# Classification of some countable descendant-homogeneous digraphs ${ }^{\text {* }}$ 

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#### Abstract

For finite $q$, we classify the countable, descendant-homogeneous digraphs in which the descendant set of any vertex is a $q$-valent tree. We also give conditions on a rooted digraph $\Gamma$ which allows us to construct a countable descendant-homogeneous digraph in which the descendant set of any vertex is isomorphic to $\Gamma$.


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## 1. Introduction

### 1.1. Background

A countable digraph is homogeneous if any isomorphism between finite (induced) subdigraphs extends to an automorphism. The digraphs with this property are classified by Cherlin in [7]. By analogy, the notion of descendanthomogeneity was introduced in [4]. A countable digraph is descendant-homogeneous if any isomorphism between finitely generated subdigraphs extends to an automorphism. Here, a subdigraph is finitely generated if its vertex set can be written as the descendant set of a finite set of vertices, that is, the set of vertices which are reachable by a directed path from the set.

Note that descendant-homogeneity can hold for trivial reasons: digraphs where the descendant set of any vertex is the whole digraph, or where no two vertices have isomorphic descendant sets are descendant-homogeneous. So it is reasonable to impose restrictions such as vertex transitivity and no directed cycles. We refer to [4] for further discussion.

In this paper we are particularly interested in vertex-transitive, descendant-homogeneous digraphs: so in this case, the descendant set of any vertex is isomorphic to some fixed digraph $\Gamma$. Examples of countable, vertex-transitive, descendanthomogeneous digraphs where $\Gamma$ is a $q$-valent directed tree (for finite $q>1$ ) were given in [9,4]. The main result of this paper is to show that the digraphs constructed in $[9,4]$ constitute a complete list of all the countable descendant-homogeneous digraphs with descendant sets of this form (Theorem 2.2). In the final section of the paper, we give general conditions on $\Gamma$ under which there is a countable, vertex-transitive, descendant-homogeneous digraph in which the descendant set of any vertex is isomorphic to $\Gamma$. In particular, these conditions are satisfied by certain 'tree-like' digraphs $\Gamma$ studied in [1]. This gives new examples of descendant-homogeneous digraphs (and indeed, highly arc-transitive digraphs).

The first (non-trivial) examples of descendant-homogeneous digraphs known to the authors arose in the context of highly arc-transitive digraphs (those whose automorphism groups are transitive on the set of $s$-arcs for all $s$ ). In answer

[^0]to a question of Cameron, Praeger, and Wormald in [6], the paper [9] gave a construction of a certain highly arc-transitive digraph $D$ having an infinite binary tree as descendant set. The digraph was constructed as an example of a highly arctransitive digraph not having the 'property $Z$ ', meaning that there is no homomorphism from $D$ onto the natural digraph on $\mathbb{Z}$ (the doubly infinite path). However, we noted in [4] that it is also descendant-homogeneous. A more systematic analysis of this notion was carried out in [4], and further examples were given. The method of [9] immediately applies to $q$-valent trees for any finite $q>1$ in place of binary trees, but in addition, it is shown in [4] that it is possible to omit certain configurations and still carry out a Fraïssé-type construction to give other examples of descendant-homogeneous digraphs whose descendant sets are $q$-valent trees. The classical Fraïssé theorem for relational structures provides a link between countable homogeneous structures (those in which any isomorphism between finite substructures extends to an automorphism) and amalgamation classes of finite structures. See [5,7,10] for instance. The analogue of Fraïssé's Theorem and the appropriate notion of amalgamation classes which applies to descendant-homogeneity is given in Section 2.1.

### 1.2. Notation and terminology

Let $D$ a digraph with vertex and edge sets $V D$ and $E D$, and let $u \in V D$. For ease, in this paper we require our digraphs to be asymmetric; that is, we never have a directed edge both from $u$ to $v$, and from $v$ to $u$, since this is true in the cases we consider anyway. For $s \geq 0$, an $s$-arc in $D$ from $u_{0}$ to $u_{s}$ is a sequence $u_{0} u_{1} \ldots u_{s}$ of $s+1$ vertices such that $\left(u_{i}, u_{i+1}\right) \in E D$ for $0 \leq i<s$. We let

$$
\operatorname{desc}^{s}(u):=\{v \in V D \mid \text { there is an } s-\operatorname{arc} \text { from } u \text { to } v\}
$$

and $\operatorname{desc}(u)=\bigcup_{s>0} \operatorname{desc}^{s}(u)$, the descendant set of $u$ (we also denote this by $\operatorname{desc}_{D}(u)$ if we need to emphasize that we are looking at descendants in $D$ ). If $X \subseteq V D$, we also let

$$
\operatorname{desc}^{s}(X):=\bigcup_{x \in X} \operatorname{desc}^{s}(x)
$$

and similarly $\operatorname{desc}(X):=\bigcup_{x \in X} \operatorname{desc}(x)$. The 'ball' of radius $s$ at $u$ is given by

$$
B^{s}(u):=\bigcup_{0 \leq i \leq s} \operatorname{desc}^{i}(u)
$$

For a digraph $D$ we often write $D$ in place of $V D$ and use the same notation for a subset of the vertices and the full induced subdigraph. Henceforth, 'subdigraph' will mean 'full induced subdigraph' and an embedding of one digraph into another will always mean a full induced subdigraph.

We say that $A \subseteq D$ is descendant-closed in $D$, written $A \leq D$ if $\operatorname{desc}_{D}(a) \subseteq A$ for all $a \in A$; and we say that an embedding $f: A \rightarrow B$ between digraphs is a $\leq$-embedding if $f(A) \leq B$. When $A, \bar{B}_{1}, B_{2}$ are digraphs we say that $\leq$-embeddings $f_{i}: A \rightarrow B_{i}$ are isomorphic if there is an isomorphism $h: B_{1} \rightarrow B_{2}$ with $f_{2}=h \circ f_{1}$.

We say that $A \leq D$ is finitely generated if there is a finite subset $X \subseteq A$ with $A=\operatorname{desc}_{D}(X)$; in this case we refer to $X$ as a generating set of $A$. If additionally no proper subset of $X$ is a generating set, then $X$ is called a minimal generating set. Clearly, in this case, no element in $X$ is a descendant of any other element of $X$.

The digraph $D$ is descendant-homogeneous if whenever $f: A_{1} \rightarrow A_{2}$ is an isomorphism between finitely generated descendant-closed subdigraphs of $D$, there is an automorphism of $D$ which extends $f$. The group of automorphisms of $D$ is denoted by $\operatorname{Aut}(D)$.

We shall mainly be concerned with digraphs $D$ where the descendant sets of single vertices are all isomorphic to a fixed digraph $\Gamma$ : in this case we refer to $\Gamma$ as 'the descendant set' of $D$. A subset of a digraph is independent if the descendant sets of any two of its members are disjoint. In any digraph in which the descendant sets are all isomorphic, for any two finite independent subsets $X$ and $Y$, any bijection from $X$ to $Y$ extends to an isomorphism from $\operatorname{desc}(X)$ to desc $(Y) \operatorname{since} \operatorname{desc}(X)$ and $\operatorname{desc}(Y)$ are both the disjoint union of $|X|$ descendant sets.

Throughout we fix an integer $q>1$ and write $T=T_{q}$ for the $q$-valent rooted tree. So $T$ has as its vertices the set of finite sequences from the set $\{0, \ldots, q-1\}$ and directed edges ( $\bar{w}, \bar{w} i$ ) (for $\bar{w}$ a finite sequence and $i \in\{0, \ldots, q-1\}$ ).

## 2. Amalgamation classes

### 2.1. The Fraïssé theorem

As in [4], the correct context for the study of descendant-homogeneous digraphs is a suitable adaptation of Fraïssé's notion of amalgamation classes. The reader who is familiar with this type of result (or with [4]) and who is mainly interested in the main classification result, Theorem 2.2, could reasonably skip to the next subsection. The extra generality which is given here is only needed in the final section of the paper.

Let $\mathscr{D}$ be a class of (isomorphism types of) digraphs. Then $\mathscr{D}$ has the $\leq$-amalgamation property if the following holds: if $A, B_{1}$ and $B_{2}$ lie in $\mathscr{D}$, and $\leq$-embeddings $f_{1}$ and $f_{2}$ of $A$ into each of $B_{1}$ and $B_{2}$ are given, then there are a structure $C \in \mathscr{D}$ and $\leq$-embeddings $g_{1}$ and $g_{2}$ of $B_{1}$ and $B_{2}$ respectively into $C$ such that $g_{1} \circ f_{1}=g_{2} \circ f_{2}$. We say that $g_{1}, g_{2}$ solve the amalgamation problem given by $f_{1}, f_{2}$.

Remark 1. Suppose $A, B_{1}$, and $B_{2}$ are digraphs and $\leq$-embeddings $f_{1}$ and $f_{2}$ of $A$ into each of $B_{1}$ and $B_{2}$ are given. We can clearly find a solution $g_{i}: B_{i} \rightarrow C$ with the property that $C=g_{1}\left(B_{1}\right) \cup g_{2}\left(B_{2}\right), g_{1}\left(B_{1}\right) \cap g_{2}\left(B_{2}\right)=g_{1}\left(f_{1}(A)\right)$ and every directed edge is contained in $g_{1}\left(B_{1}\right)$ or $g_{2}\left(B_{2}\right)$. Moreover, this solution is uniquely determined up to isomorphism by the $f_{i}$. Informally, we can regard the $f_{i}$ as inclusion maps and take $C$ to be the disjoint union of $B_{1}$ and $B_{2}$ over $A$. We make this into a digraph by taking as edge set $E C=E B_{1} \cup E B_{2}$. It is easy to see that $B_{1}, B_{2} \leq C$ and the inclusion maps $g_{i}: B_{i} \rightarrow C$ satisfy $g_{1} \circ f_{1}=g_{2} \circ f_{2}$. We say that the solution $g_{i}: B_{i} \rightarrow C$ to the problem $f_{i}: A \rightarrow B_{i}$ is the free amalgam of the $f_{i}$. When $f_{1}, f_{2}$ are inclusion maps (or are understood from the context) we shall abuse this terminology and say that $C$ is the free amalgam of $B_{1}$ and $B_{2}$ over $A$.

Note that if $B_{1}, B_{2} \leq C$ then $B_{1} \cup B_{2} \leq C$ and $B_{1} \cup B_{2}$ is the free amalgam of $B_{1}$ and $B_{2}$ over $B_{1} \cap B_{2}$ : there can be no directed edges between elements of $B_{1} \backslash B_{2}$ and $B_{2} \backslash B_{1}$ as $B_{1}, B_{2}$ are descendant-closed.

When we come to count structures and embeddings up to isomorphism (as in Lemma 4.2), it will be useful to have a more precise notation for free amalgamation. Suppose in the above that $f_{1}$ is inclusion and $f_{2}$ is an arbitrary $\leq$-embedding $f_{2}: A \rightarrow B_{2}$. The free amalgam $B_{1} *_{f_{2}} B_{2}$ has as vertex set the disjoint union of $B_{1} \backslash A$ and $B_{2}$ (and the 'obvious' directed edges). The embedding $g_{2}: B_{2} \rightarrow B_{1} *_{f_{2}} B_{2}$ is inclusion and the embedding $g_{1}: B_{1} \rightarrow B_{1} *_{f_{2}} B_{2}$ is given by $g_{1}(b)=b$ if $b \in B_{1} \backslash A$ and $g_{1}(b)=f_{2}(b)$ if $b \in A$.

We remark that in general, if $A \leq B_{1}$ and $f_{2}, f_{2}^{\prime}: A \rightarrow B_{2}$ are $\leq$-embeddings with the same image, then $B_{1} *_{f_{2}} B_{2}$ and $B_{1} *_{f_{2}^{\prime}} B_{2}$ need not be isomorphic.

The analogue of Fraïssé's Theorem which we use is the following. Recall that the notion of isomorphism of two embeddings was defined in Section 1.2.

Theorem 2.1. Suppose $M$ is a countable descendant-homogeneous digraph. Let $\mathcal{C}$ be the class of digraphs which are isomorphic to finitely generated $\leq$-subdigraphs of $M$. Then
(1) $\mathcal{C}$ is a class of countable, finitely generated digraphs which is closed under isomorphism and has countably many isomorphism types;
(2) $\mathcal{C}$ is closed under taking finitely generated $\leq$-subdigraphs;
(3) $\mathcal{C}$ has the $\leq$-amalgamation property;
(4) for all $A, B \in \mathcal{C}$ there are only countably many isomorphism types of $\leq$-embeddings from $A$ to $B$.

Conversely, if $\mathcal{C}$ is a class of digraphs satisfying (1)-(4), then there is a countable descendant-homogeneous digraph $M$ for which the class of digraphs isomorphic to finitely generated $\leq$-subdigraphs of $M$ is equal to $\mathcal{C}$. Moreover, $M$ is determined up to isomorphism by C .

We refer to a class $\mathcal{C}$ of digraphs satisfying (1)-(4) as a $\leq$-amalgamation class. The digraph $M$ determined by $\mathcal{C}$ as in the theorem is called the Fraïssé limit of $(\mathcal{C}, \leq)$.

Remark 2. It is easy to see that in place of (4) we can substitute the condition:
(4') if $A \leq B \in \mathcal{C}$ and $A$ is finitely generated, then the subgroup of the automorphism group $\operatorname{Aut}(A)$ consisting of automorphisms which extend to automorphisms of $B$ is of countable index in $\operatorname{Aut}(A)$.

Indeed, we wish to consider the number of $\leq$-embeddings $f: A \rightarrow B$ up to isomorphism. As $B$ is countable and $A$ is finitely generated there are countably many possibilities for the image $f(A)$, so it will be enough to count isomorphism types of $\leq-$ embeddings with fixed finitely generated image $Y \leq B$. Let $H$ be the subgroup of Aut $(Y)$ consisting of automorphisms which extend to automorphisms of $B$. It is straightforward to show that if $f, f^{\prime}: A \rightarrow B$ have image $Y$, then $f, f^{\prime}$ are isomorphic if and only if the map $g \in \operatorname{Aut}(Y)$ given by $g(y)=f^{\prime}\left(f^{-1}(y)\right)$ is in $H$. Thus there is a bijection between the $H$-cosets in Aut $(Y)$ and the isomorphism types.

Remark 3. The proof of Theorem 2.1 is reasonably standard, but we make some comments on the condition (4). First, suppose $M$ and $\mathcal{C}$ are as in the statement. We show that ( $4^{\prime}$ ) in Remark 2 holds. So let $A \leq B \in \mathcal{C}$ and $H \leq$ Aut $(A)$ be the automorphisms of $A$ which extend to automorphisms of $B$, as in ( $4^{\prime}$ ). We may assume $B \leq M$. Suppose $g_{1}, g_{2} \in \operatorname{Aut}(A)$ lie in different $H$-cosets. As $M$ is $\leq$-homogeneous we can extend $g_{i}$ to $k_{i} \in \operatorname{Aut}(M)$. Then $k_{1}(B) \neq k_{2}(B)$. Otherwise $h=k_{2}^{-1} k_{1}$ stabilizes $B$ and gives an automorphism of $B$ which extends $h=g_{2}^{-1} g_{1}$; this implies $h \in H$ and $g_{2} H=g_{1} H$, which is a contradiction. As there are only countably many possibilities for the image of $B$ under automorphisms of $M$, it follows that $H$ is of countable index in $\operatorname{Aut}(A)$, as required.

The converse is a fairly standard construction, and can be read off from Theorem 2.18 of [11], which in turn is adapted from Theorem 1.1 of [8]. However, we give a few details of the proof. So suppose we have a class $\mathcal{C}$ of finitely generated digraphs satisfying (1)-(4). We construct a countable chain $C_{1} \leq C_{2} \leq C_{3} \leq \ldots$ of digraphs in $\mathcal{C}$ with the property that if $A \leq C_{i}$ is finitely generated and $f: A \rightarrow B \in \mathcal{C}$ is a $\leq$-embedding, then there is $j \geq i$ and a $\leq$-embedding $g: B \rightarrow C_{j}$ with $g(f(a))=a$ for all $a \in A$. The resulting digraph $\bigcup_{i} C_{i}$ will be descendant-homogeneous, by a back-and-forth argument. Note that by (4), we have only countably many $f$ to consider (for any particular $A$ ). For if $f, g$ are as above and $f^{\prime}: A \rightarrow B$ is isomorphic to $f$ with $f^{\prime}=h \circ f$ for $h \in \operatorname{Aut}(B)$, then $g^{\prime}=g \circ h^{-1}: B \rightarrow C_{j}$ satisfies $g^{\prime}\left(f^{\prime}(a)\right)=a$ for all $a \in A$.

### 2.2. The classification result

Recall that $q \geq 2$ is an integer and $T=T_{q}$ is the $q$-valent rooted tree. We shall classify countable, descendanthomogeneous digraphs $M$ in which the descendant sets of vertices are isomorphic to $T$. Thus, by Theorem 2.1, we need to classify $\leq$-amalgamation classes of finitely generated digraphs with descendant sets isomorphic to $T$. In this case, we can replace condition (4) in Theorem 2.1 by the simpler condition:
(4") if $a_{1}, a_{2} \in B \in \mathcal{C}$, then $\operatorname{desc}_{B}\left(a_{1}\right) \cap \operatorname{desc}_{B}\left(a_{2}\right)$ is finitely generated
as in Theorem 3.4 of [4]. Indeed, if $\mathcal{C}$ satisfies $\left(4^{\prime \prime}\right)$ then $\left(4^{\prime}\right)$ is a special case of Lemma 4.2 here. Conversely, if ( $4^{\prime}$ ) holds, then to see $\left(4^{\prime \prime}\right)$ let $B=\operatorname{desc}\left(a_{1}\right) \cap \operatorname{desc}\left(a_{2}\right)$ and $A=\operatorname{desc}\left(a_{1}\right)$. Let $X$ be the minimal generating set for $B$, which just consists of its elements which are not (proper) descendants of any other members of $B$. Then $X$ is independent and any automorphism of $B$ which stabilizes $A$ must fix $X$ setwise. On the other hand, if $Z$ is an infinite independent subset of $A$ it is easy to see that the stabilizer of $Z$ in $\operatorname{Aut}(A)$ is of index continuum (as there are continuum many translates of $Z$ by automorphisms of $A$, since $A$ is a regular rooted tree).

Thus we work with the class $\mathcal{C}=\mathcal{C}_{\infty}$ consisting of all digraphs $A$ satisfying the following conditions:

- for all $a \in A, \operatorname{desc}(a)$ is isomorphic to $T$;
- $A$ is finitely generated;
- for $a, b \in A$, the intersection $\operatorname{desc}(a) \cap \operatorname{desc}(b)$ is finitely generated.

Then $\mathcal{C}$ satisfies conditions (1), (2), (4) in Theorem 2.1 (cf. the above remarks and Lemma 4.2), and we are interested in the subclasses of $\mathcal{C}$ which satisfy (3). It is easy to see that $\mathcal{C}$ satisfies (3): in fact $\mathcal{C}$ is closed under free amalgamation. It follows that $(\mathcal{C}, \leq)$ is a $\leq$-amalgamation class. The Fraïssé limit $D_{\infty}$ of $(\mathcal{C}, \leq)$ is the countable descendant-homogeneous digraph constructed in [9].

For $n \geq 2$, we now define the amalgamation classes $\mathcal{C}_{n} \subseteq \mathcal{C}$ (from [4]). Let $\mathcal{T}_{n}$ be the element of $\mathcal{C}$ generated by $n$ elements $\bar{x}_{1}, \ldots, x_{n}$, such that $\operatorname{desc}^{1}\left(x_{i}\right)=\operatorname{desc}^{1}\left(x_{j}\right)$ for all $\bar{i} \neq j$. So $\mathcal{T}_{n}$ is like the tree $T$, except that there are $n$ root vertices (all having the same out-vertices). Let $\mathcal{C}_{n}$ consist of the digraphs $A \in \mathcal{C}$ such that $\mathcal{T}_{n}$ does not embed in $A$ (as a descendant-closed subdigraph). We remark that one can analogously define $\mathcal{T}_{\infty}$, but this does not lie in $\mathcal{C}$ since it is not finitely generated.

It is clear that $\mathcal{C}_{n} \subseteq \mathcal{C}_{n+1}$ and $\mathcal{C}_{n} \subseteq \mathcal{C}$ for all $n$. In [4] it is shown that $\left(\mathcal{C}_{n}, \leq\right)$ is a $\leq$-amalgamation class, though it is clearly not a free amalgamation class. In particular, when we 'solve' an amalgamation problem $f_{i}: A \rightarrow B_{i}$ by maps $g_{i}: B_{i} \rightarrow C$, we may have $g_{1}\left(B_{1}\right) \cap g_{2}\left(B_{2}\right) \supset g_{1}\left(f_{1}(A)\right)$. Informally, this means that points of $B_{1}, B_{2}$ outside $A$ may need to become identified in the amalgam $C$.

For $n \geq 2$, let $D_{n}$ be the Fraïssé limit of $\left(\bigodot_{n}, \leq\right)$, as in Theorem 2.1. Then $D_{n}$ is a countable descendant-homogeneous digraph. Our main result is:

Theorem 2.2. Let $D$ be a countable descendant-homogeneous digraph whose descendant sets are isomorphic to T. Then $D$ is isomorphic to $D_{n}$ for some $n \in\{2, \ldots, \infty\}$.

## 3. Proof of the main theorem

We know from [4] that each $\mathcal{C}_{n} \subseteq \mathcal{C}$ is a $\leq$-amalgamation class. From now on we shall consider an arbitrary subclass $\mathfrak{D}$ of $\mathcal{C}$ which is itself a $\leq$-amalgamation class (that is, satisfies (1)-(4) of Theorem 2.1), with the goal of showing that $\mathcal{C}$ and $\mathcal{C}_{n}$ are the only possibilities for $\mathcal{D}$.

To understand the argument better, suppose that there is some integer $n \geq 2$ such that $\mathcal{T}_{n} \notin \mathscr{D}$. Choose $n$ as small as possible: so in particular, $\mathcal{T}_{n-1} \in \mathscr{D}$ (where $\mathcal{T}_{1}=T$ ) and $\mathscr{D} \subseteq \mathcal{C}_{n}$. To prove our main result it suffices to show that if $A \in \mathcal{C}_{n}$ then $A \in \mathscr{D}$, and this is done by induction on the number of generators of $A$. Let $\left\{a_{1}, \ldots, a_{k}\right\}$ be the minimal generating set of $A$ and let $A_{1}$ be the descendant-closed subdigraph of $A$ with generating set $\left\{a_{1}, \ldots, a_{k-1}\right\}$. Let $A_{0}=A_{1} \cap \operatorname{desc}\left(a_{k}\right)$. Then $A$ is the free amalgam of $A_{1}$ and $\operatorname{desc}\left(a_{k}\right)$ over $A_{0}$. By the induction hypothesis, $A_{1} \in \mathcal{D}$, and we know that $\operatorname{desc}\left(a_{k}\right) \cong T \in \mathscr{D}$. So there are $C \in \mathcal{D}$ and $\leq$-embeddings $f: A_{1} \rightarrow C$ and $g: T \rightarrow C$ such that $f(a)=g(a)$ for all $a \in A_{0}$ (identifying desc $\left(a_{k}\right)$ with $T$ ). However, a priori one cannot force $C$ to be the free amalgam. So we replace $A_{1}$ by some $B \geq A_{1}, T$ by $T^{\prime} \geq T$ and $A_{0}$ by $A_{0}^{\prime} \leq B, T^{\prime}$ in such a way that the amalgam in $\mathscr{D}$ of $B$ and $T^{\prime}$ over $A_{0}^{\prime}$ is forced to be free. This is the point of Lemmas 3.1-3.3 (which do not need the extra assumption on $\mathscr{D}$ ).

Lemma 3.1. Let $A \in \mathcal{C}$ and $X$ be a finite independent subset of $A$. Then there is a finite independent subset $Y$ of $A$ containing $X$ such that $A \backslash \operatorname{desc}(Y)$ is finite.

Proof. Let $a, x \in V A$ and let $S$ be the minimal generating set of $\operatorname{desc}(a) \cap \operatorname{desc}(x)$. Since $S$ is finite, there is $n(a, x) \in \mathbb{N}$ such that $S \subseteq B^{n(a, x)}(a)$. Let $m \geq n(a, x), y \in \operatorname{desc}^{m}(a)$ and $y \notin \operatorname{desc}(x)$. Then $\operatorname{desc}(y) \cap \operatorname{desc}(x)=\varnothing$ : if not, let $u \in \operatorname{desc}(y) \cap \operatorname{desc}(x)$. As $y \in \operatorname{desc}(a), u \in \operatorname{desc}(a) \cap \operatorname{desc}(x)$, so $u \in \operatorname{desc}(s)$ for some $s \in S$. As $\operatorname{desc}(a)$ is a tree, and by the choice of $m, y \in \operatorname{desc}(s) \subseteq \operatorname{desc}(x)$, which is a contradiction.

Let $a_{1}, \ldots, a_{r}$ be the minimal generating set for $A$. Now let $N \geq \max \left\{n\left(a_{i}, a_{j}\right), n\left(a_{i}, x\right) \mid i \neq j, x \in X\right\}$ and let $B=\bigcup_{i=1}^{r} B^{N}\left(a_{i}\right)$. So if $y \in A \backslash(B \cup \operatorname{desc}(X))$ then $\operatorname{desc}(y) \cap \operatorname{desc}(X)=\varnothing$ and if $y_{1}, y_{2} \in A \backslash(B \cup \operatorname{desc}(X))$, and neither is a descendant of the other, then $\operatorname{desc}\left(y_{1}\right) \cap \operatorname{desc}\left(y_{2}\right)=\varnothing$.

To justify this last statement, note that $y_{1} \in B^{n_{1}}\left(a_{i}\right), y_{2} \in B^{n_{2}}\left(a_{j}\right)$ for some $i, j$, and $n_{1}, n_{2}>N \geq n\left(a_{i}, a_{j}\right)$. If $i=j$ then $y_{1}, y_{2} \in \operatorname{desc}\left(a_{i}\right)$ which is a tree, so one is a descendant of the other. Thus $i \neq j$ and we may assume (by the same argument) that $y_{1} \notin \operatorname{desc}\left(a_{j}\right)$. By the first paragraph (with $a=a_{i}, x=a_{j}$ and $y=y_{1}$ ) we have $\operatorname{desc}\left(y_{1}\right) \cap \operatorname{desc}\left(a_{j}\right)=\emptyset$, so $\operatorname{desc}\left(y_{1}\right) \cap \operatorname{desc}\left(y_{2}\right)=\emptyset$.

Let $y_{1}, \ldots, y_{t}$ be the minimal generating set of $A \backslash(B \cup \operatorname{desc}(X))$. Then $Y:=X \cup\left\{y_{1}, \ldots, y_{t}\right\}$ is independent and $A \backslash \operatorname{desc}(Y) \subseteq B$ is finite.

For a finite independent subset $X$ of $A$, and $Y$ given by the lemma, we say that $Y \backslash X$ is a complement of $X$ in $A$.
For $X \subseteq D \in \mathcal{C}$, a common predecessor for $X$ in $D$ is a vertex $a \in D$ such that $(a, x)$ is a directed edge for all $x \in X$. Let $A \in \mathcal{C}$ and let $U, V$ be independent subsets of $A$ and $T$ respectively with $f: \operatorname{desc}(U) \rightarrow \operatorname{desc}(V)$ an isomorphism. Let $Q$ be the set consisting of those $q$-element subsets $p$ of $U$ such that $p$ has a common predecessor in $A$ and $f(p)$ has a common predecessor in $T$. For $p \in Q$, let $w_{p}$ and $w_{f(p)}$ be such common predecessors of $p$ and $f(p)$ respectively. We note that as $T$ is a tree, $w_{f(p)}$ is uniquely determined, but $w_{p}$ may not be. Also, as $T$ is a tree, any two members of $Q$ are disjoint. Now let

$$
U^{\prime}:=(U \backslash \bigcup Q) \cup\left\{w_{p} \mid p \in Q\right\} \quad \text { and } \quad V^{\prime}:=(V \backslash \bigcup f(Q)) \cup\left\{w_{f(p)} \mid p \in Q\right\}
$$

In words, $U^{\prime}$ is obtained from $U$ by replacing the vertices in $p \subseteq U$ by their common predecessors $w_{p}$, for all $p \in Q$. Similarly $V^{\prime}$ is obtained from $V$. Clearly $\left|U^{\prime}\right|=\left|V^{\prime}\right|, \operatorname{desc}(U) \subseteq \operatorname{desc}\left(U^{\prime}\right)$ and $\operatorname{desc}(V) \subseteq \operatorname{desc}\left(V^{\prime}\right)$. Moreover,

Lemma 3.2. (a) The sets $U^{\prime}$ and $V^{\prime}$ are independent subsets of $A$ and $T$ respectively, and the extension $F$ of $f$ which takes $w_{p}$ to $w_{f(p)}$ for each $p \in Q$ is an isomorphism from $\operatorname{desc}\left(U^{\prime}\right)$ to $\operatorname{desc}\left(V^{\prime}\right)$;
(b) if $I \subseteq A$ is disjoint from $U$ and $U \cup I$ is an independent subset of $A$, then $U^{\prime} \cup I$ is also independent.

Proof. (a) Let $u_{1}$ and $u_{2}$ be distinct members of $U^{\prime}$. If neither lies in $\left\{w_{p} \mid p \in Q\right\}$, then they are in $U$, so $\operatorname{desc}\left(u_{1}\right) \cap \operatorname{desc}\left(u_{2}\right)=$ $\varnothing$ is immediate. Next suppose that $u_{1}=w_{p}$ for $p \in Q$ and $u_{2} \notin\left\{w_{p} \mid p \in Q\right\}$. Then $\operatorname{desc}\left(u_{1}\right)=\left\{u_{1}\right\} \cup \operatorname{desc}(p)$, and as $U$ is independent, $\operatorname{desc}(u) \cap \operatorname{desc}\left(u_{2}\right)=\varnothing$ for each $u \in p$, and also $u_{1} \notin \operatorname{desc}\left(u_{2}\right)$, and it follows that $\operatorname{desc}\left(u_{1}\right) \cap \operatorname{desc}\left(u_{2}\right)=\varnothing$. Finally, if $u_{1}=w_{p_{1}}$ and $u_{2}=w_{p_{2}}$, then $\operatorname{desc}\left(u_{1}\right)=\left\{u_{1}\right\} \cup \operatorname{desc}\left(p_{1}\right)$ and $\operatorname{desc}\left(u_{2}\right)=\left\{u_{2}\right\} \cup \operatorname{desc}\left(p_{2}\right)$. Now for each $u \in p_{1}$ and $u^{\prime} \in p_{2}, \operatorname{desc}(u) \cap \operatorname{desc}\left(u^{\prime}\right)=\varnothing$ by the independence of $U$, and $u_{1} \notin \operatorname{desc}\left(p_{2}\right)$ and $u_{2} \notin \operatorname{desc}\left(p_{1}\right)$ are clear, from which it follows that $\operatorname{desc}\left(u_{1}\right) \cap \operatorname{desc}\left(u_{2}\right)=\varnothing$. This shows that $U^{\prime}$ is independent, and the proof that $V^{\prime}$ is independent is similar.

To see that $F$ is an isomorphism, note that the only new points in its domain are $w_{p}$, and $F$ maps $w_{p}$ to $w_{f(p)}$, and $f\left(\operatorname{desc}^{1}\left(w_{p}\right)\right)=\operatorname{desc}^{1}\left(w_{f(p)}\right)$.
(b) Since $\operatorname{desc}\left(w_{p}\right)=\left\{w_{p}\right\} \cup \operatorname{desc}(p)$ for each $p \in Q$, and $p \subseteq U$ and $U \cup I$ is independent, it follows that desc ( $w_{p}$ ) $\cap \operatorname{desc}(x)=\varnothing$ for all $x \in I$, so $U^{\prime} \cup I$ is also independent.

Lemma 3.3. Let $A \in \mathscr{D}$ and let $U$ be a finite independent subset of $A$. Let $M$ be the maximal number of common predecessors in $A$ of $q$-element subsets of $U$, and let $N \geq M$ be such that $\mathcal{T}_{N} \in \mathscr{D}$. Then there is $B \in \mathscr{D}$ with $A \leq B$ and such that every $q$-element subset of $U$ has at least $N$ common predecessors in $B$.
Proof. Let $P=\left\{p_{1}, \ldots, p_{t}\right\}$ be the set of all $q$-element subsets of $U$. (Note that, unlike in the previous proof, the members of $P$ need not be pairwise disjoint.) We construct a sequence $B_{0} \leq B_{1} \leq B_{2} \leq \cdots \leq B_{t}$ in $\mathscr{D}$, such that $p_{i}$ has at least $N$ common predecessors in $B_{l}$ for all $i \leq l$ and $l \leq t$. We start with $B_{0}:=A$ and assume inductively that we have constructed $B_{l}$, where $l<t$. Let $p_{l+1}=\left\{u_{1}, \ldots, u_{q}\right\}$ and consider a copy of $\mathcal{T}_{N}$ with generating set $G=\left\{g_{1}, \ldots, g_{N}\right\}$. Let desc ${ }^{1}(G)=\left\{h_{1}, \ldots, h_{q}\right\}$. Both sets $\bigcup_{j=1}^{q} \operatorname{desc}\left(u_{j}\right)$ and $\bigcup_{j=1}^{q} \operatorname{desc}\left(h_{j}\right)$ are the union of $q$ disjoint copies of $T$, so there is an isomorphism taking the first to the second such that $u_{j}$ is sent to $h_{j}$ for each $j$. Let $B_{l+1}$ be an amalgam in $\mathscr{D}$ of $B_{l}$ and $\mathcal{T}_{N}$ with $\bigcup_{j=1}^{q} \operatorname{desc}\left(u_{j}\right)$ and $\bigcup_{j=1}^{q} \operatorname{desc}\left(h_{j}\right)$ identified by this isomorphism (since $B_{l}, \mathcal{J}_{N} \in \mathscr{D}$ ). We note that $p_{l+1}$ has at least $N$ common predecessors in $B_{l+1}$ since $\left\{h_{1}, \ldots, h_{q}\right\}$ has $N$ common predecessors in $\mathcal{T}_{N}$. Hence $B=B_{t}$ is a member of $\mathscr{D}$ as required.

Proposition 3.4. Let $\mathfrak{D} \subseteq \mathcal{C}$ be $a \leq$-amalgamation class and suppose that $\mathcal{T}_{m} \notin \mathscr{D}$ for some $m \geq 2$. Then $\mathscr{D}=\mathcal{C}_{n}$ where $n$ is the least $m$ such that $\mathcal{T}_{m} \notin \mathscr{D}$.
Proof. Note that $\mathscr{D} \subseteq \mathcal{C}_{n}$ and $\mathcal{T}_{n-1} \in \mathscr{D}$. We shall show that $\mathcal{C}_{n} \subseteq \mathscr{D}$. Let $A \in \mathcal{C}_{n}$. We use induction on the number of generators of $A$ to show that $A \in \mathscr{D}$. Let $a_{1}, \ldots, a_{s}$ be the (distinct) generators of $A$. If $s=1$, or if $A$ is the disjoint union of finitely many copies of $T$, then $A$ embeds in $T$ and therefore $A \in \mathscr{D}$, since $T \in \mathscr{D}$. Now let $s \geq 2$ and suppose that $E \in \mathscr{D}$ for all $E \in \mathcal{C}_{n}$ with at most $s-1$ generators. Let $A_{1}:=\bigcup_{i=1}^{s-1} \operatorname{desc}\left(a_{i}\right)$ and let $T$ be a copy of the $q$-valent tree with $b$ its root. The digraph $A$ is the free amalgam of $A_{1}$ and $\operatorname{desc}\left(a_{s}\right)(\cong T)$ over $A_{1} \cap \operatorname{desc}\left(a_{s}\right)$ (which is finitely generated). So there are independent subsets $U=\left\{u_{1}, \ldots, u_{k}\right\}$ and $V=\left\{v_{1}, \ldots, v_{k}\right\}$ of $A_{1}$ and $T=\operatorname{desc}(b)$ respectively and an isomorphism $f$ from $\operatorname{desc}(U)$ to $\operatorname{desc}(V)$ (taking $u_{i}$ to $v_{i}$ for all $i$ ), such that $A$ is isomorphic to the free amalgam $C$ of $A_{1}$ and $T$ with desc $(U)$ and $\operatorname{desc}(V)$ identified by $f$. See Fig. 1. To prove the result it then suffices to show that there is $D \in \mathscr{D}$ embedding $C$. We shall first 'expand' $A_{1}$ to a digraph $B \in \mathscr{D}$ (using Lemma 3.3) and then amalgamate $B$ with a copy $T^{\prime} \geq T$ of $T$ over the descendant sets of some carefully chosen independent subsets. The resulting digraph is then the required digraph $D$.


Fig. 1. The digraphs $A_{1}$ and $T=\operatorname{desc}(b)$.
By the induction hypothesis, $A_{1} \in \mathscr{D}$ since $A_{1} \leq A \in \mathcal{C}_{n}$ and $A_{1}$ has $s-1$ generators. Let $P$ be the set of all $q$-element subsets of $U$.

Claim 3.5. There is $B \in \mathscr{D}$ containing $A_{1}$ such that every member of $P$ with at most $n-2$ common predecessors in $A_{1}$ has at least one common predecessor in $B$ which does not lie in $A_{1}$.

Proof. Let $M$ be the greatest number of common predecessors in $A_{1}$ of an element of $P$. Note that $M \leq n-1$ since $\mathscr{D} \subseteq \mathcal{C}_{n}$, and recall that $\mathcal{T}_{n-1} \in \mathscr{D}$. Now apply Lemma 3.3 with $N=n-1$ to obtain $B \in \mathscr{D}$ containing $A_{1}$ and such that every $p \in P$ has at least $n-1$ common predecessors in $B$. So for $p \in P$ with at most $n-2$ common predecessors in $A_{1}$, there is at least one common predecessor of $p$ in $B$ which does not lie in $A_{1}$.

Let $T^{\prime} \geq T$ be a copy of $T$ with root $z$ such that $(z, b)$ is a directed edge; let $b^{\prime} \neq b$ be another successor of $z$.
We now find independent subsets $U^{\prime} \cup I$ and $V^{\prime} \cup J$ of $B$ as given by the claim, and $T^{\prime}$ respectively, with $\operatorname{desc}(U) \subseteq \operatorname{desc}\left(U^{\prime}\right)$ and $\operatorname{desc}(V) \subseteq \operatorname{desc}\left(V^{\prime}\right)$, such that $I$ is a complement to $U$ in $A_{1}$ and $J \subseteq \operatorname{desc}\left(b^{\prime}\right)$, together with an isomorphism from $\operatorname{desc}\left(U^{\prime} \cup I\right)$ to $\operatorname{desc}\left(V^{\prime} \cup J\right)$ which takes $I$ to $J$ and extends $f$.

Indeed, if $n=2$ we let $U^{\prime}:=U$ and $V^{\prime}:=V$. Now suppose $n \geq 3$ and let $P^{\prime}$ be the subset of $P$ consisting of all $q$-element sets $p \subseteq U$ with at least one and at most $n-2$ common predecessors in $A_{1}$, and such that the image $f(p)$ in $V$ has a common predecessor in $T^{\prime}$. By Claim 3.5, $p$ has a common predecessor $w_{p}$ in $B \backslash A_{1}$. Let $w_{f(p)}$ be the common predecessor of $f(p)$ in $T^{\prime}$ and define

$$
U^{\prime}:=\left(U \backslash \bigcup P^{\prime}\right) \cup\left\{w_{p} \mid p \in P^{\prime}\right\}
$$

and

$$
V^{\prime}:=\left(V \backslash \bigcup f\left(P^{\prime}\right)\right) \cup\left\{w_{f(p)} \mid p \in P^{\prime}\right\}
$$

By Lemma 3.2, $U^{\prime}$ and $V^{\prime}$ are independent subsets of $B$ and $T^{\prime}$ respectively and the natural extension $F$ of $f$ which takes $w_{p}$ to $w_{f(p)}$ is an isomorphism from $\operatorname{desc}\left(U^{\prime}\right)$ to $\operatorname{desc}\left(V^{\prime}\right)$. In either case ( $n=2$ or $n \geq 3$ ), by Lemma 3.1 , $U$ has a complement $I$ in $A_{1}$ and by Lemma 3.2, $U^{\prime} \cup I$ is an independent set. Now let $J$ be an independent subset of $\operatorname{desc}\left(b^{\prime}\right)$ with $|J|=|I|$. Since $\operatorname{desc}(b) \cap \operatorname{desc}\left(b^{\prime}\right)=\varnothing, V^{\prime} \cup \underline{J}$ is an independent subset of $T^{\prime}$. So $U^{\prime} \cup I$ and $V^{\prime} \cup J$ are independent subsets of the same size and there is an isomorphism $\bar{F}$ from $\operatorname{desc}\left(U^{\prime} \cup I\right)$ to $\operatorname{desc}\left(V^{\prime} \cup J\right)$ extending $f$ and taking $I$ to $J$.

By $\leq$-amalgamation, there are $D \in \mathscr{D}$ and $\leq$-embeddings $g_{1}: B \rightarrow D, g_{2}: T^{\prime} \rightarrow D$ such that $g_{1}(y)=g_{2}(\bar{F}(y))$ for all $y \in \operatorname{desc}\left(U^{\prime} \cup I\right)$, where we may assume that $g_{1}$ is the identity map. As we now show, the point of the construction is that by extending before we amalgamate, we have ensured that in this amalgamation, unwanted identifications are avoided.

Claim 3.6. $A_{1} \cap g_{2}(\operatorname{desc}(b))=\operatorname{desc}(U)$.
Proof. We have $\operatorname{desc}(U)=g_{2}(\operatorname{desc}(V))$ since $\bar{F}_{\mid \operatorname{desc}(U)}=f$.As $\operatorname{desc}(V) \subseteq \operatorname{desc}(b)$, it follows that desc$(U) \subseteq A_{1} \cap g_{2}(\operatorname{desc}(b))$. Now suppose for a contradiction that there are vertices $\gamma \in A_{1} \backslash \operatorname{desc}(U), \gamma^{\prime} \in \operatorname{desc}(b) \backslash \operatorname{desc}(V)$ such that $\gamma=g_{2}\left(\gamma^{\prime}\right)$.

We first show that $\operatorname{desc}(\gamma) \backslash \operatorname{desc}(U)$ is finite. Indeed, suppose $a \in A_{1}$ is such that $\operatorname{desc}(a) \cap \operatorname{desc}(I) \neq \varnothing$. Then $a \neq g_{2}\left(\gamma^{\prime \prime}\right)$ for any $\gamma^{\prime \prime} \in \operatorname{desc}(b) \backslash \operatorname{desc}(V) \operatorname{since} \operatorname{desc}(I)=g_{2}(\operatorname{desc}(J)) \subseteq g_{2}\left(\operatorname{desc}\left(b^{\prime}\right)\right)$ and $\operatorname{desc}\left(b^{\prime}\right) \cap \operatorname{desc}(b)=\varnothing$. So $\operatorname{desc}(\gamma) \cap \operatorname{desc}(I)=\varnothing$, and $\operatorname{desc}(\gamma) \backslash \operatorname{desc}(U)=\operatorname{desc}(\gamma) \backslash \operatorname{desc}(U \cup I)$ is finite.

Now we show that there is a $q$-element subset $p$ of $U \cap \operatorname{desc}(\gamma)$ with a common predecessor in desc $(\gamma)$. Choose $u \in U \cap \operatorname{desc}(\gamma)$ at maximal distance from $\gamma$, and let $y$ be the predecessor of $u$ in $\operatorname{desc}(\gamma)$ (note that $y \in A_{1}$ ). Since $\operatorname{desc}(\gamma) \backslash \operatorname{desc}(U)$ is finite, $\operatorname{desc}(y) \backslash \operatorname{desc}(U)$ is finite. So if $u^{\prime}$ is another successor of $y, \operatorname{desc}\left(u^{\prime}\right) \backslash \operatorname{desc}(U)$ is finite and our choice of $u$ implies that $u^{\prime} \in U \cap \operatorname{desc}(\gamma)$. Thus we can take $p$ to be the set of successors of $y$.

Now we finish off the proof of Claim 3.6. Since $\gamma=g_{2}\left(\gamma^{\prime}\right)$, the $q$-element subset $f(p)$ of $\operatorname{desc}\left(\gamma^{\prime}\right) \cap V$ has a common predecessor, $y^{\prime}$ say, in $\operatorname{desc}\left(\gamma^{\prime}\right)$ and $y=g_{2}\left(y^{\prime}\right)$. If $p$ has $n-1$ predecessors in $A_{1}$, then there is a copy of $\mathcal{T}_{n}$ in $A$ because $A$ is the free amalgam of $A_{1}$ and $T$ over $\operatorname{desc}(U)=\operatorname{desc}(V)$. This is a contradiction. Therefore $p$ has at most $n-2$ common predecessors in $A_{1}$ and this means that $p \in P^{\prime}$. It follows that $y^{\prime}=w_{f(p)}$ since a $q$-element set of vertices of $T^{\prime}$ has at most one common predecessor in $T^{\prime}$ as $T^{\prime}$ is a tree. Now as $w_{p}=g_{2}\left(w_{f(p)}\right)$, we have $y=g_{2}\left(y^{\prime}\right)=g_{2}\left(w_{f(p)}\right)=w_{p}$. This is a contradiction since $w_{p} \in B \backslash A_{1}$ and $y \in A_{1}$.
We have therefore shown that, $A_{1} \cup g_{2}(\operatorname{desc}(b))$ as a subdigraph of $D$ is isomorphic to $A$. So $A$ embeds in $D$ and therefore $A \in \mathscr{D}$. This completes the proof that $\mathscr{D}=\mathcal{C}_{n}$.

Finally suppose $\mathscr{D} \subseteq \mathcal{C}$ is a $\leq$-amalgamation class and $\tau_{n} \in \mathscr{D}$ for all $n \geq 1$. A similar argument as in the above proof can be used to show that any $A \in \mathcal{C}$ lies in $\mathscr{D}$. The two important points in that proof which we need to modify slightly, are the choice of the digraph $B$ and of the subset $P^{\prime}$ of $P$. We want a digraph $B \in \mathscr{D}$ containing $A_{1}$ such that every $q$-element set of vertices of $U$ has at least $M+1$ common predecessors in $B$, where $M$ is the greatest number of common predecessors of $p$ in $A_{1}$ as $p$ ranges over $P$. For this we apply Lemma 3.3 to $A_{1}$ with $N:=M+1$. In this case it will follow that for every $p$ in $P$, there is at least one common predecessor of $p$ in $B$ which does not lie in $A_{1}$. We then take $P^{\prime}$ to be the subset of $P$ consisting of all $q$-element sets $p$ which have at least one common predecessor in $A_{1}$ and such that $f(p)$ has a common predecessor in $T^{\prime}$. The remainder of the argument follows similarly, except that when showing that $A_{1} \cap g_{2}(\operatorname{desc}(b))=\operatorname{desc}(U)$, there is only one case to consider since for every $p$ in $P^{\prime}$ there is a vertex $w_{p} \in B \backslash A_{1}$. We deduce the following.

Proposition 3.7. Let $\mathscr{D} \subseteq \mathcal{C}$ be $a \leq$-amalgamation class with $\mathcal{T}_{n} \in \mathscr{D}$ for all $n \geq 1$. Then $\mathscr{D}=\mathcal{C}$.
We have therefore shown that
Theorem 3.8. Any $\leq$-amalgamation class $\mathfrak{D} \subseteq \mathcal{C}$ is equal to $\mathcal{C}$ or to $\mathcal{C}_{n}$ for some $n \geq 2$.
This means that if $D$ is a countable descendant-homogeneous digraph whose descendant set is isomorphic to $T$, then $D \cong D_{n}$ for some $n \in\{2, \ldots, \infty\}$.

## 4. A general construction

### 4.1. Descendant sets

In this subsection we prove the following.
Theorem 4.1. Suppose $\Gamma$ is a countable digraph. Then there is a countable, vertex transitive, descendant-homogeneous digraph $M$ in which all descendant sets are isomorphic to $\Gamma$ if and only if the following conditions hold:
(C1) $\operatorname{desc}(u) \cong \Gamma$ for all $u \in \Gamma$;
(C2) If $X$ is a finitely generated subdigraph of $\Gamma$ then the subgroup of automorphisms of $X$ which extend to automorphisms of $\Gamma$ is of countable index in $\operatorname{Aut}(\Gamma)$.
For one direction of this, suppose $M$ is a vertex transitive, descendant-homogeneous digraph. The descendant sets of vertices in $M$ are all isomorphic to a fixed digraph $\Gamma$, so (C1) holds. Condition (C2) is a special case of ( $4^{\prime}$ ) in Remark 2 , so follows from Theorem 2.1 and Remark 2.

We now prove the converse. So for the rest of this subsection, suppose that $\Gamma$ is a countable digraph which satisfies conditions ( C 1 ) and ( C 2 ). Let $\mathcal{C}_{\Gamma}$ be the class of digraphs $A$ satisfying the following conditions:
(D1) $\operatorname{desc}(a)$ is isomorphic to $\Gamma$, for all $a \in A$;
(D2) $A$ is finitely generated;
(D3) for $a, b \in A$, the intersection $\operatorname{desc}(a) \cap \operatorname{desc}(b)$ is finitely generated.
Then $\mathcal{C}_{\Gamma}$ is closed under isomorphism and taking finitely generated descendant-closed substructures. Moreover, it is easy to see that if $A \leq B_{1}, B_{2} \in \mathcal{C}_{\Gamma}$ and $A$ is finitely generated, then the free amalgam of $B_{1}$ and $B_{2}$ over $A$ is in $\mathcal{C}_{\Gamma}$. Thus, Theorem 4.1 will follow once we verify that the countability conditions in (1) and (4) of Theorem 2.1 hold for $\mathcal{C}_{\Gamma}$. The following lemma will suffice.

Lemma 4.2. Suppose $A \in \mathcal{C}_{\Gamma}$. Then there are only countably many isomorphism types of $\leq$-embeddings $f: A \rightarrow B$ with $B \in \mathcal{C}_{\Gamma}$.
Once we have this, taking $A=\emptyset$ (or $A=\Gamma$ ) gives that $\mathcal{C}_{\Gamma}$ contains only countably many isomorphism types; for fixed $A, B \in \mathcal{C}_{\Gamma}$, the lemma gives condition (4) of Theorem 2.1.
Proof. We say that a $\leq$-embedding $f: A \rightarrow B$ with $A, B \in \mathcal{C}_{\Gamma}$ is an $n$-extension if $B$ can be generated by $f(A)$ and at most $n$ extra elements. We prove by induction on $n$ that for every $A \in \mathcal{C}_{\Gamma}$ there are only countably many isomorphism types of $n$-extensions of $A$.

Suppose $f: A \rightarrow B$ is a 1-extension (with $A, B \in \mathcal{C}_{\Gamma}$ ). Let $b \in B$ be such that $B$ is generated by $f(A)$ and $b$ and let $C=f(A) \cap \operatorname{desc}(b)$. It follows from property (D3) in $B$ and finite generation of $A$, that $C$ is finitely generated. Moreover, as each of $f(A)$ and $\operatorname{desc}(b)$ is descendant-closed in $B$, we have that $B$ is the free amalgam of $f(A)$ and $\operatorname{desc}(b)$ over $C$. Choose an isomorphism from $\operatorname{desc}(b)$ to $\Gamma$ and let $h$ be the restriction of this to $C$ and $g: f^{-1}(C) \rightarrow \Gamma$ be given by $g=h \circ f$. Then, in the notation of Remark 1, we have an isomorphism from $B$ to $A *_{g} \Gamma$ and therefore $f$ is isomorphic to a 1-extension $A \rightarrow A *_{g} \Gamma$ for some finitely generated $D \leq A$ and $\leq-e m b e d d i n g ~ g: D \rightarrow \Gamma$.

There are countably many possibilities for $D$ and the image $g(D)$ here (as $D$ is finitely generated), so it will suffice to show that there are only countably many isomorphism types of $A *_{g} \Gamma$ with $g: D \rightarrow \Gamma$ having fixed domain $D$ and image $E \leq \Gamma$. If $g_{1}, g_{2}: D \rightarrow \Gamma$ have image $E$ then $g_{1} \circ g_{2}^{-1}$ gives an automorphism of $E$. This extends to an automorphism of $\Gamma$ if and only if there is an isomorphism between the extensions $g_{i}: A \rightarrow A *_{g_{i}} \Gamma$. Thus, the isomorphism types here are in one-to-one correspondence with the cosets in $\operatorname{Aut}(E)$ of the subgroup of automorphisms which extend to automorphisms of $\Gamma$. So there are only countably many isomorphism types, by (C2).

This proves that there are countably many isomorphism types of 1-extensions of $A$. For the inductive step, we can take countably many representatives $f_{j}: A \rightarrow B_{j}^{\prime}($ for $j \in \mathbb{N})$ of the isomorphism types of $(n-1)$-extensions of $A$, and representatives $h_{j k}: B_{j}^{\prime} \rightarrow B_{j k}^{\prime}$ of the 1-extensions of $B_{j}^{\prime}($ for $j, k \in \mathbb{N}$ ). We claim that any $n$-extension $f: A \rightarrow B$ is isomorphic to some $h_{j k} \circ f_{j}: A \rightarrow B_{j k}^{\prime}$. Indeed, let $f(A) \leq B_{1} \leq B$ be such that $B_{1}$ is generated by $f(A)$ and $n-1$ elements, and $B$ is generated by $B_{1}$ and one extra element. So we can write $f=i \circ g$ where $g: A \rightarrow B_{1}$ is an $(n-1)$-extension and $i: B_{1} \rightarrow B$ is a 1-extension. There is $j \in \mathbb{N}$ and an isomorphism $h: B_{j}^{\prime} \rightarrow B_{1}$ with $h \circ f_{j}=g$. We can then find $k \in \mathbb{N}$ and an isomorphism $p: B_{j k}^{\prime} \rightarrow B$ with $i \circ h=p \circ h_{j k}$. Then $p \circ h_{j k} \circ f_{j}=i \circ g=f$, as required.

It then follows by Theorem 2.1 that the Fraïssé limit $D_{\Gamma}$ of $\left(\mathcal{C}_{\Gamma}, \leq\right)$ is a countable descendant-homogeneous digraph with $\mathcal{C}_{\Gamma}$ as its class of finitely generated $\leq$-subdigraphs. Vertex transitivity follows from (C1).

### 4.2. Examples and further remarks

In this subsection we show that a class of digraphs $\Gamma$ arising in [1] in the context of highly arc transitive digraphs satisfy the conditions in Theorem 4.1 and therefore arise as the descendant sets in descendant-homogeneous digraphs. We begin by reviewing some of the results of [1] and related papers.

The paper [1] studies highly arc transitive digraphs of finite out-valency and gives conditions which the descendant set $\Gamma$ of a vertex in such a digraph must satisfy. In particular:

Theorem 4.3. Suppose $\Gamma$ is the descendant set of a vertex $\alpha$ in an infinite highly arc transitive digraph $D$ of finite out-valency. Then the following properties hold:
(T1) $\Gamma=\operatorname{desc}(\alpha)$ is a rooted digraph with finite out-valency and $\operatorname{desc}^{s}(\alpha) \cap \operatorname{desc}^{t}(\alpha)=\varnothing$ whenever $s \neq t$.
(T2) $\operatorname{desc}(u) \cong \Gamma$ for all $u \in \Gamma$.
(T3) $\operatorname{Aut}(\Gamma)$ is transitive on $\operatorname{desc}^{s}(\alpha)$, for all s.
(T4) There is a natural number $N=N_{\Gamma}$ such that for $l>N$ and $x, a \in \Gamma$, if $b \in \operatorname{desc}^{l}(x) \cap \operatorname{desc}^{1}(a)$, then $a \in \operatorname{desc}(x)$.
Proof. Properties (T2) and (T3) follow immediately from high arc transitivity of $D$. Property (T1) is proved in Lemma 3.1 of [1] and (T4) is deduced from (T1, T2, T3) in [1, Propositionp 4.7(a)].

Remark 4. The paper [2] shows that there are only countably many isomorphism types of digraphs $\Gamma$ satisfying properties (T1, T2, T3). In fact, the same is true with (T3) replaced by the weaker:
(G3) There is a natural number $k$ such that if $\ell \geq k$ and $x \in \operatorname{desc}^{\ell}(\alpha)$ and $\beta \in \operatorname{desc}^{1}(\alpha)$, then $\operatorname{desc}(\beta) \cap \operatorname{desc}(x) \neq \emptyset$ implies $x \in \operatorname{desc}(\beta)$.
Moreover, these (T1, T2, G3) imply (T4). See Corollary 1.5 and Lemma 2.1 of [2] for proofs.
Explicit examples $\Gamma(\Sigma, k)$ of digraphs satisfying (T1, T2, T3) (and which are not trees) are constructed in Section 5 of [1] and constructions of highly arc transitive, but not descendant-homogeneous, digraphs with these as descendant sets are given in $[1,3]$. The construction we give here (using Theorem 4.1) gives a highly arc transitive, descendant-homogeneous digraph with descendant set $\Gamma(\Sigma, k)$ (and which does not have property $Z$ ). Indeed, it is a slightly curious corollary of the results of this section that if $\Gamma$ is a digraph of finite out-valency which is the descendant set of a vertex in an infinite, highly arc transitive digraph, then there is a descendant-homogeneous, highly arc transitive digraph which has $\Gamma$ as its descendant set.

Corollary 4.4. Suppose $\Gamma$ is a digraph of finite out-valency which satisfies conditions (T1, T2, T4). Then there is a countable, vertex transitive, descendant-homogeneous digraph in which all descendant sets are isomorphic to $\Gamma$.
Proof. We use Theorem 4.1. The digraph $\Gamma$ satisfies condition (C1) of this, by assumption (T2). So it remains to show that $\Gamma$ satisfies (C2).

Let $X$ be a finitely generated subdigraph of $\Gamma$ with minimal generating set $\left\{x_{1}, \ldots, x_{k}\right\}$. Let $N=N_{\Gamma}$ as in (T4) and $Y:=\bigcup_{i=1}^{k} B^{N}\left(x_{i}\right)$. Note that by (T1) $Y$ is finite, and it is clearly invariant under Aut $(X)$. We will show that any automorphism of $X$ fixing $Y$ pointwise extends to an automorphism of $\Gamma$. Such automorphisms form the kernel of the restriction of Aut $(X)$ to $Y$, and so form a subgroup of finite index in $\operatorname{Aut}(X)$ (the quotient group is just the group of permutations induced on the finite set $Y$ by elements of $\operatorname{Aut}(X)$ ). Thus condition (C2) will follow.

Let $a, b \in \Gamma$. We first observe that if $b \in X \backslash Y$ and $a$ is a predecessor of $b$ in $\Gamma$, then $a \in X$. Indeed, $b \in \operatorname{desc}^{l}\left(x_{i}\right)$ for some $l>N$ and $i \in\{1, \ldots, k\}$. Then by definition of $N, a \in \operatorname{desc}\left(x_{i}\right)$. Since $\operatorname{desc}\left(x_{i}\right) \subseteq X$, it follows that $a \in X$.

Let $\gamma$ be an automorphism of $X$ which fixes $Y$ pointwise. Define $\theta=\gamma \cup i d_{\Gamma \backslash X}$. To prove $\theta$ is an automorphism of $\Gamma$ we must show that $\theta$ preserves edges and non-edges. For $u \in(\Gamma \backslash X) \cup Y, \theta u=u$ and for $u \in X, \theta u=\gamma u$. So for $a, b \in(\Gamma \backslash X) \cup Y$, we have $\theta(a, b)=(\theta a, \theta b)=(a, b)$. Similarly, $\theta$ preserves edges and non-edges when $a, b \in X$ as in this case, $\theta(a, b)=\gamma(a, b)$. Now suppose $a \in \Gamma \backslash X$ and $b \in X \backslash Y$. Then $\theta(a, b)=(\theta a, \theta b)=(a, \gamma b)$. Since $\gamma$ preserves $Y, \gamma b \in X \backslash Y$. Then by the observation above, $(a, b)$ and $(a, \gamma b)$ are non-edges.

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