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Note

A Simple Way to Tell a Simple Polytope from Its Graph

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Let P be a simple d-dimensional polytope and let G(P) be the graph of P. Thus, G(P) is an abstract graph defined in the set of vertices V(P) of P. Two vertices v and u in V(P) are adjacent in G(P) if [v, u] is a 1-dimensional face of P. Perles [P] conjectured and Blind and Mani [BM] recently proved that G(P) determines the entire combinatorial structure of P. Here is a simple proof of this result. Let f denote the number of non-empty faces of P.

We consider the class of acylic orientations (i.e., edge orientations with no oriented cycles) of G(P). We will not distinguish between an acyclic orientation O of G(P) and the partial order induced by O on V(P). $(x \leq_O y)$ iff there is an O-oriented path from x to y.) Note that if O is an acyclic orientation of G(P) then the restriction of G(P) to any non-empty subset A of V(P) has a sink (=element with out-degree zero) with respect to O.

An acyclic orientation O of G(P) is good if for every non-empty face F of P, G(F) has exactly one sink. Otherwise, O is bad. The existence of good acyclic orientations of G(P) is well known. Good acyclic orientations are obtained, e.g., by orienting the edges according to the value of a linear functional on \mathbb{R}^d that is 1–1 on V(P); see [B, Sect. 15]. Our first goal is to distinguish intrinsically between good and bad orientations of G(P).

Let O be an acyclic orientation of G(P). Let h_k^O be the number of vertices of G(P) with indegree k in O. Define

$$f^{O} = h_{0}^{O} + 2h_{1}^{O} + 4h_{2}^{O} + \dots + 2^{k}h_{k}^{O} + \dots + 2^{d}h_{d}^{O}.$$

If x is a vertex of G(P) of indegree k w.r.t. O then x is a sink in 2^k faces of P. (Every *i* edges incident to x determine an *i*-face F of P which includes them.) Since each face has at least one sink we obtain that

- (I) $f^{O} \ge f$, and
- (II) O is good if and only if $f^O = f$.

To distinguish between good and bad orientations from the knowledge of G(P) only, compute f^{O} for every acyclic orientation O. The good acyclic orientations of G(P) are those having the minimal value of f^{O} .

Now we will show how to identify the faces of P. The criterion is very simple: An induced connected k-regular subgraph H of G is the graph of some k-face of P if and only if its vertices are initial w.r.t. some good acyclic orientation O of G(P). Indeed, if F is a face of P, it is well known that V(F) is an initial set with respect to some good acyclic orientation: just consider a linear functional with respect to which the vertices of F lie below all other vertices. (See [B, Sect. 18].) On the other hand, let H be a connected k-regular subgraph of G(P) and let O be a good acyclic orientation with respect to which V(H) is an initial set. Let x be a sink of H with respect to O. There are k edges containing x in H. all oriented towards x. Therefore x is a sink in a k-face F that contains these k edges. Since the orientation O is good, x is the unique sink of F, and therefore all vertices of F are $\leq x$, with respect to O. But V(H) includes the set of all vertices that are $\leq x$ with respect to O. (Remember: V(H) is an initial set with respect to O.) Thus, $V(F) \subset V(H)$. Since both H and G(F) are k-regular and connected, V(F) = V(H) and G(F) = H. This completes the proof.

Remarks. 1. We do not have a practical way to distinguish between good and bad orientations. The algorithm suggested by the proof above is exponantial in |V(P)|. We do not know of an efficient way even for computing the face numbers of P from G(P).

2. It was observed already by Perles that the 2-skeleton of P determines P up to combinatorial isomorphism. His observation is based on the following fact: Let x and y be adjacent vertices in G(P) and let F be the facet of P containing x but not y. Let z be a vertex adjacent to $x, z \neq y$. It is easy to identify the unique vertex w which is adjacent to z and does not belong to F. Let M be the (unique) 2-face of P containing x, y, and z. Then w is the vertex adjacent to z in M, different from x. This gives a quick way to identify the facets of P, hence the entire combinatorial structure of P, from the 2-skeleton of P. Perles also observed that all induced 3-gons, 4-gons, and 5-gons in G(P) correspond to 2-faces of P.

3. Perles [P] proved that simplical *d*-polytopes are determined by their $\lfloor d/2 \rfloor$ -skeleton. (Dancis [D] extended this result to a large class of simplicial manifolds.) Perles also proved that simple polytopes are determined by the incidence relations between their 1-faces and 2-faces. The proof described above can be extended to show that the combinatorial structure of a simple *d*-polytope is determined by the incidence relations

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between its *i*-faces and (i + 1)-faces, whenever $i < \lfloor d/2 \rfloor$. It is also possible to show that (d-k)-simple polytopes are determined by their k-skeleton. (P is (d-k)-simple if every (k-1)-face is included in exactly d-k+1facets.) Details will appear elsewhere. (Note that general d-polytopes are determined by their (d-2)-skeleton, and this is best possible even for quasi-simplicial polytopes, [G, Chap. 12].)

4. Perles asked whether every connected (d-1)-regular subgraph of G(P) which does not separate G(P) is the graph of a facet of P. This is still unknown.

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