o_{ij}^{*} w_{j}$ provided the series is absolutely convergent. If the series is absolutely convergent we say c_{i0} is finite. In applications w_{i} is often the cost incurred when in state i so c_{i0} is then the expected cost during a first passage from i to o.

Theorem 1 (Existence)

If the numbers m_{io} and c_{io} , i = 0, 1, ..., are finite, then the numbers

(2)
$$g = \frac{c_{oo}}{m_{oo}}$$
 and $v_i = c_{io} - gm_{io}$, $i = 0, 1, ...$

satisfy (1) and $\sum_{j=0}^{\infty} p_{ij}v_j$ converges absolutely, i = 0, 1, ... Proof:

Let
$$w_i^* = w_i - g$$
 and $Y_n^* = \sum_{j=0}^{\infty} w_j^* Z_n(j)$. Then for $i = 0, 1, ...$

and

$$\mathbf{v_{i}} = \mathbf{E} \left(\sum_{n=0}^{\infty} Y_{n}^{*} | X_{0} = \mathbf{i} \right)$$
$$= \mathbf{w_{i}^{*}} + \sum_{n=1}^{\infty} \sum_{j=0}^{\infty} \mathbf{E}(Y_{n}^{*} | X_{0} = \mathbf{i}, X_{1} = \mathbf{j}) \mathbf{p_{ij}}$$
$$= \mathbf{w_{i}^{*}} + \sum_{j=0}^{\infty} \sum_{n=1}^{\infty} \mathbf{E}(Y_{n}^{*} | X_{0} = \mathbf{i}, X_{1} = \mathbf{j}) \mathbf{p_{ij}}$$
$$= \mathbf{w_{i}^{*}} + \sum_{j=0}^{\infty} \mathbf{p_{ij}} \mathbf{v_{j}}$$

so (1) holds. The interchange of expectation and summation is justified since the finiteness of the m_{io} and c_{io} imply that $\sum_{n=0}^{\infty} E(|Y_n^*|| X_o = i) < \infty$. This in turn implies that the series above are absolutely convergent so the interchange of summations is also justified. <u>Theorem 2</u> (Uniqueness)

If the numbers m_{io} and c_{io} , i = 0, 1, ..., are finite, if $\sum_{j=0}^{\infty} o_{ij}^{*} \left(c_{jo} - \frac{c_{oo}}{m_{oo}} m_{jo} \right), i = 0, 1, ... converges absolutely, and if$ $[g, <math>v_{o}, v_{1}, ...$] is a sequence with $\sum_{j=0}^{\infty} o_{ij}^{*} v_{j}, i = 0, 1, ...,$ converging absolutely, then [g, $v_{o}, v_{1}, ...$] satisfies (1) if and only if there is a real number r such that

(3)
$$g = \frac{c_{oo}}{m_{oo}}$$
 and $v_i = c_{io} - gm_{io} + r, i = 0, 1, ...$

Proof:

It is immediate from the hypotheses and Theorem 1 that {g, v_o, v₁, ...} defined in (3) satisfies (1) and $\sum_{j=0}^{\infty} o_{ij}^{*}v_{j}$ converges absolutely as well as $\sum_{j=0}^{\infty} p_{ij}v_{j}$. Let {g', v'_o, v'₁, ...} be any other solution to (1) with $\sum_{j=0}^{\infty} p_{ij}^* v_j^i$ converging absolutely for $i = 0, 1, \dots$. Hence $\sum_{k=0}^{\infty} p_{ik} v_k^i$ is absolutely convergent. Now premultiplying both sides of (1) by $\pi_i \equiv \frac{o_i^{p_i^*}}{m_{oo}}$, summing over $i = 0, 1, \dots$, using the relations $\sum_{i=0}^{\infty} \pi_i = 1$ and $\pi_j = \sum_{k=0}^{\infty} p_{kj}\pi_k$, $j = 0, 1, \dots$, and the fact that the interchange of summations is justified, we get $g' = \sum_{i=0}^{\infty} \pi_i w_i$ which is independent of $\{v'_0, v'_1, \dots\}$. Thus since

 $\{g, v_0, v_1, \ldots\} \text{ satisfies (1) we must have } g = g^{1}.$

Letting $\triangle_i = v'_i - v_i$, $i = 0, 1, \dots$, we get from (1) on subtracting one system from the other that

Let $p_{ij}^n = \Pr(X_n = j | X_0 = i)$. Evidently for N = 1, 2, ...,

$$\sum_{n=1}^{N} p_{ij}^{n} \leq p_{ij}^{*} + (N-1) p_{oj}^{*}, j = 0, 1, \dots$$

so

(5)
$$\frac{1}{N}\sum_{n=1}^{N}p_{ij}^{n}|\Delta_{j}| \leq [p_{ij}^{*}+p_{oj}^{*}]|\Delta_{j}|, \qquad j=0, 1, \ldots.$$

Since the series on the right side of (5) converges absolutely by hypothesis, and $\lim_{N \to \infty} \frac{1}{N} \sum_{n=1}^{\infty} p^n_{ij} = \pi_j$, we get from the dominated convergence theorem that

(6)
$$\lim_{N \to \infty} \sum_{j=0}^{\infty} \frac{1}{N} \sum_{n=1}^{\infty} p_{j}^{n} \Delta_{j} = \sum_{j=0}^{\infty} \pi_{j} \Delta_{j}$$

Since from (5), $\sum_{j=0}^{\infty} p_{ij}^{n} \triangle_{j}$ converges absolutely we can iterate (4), yielding

(7)
$$\Delta_{i} = \sum_{j=0}^{\infty} p_{ij}^{n} \Delta_{j}, \quad i = 0, 1, ...; \quad n = 1, 2, ...$$

Hence on substituting (7) into (6)

$$\Delta_{\mathbf{i}} = \sum_{\mathbf{j}=0}^{\infty} \pi_{\mathbf{j}} \Delta_{\mathbf{j}}, \quad \mathbf{i} = 0, \mathbf{l}, \ldots$$

Thus \triangle_{i} is independent of i, which completes the proof. Example:

If the sequences $\{m_{i0}\}$ and $\{w_i\}$, i = 0, 1, ..., are bounded, then so is the sequence $\{c_{i0}\}$, i = 0, 1, ..., since $|c_{i0}| \leq \sup_{k,j} m_{k0} |w_j|$. Thus Theorem 1 applies and in addition the solution to (1) given in (2) is bounded. This result is used in [2].

We remark that since

$$\sum_{\mathbf{j}=\mathbf{0}}^{\infty} \mathbf{p}_{\mathbf{0}\mathbf{j}}^{\mathbf{*}} |\mathbf{u}| \geq \mathbf{p}_{\mathbf{0}\mathbf{k}\mathbf{k}}^{\mathbf{*}} \sum_{\mathbf{j}=\mathbf{0}}^{\infty} \mathbf{p}_{\mathbf{k}\mathbf{j}}^{\mathbf{*}} |\mathbf{u}_{\mathbf{j}}|$$

where

$$o^{f}ok = Pr\left(\sum_{n=0}^{\infty} Z_{n}(k) > 0 | X_{o} = 0\right) > 0$$
,

 $\sum_{j=0}^{\infty} o^{p_{kj}^{*}|u_{j}|} \text{ is absolutely convergent for every recurrent state } k$ provided that $\sum_{j=0}^{\infty} o^{p_{0j}^{*}|u_{j}|} \text{ is absolutely convergent. Thus the hypotheses of Theorems 1 and 2 could have been stated only for state 0 and the transient states.}$

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