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# Bounds for the largest Mahalanobis distance

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## Abstract

Upper and lower bounds for the magnitude of the largest Mahalanobis distance, calculated from  $n$  multivariate observations of length  $p$ , are derived. These bounds are multivariate extensions of corresponding bounds that arise for the most deviant  $Z$ -score calculated from a univariate sample of size  $n$ . The approach taken is to pose optimization problems in a mathematical context and to employ variational methods to obtain solutions. The attainability of the bounds obtained is demonstrated. Bounds for related quantities (elements of the “hat matrix”) are also derived.

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## 1. Contexts: Statistical and general

Let  $y_1, \dots, y_n$  be points in  $\mathbb{R}$ . Define  $Z$ -scores

$$z_i = \frac{(y_i - \bar{y})}{s}$$

for  $i = 1, \dots, n$ , where the sample mean is calculated as

$$\bar{y} = \sum_{i=1}^n \frac{y_i}{n}$$

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and the standard deviation is calculated as

$$s = \sqrt{\sum_{i=1}^n \frac{(y_i - \bar{y})^2}{n - 1}}.$$

Shiffler [9] showed that  $(n - 1)/\sqrt{n}$  is an upper bound for  $\max_i |z_i|$ , the magnitude of the most deviant  $Z$ -score in a univariate sample of size  $n$ . Olkin [6] noted that although this result is often attributed to Samuelson [8] historically the upper bound should identified with Thompson [7]. Various generalizations of this result have been considered. Hayes [4] showed that the magnitude of the most deviant  $Z$ -score has the lower bound  $\sqrt{(n - 1)/n}$ , and that attaining this lower bound necessarily requires that Shiffler and Harsha’s [10] upper bound for the standard deviation be also attained. Gray and Woodall [3] discussed the extension of the upper bound to residuals obtainable from the linear model with independent error terms of common variance.

In statistical analysis, the row vectors  $x'_i$  in  $\mathbb{R}^p$ , for  $i = 1, \dots, n$ , of an  $n \times p$  matrix  $X$  are often interpreted as independent realizations from a multivariate normal density  $MVN_p(\mu, \Sigma)$ . The maximum likelihood estimates of the  $p$  dimensional vector mean  $\mu$  and  $p \times p$  variance matrix  $\Sigma$  are  $\hat{\mu} = \bar{x} = X'u/n$ , and  $\hat{\Sigma} = X'CX/n$ , respectively, where  $u$  is the vector of ones in  $\mathbb{R}^n$  and the  $n \times n$  matrix  $C = I_n - \frac{1}{n}uu'$ , where  $I_n$  is the  $n \times n$  identity matrix. Observe that the matrix  $C$  is idempotent and that the matrix  $CX$  has rows  $x'_i - \bar{x}'$ , for  $i = 1, \dots, n$ . Typically, the unbiased estimate  $S = X'CX/(n - 1)$  of  $\Sigma$  is used in place of  $\hat{\Sigma}$ .

If  $S$  is non-singular, the Mahalanobis distance of the  $p$  dimensional vector  $x_i$  is defined as  $D_i^2 = (x_i - \bar{x})'S^{-1}(x_i - \bar{x})$ . This is a generalization to the multivariate setting of the quantities  $z_i^2$  of the univariate case. Two standard texts on outlier detection Barnett and Lewis [2], and Hawkins [5] provide a thorough discussion on the role of the Mahalanobis distance in the outlier literature.

Using a useful matrix identity, Olkin [6] verified that  $(n - 1)^2/n$  is an upper bound for the maximum Mahalanobis distance in the special case of bivariate data. This result follows from a more general upper bound for bivariate data derived by Olkin [6] who also noted that the same upper bound follows as a consequence of Corollary 1 to Theorem 1 of Arnold and Groeneveld [1] for higher dimensional data.

In this paper, an upper bound for the largest Mahalanobis distance obtainable from a multivariate sample is derived by posing the problem in a general mathematical context and it is then solved as a Lagrange optimization problem. The largest Mahalanobis distance obtainable in a sample of  $n$  multivariate  $p$  length vectors has upper bound  $(n - 1)^2/n$ , and when  $n > p + 1$ , the upper bound is attainable for  $k$  vectors, for each  $k = 1, \dots, p$ . Lower bounds for the largest Mahalanobis distance are also obtained in various cases, using inter alia a *dual approach* to the optimization problems; the attainability of these bounds is demonstrated. Finally, the maximum and minimum values of  $(x_i - \bar{x})'S^{-1}(x_j - \bar{x})$ , for  $i \neq j$ , are obtained.

The assumption of multivariate normality was introduced to motivate the estimates  $\bar{x}$  and  $S$ , but it is not essential and so is dispensed with for the remainder of the paper; so henceforth we need not regard the  $x_i$ s as random variables. We conclude this section by framing the problems in a generalized mathematical setting as follows. Let  $x_1, \dots, x_n$  be  $n$  vectors (data points, observations) in  $\mathbb{R}^p$ . Define the vector

$$\bar{x} = \frac{1}{n} \sum_{i=1}^n x_i \tag{1.1}$$

and the  $p \times p$  matrix

$$S = \frac{1}{n-1} \sum_{i=1}^n (x_i - \bar{x})(x_i - \bar{x})'. \tag{1.2}$$

We assume that the matrix  $S$  thus obtained is invertible and positive definite (thus requiring that  $n \geq p$ ). The  $n \times n$  matrix  $H$  is defined by its components

$$h_{ij} = \frac{1}{n-1} (x_i - \bar{x})' S^{-1} (x_j - \bar{x}). \tag{1.3}$$

$H$  is a non-negative, symmetric, idempotent matrix with trace  $p$ . Also observe that each row and column of  $H$  sums to zero, that is

$$\sum_{j=1}^n h_{ij} = \sum_{j=1}^n h_{ji} = 0.$$

For vector  $x_i$  define the Mahalanobis distance  $D_i$  by

$$D_i^2 = (n-1)h_{ii}.$$

Define, for a given set of points  $x_1, \dots, x_n$ , the maximum Mahalanobis distance

$$M = \max_{i=1, \dots, n} D_i^2.$$

Our goal is to optimize  $M$  and the  $h_{ij}$ s over points in  $\mathbb{R}^p$ , subject to the constraints that  $\bar{x}$  and  $S$  are fixed. This is achieved using the attractive and natural framework of Lagrangian methods.

## 2. The maximum Mahalanobis distance

First we find the maximum<sup>1</sup> of  $h_{kk}$ , for a given  $k$  subject to the mean and variance, as defined in (1.1) and (1.2), being fixed. Without loss of generality, we take the mean  $\bar{x} = 0$ , following an overall translation of all the (data) points. The rank of  $S$  must be  $p$ , since it is invertible. This fact, along with the mean constraint, requires that the number of points must satisfy  $n \geq p + 1$ . We will show that the maximum value of  $h_{kk}$  is  $(n-1)/n$  and that a necessary and sufficient condition for this to be achieved is that the vectors  $x_2, \dots, x_n$  be coplanar (subject to  $S$  being invertible). We also construct an example, using a suitable choice of data points, which demonstrates that it is possible to make the  $p$  largest  $h_{kk}$ s all equal to  $(n-1)/n$ ; that is the  $p$  largest Mahalanobis distances are all equal to  $(n-1)^2/n$ .

**Theorem 2.1.** *The maximum value of*

$$h_{kk} = \frac{x_k' S^{-1} x_k}{n-1}$$

is

$$\frac{n-1}{n},$$

where  $x_1, \dots, x_n$  are points in  $\mathbb{R}^p$ , constrained such that

$$\sum_{i=1}^n x_i = 0$$

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<sup>1</sup> Technically, we find the supremum  $\sup h_{kk}$  over sets of  $n$  vectors in  $\mathbb{R}^p$ , with  $\bar{x}$  and  $S$  fixed. Using the crude estimate from the trace of  $H$  that  $h_{kk} \leq p$ , the  $h_{kk}$ s are bounded above and hence the supremum exists.

and the  $p \times p$  matrix

$$S = \frac{1}{n-1} \sum_{i=1}^n x_i x_i'$$

is invertible.

**Proof.** Without loss of generality, we take  $k = 1$ . For convenience, instead of  $S$ , we use

$$\sigma = \frac{n-1}{n} S = \frac{1}{n} \sum_{i=1}^n x_i x_i'. \tag{2.1}$$

The problem is posed as a Lagrange optimization problem, with the Lagrangian function

$$L = x_1' \sigma^{-1} x_1 + \lambda' \sum_{i=1}^n x_i + \sum_{i=1}^n (x_i' \Omega x_i - \text{trace}(\Omega \sigma)),$$

where  $\lambda$  and  $\Omega$  are Lagrange multipliers and  $\Omega$  is a  $p \times p$  symmetric matrix. Taking derivatives with respect to  $x_i$  gives

$$(\sigma^{-1} + \Omega)x_1 = -\frac{1}{2}\lambda \tag{2.2}$$

and

$$\Omega x_i = -\frac{1}{2}\lambda \tag{2.3}$$

for  $i = 2, \dots, n$ . The derivative with respect to  $\lambda$  returns the constraint  $\bar{x} = 0$ . Combining these equations gives

$$\Omega x_1 = \frac{(n-1)}{2}\lambda \tag{2.4}$$

and

$$\sigma^{-1} x_1 = -\frac{n}{2}\lambda. \tag{2.5}$$

Using the definition of  $\sigma$ , multiplied on the right by  $\Omega$ , along with Eqs. (2.3), (2.4) and the constraint  $\bar{x} = 0$ , gives

$$\sigma \Omega = \frac{1}{2} x_1 \lambda'. \tag{2.6}$$

It follows from this that, for  $p \geq 2$ ,  $\det(\sigma \Omega) = 0$  and consequently,  $\Omega$  is not invertible.

Define the  $n \times p$  matrix  $X$  with row vectors  $x_i'$ , so that  $n\sigma = X'X$  and  $x_i' \sigma^{-1} x_i = n h_{ii}$ , where  $h_{ij}$  is the generic entry of the  $n \times n$  matrix

$$H = X(X'X)^{-1}X'$$

$H$  corresponds to the *hat matrix* in the language of statistical outliers. From Eq. (2.5), we obtain

$$h_{11} = \frac{1}{n} x_1' \sigma^{-1} x_1 = -\frac{1}{2} x_1' \lambda.$$

Multiplying Eq. (2.6) on the left by  $\sigma^{-1}$  and on the right by  $x_1$  gives

$$\Omega x_1 = \frac{1}{2} \sigma^{-1} x_1 (\lambda' x_1).$$

Eqs. (2.4) and (2.5) can be used to substitute for  $\Omega x_1$  and  $\sigma^{-1}x_1$  in terms of  $\lambda$  and together yield

$$\frac{(n-1)}{2}\lambda = -\frac{n}{4}\lambda(\lambda'x_1).$$

If the trivial case of  $\lambda = 0$  is excluded (which with Eq. (2.5) implies  $x_1 = 0$ ), we get

$$\lambda'x_1 = -\frac{2(n-1)}{n}$$

and hence  $h_{11} = (n-1)/n$  is the maximum possible value of a diagonal element of the matrix  $H$ . So

$$M = \frac{(n-1)^2}{n} \tag{2.7}$$

is the maximum possible Mahalanobis distance.  $\square$

**Corollary 2.2.** *The points  $x_2, \dots, x_n$  must be coplanar in  $\mathbb{R}^p$  for the maximum (2.7) to be achieved.*

**Proof.** From Eq. (2.6), it follows that

$$\sigma\Omega x_i = (x'_i\lambda)\frac{x_1}{2}$$

for  $i = 2, \dots, n$ . Substituting from Eqs. (2.3) and (2.5) then yields

$$\sigma\left(-\frac{1}{2}\lambda\right) = \left(-\frac{2}{n}x'_i\sigma^{-1}x_1\right)\frac{x_1}{2}.$$

Again using (2.5) on the left side and assuming  $x_1 \neq 0$ , we find that  $x'_i\sigma^{-1}x_1 = -1$  and so  $x'_iS^{-1}x_1 = -(n-1)/n$ , for  $i = 2, \dots, n$ . This implies that the elements of the first column (and row) of  $H$  satisfy  $h_{i1} = -1/n$ , for  $i = 2, \dots, n$ . Furthermore, we deduce that since  $S^{-1}x_1 \neq 0$ , the points  $x_2, \dots, x_n$  must be coplanar in  $\mathbb{R}^p$  for the maximum  $h_{11}$  to be achieved. (In the case of  $p = 1$  the remaining  $(n-1)$  observations must be coincident and  $h_{jj} = \frac{1}{n(n-1)}$  for  $j > 1$ .)  $\square$

**Corollary 2.3.** *The condition of Corollary 2.2 is not only necessary for achieving the maximum value (2.7), but is also sufficient.*

**Proof.** Assume that for  $i = 2, \dots, n$ , there exists a vector  $v \neq 0$  and a scalar  $\rho$  such that

$$v'x_i = \rho, \tag{2.8}$$

i.e., the points  $x_2, \dots, x_n$  lie in a common plane in  $\mathbb{R}^p$ . Using the assumption  $\bar{x} = 0$ , we also obtain  $v'x_1 = -(n-1)\rho$ . Then

$$(n-1)Sv = \sum_{n=1}^n x_i x'_i v = -(n-1)\rho x_1 + \sum_{i=2}^n \rho x_i = -n\rho x_1.$$

Since  $S$  is invertible, the right side is non-zero and in particular  $\rho \neq 0$ . Hence,

$$v = -\frac{n}{n-1}\rho S^{-1}x_1 \tag{2.9}$$

and

$$(n - 1)\rho = -v'x_1 = \frac{n}{(n - 1)}\rho x_1'S^{-1}x_1.$$

Dividing by  $\rho$ , it follows that

$$x_1'S^{-1}x_1 = \frac{(n - 1)^2}{n} = M$$

as defined in Eq. (2.7), and  $h_{11} = (n - 1)/n$ . From (2.9) we see that  $S^{-1}x_1$  is a normal vector to the plane. Also, by combining Eqs. (2.8) and (2.9), we find, as before, that  $x_i'S^{-1}x_1 = -(n - 1)/n$  for  $i = 2, \dots, n$ .  $\square$

We now demonstrate that the maximum is achievable for any given  $n$  and  $p$  with  $n > p + 1$ . In so doing, we address a more general question viz. what is the maximum number of the diagonal elements of the matrix  $H$  that can simultaneously achieve the maximum value  $(n - 1)/n$ ? From the property  $\text{trace}(H) = \sum_{i=1}^n h_{ii} = p$ , it follows that, if  $n > p + 1$ , then at most  $p$  of the  $h_{ii}$ s can be equal to the maximum value  $(n - 1)/n$ . The example below illustrates that this is achievable.

**Example 2.1.** Let  $e_\alpha$  be the standard basis of  $\mathbb{R}^p$ , i.e.,  $(e_\alpha)_\beta = \delta_{\alpha\beta}$ . Define

$$x_i = \begin{cases} e_i, & i = 2, \dots, p, \\ e_1, & i = p + 1, \dots, n. \end{cases}$$

Note that the points  $x_2, \dots, x_n$  are all coplanar, lying in the plane  $\xi_1 + \dots + \xi_p = 1$ . From the zero mean property  $x_1$  is found to be

$$x_1 = -(n - p)e_1 - \sum_{j=2}^p e_j.$$

Then  $S = (n - 1)^{-1} \sum_{i=1}^n x_i x_i'$  gives

$$S = \frac{1}{(n - 1)} \left[ (n - p)(n - p + 1)e_1 e_1' + (n - p) \sum_{i=2}^p (e_1 e_i' + e_i e_1') + \sum_{i=2}^p \sum_{j=2}^p e_i e_j' + \sum_{i=2}^p e_i e_i' \right].$$

From this, one may verify that the inverse is

$$S^{-1} = \frac{(n - 1)}{n} \left[ \frac{p}{n - p} e_1 e_1' - \sum_{i=2}^p (e_1 e_i' + e_i e_1') - \sum_{i=2}^p \sum_{j=2}^p e_i e_j' + n \sum_{i=2}^p e_i e_i' \right].$$

Using this in definition (1.3) of the generic entry  $h_{ij}$  of the matrix  $H$ , we obtain

$$H = \left( \begin{array}{cccc|ccc} \frac{n-1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ -\frac{1}{n} & \frac{n-1}{n} & \dots & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ -\frac{1}{n} & -\frac{1}{n} & \dots & \frac{n-1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ \hline -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & \frac{p}{n(n-p)} & \dots & \frac{p}{n(n-p)} \\ -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & \frac{p}{n(n-p)} & \dots & \frac{p}{n(n-p)} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & \frac{p}{n(n-p)} & \dots & \frac{p}{n(n-p)} \end{array} \right),$$

where the top left block is  $p \times p$ . One can verify that  $H$  has all the desired properties listed above (symmetric, each row sums to zero, it is idempotent and  $\text{trace}(H) = p$ ). Observe that  $h_{11} = h_{22} = \dots = h_{pp} = (n - 1)/n$ , the maximum possible number of maximum diagonal values. (Indeed if  $n = p + 1$  then  $h_{(p+1)(p+1)} = (n - 1)/n$  also.) This proves that

$$\max\{(h_{11} + \dots + h_{kk})\} = \frac{k(n - 1)}{n}$$

for each  $k = 1, \dots, p$ .

### 3. Lower bounds for the maximum Mahalanobis distance

**Theorem 3.1.** A lower bound for the quantity  $M = \max_i \{x'_i S^{-1} x_i\}$ , is

$$m_g = \frac{(n - 1)p}{n}, \tag{3.1}$$

where  $x_1, \dots, x_n$  are points in  $\mathbb{R}^p$  that are subject to the constraint

$$\bar{x} = \sum_{i=1}^n x_i = 0$$

and to the  $p \times p$  matrix  $S = (n - 1)^{-1} \sum_{i=1}^n x_i x'_i$  being invertible.

We call  $m_g$  the global lower bound for  $M$ .

**Proof.** The sum of the Mahalanobis distances for the vectors  $x_1, \dots, x_n$  is

$$\sum_{i=1}^n x'_i S^{-1} x_i = \text{trace}((n - 1)H) = (n - 1)p, \tag{3.2}$$

where the  $n \times n$  matrix  $H$  defined in Eq. (1.3). Since each term in the sum is non-negative, the maximum term  $M \geq \frac{(n-1)p}{n}$ .  $\square$

**Corollary 3.2.**  $M = m_g$  if and only if

$$x'_i S^{-1} x_i = \frac{(n - 1)p}{n} \tag{3.3}$$

for all  $i = 1, \dots, n$ .

This occurs if and only if  $h_{ii} = p/n$  for all  $i$ . We emphasize that the infimum of  $M$  will equal the global lower bound  $m_g$ , only for certain values of  $n$  and  $p$ . We explore this in some detail in the following examples.

**Example 3.1.** When  $p = 1$ , Eq. (3.3) implies that  $M = m_g$ , the global lower bound, if and only if

$$x_i^2 = \frac{(n - 1)S}{n}$$

for all  $i$ . Taking this together with the mean constraint  $\sum_{i=1}^n x_i = 0$ , we reach the following two conclusions; when  $n$  is even, the global lower bound is achieved when exactly half of the  $x_i$ s equal

$\sqrt{(n-1)S/n}$  and the other half equal  $-\sqrt{(n-1)S/n}$  (and hence  $\inf M = m_g = (n-1)/n$  for  $p = 1$  and  $n$  even), as was derived in Hayes [4]; when  $n$  is odd, the global lower bound is *never achieved* as there is no combination of  $\pm\sqrt{(n-1)S/n}$  with an odd number of terms that sums to zero. Note that this does not preclude the infimum of  $M$  being equal to  $m_g$ , although we do not show this. However, it is straightforward to put an upper bound on  $\inf M$  by taking the vector of  $x_i$ s of the type  $(a, -a, a, -a, \dots, a, -a, 0)$  thus obtaining the bounds

$$\frac{n-1}{n} \leq \inf M \leq 1$$

for  $p = 1$  and  $n$  odd.

**Example 3.2.** Next we examine the case of  $p = 2$ . Consider  $n$  (data) points uniformly distributed on the circle of radius  $a$ ; we take  $x_j = (a \cos \omega_j, a \sin \omega_j)$ , where  $\omega_j = 2\pi(j-1)/n$  for  $j = 1, \dots, n$ . Then  $\sum_{j=1}^n x_j = 0$ , for  $n > 2$  since, using complex exponential notation,

$$\sum_{j=1}^n e^{\frac{2\pi i(j-1)}{n}} = 0.$$

Also

$$x_j x'_j = a^2 \begin{pmatrix} \cos^2 \omega_j & \cos \omega_j \sin \omega_j \\ \cos \omega_j \sin \omega_j & \sin^2 \omega_j \end{pmatrix} = \frac{a^2}{2} \begin{pmatrix} 1 + \cos 2\omega_j & \sin 2\omega_j \\ \sin 2\omega_j & 1 - \cos 2\omega_j \end{pmatrix}.$$

Now

$$\sum_{j=1}^n e^{2i\omega_j} = \sum_{j=1}^n e^{\frac{4\pi i(j-1)}{n}} = 0$$

for  $n > 2$ , giving

$$(n-1)S = \sum_{j=1}^n x_j x'_j = \frac{na^2}{2} I_2,$$

and so

$$S = \frac{na^2}{2(n-1)} I_2,$$

where  $I_2$  is the  $2 \times 2$  identity matrix. Hence,

$$x'_j S^{-1} x_j = \frac{2(n-1)}{na^2} \|x_j\|^2 = \frac{2(n-1)}{n}$$

for all  $j$ . This is  $m_g$  of Theorem 3.1 with  $p = 2$ . In conclusion, for  $p = 2$  and all  $n > 2$ ,  $\inf M = m_g$  and the bound is achievable.

**Example 3.3.** Next we examine the general case, assuming that  $n \geq 2p$ . Define  $R = n \bmod 2p$ , so  $n = 2dp + R$ , where  $d \geq 1$ , and  $0 \leq R \leq 2p - 1$ . Define

$$x_i = \begin{cases} (-1)^{i+1} e_{\lfloor \frac{i+1}{2} \bmod p} & \text{if } i \leq 2dp, \\ 0 & \text{if } 2dp < i \leq n, \end{cases}$$

where we define  $e_0 = e_p$ . Then  $x_1 = e_1, x_2 = -e_1, \dots, x_{2p-1} = e_p, x_{2p} = -e_p, x_{2p+1} = e_1, x_{2p+2} = -e_1, \dots, x_{2dp-1} = e_p, x_{2dp} = -e_p, x_{2dp+1} = x_{2dp+2} = \dots = x_n = 0$ . Clearly,  $\sum_{i=1}^n x_i = 0$  and also



$$(n - 1)S = 2d \sum_{\alpha=1}^p e_{\alpha}e'_{\alpha} = 2dI_p,$$

so

$$\frac{1}{n - 1}S^{-1} = \frac{1}{2d}I_p = \frac{1}{2\lfloor \frac{n}{2p} \rfloor}I_p,$$

where  $I_p$  is the  $p \times p$  identity matrix. The diagonal elements of the  $H$  matrix are

$$h_{ii} = \begin{cases} \frac{1}{2\lfloor \frac{n}{2p} \rfloor} & \text{if } 1 \leq i \leq 2p\lfloor \frac{n}{2p} \rfloor, \\ 0 & \text{if } 2p\lfloor \frac{n}{2p} \rfloor < i \leq n. \end{cases}$$

These values satisfy the condition  $\sum_{i=1}^n h_{ii} = p$ , and it follows that

$$\max_i \{h_{ii}\} = \frac{1}{2d} = \frac{1}{2\lfloor \frac{n}{2p} \rfloor}.$$

So the maximum Mahalanobis distance for these data is

$$M_1 = \frac{n - 1}{2\lfloor \frac{n}{2p} \rfloor}. \tag{3.4}$$

Clearly,  $M_1 = m_g$ , the global lower bound, when  $2p$  divides evenly into  $n$ .

#### 4. A dual approach to optimization problems

A dual approach to the above problems and specifically the infimum of  $M$  (the minimax) problem is presented. This will also illustrate the fact that these problems are fundamentally algebraic and geometric in nature, though their origins are rooted in statistics.

We assume that  $S$  is invertible and positive definite. There then exists a rotation matrix  $\tilde{R}$  such that  $(n - 1)S = \tilde{R}'\Delta\tilde{R}$ , where  $\Delta$  is a diagonal matrix with diagonal elements  $\delta_{\alpha\alpha} > 0$ . Define  $y_i = \tilde{R}x_i$  for  $i = 1, \dots, n$ . The zero-mean condition is preserved, i.e.,  $\sum_{i=1}^n y_i = 0$  and

$$\Delta = \sum_{i=1}^n y_i y'_i. \tag{4.1}$$

The elements of the  $H$  matrix, defined in (1.3) become  $h_{ij} = y'_i \Delta^{-1} y_j$  and the various problems then consist of finding  $n$  vectors in  $\mathbb{R}^p$  with certain desired optimization properties. We now transpose the problems by defining the vectors  $z_{\alpha}$  ( $\alpha = 1, \dots, p$ ) in  $\mathbb{R}^n$  by

$$(z_{\alpha})_i = \frac{(y_i)_{\alpha}}{\sqrt{\delta_{\alpha\alpha}}}.$$

Then (4.1) becomes  $z'_{\alpha} z_{\beta} = \delta_{\alpha\beta}$ , that is  $\{z_{\alpha}\}$  forms an orthonormal set in  $\mathbb{R}^n$ . The zero mean condition becomes  $z'_{\alpha} u = 0$  where  $u$  is the vector of ones in  $\mathbb{R}^n$ . The  $H$  matrix elements are

$$h_{ij} = \sum_{\alpha=1}^p (f'_i z_{\alpha})(f'_j z_{\alpha}),$$

where  $\{f_i\}$  is the standard basis of  $\mathbb{R}^n$  and the Mahalanobis distances are

$$D_i = \sqrt{(n-1) \sum_{\alpha=1}^p (f'_i z_\alpha)^2}.$$

As before, we define, for a given set of vectors  $x_1, \dots, x_n$  in  $\mathbb{R}^p$ ,  $M = \max_i D_i^2$  which in the transposed version is defined for a given set of vectors  $z_1, \dots, z_p$  in  $\mathbb{R}^n$ . We now define, for given  $n$  and  $p$ ,

$$M_{n,p} = \sup M \tag{4.2}$$

and

$$m_{n,p} = \inf M, \tag{4.3}$$

where the supremum and infimum are taken over sets of  $p$  orthonormal vectors  $z_\alpha$  in  $\mathbb{R}^n$ , which are orthogonal to  $u$ . It follows that  $M_{n,p} = \sup D_1^2 = (n-1) \sup h_{11}$ . So the previous problems are transformed into optimization of the  $M$  over orthonormal sets of  $p$  vectors in the orthogonal space of  $u = (1, 1, \dots, 1)$  in  $\mathbb{R}^n$ . These problems depend only on  $n$  and  $p$  and not on the specific nature of any given matrix  $S$  (such as its eigenvalues). We have shown that

$$M_{n,p} = \frac{(n-1)^2}{n}.$$

We will not present a similar general formula for  $m_{n,p}$  but rather present some bounds; specifically we have the global lower bound  $m_{n,p} \geq m_g$  for all  $n, p$  and the upper bound  $M_1$  in (3.4). We now use the dual approach to obtain another upper bound. Both of these bounds give rise to an infinite number of instances when  $m_{n,p} = m_g$ .

We assume here that  $n \geq 2p$  and define

$$\mu = \lceil \log_2 p \rceil + 1. \tag{4.4}$$

Then  $p < 2^\mu \leq n$ . We write  $n = \kappa 2^\mu + \phi$ , where  $\kappa = \lfloor \frac{n}{2^\mu} \rfloor$  with  $\kappa \geq 1$  and  $0 \leq \phi < 2^\mu$ . We now construct a  $2^\mu \times 2^\mu$  orthogonal matrix  $Q$  by defining the  $(a, b)$  element, for  $a, b \in \{0, 1, \dots, 2^\mu - 1\}$  as follows: we write the binary expansions of  $a$  and  $b$ , respectively, as  $a = \sum_{l=1}^\mu 2^{l-1} a_l$  and  $b = \sum_{l=1}^\mu 2^{l-1} b_l$  (with  $a_k, b_k \in \{0, 1\}$ ) and then we define

$$Q_{ab} = 2^{-\frac{\mu}{2}} (-1)^{\sum_{k=1}^\mu a_k b_k}.$$

Note that the rows of  $2^{\frac{\mu}{2}} Q$  consist of  $\pm 1$ s; the first row is  $Q_{0\star} = 2^{-\frac{\mu}{2}} (1, 1, \dots, 1)$  and all other rows are orthogonal to this. We now choose the  $p$  vectors  $z_\alpha$  in  $\mathbb{R}^n$  as follows:

$$z_\alpha = \frac{1}{\sqrt{\kappa}} \underbrace{(Q_{\alpha\star}, Q_{\alpha\star}, \dots, Q_{\alpha\star})}_{\kappa \text{ times}}, \underbrace{(0, \dots, 0)}_{\phi \text{ times}}$$

for  $\alpha = 1, \dots, p$ , where  $Q_{\alpha\star}$  consists of the  $2^\mu$  elements of row  $\alpha$  of the matrix  $Q$  defined above. These vectors form an orthonormal set in  $\mathbb{R}^n$  and each one is orthogonal to  $Q_{0\star}$  and hence to  $u$ , the vector of ones in  $\mathbb{R}^n$ . For each  $\alpha = 1, \dots, p$ , observe that  $(f'_i z_\alpha)^2 = \frac{1}{\kappa 2^\mu}$  for  $i = 1, \dots, \kappa 2^\mu$  and  $(f'_i z_\alpha)^2 = 0$  for  $i = \kappa 2^\mu + 1, \dots, n$ . So the corresponding Mahalanobis distances  $D_i$  satisfy

$$D_i^2 = \begin{cases} \frac{(n-1)p}{\kappa 2^\mu} & \text{if } 1 \leq i \leq \kappa 2^\mu, \\ 0 & \text{if } \kappa 2^\mu < i \leq n. \end{cases}$$

So the maximum Mahalanobis distance  $M_2 = \max_i D_i^2$  for these data is given by

$$M_2 = \frac{(n - 1)p}{\kappa 2^\mu} = \frac{(n - 1)p}{2^\mu \lfloor \frac{n}{2^\mu} \rfloor} \tag{4.5}$$

with  $\mu$  as given in (4.4). Note that this produces an infinite number of instances for which  $M_2 = m_g$ , namely if there exists  $\nu$  such that  $2^\nu \geq 2p$  and  $2^\nu$  divides evenly into  $n$ . Note also that the bounds  $M_1$  in (3.4) and  $M_2$  in (4.5) are independent, as illustrated by the following examples: when  $n = 6$  and  $p = 3$ ,  $M_1 = \frac{5}{2}$  and  $M_2 = \frac{15}{4}$ , so  $M_1 < M_2$ ; when  $n = 8$  and  $p = 3$ ,  $M_1 = \frac{7}{2}$  and  $M_2 = \frac{21}{8}$  so  $M_1 > M_2$ . In conclusion, if  $n \geq 2p$ , then

$$m_g \leq m_{n,p} \leq \min \left( \frac{n - 1}{2 \lfloor \frac{n}{2p} \rfloor}, \frac{(n - 1)p}{2^\mu \lfloor \frac{n}{2^\mu} \rfloor} \right),$$

where  $\mu$  is given in (4.4).

### 5. Optimization of the elements of the $H$ matrix

We have already shown that for any given diagonal element  $h_{kk}$  of  $H$ ,

$$0 \leq h_{kk} \leq \frac{n - 1}{n}$$

and that both bounds are achievable (the lower bound being trivial).

It remains to address the optimization of the off-diagonal elements. Without loss of generality, we search for the optima of

$$h_{12} = \frac{1}{n - 1} x_1 S^{-1} x_2.$$

We again introduce the variable  $\sigma = \frac{n}{n-1} S$ , so that  $h_{12} = (x_1 \sigma^{-1} x_2) / n$ . We treat the problem as a constrained optimization with the Lagrangian function

$$L = x_1' \sigma^{-1} x_2 + \lambda' \sum_{i=1}^n x_i + \sum_{i=1}^n (x_i' \Omega x_i - \text{trace}(\Omega \sigma)),$$

where again  $\Omega$  is symmetric. Taking derivatives with respect to  $x_1, x_2$  and  $x_i$  (for  $i \geq 3$ ) gives

$$\begin{aligned} \sigma^{-1} x_2 + 2\Omega x_1 + \lambda &= 0, \\ \sigma^{-1} x_1 + 2\Omega x_2 + \lambda &= 0, \\ 2\Omega x_i + \lambda &= 0, \quad i = 3, \dots, n. \end{aligned} \tag{5.1}$$

Adding these equations gives

$$\sigma^{-1} (x_1 + x_2) = -n\lambda. \tag{5.2}$$

From the definition of  $\sigma$ , as in (2.1), multiplied on the right by  $\Omega$ , and using the fact that  $\bar{x} = 0$  along with Eqs. (5.1), we find that

$$-2n\sigma\Omega = x_1(\sigma^{-1}x_2)' + x_2(\sigma^{-1}x_1)'$$

Hence,

$$-2\sigma\Omega x_i = h_{2i}x_1 + h_{1i}x_2 \tag{5.3}$$

for  $i = 1, 2, 3, \dots, n$ . Now multiplying Eqs. (5.1) on the left by  $(-n\sigma)$ , and eliminating  $\lambda$  using (5.2), gives

$$\begin{aligned} x_1 + (1 - n)x_2 - 2n\sigma\Omega x_1 &= 0, \\ x_2 + (1 - n)x_1 - 2n\sigma\Omega x_2 &= 0, \\ x_1 + x_2 - 2n\sigma\Omega x_i &= 0, \quad i = 3, \dots, n. \end{aligned}$$

Combining these with (5.3), we obtain

$$\begin{aligned} x_1(1 + nh_{12}) + x_2(nh_{11} - (n - 1)) &= 0, \\ x_1(nh_{22} - (n - 1)) + x_2(1 + nh_{12}) &= 0, \\ x_1(1 + nh_{2i}) + x_2(1 + nh_{1i}) &= 0, \quad i = 3, \dots, n. \end{aligned} \tag{5.4}$$

We assume that  $x_1 \neq 0$  and  $x_2 \neq 0$ ; otherwise  $h_{12} = 0$ . If  $x_1$  and  $x_2$  are linearly independent, then these equations give  $h_{11} = \frac{n-1}{n} = h_{22}, h_{12} = -\frac{1}{n}$  and  $h_{1i} = h_{2i} = -\frac{1}{n}$  for  $i \geq 3$ . This case does not correspond to an extremum of  $h_{12}$ . Otherwise, there exists  $\gamma$  such that  $x_2 = \gamma x_1$ . Then  $h_{12} = \gamma h_{11}, h_{22} = \gamma^2 h_{11}$  and  $h_{2i} = \gamma h_{1i}$ . So, Eqs. (5.4) become

$$\begin{aligned} 1 + n\gamma h_{11} + \gamma(nh_{11} - n + 1) &= 0, \\ n\gamma^2 h_{11} - (n - 1) + \gamma(1 + n\gamma h_{11}) &= 0, \\ 1 + n\gamma h_{1i} + \gamma(1 + nh_{1i}) &= 0. \end{aligned} \tag{5.5}$$

Taking  $\gamma$  times the first equation minus the second equation gives  $\gamma^2 = 1$ .

If  $\gamma = 1$ , solving the equations gives  $h_{11} = \frac{n-2}{2n} = h_{12} = h_{22}$  and  $h_{1i} = h_{2i} = -\frac{1}{n}$  for  $i \geq 3$ ; if  $\gamma = -1$ , then  $h_{11} = \frac{1}{2} = h_{22}, h_{12} = -\frac{1}{2}$  and  $h_{1i} = h_{2i} = 0$  for  $i \geq 3$ . So the maximum value of  $h_{12}$  is  $\frac{1}{2} - \frac{1}{n}$  and the minimum is  $-\frac{1}{2}$ . In conclusion we find that for all  $n > 2$ ,

$$-\frac{1}{2} \leq h_{12} \leq \frac{1}{2} - \frac{1}{n}.$$

We will now demonstrate how these extreme values of  $h_{12}$  are achieved by presenting actual data that reproduce them. It is illustrative to consider first the case of one dimension,  $p = 1$ . In that instance, the  $\gamma = 1$  case corresponds to  $n$  points:  $(n - 2)r, (n - 2)r, -2r, -2r, \dots, -2r$  where  $r^2 = \frac{\sigma}{2n-4}$ ; the  $\gamma = -1$  case corresponds to  $n$  points:  $r, -r, 0, 0, \dots, 0$  where  $r^2 = \frac{n\sigma}{2}$ .

We now consider the higher dimensional cases, where  $p \geq 2$ . Observe that for  $\sigma$  to be invertible, i.e., of full rank, it is necessary that  $n \geq p + 2$ , since we now have the additional constraint on the  $x_i$ s that  $x_2 = \pm x_1$ .

**Example 5.1.** We present an example corresponding to  $\gamma = 1$ , in which one can observe that  $h_{12} = \frac{1}{2} - \frac{1}{n}$ . Let  $e_\alpha$  be the standard basis of  $\mathbb{R}^p$  and define

$$x_i = \begin{cases} -\frac{1}{2} \left( (n - p - 1)e_p + \sum_{j=1}^{p-1} e_j \right), & i = 1, 2, \\ e_{i-2}, & i = 3, \dots, p + 1, \\ e_p, & i = p + 2, \dots, n. \end{cases}$$

Then, Eq. (2.1) gives

$$\sigma = \frac{1}{2n} \left[ 2 \sum_{i=1}^{p-1} e_i e_i' + \sum_{i=1}^{p-1} \sum_{j=1}^{p-1} e_i e_j' + (n - p - 1) \sum_{i=1}^{p-1} (e_p e_i' + e_i e_p') + (n - p - 1)(n - p + 1) e_p e_p' \right].$$

From this, one may verify that the inverse is

$$\sigma^{-1} = n \sum_{i=1}^{p-1} e_i e'_i - \sum_{i=1}^{p-1} \sum_{j=1}^{p-1} e_i e'_j - \sum_{i=1}^{p-1} (e_p e'_i + e_i e'_p) + \frac{p+1}{n-p-1} e_p e'_p.$$

Using this, we then calculate the matrix

$$H = \left( \begin{array}{cc|ccc|ccc} \frac{n-2}{2n} & \frac{n-2}{2n} & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ \frac{n-2}{2n} & \frac{n-2}{2n} & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ \hline \frac{1}{n} & -\frac{1}{n} & \frac{n-1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ -\frac{1}{n} & -\frac{1}{n} & -\frac{1}{n} & \frac{n-1}{n} & \dots & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ -\frac{1}{n} & -\frac{1}{n} & -\frac{1}{n} & -\frac{1}{n} & \dots & \frac{n-1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ \hline \frac{1}{n} & -\frac{1}{n} & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & \frac{p+1}{n(n-p-1)} & \dots & \frac{p+1}{n(n-p-1)} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ -\frac{1}{n} & -\frac{1}{n} & -\frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} & \frac{p+1}{n(n-p-1)} & \dots & \frac{p+1}{n(n-p-1)} \end{array} \right),$$

where the square diagonal blocks have sides respectively of size 2,  $p - 1$  and  $n - p - 1$ .

**Example 5.2.** We present an example corresponding to  $\gamma = -1$ , in which one can observe that  $h_{12} = -\frac{1}{2}$ . Define

$$x_i = \begin{cases} (-1)^{i-1} e_1, & i = 1, 2, \\ e_{i-1}, & i = 3, \dots, p, \\ -\sum_{j=2}^{p-1} e_j - (n-p-1)e_p, & i = p+1, \\ e_p, & i = p+2, \dots, n, \end{cases}$$

where again we take  $n \geq p + 2$  and also assume  $p \geq 3$ . The case  $p = 2$  is obtained in a similar analysis by omitting the second line of this definition and the sum term from the third line. Then, Eq. (2.1) gives

$$\sigma = \frac{1}{n} \left[ 2e_1 e'_1 + \sum_{i=2}^{p-1} e_i e'_i + \sum_{i=2}^{p-1} \sum_{j=2}^{p-1} e_i e'_j + (n-p-1) \sum_{i=2}^{p-1} (e_p e'_i + e_i e'_p) + (n-p)(n-p-1)e_p e'_p \right].$$

From this, one may verify that the inverse is

$$\begin{aligned} \sigma^{-1} = & \frac{n}{2} e_1 e'_1 + n \sum_{i=2}^{p-1} e_i e'_i - \frac{n}{n-2} \sum_{i=2}^{p-1} \sum_{j=2}^{p-1} e_i e'_j - \frac{n}{n-2} \sum_{i=2}^{p-1} (e_p e'_i + e_i e'_p) \\ & + \frac{n(p-1)}{(n-2)(n-p-1)} e_p e'_p. \end{aligned}$$

Using this, we then calculate the matrix

$$H = \left( \begin{array}{cc|cccc|ccc} \frac{1}{2} & -\frac{1}{2} & 0 & 0 & \cdots & 0 & 0 & \cdots & 0 \\ -\frac{1}{2} & \frac{1}{2} & 0 & 0 & \cdots & 0 & 0 & \cdots & 0 \\ \hline 0 & 0 & \frac{n-3}{n-2} & -\frac{1}{n-2} & \cdots & -\frac{1}{n-2} & -\frac{1}{n-2} & \cdots & -\frac{1}{n-2} \\ 0 & 0 & -\frac{1}{n-2} & \frac{n-3}{n-2} & \cdots & -\frac{1}{n-2} & -\frac{1}{n-2} & \cdots & -\frac{1}{n-2} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & -\frac{1}{n-2} & -\frac{1}{n-2} & \cdots & \frac{n-3}{n-2} & -\frac{1}{n-2} & \cdots & -\frac{1}{n-2} \\ \hline 0 & 0 & -\frac{1}{n-2} & -\frac{1}{n-2} & \cdots & -\frac{1}{n-2} & \frac{p-1}{(n-2)(n-p-1)} & \cdots & \frac{p-1}{(n-2)(n-p-1)} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & -\frac{1}{n-2} & -\frac{1}{n-2} & \cdots & -\frac{1}{n-2} & \frac{p-1}{(n-2)(n-p-1)} & \cdots & \frac{p-1}{(n-2)(n-p-1)} \end{array} \right),$$

where the square diagonal blocks have sides respectively of size 2,  $p - 1$  and  $n - p - 1$ .

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