## NOTE

# ON THE NONUNIFORM FISHER INEQUALITY 

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Received 18 July 1985
Revised 30 July 1986


#### Abstract

Let $\mathscr{F}$ be a family of $m$ subsets (lines) of a set of $n$ elements (points). Suppose that each pair of lines has $\lambda$ points in common for some positive $\lambda$. The Nonuniform Fisher Inequality asserts that under these circumstances $m \leqslant n$. We examine the case when $m=n$. We give a short proof of the fact that (with the exception of a trivial case) such an $\mathscr{F}$ must behave like a geometry in the following sense: a line must pass through each pair of points. This generalizes a result of de Bruijn and Erdös.


Dedicated to the memory of H.J. Ryser
A hypergraph is a pair $\mathscr{H}=(X, \mathscr{F})$ where $X=\left\{x_{1}, \ldots, x_{n}\right\}$ is a finite set of points and $\mathscr{F}=\left\{F_{1}, \ldots, F_{m}\right\}$ is a family of subsets of $X$ called lines. We shall always use $n$ and $m$ to denote the numbers of points and lines, resp. We suggest Lovász' excellent book [7] as a general reference on hypergraphs.

Generalizing R.A. Fisher's celebrated inequality [4], Majumdar [8] and later Isbell [6] found the following result (cf. Lovász [7, Problem 13.15(b)]).

Theorem 1 (Nonuniform Fisher Inequality). Let $\mathscr{H}$ be a hypergraph with $n$ points and $m$ lines. If every pair of lines has precisely $\lambda$ points in common for some positive integer $\lambda$, then $m \leqslant n$.

Fisher's original result is equivalent to this theorem for the case of dual BIBD's. (The dual of a hypergraph is obtained by interchanging the roles of points and blocks, preserving the incidence relation.) R.C. Bose found a very elegant proof of Fisher's inquality [1] which actually works for any uniform hypergraph. ( $\mathscr{H}$ is uniform if $\left|F_{1}\right|=\cdots=\left|F_{m}\right|$ ). Bose's proof uses elementary linear algebra in a surprising way. Only a slight modification of his proof is needed in order to eliminate the uniformity conditions and thus obtain Theorem 1. No proof, avoiding this linear algebra trick, of Theorem 1 appears to be known. (Fisher's original proof [4] uses direct counting arguments and does not seem to generalize beyond the BIBD case.) We note, however, that an infinite version of Theorem 1 exists (Komjáth [15]).

The Nonuniform Fisher Inequality contains a result of de Bruijn and Erdös [3] as the particular case $\lambda=1$. For this case, a "purely combinatorial" proof exists ([3], cf. [7, Problems 13.14 and 13.15(a)]). It has been quoted as an advantage of the proof avoiding linear algebra that it permits characterization of the extremal cases $(m=n)$. As shown by de Bruijn and Erdös, in the case $\lambda=1$, equality will hold in the following cases only.
(a) One of the lines is a singleton.
(b) $\mathscr{H}$ is a possibly degenerate projective plane.

In case (a), all the other lines must have 2 points.
A possibly degenerate projective plane is a hypergraph satisfying the following axioms:
(i) Through every pair of points there is a line;
(ii) Every pair of lines intersects in precisely one point;
(iii) There exists a triangle (three points not on a line).

We obtain the axioms of a projective plane if we strengthen (iii) to
(iv) There exists a quadrilateral (four points, no three of which are on a line). A degenerate plane is one that satisfies (i), (ii), (iii) but not (iv). It is easy to show that there exists precisely one degenerate plane for every $n \geqslant 3$; it has one line with $n-1$ points and $n-1$ lines with two points each, the latter connecting the points of the long line to the point it misses.

So the crucial step in characterizing the extremal case ( $m=n$ ) in the de Bruijn-Erdös theorem is proving that unless there is a singleton among the lines, the hypergraph must satisfy (i). (If there is no singleton there, it is straightforward to verify (iii).)

Although this fact follows from the results of Ryser [11] and Woodall [14], their arguments are certainly much more involved than those of de Bruijn and Erdös. The aim of this note is to demonstrate that it is possible to use the linear algebra method to verify (i) for the extremal hypergraphs in a very simple way. This approach permits a generalization to $\lambda>1$; we shall prove that for arbitrary $\lambda$, if $m=n$ in the Nonuniform Fisher Inequality, then (i) holds.

Theorem 2 (Ryser, Woodall). Let $\mathscr{H}$ be a hypergraph with $n$ points and $m=n$ lines. Assume no line is a singleton. If every pair of lines has precisely $\lambda$ points in common for some positive $\lambda$, then through every pair of points there is a line.

If we add uniformity $\left(\left|F_{1}\right|=\cdots=\left|F_{n}\right| \stackrel{\text { def }}{=} k\right)$ to the conditions of Theorem 2, then, by a result of H.J. Ryser [9], $\mathscr{H}$ must also be regular of degree $k$ (every point must belong to precisely $k$ lines) and there must be precisely $\lambda$ lines through each pair of points. This means the hypergraph is a symmetric design [10, Ch. 8].

The nonuniform extremal configurations are called $\lambda$-designs [11]. Ryser [11] and Woodall [14] prove that the points of a $\lambda$-design have only two different degrees, $r_{1}$ and $r_{2}$, say. (The degree or "replication number" of a point is the number of lines through the point.) It is no longer true that every pair of points is
linked by $\lambda$ lines. An example of de Witte [13] (cf. [11]) shows that there exists a $\lambda$-design with $n=m=7, \lambda=2$ where some pairs of points are linked by 1 , others by 2 , still others by 3 lines.

Only one class of $\lambda$-designs is known ("point-complemented symmetric block designs" for $\lambda \geqslant 2$, cf. [2]). That all $\lambda$-designs are of this type is the $\lambda$-design conjecture. Bridges [2] shows that a $\lambda$-design is of this type if and only if every pair of points of different degrees is linked by precisely $\lambda$ lines. Seress [16] proves the $\lambda$-design conjecture assuming $r_{1} r_{2}=\lambda(n-1)$. (It is known that $r_{1} r_{2} \leqslant \lambda(n-$ 1).)

The proof of Theorem 2 is entirely self-contained and makes no reference to the results quoted above. It incorporates the (known) compact proof of Theorem 1 (see the fourth paragraph of the proof).

Proof of Theorem 2. Let us not assume for a moment that $m=n$. If there is a line with fewer than $\lambda$ points then $m=1$. If some line, say $F_{1}$, has precisely $\lambda$ points, then the sets $F_{i}-F_{1}$ must be pairwise disjoint and therefore $m \leqslant n-\lambda+1$. This is always $\leqslant n$ with equality only if $\lambda=1$ and thus one of the lines is a singleton. Henceforth we assume that the numbers $l_{i} \stackrel{\text { def }}{=}\left|F_{i}\right|-\lambda$ are all positive.

Let $M$ be the indicence matrix of $\mathscr{H} . M$ is an $m \times n$ matrix; its columns correspond to the points and its rows to the lines. The entry $M[i, j]$ is 1 if point $x_{j}$ belongs to line $F_{i}$.

The intersection matrix of $\mathscr{H}$ is the $m \times m$ matrix $A=M M^{\mathrm{T}}$. Our intersection conditions are summarized in the matrix equation

$$
\begin{equation*}
M M^{\mathrm{T}}=\lambda J+L, \tag{1}
\end{equation*}
$$

where $J$ is the $m \times m$ matrix with all entries equal to 1 and $L$ is the diagonal matrix $L=\operatorname{diag}\left(l_{1}, \ldots, l_{m}\right)$.

Since both $J$ and $L$ are positive semidefinite Hermitian matrices and $L$ is positive definite, $A$ is positive definite as well. In particular $A$ is nonsingular, therefore rank $M=m$, proving Theorem $1(m \leqslant n)$.

Let us henceforth assume $m=n$. Notice that $M$ is now a nonsingular $n \times n$ matrix.

Theorem 3. Assume $M$ is a nonnegative square matrix satisfying (1), where $\lambda>0$ and $L$ is a diagonal matrix with positive diagonal entries. Then all entries of $M^{\mathrm{T}} M$ are positive.

Our objective is to prove that all entries of the matrix $M^{\mathrm{T}} M$ are positive.
We begin with three simple observations. All matrices below are $n$ by $n$, with real entries.

Let $B$ and $C$ be two matrices. We write $B>C$ ( $B \geqslant C$, resp.) if all entries of $B-C$ are positive (nonnegative, resp.).

Observation 1. If $B>0, C \geqslant 0$ and no column of $C$ is zero then $B C>0$. If no row of $C$ is zero, then $C B>0$.

Observation 2. Any matrix of the form $B J$ can be written as $B J=D J$, where $D=\operatorname{diag}\left(d_{1}, \ldots, d_{n}\right)$ is a diagonal matrix.
(As before, $J$ is the all-ones matrix.)

Proof. Make $d_{i}$ equal to the sum of the $i$ th row of $B$.

Observation 3. Let $B$ and $C$ be arbitrary $n \times n$ matrices. If the matrices JC and $\dot{B J C}$ are positive, then so is BJ.

Proof. By Observation 2, $B J=D J$ for some diagonal matrix $D=$ $\operatorname{diag}\left(d_{1}, \ldots, d_{n}\right)$. We have to prove that all the $d_{i}$ are positive. Let $r$ denote a row of $J C$. Then the $i$ th row of BJC= DJC is $d_{i} r$. Since both $J C$ and BJC are positive, we infer $d_{i}>0$.

We proceed with the proof of Theorem 2.
By (1) we have

$$
\begin{equation*}
M^{\mathrm{T}}=M^{-1}(\lambda J+L) \tag{2}
\end{equation*}
$$

Consequently,

$$
\begin{equation*}
M^{\mathrm{T}} L^{-1}=M^{-1}\left(\lambda J L^{-1}+I\right) \tag{3}
\end{equation*}
$$

where $I$ is the identity matrix. Multiplying by $J$ from the right we obtain

$$
\begin{equation*}
M^{\mathrm{T}} L^{-1} J=M^{-1} J\left(\lambda L^{-1} J+I\right) \tag{4}
\end{equation*}
$$

The left hand side is positive by Observation 1 since $M^{\mathrm{T}} L^{-1}$ is nonnegative and non-singular. Setting $B=M^{-1}$ and $C=\lambda L^{-1} J+I$, we notice that $C$ and therefore $J C$ are positive, and $B J C$ is positive because the left hand side of (4) is. Using Observation 3 we conclude that $B J>0$, i.e.,

$$
\begin{equation*}
M^{-1} J>0 \tag{5}
\end{equation*}
$$

Multiplying both sides of (3) by $M$ from the right, we obtain

$$
\begin{equation*}
M^{\mathrm{T}} L^{-1} M=\lambda M^{-1} J L^{-1} M+I \tag{6}
\end{equation*}
$$

Here, $\lambda M^{-1} J>0$ by (5) and $L^{-1} M$ is nonnegative and nonsingular. Consequently (by Observation 1) the right hand side is positive. We conclude that

$$
M^{\mathrm{T}} M=M^{\mathrm{T}} I M \geqslant M^{\mathrm{T}} L^{-1} M>0,
$$

completing the proof of Theorem 2.

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