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# Paths in *k*-Edge-Connected Graphs

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We prove (i) if G is a 2k-edge-connected graph  $(k \ge 2)$ , s, t are vertices, and  $f_1, f_2, g$  are edges with  $f_i \ne g$  (i = 1, 2), then there exists a cycle C passing through  $f_1$  and  $f_2$  (a path P between s and t passing through  $f_1$ ) but not passing through g such that G - E(C)(G - E(P)) is (2k-2)-edge-connected, where C and P are not necessarily simple and E(C) is the set of edges of C. (ii) Every 3k-edge-connected graph  $(k \ge 1)$  is weakly (2k+1)-linked and every (3k-1)-edge-connected graph  $(k \ge 2)$  is weakly 2k-linked.  $\bigcirc$  1988 Academic Press, Inc.

### 1. INTRODUCTION

We consider finite undirected graphs possibly with multiple edges but without loops. Let G be a graph and let V(G) and E(G) be the set of vertices and edges of G, respectively. We allow repetition of vertices (but not edges) in a path and cycle. Mader [3] conjectured,

If G is a k-edge-connected graph  $(k \ge 4)$  and s, t are vertices of G, then there exists a cycle C passing through s and t such that G - E(C) is (k-2)-edge-connected,

and proved this conjecture (in fact Theorem 1(1)) for k = 4. We prove this when k is even.

THEOREM 1. Suppose that  $k \ge 4$  is an even integer, G is a k-edgeconnected graph,  $\{s, t\} \subset V(G)$ ,  $\{f_1, f_2, g\} \subset E(G)$ , and  $f_i \ne g$  (i=1, 2). Then

(1) There exists a cycle C passing through  $f_1$  and  $f_2$  but not passing through g such that G - E(C) is (k-2)-edge-connected.

(2) There exists a path P between s and t passing through  $f_1$  but not passing through g such that G - E(P) is (k-2)-edge-connected.

For odd k, the conjecture of Mader is still open, but the result of Theorem 1(1) does not always hold; Fig. 1 gives a counterexample.

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FIGURE 1

We call a graph G weakly k-linked, if for every k pairs of vertices  $(s_i, t_i)$ , there exist edge-disjoint paths  $P_1, ..., P_k$  such that  $P_i$  joins  $s_i$  and  $t_i$   $(1 \le i \le k)$ . Let

 $g(k) := \min\{m \mid \text{if } G \text{ is } m \text{-edge-connected, then } G \text{ is weakly } k \text{-linked}\}.$ 

Thomassen [6] conjectured

$$g(2k+1) = g(2k) = 2k+1 \qquad (k \ge 1).$$

The author [4] proved g(3) = 3 and Hirata, Kubota, and Saito [1] and Mader [3] proved g(4) = 5 and  $g(k) \le 2k - 3$  ( $k \ge 5$ ). We prove

THEOREM 2.  $g(2k+1) \leq 3k \ (k \geq 1)$  and  $g(2k) \leq 3k-1 \ (k \geq 2)$ .

# Notations and Definitions

 $\lambda(G)$  denotes the edge-connectivity of G. Let X, Y,  $\{x, y\} \subset V(G)$ ,  $f \in E(G)$ , and  $X \cap Y = \emptyset$ . We often denote  $\{x\}$  by x. V(f) denotes the set of end vertices of f. We denote by  $\partial(X, Y; G)$  the set of edges with one end in X and the other in Y, and set  $\partial(X; G) := \partial(X, V(G) - X; G)$ ,  $e(X, Y; G) := |\partial(X, Y; G)|$ , and  $e(X; G) := |\partial(X, V(G) - X; G)|$ .  $\lambda(X, Y; G)$  denotes the maximal number of edge-disjoint paths between X and Y. We set  $\overline{X} := V(G) - X$ ,  $N(x; G) := \{a \in V(G) - x | e(a, x) > 0\}$ ,  $N(X; G) := \bigcup_{x \in X} N(x; G)$ , and  $\Gamma(G, k) := \{Z \subset V(G) | \text{ for each } a, b \in Z, \lambda(a, b; G) \ge k\}$ . In all notations, we often omit G. G/X denotes the graph obtained from G by contracting X, and for  $a \in X$ , we denote the corresponding vertex in G/X by  $\tilde{a}$ . A path P = P[x, y] denotes a path between x and y. For a,  $b \in N(x)$  with  $a \neq b$ ,  $f \in \partial(x, a)$ , and  $g \in \partial(x, b)$ ,  $G_x^{a,b}$  denotes the graph  $(V(G), (E(G) \cup h) - \{f, g\})$ , where h is a new edge between a and b and is called a lifting of G at x arising from the lifting of f and g at x. We call  $G_x^{a,b}$  admissible if for each  $y, z \in V(G) - x$  with  $y \neq z$ ,  $\lambda(y, z; G_x^{a,b}) = \lambda(y, z; G)$ .

Throughout the paper we shall make use of the following observation:

If  $a, b \in X \subset V(G)$ , then  $\lambda(a, b; G) \ge \lambda(a, b; G/\overline{X})$  and it is easy to give examples for which the inequality is strict. However, the inequality is an equality if  $\overline{X}$  has a vertex z and a collection of e(X) edge-disjoint paths from z to X.

# 2. Proof of Theorem 1

LEMMA 1 (Okamura [5]). Suppose that  $k \ge 4$ , G is a 2-edge-connected graph, and  $\{s, t\} \subset T \in \Gamma(G, k)$ . Then

(1) If  $a \in T - \{s, t\}$  and e(a) < 2k, then there exists a path P[s, t] such that  $a \notin V(P)$ ,  $T \in \Gamma(G - E(P), k - 2)$ , and  $\{s, t, a\} \in \Gamma(G - E(P), k - 1)$ .

(2) If  $a \in V(G)$  and  $\lambda(a, s) < k$ , then there exists a path P[s, t] such that  $a \notin V(P)$ ,  $T \in \Gamma(G - E(P), k - 2)$ , and  $\lambda(a, s; G - E(P)) = \lambda(a, s; G)$ .

(3) If  $f_1, f_2 \in \partial(s)$ , then there exists a cycle C such that  $\{f_1, f_2\} \subset E(C)$  and  $T \in \Gamma(G - E(C), k - 2)$ .

LEMMA 2 (Mader [2]). If G is a graph,  $x \in V(G)$ ,  $e(x) \ge 4$ ,  $|N(x)| \ge 2$ , and x is not a cut-vertex, then there exists an admissible lifting of G at x.

LEMMA 3. If  $k \ge 3$  is an integer, G is a graph,  $V(G) = W_1 \cup W_2$ ,  $W_1 \cap W_2 = \emptyset$ ,  $W_1 \in \Gamma(G, k)$ , and each  $x \in W_2$  has even degree, then we can obtain a k-edge-connected graph  $G(W_1, k)$  from G such that  $W_1 \subset V(G(W_1, k))$  by sequences of vertex-deletions and edge-liftings.

*Proof.* We may assume G is connected. We can obtain  $G(W_1, k)$  from G as follows.

Step 1. If  $W_2 = \emptyset$ , then let  $G(W_1, k) := G$ , and stop.

Step 2. If for an  $x \in W_2$ , |N(x)| = 1, then reset G := G - x and go to Step 1.

Step 3. If for each  $x \in W_2$ , e(x) = 2, then let  $G(W_1, k)$  be the k-edgeconnected graph homeomorphic to G (that is at each  $x \in W_2$ , we lift  $\partial(x)$ and then delete x), and stop. Otherwise let  $x \in W_2$  and  $e(x) \ge 4$ .

Step 4. If x is a cut-vertex and each component of G-x has a vertex of  $W_1$ , then reset  $W_1 := W_1 \cup \{x\}$ ,  $W_2 := W_2 - x$  (note that  $W_1 \cup \{x\} \in \Gamma(G, k)$ ), and go to Step 1. If x is a cut-vertex and C is a component of G-x such that  $V(C) \cap W_1 = \emptyset$ , then reset G := G - V(C), and go to Step 1. Otherwise let  $G_x$  be an admissible lifting of G at x (see Lemma 2), reset  $G := G_x$ , and go to Step 1.

In what follows  $G \rightarrow G(W_1, k)$  denotes this operation.

LEMMA 4. Suppose that  $k \ge 4$  is an even integer, G is a k-edge-connected graph, and for each  $x \in V(G)$ ,  $e(x) \le k + 1$ . If  $f \ne g \in E(G)$ , then there exists a cycle C such that  $f \in E(C)$ ,  $g \notin E(C)$ , and  $\lambda(G - E(C)) \ge k - 2$ .

*Proof.* It is easy to see that G is 2-connected. We proceed by induction on |E(G)|. If there is an  $x \in V(f) \cap V(g)$ , then for any  $h \in \partial(x) - \{f, g\}$ , by Lemma 1(3) G has a cycle C such that  $\{f, h\} \subset E(C)$  and  $\lambda(G - E(C)) \ge k - 2$ . Then  $g \notin E(C)$ . Thus let  $V(f) \cap V(g) = \emptyset$ . Let  $V(g) = \{x, y\}$ . If e(x) = k, then let  $G_x$  be an admissible lifting at x (see Lemma 2) and let  $G_1 := G_x(V(G) - x, k)$  (see Lemma 3). In  $G_1 f \neq g$ , and by induction  $G_1$  has a cycle C such that  $f \in E(C)$ ,  $g \notin E(C)$ , and  $\lambda(G_1 - E(C)) \ge k - 2$ . Let  $C_1$  be the corresponding cycle in G and let  $C_2$  be the simple subcycle of  $C_1$  containing f, then  $C_2$  is a required cycle. Thus let e(x) = e(y) = k + 1. If  $\lambda(G - g) \ge k$ , then the result holds in G - g. Thus for some  $X \subset V(G) - y$ ,  $x \in X$  and e(X) = k. Let  $V(f) = \{a, b\}$ , we may let  $X \cap V(f) = \emptyset$  or  $\{a\}$ . If  $X \cap V(f) = \emptyset$ , then G/X has a required cycle C. If  $\tilde{x} \notin V(C)$ , then C is a required cycle for G. If  $\tilde{x} \in V(C)$ , then let  $\partial(\tilde{x}; G/X) \cap E(C) = \{h_1, h_2\}$ . By Lemma 1(3)  $G/\overline{X}$  has a cycle  $C_1$  such that  ${h_1, h_2} \subset E(C_1)$  and  $\lambda(G/\overline{X} - E(C_1)) \ge k - 2$ . Then  $g \notin E(C_1)$ . By combining C and C<sub>1</sub> in G we have a required cycle. Therefore  $X \cap V(f) = \{a\}$ . Let  $h \in \partial(X) - \{f, g\}, G_1 := G/X, \text{ and } G_2 := G/\overline{X}$ . By Lemma 1(3) for i = 1, 2, J $G_i$  has a cycle  $C_i$  such that  $\{f, h\} \subset E(C_i)$  and  $\lambda(G_i - E(C_i)) \ge k - 2$ . By combining  $C_1$  and  $C_2$  in G we have a required cycle.

The proof of Lemma 5 will be given later.

LEMMA 5. Suppose that  $k \ge 4$  is an even integer and G is a graph. If

(i)  $V(G) = \{u\} \cup A \cup W_1 \cup W_2$  (disjoint union),  $A \neq \emptyset$ , and either  $W_2 = \emptyset$  or  $W_2 = \{b\}$  and b has even degree,

(ii)  $V(G) - W_2 \in \Gamma(G, k-2)$ , and for each  $X \subset W_1 \cup W_2$  with  $X \cap W_1 \neq \emptyset$ ,  $e(X) \ge k$ ,

(iii) for each  $x \in W_1$ ,  $e(x) \leq k+1$ , and for each  $x \in V(G) - W_1$ ,  $e(x) \leq k-1$ .

(iv)  $f \in \partial(u), g \in E(G) - f, and \{f, g\} \neq \partial(b),$ 

then for some  $a \in A$ , there exists a path P[u, a] such that  $f \in E(P)$ ,  $g \notin E(P)$ , and  $V(G) - W_2 \in \Gamma(G * P, k-2)$ . Here G \* P denotes the graph  $(V(G), (E(G) \cup h) - E(P))$  and h is a new edge between u and a.

# Proof of Theorem 1

First we prove that Theorem 1(1) implies Theorem 1(2). Let h be a new edge between s and t and let  $G_1 := (V(G), E(G) \cup h)$ . Then by (1),  $G_1$  has a cycle C such that  $\{h, f_1\} \subset E(C), g \notin E(C)$ , and  $\lambda(G_1 - E(C)) \ge k - 2$ . C - h is a required path of G.

Now we prove Theorem 1(1) by induction on |E(G)|. If G is not 2-connected, then we can deduce the result by using induction on some blocks. Thus we may assume that G is 2-connected. If for an  $x \in V(G)$ ,  $e(x) \ge k+2$ , then by Lemma 2 for some  $a, b \in N(x)$  with  $a \ne b$ ,  $G_x^{a,b} =: G_1$  is admissible. If  $f_i \ne g$  (i=1,2) in  $G_1$ , then by induction the result holds in  $G_1$ . Thus let  $\partial(a, x; G) = \{f_1\}$  and  $\partial(b, x; G) = \{g\}$ .  $|N(x; G_1)| \ge 2$ , otherwise for each  $h_1, h_2 \in \partial(x; G_1) - f_2$ , the result holds in  $G - \{h_1, h_2\}$ . Thus for some  $y, z \in N(x; G_1)$  with  $y \ne z$ ,  $(G_1)_x^{y,z}$  is admissible,  $f_i \ne g$  (i=1, 2) in  $G_x^{y,z}$ , and the result holds in  $G_x^{y,z}$ . Thus we may assume that for each  $x \in V(G)$ ,  $e(x) \le k + 1$ . Let

$$F := \{h \in E(G) \mid \text{there is a cycle } C \text{ such that } \{h, f_1\} \subset E(C), \\ g \notin E(P), \text{ and } \lambda(G - E(C)) \ge k - 2\}.$$

Assume  $F \neq E(G) - g$ . By Lemma 4  $F \neq \emptyset$ . For some  $h \in F$  and  $f \in E(G) - F - g$ , h and f have a common end vertex, say u. Let C be a cycle such that  $\{h, f_1\} \subset E(C), g \notin E(C), and \lambda(G - E(P_1)) \ge k - 2$ . Let  $A := V(C) - u, W_1 := V(G) - V(C)$ , and  $W_2 = \emptyset$ . Then by Lemma 5 for some  $a \in A, G - E(C)$  has a path P[u, a] such that  $f \in E(P), g \notin E(P)$ , and (G - E(C)) \* P is (k - 2)-edge-connected. In C there are two disjoint paths joining u and a. Let  $P_1$  be one of them containing  $f_1$ , and let  $C_1 := P \cup P_1$ , then  $\lambda(G - E(C_1)) \ge k - 2$ . Thus F = E(G) - g.

To prove Lemma 5 we need some lemmas.

LEMMA 6. Suppose that  $k \ge 2$  is an even integer, G is a graph,  $V(G) = W_1 \cup W_2$ ,  $W_1 \cap W_2 = \emptyset$ , and each  $v \in W_2$  has even degree. Then

(1) If  $W_1 \in \Gamma(G, k)$ ,  $x \in V(G)$ , and e(x) is odd, then for some  $y \in W_1$ ,  $\lambda(x, y) \ge k + 1$ .

(2) If  $X \subset V(G)$ ,  $e(X) \leq k+1$ ,  $x \in X \cap W_1$ ,  $y \in \overline{X} \cap W_1$ ,  $(X \cap W_1) \cup \{\tilde{y}\} \in \Gamma(G/\overline{X}, k)$ , and  $(\overline{X} \cap W_1) \cup \{\tilde{x}\} \in \Gamma(G/\overline{X}, k)$ , then  $W_1 \in \Gamma(G, k)$ .

*Proof.* (1) Since  $x \notin W_2$ ,  $e(x) \ge k+1$ . If for some  $X \subset V(G)$  with  $x \in X$ , e(X) = k, then choose X with this property such that |X| is minimal, if not, then let X := V(G). For some  $y \in X - x$ , e(y) is odd. Then  $y \in W_1$ ,  $e(y) \ge k+1$ , and  $\lambda(y, x) \ge k+1$ .

(2) Clearly we may let e(X) = k + 1. By (1) for some  $a \in X$ ,  $\lambda(a, \overline{X}) = k + 1$  and for some  $b \in \overline{X}$ ,  $\lambda(b, X) = k + 1$ , and so  $W_1 \in \Gamma(G, k)$ .

LEMMA 7. Suppose that G is a graph and X,  $Y \subset V(G)$ . Then

(1)

$$e(X - Y) + e(Y - X) = e(X) + e(Y) - 2e(X \cap Y, \overline{X \cup Y}),$$
  
$$e(X \cap Y) + e(X \cup Y) = e(X) + e(Y) - 2e(X - Y, Y - X).$$

(2) If  $\lambda(G) \ge k$ , X - Y, Y - X,  $X \cap Y$ , and  $\overline{X \cup Y}$  are not empty and e(X) = e(Y) = k, then k is even and  $e(X - Y) = e(X \cap Y) = k$ .

*Proof.* (1) Simple counting.

(2) By (1)  $e(X - Y) = e(X \cap Y) = k$ . Thus  $k = e(X) \equiv e(X - Y) + e(X \cap Y) \equiv 0 \pmod{2}$ .

# Proof of Lemma 5

We proceed by induction on |E(G)|. Let  $\mathscr{P}(G, f, A, W_1)$  be the set of required paths of G and assume  $\mathscr{P}(G, f, A, W_1) = \emptyset$ . Note that  $e(b) \ge 4$  if  $W_2 \ne \emptyset$ . Let  $V(f) = \{u, v\}$ . Then  $v \notin A$ .

(2.1) If  $X \subset W_1 \cup W_2$  and  $|X| \ge 2$ , then  $e(X) \ge k+2$ .

*Proof.* Assume  $e(X) \leq k+1$ . Let  $x \in X$  and let  $P[u, a] \in \mathscr{P}(G/X, f, A, (W_1 - X) \cup \tilde{x})$  (g might not be in E(G/X)). If  $\tilde{x} \notin V(P)$ , then we may let e(X) = k+1 and by Lemma 6(1) for some  $y \in X$ ,  $\lambda(y, \overline{X}) = k+1$  and for some  $z \in \overline{X} - W_2$ ,  $\lambda(z, X; G * P) \geq k-1$ , and so  $V(G) - W_2 \in \Gamma(G * P, k-2)$ . Thus let  $\tilde{x} \in V(P)$  and  $h_1, h_2 \in \partial(\tilde{x}; G/X) \cap E(P)$ . If  $X \subset W_1$ , then let  $G_1 := G/\overline{X}$ , and if  $b \in X$ , then for  $Y := V(G/\overline{X}) - b$ , let  $G_1 := (G/\overline{X})(Y, k)$  (see Lemma 3). By induction for each graph H such that |E(H)| < |E(G)|, Lemma 5 holds, and so in  $G_1$  Theorem1(1) holds (see the proof of Theorem 1). Thus  $G_1$  has a cycle C such that  $\{h_1, h_2\} \subset E(C), g \notin E(C)$ , and  $\lambda(G_1 - E(C)) \geq k - 2$ . Let  $C_1$  be the corresponding cycle in  $G/\overline{X}$ . Let  $P_1[u, a]$  be the path in G obtained by combining P and  $C_1$ . Then by Lemma 6  $P_1$  is a required path.

(2.2) If  $X \subset V(G)$ ,  $|X| \ge 2$ , and  $|\overline{X}| \ge 2$ , then  $e(X) \ge k$ .

*Proof.* Assume e(X) = k-2 or k-1 and  $v \notin X$ . If  $u \in X$ , then by (ii)  $A - X \neq \emptyset$  and  $\emptyset \neq \mathscr{P}(G/X, f, A - X, W_1 - X) \subset \mathscr{P}(G, f, A, W_1)$  by induction and by Lemma 6. If  $u \notin X$ , then  $A \cap X \neq \emptyset$ . Let  $x \in X$  and  $P_1[u, a] \in \mathscr{P}(G/X, f, (A - X) \cup \tilde{x}, W_1 - X)$ . If  $a \neq \tilde{x}$ , then  $P_1$  is a required path, thus let  $a = \tilde{x}$  and  $h \in \partial(\tilde{x}; G/X) \cap E(P_1)$ . Let  $P_2 \in \mathscr{P}(G/\overline{X}, h, A \cap X, W_1 \cap X)$ . By combining  $P_1$  and  $P_2$  in G we can get a required path.

(2.3) Let x = b if  $W_2 \neq \emptyset$  and  $x \in W_1$  if  $W_2 = \emptyset$ . If |N(x)| = n and  $N(x) = \{y_1, ..., y_n\}$ , then  $n \ge 2$  and  $V(G) - x \in \Gamma(G_x^{y_1, y_i}, k-2)$   $(2 \le i \le n)$ . Moreover if  $e(x, y_1) = 1$ , then  $n \ge 3$ .

*Proof.* If n = 1, then  $\lambda(G - x) \ge k - 2$  and  $e(y_1; G - x) \le k + 1 - 4$ , a contradiction. Thus  $n \ge 2$ . By (2.2) for each  $2 \le i \le n$ ,  $V(G) - x \in \Gamma(G_x^{y_1, y_1}, k - 2)$ . Assume  $e(x, y_n) = 1$  and n = 2. Then  $e(\{x, y_1\}) \le e(y_1) - 2$ , contrary to (ii) or (iii).

(2.4)  $W_2 = \emptyset$  and for each  $x \in W_1$ , e(x) = k + 1.

*Proof.* Let x = b if  $W_2 \neq \emptyset$ , and let  $x \in W_1$  and e(x) = k if  $W_2 = \emptyset$ . By (2.3) we can obtain a lifting  $G_x$  of G at x such that  $V(G) - x \in \Gamma(G_x, k-2)$ ,  $f \neq g$  in  $G_x$ , and  $\{f, g\} \neq \partial(x; G_x)$ . By (2.1) for each  $X \subset W_1 \cup W_2$  with  $X \neq \{x\}, e(X; G_x) \ge k$ , and so there is a  $P \in \mathscr{P}(G_x, f, A, W_1 - x)$ . If x = b, then  $P \in \mathscr{P}(G, f, A, W_1)$ . If  $x \in W_1$ , then let  $P_1$  be the corresponding path in G. If  $P_1$  is not simple, then let  $P_2$  be the simple subpath of  $P_1$  between u and a, then  $P_2 \in \mathscr{P}(G, f, A, W_1)$ .

Since  $v \in W_1$ , if for some  $x \in W_1$  there is an  $h \in \partial(v, x)$ , then by (2.1), (2.2), and (2.4) the result holds in G-h. Thus  $\{u\} \neq N(v) \subset \{u\} \cup A$ . By (2.3) for some  $a \in N(v) - u$ , there is a lifting  $G_v^{u,a}$  such that  $\lambda(G_v^{u,a}) \ge k-2$ and  $f \neq g$  in  $G_v^{u,a}$ .

#### 3. PROOF OF THEOREM 2

The proof of Lemma 8 will be given later.

LEMMA 8. Suppose that  $k \ge 4$  and  $n \ge 2$  are integers, G is a 2-connected graph,  $V(G) = T \cup W_1 \cup W_2$  (disjoint union),  $T = \{s_1, ..., s_n, t_1, ..., t_n\}$ ,  $|T| = 2n, T \cup W_1 \in \Gamma(G, k), |W_2| \le 2$ , for each  $x \in W_2$ ,  $e(x) \le k - 1$  is even, and

(i) if k is odd, then for each  $x \in T \cup W_1$ , e(x) = k,

(ii) if k is even, then  $e(s_i) = e(t_i) = k$   $(1 \le i \le n)$  and for each  $x \in W_1$ , e(x) = k or k + 1.

Then there is a subgraph  $G^* \subset G$  such that

(a) for some  $1 \le i < j \le n$ ,  $G - E(G^*)$  has edge-disjoint paths  $P_1[s_i, t_i]$ and  $P_2[s_j, t_j]$ ,

(b)  $V(G^*) = K_1 \cup K_2$  and  $K_1 \cap K_2 = \emptyset$ ,

- (c)  $T \{s_i, t_i, s_i, t_i\} \subset K_1 \in \Gamma(G^*, k-3),$
- (d) for each  $x \in K_2$ ,  $e(x; G^*)$  is even.

#### Proof of Theorem 2

Let  $\alpha = 0$  or 1,  $m \ge 1$  be an integer,  $k := 3m - \alpha \ge 3$ , and  $n := 2m + 1 - \alpha$ . Assume that G is a k-edge-connected graph and  $\{s_1, ..., s_n, t_1, ..., t_n\} := T$  are vertices of G (not necessarily distinct). We prove that there are edgedisjoint paths  $P_1, ..., P_n$  such that  $P_i$  joins  $s_i$  and  $t_i$   $(1 \le i \le n)$ . We may assume that G is 2-connected. If k is odd and e(x) =: d > k for some  $x \in V(G)$ , then we replace x by d vertices of degree k (Fig. 2 gives an



example with d=8 and k=5) and assign x on one of the new vertices, producing a new graph  $G_1$ . If the result holds in  $G_1$ , then it also holds for G. Thus we may assume

# (3.1) If k is odd, then G is k-regular.

If k is even and  $e(s_1)$  is odd, then we replace  $s_1$  by two vertices of degree  $e(s_1)-1$  and  $e(s_1)$ , and assign  $s_1$  on the new vertex of degree  $e(s_1)-1$ , producing a new graph  $G_1$  (see Fig. 3). If the result holds in  $G_1$ , then it also holds for G. Thus we may assume that  $e(s_i)$  and  $e(t_i)$   $(1 \le i \le n)$  are even. We proceed by induction on |E(G)|. If k is even,  $x \in V(G)$ , and  $e(x) \ge k+2$ , then an admissible lifting  $G_x$  of G at x is k-edge-connected and the result holds in  $G_x$ . Thus

(3.2) If k is even, then  $e(s_i) = e(t_i) = k$   $(1 \le i \le n)$  and for each  $x \in V(G)$ , e(x) = k or k + 1.

By [4, 1, 3] g(3) = 3 and g(4) = 5, and so we may let  $k \ge 6$ ,  $n \ge 5$ , and  $m \ge 2$ . If  $s_1 = s_2$ , then by Theorem 1 there is a path  $P[t_1, t_2]$  such that  $s_1 \in V(P)$  and  $\lambda(G - E(P)) \ge k - 3$ . By induction G - E(P) has edge-disjoint



FIGURE 3

paths  $P_3[s_3, t_3], ..., P_n[s_n, t_n]$ . Thus let |T| = 2n. By Lemma 8 there is a subgraph  $G^* \subset G$  such that (a), (b), (c), and (d) hold. By Lemma 3  $G^*(K_1, k-3)$  is (k-3)-edge-connected, and by induction  $G^*(K_1, k-3)$  has (n-2) edge-disjoint paths joining  $(s_l, t_l)$   $(1 \le l \le n, l \ne i, j)$ . Thus the result holds in G.

#### Proof of Lemma 8

Suppose that G satisfies the hypothesis of Lemma 8, but the result does not hold. Choose G with this property such that |E(G)| is minimal.

(3.3)  $W_2 = \emptyset$ .

*Proof.* Assume  $x \in W_2$ . Then  $e(x) \ge 4$ . By Lemma 2 we have an admissible lifting  $G_x$  of G at x. The result holds in  $G_x$  and so in G.

Case 1. k is odd.

(3.4) If 
$$x, y \in W_1$$
 and  $f \in \partial(x, y)$ , then  $V(G) - \{x, y\} \notin \Gamma(G - f, k)$ .

*Proof.* Otherwise the result holds in G-f with  $V(G-f) = T \cup (W_1 - \{x, y\}) \cup \{x, y\}.$ 

(3.5) If  $a, b \in T$ , then e(a, b) = 0.

*Proof.* If  $f \in \partial(s_1, t_1)$ , then by Lemma 1(1) G has a path  $P[s_2, t_2]$  such that  $s_1 \notin V(P)$  and  $\lambda(G - E(P)) \ge k - 2$ .  $G^* := G - E(P) - f$  is a required graph. If  $f \in \partial(s_1, s_2)$ , then by Theorem 1 G - f has a path  $P[t_1, t_2]$  such that  $s_1 \in V(P)$  and  $\lambda(G - f - E(P)) \ge k - 3$ .

(3.6) If  $x \in W_1$ ,  $f \in \partial(x, s_1)$ , and  $h \in \partial(x, t_1)$ , then  $V(G) - \{x, s_1\} \notin \Gamma(G - f, k)$ .

*Proof.* Assume  $V(G) - \{x, s_1\} \in \Gamma(G - f, k)$ . Then  $V(G) - x \in \Gamma(G - \{f, h\}, k-1)$ . By Lemma 1(2)  $G - \{f, h\}$  has a path  $P[s_2, t_2]$  such that  $x \notin V(P)$  and  $\lambda(G - \{f, h\} - E(P)) \ge k - 3$ .

(3.7) If  $x \in W_1$  and  $f_i \in \partial(x, s_i)$  (i = 1, 2), then  $V(G) - \{x, s_1, s_2\} \notin \Gamma(G - \{f_1, f_2\}, k)$ .

*Proof.* Assume  $V(G) - \{x, s_1, s_2\} \in \Gamma(G - \{f_1, f_2\}, k)$ . Let  $y \in N(x) - \{s_1, s_2\}$ ,  $h \in \partial(x, y)$ , and  $G_1 := G - \{f_1, f_2, h\}$ . Then  $V(G) - x \in \Gamma(G_1, k-1)$  and  $e(x; G_1)$  is even. By Lemma 3 and Theorem 1  $G_1$  has a path  $P[t_1, t_2]$  such that  $s_1 \in V(P)$  and  $V(G_1) - x \in \Gamma(G_1 - E(P), k-3)$ . Let  $G^* := G - E(P) - \{f_1, f_2, h\}$ .

(3.8) If 
$$X \subset V(G)$$
,  $|X| \ge 2$ , and  $|\overline{X}| \ge 2$ , then  $e(X) \ge k+1$ .

*Proof.* Assume e(X) = k and for each  $Y \subsetneq X$  with  $|Y| \ge 2$ ,  $e(Y) \ge k + 1$ . Let  $u \in \overline{X}$  and  $G_1 := G/\overline{X}$ . If  $Z \subset V(G)$ ,  $|Z| \ge 2$ ,  $|\overline{Z}| \ge 2$ , and  $X \cap Z \ne \emptyset \ne X - Z$ , then  $e(Z) \ge k + 1$ . For if not, then  $Z - X \ne \emptyset \ne \overline{X \cup Z}$ , contrary to Lemma 7(2). Thus if  $x, y \in X$  and  $f \in \partial(x, y)$ , then  $V(G) - \{x, y\} \in \Gamma(G - f, k)$ , and so by (3.4)  $N(X \cap W_1: G_1) \subset T \cup \{\tilde{u}\}$ . By (3.5)  $N(X \cap T; G_1) \subset W_1 \cup \{\tilde{u}\}$ . Thus  $G_1$  is a bipartite graph with the partition  $(X \cap T, (X \cap W_1) \cup \{\tilde{u}\})$  or  $((X \cap T) \cup \{\tilde{u}\}, X \cap W_1)$  (note that  $G_1$  is k-regular and k is odd) and  $|V(G_1)| \ge 6$ . Thus for some  $x \in X \cap W_1$ ,  $e(x, \tilde{u}; G_1) = e(x, \overline{X}; G) < (k-1)/2$ . Then by Lemma 2 for some  $a, b \in X \cap T, G_x^{a,b}$  is admissible, contrary to (3.6) or (3.7).

By (3.3), (3.4), (3.5), and (3.8) G is a bipartite graph with the partition  $(T, W_1)$ . Let  $x \in W_1$ . By Lemma 2 for some  $a, b \in N(x) \subset T$ ,  $G_x^{a,b}$  is admissible, contrary to (3.6) or (3.7).

Case 2. k is even.

(3.9) If  $x \in W_1$ , then e(x) = k + 1.

*Proof.* Assume e(x) = k. By Lemma 2 there is an admissible lifting  $G_x$  of G at x. The result holds in  $G_x$  with  $V(G_x) = T \cup (W_1 - x) \cup \{x\}$ , and it also holds for G.

(3.10) If  $x, y \in W_1$  and  $f \in \partial(x, y)$ , then  $\lambda(G - f) \leq k - 1$ .

(3.11) If  $a, b \in T$ , then e(a, b) = 0.

*Proof.* We can prove  $e(s_i, t_i) = 0$   $(1 \le i \le n)$  in the same way as (3.5). If  $g \in \partial(s_1, s_2)$ , then let  $f \in \partial(s_1) - g$ . By Theorem 1 G has a path  $P[t_1, t_2]$  such that  $f \in E(P)$ ,  $g \notin E(P)$ , and  $\lambda(G - E(P)) \ge k - 2$ . Let  $G^* := G - E(P) - g$ .

(3.12) If  $x_1, x_2 \in W_1, a_1, a_2 \in T$ ,  $f_i \in \partial(x_i, a_1)$  (i = 1, 2), and  $g \in \partial(x_2, a_2)$ , then  $V(G) - a_1 \notin \Gamma(G - \{f_1, f_2\}, k)$ .

*Proof.* Assume  $V(G) - a_1 \in \Gamma(G - \{f_1, f_2\}, k)$ . If  $a_1 = s_1$  and  $a_2 = t_1$ , then by Lemma 1  $G - \{f_1, f_2\}$  has a path  $P[s_2, t_2]$  such that  $t_1 \notin V(P)$  and  $V(G) - s_1 \in \Gamma(G - \{f_1, f_2\} - E(P), k - 2)$ . Let  $G^* := G - E(P) - \{f_1, f_2, g\}$ . If  $a_1 = s_1$  and  $a_2 = s_2$ , then let  $h \in \partial(a_2) - g$ . By Theorem 1 and Lemma 3  $G - \{f_1, f_2\}$  has a path  $P[t_1, t_2]$  such that  $h \in E(P), g \notin E(P)$ , and  $V(G) - s_1 \in \Gamma(G - \{f_1, f_2\} - E(P), k - 2)$ . Let  $G^* := G - E(P) - \{f_1, f_2\}$  has a path  $P[t_1, t_2]$  such that  $h \in E(P), g \notin E(P), g \notin E(P), g \notin E(P) - \{f_1, f_2, g\}$ .

(3.13) If  $X \subset V(G)$ ,  $|X| \ge 2$ , and  $|\overline{X}| \ge 2$ , then  $e(X) \ge k + 1$ . *Proof.* Assume e(X) = k and for each  $Y \subsetneq X$  with  $|Y| \ge 2$ ,  $e(Y) \ge k + 1$ . Let  $u \in \overline{X}$ . If  $x, y \in X$  and  $f \in \partial(x, y)$ , then  $V(G) - \{x, y\} \in \Gamma(G - f, k)$ . For if not, then for some  $Z \subset V(G) - x$ ,  $y \in X$ , e(Z) = k,  $|Z| \ge 2$ , and  $|\overline{Z}| \ge 2$ , and so  $Z - X \ne \emptyset \ne \overline{X \cup Z}$  and by Lemma 7(2)  $e(X - Z) = e(X \cap Z) = k$ . Thus |X| = 2 and e(x) = e(y) = k. Then by (3.9)  $\{x, y\} \subset T$ , contrary to (3.11). Thus by (3.10)  $N(X \cap W_1; G/\overline{X}) \subset T \cup \{\widetilde{u}\}$ , and by (3.11)  $X \cap W_1 \ne \emptyset \ne X \cap T$ . By (3.9)  $|X \cap W_1| \ge 2$ , and so  $|X \cap T| \ge 2$ . Let  $a \in X \cap T$ . Since  $e(a, \overline{X}) < k/2$  (otherwise e(X - a) = k), by Lemma 2 for some  $x, y \in N(a) \cap X$ ,  $G_a^{x, y}$  is admissible. By (3.11)  $\{x, y\} \subset W_1$ . Let  $f_1 \in \partial(a, x)$  and  $f_2 \in \partial(a, y)$ , then  $V(G) - a \in \Gamma(G - \{f_1, f_2\}, k)$ . Let  $b \in ((N(x) \cup N(y)) \cap X) - a$ , then  $b \in T$ , contrary to (3.12).

By (3.3), (3.10), (3.11), and (3.13) G is a bipartite graph with the partition  $(T, W_1)$ . Let  $a \in T$ . By Lemma 2 for some x,  $y \in N(a) \subset W_1$ ,  $G_a^{x,y}$  is admissible and we can deduce a contradiction (see the proof of (3.13)).

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