



# Minimally 3-restricted edge connected graphs<sup>☆</sup>

Qinghai Liu, Yanmei Hong, Zhao Zhang<sup>\*</sup>

College of Mathematics and System Sciences, Xinjiang University, Urumqi, Xinjiang 830046, PR China

## ARTICLE INFO

### Article history:

Received 12 December 2007

Received in revised form 13 June 2008

Accepted 31 July 2008

Available online 27 August 2008

### Keywords:

Fault tolerance

Restricted edge connectivity

## ABSTRACT

For a connected graph  $G = (V, E)$ , an edge set  $S \subset E$  is a 3-restricted edge cut if  $G - S$  is disconnected and every component of  $G - S$  has order at least three. The cardinality of a minimum 3-restricted edge cut of  $G$  is the 3-restricted edge connectivity of  $G$ , denoted by  $\lambda_3(G)$ . A graph  $G$  is called minimally 3-restricted edge connected if  $\lambda_3(G - e) < \lambda_3(G)$  for each edge  $e \in E$ . A graph  $G$  is  $\lambda_3$ -optimal if  $\lambda_3(G) = \xi_3(G)$ , where  $\xi_3(G) = \max\{\omega(U) : U \subset V(G), G[U] \text{ is connected}, |U| = 3\}$ ,  $\omega(U)$  is the number of edges between  $U$  and  $V \setminus U$ , and  $G[U]$  is the subgraph of  $G$  induced by vertex set  $U$ . We show in this paper that a minimally 3-restricted edge connected graph is always  $\lambda_3$ -optimal except the 3-cube.

© 2008 Elsevier B.V. All rights reserved.

## 1. Introduction

A network can be conveniently modeled as a graph  $G = (V, E)$ . A classic measure of the fault tolerance of a network is the edge connectivity  $\lambda(G)$ . In general, the larger  $\lambda(G)$  is, the more reliable the network is [2]. A more refined measure known as restricted edge connectivity was proposed by Esfahanian and Hakimi [5], which was further generalized to  $k$ -restricted edge connectivity by Fábrega and Fiol [6] (called  $k$ -extra edge connectivity in their paper).

Let  $G$  be a connected graph. An edge set  $S \subset E(G)$  is said to be a  $k$ -restricted edge cut of  $G$  if  $G - S$  is disconnected and each component of  $G - S$  has at least  $k$  vertices. The minimum cardinality of a  $k$ -restricted edge cut is called the  $k$ -restricted edge connectivity of  $G$ , denoting by  $\lambda_k(G)$ . A  $k$ -restricted edge cut  $S$  with  $|S| = \lambda_k(G)$  is called a  $\lambda_k$ -cut. Not all graphs have  $\lambda_k$ -cuts [4,5,12,21]. Those which do have  $\lambda_k$ -cuts are called  $\lambda_k$ -connected graphs. According to current studies on  $k$ -restricted edge connectivity [10,11,13,17], it seems that the larger  $\lambda_k(G)$  is, the more reliable the network is. In [21], Zhang and Yuan proved that  $\lambda_k(G) \leq \xi_k(G)$  holds for any integer  $k \leq \delta(G) + 1$  except for a class of graphs (such a graph is constructed from a set of complete subgraphs  $K_\delta$  by adding a new vertex  $u$  and connect  $u$  to every other vertex), where  $\delta(G)$  is the minimum degree of  $G$  and  $\xi_k(G) = \min\{\omega(U) : U \subset V(G), G[U] \text{ is connected}, |U| = k\}$ ,  $\omega(U)$  is the number of edges between  $U$  and  $V \setminus U$ , and  $G[U]$  is the subgraph of  $G$  induced by  $U$ . A graph  $G$  is called  $\lambda_k$ -optimal if  $\lambda_k(G) = \xi_k(G)$ . There is much research on sufficient conditions for a graph to be  $\lambda_k$ -optimal, such as symmetric conditions [11,13,17,18], degree conditions [14,15,19], and girth-diameter conditions [1,6,16,20]. For more information on this topic, we refer the readers to the nice survey paper by Hellwig and Volkmann [7].

In this paper, we give another type of sufficient condition called a minimally restricted edge connected condition. A graph  $G$  is a *minimally  $k$ -restricted edge connected graph* (minimally  $\lambda_k$ -graph for short) if  $\lambda_k(G - e) < \lambda_k(G)$  (and thus  $\lambda_k(G - e) = \lambda_k(G) - 1$ ) for each edge  $e \in E(G)$ . It is implied in the definition that  $\lambda_k(G - e) > \lambda_k(G)$  exists for each edge  $e$ . If  $e$  is a pending edge, then  $G - e$  does not have  $\lambda_k$ -cut for  $k \geq 2$ . So, we always assume  $\delta(G) \geq 2$  when  $G$  is a minimally  $\lambda_k$ -graph for some  $k \geq 2$ . A minimally  $\lambda_1$ -graph is exactly a minimally edge connected graph, which has been shown to be  $\lambda$ -optimal ([9] Exercise 49). In [8], the authors have proved that every minimally  $\lambda_2$ -graph is  $\lambda_2$ -optimal. In this paper, we show that every minimally  $\lambda_3$ -graph is always  $\lambda_3$ -optimal except the 3-cube.

<sup>☆</sup> This research is supported by NSFC (60603003), the Key Project of Chinese Ministry of Education (208161) and XJEDU.

<sup>\*</sup> Corresponding author. Tel.: +86 13899960204; fax: +86 991 8585505.

E-mail address: zhzhao@xju.edu.cn (Z. Zhang).

## 2. Preliminaries and terminologies

Let  $G = (V, E)$  be a graph. For two disjoint vertex sets  $U_1, U_2 \subset V(G)$ , denote by  $[U_1, U_2]_G$  the set of edges of  $G$  with one end in  $U_1$  and the other end in  $U_2$ ,  $G[U]$  is the subgraph of  $G$  induced by vertex set  $U \subset V(G)$ ,  $\bar{U} = V(G) \setminus U$  is the complement of  $U$ ,  $\omega_G(U) = |[U, \bar{U}]_G|$  is the number of edges between  $U$  and  $\bar{U}$ . When the graph under consideration is obvious, we omit the subscript  $G$ . Write  $d_A(U) = |[U, A \setminus U]|$ ,  $d_A(u) = d_A(\{u\})$ . Sometimes, we use a graph itself to represent its vertex set. For example,  $\omega(C)$  is used instead of  $\omega(V(C))$ , where  $C$  is a subgraph of  $G$ ; for an edge  $e = uv$ ,  $d_A(e)$  is used instead of  $d_A(\{u, v\})$ , etc.

A  $\lambda_3$ -fragment is a subset  $U$  of  $V(G)$  with  $[U, \bar{U}]$  being a  $\lambda_3$ -cut. If  $U$  is a  $\lambda_3$ -fragment, then so is  $\bar{U}$ , and both  $G[U]$  and  $G[\bar{U}]$  are connected. A  $\lambda_3$ -fragment with minimum order is called a  $\lambda_3$ -atom. The order of a  $\lambda_3$ -atom is denoted by  $\alpha_3(G)$ . Clearly,  $\alpha_3(G) \leq |V(G)|/2$ .

A graph  $H$  is  $\lambda_3$ -independent if each component of  $H$  has at most two vertices. A connected graph of order at most two is  $\lambda_3$ -trivial. A graph is called  $\lambda_3$ -non-trivial if it has a component which contains at least three vertices.

The following two observations will be used frequently without mentioning them explicitly. The first is that if two connected subgraphs  $G_1$  and  $G_2$  have nonempty intersection, then  $G_1 \cup G_2$  is also connected. The second is that for a vertex set  $F$  of a connected graph  $G$  and a component  $C$  of  $G - F$ , if  $G[F]$  is connected, then so is  $G - C$ .

For terminologies not given here, we refer to [3] for reference.

## 3. Main result

First, it should be noted that if  $\alpha_3(G) = 3$ , then  $G$  is  $\lambda_3$ -optimal. In fact, Bonsma et al. [4] have proved that  $\lambda_3(G) \leq \xi_3(G)$  holds for any  $\lambda_3$ -connected graph  $G$ . On the other hand, considering a  $\lambda_3$ -atom  $A$  of  $G$ , we have  $\lambda_3(G) = \omega(A) \geq \xi_3(G)$ . So,  $\lambda_3(G) = \xi_3(G)$ . In view of this observation, to derive our main theorem, it suffices to show that  $\alpha_3(G) = 3$ . In the following, we prove that if  $G$  is a minimally 3-restricted edge connected graph with  $\alpha_3(G) \geq 4$ , then  $G$  is isomorphic to 3-cube.

**Lemma 1.** *Let  $G$  be a  $\lambda_3$ -connected graph with  $\delta(G) \geq 2$ ,  $F$  be a subset of  $G$  with  $G[\bar{F}]$  being connected. If one of the following conditions is satisfied:*

- (a)  $\omega(F) < \lambda_3(G)$ , or
- (b)  $\omega(F) = \lambda_3(G)$  and  $|F| < \alpha_3(G)$ ,

then  $G[F]$  is  $\lambda_3$ -independent.

**Proof.** Suppose  $F$  has a  $\lambda_3$ -non-trivial component  $C$ . Since  $G[\bar{C}]$  is connected, we have  $\omega(F) \geq \omega(C) \geq \lambda_3(G)$ , which is clearly a contradiction to condition (a). If condition (b) occurs, then  $\omega(F) = \omega(C) = \lambda_3(G)$ . Hence  $V(C)$  is a  $\lambda_3$ -fragment. But then  $|F| \geq |C| \geq \alpha_3(G)$ , contradicting  $|F| < \alpha_3(G)$ .  $\square$

**Lemma 2.** *Let  $G$  be a  $\lambda_3$ -connected graph with  $\delta(G) \geq 2$  and  $\alpha_3(G) \geq 4$ ,  $A$  be a  $\lambda_3$ -atom of  $G$ , and  $B$  be a  $\lambda_3$ -fragment of  $G$ .*

- (a) *If a subset  $U$  of  $A$  is such that  $G[U]$  is connected and  $G[A \setminus U]$  has a  $\lambda_3$ -non-trivial component, then  $d_A(U) > d_{\bar{A}}(U)$ .*
- (b) *If a subset  $U$  of  $B$  is such that  $G[U]$  is connected and  $G[B \setminus U]$  has a  $\lambda_3$ -non-trivial component, then  $d_B(U) \geq d_{\bar{B}}(U)$ .*
- (c)  $\delta(G[A]) \geq 2$ .

**Proof.** (a). Suppose  $d_A(U) \leq d_{\bar{A}}(U)$ . Then  $\omega(A \setminus U) = \omega(A) + d_A(U) - d_{\bar{A}}(U) \leq \omega(A) = \lambda_3(G)$ . By noting that  $|A \setminus U| < |A| = \alpha_3(G)$ , it follows from Lemma 1 that  $G[A \setminus U]$  is  $\lambda_3$ -independent, a contradiction.

(b). The proof of (b) is similar to that of (a); note that under the assumption  $d_B(U) < d_{\bar{B}}(U)$ , it can be deduced that  $\omega(B \setminus U) < \lambda_3(G)$ .

(c) is a consequence of (a). In fact, for each vertex  $x \in A$ , if  $A - x$  is  $\lambda_3$ -independent, then  $d_A(x) \geq 2$  since  $|A| \geq 4$ . If  $A - x$  contains a  $\lambda_3$ -non-trivial component, taking  $U = \{x\}$  in (a), we have  $d_A(x) > d_{\bar{A}}(x)$ . Then it follows from  $d_A(x) > \frac{1}{2} \cdot (d_A(x) + d_{\bar{A}}(x)) = \frac{1}{2} \cdot d_G(x) \geq 1$  that  $d_A(x) \geq 2$ .  $\square$

Similar to Lemma 2, we can prove the following lemma. The key observation to the proof, as well as some proofs after it, is that for any edge  $e \in E(G)$ , if  $\lambda_3(G - e) < \lambda_3(G)$ , then any  $\lambda_3$ -fragment of  $G - e$  contains exactly one end of  $e$ , and is a  $\lambda_3$ -fragment of  $G$ . Note that the observation is true when  $G$  is minimally 3-restricted edge connected.

**Lemma 3.** *Let  $G$  be a  $\lambda_3$ -connected graph with  $\delta(G) \geq 2$  and  $\alpha_3(G) \geq 4$ ,  $e = uv$  be an edge of  $G$ ,  $\lambda_3(G - e) < \lambda_3(G)$ ,  $A$  be a  $\lambda_3$ -atom of  $G - e$  with  $u \in A$  and  $v \notin A$ .*

- (a) *If a subset  $U$  of  $A$  is such that  $G[U]$  is connected,  $G[A] - U$  has a  $\lambda_3$ -non-trivial component, and  $e$  is not incident with  $U$ , then  $d_A(U) > d_{\bar{A}}(U)$ .*
- (b)  $d_{G[A]}(x) \geq 2$  for each  $x(\neq u) \in A$ .

**Lemma 4.** *Let  $G$  be a  $\lambda_3$ -connected graph with  $\delta(G) \geq 2$  and  $\alpha_3(G) \geq 4$ ,  $A$  be a  $\lambda_3$ -atom of  $G$ ,  $B$  be a  $\lambda_3$ -fragment of  $G$ ,  $A \cap B \neq \emptyset$ . Then for any component  $C$  of  $G[A \cap B]$ , either  $G[A] - C$  or  $G[B] - C$  is  $\lambda_3$ -independent.*

**Proof.** Suppose both  $G[A] - C$  and  $G[B] - C$  have  $\lambda_3$ -non-trivial components. Then by taking  $U = V(C)$  in Lemma 2, we have

$$d_A(U) > d_{\bar{A}}(U) \geq d_{B \setminus A}(U) = d_B(U)$$

and

$$d_B(U) \geq d_{\bar{B}}(U) \geq d_{A \setminus B}(U) = d_A(U),$$

a contradiction.  $\square$

Similar to Lemma 4, by using Lemma 3 instead of Lemma 2, it can be proved that

**Lemma 5.** Let  $G$  be a  $\lambda_3$ -connected graph with  $\delta(G) \geq 2$  and  $\alpha_3(G) \geq 4$ ,  $e = uv$  be an edge of  $G$ ,  $\lambda_3(G - e) < \lambda_e(G)$ ,  $A$  be a  $\lambda_3$ -atom of  $G - e$  with  $u \in A$  and  $v \notin A$ ,  $B$  be a  $\lambda_3$ -fragment of  $G$  such that  $A \cap B \neq \emptyset$  and  $e$  is not incident with  $B$ . For any component  $C$  of  $G[A \cap B]$ , either  $G[A] - C$  or  $G[B] - C$  is  $\lambda_3$ -independent.

**Lemma 6.** Let  $G$  be a minimally  $\lambda_3$ -graph with  $\alpha_3(G) \geq 4$ ,  $A$  be a  $\lambda_3$ -atom of  $G$ ,  $e = uv$  be an edge of  $G[A]$ ,  $B$  be a  $\lambda_3$ -atom of  $G - e$ . Then,

- (a)  $G[A \cap B]$ ,  $G[A \setminus B]$  and  $G[B \setminus A]$  are all connected;
- (b)  $|A \cap B| = 2$ ,  $|A \setminus B| = 2$ ;
- (c)  $G[A]$  is  $C_4$  (4-cycle) or  $K_4$  (complete graph on 4 vertices) or  $K_4^-$  ( $K_4$  minus one edge);
- (d)  $G[\overline{A \cup B}]$  is connected;
- (e)  $|[A \cap B, \overline{A \cup B}]| = |[A \setminus B, B \setminus A]| = 0$ , and  $d_{\bar{A}}(x) = d_A(x) - 1$  for each  $x \in A$ .

**Proof.** By the observation before Lemma 3, we may suppose, without loss of generality, that  $u \in B$  and  $v \in \bar{B}$ . Then  $A \cap B \neq \emptyset$  and  $A \setminus B \neq \emptyset$ . Because of  $|A| \leq \frac{1}{2}|V(G)|$  and  $|B| \leq \frac{1}{2}|V(G)|$ , we have  $|\overline{A \cup B}| = |V(G)| - |A| - |B| + |A \cap B| \geq |A \cap B|$ .

(a). Suppose  $G[A \cap B]$  has two components  $C_1, C_2$ . If one of  $C_1$  and  $C_2$ , say  $C_2$ , has at least two vertices, then by the connectedness of  $G[A]$  and  $G[B]$ , both  $G[A] - C_1$  and  $G[B] - C_1$  have  $\lambda_3$ -non-trivial components containing  $C_2$ , contradicting Lemma 4. Next, suppose  $V(C_1) = \{x\}$ ,  $V(C_2) = \{y\}$ , and  $y \neq u$ . By Lemma 2(c), we have  $d_A(y) \geq 2$ . By Lemma 3(b), we have  $d_B(y) \geq 2$ . So both  $G[A] - C_1$  and  $G[B] - C_1$  have  $\lambda_3$ -non-trivial components, again a contradiction. So,  $G[A \cap B]$  is connected.

Suppose  $G[A \setminus B]$  is not connected. Then a contradiction can be obtained by using Lemma 4 to  $A$  and  $\bar{B}$  (note that  $\bar{B}$  is also a  $\lambda_3$ -fragment). In fact, the case that  $G[A \setminus B]$  has a component with at least two vertices can be analyzed similar to the above paragraph. In the case that  $A \setminus B$  is an independent set with at least two vertices, we have  $|A \cap B| \geq 2$  since  $\delta(G[A]) \geq 2$ . By the connectedness of  $G[A \cap B]$ , we see that for any vertex  $y \in A \setminus B$ ,  $G[A] - y$  is  $\lambda_3$ -non-trivial. If there is a vertex  $y \in A \setminus B$  with degree one in  $G[\bar{B}]$ , then  $G[\bar{B}] - y$  is  $\lambda_3$ -non-trivial. If there exists a vertex  $x \in A \setminus B$  with  $d_{G[\bar{B}]}(x) \geq 2$ , let  $y$  be another vertex in  $A \setminus B$ . Then  $G[\bar{B}] - y$  has a  $\lambda_3$ -non-trivial component containing  $x$ . Taking  $C = \{y\}$  in Lemma 4 and using  $\bar{B}$  to take the place of  $B$ , we arrive at a contradiction.

Similar to the above deduction, by using Lemmas 3 and 5 instead of Lemmas 2 and 4, it can be proved that  $G[B \setminus A]$  is also connected.

(b). First suppose  $|A \cap B| = 1$ . By  $|B| \geq |A| \geq 4$ , we have  $|B \setminus A| \geq |A \setminus B| \geq 3$ . By (a),  $G[B \setminus A]$  and  $G[A \setminus B]$  are both  $\lambda_3$ -non-trivial, which contradicts Lemma 4 (taking  $U = A \cap B$  there).

Next suppose  $|A \cap B| \geq 3$ . By (a),  $G[A \cap B]$  is  $\lambda_3$ -non-trivial. By the connectedness of  $G[\bar{A}]$  and  $G[\bar{B}]$ , we see that  $G[\overline{A \cap B}] = G[\overline{A \cup B}]$  is also connected. Taking  $F = A \cap B$  in Lemma 1, by noting that  $|F| = |A \cap B| < |A| = \alpha_3(G)$ , we have  $\omega(A \cap B) > \lambda_3(G)$ . Combining this with  $\omega(A) = \omega(B) = \lambda_3(G)$  and the following well known submodular inequality,

$$\omega(A \cap B) + \omega(A \cup B) \leq \omega(A) + \omega(B),$$

we have  $\omega(\overline{A \cup B}) = \omega(A \cup B) < \lambda_3(G)$ . Taking  $F = \overline{A \cup B}$  in Lemma 1, we see that  $G[\overline{A \cup B}]$  is  $\lambda_3$ -independent, that is,  $G[\overline{A \cup B}]$  is composed of some singletons and complete graphs of order two.

We claim that for each component  $C$  of  $G[\overline{A \cup B}]$ ,  $d_A(C) = d_B(C)$  and  $d_{A \cap B}(C) = 0$ . It should be noted that  $C$  is connected to both  $A \setminus B$  and  $B \setminus A$ , since  $G[\bar{A}]$  and  $G[\bar{B}]$  are connected. First suppose  $C$  is a singleton  $x$ . Since  $|\overline{A \cup B}| \geq |A \cap B| \geq 3$  and  $G[B \setminus A]$  is connected, we see that  $G[\bar{A}] - x$  has a connected subgraph of order at least 3. Applying Lemma 2(b) to the  $\lambda_3$ -fragment  $\bar{A}$  and the vertex set  $\{x\}$ , we have  $d_{\bar{A}}(x) \geq d_A(x)$ . Hence

$$d_A(x) \leq d_{\bar{A}}(x) = d_{B \setminus A}(x) = d_B(x) - d_{A \cap B}(x).$$

Similarly,

$$d_B(x) \leq d_{\bar{B}}(x) = d_{A \setminus B}(x) = d_A(x) - d_{B \cap A}(x).$$

It follows that  $d_A(x) = d_B(x)$  and  $d_{A \cap B}(x) = 0$ . Next suppose  $C$  is an edge  $e = v_1 v_2$ . If  $|\overline{A \cup B}| = 3$  and  $|A \setminus B| = 1$ , then  $G[\overline{A \cup B}]$  is composed of  $e$  and a singleton  $y$ . Set  $A' = (A \setminus B) \cup \{v_1, v_2\}$ . Then both  $G[A']$  and  $G[\bar{A}']$  are connected subgraphs of order at least 3. Since  $d_{A \setminus B}(y) = d_{B \setminus A}(y)$ , we have  $\omega(A') = \omega(\bar{A}') = \omega(B \cup \{y\}) = \omega(B) + d_{A \setminus B}(y) - d_{B \setminus A}(y) = \omega(B) = \lambda_3(G)$ .

So  $A'$  is a  $\lambda_3$ -fragment. But then  $\alpha_3(G) \leq |A'| = 3$ , contradicting the assumption that  $\alpha_3(G) \geq 4$ . So,  $|\overline{A \cup B}| \geq 4$  or  $|A \setminus B| \geq 2$  (hence  $|B \setminus A| \geq 2$ ). Taking the place of  $x$  by  $e$  in the proof for  $C$  being a singleton, we have  $d_A(e) = d_B(e)$  and  $d_{A \cap B}(e) = 0$ . The claim is proved.

As a consequence of the claim, for any vertex set  $U$  which is the union of the vertex sets of some components in  $G[\overline{A \cup B}]$ , we have  $d_{A \setminus B}(U) = d_{B \setminus A}(U)$  and  $d_{A \cap B}(U) = 0$ .

Let  $U$  be the vertex set of one or two components of  $G[\overline{A \cup B}]$  such that  $|U| = 2$ . Set  $A' = (A \setminus B) \cup U$ . Then  $3 \leq |A'| < |A|$  and both  $G[A']$  and  $G[\overline{A'}]$  are connected. Since  $\omega(A') = \omega(\overline{A'}) = \omega(B) + d_{A \setminus B}(\overline{A \cup B} \setminus U) - d_{B \setminus A}(\overline{A \cup B} \setminus U) = \omega(B) = \lambda_3(G)$ , we see that  $A'$  is a  $\lambda_3$ -fragment of  $G$  with less vertices than  $A$ , a contradiction.

Hence  $|A \cap B| = 2$ . It follows that  $|A \setminus B| = |A| - |A \cap B| \geq 2$ . If  $|A \setminus B| \geq 3$ , then  $|B \setminus A| \geq 3$ . By (a), both  $G[A \setminus B]$  and  $G[B \setminus A]$  are  $\lambda_3$ -non-trivial, contradicting Lemma 4. So,  $|A \setminus B| = 2$ , and (b) is proved.

(c) follows from (a), (b) and Lemma 2(c).

Suppose  $A \setminus B = \{a, d\}$  and  $A \cap B = \{b, c\}$ . Then  $ad, bc \in E(G)$ .

**Claim 1.** For each edge  $e = xy \in E(G[A])$ ,  $d_A(e) \geq d_{\overline{A}}(e)$ , equality holds if and only if  $d_{\overline{A}}(x) = d_A(x) - 1$  and  $d_{\overline{A}}(y) = d_A(y) - 1$ .

Taking  $U = \{x\}$  in Lemma 2(a), we have  $d_A(x) > d_{\overline{A}}(x)$ . Hence  $d_{\overline{A}}(x) \leq d_A(x) - 1 = d_{A \setminus \{x,y\}}(x)$ . Similarly,  $d_{\overline{A}}(y) \leq d_A(y) - 1 = d_{A \setminus \{x,y\}}(y)$ . Then the claim follows from

$$d_A(xy) = d_{A \setminus \{x,y\}}(x) + d_{A \setminus \{x,y\}}(y) \geq d_{\overline{A}}(x) + d_{\overline{A}}(y) = d_{\overline{A}}(xy).$$

**Claim 2.**  $d_{A \setminus B}(bc) = d_{B \setminus A}(bc)$ ,  $d_{\overline{A \cup B}}(bc) = 0$ ,  $d_{\overline{A}}(b) = d_A(b) - 1$  and  $d_{\overline{A}}(c) = d_A(c) - 1$ .

By Claim 1, we have

$$d_{A \setminus B}(bc) = d_A(bc) \geq d_{\overline{A}}(bc) = d_{B \setminus A}(bc) + d_{\overline{A \cup B}}(bc). \tag{1}$$

If  $|B \setminus A| = 2$ , then  $B$  is also a  $\lambda_3$ -atom of  $G$ , and similar to (1), we have

$$d_{B \setminus A}(bc) = d_B(bc) \geq d_{\overline{B}}(bc) = d_{A \setminus B}(bc) + d_{\overline{A \cup B}}(bc). \tag{2}$$

If  $|B \setminus A| \geq 3$ , then by (a),  $G[B \setminus A]$  is  $\lambda_3$ -non-trivial. By Lemma 2(b), inequality (2) is still valid. Combining (1) and (2), we have the first and the second equality of the claim. As a consequence, equality holds for (1), and thus the third and the fourth equality of the claim follow from Claim 1.

(d). Suppose  $C_1$  and  $C_2$  are two components of  $G[\overline{A \cup B}]$ . Since  $d_A(ad) \geq d_{\overline{A}}(ad)$  by Claim 1, we have

$$\begin{aligned} \omega(C_1) < \omega(\overline{A \cup B}) &= \omega(A \cup B) = \omega(B) + d_{\overline{A \cup B}}(ad) - d_{A \cup B}(ad) \\ &= \omega(B) + (d_{\overline{A}}(ad) - d_{B \setminus A}(ad)) - (d_{B \setminus A}(ad) + d_A(ad)) \\ &= \omega(B) - d_A(ad) + d_{\overline{A}}(ad) - 2d_{B \setminus A}(ad) \\ &\leq \omega(B) = \lambda_3(G). \end{aligned}$$

By Lemma 1,  $C_1$  is  $\lambda_3$ -independent, and thus  $|V(C_1)| \leq 2$ . Since  $G[B \setminus A]$  is connected,  $|B \setminus A| \geq 2$ , and  $C_2$  is connected to  $B \setminus A$ , we see that  $G[\overline{A}] - C_1$  has a  $\lambda_3$ -non-trivial component. By Lemma 2(b) and Claim 2, we have

$$d_A(C_1) \leq d_{\overline{A}}(C_1) = d_{B \setminus A}(C_1) = d_B(C_1).$$

Similarly,

$$d_B(C_1) \leq d_{\overline{B}}(C_1) = d_{A \setminus B}(C_1) = d_A(C_1).$$

It follows that  $d_B(C_1) = d_A(C_1)$ . Set  $A' = (A \setminus B) \cup C_1$ . Then both  $G[A']$  and  $G[\overline{A'}]$  are connected. By Claim 2, we have  $\omega(A') = \omega(A \setminus B) + d_B(C_1) - d_{A \setminus B}(C_1) = \omega(A \setminus B) + d_B(C_1) - d_A(C_1) = \omega(A \setminus B) = \omega(A) + d_{A \setminus B}(bc) - d_{B \setminus A}(bc) = \omega(A) = \lambda_3(G)$ . Hence  $A'$  is a  $\lambda_3$ -fragment of  $G$ . since  $|A'| \leq |A|$ , we see that  $A'$  is also a  $\lambda_3$ -atom of  $G$ . Applying Claim 1 to  $A'$  and  $A$  respectively, we have

$$d_{A'}(ad) \geq d_{\overline{A'}}(ad) \geq d_A(ad) + d_{C_2}(ad) > d_A(ad),$$

and

$$d_A(ad) \geq d_{\overline{A}}(ad) \geq d_{C_1}(ad) = d_{A'}(ad),$$

a contradiction. So,  $G[\overline{A \cup B}]$  is connected.

(e). Noting that in proving Claim 2, it suffices for  $B$  to be a  $\lambda_3$ -fragment (it needs not to be a  $\lambda_3$ -atom of  $G - e$ ). Hence, by taking the place of  $B$  by  $\overline{B}$ , we have  $d_{B \setminus A}(ad) = 0$ , and  $d_{\overline{A}}(a) = d_A(a) - 1$ ,  $d_{\overline{A}}(d) = d_A(d) - 1$ . Together with Claim 2, the results follow immediately.  $\square$

**Theorem 1.** The only minimally  $\lambda_3$ -connected graph which is not  $\lambda_3$ -optimal is the 3-cube.

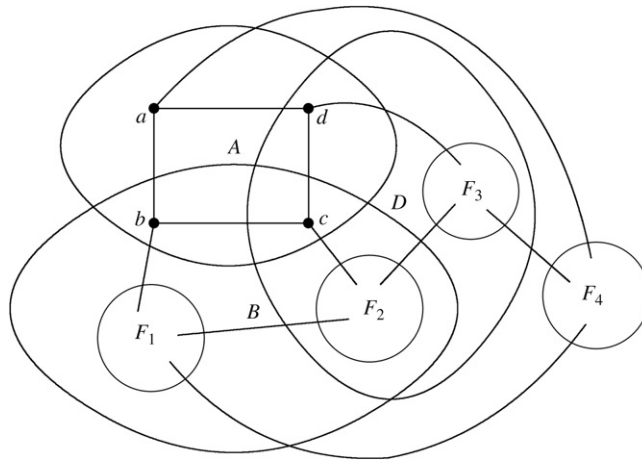


Fig. 1. An illustration for the proof of Theorem 1.

**Proof.** Suppose  $G$  is a minimally  $\lambda_3$ -connected graph which is not  $\lambda_3$ -optimal. Then  $\alpha_3(G) \geq 4$ . Let  $A$  be a  $\lambda_3$ -atom of  $G$ ,  $e = uv$  be an edge of  $G[A]$ , and  $B$  be a  $\lambda_3$ -atom of  $G - e$ . Then Lemma 6 is applicable to  $A$  and  $B$ . Suppose  $A \cap B = \{b, c\}$ ,  $A \setminus B = \{a, d\}$ . Then  $ad, bc \in E(G)$ . By Lemma 6(c), we may assume, without loss of generality, that  $ab, cd \in E(G)$ .

By Lemma 6(e),  $d_{\bar{A}}(x) = d_A(x) - 1$  for each  $x \in \{a, b, c, d\}$ . Hence

$$\begin{aligned} \lambda_3(G) = \omega(A) &= d_{\bar{A}}(a) + d_{\bar{A}}(b) + d_{\bar{A}}(c) + d_{\bar{A}}(d) \\ &= d_A(a) + d_A(d) - 2 + d_A(b) + d_A(c) - 2 \\ &= d_A(ad) + d_A(bc) \\ &= 2|[ad, bc]|. \end{aligned}$$

Let  $D$  be a  $\lambda_3$ -atom of  $G - bc$ . Then Lemma 6 is also applicable to  $A$  and  $D$ . Suppose, without loss of generality, that  $A \cap D = \{c, d\}$ . Set  $F_1 = (B \setminus A) \setminus D$ ,  $F_2 = (B \setminus A) \cap D$ ,  $F_3 = (A \cup B) \cap D$ ,  $F_4 = (\bar{A} \cup B) \setminus D$  (see Fig. 1).

By Lemma 6(e), we have

$$|[a, d], F_1 \cup F_2| = |[b, c], F_3 \cup F_4| = |[a, b], F_2 \cup F_3| = |[c, d], F_1 \cup F_4| = 0. \tag{3}$$

It follows that  $|[a, F_1 \cup F_2 \cup F_3]| = 0$ , and thus  $|[a, F_4]| = |[a, \bar{A}]| = d_{\bar{A}}(a) = d_A(a) - 1 \geq \delta(G[A]) - 1 \geq 1$ . By the same argument, we have

$$\begin{aligned} |[a, F_4]| &= d_{\bar{A}}(a) \geq 1, & |[b, F_1]| &= d_{\bar{A}}(b) \geq 1, \\ |[d, F_3]| &= d_{\bar{A}}(d) \geq 1, & |[c, F_2]| &= d_{\bar{A}}(c) \geq 1. \end{aligned} \tag{4}$$

As a consequence,

$$[A, \bar{A}] = [a, F_4] \cup [b, F_1] \cup [c, F_2] \cup [d, F_3], \tag{5}$$

and none of  $F_1, F_2, F_3, F_4$  is empty. Since the four subgraphs  $G[F_1 \cup F_2] = G[B \setminus A]$ ,  $G[F_2 \cup F_3] = G[D \setminus A]$ ,  $G[F_1 \cup F_4] = G[\bar{A} \cup D]$  and  $G[F_3 \cup F_4] = G[\bar{A} \cup B]$  are all connected, we have

$$|[F_1, F_2]| \geq 1, \quad |[F_2, F_3]| \geq 1, \quad |[F_1, F_4]| \geq 1, \quad |[F_3, F_4]| \geq 1. \tag{6}$$

By equality (3), we have

$$\begin{aligned} \omega(B) &= |[bc, ad]| + |[F_1 \cup F_2, F_3 \cup F_4]|, \\ \omega(D) &= |[dc, ab]| + |[F_2 \cup F_3, F_1 \cup F_4]|. \end{aligned} \tag{7}$$

Because of Lemma 6(c), we consider the following three cases.

Case 1.  $G[A] \cong C_4$ .

In this case,  $\lambda_3(G) = 2|[ad, bc]| = 4$ . Then equalities hold in (4). By the equalities (7) and the inequalities (6), we have

$$\omega(B) \geq 4 \quad \text{and} \quad \omega(D) \geq 4. \tag{8}$$

Since  $\omega(B) = \omega(D) = \lambda_3(G) = 4$ , all inequalities in (8) become equalities. In particular,  $|[F_1, F_2]| = |[F_1, F_4]| = 1$  and  $|[F_1, F_3]| = 0$ . Combining these with  $|[F_1, A]| = |[b, F_1]| = 1$ , we have  $\omega(F_1) = 3 < \lambda_3(G)$ . Then by Lemma 1,  $G[F_1]$  is  $\lambda_3$ -independent. Since  $G[F_1 \cup F_2]$  is connected and  $|[F_1, F_2]| = 1$ , we see that  $G[F_1]$  is connected. Thus  $|F_1| \leq 2$ .

If  $|F_1| = 2$ , then  $G[F]$  is  $\lambda_3$ -non-trivial, where  $F = F_1 \cup \{b\}$ . On the other hand, by noting that  $|F| = 3 < \alpha_3(G)$  and  $\omega(F) = \omega(F_1) - |[b, F_1]| + d_A(b) = 3 - 1 + 2 = 4 = \lambda_3(G)$ , it follows from Lemma 1 that  $G[F]$  is  $\lambda_3$ -independent, a contradiction. So  $|F_1| = 1$ . Similarly, it can be deduced that  $|F_2| = |F_3| = |F_4| = 1$ . Then  $G$  is a 3-cube.

Case 2.  $G[A] \cong K_4^-$ .

In this case,  $\lambda_3(G) = 2|[ad, bc]| = 6$ . We assume, without loss of generality, that  $bd \in E(G)$ . By equality (7) and  $\omega(B) = \omega(D) = \lambda_3(G) = 6$ , we have

$$|[F_1, F_3 \cup F_4]| + |[F_2, F_3 \cup F_4]| = |[F_1 \cup F_2, F_3 \cup F_4]| = 3, \quad (9)$$

and

$$|[F_2 \cup F_3, F_1]| + |[F_2 \cup F_3, F_4]| = |[F_2 \cup F_3, F_1 \cup F_4]| = 3. \quad (10)$$

By (6) and (9), one of  $[F_1, F_3 \cup F_4]$  and  $[F_2, F_3 \cup F_4]$  is 1. Suppose, without loss of generality, that  $[F_1, F_3 \cup F_4] = 1$ . Then  $[F_1, F_4] \leq 1$ . By (6) and (10),  $[F_1, F_2 \cup F_3] \leq 2$ . By equality (5) and Lemma 6(e),  $[b, F_1] = d_A(b) - 1 = 2$ . Hence  $\omega(F_1) = |[b, F_1]| + |[F_1, F_2 \cup F_3]| + |[F_1, F_4]| \leq 5 < \lambda_3(G)$ . By Lemma 1,  $G[F_1]$  is  $\lambda_3$ -independent. Since  $G[F_1 \cup F_4]$  is connected and  $[F_1, F_4] = 1$ , we see that  $G[F_1]$  is connected and thus  $|F_1| \leq 2$ . Combining this with  $[b, F_1] = 2$ , we have  $|F_1| = 2$ . Set  $F = F_1 \cup \{b\}$ . Then both  $G[F]$  and  $G[\bar{F}]$  are connected. By noting that  $|F| = 3 < \alpha_3(G)$  and

$$\omega(F) = \omega(F_1) + d_A(b) - |[b, F_1]| \leq 5 + 3 - 2 = 6 = \lambda_3(G),$$

it follows from Lemma 1 that  $G[F]$  is  $\lambda_3$ -independent, a contradiction.

Case 3.  $G[A] \cong K_4$ .

In this case,  $\lambda_3(G) = 2|[ad, bc]| = 8$ . For each vertex  $x \in A$ , we have  $d_A(x) = d_A(x) - 1 = 2$  by Lemma 6(e). It follows from (5) that  $[a, F_4] = [b, F_1] = [c, F_2] = [d, F_3] = 2$ . Similar to the deduction of (9) and (10), we have

$$|[F_1, F_3 \cup F_4]| + |[F_2, F_3 \cup F_4]| = |[F_2, F_1 \cup F_4]| + |[F_3, F_1 \cup F_4]| = 4.$$

Then  $[F_1, F_3 \cup F_4] \leq 3$ . Suppose, without loss of generality, that  $[F_2, F_1 \cup F_4] \leq 2$ . Then  $[F_2, F_1] \leq 2$ . If  $[F_1, F_3 \cup F_4] = 3$ , then  $[F_2, F_3 \cup F_4] = 1$ . Set  $F = F_2 \cup \{c\}$ . Then

$$\omega(F) = d_A(c) + |[F_2, F_1]| + |[F_2, F_3 \cup F_4]| \leq 3 + 2 + 1 = 6 < \lambda_3(G).$$

By Lemma 1,  $F$  is  $\lambda_3$ -independent. On the other hand, since  $[c, F_2] = 2$ , we see that  $G[F]$  contains a  $\lambda_3$ -non-trivial component, a contradiction. If  $[F_1, F_3 \cup F_4] \leq 2$ , set  $F = F_1 \cup \{b\}$ . Then

$$\omega(F) = d_A(b) + |[F_1, F_2]| + |[F_1, F_3 \cup F_4]| \leq 3 + 2 + 2 = 7 < \lambda_3(G).$$

Similar to the above, we arrive at a contradiction.  $\square$

## References

- [1] C. Balbuena, P. García-Vázquez, X. Marcote, Sufficient conditions for  $\lambda'$ -optimality in graphs with girth  $g$ , *J. Graph Theory* 52 (2006) 73–86.
- [2] D. Bauer, F. Boesch, C. Suffel, R. Van Syke, On the validity of a reduction of reliable network design to a graph extremal problem, *IEEE Trans. Circuits Syst.* 34 (1989) 1579–1581.
- [3] J.A. Bondy, U.S.R. Murty, *Graph Theory with Applications*, Macmillan, London, 1976.
- [4] P. Bonsma, N. Ueffing, L. Volkmann, Edge-cuts leaving components of order at least three, *Discrete Math.* 256 (2002) 431–439.
- [5] A.H. Esfahanian, S.L. Hakini, On computing a conditional edge connectivity of a graph, *Inform. Process. Lett.* 27 (1988) 195–199.
- [6] J. Fábrega, M.A. Fiol, Extraconnectivity of graphs with large girth, *Discrete Math.* 127 (1994) 163–170.
- [7] A. Hellwig, L. Volkmann, Maximally edge-connected and vertex-connected graphs and digraphs: A survey, *Discrete Math.* 308 (2008) 3265–3296.
- [8] Y. Hong, Q. Liu, Z. Zhang, Minimally restricted edge connected graphs, *Appl. Math. Lett.* (2007) doi:10.1016/j.aml.2007.09.004.
- [9] L. Lovász, *Combinatorial Problems and Exercises*, North-Holland Publishing Company, 1979.
- [10] J.X. Meng, Y.H. Ji, On a kind of restricted edge connectivity of graphs, *Discrete Appl. Math.* 117 (2002) 183–193.
- [11] J.X. Meng, Optimally super-edge-connected transitive graphs, *Discrete Math.* 260 (2003) 239–248.
- [12] J.P. Ou, Edge cuts leaving components of order at least  $m$ , *Discrete Math.* 305 (2005) 365–371.
- [13] M. Wang, Q. Li, Conditional edge connectivity properties, reliability comparisons and transitivity of graphs, *Discrete Math.* 258 (2002) 205–214.
- [14] Ying-Qian Wang, Qiao Li, Super-edge-connectivity properties of graphs with diameter 2, *J. Shanghai Jiaotong Univ.* 33 (6) (1999) 646–649 (in Chinese).
- [15] Ying-Qian Wang, Qiao Li, Sufficient conditions for a graph to be maximally restricted edge-connected, *J. Shanghai Jiaotong Univ.* 35 (8) (2001) 1253–1255 (in Chinese).
- [16] Ying-Qian Wang, Li Qiao, A sufficient condition for the equality between the restricted edge-connectivity and minimum edge-degree of graphs, *Appl. Math. J. Chinese Univ.* 16A (3) (2001) 269–275 (in Chinese).
- [17] J.M. Xu, K.L. Xu, On restricted edge-connectivity of graphs, *Discrete Math.* 243 (2002) 291–298.
- [18] Z. Zhang, J.X. Meng, Restricted edge connectivity of edge transitive graphs, *Ars Combinatoria LXXVIII* (2006) 297–308.
- [19] Z. Zhang, J.J. Yuan, Degree conditions for restricted-edge-connectivity and isoperimetric-edge-connectivity to be optimal, *Discrete Math.* 307 (2007) 293–298.
- [20] Z. Zhang, Q.H. Liu, Sufficient conditions for a graph to be  $\lambda_k$ -optimal with given girth and diameter (submitted for publication).
- [21] Z. Zhang, J.J. Yuan, A proof of an inequality concerning  $k$ -restricted edge connectivity, *Discrete Math.* 304 (2005) 128–134.