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# Completeness theory for the product of finite partial algebras

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## Abstract

A general completeness criterion for the finite product  $\prod \mathbb{P}(k_i)$  of full partial clones  $\mathbb{P}(k_i)$  (composition-closed subsets of partial operations) defined on finite sets  $E(k_i)$  ( $|E(k_i)| \geq 2$ ,  $i = 1, \dots, n$ ,  $n \geq 2$ ) is considered and a Galois connection between the lattice of subclones of  $\prod \mathbb{P}(k_i)$ , called partial  $n$ -clones, and the lattice of subalgebras of multiple-base invariant relation algebra, with operations of a restricted quantifier free calculus, is established. This is used to obtain the full description of all maximal partial  $n$ -clones via multiple-base invariant relations and, thus, to solve the general completeness problem in  $\prod \mathbb{P}(k_i)$ .

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## 1. Introduction and basic definitions

Let  $k \geq 2$  be an integer and  $E(k) = \{0, 1, \dots, k - 1\}$ . For an integer  $m \geq 1$  an  $m$ -ary partial operation  $f$  on  $E(k)$  (an  $m$ -ary partial function of  $k$ -valued logic) is a one-to-one map from a subset  $D_f = \text{Dom}(f)$  of  $E^m(k)$  (called the domain of  $f$ ) into  $E(k)$ ,  $f : D_f \rightarrow E(k)$ . Denote  $P^m(k)$  the set of all partial  $m$ -ary operations on  $E(k)$  including the empty operation  $p_m$  having an empty domain. Set  $P(k) = \bigcup_{m \geq 1} P^m(k)$ .

The notion of a composition of partial operations from  $P(k)$  is formally equivalent to the operations of iterative Post algebra  $\mathbb{P}(k) = \langle P(k); \zeta, \tau, \Delta, *, e_1^2 \rangle$  (see [11]), where  $e_1^2(x_1, x_2) = x_1$  is a binary selector (projection) and for any  $n > 1$  and  $f \in P^n(k)$  we

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have

$$(\zeta f)(x_1, x_2, x_3, \dots, x_n) = f(x_2, x_3, \dots, x_n, x_1),$$

$$(\tau f)(x_1, x_2, x_3, \dots, x_n) = f(x_2, x_1, x_3, \dots, x_n),$$

$$(\Delta f)(x_1, x_2, x_3, \dots, x_n) = f(x_1, x_1, x_3, \dots, x_{n-1}),$$

where the left sides of identities are defined whenever the right sides are defined. For  $n = 1$  we put  $\zeta f = \tau f = \Delta f = f$ .

Next for  $f \in P^n(k)$  and  $g \in P^m(k)$  ( $n, m \geq 1$ ) we set

$$(f * g)(x_1, \dots, x_{m+n-1}) = f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{m+n-1}),$$

where again the left side is defined whenever the right side is defined.

In universal algebra terminology  $\mathbb{P}(k)$  is called the *full partial clone* [7] and each subalgebra of it is called a *partial clone* on  $E(k)$ . A set  $S$  of partial operations is *complete* in  $\mathbb{P}(k)$  when it is a generating set in  $P(k)$  with respect to operations of the iterative Post algebra (or, equivalently, with respect to any compositions of partial operations). A general completeness criterion establishes the necessary and sufficient conditions for a given set  $S \subset P(k)$  to be complete. Since  $\mathbb{P}(k)$  is finitely generated this criterion is known (see, e.g., [2] or [4]) to be based on the knowledge of the full list of all maximal subalgebras of  $\mathbb{P}(k)$  or *maximal partial clones* on  $E(k)$  ( $k \geq 2$ ).

For  $k = 2$  this problem was introduced and solved by Freivald [3,4] who listed all 8 maximal partial clones on  $E(2)$ . The case  $k \geq 3$  was considered in [15], where the list of maximal partial clones on  $E(3)$  was presented (3 clones were inadvertently omitted, see [6,20]),<sup>2</sup> and the Slupecki-type criterion for  $k \geq 3$  was given (completeness with all unary partial operations), as well as some series of maximal partial clones on  $E(k)$ ,  $k \geq 4$ , were found. The full description of all maximal partial clones on  $E(k)$ ,  $k \geq 4$ , was provided independently by Lo Czukai [9,10] (see also comments on these results in [20]), Haddad and Rosenberg [5,7] and the author [20]. All of the variants of a final solution were grounded on the fact [15] that, with one exception, each maximal partial clone is determined by a relation of arity less or equal  $k$  defined on the same set  $E(k)$ ,  $k \geq 4$ .

**Remark.** In the case of an infinite base set  $E$  the general completeness criterion cannot be formulated entirely in terms of maximal partial clones (see, e.g., [16,24]), although the knowledge of these clones is still of a great importance. We'll mention only three results in this field: (1) Slupecki-type criterion for local completeness in  $P(E)$  [17]; (2) the full description of all maximal local partial clones [22]; (3) the full description of maximal partial clones which can be determined by a finite arity relation on  $E$  [24].

<sup>2</sup> The list of all 58 maximal partial clones on  $E(3)$  was also presented in the thesis: D. Lau, "Eigenschaften gewisser abgeschlossener Klassen in Postschen Algebren", University Rostock, GDR, 1977.

In this paper we consider the completeness problem for vectors of partial operations defined on finite sets. For integers  $k_1, \dots, k_n$  greater than 1 and  $m \geq 1$  consider the set:

$$A(m) = P^m(k_1) \times \dots \times P^m(k_n) \tag{1}$$

of all  $n$ -vectors ( $n \geq 2$ ) of partial  $m$ -ary operations defined on the sets  $E(k_1), \dots, E(k_n)$  resp. Denote  $\mathbf{e}_1^2 = \langle e_1^2(x, y), \dots, e_1^2(x, y) \rangle \in A(2)$  the  $n$ -vector produced from the projection  $e_1^2(x, y) = x$ . We introduce the arity-calibrated product of full partial clones as follows:

$$\begin{aligned} \prod \mathbb{P}(k_i) &:= \prod_{i=1}^n \mathbb{P}(k_i) = \mathbb{P}(k_1) \times \dots \times \mathbb{P}(k_n) \\ &= \left\langle \bigcup_{m \geq 1} A(m); \zeta, \tau, \Delta, *, \mathbf{e}_1^2 \right\rangle, \end{aligned} \tag{2}$$

where the operations  $\zeta, \tau, \Delta$ , and  $*$  are applied coordinatewise.

So if  $\mathbf{f} = \langle f_1, \dots, f_n \rangle \in A(m)$  and  $\mathbf{g} = \langle g_1, \dots, g_n \rangle \in A(s)$  ( $m, s \geq 1$ ), then  $\mathbf{f} * \mathbf{g} = \langle f_1 * g_1, \dots, f_n * g_n \rangle$  and  $\varepsilon \mathbf{f} = \langle \varepsilon f_1, \dots, \varepsilon f_n \rangle$ , where  $\varepsilon \in \{\zeta, \tau, \Delta\}$ . The  $n$ -vector  $\mathbf{e}_1^2$  is a constant operation. This formalism represents all compositions of  $n$ -vectors of partial algebraic operations. The product  $\prod \mathbb{P}(k_i)$  is called the *full partial  $n$ -clone*. Any its subalgebra is called a *partial  $n$ -clone*, which is exactly a subdirect product of  $n$  partial clones defined on the sets  $E(k_i)$  ( $i = 1, \dots, n$ ). Next a partial  $n$ -clone is called *maximal* if there is no partial  $n$ -clone, other than the full  $n$ -clone, covering it.

Similarly to its factors  $\mathbb{P}(k_i)$  ( $i = 1, \dots, n$ ) the full partial  $n$ -clone  $\prod \mathbb{P}(k_i)$  is finitely generated (e.g. it is easy to verify that  $A(2)$  is a finite generating set in it). Hence, from the common algebraic results (see [2]) it follows that each proper partial  $n$ -clone is contained in a maximal partial  $n$ -clone and, therefore, a set  $S$  is complete in  $\prod \mathbb{P}(k_i)$  if and only if it is not contained in any maximal partial  $n$ -clone. So the description of all maximal partial  $n$ -clones (dual atoms in the lattice of all partial  $n$ -clones) provides the solution of the general completeness problem in  $\prod \mathbb{P}(k_i)$ .

We will explore the properties of the lattice of partial  $n$ -clones via multiple-base invariant relations defined on the same base sets  $E(k_i)$  ( $i = 1, \dots, n$ ), similar to the case of products of the full clones of everywhere defined operations  $\mathbb{Q}(k_1) \times \dots \times \mathbb{Q}(k_n)$  (see e.g., [14,18,19,21]), where  $\mathbb{Q}(k) = \langle Q(k); \zeta, \tau, \Delta, *, \mathbf{e}_1^2 \rangle$  is the full clone of algebraic operations and  $Q(k)$  is the set of all everywhere defined operations on  $E(k)$  ( $k \geq 2$ ).

We will follow a traditional way (see [1,14,16]) in providing the relational description of dual atoms in the lattice of partial  $n$ -clones. First we establish a Galois connection between the lattice of partial  $n$ -clones closed under all restrictions of their elements and the lattice of multiple-base relations sets closed under the formation of  $(\&, =_{1, \dots, n})$ -formulas of the restricted quantifier free first order calculus. Then we prove that each maximal partial  $n$ -clone, with  $n$  exceptions, is determined by a multiple-base relation, which is minimal under the expressibility by these formulas. Next starting with the Slupecki criterion we find all those multiple-base relations for the general case  $\mathbb{P}(k_1) \times \dots \times \mathbb{P}(k_n)$  using predicative descriptions and also combinatorial considerations as well as for the case  $\mathbb{P}(2) \times \dots \times \mathbb{P}(2)$  which requires only predicative descriptions of relations. The short version of these results, without proofs, was published in [26].

## 2. Multiple-base relations

We consider multiple-base relations on  $n$  base sets  $E(k_1), \dots, E(k_n)$  ( $n \geq 1$ ), each of them corresponds to its own *sort* of variables from the set  $I = \{1, \dots, n\}$ . In what follows we denote  $x^i$  or  $y^j$  variables of  $i$ th sort in both function and relation taking on values from  $E(k_i)$  ( $i = 1, \dots, n$ ). Let  $m_1, \dots, m_n$  be nonnegative integers. A *multiple-base* relation  $R(x_1^1, \dots, x_{m_1}^1, x_1^2, \dots, x_{m_2}^2, \dots, x_1^n, \dots, x_{m_n}^n)$  of arity  $(m_1, \dots, m_n)$  is a relation with  $m_i$  coordinates from the set  $E(k_i)$ , where  $m_i \geq 0$  ( $i = 1, \dots, n$ ). In case  $m_j > 0$ , while  $m_i = 0$  for all  $1 \leq i \leq n$ ,  $i \neq j$ , we identify this relation with an ordinary single-base relation on the set  $E(k_j)$ . The set  $J(R)$  of all indices  $j$  for which  $m_j > 0$  is called *type* of  $R$ ,  $J(R) \subseteq I$ .

**Example 2.1.** Let  $n = 3$  and  $k_i = 2$  ( $i = 1, 2, 3$ ). Then  $R \equiv (x_1^1 = x_2^1) \& (x_1^2 = x_2^2)$ , where  $\&$  is a conjunction of multi-sorted predicates, is a multiple-base relation of arity  $(2, 2, 0)$  and type  $J(R) = \{1, 2\}$ . Notice that in order to present  $R$  as a set of  $(2, 2)$ -tuples one has to distinguish each base set from the others. Namely, one way is to put semicolon to separate coordinates of different sorts. So we have  $R = \{(0, 0; 0, 0), (0, 0; 1, 1), (1, 1; 0, 0), (1, 1; 1, 1)\}$ . Another way [14] is to assume that all  $E(k_i)$  ( $i = 1, \dots, n$ ) are distinct pairwise disjoint sets (this assumption in no way affects further results). So we may rewrite  $R = \{(0, 0, a, a), (1, 1, a, a), (0, 0, b, b), (1, 1, b, b)\}$ , where  $E(k_1) = \{0, 1\}$ ,  $E(k_2) = \{a, b\}$ . In the sequel we will use (whenever it is possible) different letters for variables from different sorts, so we may put in our case  $R(x_1, x_2, y_1, y_2) \equiv (x_1 = x_2) \& (y_1 = y_2)$ .

**Definition 2.1.** A vector of partial operations  $\mathbf{f} = \langle f_1(x_1, \dots, x_m), f_2(y_1, \dots, y_m), \dots, f_n(z_1, \dots, z_m) \rangle$  ( $m \geq 1$ ) preserves a multiple-base relation  $R(x_1, \dots, x_k, y_1, \dots, y_p, \dots, z_1, \dots, z_s)$  of arity  $(k, p, \dots, s)$  if

$$\begin{aligned}
 &R(x_{11}, \dots, x_{1k}, y_{11}, \dots, y_{1p}, \dots, z_{11}, \dots, z_{1s}) \& \dots \\
 &\& R(x_{m1}, \dots, x_{mk}, y_{m1}, \dots, y_{mp}, \dots, z_{m1}, \dots, z_{ms}) \\
 &\& f_1(x_{11}, \dots, x_{m1}) = x_1 \& \dots \\
 &\& f_1(x_{1k}, \dots, x_{mk}) = x_k \& f_2(y_{11}, \dots, y_{m1}) = y_1 \& \dots \\
 &\& f_2(y_{1p}, \dots, y_{mp}) = y_p \& \dots \\
 &\& f_n(z_{11}, \dots, z_{m1}) = z_1 \& \dots \\
 &\& f_n(z_{1s}, \dots, z_{ms}) = z_s \rightarrow R(x_1, \dots, x_k, y_1, \dots, y_p, \dots, z_1, \dots, z_s) \quad (3)
 \end{aligned}$$

holds for all values of all sorts of variables  $x, y, \dots, z$  involved.

Notice that a predicate  $f(x_1, \dots, x_m) = x$  ( $f \in P^m(k)$ ) is valid in (3) whenever  $f(x_1, \dots, x_m)$  is defined and equals  $x$ . Hence each  $\mathbf{f}$  that contains a void (empty) function as its coordinate preserves any relation  $R$ . Denote  $F = \bigcup_{m \geq 1} \{ \langle f_1, \dots, f_n \rangle \in A(m) : \exists i \in \{1, \dots, n\} f_i = p_m \}$  the set of all vector-functions having at least one empty coordinate.

Definition 2.1 can be interpreted in terms of constructing of all possible  $m \times (k + p + \dots + s)$  matrices over the sets  $E(k_1), \dots, E(k_n)$  with rows that are tuples from  $R$  and then applying  $\mathbf{f}$  coordinatewise to these matrices according to each sort of variables. Namely,  $f_1$  is applying to  $k$  coordinates of the 1st sort,  $\dots$ ,  $f_n$  is applying to  $s$  coordinates of the  $n$ th sort. Finally, if the result of each application of  $\mathbf{f}$  to any matrix constructed above (while existed) is also a tuple of  $R$ , then  $\mathbf{f}$  preserves  $R$ .

For everywhere defined vector-operations from  $\mathbb{Q}(k_1) \times \dots \times \mathbb{Q}(k_n)$ , the expression (1) coincides with the definition given in [14,19]. If  $n = 1$ , then we obtain partial operations and relations on  $E(k)$ ,  $k \geq 2$  (see, e.g., [16]). And, finally, for  $f \in Q(k)$  we get the conventional definition of an algebraic operation preserving a relation on the same set  $E(k)$ .

Let  $\text{Pol}(R) = \{\mathbf{f} \in \prod \mathbb{P}(k_i) : \mathbf{f} \text{ preserves } R\}$  and  $\text{Pol}'(R) = \{\mathbf{f} \in \prod \mathbb{Q}(k_i) : \text{preserves } R\}$ . Clearly  $\text{Pol}(R)(\text{Pol}'(R))$  is a partial  $n$ -clone ( $n$ -clone, respectively) and  $F \subset \text{Pol}(R)$ . Set  $\text{Pol}(\mathfrak{R}) = \bigcap_{R \in \mathfrak{R}} \text{Pol}(R)$  for any set  $\mathfrak{R}$  of multiple-base relations.

**Example 2.2.** Let  $R$  be the relation of Example 2.1. Then it is easy to verify that  $\text{Pol}(R) = \prod \mathbb{P}(k_i)$  and also  $\text{Pol}'(R) = \prod \mathbb{Q}(k_i)$  for any  $n \geq 2$ .

Let  $\mathbf{f}, \mathbf{g} \in A(m)$  ( $m \geq 1$ ) be such that  $\text{Dom}(g_i) \subseteq \text{Dom}(f_i)$  and  $g_i = f_i|_{\text{Dom}(g_i)}$  ( $i = 1, \dots, n$ ). We call  $\mathbf{g}$  a restriction of  $\mathbf{f}$  and in turn  $\mathbf{f}$  is called an extension of  $\mathbf{g}$ . Clearly if  $\mathbf{f}$  preserves  $R$ , then  $\mathbf{g}$  also preserves  $R$  and so each partial  $n$ -clone  $\text{Pol}(\mathfrak{R})$  is restriction-closed. The converse is also true.

**Proposition 2.1.** Any partial  $n$ -clone can be presented in the form  $\text{Pol}(\mathfrak{R})$  if and only if it is restriction-closed and also contains  $F$ .

**Proof.** Let  $\mathbf{A}$  be a restriction-closed partial  $n$ -clone and  $F \subset \mathbf{A}$ . Similar to the case  $n = 1$  (see [16]) we introduce  $m$ -graphs of  $\mathbf{A}$  ( $m = 1, 2, \dots$ ) as follows: for each set  $D \subseteq E^m(k_1) \cup \dots \cup E^m(k_n)$ ,  $D \neq \emptyset$ ,  $1 \leq |D| \leq k_1^m + \dots + k_n^m$ , which is considered as  $m$  multiple-base tuples  $r_1, \dots, r_m$  of the same arity  $(s_1, \dots, s_n)$  ( $0 \leq s_i \leq k_i^m$ ,  $i = 1, \dots, n$ ) and presented as a  $m \times (s_1 + \dots + s_n)$  matrix  $D = [r_1, \dots, r_m]$  over  $E(k_1), \dots, E(k_n)$ , we define the relation of arity  $(s_1, \dots, s_n)$ :

$$G_m(\mathbf{A}, D) = \{r : \mathbf{f}(r_1, \dots, r_m) = r \text{ for some } \mathbf{f} \in \mathbf{A} \text{ of arity } m \geq 1\}, \tag{4}$$

where  $\mathbf{f}(r_1, \dots, r_m)$  is a  $(s_1, \dots, s_n)$ -tuple obtained by column-wise application of  $\mathbf{f}$  to  $[r_1, \dots, r_m]$ .

Notice that in this case we have  $D \subseteq \text{Dom}(f_1) \cup \dots \cup \text{Dom}(f_n)$ . Then we introduce the set of relations  $\mathbf{G} = \{G_m(\mathbf{A}, D) : \text{for all non-void subsets } D \text{ and } m \geq 1\}$ . Next we prove:

$$\mathbf{A} = \text{Pol}(\mathbf{G}). \tag{5}$$

It is easy to verify that  $\mathbf{A}$  preserves each relation (4) and so we have  $\mathbf{A} \subseteq \text{Pol}(\mathbf{G})$ . Now assume that there exists  $\mathbf{f} \in \text{Pol}(\mathbf{G}) \setminus \mathbf{A}$  of arity  $m \geq 1$ . Consider  $G_m(\mathbf{A}, D)$ , where

$D = \text{Dom}(f_1) \cup \dots \cup \text{Dom}(f_n)$ . Then by (4) we have  $\mathbf{f}(r_1, \dots, r_m) = r \notin G_m(\mathbf{A}, D)$  (otherwise  $\mathbf{f} \in \mathbf{A}$ ). Hence  $\mathbf{f}$  does not preserve this relation. On the other hand,  $\mathbf{f}$  preserves each relation from  $\mathbf{G}$ . This contradiction proves (5).  $\square$

For any nonempty system  $\mathbf{A}$  of partial  $n$ -operations let  $\text{Inv}(\mathbf{A})$  be the set of all multiple-base relations that are preserved by each element of  $\mathbf{A}$ :  $\text{Inv}(\mathbf{A}) = \{R: \mathbf{A} \subseteq \text{Pol}(R)\}$ . The functors  $\text{Pol}$  and  $\text{Inv}$  establish the *Galois connection* (see, e.g., [1]) between the sets of partial  $n$ -operations and multiple-base relations. The sets having the form  $\text{Pol}(\mathfrak{R})$  and  $\text{Inv}(\mathbf{A})$  are called *Galois-closed* and consequently  $\text{Pol}(\text{Inv}(\mathbf{A}))$  ( $\text{Inv}(\text{Pol}(\mathfrak{R}))$ ) is called the *Galois closure* on sets of partial  $n$ -operations (sets of  $n$ -base relations, respectively).

Notice that Proposition 2.1 gives us the description of Galois-closed sets on the side of partial  $n$ -operations. In order to produce similar description on another side we consider some operations on  $n$ -base relations. Let  $=_i$  be the equality relation on  $E(k_i)$  ( $i=1, \dots, n$ ). We introduce  $(\&, =_1, \dots, =_n)$ -formulas of the restricted multi-sorted first order calculus over the set of relations  $\mathfrak{R}$  which are constructed by the operation  $\&$  from  $=_i$  ( $i=1, \dots, n$ ) and the symbols of relations from  $\mathfrak{R}$  with arbitrary permutations and identifications of variables. Operations  $\pi_i$  ( $i=1, \dots, n$ ), peculiar to the case of partial  $n$ -operations, are used to obtain relations of the smaller type. Namely, if  $R$  can be presented in the form  $(x^i = x^i) \& R'$ , then  $\pi_i(R) = R'$ , otherwise  $\pi_i(R) = R$  ( $i=1, \dots, n$ ).

**Example 2.3.** If  $\mathfrak{R}$  is the empty set, then applying  $\&$ -formulas to  $=_i$  ( $i=1, \dots, n$ ) we obtain *multiple-base diagonals* [14], which can be presented in the form  $D = D_1 \& \dots \& D_n$ , where each  $D_i$  is a single-base diagonal on  $E(k_i)$  constructed by a  $\&$ -formula from  $=_i$  ( $i=1, \dots, n$ ). Denote  $\mathbf{D}$  the set of all  $n$ -base diagonals including empty relations. It is easy to check that  $\text{Pol}(D) = \prod \mathbb{P}(k_i)$  and also  $\text{Pol}'(D) = \prod \mathbb{Q}(k_i)$  for any  $D \in \mathbf{D}$  ( $n \geq 2$ ).

Clearly  $\text{Pol}(R) = \text{Pol}(\pi_i R)$  ( $i=1, \dots, n$ ), and if a relation  $Q$  is constructed by some  $(\&, =_1, \dots, =_n)$ -formula from  $\mathfrak{R}$ , then  $\text{Pol}(\mathfrak{R}) \subseteq \text{Pol}(Q)$ . Applying antimonotone property of the functor  $\text{Inv}$  we obtain  $\text{Inv}(\text{Pol}(\mathfrak{R})) \supseteq \text{Inv}(\text{Pol}(Q))$ , which with  $Q \in \text{Inv}(\text{Pol}(\mathfrak{R}))$ , gives us  $Q \in \text{Inv}(\text{Pol}(\mathfrak{R}))$ . Thus, we proved the property:

Any set of the form  $\text{Inv}(\mathbf{A})$ ,  $\mathbf{A} \subseteq \prod \mathbb{P}(k_i)$ , is closed under application of  $(\&, =_1, \dots, =_n)$ -formulas and also operations  $\pi_i$  ( $i=1, \dots, n$ ).

The converse is also true, and in this way we obtain the characteristics of Galois-closed sets of multiple-base relations.

**Theorem 2.1.** *Any system of  $n$ -base relations has the form  $\text{Inv}(\mathbf{A})$ ,  $\mathbf{A} \subseteq \prod \mathbb{P}(k_i)$ , if and only if it is closed under formation of  $(\&, =_1, \dots, =_n)$ -formulas and application of  $\pi_i$  ( $i=1, \dots, n$ ).*

**Proof.** ( $\Rightarrow$ ) See the property from the above.

( $\Leftarrow$ ) Without the loss of generality we consider  $n=2$ . The common case can be obtained by using the same technique.

**Lemma 2.1.** *Let  $\mathfrak{R}$  be a set of 2-base relations which is closed under formation of  $(\&, =_1, =_2)$ -formulas and  $\pi_i$  ( $i=1, 2$ ). Then for every  $R \in \text{Inv}(\text{Pol}(\mathfrak{R}))$  we have  $R \in \mathfrak{R}$ .*

**Proof.** Clearly  $\mathbf{D} \subseteq \mathfrak{R}$ . Let  $R(x_1, \dots, x_s, y_1, \dots, y_m)$ ,  $R \in \text{Inv}(\text{Pol}(\mathfrak{R}))$ , be a 2-base non-diagonal relation of arity  $(s, m)$ ,  $s, m \geq 1$ . Consider the set  $N = \{Q_1, \dots, Q_t\}$  of all 2-base relations  $Q_i$  from  $\mathfrak{R}$  such that  $R \subseteq Q_i$  ( $i=1, \dots, t$ ) (inclusion of 2-base relations as sets of  $(s, m)$ -tuples). It is obvious that this set is non-void, since it contains at least the full relation of arity  $(s, m)$ . Then we construct the relation  $T$  of arity  $(s, m)$ :

$$T(x_1, \dots, x_s, y_1, \dots, y_m) \equiv \&_{i=1}^t Q_i(x_1, \dots, x_s, y_1, \dots, y_m). \tag{6}$$

Since  $T$  itself is constructed by a  $(\&, =_1, =_2)$ -formula we have  $T \in \mathfrak{R}$  and, therefore,  $R \subseteq T$ . Our goal is to show that  $R \equiv T$  which proves the lemma.

Let  $R \subset T$  (strict inclusion) and  $R = \{r_1, \dots, r_n\}$  be presented as a set of  $n$   $(s, m)$ -tuples,  $n = |R| \geq 1$ . Choose an  $(s, m)$ -tuple  $r \in T \setminus R$ . Then we define a 2-mapping  $\mathbf{f} = \langle f_1, f_2 \rangle$  of arity  $n$ :  $\text{Dom}(\mathbf{f}) = [r_1, \dots, r_n] = \{\langle r_1(i), \dots, r_n(i) \rangle : i = 1, \dots, s+m\}$  and  $\mathbf{f}(r_1, \dots, r_n) = \langle f_1(r_1(1), \dots, r_n(1)), \dots, f_1(r_1(s), \dots, r_n(s)), f_2(r_1(s+1), \dots, r_n(s+1)), \dots, f_2(r_1(s+m), \dots, r_n(s+m)) \rangle = \langle r(1), \dots, r(s+m) \rangle = r$ .

Since  $R$  is a non-full relation  $\mathbf{f}$  is not everywhere defined. In addition,  $\mathbf{f}$  is a partial 2-operation, i.e., both components  $f_1$  and  $f_2$  are one-to-one partial operations. In other words, for every equal columns  $\langle r_1(i), \dots, r_n(i) \rangle$  and  $\langle r_1(j), \dots, r_n(j) \rangle$  from  $\text{Dom}(\mathbf{f})$  we have  $r(i) = r(j)$  ( $1 \leq i, j \leq s$  or  $s+1 \leq i, j \leq s+m$ ). It is true because in this case  $R \subset D$ , where  $D(x_1, \dots, x_s, y_1, \dots, y_m) \equiv (x_i = x_j)$  ( $1 \leq i, j \leq s$ ) is a 2-base diagonal of arity  $(s, m)$ , and hence  $D$  is involved in formula (6) which gives us  $T \subseteq D$  and  $r \in D$ .

Next we need three facts about  $\mathbf{f}$ .

**Fact 1.**  $\mathbf{f} \notin \text{Pol}(R)$  ( $\mathbf{f}$  does not preserve  $\text{Pol}(R)$ ).

Holds straightforward from the definition of  $\mathbf{f}$ .

**Fact 2.**  $\mathbf{f} \notin \text{Pol}(\mathfrak{R})$ .

Since  $R \in \text{Inv}(\text{Pol}(\mathfrak{R}))$  we obtain  $\text{Pol}(\mathfrak{R}) \subseteq \text{Pol}(R)$  by using antimonotone property of the functor  $\text{Pol}$  (see, e.g., [1]). Then we apply Fact 1.

**Fact 3.** *There exists such relation  $Q \in \mathfrak{R}$  that  $\mathbf{f}$  does not preserve  $Q$ .*

Follows straight from the Fact 2.

First let  $Q$  be a 2-base relation of arity  $(p, t)$  ( $p, t \geq 1$ ). Then from the Fact 3 there exist  $n$  2-base  $(p, t)$ -tuples  $q_1, \dots, q_n \in Q$  such that  $\mathbf{f}(q_1, \dots, q_n) = q \notin Q$ . In addition, since  $\text{Dom}(\mathbf{f}) = [r_1, \dots, r_n]$  we have  $[q_1, \dots, q_n] \subseteq [r_1, \dots, r_n]$  (inclusion as sets of  $n$ -tuples  $[q_1, \dots, q_n] = \{\langle q_1(i), \dots, q_n(i) \rangle : i = 1, \dots, p+t\}$  and  $[r_1, \dots, r_n] = \{\langle r_1(j), \dots, r_n(j) \rangle : j = 1, \dots, s+m\}$ ). Notice that by identification of equal coordinates in  $Q$  one can reduce its arity to  $p \leq s$  and  $t \leq m$  still satisfying the Fact 3.

We introduce two everywhere defined one-to-one mappings  $\varphi : \{1, \dots, p\} \rightarrow \{1, \dots, s\}$ ,  $i \mapsto \varphi i$ , and  $\psi : \{1, \dots, t\} \rightarrow \{1, \dots, m\}$ ,  $j \mapsto \psi j$ , between the numbers of  $n$ -tuples from

$[q_1, \dots, q_n]$  and  $[r_1, \dots, r_n]$ :

$$\begin{aligned} \langle q_1(i), \dots, q_n(i) \rangle &= \langle r_1(\varphi i), \dots, r_n(\varphi i) \rangle \quad \text{for all } n\text{-tuples on } E(k_1), \\ \langle q_1(j), \dots, q_n(j) \rangle &= \langle r_1(\psi j), \dots, r_n(\psi j) \rangle \quad \text{for all } n\text{-tuples on } E(k_2). \end{aligned} \quad (7)$$

Now we define the relation  $S$  of arity  $(s, m)$  as follows:

$$S(x_1, \dots, x_s, y_1, \dots, y_m) \equiv Q(x_{\varphi 1}, \dots, x_{\varphi p}, y_{\psi 1}, \dots, y_{\psi t}), \quad (8)$$

where all coordinates, other than explicitly shown on the right side, are free or complete.

Next we establish several properties of  $S$ :

(i)  $R \subseteq S$ .

According to (7) we have  $r_1, \dots, r_n \in S$  and so  $R = \{r_1, \dots, r_n\} \subseteq S$ .

(ii)  $T \subseteq S$ .

Since  $S$  is constructed via  $\&$ -formula from  $Q \in \mathfrak{R}$  we get that  $S$  is involved in the formula (6) and so  $T \subseteq S$ .

(iii)  $r \in S$

Follows straight from (ii).

Since  $f_1$  and  $f_2$  are one-to-one operations we obtain from (7) that:

$$\begin{aligned} \langle q(i) \rangle &= \langle r(\varphi i) \rangle \quad \text{for elements from } E(k_1), \\ \langle q(j) \rangle &= \langle r(\psi j) \rangle \quad \text{for elements from } E(k_2). \end{aligned} \quad (9)$$

Next we define two mappings  $\alpha: \{1, \dots, s\} \rightarrow \{1, \dots, p\}$ ,  $i \mapsto \alpha$ , and  $\beta: \{1, \dots, m\} \rightarrow \{1, \dots, t\}$ ,  $j \mapsto \beta j$  such that:  $\alpha i = j$ , when  $\varphi j = i$  and  $\alpha i = 1$  otherwise ( $i = 1, \dots, s$ );  $\beta i = j$ , when  $\psi j = i$ , and  $\beta i = 1$  otherwise ( $i = 1, \dots, m$ ).

Finally, from the formula (8) we obtain:

$$Q(x_1, \dots, x_p, y_1, \dots, y_t) \equiv S(x_{\alpha 1}, \dots, x_{\alpha s}, y_{\beta 1}, \dots, y_{\beta m}). \quad (10)$$

Moreover, from (iii) ( $r \in S$ ) and (9) we obtain that in formula (10)  $q \in Q$  that contradicts our previous assumptions.

In the case, when  $Q$  is single-sorted, we use only one mapping  $\varphi: \{1, \dots, p\} \rightarrow \{1, \dots, s\}$  and obtain  $S(x_1, \dots, x_s, y_1, \dots, y_m) \equiv Q(x_{\varphi 1}, \dots, x_{\varphi p})$  with the converse identification (instead of (10)):  $\pi_2 S(x_{\alpha 1}, \dots, x_{\alpha s}, y, \dots, y) \equiv Q(x_1, \dots, x_p)$ .

So there is no  $r \in T \setminus R$  and  $R \equiv T$ , which proves the lemma.

Applying Lemma 2.1 we get  $\mathfrak{R} = \text{Inv}(\text{Pol}(\mathfrak{R}))$  for any set  $\mathfrak{R}$  closed under formation of  $(\&, =_1, \dots, =_n)$ -formulas and application of  $\pi_i$  ( $i = 1, \dots, n$ ). This proves the theorem.  $\square$

If  $\mathbf{A} = \text{Pol}(R)$ , then we say that  $R$  determines partial  $n$ -clone  $\mathbf{A}$ . Using Galois connection properties we obtain that in this case  $R$  is a generating relation for the set  $\text{Inv}(\mathbf{A})$  with respect to operations mentioned in Theorem 2.1.

**Corollary 2.1.** *A relation  $R$  determines  $\prod \mathbb{P}(k_i)$  if and only if  $R$  is a multiple-base diagonal.*



Let  $\Phi(k) = Q(k) \cup \{p_m: m \geq 1\}$  be a partial clone on  $E(k)$ ,  $k \geq 2$ , consisting of all everywhere defined and empty operations. It is known [4] that  $\Phi(k)$  is a maximal partial clone (moreover, in [24] this result was extended to an infinite base set  $E$ ). Consider  $n$  ( $n \geq 2$ ) partial  $n$ -clones:

$$\begin{aligned} \Phi_1 &= (\Phi(k_1) \times P(k_2) \times \dots \times P(k_n)) \cup F, \\ \Phi_2 &= (P(k_1) \times \Phi(k_2) \times \dots \times P(k_n)) \cup F, \dots, \\ \Phi_n &= (P(k_1) \times P(k_2) \times \dots \times \Phi(k_n)) \cup F. \end{aligned} \tag{11}$$

**Proposition 2.2.**  $\Phi_i$  ( $i=1, \dots, n$ ) are the only maximal partial  $n$ -clones containing the  $n$ -clone  $\prod Q(k_i)$ .

**Proof.** Consider  $n$ -clone  $\text{Sel} = \text{Sel}(k_1) \times \text{Sel}(k_2) \times \dots \times \text{Sel}(k_n)$ , which is the direct arity-calibrated product of  $n$  clones of all projections (selectors)  $\text{Sel}(k_i)$  on  $E(k_i)$  ( $i = 1, \dots, n$ ). In what follows, we will use the fact, which is based on the properties of  $\text{Sel}$ .

**Fact.** If  $\mathbf{A}$  is a partial  $n$ -clone with  $\text{Sel} \subset \mathbf{A}$ , then  $\mathbf{A} \setminus F$  can be presented in the form  $A_1 \times A_2 \times \dots \times A_n$  of an arity-calibrated direct product of  $n$  partial clones  $A_i$  on  $E(k_i)$  ( $i = 1, \dots, n$ ).

First it is easy to prove maximality of each  $\Phi_i$  using that  $\Phi(k_i)$  is maximal in  $P(k_i)$  ( $i=1, \dots, n$ ). Next from  $\text{Sel} \subset \prod Q(k_i)$  we get that each maximal partial  $n$ -clone containing  $\prod Q(k_i)$  can be presented as a direct product. This proves the second part of the proposition.  $\square$

Denote  $[A]$  the partial  $n$ -clone generated by a set of  $n$ -operations  $A$ .

**Corollary 2.2.**  $\prod P(k_i)$  is generated by the set  $A(2)$ .

**Proof.** Since all binary  $n$ -selectors  $\text{Sel}^{(2)}$  are contained in  $A(2)$  and also  $\text{Sel}^{(2)}$  generates  $\text{Sel}$  the partial  $n$ -clone  $[A(2)]$  generated by  $A(2)$  is presented as a direct product. Next we apply the result that the set of all partial binary operations generates  $P(k_i)$  ( $i = 1, \dots, n$ ) (see [4]).  $\square$

Hence from common algebraic results (see, e.g., [2]) it follows that each proper partial  $n$ -clone is contained in a maximal partial  $n$ -clone and, therefore, a set of partial  $n$ -operations is complete in  $\prod P(k_i)$  if and only if it is not contained in any maximal partial  $n$ -clone ( $n \geq 2$ ).

**Theorem 2.2.** Each maximal partial  $n$ -clone, with the exception of  $\Phi_i$  ( $i = 1, \dots, n$ ), is determined by a multiple-base relation that is minimal under the expressibility by  $\&$ -formulas and distinct from a multiple-base diagonal.

**Proof.** Without the loss of generality consider  $n = 2$ . Let  $\mathbf{A}$  be a maximal partial 2-clone, other than  $\Phi_i$  ( $i = 1, 2$ ). Then applying Proposition 2.2 we obtain  $\mathbf{B} = \mathbf{A} \cap \mathcal{Q}(k_1) \times \mathcal{Q}(k_2) \subset \mathcal{Q}(k_1) \times \mathcal{Q}(k_2)$ , where  $\mathbf{B}$  is a proper 2-clone. Next for binary operations we have  $\mathbf{B}^{(2)} = \mathbf{A}^{(2)} \cap \mathcal{Q}^{(2)}(k_1) \times \mathcal{Q}^{(2)}(k_2)$ . Clearly  $\mathbf{B}^{(2)}$  is included properly in  $\mathcal{Q}^{(2)}(k_1) \times \mathcal{Q}^{(2)}(k_2)$ , otherwise  $[\mathbf{B}^{(2)}] = [\mathcal{Q}^{(2)}(k_1) \times \mathcal{Q}^{(2)}(k_2)] = \mathcal{Q}(k_1) \times \mathcal{Q}(k_2)$ , a contradiction to Proposition 2.2.

Consider a 2-graph  $G_2(\mathbf{B})$  of the  $n$ -clone  $\mathbf{B}$ . We choose the set  $D = E^2(k_1) \cup E^{(2)}(k_2)$ ,  $|D| = k_1^2 + k_2^2$ , where  $D = [r_1, r_2]$  consists of two 2-base tuples  $r_1$  and  $r_2$  of arity  $(k_1^2, k_2^2)$  over  $E(k_1)$  and  $E(k_2)$ . Next we define the relation  $G_2(\mathbf{B})$  of arity  $(k_1^2, k_2^2)$  as follows:

$$G_2(\mathbf{B}) = \{r: \mathbf{f}(r_1, r_2) = r \text{ for some } \mathbf{f} \in \mathbf{B}^{(2)}\}.$$

Clearly  $G_2(\mathbf{B})$  is a non-full relation hence it is not a 2-base diagonal as well (the  $2 \times (k_1^2 + k_2^2)$  matrix  $[r_1, r_2]$  does not have equal columns and so no non-full diagonal contains  $G_2(\mathbf{B})$ ). Finally, it is easy to verify, applying the maximality of  $\mathbf{A}$ , that  $\mathbf{A} = \text{Pol}(G_2(\mathbf{B}))$ .

Hence we proved that each maximal partial  $n$ -clone, other than  $\Phi_i$  ( $i = 1, \dots, n$ ), is determined by a multiple-base relation (in the common case of arity  $(k_1^2, \dots, k_n^2)$ ). Now from Proposition 2.1 we get that maximal partial  $n$ -clones of this type are precisely maximal restriction-closed partial  $n$ -clones. So applying properties of the Galois connection we obtain that  $G_2(\mathbf{B})$  is a generating relation with the minimal expressibility property in the atom  $\text{Inv}(\mathbf{A})$  of the lattice of Galois-closed sets of multiple-base relations, i.e., every non-diagonal  $Q$ ,  $Q \in \text{Inv}(\mathbf{A})$ , can be obtained from  $G_2(\mathbf{B})$  by using operations of the Galois closure on the set of relations and, conversely,  $G_2(\mathbf{B})$  is constructed from  $Q$  via the same operations. Notice that  $G_2(\mathbf{B})$  has no equal or fictitious (dummy) coordinates. Moreover, if we also consider  $Q$  without equal or fictitious coordinates, then  $Q$  can be obtained from  $G_2(\mathbf{B})$  via a  $\&$ -formula and vice versa.  $\square$

In the sequel, we call relations without equal or fictitious coordinates satisfying Theorem 2.2 *minimal*. Straight from the definition of minimal relations we obtain the corollary which enables us to incorporate minimal  $m$ -base relations into  $n$ -base relations, i.e., partial  $m$ -clones into partial  $n$ -clones ( $1 \leq m \leq n$ ).

**Corollary 2.3.** *Every minimal relation over the type  $J$ ,  $|J| \geq 1$ , is also minimal over any type  $I$ ,  $J \subset I$ .*

### 3. Slupecki-type criterion

In order to find the exact estimates of minimal relations arities we will establish a Slupecki-type criterion, i.e., a completeness criterion for systems of partial  $n$ -operations, containing the set  $\Omega(k_1, \dots, k_n) = P^{(1)}(k_1) \times P^{(1)}(k_2) \times \dots \times P^{(1)}(k_n)$  of all unary partial  $n$ -operations.

Namely, we will find all maximal partial  $n$ -clones containing  $\Omega(k_1, \dots, k_n)$ , called *Slupecki* partial  $n$ -clones, via  $n$ -base relations determining them. Notice that  $\Omega(k_1, \dots, k_n)$  is a direct product of  $n$  semigroups  $\Omega(k_i)$  ( $i = 1, \dots, n$ ) of all partial unary

operations defined on  $n$  base sets. At the same time we may also consider  $\Omega(k_1, \dots, k_n)$  as a partial  $n$ -clone by applying  $n$ -selectors [14] (or constant operation  $\mathbf{e}_1^2$ ) to it.

We will describe the structure of  $\text{Inv}(\mathbf{A})$  in the case of unary partial  $n$ -operations (for  $n = 1$  see [16]).

**Proposition 3.1.** *Let  $\mathbf{A}$  be a restriction-closed partial  $n$ -clone. Then  $\mathbf{A}$  is a subsemigroup of  $\Omega(k_1, \dots, k_n)$  (consists of only unary partial  $n$ -operations,  $n$ -selectors and  $F$ ) if and only if  $\text{Inv}(\mathbf{A})$  is closed under application of any disjunction of relations.*

The proof basically follows the case  $n = 1$  (see, e.g., [1]).

**Corollary 3.1.** *The set  $\text{Inv}(\Omega(k_1, \dots, k_n))$  consists of any disjunction of  $n$ -base diagonals.*

Denote  $\mathcal{J}$  the set consisted of any disjunction of  $n$ -base diagonals ( $n \geq 1$ ). Applying Proposition 2.1 we get the corollary.

**Corollary 3.2.** *Each restriction-closed partial  $n$ -clone, containing  $\Omega(k_1, \dots, k_n)$ , is determined by a set of relations from  $\mathcal{J}$ .*

Then applying Theorem 2.2 we obtain the following corollary.

**Corollary 3.3.** *Each Slupecki partial  $n$ -clone is determined by a minimal relation from the set  $\mathcal{J}$ .*

Now it suffices to find all minimal relations in the set  $\mathcal{J} \setminus \mathbf{D}$ , which determine distinct partial  $n$ -clones.

**Definition 3.1.** A non-diagonal  $n$ -base relation  $S$  ( $n \geq 2$ ) is called irreducible if by applying to  $S$  intersections with permutations of coordinates, identifications of coordinates of the same sort and also  $\pi_i$  ( $i = 1, \dots, n$ ) one cannot obtain a non-diagonal relation of either less arity, or less type, or less number of tuples.

For any  $(h_1, \dots, h_n)$ -tuple  $r$  ( $h_1, \dots, h_n \geq 1$ ) denote  $\varepsilon(r)$  the equivalence relation on numbers of coordinates induced by equal coordinates in  $r$ , e.g., for a  $(2, 2)$ -tuple  $r = (0, 0; 1, 1)$  we have  $\varepsilon(r) = \{(1, 2), (3, 4)\}$  and  $\varepsilon(r) = \varepsilon(D)$ , where  $D \equiv x_1 = x_2 \ \& \ y_1 = y_2$  is a 2-base diagonal corresponding to  $\varepsilon(r)$ . If  $r$  has no equal coordinates, then  $\varepsilon(r)$  is the trivial equivalence which represents the full  $n$ -base diagonal of arity  $(h_1, \dots, h_n)$ .

**Lemma 3.1.** *Let  $S$  be an irreducible  $n$ -base relation. Then for every  $r \in S$  such that  $\varepsilon(r)$  is non-trivial we have  $D \subset S$ , where  $\varepsilon(D) = \varepsilon(r)$  and  $D \in \mathbf{D}$ .*

**Proof.** Assume  $D \not\subset S$  for some  $r \in S$  and  $\varepsilon(D) = \varepsilon(r)$ . Then applying to  $S$  identifications of coordinates according to all blocks of  $\varepsilon(D)$  we obtain a non-diagonal relation

which contradicts the fact that  $S$  is irreducible.  $\square$

$$\text{Set } T(h) \equiv \bigvee_{1 \leq i < j \leq h} (x_i = x_j), \quad h \geq 2.$$

**Proposition 3.2.** *Each irreducible relation  $S$ ,  $S \in \mathfrak{J} \setminus \mathbf{D}$ , of arity  $(h_1, \dots, h_n)$  and type  $\{1, \dots, n\}$  ( $2 \leq h_1 \leq k_1, \dots, 2 \leq h_n \leq k_n$ ,  $n \geq 2$ ) is presented as a disjunction  $T(h_1) \vee \dots \vee T(h_n)$  of  $n$  single-base relations defined on sets  $E(k_1), \dots, E(k_n)$ , respectively.*

**Proof.** We consider the proof for the case  $n=2$ . The same idea is applicable to  $n \geq 2$ .

Let  $S(x_1, \dots, x_n, y_1, \dots, y_m)$  ( $n, m \geq 2$ ) be a 2-base irreducible relation of arity  $(n, m)$  (if  $n=1$ , then using  $\pi_1$  we obtain a single-base non-diagonal relation). So  $S$  can be presented in the form:

$$S \equiv \bigvee_{i=1}^t D_1^i(x_1, \dots, x_n) \& D_2^i(y_1, \dots, y_m), \quad (12)$$

where  $D_1^i$  are diagonals of the 1st sort and  $D_2^i$  are diagonals of the 2nd sort ( $i=1, \dots, t$ ).

Now consider the relation  $D(y_1, \dots, y_m) \equiv \pi_1 S(x, \dots, x, y_1, \dots, y_m)$ , which is a diagonal due to Definition 2.1. If  $D$  is a non-full diagonal, then from (12) we get  $D_2^i \subseteq D$  ( $i=1, \dots, t$ ). Hence  $S \subset \bigvee_{1 \leq i \leq t} D_2^i \subseteq D$  and so  $S \cap D = S$  (here  $D$  has  $n$  fictitious variables of the 1st sort). Then by Lemma 3.1  $D \subseteq S$  and  $D = S$ . Contradiction.

So  $D$  is the full diagonal and, therefore, there exists  $a \in E(k_1)$  such that  $(a, \dots, a; b_1, \dots, b_m) \in S$ , where  $(b_1, \dots, b_m)$  are all possible  $m$ -tuples from  $E^m(k_2)$ . Then applying Lemma 3.1 we have  $x_1 = \dots = x_n \subset S$ . Similarly we obtain  $y_1 = \dots = y_m \subset S$ . Hence we proved that  $S$  can be presented in a form of separated single-base disjunctive components:

$$S \equiv R_1(x_1, \dots, x_n) \vee T(x_1, \dots, x_n, y_1, \dots, y_m) \vee R_2(y_1, \dots, y_m), \quad (13)$$

where  $R_1$  and  $R_2$  are non-full single-base diagonals and  $R_1, R_2, T \in \mathfrak{J}$ .

In addition, we choose  $R_1$  and  $R_2$  as the greatest single-base disjunctive components, i.e., if a single-base diagonal  $D^1 \subset S$  ( $D^2 \subset S$ ), then  $D^1 \subseteq R_1$  ( $D^2 \subseteq R_2$ , respectively). At the same time, we assume that  $T$  does not contain any single-base diagonals with fictitious coordinates.

**Fact 1.** *Relations  $R_1$  and  $R_2$  in the expression (13) are totally symmetric, i.e., stable under any permutations of coordinates.*

**Proof.** Let  $\alpha$  be a permutation of  $n$  variables in  $R_1$ :  $R_1^\alpha(x_1, \dots, x_n) \equiv R_1(x_{\alpha 1}, \dots, x_{\alpha n})$ . Then from (13) we get  $S^\alpha \equiv R_1^\alpha(x_1, \dots, x_n) \vee T(x_{\alpha 1}, \dots, x_{\alpha n}, y_1, \dots, y_m) \vee R_2(y_1, \dots, y_m)$ . Hence by using properties of operations  $\&$  and  $\vee$  we have

$$S \& S^\alpha \equiv R_1^\alpha(x_1, \dots, x_n) \& R_1(x_1, \dots, x_n) \vee T_1(x_1, \dots, x_n, y_1, \dots, y_m) \vee R_2(y_1, \dots, y_m),$$

where  $T_1 \equiv R_1 \& T^\alpha \vee T \& T^\alpha \vee R_1^\alpha \& T$  is a 2-base relation from  $\mathfrak{J}$ .

Since  $x_1 = \dots = x_n \subseteq R_1^\alpha \& R_1$ , the relation  $S \& S^\alpha$  is not a diagonal (it has single-base disjunctive components for each sort of variables). Hence from Definition 3.1  $S \& S^\alpha \equiv R_2$  and so  $R \equiv R_1^\alpha \& R_1$ , which implies  $R_1^\alpha \equiv R_1$ .  $\square$

Now consider  $S$  in formula (13) in two different cases.

*Case  $n = 2$  ( $m = 2$ ):* Here it is easy to verify that  $T(x_1, x_2, y_1, \dots, y_m) \subseteq x_1 = x_2$  and so  $S$  can be presented in the form:  $S \equiv x_1 = x_2 \vee R_2(y_1, \dots, y_m)$ . If  $m = 2$ , then  $R_2 \equiv y_1 = y_2$  and so  $S \equiv x_1 = x_2 \vee y_1 = y_2$ .

Let  $m \geq 3$ . Then the relation  $S(x_1, x_2, y_1, y_1, y_3, \dots, y_m)$  is the full 2-base diagonal of arity  $(2, m - 1)$  (because of the disjunctive component  $x_1 = x_2$  it cannot be a non-full diagonal). Hence applying Lemma 3.1 we get  $y_1 = y_2 \subseteq S$  and also  $y_1 = y_2 \subseteq R_2$  (greatest disjunctive component property). Next from the Fact 1 we conclude that  $R_2 \equiv T(m) \equiv \bigvee_{1 \leq i < j \leq m} (x_i = x_j)$ . Note that here we have  $3 \leq m \leq k_2$ , since for  $m > k_2$   $T(m)$  is the full relation and  $S$  is also full relation. Finally, we get  $S \equiv x_1 = x_2 \vee T(m)$ ,  $2 \leq m \leq k_2$ . (For the case  $m=2$  we have  $S \equiv T(n) \vee x_1 = x_2$ ,  $2 \leq n \leq k_1$ ).

*Case  $n, m > 2$ :* Consider the relation  $S' \equiv S(x_1, \dots, x_n, y_1, y_1, y_3, \dots, y_m)$  which is a 2-base diagonal of arity  $(n, m - 1)$  (see Definition 3.1). If  $S'$  is a non-full diagonal, then this contradicts the inclusion  $R_1(x_1, \dots, x_n) \subseteq S'$  which follows straight from (13). So  $S'$  is the full diagonal and applying Lemma 3.1 we obtain  $y_1 = y_2 \subseteq S$  and from the greatest disjunctive component property we have  $y_1 = y_2 \subseteq R_2$ . Then from the Fact 1 we get  $R_2 \equiv T(m)$ ,  $2 \leq m \leq k_2$ . Next by repeating the same steps we obtain  $R_1 \equiv T(n)$ ,  $2 \leq n \leq k_1$ . Finally,  $S \equiv T(n) \vee T(m)$ ,  $n, m \geq 2$ .  $\square$

It is obvious that every maximal partial  $n$ -clone, with  $n$  exceptions, can be determined by an irreducible relation. Hence from Corollary 3.3 and Proposition 3.2 we get corollary.

**Corollary 3.4.** *Each Slupecki partial  $n$ -clone, that is a subdirect product of  $m$  ( $2 \leq m \leq n$ ) factors, is determined by a relation which is contained among the relations having the form:  $T(h_1) \vee \dots \vee T(h_m)$  ( $2 \leq h_1 \leq k_1, \dots, 2 \leq h_m \leq k_m$ ), where  $T(h_i)$  has the type  $\{i\}$  ( $1 \leq i \leq m$ ), or by a relation obtained from them by a permutation of numbers of base sets.*

The converse is also true.

**Proposition 3.3.** *Each relation  $T(h_1) \vee \dots \vee T(h_m)$  ( $2 \leq h_1 \leq k_1, \dots, 2 \leq h_m \leq k_m$ ,  $2 \leq m \leq n$ ) is a minimal  $m$ -base relation.*

The proof for the general case will be presented in the next section (Proposition 4.2).

Recall that all  $k$  Slupecki partial clones on  $E(k)$ ,  $k \geq 3$ , were described in [15] by  $k$  invariant relations:  $H_1 \equiv x_1 = x_2 \& x_3 = x_4 \vee x_1 = x_3 \& x_2 = x_4$ ,  $H_2 \equiv x_1 = x_2 \& x_3 = x_4 \vee x_1 = x_3 \& x_2 = x_4 \vee x_1 = x_4 \& x_2 = x_3$ ,  $T(h)$  ( $h = 3, \dots, k$ ). If  $k = 2$ , then there exist 2 Slupecki partial clones  $\text{Pol}(H_1)$  and  $\text{Pol}(H_2)$  (see [4]). We define the set  $G_i$  of  $k_i$  single-base relations of type  $\{i\}$  on  $E(k_i)$  ( $i = 1, \dots, n$ ) as follows: if  $k_i \geq 3$ ,  $G_i = \{H_1, H_2, T(h) \ (3 \leq h \leq k_i)\}$  and if  $k_i = 2$ , then  $G_i = \{H_1, H_2\}$  ( $i = 1, \dots, n$ ).

Finally, from the results of this section we obtain the theorem.

**Theorem 3.1.** *Each Slupecki partial  $n$ -clone ( $n \geq 2$ ) is defined by a relation such that*

- (1)  $R \in G_1 \cup \dots \cup G_n$  or
- (2)  $R$  is represented as a disjunction  $R_1 \vee \dots \vee R_n$ , where each  $R_i$  ( $i=1, \dots, n$ ) is either one of  $T(h)$  ( $2 \leq h \leq k_i$ ) with the type  $J(R_i) = \{i\}$ , or empty and, moreover, at least two of disjunctive components  $R_i$  are nonempty.

Similarly to the case  $n = 1$  [15] each maximal restriction-closed partial  $n$ -clone  $\mathbf{A}$ , except for Slupecki partial  $n$ -clones, is determined by its 1-graph  $G_1(\mathbf{A})$  (the graph of all unary  $n$ -operations  $\mathbf{A} \cap \Omega(k_1, \dots, k_n)$ ) which has an arity  $(k_1, \dots, k_n)$ .

**Corollary 3.5.** *Each maximal partial  $n$ -clone, other than  $\Phi_i$  ( $i=1, \dots, n$ ), is determined by a minimal multiple-base relation of arity  $(k_1, \dots, k_n)$ .*

*Slupecki criterion for  $\mathbb{Q}(k_1) \times \dots \times \mathbb{Q}(k_n)$ .* We will apply results of this section to the description of all Slupecki  $n$ -clones ( $n \geq 2$ ) (maximal  $n$ -clones including all unary  $n$ -operations). Similar to partial  $n$ -clones by establishing analogues of Proposition 3.1 and Corollary 3.1 we obtain the fact: every non-full  $n$ -clone  $\mathbf{B}$ , which contains all unary  $n$ -operations, has the form  $\mathbf{B} = \text{Pol}^l(\mathfrak{R})$ , where  $\mathfrak{R} \subseteq \mathcal{J}$  and  $\mathfrak{R} \cap (\mathcal{J} \setminus \mathbf{D}) \neq \emptyset$ . Hence we get the following proposition.

**Proposition 3.4.** *Each Slupecki  $n$ -clone ( $n \geq 2$ ) is determined by a non-diagonal relation  $R$ ,  $R \in \mathcal{J}$ , such that  $\text{Pol}^l(R)$  is a maximal element by inclusion among all  $n$ -clones of the form  $\text{Pol}^l(S)$ ,  $S \in \mathcal{J} \setminus \mathbf{D}$ .*

Next it suffices to investigate only irreducible relations described in Proposition 3.2, because if  $S$  is reduced by intersections and identifications to irreducible  $R$ , then it is obvious that  $\text{Pol}^l(S) \subseteq \text{Pol}^l(R)$ . Further we will need the lemma.

**Lemma 3.2.** *An  $n$ -operation  $\mathbf{f} = \langle f_1, \dots, f_n \rangle$ ,  $\mathbf{f} \notin \text{Sel}$ , belongs to  $\text{Pol}^l(T(h_1) \vee \dots \vee T(h_m))$  ( $2 \leq h_1 \leq k_1, \dots, 2 \leq h_m \leq k_m, 2 \leq m \leq n$ ) if and only if there exists  $i$ ,  $1 \leq i \leq n$ , such that the range of  $f_i$  is less or equal  $h_i - 1$  ( $2 \leq h_i \leq k_i$ ).*

The proof of this lemma is based on the case  $n = 1$  (Slupecki criterion for  $k$ -valued logic, see, e.g., [8]). Recall that Slupecki  $n$ -clones determined by single-base relations for each type  $\{i\}$  are: the Slupecki clone  $\text{Pol}^l(T(k_i))$ , when  $k_i \geq 3$ , or the clone of all linear Boolean functions [13] having the form  $\text{Pol}^l(H_2)$ , when  $k_i = 2$  ( $i = 1, \dots, n$ ).

Applying Lemma 3.2 we get  $\text{Pol}^l(T(h_1) \vee \dots \vee T(h_m)) \subset \text{Pol}^l(T(t_1) \vee \dots \vee T(t_m))$ , where  $h_1 \leq t_1 \leq k_1, \dots, h_m \leq t_m \leq k_m$ , and there is at least one strict inequality ( $2 \leq m \leq n$ ). So all maximal elements by inclusion satisfying Proposition 3.4 are exactly  $\text{Pol}^l(T(k_1) \vee \dots \vee T(k_m))$  and the ones obtained from them by permutations of the numbers of base sets. Finally, we obtain the description of all Slupecki  $n$ -clones (see also [12,19,25]).

**Theorem 3.2.** *There are exactly  $2^n - 1$  Slupecki  $n$ -clones that are determined by multiple-base relations having the form*

- (a)  $R \equiv R_1 \vee \dots \vee R_n$ , where  $R_i \in \{\emptyset, T(k_i)\}$ ,  $J(R_i) = \{i\}$ , and at least two  $R_i$  are nonempty;
- (b) single-sorted relations of the type  $\{i\}$  ( $i = 1, \dots, n$ ), namely,  $R \equiv T(k_i)$ , when  $k_i \geq 3$ , or  $R \equiv H_2$ , when  $k_i = 2$ .

We call an  $n$ -operation  $\mathbf{f} = \langle f_1, \dots, f_n \rangle$  essential over type  $\{i\}$  ( $1 \leq i \leq n$ ), if either  $f_i$  is essential (has the full range  $k_i$  and is a non-selector), when  $k_i \geq 3$ , or  $f_i$  is a non-linear Boolean function, when  $k_i = 2$ . Then  $\mathbf{f}$  is essential over type  $J = \{i_1, \dots, i_m\}$ ,  $J \subseteq I$ ,  $2 \leq |J| \leq n$ , if for every  $i \in J$   $f_i$  has the full range and  $\langle f_{i_1}, \dots, f_{i_m} \rangle$  is not equal to an  $m$ -vector of unary partial operations  $\langle \alpha_{i_1}, \dots, \alpha_{i_m} \rangle \in \Omega(k_{i_1}) \times \dots \times \Omega(k_{i_m})$  (up to fictitious coordinates). Next  $\mathbf{f}$  is essential if for every  $i \in \{1, \dots, n\}$  either  $f_i$  has the full range and is a non-selector, when  $k_i \geq 3$ , or  $f_i$  is a non-linear Boolean function, when  $k_i = 2$ . Finally, we obtain Slupecki criterion for  $n$ -clones (see also [23]).

**Proposition 3.5.** *A set  $B$  of  $n$ -operations is complete in  $\mathbb{Q}(k_1) \times \dots \times \mathbb{Q}(k_n)$  with all unary  $n$ -operations if and only if for every type  $J$ ,  $J \subseteq \{1, \dots, n\}$ ,  $1 \leq |J| \leq n$  ( $n \geq 2$ ), there exists an  $n$ -operation  $\mathbf{f} \in B$  which is essential over  $J$ .*

**Corollary 3.6.** *An  $n$ -operation  $\mathbf{f}$  is complete in  $\mathbb{Q}(k_1) \times \dots \times \mathbb{Q}(k_n)$  with all unary  $n$ -operations if and only if  $\mathbf{f}$  is essential.*

**Corollary 3.7.** *Each maximal  $n$ -clone is determined by a multiple-base relation of arity  $(k_1, \dots, k_n)$  (with the exception of a single-base relation  $H_2$  of arity 4 on  $E(2)$ ).*

#### 4. Maximal partial $n$ -clones

In what follows we explore irreducible relations of arity less or equal  $(k_1, \dots, k_n)$  which do not belong to  $\mathfrak{I}$ . Without loss of generality we consider  $n$ -base relations of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$ , where  $0 \leq m \leq n$  and  $2 \leq h_i \leq k_i$  ( $i = 1, \dots, m$ ) (one can pass to the general case by changing numbers of base sets). We also need definitions extending case  $n = 1$ .

1. A multiple-base relation  $R$  is called *areflexive* if it contains no tuples with equal coordinates, i.e.,  $R \cap (T(h_1) \vee \dots \vee T(h_m)) = \emptyset$ . Denote  $\mathbf{R}$  the set of all areflexive relations.

2. A multiple-base relation  $H$  is called *totally symmetric*, when it is stable under each permutation of coordinates of the same  $i$ th sort ( $1 \leq i \leq m$ ) and *totally reflexive*, when  $T(h_1) \vee \dots \vee T(h_m) \subseteq H$ . Denote  $\mathbf{H}$  the set of all totally reflexive and totally symmetric non-full relations (for  $n = 1$  see [27]).

**Example 4.1.** Let  $E(k_1) = k \geq 3$ ,  $E(k_2) = 2$  and a 2-base relation of arity  $(h, 1)$ ,  $2 \leq h \leq k$ , be defined as follows:  $H(x_1, \dots, x_h, y) \equiv \{(x_1, \dots, x_h, y) : (x_1, \dots, x_h) \in T(h)\}$  or

$(x_1 + \dots + x_h) = 0 \pmod{k} \ \& \ y = 1 \} \equiv T(h)(x_1, \dots, x_h) \vee \langle \langle (x_1 + \dots + x_h) = 0 \rangle \rangle \ \& \ y = 1$ .  
Then  $H \in \mathbf{H}$ .

3. For every non-single type  $J$ ,  $2 \leq |J| \leq n$ , the set  $\mathbf{K}$  consists of all nonempty, non-full relations of arity  $(1, \dots, 1)$  and type  $J$ .

Note that if a relation  $T$  is obtained by a  $\&$ -formula from irreducible  $Q$  of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$ , then by identification of coordinates of types  $s > m$  in  $T$  we also get a non-diagonal relation  $S$  of arity  $\langle s_1, \dots, s_m, 1, \dots, 1, 0, \dots, 0 \rangle$  and  $s_i \geq h_i$  ( $i = 1, \dots, m$ ). Now we will consider a special presentation of  $S$  by a  $\&$ -formula from  $Q$ . Without loss of generality  $Q$  has arity  $(h_1, \dots, h_m)$  and  $S - (s_1, \dots, s_m)$  respectively  $(h_i \leq s_i, i = 1, \dots, m)$ . Then we introduce an *index  $m$ -base relation*  $M$  of arity  $(h_1, \dots, h_m)$  on base sets  $E(s_1), \dots, E(s_m)$ . An index relation  $M$  represents any  $S$  constructed by a  $\&$ -formula from  $Q$ :

$$S(x_0, \dots, x_{s_1-1}, y_0, \dots, y_{s_2-1}, \dots, z_0, \dots, z_{s_m-1}) \equiv \&_{r \in M} Q^r, \quad (14)$$

where  $r = (r(1, 1), \dots, r(1, h_1); r(2, 1), \dots, r(m, h_m)) \in M$  is a  $(h_1, \dots, h_m)$ -tuple over  $E(s_1), \dots, E(s_m)$  and  $Q^r \equiv Q(x_{r(1,1)}, \dots, y_{r(2,1)}, \dots, z_{r(m,1)}, \dots)$ .

Next if  $Q$ , in turn, can be obtained by a  $\&$ -formula from  $S$ , then clearly it can be done by using intersections with identifications and permutations of coordinates. So we get.

**Lemma 4.1.** *Any irreducible multiple-base relation  $Q$  of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$  is minimal if and only if from every non-diagonal relation  $S$  of arity  $\langle s_1, \dots, s_m, 1, \dots, 1, 0, \dots, 0 \rangle$  constructed by the formula (14) one can obtain  $Q$  using intersections with identifications and permutations of coordinates in  $S$ .*

**Proposition 4.1.** *Each  $Q \in \mathbf{K}$  is a minimal relation.*

**Proof.** Let  $S$  be constructed from  $Q \in \mathbf{K}$  by the formula (14). We consider an identification  $\Delta$  of coordinates in  $S$  as follows: for all  $r \in M$   $r(1, 1) \rightarrow 1, \dots, r(m, 1) \rightarrow 1$ . Hence we get  $\Delta S \equiv Q$ . Then we apply Lemma 4.1.  $\square$

**Proposition 4.2.** *Each  $Q \in \mathbf{H}$  is a minimal relation.*

**Proof.** Clearly that in this case if  $S$  in the formula (14) is not a diagonal, then there exists a point  $q \in M$  with all pairwise distinct coordinates of the same sort. Consider identification  $\Delta$  of coordinates in  $S: q(i, j) \rightarrow q(i, j)$  ( $i = 1, \dots, m; j = 1, \dots, h_i$ ) and  $r(i, j) \rightarrow q(i, j)$  for any  $r \in M \setminus \{q\}$ . Hence we have  $\Delta S \equiv Q$ . Next see Lemma 4.1.  $\square$

Note that all minimal relations from Proposition 3.3 are included into the set  $\mathbf{H}$ . So the above proof also covers that case.

**Lemma 4.2.** *For each irreducible non-single sort relation  $Q$  of arity less or equal  $(k_1, \dots, k_n)$  we have either:*

- (1)  $Q$  belongs to  $\mathbf{K}$  ( $Q \in \mathbf{K}$ );
- (2)  $Q$  belongs to  $\mathbf{H}$  ( $Q \in \mathbf{H}$ );



- (3)  $Q$  is areflexive ( $Q \in \mathbf{R}$ );
- (4)  $Q$  has the form  $R \vee D$ , where  $R \in \mathbf{R}$  and  $D$  is a multiple-base non-full diagonal of the same arity as  $R$ .

**Proof.** Let  $Q$  be an irreducible relation of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$  ( $m \geq 1$ ) and  $Q \notin \mathbf{K}$ . Then either  $Q$  is areflexive or  $Q \cap D \neq \emptyset$ , where  $D$  is a multiple-base non-full diagonal. Applying Lemma 3.1 we obtain that  $Q \equiv R \vee S$  or  $Q \equiv S$ , where  $R \in \mathbf{R}$  and  $S \in \mathcal{J}$ . If  $S$  is a diagonal, then  $Q \equiv R \vee D$  (case 4). Next if  $S$  is a non-diagonal, then according to Proposition 3.2,  $S$  has the form  $T(h_1) \vee \dots \vee T(h_m)$  and, moreover,  $R$  admits all permutations, since  $Q$  is irreducible. Hence  $Q \in \mathbf{H}$ .  $\square$

Now it suffices to clear cases (3) and (4) in the previous lemma. In what follows  $Q$  will be of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$  with  $s$  ( $s \geq 2$ ) non-void sorts of coordinates ( $m \leq s \leq n$ ).

**Lemma 4.3.** *An  $s$ -base irreducible relation of the form  $Q \equiv R$  ( $R \in \mathbf{R}$ ) or  $Q \equiv R \vee D$  ( $R \in \mathbf{R}, D \in \mathbf{D}$ ) is minimal if and only if every relation  $T \equiv \&_{r \in M} Q^r$  ( $M \subseteq R$ ) of arity  $\langle k_1, \dots, k_m, k_{m+1}, \dots, k_s, 0, \dots, 0 \rangle$  can be reduced by some identification of coordinates to  $Q$ .*

**Proof.** First it easy to verify that  $\text{Pol}(Q)$  is not included in any Slupecki partial  $n$ -clone, i.e., using any  $\&$ -formula one cannot obtain from  $Q$  an  $s$ -base relation of the form  $T(h_1) \vee \dots \vee T(h_t)$  ( $2 \leq t \leq s$ ). The proof of this fact is similar to the case  $n = 1$  (see [22]). Hence from the results of the previous section each maximal partial  $n$ -clone  $\mathbf{A}$ , such that  $\text{Pol}(Q) \subseteq \mathbf{A}$ , satisfy the condition  $\Omega(k_1, \dots, k_n) \notin \mathbf{A}$ . Moreover, there exists such  $\mathbf{A}$  that it is determined by a non-diagonal  $s$ -base relation. To construct such relation consider 1-graph of any  $\mathbf{A} = \mathbf{A}' \times \mathbb{P}(k_{s+1}) \times \dots \times \mathbb{P}(k_n)$ , where  $\mathbf{A}'$  is a subdirect product of  $s$  factors  $\mathbb{P}(k_i)$  ( $1 \leq i \leq s$ ). Namely, we have a relation  $G_1(\mathbf{A}) = \{ \mathbf{f} p : \mathbf{f} \in A \cap \Omega(k_1, \dots, k_n) \}$  of arity  $(k_1, \dots, k_s)$ , where  $p = \langle \mathbf{E}(k_1, ); \dots; \mathbf{E}(k_s) \rangle$  is a  $(k_1, \dots, k_s)$ -tuple,  $\mathbf{E}(k_i) = (0, 1, \dots, k_i - 1)$  ( $1 \leq i \leq s$ ). From the fact that  $\mathbf{A}$  is maximal we get  $\mathbf{A} = \text{Pol}(G_1(\mathbf{A}))$ . Hence  $G_1(\mathbf{A}) \in \text{Inv}(\text{Pol}(Q))$  and so  $G_1(\mathbf{A})$  can be obtained by a  $\&$ -formula from  $Q$ . Therefore, grounding on Lemma 4.1 it is sufficient to consider in the formula (14) only relations  $T \equiv \&_{r \in M} Q^r$  of arity  $(k_1, \dots, k_s)$ , which contain the point  $p$ . It is easy to prove two facts about such relations:

- (1) if  $M \subseteq R$ , then  $p \in T$ ;
- (2) if there exists  $r \in M$  and  $r \notin R$ , then  $p \notin T$ .

So we may consider only index relations  $M$ ,  $M \subseteq R$ . Moreover, since  $Q$  is irreducible each identification of  $T$  to arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$  is either  $Q$  or a diagonal.  $\square$

**Example 4.2.** Consider 2-base irreducible relation  $R = \{(0, 1, a), (1, 0, b)\}$  of arity  $(2, 1)$  on the sets  $E(k_1) = \{0, 1\}$  and  $E(k_2) = \{a, b\}$ . Then by Lemma 4.3 we need to investigate only three relations containing the point  $p$ :  $T_1(x_0, x_1, y_0, y_1) \equiv R(x_0, x_1, y_0)$ ,

$T_2(x_0, x_1, y_0, y_1) \equiv R(x_1, x_0, y_1)$  and  $T_3(x_0, x_1, y_0, y_1) \equiv R(x_0, x_1, y_0) \& R(x_1, x_0, y_1)$ , where  $T_3 = \{(0, 1, a, b), (1, 0, b, a)\}$ . So there is no identification of  $T_3$  to arity  $(2, 1)$  other than empty. Hence applying Lemma 4.3 we obtain that  $R$  is not a minimal relation. At the same time, a single-base projection of  $R$  on the type  $\{1\}$   $R' = \{(0, 1), (1, 0)\}$  is a minimal relation [4].

Let  $G(R)$  be a symmetry group of  $R$ , i.e.,  $G(R)$  is a subgroup of the product  $S(h_1) \times \dots \times S(h_m)$  ( $m \geq 1$ ) of the symmetric permutation groups on numbers of coordinates of each sort  $i$ ,  $1 \leq i \leq m$ , for which  $R$  contains at least two coordinates, such that for each  $\alpha \in G(R)$  the resulting relation  $R^\alpha(x_1, \dots, y_1, \dots, z_1, \dots) \equiv R(x_{\alpha_1}, \dots, y_{\alpha_1}, \dots, z_{\alpha_1}, \dots)$  under application of  $\alpha$  to the numbers of its coordinates equals  $R$  and for each  $\beta \notin G(R)$  we have  $R^\beta \neq R$ . We call  $R$  normal [20] if for each  $\beta \notin G(R)$  we have  $R \cap R^\beta = \emptyset$ . It is obvious, that areflexive  $R$  is normal if and only if it is irreducible, e.g., for  $R$  from Example 4.2 we have  $R(x_0, x_1, y_0) \& R(x_1, x_0, y_1) = \emptyset$  and so  $R$  is a normal relation. Also notice that in this case  $G(R) = \{e\}$  is the identity group.

Denote  $\text{Orb}(G(R))$  the  $(h_1, \dots, h_m)$ -orbit of the group  $G(R)$  (a generalization of the notion of the  $h$ -orbit of a permutation group) that consists of the images of all vector-permutations  $\alpha = \langle \alpha_1, \dots, \alpha_m \rangle \in G(R)$  applied to the  $(h_1, \dots, h_m)$ -tuple  $p = \langle \mathbf{E}(h_1); \dots; \mathbf{E}(h_m) \rangle$ , i.e.,  $\text{Orb}(G(R)) = \{ \langle \alpha_1 \mathbf{E}(h_1); \dots; \alpha_m \mathbf{E}(h_m) \rangle : \langle \alpha_1, \dots, \alpha_m \rangle \in G(R) \}$ , where  $\alpha_i \mathbf{E}(h_i) = (\alpha_i 0, \dots, \alpha_i (h_i - 1))$ ,  $1 \leq i \leq m$ . Hence  $\text{Orb}(G(R))$  is an  $m$ -base relation of arity  $(h_1, \dots, h_m)$  and type  $J = \{1, \dots, m\}$ .

Let  $\Psi_i : E(k_i) \rightarrow E(h_i)$  ( $2 \leq h_i \leq k_i$ ,  $i = 1, \dots, m$ ) be epimorphisms (one-to-one onto mappings) and  $\Psi = \langle \Psi_1, \dots, \Psi_m \rangle$  be the corresponding vector-epimorphism. Also denote  $\Psi R$  the  $m$ -base relation defined on the sets  $E(h_1), \dots, E(h_m)$  that is obtained from the restriction of  $R$  on the coordinates of type  $J = \{1, \dots, m\}$  (each sort of  $J$  contains at least two coordinates in  $R$ ) by application  $\Psi$  to all its points, while  $\Psi_i$  is applied to coordinates of sort  $i$  ( $i = 1, \dots, m$ ). For example, let  $R(x_1, x_2, y_1, y_2, z) = \{(0, 1; 0, 1; 0), (1, 0; 1, 0; 1)\}$  be a relation of arity  $(2, 2, 1)$  over three two-element base sets  $E(2) = \{0, 1\}$ . Then for any 2-epimorphism  $\Psi = \langle \Psi_1, \Psi_2 \rangle (\Psi_1 : E(2) \rightarrow E(2), \Psi_2 : E(2) \rightarrow E(2))$  we have a  $(2, 2)$ -relation  $\Psi R = \{ (\Psi_1 x_1, \Psi_1 x_2, \Psi_2 y_1, \Psi_2 y_2) : (x_1, x_2, y_1, y_2) \in R \}$ . Notice that in this case  $G(R)$  is the identity group and so  $\text{Orb}(G(R)) = \{(0, 1; 0, 1)\}$ .

**Proposition 4.3.** *Each areflexive  $s$ -base relation  $R$  of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$ , where  $1 \leq m \leq s \leq n$ ,  $2 \leq h_i \leq k_i$  ( $i = 1, \dots, m$ ),  $s \geq 2$ , is minimal if and only if:*

- (1)  $R$  is normal (sufficient condition for arity  $(k_1, \dots, k_s)$ );
- (2) there exists a vector-epimorphism  $\Psi = \langle \Psi_1, \dots, \Psi_m \rangle$  such that  $\Psi R = \text{Orb}(G(R))$ .

**Proof.** Straight from the Lemma 4.3 we get that  $R$  of arity  $(k_1, \dots, k_s)$  is minimal if and only if it is normal (the case  $n = 1$  see in [20]). Now consider the common case.

First we show that the part (1) of this proposition is the necessary condition for a relation to be minimal (this condition is absent in the results [9,10] for  $n = 1$ ). Indeed, if  $R \cap R^\alpha = R'$ ,  $\emptyset \neq R' \subset R$ , for some vector-permutation  $\alpha$ , then we have

$\text{Pol}(R) \subset \text{Pol}(R')$ , since one cannot obtain  $R$  via  $\&$ -formula from areflexive  $R'$  of the same arity, which is included in  $R$ .

Then it is obvious that any identification of  $T$  in Lemma 4.3 to a relation of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$  corresponds to application of some vector-epimorphism  $\Psi$  to the index relation  $M$ ,  $M \subseteq R$ , provided that all variables of each sort  $i$ ,  $m \leq i \leq s$ , are identified with a single variable of the same sort.

Next since  $R$  is a normal relation each identification of  $T \equiv \&_{r \in M} R^r$  ( $M \subseteq R$ ) in Lemma 4.3 to the arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$  is either  $R$  or empty. Hence if there exists  $\Psi$  such that  $T^\Psi \equiv \&_{r \in \Psi R} R^r = R$ , then the same  $\Psi$ , while applied to any non-void  $M \subseteq R$ , gives us  $T^\Psi \equiv \&_{r \in \Psi M} R^r = R$ . So for the case  $Q \equiv R$  it is sufficient to consider in Lemma 4.3 only relations of the form  $T \equiv \&_{r \in R} R^r$ .

It is easy to verify that  $\Psi R \subseteq \text{Orb}(G(R))$  implies  $T^\Psi = R$ . Moreover, in this case we have  $\Psi R = \text{Orb}(G(R))$ , since  $\Psi$  is a vector-epimorphism and  $R$  is normal. Next if there exists  $\Psi$  such that  $T^\Psi \equiv R$  and  $p \notin \Psi R$ , then we can find such vector-permutation  $\alpha = \langle \alpha_1, \dots, \alpha_m \rangle$  on  $E(h_1) \times \dots \times E(h_m)$ ,  $\alpha \notin G(R)$ , that for  $\alpha\Psi = \langle \alpha\Psi_1, \dots, \alpha\Psi_m \rangle$  we have  $p \in (\alpha\Psi)R$  and also  $T^{\alpha\Psi} \equiv R$ .

The class of relations established in the previous proposition, including those obtained by arbitrary permutations of numbers of base sets, is denoted by  $\mathbf{R}_1$  (similarly to  $n = 1$  [20]).

Consider incomplete  $s$ -base relations of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$ ,  $1 \leq m \leq s \leq n$ ,  $s \geq 2$ ,  $2 \leq h_i \leq k_i$  ( $i = 1, \dots, m$ ), having the form  $Q \equiv R \vee D_1 \& \dots \& D_m$ , where  $R$  is non-empty areflexive relation of the same arity as  $Q$ ,  $D_i$  is a single-base diagonal of arity  $h_i$  and sort  $i$  ( $1 \leq i \leq m$ ). Let  $G(D_i)$  be the symmetry group of  $D_i$ , i.e., the group of all permutations of coordinates preserving the equivalence relation  $\varepsilon(D_i)$  on the set of numbers of coordinates  $E(h_i)$  induced by equal, non-dummy coordinates in  $D_i$  ( $i = 1, \dots, m$ ). Denote,  $D_i(h_i)$  the diagonal on  $E(h_i)$  induced by the same equivalence relation:  $\varepsilon(D_i) \equiv \varepsilon(D_i(h_i))$ .

**Proposition 4.4.** *An  $s$ -base incomplete relation  $Q \equiv R \vee D_1 \& \dots \& D_m$  of arity  $\langle h_1, \dots, h_m, 1, \dots, 1, 0, \dots, 0 \rangle$ ,  $1 \leq m \leq s \leq n$ ,  $s \geq 2$ ,  $2 \leq h_i \leq k_i$  ( $i = 1, \dots, m$ ), is minimal if and only if:*

- (1)  $R$  is normal and  $G(R) \subseteq G(D_1) \times \dots \times G(D_m)$  (a sufficient condition for arity  $(k_1, \dots, k_s)$ );
- (2) For each non-empty subrelation  $M \subseteq R$  there exists a vector-epimorphism  $\Psi = \langle \Psi_1, \dots, \Psi_m \rangle$  such that  $\Psi M \subseteq \text{Orb}(G(R)) \cup D(h_1) \times \dots \times D(h_m)$  and  $\Psi M \cap \text{Orb}(G(R)) \neq \emptyset$ .

**Proof.** Part 1. Clearly the condition of Part 1 is equivalent to the fact that  $Q$  is irreducible. Next similarly to Proposition 4.3 one can show that this condition is the necessary for  $Q$  to be minimal. From Lemma 4.3 we obtain that it is also a sufficient condition for the arity  $(k_1, \dots, k_s)$ .

Part 2. Notice that each identification of  $T$  in Lemma 4.3 gives us either  $Q$  (up to permutations of coordinates of the same sort) or a diagonal (since  $Q$  is irreducible)

and it is equivalent to application of some vector-epimorphism  $\Psi$  to the index relation  $M$ . Next if  $r \in M$  and  $\Psi r \in D(h_1) \times \cdots \times D(h_m)$ , then  $Q^{\Psi r}$  is a full diagonal, which does not affect the result of identification. But if  $\Psi r \in D \setminus D(h_1) \times \cdots \times D(h_m)$ , where  $D$  is an  $m$ -base incomplete diagonal, then  $D^{\Psi r}$  is an incomplete diagonal itself and so  $T^\Psi \neq Q$ . Therefore, any reflexive part of  $\Psi M$  leading to  $T^\Psi = Q$  is included in  $D(h_1) \times \cdots \times D(h_m)$ .

It is obvious that the requirements of Part 2 imply the minimality of an irreducible relation  $Q$ . On the other side, if there exists  $\Psi$  such that  $T^\Psi = Q$  and  $p \notin \Psi R$ , then by using some vector-permutation  $\alpha$  on  $E(h_1) \times \cdots \times E(h_m)$ ,  $\alpha \notin G(R)$ , one can prove (similar to Proposition 4.3) that  $\alpha\Psi$  satisfies conditions of Part 2 and we have  $T^{\alpha\Psi} = Q$ .  $\square$

Denote  $\mathbf{R}_2$  the class of relations established in Proposition 4.4 including the ones obtained by permutations of numbers of base sets. Let  $\mathbf{B}(k_i)$  be the set of single-base relations of sort  $i$  determining all maximal partial clones on  $E(k_i)$ , except  $\Phi(k_i)$  ( $i = 1, \dots, n$ ) (see [20] or [7] and also [3]). Set  $\mathbf{B} = \mathbf{B}(k_1) \cup \cdots \cup \mathbf{B}(k_n)$ .

Finally, summarizing the results of the three sections we obtain the theorem.

**Theorem 4.1.** *Every maximal partial  $n$ -clone ( $n \geq 2$ ), except  $\Phi_i$  ( $i = 1, \dots, n$ ), is determined by a relation from classes  $\mathbf{K}$ ,  $\mathbf{H}$ ,  $\mathbf{R}_1$ ,  $\mathbf{R}_2$  and  $\mathbf{B}$ .*

**Corollary.** *A system of partial  $n$ -operations  $S$  is complete in  $\prod \mathbb{P}(k_i)$  if and only if:*

- (1) for every  $i$  ( $1 \leq i \leq n$ ) the set  $(S \setminus F)^i$  of all restrictions  $S \setminus F$  on its  $i$ th coordinate is complete in  $\mathbb{P}(k_i)$ ;
- (2) for every relation from classes  $\mathbf{K}$ ,  $\mathbf{H}$ ,  $\mathbf{R}_1$ , and  $\mathbf{R}_2$  the set  $S$  contains a partial  $n$ -operation not preserving it.

**Remark.** Note that the elements of  $F$  play the same role as empty operations for the case  $n=1$ . So, if we consider completeness criteria for  $\prod \mathbb{P}(k_i)$ , then elements of  $F$  are not supposed to be produced by compositions of partial  $n$ -operations from a complete set.

*Note:* (1) all relations from the above listed classes determine distinct partial  $n$ -clones unless they could be transposed to one another by some permutation of coordinates; (2) minimal relations have the minimum arity (comparing coordinatewise) among all relations determining the same maximal partial  $n$ -clone (for  $n=1$  see [22]).

## 5. Completeness in $\mathbb{P}(2) \times \cdots \times \mathbb{P}(2)$

We apply the previous results to vectors of partial Boolean functions (partial Boolean  $n$ -operations). In this case the description of maximal  $n$ -clones has a special simplified form that avoids the usage of epimorphic images. We introduce all these classes of minimal relations defined on  $n$  base sets  $E(2) = \{0, 1\}$ .

(1) Class **K** is the set of all nonempty, incomplete relations of arity  $(1, \dots, 1)$  having an arbitrary non-single type  $J \subseteq \{1, \dots, n\}$ ,  $2 \leq |J| \leq n$ , which cannot be reduced by  $\pi_i$  ( $i \in J$ ) to relations of smaller type.

Next according to Corollary 3.5 we may consider only relations of arity  $(2, \dots, 2, 1, \dots, 1, 0, \dots, 0)$  with the first  $m$  sorts having arity 2 and the next  $s - m$  sorts having arity 1 ( $1 \leq m \leq s \leq n$ ,  $s \geq 2$ ). Let  $H \in \mathbf{H}$  be a relation of the above arity. By Lemma 3.1 being totally reflexive in the case  $k_1 = 2$  means  $x_1 = x_2 \subset H$ . Moreover, if there exist two points  $(0, 1; q)$ ,  $(1, 0; q) \in H$ , where  $q$  is a tuple over the type  $\{2, \dots, s\}$ , then clearly that together with  $x_1 = x_2 \subset H$  we get  $x_1 = x_2 \vee q \subseteq H$ . Hence  $H \equiv x_1 = x_2 \vee H'$ , where  $H'$  has the type  $\{2, \dots, s\}$ .

(2) Class **H** consists of all relations having the form (as well as ones obtained from them by permutations of base sets):

$$x_1^1 = x_2^1 \vee \dots \vee x_1^m = x_2^m \vee K(x^{m+1}, \dots, x^s), \tag{15}$$

where either  $K \in \mathbf{K}$  of type  $\{m + 1, \dots, s\}$ , when  $s \geq m + 2$ , or  $K \in \{x^s = 0, x^s = 1\}$ , when  $s = m + 1$ , or  $K$  is void, when  $s = m$ .

(3) Note that in the Boolean case a vector-epimorphism  $\Psi$  from the Proposition 4.3 becomes a vector-isomorphism. Hence here each minimal  $R \in \mathbf{R}_1$  consists of only one block (orbit) of its group  $G(R)$  which in this case is a subgroup of the direct product  $S_2 \times \dots \times S_2$  of  $m$  symmetric groups  $S_2 = \{e, \alpha\}$  on  $E(2): \alpha: 0 \rightarrow 1, 1 \rightarrow 0$  and  $\alpha^2 = e$ . Notice that  $G(R)$  consists of vectors  $\langle \alpha_1, \dots, \alpha_m \rangle$ , where either  $\alpha_i = \alpha$  or  $\alpha_i = e$  ( $i = 1, \dots, m$ ). Next if  $G(R) = S_2 \times A$ , where  $A$  is a group over the type  $\{2, \dots, m\}$ , then  $R \equiv (x_1^1 \neq x_2^1) \& T$ , where  $T$  has the type  $\{2, \dots, m\}$ . Moreover, if we have  $\alpha_i = e$  for all elements  $\alpha \in G(R)$ , then  $R \equiv (x_1^i = 0 \& x_2^i = 1) \& T$ , where  $T$  has the type  $\{2, \dots, m\} \setminus \{i\}$ .

So for a group  $G(R)$  which is the direct product of  $S_2$  and the unit group  $\{e\}$ , i.e.,  $G(R) = S_2(1) \times \dots \times S_2(m)$ , where  $S_2(i) \in \{\{e\}, S_2\}$  ( $i = 1, \dots, m$ ), we have the presentation of the corresponding  $R \in \mathbf{R}_1$  (up to arbitrary permutations of base sets):

$$R \equiv R_1 \& \dots \& R_m \& K, \tag{16}$$

where  $R \in \{x_1^i = 0 \& x_2^i = 1, x_1^i \neq x_2^i\}$  ( $i = 1, \dots, m$ );  $K \in \mathbf{K}$  is of the type  $\{m + 1, \dots, s\}$ , when  $s \geq m + 2$ , or  $K \in \{x^s = 0, x^s = 1\}$ , when  $s = m + 1$ , or  $K$  is the full relation of the type  $\{1, \dots, m\}$ , when  $s = m$ .

Now consider the common case  $G(R) = S_2[t] \times S_2(t+1) \times \dots \times S_2(m)$ , where  $S_2[t]$  is a subdirect product of  $t$  groups  $S_2$  ( $2 \leq t \leq m$ ). Let  $\text{Orb}(S_2[t]) = \{(\alpha_1 0, \alpha_1 1; \dots; \alpha_t 0, \alpha_t 1): \langle \alpha_1, \dots, \alpha_t \rangle \in S_2[t]\}$  be the  $(2, \dots, 2)$ -orbit of this group.

Hence  $\mathbf{R}_1$  consists of all relations defined in (16) and also relations having the form (including those obtained by permutations of numbers of base sets):

$$\text{Orb}(S_2[t]) \& R, \tag{17}$$

where either  $R$  is a relation from (16) over the type  $\{t + 1, \dots, s\}$ , when  $t < m$ , or  $R \in \mathbf{K}$  over the type  $\{m + 1, \dots, s\}$ , when  $t = m < s$  and  $s \geq m + 2$ , or  $R \in \{x^s = 0, x^s = 1\}$ , when  $s = m + 1$  and  $t = m$ , or  $R$  is the full relation over the type  $\{1, \dots, m\}$ , when  $t = m = s$ .

**Example.** Let  $t = 3, m = 5, s = 7, S_2[3] = \{\langle e, e, e \rangle, \langle \alpha, e, \alpha \rangle, \langle e, \alpha, e \rangle, \langle \alpha, \alpha, \alpha \rangle\}$  and  $G(R) = S_2[3] \times S_2 \times \{e\}$  is the symmetry group of  $R$ . We have  $\text{Orb}(S_2[3]) = \{(0, 1; 0, 1; 0, 1), (0, 1; 1, 0; 0, 1), (1, 0; 0, 1; 1, 0), (1, 0; 1, 0; 1, 0)\}$ . Next we construct relations  $R \in \mathbf{R}_1$ :

$$R \equiv \text{Orb}(S_2[3]) \& x_1^4 \neq x_2^4 \& x_1^5 = 0 \& x_2^5 = 1 \& K,$$

where  $K \in \mathbf{K}$  has the type  $\{6, 7\}$ .

(4) From Proposition 4.4 we get that each  $Q \in \mathbf{R}_2$  is obtained from  $R \in \mathbf{R}_1$  using disjunction with an incomplete  $m$ -base diagonal of the same arity:

$$Q \equiv R \vee D_1 \& \dots \& D_m, \quad (18)$$

where  $R \in \mathbf{R}_1$ ,  $D_i$  is a single-base diagonal of the sort  $i$  ( $i = 1, \dots, m$ ) and at least one of  $D_i$  is the equality relation ( $1 \leq m \leq n$ ).

In total each  $R \in \mathbf{R}_1$  produces  $2^m - 1$  different relations  $Q \in \mathbf{R}_2$ .

(5) Recall that we have 7 single-base minimal relations (see [3]) over the type  $i$  ( $i = 1, \dots, n$ ):  $x = 0, x = 1, x_1 \neq x_2, x_1 \leq x_2, x_1 = 0 \& x_2 = 1, H_1 \equiv x = y \& u = z \vee x = u \& y = z$  and  $H_2 \equiv x = y \& u = z \vee x = u \& y = z \vee x = z \& y = u$ . The 8th maximal partial Boolean clone  $\Phi(2)$ , consisting of  $Q(2)$  and empty operations, produces the maximal partial  $n$ -clone  $\Phi_i$  ( $i = 1, \dots, n$ ).

Thus, we obtained the following theorem.

**Theorem 5.1.** *A system of partial Boolean  $n$ -operations  $S$  is complete in  $\mathbb{P}(2) \times \dots \times \mathbb{P}(2)$  if and only if:*

- (1) each coordinate set  $(S \setminus F)^i$  ( $i = 1, \dots, n$ ) is complete in  $\mathbb{P}(2)$ ;
- (2) for each relation from the classes (1)–(4)  $S$  contains a partial  $n$ -operation not preserving it.

Recall that all maximal  $n$ -clones of  $\mathbb{Q}(2) \times \dots \times \mathbb{Q}(2)$  were described in [19] by the following relations (another approach see in [28]):

- (a) single-base relations determining all 5 maximal clones on  $E(2)$  (see [13]):  $x = 0, x = 1, x_1 \neq x_2, x_1 \leq x_2$ , and  $H_2$  of the sort  $\{i\}$  ( $i = 1, \dots, n$ );
- (b) 2-base relations  $x = 0 \& y = 0 \vee x = 1 \& y = 1, x = 0 \& y = 1 \vee x = 1 \& y = 0$  for all pairs of different sorts from  $\{1, \dots, n\}$ ;
- (c)  $s$ -base relations ( $1 \leq h \leq s \leq n, s \geq 2$ ) of the form (including the ones obtained by permutations of numbers of base sets):

$$x_1^1 = x_2^1 \vee \dots \vee x_1^h = x_2^h \vee R_{h+1} \vee \dots \vee R_s,$$

where  $R_i \in \{x^i = 0, x^i = 1\}$  ( $i = h + 1, \dots, s$ ).

Clearly class (a) is included in (5) and relations from (b) and (c) are contained in  $\mathbf{K} \cup \mathbf{H}$ .

**Corollary 5.1.** *Each relation from classes (a), (b) and (c) determining maximal  $n$ -clone of Boolean functions also determines maximal partial  $n$ -clone of partial Boolean functions.*

Case  $\mathbb{P}(2) \times \mathbb{P}(2)$ :

Applying the results of this section we describe all 67 maximal partial 2-clones of Boolean operations, i.e., all maximal iterative Post subalgebras in the system of all pairs of partial Boolean functions.

(1) Considering class **K** we get 10 minimal double-base relations:  $x = a \& y = b, x = a \& y = b, \text{ where } a, b \in \{0, 1\}, x = 0 \& y = 0 \vee x = 1 \& y = 1, x = 0 \& y = 1 \vee x = 1 \& y = 0.$

(2) Class **H** contributes 5 relations:  $x_1 = x_2 \vee y_1 = y_2, x = 0 \vee y_1 = y_2, x = 1 \vee y_1 = y_2, x_1 = x_2 \vee y = 0, x_1 = x_2 \vee y = 1.$

(3) Classes **R**<sub>1</sub> and **R**<sub>2</sub> give 20 relations of arity (2,2):  $R_1 \equiv x_1 = 0 \& x_2 = 1 \& y_1 = 0 \& y_2 = 1, R_2 \equiv x_1 = 0 \& x_2 = 1 \& y_1 \neq y_2, R_3 \equiv x_1 \neq x_2 \& y_1 = 0 \& y_2 = 1, R_4 \equiv x_1 \neq x_2 \& y_1 \neq y_2, R_5 \equiv x_1 = 0 \& x_2 = 1 \& y_1 = 0 \& y_2 = 1 \vee x_1 = 1 \& x_2 = 0 \& y_1 = 1 \& y_2 = 0.$  And also  $R_i \vee D$  ( $i = 1, \dots, 5$ ), where  $D \in \{x_1 = x_2, y_1 = y_2, x_1 = x_2 \& y_1 = y_2\}.$

(4) There are also 16 relations from **R**<sub>1</sub>  $\cup$  **R**<sub>2</sub> of arity (2,1) and (1,2):  $Q_1 \equiv x_1 = 0 \& y_1 = 0 \& y_2 = 1, Q_2 \equiv x_2 = 1 \& y_1 = 0 \& y_2 = 1, Q_3 \equiv x = 0 \& y_1 \neq y_2, Q_4 \equiv x = 1 \& y_1 \neq y_2, Q_i \vee y_1 = y_2$  ( $i = 1, \dots, 4$ )-yields 8 relations. Interchanging  $x$  and  $y$  we obtain 8 relations of arity (2,1).

(5) Add 16 partial 2-clones of the form  $(A \times P(2)) \cup F$  and  $(P(2) \times A) \cup F$ , where  $A$  is maximal partial clone on  $E(2)$  (see [4]).

Finally, we get in total 67 maximal partial 2-clones of Boolean functions.

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