# A SOLUTION TO A COUNTABLE SYSTEM OF EQUATIONS ARISING IN MARKOVIAN DECISION PROCESSES 

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## Nontechnical Summary

Let $X_{o}, X_{1}, \ldots$ be a sequence of non-negative integer valued random variables with the property that

$$
\operatorname{Pr}\left(x_{n+1}=j \mid x_{0}=x_{0}, \ldots, X_{n-1}=x_{n-1}, x_{n}=i\right)=p_{i j}
$$

for all $i, j, x_{0}, \ldots, x_{n}, n$. The collection of random variables $\left\{X_{n}\right\}$ is called a Markov chain and the $p_{i j}$ are called transition probabilities. We refer to $X_{n}$ as the state of the process at time $n_{\text {. }}$ Let $w_{i}$ be the cost incurred at time $n$ if the process is in state $i$ at that time. Consider the system of equations

$$
\begin{equation*}
g+v_{i}=w_{i}+\sum_{j=0}^{\infty} p_{i j} v_{j}, i=0,1, \ldots \tag{I}
\end{equation*}
$$

in the unknown variables $E$, $v_{0}, v_{1}, \ldots$. Such a system arises in connection with constructing optimal rules for controlling Markovian decision processes. Also the numbers $g, v_{0}, v_{1}, \ldots$ are of interest in their own right. Often $g$ is the long run expected average cost and $v_{i}-v_{j}$ is the limit, as $n \rightarrow \infty$, of the difference between expected total cost during times $0,1, \ldots, n$ given that the process siaris iftsiaies i anā $j$ respecinveiy.

We show in this paper that one solution to the system (1) is given by
(2) $\quad g=\frac{c_{00}}{m_{00}}$ and $v_{i}=c_{i 0}-g m_{i 0}, i=0,1, \ldots$
provided that the expected time $m_{i o}$ required to go from state i to state 0 is finite and that the expected cost $c_{\text {io }}$ incurred during that time is also finite, $i=0,1, \ldots$ Notice that $v_{0}=0$. As an illustration of the above ideas, consider a single item inventory model in which the demands in periods $1,2, \ldots$ are independent. A demand of size one occurs with probability $p, 0<p<1$, and a demand of size zero occurs with probability 1 - $p$. Let $X_{n}$ denote. the stock on hand at the beginning of period $n$. An order for one unit is placed in period $n$ with immediate delivery if $X_{n}=0$; otherwise, no order is placed in period $n$. There is a unit cost $h$ for each unit of stock on hand after ordering in a period. There is a cost $K$ for placing an order in a period. Under these assumptions the nonzero transition probabilities are $p_{00}=p, p_{o l}=1-p, p_{i i}=1-p$, and $p_{i, i-1}=p, i=1,2, \ldots$. Also $w_{o}=K+h$ and $w_{i}=h i, i=1,2, \ldots$. Thus the system (1) becomes

$$
\begin{aligned}
& g+v_{0}=K+h+p v_{0}+(1-p) v_{1} \\
& g+v_{i}=i h+p v_{i-1}+(1-p) v_{i}, i=1,2, \ldots 0
\end{aligned}
$$

The solution given in (2) is

$$
\begin{gathered}
g=p K+h, \\
v_{i}=\frac{h i(i-1)}{2 p}-K i, i=0, l, \ldots
\end{gathered}
$$

Thus $g$ is here the long run expected average cost under the indicated ordering policy. Also $v_{i}$ is the limit, as $n \rightarrow \infty$, of the amount by which the expected cost in periods $0,1, \ldots, n$ starting with $i$ units of stock on hand exceeds that starting with no stock on hand.

# A SOLUTION TO A COUNTABLE SYSTEM OF EQUATIONS ARISING IN MARKOVIAN DECISION PROCESSES 

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Let $\left\{X_{n}\right\}, n=0,1, \ldots$, be a Markov chain having a state space consisting of the non-negative integers and having stationary transition probabilities $\left\{p_{i j}\right\}$. Let $\left\{w_{1}\right\}, i=0,1, \ldots$, be a sequence of real numbers. Consider the system of equations

$$
\begin{equation*}
g+v_{i}=w_{i}+\sum_{j=0}^{\infty} p_{i j} v_{j}, \quad i=0, l, \ldots, \tag{1}
\end{equation*}
$$

in the unknown variables $\left\{\mathrm{g}, \mathrm{v}_{0}, \mathrm{v}_{1}, \ldots\right\}$. In [2], the system (1) arises in connection with conditions for the existence and construction of optimal rules for controlling a Markovian decision process. For a finite state space existence of solutions to (i) is guaranteed by the condition that the Markov chain have at most one ergoaic class of states. (See [3].) In this note we give conditions ensuring the existence (Theorem 1) and uniqueness (Theorem 2) of solutions to (1).

Let

$$
\begin{aligned}
& z_{n}(j)=\left\{\begin{array}{l}
1, \quad \text { if } X_{n}=j \text { and if } x_{m} \neq 0 \text { for } 0<m \leq n \\
0, \text { otherwise }
\end{array}\right. \\
& \quad j, n=0,1, \ldots, \\
& o^{p_{i j}^{*}=}=E\left(\sum_{n=0}^{\infty} z_{n}(j) \mid x_{o}=i\right), i, j=0,1, \ldots,
\end{aligned}
$$

and

$$
m_{i o}=\sum_{j=0}^{\infty}{ }_{o} p_{i j}^{*}, \quad i=0,1, \ldots .
$$

If the last series converges absolutely, then $m_{i o}$ is the mean first passage time from $i$ to 0 and we say $m_{i o}$ is finite. If the $m_{i o}$ are all finite, as we assume throughout, then state 0 is positive recurrent and there is only one recurrent class.

$$
\text { Let } Y_{n}=\sum_{j=0}^{\infty} w_{j} Z_{n}(j) \text { and } c_{i o}=E\left(\sum_{n=0}^{\infty} Y_{n} \mid X_{o}=i\right)
$$

By an obvious generalization of Theorem 5 in [1, p. 81] we get $c_{i o}=\sum_{j=0}^{\infty} o^{p_{i j}^{*}}{ }^{w} j$ provided the series is absolutely convergent. If the series is absolutely convergent we say $c_{i o}$ is finite. In applications $w_{i}$ is often the cost incurred when in state $i$ so $c_{i o}$ is then the expected cost during a first passage from $i$ to 0.

Theorem 1 (Existence)
If the numbers $m_{i o}$ and $c_{i o}, i=0, l, \ldots$, are finite, then the numbers
(2)

$$
g=\frac{c_{00}}{m_{00}} \text { and } v_{i}=c_{i o}-g m_{i 0}, \quad i=0, I, \ldots
$$

satisfy (1) and $\sum_{j=0}^{\infty} p_{i j}{ }_{j}$ converges absolutely, $i=0,1, \ldots$. Proof:

Let $w_{i}^{*}=w_{i}-g$ and $Y_{n}^{*}=\sum_{j=0}^{\infty} w_{j}^{*} Z_{n}(j)$. Then for $i=0,1, \ldots$

$$
\begin{aligned}
v_{i} & =E\left(\sum_{n=0}^{\infty} Y_{n}^{*} \mid X_{o}=i\right) \\
& =w_{i}^{*}+\sum_{n=1}^{\infty} \sum_{j=0}^{\infty} E\left(Y_{n}^{*} \mid X_{o}=i, X_{l}=j\right) p_{i j} \\
& =w_{i}^{*}+\sum_{j=0}^{\infty} \sum_{n=1}^{\infty} E\left(Y_{n}^{*} \mid X_{o}=i, X_{1}=j\right) p_{i j} \\
& =w_{i}^{*}+\sum_{j=0}^{\infty} p_{i j} v_{j}
\end{aligned}
$$

so (1) holds. The interchange of expectation and summation is justified since the finiteness of the $m_{i o}$ and $c_{i o}$ imply that $\sum_{n=0}^{\infty} E\left(\left|Y_{n}^{*}\right| \mid X_{o}=i\right)<\infty$. This in turn implies that the series above are absolutely convergent so the interchange of summations is also justified. Theorem 2 (Uniqueness)

If the numbers $m_{i o}$ and $c_{i o}, i=0,1, \ldots$, are finite, if $\sum_{j=0}^{\infty} o^{p_{i j}^{*}}\left(c_{j o}-\frac{c_{00}}{m_{o o}} m_{j o}\right)^{10}, i=0,1, \ldots$ converges absolutely, and if $\left\{g, v_{o}, v_{1}, \ldots\right\}$ is a sequence with $\sum_{j=0}^{\infty} o^{p_{i j}^{*}} v_{j}, i=0,1, \ldots$, converging absolutely, then $\left\{g, v_{o}, v_{l}, \ldots\right\}$ satisfies (l) if and only if there is a real number $r$ such that

$$
\begin{equation*}
g=\frac{c_{o o}}{m_{o o}} \text { and } v_{i}=c_{i o}-{g m_{i o}}+r, i=0, l, \ldots . \tag{3}
\end{equation*}
$$

Proof:
It is immediate from the hypotheses and Theorem 1 that $\left\{g, v_{o}, v_{l}, \ldots\right\}$ defined in (3) satisfies (1) and $\sum_{j=0}^{\infty}{ }_{o} p_{i j}^{*} v_{j}$ converges absolutely as well as $\sum_{j=0}^{\infty} p_{i j} v_{j}$. Let $\left\{g^{\prime}, v_{o}^{\prime}, v_{1}^{\prime}, \ldots\right\}$ be
any other solution to (1) with $\sum_{j=0}^{\infty} \circ^{p_{i j}^{*}} v_{j}^{\prime}$ converging absclutely for $i=0,1, \ldots$ Hence $\sum_{k=0}^{\infty} p_{i k} v_{k}^{\prime}$ is absolutely convergent. Now premultiplying both sides of (1) by $\pi_{i} \equiv \frac{o^{p}}{m_{00}}$, summing over $i=0,1, \ldots$, using the relations $\sum_{i=0}^{\infty} \pi_{i}=1$ and $\pi_{j}=\sum_{k=0}^{\infty} p_{k j}{ }^{\top} k, j=0,1, \ldots$, and the fact that the interchange of summations is justified, we get $g^{\prime}=\sum_{i=0}^{\infty} \pi_{i} W_{i}$ which is independent of $\left\{v_{o}^{\prime}, v_{l}^{\prime}, \ldots\right\}$. Thus since $\left\{g, v_{o}, v_{1}, \ldots\right\}$ satisfies (1) we must have $g=g^{\prime}$.

Letting $\Delta_{i}=v_{i}^{\prime}-v_{i}, i=0,1, \ldots$, we get from (l) on subtracting one system from the other that

$$
\begin{equation*}
\Delta_{i}=\sum_{j=0}^{\infty} p_{i j} \Delta_{j}, \quad i=0,1, \ldots \tag{4}
\end{equation*}
$$

Let $p_{i j}^{n}=\operatorname{Pr}\left(X_{n}=j \mid X_{o}=i\right)$. Evidently for $N=1,2, \ldots$,

$$
\sum_{n=1}^{N} p_{i j}^{n} \leq{ }_{o} p_{i j}^{*}+(N-1)_{o} p_{o j}^{*}, \quad j=0,1, \ldots
$$

so

$$
\begin{equation*}
\frac{1}{\mathbb{N}} \sum_{n=1}^{N} p_{i j}^{n}\left|\Delta_{j}\right| \leq\left[{ }_{o} p_{i j}^{*}+{ }_{o} p_{o j}^{*}\right]\left|\Delta_{j}\right|, \quad j=0,1, \ldots \tag{5}
\end{equation*}
$$

Since the series on the right side of (5) converges absolutely by hypothesis, and $\lim _{N \rightarrow \infty} \frac{1}{N} \sum_{n=1}^{\infty} p_{i j}^{n}=\pi_{j}$, we get from the dominated convergence theorem that

$$
\begin{equation*}
\lim _{N \rightarrow \infty} \sum_{j=0}^{\infty} \frac{1}{N} \sum_{n=1}^{\infty} p_{i j}^{n} \Delta_{j}=\sum_{j=0}^{\infty} \pi_{j} \Delta_{j} . \tag{6}
\end{equation*}
$$

Since from (5), $\sum_{j=0}^{\infty} p_{i j}^{n} \Delta_{j}$ converges absclutely we can iterate (4), yielding

$$
\begin{equation*}
\Delta_{i}=\sum_{j=0}^{\infty} p_{i j}^{n} \Delta_{j}, \quad i=0,1, \ldots ; \quad n=1,2, \ldots . \tag{7}
\end{equation*}
$$

Hence on substituting (7) into (6)

$$
\Delta_{i}=\sum_{j=0}^{\infty} \pi_{j} \triangle_{j}, \quad i=0,1, \ldots .
$$

Thus $\Delta_{i}$ is independent of $i$, which completes the proof.
Example:
If the sequences $\left\{m_{i 0}\right\}$ and $\left\{w_{i}\right\}, i=0,1, \ldots$, are bounded, then so is the sequence $\left\{c_{i o}\right\}, i=0,1, \ldots$, since $\left|c_{i o}\right| \leq \sup _{k, j} m_{k o}\left|w_{j}\right|$. Thus Theorem $I$ applies and in addition the solution to (1) given in (2) is bounded. This result is used in [2].

We remark that since

$$
\sum_{j=0}^{\omega} o_{o}^{*}|u| \geq o_{o k}^{f} \sum_{j=0}^{m} o^{p_{k j}^{*}}\left|u_{j}\right|
$$

where

$$
o_{o k}=\operatorname{Pr}\left(\sum_{n=0}^{\infty} z_{n}(k)>0 \mid X_{0}=0\right)>0
$$

$\sum_{j=0}^{\infty} o^{p_{k j}^{*}}\left|u_{j}\right|$ is absolutely convergent for every recurrent state $k$ provided that $\sum_{j=0}^{\infty} o_{o j} p_{j}^{*}\left|u_{j}\right|$ is absolutely convergent. Thus the hypotheses of Theorems 1 and 2 could have been stated only for state 0 and the transient states.

## References

[1] Chung, K. L. (1960), Markov Chains with Stationary Transition Probabilities, Springer, Berlin.
[2] Derman, Cyrus (1966), "Denumerable State Markovian Decision Processes - Average Cost Criterion," (To Appear in Ann. Math. Stat.).
[3] Howard, Ronald (1960), Dynamic Programming and Markov Processes, John Wiley, New York.

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