provided by Elsevier - Publisher



Available online at www.sciencedirect.com



Discrete Mathematics 306 (2006) 1969 – 1974

Note



[www.elsevier.com/locate/disc](http://www.elsevier.com/locate/disc)

# Kernels in quasi-transitive digraphs

Hortensia Galeana-Sánchez<sup>a</sup>, Rocío Rojas-Monroy<sup>b</sup>

<sup>a</sup>*Instituto de Matemáticas, UNAM, Ciudad Universitaria, Circuito Exterior, 04510 México D.F., Mexico*

<sup>b</sup>*Facultad de Ciencias, Universidad Autónoma del Estado de México, Instituto Literario No. 100, Centro 50000, Toluca,*

*Edo. de México, Mexico*

Received 18 June 2004; received in revised form 25 November 2005; accepted 3 February 2006 Available online 14 June 2006

#### **Abstract**

Let *D* be a digraph,  $V(D)$  and  $A(D)$  will denote the sets of vertices and arcs of *D*, respectively.

A kernel *N* of *D* is an independent set of vertices such that for every  $w \in V(D) - N$  there exists an arc from *w* to *N*. A digraph is called *quasi-transitive* when  $(u, v) \in A(D)$  and  $(v, w) \in A(D)$  implies  $(u, w) \in A(D)$  or  $(w, u) \in A(D)$ . This concept was introduced by Ghouilá–Houri [Caractérisation des graphes non orientés dont on peut orienter les arrêtes de maniere à obtenir le graphe d' un relation d'ordre, C.R. Acad. Sci. Paris 254 (1962) 1370–1371] and has been studied by several authors. In this paper the following result is proved: Let *D* be a digraph. Suppose  $D = D_1 \cup D_2$  where  $D_i$  is a quasi-transitive digraph which contains no asymmetrical infinite outward path (in  $D_i$ ) for  $i \in \{1, 2\}$ ; and that every directed cycle of length 3 contained in *D* has at least two symmetrical arcs, then *D* has a kernel. All the conditions for the theorem are tight.

© 2006 Elsevier B.V. All rights reserved.

*Keywords:* Kernel; Kernel-perfect digraph; Quasi-transitive digraph

#### **1. Introduction**

For general concepts we refer the reader to [\[4\].](#page-5-0) In the paper we write digraph to mean 1-digraph in the sense of Berge [\[4\].](#page-5-0) In this paper *D* will denote a possibly infinite digraph with  $V(D)$  and  $A(D)$  being the sets of vertices and arcs of *D*, respectively. Often we shall write  $u_1u_2$  instead of  $(u_1, u_2)$ . An arc  $u_1u_2 \in A(D)$  is called asymmetrical (resp. symmetrical) if  $u_2u_1 \notin A(D)$  (resp.  $u_2u_1 \in A(D)$ ). If *S* is a nonempty subset of  $V(D)$  then the subdigraph  $D[S]$ induced by *S* is the digraph with vertex set *S* and whose arcs are those arcs of *D* which join vertices of *S*.

A directed path is a finite or infinite sequence  $(x_1, x_2, \ldots)$  of distinct vertices of *D* such that  $(x_i, x_{i+1}) \in A(D)$  for each *i*. When *D* is infinite and the sequence is infinite we call the directed path an infinite outward path. Let  $S_1$  and  $S_2$  be subsets of  $V(D)$ . A finite directed path  $(x_1, \ldots, x_n)$  will be called an  $S_1S_2$ -directed path whenever  $x_1 \in S_1$  and  $x_2 \in S_2$ , in particular when the directed path is an arc, we will call it an  $S_1S_2$ -arc.

**Definition 1.1.** A set  $I \subseteq V(D)$  is independent if  $A(D[I]) = \emptyset$ . A kernel *N* of *D* is an independent set of vertices such that for each  $z \in V(D) - N$  there exists a *zN*-arc in *D*.

A digraph *D* is called a kernel-perfect digraph when every induced subdigraph of *D* has a kernel.

*E-mail address:* [hgaleana@matem.unam.mx](mailto:hgaleana@matem.unam.mx) (H. Galeana-Sánchez).

<sup>0012-365</sup>X/\$ - see front matter © 2006 Elsevier B.V. All rights reserved. doi:10.1016/j.disc.2006.02.015

The concept of kernel was introduced by Von Neumann and Morgenstern [\[15\]](#page-5-0) in the context of Game Theory. The problem of the existence of a kernel in a given digraph has been studied by several authors in particular by Richardson [16,17], Duchet and Meyniel [\[9\],](#page-5-0) Duchet [7,8], Galeana-Sánchez and Neumann-Lara [\[10\].](#page-5-0)

A digraph *D* is transitive whenever  $(u, v) \in A(D)$  and  $(v, w) \in A(D)$  implies  $(u, w) \in A(D)$ . A digraph is called *quasi-transitive* if whenever  $(u, v) \in A(D)$  and  $(v, w) \in A(D)$ , then  $(u, w) \in A(D)$  or  $(w, u) \in A(D)$ .

Quasi-transitive digraphs were introduced by Ghouilá-Houri [\[12\]](#page-5-0) and have been studied by several authors for example Bang-Jensen and Huang  $[1-3]$ , Huang  $[13]$ , Skrien  $[19]$ . It was proved by Ghouilá-Houri  $[12]$  that an undirected graph can be oriented as a quasi-transitive digraph if and only if it can be oriented as a transitive digraph, namely a comparability graph. More information about comparability graphs can be found in [11,14].

In [\[6\]](#page-5-0) Boros and Gurvich proved that if *G* is a perfect graph then any orientation of *G* in which each complete subdigraph has a kernel is kernel-perfect. It is well known that comparability graphs are perfect graphs (see for example [\[5\]\)](#page-5-0). Meyniel [\[9\]](#page-5-0) observed that if *D* is a digraph such that every directed cycle of length 3 has at least two symmetrical arcs, then each complete subdigraph of *D* has a kernel.

We can conclude the following result.

**Theorem 1.2.** *If D is a finite quasi-transitive digraph such that every directed cycle of length* 3 *has at least two symmetrical arcs*, *then D is a kernel-perfect digraph*.

The result proved in this paper generalizes Theorem 1.2 and the following result of Sands et al. [\[18\].](#page-5-0)

**Theorem 1.3** (*Sands et al. [\[18\]](#page-5-0)*). *Let D be a digraph whose arcs are colored with two colors. If D contains no monochromatic infinite outward path*, *then there exists a set S of vertices of D such that no two vertices of S are connected by a monochromatic directed path and for every vertex x not in S there is a monochromatic directed path from x to a vertex in S*.

We include the following definitions in order to understand Theorem 1.3 in terms of kernels.

We call the digraph *D* an *m*-colored digraph if the arcs of *D* are colored with *m* colors. A directed path is called monochromatic if all of its arcs are colored alike. A kernel by monochromatic paths in an *m*-colored digraph *D* is a set of vertices *N* which satisfies the following two conditions: (i) for every pair of different vertices  $u, v \in N$  there is no monochromatic directed path between them; and (ii) for every vertex  $x \in V(D) - N$  there is a vertex  $y \in N$  such that there is an *xy*-monochromatic directed path.

If *D* is an *m*-colored digraph then the closure of *D*, denoted  $\mathcal{C}(D)$  is the digraph defined as follows:  $V(\mathcal{C}(D))=V(D)$ and  $(u, v) \in A(\mathscr{C}(D))$  iff there exists a uv-monochromatic directed path contained in *D*.

Note that for any *m*-colored *D*, *D* has a kernel by monochromatic paths if and only if  $\mathcal{C}(D)$  has a kernel.

In this terminology Theorem 1.3 asserts that if *D* is a 2-colored digraph, which contains no monochromatic infinite outward path, then  $\mathcal{C}(D)$  has a kernel.

Now it is clear that Theorem 1.3 is equivalent to the following assertion. Let *D* be a digraph;  $D_1$  and  $D_2$  transitive subdigraphs of *D* such that  $D = D_1 \cup D_2$  (recall that  $D_1 \cup D_2$  is defined as follows:  $V(D_1 \cup D_2) = V(D_1) \cup V(D_2)$  and  $A(D_1 \cup D_2) = A(D_1) \cup A(D_2)$  and  $A(D_1) \cap A(D_2) = \emptyset$ . If *D* has no infinite outward path contained in  $D_i$  (i = 1, 2), then *D* has a kernel.

Finally, we will introduce some notation. Two subdigraphs  $D_1$  and  $D_2$  of  $D$  are given (possibly  $A(D_1) \cap A(D_2) \neq \emptyset$ ). For distinct vertices x, y of D,  $x \stackrel{i}{\rightarrow} y$  will mean that the arc  $(x, y) \in A(D_i)$  and  $x \stackrel{i}{\rightarrow} S$  will mean that there exists an arc in  $D_i$  from *x* to a vertex in *S*, the negation of  $x \stackrel{i}{\rightarrow} y$  (resp.  $x \stackrel{i}{\rightarrow} S$ ) will be denoted by  $x \stackrel{i}{\rightarrow} y$  (resp.  $x \stackrel{i}{\rightarrow} S$ ), for  $i = 1, 2$ . When we do not know if the arc is in  $D_1$  or in  $D_2$  we write simply  $x \to y$ ; and  $x \to y$  will mean that  $(x, y) \notin A(D)$ . A directed cycle of length 3 will be called a triangle.

## **2. Kernels in the union of two quasi-transitive digraphs**

The main result of this section is Theorem 2.3. The proof is similar to that in Sands et al. [\[18\].](#page-5-0)

**Lemma 2.1.** Let D be a digraph such that every triangle has at least two symmetrical arcs. If  $D_1$  is a quasi*transitive subdigraph of* D *and*  $(v_1, v_2, \ldots, v_n)$  *is a sequence of vertices of*  $D_1$  *such that*  $(v_i, v_{i+1}) \in A(D_1)$  *and*   $(v_{i+1}, v_i) \notin A(D)$ , then the sequence is an asymmetrical directed path of D contained in  $D_1$ , and for each  $i \in$  ${1, \ldots, n-1}, (v_i, v_j) \in A(D_1)$  *and*  $(v_j, v_i) \notin A(D)$  *for every*  $j \in \{i+1, \ldots, n\}.$ 

**Proof.** We proceed by induction on *n*. The result is obvious for  $n \le 2$ . Assume the result is true for a sequence  $(v_1,\ldots,v_n)$  which satisfies the hypothesis of Lemma 2.1. Consider a sequence  $T = (v_1,\ldots,v_n,v_{n+1})$  such that for each  $i \in \{1,\ldots,n\}$ ,  $(v_i, v_{i+1}) \in A(D_1)$  and  $(v_{i+1}, v_i) \notin A(D)$ . Since  $T' = (v_1,\ldots,v_n)$  satisfies the inductive hypothesis, we have that T' is an asymmetrical directed path contained in  $D_1$  and for each  $i \in \{1, \ldots, n-1\}$   $(v_i, v_j) \in A(D_1)$  and  $(v_j, v_i) \notin A(D)$  for every  $j \in \{i+1, \ldots, n\}$ . So we only need to prove that for each  $i \in \{1, \ldots, n-1\}$ ,  $v_i \neq v_{n+1}$ ,  $(v_i, v_{n+1}) \in A(D_1)$  and  $(v_{n+1}, v_i) \notin A(D)$ .

First assume by contradiction that  $v_{n+1} = v_i$  for some  $i \in \{1, ..., n-1\}$ . It follows from the inductive hypothesis on T' that  $(v_i, v_n) = (v_{n+1}, v_n) \in A(D_1)$  and thus  $(v_{n+1}, v_n) \in A(D)$  contradicting our hypothesis on T. We conclude that *T* is an asymmetrical directed path of *D* contained in  $D_1$ . Now, we have from the inductive hypothesis on  $T'$  that for each  $i \in \{1, ..., n-1\}$ ,  $(v_i, v_n) \in A(D_1)$  and since  $(v_n, v_{n+1}) \in A(D_1)$  and  $D_1$  is a quasi-transitive digraph, we have that  $(v_i, v_{n+1}) \in A(D_1)$  or  $(v_{n+1}, v_i) \in A(D_1)$ . If  $(v_{n+1}, v_i) \in A(D)$  then  $C_3 = (v_i, v_n, v_{n+1}, v_i)$  is a triangle and from the hypothesis on *D*, C<sub>3</sub> has at least two symmetrical arcs which is impossible as  $(v_{n+1}, v_n) \notin A(D)$  (hypothesis on *T*) and  $(v_n, v_i) \notin A(D)$  (inductive hypothesis). Thus  $(v_{n+1}, v_i) \notin A(D)$  and  $(v_i, v_{n+1}) \in A(D_1)$ .  $\Box$ 

Lemma 2.2. Let D be a digraph such that every triangle has at least two symmetrical arcs, and  $D_1$  be a quasi-transitive *subdigraph of D which contains no asymmetrical (in D) infinite outward path. If*  $\emptyset \neq U \subseteq V(D)$  *then there exists*  $x \in U$  such that for all  $y \in U(x, y) \in A(D_1)$  implies  $(y, x) \in A(D)$ .

**Proof.** Suppose by contradiction that for each  $x \in U$ , there exists  $y \in U$  such that  $(x, y) \in A(D_1)$  and  $(y, x) \notin A(D)$ . Consider some  $x_1 \in U$ . Then there exists  $x_2 \in U$  such that  $(x_1, x_2) \in A(D_1)$  and  $(x_2, x_1) \notin A(D)$ . So for each  $n \in \mathbb{N}$ , given  $x_n \in U$ , there exists  $x_{n+1} \in U$  such that  $(x_n, x_{n+1}) \in A(D_1)$  and  $(x_{n+1}, x_n) \notin A(D)$ . It follows from Lemma 2.1 that  $T_{n+1} = (x_1, x_2, \ldots, x_{n+1})$  is an asymmetrical directed path of *D* contained in  $D_1$ . Consider the sequence  $T = (x_n)_{n \in \mathbb{N}}$ ; for each  $n \in \mathbb{N}$ ,  $(x_n, x_{n+1}) \in A(D_1)$ , and for  $n < m$  we have  $\{x_n, x_m\} \subseteq V(T_m)$  and since  $T_m$  is a directed path we obtain  $x_n \neq x_m$ ; hence *T* is an asymmetrical infinite outward path of *D* contained in  $D_1$ , a contradiction.  $\Box$ 

**Theorem 2.3.** *Let D be a digraph such that*  $D = D_1 \cup D_2$  (*possibly*  $A(D_1) \cap A(D_2) \neq \emptyset$ *), where*  $D_i$  *is a quasi-transitive subdigraph of D which contains no asymmetrical* (*in D*) *infinite outward path. If every triangle contained in D has at least two symmetrical arcs*, *then D is a kernel-perfect digraph*.

**Proof.** It suffices to prove that *D* has a kernel, as any induced subdigraph of *D* satisfies the hypothesis of Theorem 2.3. For independent sets *S*, *T* of *D*, we write  $S \le T$  if and only if for each  $s \in S$  there exists  $t \in T$  such that either  $s = t$ 

or  $(s \to t$  and  $t \to s)$ . Note that in particular  $S \subseteq T$  implies  $S \leq T$ .

(1) The collection of all independent sets of vertices of *D* is partially ordered by  $\leq$ .

 $(1.1) \leq$  is reflexive.

This follows from the fact  $S \subseteq S$ .

 $(1.2) \leqslant$  is transitive.

Let S, T and R be independent sets of vertices of D, such that  $S \le T$  and  $T \le R$ , and let  $s \in S$ . Since  $S \le T$  there exists  $t \in T$  such that either  $s = t$  or  $(s \to t$  and  $t \to s$ ) and  $T \le R$  implies that there exists  $r \in R$  such that either  $t = r$  or  $(t \to r \text{ and } r \to t)$ . If  $s = t \text{ or } t = r$ , then  $s = r \text{ or } (s \to r \text{ and } r \to s)$  with  $r \in R$ . So we can assume  $s \neq t$ ,  $t \neq r$ ,  $(s \to t \text{ and } t \to s)$ and  $t\to s$ ) and  $(t\to r$  and  $r\to t)$ . Since  $D_1$  is a quasi-transitive digraph it follows from Lemma 2.1 on the sequence  $(s, t, r)$  that  $(s \stackrel{1}{\rightarrow} r$  and  $r \rightarrow s$ ).

 $(1.3) \leq$  is antisymmetrical.

Let *S* and *T* be independent sets of vertices of *D* such that  $S \le T$  and  $T \le S$ , and let  $s \in S$ . Since  $S \le T$  there exists  $t \in T$  such that either  $s = t$  or  $(s \to t$  and  $t \to s)$ . Suppose that  $s \neq t$ . The fact  $T \leq S$  implies that there exists  $s' \in S$ such that either  $t = s'$  or  $(t \to s'$  and  $s' \to t)$ . When  $t = s'$  we obtain  $s \to s'$  contradicting that *S* is an independent set; so  $t \neq s'$  and  $(t \stackrel{1}{\rightarrow} s'$  and  $s' \nightharpoonup t$ ). Now applying Lemma 2.1 on the sequence  $(s, t, s')$ , we have  $s \stackrel{1}{\rightarrow} s'$  contradicting that *S* is an independent set. We conclude  $s = t$  and consequently  $s \in T$  and  $S \subseteq T$ . Analogously it can be proved  $T \subseteq S$ .

Let F be the family of all nonempty independent sets S of vertices of D such that  $S \to y$  implies  $y \to S$  for all vertices *y* of *D*.

**(2)** ( $\mathcal{F}, \leq$ ) has maximal elements.

 $(2.1) \mathscr{F} \neq \emptyset$ .

Since  $D_2$  is a quasi-transitive digraph which contains no asymmetrical infinite outward path, it follows from Lemma 2.2 (taking  $U = V(D)$  and  $D_2$  instead of  $D_1$ ) that there exists a vertex  $x \in V(D)$  such that  $x \stackrel{2}{\rightarrow} y$  implies  $y \rightarrow x$ , for all vertices *y* of *D*, so  $\{x\} \in \mathcal{F}$ .

**(2.2)** Every chain in  $(\mathcal{F}, \leq)$  is upper bounded.

Let C be a chain in  $(\mathscr{F}, \leq)$ , and define  $S^{\infty} = \{s \in \bigcup_{S \in \mathscr{C}} S | \text{ there exists } S \in \mathscr{C} \text{ such that } s \in T \text{ whenever } T \in \mathscr{C} \text{ such that } S \in \mathscr{C} \text{ such$ and  $T \ge S$ . (S<sup>∞</sup> consists of all vertices of *D* that belong to every member of *C* from some point on.)

We will prove that  $S^{\infty}$  is an upper bound of  $\mathscr{C}$ .

**(2.2.1)**  $S^{\infty} \neq \emptyset$ , and for each  $S \in \mathscr{C}$ ,  $S^{\infty} \geq S$ .

Let  $S \in \mathscr{C}$  and  $t_0 \in S$ . We will prove that there exists  $t \in S^{\infty}$  such that either  $t_0 = t$  or  $(t_0 \to t$  and  $t \to t_0)$ . If  $t_0 \in S^{\infty}$ we are done. So assume  $t_0 \notin S^{\infty}$ . We proceed by contradiction; suppose that if  $t \in V(D)$  with  $(t_0 \stackrel{1}{\to} t$  and  $t \to t_0$ ), then  $t \notin S^{\infty}$ . Take  $T_0 = S$ . Since  $t_0 \notin S^{\infty}$  we have that there exists  $T_1 \in \mathscr{C}$ ,  $T_1 \geq T_0$  such that  $t_0 \notin T_1$ . Hence there exists  $t_1 \in T_1$  such that  $t_0 \to t_1$  and  $t_1 \to t_0$ . And our assumption implies  $t_1 \notin S^{\infty}$ . The fact  $t_1 \notin S^{\infty}$  implies  $t_1 \notin T_2$  for some  $T_2 \in \mathscr{C}, T_2 \ge T_1$ . Hence there exists  $t_2 \in T_2$  such that  $t_1 \stackrel{1}{\to} t_2$  and  $t_2 \nightharpoonup t_1$ . Since  $D_1$  is a quasi-transitive digraph, it follows from Lemma 2.1 on the sequence  $\tau_2 = (t_0, t_1, t_2)$  that  $\tau_2$  is an asymmetrical directed path of *D* contained in  $D_1$ ,  $(t_0 \to t_2 \text{ and } t_2 \to t_0)$ ; and  $t_2 \notin S^{\infty}$ . We may continue this way and we obtain, for each  $n \in \mathbb{N}$ ,  $T_n \in \mathscr{C}$ ,  $t_n \in T_n$ ,  $(t_0 \to t_n)$ and  $t_n \to t_0$ ) and  $t_n \notin S^\infty$ . Hence there exists  $T_{n+1} \in \mathscr{C}$  such that  $T_{n+1} \geq T_n$  and  $t_n \notin T_{n+1}$ . So there exists  $t_{n+1} \in T_{n+1}$ . with  $(t_n \stackrel{1}{\rightarrow} t_{n+1}$  and  $t_{n+1} \rightarrow t_n$ ).

Since  $D_1$  is a quasi-transitive digraph, and  $(t_n \to t_{n+1}$  and  $t_{n+1} \to t_n)$  for each  $n \in \mathbb{N}$ , it follows from Lemma 2.1 (on the sequence  $\tau_{n+1} = (t_0, t_1, \ldots, t_{n+1})$  that  $\tau_{n+1}$  is an asymmetrical directed path contained in  $D_1$  and in particular  $(t_0 \to t_{n+1}$  and  $t_{n+1} \to t_0$ ). Our assumption implies  $t_{n+1} \notin S^{\infty}$ . Now consider the sequence  $\tau = (t_n)_{n \in \mathbb{N}}$ . For each  $n \in \mathbb{N}$ we have  $(t_n \to t_{n+1}$  and  $t_{n+1} \to t_n)$ , and observe that for  $n < m$ ,  $\{t_n, t_m\} \subseteq V(\tau_m)$ , and since  $\tau_m$  is a directed path we have  $t_n \neq t_m$ . Hence  $\tau$  is an asymmetrical infinite outward path contained in  $D_1$ , a contradiction. We conclude that there exists  $t \in S^{\infty}$  such that  $(t_0 \stackrel{1}{\to} t$  and  $t \to t_0)$ .

**(2.2.2)**  $S^{\infty}$  is an independent set.

Let  $s_1, s_2 \in S^{\infty}$  and suppose without loss of generality that  $S_1, S_2 \in \mathscr{C}$  are such that  $s_1 \in S_1$ ,  $s_2 \in S_2$ ,  $S_1 \le s_2$ , since  $s_1 \in S^{\infty}$  we have  $s_1 \in S$  whenever  $S \in \mathscr{C}$  and  $S \ge S_1$ , so  $s_1 \in S_2$ , and since  $S_2$  is independent, there is no arc in *D* between  $s_1$  and  $s_2$ .

 $(2.2.3)$   $S^{\infty} \in \mathscr{F}$ .

Suppose  $S^{\infty} \xrightarrow{2} y$  with  $y \in V(D)$ , so there exists  $s \in S^{\infty}$  with  $s \xrightarrow{2} y$ . Let  $S \in \mathscr{C}$  such that  $s \in T$  for all  $T \in \mathscr{C}$ ,  $T \ge S$ . Since  $S \in \mathcal{F}$  we have  $y \to S$ , so there exists  $s' \in S$  with  $y \to s'$ . When  $s' \in S^{\infty}$  we are done. When  $s' \notin S^{\infty}$  we analyze the two possibilities;  $y \stackrel{1}{\rightarrow} s'$  or  $y \stackrel{2}{\rightarrow} s'$ . First suppose  $y \stackrel{2}{\rightarrow} s'$ . Since  $s \stackrel{2}{\rightarrow} y$  and  $D_2$  is a quasi-transitive digraph it follows that  $s \stackrel{2}{\rightarrow} s'$  or  $s' \stackrel{2}{\rightarrow} s$  which is impossible as *S* is an independent set and  $\{s, s'\} \subseteq S$ . Now suppose  $y \stackrel{1}{\rightarrow} s'$ . Since  $s' \in S$ ,  $S \le S^{\infty}$  by (2.2.1) and  $s' \notin S^{\infty}$ , there exists  $t \in S^{\infty}$  such that  $s' \stackrel{1}{\to} t$  and  $t \to s'$ . So we obtain  $y \stackrel{1}{\to} t$  or  $t \to y$  (as  $y \to s'$ ,  $s' \to t$  and  $D_1$  is a quasi-transitive digraph). If  $y \to t$  then  $y \to s^{\infty}$  and we are done. If  $t \to y$  then we obtain the triangle  $(y, s', t, y)$  and it follows from the hypothesis that it has two symmetrical arcs and since  $t \rightarrow s'$ we have  $s' \rightarrow y$  and  $y \rightarrow t$ , so  $y \rightarrow S^{\infty}$ .

We have proven that any chain in  $\mathcal F$  has an upper bound in  $\mathcal F$ , and so by Zorn's Lemma,  $(\mathcal F, \leqslant)$  contains maximal elements. Let *S* be a maximal element of  $(\mathscr{F}, \leq)$ .

**(3)** *S* is a kernel of *D*.

Since  $S \in \mathcal{F}$ , *S* is an independent set of vertices of *D*.

**(3.1)** For each  $x \in (V(D) - S)$  there exists an *xS*-arc.

Suppose by contradiction there exists  $x \in (V(D) - S)$  such that  $x \rightarrow S$ .

**(3.1.1)** There exists a vertex  $x_0 \in V(D)$  such that  $x_0 \to S$  and  $x_0$  satisfies:  $x_0 \stackrel{2}{\to} y$  and  $y \to S$  imply  $y \to x_0$  for all vertices  $y \in V(D)$ . Let  $U = \{z \in V(D_2) - S | z \rightarrow S \}$ . When  $U \neq \emptyset$ , it follows from Lemma 2.2 (applied on  $D_2$  and *U*) that there exists  $x_0$  with the required properties. When  $U = \emptyset$  it follows from our assumption that  $z \rightarrow S$ , for some vertex *z* in  $V(D_1) - (S \cup V(D_2))$ , and we take  $x_0$  to be any such vertex.

Note that the choice of  $x_0$  implies  $x_0 \to S$  and since  $S \in \mathcal{F}$ , we also have  $S \stackrel{2}{\to} x_0$ . Let  $T = \{s \in S | s \stackrel{1}{\to} x_0\}$ , it follows from above that  $T \cup \{x_0\}$  is an independent set of vertices of *D*.

 $(3.1.2)$   $T \cup \{x_0\} \in \mathscr{F}$ .

Suppose  $T \cup \{x_0\} \stackrel{2}{\rightarrow} y$  and  $y \nrightarrow T$ . We will prove  $y \rightarrow x_0$ . First we make the following observation.

(3.1.2.1) If 
$$
y \stackrel{1}{\rightarrow} (S - T)
$$
 then  $y \rightarrow x_0$ .

Let  $s \in (S-T)$  such that  $y \stackrel{1}{\to} s$ . Since  $s \in (S-T)$  we have  $s \stackrel{1}{\to} x_0$ . Now the fact that  $D_1$  is a quasi-transitive digraph implies  $y \to x_0$  or  $x_0 \to y$ . If  $x_0 \to y$  then  $(y, s, x_0, y)$  is a triangle which by the hypothesis has two symmetrical arcs, and since  $x_0 \rightarrow s$  it follows that  $y \rightarrow x_0$ .

We proceed to prove  $(3.1.2)$  by considering the two following cases:

*Case* a:  $T \stackrel{2}{\rightarrow} y$ .

Since  $T \subset S$  we have  $S \stackrel{2}{\to} y$  and the fact  $S \in \mathcal{F}$  implies  $y \to S$ . So  $y \to (S - T)$  (as we are assuming  $y \to T$ ).

When  $y \rightarrow (S - T)$  it follows from (3.1.2.1) that  $y \rightarrow x_0$ .

When  $y \stackrel{2}{\rightarrow} (S-T)$ , since we have  $T \stackrel{2}{\rightarrow} y$  and  $D_2$  is a quasi-transitive digraph, we obtain  $T \stackrel{2}{\rightarrow} (S-T)$  or  $(S-T) \stackrel{2}{\rightarrow} T$ and this is impossible as  $T \subseteq S$  and *S* is an independent set.

*Case* b:  $x_0 \stackrel{2}{\rightarrow} y$ .

We consider two possible subcases:

*Case* b.1:  $v \rightarrow S$ .

Since  $x_0 \stackrel{2}{\rightarrow} y$  and  $y \rightarrow S$ , the choice of  $x_0$  (see (3.1.1)) implies  $y \rightarrow x_0$ .

*Case* b.2:  $y \rightarrow S$ .

In this case we have  $y \rightarrow (S - T)$  (as we are assuming  $y \rightarrow T$ ).

When  $y \stackrel{2}{\rightarrow} (S - T)$ , since  $x_0 \stackrel{2}{\rightarrow} y$  and  $D_2$  is a quasi-transitive digraph, we have  $x_0 \stackrel{2}{\rightarrow} (S - T)$  or  $(S - T) \stackrel{2}{\rightarrow} x_0$ . Now recalling  $x_0 \to S$ , we obtain  $(S - T) \stackrel{2}{\to} x_0$  and since  $S \in \mathcal{F}$  it follows  $x_0 \to S$  which is impossible.

When  $y \rightarrow (S - T)$  it follows from (3.1.2.1) that  $y \rightarrow x_0$ . **(3.1.3)**  $S < T \cup \{x_0\}.$ 

For  $s \in (S-T)$  we have  $s \stackrel{1}{\to} x_0$  and we have noted  $x_0 \nrightarrow S$ ; hence  $S \leq T \cup \{x_0\}$ . Moreover since  $x_0 \notin S$  (by construction in (3.1.1)) we have  $S < T \cup \{x_0\}$ .

Clearly propositions (3.1.2) and (3.1.3) contradict that *S* is a maximal element of  $(\mathscr{F}, \leqslant)$ .  $\Box$ 

**Remark 2.4.** The condition that  $D_i$  has no infinite outward path in Theorem 2.3 is necessary.

Consider the following digraph D' with  $V(D') = \{u_n | n \in \mathbb{N}\}\$  and  $A(D') = \{(u_n, u_m) | n, m \in \mathbb{N}\}\$  and  $n < m\}$ ,  $D_1 = D', D_2 = D'$  and  $D = D_1 \cup D_2$ .

**Remark 2.5.** The following digraph *D* is the union of two quasi-transitive finite digraphs; each triangle in *D* has at least one symmetrical arc and *D* has no kernel.

 $V(D_1) = \{u_0, u_1, u_2, u_3\},\$  $V(D_2) = V(D_1) \cup \{w\},\$  $A(D_1) = \{(u_i, u_{i+1}) | i \in \{0, 1, 2, 3\} \pmod{4}\} \cup \{(u_0, u_2), (u_2, u_0), (u_1, u_3), (u_3, u_0)\},\$  $A(D_2) = \{(w, u_i)|i \in \{0, 1, 2, 3\}\},\$  $D = D_1 \cup D_2$ .

**Remark 2.6.** Clearly  $\vec{C}_5$  the directed cycle of length 5 is the union of two finite digraphs,  $\vec{C}_5$  has no triangle and  $\vec{C}_5$ has no kernel.

We conclude that the conditions on Theorem 2.3 are tight.

## <span id="page-5-0"></span>**Acknowledgments**

The authors wish to thank the referees for many suggestions which improved the rewriting of this paper.

### **References**

- [1] J. Bang-Jensen, J. Huang, Quasi-transitive digraphs, J. Graph Theory 20 (2) (1995) 141–161.
- [2] J. Bang-Jensen, J. Huang, Kings in quasi-transitive digraphs, Discrete Math. 185 (1–3) (1998) 19–27.
- [3] J. Bang-Jensen, J. Huang, Strongly connected spanning subdigraphs with the minimum number of arcs in quasi-transitive digraphs, SIAM J. Discrete Math. 16 (2) (2003) 335–343.
- [4] C. Berge, Graphs, North-Holland Mathematical Library, vol. 6, North-Holland, Amsterdam, 1985.
- [5] C. Berge, V. Chvatal, Editors, Topics on perfect graphs, Ann. Discrete Math. vol. 21, North-Holland, Amsterdam, 1985. ´
- [6] E. Boros, V. Gurvich, Perfect graphs are kernel solvable, Discrete Math. 159 (1996) 35–55.
- [7] P. Duchet, Graphes Noyau-Parfaits, Ann. Discrete Math. 9 (1980) 93–101.
- [8] P. Duchet, A sufficient condition for a digraph to be kernel-perfect, J. Graph Theory 11 (1) (1987) 81–85.
- [9] P. Duchet, H. Meyniel, A note on kernel–critical graphs, Discrete Math. 33 (1981) 103–105.
- [10] H. Galeana-Sánchez, V. Neumann-Lara, On kernels and semikernels of digraphs, Discrete Math. 48 (1984) 67–76.
- [11] T. Gallai, Transitiv orientierbare graphen, Acta Math. Acad. Sci. Hungar. 18 (1967) 25–66.
- [12] A. Ghouilá-Houri, Caractérisation des graphes non orientés dont on peut orienter les arrêtes de maniere à obtenir le graphe d' un relation d'ordre, C.R. Acad. Sci. Paris 254 (1962) 1370–1371.
- [13] J. Huang, Tournament-like oriented graphs, Ph.D. Thesis, Simon Fraser University, 1992.
- [14] D. Kelly, Comparability graphs, in: I. Rival (Ed.), Graphs and Order, Nato ASI Series C 147, D. Reidel, Dordrecht, 1985, pp. 3–40.
- [15] J. Von Neumann, O. Morgenstern, Theory of Games and Economic Behavior, Princeton University Press, Princeton, NJ, 1944.
- [16] M. Richardson, Solutions of irreflexive relations, Ann. Math. 58 (2) (1953) 573-580.
- [17] M. Richardson, Extensions theorems for solutions of irreflexive relations, Proc. Natl. Acad. Sci. USA 39 (1953) 649–651.
- [18] B. Sands, N. Sauer, R. Woodrow, On monochromatic paths in edge-colored digraphs, J. Combin. Theory B 33 (1982) 271–275.
- [19] D.J. Skrien, A relationship between triangulated graphs, comparability graphs, proper interval graphs proper circulant graphs and nested interval graphs, J. Graph Theory 6 (1980) 309–316.