# Extensions of barrier sets to nonzero roots of the matching polynomial 

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#### Abstract

In matching theory, barrier sets (also known as Tutte sets) have been studied extensively due to their connection to maximum matchings in a graph. For a root $\theta$ of the matching polynomial, we define $\theta$-barrier and $\theta$-extreme sets. We prove a generalized Berge-Tutte formula and give a characterization for the set of all $\theta$-special vertices in a graph.


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## 1. Introduction

All the graphs in this paper are simple and finite.
Definition 1.1. An $r$-matching in a graph $G$ is a set of $r$ edges, no two of which have a vertex in common. The number of $r$-matchings in $G$ will be denoted by $p(G, r)$. Set $p(G, 0)=1$. The matching polynomial of $G$ is defined by

$$
\mu(G, x)=\sum_{r=0}^{\lfloor n / 2\rfloor}(-1)^{r} p(G, r) x^{n-2 r} .
$$

In [1], Chen and Ku developed a Gallai-Edmonds decomposition associated to a root $\theta$ of the matching polynomial, generalizing the usual one which is the special case where $\theta=0$. Note that 0 is a root of the matching polynomial if and only if the graph has no perfect matching. In this paper, we extend the notions of barrier and extreme sets to $\theta$-barrier and $\theta$-extreme sets and show connections with the Gallai-Edmonds decomposition for general $\theta$.

We shall denote the multiplicity of $\theta$ as a root of $\mu(G, x)$ by mult $\theta, G)$. In particular, mult $(\theta, G)=0$ if and only if $\theta$ is not a root of $\mu(G, x)$.

The following are properties of $\mu(G, x)$.
Theorem 1.2 (Theorem 1.1 on $p .2$ of [2]).
(a) $\mu(G \cup H, x)=\mu(G, x) \mu(H, x)$ when $G$ and $H$ are disjoint graphs,
(b) $\mu(G, x)=\mu(G-e, x)-\mu(G \backslash u v, x)$ if $e=\{u, v\}$ is an edge of $G$,
(c) $\mu(G, x)=x \mu(G \backslash u, x)-\sum_{i \sim u} \mu(G \backslash u i, x)$, where $i \sim u$ means $i$ is adjacent to $u$,
(d) $\frac{\mathrm{d}}{\mathrm{d} x} \mu(G, x)=\sum_{i \in V(G)} \mu(G \backslash i, x)$, where $V(G)$ is the vertex set of $G$.

[^0]It is well known that all roots of $\mu(G, x)$ are real (see [5] and in particular [2, Corollary 1.2]). By Theorem 5.3 on p. 29 and Theorem 1.1 on p. 96 of [2], one can easily deduce the following lemma (see also [4]).

Lemma 1.3. If $G$ is a graph and $u \in V(G)$, then

$$
\operatorname{mult}(\theta, G)-1 \leq \operatorname{mult}(\theta, G \backslash u) \leq \operatorname{mult}(\theta, G)+1
$$

As a consequence of Lemma 1.3, we can classify the vertices in a graph with respect to $\theta$ as follows.
Definition 1.4 (See [3, Section 3]). For any $u \in V(G)$,
(a) $u$ is $\theta$-essential if $\operatorname{mult}(\theta, G \backslash u)=\operatorname{mult}(\theta, G)-1$,
(b) $u$ is $\theta$-neutral if $\operatorname{mult}(\theta, G \backslash u)=\operatorname{mult}(\theta, G)$,
(c) $u$ is $\theta$-positive if $\operatorname{mult}(\theta, G \backslash u)=\operatorname{mult}(\theta, G)+1$.

Furthermore, when $u$ is not $\theta$-essential but is adjacent to some $\theta$-essential vertex, we say that $u$ is $\theta$-special.
It turns out that $\theta$-special vertices play an important role in the Gallai-Edmonds decomposition of a graph (see [1]). One of the results in this paper is a characterization of the set of these vertices in terms of $\theta$-barriers.

Note that, if $\operatorname{mult}(\theta, G)=0$, then, for any $u \in V(G), u$ is either $\theta$-neutral or $\theta$-positive, and no vertices in $G$ can be $\theta$-special. By Corollary 4.3 of [3], a $\theta$-special vertex is $\theta$-positive. Let $D_{\theta}(G), A_{\theta}(G)$, and $N_{\theta}(G)$, respectively, denote the sets of $\theta$-essential, $\theta$-special, and $\theta$-neutral vertices, and let $P_{\theta}(G)$ denote the set of vertices that are $\theta$-positive but not $\theta$-special. These four sets partition $V(G)$.

Note that there are no 0 -neutral vertices. If there were, then there would be a vertex, say $u$, with mult $(0, G)=$ mult $(0, G \backslash u)$. There is then a maximum matching that does not cover $u$, and so $u \in D_{0}(G)$, a contradiction, for $D_{0}(G)$ is the set of all points in $G$ which are not covered by at least one maximum matching of $G$ (see [11, Section 3.2 on $p$. 93] for the details). Thus $N_{0}(G)=\varnothing$ and $V(G)=D_{0}(G) \cup A_{0}(G) \cup P_{0}(G)$.

Definition 1.5 (See [3, Section 3]). A graph $G$ is said to be $\theta$-critical if all vertices in $G$ are $\theta$-essential and mult $\theta, G)=1$.
The Gallai-Edmonds structure theorem describes a certain canonical decomposition of $V(G)$ with respect to the zero root of $\mu(G, x)$.

Theorem 1.6 (Theorem 1.5 of [1]). Let $G$ be a graph with $\theta$ a root of $\mu(G, x)$. If $u \in A_{\theta}(G)$, then
(i) $D_{\theta}(G \backslash u)=D_{\theta}(G)$,
(ii) $P_{\theta}(G \backslash u)=P_{\theta}(G)$,
(iii) $N_{\theta}(G \backslash u)=N_{\theta}(G)$,
(iv) $A_{\theta}(G \backslash u)=A_{\theta}(G) \backslash\{u\}$.

Theorem 1.7 (Theorem 1.7 of [1]). If $G$ is connected and every vertex of $G$ is $\theta$-essential, then $\operatorname{mult}(\theta, G)=1$.
By Theorems 1.6 and 1.7, it is not hard to deduce the following, whose proof is omitted.

## Corollary 1.8.

(i) $A_{\theta}\left(G \backslash A_{\theta}(G)\right)=\varnothing, D_{\theta}\left(G \backslash A_{\theta}(G)\right)=D_{\theta}(G), P_{\theta}\left(G \backslash A_{\theta}(G)\right)=P_{\theta}(G)$, and $N_{\theta}\left(G \backslash A_{\theta}(G)\right)=N_{\theta}(G)$.
(ii) $G \backslash A_{\theta}(G)$ has exactly $\left|A_{\theta}(G)\right|+$ mult $(\theta, G) \theta$-critical components.
(iii) If $H$ is a component of $G \backslash A_{\theta}(G)$, then either $H$ is $\theta$-critical or mult $(\theta, H)=0$.
(iv) The subgraph induced by $D_{\theta}(G)$ consists of all the $\theta$-critical components in $G \backslash A_{\theta}(G)$.

Consider the Gallai-Edmonds decomposition of the graph $G$ in Fig. 1 for $\theta=0$ and $\theta=1$. For $\theta=0$, it is the usual Gallai-Edmonds decomposition (see [11, Section 3.2 on p. 93]). First note that mult $(1, G)=1=\operatorname{mult}(0, G)$.

For $\theta=1$, we have $A_{1}(G)=\left\{u_{1}\right\}, D_{1}(G)=\left\{u_{2}, u_{3}, u_{4}, u_{5}\right\}, P_{1}(G)=\left\{u_{7}, u_{10}\right\}$, and $N_{1}(G)=\left\{u_{6}, u_{8}, u_{9}, u_{11}, u_{12}, u_{13}\right\}$. Now $C_{1}, C_{2}, C_{3}, C_{4}$ are the only components in $G \backslash A_{1}(G)$. Note that $C_{1}$ and $C_{2}$ are 1-critical, and mult $\left(1, C_{3}\right)=0=\operatorname{mult}\left(1, C_{4}\right)$.

For $\theta=0$, we have $A_{0}(G)=\left\{u_{2}, u_{4}, u_{7}, u_{8}, u_{10}\right\}, D_{0}(G)=\left\{u_{1}, u_{3}, u_{5}, u_{6}, u_{9}, u_{12}\right\}$, and $P_{0}(G)=\left\{u_{11}, u_{13}\right\}$. Now all components in $G \backslash A_{0}(G)$ consist of a single vertex except $H$ (see Fig. 2). The single vertex is 0 -critical, and mult $(0, H)=0$.

Let $G$ be a graph. The deficiency of $G$, denoted by $\operatorname{def}(G)$, is defined to be the number of points left uncovered by any maximum matching. Let the number of odd components in $G$ be denoted by $o(G)$. Then $\operatorname{def}(G)=\max _{X \subseteq V(G)} o(G \backslash X)-|X|$ (see [11, Theorem 3.1.14 on p. 90]), and this is called the Berge-Tutte formula. Note that the multiplicity of 0 as a root of $\mu(G, x)$ is $|V(G)|$ minus the largest $r$ for which there is a matching of size $r$. Therefore mult $(0, G)=\operatorname{def}(G)$, and the following theorem follows.

Theorem 1.9. $\operatorname{mult}(0, G)=\max _{X \subseteq V(G)} o(G \backslash X)-|X|$.


Fig. 1.

Fig. 2.
Definition 1.10. Motivated by the Berge-Tutte formula, a barrier set is defined to be a set $X \subseteq V(G)$ for which mult $(0, G)=$ $o(G \backslash X)-|X|$. An extreme set is defined to be a set for which $\operatorname{mult}(0, G \backslash X)=\operatorname{mult}(0, G)+|X|$.

It should be noted that the standard terminology for a barrier set is a Tutte set in the classical matching theory.
Properties of extreme and barrier sets can be found in [11, Section 3.3]. In fact a barrier set is an extreme set. An extreme set is not necessarily a barrier set, but it can be shown that an extreme set is contained in some barrier set. In general, the union or intersection of two barrier sets is not a barrier set. However, it can be shown that the intersection of two (inclusionwise) maximal barrier sets is a barrier set. The $A_{0}(G)$ is both a barrier set and an extreme set. It can be shown that $A_{0}(G)$ is in fact the intersection of all the maximal barrier sets in $G$. We shall extend this fact to $A_{\theta}(G)$ (see Theorem 3.6).

In the next section, we prove a version of the Berge-Tutte formula extended to general $\theta$. Let the number of $\theta$-critical components in $G$ be denoted by $c_{\theta}(G)$.

Theorem 2.1 (Generalized Berge-Tutte Formula).

$$
\operatorname{mult}(\theta, G)=\max _{X \subseteq V(G)} c_{\theta}(G \backslash X)-|X| .
$$

Definition 1.11. Motivated by the generalized Berge-Tutte formula, we define a $\theta$-barrier set to be a set $X \subseteq V(G)$ for which $\operatorname{mult}(\theta, G)=c_{\theta}(G \backslash X)-|X|$.

We define a $\theta$-extreme set to be a set $X \subseteq V(G)$ for which $\operatorname{mult}(\theta, G \backslash X)=\operatorname{mult}(\theta, G)+|X|$.
The main theorem of this paper, which is proved in Section 3, is the following.
Theorem 3.6. If $N_{\theta}(G)=\varnothing$, then $A_{\theta}(G)$ is the intersection of all maximal $\theta$-barrier sets in $G$.
We emphasize that this paper is built up by generalizing some of the statements given in Chapter 3 of Lovász and Plummer's book [11] to the roots of the matching polynomial. Almost all proofs here have a resemblance to those found in [11]. The novelty of this paper is to merge the tools developed by Godsil [3] with the Lovász-Plummer investigations. This paper also fits into a series of papers [6-10] by the authors about the generalization of the results of classical matching theory with respect to the roots of the matching polynomial.

## 2. Properties of $\boldsymbol{\theta}$-barrier sets

An immediate consequence of part (a) of Theorems 1.2 and 1.7 is the following inequality, which we use frequently.

$$
\begin{equation*}
\operatorname{mult}(\theta, G) \geq c_{\theta}(G) \quad \text { for any graph } G . \tag{1}
\end{equation*}
$$

We prove the following analogue of the Berge-Tutte formula. The proof is similar to that of the generalization of Tutte's theorem due to the authors in [7]. For the sake of completeness, we repeat the statement.

Theorem 2.1 (Generalized Berge-Tutte Formula).

$$
\operatorname{mult}(\theta, G)=\max _{X \subseteq V(G)} c_{\theta}(G \backslash X)-|X| .
$$



Fig. 3.
Proof. We claim that $c_{\theta}(G \backslash X) \leq|X|+\operatorname{mult}(\theta, G)$ for all $X \subseteq V(G)$. If not, then $c_{\theta}(G \backslash X)>|X|+\operatorname{mult}(\theta, G)$ for some $X \subseteq V(G)$. Recall that mult $(\theta, G \backslash X) \geq c_{\theta}(G \backslash X)$. Together with Lemma 1.3, we have mult $(\theta, G) \geq \operatorname{mult}(\theta, G \backslash X)-|X|>$ $\operatorname{mult}(\theta, G)$, a contradiction.

Now it suffices to show that there is a set $X \subseteq V(G)$ for which mult $(\theta, G)=c_{\theta}(G \backslash X)-|X|$. Take $X=A_{\theta}$ ( $G$ ); by (ii) of Corollary 1.8 we are done.

Note that the definitions of 0-extreme set and extreme set coincide, but the definitions of 0-barrier set and barrier set are different. Our next proposition shows that a 0 -barrier set is a barrier set.

Proposition 2.2. A 0-barrier set is a barrier set.
Proof. If $X$ is a 0 -barrier set, then $c_{0}(G \backslash X)=\operatorname{mult}(0, G)+|X|$. Note that $c_{0}(G \backslash X) \leq o(G \backslash X)$. Using Theorem 1.9, we conclude that $o(G \backslash X)=\operatorname{mult}(0, G)+|X|$. Hence $X$ is a barrier set.

The converse of Proposition 2.2 is not true. The graph $G$ in Fig. 3 is well known (see [11, Figure 3.3 .1 on p. 105]). Note that $X=\{u, v\}$ is a barrier set in $G$, but it is not a 0 -barrier set.

A weak converse of Proposition 2.2 can be easily proved by using part (b) of Exercise 3.3.18 on p. 109 of [11].
Proposition 2.3. A (inclusionwise) maximal barrier set is a maximal 0-barrier set.
Now we shall study the properties of $\theta$-barrier and $\theta$-extreme sets.
Lemma 2.4. A subset of a $\theta$-extreme set is $a \theta$-extreme set.
Proof. Let $X$ be an $\theta$-extreme set, and consider $Y \subseteq X$. Now mult $(\theta, G \backslash X)=\operatorname{mult}(\theta, G)+|X|$. By Lemma 1.3, mult $(\theta, G \backslash Y)$ $\leq \operatorname{mult}(\theta, G)+|Y|$. If $Y$ is not $\theta$-extreme, then $\operatorname{mult}(\theta, G \backslash Y)<\operatorname{mult}(\theta, G)+|Y|$, and by Lemma 1.3 again, mult $(\theta, G \backslash X) \leq$ $\operatorname{mult}(\theta, G \backslash Y)+|X \backslash Y|<\operatorname{mult}(\theta, G)+|X|$, a contradiction. Hence a subset of an $\theta$-extreme set is $\theta$-extreme.

Lemma 2.5. If $X$ is a $\theta$-barrier [ $\theta$-extreme] set and $Y \subseteq X$, then $X \backslash Y$ is a $\theta$-barrier [ $\theta$-extreme] set in $G \backslash Y$.
Proof. Note that $c_{\theta}(G \backslash X)=|X|+\operatorname{mult}(\theta, G)$. By Theorem 2.1 and Lemma 1.3, $c_{\theta}(G \backslash X) \leq|X \backslash Y|+\operatorname{mult}(\theta, G \backslash Y) \leq$ $|X \backslash Y|+\operatorname{mult}(\theta, G)+|Y|=|X|+\operatorname{mult}(\theta, G)$. Hence $c_{\theta}(G \backslash X)=|X \backslash Y|+\operatorname{mult}(\theta, G \backslash Y)$, and $X \backslash Y$ is a $\theta$-barrier set in $G \backslash Y$.

Lemma 2.6. Every $\theta$-extreme set of $G$ lies in $a \theta$-barrier set.
Proof. If $X$ is a $\theta$-extreme set and $T=A_{\theta}(G \backslash X) \cup X$, then

$$
\begin{aligned}
c_{\theta}(G \backslash T) & =c_{\theta}\left(G \backslash\left(A_{\theta}(G \backslash X) \cup X\right)\right) \\
& =c_{\theta}\left((G \backslash X) \backslash A_{\theta}(G \backslash X)\right) \\
& =\left|A_{\theta}(G \backslash X)\right|+\operatorname{mult}(\theta, G \backslash X) \quad \text { (by (ii) of Corollary 1.8) } \\
& =\left|A_{\theta}(G \backslash X)\right|+\operatorname{mult}(\theta, G)+|X| \quad(X \text { is } \theta \text {-extreme) } \\
& =|T|+\operatorname{mult}(\theta, G),
\end{aligned}
$$

and hence $T$ is a $\theta$-barrier set.
Lemma 2.7. If $X$ is a $\theta$-barrier set, then $X$ is a $\theta$-extreme set.
Proof. Recall from (1) that $\operatorname{mult}(\theta, G \backslash X) \geq c_{\theta}(G \backslash X)$. Since $c_{\theta}(G \backslash X)=|X|+\operatorname{mult}(\theta, G)$, by Lemma 1.3, we have $\operatorname{mult}(\theta, G) \geq \operatorname{mult}(\theta, G \backslash X)-|X| \geq c_{\theta}(G \backslash X)-|X|=\operatorname{mult}(\theta, G)$.

Hence $\operatorname{mult}(\theta, G \backslash X)=\operatorname{mult}(\theta, G)+|X|$, and $X$ is a $\theta$-extreme set.
Note that in general a $\theta$-extreme set is not a $\theta$-barrier set. In Fig. $3, X_{1}=\{u\}$ is a 0 -extreme set but is not a 0 -barrier set. Furthermore, in Fig. 1, $X_{2}=\left\{u_{1}, u_{10}\right\}$ is a 1-extreme set but is not a 1-barrier set.

Lemma 2.8. If $X$ is a $\theta$-barrier set and $H$ is a component of $G \backslash X$, then either $H$ is $\theta$-critical or mult $(\theta, H)=0$.

Proof. Note that $c_{\theta}(G \backslash X)=|X|+\operatorname{mult}(\theta, G)$. By Lemma 2.7, $X$ is a $\theta$-extreme set. Therefore mult $(\theta, G \backslash X)=\operatorname{mult}(\theta, G)+$ $|X|=c_{\theta}(G \backslash X)$. Now, if $H$ is not $\theta$-critical and $\operatorname{mult}(\theta, H)>0$, then, by part (a) of Theorem 1.2, mult $(\theta, G \backslash X)>c_{\theta}(G \backslash X)$, a contradiction. Hence either $H$ is $\theta$-critical or $\operatorname{mult}(\theta, H)=0$.

Lemma 2.9. Let $X$ be a maximal $\theta$-barrier set. If $H$ is a component of $G \backslash X$ and mult $(\theta, H)=0$, then, for all $u \in V(H)$, $u$ is $\theta$-neutral in $H$. Furthermore, if $Y \subseteq V(H)$ and $Y \neq \varnothing$, then $c_{\theta}(H \backslash Y) \leq|Y|-1$.

Proof. If $H$ has a $\theta$-positive vertex, say $u$, then $\operatorname{mult}(\theta, H \backslash u)=1$. By (ii) of Corollary 1.8, $c_{\theta}\left((H \backslash u) \backslash A_{\theta}(H \backslash u)\right)=$ $\left|A_{\theta}(H \backslash u)\right|+\operatorname{mult}(\theta, H \backslash u)=\left|A_{\theta}(H \backslash u)\right|+1$. Now

$$
\begin{aligned}
c_{\theta}\left(G \backslash\left(X \cup\{u\} \cup A_{\theta}(H \backslash u)\right)\right) & =c_{\theta}(G \backslash X)+c_{\theta}\left((H \backslash u) \backslash A_{\theta}(H \backslash u)\right) \\
& =|X|+\operatorname{mult}(\theta, G)+\left|A_{\theta}(H \backslash u)\right|+1 \\
& =\left|X \cup\{u\} \cup A_{\theta}(H \backslash u)\right|+\operatorname{mult}(\theta, G),
\end{aligned}
$$

and so $X \cup\{u\} \cup A_{\theta}(H \backslash u)$ is a $\theta$-barrier in $G$, a contradiction to the maximality of $X$. Hence, for all $u \in V(H)$, $u$ is $\theta$-neutral in $H$.

Since $Y \neq \varnothing$, we may choose $y \in Y$. Let $Y^{\prime}=Y \backslash y$ and $H^{\prime}=H \backslash y$. Note that mult $(\theta, H \backslash y)=0$ since $y$ is $\theta$-neutral in H. By Theorem 2.1, $c_{\theta}\left(H^{\prime} \backslash Y^{\prime}\right) \leq\left|Y^{\prime}\right|$. Since $H \backslash Y=H^{\prime} \backslash Y^{\prime}$, we have $c_{\theta}(H \backslash Y) \leq|Y|-1$.

Lemma 2.10. If $G$ is $\theta$-critical, then, for all $Y \subseteq V(G)$ and $Y \neq \varnothing, c_{\theta}(G \backslash Y) \leq|Y|-1$.
Proof. Since $Y \neq \varnothing$, we may choose $y \in Y$. Let $Y^{\prime}=Y \backslash y$ and $G^{\prime}=G \backslash y$. Note that mult $(\theta, G \backslash y)=0$ since $y$ is $\theta$-essential in $G$. By Theorem 2.1, $c_{\theta}\left(G^{\prime} \backslash Y^{\prime}\right) \leq\left|Y^{\prime}\right|$. Since $G \backslash Y=G^{\prime} \backslash Y^{\prime}$, we have $c_{\theta}(G \backslash Y) \leq|Y|-1$.

In general, the union of two $\theta$-barrier sets is not necessarily a $\theta$-barrier set. In Fig. 3, $X_{3}=\{u, v, w\}$ and $X_{4}=\{v, w, z\}$ are two 0-barrier sets, but $X_{3} \cup X_{4}$ is not a 0-barrier set. In Fig. $1, X_{5}=\left\{u_{1}, u_{7}\right\}$ and $X_{6}=\left\{u_{1}, u_{10}\right\}$ are 1-barrier sets and $X_{5} \cup X_{6}$ is a 1-barrier set. Let $C_{3}$ be a cycle with three vertices. Every set containing a single vertex of $C_{3}$ is a 1-barrier set, but the union of two such sets is not 1-barrier set.

However, the intersection of two $\theta$-barrier sets is a $\theta$-barrier set. We shall prove this fact in Theorem 3.10. At present, let us use the results in this section to prove a weaker version.

Theorem 2.11. The intersection of two maximal $\theta$-barrier sets is a $\theta$-barrier set.
Proof. Let $X$ and $Y$ be two maximal $\theta$-barrier sets. Let $G_{1}, G_{2}, \ldots, G_{k}$ be the $\theta$-critical components of $G \backslash X$ and $H_{1}, H_{2}, \ldots, H_{m}$ be the components of $G \backslash Y$. Note that $k=|X|+\operatorname{mult}(\theta, G)$. Let $X_{i}=X \cap V\left(H_{i}\right), Y_{i}=Y \cap V\left(G_{i}\right)$, and $Z=X \cap Y$. By relabelling if necessary, we may assume that $X_{1}, \ldots, X_{m_{1}} \neq \varnothing$ and $Y_{1}, \ldots, Y_{k_{1}} \neq \varnothing$, but $X_{m_{1}+1}=\cdots=X_{m}=Y_{k_{1}+1}=\cdots=Y_{k}=\varnothing$, and also that $k_{1} \leq m_{1}$. Note that $G_{k_{1}+1}, \ldots, G_{k}$ are $\theta$-critical components in $(G \backslash X) \backslash Y$, so each is contained in a component of $G \backslash Y$.

Next we count the indices $i$ with $k_{1}+1 \leq i \leq k$ such that $G_{i}$ is contained in some $H_{j}$. If $m_{1}+1 \leq j \leq m$, then $H_{j}$ is a component in $(G \backslash X) \backslash Y$. So, if $G_{i} \subseteq H_{j}$, then $G_{i}=H_{j}$. Furthermore, $G_{i}$ is a component of $G \backslash Z$. By Theorem 2.1, the number of such $G_{i}$ 's is at most $c_{\theta}(G \backslash Z) \leq|Z|+\operatorname{mult}(\theta, G)$.

Suppose that $1 \leq j \leq m_{1}$. If $G_{i_{1}}, \ldots, G_{i_{t}}$ are contained in $H_{j}$, then they are $\theta$-critical components in $H_{j} \backslash X_{j}$. By Lemma 2.8, either $H_{j}$ is $\theta$-critical or $\operatorname{mult}(\theta, H)=0$. If $\operatorname{mult}(\theta, H)=0$, then, by Lemma 2.9, $c_{\theta}\left(H_{j} \backslash X_{j}\right) \leq\left|X_{j}\right|-1$. If $H_{i}$ is $\theta$-critical, then, by Lemma 2.10, $c_{\theta}\left(H_{j} \backslash X_{j}\right) \leq\left|X_{j}\right|-1$. Therefore, in either case, $t \leq\left|X_{j}\right|-1$.

The number of $G_{i}^{\prime}$ s where $k_{1}+1 \leq i \leq k$ that are disjoint from $Y$ is at most

$$
\begin{aligned}
c_{\theta}(G \backslash Z)+\sum_{j=1}^{m_{1}}\left(\left|X_{j}\right|-1\right) & \leq|Z|+\operatorname{mult}(\theta, G)+|X \backslash Z|-m_{1} \\
& =|X|+\operatorname{mult}(\theta, G)-m_{1} \\
& =k-m_{1} \\
& \leq k-k_{1} .
\end{aligned}
$$

Since this number is exactly $k-k_{1}$, we infer that equality must hold throughout. Hence $c_{\theta}(G \backslash Z)=|Z|+\operatorname{mult}(\theta, G)$, and $Z$ is a $\theta$-barrier set.

## 3. Characterizations of $\boldsymbol{A}_{\boldsymbol{\theta}}(\boldsymbol{G})$

A characterization of $A_{\theta}(G)$ is that it is the unique inclusion-minimal $\theta$-barrier set (see Theorem 3.5). If $N_{\theta}(G)=\varnothing$, then another characterization of $A_{\theta}(G)$ is that it is the intersection of all maximal $\theta$-barrier sets in $G$ (see Theorem 3.6).

Lemma 3.1. If $X$ is a $\theta$-barrier set or a $\theta$-extreme set, then $X \subseteq A_{\theta}(G) \cup P_{\theta}(G)$.


Fig. 4.
Proof. By Lemma 2.7, we may assume that $X$ is $\theta$-extreme. Let $x \in X$. By Lemma $2.4,\{x\}$ is a $\theta$-extreme set. Therefore $\operatorname{mult}(\theta, G \backslash x)=\operatorname{mult}(\theta, G)+1$, and $x$ is $\theta$-positive. Hence $x \in A_{\theta}(G) \cup P_{\theta}(G)$, and $X \subseteq A_{\theta}(G) \cup P_{\theta}(G)$.

Lemma 3.2. Let $X$ be a $\theta$-barrier set. If $X \subseteq A_{\theta}(G)$, then $X=A_{\theta}(G)$.
Proof. Note that $c_{\theta}(G \backslash X)=\operatorname{mult}(\theta, G)+|X|$. By Lemma 2.8, we conclude that $A_{\theta}(G \backslash X)=\varnothing$. By Theorem 1.6, $A_{\theta}(G \backslash X)$ $=A_{\theta}(G) \backslash X$. Hence $X=A_{\theta}(G)$.

We shall need the following result of Godsil [3].
Theorem 3.3 (Theorem 4.2 of [3]). If $\theta$ is a root of $\mu(G, x)$ with non-zero multiplicity $k$ and we let $u$ be a $\theta$-positive vertex in $G$, then
(a) if $v$ is $\theta$-essential in $G$, then it is $\theta$-essential in $G \backslash u$;
(b) if $v$ is $\theta$-positive in $G$, then it is $\theta$-essential or $\theta$-positive in $G \backslash u$;
(c) if $u$ is $\theta$-neutral in $G$, then it is $\theta$-essential or $\theta$-neutral in $G \backslash u$.

Lemma 3.4. If $u \in P_{\theta}(G)$, then $A_{\theta}(G) \subseteq A_{\theta}(G \backslash u)$.
Proof. If $A_{\theta}(G)=\varnothing$, then we are done. Suppose that $A_{\theta}(G) \neq \varnothing$. If $v \in A_{\theta}(G)$, then $v$ is adjacent to a $\theta$-essential vertex $w$. By Theorem 3.3, $w$ is $\theta$-essential in $G \backslash u$, and $v$ is either $\theta$-positive or $\theta$-essential in $G \backslash u$. If $v$ is $\theta$-essential in $G \backslash u$, then $\operatorname{mult}(\theta, G \backslash u v)=\operatorname{mult}(\theta, G)$. By Theorem 1.6, $u \in P_{\theta}(G)=P_{\theta}(G \backslash v$ ). Since $v$ is $\theta$-special in $G, v$ is $\theta$-positive in $G$ (see Corollary 4.3 of [3]). Hence $\operatorname{mult}(\theta, G \backslash u v)=\operatorname{mult}(\theta, G)+2$, a contradiction. Therefore $v$ is $\theta$-positive in $G \backslash u$. Since $v$ is adjacent to $w$, we must have $v \in A_{\theta}(G \backslash u)$. Hence $A_{\theta}(G) \subseteq A_{\theta}(G \backslash u)$.

Theorem 3.5. If $X$ is a $\theta$-barrier set in $G$, then $A_{\theta}(G) \subseteq X$. In particular, $A_{\theta}(G)$ is the unique minimal $\theta$-barrier set.
Proof. By Lemma 3.1, $X \subseteq A_{\theta}(G) \cup P_{\theta}(G)$. We shall prove the result by induction on $\left|X \cap P_{\theta}(G)\right|$. If $\left|X \cap P_{\theta}(G)\right|=0$, then $X \subseteq A_{\theta}(G)$, and, by Lemma 3.2, $X=A_{\theta}(G)$. Suppose that $\left|X \cap P_{\theta}(G)\right| \geq 1$. We may assume that, if $X^{\prime}$ is a $\theta$-barrier set in $G^{\prime}$ with $\left|X^{\prime} \cap P_{\theta}\left(G^{\prime}\right)\right|<\left|X \cap P_{\theta}(G)\right|$, then $A_{\theta}\left(G^{\prime}\right) \subseteq X^{\prime}$.

Let $x \in X \cap P_{\theta}(G)$. By Lemma 2.5, $X^{\prime}=X \backslash x$ is a $\theta$-barrier set in $G^{\prime}=G \backslash x$. By Lemmas 3.1 and 3.4, we have $X^{\prime} \subseteq A_{\theta}\left(G^{\prime}\right) \cup P_{\theta}\left(G^{\prime}\right)$ and $A_{\theta}(G) \subseteq A_{\theta}\left(G^{\prime}\right)$. Therefore $\left|X^{\prime} \cap P_{\theta}\left(G^{\prime}\right)\right|<\left|X \cap P_{\theta}(G)\right|$. By the induction hypothesis, $A_{\theta}\left(G^{\prime}\right) \subseteq X^{\prime}$. Hence $A_{\theta}(G) \subseteq X$.

In general, $A_{\theta}(G)$ is not the intersection of all maximal $\theta$-barrier sets in $G$. For instance, in Fig. 4 , mult $(\sqrt{3}, G)=0$ and $A_{\sqrt{3}}(G)=\varnothing$. Now $\{u\}$ is the only maximal $\sqrt{3}$-barrier set, but $A_{\sqrt{3}}(G) \neq\{u\}$. However, we can show that $A_{\theta}(G)$ is the intersection of all maximal $\theta$-barrier sets in $G$ if $N_{\theta}(G)=\varnothing$.

Theorem 3.6. If $N_{\theta}(G)=\varnothing$, then $A_{\theta}(G)$ is the intersection of all maximal $\theta$-barrier sets in $G$.
Proof. By Theorem 3.5, $A_{\theta}(G)$ is contained in the intersection of all maximal $\theta$-barriers in $G$. It is sufficient to show that for each $x \in V(G) \backslash A_{\theta}(G)$ there is a maximal barrier that does not contain $x$. If $x \in D_{\theta}(G)$, then, by Lemma 3.1, $x$ is not contained in any $\theta$-barriers and thus any maximal $\theta$-barriers. If $x \in P_{\theta}(G)$, then $x$ is contained in a component $H$ in $G \backslash A_{\theta}(G)$ with $\operatorname{mult}(\theta, H)=0$. Note that $|V(H)| \geq 2$ for $x \in P_{\theta}(G)=P\left(G \backslash A_{\theta}(G)\right.$ ), and $\operatorname{mult}(\theta, H \backslash x)=1$ (see Theorem 1.6). By (c) of Theorem 1.2 and the fact that $\operatorname{mult}(\theta, H)=0$, we deduce that there is a vertex $y \in V(H \backslash x)$ for which mult $(\theta, H \backslash x y)=0$. Now $y \in P_{\theta}(G)$ for $N_{\theta}(G)=\varnothing$. Furthermore, $x$ is $\theta$-essential in $H \backslash y$. Therefore $x \notin A_{\theta}(H \backslash y$ ) and, by (ii) of Corollary 1.8, $c_{\theta}\left((H \backslash y) \backslash A_{\theta}(H \backslash y)\right)=\left|A_{\theta}(H \backslash y)\right|+1$. Hence

$$
\begin{aligned}
c_{\theta}\left(G \backslash\left(A_{\theta}(G) \cup\{y\} \cup A_{\theta}(H \backslash y)\right)\right) & =c_{\theta}\left(G \backslash A_{\theta}(G)\right)+c_{\theta}\left((H \backslash y) \backslash A_{\theta}(H \backslash y)\right) \\
& =\left|A_{\theta}(G)\right|+\operatorname{mult}(\theta, G)+\left|A_{\theta}(H \backslash y)\right|+1 \\
& =\left|A_{\theta}(G) \cup\{y\} \cup A_{\theta}(H \backslash y)\right|+\operatorname{mult}(\theta, G),
\end{aligned}
$$

and so $A_{\theta}(G) \cup\{y\} \cup A_{\theta}(H \backslash y)$ is a $\theta$-barrier set not containing $x$. Let $Z$ be a maximal $\theta$-barrier set containing $Y=$ $A_{\theta}(G) \cup\{y\} \cup A_{\theta}(H \backslash y)$. By Lemma $2.5, Z \backslash Y$ is a $\theta$-barrier set in $G \backslash Y$. Using Theorem 1.6 and the fact that $x$ is $\theta$-essential in $H \backslash y$, we can deduce that $x \in D_{\theta}(G \backslash Y)$. By Lemma 3.1, we conclude that $x \notin Z \backslash Y$, and hence $x \notin Z$. The proof of the theorem is completed.

Since $N_{0}(G)=\varnothing$, by Theorem 3.6 and Proposition 2.3, we deduce the following classical result.

## Corollary 3.7 (Theorem 3.3 .15 of [11]). $A_{0}(G)$ is the intersection of all maximal barrier sets in $G$.

Finally, we prove that the intersection of two $\theta$-barrier sets is a $\theta$-barrier set. We shall need the following two lemmas.
Lemma 3.8. A set $X \subseteq V(G)$ is a $\theta$-barrier set in $G$ if and only if $X \cap H$ is a $\theta$-barrier set in $H$ for each component $H$ of $G$.
Proof. Let $H_{1}, \ldots, H_{m}$ be the components of $G$. Note that $c_{\theta}(G \backslash X)=\sum_{i=1}^{m} c_{\theta}\left(H_{i} \backslash X\right)$. By part (a) of Theorem 1.2, mult $(\theta, G)$ $=\sum_{i=1}^{m} \operatorname{mult}\left(\theta, H_{i}\right)$ and $\operatorname{mult}(\theta, G \backslash X)=\sum_{i=1}^{m} \operatorname{mult}\left(\theta, H_{i} \backslash X\right)$.
$(\Leftarrow)$ If $X \cap H_{i}$ is a $\theta$-barrier set in $H_{i}$ for all $i$, then $\operatorname{mult}\left(\theta, H_{i}\right)=c_{\theta}\left(H_{i} \backslash X\right)-\left|H_{i} \cap X\right|$. Therefore mult $(\theta, G)=\sum_{i=1}^{m}$ $\left(c_{\theta}\left(H_{i} \backslash X\right)-\left|H_{i} \cap X\right|\right)=c_{\theta}(G \backslash X)-|X|$, and $X$ is a $\theta$-barrier set in $G$.
$(\Rightarrow)$ If $X$ is a $\theta$-barrier set in $G$, then mult $(\theta, G)=c_{\theta}(G \backslash X)-|X|$. So $\sum_{i=1}^{m} \operatorname{mult}\left(\theta, H_{i}\right)=\sum_{i=1}^{m}\left(c_{\theta}\left(H_{i} \backslash X\right)-\left|H_{i} \cap X\right|\right)$ and $\sum_{i=1}^{m}\left(\operatorname{mult}\left(\theta, H_{i}\right)-\left(c_{\theta}\left(H_{i} \backslash X\right)-\left|H_{i} \cap X\right|\right)\right)=0$. By Theorem 2.1, each summand on the left in the last equation must be non-negative. We thus conclude that mult $\left.\theta, H_{i}\right)=c_{\theta}\left(H_{i} \backslash X\right)-\left|H_{i} \cap X\right|$, and $X \cap H_{i}$ is a $\theta$-barrier set in $H_{i}$ for all $i$.

Lemma 3.9. If $B$ is a $\theta$-barrier set in $G$ with $X=B \cup T$ for some $T \subseteq V(G \backslash B)$, then $X$ is $a \theta$-barrier set in $G$ if and only if $T$ is $a$ $\theta$-barrier set in $G \backslash B$.
Proof. First note that, by Lemma 2.7, $\operatorname{mult}(\theta, G \backslash B)=\operatorname{mult}(\theta, G)+|B|$.
$(\Leftarrow)$ If $T$ is a $\theta$-barrier set in $G \backslash B$, then $\operatorname{mult}(\theta, G \backslash B)=c_{\theta}(G \backslash(B \cup T))-|T|$, and so mult $(\theta, G)=c_{\theta}(G \backslash(B \cup T))-|B \cup T|$. Hence $X$ is a $\theta$-barrier set in $G$.
$(\Rightarrow)$ If $X$ is a $\theta$-barrier set in $G$, then $\operatorname{mult}(\theta, G)=c_{\theta}(G \backslash(B \cup T))-|B \cup T|$, and therefore mult $(\theta, G \backslash B)=c_{\theta}(G \backslash(B \cup T))-|T|$. Hence $T$ is a $\theta$-barrier set in $G \backslash B$.

Theorem 3.10. The intersection of two $\theta$-barrier sets is a $\theta$-barrier set.
Proof. Let $B_{1}$ and $B_{2}$ be two $\theta$-barrier sets. By Theorem 3.5, $B_{1}=A_{\theta}(G) \cup T_{1}$ for some $T_{1} \subseteq V\left(G \backslash A_{\theta}(G)\right)$. By Lemma 3.9, $T_{1}$ is a $\theta$-barrier set in $G \backslash A_{\theta}(G)$. Similarly, $B_{2}=A_{\theta}(G) \cup T_{2}$ for some $\theta$-barrier set $T_{2}$ in $G \backslash A_{\theta}(G)$. Now $B_{1} \cap B_{2}=A_{\theta}(G) \cup\left(T_{1} \cap T_{2}\right)$. By Lemma 3.9, it suffices to show that $T_{1} \cap T_{2}$ is a $\theta$-barrier set in $G \backslash A_{\theta}(G)$.

If $\left|T_{1} \cap T_{2}\right|=0$, then $T_{1} \cap T_{2}=\varnothing$ and we are done (for an empty set is a $\theta$-barrier set). Suppose $\left|T_{1} \cap T_{2}\right| \geq 1$. Assume that, if $T_{3}$ and $T_{4}$ are $\theta$-barrier sets in $G \backslash A_{\theta}(G)$ and $\left|T_{3} \cap T_{4}\right|<\left|T_{1} \cap T_{2}\right|$, then $T_{3} \cap T_{4}$ is a $\theta$-barrier set in $G \backslash A_{\theta}(G)$.

Let $x \in T_{1} \cap T_{2}$. Since $T_{1}$ is a $\theta$-barrier set in $G \backslash A_{\theta}(G)$, by part (iii) of Corollary 1.8 and Lemma 3.1, we deduce that $x \in H$, where $H$ is a component of $G \backslash A_{\theta}(G)$ with $\operatorname{mult}(\theta, H)=0$. Furthermore, mult $(\theta, H \backslash x)=1$, and we deduce that $c_{\theta}\left((H \backslash x) \backslash A_{\theta}(H \backslash x)\right)=1+\left|A_{\theta}(H \backslash x)\right|=\left|\{x\} \cup A_{\theta}(H \backslash x)\right|$ (by Corollary 1.8 and part (a) of Theorem 1.2). Therefore $\{x\} \cup A_{\theta}(H \backslash x)$ is a $\theta$-barrier set in $H$.

On the other hand, $T_{1} \cap H$ is a $\theta$-barrier set in $H$ by Lemma 3.8. By Lemma 2.5, $\left(T_{1} \cap H\right) \backslash x$ is a $\theta$-barrier set in $H \backslash x$, which yields $A_{\theta}(H \backslash x) \subseteq\left(T_{1} \cap H\right) \backslash x$ (Theorem 3.5). Therefore, we may let $T_{1} \cap H=T_{3} \cup\left(\{x\} \cup A_{\theta}(H \backslash x)\right.$ ) for some $T_{3} \subseteq V\left(H \backslash\left(\{x\} \cup A_{\theta}(H \backslash x)\right)\right)$. Moreover, since $T_{1} \cap H$ is a $\theta$-barrier set in $H, T_{3}$ is a $\theta$-barrier set in $H$ (Lemma 3.9). In fact, it is not hard to see that $T_{3}$ is a $\theta$-barrier set in $G \backslash A_{\theta}(G)$ (because $H$ is a component of $G \backslash A_{\theta}(G)$ ). Similarly, $T_{2} \cap H=T_{4} \cup\left(\{x\} \cup A_{\theta}(H \backslash x)\right)$ for some $\theta$-barrier set $T_{4}$ in $H$ that is also a $\theta$-barrier set in $G \backslash A_{\theta}(G)$. Clearly $\left|T_{3} \cap T_{4}\right|<\left|T_{1} \cap T_{2}\right|$. By the induction hypothesis, we conclude that $T_{3} \cap T_{4}$ is a $\theta$-barrier set in $G \backslash A_{\theta}(G)$, and thus is also a $\theta$-barrier set in $H$. Since $T_{1} \cap T_{2} \cap H=\left(T_{3} \cap T_{4}\right) \cup\left(\{x\} \cup A_{\theta}(H \backslash x)\right.$ ), we deduce from Lemma 3.9 that $T_{1} \cap T_{2} \cap H$ is a $\theta$-barrier set in $H$.

Now if $H^{\prime}$ is a $\theta$-critical component of $G \backslash A_{\theta}(G)$, then $T_{1} \cap T_{2} \cap H^{\prime}=\varnothing$, and so $T_{1} \cap T_{2} \cap H^{\prime}$ is a $\theta$-barrier set in $H^{\prime}$. Thus $T_{1} \cap T_{2} \cap H^{\prime \prime}$ is $\theta$-barrier set in $H^{\prime \prime}$ for any component $H^{\prime \prime}$ of $G \backslash A_{\theta}(G)$, and so $T_{1} \cap T_{2}$ is a $\theta$-barrier set in $G \backslash A_{\theta}(G)$ (Lemma 3.8).

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