A-Matrix Optimization Problem

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ABSTRACT

A minimization problem for a matrix-valued matrix function is considered. A duality theorem is proved. Some examples illustrate its applicability.

Let Σ_n be the set of real symmetric $n \times n$ matrices. In Σ_n a partial ordering—the Löwner ordering—is defined by $A \ge B$ iff A - B is nonnegative definite. Using this ordering, investigations of monotonicity, convexity, and extremum properties of matrix functions with values in Σ_n have recently aroused considerable interest, mainly due to applications in statistics and electrical network theory; see e.g. [7], [5], [2], [8]. In this paper we will derive a primal-dual relation for a quadratic minimization problem under a linear constraint, and we will show that some results from the literature can be subsumed under this problem.

We will use the following terminology: For a real $m \times n$ matrix A we write $A \in \mathbb{R}^{m \times n}$; A^T is the transpose of A, and A^+ the Moore-Penrose inverse of A. I_n is the unit matrix in $\mathbb{R}^{n \times n}$. If $A \in \mathbb{R}^{n \times n}$, $A \ge 0$, $B \in \mathbb{R}^{q \times n}$, then \overline{B}_A denotes a minimum-A-seminorm g-inverse of B, i.e., $\overline{B}_A \in \mathbb{R}^{n \times q}$ satisfies

$$B\overline{B}_{A}B = B$$
 and $(\overline{B}_{A}B)^{T}A = A\overline{B}_{A}B$ (1)

[10, p. 46]. The matrix

$$S = S(A, B) = A\overline{B}_A B \tag{2}$$

is invariant under the choice of \overline{B}_A and satisfies

$$S = B^{T} (\overline{B}_{A})^{T} A \overline{B}_{A} B; \qquad (3a)$$

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in particular

$$S \ge 0$$
 and $A \ge S$ (3b)

(see [8]). S is called the shorted matrix A with respect to B.

Using the shorted operator S for the present problem is suggested by the fact that $\Re(A) \cap \Re(B^T) = \Re(S)$, where $\Re(D)$ denotes the range of the matrix D. This together with $S \ge 0$ allows us to treat the p-dimensional problem, $p \ge 2$, with Löwner ordering in the same way as the case p = 1.

THEOREM. Let $A \in \mathbb{R}^{n \times n}$, $A \ge 0$, $B \in \mathbb{R}^{q \times n}$, $C \in \mathbb{R}^{q \times p}$ be given, and let S be defined by (2). Consider the following pair of primal and dual problems (with respect to the Löwner ordering):

$$g(X) = XAX^{T} = \min$$
subject to
$$BX^{T} = C, \quad X \in \mathbb{R}^{p \times n}$$
(4)

and

$$h(Y) = C^T B^{+T} Y^T + Y B^+ C - Y S^+ Y^T = \max$$
subject to
$$(I_n - SS^+) Y^T = 0, \quad Y \in \mathbb{R}^{p \times n}.$$
(5)

If there exists a feasible solution of (4), then:

(a) There exist feasible solutions \hat{X} and \hat{Y} of (4) and (5), respectively, such that for all feasible X and Y

$$g(X) \ge g(\hat{X}) = C^T B^{+T} S B^+ C = h(\hat{Y}) \ge h(Y).$$

(b) The general solution of (4) is

$$\hat{X} = C^T \overline{B}_A^T + Z (I - (A - S)\overline{A}) (I - B^T \overline{B}_A^T),$$

where Z is arbitrary, $Z \in \mathbb{R}^{p \times n}$, \overline{A} is an arbitrary g-inverse of A, and \overline{B}_A is an arbitrary, but fixed, minimum-A-seminorm g-inverse of B.

(c) The solution of (5) is unique, namely

$$\hat{Y} = C^T B^+ {}^T S.$$

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Proof. Note first that a solution X of $BX^T = C$ exists iff $C = B \overline{B} C$ for a g-inverse \overline{B} of B.

(a): Let X and Y be feasible. Then using (2) and (3)

$$h(Y) = C^{T}B^{+T}SS^{+}Y^{T} + YS^{+}SB^{+}C - YS^{+}Y^{T}$$

= $-(Y^{T} - SB^{+}C)^{T}S^{+}(Y^{T} - SB^{+}C) + C^{T}B^{+T}SB^{+}C$
 $\leq C^{T}B^{+T}SB^{+}C = XB^{T}B^{+T}SB^{+}BX^{T} = XSX^{T}$
 $\leq XSX^{T} + (X^{T} - A^{+}SX^{T})^{T}A(X^{T} - A^{+}SX^{T}) = XAX^{T} = g(X).$

For $\hat{X} = C^T \overline{B}_A^T$ we have $B\hat{X}^T = B\overline{B}_A C = C$, i.e., \hat{X} is feasible; and from (3)

$$\mathbf{g}(\hat{X}) = C^T \overline{B}_A^T A \overline{B}_A C = C^T B^{+T} S B^+ C.$$

Further, $\hat{Y} = C^T B^{+T} S$ is feasible and

$$h(\hat{Y}) = 2C^{T}B^{+T}SB^{+}C - C^{T}B^{+T}SS^{+}SB^{+}C = C^{T}B^{+T}SB^{+}C.$$

(b): If \overline{B}_A is fixed, the general solution of $BX^T = C$ is

$$X = C^T \overline{B}_A^T + U (I - B^T \overline{B}_A^T),$$

where $U \in \mathbb{R}^{p \times n}$ is arbitrary [10, p. 24]. The proof of (a) shows that we have $g(X) = g(\hat{X})$ iff $XAX^T = XSX^T$, i.e. X(A - S) = 0. Using $C = B\bar{B}_AC$, (1), and (2), this implies U(A - S) = 0; hence

$$U = Z(I - (A - S)\overline{(A - S)}),$$

where $Z \in \mathbb{R}^{p \times n}$ is arbitrary. Now from (1)–(3) it follows that \overline{A} is a g-inverse of A - S. The converse is straightforward, and (c) can be proved similarly.

REMARK. The dual problem (5) has been obtained by a Lagrangian-multiplier approach to (4). Therefore, the result admits the following version as a saddle-point theorem: Let the operator $F: \mathbb{R}^{p \times n} \times \mathbb{R}^{p \times q} \to \mathbb{R}^{p \times p}$ be defined by

$$F(X,Y) = XAX^{T} - XB^{T}Y^{T} - YBX^{T} + YC + C^{T}Y^{T}.$$

Then for $\hat{X} = C^T \overline{B}_A^T$, $\hat{Y} = C^T B^+ {}^T S B^+$, and all $X \in \mathbb{R}^{p \times n}$, $Y \in \mathbb{R}^{p \times q}$ it holds that

$$F(\hat{X}, Y) \leqslant F(\hat{X}, \hat{Y}) \leqslant F(X, \hat{Y}).$$

In the case p = 1 this approach has been used in [6] to give a statistical interpretation of \hat{Y} , if the primal problem is that of finding a BLUE in the Gauss-Markoff model.

Let us consider now some special cases of the theorem.

EXAMPLE 1. If A is positive definite, then $\overline{B}_A = A^{-1}B^T(BA^{-1}B^T)^+$ satisfies (1); hence $S = B^T(BA^{-1}B^T)^+ B$. From this one easily obtains that \hat{X} is uniquely determined. If additionally $B = CC^T$, then $\overline{B}_A = A^{-1}C(C^TA^{-1}C)^+C^+$ and

$$\hat{X} = C^+ C (C^T A^{-1} C) C^T A^{-1}.$$

 \hat{X} satisfies the conditions (3.3.5) in [10]. Thus we obtain a characterization of the I_q -norm A^{-1} -least-squares g-inverse of C as solution of a matrix optimization problem. In the special case $A = I_q$ we get with $\hat{X} = C^+$ the well-known characterization of the Moore-Penrose inverse; cf. [9] and [4].

EXAMPLE 2. Let $A_1, A_2 \in \mathbb{R}^{n \times n}$, $A_1 \ge 0$, $A_2 \ge 0$; put

$$A = \begin{bmatrix} A_1 & 0 \\ 0 & A_2 \end{bmatrix}, \qquad B = \begin{bmatrix} I_n, I_n \end{bmatrix}.$$

Then

$$\overline{B}_{A} = \frac{1}{2} \begin{bmatrix} I_{n} + (A_{1} + A_{2})^{+} (A_{1} - A_{2}) \\ I_{n} + (A_{1} + A_{2})^{+} (A_{2} - A_{1}) \end{bmatrix}$$

satisfies the condition (1); hence

$$\mathbf{S} = \begin{bmatrix} D & D \\ D & D \end{bmatrix},$$

where

$$D = A_1 - A_1(A_1 + A_2)^+ A_1 = A_2 - A_2(A_1 + A_2)^+ A_2^-$$
$$= A_1(A_1 + A_2)^+ A_2 = A_2(A_1 + A_2)^+ A_1^-$$

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is the parallel sum of A_1 and A_2 ; cf. [1]. Thus the optimum value of the problem

$$X_1 A_1 X_1^T + X_2 A_2 X_2^T = \min$$
(6)

subject to

$$X_1^T + X_2^T = C, \qquad X_1, X_2 \in \mathbb{R}^{p \times n}$$

is $C^T DC$. A solution is given by

$$\hat{X}_1 = C^T - C^T A_1 (A_1 + A_2)^+, \qquad \hat{X}_2 = C^T A_1 (A_1 + A_2)^+.$$

The solution of (5) is $\hat{Y} = [\hat{Y}_1, \hat{Y}_2] = [C^T D, C^T D]$. This is a generalization of problem (14) in [2]. Moreover, (6) provides a characterization of the parallel sum as the optimum value of (6) when $C = I_n$.

EXAMPLE 3. Considering the problem (4) for the case that B = C, we get $\min\{XAX^T: BX^T = B, X \in \mathbb{R}^{p \times n}\} = B^T B^{+T} S B^+ B = S$. This is a characterization of the shorted matrix S which is different from those given in [2, Theorems 1, 5] and in [8, Theorem 2.2].

EXAMPLE 4. Let p = q = 1, A be positive definite, $B = e_k^T$, where e_k is the kth unit vector in \mathbb{R}^n , C = 1. In this case $\overline{B}_A = A^{-1}e_k(e_k^T A^{-1}e_k)^{-1}$ satisfies (1), and hence

$$\mathbf{S} = \mathbf{S}(\mathbf{A}, \mathbf{e}_k^T) = \left(\mathbf{e}_k^T \mathbf{A}^{-1} \mathbf{e}_k\right)^{-1} \mathbf{e}_k \mathbf{e}_k^T;$$

see also [3, Theorem 7]. If A_k denotes the matrix obtained when deleting the kth row and column of A, and a^{kk} is the kth diagonal element of A^{-1} , then $e_k^T A^{-1} e_k = a^{kk} = \det A_k / \det A$. Thus we get from the theorem Bergström's inequality [7, p. 478],

$$\begin{aligned} \frac{\det(A^{(1)} + A^{(2)})}{\det(A^{(1)}_k + A^{(2)}_k)} &= \min\{x^T (A^{(1)} + A^{(2)}) x : e_k^T x = 1\} \\ &\ge \min\{x^T A^{(1)} x : e_k^T x = 1\} \\ &+ \min\{x^T A^{(2)} x : e_k^T x = 1\} \\ &= \frac{\det A^{(1)}}{\det A^{(1)}_k} + \frac{\det A^{(2)}}{\det A^{(2)}_k} \end{aligned}$$

for any two positive definite matrices $A^{(1)}$, $A^{(2)} \in \mathbb{R}^{n \times n}$. Also from $A \ge S$ we get the well-known relation $a_{kk} \ge (a^{kk})^{-1}$.

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