



# On fast algorithms for the evaluation of Legendre coefficients

Shuhuang Xiang

Department of Applied Mathematics and Software, Central South University, Changsha, Hunan 410083, PR China

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## ABSTRACT

In this paper, we present formulas between Legendre and Chebyshev expansion coefficients for a piecewise smooth (or Dini–Lipschitz) function, analyze the error bounds for the Piessens' algorithm and present a new algorithm with  $O(N \log N)$  operations for computation of the first  $N + 1$  coefficients of the Legendre expansion. Finally, we show the identity to the formulas given by Iserles [5] for analytic functions in a neighborhood of  $[-1, 1]$ .

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## 1. Introduction

Legendre expansions are widely used in approximation theory, numerical integration, solutions of partial differential equations, analysis of pseudospectral methods, special function theory, etc. Given a function  $f \in L^2[-1, 1]$ , it takes  $O(N^2)$  operations to compute the first  $N + 1$  terms of the Legendre expansion

$$f(x) = \sum_{m=0}^{\infty} a_m P_m(x) \quad (1.1)$$

by the standard discrete expansion [1]

$$a_m = \left(m + \frac{1}{2}\right) \int_{-1}^1 f(x) P_m(x) dx \approx \left(m + \frac{1}{2}\right) \sum_{n=0}^N w_n f(x_n) P_m(x_n), \quad m = 0, 1, \dots, N, \quad (1.2)$$

where  $\{x_n\}_{n=0}^N$  is the Gauss-point set (the roots of  $P_{N+1}(x) = 0$ ) and  $w_n$  the corresponding weights in the Gauss–Legendre quadrature [2].

Piessens [3] presented another efficient algorithm with order  $O(N^2)$  to approximate the first  $N + 1$  Legendre coefficients by using the Chebyshev expansion:

$$f(x) = \sum_{m=0}^{\infty} {}'c_m T_m(x), \quad c_m = \frac{2}{\pi} \int_{-1}^1 \frac{f(x) T_m(x)}{\sqrt{1-x^2}} dx, \quad m \in \mathbb{Z}_+, \quad (1.3)$$

where  $\sum'$  denotes a sum whose first term is halved. Then  $a_m$  can be approximated [3] by

$$a_m \approx \left(m + \frac{1}{2}\right) \sum_{k=0}^N {}'c_k J_{m,k}, \quad m = 0, 1, \dots, N$$

E-mail address: [xiangsh@mail.csu.edu.cn](mailto:xiangsh@mail.csu.edu.cn).

which yields a fast algorithm [3]

$$\bar{a}_m = \left(m + \frac{1}{2}\right) \sum_{j=0}^N {}''I_{m,j} \tilde{c}_j, \quad m = 0, 1, \dots, N. \tag{1.4}$$

Here  $I_{m,k} = \int_{-1}^1 P_m(x) T_k(x) dx$  can be calculated by recursion in  $O(N^2)$  operations

$$I_{m,k} = \begin{cases} \frac{2}{1-k^2}, & m = 0, k \text{ is even} \\ \frac{2^{2m}(m!)^2}{(2m+1)!}, & k = m \geq 1 \\ \frac{[(m+2j-3)(m+2j-2) - m(m+1)](m+2j)}{[(m+2j+1)(m+2j) - m(m+1)](m+2j-2)} I_{m,m+2j-2}, & m \geq 1, k = m+2j, j \geq 1 \\ 0, & \text{otherwise,} \end{cases} \tag{1.5}$$

$\tilde{c}_m$  ( $m = 0, 1, \dots, N$ ) can be evaluated by FFT in  $O(N \log N)$  operations [4]:

```
function b = coefficients(f, n) % (n + 1) coefficients for f
a0 = -1; b0 = 1 % approximate interval of f
x = (a0 + b0)/2 + (b0 - a0)/2 * cos(pi * (0 : n)'/n); % Chebyshev points of the second kind
fx = feval(f, x)/(2 * n); % f evaluated at these points
g = fft(fx([1 : N0 + 1N0 : -1 : 2])); % FFT
c_tilde = [2 * g(1); g(2 : N0) + g(2 * N0 : -1 : N0 + 2); 2 * g(N0 + 1)]; % Chebyshev coefficients
```

and  $\sum''$  denotes a sum whose first and last terms are halved.

More recently, Iserles [5] proposed a new improvement by using the Cauchy theorem and Taylor expansions of the Legendre and hypergeometric functions  $\Phi_m$  and  ${}_2F_1$ , respectively, for  $f(x)$  analytic on and inside the Bernstein ellipse  $\mathcal{E}_\rho$  with foci  $\pm 1$  and major and minor semiaxis lengths summing to  $\rho (> 1)$ . The derivation is rather convoluted but the outcome is a surprisingly simple numerical algorithm [5]. The total error bound on the algorithm remains open there [5].

In this paper, we will focus on fast algorithms for the evaluation of Legendre expansion for non-analytic functions of finite regularity and their error analysis. We will show that the Legendre coefficient  $a_m$  for piecewise smooth (or Dini-Lipschitz)  $f(x)$  on  $[-1, 1]$  can be represented by the Chebyshev coefficients as

$$a_0 = \frac{1}{2}c_0 + \frac{1}{2} \sum_{n=1}^{\infty} I_{0,2n} c_{2n}, \quad a_m = \left(m + \frac{1}{2}\right) \sum_{n=0}^{\infty} I_{m,m+2n} c_{m+2n}, \quad m = 1, 2, \dots, \tag{1.6}$$

from which we can construct an efficient algorithm to compute  $a_m$

$$\begin{aligned} \tilde{a}_0 &= \frac{1}{2}c_0 + \frac{1}{2} \sum_{n=1}^{N_0} I_{0,2n} \tilde{c}_{2n}, & \tilde{a}_m &= \left(m + \frac{1}{2}\right) \sum_{n=0}^{N_0} I_{m,m+2n} \tilde{c}_{m+2n}, \quad m = 1, \dots, N-1, \\ \tilde{a}_N &= \left(N + \frac{1}{2}\right) \sum_{n=0}^{N_0-1} I_{N,N+2n} \tilde{c}_{N+2n} + \frac{1}{2} \left(N + \frac{1}{2}\right) I_{N,N+2N_0} \tilde{c}_{N+2N_0}, \end{aligned} \tag{1.7}$$

where  $N_0$  is a fixed positive integer depending on  $f(x)$  and independent of  $N$  for a given tolerance  $\delta > 0$ ,  $\tilde{c}_m$  ( $m = 0, 1, \dots, N + 2N_0$ ) can be evaluated by FFT in  $O((N + 2N_0) \log(N + 2N_0))$  operations, and  $I_{m,j}$  can be obtained by (1.5) in  $O(NN_0)$  operations. Therefore, the total cost is  $O(N \log N)$ . In Section 2, we will derive the error bounds for Piessens' algorithm (1.4) and the new algorithm (1.7). In the final remark, we will show that (1.6) is identity to that given in [5] for analytic functions in a neighborhood of  $[-1, 1]$ .

## 2. Main results

**Lemma 2.1.** Suppose  $f(x) \in L^2[-1, 1]$  is piecewise smooth (or Dini-Lipschitz) on  $[-1, 1]$ , then

$$a_0 = \frac{1}{2}c_0 + \frac{1}{2} \sum_{n=1}^{\infty} I_{0,2n} c_{2n}, \quad a_m = \left(m + \frac{1}{2}\right) \sum_{n=0}^{\infty} I_{m,m+2n} c_{m+2n}, \quad m = 1, 2, \dots$$

**Proof.** Recall that a function  $h(x)$  is called piecewise continuous on  $[-1, 1]$  provided (i)  $h$  is continuous on  $[-1, 1]$  except at finitely many points  $y_1, \dots, y_k$ ; and (ii) at each of the points  $y_1, \dots, y_k$ , the left-hand and right-hand limits of  $h$  exist.  $h$  is

called piecewise smooth if  $h$  and  $h'$  are both piecewise continuous on  $[-1, 1]$  [6, p. 32]. Since  $f(x)$  is piecewise smooth on  $[-1, 1]$ , then  $f(\cos(\theta))$  is piecewise smooth on  $[-\pi, \pi]$  and

$$f(\cos(\theta)) = \sum_{n=0}^{\infty} c_n \cos(n\theta), \quad \cos(\theta) \neq y_j, \quad j = 1, \dots, k$$

(see [6, p. 35]), where  $c_n$  is the same as that in (1.3). Consequently we have

$$f(x) = \sum_{n=0}^{\infty} c_n T_n(x), \quad x \neq y_j, \quad j = 1, \dots, k. \tag{2.1}$$

Thus, from (2.1) and Theorem 1.27 [7, p. 22] it follows

$$\begin{aligned} a_m &= \left(m + \frac{1}{2}\right) \int_{-1}^1 f(x) P_m(x) dx = \left(m + \frac{1}{2}\right) \int_{-1}^1 \sum_{n=0}^{\infty} c_n T_n(x) P_m(x) dx \\ &= \left(m + \frac{1}{2}\right) \sum_{n=0}^{\infty} c_n \int_{-1}^1 T_n(x) P_m(x) dx, \quad m \in \mathbb{Z}_+, \end{aligned}$$

which directly leads to (1.6).

In the case that  $f$  is Dini–Lipschitz, from the Dini–Lipschitz Theorem [8, p. 129], the Chebyshev expansion uniformly converges to  $f(x)$ . By the above same proof, (1.6) is also satisfied.  $\square$

**Lemma 2.2.** For the Legendre expansion (1.1) and Chebyshev expansion (1.3), the coefficients satisfy

$$|a_n| \leq \frac{2\sqrt{nM}}{\rho^{n-1}(\rho^2 - 1)} \quad (\text{see [9]}), \quad |c_n| \leq \frac{2M}{\rho^n} \quad (\text{see [10]}), \quad n = 1, 2, \dots \tag{2.2}$$

if  $f$  is analytic with  $|f(z)| \leq M$  in the region bounded by the ellipse  $\mathcal{E}_\rho$  with foci  $\pm 1$  and major and minor semiaxis lengths summing to  $\rho > 1$ ; or

$$|a_n| \leq \frac{\sqrt{\pi} V_p}{\sqrt{2(n-p-1)} \left(n - \frac{1}{2}\right) \left(n - \frac{3}{2}\right) \cdots \left(n - \frac{2p-1}{2}\right)} \quad (\text{see [11]}), \tag{2.3a}$$

$$|c_n| \leq \frac{2V_p}{\pi n(n-1) \cdots (n-p)} \quad (\text{see [12]}) \tag{2.3b}$$

if  $f, \dots, f^{(p-1)}$  are absolutely continuous on  $[-1, 1]$  and  $f^{(p)}$  has bounded variation  $V_p$  for some  $p \geq 1$ .

**Theorem 2.1.** The error bound for the Piessens' algorithm (1.4) for evaluating  $a_m$  can be given as

$$|\bar{a}_m - a_m| \leq \begin{cases} \frac{(2m+1)I_{m,m}V_p}{p\pi N(N-1)\cdots(N-p+1)} + \frac{V_p}{(p-1)\left(N - \frac{1}{2}\right)\left(N - \frac{3}{2}\right)\cdots\left(N - \frac{2p-3}{2}\right)} \sqrt{\frac{\pi}{2(N-p)}}, & V_p < \infty \\ \frac{(4m+2)I_{m,m}M}{(\rho-1)\rho^N} + \frac{2M[(N+1)(\rho-1)+1]}{\sqrt{N+1}\rho^{N-1}(\rho-1)^3(\rho+1)}, & f(x) \text{ analytic in } \mathcal{E}_\rho \end{cases} \tag{2.4}$$

for  $m = 0, 1, 2, \dots, N$ .

**Proof.** From the definitions of  $I_{m,k}$ s, it follows for  $m \geq 0$  that

$$|a_m| \leq \left(m + \frac{1}{2}\right) \sum_{n=0}^{\infty} |c_n I_{m,n}| = \left(m + \frac{1}{2}\right) |c_m I_{m,m}| + \left(m + \frac{1}{2}\right) \sum_{j=1}^{\infty} |c_{m+2j} I_{m,m+2j}|,$$

and for  $j \geq 0$  that

$$|I_{m,m+2j+2}| = \frac{|(m+2j-3)(m+2j-2) - m(m+1)|(m+2j)}{|(m+2j+1)(m+2j) - m(m+1)|(m+2j-2)} |I_{m,m+2j-2}| \leq \frac{|2j-1|}{2j+1} |I_{m,m+2j}|,$$

which implies

$$|I_{m,m+2j}| \leq \frac{1}{2j-1} |I_{m,m}|, \quad j = 0, 1, 2, \dots$$

Moreover, from (1.4)–(1.6), the error bound for evaluation of  $a_m$  by (1.4) can be given as

$$|\bar{a}_m - a_m| \leq \left(m + \frac{1}{2}\right) I_{m,m} \left( \sum_{j=0}^{N-1} |\tilde{c}_j - c_j| + \frac{1}{2} |\tilde{c}_N - c_N| \right) + \sum_{j=N+1}^{\infty} |a_j|, \quad m = 0, 1, 2, \dots, N. \quad (2.5)$$

Recalling (2.13.1.11) in [2] (also see [13, p. 96])

$$\tilde{c}_j - c_j = \sum_{\ell=1}^{\infty} (c_{2\ell n-j} + c_{2\ell n+j}), \quad j = 0, 1, \dots, n, \quad (2.6)$$

which, together with (2.5) and (2.6), gives

$$|\bar{a}_m - a_m| \leq \left(m + \frac{1}{2}\right) I_{m,m} \sum_{j=N+1}^{\infty} |c_j| + \sum_{j=N+1}^{\infty} |a_j|. \quad (2.7)$$

Furthermore, note that  $\sum_{j=N+1}^{\infty} |c_j| \leq \frac{2V_p}{p\pi N(N-1)\dots(N-p+1)}$  [14, Theorem 2.1] for  $V_p < \infty$  and

$$\begin{aligned} \sum_{j=N+1}^{\infty} |a_j| &\leq V_p \sqrt{\frac{\pi}{2(N-p)}} \sum_{j=N+1}^{\infty} \frac{1}{(j - \frac{1}{2})(j - \frac{3}{2}) \dots (j - \frac{2p-1}{2})} \\ &= V_p \sqrt{\frac{\pi}{2(N-p)}} \frac{1}{p-1} \sum_{j=N+1}^{\infty} \left[ \frac{1}{(j - \frac{3}{2}) \dots (j - \frac{2p-1}{2})} - \frac{1}{(j - \frac{1}{2}) \dots (j - \frac{2p-3}{2})} \right] \\ &= \frac{V_p}{(p-1)(N - \frac{1}{2})(N - \frac{3}{2}) \dots (N - \frac{2p-3}{2})} \sqrt{\frac{\pi}{2(N-p)}}. \end{aligned}$$

These together lead to (2.4) for  $V_p < \infty$ .

Similarly, by (2.2) together with Theorem 2.1 [14] and Corollary 2.1 [9], it obtains (2.4) for  $f(x)$  analytic in  $\mathcal{E}_\rho$ .  $\square$

**Theorem 2.2.** The error bound for the algorithm (1.7) for evaluating  $a_m$  can be given as

$$|\tilde{a}_m - a_m| \leq \begin{cases} \frac{(4m+2)I_{m,m}V_p}{p\pi(N_0+2m)(N_0+2m-1)\dots(N_0+2m-p+1)}, & V_p < \infty \\ \frac{(4m+2)I_{m,m}M}{(\rho-1)\rho^{m+2N_0}}, & f(x) \text{ analytic in } \mathcal{E}_\rho \end{cases} \quad (2.8)$$

for  $m = 0, 1, 2, \dots, N$ .

**Proof.** In a similar way, the error for computation of each  $a_m$  by  $\tilde{a}_m$  can be estimated from (1.6) and (1.7) for  $m = 0, 1, \dots, N$  by

$$\begin{aligned} |\tilde{a}_m - a_m| &\leq \left(m + \frac{1}{2}\right) I_{m,m} \left( \sum_{n=0}^{m+2N_0-1} |\tilde{c}_n - c_n| + \left| \frac{1}{2} \tilde{c}_{N+2N_0} - c_{N+2N_0} \right| + \sum_{n=N_0+1}^{\infty} |c_{m+2n}| \right) \\ &\leq \left(m + \frac{1}{2}\right) I_{m,m} \left( \sum_{n=m+2N_0+1}^{\infty} |c_n| + \sum_{n=N_0+1}^{\infty} |c_{m+2n}| \right) \quad \text{by (2.6)} \\ &\leq (2m+1) I_{m,m} \sum_{n=m+2N_0+1}^{\infty} |c_n|, \end{aligned}$$

which together with  $\sum_{j=N+1}^{\infty} |c_j| \leq \frac{2V_p}{p\pi N(N-1)\dots(N-p+1)}$  for  $V_p < \infty$  and  $\sum_{j=N+1}^{\infty} |c_j| \leq \frac{2M}{(\rho-1)\rho^N}$  [14, Theorem 2.1] yields the desired result.  $\square$

**Remark 1.** From the definition of  $I_{m,m}$ s, we see that  $\frac{1}{2}I_{0,0} = 1$  and  $(m + \frac{1}{2})I_{m,m} \leq \sqrt{m}$  for  $m \geq 1$ .

**Remark 2.** Both algorithms (1.2) and (1.5) can achieve higher accuracy as  $N$  increases. However, these methods cannot improve the accuracy of the relative errors at the last few coefficients as reported in [3]. But from (2.8), it shows that the relative errors by (1.7) at the last few coefficients can achieve higher accuracy as  $N_0$  increases.

**Example 1.** We illustrate the convergence rates of (1.7) and compare with that evaluated by the discrete expansion (1.2) and (1.5) with  $N = 100$  (Tables 1 and 2).

**Table 1**  
Approximation of coefficients of the Legendre expansions by (1.7).

$f(x)$	$a_k$	$N_0 = 2$	$N_0 = 10$	$N_0 = 10^2$	$N_0 = 10^3$	Exact value
$e^x$	$a_0$	1.17520248187784	1.17520119364380	1.17520119364380	1.17520119364380	1.17520119364380
	$a_{100}$	0.00000000000000	0.00000000000000	0.00000000000000	0.00000000000000	0.00000000000000
$\frac{1+x}{4+x^2}$	$a_0$	0.23182165921697	0.23182380450040	0.23182380450040	0.23182380450040	0.23182380450040
	$a_{100}$	0.00000000000000	0.00000000000000	0.00000000000000	0.00000000000000	0.00000000000000
$e^{-1/x^2}$	$a_0$	0.08855032961820	0.08907377800247	0.08907385589078	0.08907385589078	0.08907385589078
	$a_{100}$	0.00000000037222	0.00000000032330	0.00000000032424	0.00000000032424	0.00000000032424
$ x ^3$	$a_0$	0.24979747886432	0.24999996006293	0.25000000009753	0.25000000000004	0.25000000000000
	$a_{100}$	0.00000154312226	0.00000117911069	0.00000094424591	0.00000094223996	0.00000094223929

**Table 2**  
Approximation of coefficients of the Legendre expansions by (1.2) and (1.5) respectively.

$f(x)$	$a_k$	By (1.2)	By (1.5)	Exact value
$e^x$	$a_0$	1.17520119364380	1.17520119364380	1.17520119364380
	$a_{100}$	0.00000000000000	0.00000000000000	0.00000000000000
$\frac{1+x}{4+x^2}$	$a_0$	0.23182380450040	0.23182380450040	0.23182380450040
	$a_{100}$	0.00000000000001	0.00000000000000	0.00000000000000
$e^{-1/x^2}$	$a_0$	0.08907385589078	0.08907385589078	0.08907385589078
	$a_{100}$	0.00000000091831	0.00000000016170	0.00000000032424
$ x ^3$	$a_0$	0.25000000765123	0.25000000812140	0.25000000000000
	$a_{100}$	0.00000183956737	0.00000068850364	0.00000094223929

**3. Final remarks**

Iserles [5] proposed fast algorithms for analytic function  $f(x)$  on and inside the Bernstein ellipse  $\mathcal{E}_\rho$ . By using the Cauchy theorem and Taylor expansions of the Legendre and hypergeometric functions  $\Phi_m$  and  ${}_2F_1$ , the Legendre coefficient  $a_m$  can be represented by [5] in the form of

$$a_m = \frac{2^{2m}(m!)^2}{2(2m)!\rho^m} \sum_{j=0}^{\infty} \int_{-\pi}^{\pi} (1 - \rho^{-2}e^{2i\theta})f\left(\frac{1}{2}(\rho^{-1}e^{i\theta} + \rho e^{-i\theta})\right) {}_2F_1\left(m + 1, \frac{1}{2}; m + \frac{3}{2}; \rho^{-2}e^{2i\theta}\right) e^{im\theta} d\theta.$$

Letting  $\rho \rightarrow 1$  it follows [5]

$$a_m = \frac{1}{2\pi} \sum_{j=0}^{\infty} \tilde{g}_{m,j} \int_{-\pi}^{\pi} (1 - e^{2i\theta})f(\cos(\theta))e^{i(m+2j)\theta} d\theta = \frac{1}{2} \sum_{j=0}^{\infty} \tilde{g}_{m,j}(c_{m+2j} - c_{m+2j+2}), \tag{3.1}$$

where  $\tilde{g}_{m,j} = \frac{2^{2m}(m!)^2(m+1)j(\frac{1}{2})_j}{(2m)j!(m+\frac{1}{2})_j}$  can be computed by

$$\tilde{g}_{0,0} = 1, \quad \tilde{g}_{m,0} = \frac{m}{m - \frac{1}{2}}\tilde{g}_{m-1,0}, \quad \tilde{g}_{m,j} = \frac{(m+j)(j - \frac{1}{2})}{j(m+j + \frac{1}{2})}\tilde{g}_{m,j-1}, \quad m = 1, 2, \dots, j = 1, 2, \dots,$$

which together leads to a fast algorithm [5]

$$\hat{a}_m = \frac{1}{2} \sum_{j=0}^{N_0-1} \tilde{g}_{m,j}(\tilde{c}_{m+2j} - \tilde{c}_{m+2j+2}), \quad m = 0, 1, \dots, N - 1$$

$$\hat{a}_N = \frac{1}{2} \sum_{j=0}^{N_0-2} \tilde{g}_{N,j}(\tilde{c}_{N+2j} - \tilde{c}_{N+2j+2}) + \frac{1}{2}\tilde{g}_{N,j}\left(\tilde{c}_{N+2N_0-2} - \frac{1}{2}\tilde{c}_{N+2N_0}\right). \tag{3.2}$$

In the following, we will show that (3.1) and (1.6) are identity, that is,

$$a_0 = \frac{1}{2}c_0 + \frac{1}{2} \sum_{n=1}^{\infty} I_{0,2n}c_{2n} = \frac{1}{2} \sum_{j=0}^{\infty} \tilde{g}_{0,j}(c_{2j} - c_{2j+2})$$

$$a_m = \left(m + \frac{1}{2}\right) \sum_{n=0}^{\infty} I_{m,m+2n}c_{m+2n} = \frac{1}{2} \sum_{j=0}^{\infty} \tilde{g}_{m,j}(c_{m+2j} - c_{m+2j+2}), \quad m \geq 1.$$

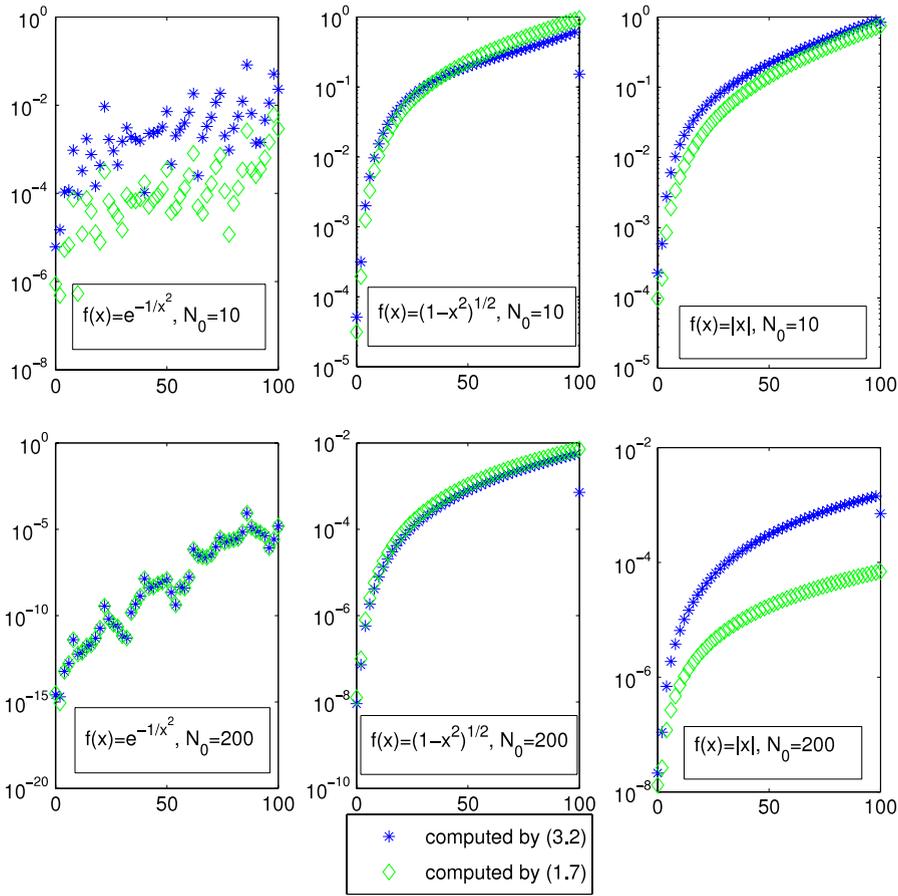


Fig. 1. The relative errors computed by (3.2) and (1.7) respectively:  $m = 0 : 2 : 100$ .

Notice that (3.1) can be rewritten as

$$a_m = \frac{1}{2} \tilde{g}_{m,0} c_m + \frac{1}{2} \sum_{j=1}^{\infty} (\tilde{g}_{m,j} - \tilde{g}_{m,j-1}) c_{m+2j} = \frac{1}{2} \tilde{g}_{m,0} c_m - \frac{1}{2} \sum_{j=1}^{\infty} \frac{\frac{1}{2}m + j}{j(m + j + \frac{1}{2})} \tilde{g}_{m,j-1} c_{m+2j}. \quad (3.3)$$

Comparing the first coefficient in (3.3) with that in (1.6), it is obvious by the definitions of  $I_{m,m}$  and  $\tilde{g}_{m,0}$  that

$$\left(m + \frac{1}{2}\right) I_{m,m} = \frac{1}{2} \tilde{g}_{m,0}, \quad m = 0, 1, 2, \dots$$

Hence, it is only necessary to prove that for each fixed  $m$

$$\left(m + \frac{1}{2}\right) I_{m,m+2j} = -\frac{1}{2} \frac{\frac{1}{2}m + j}{j(m + j + \frac{1}{2})} \tilde{g}_{m,j-1}, \quad j = 1, 2, \dots, \quad (3.4)$$

which can be proved by induction on  $j$ . Suppose (3.4) is true for  $k = j$ , then for  $k = j + 1$ , by (1.6) it yields

$$\begin{aligned} & \left(m + \frac{1}{2}\right) I_{m,m+2j+2} \\ &= \left(m + \frac{1}{2}\right) I_{m,m+2j} \frac{[(m + 2j - 1)(m + 2j) - m(m + 1)](m + 2j + 2)}{[(m + 2j + 3)(m + 2j + 2) - m(m + 1)](m + 2j)} \\ &= -\frac{1}{2} \frac{\frac{1}{2}m + j}{j(m + j + \frac{1}{2})} \tilde{g}_{m,j-1} \frac{[(m + 2j - 1)(m + 2j) - m(m + 1)](m + 2j + 2)}{[(m + 2j + 3)(m + 2j + 2) - m(m + 1)](m + 2j)} \\ &= -\frac{1}{2} \frac{\frac{1}{2}m + j + 1}{(j + 1)(m + j + 1 + \frac{1}{2})} \cdot \frac{(m + j)(j - \frac{1}{2})}{j(m + j + \frac{1}{2})} \tilde{g}_{m,j-1} \end{aligned}$$

$$\begin{aligned} & \times \left\{ \frac{j(j+1)(m+j+\frac{1}{2})(m+j+\frac{3}{2})}{(\frac{1}{2}m+j+1)(m+j)(j-\frac{1}{2})} \frac{(\frac{1}{2}m+j)[(m+2j-1)(m+2j)-m(m+1)](m+2j+2)}{j(m+j+\frac{1}{2})[(m+2j+3)(m+2j+1)-m(m+1)](m+2j)} \right\} \\ & = -\frac{1}{2} \frac{\frac{1}{2}m+j+1}{j(m+j+\frac{3}{2})} \tilde{g}_{m,j} \end{aligned}$$

where we used MAPLE 11 to verify the last factor in the right-hand side of the third equality is equal to 1.

Therefore, the fast algorithm (3.2) can also be applied to piecewise smooth (or Dini–Lipschitz) functions. Comparing (3.3) with (1.7) ( $N_0 - 1$  instead of  $N_0$  in (1.7)), we see that for fixed  $m = 0, 1, \dots, N$  in (3.2) and (1.7)

$$(3.2) = (1.7) - \frac{1}{2} \tilde{g}_{m, N_0} \tilde{c}_{m+2N_0+2}.$$

**Example 2.** We use  $f(x) = e^{-1/x^2}, \sqrt{1-x^2}, |x|$  to compare (3.2) with (1.7) to show the efficiency of these two algorithms (Fig. 1).

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