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Building counterexamples

Irena Rusu*

L.R.I., U.R.A. 410 du C.N.R.S., bât. 490, 91405 Orsay-cedex, France

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Abstract

A conjecture concerning perfect graphs asserts that if for a Berge graph G the following three conditions hold: (1) neither G, nor \overline{G} has an even pair; (2) neither G, nor \overline{G} has a stable cutset; (3) neither G, nor \overline{G} has a star-cutset, then G or \overline{G} is diamond-free. We show that this conjecture is not valid and that, in a way, every weaker version is false too. To this end, we construct a class of perfect graphs satisfying the hypothesis above and indicate counterexamples within this class for the instances of the conjecture obtained by replacing the diamond with any graph H which is the join of a clique and a stable set.

1. Introduction

For a graph G = (V, E), let us call a *clique* any set of pairwise adjacent vertices in G. The *clique number* $\omega(G)$ of G represents the cardinality of a largest clique in G, while the *chromatic number* $\chi(G)$ of G is the minimum number of colours necessary to colour the vertices of G in such a way that any two adjacent vertices have different colours. Using these two parameters, Berge defined a graph G = (V, E) to be *perfect* if for each of its subgraphs the clique number equals the chromatic number. It was proved by Lovász [5] that a graph G is perfect if and only if its complement graph \overline{G} is perfect too. No characterization with minimal forbidden subgraphs is known for perfect graphs, although a conjecture of Berge (called the strong perfect graph conjecture, abbreviated SPGC) has been formulated thirty years ago and it is neither proved, nor invalidated.

Conjecture 1 (SPGC). A graph is perfect if and only if it contains no odd hole and no odd antihole.

A hole is a chordless cycle with at least five vertices, while an *antihole* is the complement graph of a hole. A hole is *odd* if it has odd number of edges. Graphs without odd holes and odd antiholes are usually called *Berge graphs*. An equivalent

^{*} E-mail: rusu@chambord.univ-orleans.fr.

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version of the SPGC can be formulated using the notion of *minimal imperfect graph*, which designates an imperfect graph such that every proper induced subgraph is perfect.

Conjecture 2 (SPGC'). Every minimal imperfect graph is an odd hole or an odd antihole.

As a first step toward proving (SPGC'), some properties of minimal imperfect graphs have been investigated.

A subset C of V is called a *cutset* of G if the subgraph of G induced by V - C is not connected. As it was shown by Tucker [8], a minimal imperfect Berge graph cannot contain a *stable cutset*, i.e., a cutset C such that the subgraph $[C]_G$ induced by C in G is edgeless. Moreover, in [1], Chvátal proved that no minimal imperfect graph has a *star-cutset*, that is, a cutset C containing a vertex adjacent to all the other vertices in C.

In a graph G two nonadjacent vertices x, y form an *even pair* if there is no odd chordless path joining x and y in G. It was proved by Meyniel [6] that no minimal imperfect graph has an even pair.

Unfortunately, the exact importance of these three properties of minimal imperfect graphs is not known. The following conjecture, due to Bruce Reed [7] and involving a well-known class of perfect graphs, namely the line-graphs of bipartite graphs, has been disproved by Hougardy [4]:

Conjecture 3 (*Reed*). Let G be a Berge graph satisfying the conditions below:

1. neither G, nor \overline{G} has an even pair;

2. neither G, nor \overline{G} has a star cutset.

Then G or \overline{G} is the line-graph of a bipartite graph.

Several authors proposed a weaker conjecture by substituting the conclusion "G or \overline{G} is the line-graph of a bipartite graph" with "G or \overline{G} is diamond-free", where a diamond is the graph obtained by deleting an edge of a clique on four vertices. Still weaker than the conjecture above is the one we shall examine along this paper.

Conjecture 4. Let G be a Berge graph such that

- 1. neither G, nor \overline{G} has an even pair;
- 2. neither G, nor \overline{G} has a star cutset;
- 3. neither G, nor \overline{G} has a stable cutset.
- Then G or \overline{G} is diamond-free.

Obviously, since the diamond-free Berge graphs are perfect [9], a proof of this conjecture would also be a proof of the SPGC. Our purpose here is to give a counterexample not only for this conjecture, but also for every weaker version of it obtained by replacing the diamond with any graph H which is the join of a clique and a stable set (that is, all possible edges between the two sets are present in H). The counter-

examples will be iteratively built using a composition of two graphs G_1 and G_2 into another graph G with the same properties as the initial graphs.

2. The 2-join

Let G = (V, E) be a graph and consider a partition of its vertex set V in p subsets V_1, V_2, \ldots, V_p . This is a natural way to indicate the classes of vertices in G having the same behaviour relative to a certain property, whatever is this property. In our case, the vertices in a class will have the same neighbours in the graph joined to G by the operation defined below. Some reasons of simplicity decided us to limit the approach to the case p = 3. The partition sets will be denoted by R, W, B, each of them representing the set of vertices in G coloured respectively in red, white and blue. We shall also use the term of Q-vertex for a vertex coloured in Q, where $Q \in \{R, W, B\}$.

Given two graphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ coloured in R, W and B, denote by $G = G_1 \phi G_2$ the graph G = (V, E) defined as follows:

$$V = V_1 \cup V_2$$

$$E = E_1 \cup E_2 \cup \{xy \mid (x, y) \in R_1 \times R_2\} \cup \{zt \mid (z, t) \in W_1 \times W_2\},\$$

where R_i (resp. W_i) is the set of red (resp. white) vertices in G_i , for i=1, 2. The blue vertices in G have precisely the same neighbours as they had in the two initial graphs. Notice that the new graph G may be also obtained from the graphs

$$G'_1 = (V(G_1) \cup \{x_1, y_1\}, E(G_1) \cup \{x_1v, v \in R_1\} \cup \{y_1w, w \in W_1\} \cup \{x_1y_1\})$$

and

$$G'_{2} = (V(G_{2}) \cup \{x_{2}, y_{2}\}, E(G_{2}) \cup \{x_{2}v, v \in R_{2}\} \cup \{y_{2}w, w \in W_{2}\} \cup \{x_{2}y_{2}\})$$

using the 2-join operation defined by Cornuéjols and Cunningham [3]. According to that definition, the 2-join of G'_1 and G'_2 is the graph H resulted by eliminating x_i, y_i (i = 1, 2) and joining every neighbour of x_1 (resp. of y_1) in G_1 to every neighbour of x_2 (resp. of y_2) in G_2 . A brief verification shows that G and H are in fact the same graph. That is why, for convenience, we shall say along this paper that $G = G_1 \phi G_2$ is the 2-join of G_1 and G_2 .

For every $n \ge 2$, consider now the graph F_n (already coloured) with the vertex set

$$V(F_n) = \{x_1, x_2, \dots, x_n\} \cup \{y_1, y_2, \dots, y_n\} \cup \{z_1, z_2, \dots, z_n\},\$$

where $R = \{x_1, x_2, ..., x_n\}$, $W = \{y_1, y_2, ..., y_n\}$, $B = \{z_1, z_2, ..., z_n\}$ are the sets of vertices coloured, respectively, in red, white and blue.

The edge set $E(F_n)$ is defined such that $[R]_{F_n}$, $[W]_{F_n}$, $[B]_{F_n}$ are *n*-cliques and for all $i \in \{1, 2, ..., n\}$, $[x_i, y_i, z_i]_{F_n}$ is a 3-clique (see Fig. 1).



Fig. 1. The graph F_n .

It is easy to verify that, for $n \ge 3$, F_n satisfies the hypothesis of the conjecture (the case n = 2 is particular). In order to define a class of graphs using F_n $(n \ge 2)$ and the 2-join operation, let us see which general conditions are sufficient to preserve the hypothesis.

In the four lemmas below, we shall say that a graph is *coloured* in R, W, B (or simply *coloured*) if its vertex set is partitioned in R, W, B such that the following condition holds (Q is an arbitrary colour of the set $\{R, W, B\}$):

(C1): For every Q-vertex $x \in V$, the neighbourhood $N_G(x)$ intersects each of R, W, B, but contains none of them, except possibly for Q.

As we can easily notice, the 2-join of two coloured graphs is a coloured graph too.

The proofs of the lemmas are symmetrical for R, W and for the two graphs G_1 , G_2 , therefore we shall analyse only the nonsymmetrical cases.

Lemma 1. Let G_i (i = 1, 2) be two coloured Berge graphs satisfying the following conditions:

- (C2): If $x \in B_i$, then $N_{G_i}(x) = K_1 \cup K_2$, where $K_1 \subset B_i$ and $K_2 \subset R_i \cup W_i$ are two disjoint cliques with no edge between them;
- (C3): For every odd chordless path $Ru_1 \dots u_{2k}R$ (resp. $Wu_1 \dots u_{2k}W$) in G_i , there is some $p \in \{1, 2, \dots, 2k\}$ such that $u_p \in R_i$ (resp. W_i);
- (C4): There is no odd chordless path RWW...WR (resp. WRR...RW) in \overline{G}_i .

Then $G = G_1 \phi G_2$ is a coloured Berge graph and satisfies (C2), (C3) and (C4).

Proof. We firstly show that G has no odd holes and no odd antiholes. Suppose there is an odd hole C induced in G and let V(C) be its vertex set. Obviously, C is not entirely contained in a graph G_i since these graphs are Berge. To pass from a graph into the other one, an edge xy such that x and y have the same colour is needed. Without loss of generality we may suppose that $x \in R_1$ and $y \in R_2$. Let z be the other neighbour of x along the cycle C. If $z \in R_1$ then $yz \in E(G)$ is a chord in C, a contradiction. Thus $z \in R_2$ or z is not a red vertex.

Three positions of the cycle are then possible with respect to the initial graphs.

- $C = xyP_2zx$, where $x \in R_1$, $z, y \in R_2$ and P_2 is a chordless path in $G_2 R_2$. In this case, yP_2z is an odd chordless path in G_2 in which only the extremities are red vertices. By (C3), such a path is not induced in G_2 .
- $C = xyP_2tuvP'_2zx$, where $x \in R_1$, $z, y \in R_2$, $u \in W_1$, $t, v \in W_2$ and P_2 , P'_2 are chordless paths in B_2 . Suppose that yP_2t is even. Then $V(P_2)$ must contain an odd number of *B*-vertices. But this is not possible because of (C2).
- $C = xyP_2tvP_1x$, where $x \in R_1$, $y \in R_2$, $v \in W_1$, $t \in W_2$ and P_i (i = 1, 2) is a chordless path in B_i . Again, one of the two paths must contain an odd number of *B*-vertices and that is not possible.

We deduce that G contains no odd holes.

Suppose now that G contains some odd antihole \overline{C} and that \overline{C} has a *B*-vertex. If \overline{C} has more than five vertices, then in \overline{G} the *B*-vertex is nonadjacent to at least one path on four vertices. Therefore, in G the *B*-vertex has such a path in its neighbourhood. But that contradicts (C2). Consequently, \overline{C} induces a 5-cycle in \overline{G} , thus its complement is a 5-cycle in G. But G has no odd holes, a contradiction.

We may then suppose that the antihole \overline{C} contains no *B*-vertices. By (C4), a reasoning similar to the one we used to deduce that *G* has no odd holes proves that in fact there is no odd hole in \overline{G} . Then *G* is a Berge graph.

We subsequently prove that G satisfies the conditions (C2), (C3), (C4).

(C2): By definition, the operation ϕ does not change the neighbour set of a *B*-vertex.

(C3): Assume the contrary and let x, y be two nonadjacent *R*-vertices in *G* joined by an odd chordless path *P* with no *R*-vertices. Then x and y are both in G_1 or both in G_2 (say they are in G_1). By (C2), if there are any *B*-vertices on *P*, their number is even. An even number of *W*-vertices is then required in order to have an odd chordless path (notice that if there are no *W*-vertices, *P* is entirely contained in G_1 , a contradiction). Whatever would be the repartition of the *W*-vertices in G_1 and G_2 , there is a chord in *P*.

(C4): Assume the contrary and let x, y be two nonadjacent *R*-vertices in \tilde{G} joined by an odd chordless path P' containing only *W*-vertices. Notice that in \bar{G} , the red vertices in a graph are joined to the white and to the blue vertices in the other one. Consequently, if x, y are both in G_1 , then every *W*-vertex of P' in G_2 (and there exists at least one) would be adjacent to both x and y, a contradiction. The only possible case is $x \in R_1, y \in R_2$ such that their neighbours along P' (u and, respectively, v) are in W_2 , respectively in W_1 . The *W*-vertices u and v are nonadjacent in \bar{G} , therefore at least another *W*-vertex occurs on P'. But then P' would have chords. \Box

Remark 1. Notice that the condition (C2) could be relaxed, but such a modification would also be relaxing the conclusion of Lemma 1. Later reasonings will involve precisely the indicated form of (C2).

Lemma 2. Let G_i (i = 1, 2) be two coloured graphs such that neither G_i , nor $\tilde{G_i}$ has an even pair and

- (C2): If $x \in B_i$, then $N_{G_i}(x) = K_1 \cup K_2$, where $K_1 \subset B_i$ and $K_2 \subset R_i \cup W_i$ are two disjoint cliques with no edge between them;
- (C5): For every R-vertex (resp. W-vertex) x in G_i , there is an even chordless path xRR...RW (resp xWW...WR) in G_i .

Then neither $G = G_1 \phi G_2$, nor \overline{G} has an even pair and G satisfies (C2) and (C5).

Proof. For every pair of vertices u, v in the same graph G_i , if u, v are nonadjacent (resp. adjacent) then there is an odd chordless path joining them in the same graph (resp. in its complement graph). This path is also an odd chordless path in G (resp. in \overline{G}).

Let now u, v be a pair of vertices in G such that $u \in V(G_1)$ and $v \in V(G_2)$. Four nonsymmetrical cases can occur:

- $u \in R_1$, $v \in R_2$; by (C1), there is a *W*-vertex *t* in G_1 adjacent to *u* and, obviously, nonadjacent to *v*. Also, there is a *B*-vertex *w* in G_2 adjacent to *v* and nonadjacent to *u*. Moreover, *t* and *w* are nonadjacent, so *tuvw* is a chordless path on four vertices in *G*. The complement of this graph is the odd chordless path on four vertices *uwtv* joining *u* and *v* in \overline{G} . We deduce that (u, v) is not an even pair in \overline{G} .
- $u \in R_1$, $v \in W_2$; by (C5) there is an even chordless path uRR...RW in G_1 ; then uRR...Wv is an odd chordless path in G.
- $u \in R_1$, $v \in B_2$; by (C1) there exists a *B*-neighbour *w* of *v* in G_2 and by (C2) *w* has an R-neighbour *t* that is nonadjacent to *v*. Then *utwv* is an odd chordless path joining *u* and *v* in *G*.
- $u \in B_1$, $v \in B_2$; if t is a W-vertex in G_1 adjacent to u (according to (C1)) and w a W-vertex in G_2 adjacent to v, then utwv is an odd chordless path joining u and v in G.

Thus G and \overline{G} have no even pair. Obviously, the conditions (C2) and (C5) are also valid for G. \Box

Remark 2. In fact, the *RW*-, *RB*- and *WB*-edges in G_i may be even pairs in \overline{G}_i and the conclusion is still valid. While 2-joining G_1 and G_2 , by (C1) we can find for a *WR*-edge xy in G_1 a *W*-vertex u and an *R*-vertex v in G_2 such that $uv \notin E(G)$, so xvuy is an odd chordless path joining x and y in \overline{G} . For a *WB*-edge zt in G_1 , since (C2) is true the *B*-vertex t has a *B*-neighbour a in G_1 nonadjacent to z. As before, zhas a *W*-neighbour b in G_2 and atzb is a P_4 in G that induces an odd chordless path in \overline{G} joining z and t.

Lemma 3. Let G_i (i = 1, 2) be two coloured graphs such that neither G_i , nor \overline{G}_i has a stable cutset. Then the same property holds for $G = G_1 \phi G_2$.

Proof. Assume the contrary and let S be a stable cutset in G. Since neither G_1 , nor G_2 has a stable cutset, $G'_1 = G_1 - S$ and $G'_2 = G_2 - S$ are connected. Moreover, they can have at most one colour in common and that one is B (otherwise G - S is connected). Consequently, S contains at least one of R_1, R_2 , so by (C1) it is not a stable.

Let now S' be a stable cutset in \overline{G} and let $\overline{G'_1}$, $\overline{G'_2}$ be the connected components of $\overline{G} - S'$. Neither $\overline{G'_1}$, nor $\overline{G'_2}$ contains a *B*-vertex (otherwise the two subgraphs are connected), therefore all the *B*-vertices are in S'. Since none of B_1 , B_2 is empty, S' cannot be a stable set, a contradiction. \Box

Remark 3. We can notice that in the proof we do not really need to use the hypothesis that \overline{G}_i (i = 1, 2) has no stable cutset. Indeed, suppose that \overline{G}_i may have stable cutsets, let S' be a stable cutset of \overline{G} and consider A_1, A_2, \ldots, A_s the connected components of $\overline{G} - S'$. The graph $\overline{G} - S'$ contains at least one *B*-vertex v. Without loss of generality we may assume that v is a *B*-vertex in \overline{G}_1 contained in the connected component A_1 . Then $V(G_2) \subseteq A_1 \cup S'$, so $A_2, \ldots, A_s \subseteq V(G_1)$. This implies that B_2 and one of R_2, W_2 are included in S'. But since S' is a stable set we get a contradiction to (C1).

Lemma 4. Let G_i (i = 1, 2) be two coloured graphs such that neither G_i , nor G_i has a star-cutset. Then the same property holds for $G = G_1 \phi G_2$.

Proof. We shall use a result of Chvátal [1] stating that a graph G has a star-cutset if and only if it has at least one of the following properties:

- (i) G has a vertex w such that the set of all the vertices distinct from w and not adjacent to w induces a disconnected subgraph of G;
- (ii) G has at least two nonadjacent vertices, and it has adjacent vertices v, w such that w dominates v.

Assume now that the lemma is not true and let S be a star-cutset of G. We denote by x the vertex of S adjacent to all the other vertices in S. Two nonsymmetrical cases are possible:

- $x \in B_1$; then $S \subset V(G_1)$ and $G_1 S$ is connected. The two connected components are precisely $G_1 S$ and G_2 , therefore $G_1 S$ contains no *R*-vertices and no *W*-vertices. The condition (C1) is violated.
- $x \in R_1$; denote $N'_G(x) = N_G(x) \cup \{x\}$ and $G'_1 = G_1 N'_G(x)$. Then $V(G) N'_G(x) = W_2 \cup B_2 \cup V(G'_1)$. By (C1), every *B*-vertex of G_2 has a neighbour in G_2 coloured *W*. This one is adjacent to all the *W*-vertices in G'_1 (at least one such vertex exists, by (C1)). Moreover, G'_1 is a connected graph (otherwise $N'_G(x) \cap V(G'_1)$ would be a star-cutset for G_1 , a contradiction). Conclusion: $G N'_G(x)$ is connected.

To prove that no star-cutset S exists, it is sufficient to show that every neighbour of x in G is adjacent to at least one nonneighbour of x. For the neighbours of x in G_1 , this property is insured by the fact that G_1 has no star-cutset. The neighbours of x in G_2 are R-vertices, which are adjacent to some of the W-vertices in G_2 .

Let us now prove that \overline{G} has no star-cutset. Suppose the contrary.

• $x \in B_1$; as before, denote $N'_{\bar{G}}(x) = \{x\} \cup N_{\bar{G}}(x)$ and let $\bar{G'_1} = \bar{G_1} - N'_{\bar{G}}(x)$ (connected) be the graph induced in $\bar{G_1}$ by the vertices nonadjacent to x. One has $\bar{G'_1} = \bar{G} - N'_{\bar{G}}(x)$ (since all vertices in $\bar{G_2}$ are adjacent to x), so $N'_{\bar{G}}(x)$ is not a star-cutset of \bar{G} . By (C1), there is at least one W-vertex and at least one R-vertex in G_1 adjacent to x. Therefore, in $\bar{G'_1}$ there exist at least one R-vertex and at least one W-vertex. Consequently, every neighbour $y \in \overline{G}_2$ of x is adjacent to a nonneighbour of x (the R-vertex or the W-vertex in \overline{G}_1). Since $N'_{\overline{G}}(x) \cap V(\overline{G}_1)$ is not a star-cutset in \overline{G}_1 , the same holds for the neighbours of x in \overline{G}_1 .

• $x \in R_1$; the graph \bar{G}'_1 is connected and contains at least one *W*-vertex, by (C1). This vertex is adjacent to all the *R*-vertices in \bar{G}_2 , so $\bar{G} - N'_{\bar{G}}(x)$ is connected (since $\bar{G} - N'_{\bar{G}}(x) = \bar{G}'_1 \cup R_2$). Every neighbour $y \in V(G_2)$ of x is a *W*-vertex or a *B*-vertex and, again by (C1), is adjacent to at least one *R*-vertex z in \bar{G}_2 (nonadjacent to x). Also, every neighbour $y \in V(G_1)$ of x is adjacent to a nonneighbour of x in \bar{G}_1 , since \bar{G}_1 has no star-cutset. \Box

The five conditions identified in the lemmas above are now sufficient to define the class of graphs generated by F_n $(n \ge 2)$ using the 2-join operation. Let us firstly say that a graph G is *finely-colourable* if its vertex set may be partitioned into three sets R, W, B such that the conditions (C1)–(C5) hold. A colouring of G with these properties is called a *fine-colouring* and a graph G provided with a fine-colouring is said to be *finely-coloured*.

We define the class Γ generated by the basic graphs F_n $(n \ge 2)$ as following:

- for every $n \ge 2$, $F_n \in \Gamma$;
- if G_1 and G_2 in Γ are finely-coloured in R, W, B, then $G_1\phi G_2 \in \Gamma$.

The theorem below uses the four previous lemmas to prove that Γ is correctly defined and, moreover, that every graph in $\Gamma \setminus \{F_2\}$ satisfies the hypothesis of the conjecture:

Theorem 1. The graphs in $\Gamma \setminus \{F_2\}$ are finely-colourable Berge graphs with the following properties:

- 1. neither G, nor \overline{G} has an even pair;
- 2. neither G, nor \overline{G} has a star cutset;
- 3. neither G, nor \overline{G} has a stable cutset.

Proof. It is a routine matter to verify that, for $n \ge 3$, F_n satisfies the hypothesis of Lemmas 1, 2, 3, 4.

For n = 2 not all the properties hold. Namely, there are two exceptions, both in $\overline{F_2}$:

- the nonadjacent pairs of vertices coloured in RW, RB or WB are even pairs in $\bar{F_2}$;
- the two 3-stables are stable cutsets in $\overline{F_2}$.

According to Remarks 2 and 3, the two exceptions above do not disturb the proofs of Lemmas 2 and 3.

Then we can easily prove by induction that every graph in Γ has the indicated properties. \Box

3. A class of perfect graphs

The purpose of this section is to show that the class Γ defined in Section 2 is a class of perfect graphs. To this end we prove that the fine-colourings of the graphs in

 Γ have a common structure and use this property to deduce that no P_4 of type *RRRR*, *WWWW* exists in a fine-colouring. The perfection results as a simple application of a theorem due to Chvátal et al. [2].

Since all the graphs in Γ are obtained from the basic graphs by repeated 2-joins, we can regard every $G \in \Gamma$ as a set of basic graphs whose vertices are joined accordingly to the indicated rule. Apparently, in a fine-colouring of G it is no need for a basic graph to be coloured as described before. Actually, they are coloured in this way.

Lemma 5. In a fine-colouring of $G \in \Gamma$, every basic graph is finely-coloured.

Proof. A brief verification shows that a fine-colouring of a basic graph F_n is a one-toone application from the set of *n*-cliques to the set of colours; therefore, for this easy case, the lemma is proved.

Let $G = G_1 \phi G_2$ be a graph of Γ obtained by 2-joining the finely-coloured graphs G_1 and G_2 . We denote by R_i^{old} (resp. $W_i^{\text{old}}, B_i^{\text{old}}$) the sets of *R*-vertices (resp. *W*-, *B*-vertices) in the fine-colouring of G_i used to obtain *G* and by R_i^{new} (resp. $W_i^{\text{new}}, B_i^{\text{new}}$) the vertices in G_i (i = 1, 2) coloured in *R* (resp. *W*, *B*) in an arbitrary fine-colouring of *G*. Then every vertex in R_1^{old} is adjacent to every vertex in R_2^{old} and the same holds for the *W*-vertices. The vertices in B_i^{old} (i = 1, 2) have neighbours only in G_i .

Claim 1. No R- or W-vertex in G_i (old) becomes a B-vertex in G (new).

Proof. Suppose there is such a vertex $x \in B_2^{\text{new}} \cap W_2^{\text{old}}$. Then x has at least one neighbour y in B_2^{old} . We have to consider the following three cases:

- x has a neighbour $y \in B_2^{old} \cap R_2^{new}$. Then all the neighbours of x in G_1 must be *B*-vertices (new), otherwise the *R* and *W*-neighbours of x in G do not induce a clique and this contradicts (C2). Consequently, all the vertices in W_1^{old} are now *B*-vertices and form a clique. Let z be a vertex in this set. By (C1), z has an R-neighbour t in G. If t is in G_1 , then tzxy is a P_4 coloured *RBBR* and (C3) is violated. Then all the R-neighbours of z in G are in fact in W_2^{old} . We deduce that no W-neighbour (new) u of z occurs in G_1 . Otherwise u is in R_1^{old} or in B_1^{old} and it is not adjacent to t, therefore z contradicts (C2). Now, the neighbourhood of z coloured in B contains x and at least a vertex in B_1^{old} (according to (C1) in G_1), thus it is not a clique, a contradiction.
- x has a neighbour $y \in B_2^{\text{old}} \cap W_2^{\text{new}}$. The reasoning is similar to the preceding one.
- every neighbour of x in B_2 (old) is also coloured with B in G. Then all the vertices in W_1^{old} are now coloured in R, W and form a clique. Also, every neighbour of x in R_2^{old} is a B-vertex in G.

We shall firstly prove that in W_1^{old} there exist both new *W*-vertices and new *R*-vertices. Suppose this is not the case and all the vertices in W_1^{old} are *W* in the new colouring. We then consider an old *R*-neighbour *y* of *x* (which must be a *B*-vertex

in the new colouring) and a neighbour q of y in R_1^{old} , which is not adjacent to a fixed vertex s of W_1^{old} (such a vertex exists, otherwise s would be adjacent to the entire set R_1^{old} , a contradiction). The P_4 qyxs implies $q \in R_1^{\text{new}}$. Let r be an R-neighbour (new) of x, so $r \in W_2^{\text{old}}$ (otherwise the new W- and R-neighbours of x would not form a clique). We deduce that qyxr is a P_4 coloured RBBR, except if $yr \in E$; but then r and y are two adjacent neighbours of the B-vertex x coloured in R and B, a contradiction. The reasoning is similar if we suppose that all the vertices in W_1^{old} are R in the new colouring.

Consequently, in W_1^{old} there exist vertices of both colours R (new) and W (new). Let z be a neighbour of x in B_2^{old} . Then z is a B-vertex in G too and its R- or Wneighbours are all in G_2 . Suppose it has an R-neighbour v in B_2^{old} or in R_2^{old} . For every vertex $w \in W_1^{\text{old}} \cap R_1^{\text{new}}$ we have then the P_4 wxzv which is coloured RBBR, a contradiction. We deduce that every R-neighbour of z was coloured with W in the fine-colouring of G_2 . In the same way, we deduce that every W-neighbour of zwas coloured with W in the fine-colouring of G_2 .

Consider now a neighbour t of x in R_2^{old} . Then t is coloured in B (new) and its neighbours in G_1 cannot be B-vertices (since x is also a B-neighbour of t, the Bvertices adjacent to t in G would not form a clique). Let $u \in R_1^{\text{new}} \cap R_1^{\text{old}}$ (if there is one) be a neighbour of t and $v \in W_2^{\text{old}}$ an R-neighbour of z. Since $tz \in E(G)$ (both are B-vertices in $N_G(x)$), then utzv is a P_4 coloured RBBR, again in contradiction with (C3). We have necessarily $tv \in E(G)$ and this contradicts the fact that the Rand W-vertices adjacent to t form a clique (since v and u are nonadjacent). The same reasoning is valid for the W-neighbours of t in R_1^{old} , so the former R-vertices in G_1 cannot be coloured in R, W or B, a contradiction. \Box

Claim 2. No B-vertex in G_i (old) can change its colour in a fine-colouring of G.

Proof. In the basic graphs F_n , every R- or W-vertex is adjacent to exactly one B-vertex. By Claim 1, no B-vertex can be added to the initial class, so by induction we may suppose that in G_1 and G_2 every R- or W-vertex is also adjacent to exactly one B-vertex. If a B-vertex (old) in G_i (i=1 or 2) changes its colour, every former R-neighbour of it would have no more B-neighbours in G and that would contradict (C1). \Box

Proof of Lemma 5 (*continued*). We show that in every fine-colouring of G, each basic graph has the vertices in an *n*-clique of the same colour. By Claims 1 and 2, the set of *B*-vertices in every basic graph is unchanged while using the 2-join operation. Consequently, one of the three *n*-cliques of each F_n (say $\{z_1, z_2, ..., z_n\}$) is always coloured in B and all its neighbours are vertices of the same F_n (see Fig. 1). By (C1) and (C2) for G, any *B*-vertex z_i has the property that its neighbours x_i, y_i are coloured in R and W. Suppose that for two indices *i* and *j*, x_i and x_j have not the same colour. If x_i is an R-vertex and x_j a W-vertex, then y_j is an R-vertex and $x_i z_i z_j y_j$ is a P_4 coloured *RBBR*, a contradiction. Therefore, all the vertices in an *n*-clique have the

same colour and this is a fine-colouring of F_n . In fact, every fine-colouring of G may be obtained from another fine-colouring by interchanging the colours R and W in some basic graphs. \Box

Remark 4. Lemma 5 insures that the vertices of an *n*-clique in a basic graph F_n are always of the same colour in a fine-colouring of G and have the same neighbour set in $G-F_n$. Also, the neighbourhood in G of each B-vertex is precisely its neighbourhood in the basic graph containing it. Moreover, given two basic subgraphs in a graph $G \in \Gamma$, either the connection between them is R-R, W-W, or it is W-R, R-W.

Lemma 6. No P_4 of type RRRR or WWWW is induced by a fine-coloring of $G \in \Gamma$.

Proof. Suppose the contrary and let xyzt be a P_4 of type RRRR.

If any two vertices of this path are contained in the same basic graph F_n , then they are contained in the same *n*-clique of this graph, so they have precisely the same neighbours in $G - F_n$, according to Remark 4. Consequently, the other two vertices must be in the same basic graph, and this contradicts the fact that in a basic graph the vertices of the same colour form a clique.

Then every two vertices are in different basic graphs. Consider now F^x and F^z the basic graphs containing x and z, respectively. The *R*-vertices in F^z are not adjacent to the *R*-vertices in F^x , therefore, by Remark 4, they must be adjacent to the *W*-vertices in F^x . If x' is a *W*-vertex in F^x , then $x'z \in E(G)$. The same holds for y and t, so if t' is a *W*-vertex in F^t then $yt' \in E(G)$. Moreover, x' and t' are nonadjacent (otherwise the edges between F^x and F^t would be *WW*-edges or *RR*-edges and x would be adjacent to t since they are both *R*-vertices), and the same holds for x' and y, respectively t' and z. We deduce that x'zyt' is a P_4 coloured in *WRRW*, a contradiction. \Box

Theorem 2. The graphs in Γ are perfect graphs.

Proof. Let $G \in \Gamma$ be a graph provided with a fine-colouring. By Lemma 6, the fine-colouring is a partition of V(G) into three sets R, W, B which satisfies:

- no P_4 induced in G is coloured RRRR, WRRW, RWWR or WWWW.

- for each *B*-vertex x in G, $N_G(x) = K_1 \cup K_2$, where K_1 and K_2 are two disjoint cliques with no edge between them.

If G was not perfect, then it would contain a minimal imperfect subgraph G'. Two cases occur:

Case 1: G' contains no B-vertex. If all the vertices of G' have the same colour, then G' is P_4 -free, thus it is perfect. If both colours R and W are present in G', we have a partition of V(G') into two sets R, W such that no induced P_4 is coloured RRRR, RWWR, WRRW or WWWW. By a theorem of Chvátal et al. [2], the graph G' is perfect if and only if the subgraphs induced by the vertices coloured in R and, respectively, in W are perfect. In our case, the two subgraphs are P_4 -free, thus they are perfect. Conclusion: G' is also perfect. *Case* 2: G' contains at least one *B*-vertex. Since G' is minimal imperfect, the *B*-vertex is contained in exactly ω ω -cliques, where $\omega = \omega(G')$ is the clique number of G'. But every *B*-vertex is contained in exactly two cliques in G, so it is contained in at most two cliques in G', therefore $\omega = 1$ or 2. In the first case, G' is perfect; in the second one, it is an odd hole and this contradicts the fact that G is a Berge graph.

In both cases we obtain that every subgraph of G is perfect, so G itself is a perfect graph. \Box

4. Counterexamples

For an arbitrary graph H, consider the following conjecture:

Conjecture (C_H) . If G is a Berge graph such that

- 1. neither G, nor \overline{G} has an even pair;
- 2. neither G, nor \overline{G} has a star cutset;
- 3. neither G, nor \overline{G} has a stable cutset

then G or \overline{G} is H-free.

For all graphs H such that there is a graph G_H in Γ containing both H and \overline{H} , the corresponding conjecture (C_H) is false. We don't know exactly which are these graphs, but we indicate a constructive method to found counterexamples for the cases when H is a join of a clique and a stable set.

If $G = G_1 \phi G_2$, let $(G_1 \phi G_2)_{(R \leftrightarrow W)}$ or $G_{(R \leftrightarrow W)}$ be the graph G with the colouring obtained from the initial one by interchanging the colours R and W in G_2 . Consequently, the edges between G_1 and G_2 are no more RR-edges or WW-edges, but RWor WR-edges. Obviously, the new colouring is a fine-colouring for the graphs G_1 and G_2 . Is it a fine-colouring for the graph G? A reasoning similar to the one in Lemma 1 gives the affirmative answer for the conditions (C3), (C4) (notice that in $G_{(R \leftrightarrow W)}$ the adjacency is R-W, W-R, as in \overline{G} before switching the colours R and W). The other conditions are insured by the internal properties of G_i (i=1, 2).

For a fixed $n \ge 2$, consider now the sequence of graphs defined by induction as follows:

 $G_n^1 = F_n$ $G_n^2 = F_n \phi F_n$ \dots $G_n^k = G_{n,(R \leftrightarrow W)}^{k-1} \phi F_n$ \dots

The graph G_n^k is obtained from G_n^{k-1} by interchanging the colours R and W in *its* subgraph F_n (i.e. the subgraph used to build G_n^{k-1}) and 2-joining with a new graph F_n ; therefore $G_n^k \in \Gamma$.

The iterative construction above allows us to identify a certain structure of these graphs. More precisely, if we denote by F_n^i the graph F_n used during the step *i* of the composition (that is, the graph composed with $G_{n,(R \leftrightarrow W)}^{i-1}$ to obtain G_n^i) and by R_i , W_i , B_i the *n*-cliques of F_n^i coloured in *R*, *W*, *B* at the present moment of the composition, then the graph G_n^k ($k \ge 2$) has the structure in Fig. 2 (The *n*-cliques are represented by points, the join of two subgraphs by continuous line and the already known connections in F_n^i by dashed line.)

Indeed, the graph G_n^2 has precisely this structure. By induction, suppose that G_n^k has the configuration in Fig. 2. Its successor G_n^{k+1} is obtained from G_n^k by interchanging the colours R, W in F_n^k (that is, in Fig. 2, W_k becomes R_k and vice versa) and by composing with F_n^{k+1} according to the new colouring. The structure obtained in this way is the one in Fig. 2.

In G_n^{k+1} every vertex in the clique W_{k+1} is adjacent to every vertex in the cliques $W_1, W_2, ..., W_k$, for all $k \ge 1$. We deduce that G_n^{n+1} contains as an induced subgraph the graph H_n which is the join of an *n*-clique and an *n*-stable. Also, G_n^{n+1} contains \bar{H}_n , since the *n*-clique W_{n+1} is nonadjacent to any of the cliques $R_1, R_2, ..., R_n$. So neither G_n^{n+1} , nor \bar{G}_n^{n+1} is H_n -free.

We can then state that the conjecture (C_H) is false for every H which is the join of a clique K and a stable set S. A counterexample is G_p^{p+1} , where $p = max\{|K|, |S|\}$.

Remark 5. Another conjecture says that a minimal imperfect graph cannot contain an *odd pair*, i.e., a pair of nonadjacent vertices joined by no even chordless path. Similarly to Lemma 2 and Remark 2, one can show that if the only odd pairs in G_i , \bar{G}_i (i = 1,



Fig. 2.



2) correspond to the *RR*-, *BB*-, *WW*- edges in G_i , then the 2-join G has the property that neither G, nor \overline{G} has an odd pair. We deduce that all G_n^k (except for $G_2^1 = F_2$) have this property and that even if we add to (C_H) the following hypothesis:

4. Neither G, nor \overline{G} has an odd pair,

the conjecture is false.

For n = 2, the graph $H_n = H_2$ is precisely the diamond and the counterexample G_2^3 is the 2-join of the two finely-coloured Berge graphs in Fig. 3.

Chvátal noticed that the graph G_2^3 can be described in the following easier manner: in the induced cycle on six vertices w_1, w_2, \ldots, w_6 , substitute the vertex w_i with the 2-clique induced by the vertices u_i, v_i ; then, for any i = 1, 2, 3, add the vertices x_i, y_i and the edges $x_i y_i, x_i u_i, x_i u_{i+3}, u_i u_{i+3}, y_i v_i, y_i v_{i+3}, v_i v_{i+3}$.

5. Conclusions

It is quite hard to believe that a conjecture of type (C_H) may be true, although the counterexamples we have built do not exhaust the subject. The existence of such a large class of "nonconventional" perfect graphs contributes to confirm the idea that we really need some new properties of minimal imperfect graphs before being able to solve the difficult perfect graph problems.

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