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## ON CRITICAL SUBGRAPHS OF COLOUR-CRITICAL GRAPHS

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**Abstract.** Some results on the distribution of critical subgraphs in colour-critical graphs are obtained. Characterizations of  $k$ -critical graphs in terms of their  $(k - 1)$ -critical subgraphs are given. The special case  $k = 4$  is considered, and it is proved that if a 4-critical graph  $\Gamma$  has a vertex  $x$  of large valency (compared to the number of vertices of  $\Gamma$  not adjacent to  $x$ ), then  $\Gamma$  contains vertices of valency 3. Finally, a list of all 4-critical graphs with  $\leq 9$  vertices is exhibited.

## 1. Terminology and introduction.

We consider *finite graphs without loops and multiple edges*. The set of *vertices* and the set of *edges* of a graph  $\Gamma$  are denoted  $V(\Gamma)$  and  $E(\Gamma)$ , respectively. The *complete  $k$ -graph* is denoted  $\langle k \rangle$ . The terms *path* and *circuit* are used in the sense in which the corresponding terms *Weg* and *Kreis* are used in [6]. The *length* of a path or circuit is the number of edges contained in it. We allow a path to have length 0, but a circuit has length  $\geq 3$ . A path or a circuit is odd or even according to whether its length is odd or even.

If  $\Delta$  and  $\Gamma$  are graphs satisfying  $V(\Delta) \subseteq V(\Gamma)$  and  $E(\Delta) \subseteq E(\Gamma)$ , then  $\Delta$  is a *subgraph* of  $\Gamma$ , denoted  $\Delta \subseteq \Gamma$ . If  $T \subseteq V(\Gamma)$ , then  $\Gamma(T)$  denotes the subgraph of  $\Gamma$  *spanned by  $T$* , i.e.,  $V(\Gamma(T)) = T$  and  $E(\Gamma(T))$  consists of all edges of  $E(\Gamma)$  having both endvertices in  $T$ . If  $E(\Gamma(T)) = \emptyset$ , then  $T$  is an *independent set* of vertices of  $\Gamma$ , and  $T$  is *maximal independent* if it is not a proper subset of any other independent set of vertices of  $\Gamma$ . The *valency* of a vertex  $x$  of  $\Gamma$ , denoted  $\text{val}(x, \Gamma)$ , is the number of edges of  $\Gamma$  incident with  $x$ . If  $V' \subseteq V(\Gamma)$  and  $E' \subseteq E(\Gamma)$ , then  $\Gamma - V' - E'$  denotes the subgraph of  $\Gamma$  obtained by deleting from  $\Gamma$  all vertices of  $V'$

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and all edges of  $\Gamma$  incident with vertices of  $V'$ , and then deleting all edges of  $E'$  that remain, without deleting any more vertices. If  $\Gamma$  is connected and  $\Gamma - E'$  is disconnected, then  $E'$  is a *separating* set of edges of  $\Gamma$ .

A graph  $\Gamma$  is *k-colourable* if  $V(\Gamma)$  can be partitioned into at most  $k$  mutually disjoint (colour) classes in such a way that each class is an independent set of vertices. If  $k$  is the least integer for which  $\Gamma$  is *k-colourable*, then  $k$  is the *chromatic number* of  $\Gamma$  and  $\Gamma$  is *k-chromatic*. An element  $t \in V(\Gamma) \cup E(\Gamma)$  of a *k-chromatic* graph  $\Gamma$  is *critical* if  $\Gamma - t$  is  $(k - 1)$ -colourable. A connected graph  $\Gamma$  is *critical k-chromatic* (or simply *k-critical*) if it has chromatic number  $k$  and all edges – and consequently all vertices – of  $\Gamma$  are critical. Each vertex of a *k-critical* graph has valency  $\geq k - 1$ .

The 1- and 2-critical graphs are the  $\langle 1 \rangle$  and the  $\langle 2 \rangle$ , respectively. The 3-critical graphs are the odd circuits ([6, p. 151, Satz 12]). Hence each vertex of a *k-critical* graph with  $k \leq 3$  has valency  $k - 1$ . It seems hopeless to determine the structure of all 4-critical graphs. T. Gallai constructed an infinite class of regular 4-critical graphs of valency 4, thus proving that a 4-critical graph need not contain a vertex of valency 3 (see [4, p. 172, (2.3)] and [5]). M. Simonovits and the present author even proved that for any natural number  $h$ , there exist 4-critical graphs in which all vertices have valency  $\geq h$  (see [8, Chapter 6] and [9]). However, it is still unknown whether any planar 4-critical graph necessarily contains a vertex of valency 3 (see [5]).

In Section 4 of this paper, we shall prove that if  $\Gamma$  is a 4-critical graph having a vertex adjacent to all except  $\alpha$  vertices of  $\Gamma$ , then there is an upper bound  $m(\alpha)$  depending only on  $\alpha$  for the number of vertices of  $\Gamma$  of valency  $\geq 4$ , i.e., if  $|V(\Gamma)| > m(\alpha)$ , then  $\Gamma$  contains at least  $|V(\Gamma)| - m(\alpha)$  vertices of valency 3. The possible structures of  $\Gamma$  in the cases  $\alpha = 2$  and  $\alpha = 3$  are determined. This enables us to establish a catalogue of all 4-critical graphs with at most 8 vertices. Also the case of 9 vertices is mentioned. The proofs in Section 4 are based on a characterization in Section 3 of *k-critical* graphs in terms of their  $(k - 1)$ -critical subgraphs. In Section 2, a more elementary result is obtained and also a new, simple proof of the result, that any separating set of edges in a *k-critical* graph contains  $\geq k - 1$  edges, is presented.

## 2. Critical subgraphs

By deleting any vertex or any edge from a  $k$ -critical graph, the remaining graph has chromatic number  $k - 1$  and therefore contains at least one  $(k - 1)$ -critical subgraph. Theorem 2.1 gives extensions of this statement.

**Theorem 2.1.** *Let  $\Gamma$  be a  $k$ -critical graph ( $k \geq 3$ ), and let  $\alpha$  be an integer satisfying  $1 \leq \alpha \leq k - 2$ .*

(a) *By deleting at most  $\alpha$  vertices from  $\Gamma$ , there exists for each vertex  $x$  of the remaining graph  $\Delta$ , a  $(k - \alpha)$ -critical subgraph of  $\Delta$  containing  $x$ .*

(b) *By deleting at most  $\alpha$  edges from  $\Gamma$ , there exists for each edge  $e$  of the remaining graph  $\Delta$ , a  $(k - \alpha)$ -critical subgraph of  $\Delta$  containing  $e$ , but not containing both endvertices of any of the deleted edges.*

Theorem 2.1(a) follows easily by induction from the case  $\alpha = 1$ , and the case  $\alpha = 1$  is an immediate corollary of a characterization of  $k$ -critical graphs given in Section 3. Theorem 2.1(a) was obtained earlier by Dirac [3, p. 45, (4)], and the following proof of Theorem 2.1(b) is similar to the proof in [3] of Theorem 2.1(a).

**Proof of Theorem 2.1(b).** Let  $e_1, \dots, e_\nu$  be a set of  $\nu$  edges of  $\Gamma$ , where  $\nu \leq \alpha$ , and let  $e$  be an edge of  $\Gamma - e_1 - \dots - e_\nu (= \Delta)$ .  $\Gamma - e$  has a  $(k - 1)$ -colouring  $K$  with colours  $1, 2, \dots, k - 1$  such that the two endvertices of  $e$  both have the colour  $k - 1$  and such that for  $i = 1, \dots, \nu$ , at least one endvertex of  $e_i$  has a colour among  $1, \dots, \nu$ . Delete from  $\Gamma$  all vertices having colours  $1, \dots, \alpha$  in  $K$  and call the remaining graph  $\Delta'$ .  $e_i \notin E(\Delta')$  for  $i = 1, \dots, \nu$ , but since  $\alpha \leq k - 2$ ,  $e \in E(\Delta')$ .  $\Delta'$  has chromatic number  $\geq k - \alpha$ , because if  $\Delta'$  were  $(k - \alpha - 1)$ -colourable, then  $\Gamma$  would be  $(k - 1)$ -colourable. However,  $K$ 's restriction to  $\Delta' - e$  shows that  $\Delta' - e$  is  $(k - \alpha - 1)$ -colourable, hence  $e$  is contained in any  $(k - \alpha)$ -critical subgraph of  $\Delta'$ . This proves Theorem 2.1(b).

For  $\alpha = k - 2$ , Theorem 2.1(a) is equivalent to the statement that each vertex of a  $k$ -critical graph has valency  $\geq k - 1$ , and (b) is trivial. The case  $\alpha = k - 3$  of (b) gives a simple proof of the following well-known result [2, p. 45, Theorem 1]:

- (1) *A separating set of edges of a  $k$ -critical graph ( $k \geq 3$ ) contains at least  $k - 1$  edges.*

**Proof.** For  $k = 3$ , the statement is true. Suppose that  $k \geq 4$  and let  $E$  be a separating set of edges of the  $k$ -critical graph  $\Gamma$ . If we delete from  $\Gamma$  all edges of  $E$  except one, then the remaining edge of  $E$  is not contained in any circuit of the remaining graph. Then by Theorem 2.1(b) with  $\alpha = k - 3$ , at least  $k - 2$  edges have been deleted, i.e.,  $|E| \geq k - 1$ , and (1) has been proved. Another version of this proof based on the case  $\alpha = 1$  of Theorem 2.1(b) and by induction over  $k$  can be given.

The negation of the following statement (2) would – if true – have generalized both parts of Theorem 2.1 in the case  $\alpha = 1$ .

- (2) *An edge of a graph  $\Gamma - x$ , where  $\Gamma$  is  $k$ -critical ( $k \geq 4$ ) and  $x \in V(\Gamma)$ , is not necessarily contained in a  $(k - 1)$ -critical subgraph of  $\Gamma - x$ .*

**Proof.** Let  $\Gamma$  be a graph obtained from two disjoint,  $k$ -critical graphs  $\Delta_1$  and  $\Delta_2$  ( $k \geq 4$ ) by Hajós' construction, i.e., delete an edge  $(x_1, y_1)$  from  $\Delta_1$  and an edge  $(x_2, y_2)$  from  $\Delta_2$ , identify  $x_1$  and  $x_2$  to a new vertex  $x$  and join  $y_1$  and  $y_2$  by a new edge  $e$ .  $\Gamma$  is  $k$ -critical and the edge  $e$  is not contained in any  $(k - 1)$ -critical subgraph of  $\Gamma - x$ . This proves (2).

The above considerations give rise to an unsolved problem: For  $\Gamma$   $k$ -critical ( $k \geq 4$ ) and  $x \in V(\Gamma)$ , characterize the set  $E_x$  of those edges of  $\Gamma - x$  that are not contained in any  $(k - 1)$ -critical subgraph of  $\Gamma - x$ . By (1), any edge contained in a separating set of  $\leq k - 3$  edges of  $\Gamma - x$  belongs to  $E_x$ . A result of Dirac [3, p. 48, Corollary to Theorem 3] implies that for  $k = 4$ , also the converse of this statement is true, i.e., for  $k = 4$ ,  $E_x$  consists of precisely the separating edges (also called the bridges) of  $\Gamma - x$ . This implies that the above examples showing (2) are the only such examples for  $k = 4$ . However, the situation changes for  $k \geq 5$ . Fig. 1 shows a 5-critical graph  $\Gamma$  (we leave it to the reader to check this) in which the edge  $e_1$  belongs to  $E_x$ , but where it is not contained in any separating set of  $\leq 2$  edges of  $\Gamma - x$ . To see that  $e_1 \in E_x$  assume on the contrary that there exists a 4-critical subgraph  $\Delta$  of  $\Gamma - x$

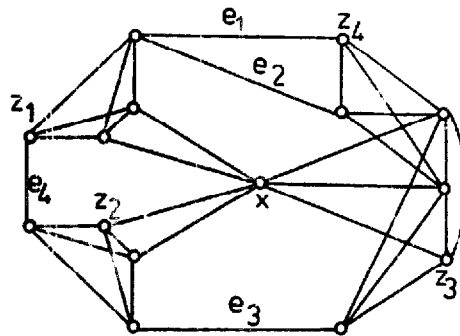


Fig. 1.

containing  $e_1$ . Then by (1), also  $e_2$  and  $e_3$  are contained in  $\Delta$ . However,  $\{e_3, e_4\}$  is a separating set of two edges of  $\Gamma - x$ , and  $e_3$  is therefore not contained in  $\Delta$ . This is a contradiction. The example of Fig. 1 may be generalized to larger values of  $k$  by replacing each of the vertices  $z_1, z_2, z_3$  and  $z_4$  by a  $(k - 4)$ , showing the existence of a  $k$ -critical graph ( $k \geq 5$ ) with an edge  $e \in E_x$ , where  $e$  is not contained in any separating set of  $\leq k - 3$  edges of  $\Gamma - x$ .

### 3. Characterizations of critical graphs

Let for a graph  $\Gamma$  and any vertex  $x$  of  $\Gamma$ ,  $A_x$  denote the set of vertices of  $\Gamma$  not adjacent to  $x$  and let  $B_x$  denote the set of vertices of  $\Gamma$  adjacent to  $x$ .  $V(\Gamma)$  is thus the disjoint union of  $\{x\}$ ,  $A_x$  and  $B_x$ . Let  $\Delta_x$  denote the subgraph  $\Gamma - x - E(\Gamma(A_x))$ .

**Theorem 3.1.** *Let  $\Gamma$  be a non-empty graph and let  $k \geq 3$ . The following four statements are equivalent.*

- (a)  $\Gamma$  is  $k$ -critical.
- (b) For all  $x \in V(\Gamma)$ , the following statement holds:

(\*) *Let the maximal independent sets of vertices of  $\Gamma(A_x)$  be  $I_1^x, \dots, I_\mu^x$ . Then for  $j = 1, \dots, \mu$ ,  $\Gamma - x - I_j^x$  contains a  $(k - 1)$ -critical subgraph. Let for  $j = 1, \dots, \mu$ ,  $\theta_j$  denote any such  $(k - 1)$ -critical subgraph of  $\Gamma - x - I_j^x$ . Then  $\Delta_x \subseteq \bigcup_{j=1}^{\mu} \theta_j$ .*

(c) *Let  $I_1, \dots, I_\nu$  be all maximal independent sets of vertices of  $\Gamma$ . Then (\*\*) holds.*

For  $j = 1, \dots, \nu$ ,  $\Gamma - I_j$  contains a  $(k - 1)$ -critical subgraph. Let  
 (\*\*)  $\theta_j$  be any such  $(k - 1)$ -critical subgraph of  $\Gamma - I_j$ . Then  $\Gamma = \bigcup_{j=1}^{\nu} \theta_j$ .

(d) There exist independent sets  $I_1, \dots, I_{\nu}$  of vertices of  $\Gamma$  such that (\*\*) holds and such that for at least one vertex  $x$  of  $\Gamma$  all maximal independent sets of vertices of  $\Gamma$  containing  $x$  are among  $I_1, \dots, I_{\nu}$ .

**Remark 3.2.** In Section 4, only the statement (a)  $\Rightarrow$  (b) will be used.

**Remark 3.3.** The  $\theta_j$ 's in (\*) and (\*\*) are not necessarily different and not necessarily uniquely determined.

**Remark 3.4.** The assertion of Theorem 2.1(a) in the case  $\alpha = 1$  is an immediate consequence of (b).

**Proof of Theorem 3.1.** (a)  $\Rightarrow$  (b). Let  $\Gamma$  be  $k$ -critical and let  $x \in V(\Gamma)$ . We shall prove that (\*) holds. For each  $j$ ,  $1 \leq j \leq \mu$ ,  $\{x\} \cup I_j^x$  is an independent set of vertices of  $\Gamma$ , hence  $\Gamma - x - I_j^x$  is  $(k - 1)$ -chromatic and therefore contains a  $(k - 1)$ -critical subgraph  $\theta_j$ . Let  $t \in V(\Delta_x) \cup E(\Delta_x)$ . In order to finish the proof of (\*) we shall prove that  $t \in \theta_1 \cup \dots \cup \theta_{\mu}$ .  $\Gamma$  is  $k$ -critical, hence  $\Gamma - t$  has a  $(k - 1)$ -colouring  $K$ . Let  $I$  denote the set of vertices of  $A_x$  having the same colour as  $x$  in  $K$ .  $\{x\} \cup I$  is one of the  $k - 1$  colour-classes of  $K$ , hence  $\Gamma - x - I - t$  is  $(k - 2)$ -colourable and since  $t \in \Delta_x$ ,  $\{x\} \cup I$  is an independent set of vertices of  $\Gamma$ . It follows that  $\Gamma - x - I$  is  $(k - 1)$ -chromatic and that  $t$  is contained in any  $(k - 1)$ -critical subgraph of  $\Gamma - x - I$ . There exists a  $j$  such that  $I \subseteq I_j^x$ , hence either  $t \in I_j^x$  or  $t \in \theta_j$ . If  $t \in I_j^x$ , then a  $(k - 1)$ -colouring of  $\Gamma$  may be obtained from  $K$  by giving to  $t$  the colour of  $x$  in  $K$ . But  $\Gamma$  is  $k$ -chromatic, hence  $t \in \theta_j$ . This proves (a)  $\Rightarrow$  (b).

(b)  $\Rightarrow$  (c). Let  $\Gamma$  satisfy (b). For each integer  $j$ ,  $1 \leq j \leq \nu$ , and  $x \in I_j$ , there exists an integer  $h$  such that  $I_j = \{x\} \cup I_h^x$ , hence by (\*),  $\Gamma - I_j = \Gamma - x - I_h^x \supseteq \theta_j$ , where  $\theta_j$  is  $(k - 1)$ -critical. This proves the first part of (\*\*) and implies  $|V(\Gamma)| \geq k$ . By Remark 3.4, we have that for any vertex  $x$  of  $\Gamma$ , each vertex of  $\Gamma - x$  is contained in a  $(k - 1)$ -critical subgraph of  $\Gamma - x$ , hence no connected component of  $\Gamma$  is a  $\langle 1 \rangle$  or a  $\langle 2 \rangle$ , which implies that  $\bigcup_{x \in V(\Gamma)} \Delta_x = \Gamma$ . For each  $x \in V(\Gamma)$  and each  $I_h^x$ , there exists a  $j$  such that  $\{x\} \cup I_h^x = I_j$ , hence by (\*),

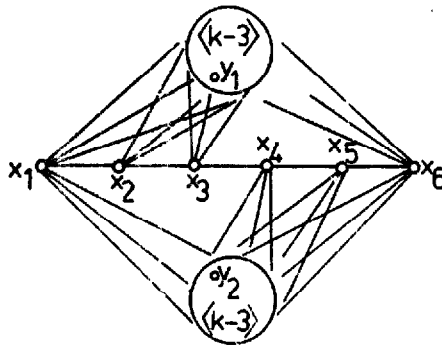


Fig. 2.

$$\Gamma = \bigcup_{x \in V(\Gamma)} \Delta_x \subseteq \bigcup_{j=1}^p \theta_j \subseteq \Gamma,$$

which proves (b)  $\Rightarrow$  (c).

(c)  $\Rightarrow$  (d). Trivial.

(d)  $\Rightarrow$  (a). Let  $\Gamma$  satisfy (d). Suppose that  $\Gamma$  has a  $(k - 1)$ -colouring  $K$  and let  $C$  be the colour-class of  $K$  containing  $x$ . Then  $\Gamma - C$  is  $(k - 2)$ -colourable and since  $x \in C$ , there exists a  $j$  such that  $C \subseteq I_j$ , hence also  $\Gamma - I_j$  is  $(k - 2)$ -colourable. This contradicts (\*\*). Hence  $\Gamma$  has chromatic number  $\geq k$ . In order to prove that  $\Gamma$  is  $k$ -critical it is therefore sufficient to prove that for any element  $t \in V(\Gamma) \cup E(\Gamma)$ ,  $\Gamma - t$  is  $(k - 1)$ -colourable. Let  $t \in V(\Gamma) \cup E(\Gamma)$ . By (\*\*), there exists a  $j$  such that  $t$  is necessarily contained in any  $(k - 1)$ -critical subgraph of  $\Gamma - I_j$ , because otherwise each  $\theta_j$  could be chosen such that  $\theta_j \subseteq \Gamma - t$ , contradicting  $\Gamma = \bigcup_{j=1}^p \theta_j$ . It follows that there exists a  $j$  such that  $\Gamma - I_j - t$  is  $(k - 2)$ -colourable, but then  $\Gamma - t$  is  $(k - 1)$ -colourable, because  $I_j$  is an independent set of vertices. This proves (d)  $\Rightarrow$  (a).

Theorem 3.1 has then been proved.

Let us remark that the proof of (a)  $\Rightarrow$  (b) may be applied also in the case where  $t \in E(\Gamma(A_x))$  if  $\Gamma - t$  has a  $(k - 1)$ -colouring in which the two endvertices of  $t$  and  $x$  have not all three the same colour. This implies that if  $t \in E_x$  (defined in Section 2), then in any  $(k - 1)$ -colouring of  $\Gamma - t$  the two endvertices of  $t$  and  $x$  have necessarily all three the same colour.

\* The condition of (d), that there exists a vertex  $x$  such that all maximal independent sets containing  $x$  are among  $I_1, \dots, I_p$ , cannot be omitted, for  $k \geq 4$  not even if  $I_1, \dots, I_p$ , are required all to be maximal independent. This is shown, for  $k \geq 4$ , by the graph  $\Gamma$  of Fig. 2. For this

graph  $\Gamma$ , the maximal independent sets  $I_1 = \{x_2, v_2\}$ ,  $I_2 = \{x_3, y_2\}$ ,  $I_3 = \{x_4, y_1\}$  and  $I_4 = \{x_5, y_1\}$  satisfy (\*\*), but  $\Gamma$  is  $(k - 1)$ -colourable. It can, however, be proved that a non-empty graph  $\Gamma$  is 3-critical, i.e., an odd circuit, if and only if there exist maximal independent sets  $I_1, \dots, I_p$  of vertices of  $\Gamma$  such that (\*\*) with  $k = 3$  is satisfied.

The following well-known results on  $k$ -critical graphs are easy consequences of Theorem 3.1(b).

- (3) *If  $A_x \neq \emptyset$ , then each vertex of  $\Gamma(A_x)$  has valency  $\geq 1$  in  $\Gamma(A_x)$ , especially  $|A_x| \geq 2$ .*

**Proof.** If  $y \in A_x$ , then  $y \in V(\Delta_x)$ , hence by (\*),  $y$  is not contained in all maximal independent sets of vertices of  $\Gamma(A_x)$ . This proves (3).

- (4) *If  $A_x \neq \emptyset$ , then  $\Gamma(B_x)$  does not contain a  $(k - 1)$ -critical subgraph.*

**Proof.** If  $\theta$  is a  $(k - 1)$ -critical subgraph of  $\Gamma(B_x)$ , then we may choose  $\theta_j = \theta$  for all  $j$ , contradicting  $\Delta_x \subseteq \bigcup_{j=1}^{\mu} \theta_j = \theta$ .

#### 4. 4-critical graphs

Let now  $\Gamma$  be a 4-critical graph. By Theorem 3.1,  $\Gamma$  satisfies Theorem 3.1(b) in the case  $k = 4$  in which the graphs  $\theta_j$  are odd circuits. M. Simonovits made me aware of this extension (for  $k = 4$ ) of (4):

- (5) *If  $A_x \neq \emptyset$ , then  $\Gamma(B_x)$  does not contain any circuits.*

**Proof.** Suppose that  $\theta$  is a circuit of  $\Gamma(B_x)$ . By (4),  $\theta$  is even and therefore the two endvertices of an edge  $e \in E(\theta)$  have different colours in a 3-colouring of  $\Gamma - e$ . But then  $\Gamma$  is 3-colourable, which is a contradiction. This proves (5).

By (5),  $\Gamma(B_x) \cap \theta_j$  is either empty or it consists of a set of mutually disjoint paths. If  $\Gamma(B_x) \cap \theta_j \neq \emptyset$ , then let these paths be  $P_1^j, P_2^j, \dots, P_{f(j)}^j$ , where  $f(j) \geq 1$ , otherwise let  $f(j) = 0$ .

- (6) *Let  $1 \leq j < m \leq \mu$  and let  $1 \leq i \leq f(j)$  and  $1 \leq h \leq f(m)$ . Then  $P_i^j \cap P_h^m$  is either empty or a path.*



**Proof.** if not, then  $\Gamma(B_x)$  would contain a circuit, contradicting (5). This proves (6).

On the basis of Theorem 3.1(b) and (5) we get:

**Theorem 4.1.** *Let  $\alpha$  be an integer  $\geq 2$ . There exist integers  $n(\alpha)$  and  $m(\alpha)$  depending only on  $\alpha$  such that if  $\Gamma$  is a 4-critical graph containing a vertex  $x$  of valency  $|V(\Gamma)| - \alpha - 1$ , then:*

(a)  $|E(\Gamma)| \leq 2 \cdot |V(\Gamma)| + n(\alpha)$ ;

(b)  $\Gamma$  contains at most  $m(\alpha)$  vertices of valency  $\geq 4$ , i.e., if  $|V(\Gamma)| > m(\alpha)$ , then  $\Gamma$  contains at least  $|V(\Gamma)| - m(\alpha)$  vertices of valency 3.

**Remark 4.2.** The statement for 5-critical graphs corresponding to (b) of Theorem 4.1 (i.e., the statement obtained from (b) by replacing 3 and 4 by 4, 5, respectively) is not true for any value of  $\alpha$ . Counterexamples of infinitely many 5-critical graphs  $\Gamma$  containing vertices  $x$  of valencies  $|V(\Gamma)| - \alpha - 1$  in which all vertices have valencies  $\geq h$ , where  $h$  is a given integer  $\geq 5$ , may be constructed from the 4-critical graphs having large minimal valency (they may be constructed such that the resulting graphs have separating sets of two vertices of which  $x$  is one).

**Proof of Theorem 4.1.** Let  $\Gamma$  be 4-critical and let  $x \in V(\Gamma)$  have valency  $|V(\Gamma)| - \alpha - 1$  in  $\Gamma$ , i.e.,  $|A_x| = \alpha$ . The number of edges of  $\Gamma(B_x)$  is at most  $|B_x| - 1$  by (5). The number of edges of  $\Gamma(A_x)$  is at most  $\frac{1}{2}\alpha(\alpha - 1)$ . The number of edges having one endvertex in  $A_x$  and the other in  $B_x$  is  $\leq 2 \cdot \sum_{j=1}^{\mu} f(j)$  by Theorem 3.1(b).  $\mu$  is bounded by the number of different non-empty subsets of  $A_x$ , i.e.,  $\mu \leq 2^\alpha - 1$ , in fact,  $\mu \leq 2^\alpha - 2$  since the vertices of  $A_x$  are not independent by (3).  $f(j) \leq \alpha - 1$  since the circuit  $\theta_j$  is "passing through"  $\Gamma(A_x)$  at most  $\alpha - 1$  times. It follows that

$$\begin{aligned} |E(\Gamma)| &\leq |V(\Gamma)| - \alpha - 1 + |B_x| - 1 + \frac{1}{2}\alpha(\alpha - 1) + 2(2^\alpha - 2)(\alpha - 1) \\ &= 2 \cdot |V(\Gamma)| - 2\alpha - 3 + \frac{1}{2}\alpha(\alpha - 1) + 2(2^\alpha - 2)(\alpha - 1) \\ &= 2 \cdot |V(\Gamma)| + n(x). \end{aligned}$$

This proves (a).

By the formula  $\sum_{y \in V(\Gamma)} \text{val}(y, \Gamma) = 2|E(\Gamma)|$ , we get

$$|V(\Gamma)| - \alpha - 1 + 4(|V(\Gamma)| - 1 - r) + 3r \leq 2|E(\Gamma)|,$$

where  $r$  denotes the number of vertices of  $A_x \cup B_x$  having valency 3 in  $\Gamma$ . By (a), this implies

$$r \geq |V(\Gamma)| - 2n(\alpha) - \alpha - 5,$$

hence  $\Gamma$  contains at most  $2n(\alpha) + \alpha + 5$  vertices of valency  $\geq 4$ . This proves (b), hence Theorem 4.1 has been proved.

Unfortunately, I know next to nothing on exact values and orders of magnitudes of the best possible  $n(\alpha)$  and  $m(\alpha)$  that may be used in Theorem 4.1. The estimation in the proof is not best possible. Thus the number of edges of  $\Gamma(A_x)$  is  $\leq \frac{1}{3}\alpha^2$  by Turán's theorem [11], and  $\mu \leq 3^{\alpha/3}$  by a theorem of Moon and Moser [7]. However, even this gives an estimation which is probably far from the correct order of magnitude. The estimation of the proof gives for  $\alpha = 2$ ,  $n(2) = -2$  and  $m(2) = 3$ , which is best possible as we shall see in the following, where the special cases  $\alpha = 2$  and  $\alpha = 3$  are treated. Dirac and Gallai conjectured [5, p. 44, Conjecture] that if  $\Gamma$  is assumed also to be planar, then  $n(\alpha) = -2$  may be used for all values of  $\alpha$ .

#### 4.1. $\alpha = 2$

Let  $A_x = \{y, z\}$ . By (3),  $\Gamma(A_x) = \langle 2 \rangle$  and  $\mu = 2$ . We may assume that  $I_1^x = \{y\}$  and  $I_2^x = \{z\}$ . By (5), the odd circuit  $\theta_1$  contains  $z$ , but not  $y$ , and the odd circuit  $\theta_2$  contains  $y$ , but not  $z$ . By Theorem 3.1(b),

$$A_x = \Gamma - x - (y, z) = \theta_1 \cup \theta_2,$$

and by (c)  $\theta_1 \cap \theta_2$  is either empty or a path. If  $\theta_1 \cap \theta_2$  is a path with  $\geq 2$  vertices, then none of the endvertices of that path are adjacent to both vertices of  $A_x$  since otherwise the edge from  $x$  to such an endvertex would not be critical. (The case  $\alpha = 2$  was also considered in [10, §6].)

#### 4.2. $\alpha = 3$

Let  $A_x = \{p, y, z\}$ . By (3),  $\Gamma(A_x)$  is either

- (i) a path, or
- (ii) complete.

Let us consider the two cases in turn.

(i)  $E(\Gamma(A_x)) = \{(p, y), (y, z)\}$  .

$\mu = 2$  and we may assume that  $I_1^x = \{p, z\}$  and  $I_2^x = \{y\}$ . By (5), the odd circuit  $\theta_1$  contains  $y$ , but not  $p$  nor  $z$ . By Theorem 3.1(b), the odd circuit  $\theta_2$  contains  $p$  and  $z$ , but not  $y$ , and it consists therefore of an even path  $P_1$  of length  $\geq 2$  and an odd path  $P_2$  of length  $\geq 3$  both joining  $p$  and  $z$  and with no interior vertices in common. By Theorem 3.1(b),

$$\Delta_x = \Gamma - x - (p, y) - (y, z) = \theta_1 \cup P_1 \cup P_2 .$$

If  $P_i \cap \theta_1 \neq \emptyset$ , then  $P_i \cap \theta_1$  is a path by (6) and  $P_j \cap \theta_1 = \emptyset$  ( $j = 1$  if  $i = 2$  and conversely) since otherwise the edges  $(p, y)$  and  $(y, z)$  would not both be critical. If  $P_1$  has length 2 and if the interior vertex  $v$  of  $P_1$  is on  $\theta_1$ , then  $v$  is not adjacent to  $y$  since otherwise  $(x, v)$  would not be critical.

(ii)  $\Gamma(A_x) = \langle 3 \rangle$  .

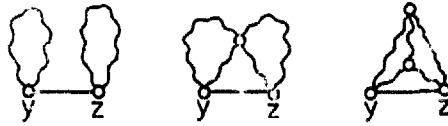
$\mu = 3$  and we may assume that  $I_1 = \{p\}$ ,  $I_2 = \{y\}$  and  $I_3 = \{z\}$ . By (5), each of  $\theta_1, \theta_2$  and  $\theta_3$  contains at least one vertex of  $A_x$ . We shall consider two cases.

(ii.1) Suppose that one of  $\theta_1, \theta_2$  and  $\theta_3$  may be chosen such that it contains only one vertex of  $A_x$ , say  $\theta_1$  contains  $y$ , but not  $p$  nor  $z$ . Then we may choose  $\theta_3 = \theta_1$ . By Theorem 3.1(b),  $\theta_2$  contains necessarily both  $p$  and  $z$ , but not  $y$ , hence there is in  $\Delta_x$  an even path  $P$  of length  $\geq 2$  joining  $p$  and  $z$  and not containing  $y$ .  $\theta_2$  may therefore be chosen as  $P$  together with the edge  $(p, z)$ . By Theorem 3.1(b),

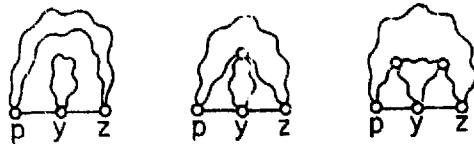
$$\Delta_x = \theta_1 \cup P ,$$

where  $\theta_1 \cap P = \emptyset$  since otherwise one of the two edges of  $\theta_1$  incident with  $y$  would not be critical.

$\alpha = 2 :$



$\alpha = 3 . (i) :$



$\alpha = 3 . (ii) :$

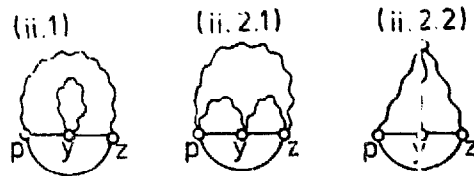


Fig. 3.

(ii.2) The alternative to consider is the case where each of  $\theta_1, \theta_2$  and  $\theta_3$  necessarily contains precisely two vertices of  $A_x$ . In this case,  $\theta_1, \theta_2$  and  $\theta_3$  may be chosen as even paths  $P_1, P_2$  and  $P_3$  of length  $\geq 2$  joining the vertices  $y$  and  $z, p$  and  $z, p$  and  $y$  respectively, together with the edges  $(y, z), (p, z)$  and  $(p, y)$ , respectively. By Theorem 3.1(b),

$$\Delta_x = P_1 \cup P_2 \cup P_3 .$$

One possibility is that  $P_1, P_2$  and  $P_3$  are mutually disjoint outside  $A_x$ . The alternative is that, say,  $V(P_1) \cap V(P_2) \cap B_x \neq \emptyset$ . Let  $q$  be the first vertex of  $P_1$  (going from  $y$  towards  $z$ ) belonging also to  $P_2$ . Let  $R_1, R_2$  and  $R_3$  be three paths joining  $q$  with  $p, y$  and  $z$  respectively, where  $R_1$  and  $R_3$  are parts of  $P_2$  and  $R_2$  is a part of  $P_1$ .  $R_1, R_2$  and  $R_3$  are mutually disjoint except for the vertex  $q$ , and since  $P_2$  is even,  $R_1$  and  $R_3$  have the same parity. By (6),  $P_1 \cap P_2$  is a path and since (ii.1) is not the case, also  $R_2$  and  $R_3$  have the same parity, i.e.,  $R_1, R_2$  and  $R_3$  have all the same parity. Then  $\theta_1, \theta_2$  and  $\theta_3$  may be chosen as  $R_2 \cup R_3 \cup (y, z), R_1 \cup R_3 \cup (p, z)$  and  $R_1 \cup R_2 \cup (p, y)$ , respectively. Hence by Theorem 3.1(b),

$$\Delta_x = R_1 \cup R_2 \cup R_3 .$$

Since  $\Gamma - x$  is 3-colourable, the lengths of  $R_1$ ,  $R_2$  and  $R_3$  are not all 1.

This completes the treatment of the cases  $\alpha = 2$  and  $\alpha = 3$ . Fig. 3 shows the various possibilities for  $\Gamma - x$ . It can be proved that the conditions on the structure of  $\Gamma$  that we have established are not only necessary, but also sufficient conditions for  $\Gamma$  to be 4-critical. The analysis gives as a corollary:

**Theorem 4.3.** *Let  $\Gamma$  be a 4-critical graph with a vertex  $x$  of valency  $|V(\Gamma)| - \alpha - 1$ .*

(a) *If  $\alpha = 2$ , then all vertices of  $\Gamma - x$  have valency 3 in  $\Gamma$  except possibly either one vertex of valency 5 or two vertices each of valency 4.*

*All three cases occur.*

(b) *If  $\alpha = 3$ , then all vertices of  $\Gamma - x$  have valency 3 in  $\Gamma$  except either one vertex of valency 4 or two vertices of valencies 4 and 5 respectively, or three vertices each of valency 4. All three cases occur.*

It follows from Theorem 4.3 that the best possible values of  $n(\alpha)$  and  $m(\alpha)$  in the cases  $\alpha = 2$  and  $\alpha = 3$  are  $n(2) = n(3) = -2$ ,  $m(2) = 3$  and  $m(3) = 4$ .

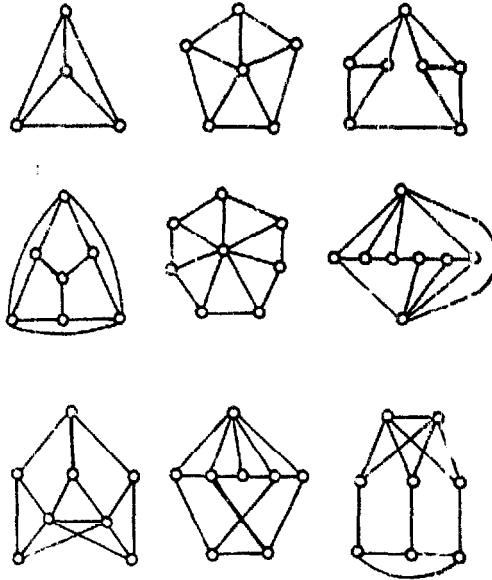
The case  $\alpha = 4$  has also been considered, and by the above method it is in this case possible to prove:

(7) *There exists precisely one 4-critical graph with 9 vertices in which each vertex has valency 3 or 4, namely the last graph of Fig. 4.*

I shall leave out my rather cumbersome proof of (7).

If  $\Gamma$  is a 4-critical graph with at most 8 vertices, then by the theorem of Brooks [1] it contains a vertex  $x$  of valency  $|V(\Gamma)| - \alpha - 1$  for either  $\alpha = 0$ ,  $\alpha = 2$  or  $\alpha = 3$ . By (7), this holds also if  $|V(\Gamma)| = 9$  and  $\Gamma =$  the graph of (7). If  $\alpha = 0$ , then it is well known that  $\Delta_x$  is an odd circuit, hence the above analysis provides us with a complete list of all 4-critical graphs with  $\leq 9$  vertices. There are 30 such graphs and they are exhibited in Fig. 4. One of the graphs has 4 vertices, one has 6 vertices, 2 have 7 vertices, 5 have 8 vertices and 21 have 9 vertices.

$|V(\Gamma)| \cong 8 :$



$|V(\Gamma)| = 9$  and  $\max val(x, \Gamma) = 6 :$

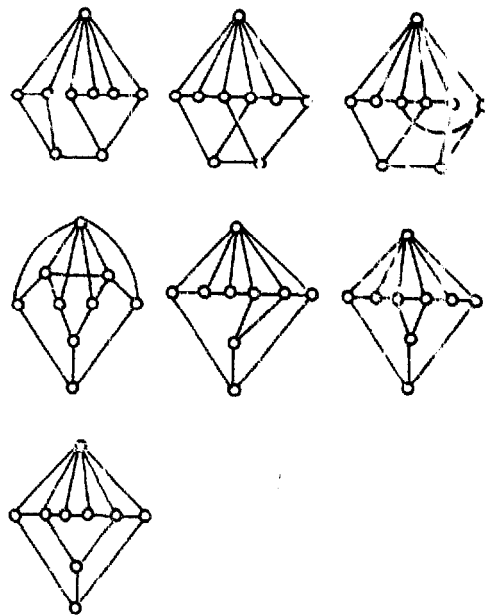
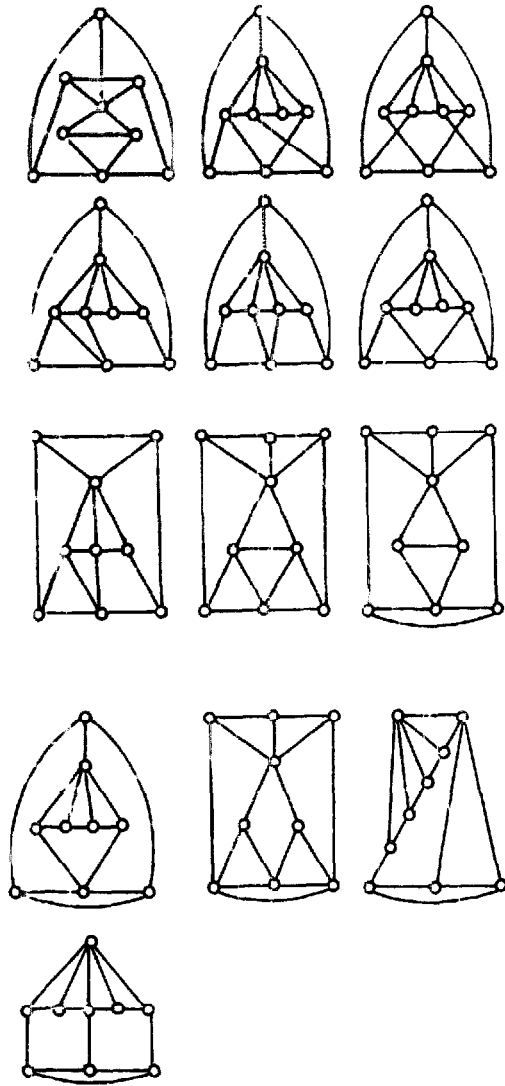


Fig. 4.

$|V(\Gamma)| = 9$  and  $\maxval(x, \Gamma) = 5$  :



$|V(\Gamma)| = 9$  and  $\maxval(x, \Gamma) = 4$  :

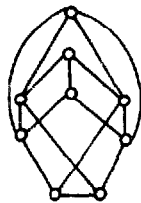


Fig. 4 (continued).

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