

Least Fixed Point of a Functor

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INTRODUCTION

Fixed points of a functor $F: \mathcal{K} \rightarrow \mathcal{K}$ are objects X of \mathcal{K} such that $FX \cong X$. They arise, e.g., in Scott's approach to data types as lattices D satisfying $D \cong R \times [S \rightarrow D]$: These are fixed points of $F(?) = R \times [S \rightarrow ?]$. Scott remarks in [13] that a general theory of fixed points of functors would be of a value; the aim of the present paper is to lay foundations for such a theory.

We introduce the least fixed point (LFP) of a functor $F: \mathcal{K} \rightarrow \mathcal{K}$. We exhibit a construction of the LFP, generalizing the Knaster–Tarski formula $\text{lub}\{F^n(0)\}_{n \in \omega}$: lubs are substituted by well-ordered colimits and n is allowed to be an arbitrary ordinal. Related LFP constructions, always restricted to $n \in \omega$, have been considered by various authors [6, 11, 13, 14]. The advantage of the present approach is its effectiveness: Whenever a functor F has a fixed point then our construction stops, yielding the LFP of F .

Least fixed points are closely related to free algebras. In fact, the LFP construction, studied in the present paper, is a result of application of the free-algebra construction, exhibited by the first author [1], to initial objects. (This application is an idea of Arbib [5].) We proceed in a converse order here, introducing free algebras as special fixed points.

We further investigate categories with the fixed-point property (with respect to endofunctor). Example: Sets of power $\leq \alpha$ and vector spaces of dimension $\leq \alpha$ form categories with the fixed-point property (for any cardinal α). On the other hand, we show that a complete category with the fixed-point property must be a preordered class. The problem of characterizing categories with the fixed-point property is difficult: already for posets it is an open problem of long standing.

Finally, we show that a result, concerning the construction of free algebras for set functors, is undecidable, since it is equivalent to the nonexistence of measurable cardinals.

Results of the present paper have been announced in [4].

A. FIXED POINTS AND LFP

1. Throughout this paper we assume that a category \mathcal{K} and a functor $F: \mathcal{K} \rightarrow \mathcal{K}$ are given. An object X , isomorphic to its image FX , is called a *fixed point of F* (the first one to use this notion was probably Lambek [10]). We shall, sometimes, adopt a more rigorous attitude and we shall call a pair (X, v) a fixed point, if X is an object and $v: FX \rightarrow X$ is an isomorphism.

The *least fixed point (LFP)* of F is then defined as a fixed point (X_0, v_0) with the following universal property: For every fixed point (X, v) there exists a unique morphism $f: X_0 \rightarrow X$ such that $Ff = v^{-1} \cdot f \cdot v_0$.

$$\begin{array}{ccc}
 FX_0 & \xrightarrow{v_0} & X_0 \\
 Ff \downarrow & & \downarrow f \\
 FX & \xrightarrow[v^{-1}]{} & X
 \end{array}$$

This notion has been used by Arbib [6] and Wand [19].

Notice that, while a functor may have a number of distinct fixed points, the LFP is unique up to a (natural) isomorphism. For example, every object is a fixed point of the identity functor I of \mathcal{K} . But the LFP is precisely the *initial object* 0 characterized (up to isomorphism) by the property that each object X has precisely one morphism $f: 0 \rightarrow X$.

2. A category is said to be *chain-cocomplete* if each well-ordered diagram has a colimit. This is a property shared by a majority of “everyday” categories. For example the categories of ω -complete posets or of complete posets (with inf-preserving functions) are chain-cocomplete, in fact cocomplete. For chain-cocomplete categories we present a fixed-point construction. This construction is not applicable to categories of complete lattices (with sup- and inf-preserving functions) since they fail to be chain-cocomplete. For such categories the technique of D. Scott, based on the coincidence of certain limits and colimits, is quite different.

Let us remark that each chain-cocomplete category has an initial object 0 —the colimit of the void diagram.

3. *LFP construction.* For every endofunctor $F: \mathcal{K} \rightarrow \mathcal{K}$ of a chain-cocomplete category \mathcal{K} we define a transfinite chain of objects $W_i \in \mathcal{K}$ (for each ordinal i) and morphisms $W_{i,j}: W_i \rightarrow W_j$ ($i \leq j$) subject to

$$W_{i,i} = 1_{W_i} \text{ and } W_{i,k} = W_{i,k} \cdot W_{i,j} \quad \text{whenever } i \leq j \leq k.$$

$$0 \xrightarrow{W_{0,1}} F0 \xrightarrow{FW_{0,1}} F^20 \xrightarrow{F^2W_{0,1}} \dots \rightarrow W_\omega = \operatorname{colim}_{n < \omega} F^n0 \xrightarrow{W_{\omega,\omega+1}} FW_\omega \xrightarrow{FW_{\omega,\omega+1}} \dots$$

First,

$$W_0 = 0, \quad W_1 = F0, \quad W_{0,1}: 0 \rightarrow F0 \text{ canonical.}$$

Then for each isolated step we have $W_{i,i+1} : W_i \rightarrow W_{i+1} = FW_i$ and we put

$$W_{i+2} = FW_{i+1} \quad \text{and} \quad W_{i+1,i+2} = FW_{i,i+1} : FW_i \rightarrow FW_{i+1};$$

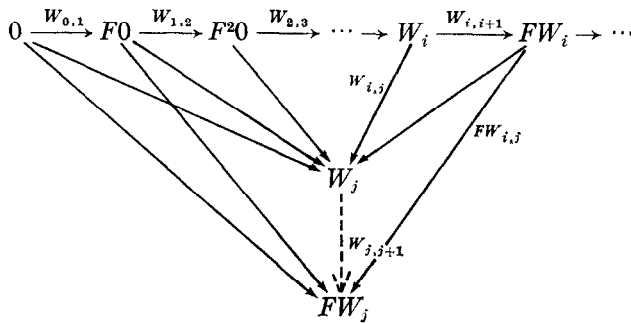
finally, given a limit ordinal j put

$$W_j = \operatorname{colim}_{i < j} W_i \text{ with } W_{i,j} : W_i \rightarrow W_j \quad (i < j) \text{ canonical}$$

and

$$W_{j+1} = FW_j;$$

we have a compatible family of morphisms $FW_{i,j} \cdot W_{i,i+1} : W_i \rightarrow FW_j$ ($i < j$)



which yields a unique morphism $W_{j,j+1} : W_j \rightarrow FW_j = W_{j+1}$ with

$$W_{j,j+1} \cdot W_{i,j} = FW_{i,j} \cdot W_{i,i+1} \quad ((i < j)).$$

The LFP construction is said to *stop after α steps* if $W_{\alpha,\alpha+1}$ is an isomorphism.

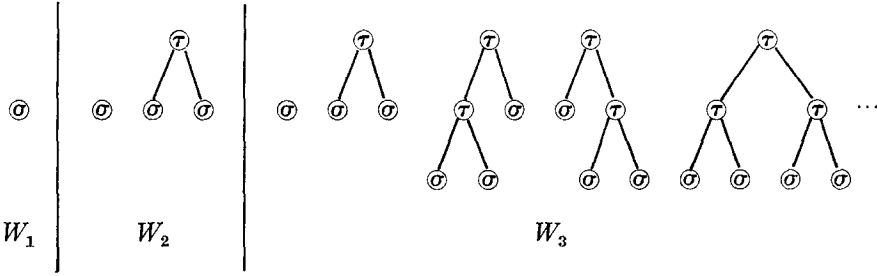
4. EXAMPLE. Let Σ be a ranked alphabet, i.e., a set, equipped by an arity function $\text{ar} : \Sigma \rightarrow \{0, 1, 2, \dots\}$. Define a set functor F_Σ , assigning to each set X the disjoint union of its Cartesian powers, each power related to a letter $\sigma \in \Sigma$ (the 0th power X^0 means a singleton set):

$$F_\Sigma X = \coprod_{\sigma \in \Sigma} X^{\text{ar}(\sigma)}.$$

It is clear how to define (naturally) F_Σ on morphisms to obtain a functor. Denoting $\Sigma_0 = \{\sigma \in \Sigma; \text{ar}(\sigma) = 0\}$, we have

$$W_0 = \emptyset; \quad W_1 = \coprod_{\sigma \in \Sigma} \emptyset^{\text{ar}(\sigma)} = \Sigma_0; \quad W_2 = \coprod_{\sigma \in \Sigma} \Sigma_0^{\text{ar}(\sigma)}; \dots$$

The n th step W_n can be clearly interpreted as the set of all Σ trees (i.e., trees labeled by Σ so that arities agree with the number of successors) of length smaller than n . For example, for $\Sigma = \{\sigma, \tau\}$ with $\text{ar}(\sigma) = 0, \text{ar}(\tau) = 2$:



Then $W_\omega \cong \bigcup_{n=0}^\omega W_n$ is the set of all Σ trees and $W_{\omega, \omega+1}$ is a bijection, hence the LFP construction stops after ω steps.

Allowing infinite arities, the LFP construction for F_Σ always stops but later than after ω steps; W_i is the set of Σ trees of length smaller than i .

5. PROPOSITION *If the LFP construction stops after α steps then W_α is the LFP with respect to $W_{\alpha, \alpha+1}^{-1} : FW_\alpha \rightarrow W_\alpha$.*

Proof. Given a fixed point (X, v) , define $f_i : W_i \rightarrow X$ by induction to yield a compatible family ($f_i = f_j \cdot W_{i,j}$ for each $i \leq j$).

$f_0 : 0 \rightarrow X$ is canonical.

$f_{i+1} = v \cdot Ff_i : FW_i = W_{i+1} \rightarrow X$.

f_γ (γ a limit ordinal) is determined by the fact that f_i ($i \leq \gamma$) form a compatible family.

Particularly, for $f = f_\alpha$ we have $v \cdot Ff = f_{\alpha+1}$ and so

$$v \cdot Ff = fW_{\alpha, \alpha+1}^{-1},$$

hence $Ff = v^{-1} \cdot f \cdot W_{\alpha, \alpha+1}^{-1}$.

The uniqueness of f is easy: Given $f : W_\alpha \rightarrow X$ with $Ff = v^{-1} \cdot f \cdot W_{\alpha, \alpha+1}^{-1}$, we have $f \cdot W_{\alpha, 0} = f_0$ (because 0 is an initial object) and it suffices to show that $f \cdot W_{i, \alpha} = f_i$ implies $f \cdot W_{i+1, \alpha} = f_{i+1}$ (then, also $f \cdot W_{\gamma, \alpha} = f_\gamma$ for a limit ordinal γ and, by induction, $f = f \cdot W_{\alpha, \alpha} = f_\alpha$):

$$\begin{aligned} f \cdot W_{i+1, \alpha} &= f \cdot W_{\alpha, \alpha+1}^{-1} \circ W_{i+1, \alpha+1} = f \cdot W_{\alpha, \alpha+1}^{-1} \cdot FW_{i, \alpha} \\ &= v \cdot Ff \cdot FW_{i, \alpha} = v \cdot Ff_i \\ &= f_{i+1}. \end{aligned}$$

Thus, if the LFP construction stops then F has LFP. But not conversely.

6. EXAMPLE. Let \mathcal{X} be the category of sets and let $F = P^*$ be the “infinite-power-set” functor defined as follows:

$P^* \emptyset = N$, the set of natural numbers;

$P^*X = \{T \subset X; T \text{ infinite}\} \cup \{*\}$ for $X \neq \emptyset$.

Given a mapping $f: X \rightarrow Y$, we have three possibilities:

- $X \neq \emptyset$, then $P^*f(T) = f(T)$ if $f(T)$ is infinite, $P^*f(T) = * = P^*f(*)$ if $f(T)$ is finite;
- $X = Y = \emptyset$, then $P^*f = \text{id}_N$;
- $X = \emptyset \neq Y$, then $P^*f(n) = *$ for $n \in N$.

This functor has a unique fixed point: singleton set. Hence, this is the LFP of P^* . Yet, the LFP construction does not stop: We have $W_1 = P^*\emptyset = N$ and, whenever W_i is an infinite set then P^*W_i has a greater power than W_i —thus, W_i cannot be a fixed point of P .

The reason why LFP construction fails for the functor P^* is that P^* does not preserve monos (the void map $\emptyset \rightarrow \{*\}$ is lifted to the constant map $N \rightarrow \{*\}$). Before discussing how preservation of monos influences the situation we shall introduce a broader view of F -algebras.

B. FIXED POINTS AND FREE ALGEBRAS

7. Given a functor $F: \mathcal{K} \rightarrow \mathcal{K}$, an F -algebra is a pair (Q, d) consisting of an object Q of \mathcal{K} and a morphism $d: FQ \rightarrow Q$. A homomorphism between F -algebras $f: (Q, d) \rightarrow (Q', d')$ is a morphism $f: Q \rightarrow Q'$ in \mathcal{K} such that the square

$$\begin{array}{ccc} FQ & \xrightarrow{d} & Q \\ Ff \downarrow & & \downarrow f \\ FQ' & \xrightarrow{d'} & Q' \end{array}$$

commutes. For example, let $F = F_\Sigma$ be as in Example 4, then an F_Σ -algebra is a set Q together with a map $d: \coprod_{\sigma \in \Sigma} Q^{\text{ar}(\sigma)} \rightarrow Q$, which can be viewed as a collection of operations

$$Q^{\text{ar}(\sigma)} \rightarrow Q \quad (\sigma \in \Sigma).$$

In that sense, F_Σ -algebras are just universal algebras of type Σ . The above notion of homomorphism coincides with that from universal algebra.

8. A free F -algebra, generatee by an object I , is an F -algebra $(I^\#, \varphi)$ together with a morphism $s: I \rightarrow I^\#$ (insertion of generators) which is universal in the following sense: Every diagram

$$\begin{array}{ccc} FI^\# & \xrightarrow{\sigma} & I^\# \xleftarrow{s} I \\ & & \searrow \varphi \\ FQ & \xrightarrow{d} & Q \end{array}$$

has a unique commutative fill-in $f^\#$:

$$\begin{array}{ccc}
 FI^\# & \xrightarrow{\varphi} & I^\# \xleftarrow{s} I \\
 Ff^\# \downarrow & & \downarrow f^\# \swarrow f \\
 FQ & \xrightarrow{a} & Q
 \end{array}$$

Free F -algebras play an important role in the theory of machines in a category due to Arbib and Manes [7]. We call F a *variator* (input process in [7]) provided that for every object I there exists the free algebra $(I^\#, \varphi)$.

9. The free algebra over the initial object 0 (provided that both exist) is the algebra having exactly one homomorphism into any other F -algebra. Thus, this is the *initial algebra*, i.e., the initial object in the category of F -algebras. If $(0^\#, \varphi)$ is the initial algebra then $0^\#$ is a fixed point, indeed the LFP, see [6]; thus, it is natural to ask:

- (a) Does the existence of the LFP guarantee the existence of an initial algebra?
- (b) Does the existence of a fixed point guarantee the existence of the LFP (and the stop of the LFP construction)?

In general, the answer is negative:

EXAMPLE. LFP exists but initial algebra does not. Let \mathcal{K} be the category of partial semigroups. Denote by $(I, *)$ the (total) one-element semigroup and define $F: \mathcal{K} \rightarrow \mathcal{K}$ on objects by

$$\begin{aligned}
 F(X, \cdot) &= (I, *) && \text{if the operation } \cdot \text{ is anywhere defined (i.e. } \cdot \neq \emptyset); \\
 F(X, \emptyset) &= (\exp X, \emptyset).
 \end{aligned}$$

There is a unique way of defining F on morphisms provided that

$$Ff = \exp f \quad \text{for each } f: (X, \emptyset) \rightarrow (Y, \emptyset).$$

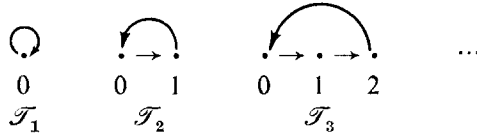
Then F has a unique (up to isomorphism) fixed point, which is its LFP, viz., $(I, *)$. Yet F has no initial algebra. Indeed, if an initial algebra existed it would have to coincide with the LFP. But, given any set X and any map $d: \exp X \rightarrow X$ we have an F algebra

$$((X, \emptyset), d) \quad [\text{since } d: F(X, \emptyset) \rightarrow (X, \emptyset)]$$

and there exists no homomorphism from $(I, *)$ to (X, \emptyset) , of course.

EXAMPLE. Fixed points exist but LFP does not. Let \mathcal{K} be the category of graphs (i.e., pairs $\mathcal{G} = (V, E)$, where V is a set and $E \subset V \times V$) and graph morphisms (i.e., maps $f: (V, E) \rightarrow (V', E')$ such that $(x, y) \in E$ implies $(f(x), f(y)) \in E'$). Denote by \mathcal{K}_n

the class of all graphs which contain a cycle of length n , i.e., which have a subgraph, isomorphic to \mathcal{T}_n :

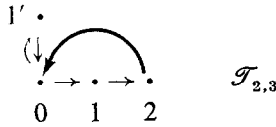


It is easy to verify that given a graph morphism $f: \mathcal{G} \rightarrow \mathcal{G}'$ then $\mathcal{G} \in \mathcal{K}_1$ implies $\mathcal{G}' \in \mathcal{K}_2$, $\mathcal{G} \in \mathcal{K}_2$ implies $\mathcal{G}' \in \mathcal{K}_1 \cup \mathcal{K}_2$, and $\mathcal{G} \in \mathcal{K}_3$ implies $\mathcal{G}' \in \mathcal{K}_1 \cup \mathcal{K}_3$.

Define $F: \mathcal{K} \rightarrow \mathcal{K}$ on objects as follows:

$$\begin{aligned}
 F\mathcal{G} &= \mathcal{T}_1 && \text{if } \mathcal{G} \in \mathcal{K}_1, \\
 F\mathcal{G} &= \mathcal{T}_2 && \text{if } \mathcal{G} \in \mathcal{K}_2 - (\mathcal{K}_1 \cup \mathcal{K}_3), \\
 F\mathcal{G} &= \mathcal{T}_3 && \text{if } \mathcal{G} \in \mathcal{K}_3 - (\mathcal{K}_1 \cup \mathcal{K}_2); \\
 F\mathcal{G} &= \mathcal{T}_{2,3} && \text{if } \mathcal{G} \in \mathcal{K}_2 \cap \mathcal{K}_3,
 \end{aligned}$$

where $\mathcal{T}_{2,3}$ is the graph



$$F\mathcal{G} = (\exp V, \emptyset) \quad \text{if } \mathcal{G} = (V, E) \in \mathcal{K} - (\mathcal{K}_1 \cup \mathcal{K}_2 \cup \mathcal{K}_3).$$

Given a graph morphism $f: \mathcal{G} \rightarrow \mathcal{G}'$, define Ff as follows:

- (a) $\mathcal{G} \in \mathcal{K} - (\mathcal{K}_1 \cup \mathcal{K}_2 \cup \mathcal{K}_3)$,
 $Ff = \exp f$ if $\mathcal{G}' \in \mathcal{K} - (\mathcal{K}_1 \cup \mathcal{K}_2 \cup \mathcal{K}_3)$,
 $Ff = \text{const } 0$ if $\mathcal{G}' \in \mathcal{K}_1 \cup \mathcal{K}_2 \cup \mathcal{K}_3$,
- (b) $\mathcal{G} \in \mathcal{K}_1$ [implies $\mathcal{G}' \in \mathcal{K}_1$],
 $Ff = \text{id}_{\{0\}}$,
- (c) $\mathcal{G} \in \mathcal{K}_2$ [implies $\mathcal{G}' \in \mathcal{K}_1 \cup \mathcal{K}_2$],
 $Ff = \text{const } 0$ if $\mathcal{G}' \in \mathcal{K}_1$,
 $Ff = \text{id}_{\{0,1\}}$ if $\mathcal{G}' \in \mathcal{K}_2 - (\mathcal{K}_1 \cup \mathcal{K}_3)$,
 $Ff = \text{embedding } \mathcal{T}_2 \rightarrow \mathcal{T}_{2,3}$ if $\mathcal{G}' \in \mathcal{K}_3$,
- (d) analogously for $\mathcal{G} \in \mathcal{K}_3$.

A mechanical checking shows that F is a correctly defined functor. It has, up to isomorphism, exactly four fixed points, viz.,

$$\mathcal{F}_1, \quad \mathcal{F}_2, \quad \mathcal{F}_3, \quad \text{and} \quad \mathcal{F}_{2,3}.$$

Since none of these admits a homomorphism into each of the other three, F has no LFP.

10. DEFINITION. A class C of morphisms is said to be *chain-cocomplete* if for every ordinal γ and every chain of C -morphisms

$$W_{i,j} : W_i \rightarrow W_j \text{ in } C \quad (i \leq j < \gamma)$$

there is a colimit, say W_γ and $W_{i,\gamma} : W_i \rightarrow W_\gamma$, such that

- (i) $W_{i,\gamma} \in C$ for each $i < \gamma$;
- (ii) for every compatible family $U_i : W_i \rightarrow U$ in C the factorizing morphism $h : W_\gamma \rightarrow U$ [$h \cdot W_{i,\gamma} = U_i$] is in C , too.

EXAMPLES. All monos form a chain-cocomplete class in a number of categories, including sets, complete posets, topological spaces, varieties of finitary algebras, etc. They fail to form a chain-cocomplete class, e.g., in the category of compact Hausdorff spaces. See [18] for details.

Note. If C is a chain-cocomplete class then

- (0) \mathcal{K} has an initial object 0 with each $0 \rightarrow X$ in C .

Proof. Apply (ii) in the definition to the void chain.

- (1) C is closed to isomorphisms.

Proof. Apply (i) to the chain $1_A : A \rightarrow A$ of length 1: Given an isomorphism $v : A \rightarrow B$ then B together with v is the colimit of this chain.

- (2) C is closed to composition.

Proof. Apply (i) to the chain $f : A \rightarrow B, g : B \rightarrow C$ of length 2. Its colimit is C with colimit injections $g \cdot f : A \rightarrow C$ (1), $g : B \rightarrow C$, and $1 : C \rightarrow C$.

11. THEOREM. Let \mathcal{M} be a chain-cocomplete class of monos in a category \mathcal{K} and let \mathcal{K} be \mathcal{M} -well powered. The following conditions are equivalent for any functor $F : \mathcal{K} \rightarrow \mathcal{K}$, preserving \mathcal{M} (i.e., such that $m \in \mathcal{M}$ implies $Fm \in \mathcal{M}$):

- (i) F has a fixed point;
- (ii) F has the LFP;
- (iii) the initial F -algebra exists;
- (iv) the LFP construction stops.

Proof. Only (i) \Rightarrow (iii) & (iv) has to be proved, of course. Thus, let (X, ν) be a fixed point of F . To see that the LFP construction stops, we shall construct a compatible

family $h_i : W_i \rightarrow X$ of \mathcal{M} -monos. Since \mathcal{K} is \mathcal{M} -well powered, this will guarantee that two of these subobjects of X , say $h_\alpha, h_{\alpha'}$ ($\alpha < \alpha'$) are isomorphic, from which it follows that $W_{\alpha, \alpha'}$ (and $W_{\alpha, \alpha+1}$) is an isomorphism—hence, the construction stops and yields the LFP. It is proved in [1] that if this is the case, then the LFP is the initial F -algebra.

We shall construct h_i by induction in i , proving also, as we proceed, that $W_{t, i+1}$ is also in \mathcal{M} for each $t \leq i + 1$. First, both $h_0 : 0 \rightarrow X$ and $W_{0,1} : 0 \rightarrow F0$ are in \mathcal{M} . Second, given $h_i \in \mathcal{M}$ we define $h_{i+1} = v \cdot Fh_i$.

$$\begin{array}{ccc}
 W_i & \xrightarrow{W_{i, i+1}} & W_{i+1} = FW_i \xrightarrow{FW_{i, i+1}} FW_{i+1} \\
 h_i \downarrow & & \downarrow Fh_i \\
 X & \xleftarrow{v} & FX
 \end{array}$$

Since $h_i \in \mathcal{M}$ implies $Fh_i \in \mathcal{M}$ we see that $h_{i+1} \in \mathcal{M}$ (for, by the above note, \mathcal{M} is closed to composition with isomorphisms). Further,

$$W_{t, i+1} = (FW_{t, i}) \cdot W_{t, i+1} \in \mathcal{M}.$$

Finally, given a limit ordinal γ and given $h_i : W_i \rightarrow X$ in \mathcal{M} for each $i < \gamma$, we define $h_\gamma : W_\gamma \rightarrow X$ by $h_\gamma \cdot W_{i, \gamma} = h_i$ for each $i < \gamma$. Since \mathcal{M} is a chain-cocomplete class, we have both $h_\gamma \in \mathcal{M}$ and $W_{\gamma, \gamma+1} \in \mathcal{M}$.

12. In a category \mathcal{K} with finite coproducts the above theorem yields a criterion on a functor $F : \mathcal{K} \rightarrow \mathcal{K}$ to be a variator. For each object I in \mathcal{K} denote by $F_{(I)} : \mathcal{K} \rightarrow \mathcal{K}$ the functor, defined by

$$F_{(I)}X = FX + I \quad \text{and} \quad F_{(I)}f = Ff + \text{id}_I$$

(i.e., $F_{(I)}$ is a coproduct of F and the constant functor to I).

Observation. The free F -algebra, generated by an object I , is precisely the initial $F_{(I)}$ -algebra.

Now, if F preserves a class \mathcal{M} which is additive (i.e., given $m : A \rightarrow B$ and $m' : A' \rightarrow B'$ in \mathcal{M} then also $m + m' : (A + A') \rightarrow (B + B')$ is in \mathcal{M}) then each of the functors $F_{(I)}$ preserves \mathcal{M} . Hence, we obtain the following result, originally proved in [18]:

COROLLARY. *Let \mathcal{K} be a category with finite coproducts and let \mathcal{M} be its additive, chain-cocomplete class of monos such that \mathcal{K} is \mathcal{M} -well powered. A functor $F : \mathcal{K} \rightarrow \mathcal{K}$, preserving \mathcal{M} , is a variator iff for every object I there exists an object X , isomorphic to $FX + I$ (i.e., a fixed point of $F_{(I)}$).*

13. An important feature of Dana Scott’s fixed points is that every object has a “nice” embedding into a fixed point. Let us make a remark concerning this situation.

DEFINITION. We say that an object X is a *canonical subobject* of an object T if T is isomorphic to $X + Y$ for some object Y .

THEOREM. *Let \mathcal{K}, \mathcal{M} be as in the above corollary and let $F: \mathcal{K} \rightarrow \mathcal{K}$ be a functor, preserving \mathcal{M} and such that each object of \mathcal{K} is a canonical subobject of a fixed point of F . If \mathcal{K} has countable coproducts then F is a variator.*

Proof. We shall use the above corollary: Given an object I , we form the coproduct I_ω of countably many copies of I ; by hypothesis, there exists a fixed point X , isomorphic to $I_\omega + Y$ (for some object Y). Then

$$X \cong FX$$

and also, via $I_\omega + I \cong I_\omega$:

$$X + I \cong Y + I_\omega + I \cong Y + I_\omega \cong Y.$$

Hence, $FX \cong FX + I$ and we have

$$X \cong FX + I.$$

The converse fails even for very reasonable variators. For example, let \mathcal{K} be the category of partial commutative groupoids and let FX be the subgroupoid of X over all idempotents, $FX = \{y \in X; y \cdot y = y\}$. Then F preserves coproducts (hence it is a variator with $I^\# = \coprod_{n=0}^\infty F^n I$) and preserves limits. Yet, no object other than a fixed point of F has a canonical embedding into a fixed point!

C. FIXED-POINT PROPERTY OF CATEGORIES

14. EXAMPLE. Let $\text{Set}(\alpha)$ denote the category of sets of power $\leq \alpha$ (and mappings between them). For every α , $\text{Set}(\alpha)$ has the *fixed-point property*: Every endofunctor has a fixed point. Indeed, for any endofunctor F put $C_F = \{\beta; \text{if } \text{card } X = \beta \text{ then } \text{card } FX \leq \beta\}$. Then $\alpha \in C_F$ hence $C_F \neq \emptyset$, and we can choose X_0 with $\text{card } X_0 = \min C_F$. Then $\text{card } FX_0 \leq \text{card } X_0 (\in C_F)$ and so either $FX_0 = \emptyset$ (in which case necessarily $X_0 = \emptyset$ is the LFP) or there is a retraction $r: X_0 \rightarrow FX_0$. Since Fr is also a retraction, there follows $\text{card } FFX_0 \leq \text{card } FX_0$, i.e., $\text{card } FX_0 \in C_F$. Then $\text{card } FX_0 = \text{card } X_0$, so that X_0 is a fixed point of F .

EXAMPLE. The category $T\text{-Vect}(\alpha)$ of all vector spaces over a field T of dimension $\leq \alpha$ (and all linear mappings between them) has also the fixed-point property. The proof is the same.

Note. In $T\text{-Vect}(\alpha)$ the monomorphisms all split, hence they are preserved by every endofunctor. The class of all monos has all properties, required in Section 12 above. Hence, *every functor is a variator*. In fact, with the exception of functors, equivalent to constant functors, every functor F has the property, mentioned in Section 13: Each space is a subspace of a fixed point.

The same is true for $\text{Set}(\alpha)$ but the argument must be more careful, see [12], because the monos from the void set do not split and are, generally, mapped by functors to nonmonos.

Our study of the fixed-point property of categories has been inspired by the result of Trnková [17] that every endofunctor of $\text{Set}(\aleph_0)$ is a variator.

15. We have seen in Section 14 a type of categories with the fixed-point property. Another type is ω -complete posets. We shall show now that the latter are the only "reasonable" categories with the fixed-point property. We denote by Set^+ the category of nonvoid sets and mappings.

PROPOSITION. *For every nonconstant set functor F there exists a cardinal α such that $\text{card } FX \geq \text{card } X$ for every set of power greater than α .*

In particular, the image $\{FX; X \text{ a set}\}$ of a nonconstant set functor F contains a proper class of pairwise nonisomorphic sets.

Proof. See [8].

LEMMA. *Let \mathcal{K} be a category, for which a functor $\Phi: \text{Set}^+ \rightarrow \mathcal{K}$ exists, which is not naturally equivalent to a constant functor. Then \mathcal{K} has an endofunctor without fixed points.*

Proof. (1) There exists a map $f: X \rightarrow X$ with $\Phi f \neq 1_{\Phi X}$. *Proof:* Since Φ is not equivalent to a constant functor, there are two possibilities:

(a) Φh is not an isomorphism, for some $h: A \rightarrow B$ in Set^+ . Since $A \neq \emptyset$, we can write $h = m \cdot e$, where $e: A \rightarrow C$ is a split epi and $m: C \rightarrow B$ is a split mono. If Φe is not an isomorphism, then we choose $r: C \rightarrow A$ with $e \cdot r = 1$ and we find out that $\Phi(r \cdot e) \neq 1$ (else $\Phi e = (\Phi r)^{-1}$); then we put $f = r \cdot e$. If Φe is an isomorphism, then Φm cannot be an isomorphism. We choose $p: B \rightarrow C$ with $p \cdot m = 1$ and we find out that $\Phi(m \cdot p) \neq 1$; then we put $f = m \cdot p$.

(b) $\Phi h \neq \Phi k$ for some $h, k: A \rightarrow B$ in Set^+ . Then for the coequalizer e of h, k we have: Φe is not an isomorphism (else, $\Phi e \cdot \Phi h = \Phi e \cdot \Phi k$ implies $\Phi h = \Phi k$). Then we return to a).

(2) The image of Φ contains a proper class of pairwise nonisomorphic objects of \mathcal{K} . *Proof:* Consider the set functor $F = \text{hom}(\Phi X, \Phi -)$, where $\Phi f \neq 1$ for some $f: X \rightarrow X$. Then F is nonconstant, because $Ff(1_{\Phi X}) = \Phi f$ while $F1_{\Phi X}(1_{\Phi X}) = 1_{\Phi X}$. By the preceding proposition we can choose a class \mathcal{D} of sets such that the sets $\{FX; X \in \mathcal{D}\}$ are pairwise nonisomorphic. Hence, the objects $\{\Phi X; X \in \mathcal{D}\}$ are pairwise nonisomorphic in \mathcal{K} .

(3) For every cardinal t there exists a cardinal t^* such that

$$|\text{hom}(\Phi X, \Phi M)| = t \text{ implies } |M| \leq t^*$$

for every set M . *Proof:* By the preceding proposition it is clear that for every nonconstant set functor F and every cardinal t , $\text{card } FX > t$ whenever X is sufficiently great (card

$X > \alpha + t$) and we can choose t^* such that $\text{card } X > t^*$ implies $\text{card } FX \neq t$. Apply this rule to F above: For every t there is t^* such that

$$|FM| = |\text{hom}(\Phi X, \Phi M)| = t \text{ implies } |M| \leq t^*.$$

(4) There exists a set functor T such that the functor $\Phi \cdot T \cdot \text{hom}(\Phi X, -)$ has no fixed point. For each cardinal t we can choose an infinite cardinal $\lambda(t)$, greater than t^* above, in such a way that

$$t_1 < t_2 \quad \text{implies} \quad \lambda(t_2) > 2^{\lambda(t_1)} + 2^{t_2}. \tag{*}$$

By (*) it follows from [8] that a set functor T exists for which $\text{card } Z = t$ implies $\text{card } TZ = \lambda(t)$. Now, given an object Y in \mathcal{X} , put $M = \text{hom}(\Phi X, Y)$, $t = |M|$. Then $|TM| = \lambda(t)$ implies

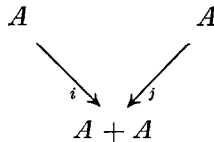
$$|\text{hom}(\Phi X, \Phi TM)| \neq t$$

via (3) above. This proves that Y is not a fixed point of $\Phi \cdot T \cdot \text{hom}(\Phi X, -)$: else, $t = |\text{hom}(\Phi X, Y)| = |\text{hom}(\Phi X, \Phi \cdot T \cdot \text{hom}(\Phi X, Y))| = |\text{hom}(\Phi X, \Phi M)|$. This concludes the proof.

THEOREM. *If a category \mathcal{X} has the fixed-point property (i.e., if every functor $F: \mathcal{X} \rightarrow \mathcal{X}$ has a fixed point) and has powers or copowers of all objects, then \mathcal{X} is a preordered class.*

Proof. Assuming that \mathcal{X} is not a preordered class, it suffices to find a functor $\Phi: \text{Set}^+ \rightarrow \mathcal{X}$, nonequivalent to a constant functor. We shall assume that \mathcal{X} has copowers; if it has powers we can work with its dual.

Thus, let $a, b: A \rightarrow B$ be distinct morphisms. We define $\Phi: \text{Set}^+ \rightarrow \mathcal{X}$ as the copowers of A : $\Phi M = \coprod_{m \in M} A_m$ with $A_m = A$; for $f: M \rightarrow N$ there is an induced morphism $\Phi f: \coprod_{m \in M} A_m \rightarrow \coprod_{n \in N} A_n$. To prove that Φ is indeed nonequivalent to a constant functor, it suffices to verify that in the copower $A + A$



i is not an isomorphism. Assume that, to the contrary, i is an isomorphism. Since we have a pair $1_A, 1_A: A \rightarrow A$, factorizing as $1_A = k \cdot i, 1_A = k \cdot j$ for a unique $k: A + A \rightarrow A$, there follows that $k = i^{-1}$ and $j = k^{-1} = i$. This is a contradiction: For $a, b: A \rightarrow B$ there exists $c: A + A \rightarrow B$ with $a = c \cdot i = c \cdot j = b$, though $a \neq b$.

Note. Although only finite copowers were explicitly used in the argument above, the existence of finite copowers is not enough for the theorem to hold. Notice that $\text{Set}(2^\alpha)$ has all limits and colimits of diagrams with less than α objects; yet, $\text{Set}(2^\alpha)$ has the fixed-point property.

D. FREE ALGEBRAS AND MEASURABLE CARDINALS

16. The most important part of the LFP construction is the finitary part (W_i for $i \leq \omega$). In [1], a variator F is called *algorithmic* if free algebras are obtained by the finitary part of the free-algebra construction, in other words, if for each $F_{(j)}$ the LFP construction stops after ω steps.

Every *finitary functor*, i.e., functor preserving filtered colimits, is an algorithmic variator. In the case where \mathcal{X} is the category of sets the converse seemed to be indicated by concrete examples: Every algorithmic variator is finitary. Is it true?

For the case of sets the following conditions on $F: \text{Set} \rightarrow \text{Set}$ are equivalent (see [2]),

- (i) F is finitary;
- (ii) F preserves well-ordered unions (i.e., given a chain of sets W_i ($i < \alpha$), then $F(\bigcup_{i < \alpha} W_i) = \bigcup_{i < \alpha} Ft_i(FW_i)$, where $t_{i_0}: W_{i_0} \rightarrow \bigcup_{i < \alpha} W_i$ denotes the inclusion map);
- (iii) for every set X and every point $x \in FX$ there exists a finite set $Y \subset X$ and a point $y \in FY$ with $x = Ft(y)$ ($t: Y \rightarrow X$ is the inclusion).

17. Recall that a cardinal α is measurable iff there exists a σ -additive $\{0, 1\}$ -measure on a set X of power α ; equivalently, if there exists an ultrafilter \mathfrak{F} on X , which is free (i.e., not containing any singleton set) and is closed to countable intersections (i.e., $\bigcap_{n=1}^\infty A_n \in \mathfrak{F}$ whenever $A_n \in \mathfrak{F}$ for each n). The following lemma is folklore and we prove it here only because we do not know any reference.

LEMMA. *Let there exist no measurable cardinal. Let \mathfrak{F} be a filter on a set X , not containing any finite subset of X . Then there exist sets $X_0 \subset X_1 \subset X_2 \subset \dots$ such that $X = \bigcup_{n=0}^\infty X_n$ and $X_n \notin \mathfrak{F}$ for each n .*

Proof. (A) Let \mathfrak{F} be a free filter. Then \mathfrak{F} is contained in a free ultrafilter \mathfrak{G} . By hypothesis, \mathfrak{G} is not closed under countable intersections, hence there exist sets $A_n \in \mathfrak{G}$ with $\bigcap_{n=1}^\infty A_n \notin \mathfrak{G}$. Put $X_0 = \bigcap_{n=1}^\infty A_n$ and $X_k = X_0 \cup (X - \bigcap_{n=1}^k A_n)$, $k = 1, 2, 3, \dots$. Then $X_0 \notin \mathfrak{G}$ and $X - \bigcap_{n=1}^k A_n \notin \mathfrak{G}$ (because $\bigcap_{n=1}^k A_n \in \mathfrak{G}$); since \mathfrak{G} is an ultrafilter, necessarily $X_k \notin \mathfrak{G}$. Hence, for each k , $X_k \notin \mathfrak{F}$. And

$$\bigcup_{k=0}^\infty X_k = \left(\bigcap_{n=1}^\infty A_n \right) \cup \bigcup_{k=1}^\infty \left(X - \bigcup_{n=1}^k A_n \right) = X.$$

(B) \mathfrak{F} is arbitrary. Let F be the intersection of all elements of \mathfrak{F} . If $F \in \mathfrak{F}$, then F cannot be finite and we can choose a sequence f_0, f_1, f_2, \dots of pairwise distinct elements of F . Put $X_n = X - \{f_n, f_{n+1}, f_{n+2}, \dots\}$: since $F = \bigcap \mathfrak{F}$ but X_n does not contain F , clearly $X_n \notin \mathfrak{F}$. Yet, $X = \bigcup_{n=0}^\infty X_n$ and $X_0 \subset X_1 \subset X_2 \subset \dots$.

If $F \notin \mathfrak{F}$, put $X' = X - F$ ($\neq \emptyset$) and let $\mathfrak{F}' = \{T \subset X'; T = S \cap X' \text{ for some } S \in \mathfrak{F}\}$. Then \mathfrak{F}' is a free filter on X' . By part (A) there exist sets $X'_0 \subset X'_1 \subset X'_2 \subset \dots$ with $X' = \bigcup_{n=0}^\infty X'_n$ and $X'_n \notin \mathfrak{F}'$. Then for $X_n = X'_n \cup F$ we have $X_0 \subset X_1 \subset X_2 \subset \dots$, $X = \bigcup_{n=0}^\infty X_n$, and $X_n \notin \mathfrak{F}$ (the last follows from $X_n \cap X' = X'_n \notin \mathfrak{F}'$).

PROPOSITION. (Set functors preserve finite nonvoid intersections.) Let F be an arbitrary set functor and let X be a set, x an element of FX . Then the collection

$$\mathfrak{F}_x = \{Y; \emptyset \neq Y \subset X; x \in Ft(FY), t: Y \rightarrow X \text{ the inclusion map}\}$$

is either a filter on X or $\mathfrak{F}_x = \exp X - \{\emptyset\}$.

Proof. See [15].

18. THEOREM. The nonexistence of measurable cardinals is equivalent to the validity of the following statement:

(*) A functor $F: \text{Set} \rightarrow \text{Set}$ is an algorithmic variator iff it is finitary.

Proof. Let there exist measurable cardinals. Then the following functor $B: \text{Set} \rightarrow \text{Set}$ is clearly nonfinitary:

BX is the set of all ultrafilters on X , closed to countable intersections (free or fixed);

given $f: X \rightarrow Y$ and an ultrafilter $\mathfrak{F} \in BX$ then

$$Bf(\mathfrak{F}) = \{R \subset Y; f^{-1}(R) \in \mathfrak{F}\}.$$

It is proved in [16] that B preserves countable colimits; then B is an algorithmic variator Hence (*) fails.

Conversely, assuming that measurable cardinals do not exist, we shall prove (*), i.e., we shall verify that any functor F , which is not finitary, fails to be an algorithmic variator. By (iii) in 16 there exists a point $x \in FX$ such that the collection

$$\mathfrak{F}_x = \{Y \subset X; x \in Ft(FY), t: Y \rightarrow X \text{ inclusion}\}$$

contains no finite set. By the proposition in Section 17, \mathfrak{F}_x is a filter and, since measurable cardinals do not exist, we have an increasing sequence of sets $X_n \subset X$, $X_n \notin \mathfrak{F}_x$ with $X = \bigcup_{n=0}^{\infty} X_n$.

Let α be a cardinal with $\text{card } FT \geq \text{card } T$ whenever T has a power greater than α (see the proposition in Section 15; F is clearly nonconstant). Choose a set I whose power is greater than both α and all $\text{card } X_n$ (e.g., $I = X \times \alpha$). We shall verify that the FLP construction does not stop after ω_0 steps for $F_{(I)}$. In other words, we put

$$\begin{aligned} W_0 &= F_{(I)}\emptyset = I, \\ W_{n+1} &= F_{(I)}W_n = I + FW_n, \\ c_0 &: I \rightarrow I + FW_n \text{ canonical,} \\ c_{n+1} &= 1 + Fc_n : I + FW_n \rightarrow I + FW_{n+1}, \end{aligned}$$

and we shall verify that the construction does not stop after ω_0 steps. To do so, we shall exhibit a point in FW_{ω_0} which is not in the image of Fw_n for any n , where $w_n : W_n \rightarrow W_{\omega_0}$ ($n = 0, 1, 2, \dots$) is the colimit of $\{W_n\}$.

Since $\text{card } I > \text{card } X_0$, there exists a one-to-one mapping $t_0 : X_0 \rightarrow I = W_0$. Since $\text{card } I > \alpha$, it is easy to see by induction

$$\begin{array}{ccccccc}
 X_0 & \longrightarrow & X_1 & \longrightarrow & X_2 & \longrightarrow & \cdots & \longrightarrow & X \\
 \downarrow t_0 & & \downarrow t_1 & & \downarrow t_2 & & & & \downarrow t \\
 W_0 & \xrightarrow{c_0} & W_1 & \xrightarrow{c_1} & W_2 & \xrightarrow{c_2} & \cdots & \longrightarrow & W_{\omega_0}
 \end{array}$$

that $\text{card}(W_n - \text{Im } c_n) \geq \text{card } I > \text{card } X_n$ for each n . Therefore there exist one-to-one mappings $t_n : X_n \rightarrow W_n$ such that

- (a) for $x \in X_{n-1}$ always $t_n(x) = c_{n-1} \cdot t_{n-1}(x)$ and
- (b) for $x \in X_n - X_{n-1}$ always $t_n(x)$ is out of the image of c_{n-1} .

Hence, we get a one-to-one mapping $t : X \rightarrow W_{\omega_0}$ with restrictions to X_n equal to $w_n \cdot t_n$. Supposing $Ft(x) \in \text{Im } Fw_n$ for some n , we shall derive a contradiction, thus concluding the whole proof.

Let us apply the proposition in Section 17 to the point $y = Ft(x)$. First, we remark that $c_n : W_n \rightarrow W_{n+1}$ are all one to one (this is clear for c_0 ; since $I \neq \emptyset$, c_0 is then a split mono, hence $c_1 = 1 + Fc_0$ is a mono, etc.) and, therefore, so are $w_n : W_n \rightarrow W_{\omega_0}$. Since both $t(x)$ and $w_n(W_n)$ are elements of \mathfrak{F}_y , so is $t(x) \cap w_n(W_n) = t \cdot i(X_n)$, where $i : X_n \rightarrow X$ is the inclusion map. Hence, $y = F(t \cdot i)(x_0)$ for some $x_0 \in FX_n$. Since t is a split mono, Ft is one to one and so $Fi(x_0) = x$, because

$$Ft \cdot (Fi(x_0)) = y = Ft(x).$$

Thus, $x \in Fi(FX_n)$. This is a contradiction to $X_n \notin \mathfrak{F}_x$.

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