

On $(1 - u)$ -cyclic codes over $\mathbb{F}_{p^k} + u\mathbb{F}_{p^k}$

Maria Carmen V. Amarra*, Fidel R. Nemenzo

Institute of Mathematics, University of the Philippines, Diliman, Quezon City, Philippines

Received 27 July 2007; accepted 27 July 2007

Abstract

We extend the results of [J.F. Qian, L.N. Zhang, S.X. Zhu, $(1 + u)$ -constacyclic and cyclic codes over $\mathbb{F}_2 + u\mathbb{F}_2$, Appl. Math. Lett. 19 (2006) 820–823. [3]] to codes over the commutative ring $R = \mathbb{F}_{p^k} + u\mathbb{F}_{p^k}$, where p is prime, $k \in \mathbb{N}$ and $u^2 = 0$. In particular, we prove that the Gray image of a linear $(1 - u)$ -cyclic code over R of length n is a distance-invariant quasicyclic code of index p^{k-1} and length $p^k n$ over \mathbb{F}_{p^k} . We also prove that if $(n, p) = 1$, then every code of length $p^k n$ over \mathbb{F}_{p^k} which is the Gray image of a linear cyclic code of length n over R is permutation-equivalent to a quasicyclic code of index p^{k-1} .

© 2008 Elsevier Ltd. All rights reserved.

Keywords: Cyclic and quasicyclic codes; Gray map; Finite rings

1. Preliminaries

Let p be prime and $k \in \mathbb{N}$. Let R be the ring $\mathbb{F}_{p^k} + u\mathbb{F}_{p^k}$, where $u^2 = 0$ and $\mathbb{F}_{p^k} = GF(p^k)$. Then R is a finite chain ring with maximal ideal uR and residue field \mathbb{F}_{p^k} . Let \mathcal{C} be a code of length n over R , and $P(\mathcal{C})$ be its polynomial representation, i.e.,

$$P(\mathcal{C}) = \left\{ \sum_{i=0}^{n-1} r_i x^i \mid (r_0, \dots, r_{n-1}) \in \mathcal{C} \right\}.$$

Let σ and ν be maps from R^n to R^n given by

$$\sigma(r_0, r_1, \dots, r_{n-1}) = (r_{n-1}, r_0, \dots, r_{n-2})$$

and

$$\nu(r_0, r_1, \dots, r_{n-1}) = ((1 - u)r_{n-1}, r_0, \dots, r_{n-2}).$$

Then \mathcal{C} is said to be *cyclic* if $\sigma(\mathcal{C}) = \mathcal{C}$, and $(1 - u)$ -*cyclic* if $\nu(\mathcal{C}) = \mathcal{C}$. It is known that:

Theorem 1.1. *A code \mathcal{C} of length n over R is cyclic if and only if $P(\mathcal{C})$ is an ideal of $R[x]/\langle x^n - 1 \rangle$.*

* Corresponding address: School of Mathematics and Statistics, The University of Western Australia, Perth, Australia.

E-mail addresses: mcamarra@maths.uwa.edu.au, carmen_amarra@yahoo.com (M.C.V. Amarra), fidel@math.upd.edu.ph (F.R. Nemenzo).

Theorem 1.2. A code \mathcal{C} of length n over R is $(1 - u)$ -cyclic if and only if $P(\mathcal{C})$ is an ideal of $R[x]/\langle x^n - (1 - u) \rangle$.

Let $\mathbf{a} \in \mathbb{F}_{p^k}^{pn}$, with $\mathbf{a} = (a_0, \dots, a_{p^k n-1}) = (a^{(0)} \mid \dots \mid a^{(p^k-1)})$, $a^{(i)} \in \mathbb{F}_{p^k}^{pn}$ for all $i = 0, \dots, p^k - 1$. Let $\sigma^{\otimes p^k-1}$ be the map from $\mathbb{F}_{p^k}^{pn}$ to $\mathbb{F}_{p^k}^{pn}$ given by

$$\sigma^{\otimes p^k-1}(\mathbf{a}) = (\tilde{\sigma}(a^{(0)}) \mid \dots \mid \tilde{\sigma}(a^{(p^k-1)})),$$

where $\tilde{\sigma}$ is the usual cyclic shift $(c_0, c_1, \dots, c_{pn-1}) \mapsto (c_{pn-1}, c_0, \dots, c_{pn-2})$ on $\mathbb{F}_{p^k}^{pn}$. A code $\tilde{\mathcal{C}}$ of length $p^k n$ over \mathbb{F}_{p^k} is said to be *quasicyclic of index p^{k-1}* if $\sigma^{\otimes p^k-1}(\tilde{\mathcal{C}}) = \tilde{\mathcal{C}}$.

In [1], a homogeneous weight on arbitrary finite chain rings is defined; we give it here for the case of the ring $\mathbb{F}_{p^k} + u\mathbb{F}_{p^k}$. The homogeneous weight $w_{\text{hom}}(r)$ of $r \in R$ is given by

$$w_{\text{hom}}(r) = \begin{cases} p^k - 1 & \text{if } r \in R \setminus Ru \\ p^k & \text{if } r \in Ru \setminus \{0\} \\ 0 & \text{otherwise.} \end{cases}$$

This extends to a weight function in R^n : if $\mathbf{r} = (r_0, \dots, r_{n-1}) \in R^n$, then

$$w_{\text{hom}}(\mathbf{r}) := \sum_{i=0}^{n-1} w_{\text{hom}}(r_i).$$

The homogeneous distance $d_{\text{hom}}(\mathbf{x}, \mathbf{y})$ between any distinct vectors $\mathbf{x}, \mathbf{y} \in R^n$ is defined to be $w_{\text{hom}}(\mathbf{x} - \mathbf{y})$.

2. Gray images of $(1 - u)$ -cyclic codes over $\mathbb{F}_{p^k} + u\mathbb{F}_{p^k}$

Any element $\epsilon \in \mathbb{Z}_{p^k}$ has p -adic representation

$$\gamma_{0,\epsilon} + \gamma_{1,\epsilon}p + \dots + \gamma_{(k-1),\epsilon}p^{k-1},$$

where $\gamma_{i,\epsilon} \in \{0, 1, 2, \dots, p - 1\}$. If α is a fixed primitive element of \mathbb{F}_{p^k} , then corresponding to every $\epsilon \in \mathbb{Z}_{p^k}$ is an element $\alpha_\epsilon \in \mathbb{F}_{p^k}$, given by

$$\alpha_\epsilon := \gamma_{0,\epsilon} + \gamma_{1,\epsilon}\alpha + \dots + \gamma_{(k-1),\epsilon}\alpha^{k-1}.$$

The *Gray map* Φ on R , which is a special case of the Gray map defined in [1], is given by

$$\begin{aligned} \Phi : R^n &\longrightarrow \mathbb{F}_{p^k}^{p^k n} \\ x + yu &\longmapsto (y, \alpha_1 x \oplus y, \dots, \alpha_{p^k-1} x \oplus y), \end{aligned}$$

where \oplus is componentwise addition in \mathbb{F}_{p^k} . The Gray map Φ is an isometry from (R^n, d_{hom}) to $\mathbb{F}_{p^k}^{p^k n}$ under the Hamming distance. It is also clear that Φ preserves linearity of codes.

From the above definitions, we have:

Proposition 2.1.

$$\Phi \circ \nu = \sigma^{\otimes p^k-1} \circ \Phi.$$

Proof. Let $\mathbf{r} = (r_0, \dots, r_{n-1}) \in R^n$ and $\mathbf{x} = (x_0, x_1, \dots, x_{n-1})$, $\mathbf{y} = (y_0, y_1, \dots, y_{n-1}) \in \mathbb{F}_{p^k}^n$ such that $\mathbf{r} = \mathbf{x} + \mathbf{y}u$.

Let $\Phi(\mathbf{r}) = (a_0, \dots, a_{p^k n-1})$. Then $\sigma^{\otimes p^k-1}(\Phi(\mathbf{r})) = (b_0, \dots, b_{p^k n-1})$, where

$$b_{\epsilon n+j} = \begin{cases} a_{(\epsilon+p-1)n+(n-1)} & \text{if } j = 0, \gamma_{0,\epsilon} = 0 \\ a_{(\epsilon-1)n+(n-1)} & \text{if } j = 0, \gamma_{0,\epsilon} \neq 0 \\ a_{\epsilon n+j-1} & \text{if } j \neq 0. \end{cases}$$

On the other hand,

$$v(\mathbf{r}) = ((1 - u)r_{n-1}, r_0, \dots, r_{n-2}),$$

where $(1 - u)r_{n-1} = x_{n-1} + u(-x_{n-1} \oplus y_{n-1})$. Then $\Phi(v(\mathbf{r})) = (c_0, \dots, c_{p^k n-1})$, where

$$c_{\epsilon n+j} = \begin{cases} \left(\sum_{i=0}^{k-1} \gamma_{i,\epsilon} \alpha^i - 1 \right) x_{n-1} \oplus y_{n-1}, & \text{if } j = 0 \\ \left(\sum_{i=0}^{k-1} \gamma_{i,\epsilon} \alpha^i \right) x_{j-1} \oplus y_{j-1}, & \text{if } j \neq 0 \end{cases}$$

$$= \begin{cases} \left(\sum_{i=0}^{k-1} \gamma_{i,\epsilon} \alpha^i + p - 1 \right) x_{n-1} \oplus y_{n-1}, & \text{if } j = 0, \gamma_{0,\epsilon} = 0 \\ \left(\sum_{i=0}^{k-1} \gamma_{i,\epsilon} \alpha^i - 1 \right) x_{n-1} \oplus y_{n-1}, & \text{if } j = 0, \gamma_{0,\epsilon} \neq 0 \\ \left(\sum_{i=0}^{k-1} \gamma_{i,\epsilon} \alpha^i \right) x_{j-1} \oplus y_{j-1}, & \text{if } j \neq 0. \end{cases}$$

The conclusion follows. \square

As a consequence:

Theorem 2.2. A code \mathcal{C} of length n over R is $(1 - u)$ -cyclic if and only if $\Phi(\mathcal{C})$ is quasicyclic of index p^{k-1} and length $p^k n$ over \mathbb{F}_{p^k} .

Proof. Suppose \mathcal{C} is $(1 - u)$ -cyclic. Then

$$\sigma^{\otimes p^{k-1}}(\Phi(\mathcal{C})) = \Phi(v(\mathcal{C})) = \Phi(\mathcal{C}),$$

so $\Phi(\mathcal{C})$ is quasicyclic of index p^{k-1} . Conversely, if $\Phi(\mathcal{C})$ is quasicyclic of index p^{k-1} , then

$$\Phi(v(\mathcal{C})) = \sigma^{\otimes p^{k-1}}(\Phi(\mathcal{C})) = \Phi(\mathcal{C}).$$

Since Φ is injective, it follows that $v(\mathcal{C}) = \mathcal{C}$. \square

3. Gray images of cyclic codes over $\mathbb{F}_{p^k} + u\mathbb{F}_{p^k}$

Suppose that $(n, p) = 1$. Following [2], let us have $n' \in \{0, 1, \dots, p - 1\}$ such that $nn' \equiv 1 \pmod{p}$ and $\beta = 1 + n'u$. If μ is the map

$$\begin{aligned} \mu : R[x]/(x^n - 1) &\longrightarrow R[x]/(x^n - (1 - u)) \\ r(x) &\longmapsto r(\beta x) \end{aligned}$$

then μ is a ring isomorphism. Hence I is an ideal of $R[x]/(x^n - 1)$ if and only if $\mu(I)$ is an ideal of $R[x]/(x^n - (1 - u))$. If $\bar{\mu}$ is the map

$$\begin{aligned} \bar{\mu} : R^n &\longrightarrow R^n \\ r &\longmapsto (r_0, \beta r_1, \beta^2 r_2, \dots, \beta^{n-1} r_{n-1}) \end{aligned}$$

then it also follows that:

Proposition 3.1. The set $\mathcal{C} \subseteq R^n$ is a linear cyclic code if and only if $\bar{\mu}(\mathcal{C})$ is a linear $(1 - u)$ -cyclic code.

We now define the permutation $\pi^{\otimes p^{k-1}}$, which is an extension of the Nechaev permutation introduced in [4], as follows:

For $\mathbf{c} = (c^{(p^0)} | \dots | c^{(p^{k-1})}) \in \mathbb{F}_{p^k}^{p^k n}$,

$$\pi^{\otimes p^{k-1}}(\mathbf{c}) := \left(\pi \left(c^{(p^0)} \right) \mid \dots \mid \pi \left(c^{(p^{k-1})} \right) \right),$$

where

$$\pi(\mathbf{a}) = (a_{\tau(0)}, \dots, a_{\tau(pn-1)})$$

for $\mathbf{a} = (a_0, \dots, a_{pn-1}) \in \mathbb{F}_{p^k}^{pn}$, where

$$\tau(\gamma n + j) = (\gamma + jn')_p n + j,$$

$0 \leq \gamma \leq p - 1, 0 \leq j \leq n - 1$, and $(\gamma + jn')_p$ is the least residue of $\gamma + jn'$ modulo p .

Proposition 3.2.

$$\Phi \circ \bar{\mu} = \pi^{\otimes p^{k-1}} \circ \Phi.$$

Proof. Let $\mathbf{r} = (r_0, \dots, r_{n-1}) \in R^n$. If $\Phi(\mathbf{r}) = (a_0, \dots, a_{p^k n-1})$, then $\pi^{\otimes p^{k-1}}(\Phi(\mathbf{r})) = (b_0, \dots, b_{p^k n-1})$, where

$$\begin{aligned} b_{\epsilon n+j} &= a_{[\gamma^{(\epsilon)}+(\gamma_{0,\epsilon}+jn')_p]n+j} \\ &= \left[\sum_{i=1}^{k-1} \gamma_{i,\epsilon} \alpha^i + (\gamma_{0,\epsilon} + jn')_p \right] x_j \oplus y_j. \end{aligned}$$

On the other hand, $\bar{\mu}(\mathbf{r}) = (r_0, \beta r_1, \dots, \beta^{n-1} r_{n-1})$, where

$$\begin{aligned} \beta^j r_j &= (1 + n'u)^j r_j \\ &= x_j + u(jn'x_j \oplus y_j). \end{aligned}$$

Hence $\Phi(\bar{\mu}(\mathbf{r})) = (c_0, c_1, \dots, c_{p^k n-1})$, where

$$\begin{aligned} c_{\epsilon n+j} &= \left(\sum_{i=0}^{k-1} \gamma_{i,\epsilon} \alpha^i \right) x_j \oplus (jn'x_j \oplus y_j) \\ &= \left[\sum_{i=1}^{k-1} \gamma_{i,\epsilon} \alpha^i + (\gamma_{0,\epsilon} + jn')_p \right] x_j \oplus y_j. \end{aligned}$$

The conclusion follows. \square

Corollary 3.3. *If \tilde{C} is the Gray image of a linear cyclic code of length n over R , then \tilde{C} is equivalent to a quasicyclic code of index p^{k-1} and length $p^k n$ over \mathbb{F}_{p^k} .*

Proof. From Proposition 3.1, a code \mathcal{C} of length n over R is linear cyclic if and only if $\bar{\mu}(\mathcal{C})$ is linear $(1 - u)$ -cyclic. From Theorem 2.2, this is so if and only if $\Phi(\bar{\mu}(\mathcal{C}))$ is a linear quasicyclic code of index p^{k-1} over \mathbb{F}_{p^k} , that is, if and only if $\pi^{\otimes p^{k-1}}(\Phi(\mathcal{C}))$ is linear quasicyclic of index p^{k-1} over \mathbb{F}_{p^k} . \square

The following example illustrates the above results. Computations were done using MAGMA.

Example. Let $\mathbb{F}_4 = \{0, 1, \omega, 1 + \omega = \omega^2\}$. Consider $x^3 - 1$. In $\mathbb{F}_4 + u\mathbb{F}_4$,

$$x^3 - 1 = (x + 1)(x + \omega)(x + \omega^2).$$

Applying the map $\mu : x \mapsto (1 + u)x$, we obtain

$$x^3 - (1 - u) = [x + (u + 1)][x + (\omega u + \omega)][x + (\omega^2 u + \omega^2)] = f_1(x)f_2(x)f_3(x).$$

1. Let \mathcal{C}_1 be the $(1 - u)$ -cyclic code generated by a , where

$$a(x) = f_1(x)f_2(x) = x^2 + (\omega^2u + \omega^2)x + 1.$$

The code $\Phi(\mathcal{C}_1)$ is a $[12, 2, 9]$ quaternary code, which is an optimal code.

2. Let $a(x)$ be as in the above,

$$b(x) = uf_1(x) = ux + u,$$

and \mathcal{C}_2 be the $(1 - u)$ -cyclic code over $\mathbb{F}_4 + u\mathbb{F}_4$ generated by a and b . The code $\Phi(\mathcal{C}_2)$ is a $[12, 3, 8]$ quaternary code, which is an optimal code.

References

- [1] M. Greferath, S.E. Schmidt, Gray isometries for finite chain rings and a nonlinear ternary $(36, 3^{12}, 15)$ code, IEEE Trans. Inform. Theory 45 (1999) 2522–2524.
- [2] S. Ling, J. Blackford, $\mathbb{Z}_{p^{k+1}}$ -linear codes, IEEE Trans. Inform. Theory 48 (2002) 2592–2605.
- [3] J.F. Qian, L.N. Zhang, S.X. Zhu, $(1 + u)$ -constacyclic and cyclic codes over $\mathbb{F}_2 + u\mathbb{F}_2$, Appl. Math. Lett. 19 (2006) 820–823.
- [4] J. Wolfmann, Negacyclic and cyclic codes over \mathbb{Z}_4 , IEEE Trans. Inform. Theory 45 (1999) 2527–2532.