

# Switching Reconstruction and Diophantine Equations

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*Communicated by the Editors*

Received December 19, 1989

Based on a result of R. P. Stanley (*J. Combin. Theory Ser. B* **38**, 1985, 132–138) we show that for each  $s \geq 4$  there exists an integer  $N_s$  such that any graph with  $n > N_s$  vertices is reconstructible from the multiset of graphs obtained by switching of vertex subsets with  $s$  vertices, provided  $n \not\equiv 0 \pmod{4}$  if  $s$  is odd. We also establish an analog of P. J. Kelly's lemma (*Pacific J. Math.*, 1957, 961–968) for the above  $s$ -switching reconstruction problem. © 1992 Academic Press, Inc.

## 1. INTRODUCTION

Let  $G = G(V, E)$  be a graph. Given a subset  $W \subseteq V$ , the switching  $G_W$  of  $G$  at  $W$  is the graph obtained from  $G$  by replacing all edges between  $W$  and  $V \setminus W$  by the nonedges. The multiset of unlabelled graphs  $D_s(G) = \{G_W : |W| = s\}$  is called the  $s$ -switching deck of  $G$ . A graph  $G$  is called  $s$ -switching reconstructible if it is uniquely defined, up to isomorphism, by  $D_s(G)$ .

A question concerning  $s$ -switching reconstruction of graphs was proposed by Stanley, who established the following result [11]:

**THEOREM 1.** *Suppose that the Krawtchouk polynomial*

$$K_s^n(x) = \sum_{i=0}^s (-1)^i \binom{x}{i} \binom{n-x}{s-i} \quad (1)$$

*has no even integer roots in the interval  $[0, n]$ . Then any graph with  $n$  vertices is  $s$ -switching reconstructible.*

Other, probably equivalent conditions, were given in [4] (for sufficient conditions of different types see [4, 5]).

For  $s \leq 3$  a direct calculation of the Krawtchouck polynomials [11] yields that a graph is reconstructible if

$$s = 1 \text{ and } n \neq 0 \pmod{4};$$

$$s = 2 \text{ and } n \neq t^2, \text{ where } t = 0, 1 \pmod{4};$$

$$s = 3 \text{ and } n \neq 0 \pmod{4}, n \neq (t^2 + 2)/3, \text{ where } t = 1, 2, 5, 10 \pmod{12}.$$

Note that  $t$  can be also negative.

It turned out that for  $s \geq 4$  the situation is quite different. Namely, we show that for any fixed  $s \geq 4$  and  $n \neq 0 \pmod{4}$  if  $s$  odd, then for all sufficiently large  $n$  the corresponding Krawtchouck polynomial has no integer roots. Thus, for each fixed  $s \geq 4$  and  $s \neq 0 \pmod{4}$  if  $s$  odd, all but maybe a finite number of graphs are  $s$ -vertex switching reconstructible (Theorem 4).

Our proof is based on the two following theorems on diophantine equations:

**THEOREM 2** [10]. *The equation with integer coefficients*

$$y^2 = a_0 x^n + a_1 x^{n-1} + \dots + a_n$$

*has only finitely many integer solutions if the RHS has at least three different roots over  $C$ , where  $C$  is the complex field.*

**THEOREM 3** [2, 8]. *Let  $Z[X, Y]$  be the set of polynomials in two variables with integer coefficients. If  $f \in Z[X, Y]$  is an irreducible binary form of degree at least three and  $g \in Z[X, Y]$  has degree less than the degree of  $f$  then  $f(x, y) = g(x, y)$  has only finitely many integer solutions.*

Note that Theorem 3 is ineffective whenever an effective version of Theorem 2 was established by Baker [1], although, as far as we know, these bounds are too large to solve completely diophantine equations arising in this paper (further information can be found in [9]).

Since Theorem 1 provides only a partial answer to the switching reconstruction problem one can look for parameters of a graph which are defined by  $D_s(G)$ . For  $s = 1$  some results in this direction can be found in [4, 5]. Here we establish an analog of Kelly's Lemma [3] for  $s$ -switching reconstruction. Namely, we show that the number of induced subgraphs isomorphic to a given graph  $H$  on  $m$  vertices is  $s$ -vertex switching reconstructible if  $\binom{n-m}{s} + \binom{n-m}{s-m} > (1/2)\binom{n}{s}$  (Theorem 5). A stronger result will be given for  $m = 2$  and 3 (Theorem 6).

## 2. PROOFS

**THEOREM 4 (Main Theorem).** *A graph with  $n$  vertices is  $s$ -vertex switching reconstructible if*

- (i)  $s = 1$  and  $n \not\equiv 0 \pmod{4}$ ;
- (ii)  $s = 2$  and  $n \neq t^2$ , where  $t = 0, 1 \pmod{4}$ ;
- (iii)  $s = 3$  and  $n \not\equiv 0 \pmod{4}$ ,  $n \not\equiv (t^2 + 2)/3$ , where  $t = 1, 2, 5, 10 \pmod{12}$ .

*Moreover, for each  $s \geq 4$  there exists an integer  $N_s$  such that a graph is  $s$ -switching reconstructible provided*

- (iv)  $n > N_s$ , for  $s$  even,
- (v)  $n > N_s$  and  $s \not\equiv 0 \pmod{4}$ , for  $s$  odd.

*Proof.* Throughout the proof we set  $P_s^n(y) = K_s^n((n-y)/2)$ . Thus,  $x$  is even iff  $y \equiv n \pmod{4}$ . Note that for the cases  $s = 4$  and  $s = 5$  the proof is based on Theorem 2, whenever for  $s \geq 6$  we use Theorem 3.

*Case  $s < 4$ .* Suppose that  $G$  is not reconstructible. By Theorem 1, for  $s = 1$  we have,  $P_1^n(y) = y = 0$ , hence  $n \equiv 0 \pmod{4}$ .

For  $s = 2$  we have  $P_2^n(y) = (1/2)(y^2 - n) = 0$ . Since  $n \equiv y = y^2 \pmod{4}$  then  $y \equiv 0, 1 \pmod{4}$  and (ii) follows.

For  $s = 3$  we have  $P_3^n(y) = (y/6)(y^2 - 3n + 2)$ , hence  $n \equiv (y^2 + 2)/3$  and  $y \equiv 1, 2 \pmod{3}$ . Now,  $y^2 \equiv 3n - 2 \equiv 3y - 2 \pmod{4}$ , hence  $y \equiv 1, 2 \pmod{4}$  and so,  $y \equiv 1, 2, 5, 10 \pmod{12}$ .

*Case  $s = 4$ .* For  $s = 4$  the Krawtchouck polynomial just is

$$P_4^n(y) = \frac{1}{4!} (3n^2 - 6n(y^2 + 1) + y^4 + 8y^2),$$

and has exactly four different roots for any integer  $n$ . Hence  $P_4^n(y) = 0$  yields

$$n = y^2 + 1 \pm \left( \frac{6y^4 - 6y^2 + 9}{9} \right)^{1/2}.$$

Thus  $6y^4 - 6y^2 + 9 = z^2$  for some integer  $z$ . But, by Theorem 2, this equation has only finitely many solutions, hence, by Theorem 1, all but a finite number of graphs are reconstructible from  $D_4$ .

*Case  $s = 5$ .*

$$P_5^n(y) = \frac{y}{5!} (15n^2 - 10n(y^2 + 5) + y^4 + 20y^2 + 24),$$

and the second factor again has four different roots. Thus, either  $y = 0$  and, by Theorem 1,  $n = 0 \pmod{4}$ , or

$$n = \frac{y^2 + 5}{3} \pm \frac{(10y^4 - 50y^2 + 265)^{1/2}}{15}.$$

In the last case  $(10y^4 - 50y^2 + 265)^{1/2}$  must be an integer. But, by Theorem 2, there are only finitely many such  $y$ 's.

Case  $s \geq 6$ . It is known (see, e.g., [6]) that the Krawtchouck polynomials satisfy the following recurrence relation

$$\begin{aligned} (s + 1)P_{s+1}^n(y) &= yP_s^n(y) - (n - s + 1)P_{s-1}^n(y), \\ P_0^n(y) &= 1, \quad P_1^n(y) = y. \end{aligned} \tag{2}$$

Putting  $z = y^2$  and using induction on  $s$  one obtains

$$\begin{aligned} P_{2s}^n(y) &= f_{2s}(z, n) + g_{2s}(z, n) \\ P_{2s+1}^n(y) &= z^{1/2}(f_{2s+1}(z, n) + g_{2s+1}(z, n)), \end{aligned} \tag{3}$$

where  $f_i(z, n)$  is a binary form of degree  $\lfloor i/2 \rfloor$  and  $g_i(z, n)$  is a polynomial of degree less than that of  $f_i(z, n)$ . Indeed, since  $P_s^n(y) = (1/s!)(y^s + a_1(n)y^{s-1} + \dots + a_s(n))$  has degree exactly  $s$  then  $f_s(z, n)$  is not identically zero. Rewriting (2) as

$$(s + 1)P_{s+1}^n(y) = (yP_s^n(y) - nP_{s-1}^n(y)) + (s - 1)P_{s-1}^n(y)$$

and using the induction hypothesis we convince that the first term in the RHS is a binary form of degree  $\lfloor i/2 \rfloor$  plus a polynomial of degree less than  $\lfloor i/2 \rfloor$ , whenever the second term is a polynomial of degree less than  $\lfloor i/2 \rfloor$ .

Thus, in view of Theorem 3, it is enough to show that  $f_{2s}(z, n)$  is irreducible.

For set  $Q_{2s}(z, n) = (2s)!f_{2s}(z, n)$ ,  $Q_{2s+1}(z, n) = (2s + 1)!z^{1/2}f_{2s+1}(z, n)$ . Then, by (2) and (3),  $Q_i(z, n)$  satisfy the recurrence relation

$$\begin{aligned} Q_{s+1}(z, n) &= z^{1/2}Q_s(z, n) - nsQ_{s-1}(z, n), \\ Q_0(z, n) &= 1, \quad Q_1(z, n) = z^{1/2}. \end{aligned} \tag{4}$$

By induction on  $s$  one easily gets

$$\begin{aligned} Q_{2s}(z, n) &= \sum_{i=0}^s az^{s-i}n^i = \sum_{i=0}^s (-1)^i \binom{2s}{2i} (2i - 1)!! z^{s-i}n^i \\ Q_{2s+1}(z, n) &= z^{1/2} \sum_{i=0}^s b_i z^{s-i}n^i = z^{1/2} \sum_{i=0}^s (-1)^i \binom{2s+1}{2i} (2i - 1)!! z^{s-i}n^i, \end{aligned} \tag{5}$$

where  $(2m - 1)!! = \prod_{j=1}^m (2j - 1)$ ,  $(-1)!! = 1$ .

Let now  $p$  be the largest prime less than  $2s$ . Then  $a_0 = 1$ ,  $a_s = (2s - 1)!!$  and so,  $p \nmid a_0$ ,  $p^2 \mid a_s$ . On the other hand, if  $2i - 1 \geq p$  then  $p \mid (2i - 1)!!$ , hence  $p \mid a_i$ . If  $2i - 1 < p$  then  $p \mid \binom{2s}{2i}$ , hence, again  $p \mid a_i$ . Thus,  $Q_{2s}(z, n)$  is irreducible by Eisenstein criterion (see, e.g., [12, p. 161]), and hence, the theorem is proved  $s$  even.

Similarly, for  $s$  odd  $z^{-1/2}Q_s(z, n)$  is also irreducible. Thus, for sufficiently large  $n$  the only integer root of  $P_s^n$  arises from  $z = 0$ . Then  $n - 2x = 0$  and, by Theorem 1, for a non-reconstructible graph we get  $n = 0 \pmod{4}$ . Hence the proof is completed. ■

*Remark 1.* For  $s = 4$  and  $n \leq 10^8$  the polynomial  $P_4^n(x)$  has an even root in the interval  $[0, n]$  only for  $n = 17, 66, 1521, 15043$ .

For  $s = 5$  and  $n \leq 10^8$  the corresponding exceptional values of  $n$ ,  $n \neq 0 \pmod{4}$ , are 17, 67, 289, 10882.

A question of whether 15043 is sufficiently large remains open.

**THEOREM 5.** *Let  $\mu(H \rightarrow G)$  be the number of subgraphs of  $G$  isomorphic to  $H$ . Then  $\mu(H \rightarrow G)$  is reconstructible from  $D_s(G)$ , provided  $\binom{n-m}{s-m} + \binom{n-m}{s-m} > (1/2)\binom{n}{s}$ , where  $m = |V(H)|$  and  $\binom{a}{b} = 0$  if  $a < b$  or  $b < 0$ .*

*Proof.* A switching  $G_W$  with  $|W| = k$  will be called a  $k$ -switching. Given a graph  $G$ ,  $|V(G)| = n$ , and integers  $s, m$  satisfying

$$\binom{n-m}{s} + \binom{n-m}{s-m} > \frac{1}{2} \binom{n}{s}. \tag{6}$$

Let  $L_m = \{H^1, H^2, \dots\}$  be the set of all unlabelled graphs on  $m$  vertices. Let  $A_k^m(ij)$  be a matrix whose rows and column are indexed by elements of  $L_m$  and the entries  $a_{ij} = |\{W \subseteq V(H^j) : H_W^j \simeq H^i, |W| = k\}|$ . Note that  $A_0^m = A_m^m$  is a unite matrix, since a switching of an empty set as well as of the whole set of vertices is the identity.

Consider the matrix  $B = B_s^m = \sum_{k=0}^m \binom{n-m}{s-k} A_k^m$ . Observe that any column sum of  $A_k^m$  is  $\binom{m}{k}$ , hence, for column sums of  $B$  we have

$$\sum_{k=0}^m \binom{n-m}{s-k} \binom{m}{k} = \binom{n}{s}.$$

Moreover, each diagonal element  $b_{ii}$  of  $B$  is at least  $\binom{n-m}{s-m} + \binom{n-m}{s-m}$ , the contribution of  $\binom{n-m}{s-m} A_0^m + \binom{n-m}{s-m} A_m^m$  in  $B$ . Hence, by (6),  $b_{ii} > \sum_{j \neq i} a_{ji}$  and thus,  $B$  is invertible.

Now, define a vector  $\mu_m(G) = \mu(G) = (\mu_1, \mu_2, \dots)$  where  $\mu_i = \mu(H^i \rightarrow G)$ . We also set  $\mu(D_s(G)) = \sum_{F \in D_s(G)} \mu(F)$ .

Fix  $F \subset G$ ,  $|V(F)| = m$  and  $Z \subset V(F)$ ,  $|Z| = k$ . Consider an  $s$ -switching  $G_W$  such that  $W \cap V(F) = Z$ . There are  $\binom{n-m}{s-k}$  possible choices of such a  $W$

each of which transforms  $F$  into  $F_Z$ . Therefore, the  $l$ th component of the vector  $\binom{n-m}{s-k} A_k^m \mu(G)$  is just the number of subgraphs isomorphic to  $H^l$  in  $D_s(G)$  which were obtained by a  $k$ -switching of the  $m$  vertices subgraphs of  $G$ . Therefore we have the equation  $B \mu(G) = \mu(D_s(G))$ .

Here the RHS is known, the matrix  $B$  is invertible and so, one can find  $\mu(G)$ . ■

For  $m=2$ , i.e., when  $H$  is a single edge, and  $m=3$  we will show a little more, namely,

**THEOREM 6.** *If  $m=2, 3$  then  $\mu(H \rightarrow G)$  is reconstructible from  $D_s(G)$  except, possibly, the cases  $s = \binom{t}{2}$  and  $n = t^2$  or  $n = (t-1)^2$ ,  $t = 2, 3, \dots$*

*Proof.* For  $m=2$  or 3 the matrix  $B_s^m$  can be easily calculated, namely,

$$A_1^2 = \begin{pmatrix} 0 & 2 \\ 2 & 0 \end{pmatrix}, \quad B_s^2 = \begin{pmatrix} a & 2b \\ 2b & a \end{pmatrix},$$

$$A_1^3 = A_2^3 = \begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & 2 & 0 & 3 \\ 3 & 0 & 2 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix}, \quad B_s^3 = \begin{pmatrix} c & 0 & d & 0 \\ 0 & c+2d & 0 & 3d \\ 3d & 0 & c+2d & 0 \\ 0 & d & 0 & c \end{pmatrix},$$

where

$$a = \binom{n-2}{s} + \binom{n-2}{s-2}, \quad b = \binom{n-2}{s-1},$$

$$c = \binom{n-3}{s} + \binom{n-3}{s-3}, \quad d = \binom{n-3}{s-1} + \binom{n-3}{s-2},$$

and the graphs are listed by increasing the number of edges.

Hence,  $\det B_s^2 = (a-2b)(a+2b)$ ,  $\det B_s^3 = (a+3b)^2(a-b)^2$ . One can see that in both cases  $B$  is not invertible only if  $s = \binom{t}{2}$  and  $n = t^2$ ,  $t = 2, 3, \dots$ . We omit the details. ■

*Remark 2.* One can check that the matrices  $A_k^m$  satisfy the recurrence

$$(k+1)A_{k+1}^m = (k+1)A_k^m A_1^m - (m-k+1)A_{k-1}^m, \tag{7}$$

i.e., precisely the recurrence (2) for the Krawtchouck polynomials. Indeed, the entry  $ij$  of  $A_k^m A_1^m$  is just the number of ways to obtain  $H^i$  from  $H^j$  by a two-steps switching: first  $k$  vertices and then one vertex. Thus, the result will be either  $(k+1)$ - or  $(k-1)$ -switching, and in the first case we have  $(k+1)$  choices for the first step, while in the second case there are  $(n-k+1)$  choices.

This observation shows that  $B_s^m$  is invertible iff no eigenvalue of  $A_1^m$  is the root of the polynomial

$$R_k^m(y) = \sum_{k=0}^m \binom{n-m}{s-k} P_k^m(y).$$

*Remark 3.* Note that two graphs are not  $s$ -switching reconstructible iff the corresponding columns of  $A_s^n$  are equal. It easily follows from (7) and (3) that  $A_{2s+1}^n = A_1^n C$  for some matrix  $C$ . Hence, if  $A_1^n$  has two equal columns then  $A_{2s+1}^n$  has two also. Thus if  $G$  is not 1-switching reconstructible then it is not  $(2s+1)$ -switching reconstructible for all  $s$ .

It is natural to ask whether the degree sequence of a graph is reconstructible? Stanley proved that the answer is "yes" for  $s=1$  and  $n \neq 4$  [11]. As far as we know, the question remains open even for  $s=2$ .

In conclusion let us formulate the following conjecture which can be considered as an analog of the Nah-Williams Lemma [7] for the 1-switching reconstruction problem:

*Conjecture.* Let  $D_1(G) = D_1(H)$  but  $G \not\cong H$ , then there is a pairing  $(v, \sigma(v))$ ,  $v \neq \sigma(v)$ , of the vertices of  $G$  such that the switching of any  $t$  pairs results in  $H$  for  $t$  odd and in  $G$  for  $t$  even.

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