

## Nearly-Neighborly Families of Tetrahedra and the Decomposition of Some Multigraphs

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A family of  $d$ -polyhedra in  $E^d$  is called nearly-neighborly if every two members are separated by a hyperplane which contains facets of both of them. Reducing the known upper bound by 1, we prove that there can be at most 15 members in a nearly-neighborly family of tetrahedra in  $E^3$ . The proof uses the following statement: "If the graph, obtained from  $K_{16}$  by duplicating the edges of a 1-factor, is decomposed into  $t$  complete bipartite graphs, then  $t \geq 9$ ." Similar results are derived for various graphs and multigraphs. © 1988 Academic Press, Inc.

A  $d$ -polyhedron is the finite intersection of closed half-spaces in  $E^d$ , having an interior point. A family  $F$  of  $d$ -polyhedra in  $E^d$  is called *nearly-neighborly* [13] if for every two members there exists a hyperplane, which separates them and contains a facet of each. This notion is closely related to the notion of neighborliness, where a family of  $d$ -polyhedra in  $E^d$  is called *neighborly* [4, 6, 7, 13–16] if every two members meet in a  $(d-1)$ -dimensional set; this set lies in a hyperplane which separates the two members and which contains a facet of each one of them. Thus a neighborly family is also nearly-neighborly.

Following [13], and slightly changing the notation, let  $g_u(d, k)$  ( $f_u(d, k)$ ) denote the maximum number of  $d$ -polyhedra in a nearly-neighborly (neighborly, respectively) family in  $E^d$ , in which every member has at most  $k$  facets. Let  $g_b(d, k)$  and  $f_b(d, k)$  denote the corresponding maxima, when restricted to *bounded*  $d$ -polyhedra (i.e., to convex  $d$ -polytopes), having at most  $k$  facets.

Clearly,  $f_b(d, k) \leq g_b(d, k) \leq g_u(d, k)$  and  $f_b(d, k) \leq f_u(d, k) \leq g_u(d, k)$ .

Tietze [11] and Besicovitch [3] gave examples of infinite neighborly families in  $E^3$ ; these examples show that  $f_b(3, k)$  tends to  $\infty$  as  $k$  tends to  $\infty$ ; the same is true for all  $d \geq 3$ .

The first proof of the finiteness of  $g_b(d, k)$ , hence (as is easily seen) the finiteness of all the other functions as well, conjectured in [4], was given in [13]; the best known upper bound for  $g_b(d, k)$  is  $2^k$ , due to Perles [8].

Considering neighborly families of tetrahedra in  $E^3$ , Bagemihl [1] showed that  $8 \leq f_b(3, 4) \leq 17$ ; Baston [2] reduced it to  $8 \leq f_b(3, 4) \leq 9$ ; both of them conjectured that  $f_b(3, 4) = 8$  (and similarly that  $f_b(d, d+1) = 2^d$  for all  $d$ ). We [15, 16] have recently proved this conjecture, showing that no neighborly families consisting of nine tetrahedra in  $E^3$  exist. The current situation with  $f_b(d, d+1)$  is given by  $2^d \leq f_b(d, d+1) \leq 2^{d+1}$ , where the upper bound is due to Perles [8] and the lower bound is due to [14].

We wish to remark that Perles' upper bound  $2^k$  for  $g_u(d, k)$  is best in case  $k = d + 1$  for all  $d \geq 2$ , i.e.,  $g_u(d, d+1) = 2^{d+1}$  for all  $d \geq 2$  (for details, see Remark 1 at the end of the paper). In addition,  $f_u(3, 4) = g_u(3, 4) = 16$  (see Remark 2).

$g_u(3, 4) = 16$  implies that  $8 \leq g_b(3, 4) \leq 16$ . We make the following

*Conjecture.* There can be at most eight nearly-neighborly tetrahedra in  $E^3$ .

A stronger conjecture would be that  $g_b(d, d+1) = f_b(d, d+1)$  for all  $d \geq 3$ .

One of the purposes of this paper is to reduce the upper bound of  $g_b(3, 4)$  from 16 to 15, which is expressed as

**THEOREM 1.** *There can be at most fifteen nearly-neighborly tetrahedra in  $E^3$ .*

The other purpose of this paper is to extend a theorem, due to R. L. Graham and H. O. Pollak [5]; this theorem states that  $K_n$ , the complete graph on  $n$  vertices, cannot be decomposed into fewer than  $n - 1$  complete bipartite graphs. Let  $b(G)$  denote the minimum number of complete bipartite graphs into which the multigraph  $G$  can be decomposed;  $b(G)$  is well defined, and it is at most equal to the number of edges in  $G$ . The Graham–Pollak theorem states that  $b(K_n) \geq n - 1$ ; in fact, it follows easily that  $b(K_n) = n - 1$ . For extensions of this theorem, see [9, 10].

Let  $M_m$  denote a matching in  $K_n$ , consisting of  $m$  disjoint edges;  $2m \leq n$ . Let  $K_n + M_m$  denote the multigraph, obtained from  $K_n$  by taking all the edges of  $M_m$  as double edges.

We have the following results.

**THEOREM 2.**  $b(K_n + M_m) \geq n - m$  for all  $m \geq 1$ .

**THEOREM 3.** *Let  $K_n + M_m$  ( $m \geq 2$ ) have a decomposition into  $n - m$  complete bipartite graphs  $K_{A_j, B_j}$ , where  $|A_j| \leq |B_j|$ . Then, for each  $j$ ,  $|A_j| \geq m$  or  $|A_j| \leq n - 2m$ .*

COROLLARY 1.  $b(K_{2m} + M_m) \geq m$ , and equality holds only for  $m = 2$ .

*Proof of Theorem 2.* The proof uses Tverberg's [12] proof of the Graham–Pollak theorem, in a form due to R. L. Graham (private communication). Let the vertex set of  $K_n$  be  $\{1, 2, \dots, n\} = N$  and let  $M_m = \{(2j-1, 2j) \mid 1 \leq j \leq m\}$ . Suppose  $K_n + M_m$  has a decomposition into  $t$  complete bipartite graphs; denote these  $t$  graphs by  $K_{A_j, B_j}$ ,  $1 \leq j \leq t$ ,  $\emptyset \neq A_j$ ,  $B_j \subset N$ ,  $A_j \cap B_j = \emptyset$ , and  $|A_j| \leq |B_j|$  for all  $j$ . Thus we have

$$K_n + M_m = \sum_{j=1}^t K_{A_j, B_j} \quad (\text{edge-disjoint sum}). \tag{1}$$

Consider the following system of homogeneous linear equations in the  $n$  variables  $x_1, \dots, x_n$ :

$$\sum_{i \in A_j} x_i = 0 \quad \text{for all } j = 1, \dots, t. \tag{2}$$

$$\sum_{i \in N} x_i = 0. \tag{3}$$

By squaring (3) we get

$$\begin{aligned} 0 &= \left( \sum_{i \in N} x_i \right)^2 = \sum_{i \in N} x_i^2 + 2 \sum_{1 \leq i < k \leq n} x_i x_k \\ &= \sum_{i \in N} x_i^2 + 2 \sum_{\substack{(i,k) \in E(K_n) \\ i < k}} x_i x_k \\ &= \sum_{i \in N} x_i^2 + 2 \left[ \sum_{\substack{(i,k) \in E(K_n + M_m) \\ i < k}} x_i x_k - \sum_{\substack{(i,k) \in E(M_m) \\ i < k}} x_i x_k \right] \\ &= \sum_{i \in N} x_i^2 + 2 \left[ \sum_{j=1}^t \left( \sum_{i \in A_j} x_i \right) \left( \sum_{i \in B_j} x_i \right) - \sum_{j=1}^m x_{2j-1} x_{2j} \right] \\ &= \sum_{j=1}^m (x_{2j-1} - x_{2j})^2 + \sum_{i=2m+1}^n x_i^2. \end{aligned}$$

Remark that (2) has been used in the last step, to cancel the middle term. It follows that the system (2), (3) satisfies

$$\begin{aligned} x_{2j-1} &= x_{2j} & \text{for all } j, 1 \leq j \leq m, \\ x_i &= 0 & \text{for all } i, i \geq 2m + 1. \end{aligned} \tag{4}$$

The case  $m = 0$  is just Tverberg's proof of the Graham–Pollak theorem, since (4) means that (2), (3) has only the trivial solution, thus  $t + 1 \geq n$  or  $t \geq n - 1$ .

If  $m = 1$ , then (4) means that  $x_1 = x_2$  and  $x_i = 0$  for all other values of  $i$ ; by (3),  $x_1 + x_2 = 0$ , therefore  $x_1 = x_2 = 0$  as well. It follows that in this case, too, there exists only the trivial solution, hence  $t \geq n - 1$ , which for  $m = 1$  means also that  $t \geq n - m$ .

Suppose  $m \geq 2$ ; from (4) it follows that  $\sum_{j=1}^m x_{2j-1} = \sum_{j=1}^m x_{2j}$ , while (3) implies that  $\sum_{j=1}^m x_{2j-1} + \sum_{j=1}^m x_{2j} = 0$ ; therefore each one of these sums is equal to 0, and we get

$$\begin{aligned} x_{2j-1} &= x_{2j} && \text{for all } j, 1 \leq j \leq m-1, \\ x_{2m-1} &= x_{2m} = - \sum_{j=1}^{m-1} x_{2j} && (5) \\ x_i &= 0 && \text{for all } i, i \geq 2m+1. \end{aligned}$$

Thus, the dimension of the solution set of (2), (3) is at most  $m - 1$ , and the rank of the system (2), (3) is at most  $t + 1$ ; it follows that

$$n = \text{rank of system} + \text{dimension of solution} \leq (t + 1) + (m - 1) = t + m,$$

therefore  $t \geq n - m$ .

*Proof of Theorem 3.* Suppose that for some  $m \geq 2$  and some  $n, n \geq 2m$ ,  $K_n + M_m$  has a decomposition into  $n - m$  complete bipartite graphs  $K_{A_j, B_j}$ ,  $1 \leq j \leq n - m$ . Applying the procedure of the proof of Theorem 2, we get a system (2), (3) which has the solution (5), in terms of the  $m - 1$  parameters  $\{x_{2j} \mid 1 \leq j \leq m - 1\}$ , and so that the solution set has dimension exactly  $m - 1$ ; thus the parameters are linearly independent. In particular, for each  $j, 1 \leq j \leq n - m (=t)$ , the  $j$ th equation in (2) does not represent a linear dependence of the parameters  $\{x_{2j} \mid 1 \leq j \leq m - 1\}$ . Due to the special coefficients in the equations in (2), it follows that

either  $A_j \cap \{1, 2, \dots, 2m\}$  contains at least one of the two numbers  $2m - 1$  and  $2m$ , and for each one of them appearing in  $A_j$  there must be  $m - 1$  other integers in  $A_j$ , one of  $x_{2j-1}$  and  $x_{2j}$  for all  $j, 1 \leq j \leq m - 1$ ,

or else  $A_j \cap \{1, 2, \dots, 2m\} = \emptyset$ , implying that  $A_j \subset \{2m + 1, \dots, n\}$ .

In the first case  $|A_j| \geq m$ , therefore  $m \leq |A_j| \leq |B_j|$ , and in the latter case  $|A_j| \leq n - 2m$ .

*Proof of Corollary 1.* For all  $m \geq 2$ ,  $b(K_{2m} + M_m) \geq m$ , by Theorem 1. Trivially,  $b(K_2 + M_1) = 2$ .

Suppose that for some  $m \geq 2$ ,  $b(K_{2m} + M_m) = m$ , say  $K_{2m} + M_m = \sum_{j=1}^m K_{A_j, B_j}$ , where  $|A_j| \leq |B_j|$ . Now  $2m \geq |A_j| + |B_j| \geq 2|A_j| \geq 2m$  (by Theorem 3), so that  $|A_j| = |B_j| = m$ . Thus  $K_{2m} + M_m$  has a decomposition into  $m$  copies of  $K_{m,m}$ , which implies that  $m = 2$ .

To see that  $b(K_4 + M_2) = 2$ , we observe that  $K_4 + M_2$  has the following decomposition into  $K_{\{1,3\},\{2,4\}} + K_{\{1,4\},\{2,3\}}$ .

It is not hard to show that  $b(K_6 + M_3) = 4$ , using the inequality  $\geq 4$ , due to Corollary 1, and the decomposition

$$K_6 + M_3 = K_{\{1,3\},\{2,4\}} + K_{\{1,4\},\{2,3\}} + K_{\{5\},\{1,2,3,4,6\}} + K_{\{6\},\{1,2,3,4,5\}}.$$

In fact, the following recursive relation holds.

**THEOREM 4.** *If  $p$  and  $q$  are natural numbers and  $m = p + q$ , then  $b(K_{2m} + M_m) \leq b(K_{2p} + M_p) + b(K_{2q} + M_q) + 1$ .*

The proof of Theorem 4 follows easily from the decomposition of  $K_{2m} + M_m$  into  $K_{2p,2q} + (K_{2p} + M_p) + (K_{2q} + M_q)$ .

Theorem 4 and Corollary 1 yield  $b(K_8 + M_4) = 5$ . The value of  $b(K_{2m} + M_m)$  for  $m \geq 5$  can be estimated:  $6 \leq b(K_{10} + M_5) \leq 7$ ,  $7 \leq b(K_{12} + M_6) \leq 9$ ,  $8 \leq b(K_{14} + M_7) \leq 10$ , and  $9 \leq b(K_{16} + M_8) \leq 11$ . In general,  $2m + 1 \leq b(K_{4m} + M_{2m}) \leq 3m$ ,  $2m + 2 \leq b(K_{4m+2} + M_{2m+1}) \leq 3m + 2$ , and  $4m + 1 \leq b(K_{8m} + M_{4m}) \leq 6m - 1$ , for all  $m \geq 2$ .

Additional relations on  $b(K_n + M_m)$  can be derived from a particular decomposition of  $K_n$  which starts with a spanning  $K_{1,n-1}$ ; thus, in general,  $b(K_n + M_m) \leq b(K_{n-1} + M_m) + 1$ . It follows that for a fixed  $m$ ,  $b(K_n + M_m) - n$  is fixed for large values of  $n$ .

We return to deal with Theorem 1, which states that  $g_b(3, 4)$  can be at most 15, i.e., that there can be at most fifteen nearly-neighborly tetrahedra in  $E^3$ . We present the

*Proof of Theorem 1.* The proof uses the idea of the proof in [8], as follows. Suppose there exists a nearly-neighborly family  $F$  in  $E^3$ , consisting of 16 tetrahedra  $P_1, \dots, P_{16}$ . Let  $H_1, \dots, H_s$  be the collection of all the planes in  $E^3$  which contain facets of some  $P_i$ , and let  $H_j^+$  and  $H_j^-$  denote the two closed half-spaces determined by  $H_j$ ,  $1 \leq j \leq s$ .

The Baston matrix  $B(F) = (b_{ij})$  of  $F$  is defined (see [2, 13-15]) by

$$b_{ij} = \begin{cases} 1 & \text{if } H_j \text{ contains a facet of } P_i \text{ and } P_i \subset H_j^+, \\ -1 & \text{if } H_j \text{ contains a facet of } P_i \text{ and } P_i \subset H_j^-, \\ 0 & \text{otherwise, } 1 \leq i \leq 16, 1 \leq j \leq s. \end{cases}$$

Each row of  $B(F)$  contains precisely four non-zero terms, corresponding to the four facets of the tetrahedron; the nearly-neighborliness of  $F$  translates into the following property of  $B(F)$ : for every two row indices  $i$  and  $k$ ,

$1 \leq i < k \leq 16$ , there exists (at least one) column index  $j$ ,  $1 \leq j \leq s$ , such that  $b_{ij} \cdot b_{kj} = -1$ , i.e.,  $\{b_{ij}, b_{kj}\} = \{1, -1\}$ .

Let  $C$  be the  $\pm 1$ -matrix, obtained from  $B(F)$  by replacing each row of  $B(F)$  with  $2^{s-4}$  rows, so that all the zero terms in the row of  $B(F)$  are replaced by either 1 or  $-1$ , in all the  $2^{s-4}$  different ways.

It follows easily that all the rows of  $C$  are different;  $C$  has  $16 \cdot 2^{s-4} = 2^s$  rows of 1 or  $-1$ , and it has  $s$  columns; therefore the matrix  $C$  is full, in the sense that every  $\pm 1$  vector on  $s$  coordinates appears exactly once in  $C$ . It follows therefore that in each *column* of  $C$  there are equal numbers of terms of each sign. This can happen only when each column of  $B(F)$  has the same number of non-zero terms of each sign.

Following [13], let  $x_{ij}$ ,  $i \leq j$ , denote the number of columns of  $B(F)$  in which there are precisely  $i$  non-zero terms of one sign and  $j$  non-zero terms of the opposite sign.

The property of  $B(F)$  which was found can be stated:  $x_{ij} \neq 0$  implies  $i = j$ .

Using Lemmas 9 and 10 of [13] it follows that  $x_{ij} \neq 0$  implies that that  $i, j \leq 4$ ; by Lemma 5 of [13] the following hold

$$2x_{1,1} + 4x_{2,2} + 6x_{3,3} + 8x_{4,4} = 64 \quad (= 16 \cdot 4), \quad (6)$$

$$x_{1,1} + 4x_{2,2} + 9x_{3,3} + 16x_{4,4} \geq 120 \quad \left( = \binom{16}{2} \right). \quad (7)$$

This Diophantine system has three possible solutions, as given in the following table:

	$x_{1,1}$	$x_{2,2}$	$x_{3,3}$	$x_{4,4}$
1.	1	0	1	7
2.	0	2	0	7
3.	0	0	0	8

A member  $P_i$  of  $F$  is said to be of *type*  $(a, b, c, d)$ ,  $a \geq b \geq c \geq d$ , if there exist precisely  $a, b, c$ , and  $d$  members of  $F$ , having one facet on any one (or more) of the four planes containing facets of  $P_i$ , such that these other members of  $F$  are separated from  $P_i$  by these (four) planes. By the near-neighborliness of  $F$  it follows that  $a + b + c + d \geq 15$ , and by Lemmas 9 and 10 of [13] it follows that  $a, b, c, d \leq 4$ . Thus members of  $F$  can be of type  $(4, 4, 4, 4)$  or  $(4, 4, 4, 3)$ .

The solutions 1 and 2 are impossible, since in these solutions  $x_{i,i} = 1$  for some  $i \leq 2$ , implying that there should be a member of  $F$  of type  $(a, b, c, d)$ , where  $\{a, b, c, d\} \cap \{1, 2\} \neq \emptyset$ .

In the case of solution 3, it follows that all the 16 members of  $F$  are of type  $(4, 4, 4, 4)$ . It means that for every member of  $F$ , the fifteen other

members of  $F$  “appear” altogether 16 times in the expression  $a + b + c + d$  of the (common) type; hence the following property holds:

For each member  $P_i$  of  $F$  there exists precisely one other member  $P_j$  of  $F$ ,  $i \neq j$ , such that  $P_i$  and  $P_j$  are separated by exactly two planes which contain facets of both of them; for all other members  $P_r$ ,  $r \neq i, j$ , of  $F$ ,  $P_i$  and  $P_r$  are separated by exactly one plane which contains facets of both  $P_i$  and  $P_r$ . (8)

Property (8) of  $F$  can be translated to the following property of  $B(F)$ :

For each row index  $i$ ,  $1 \leq i \leq 16$ , there exists a *unique* row index  $j$ ,  $1 \leq j \leq 16$ ,  $j \neq i$ , for which there exist precisely two column indices  $p$  and  $q$ ,  $1 \leq p < q \leq s$ , such that  $\{b_{i,p}, b_{i,p}\} = \{b_{i,q}, b_{i,q}\} = \{1, -1\}$ ; for all other row indices  $k$ ,  $k \neq i, j$ , there exists a unique column index  $r$  for which  $\{b_{i,r}, b_{k,r}\} = \{1, -1\}$ . (9)

Let us define the multigraph  $G(F)$  as follows:  $G(F)$  has the 16 vertices  $\{1, 2, \dots, 16\}$ ; two vertices  $n$  and  $m$  of  $G(F)$  are connected by as many edges as there are column-indices  $r$  (in  $B(F)$ ), for which  $\{b_{n,r}, b_{m,r}\} = \{1, -1\}$ .

It follows from (9) that  $G(F)$  is equal to the multigraph, obtained from  $K_{16}$  by duplicating the edges of some 1-factor (=maximal matching) of  $K_{16}$ ; i.e.,  $G(F) = K_{16} + M_8$ .

The collection of the edges of  $G(F)$  which are contributed by any one column of  $B(F)$ , a column counted by  $x_{i,j}$ , form a complete bipartite graph of the form  $K_{i,j}$ . It follows that  $G(F) = K_{16} + M_8$  has a decomposition into eight  $K_{4,4}$ , since in the solution under consideration  $x_{4,4} = 8$  and  $x_{i,j} = 0$  otherwise. However, this contradicts the inequality  $b(K_{16} + M_8) \geq 9$ , proved earlier (following Theorem 4).

Therefore there exist no nearly-neighborly families in  $E^3$  consisting of sixteen tetrahedra.

In the first few steps of the proof of Theorem 1 we have actually proved the following.

**COROLLARY 2.** *If  $F$  is a nearly-neighborly family in  $E^d$ , in which every member has at most  $k$  facets, and if  $|F| = 2^k$ , then each member of  $F$  has precisely  $k$  facets and  $B(F)$  has the property that  $x_{i,j} \neq 0$  implies that  $i = j$ .*

A similar counting argument yields the following.

**COROLLARY 3.** *If  $F$  is a nearly-neighborly family in  $E^d$ , in which every member has at most  $k$  facets, and if  $|F| = 2^k - p$ , then  $B(F)$  has the following property:  $x_{i,j} \neq 0$  implies  $j - i \leq p$ .*

We were unable to prove that  $g_b(3, 4) \neq 15$ ;\* using Corollary 3, and assuming there exists a nearly-neighborly family consisting of fifteen tetrahedra in  $E^3$ , the analogous system to (6, 7) is  $\sum_{i \leq j} (i+j)x_{i,j} = 60$  and  $\sum_{i \leq j} ijx_{i,j} \geq 105$ , where the variables are  $x_{i,j}$ , for  $0 \leq i \leq j \leq 4$  and  $j-i \leq 1$ . So far we are unable to refute some of the solutions of this system.

*Remarks.* 1. It is very easy to show that  $g_u(d, d+1) = 2^{d+1}$  for all  $d \geq 2$ ; merely observe that the following family of  $2^{d+1}$   $d$ -polyhedra in  $E^d$  is nearly-neighborly for all  $d \geq 2$ . Take in each one of the orthant the following two sets: a  $d$ -simplex occupying the corner (i.e., spanned by the origin and  $d$  points, one on each one of the semi-axes in that orthant) and the closure of the complement of this  $d$ -simplex, taken relative to the orthant.

2. It is slightly harder and less trivial to show that  $f_u(3, 4) = 16$ ; in [13, p. 280, 1-11 to p. 282, 1-15], we gave an example of a neighborly family in  $E^3$ , consisting of 16 pyramids (having quadrangular bases); each one of these pyramids has one facet which is free (see [13]; a facet of a member of a neighborly family is called *free* if it contains no one of the intersections of pairs of members). By deleting the free facet from each pyramid (i.e., if the pyramid is  $\bigcap_{i=1}^5 H_i^+$ , and  $H_5$  is the hyperplane containing a free facet, then consider  $\bigcap_{i=1}^4 H_i^+$ ), we get a neighborly family in  $E^3$ , consisting of 16 3-polyhedra, each one having four facets, thus  $f_u(3, 4) \geq 16$ ; equality follows from  $f_u(3, 4) \leq g_u(3, 4) = 16$ .

\* *Note added in proof.* Using a similar yet more detailed approach, S. Furino, B. Gamble, and J. Zako proved that there can be at most 14 nearly-neighborly tetrahedra in  $E^3$ .

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