I have one general concern: If we consider a certain entity of Gödel numbers $T$ (any kind of Gödelization) of closed terms in the language of BA and another entity $T_x$ of Gödel numbers of terms in one fixed variable $x$, and if we further consider an evaluation function $\text{val} : T \to \mathbb{N}$ and a substitution function $\text{sub} : T_x \times \mathbb{N} \to T$, then faithfulness of the arithmetization implies that all polynomial growth rates can be represented in the system $(T, T_x, \text{val}, \text{sub})$. For proving Gödel’s diagonalization lemma we need the substitution function $\text{sub}$ to have a polynomial growth rate, and for showing the correctness of derivations in BA by the evaluation strategy we need the evaluation function $\text{val}$ to have a polynomial growth rate. But this is impossible, because then $f(a, b) := \text{val}(\text{sub}(a, b))$ has polynomial growth rate, i.e., $|f(a, b)| \leq \max(a, b)^c$ for some $c \in \mathbb{N}$. By faithfulness there is some $t_{c+1} \in T_x$ such that $|m|^c+1 \leq |\text{val}(\text{sub}(t_{c+1}, m))| \leq \max(t_{c+1}, m)^c = |m|^c$ for $m$ big enough, a contradiction. I think this indicates that in order to obtain a separation via Gödel sentences we need further ideas that differ strongly from the evaluation strategy.

The first reviewed paper summarizes all necessary notions and results on strictly $i$-normal proofs and partial truth definitions for them. Let $\text{Prf}^f(w, [\varphi^f])$ denote “$w$ is the Gödel number of a strictly $i$-normal proof of $\varphi$,” and let $|x|_k$ denote the $k$-times iterated “integer-logarithm” $\lfloor \cdots \lfloor x \rfloor \rfloor$ applied to $x$. Takeuti shows that BA can prove weak consistency statements like $\forall w=\text{Prf}^f([w], \Gamma \to \Delta)$ (the empty sequent “$\cdots$” representing contradiction) and weak forms of reflection like $\forall w \text{Prf}^f([|w|_1, \forall x \varphi(x)]) \to \forall x \varphi(|x|)$ for simple $\varphi$. Of course, Gödel’s second incompleteness theorem can only be proved for the “usual” consistency statement $S_1^f \not\vdash \forall w=\text{Prf}^f([w], \Gamma \to \Delta)$. Takeuti defines Gödel sentences $\varphi_k^f$ satisfying $S_1^f \vdash \varphi_k^f \iff \forall x=\text{Prf}^f([|x|_k, \eta \varphi_k^f])$. They have similar properties to the above-described variations of consistency. By Gödel’s diagonalization lemma, Takeuti obtains $S_1^f \not\vdash \varphi_k^f$ for all $k$, but on the other hand $S_{k+1}^f \vdash \forall x=\text{Prf}^f([|x|_{k+1}, \eta \varphi_k^f])$ for $k \geq 2$. He formulates two conjectures in the above-described form that would imply $P \neq NP$.

The second reviewed paper extends this line of research and contrasts it with Gödel sentences based on the usual (unrestricted) derivability. That is, let $\text{Prf}^F(w, [\varphi^F])$ denote the formalized notion of “$w$ is the Gödel number of an (unrestricted) $S_2^F$-proof of $\varphi$,” and define Gödel sentences $\Phi_k$ satisfying $S_1^F \vdash \Phi_k \iff \forall x=\text{Prf}^F([|x|_k, \Phi_k^F])$. Takeuti compares pairs $(\text{Prf}^F, \Phi_k^F)$ with $(\text{Prf}^f, \varphi_k^F)$, obtaining results like $\Delta_0 + \exp \not\vdash \forall x=\text{Prf}^F([|x|_{k+1}, \Gamma \varphi_k^F])$ for all $k$, in contrast to the above-obtained results for $(\text{Prf}^f, \varphi_k^F)$. He further establishes relationships between different Gödel sentences, and between Gödel sentences and weak consistency statements.

The reviewed papers are top-level research papers in the field of applications of Gödel sentences to the separation problem of BA. This implies that for the study of them one needs some background in BA (at least parts of the Buss book). Like all papers of Takeuti that I have read, I found the ones under review inspiring, showing new ideas to the separation problem of BA.

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The paper under review is concerned with Gödel’s system $T$ of primitive recursive functionals of finite type, which is sufficient for providing a functional interpretation of first-order Peano arithmetic PA. The main technical achievement of the paper is a novel proof of the
strong normalization of $T$ by defining a new assignment $\sem{\ }_0 : T \to \omega$ which decreases under one-step reduction $\succ^1$ of terms in $T$, i.e., $t \succ^1 t'$ entails $\sem{t'}_0 < \sem{t}_0$. Moreover, $\sem{\ }_0$ is an $\varepsilon_0$-recursive function. As a byproduct, optimal bounds for the natural fragments $T_n$ of $T$ are obtained.

As the title of the paper suggests, this new analysis of Gödel’s $T$ has to be seen as a specific case study of a more general program, which has its roots in infinitary proof theory. There the general concept of miniaturization and concretion of certain (large) cardinals has long been central in the development of ordinal notation systems. Further, as has been shown by Weiermann, it is often possible to miniaturize or project down the ordinal analysis of a (strong) system in order to obtain a perspicuous proof-theoretic analysis of a corresponding weaker system. For example, he has shown in *How to characterize provably total functions by local predicativity* (*The journal of symbolic logic*, vol. 61 (1996), pp. 52–69) how to pin down the standard local predicativity treatment (due to Pohlers) of the theory of one inductive definition $ID_1$ to a technically smooth analysis of Peano arithmetic $PA$. In this miniaturization, the collapsing function $D : \varepsilon_{\Omega+1} \to \Omega$ is replaced by the miniaturized collapsing function $\psi : \varepsilon_0 \to \omega$. The treatment of $PA$ thus obtained indeed also produces optimal bounds for the fragments $I \Sigma^+_n$ of $PA$.

The treatment of Gödel’s $T$ in the paper under review can be seen as a miniaturization of Howard’s analysis of bar recursion of type zero and is inspired by the local predicativity approach to pure proof theory. In particular, the definition of the assignment $\sem{\ }_0$ mentioned above makes crucial use of the miniaturized collapsing function $\psi$. Whereas previous treatments of the subsystems $T_n$ of $T$ (in which the recursors have type level less than or equal to $n + 2$) used the ordinal bound $\omega_{n+3}$ (e.g. in Weiermann’s paper *A proof of strongly uniform termination for Gödel’s $T$ by methods from local predicativity*, *Archive for mathematical logic*, vol. 36 (1997), pp. 445–460—a precursor of the article under review), the present paper yields the optimal strong normalization bound $\omega_{n+2}$ for $T_n$.

This is a very important paper both from the technical as well as from the conceptual point of view. In a long and sparkling introduction, the author outlines his vision of how to apply a certain kind of infinitary methods to questions of finitary mathematics, and he discusses exciting connections of his results with term rewriting theory, hierarchy theory, and computational complexity. The paper is largely self-contained and, due to its extensive motivation and conceptual discussion, it should be accessible to a wide readership.

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In this article, the author continues his investigations of the rate of growth of functions definable in certain subsystems of analysis in finite types. In previous papers (e.g. *Mathematically strong subsystems of analysis with low rate of growth of provably recursive functionals*, BSL VII 280) the classical systems $G_nA^\omega$ were analyzed and proof-theoretic tools such as monotone functional interpretation (introduced by the author) were used. In the paper under review, Kohlenbach studies the intuitionistic versions $G_nA^\omega_i$ of these systems and analyzes them via a new monotone realizability interpretation. He shows that the addition of certain strong non-constructive and even classically refutable analytical and logical principles has no impact on the rate of growth of definable functions.

The main result of the paper is summarized in the following.

**Main theorem.** Let $E - G_nA^\omega$ be the system $G_nA^\omega$ plus extensionality axioms in all finite types; the full axiom of choice AC in all finite types, and the independence of premiss principle $IP_-$ for negated formulas. Furthermore, let $A$ consist of the following (classically