



Note

An eigenvalue bound for the Laplacian of a graph

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Abstract

We present a lower bound for the smallest non-zero eigenvalue of the Laplacian of an undirected graph. The bound is primarily useful for graphs with small diameter.

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1. Introduction

We use the following definition for the Laplacian matrix of a graph, consistent with [3]:

Definition. Let G be an undirected graph with adjacency matrix \mathbf{A} , and let \mathbf{D} be the diagonal degree matrix defined by $d_{ii} = \deg(v_i)$ and $d_{ik} = 0$ for $i \neq k$. The *Laplacian* of G is the matrix $\mathcal{L} = \mathbf{I} - \mathbf{D}^{-1/2}\mathbf{A}\mathbf{D}^{-1/2}$.

The eigenvalues of \mathcal{L} are in the range $[0, 2]$. Zero is always an eigenvalue with multiplicity equal to the number of connected components of G , and 2 occurs as an eigenvalue if and only if G is bipartite. The eigenvalues of \mathcal{L} contain additional information regarding the structure of G . They can be used to establish bounds on the diameter of G as well as distances between subsets of G [1,4,2,5]. The magnitudes of the eigenvalues also determine the rate of convergence of various iterative computations such as those described in [6,7]; it is therefore desirable to find bounds on the eigenvalues themselves. One of the best lower bounds for

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the smallest non-zero eigenvalue λ_1 is established in [3]

$$\lambda_1 \geq \frac{1}{D \operatorname{vol}(G)}. \quad (1)$$

Here D is the diameter of G , and $\operatorname{vol}(G)$ is the sum of the degrees of all vertices. In this paper, we present a lower bound on λ_1 which is easy to compute and is tighter than (1) for certain graphs with low diameter.

2. A lower bound for λ_1

Consider the similar matrix $\mathbf{L} = \mathbf{D}^{-1/2} \mathcal{L} \mathbf{D}^{1/2} = \mathbf{I} - \mathbf{D}^{-1} \mathbf{A}$ which has the same eigenvalues as \mathcal{L} . If f is an eigenfunction of \mathbf{L} corresponding to eigenvalue λ then for any vertex v

$$(1 - \lambda)f(v) = \frac{1}{\deg(v)} \sum_{u \sim v} f(u), \quad (2)$$

where $u \sim v$ denotes that the vertices u, v are connected. Let v_1, v_2, \dots, v_{m+1} be a sequence of connected vertices such that $f(v_1)$ is maximal and $f(v_{m+1}) \leq 0$. For convenience, set $x_i = f(v_i)$. Let $\alpha = 1 - \lambda$, and let d be the maximum degree of any vertex of G . Since $v_1 \sim v_2$ and $f(v_1)$ is maximal, Eq. (2) gives us

$$\alpha x_1 = \frac{1}{\deg(v_1)} \sum_{u \sim v_1} f(u) \leq \frac{x_2}{\deg(v_1)} + \frac{(\deg(v_1) - 1)x_1}{\deg(v_1)} \leq \frac{x_2}{d} + \frac{(d - 1)x_1}{d}. \quad (3)$$

Similarly, since $v_i \sim v_{i-1}$ and $v_i \sim v_{i+1}$ for $2 \leq i \leq m$, we have

$$\alpha x_i \leq \frac{x_{i-1} + x_{i+1}}{d} + \frac{(d - 2)x_i}{d}. \quad (4)$$

Scaling f if necessary, we may assume that $x_1 = 1$ and rewrite inequalities (3) and (4) as

$$\begin{aligned} x_2 &\geq 1 - \lambda d, \\ x_{i+1} &\geq \alpha d x_i - x_{i-1} - (d - 2). \end{aligned} \quad (5)$$

We assume that $\lambda < 1$ and $\alpha d \geq 1$; otherwise we have the bound $\lambda > (d - 1)/d$ which is much better than the one we will derive.

Lemma. For $3 \leq k \leq m + 1$ we have $x_k \geq 1 - \lambda \alpha^{k-3} d^{k-2} - \lambda \alpha^{k-2} d^{k-1}$.

Proof. Our proof is by induction. Setting $i = 2$ in inequality (5) establishes the base case:

$$x_3 \geq \alpha d x_2 - 1 - (d - 2) \geq \alpha d(1 - \lambda d) + 1 - d = 1 - \lambda d - \lambda \alpha d^2.$$

Now suppose the inequality holds for $k \leq i$, where $i \geq 3$. If we add the inequalities (5) up to $i + 1$ we obtain

$$\begin{aligned} x_2 &\geq 1 - \lambda d, \\ x_3 &\geq \alpha d x_2 - 1 - (d - 2), \\ x_4 &\geq \alpha d x_3 - x_2 - (d - 2), \\ &\vdots \\ x_i &\geq \alpha d x_{i-1} - x_{i-2} - (d - 2), \\ &\quad + x_{i+1} \geq \alpha d x_i - x_{i-1} - (d - 2), \\ x_{i+1} &\geq (\alpha d - 2)(x_2 + x_3 + \dots + x_{i-1} + x_i) + x_i - \lambda d - (i - 1)(d - 2), \end{aligned}$$

so by induction

$$\begin{aligned} x_{i+1} &\geq (\alpha d - 2)((i - 1) - 2\lambda d - 2\lambda \alpha d^2 - \dots - 2\lambda \alpha^{i-3} d^{i-2} - \lambda \alpha^{i-2} d^{i-1}) \\ &\quad + (1 - \lambda \alpha^{i-3} d^{i-2} - \lambda \alpha^{i-2} d^{i-1}) - \lambda d - (i - 1)(d - 2) \\ &= 1 - (i - 4)\lambda d + 2\lambda \alpha d^2 + 2\lambda \alpha^2 d^3 + \dots + 2\lambda \alpha^{i-4} d^{i-3} + \lambda \alpha^{i-3} d^{i-2} \\ &\quad - \lambda \alpha^{i-2} d^{i-1} - \lambda \alpha^{i-1} d^i \\ &\geq 1 - (i - 4)\lambda d + 2\lambda d + 2\lambda d + \dots + 2\lambda d + \lambda d - \lambda \alpha^{i-2} d^{i-1} - \lambda \alpha^{i-1} d^i \\ &= 1 + (i - 3)\lambda d - \lambda \alpha^{i-2} d^{i-1} - \lambda \alpha^{i-1} d^i \geq -\lambda \alpha^{i-2} d^{i-1} - \lambda \alpha^{i-1} d^i. \quad \square \end{aligned}$$

Theorem. Let G be a graph with diameter D and maximum vertex degree d . Then

$$\lambda_1 \geq \frac{1}{(d + 1)d^{\lceil D/2 \rceil - 1}}. \tag{6}$$

Proof. Since the distance from a vertex which maximizes f to one which minimizes f is at most D , we can obtain the described sequence of connected vertices with $m \leq \lceil D/2 \rceil$, negating f if necessary. Then applying the lemma to x_{m+1} we have

$$0 \geq x_{m+1} \geq 1 - \lambda_1 \alpha^{m-2} d^{m-1} - \lambda_1 \alpha^{m-1} d^m \geq 1 - \lambda_1 d^{m-1} - \lambda_1 d^m$$

from which the result follows. \square

3. Discussion

Since the size of the denominator of (6) is exponential in D , the bound is primarily useful for graphs of low diameter. Note that in the special case $D = 2$ we have $0 \geq x_2 \geq 1 - \lambda_1 d$ which gives us the improved bound

$$\lambda_1 \geq \frac{1}{d}.$$

Table 1
Maximum value of d for which the bound (6) is necessarily tighter on d -regular graphs

D	5	6	7	8	9	10
d	10	18	7	5	3	3

In general, the bound (6) will be tighter than (1) when $(d + 1)d^{\lceil D/2 \rceil - 1} \leq D \text{vol}(G)$. For d -regular graphs in which all vertices have degree d , this condition becomes

$$(d + 1)d^{\lceil D/2 \rceil - 2} \leq D|G|. \quad (7)$$

This is always true for $D = 3, 4$. For $D = 5$, a d -regular graph has at least $2d + 4$ vertices, so condition (7) will certainly be satisfied if $(d + 1)d \leq 5(2d + 4)$ which is true for $d \leq 10$; for $D = 6$ a d -regular graph has at least $3d + 3$ vertices, so (7) will be satisfied if $(d + 1)d \leq 6(3d + 3)$, i.e. if $d \leq 18$. Continuing in this manner we obtain Table 1 which lists, for $5 \leq D \leq 10$, the maximum value of d for which condition (7) is guaranteed to be satisfied.

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