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Computing in algebraic geometry and commutative algebra using *Macaulay 2*

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Abstract

We present recent research of Eisenbud, Fløystad, Schreyer, and others, which was discovered with the help of experimentation with *Macaulay* 2. In this invited, expository paper, we start by considering the exterior algebra, and the computation of Gröbner bases. We then present, in an elementary manner, the explicit form of the Bernstein–Gelfand–Gelfand relationship between graded modules over the polynomial ring and complexes over the exterior algebra, that Eisenbud, Fløystad and Schreyer found. We present two applications of these techniques: cohomology of sheaves, and the construction of determinantal formulae for (powers of) Macaulay resultants. We show how to use *Macaulay* 2 to perform these computations. © 2003 Published by Elsevier Ltd.

1. Introduction

This invited talk at ISSAC 2002 has three goals. We wish to present some exciting new research of Eisenbud, Fløystad, Schreyer, and others, which was discovered with the help of experimentation with *Macaulay 2*. In the process, we hope to convince the reader that it is possible to compute with relatively abstract notions in algebraic geometry, and finally, we show how these computations can be performed using *Macaulay 2*.

Macaulay 2 is computer software for algebraic geometry, commutative algebra and related fields. Grayson and I have been working on *Macaulay 2* since we started the project in 1993. *Macaulay 2* is freely available (Grayson and Stillman, 1993–2003).

This paper is an introduction to the work of Eisenbud, Fløystad, and Schreyer. More details and proofs can be found in the papers (Eisenbud et al., 2001; Eisenbud and Schreyer, 2001; Decker and Eisenbud, 2001) and in the book (Eisenbud, 2003). In particular, they prove considerably more than we explain here. What we do instead is explain their constructions and apply them to examples. We present two applications of these techniques: cohomology of sheaves, and the construction of determinantal formulae for

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(powers of) Macaulay resultants. We culminate with finding an explicit skew symmetric 8×8 matrix whose pfaffian (square root of the determinant) is the Macaulay resultant of three ternary quadratic forms. This is one of the new determinantal-like formulae that Eisenbud and Schreyer found in Eisenbud and Schreyer (2001).

2. The exterior algebra

Fix a field k. The exterior algebra E on n letters e_1, \ldots, e_n is defined to be the free associative k-algebra on e_1, \ldots, e_n , modulo the relations

$$e_1^2=\cdots=e_n^2=0,$$

and

$$e_i e_j = -e_j e_i,$$

for all *i* and *j*. There is no truly standard notation for this algebra. We will denote this exterior algebra by $E = k \langle e_1, \ldots, e_n \rangle$. A *k*-basis for *E* consists of all square-free monomials, so has dimension 2^n as a vector space over *k*.

Since multiplication is almost commutative, Gröbner bases and Buchberger's algorithm both extend, with simple modifications, to ideals in E, as well as to modules over E. (We always consider right-ideals and right modules.) The only difference is that the notion of "S-pair" must be modified. An example should make it clear.

Example 1. Let $E = \mathbb{Q}(a, b, c, d)$ be the exterior algebra on letters a, \ldots, d . Let I be the ideal generated by F = ac - bd and G = bc - ad.

Over the usual polynomial ring, the Buchberger algorithm works by selecting pairs of polynomials, cancelling their lead terms by taking a linear combination of the two polynomials, and computing its remainder.

We start to compute a Gröbner basis by considering the pair (F, G). The combination that cancels lead terms is

bF + aG = b(ac - bd) + a(bc - ad) = 0.

The one difference, other than arithmetic, is that we may uncover new lead terms by multiplying a polynomial such as F by any variable which occurs in its lead term. Thus,

$$aF = a(ac - bd) = -abd$$

and

$$cF = c(ac - bd) = bcd$$

H = abd is not divisible by the lead terms of F or G, and so we add it to the Gröbner basis. The element *bcd* does reduce to zero, by subtracting *dG*.

All other "S-pairs" on each of F, G and H reduce to zero, and so $\{F, G, H\}$ is a Gröbner basis of I.

By using this straightforward extension of Buchberger's algorithm, we can compute Gröbner bases. As with polynomial rings, if we extend the algorithm to modules, and keep track of how Gröbner basis elements are expressed in terms of the original generators, we can also compute *syzygies*, i.e. kernels of $r \times s$ matrices

 $\phi: E^s \longrightarrow E^r.$

For example, let's compute with the above ideal I using Macaulay 2.

i1 : E = QQ[a..d, SkewCommutative=>true];

Multiplication is as defined above.

i2 : c*b o2 = -b*c o2 : E i3 : b^2 o3 = 0 o3 : E i4 : I = ideal(a*c-b*d, b*c-a*d) o4 = ideal (a*c - b*d, b*c - a*d) o4 : Ideal of E i5 : transpose gens gb I o5 = {-2} | bc-ad | {-2} | ac-bd | {-3} | abd | 3 1 o5 : Matrix E <--- E i6 : m = generators I o6 = | ac-bd bc-ad | 1 2 06 : Matrix E <--- E 2 i7 : m1 = syz m o7 = {2} | c -d a b 0 | {2} | -d c b a bd | 2 5 o7 : Matrix E <--- E

We can compute the kernel of this map.

i8 : syz m1 $o8 = \{3\} | c -d 0 0 a b 0 0 0 0 | \\ \{3\} | -d c 0 0 b a 0 0 0 b d | \\ \{3\} | 0 0 a b c -d 0 0 0 0 0 | \\ \{3\} | 0 0 b a -d c 0 0 b d 0 | \\ \{4\} | 0 0 0 0 0 0 0 d b -a -c | \\ 5 10 \\ o8 : Matrix E < --- E$

The resolution routine iterates this process as far as we want, to compute a (minimal) free resolution for some number of steps.

i9 : resolution(cokernel m, LengthLimit=>7)

 $o9 = E \xrightarrow{1} - E \xrightarrow{2} E \xrightarrow{5} 10 \xrightarrow{18} 30 \xrightarrow{47} 70$ $o9 = E \xrightarrow{1} - E \xrightarrow{10} E \xrightarrow{10} E \xrightarrow{18} - E \xrightarrow{10} E \xrightarrow{10}$

o9 : ChainComplex

There are two major differences between computing over the usual polynomial ring (the symmetric algebra) and the exterior algebra: first, Gröbner bases and syzygies are *much* easier to compute over the exterior algebra. This is in large part due to the small number of monomials in the exterior algebra. Second, Hilbert proved in 1890 that minimal free resolutions over the polynomial ring are always finite. Over the exterior algebra, free resolutions are almost never finite. However, because of the small number of monomials in E, finite parts of these resolutions can often be found quickly.

3. The link between the exterior algebra and the symmetric algebra

Let $S = k[x_1, ..., x_n]$, and let $E = k \langle e_1, ..., e_n \rangle$ be the corresponding exterior algebra. Let

$$M = \bigoplus_{d \in \mathbb{Z}} M_d,$$

be a graded *S*-module, where each degree *d* piece M_d is a finite dimensional vector space over *k*. Once and for all, choose a basis of each vector space M_d . If $m \in M_d$, denote by [m] its representation in this basis. Throughout this section, we set $m_d := \dim M_d$, for all *d*.

The whole theory rests on the following particular method for encoding the data of multiplication by linear forms from M_d to M_{d+1} .

Definition 2 (Bernstein–Gelfand–Gelfand Maps). For a given d, define the d th BGG map

 $\phi_d: E^{m_d} \longrightarrow E^{m_{d+1}}$

of *M* by

$$[m]\mapsto \sum_{i=1}^n e_i[x_im].$$

We could do everything without coordinates, and then this would be the adjoint map to the multiplication map. Given a vector space V of dimension n, define $W = V^*$ to be the dual vector space, and then set $S = S_*V$ and $E = \Lambda^*W$ to be the symmetric and exterior algebras respectively. If V is the span of x_1, \ldots, x_n , then this is exactly what we have considered already. The multiplication map is $V \otimes_k M_d \longrightarrow M_{d+1}$. The adjoint is $M_d \longrightarrow V^* \otimes_k M_{d+1}$. Considering the elements of $V^* = W$ as linear elements of E, the adjoint gives a matrix of linear forms ϕ_d defined by the above formula.

Example 3. As a simple example, consider $S = k[x_1, x_2]$ and the corresponding exterior algebra $E = k \langle e_1, e_2 \rangle$. Let

$$M = S/(x_1^2, x_2^3) = M_0 \oplus M_1 \oplus M_2 \oplus M_3,$$

where $M_0 = k$, $M_1 = k^2$, $M_2 = k^2$, and $M_3 = k$. We choose monomials as the basis elements of these vector spaces. The basis of M_0 is {[1]}, the basis of M_1 is {[x_1], [x_2]}, the basis of M_2 is {[x_1x_2], [x_2^2]} and the basis of M_3 is {[$x_1x_2^2$]}.

Setting d = 1 for example, $\phi_1 : E^2 \longrightarrow E^2$ is defined by

$$[x_1] \mapsto e_1[x_1^2] + e_2[x_1x_2],$$

and

$$[x_2] \mapsto e_1[x_1x_2] + e_2[x_2^2]$$

Since $[x_1^2] = 0$, $\phi_1 = \begin{pmatrix} e_2 & e_1 \\ 0 & e_2 \end{pmatrix}$.

Computing ϕ_0 and ϕ_2 , we obtain a sequence

$$0 \longrightarrow k \xrightarrow{\begin{pmatrix} e_1 \\ e_2 \end{pmatrix}} k^2 \xrightarrow{\begin{pmatrix} e_2 & e_1 \\ 0 & e_2 \end{pmatrix}} k^2 \xrightarrow{\begin{pmatrix} e_2 & e_1 \\ 0 & e_2 \end{pmatrix}} k^2 \xrightarrow{\begin{pmatrix} e_2 & e_1 \end{pmatrix}} k \longrightarrow 0$$

Note that this is a complex: applying two maps in a row gives zero.

In the general case, if we apply the construction for each d to a graded S-module M, we get a (possibly infinite) sequence of maps:

 $\cdots \longrightarrow E^{m_0} \xrightarrow{\phi_0} E^{m_1} \xrightarrow{\phi_1} E^{m_2} \xrightarrow{\phi_2} \cdots$

There are two basic facts about this sequence.

- This is a *complex*, i.e. $\phi_{i+1}\phi_i = 0$ for all *i*. This is a simple exercise, using the fact that multiplication in *S* is commutative.
- This complex is eventually exact, i.e. for $i \gg 0$, ker $(\phi_{i+1}) = im(d_i)$. This theorem is proved in Eisenbud et al. (2001). They also show that if the Castelnuovo–Mumford regularity of *M* is *r*, then this sequence is exact after the *r*th step.

The cohomology of this sequence would be interesting to investigate, but for now, we take a tail of this complex which is exact:

 $\cdots \longrightarrow E^{m_r} \xrightarrow{\phi_r} E^{m_{r+1}} \xrightarrow{\phi_{r+1}} E^{m_{r+2}} \xrightarrow{\phi_{r+2}} \cdots$

The crucial link is the following exact sequence.

Definition 4. The **Tate resolution** T(M) of M is the exact complex (possibly infinite in both directions)

 $\cdots \longrightarrow E^{m_{r-2}} \xrightarrow{\psi_{r-2}} E^{m_{r-1}} \xrightarrow{\psi_{r-1}} E^{m_r} \xrightarrow{\psi_r} E^{m_{r+1}} \xrightarrow{\psi_{r+1}} \cdots$

obtained by computing a *free resolution* of ker(ϕ_r), where *r* is chosen large enough so that the tail of the complex is exact, and $\psi_i = \phi_i$, for $i \ge r$.

In the example above, the Tate resolution is the zero complex, since the complex is eventually zero.

Eisenbud, Fløystad, and Schreyer prove that the Tate resolution is independent (in a suitable sense) of the truncation location *r*, as long as it is chosen so that the tail is exact. This implies that if $M_{\geq e} := \bigoplus_{d \geq e} M_d$ is the *e*th truncation of *M*, then $\mathbf{T}(M) = \mathbf{T}(M_{\geq e})$.

This is reminiscent of a property of sheaves on projective space. So, before continuing, let us brush up on sheaves.

3.1. An aside: a crash course on implementing coherent sheaves on \mathbb{P}^{n-1}

Serre's famous FAC paper (Serre, 1955) introduced sheaves to algebraic geometry. What is perhaps less well known is that he describes (in essence) how to represent sheaves on projective space as modules, and how to compute their cohomology. In this sense, his paper is perhaps the first paper in computational abstract algebraic geometry.

A graded *S*-module *M* determines in a canonical manner a coherent sheaf \widetilde{M} on \mathbb{P}^{n-1} , and all coherent sheaves on \mathbb{P}^{n-1} arise in this manner. Unfortunately, the correspondence is not one to one. For example, if $M_{\geq e}$ is the *e*th truncation of the module *M*, then $\widetilde{M} = \widetilde{M_{\geq e}}$. In fact, two coherent sheaves \widetilde{M} and \widetilde{N} are isomorphic if and only if there is an integer *e* such that $M_{\geq e} \simeq N_{\geq e}$.

If $X \subset \mathbb{P}^{n-1}$ is a projective variety defined by an ideal $I_X \subset S$, then $\mathcal{O}_X := \widetilde{S/I_X}$ is called the sheaf of regular functions on X.

Another important construction is the twist of a sheaf. If *e* is an integer, then $\widetilde{M}(e) := \widetilde{M(e)}$, where M(e) is the same module as *M*, but with a shift in the grading: $M(e)_d := M_{e+d}$.

Cohomology of sheaves can be computed using the representation of the sheaf as a graded *S*-module. We will not describe these algorithms here (see Eisenbud's chapter in Vasconcelos, 1998). Instead, we use *Macaulay 2* to compute some of the cohomology groups of the sheaves $\mathcal{O}_C(d)$, where $C \subset \mathbb{P}^3$ is the twisted cubic curve. Note that $h^i(\widetilde{M}) := \dim_k H^i(\widetilde{M})$ is the notation often used for the dimensions of the cohomology groups (and are also *k* vector spaces).

 $\mathbb{P}^1 \longrightarrow \mathbb{P}^3$

which sends

 $(s, t) \mapsto (W, X, Y, Z) = (s^3, s^2t, st^2, t^3).$

The ideal I_C of the image is generated by $\{X^2 - WY, Y^2 - XZ, WZ - XY\}$.

```
i10 : S = QQ[W,X,Y,Z];
i11 : IC = ideal(X^2-W*Y, Y^2-X*Z, W*Z-X*Y);
o11 : Ideal of S
i12 : C = variety IC
o12 = C
o12 : ProjectiveVariety
i13 : (HH^0(00_C), HH^0(00_C(1)), HH^1(00_C(-5)))
o13 = (QQ, QQ, QQ)
o13 : Sequence
```

So $h^0(\mathcal{O}_C) = 1$, $h^0(\mathcal{O}_C(1)) = 4$, and $h^1(\mathcal{O}_C(-5)) = 14$. Here are more cohomology groups:

3.2. Tate resolutions and cohomology of coherent sheaves

What have we done so far? We start with a graded *S*-module *M*, or its associated sheaf \tilde{M} and obtain an exact complex $\mathbf{T}(M)$ of free *E*-modules, which is eventually linear. This is cute, and pretty, but so what? What good is it? Well, it turns out to be amazingly useful. We will see two completely different applications below. In addition, there are several other applications that we do not have the time or space to describe (see Eisenbud et al., 2001; Eisenbud and Schreyer, 2001; Decker and Eisenbud, 2001).

Let us compute the Tate resolution of the twisted cubic curve in projective 3-space. We'll use *Macaulay 2* to do the computations for us.

```
i16 : load "bgg.m2"; -- described in the appendix
i17 : E = QQ[w,x,y,z, SkewCommutative => true];
i18 : M = cokernel generators IC;
```

M is the homogeneous coordinate ring of the twisted cubic curve. The degree one and two parts have bases consisting of the following sets of monomials.

```
i19 : basis(1,M)
o19 = | W X Y Z |
o19 : Matrix
i20 : basis(2,M)
o20 = | W2 WX WY WZ XZ YZ Z2 |
o20 : Matrix
```

The routine bgg computes the map ϕ_d defined above.

```
i21 : phi1 = bgg(1,M,E)
o21 = {-2} | w 0 0 0 |
      {-2} | x w 0 0 |
      {-2} | y x w 0 |
      {-2} | z y x w |
     {-2} | 0 z y x |
      {-2} | 0 0 z y |
      {-2} | 0 0 0 z |
                      4
              7
o21 : Matrix E <--- E
i22 : phi2 = bgg(2,M,E)
o22 = {-3} | w 0 0 0 0 0 0 |
      {-3} | x w 0 0 0 0 0
                           1
      {-3} | y x w 0 0 0 0 |
      {-3} | z y x w 0 0 0 |
      {-3} | 0 z y x w 0 0
      {-3} | 0 0 z y x w 0 |
      {-3} | 0 0 0 z y x w |
      {-3} | 0 0 0 0 z y x |
      {-3} | 0 0 0 0 0 z y |
      {-3} | 0 0 0 0 0 0 z |
              10
                       7
o22 : Matrix E <--- E
i23 : phi2 * phi1
```

o23 = 0 023 : Matrix E <---- E

Here is a *Macaulay 2* routine for computing (a part of) the Tate resolution. The routine first truncates M at the regularity of M, calls bgg, and then computes several steps of a free resolution. Notice that the second matrix from the left is $\phi_2 : E^7 \longrightarrow E^{10}$ (up to change of basis).

```
i24 : Ta = tateResolution(presentation M,E,-3,4)

o24 = E^{13} <br/>
-4^{-3} <br/>
-2^{-1} <br/>
-1^{-1} <br/>
-2^{-1} <br/>
-1^{-1} <br/>
-2^{-1} <br/>
-1^{-1} <br/>
-1^{-
```

Caution! Macaulay 2 displays maps from right to left, so the eventually linear part here is the leftmost displayed part of the Tate resolution. For example, $\psi_{-1} : E^5 \longrightarrow E^3$ is

```
i25 : Ta.dd_1

o25 = {0} | 0 0 -wz wy wx |

{1} | z y x -w 0 |

{1} | 0 z y -x -w |

3 5

o25 : Matrix E <--- E
```

The graded pieces of each module are displayed using the betti command.

```
i26 : betti Ta

o26 = total: 13 10 7 4 3 5 8 11

0: 13 10 7 4 1 . . .

1: . . . 2 5 8 11
```

The entry in row d: and column c (where the first column displayed is c = -4) is the number of generators of degree d + c in the cth free module, where each variable in the exterior algebra has degree 1. For example, ψ_{-1} , as a map of graded free modules, has the form ψ_{-1} : $E(-2)^5 \longrightarrow E(-1)^2 \oplus E$, where E(-d) is the graded free module of rank one, having its generator in degree d.

Wait! These numbers in the betti diagram are the same numbers we encountered when computing the cohomology of \mathcal{O}_C and its twists. Eisenbud, Fløystad, and Schreyer observed this, and then were able to prove in general that the graded pieces of the Tate resolution $\mathbf{T}(M)$ are exactly the cohomology modules of \widetilde{M} and its twists. In terms of this betti diagram, the statement is:

Theorem 6 (Eisenbud et al., 2001). Let M be a graded $S = k[x_1, ..., x_n]$ module. Let \widetilde{M} be the corresponding sheaf. Then: the betti diagram of the Tate resolution of M has the form

4. Resultants, Chow forms, and the Tate resolution

There has been a great deal of interest in finding determinantal formulae for multivariate resultants (the *Macaulay* resultant), and for sparse resultants. For one such result from these proceedings, with pointers to the literature for others, see Khetan (2002).

Khetan (2002) discovered that their exterior algebra methods can be used to construct determinantal formulae for some of these resultants. In this section, we present a part of this work, leading up to an explicit Bezout formula for the Macaulay resultant of three quadratic forms in three variables.

Macaulay resultants are the Chow forms of Veronese varieties, and Eisenbud and Schreyer find formulae for Chow forms. Thus, our story starts with Chow forms.

4.1. The Chow divisor and Chow form of V

Let $X \subset \mathbb{P}^{n-1}$ be a projective variety of dimension d. Let G be the set of all codimension d + 1 planes L in \mathbb{P}^{n-1} . This is a Grassmann variety, and has dimension (d+1)(n-d-1).

The **Chow divisor** D_X of X is

 $D_X = \{ L \in G \mid X \cap L \neq \emptyset \}.$

It is an exercise in dimension theory to show that D_X has codimension one in G.

An element L of G is represented by a $(d + 1) \times n$ matrix M such that if $H_i = M_{i1}x_1 + \cdots + M_{in}x_n$, then $L = (H_1 = 0) \cap \cdots \cap (H_{d+1} = 0)$. The Chow divisor D_X is defined by a single equation Ch_X (the **Chow form**) in the indeterminates M_{ij} . This polynomial may also be expressed as a polynomial in the Plücker coordinates

$$[i_1, \dots, i_{d+1}] = \det \begin{pmatrix} M_{1,i_1} & M_{1,i_2} & \dots & M_{1,i_{d+1}} \\ \vdots & & \ddots & \vdots \\ M_{d+1,i_1} & M_{d+1,i_2} & \dots & M_{d+1,i_{d+1}} \end{pmatrix}.$$

The degree of the polynomial in the Plücker coordinates is the degree of X.

Let us now specialize to the Veronese surface $V \subset \mathbb{P}^5$. This surface is the image of the map

 $\mathbb{P}^2 \longrightarrow \mathbb{P}^5$

.

given by

$$(r, s, t) \mapsto (A, B, C, D, E, F) = (r^2, rs, rt, s^2, st, t^2).$$

The variety V has dimension 2 and degree 4.

What is the Chow form Ch_V of the Veronese surface? In this situation, G is the Grassmannian of codimension 3 subspaces of \mathbb{P}^5 . Each element $L \in G$ is determined by a 3×6 matrix

$$M = \begin{pmatrix} a_1 & \dots & a_6 \\ b_1 & \dots & b_6 \\ c_1 & \dots & c_6 \end{pmatrix}$$

where if

$$H_a = a_1 x_1 + \dots + a_6 x_6, H_b = b_1 x_1 + \dots + b_6 x_6, H_c = c_1 x_1 + \dots + c_6 x_6,$$

then $L = (H_a = 0) \cap (H_b = 0) \cap (H_c = 0)$. The Plücker coordinates are

$$[i, j, k] := \det \begin{pmatrix} a_i & a_j & a_k \\ b_i & b_j & b_k \\ c_i & c_j & c_k \end{pmatrix}.$$

 Ch_V is a polynomial in the 18 variables a, b, c. Ch_V can also be expressed in terms of the Plücker coordinates. Since the degree of the Veronese surface is four, general theory tells us that Ch_V is a polynomial of degree four in the Plücker coordinates (and so of degree 12 in the *a*, *b*, *c* variables).

Let

$$F_a = a_1r^2 + a_2rs + a_3rt + a_4s^2 + a_5st + a_6t^2$$

$$F_b = b_1r^2 + b_2rs + b_3rt + b_4s^2 + b_5st + b_6t^2$$

$$F_c = c_1r^2 + c_2rs + c_3rt + c_4s^2 + c_5st + c_6t^2.$$

We now identify the Chow form of V.

$$Ch_V(a, b, c) = 0$$

$$\iff V \cap (H_a = 0) \cap (H_b = 0) \cap (H_c = 0) \neq \emptyset$$

$$\iff F_a(r, s, t) = F_b(r, s, t) = F_c(r, s, t) = 0 \quad \text{for some } (r, s, t) \in \mathbb{P}^2$$

$$\iff Res_{2,2,2}(F_a, F_b, F_c) = 0.$$

Therefore (since both are irreducible polynomials)

 $Ch_V(a, b, c) = Res_{2,2,2}(F_a, F_b, F_c)$

is the Macaulay resultant of three ternary quadratic forms.

The goal is to find determinantal formulae for resultants such as this. Here is one example, which appears in the list in Gelfand et al. (1994). The Chow form Ch_V is the determinant of the 6×6 matrix

| | a_1 | - | - | [1, 2, 6] | 0 | [1, 2, 3] | |
|---|-------|-------|-------|-----------------------|-----------------------|-----------------------|---|
| | a_2 | b_2 | c_2 | [1, 4, 6] | [1, 4, 5] | [1, 2, 5] - [1, 3, 4] | |
| | a_3 | b_3 | с3 | [1, 5, 6] - [2, 3, 6] | [1, 4, 6] | [1, 2, 6] | |
| | a_4 | b_4 | С4 | 0 | [2, 4, 5] | [1, 4, 5] | ŀ |
| | a_5 | b_5 | c_5 | [3, 4, 6] | [3, 4, 5] + [2, 4, 6] | [1, 4, 6] | |
| 1 | a_6 | b_6 | c_6 | [3, 5, 6] | [3, 4, 6] | 0 / | |

If we use a Laplace expansion with the first three columns, we see that the determinant is a polynomial of degree 4 in the cubics [i, j, k].

Question: Can the polynomial Ch_V be expressed as the determinant of a 4×4 matrix whose entries are linear in the Plücker coordinates [i, j, k]? Or, if not, is there *any* nice formula involving only the [i, j, k]'s? In fact, there is no such 4×4 determinant, but Eisenbud and Schreyer construct an 8×8 skew symmetric matrix whose pfaffian is $Ch_V = Res_{2,2,2}$. In the rest of this section, we describe their construction, and at the end we obtain the 8×8 matrix explicitly.

4.2. The Eisenbud–Schreyer construction

For the general construction, see Eisenbud and Schreyer (2001). Here we present an important special case, which works for the Veronese and many other cases.

Start with a variety $X \subset \mathbb{P}^{n-1}$. Let M be a graded $S = k[x_1, \ldots, x_n]$ module, which is supported on X (i.e. $I_X \subset ann(M)$). Assume that the sheaf associated to M is locally free on X of rank r. There is an additional assumption on M (M is "Ulrich") for the formula below to work as nicely as it does, but we will not get into that here. See Eisenbud and Schreyer (2001) for the specific condition.

From this module M, find the Tate resolution of M. In the case when M is suitably nice, i.e. is "Ulrich", the resolution has the form

 $\cdots \longrightarrow E^{\gamma} \longrightarrow E^{\alpha} \xrightarrow{\psi} E^{\alpha} \longrightarrow E^{\beta} \longrightarrow \cdots$

The entries of all of the matrices except ψ are linear in e_1, \ldots, e_n , and the non-zero entries of the matrix ψ all have degree d+1. The final step of the construction is to create the $\alpha \times \alpha$ matrix $\mathbf{U}(\psi)$ whose entries are obtained from those of ψ by setting $T(e_i e_j e_k) = [i, j, k]$, and extending via *k*-linearity. For example, $T(e_1 e_2 e_3 + e_1 e_2 e_4) = [1, 2, 3] + [1, 2, 4]$.

Theorem 7. If \widetilde{M} is locally free on X of rank r, and M is "Ulrich", then

 $\det \mathbf{U}(\psi) = (Ch_X)^r.$

In particular, if \widetilde{M} is a line bundle (r = 1), then the construction provides a determinantal formula for Ch_X .

To summarize,

Construction 8 (A Power of the Chow Form).

input: An S-module M, such that the sheaf M is locally free of rank r on a variety $X \subset \mathbb{P}^{n-1}$, and M is "Ulrich".

output: A square matrix $\mathbf{U}(\psi)$ whose entries are linear forms in the Plücker coordinates, such that $(Ch_X)^r = \det \mathbf{U}(\psi)$.

begin

Compute the matrix $\psi = \psi_{-1}$ in the Tate resolution.

if the matrix ψ is not square or has entries not of degree d + 1

then error *M* is not Ulrich.

Return $\mathbf{U}(\psi)$.

end.

In practice, such an Ulrich module M may not exist, or at least might be difficult to find. Eisenbud and Schreyer (2001) give criteria for when such a module exists, and how to find one. In particular, with their construction they can reproduce all of the known determinantal formulae (at least the ones that appear in Gelfand et al., 1994 for Macaulay resultants).

4.3. The resultant $Res_{2,2,2}(F_a, F_b, F_c)$

In our example of the Veronese surface $V \subset \mathbb{P}^5$, we choose M = TV to be a graded S-module which corresponds to the (rank two) tangent bundle of V. In this case, the module M is "Ulrich". When we apply the above construction, we obtain an 8×8 skew symmetric matrix whose determinant is the square of $Ch_V = Res_{2,2,2}(F_a, F_b, F_c)$. Recall that the determinant of a skew symmetric matrix (even a matrix of polynomials) is a square. Its square root is called the *pfaffian* of the matrix. Therefore the resultant $Res_{2,2,2}$ is the pfaffian of an 8×8 matrix of linear forms in the Plücker coordinates [i, j, k]. We will construct this matrix using *Macaulay 2*.

```
i27 : S3= QQ[r,s,t];
```

We use the variables A, \ldots, F , and a, \ldots, f instead of x_i and e_i , to improve the readability of the *Macaulay 2* output.

```
i28 : S6 = QQ[A..F,Degrees=>{2,2,2,2,2,2}];
i29 : E6 = QQ[symbol a..symbol f,SkewCommutative=>true]
o29 = E6
o29 : PolynomialRing
i30 : FV = map(S3,S6,{r^2, r*s, r*t, s^2, s*t, t^2})
o30 = map(S3,S6,{r , r*s, r*t, s , s*t, t })
o30 : RingMap S3 <--- S6</pre>
```

The following three lines of *Macaulay 2* code is one way to compute TV, a module in S6 corresponding to the tangent sheaf of V. This method starts with the tangent bundle

of \mathbb{P}^2 , truncates it so that all of the generators are in degree zero, and pushes it forward to a bundle on the image *V*.

```
i31 : TP2 = coker transpose vars S3
o31 = cokernel {-1} | r |
            {-1} | s |
            {-1} | t |
                          3
o31 : S3-module, quotient of S3
i32 : MO = prune truncate(0, TP2)
o32 = cokernel | 0 0 s 0 0 0 0 r t |
            | 0 0 -t 0 r t 0 0 0 |
            | 0 0 0 0 -s 0 0 -t 0 |
            | 0 s 0 0 0 0 r 0 0 |
            | 0 -t 0 r 0 s 0 0 t |
            | 0 0 0 -s 0 0 -t 0 0 |
            | s 0 0 0 0 0 0 0 r |
            | -t 0 0 0 0 r 0 0 |
                          8
o32 : S3-module, quotient of S3
i33 : TV = prune coker pushForward1(FV,MO)
o33 = cokernel | -F 0 0 -E 0 0 C 0 0 0 0 0 B 0 0 0 0 A 0 0
                                                              . . .
            | E 0 0 D 0 0 C 0 0 0 0 B 0 0 0 0 A 0
                                                               . . .
            0
                0 0 0 0 0 0 0 -E -F 0 0 0 0 -D -E 0 0 D E -B ···
            | 0 -F 0 0 -E 0 E F C 0 0 0 D E B 0 0 0 0 0 A ···
            | 0 E 0 0 D 0 000 C 0 0 000 B 0 0 000 ···
            | 0 0 0 0 0 0 0 0 0 0 -E -F 0 0 0 0 -D -E E F 0
                                                               . . .
            | 0 0 -F 0 0 -E F 0 0 0 C 0 E F 0 0 B 0 0 0 0 ···
            | 0 0 E 0 0 D 0 F 0 0 0 C 0 0 0 0 B 0 0 0 ···
                          8
```

```
o33 : S6-module, quotient of S6
```

The degrees of the ring S6 were chosen so that FV would be homogeneous. We must adjust the degrees of the ring to be all of degree 1, before computing the Tate resolution.

```
i34 : R6 = QQ[A..F];
i35 : TV = coker substitute(presentation TV, R6);
i36 : Ta = tateResolution(presentation TV,E6,-3,4)
                            24
                                   8
                     48
               80
                                            8
                                                    24
      120
                                                             48
         <-- E6 <-- E6
o36 = E6
              -3
                     -2
                                   0
                                           1
                                                   2
                                                            3
     -4
                            -1
o36 : ChainComplex
i37 : betti Ta
```

```
o37 = total: 120 80 48 24 8 8 24 48
          0: 120 80 48 24 8 . . .
          1:
               .
                  •
                     •
                        . .
                  . . . . 8 24 48
          2:
i38 : Ta.dd_1
o38 = | -aef -acf
                       -adf -ace
                                     -ace+abf 0
                                                        -acd
                                                                  abc
       -bef -bcf
                       -bdf -bce+adf -bce+adf abf
                                                        -bcd+ade
                                                                 -acd+abe
       -cef 0
                       -cdf aef
                                     -bcf+aef acf
                                                        adf
                                                                  abf
      1
       adf
                       ade
                                                        abd
                                                                  0
            abf
                            abe
                                     acd
                                              abc
            adf
       bdf
                      bde
                            ade
                                                        0
      1
                                     bcd
                                               -acd
                                                                  -abd
       cdf
             -bcf+aef cde
                            -bce+adf 0
                                              -ace+abf
                                                        -bcd
                                                                  -acd
                      0
                            cde+bdf
                                              adf
       -def cdf
                                     cde
                                                        bde
                                                                  ade
       0
                      def
                           cdf+bef
                                              aef
                                                        bdf
                                                                  adf
      L
             cef
                                     cdf
               8
                        8
o38 : Matrix E6 <--- E6
```

This is the map ψ_{-1} . This matrix is almost the desired 8×8 skew symmetric matrix, except for one problem: it is not skew symmetric! This is because of the choices made by *Macaulay 2* in computing the resolution. By row and column operations over k, it is straightforward to produce the desired skew-symmetric matrix. The following lines of *Macaulay 2* code perform these row and column operations.

```
i39 : load "sparsemat.m2";
i40 : (m = sparseMutableMatrix Ta.dd_1;
          rflip(m,0,7);rflip(m,1,2);rflip(m,6,2);cflip(m,3,6);
          rflip(m,3,4);rscale(m,-1_E6,3);cflip(m,4,5);rflip(m,7,4);
          cflip(m,5,7);rflip(m,7,5);rscale(m,-1_E6,5);caxy(m,-1_E6,7,6);
          rscale(m,-1_E6,7);
          matrix m)
o40 = | 0
                     def bdf
                                            adf
                                                      bef
                                                              cdf
             cef
                                    aef
       -cef O
                     -cdf adf
                                    acf
                                            abf
                                                      bcf
                                                              -bcf+aef
                     0
                                                      bdf
       -def cdf
                          bde
                                    adf
                                            ade
                                                              cde
       -bdf -adf
                     -bde 0
                                            abd
                                                      bcd-ade
                                                              -bcd
                                    acd
       -aef -acf
                     -adf -acd
                                    0
                                            abc
                                                      -abf
                                                              -ace+abf
       -adf -abf
                      -ade -abd
                                            0
                                                      acd-abe
                                                              -acd
                                    -abc
       -bef -bcf
                      -bdf -bcd+ade abf
                                            -acd+abe
                                                      0
                                                              -bce+adf
       -cdf bcf-aef -cde bcd
                                    ace-abf acd
                                                      bce-adf 0
      L
                        8
               8
o40 : Matrix E6 <--- E6
```

Theorem 9 (Eisenbud and Schreyer, 2001). The Macaulay resultant

 $Res_{2,2,2}(F_a, F_b, F_c)$

is the pfaffian of the 8×8 matrix U(m) where m is the skew symmetric 8×8 matrix in the Macaulay 2 code above.

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Appendix A. Macaulay 2 code used in this paper

The code that we use here appeared in Decker and Eisenbud (2001). They also include *Macaulay 2* code to compute the Beilinson monad, which is another interesting and important application of these exterior algebra techniques. All of these make up the file "bgg.m2".

The routine bgg computes the matrix ϕ_i corresponding to multiplication from M_i to M_{i+1} . The "BGG" stands for Bernstein–Gelfand–Gelfand.

```
i41 : code bgg

o41 = -- bgg.m2:10-20

bgg = (i,M,E) ->(

        S :=ring(M);

        numvarsE := rank source vars E;

        ev:=map(E,S,vars E);

        f0:=basis(i,M);

        f1:=basis(i+1,M);

        g :=((vars S)**f0)//f1;

        b:=(ev g)*((transpose vars E)**(ev source f0));

        --correct the degrees (which are otherwise

        --wrong in the transpose)

        map(E^{(rank target b):i+1},E^{(rank source b):i}, b));
```

The routine symExt is a subroutine of tateResolution. The input is a presentation matrix for the module M above, and it is a method to obtain bgg(coker m, 0, E) with less computation.

i42 : code symExt

The routine tateResolution takes as input a presentation matrix for the module M, the corresponding exterior algebra E, and a low and high degree loDeg and hiDeg which determines the part of the Tate resolution to return. If r is the regularity of the module M, the piece that is returned is

 $\mathbf{T}^{\max(r+2,hi\mathrm{Deg})}(M) \longleftarrow \cdots \longleftarrow \mathbf{T}^{lo\mathrm{Deg}}(M).$

We have modified the code slightly from Decker and Eisenbud (2001) by shifting the cohomological degrees of the result, so that if T is the result, then $T_{-}(-d)$ is $T^{d}(M)$.

```
i43 : code tateResolution
o43 = -- bgg.m2:21-32
     tateResolution = (m,E,loDeg,hiDeg)->(
          M := coker m;
          reg := regularity M;
          bnd := max(reg+1,hiDeg-1);
          mt := presentation truncate(bnd,M);
              := symExt(mt,E);
          0
           --adjust degrees, since symExt forgets them
          ofixed
                   := map(E^{(rank target o):bnd+1},
                     E^{(rank source o):bnd},
                      o):
          C := res(coker ofixed, LengthLimit=>max(1,bnd-loDeg+1));
          C[bnd+1])
```

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