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## Optimal orientations of products of paths and cycles

K.M. Koh\*, E.G. Tay

*Department of Mathematics, National University of Singapore, Lower Kent Ridge Road,  
Singapore 119260*

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### Abstract

For a graph  $G$ , let  $\mathcal{O}(G)$  be the family of strong orientations of  $G$ ,  $\mathbf{d}(G) = \min\{d(D) \mid D \in \mathcal{O}(G)\}$  and  $\rho(G) = \mathbf{d}(G) - d(G)$ , where  $d(G)$  and  $d(D)$  are the diameters of  $G$  and  $D$  respectively. In this paper we show that  $\rho(G) = 0$  if  $G$  is a cartesian product of (1) paths, and (2) paths and cycles, which satisfy some mild conditions.

*Keywords:* Path; Cycle; Bipartite graph; Diameter; Strong orientation

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### 1. Introduction

Let  $G$  (resp.,  $D$ ) be a graph (resp., digraph) with vertex set  $V(G)$  (resp.,  $V(D)$ ) and edge set  $E(G)$  (resp.,  $E(D)$ ). For  $v \in V(G)$ , the *eccentricity*  $e(v)$  of  $v$  is defined as  $e(v) = \max\{d(v, x) \mid x \in V(G)\}$ , where  $d(v, x)$  denotes the distance from  $v$  to  $x$ . The notion  $e(v)$  in  $D$  is similarly defined. The *diameter* of  $G$  (resp.,  $D$ ), denoted by  $d(G)$  (resp.,  $d(D)$ ), is defined as  $d(G) = \max\{e(v) \mid v \in V(G)\}$  (resp.,  $d(D) = \max\{e(v) \mid v \in V(D)\}$ ).

An *orientation* of a graph  $G$  is a digraph obtained from  $G$  by assigning to each edge in  $G$  a direction. An orientation  $D$  of  $G$  is *strong* if every two vertices in  $D$  are mutually reachable in  $D$ . An edge  $e$  in a connected graph  $G$  is a *bridge* if  $G - e$  is disconnected. Robbins' celebrated one-way street theorem [15] states that a *connected graph  $G$  has a strong orientation if and only if no edge of  $G$  is a bridge*. As a possible way of extending Robbins' theorem, Boesch and Tindell [1] introduced the notion  $\rho(G)$  given below. For a connected graph  $G$  containing no bridges, let  $\mathcal{O}(G)$  be the family

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\* Corresponding author. E-mail: [matkohkm@leonis.nus.sg](mailto:matkohkm@leonis.nus.sg).

of strong orientations of  $G$ . Define

$$d(G) = \min\{d(D) \mid D \in \mathcal{D}(G)\} \quad \text{and} \quad \rho(G) = d(G) - d(G).$$

The problem of evaluating  $\rho(G)$  for an arbitrary connected graph  $G$  is very difficult. As a matter of fact, Chvátal and Thomassen [2] showed that the problem of deciding whether a graph admits an orientation of diameter two is NP-hard.

On the other hand, the parameter  $\rho(G)$  has been studied in various classes of graphs including complete graphs [1, 11, 14], complete bipartite graphs [1, 3, 20], complete  $k$ -partite ( $k \geq 3$ ) graphs [4, 6, 7, 13], and  $n$ -cubes [12, 20]. Let  $G \times H$  denote the cartesian product of two graphs  $G$  and  $H$  (see Section 2 for the definition), and  $P_r, C_r$  and  $K_r$ , respectively, the path, cycle and complete graph of order  $r$ . Roberts and Xu [16–19], and independently Koh and Tan [5], evaluated the quantity  $\rho(P_m \times P_n)$ . Very recently, Koh and Tay have further determined the quantities  $\rho(C_{2m} \times P_k)$  [8],  $\rho(K_m \times P_k)$ ,  $\rho(K_m \times C_{2k+1})$  and  $\rho(K_m \times K_n)$  [9] and  $\rho(C_{2m} \times K_n)$  [10]. In this paper, we shall evaluate  $\rho(G_1 \times G_2 \times \dots \times G_m)$ , where  $m \geq 2$  and  $\{G_i \mid 1 \leq i \leq m\}$  is a combination of paths and cycles.

## 2. Cartesian product of paths

The cartesian product of a family of graphs  $G_1, G_2, \dots, G_n$ , denoted by  $G_1, G_2 \times \dots \times G_n$  or  $\prod_{i=1}^n G_i$ , where  $n \geq 2$ , is the graph  $G$  having  $V(G) = V(G_1) \times V(G_2) \times \dots \times V(G_n)$  and two vertices  $(u_1, u_2, \dots, u_n)$  and  $(v_1, v_2, \dots, v_n)$  are adjacent if and only if there exists  $r \in \{1, 2, \dots, n\}$  such that  $u_r v_r \in E(G_r)$  and  $u_i = v_i$  for all  $i = 1, 2, \dots, n$  with  $i \neq r$ . In this section, we shall evaluate  $\rho(G)$ , where  $G$  is of the form  $\prod_{i=1}^n P_{k_i}$  with  $n \geq 2$  and  $k_i \geq 2$  for each  $i = 1, 2, \dots, n$ . For convenience, the vertices in the graph are labelled  $(x_1, x_2, \dots, x_n)$ , where  $1 \leq x_i \leq k_i$  for each  $i = 1, 2, \dots, n$ , such that the vertices  $(a_1, a_2, \dots, a_n)$  and  $(b_1, b_2, \dots, b_n)$  are adjacent iff  $|a_r - b_r| = 1$  for exactly one  $r \in \{1, 2, \dots, n\}$ , and  $a_i = b_i$  for all  $i$  with  $i \neq r$ .

Let  $D$  be a digraph. A dipath (resp., dicycle) in  $D$  is simply called a path (resp., cycle) in  $D$ . A path from  $u$  to  $v$  in  $D$  is simply called a  $u$ - $v$  path. For  $X \subseteq V(D)$ , the subdigraph of  $D$  induced by  $X$  is denoted by  $D[X]$ . For  $x, y \in V(D)$  and  $A \subseteq V(D)$ , we write ' $x \rightarrow y$ ' if  $x$  is adjacent to  $y$  in  $D$ , and write ' $x \rightarrow A$ ' (resp., ' $A \rightarrow y$ ') if  $x \rightarrow y$  for each  $y \in A$  (resp., for each  $x \in A$ ).

Our first main result is as follows:

**Theorem 1.**  $\rho(\prod_{i=1}^n P_{k_i}) = 0$ , where  $n \geq 2$ ,  $k_1 \geq 3$ ,  $k_2 \geq 6$  with  $(k_1, k_2) \neq (3, 6)$ .

Let

$$G_n = \overbrace{P_2 \times P_2 \times \dots \times P_2}^n$$

(i.e., the  $n$ -cube). In proving that  $\rho(G_n) = 0$  for  $n \geq 4$ , McCanna [12] made use of the following subtle observation due to C. Thomassen.

**Lemma 1.** *If a bipartite graph  $G$  admits an orientation of diameter at most  $k$ , where  $k \geq 3$ , such that every vertex is in a cycle of length at most  $k$ , then the graph  $G \times P_2$  admits an orientation of diameter at most  $k + 1$  such that every vertex is in a cycle of length at most  $k$ .*

We shall now extend Thomassen’s observation from  $P_2$  to  $\prod_{i=1}^n P_{k_i}$ , and shall make use of the extension to prove some of our main results in this paper.

**Lemma 2.** *If a bipartite graph  $G$  admits an orientation of diameter at most  $k$ , where  $k \geq 3$ , such that every vertex is in a cycle of length at most  $k$ , then the graph  $G \times \prod_{i=1}^n P_{k_i}$ , where  $n \geq 1$ , admits an orientation of diameter at most  $k - n + \sum_{i=1}^n k_i$  such that every vertex is in a cycle of length at most  $k$ .*

**Proof.** Let  $V_1$  and  $V_2$  be the partite sets of  $G$ . Let  $F \in \mathcal{O}(G)$  with  $d(F) \leq k$  such that every vertex is in a cycle of length at most  $k$  in  $F$ . We shall now orient  $G \times \prod_{i=1}^n P_{k_i}$  inductively as follows:

- (i) In  $G \times P_{k_i}$ , for  $1 \leq i \leq k_1 - 1$ , orient  $(x, i) \rightarrow (x, i + 1)$  iff  $x \in V_1$ ; and for  $1 \leq i \leq k_1$ , orient  $(x, i) \rightarrow (y, i)$  iff  $xy \in E(F)$ .
- (ii) Suppose  $G \times \prod_{i=1}^r P_{k_i}$ , where  $1 \leq r \leq n - 1$ , has been oriented. Orient  $G \times \prod_{i=1}^{r+1} P_{k_i}$  so that the orientation of  $G \times \prod_{i=1}^r P_{k_i} \times \{j\}$  is isomorphic to that of  $G \times \prod_{i=1}^r P_{k_i}$  for each  $j = 1, 2, \dots, k_{r+1}$ , and for  $1 \leq i \leq k_{r+1} - 1$ , orient  $(x, a_1, a_2, \dots, a_r, i) \rightarrow (x, a_1, a_2, \dots, a_r, i + 1)$  iff  $x \in V_1$ .

Let  $F^*$  be the resulting orientation of  $G \times \prod_{i=1}^n P_{k_i}$ .

**Claim.**  $e(u) \leq k - n + \sum_{i=1}^n k_i$  for each vertex  $u$  in  $F^*$ .

Let  $u = (x, a_1, a_2, \dots, a_n)$  and assume that  $x \in V_1$ , say. Take an arbitrary vertex  $v = (y, b_1, b_2, \dots, b_n)$  in  $F^*$ . As the cartesian product is commutative, we may assume that  $a_i \leq b_i$  for  $1 \leq i \leq m$  and  $a_i > b_i$  for  $m + 1 \leq i \leq n$ , where  $m \leq n$ .

(1) Let  $w = (x, b_1, b_2, \dots, b_m, a_{m+1}, \dots, a_n)$ . Observe that there is a  $u-w$  path of length at most  $\sum_{i=1}^m k_i - m$  in  $F^*$ .

(2) If  $x \neq y$ , let  $w' = (x', b_1, b_2, \dots, b_m, a_{m+1}, \dots, a_n)$ , where  $x'$  is adjacent from  $x$  in an  $x-y$  path of length at most  $k$  in  $F$ . Then  $w \rightarrow w'$  in  $F^*$ . (Note that  $x' \in V_2$ .) If  $x = y$ , take a cycle of length at most  $k$  containing  $x$  in  $F$ .

(3) There is a path of length at most  $\sum_{i=m+1}^n k_i - (n - m)$  from  $w'$  to  $(x', b_1, b_2, \dots, b_n)$  in  $F^*$ .

(4) There is a path of length at most  $k - 1$  from  $(x', b_1, b_2, \dots, b_n)$  to  $v$  in  $F^*$ .

Combining (1)–(4),  $d(u, v) \leq \sum_{i=1}^m k_i - m + 1 + \sum_{i=m+1}^n k_i - (n - m) + k - 1 = k - n + \sum_{i=1}^n k_i$ . This proves the claim.

Thus  $d(F^*) \leq k - n + \sum_{i=1}^n k_i$ . The second part of Lemma 2 is obvious as each vertex in  $F^*$  is contained in a cycle of length at most  $k$  in  $F$ .  $\square$

We need also the following result.

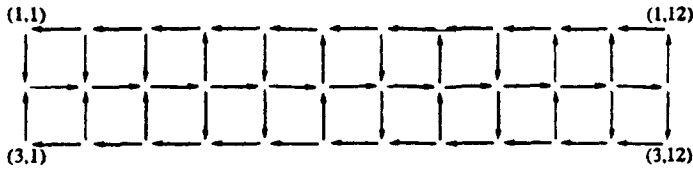


Fig. 1. Orientation of  $P_3 \times P_{12}$ .

**Lemma 3.** For  $m \geq 3, n \geq 6$  with  $(m,n) \neq (3,6)$ , there exists  $F \in \mathcal{D}(P_m \times P_n)$  such that

- (i)  $d(F) = d(P_m \times P_n) = m + n - 2$  and
- (ii) every vertex in  $P_m \times P_n$  is in a cycle of length at most  $m + n - 2$  in  $F$ .

**Note.** (1)  $d(P_m \times P_n) = m + n - 2$  for all  $m, n \geq 1$ .

(2) It was shown in [5] that  $d(P_3 \times P_6) = 8 (= m + n - 1)$ .

**Proof of Lemma 3.** Part (i) (except some isolated cases) was first obtained by Roberts and Xu [16–19]. Here, we shall use the orientations of  $P_m \times P_n$  introduced by Koh and Tan [5] to prove part (ii). Following [5], we have seven cases to consider.

*Case A:*  $m = 3$  and  $n \equiv 0 \pmod{2}$  with  $n \geq 8$ . Define  $F \in \mathcal{D}(P_m \times P_n)$  as follows (see Fig. 1):

- (1) For  $i = 1, 3$  and  $j = 1, 2, \dots, n - 1$ , orient  $(i, j + 1) \rightarrow (i, j)$ ;
- (2) For  $j = 1, 2, \dots, n - 1$ , orient  $(2, j) \rightarrow (2, j + 1)$ ;
- (3) For  $j = 1, 2, 3$ , orient  $\{(1, j), (3, j)\} \rightarrow (2, j)$ ;
- (4) Orient  $(2, 4) \rightarrow \{(1, 4), (3, 4)\}$ ;
- (5) Orient  $(2, n) \rightarrow \{(1, n), (3, n)\}$  and  $(2, n - 1) \rightarrow \{(1, n - 1), (3, n - 1)\}$ ;
- (6) For  $j = 5, 6, \dots, n - 2$ , orient  $(3, j) \rightarrow (2, j) \rightarrow (1, j)$  if  $j \equiv 0 \pmod{2}$ ; and  $(1, j) \rightarrow (2, j) \rightarrow (3, j)$  if  $j \equiv 1 \pmod{2}$ .

Note that  $d(F) = m + n - 2$ . Now, consider the following cycles (see also Fig. 1):

- $(A_1)$   $(1, 1)(2, 1)(2, 2)(2, 3)(2, 4)(1, 4)(1, 3)(1, 2)(1, 1)$ ,
- $(A_2)$   $(3, 1)(2, 1)(2, 2)(2, 3)(2, 4)(3, 4)(3, 3)(3, 2)(3, 1)$ ,
- $(A_3)$   $(3, 5)(3, 4)(3, 3)(2, 3)(2, 4)(2, 5)(3, 5)$ ,
- $(A_4)$   $(1, n)(1, n - 1)(1, n - 2)(1, n - 3)(2, n - 3)(2, n - 2)(2, n - 1)(2, n)(1, n)$ ,
- $(A_5)$   $(3, n)(3, n - 1)(3, n - 2)(2, n - 2)(2, n - 1)(2, n)(3, n)$ .

It can be checked that each of the above cycles is of length at most  $m + n - 2$ , and that the cycles cover vertices  $(3, 5), (1, n - 3), (2, n - 3)$  and  $(i, j)$ , where  $i = 1, 2, 3$  and  $j = 1, 2, 3, 4, n - 2, n - 1, n$ . On the other hand, each of the remaining vertices lies in a cycle of length 4 in  $F$ .

*Case B:*  $m = 3$  and  $n = 7$ . Define  $F \in \mathcal{D}(P_3 \times P_7)$  as shown in Fig. 2. It can be checked that  $d(F) = 8 = m + n - 2$  and that (ii) is satisfied as shown in Fig. 2.

*Case C:*  $m = 3$  and  $n \equiv 1 \pmod{2}$  with  $n \geq 9$ . Define  $F \in \mathcal{D}(P_m \times P_n)$  as follows (see Fig. 3):

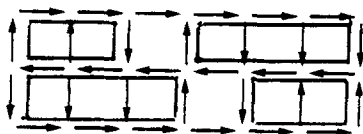


Fig. 2. Orientation of  $P_3 \times P_7$ .

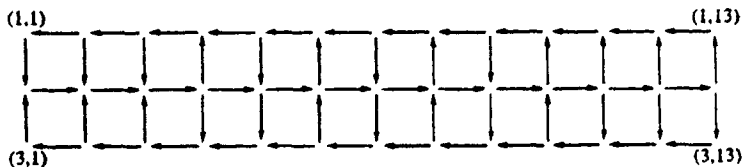


Fig. 3. Orientation of  $P_3 \times P_{13}$ .

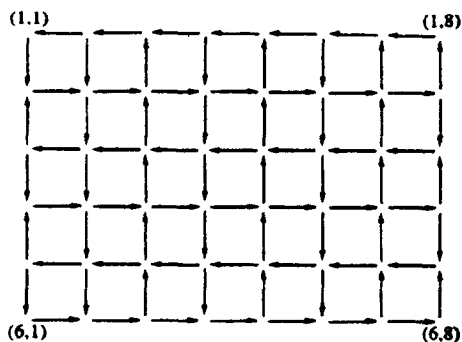


Fig. 4. Orientation of  $P_6 \times P_8$ .

(1)  $F[P_m \times P_{n-1}]$  is identical with the orientation in Case A;

(2) Orient  $(2, n) \rightarrow \{(1, n), (3, n)\}$ ;

(3) Orient  $(1, n) \rightarrow (1, n-1)$ ,  $(2, n-1) \rightarrow (2, n)$ , and  $(3, n) \rightarrow (3, n-1)$ .

Note that  $d(F) = m + n - 2$ . Now, consider the following cycles (see also Fig. 3):

( $C_1$ )  $(1, 1)(2, 1)(2, 2)(2, 3)(2, 4)(1, 4)(1, 3)(1, 2)(1, 1)$ ,

( $C_2$ )  $(3, 1)(2, 1)(2, 2)(2, 3)(2, 4)(2, 5)(3, 5)(3, 4)(3, 3)(3, 2)(3, 1)$ ,

( $C_3$ )  $(1, n)(1, n-1)(1, n-2)(1, n-3)(1, n-4)(2, n-4)(2, n-3)(2, n-2)(2, n-1)(2, n)(1, n)$ ,

( $C_4$ )  $(3, n)(3, n-1)(3, n-2)(3, n-3)(2, n-3)(2, n-2)(2, n-1)(2, n)(3, n)$ .

Each of these cycles is of length at most  $m + n - 2$  and they cover vertices  $(3, 5)$ ,  $(1, n-4)$ ,  $(2, n-4)$  and  $(i, j)$ , where  $i = 1, 2, 3$  and  $j = 1, 2, 3, 4, n-3, n-2, n-1, n$ . On the other hand, each of the remaining vertices lies in a cycle of length 4 in  $F$ .

Case D:  $m \equiv n \equiv 0 \pmod{2}$  with  $m \geq 4$  and  $n \geq 6$ . Define  $F \in \mathcal{Q}(P_m \times P_n)$  as follows (see Fig. 4):

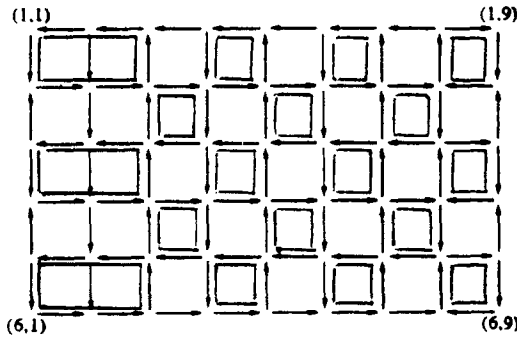


Fig. 5. Orientation of  $P_6 \times P_9$ .

(1) For  $i = 1, 2, \dots, m$  and  $j = 1, 2, \dots, n - 1$ , orient

$$\begin{aligned} (i, j) &\rightarrow (i, j + 1) && \text{if } i \equiv 0 \pmod{2}, \\ (i, j + 1) &\rightarrow (i, j) && \text{if } i \equiv 1 \pmod{2}; \end{aligned}$$

(2) For  $i = 1, 2, \dots, m - 1$  and  $j = 2, 3, \dots, n - 1$ , orient

$$\begin{aligned} (i, j) &\rightarrow (i + 1, j) && \text{if } j \equiv 0 \pmod{2}, \\ (i + 1, j) &\rightarrow (i, j) && \text{if } j \equiv 1 \pmod{2}; \end{aligned}$$

(3) Orient  $(1, 1) \rightarrow (2, 1)$  and  $(i, 1) \rightarrow \{(i - 1, 1), (i + 1, 1)\}$  for each  $i = 3, 5, \dots, m - 1$ ;

(4) Orient  $(i, n) \rightarrow \{(i - 1, n), (i + 1, n)\}$  for each  $i = 2, 4, \dots, m - 2$ ; and

(5) Orient  $(m, n) \rightarrow (m - 1, n)$ .

Note that  $d(F) = m + n - 2$ . Now, consider the following cycles:

( $D_1$ ) For  $i = 1, 3, \dots, m - 1$ ,

$$(i, 1)(i + 1, 1)(i + 1, 2)(i + 1, 3)(i, 3)(i, 2)(i, 1);$$

( $D_2$ ) For  $i = 1, 3, \dots, m - 1$ ,

$$(i, n)(i, n - 1)(i, n - 2)(i + 1, n - 2)(i + 1, n - 1)(i + 1, n)(i, n).$$

Each of these cycles is of length not exceeding  $m + n - 2$ , and the cycles cover vertices  $(i, j)$ , where  $i = 1, 2, \dots, m$  and  $j = 1, 2, 3, n - 2, n - 1, n$ . On the other hand, each of the remaining vertices lies in a cycle of length 4 in  $F$ .

Case E:  $m \equiv 0 \pmod{2}$  and  $n \equiv 1 \pmod{2}$  with  $m \geq 4$  and  $n \geq 7$ . Define  $F \in \mathcal{D}(P_m \times P_n)$  as follows (see Fig. 5):

- (1)  $F[P_m \times P_{n-1}]$  is identical with the orientation in Case D;
- (2) For  $i = 1, 3, \dots, m - 1$ , orient  $(i, n - 1) \rightarrow (i, n)$  and  $(i + 1, n) \rightarrow (i + 1, n - 1)$ ;
- (3) For  $i = 3, 5, \dots, m - 1$ , orient  $(i, n) \rightarrow \{(i - 1, n), (i + 1, n)\}$ ;
- (4) Orient  $(1, n) \rightarrow (2, n)$ .

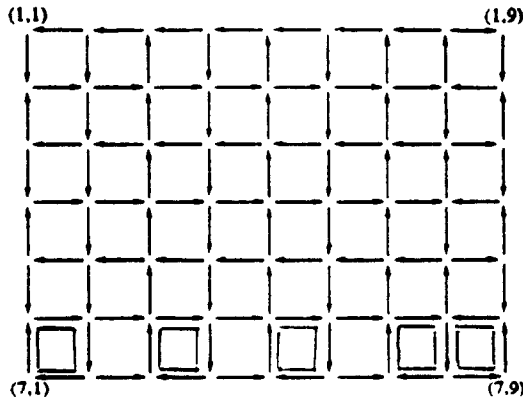


Fig. 6. Orientation of  $P_7 \times P_9$ .

Note that  $d(F) = m + n - 2$ . Also, it can be shown (see Fig. 5 as an illustration) that each vertex is in a cycle of length not exceeding  $m + n - 2$  in  $F$ .

Case F:  $m \equiv n \equiv 1 \pmod{2}$  with  $m \geq 5$  and  $n \geq 7$ . Define  $F \in \mathcal{Q}(P_m \times P_n)$  as follows (see Fig. 6):

- (1)  $F[P_{m-1} \times P_n]$  is identical with the orientation in Case E;
- (2) For each  $j = 2, 4, \dots, n - 1$ , orient

$$(m, j) \rightarrow \{(m, j - 1), (m, j + 1)\},$$

$$(m - 1, j) \rightarrow (m, j) \text{ and}$$

$$(m, j - 1) \rightarrow (m - 1, j - 1);$$

- (3) Orient  $(m, n) \rightarrow (m - 1, n)$ .

Note that  $d(F) = m + n - 2$ . Also, it can be checked (see Fig. 6 as an illustration) that each vertex is in a cycle of length not exceeding  $m + n - 2$  in  $F$ .

Finally, we consider the case when  $m \equiv 1 \pmod{2}$  and  $n \equiv 0 \pmod{2}$  with  $m \geq 5$  and  $n \geq 6$ . By symmetry and the result in Case E, we need only consider the following:

Case G:  $m = 5$  and  $n \equiv 0 \pmod{2}$  with  $n \geq 6$ . Let  $n = 2k$  and define  $F \in \mathcal{Q}(P_5 \times P_n)$  as follows (see Fig. 7):

- (1) For  $i = 1, 2, 4, 5$  and  $j = 1, 2, \dots, k - 1$ , orient  $(i, j) \rightarrow (i, j + 1)$  and  $(3, j + 1) \rightarrow (3, j)$ ;
- (2) For  $i = 1, 2, 4, 5$  and  $j = k + 1, k + 2, \dots, 2k - 1$ , orient  $(i, j + 1) \rightarrow (i, j)$  and  $(3, j) \rightarrow (3, j + 1)$ ;
- (3) For  $j \neq k, k + 1$ , orient  $(1, j) \leftarrow (2, j) \leftarrow (3, j) \rightarrow (4, j) \rightarrow (5, j)$ ;
- (4) For  $j = k, k + 1$ , orient  $(1, j) \rightarrow (2, j) \rightarrow (3, j) \leftarrow (4, j) \leftarrow (5, j)$ ;
- (5) Orient  $(2, k) \rightarrow (2, k + 1), (3, k) \leftarrow (3, k + 1)$  and  $(4, k) \rightarrow (4, k + 1)$ . The edges  $(1, k)(1, k + 1)$  and  $(5, k)(5, k + 1)$  may be arbitrarily oriented.

Note that  $d(F) = n + 3$  and each vertex is in a cycle of length not exceeding  $n + 3$  (see Fig. 7 as an illustration).

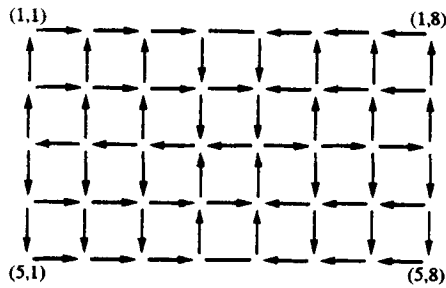


Fig. 7. Orientation of  $P_5 \times P_8$ .

The proof of Lemma 3 is now complete.  $\square$

**Proof of Theorem 1.** Let  $G = P_{k_1} \times P_{k_2}$ . By Lemma 3, the bipartite graph  $G$  admits an orientation  $F$  with  $d(F) = k_1 + k_2 - 2$  such that each vertex in  $G$  is in a cycle of length at most  $k_1 + k_2 - 2$  in  $F$ . By Lemma 2, the graph  $\prod_{i=1}^n P_{k_i}$  admits an orientation  $F^*$  with

$$\begin{aligned} d(F^*) &\leq (k_1 + k_2 - 2) - (n - 2) + \sum_{i=3}^n k_i \\ &= \sum_{i=1}^n k_i - n \\ &= d\left(\prod_{i=1}^n P_{k_i}\right). \end{aligned}$$

The result thus follows.  $\square$

### 3. Cartesian product of paths and cycles

The main aim in this section is to prove the following results.

**Theorem 2.** (i)  $\rho(C_{2n} \times \prod_{i=1}^m P_{k_i}) = 0$  for  $m \geq 1$ ,  $n \geq 2$  and  $k_1 \geq 4$ .

(ii)  $\rho(\prod_{i=1}^m P_{k_i} \times \prod_{i=1}^r C_{n_i}) = 0$  for  $m \geq 2$ ,  $r \geq 0$ ,  $k_1 \geq 3$  and  $k_2 \geq 6$  with  $(k_1, k_2) \neq (3, 6)$ .

(iii)  $\rho(C_{2n} \times \prod_{i=1}^m P_{k_i} \times \prod_{i=1}^r C_{n_i}) = 0$  for  $m \geq 1$ ,  $r \geq 0$ ,  $n \geq 2$  and  $k_1 \geq 4$ .

Note that results (ii) and (iii) are overlapping with (ii) requiring stronger conditions on two paths whereas (iii) requiring a cycle to be even and the length of a path at least four.

In what follows, the vertices of  $\prod_{i=1}^r C_{n_i}$  are labelled  $(x_1, x_2, \dots, x_r)$ , where  $1 \leq x_i \leq n_i$ ,  $1 \leq i \leq r$  so that  $(a_1, a_2, \dots, a_r)$  and  $(b_1, b_2, \dots, b_r)$  are adjacent iff  $|a_k - b_k| = 1 \pmod{n_k - 2}$



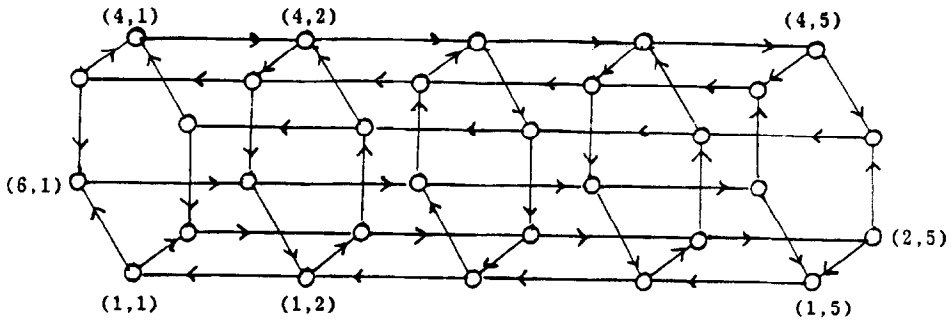


Fig. 8. Orientation of  $C_6 \times P_5$ .

for exactly one  $k$ ,  $1 \leq k \leq r$ , and  $a_i = b_i$  for all  $i$  with  $i \neq k$ . For a real  $x$ , we shall denote by  $[x]$  the greatest integer not exceeding  $x$ .

To prove Theorem 2(i), we need the following result.

**Lemma 4.** For  $n \geq 2$  and  $k \geq 4$ , there exists  $F \in \mathcal{L}(C_{2n} \times P_k)$  such that

- (i)  $d(F) = d(C_{2n} \times P_k) = n + k - 1$  and
- (ii) every vertex in  $C_{2n} \times P_k$  is in a cycle of length at most  $n + k - 1$  in  $F$ .

**Proof.** In [8], the following orientation  $F$  of  $C_{2n} \times P_k$  was introduced (see Fig. 8):

- (i) For  $i \equiv 1 \pmod{2}$  and  $1 \leq i \leq 2n - 1$ , orient  $(i, 1) \rightarrow \{(i + 1, 1), (i - 1, 1)\}$ ;
- (ii) For  $j \equiv 0 \pmod{2}$ ,  $2 \leq j \leq k - 1$ , and  $1 \leq i \leq 2n$ , orient  $(i, j) \rightarrow (i + 1, j)$ ;
- (iii) For  $j \equiv 1 \pmod{2}$ ,  $3 \leq j \leq k - 1$ , and  $1 \leq i \leq 2n$ , orient  $(i, j) \rightarrow (i - 1, j)$ ;
- (iv) For  $i \equiv 0 \pmod{2}$  and  $2 \leq i \leq 2n$ , orient  $(i, k) \rightarrow \{(i + 1, k), (i - 1, k)\}$ ;
- (v) For  $i \equiv 0 \pmod{2}$ ,  $2 \leq i \leq 2n$  and  $1 \leq j \leq k - 1$ , orient  $(i, j) \rightarrow (i, j + 1)$ ;
- (vi) For  $i \equiv 1 \pmod{2}$ ,  $1 \leq i \leq 2n - 1$  and  $2 \leq j \leq k$ , orient  $(i, j) \rightarrow (i, j - 1)$ .

It was shown in [8] also that  $d(F) = d(C_{2n} \times P_k) = n + k - 1$  for  $k \geq 4$  and  $n \geq 2$ . It remains to prove (ii).

Consider the following cycles (see also Fig. 8):

- (A<sub>1</sub>) For  $j \equiv 0 \pmod{2}$  with  $2 \leq j \leq k - 3$  for odd  $k$  or  $2 \leq j \leq k - 2$  for even  $k$ , and  $i \equiv 1 \pmod{2}$  with  $1 \leq i \leq 2n - 1$ ,  $(i, j)(i + 1, j)(i + 1, j + 1)(i, j + 1)(i, j)$ ;
- (A<sub>2</sub>) For  $i \equiv 1 \pmod{2}$ ,  $1 \leq i \leq 2n - 1$ ,  $(i, 1)(i - 1, 1)(i - 1, 2)(i, 2)(i, 1)$ ; and
- (A<sub>3</sub>) For  $i \equiv 1 \pmod{2}$ ,  $1 \leq i \leq 2n - 1$ ,  $(i + 1, k)(i, k)(i, k - 1)(i + 1, k - 1)(i + 1, k)$  if  $k$  is odd, or  $(i - 1, k)(i, k)(i, k - 1)(i - 1, k - 1)(i - 1, k)$  if  $k$  is even.

Clearly, all the above cycles are of length 4 ( $\leq n + k - 1$ ) and they cover  $V(C_{2n} \times P_k)$ . □

**Proof of Theorem 2(i).** Let  $G = C_{2n} \times P_{k_1}$ . By Lemma 4, the bipartite graph  $G$  admits an orientation  $F$  with  $d(F) = n + k_1 - 1$  such that each vertex in  $G$  is in a cycle of length at most  $n + k_1 - 1$  in  $F$ . By Lemma 2, the graph  $H = C_{2n} \times \prod_{i=1}^m P_{k_i}$  admits an

orientation  $F^*$  with

$$\begin{aligned} d(F^*) &\leq n + k_1 - 1 - (m - 1) + \sum_{i=2}^m k_i \\ &= n - m + \sum_{i=1}^m k_i \\ &= d(H). \end{aligned}$$

This proves Theorem 2(i).  $\square$

To prove Theorem 2(ii), we shall extend Thomassen’s observation from  $P_2$  to  $\prod_{i=1}^n C_{k_i}$ .

**Lemma 5.** *If a bipartite graph  $G$  admits an orientation of diameter at most  $k$ , where  $k \geq 3$ , such that each vertex is in a cycle of length at most  $k$ , then the graph  $G \times \prod_{i=1}^n C_{k_i}$  admits an orientation of diameter not exceeding  $k + \sum_{i=1}^n \lceil k_i/2 \rceil$  such that each vertex is in a cycle of length at most  $k$ .*

**Proof.** Let  $V_1$  and  $V_2$  be the partite sets of  $G$ . Let  $F \in \mathcal{D}(G)$  with  $d(F) \leq k$  such that every vertex is in a cycle of length at most  $k$  in  $F$ .

Orient  $G \times \prod_{i=1}^n C_{k_i}$  inductively as follows:

(i) In  $G \times C_{k_1}$ , for  $1 \leq i \leq k_1$ , orient

$$\begin{aligned} (x, i) &\rightarrow (x, i + 1) \quad \text{iff } x \in V_1, \\ (x, i) &\rightarrow (y, i) \quad \text{iff } xy \in E(F). \end{aligned}$$

(Note that the second coordinate of  $(x, i + 1)$  is taken modulo  $k_1$ .)

(ii) Suppose  $G \times \prod_{i=1}^r C_{k_i}$ , where  $1 \leq r \leq n - 1$ , has been oriented. Orient  $G \times \prod_{i=1}^{r+1} C_{k_i}$  so that the orientation of  $G \times \prod_{i=1}^r C_{k_i} \times \{j\}$  is isomorphic to that of  $G \times \prod_{i=1}^r C_{k_i}$  for each  $j = 1, 2, \dots, k_{r+1}$ , and for  $1 \leq i \leq k_{r+1}$ , orient  $(x, a_1, a_2, \dots, a_r, i) \rightarrow (x, a_1, a_2, \dots, a_r, i + 1)$  iff  $x \in V_1$  (note that the last coordinate is taken modulo  $k_{r+1}$ ).

Let  $F^*$  be the resulting orientation of  $G \times \prod_{i=1}^n C_{k_i}$ . We shall now show that there is a path of length at most  $k + \sum_{i=1}^n \lceil k_i/2 \rceil$  from an arbitrary vertex  $u$  to any other vertex  $v$  in  $F^*$ . Let  $u = (x, a_1, a_2, \dots, a_n)$  and  $v = (y, b_1, b_2, \dots, b_n)$ . We may assume that  $x \in V_1$ . As the cartesian product is commutative, we further assume that

$$0 \leq b_i - a_i \leq \left\lfloor \frac{k_i}{2} \right\rfloor \pmod{k_i} \quad \text{for } i = 1, 2, \dots, m$$

and

$$0 \leq a_i - b_i \leq \left\lfloor \frac{k_i}{2} \right\rfloor \pmod{k_i} \quad \text{for } i = m + 1, \dots, n.$$

(1) Let  $w = (x, b_1, b_2, \dots, b_m, a_{m+1}, \dots, a_n)$ . Clearly, there is a  $u - w$  path in  $F^*$  of length at most  $\sum_{i=1}^m \lceil k_i/2 \rceil$ .

- (2) If  $x \neq y$ , let  $w' = (x', b_1, b_2, \dots, b_m, a_{m+1}, \dots, a_n)$ , where  $x'$  is adjacent from  $x$  in an  $x$ - $y$  path of length at most  $k$  in  $F$ . Observe that  $x' \in V_2$ . If  $x = y$ , take a cycle of length at most  $k$  containing  $x$  in  $F$ .
- (3) Let  $w^* = (x', b_1, b_2, \dots, b_n)$ . Clearly, there is a  $w'$ - $w^*$  path of length at most  $\sum_{i=m+1}^n \lceil k_i/2 \rceil$  in  $F^*$ .
- (4) There is a  $w^*$ - $v$  path of length at most  $k - 1$  in  $F^*$ .

Combining (1)–(4), we have

$$\begin{aligned}
 d(u, v) &\leq \sum_{i=1}^m \left\lceil \frac{k_i}{2} \right\rceil + 1 + \sum_{i=m+1}^n \left\lceil \frac{k_i}{2} \right\rceil + k - 1 \\
 &= k + \sum_{i=1}^n \left\lceil \frac{k_i}{2} \right\rceil.
 \end{aligned}$$

This shows that  $d(F^*) \leq k + \sum_{i=1}^n \lceil k_i/2 \rceil$ . The second part of Lemma 5 is obvious as each vertex in  $F^*$  is contained in a cycle of length at most  $k$  in  $F$ .  $\square$

**Proof of Theorem 2(ii).** Let  $G = \prod_{i=1}^m P_{k_i}$ . By Theorem 1 and Lemma 2, the bipartite graph  $G$  admits an orientation  $F$  with  $d(F) = \sum_{i=1}^m k_i - m$ , and every vertex in  $G$  lies in a cycle of length not exceeding  $k_1 + k_2 - 2$  ( $\leq \sum_{i=1}^m k_i - m$ ) in  $F$ . Thus by Lemma 5, the graph  $H = G \times \prod_{i=1}^r C_{n_i}$  admits an orientation  $F^*$  with

$$\begin{aligned}
 d(F^*) &\leq \sum_{i=1}^m k_i - m + \sum_{i=1}^r \left\lceil \frac{n_i}{2} \right\rceil \\
 &= d(H).
 \end{aligned}$$

This proves Theorem 2(ii).  $\square$

**Proof of Theorem 2(iii).** Let  $G = C_{2n} \times \prod_{i=1}^m P_{k_i}$ . By Theorem 2(i) and Lemma 2, the bipartite graph  $G$  admits an orientation  $F$  with  $d(F) = n + \sum_{i=1}^m k_i - m$ , and every vertex in  $G$  lies in a cycle of length at most  $4$  ( $\leq n + \sum_{i=1}^m k_i - m$ ) in  $F$  (see the proof of Lemma 4(ii)). Thus by Lemma 5, the graph  $H = G \times \prod_{i=1}^r C_{n_i}$  admits an orientation  $F^*$  with

$$\begin{aligned}
 d(F^*) &\leq n + \sum_{i=1}^m k_i - m + \sum_{i=1}^r \left\lceil \frac{n_i}{2} \right\rceil \\
 &= d(H).
 \end{aligned}$$

This proves Theorem 2(iii).  $\square$

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