Asymptotic arc-components of unimodal inverse limit spaces

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Abstract

We consider the inverse limit space \((I, f)\) of a unimodal bonding map \(f\) as fixed bonding map. If \(f\) has a periodic turning point, then \((I, f)\) has a finite non-empty set of asymptotic arc-components. We show how asymptotic arc-components can be determined from the kneading sequence of \(f\). This gives an alternative to the substitution tiling space approach taken by Barge and Diamond [Ergodic Theory Dynamical Systems 21 (2001) 1333].

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1. Introduction

Inverse limit spaces of endomorphisms appear as the global attractors of many dynamical systems [21]. For instance, the relevant example for this paper is the global attractor of Hénon maps, cf. [5]. We will study inverse limit spaces for which a unimodal map of the interval is the (fixed) bonding map. Since two conjugate maps gives rise to homeomorphic inverse limit spaces, it suffices to consider quadratic maps \(f(x) = 1 - ax^2\) with \(a \in [0, 2]\). Then \(f\) has a unique critical point \(c = 0\) and maps the interval \(I := [1 - a, 1]\) onto itself. Write \(X := (I, f)\) for the inverse limit space with \(f\) as single bonding map. For values of \(a\) below the Feigenbaum parameter \((a < a_{\text{Feig}} \approx 1.40115\ldots)\), the structure
of the inverse limit space $X$ is relatively simple, see [6]. For $a > a_{\text{Feig}}$, the known results are largely restricted to parameters for which the critical orbit is finite. In this case, $X$ consists of uncountably many arc-components which (assuming $f$ is non-renormalizable) lie dense in $X$. If $c$ is periodic of period $N$, then $X$ contains $N$ endpoints, cf. [7], and the arc-component of these endpoints are continuous images of the half-line $[0, \infty)$. All other arc-components are continuous images of $\mathbb{R}$. If $c$ is strictly preperiodic, no endpoints exists [8]; instead there are turnlink points at which $X$ is not homogeneous, [10]. Each point that is neither endpoint nor turnlink point has a neighborhood homeomorphic to the product of a Cantor set and an interval.

There are still more inhomogeneities. In [1], Barge and Diamond point out the existence of asymptotic arc-components. Two arc-components $C$ and $\tilde{C}$ are asymptotic if there exists parametrizations $\varphi, \tilde{\varphi}: \mathbb{R} \to C, \tilde{C}$ such that $\lim_{t \to \infty} d(\varphi(t), \tilde{\varphi}(t)) = 0$. Here $d$ is a metric on $X$ compatible with the topology. In this paper we use symbolic dynamics to describe the asymptotic arc-components. Depending on the kneading sequence, we can algorithmically determine the pattern of asymptotic arc-components. They appear in $k$-fans (i.e., $k$ arc-components which are all asymptotic in one direction), or $k$-cycles (i.e., $k$ arc-components each of which is asymptotic in either direction to a neighboring arc-component), or more complicated combinations of these two, see Fig. 1.

The induced homeomorphism permutes the asymptotic arc-components and from the structure of this permutation, the pattern of the asymptotic arc-components can be further analyzed. The pattern gives some visualization of why inverse limit spaces of non-conjugate “periodic” unimodal maps are non-homeomorphic, contributing to the partial results on the classification in [3,11,19,15]. For example, it seems that for any given period $N \geq 6$, the four $N$-periodic kneading sequences appearing last in the parity-lexicographical order, all lead to non-homeomorphic inverse limit spaces, cf. [19, Theorem 5.1]. However, our results provide no complete classification, and hence cannot replace the claim of [16, 18] that all “periodic” inverse limit spaces are non-homeomorphic.

The paper is organized as follows. The next section gives preliminaries on inverse limit spaces of unimodal maps and how to track their arc-components using kneading theory. In Section 3, we determine when two arc-components are asymptotic. Section 4 presents the results for the unimodal maps with periodic critical point up to period 8, and gives the theorems that led to this classification. The proofs are given in the final section.
2. Preliminaries

Let \( f(x) = 1 - ax^2 \) be a quadratic map on the interval \( I := [1 - a, 1] \). The critical point is \( c = 0 \), and we write \( c_i = f^i(c) \) for the \( i \)th image of \( c \). Therefore the interval \( [1 - a, 0] = [c_2, c_1] \) is the core of the map; throughout this paper, we will always restrict unimodal maps to their cores. The inverse limit space is

\[
X := (I, f) = \{(\ldots, x_{-3}, x_{-2}, x_{-1}) \mid x_i = f(x_{i-1}) \in I \text{ for all } i < 0 \}.
\]

Endow \( X \) with product topology and metric \( d(x, y) = \sum_{i<0} 2^{|x_i - y_i|} \). Let \( \pi_i(x) = x_i \) be the projection on the \( i \)th coordinate for \( i < 0 \) and \( \pi(x) := f(x_{-1}) \) for the projection on the 0th coordinate. The induced homeomorphism is

\[
\hat{f}(\ldots, x_{-3}, x_{-2}, x_{-1}) = (\ldots, x_{-3}, x_{-2}, x_{-1}, f(x_{-1}))
\]

with the right-shift as inverse \( \hat{f}^{-1} \).

We use the standard symbolic dynamics known as kneading theory for \( f \). Given \( x \in I \), the itinerary of \( x \) is the sequence \( e(x) = e_0 e_1 e_2 \ldots \), with

\[
e_i = \begin{cases} 1 & \text{if } f^i(x) > c, \\ * & \text{if } f^i(x) = c, \\ 0 & \text{if } f^i(x) < c. \end{cases}
\]

The itinerary of the critical value \( e_1 \) is called the kneading sequence and will be denoted as \( v = v_1 v_2 v_3 \ldots \). Let \( \vartheta(e_1 \ldots e_n) \) denote the number of ones in \( e_1 \ldots e_n \). We use the shorthand \( \vartheta(n) = \vartheta(v_1 \ldots v_n) \). If \( c \) is periodic, say of period \( N \), then let by convention

\[
v_N \in \{0, 1\} \quad \text{be such that } \vartheta(N) \text{ is even.}
\]

We need the parity-lexicographical ordering \( \preceq \) if \( e = e_1 e_2 e_3 \ldots \) and \( \tilde{e} = \tilde{e}_1 \tilde{e}_2 \tilde{e}_3 \ldots \) are finite or infinite words of the symbols \( 0, *, 1 \) and their first difference is at entry \( k \), then

\[
e \prec \tilde{e} \quad \text{if } \begin{cases} e_k < \tilde{e}_k & \text{and } \vartheta(e_1 \ldots e_{k-1}) \text{ is even, or} \\ e_k > \tilde{e}_k & \text{and } \vartheta(e_1 \ldots e_{k-1}) \text{ is odd,} \end{cases}
\]

where \( 0 < * < 1 \). Assuming that \( e_i = * \) implies that \( e_{i+1} e_{i+2} \ldots = v_1 v_2 \ldots \), this is a total ordering on the set of 0 1-words. A kneading sequence \( v \in \{0, 1\}^\mathbb{N} \) with \( v_1 = 1 \) is called admissible if there exists a unimodal map with \( v \) as kneading sequence. Admissible sequences can be characterized by the fact that they are shift-maximal with respect to the left-shift \( \sigma \) and the parity-lexicographical ordering:

\[
\sigma v \leq \sigma^n v \leq v \quad \text{for all } n \geq 0.
\]

Equivalent admissibility conditions for kneading sequences are given in Section 5. Given a kneading sequence \( v \), we say that a 0 1-word \( A \) is admissible (cf. [12, Chapter 1.18]) if

\[
\sigma v \leq B v \leq v \quad \text{for all subwords } B \text{ of } A.
\]

Note that a finite subword \( B \) corresponds to an interval in \( I \) whose itinerary starts with \( B \). We extend the symbolic dynamics to the inverse limit space, by giving \( x \in X \) the backward itinerary \( e(x) = \ldots e_3 e_2 e_1 \), where

\[
e_i = \begin{cases} 1 & \text{if } x_i > c, \\ * & \text{if } x_i = c, \\ 0 & \text{if } x_i < c. \end{cases}
\]
Lemma 1 (Cf. [12, Chapter 1.18]). The sequence $e$ is the backward itinerary of some point $x \in X$ if and only if

1. if $e_{-k} = \star$, then $e_{-k+1} \ldots e_{-1} = v_1 \ldots v_k$;
2. $\sigma v \preceq Bv \preceq v$ for all subwords $B$ of $\ldots e_{-3} e_{-2} e_{-1}$.

Lemma 2. Let $f$ have a periodic critical point. Let $p, \tilde{p} \in X$ have backward itineraries $e, \tilde{e}$. Then $p$ and $\tilde{p}$ belong to the same arc-component if and only if $e$ and $\tilde{e}$ have the same tail, i.e., $e_{-i} = \tilde{e}_{-i}$ for all $i$ sufficiently large.

Proof. This is contained [8]. In fact, assuming that $f$ is long-branched (or equivalently, its kneading map if bounded, see Lemma 4) is sufficient for this result. In general, the “if” direction is true: if backward itineraries $e$ and $\tilde{e}$ agree from entry $-N$ onwards, then the arc-component can be parametrized by the $-N$th coordinate. The “only if” direction fails, however. There can be arcs whose endpoint(s) have backward itineraries with different tails than the backward itineraries of the rest of the arc, cf. the endpoint characterizations in [10]. □

To describe the folding pattern of an arc-component, we define

$$\tau_R(e) = \sup \{ n \geq 1; e_{-(n-1)} \ldots e_{-1} = v_1 \ldots v_n, \vartheta(n-1) \text{ is even} \}$$

and

$$\tau_L(e) = \sup \{ n \geq 1; e_{-(n-1)} \ldots e_{-1} = v_1 \ldots v_n, \vartheta(n-1) \text{ is odd} \}.$$

It was shown in [10] (cf. [8]), that the set of points $x \in X$ with given sequence $e$ as itinerary is an arc $A$ which projects to

$$\pi(A) = [c_{\tau_L(e)}, c_{\tau_R(e)}].$$

(2)

(Here we assume that $\tau_L$ and $\tau_R$ both are finite; for examples otherwise, see [10].) Moreover, if $A$ and $\tilde{A}$ are two such adjacent arcs in the same arc-component of $X$, then for the corresponding backward itineraries $e$ and $\tilde{e}$, $e_i = \tilde{e}_i$, except for a single $i < 0$ which is either $i = \tau_L(e) = \tau_L(\tilde{e})$ or $i = \tau_R(e) = \tau_R(\tilde{e})$. This gives an algorithm to compute the folding pattern of an arc-component.

Define $(a_n)_{n \in \mathbb{Z} \setminus \{0\}}, \alpha_n = \alpha_n(e) \in \mathbb{Z}$ as follows.

$$\alpha_1(e) = \tau_R(e) \quad \text{and} \quad \alpha_{-1}(e) = \tau_L(e).$$

The coordinate $\alpha_0$ remains undefined. Next define $Re$ by

$$(Re)_i = \begin{cases} e_i & \text{if } i \neq \tau_R(e), \text{ i.e., } i \neq \alpha_1(e), \\ 1 - e_i & \text{if } i = \tau_R(e), \text{ i.e., } i = \alpha_1(e). \end{cases}$$

Similarly, $R^{-1}e$ is defined as

$$(R^{-1}e)_i = \begin{cases} e_i & \text{if } i \neq \tau_L(e), \text{ i.e., } i \neq \alpha_{-1}(e), \\ 1 - e_i & \text{if } i = \tau_L(e), \text{ i.e., } i = \alpha_{-1}(e). \end{cases}$$
Fig. 2. Embedding in the plane of the arc-component through the fixed point $p$ for $v = 101$. The horizontal lines denote adjacent arcs $A$ with constant backward itinerary. To the left, at the same horizontal line, these backward itineraries are given, and the projection $\pi(A)$ according to Eq. (2).

We continue to define $\alpha_2(e) = \tau_L(Re)$, $\alpha_3(e) = \tau_R(R^2e)$, etc., and for negative subscripts $\alpha_{-2}(e) = \tau_R(R^{-1}e)$, $\alpha_{-3}(e) = \tau_L(R^{-2}e)$, etc. In other words, the numbers $\alpha_1, \alpha_2, \ldots$ record at which entries the backward itineraries change as we follow the arc-component through $x$ to the right, while $\alpha_{-1}, \alpha_{-2}, \ldots$ record changes as we follow the arc-component through $x$ to the left. Fig. 2 gives a suggestive embedding into the plane of the arc-component through the fixed point $p$ (with backward itinerary ...111) of the induced homeomorphism for $v = 101$.

If $p \neq \tilde{p} \in X$ lie on the same arc-component $C$ (and hence their backward itineraries $e$ and $\tilde{e}$ have the same tail), then we say that $p \prec \tilde{p}$ if $\tilde{p}$ is reached from $p$ by following $C$ to the right. This means that

$$
\begin{align*}
\pi(p) &< \pi(\tilde{p}) & \text{if } e = \tilde{e}, \\
\tilde{e} & = R^k e & \text{for some } k > 0 & \text{if } e \neq \tilde{e}.
\end{align*}
$$

By abuse of notation, we say that $e \prec \tilde{e}$ in this case. Note that $\prec$ is not a transitive order relation: Since $e$ and $Re$ meet at their ‘right’ endpoints (and hence both $e \prec Re$ and $Re \prec e$ are true), what is to the right of $Re$ turns out to be to the left of $e$. We use this repeatedly in the following form:

$$Re \prec R^m e \implies e \not\prec R^m e \text{ and hence } m < 0. \quad (3)$$

Define the last discrepancy of $e$ and $\tilde{e}$ as

$$\bar{d} = \bar{d}(e, \tilde{e}) = \sup\{i; \ e_{-i} \neq \tilde{e}_{-i}\}.$$ 

Obviously, $\bar{d} = \infty$ if $e$ and $\tilde{e}$ have different tails.

---

1 Note that $(R^n)_{n \in \mathbb{Z}}$ fails the usual group property of dynamical systems, since $R^{i+m}(e) \neq (R^n)(R^m(e))$ if $m$ is odd, see (3).
Lemma 3. Let $e, \tilde{e}$ be two backward itineraries with the same tail. Then $e \prec \tilde{e}$ if and only if $\sum_{i=1}^{d-1} e_{-i}$ is even.

Proof. Assume that $p, \tilde{p} \in X$ have backward itineraries $e$ and $\tilde{e}$. Without loss of generality, $\pi_i(p), \pi_i(\tilde{p}) \neq c$ for all $i$. Let $B$ be the arc connecting $p$ and $\tilde{p}$. It is easy to see that $\pi_\alpha : B \to I$ is a homeomorphism, and that the critical point lies between $p_{-\tilde{d}} = \pi_{-\tilde{d}}(p)$ and $\tilde{p}_{-\tilde{d}} = \pi_{-\tilde{d}}(\tilde{p})$.

Assume $p_{-\tilde{d}} < c < \tilde{p}_{-\tilde{d}}$. Take $q$ in the interior of $B$ such that $q$ and $p$ have the same itinerary. Then $\pi(p) < \pi(q)$ if and only if $f^{\tilde{d}}$ preserves orientation on $(p_{-\tilde{d}}, q_{-\tilde{d}})$ if and only if $\sum_{i=1}^{d-1} e_{-i}$ is even. In this case we have indeed $\tilde{p} > p$.

The other case $\tilde{p}_{-\tilde{d}} < c < p_{-\tilde{d}}$ is treated in the same way. □

Corollary 1. Suppose two backward itineraries $e$ and $R^ne$ coincide on the rightmost $k$ entries: $e_{-k} \cdots e_{-1} = (R^n e)_{-k} \cdots (R^n e)_{-1}$. Then there exists $m < n$ such that $\alpha_m(e) > k$.

3. Asymptotic arc-components

Definition 1. Two arc-components $C$ and $\tilde{C}$ of $X$ are asymptotic if there exists parametrizations $\phi, \tilde{\phi}$ of $C, \tilde{C}$ such that $\lim_{t \to \infty} d(\phi(t), \tilde{\phi}(t)) = 0$.

Given two itineraries $e$ and $\tilde{e}$, let

$$d_\alpha(e, \tilde{e}) = \min\{n; e_{-n} \neq \tilde{e}_{-n}\}$$

be the first discrepancy. Similarly, $d_\alpha = d(R^n(e), R^n(\tilde{e}))$.

Proposition 1. Let $C$ and $\tilde{C}$ be two arc-components, containing points with backward itineraries $e$ and $\tilde{e}$. Then $C$ and $\tilde{C}$ are asymptotic if and only if the following holds:

1. There exists $k, \tilde{k} \in \mathbb{Z}$ such that $d_\alpha(R^k(e), R^k(\tilde{e})) \to \infty$ as $n \to \infty$.
2. Let $\alpha$ and $\tilde{\alpha}$ be the folding patterns of $R^k(e)$ and $R^k(\tilde{e})$, respectively. Then $|c_\alpha_n - c_{\tilde{\alpha}_n}| \to 0$ as $n \to \infty$.

Proof. We carry out the proof for the case that $c$ is periodic of period $N$. This means that Condition (2) is equivalent to $c_\alpha_n = c_{\tilde{\alpha}_n}$ for all $n$ sufficiently large.

Let $\varepsilon = \frac{1}{K} |c_i - c_j|; 0 < i < j \leq N$. If $C$ and $\tilde{C}$ are indeed asymptotic, then they have parametrizations $\phi$ and $\tilde{\phi}$ such that $d(\phi(t), \tilde{\phi}(t)) \to 0$ as $t \to \infty$. Take $t_0$ so large that $d(\phi(t), \tilde{\phi}(t)) \leq 4 |\phi \circ \psi(t) - \tilde{\phi} \circ \tilde{\psi}(t)| < \varepsilon$ for all $t \geq t_0$. Then the right folding patterns of $\phi(t_0)$ and $\tilde{\phi}(t_0)$ must be the same modulo $N$. For any backward itinerary $e$ one can find $k \in \mathbb{Z}$ such that the backward itinerary of $\phi(t_0)$ coincides with $R^k(e)$. The analogous statement holds for $\tilde{e}$. Condition (2) follow immediately. Suppose by contradiction that $d_\alpha(R^n(e), R^n(\tilde{e})) = K < \infty$ infinitely often. For each such $n$, the middle points on the arcs corresponding to $R^{k+n}(e)$ and $R^{k+n}(\tilde{e})$ are some definite distance away. This contradicts that $d(\phi(t), \tilde{\phi}(t)) \to 0$. 
 Conversely, if the folding patterns of $R^k(e)$ and $R^k(\hat{e})$ are the same modulo $N$, then there are parametrizations $\varphi$ of $C$ and $\hat{\varphi}$ of $\hat{C}$ such that $\varphi(0)$ has backward itinerary $e$, $\varphi(0)$ has backward itinerary $\hat{e}$, and $\pi \circ \varphi(t) = \pi \circ \hat{\varphi}(t)$ for all $t \geq 0$. Condition (1) implies that also $d(\varphi(t), \hat{\varphi}(t)) \to 0$.

The proof for when $e$ is not periodic is similar. □

**Example 1.** Let $\nu = 101$. The arc-component $C$ through the fixed point $p$ of $\hat{f}$ is shown Fig. 2, whereas Fig. 3 shows its folding pattern; more precisely, points $x$ with backward itinerary $e = \ldots 11111101$ and $\hat{x}$ with backward itinerary $\hat{e} = \ldots 11110110$ are compared. (If $\pi(x) < c$, then we could even take $\hat{x} = \hat{f}(x)$, which lies in the same arc-component because $\hat{f}(C) = C$.) The table shows that the two components of $C \setminus \{p\}$ are asymptotic to each other. For this reason, we want to call the arc-component $C$ self-asymptotic. (This is stronger than $C$ and $C$ being asymptotic. In the sense of Definition 1, every arc-component is trivially asymptotic to itself.) To prove that $C$ is self-asymptotic, observe that $\theta(\alpha_n) - \theta(\hat{\alpha}_n)$ is always a multiple of 3 and $|d_n(e, \hat{e})| \to \infty$. Note that $R^2e = e11$ and $R^2\hat{e} = \hat{e}11$ (and $R^3e = e111$, $R^3\hat{e} = \hat{e}1111$, and also $Re = \hat{e}1$, $R\hat{e} = e1$ and $R^6e = \hat{e}111$, $R^6\hat{e} = e111$). This similarity is used to prove the above statements, see case I in Section 5.

**Corollary 2.** If $f$ has a strictly preperiodic critical point, then $X$ has no asymptotic pair of arc-components.

**Proof.** Let $f$ have period $N$ and preperiod $M$. Suppose that $C \neq \hat{C}$ are asymptotic. Choose corresponding itineraries $e$ and $\hat{e}$ such that for the folding patterns $|c_{\alpha_n} - c_{\hat{\alpha}_n}| \to 0$ as $n \to \infty$. Hence for some $n_0 > 0$, $N$ divides $\alpha_n - \hat{\alpha}_n$ for all $n \geq n_0$. At the same time, the discrepancies $d_n \to \infty$. Fix some $n_1 \geq n_0$ such that $d_{n_1} \gg N + M$. Let $n_2 \geq n_1$ be the smallest integer such that $\alpha_{n_2} = d_{n_1} < \alpha_{n_1}$. (The case $\alpha_{n_2} = d_{n_1} < \alpha_{n_1}$ can be treated by the same argument.) Note that such $n_2$ must exists, because otherwise the discrepancy $d_{n_1}$ is never resolved, contradicting that $d_n \to \infty$.

Now we have $\alpha_{n_2} - \alpha_{n_1} = iN$ for some $i \geq 1$. Let $\eta = R^{n_2}(e)$ and $\hat{\eta} = R^{n_2}(\hat{e})$. Then

\[ \eta - \alpha_{n_2}, \ldots, \eta - 1 = \bar{\eta} - \alpha_{n_2}, \ldots, \eta - 1 = v_1 \ldots v_{n_2} - 1. \]
On the other hand,

\[ \tilde{\eta} - \tilde{\alpha}_n - 2 \cdots \tilde{\eta} - 1 = v_1 \cdots v_{\tilde{\alpha}_n - 1} = v_1 \cdots v_N v_1 \cdots v_{\tilde{\alpha}_n - 1}. \]

This contradicts that \( \nu \) is preperiodic. \( \square \)

A similar proof shows that \( X \) has no asymptotic arc-components when \( f \) has a non-recurrent critical point.

### 4. Periodic kneading sequences

Let from now on \( c \) be periodic with period \( N \). Recall that by convention we write \( \nu = v_1 \cdots v_N \), where \( v_N \in \{0, 1\} \) is such that \( \varphi(N) \) is even. In Fig. 4, we list the admissible periodic kneading sequences up to period 8, the tails of their asymptotic arc-components and the pattern these arc-components make. The kneading sequences for \( a < a_{\text{reg}} \) are left out, because their inverse limit spaces possess no asymptotic arc-components, see [6].

Suspensions of substitution shifts (also called substitution tilings spaces) have frequently been studied in the 20th century, see, e.g., monographs of Hedlund and Gottschalk [13], and Queffelec [17]. It is in this context that asymptotic arc-components (their existence and finiteness) were noticed first, cf. [17, Theorem V.21]. Since such spaces appear as orientable 2-to-1 coverings of unimodal inverse limit spaces, one can conclude that the collection \( A \) of asymptotic arc-components is non-empty and finite. In [4, Proposition 4], it is shown (after subtracting the \( N \) arc-components with endpoints), that the cardinality \( \#A \leq 2(N - 2) \). As \( \hat{f} \) permutes the asymptotic arc-components, they can be viewed as the unstable manifolds of periodic points of \( \hat{f} \), and their backward itineraries have periodic tails. The maximal (in parity-lexicographical order) shifts of these tails are shown in the fourth column. Let us write \( s \sim t \) if the tails \( s \) and \( t \) are tails of asymptotic arc-components. If only one tail \( t \) is given, it means that \( t \sim \sigma^nt \) for some \( n \geq 1 \): the tail is asymptotic to a shift of itself. For instance, in line 2, all three shifts of \( t = 101 \) arc simultaneously asymptotic, see case I of Theorem 1. This leads to a 3-fan of asymptotic arc-components (cf. Fig. 1).

In line 4, the shift of the same tail \( t = 101 \) are only pairwise asymptotic: \( t \sim \sigma t, \sigma t \sim \sigma^2 t \) and \( \sigma^2 t \sim t \). The resulting arrangement of asymptotic arc-components is called a 3-cycle, see Fig. 1.

In line 12, the tail \( t = 100010001 \) is asymptotic to its fifth shift: \( t \sim \sigma^5 t \), and similarly \( \sigma t \sim \sigma^6 t, \sigma^2 t \sim \sigma^7 t, \sigma^3 t \sim \sigma^8 t \) and \( \sigma^4 t \sim \sigma^9 t \). Hence we see five 2-fans, and \( \hat{f} \) permutes all of the 10 strands in one cycle. In line 15, we see two tails \( t = 100100 \) and \( s = 1011 \). Here \( t \sim s, \sigma t \sim \sigma s \), etc., so again we see five 2-fans, but this time \( \hat{f} \) permutes them in two separate cycles. We do not know if all homeomorphisms of \( X \) into itself are homotopic to iterates of \( \hat{f} \), but if this is true, it would follow that the inverse limit spaces of line 12 and line 15 are non-homeomorphic.\(^2\) In any case, a homeomorphism between these spaces cannot commute with the induced homeomorphisms.

\(^2\) Swanson and Volkmer [19] showed that these inverse limit spaces are indeed non-homeomorphic, but for a different reason.
In line 13, we find two different tails $t = 1010$ and $s = 10010$. Since $t \sim s$, $\sigma t \sim \sigma s$ and also $t = \sigma^2 t \sim \sigma^2 s$, $\sigma t = \sigma^3 t \sim \sigma^3 s$, we find that the four resulting 2-fans must actually be linked in pairs. The pattern consists of two copies of a linked pair, see Fig. 1.

In line 31, we have $t = 101111010$ which is asymptotic to $\sigma^2 t$. Moreover, it turns out that $\sigma^3 t \sim \sigma^6$ and $\sigma^6 t \sim t$. The same happens for $\sigma t \sim \sigma^4 t$, $\sigma^4 t \sim \sigma^7 t$, $\sigma^7 t \sim \sigma t$ and $\sigma^2 t \sim \sigma^5 t$, $\sigma^5 t \sim \sigma^8 t$, $\sigma^8 t \sim \sigma^2 t$. This gives three copies of 3-cycles.

The following theorem predicts the asymptotic arc-components and their tails. We write $\rho : \mathbb{N} \to \mathbb{N}$,

$$\rho(n) = \min\{k > n; \nu_k \neq \nu_{k-n}\}.$$  

and we use the notation $a' := 1 - a$ for $a \in [0, 1]$.
Theorem 1. Given a periodic kneading sequence $\nu = \nu_1 \ldots \nu_N$ which is not of $\ast$-product structure (i.e., $f$ is not renormalizable, see below),

Case I: $\nu = 10^{N-2}$. Then
\[
\begin{align*}
  e &= 10^{N-3}10^{N-2}1 \\
  \tilde{e} &= 10^{N-3}10^{N-2}10^l \quad \text{for } l = 1, \ldots, N - 2
\end{align*}
\]
are asymptotic to the right.

Assume that $k < N$ is such that $\rho(k) \geq N$.

Case II: $\vartheta(k)$ is odd. Then
\[
\begin{align*}
  e &= \nu_1 \ldots \nu_k \nu_1 \ldots \nu_N \\
  \tilde{e} &= \nu_1 \ldots \nu_N - k \nu_N - k + 1 \ldots \nu_k \nu_{N-k} \nu_{N-k+1} \ldots \nu_N
\end{align*}
\]
are asymptotic to the right (if admissible).

Case III: $\vartheta(k)$ is even and $N = ak + r$ for $0 \leq r < N$. Then
\[
\begin{align*}
  e &= v_{r+1} \ldots v_k v_1 \ldots v_N \\
  \tilde{e} &= v_{r+1} \ldots v_k v_1 \ldots v_N v_{N-k+1} \ldots v_N
\end{align*}
\]
are asymptotic to the right (if admissible).

The resulting basic pattern of asymptotic arc-components is a $(k - 1)$-fan in case I, $k$-fans in cases II and III, provided $a = 1$. However, if the backward itineraries $e$ and $\tilde{e}$ have same additional symmetry, a more complicated pattern may arise.

Theorem 1 does not predict, for cases II and III, whether the backward itineraries $e$ and $\tilde{e}$ are admissible. Corollary 3 shows that in each case, only one value of $k$ is possible. For most lines in Fig. 4, the choice $k = \min\{i < N; \rho(i) \geq N\}$ works, and other choices of $k$ give non-admissible backward itineraries. An exception is line 17. Here $k = 5$ is taken, whereas the minimal value $k = 3$ leads to a non-admissible backward itinerary.

Cases I–III are mutually exclusive, see Proposition 2 in Section 5. Note that, if an $N$-periodic kneading sequence $\nu$ satisfies case I, then it also satisfies case II, which leads to one more asymptotic arc-component with tail $\ldots 0000$. However, this tail is only admissible if we enlarge the interval $I$ on which $f$ is defined to $[q, 1]$, where $q = \frac{1}{2}(1 - \sqrt{4a + 1})$ is the orientation preserving fixed point of $f$. The arc-component with tail $\ldots 0000$ is the arc-component of $(\ldots, q, q, q)$, which coils into $(I, f)$ in such a way that it is asymptotic to the $(N - 1)$-fan.

A unimodal map is renormalizable if there exists a closed neighborhood $J \ni c$ and an integer $m > 1$ such that $J$, $f(J)$, $f^{m-1}(J)$ have pairwise disjoint interiors, while $f^m(J) \subset J$. Let $\mu = \mu_1 \ldots \mu_m$ be the itinerary of $f(J)$, where $\mu_m \in \{0, 1\}$, using the same convention as (1). Also let $\tilde{\mu}$ be the kneading sequence of $f^m|J$. The kneading sequence of $\nu$ can be written as $\ast$-product $\nu = \mu \ast \tilde{\mu}$, defined as
\[
\nu_i = \begin{cases}
  \mu_i \mod m & \text{if } i \mod m \neq 0, \\
  \mu_m & \text{if } i \mod m = 0, \text{ and } \vartheta(\mu_1 \ldots \mu_m) \text{ is even}, \\
  \mu'_m & \text{if } i \mod m = 0, \text{ and } \vartheta(\mu_1 \ldots \mu_m) \text{ is odd}.
\end{cases}
\]
The following theorem takes care of these renormalizable examples.

**Theorem 2.** Let \( v \) be a periodic kneading sequence.

Case IV: \( v = \mu \ast \tilde{\mu} \) has \( \ast \)-product structure, where \( \mu \) has period \( m \). Then \( X \) has \( m \) subcontinua, each of them homeomorphic to the inverse limit space of the unimodal map with kneading sequence \( \tilde{\mu} \), having corresponding asymptotic arc-components.

**Proof.** This is standard. Since \( f \) is renormalizable, the set
\[
\{ x \in \mathbb{I}^m : x^{im+j} \in f^j(J) \text{ for all } i \geq 1 \text{ and } j = 0, 1, \ldots, m-1 \}
\]
forms a subcontinuum \( H \) homeomorphic to \((J, f^m|J)\). The subcontinua \( H, \hat{f}(H), \ldots, \hat{f}^{m-1}(H) \) are pairwise disjoint, and connected to each other by additional arc-components, see [2]. Each \( H \) has the same asymptotic arc-components as \((J, f^m|J)\).

As an example, in line 32 from Fig. 4, \( v = 11 \ast 1001 \). The corresponding inverse limit space has two copies of the inverse limit space of line 2 as subcontinua, and hence two 3-fans. In addition, there is the arc-component with tail \( \ldots 1111 \), which turns out to entwine in the these subcontinua asymptotically to the two 3-fans, rendering them two linked 4-fans. The same mechanism in line 9 gives two linked 3-fans where the non-linking strands of each 3-fan are in fact asymptotic to themselves.

We think that Theorems 1 and 2 give the complete picture of asymptotic arc-components. They lead to at most \( 2(N-2) \) asymptotic arc-components (obtained in cases II and III, when \( k = N-2 \) and no additional symmetries in the asymptotic backward itineraries are present), which confirms the bound from [4].

**Conjecture 1.** Given a periodic kneading sequence, let \( \mathcal{A} \) be the collection of asymptotic arc-components of the corresponding inverse limit space. If \( C, \tilde{C} \in \mathcal{A} \), then \( C \) is asymptotic to, or coincides with, \( \hat{f}^n \tilde{C} \) for some \( n \in \mathbb{Z} \).

5. Proofs

In this section we give the proof of Theorem 1. Case I expresses the scheme of the proof in its simplest form. For the other cases, several more technical arguments are necessary to verify the basic steps. This involves a discussion of admissible words for a given kneading sequence, and of the structure of cutting and co-cutting times.

**Proof. Case I:** We have
\[
\begin{align*}
e &= 10^{N-3}110^{N-2}1, \\
\bar{e} &= 10^{N-3}110^{N-2}10',
\end{align*}
\]
so \( \alpha_1(e) = N + 1 \) and \( \alpha_1(\bar{e}) = 1 \). Consequently,
\[
\begin{align*}
Re &= 10^{N-3}110^{N-2}10^{N-2}1, \\
R\bar{e} &= 10^{N-3}110^{N-2}10^{N-1}11.
\end{align*}
\]
and \( d(Re, R\tilde{e}) = l + 1 \). Let \( n > 1 \) be minimal such that \( R^n e \) ends with \( 10^{l-1} \). We claim that
\[
\alpha_m(e), \alpha_m(\tilde{e}) \leq l + 1 \quad \text{for } m < n.
\]
Indeed, if \( \alpha_m(e) > l + 1 \) for the first time, then \( \alpha_m = N + 1 \) and \( R^m e = \ldots 10^{N-2} 1 \) ends in the same way as \( Re \). By Corollary 1 this cannot happen. If, on the other hand, \( \alpha_m(\tilde{e}) > l + 1 \), then \( \alpha_m(\tilde{e}) = N + l + 1 \) and \( R^m \tilde{e} = 10^{l-1} \). But then also \( R^n e = \ldots 10^{l-1} \) and \( \alpha_m(e) = l + 1 \). This proves the claim, and we find \( \alpha_n(e) = l + 1 \) and \( \alpha_n(\tilde{e}) = N + l + 1 \).

Therefore
\[
\begin{align*}
R^{n+1} e &= \ldots 10^{N-3} 10^{N-2} 10^{N-l-2} 10^{l-1} = e 0^{N-l-2} 110^{l-1}, \\
R^{n+1} \tilde{e} &= \ldots 10^{N-3} 110^{N-2} 10^{N-2} 110^{l-1} = \tilde{e} 0^{N-l-2} 110^{l-1},
\end{align*}
\]
which are left-shifted copies of the original backward itineraries. Since \( e \) and \( R^{n+1} e \) differ by an even number of entries, \( n + 1 \) is even. Therefore, we can repeat the argument and find that \( d_n(e, \tilde{e}) \to \infty \) while \( \alpha_n(e) - \alpha_n(\tilde{e}) \) is always a multiple of \( N \). This proves case I. \( \square \)

Define the cutting times \( (S_k)_{k \geq 0} \) and co-cutting times \( (T_l)_{l \geq 0} \) of a kneading sequence \( \nu \) as:
\[
S_0 = 1, \quad S_k = \rho(S_{k-1})
\]
and
\[
T_0 = \min \{ i > 1; \nu_i = 1 \}, \quad T_l = \rho(T_{l-1}).
\]
We list some facts of (co-)cutting times and admissibility from \([9,14,20]\).

**Lemma 4.** The following statements are equivalent:

1. A kneading sequence, \( \nu \) is admissible.
2. \( \sigma^* \nu \leq \nu \) for all \( n \geq 0 \).
3. \( \rho(n) - n \) is a cutting time for each \( n \geq 1 \).
4. \( \rho(n) - n \) is a cutting time for each cutting or co-cutting time \( n \), and \( (S_k)_{k \geq 0} \) and \( (T_l)_{l \geq 0} \) are disjoint sequences.
5. \( S_k - S_{k-1} = S_Q(k) \) for some \( Q : \mathbb{N} \to \mathbb{N} \cup \{ 0 \} \) (this map \( Q \) is called the kneading map of \( \nu \)) and
\[
\{ Q(k + j) \}_{j \geq 1} \leq_{\text{lex}} Q \{ Q^2(k) + j \}_{j \geq 1} \quad (4)
\]
for all \( k \) (\( \leq_{\text{lex}} \) is lexicographical order).

The (co-)cutting times serve as natural dividers of the kneading sequence, and they determine which subwords of \( \nu \) are admissible. We list some of these properties more precisely, writing \( \overline{S}(n) = \min \{ S_k; \ S_k > n \} \) and \( \overline{T}(n) = \min \{ T_l; \ T_l > n \} \).

**Lemma 5.**

(a) \( \vartheta(S_k) \) is odd for all \( k \) and \( \vartheta(T_l) \) is even for all \( l \).
(b) If \( \rho(n) > S(n) \), then \( \vartheta(v_{n+1} \ldots v_{S(n)}) \) is odd. If \( \rho(n) > T(n) \), then \( \vartheta(v_{n+1} \ldots v_{T(n)}) \) is odd.

(c) The word \( v_1 \ldots v_{n-1}v'_n \) is admissible if and only if \( n \) is a cutting time.

(d) If \( v \) is \( N \)-periodic (with convention (1)), then \( N \) is a co-cutting time and no cutting time is a multiple of \( N \).

**Proof.** (a) This follows by induction and the fact that \( \rho(n) - n \) is a cutting time.

(b) We prove the following statement by induction

\[
\begin{align*}
\rho(k) > S(k) & \quad \text{implies} \quad \vartheta(k) \text{ is even.} \\
\rho(k) = S(k) & \quad \text{implies} \quad \vartheta(k) \text{ is odd.} \\
\rho(k) > T(k) & \quad \text{implies} \quad \vartheta(k) \text{ is odd.} \\
\rho(k) = T(k) & \quad \text{implies} \quad \vartheta(k) \text{ is even.}
\end{align*}
\]

This is obviously true for \( k \leq \min\{i > 1 : v_i = 1\} \). Suppose (5) is true for all integers less than \( k \). Assume \( \rho(k) > S(k) \) and \( S(k) = S_i \), so \( S_{i-1} \leq k < S(k) < \rho(k) \). Then \( k > S_{i-1} \) because otherwise \( \rho(k) = S_i \). It follows that \( \rho(k - S_{i-1}) = S_i - S_{i-1} \) is a cutting time, and by induction, \( \vartheta(k - S_{i-1}) \) is odd. Therefore \( \vartheta(k) = \vartheta(S_{i-1}) + \vartheta(k - S_{i-1}) \) is even, as asserted. The remaining three statements of (5) are proven in the same way. Now statement (b) follows immediately from statement (a).

(c) The points in \( I \) with itinerary starting with \( v_1 \ldots v_n \) or \( v_1 \ldots v'_n \) form a one-sided neighborhood \( U_n \) of \( c_1 \); this is the largest neighborhood in \( I \) on which \( f^{n-1} \) is monotone. Both \( v_1 \ldots v_n \) and \( v_1 \ldots v'_n \) are admissible if and only if \( f^{n-1}(U_n) \ni c \) if and only if \( n \) is a cutting time.

(d) If \( f \) is a quadratic map with an \( N \)-periodic critical point, then varying the parameter \( a \) shows that both \( v_1 \ldots v_{N-1}0 \) and \( v_1 \ldots v_{N-1}1 \) can be realized as kneading sequences. The sequences of cutting and co-cutting times are disjoint (see Lemma 4); therefore using convention (1), we conclude that \( N \) is a co-cutting time. Since \( \vartheta(v_1 \ldots v_{N}) \) is even for all \( i \geq 1, iN \) is never a cutting time.

\[ \square \]

**Corollary 3.** The only possible admissible choices of \( k \) are in case II: \( k = \min\{i < N; \rho(i) > N\} \), and in case III: \( k = \min\{i < N; \rho(i) = N\} \).

**Proof.** In case II, the minimal \( k \) such that \( \rho(k) > N \) is the last cutting time before \( N \). If \( k' > k \) is taken such that \( \rho(k') > N \), then \( \vartheta(k') \) is odd, so \( k' \) is neither cutting nor co-cutting time. Then, for the backward itinerary \( e \) based on \( k' \), the subword \( v_1 \ldots v'_{k'} \) is not admissible.

In case III, if \( k' > k \) is such that \( \rho(k') = N \) and hence \( \vartheta(k') \) is even, then also \( \vartheta(v_{k+1} \ldots v_{k'}) \) is even. If the backward itinerary \( e \) were based on \( k' \), then it contains \( v_{k+1} \ldots v'_{k'} \) as a subword, but \( v_{k+1} \ldots v'_{k'} = v_{k+1} \ldots v'_{k'-k} \geq v_{1} \ldots v'_{k'-k} \), so it is not admissible with respect to \( v \).

\[ \square \]

Now we continue the proof of Theorem 1.
**Proof.** Case II: Since \( \vartheta (k) \) is odd, Lemma 5(b) implies that \( k \geq S_i \), the last cutting before \( N \). If \( k > S_i \), then by Lemma 5(c), the word \( v_1 \ldots v'_k \) is not admissible, so \( k = S_i \). It follows that

\[
v_1 \ldots v_k v_1 \ldots v_{N-k} = v_1 \ldots v_N
\]

and \( N - k \) is a cutting time too. In addition, \( k > N/2 \) because otherwise either \( S_{i+1} - S_i \) is not a cutting time, or \( v \) is \( k \)-periodic, contradicting Lemma 5(d). We have the backward itineraries

\[
\begin{aligned}
e &= v_1 \ldots v_k v_1 \ldots v_N, \\
\tilde{e} &= v_1 \ldots v_{N-k} v_{N-k+1} \ldots v_k v_1 \ldots v'_N v_{N-k+1} \ldots v'_N.
\end{aligned}
\]

Since \( N \) is the minimal period of \( v \), we find that \( \alpha (e) = e \). Now for \( \tilde{e} \), let \( \alpha := \alpha (\tilde{e}) = \tau_R (\tilde{e}) \), so \( \vartheta (\alpha) \) is even. If \( 1 < \alpha \leq k \), then

\[
v_1 \ldots v_{N-\alpha} v_{N-\alpha+1} \ldots v'_N = v_1 \ldots v_{N-\alpha} v_1 \ldots v_N,
\]

and by Lemma 5(b), \( \vartheta (N - \alpha) \) is even. Therefore \( \vartheta (v_1 \ldots v'_N) = \vartheta (\alpha) + \vartheta (N - \alpha) \) is both even and odd, a contradiction. If \( k < \alpha \leq N \), then

\[
v_1 \ldots v_{N-\alpha} v_{N-\alpha+1} \ldots v'_{N-k} v_{N-k+1} \ldots v'_N = v_1 \ldots v_{N-\alpha} v_1 \ldots v_N.
\]

Since \( N - k \) is a cutting time, and \( \rho (N - \alpha) = N - k \), we have by Lemma 5(b) that \( \vartheta (N - \alpha) \) is odd. Therefore \( \vartheta (v_1 \ldots v'_{N-k} v_{N-k+1} \ldots v'_N) = \vartheta (v_1 \ldots v_N) \) is even, again a contradiction. Obviously \( \alpha \neq N + 1 \), and if \( \alpha = iN + \beta \) for some \( i \geq 0 \), \( 1 < \beta \leq N + 1 \), then \( \tilde{e} \) ends with \( v_1 \ldots v_\beta \) and \( \vartheta (\beta) \) is even, which was excluded by the above arguments.

This shows that \( \alpha (\tilde{e}) = 1 \), and hence

\[
\begin{aligned}
Re &= v_1 \ldots v_k v_1 \ldots v_k v_1 \ldots v_N \\
&= v_1 \ldots v_k v_1 \ldots v_N v_{N-k+1} \ldots v_N, \\
\tilde{R} \tilde{e} &= v_1 \ldots v'_{N-k} v_{N-k+1} \ldots v_k v_1 \ldots v'_{N-k} v_{N-k+1} \ldots v_N \\
&= v_1 \ldots v'_{N-k} v_{N-k+1} \ldots v_k v_1 \ldots v'_{N-k} v_{N-k+1} \ldots v'_N v_{N-k+1} \ldots v'_N.
\end{aligned}
\]

As a result, \( d_k (e, \tilde{e}) = k + 1 \). Let \( n \geq 1 \) be minimal such that \( R^n \tilde{e} \) ends with \( v_1 \ldots v_k \). We claim:

There is no \( 1 < m < n \) such that

\[
\alpha_m (\tilde{e}) \geq k + 1 \quad \text{or} \quad \alpha_m (\tilde{e}) \geq k + 1.
\]

Let us start with \( \tilde{e} \). Take \( 1 \leq m \leq n \) minimal such that \( \alpha := \alpha_m (\tilde{e}) \geq k + 1 \). If \( k < \alpha < 2k \), then by Lemma 5(b) applied to co-cutting time \( N \), \( \vartheta (v_{N+k-\alpha+1} \ldots v'_{N-k}) \) is odd. Because \( N - k \) is a cutting time, \( \vartheta (v_{N+k-\alpha+1} \ldots v_N) \) is odd. Hence \( \vartheta (v_{N+k-\alpha+1} \ldots v'_{N-k} v_{N-k+1} \ldots v_N) \) is even. Comparing \( R \tilde{e} \) and \( \tilde{R} \tilde{e} \) using Lemma 3 gives that \( R \tilde{e} \neq \tilde{R} \tilde{e} \), a contradiction. If \( 2k < \alpha < N + k \), then Lemma 5(b) applied to cutting time \( N - k \) shows that \( \vartheta (v_{N+k-\alpha+1} \ldots v'_{N-k}) \) is odd. Therefore
\[ \partial(v_{N+k-\alpha+1} \ldots v_{N-k} v_{N-k+1} \ldots v_N v_{N-k+1} \ldots v_N) \] is even. The above argument shows that \( \bar{e} \neq R^{m+1} \bar{e} \), so again we get a contradiction.

Since \( k \) is the last cutting time before \( N \), \( \rho(k) > N \), and \( v_{k+1} \ldots v_N = v_1 \ldots v_{N-k} \). Hence \( v_{N-k} = v_N \) and \( \partial(v_1 \ldots v_{N-k} v_{N-k+1} \ldots v_N) = \nu(N) \pm 2 \). Since \( \partial(v_{i+1} \ldots v_{N+i}) = \partial(N) \) for all \( i \geq 0 \), also \( \alpha > N + k \) is not possible. This proves the first half of the claim.

We use similar arguments for \( \bar{e} \): Suppose that \( m > 0 \) is minimal such that \( \alpha := \alpha_m(e) \geq k + 1 \). If \( k < \alpha \leq N + k \), then by Lemma 5(b), \( \partial(v_{N+k-\alpha+1} \ldots v_N) \) is odd. We already know that \( \partial(v_{N-k+1} \ldots v_N) \) is odd, so \( \partial(v_{N+k-\alpha+1} \ldots v_N v_{N-k+1} \ldots v_N) \) is even. By Lemma 3, we find \( R^n e \neq R^{m+1} e \), so again by (3), \( e \neq R^{m+1} e \). The case \( \alpha = N + k + 1 \) is possible, but then \( R^n e \) ends with \( v_1 \ldots v_k \), so \( m = n \). Finally, \( \alpha > N + k + 1 \) is impossible, since \( N \) is the smallest period of \( \nu \) and \( v_1 \ldots v_k \neq v_{N-k+1} \ldots v_N \) (they have different parity). This proves the claim, and hence \( \alpha_m(e) = \alpha_m(\bar{e}) \) for all \( 1 < m < n \). We find

\[
\begin{align*}
Re &= v_1 \ldots v_k v_1 \ldots v_k v_1 \ldots v_N v_1 \ldots v_k \\
&= v_1 \ldots v_k v_1 \ldots v_N v_1 \ldots v_k \\
R\bar{e} &= v_1 \ldots v_{N-k} v_{N-k+1} \ldots v_k v_1 \ldots v_{N-k+1} v_1 \ldots v_k.
\end{align*}
\]

By assumption, \( \partial(k) \) is odd, while \( \partial(v_{N-k+1} \ldots v_N) \) is even as we saw before. Therefore \( \bar{e} \) and \( R^n \bar{e} \) differ at an odd number of entries, and hence \( n \) is odd. Using the same arguments as for \( \alpha_1 \) above, we find \( \alpha_n(e) = \tau_L(R^n e) = N + k + 1 \) and \( \alpha_n(\bar{e}) = \tau_L(R^n \bar{e}) = k + 1 \), giving a difference of \( N \). Taking one more iterate, we obtain

\[
R^{n+1} e = v_1 \ldots v_k v_1 \ldots v_k v_1 \ldots v_N v_1 \ldots v_k
\]

and

\[
R^{n+1} \bar{e} = v_1 \ldots v_{N-k} v_{N-k+1} \ldots v_k v_1 \ldots v_{N-k} v_{N-k+1} v_1 \ldots v_k
\]

Hence \( R^{n+1} e \) and \( R^{n+1} \bar{e} \) are left-shifted copies of \( e \) and \( \bar{e} \). Since \( n + 1 \) is even, we can repeat the argument and find that \( d_n(e, \bar{e}) \to \infty \) while \( \alpha_n(e) - \alpha_n(\bar{e}) \) is always a multiple of \( N \). This proves case II.

Case III: We will use the same arguments as for case II. Since \( r = \rho(ak) - ak, r \) is a cutting time and hence \( \partial(r) \) is odd. We started with

\[
\begin{align*}
e &= v_{r+1} \ldots v_N v_1 \ldots v_N \\
\bar{e} &= v_{r+1} \ldots v_N v_1 \ldots v_{N-k+1} \ldots v_N.
\end{align*}
\]
Since $N$ is the smallest period of $\nu$, $\alpha_1(\nu) = iN + 1$ for some $i \geq 1$. If $i > 1$, then $v_1 \cdots v_N v_r v_{r+1} \cdots v_{k} v_1 \cdots v_N = v_1 \cdots v_k v_1 \cdots v_N v_1 \cdots v_N$, so $\rho(k - r) > N$. Note that $\vartheta(k - r) = \vartheta(k) - \vartheta(r)$ is odd. Obviously, $k - r < N/2$ and there is always a cutting time $S$ between $N/2$ and $N$. By Lemma 5(b), $\rho(k - r) \leq S$, giving a contradiction. Therefore $\alpha_1(\nu) = N + 1$.

Now for $\tilde{\nu}$, let $\nu := \alpha_1(\tilde{\nu}) = \tau_R(\tilde{\nu})$, so $\vartheta(\alpha - 1)$ is even. We have

$$\vartheta(\nu_{N-k+1} \cdots v_N') = \vartheta(v_1 \cdots v_N') - \vartheta(N-k) = \vartheta(N) \pm 1 - (a - 1)\vartheta(k) - \vartheta(r) \text{ is even.}$$

If $\nu = k + 1$, then $v_{N-k+1} \cdots v_N = v_1 \cdots v_k$. Therefore $\rho(N-k) = N$ whence $\rho((N-k) - (N-k)) = k$ is a cutting time by Lemma 4. But $\vartheta(k) = \vartheta(\alpha - 1)$ cannot be both odd and even. If $1 < \alpha \leq k$, then since $v_{N-a+1} \cdots v_N = v_1 \cdots v_a$ and $\vartheta(\alpha)$ is even, $\vartheta(N - \alpha)$ is also even. But then $\vartheta(N) = \vartheta(N - \alpha) + \vartheta(v_{N-a+1} \cdots v_N')$ is odd, contradicting our convention (1).

If $k + 1 < \alpha \leq k + N$, then by Lemma 5 applied to co-cutting time $N$, $\vartheta(v_{N+k-a+1} \cdots v_N)$ is odd. Hence $\vartheta(v_{N+k-a+1} \cdots v_N) = \vartheta(v_{N+k-a+1} \cdots v_N') + \vartheta(v_{N-k+1} \cdots v_N')$ is odd, and $\alpha \neq \tau R(\tilde{\nu})$. If $\alpha > k + N$, then we repeat the above arguments on $\beta = \alpha - iN$ such that $k + 1 \leq \beta \leq k + N$.

The remaining possibility is $\alpha_1(\tilde{\nu}) = 1$, so

$$\begin{align*}
Re &= v_{r+1} \cdots v_k v_1 \cdots v_N v_r v_{r+1} \cdots v_{k-1} \cdots v_N, \\
\tilde{R} &= v_{k+1} \cdots v_k v_1 \cdots v_N v_{k-1} \cdots v_N.
\end{align*}$$

Take $n$ minimal such that $R^n e$ ends with $v_1 \cdots v_k$. We claim:

There is no $1 < m < n$ such that $\alpha_m(e) \geq k + 1$ or $\alpha_m(\tilde{e}) \geq k + 1$.

First take $m$ is minimal such that $\alpha := \alpha_m(e) > k$. If $\alpha = k + 1$, then $m = n$ by definition of $n$. If $k + 1 < \alpha \leq N + k - r$, then by Lemma 5(b) applied to co-cutting time $N$, $\vartheta(v_{N+k-a+1} \cdots v_N')$ is odd. From this it follows that $\vartheta(v_{N+k-a+1} \cdots v_N v_{N-k+1} \cdots v_N)$ is even, and by Lemma 3, $Re \in R^{m+1} e$. Therefore (3) shows that $e \notin R^{m+1} e$, a contradiction.

If $N + k - r < \alpha \leq 2N + k - r$, then by Lemma 5(b) applied to co-cutting time $N$ we get $\vartheta(v_{N+k-r-a+1} \cdots v_N)$ is odd. Since $\vartheta(r)$ is odd, also $\vartheta(v_{r+1} \cdots v_N') = \vartheta(N) \pm 1 - \vartheta(r)$ is even and as shown before, $\vartheta(v_{N-k+1} \cdots v_N)$ is odd. Therefore $\vartheta(v_{N+k-r-a+1} \cdots v_N v_{r+1} \cdots v_N' v_{N-k+1} \cdots v_N)$ is even, and we apply Lemma 3 to conclude that $Re \in R^{m+1} e$ and hence $e \notin R^{m+1} e$. Finally, if $\alpha > 2N + k - r$, then $\alpha = iN + k - r + 1$ for some $i \geq 2$. However, since $r$ is a cutting time, say $r = S_i < N/2$, and since there is a next cutting time $S_j < N/2$, we obtain that $v_{S_j+1} \cdots v_{S_j+1} = v_1 \cdots v_{S_j+1} - S_i \neq v_1 \cdots v_{S_j+1} - S_i$. This shows that $\alpha \neq iN + k - r + 1$, so $\alpha > 2N + k - r$ is also impossible.

Now for $\tilde{e}$, assume that $m$ is minimal such that $\alpha := \alpha_m(\tilde{e}) > k$. If $\alpha = k + 1$, then $R^m e$ and $R^m \tilde{e}$ both end with $v_1 \cdots v_k$, so $m = n$. If $k + 1 < \alpha \leq k + N$, then the above argument
most one of cases II and III is admissible. Hence at an odd number of entries, so that $\alpha N = \alpha R^N e = e$. Therefore $\alpha N$ are left-shifted copies of $e$ and $\alpha R^N e$. Since $n + 1$ is odd and $\alpha N = \alpha N$ is odd, we can repeat the argument and find that $\alpha N = \alpha N$ is always a multiple of $N$. This proves case III. □

Proposition 2. Cases II and III of Theorem 1 are mutually exclusive.

Proof. The idea of the proof is to find subwords of $\alpha$ for case II and $e$ for case III that cannot be simultaneously majorized, in the parity-lexicographical order, by $\nu$. Hence at most one of cases II and III is admissible.

Let $k$ be as in case II, i.e., $k = S_i$ is the last cutting time before $N$ and $N - k =: S_i$ is again a cutting time. Let $l$ serve the role of $k$ in case III. Then $r = N - al$ is also a cutting time, say $r =: S_i$. We distinguish two cases.

(1) First assume that $k < l$, so $j < i$. If in case II,

$$\nu_1 \ldots \nu_{N-k} \nu_{N-k+1} \ldots \nu_k \nu_1 \ldots \nu_{N-k} \nu_{N-k+1} \ldots \nu_N$$

is admissible, then

$$\nu_{N-k} \nu_1 \ldots \nu_{N-k+1} \ldots \nu_k \nu_1 \ldots \nu_{N-k} \nu_{N-k+1} \ldots \nu_N$$

is admissible, and

$$\nu_{N-k} \nu_1 \ldots \nu_{N-k+1} \ldots \nu_k \nu_1 \ldots \nu_{N-k} \nu_{N-k+1} \ldots \nu_N$$
If in case III, if \( v_{r+1} \ldots v_l v_{l+1} \ldots v_N \) is admissible, then
\[
v_{N-k+1} \ldots v_N v_{r+1} \ldots v_l = v_1 \ldots v_S v_{r+1} \ldots v_l \leq v.
\] (7)

Recall the kneading map \( Q : \mathbb{N} \to \mathbb{N} \cup \{0\} \) from Lemma 4. By construction \( Q(k) < k \) for each \( k \). Condition (4) is in fact a way to express the paritalexicographic order. If \( v \) and \( \tilde{v} \) are two 01-words (or kneading sequences) starting with 1, then for both cutting times and kneading maps \( Q \) respectively \( \tilde{Q} \) can be denned. By induction, it is not hard to show that \( v \leq \tilde{v} \) if and only if \( \{ Q(t) \}_{t \geq 1} \succeq_{\text{lex}} \{ \tilde{Q}(t) \}_{t \geq 1} \). Applying this to (6) and (7), we obtain
\[
\{ Q(t) \}_{t=i+1}^b \succeq_{\text{lex}} \{ Q(t) \}_{t=j+1}^{b+j-i} \quad \text{and} \quad \{ Q(t) \}_{t=i+1}^b \succeq_{\text{lex}} \{ Q(t) \}_{t=j+1}^{b+i-j}.
\] (8)

Therefore \( Q(i+t) = Q(j+t) \) for \( t = 1, \ldots, b - i \), which is still possible, but it implies that \( Q(t) < i \) for \( t \leq b \). A closer look at (6) shows that to fulfill the second condition of (8), we also need \( N - S_b = S_i = S_{Q(b+j-i)} \), but this is impossible.

(2) Assume that \( l < k \). Recall that \( N = al + r \) and \( r = S_j \). Let \( S_b \) be the last cutting time before \( l \). Thus \( S_i, S_j \leq S_b < l < k = S_b \). If, in case II
\[
v_1 \ldots v_{N-k} v_{N-k+1} \ldots v_{k} v_1 \ldots v_{N-k} v_{N-k+1} \ldots v_N \]
is admissible, then since \( v_{r+1} \ldots v_k v_1 \ldots v_{N-k} = v_{r+1} \ldots v_k v_{k+1} \ldots v_N \) and \( v_{a+l} \ldots v_N = v_1 \ldots v_r \), we find the subword
\[
v_{a+l} \ldots v_k v_{N-k} \ldots v_N = v_1 \ldots v_S v_{S_j+1} \ldots v_{S_b} \leq v.
\] (9)

If in case III: \( v_{r+1} \ldots v_l v_1 \ldots v_N \), is admissible, then
\[
v_{N-k+1} \ldots v_N v_{r+1} \ldots v_l = v_1 \ldots v_S v_{S_j+1} \ldots v_l \leq v.
\] (10)

Combining (9) and (10) gives
\[
\{ Q(t) \}_{t=i+1}^b \succeq_{\text{lex}} \{ Q(t) \}_{t=j+1}^{b+j-i} \quad \text{and} \quad \{ Q(t) \}_{t=i+1}^b \succeq_{\text{lex}} \{ Q(t) \}_{t=j+1}^{b+i-j}.
\] (11)

Therefore \( Q(i+t) = Q(j+t) \) for \( t = 1, \ldots, \tilde{b} - j \). Combining (10) with the second part of (11), we get that \( S_{Q(b+1)} > l - S_b \geq S_{Q(b+i-j+1)} \). The first part of (11), however, yields \( Q(i + \tilde{b} - j + 1) \geq Q(\tilde{b} + 1) \). This contradiction completes the proof. \( \square \)

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References


Further reading