

Unordered Canonical Ramsey Numbers

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We define a weak form of canonical colouring, based on that of P. Erdős and R. Rado (1950, *J. London Math. Soc.* **25**, 249–255). This yields a class of unordered canonical Ramsey numbers $CR(s, t)$, again related to the canonical Ramsey

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known corresponding bounds for $ER(2, s)$. © 2000 Academic Press

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Erdős and Rado extended Ramsey's original theory of monochromatic graphs [6] to allow infinitely many colours [4]. This extension created the theory of canonical colourings. The Erdős–Rado theorem assumes that there is a pre-existing ordering on the vertices of the graph. It states that, for sufficiently large n , an arbitrary colouring of the edges of K_n contains some K_s which is coloured according to one of the following methods (where $c(e)$ is the colour of the edge e and uv is an edge with $u < v$):

- K_s is monochromatic;
- $c(uv) = f(u)$, where f injective, i.e., minimum-coloured;
- $c(uv) = f(v)$, where f injective, i.e., maximum-coloured; or
- $c(uv) = f(uv)$, where f injective, i.e., K_s is *distinctly edge coloured*.

Let $ER(2; s)$ be the smallest n for which the above holds. The best bounds known for $ER(2; s)$ are those given by Lefmann and Rödl in [5], $2^{c_1 s^2} \leq ER(2; s) \leq 2^{c_2 s^2 \log s}$ for some constants c_1, c_2 . When considering k -uniform hypergraphs, the best upper bounds for $ER(k; s)$ for general fixed k are due to Shelah [7]. The upper bounds obtained by Shelah are towers of height k , the same height as the towers of the lower bounds of the corresponding Ramsey numbers, and therefore in a sense these bounds are best possible.

Here we present a weaker definition of canonical colourings, which does not require a fixed ordering on the vertices of our graph. In our definition, we combine the first three categories into a single type. Take some $G = K_s$ for arbitrary s . Colour the edges of G with arbitrary colours. We call such a colouring *orderable* if we can order $V(G)$ such that the colour of each edge is completely determined by the smaller of its two vertices. That is, if $c(e)$ denotes the colour of the edge e , and uv is an edge with $u < v$, then $c(uv) = f(u)$. Note that this definition does not require that f be injective, and hence monochromatic graphs are orderable. Also note that the minimum- and maximum-colourings specified by Erdős and Rado are orderable, as we can select either the given ordering of the vertices (to satisfy a minimum-colouring) or reverse ordering (to satisfy a maximum-colouring).

Let $CR(s, t)$ be the smallest N such that every colouring of the edges of K_N contains either an orderable K_s or a distinctly edge coloured K_t . Certainly $CR(s, t) \leq ER(2; \max(s, t))$, and so in particular $CR(s, t)$ exists. In this paper we prove both an upper bound (in Theorem 0.1) and lower bound (Theorem 0.2) for $CR(s, t)$, having the same order of magnitude.

The notion of orderable colourings was stimulated originally by the proof of Ramsey's theorem given by Erdős and Szekeres [2]. Restricting ourselves to two colours, we have an analogue of the standard Ramsey numbers, rather than the Erdős–Rado numbers. Denote by $\rho(n)$ the smallest k such that any 2-colouring of K_k must contain an orderable K_n . Now a monochromatic K_n is orderable, and conversely any 2-coloured orderable K_{2n-2} must contain a monochromatic K_n . (This is why we do not require f to be injective in our definition of orderable graphs.) Thus $\rho(n) \leq r(n) \leq \rho(2n-2)$, and so the standard bounds on $r(n)$ [2, 3] show that $\rho(n)$ grows exponentially in n . Moreover, the argument of Erdős and Szekeres [2] shows $\rho(n+1) \leq 2\rho(n)$, and since it is easily checked that $\rho(3) = 3$ and $\rho(4) = 6$ we obtain the upper bound $\rho(n) \leq 3 \cdot 2^{n-3}$.

We now return to our main purpose, dealing with arbitrarily many colours and thereby obtaining bounds on $CR(s, t)$. We begin with an upper bound.

THEOREM 0.1. $CR(s, t) \leq 7^{3-s} t^{4s-4}$.

Proof. Let $n = CR(s, t) - 1$. Then there exists some colouring G of K_n such that G contains no orderable K_s or distinctly edge coloured K_t . Hence every $K_t \subset G$ contains a pair of like-coloured edges. We say that a pair of edges of the same colour which lie in the same K_t block that K_t (as they prevent it from being distinctly edge coloured). As each K_t is blocked by at least two edges, some edge e blocks at least $2\binom{n}{t}/\binom{n}{2}$ distinct K_t .

Two disjoint edges of like colour block $\binom{n-4}{t-4}$ distinct K_t , while two adjacent edges of like colour block $\binom{n-3}{t-3}$ distinct K_t . If we denote by k the

number of edges incident with e of the same colour as e , and by l the number of other edges of the same colour as e , then

$$l \binom{n-4}{t-4} + k \binom{n-3}{t-3} \geq 2 \frac{\binom{n}{t}}{\binom{n}{2}}$$

So there must be a large number of edges of the same colour as e . Note, however, that G is also free of orderable K_s graphs, and this restricts the number of edges of any single colour. Let $c = CR(s-1, t)$. If some vertex $v \in G$ had degree at least c in one colour, then either G contains a K_t distinctly edge coloured, or the neighbourhood of v in that colour contains an orderable K_{s-1} , and hence G contains an orderable K_s . Hence each vertex incident with e has at most $c-2$ edges of the same colour as e besides e itself. Hence $k \leq 2c-4$.

Similarly, none of the $n-2$ vertices disjoint from e can have c incident edges of the same colour as e . The maximum number of edges coloured similarly to e but disjoint from e is therefore $\frac{1}{2}(c-1)(n-2)$. Hence

$$\binom{n-4}{t-4} \binom{n-2}{2} (c-1) + (2c-4) \binom{n-3}{t-3} \geq 2 \frac{\binom{n}{t}}{\binom{n}{2}}$$

So, writing $t_{(4)}$ for $t(t-1)(t-2)(t-3)$, we have

$$\frac{n-2}{2(n-3)} (c-1) + \frac{2}{t-3} (c-2) \geq \frac{4(n-2)}{t_{(4)}}.$$

Note that $CR(s, t) = n+1 \geq c$. The above inequality therefore holds whenever

$$\frac{c}{2} + \frac{3}{2(n-3)} + \frac{2}{t-3} (c-2) \geq \frac{4(n-2)}{t_{(4)}}$$

and so

$$\begin{aligned} CR(s, t) &\leq 3 + \frac{t_{(4)}}{8} \left(c + \frac{3}{CR(s, t) - 4} \right) + \frac{t_{(3)}(c-2)}{2} \\ &\leq \frac{t^4}{8} \left(c + \frac{3}{CR(s, t) - 4} \right) \end{aligned}$$

whenever $n \geq 4$, $t \geq 4$.

If $CR(s, t) < 10$, then as $t \geq 2$ $CR(s, t) \leq 3t^4/2$. Otherwise, if $CR(s-1, t) < 6$, then from the equation above we have that

$$CR(s, t) \leq \frac{t^4}{8} \left(5 + \frac{3}{CR(s, t) - 4} \right),$$

and as $CR(s, t) \geq 10$ this means $CR(s, t) \leq 3t^4/2$. If neither of these bounds holds, then we instead bound $CR(s, t)$ above by

$$\frac{t^4}{8} \left(CR(s-1, t) + \frac{1}{2} \right),$$

so

$$CR(s, t) \leq \frac{t^4}{7} CR(s-1, t).$$

Note also that $CR(2, t) = 2$ for every $t \in \mathbb{N}$ (as any single edge is an orderable K_2). Hence we can obtain an upper bound for any $CR(s, t)$ by using the third bound,

$$CR(s, t) \leq \frac{t^4}{7} CR(s-1, t),$$

at most $s-2$ times in succession, until either $CR(s, t) < 10$ or $CR(s-1, t) < 6$, both of which then yield a fixed bound of at most $3t^4/2$.

So, for all $s \geq 2, t \geq 4$,

$$CR(s, t) \leq \frac{3t^4}{2} \frac{t^{4s-8}}{7^{s-2}} = \frac{3}{2} 7^{2-s} t^{4s-4} \leq 7^{3-s} t^{4s-4}. \blacksquare$$

We now consider lower bounds for $CR(s, t)$. Like the original Ramsey numbers we can obtain a lower bound for $CR(s, t)$ probabilistically. Random colourings using $\binom{t}{2} - 1$ colours give $CR(s, t) \geq ((\binom{t}{2} - 1)^{(s-1)(s-2)/2s})$. But, unlike the original Ramsey numbers, for the unordered canonical Ramsey numbers $CR(s, t)$ we can give a lower bound by a simple construction, which is better than the probabilistic bound.

THEOREM 0.2. $CR(s, t) \geq ((\binom{t}{2} - 1)(CR(s-1, t) - 1) + 1$.

Proof. Let G be a 2-coloured graph of order $CR(s-1, t) - 1$. Take $\binom{t}{2} - 1$ copies of G , labelled

$$G_0, \dots, G_{\binom{t}{2}-2}.$$

Let $c_k, 0 \leq k \leq \binom{t}{2} - 2$ be distinct colours not used in G . Colour all edges between any pair of copies G_i and G_j with colour c_l , where $l \equiv i + j \pmod{(\binom{t}{2} - 1)}$.

As a result, the colours connecting each G_j to the $\binom{t}{2} - 2$ other copies of G are distinct for every j . Call the resultant complete graph H . We claim that H contains no distinctly edge coloured K_t , and no orderable K_s .

Suppose that there exists $H' \subset H$, distinctly edge coloured and of order t . We know that $H' \not\subset G_k$ for every $1 \leq k \leq \binom{t}{2} - 1$, as each G_k is free of distinctly edge coloured K_t 's. Alternatively, if no two vertices of H' lie in the same $G_k \subset H$, then all edges in H' have colours c_k for some $0 \leq k \leq \binom{t}{2} - 2$, and hence at least two edges in H' must have the same colour. So there must be some $G_k \subset H$ such that $H' \not\subset G_k$, but $|H' \cap G_k| \geq 2$. Then there exists G_l , $l \neq k$, such that at least one vertex of H' lies in G_l . But then H' contains two edges between G_k and G_l , which must be of the same colour. We therefore know that there is no distinctly edge coloured K_t inside H .

We now show that H contains no orderable K_s . For suppose such a K_s exists. G contains no orderable K_{s-1} , hence any orderable $K_s \subset H$ must contain either (i) vertices from three different copies of G , or (ii) at least two vertices from each of two distinct copies of G . In case (i) the three vertices from distinct copies will form a triangle in which each edge is a distinct colour, which prevents the K_s from being orderable. In case (ii), regardless of how we order the vertices, the first vertex v_1 has at least one neighbour in the same copy of G , and one neighbour in a different copy of G . These two vertices must be connected to v_1 , by edges of different colours, ensuring that G is not orderable. Hence $CR(s, t) \geq |H| + 1 = ((\binom{t}{2} - 1)(CR(s-1, t) - 1) + 1$. ■

As $CR(2, t) = 2$ for every $t \geq 2$, we can use Theorem 0.2 recursively to obtain the following general lower bound.

COROLLARY 0.1. $CR(s, t) \geq ((\binom{t}{2} - 1)^{s-2} + 1$.

Further, a theorem of Babai [1] gives us a way to improve the exponent in this lower bound, but in a non-constructive way. Babai proved that there exists an edge-coloured complete graph of order $\Theta(t^3/\log t)$ in which no two edges of the same colour are adjacent, but which does not contain any distinctly edge-coloured K_t . So by using the nesting argument of Theorem 0.2 with graphs of order $\Theta(t^3/\log t)$, rather than order $\binom{t}{2} - 1$ as used in Theorem 0.2, we obtain a graph of order $\Theta(t^{3s}/(\log t)^s)$. This graph is free of distinctly edge coloured K_t by [1], and is so free of orderable graphs of order s , hence $CR(s, t) = \Omega(t^{3s}/(\log t)^s)$.

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