

Two-Dimensional Binary Arrays with Good Autocorrelation

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Calabro and Wolf (1968, *Inform. Contr.* 11) investigate the autocorrelation properties of certain periodic two-dimensional arrays. This note points out a relationship between periodic $p \times q$ -arrays with two-level autocorrelation and difference sets in the group $C(p) \times C(q)$, where $C(n)$ denote the cyclic group of order n . This observation enables us to construct several families of such arrays, some of which are perfect.

1. Calabro and Wolf (1968) investigate the autocorrelation properties of certain periodic two-dimensional arrays. This note points out a relationship between periodic $p \times q$ -arrays with two-level autocorrelation and difference sets in the group $C(p) \times C(q)$, where $C(n)$ denotes the cyclic group of order n . This observation enables us to construct several families of such arrays, some of which are perfect, i.e. with zero autocorrelation for all out-of-phase shifts.

A periodic binary $p \times q$ -array is an infinite matrix $A = (a(i, j))$ with $i, j = 0, 1, 2, \dots$, where $a(i, j)$ is either 0 or 1 and where p, q are the smallest positive integers such that

$$a(i, j) = a(i, j + q) = a(i + p, j) \quad \text{for all } i, j.$$

When no confusion can arise we sometimes regard it as a $p \times q$ -matrix $A = (a(i, j))$ with $i = 0, 1, 2, \dots, p - 1, j = 0, 1, 2, \dots, q - 1$. Since this is the only type of arrays we shall be looking at in this note, we simply call them $p \times q$ -arrays. The autocorrelation function of A is given by

$$R_A(r, s) = \sum_{i=0}^{p-1} \sum_{j=0}^{q-1} b(i, j) b(i + r, j + s)$$

where $b(i, j) = 1 - 2a(i, j)$. We say A has two-level autocorrelation if R_A is of the form

$$R_A(r, s) = \begin{cases} pq & \text{if } r \equiv 0 \pmod{p} \text{ and } s \equiv 0 \pmod{q} \\ c & \text{if otherwise.} \end{cases} \quad (1)$$

A is said to be *perfect* if c (in (1)) is equal to zero.

A subset $D = \{g_1, \dots, g_k\}$ of a group G of order v is called a (v, k, λ) -difference set if the equation $g_i g_j^{-1} = g$ has exactly λ solutions for every $g \neq 1$ in G . In what follows we shall be looking at the case $G = C(p) \times C(q)$, and we shall use the additive notation.

2. It is well-known (see Baumert, 1971, Chap. I) that periodic binary sequences with two-level autocorrelation correspond to cyclic difference sets (i.e. difference sets in cyclic groups). The following result is a natural generalization.

PROPOSITION 1. (a) Let D be a (pq, k, λ) -difference set in $C(p) \times C(q)$, and let $A = (a(i, j))$ be defined by

$$a(i, j) = \begin{cases} 0 & \text{if } (i, j) \in D \\ 1 & \text{if otherwise,} \end{cases}$$

then A has two-level autocorrelation with c (in (1)) equal to $pq - 4(k - \lambda)$.

(b) Let $A = (a(i, j))$ be a $p \times q$ -array with two-level autocorrelation (see (1)), then $pq + pqc - c$ is a perfect square, say m^2 , and $D = \{(i, j) \in C(p) \times C(q) \mid a(i, j) = 0\}$ is a difference set with parameters $(pq, \frac{1}{2}(pq - m), \frac{1}{4}(pq + c - 2m))$ or $(pq, \frac{1}{2}(pq + m), \frac{1}{4}(pq + c + 2m))$. If we restrict our attention to $p \times q$ -arrays with more 1's than 0's, then the former is the case.

Proof. (a) Let $\lambda(r, s)$ be the number of matching 0's between A and the (r, s) -translate of A (i.e. with (i, j) th entry equal to $a(i + r, j + s)$). It is not hard to see that, for $r \not\equiv 0 \pmod{p}$ or $s \not\equiv 0 \pmod{q}$, $\lambda(r, s)$ is equal to the number of solution pairs (d_i, d_j) with $d_i, d_j \in C(p) \times C(q)$ and $d_i - d_j = (r, s)$, which is a constant ($= \lambda$) by hypothesis. From this we can compute c , which turns out to be $pq - 4(k - \lambda)$.

(b) Note that

$$pq + (pq - 1)c = \sum_{r=0}^{p-1} \sum_{s=0}^{q-1} R_A(r, s) = m^2,$$

where $m = \sum_{i=0}^{p-1} \sum_{j=0}^{q-1} b(i, j)$ is the difference between the number of 0's and 1's. Again, by similar argument as in (a), we have $c = pq - 4(k - \lambda(r, s))$ where k is the number of 0's. This implies $\lambda(r, s)$ is a constant for $r \not\equiv 0 \pmod{p}$ or $s \not\equiv 0 \pmod{q}$. Finally, k can be computed from m and pq . Q.E.D.

EXAMPLE I. Let $f: G \rightarrow G'$ be an isomorphism of groups. It is clear that D is a difference set in G if and only if $f(D)$ is a difference set in G' . We thus obtain $p \times q$ -arrays with two-level autocorrelation when $(p, q) = 1$ by choosing cyclic difference sets with parameters (pq, k, λ) , using the isomorphism $f: C(pq) \rightarrow C(p) \times C(q)$ given by $f(s) = (xs, ys)$ where $xq + yp = 1$. In particular, case (b) of Theorem 2 in Calabro and Wolf (1968) is obtained by choosing the twin-prime set (see Baumert (1971), Theorem 5.27) when p, q are twin primes.

If D is a difference set in a group G , then for $g \in G$, its g -translate $D + g = \{d + g \mid d \in D\}$ is also a difference set. Two difference sets D_1, D_2 are said to be equivalent if $D_2 = f(D_1) + g$ for some g in G and some automorphism $f: G \rightarrow G$.

When $G = C(p) \times C(q)$, $D + (r, s)$ corresponds to the $(-r, -s)$ -translate of the array corresponding to D . For $(p, q) = 1$, any automorphism of $C(p) \times C(q)$ sends (x, y) to (rx, sy) where $(r, p) = (s, q) = 1$, and its effect on the corresponding array amounts to forming the new array whose (i, j) th entry is $a(r'i, s'j)$ with $rr' \equiv 1 \pmod{p}$ and $ss' \equiv 1 \pmod{q}$. For $(p, q) \neq 1$, the effect of an automorphism of $C(p) \times C(q)$ on the corresponding array may be more complicated.

EXAMPLE II. McFarland (1973) constructs $\frac{1}{2}(q^s + 1)$ inequivalent difference sets in the group $E \times K$ where E is an elementary abelian group of order q^{s+1} and K is any group of order $[(q^{s+1} - 1)/(q - 1)] + 1$, where q is an odd prime power. By taking q to be an odd prime, $s = 1$ and K to be $C(q + 2)$, and observing that $C(q) \times C(q) \times C(q + 2)$ is isomorphic to $C(q) \times C(q(q + 2))$, we obtain $\frac{1}{2}(q + 1)$ inequivalent difference sets in $C(q) \times C(q(q + 2))$ and hence $\frac{1}{2}(q + 1) q \times q(q + 2)$ -arrays with two-level autocorrelation with c (in (1)) equal to $q^2(q - 2)$.

The explicit construction is as follows. It is obtained by tracing the isomorphism $f: C(q) \times C(q) \times C(q + 2) \rightarrow C(q) \times C(q(q + 2))$ given by $f(x, y, z) = (x, y(q + 2) + zq)$. For a given $t = -1, 0, 1, 2, \dots, \frac{1}{2}(q - 3)$, we put $A_t = (a_t(i, j))$ where

$$a_t(i, j) = \begin{cases} 0 & \text{if } (i, j) = (sl + 1, (q + 2)s + ql) \\ & \text{or } (sk, (q + 2)s + qk) \\ & \text{or } (s, q^2) \\ & \text{with } l = 0, 1, \dots, t, k = t + 1, \dots, q - 1, \\ & \quad s = 0, 1, \dots, q - 1 \\ 1 & \text{if otherwise.} \end{cases}$$

(For $t = -1$, it is to be understood that l does not appear in the description of (i, j) and k ranges from 0 to $q - 1$). Altogether this gives $\frac{1}{2}(q + 1)$ arrays with two-level autocorrelation.

When q is an odd prime with $q + 2$ a perfect square, say n^2 (e.g. $(q, n) = (7, 3), (23, 5), (47, 7), \dots$), then by choosing K to be $C(n) \times C(n)$ we obtain in a similar manner $\frac{1}{2}(q + 1) qn \times qn$ -arrays with two level autocorrelation with $c = q^2(q - 2)$.

3. It is desirable to obtain perfect $p \times q$ -arrays, which corresponds to $(4N^2, 2N^2 - N, N^2 - N)$ -difference sets in $C(p) \times C(q)$ by Proposition 1. Again we restrict our attention to arrays with more 1's than 0's.

For $(p, q) = 1$, the group $C(p) \times C(q)$ is isomorphic to $C(pq)$. It is conjectured that no cyclic $(4N^2, 2N^2 - N, N^2 - N)$ -difference set exists except for $N = 1$. Turyn (1968) has verified the conjecture for $N < 55$. Thus it seems that $(1 \ 1 \ 1 \ 0)$ is the only possible perfect $p \times q$ -array with $(p, q) = 1$.

For $(p, q) \neq 1$, Calabro and Wolf (1968) give the two perfect arrays

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \quad \text{and} \quad \begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

which are obtained from $(1 \ 1 \ 1 \ 0)$. We shall now produce two more families of perfect arrays.

EXAMPLE III. Spence (1977) constructs three inequivalent $(36, 15, 6)$ -difference sets in $C(6) \times C(6)$ and four inequivalent $(36, 15, 6)$ -difference sets in $C(3) \times C(12)$. Thus we obtain three corresponding perfect 6×6 -arrays and four corresponding perfect 3×12 -arrays. An example of each type is given below.

$$\begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}, \quad \begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \end{bmatrix}.$$

EXAMPLE IV. Let D_1 be a $(36, 15, 6)$ -difference set in $C(3) \times C(12)$, and let D_2 be the trivial $(4, 1, 0)$ -difference set in $C(4)$. The subset $D = (D_1, \bar{D}_2) \cup (\bar{D}_1, D_2)$, where \bar{D}_1 and \bar{D}_2 denote the complementary difference sets of D_1 and D_2 respectively, is a $(144, 66, 30)$ -difference set in $C(3) \times C(12) \times C(4)$ by Theorem 2 of Menon (1962). As $C(3) \times C(12) \times C(4)$ is isomorphic to $C(12) \times C(12)$, we obtain four perfect 12×12 -arrays from the four inequivalent difference sets in $C(3) \times C(12)$ in Spence (1977) by this method. Again we just give one example below.

$$\begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \end{bmatrix}.$$

If we take both D_1 and D_2 to be the trivial $(4, 1, 0)$ -difference set in $C(4)$, we obtain via this method the perfect 4×4 -array given by Calabro and Wolf (1968) mentioned at the beginning of this section.

4. We summarize the preceding results in Table I (for $v \leq 200$ and $p \leq q$). Most of the entries come from Table 6.1 of Baumert (1971). We shall omit all $1 \times v$ -arrays (which are simply sequences of period v) and all trivial cases (which correspond to trivial difference sets with $k - \lambda = 0$ or 1). The last column in Table 1 refers to the type of example described in this note.

TABLE I

v	p	q	c	Type
4	2	2	0	Calabro and Wolf (1968)
15	3	5	-1	I
16	4	4	0	IV
21	3	7	5	I
35	5	7	-1	I
36	3	12	0	III
36	6	6	0	III
40	5	8	4	I
45	3	15	9	II
57	3	19	29	I
63	7	9	-1	I
85	5	17	21	I
91	7	13	55	I
133	7	19	89	I
133	7	19	33	I
143	11	13	-1	I
144	12	12	0	IV
156	3	52	56	I
156	4	39	56	I
156	12	13	56	I
175	5	35	75	II
183	3	61	131	I

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