



## Original Article

# Domain Decomposition Strategy for Pin-wise Full-Core Monte Carlo Depletion Calculation with the Reactor Monte Carlo Code

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## ABSTRACT

Because of prohibitive data storage requirements in large-scale simulations, the memory problem is an obstacle for Monte Carlo (MC) codes in accomplishing pin-wise three-dimensional (3D) full-core calculations, particularly for whole-core depletion analyses. Various kinds of data are evaluated and quantificational total memory requirements are analyzed based on the Reactor Monte Carlo (RMC) code, showing that tally data, material data, and isotope densities in depletion are three major parts of memory storage. The domain decomposition method is investigated as a means of saving memory, by dividing spatial geometry into domains that are simulated separately by parallel processors. For the validity of particle tracking during transport simulations, particles need to be communicated between domains. In consideration of efficiency, an asynchronous particle communication algorithm is designed and implemented. Furthermore, we couple the domain decomposition method with MC burnup process, under a strategy of utilizing consistent domain partition in both transport and depletion modules. A numerical test of 3D full-core burnup calculations is carried out, indicating that the RMC code, with the domain decomposition method, is capable of pin-wise full-core burnup calculations with millions of depletion regions.

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## 1. Introduction

With the higher requirements for the safety and economy of nuclear reactors, as well as the developments of new types of nuclear systems, traditional methods and tools for reactor analysis are being challenged. The Monte Carlo (MC) method

is becoming an important area of research for the next generation of methods for reactor physics calculations. With the development of parallel computing technology, the expectations are rising to see the MC method being truly applied in nuclear reactor engineering design practices [1]. However, a prohibitive amount of data is required for storage in large-

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scale calculations in MC codes. Such excessive memory demands turn into a key obstacle for the application of MC method in accomplishing pin-wise three-dimensional (3D) full-core calculations. In particular, for whole-core burnup calculations with millions of burnup regions, the data storage reaches up to hundreds of gigabytes or even terabytes, which far exceed the capacity of current computers.

Data decomposition [2,3] and domain decomposition [4,5] are two feasible ways to solve the memory problem of MC. Through the former method, specific types of data are decomposed and distributed on different processors and parallel communications are called for data operation in neutron-simulating processes. For the latter method, the idea is to divide the problem model into smaller geometry domains, which are assigned to different processors. The domain-related data are, meanwhile, decomposed. The particles are communicated between processors in the domain decomposition method as the tracks of particles are cut into pieces.

In previous studies, tally data decomposition (TDD) algorithms [6] have been designed and implemented based on the Reactor Monte Carlo (RMC) code [7]. Thereafter, a combination of TDD and depletion isotope data decomposition [8] is utilized to alleviate the memory problem, enabling simple 3D whole-core MC burnup calculations with hundreds of thousands of depletion regions [9]. However, the memory problem still exists for larger scale or fine 3D whole-core burnup, because the material data cannot be decomposed in the TDD method. In this paper, the domain decomposition method is investigated to solve the memory problem thoroughly. Through this work, fine 3D whole-core burnup calculations with millions of depletion region are achieved.

## 2. Memory evaluation of MC codes

For an in-depth knowledge of the memory problem of MC, it is necessary to classify data and analyze each data class quantitatively. Taking RMC as the reference, normally suitable to other MC codes, the data can be classified into six categories: geometry, material, nuclear data, particles, tallies, and burnup. The memory model can be constructed by going deep into each data type and evaluating their memory sizes in detail, as shown in Eqs. (1) and (2):

$$M = M_{\text{geo}} + M_{\text{mat}} + M_{\text{cs}} + M_{\text{part}} + M_{\text{tally}} + M_{\text{burn}} + M_{\text{temp}} \quad (1)$$

$$M \approx N_{\text{cell}} \bar{m}_{\text{cell}} + N_{\text{mat}} \bar{m}_{\text{mat}} + N_{\text{tot.nuc}} \bar{m}_{\text{nuc.cs}} + N_{\text{part}} m_{\text{part}} + N_{\text{tally}} \bar{m}_{\text{tally}} + N_{\text{burncell}} m_{\text{burncell}} \quad (2)$$

where total memory usage of a code,  $M$ , is the sum of the memory of different data types,  $M_x$ , with  $x$  as the data category. For example,  $M_{\text{geo}}$  defines memory of geometry data.  $M_{\text{temp}}$  represents all other temporary and supporting data, which are generally negligible in the memory footprint. Furthermore, all data of concern have a vector structure, and their memory sizes are proportional to the amount of unit data. Total memory approximates into Eq. (2), where  $N_y$  and  $m_y/\bar{m}_y$  are the number and unit or average storage size of specific data structure  $y$ , respectively.

Specifically, for the RMC code, the unit or average storage of each data type can be estimated. For example, data of one material contain names (12 bytes), ID (4 bytes), and atom/mass

densities (16 bytes) of all nuclides in the material. For a depletion calculation, assuming there is an average of 150 nuclides/region, the memory storage of one material is about  $32 \times 150 = 4.8 \times 10^3$  bytes. Particle data contain eight double-precision floating variables (3 for coordinates, 3 for direction, 1 for energy, and 1 for weight) to record particle state information, and therefore, its unit size is 64 bytes. Similarly, unit storage of tally, which is composed of statistics and filter data, is about 70 bytes. For burnup data, RMC accounts for 1500 nuclides in the depletion chain and the predictor–corrector method is used, and therefore, the unit datum is  $3.6 \times 10^4$  bytes. Finally, Eq. (3) is obtained to describe the memory consumption in burnup simulations using RMC.

$$M_{\text{RMC}} \approx (100 \times N_{\text{cell}} + 4.8 \times 10^3 \times N_{\text{mat}} + 2 \times 10^6 \times N_{\text{tot.nuc}} + 64 \times N_{\text{part}} + 70 \times N_{\text{tally}} + 3.6 \times 10^4 \times N_{\text{burncell}}) \text{ bytes} \quad (3)$$

Table 1 summarizes unit storage, scale of unit, and maximum storage of each data type. It can be seen that three types of data (i.e., tally data, burnup, and material) are the main sources of memory problems of MC codes.

The Hoogenboom–Martin whole core [10] was chosen as a case study of large-scale MC burnup calculations. There are a total of 241 assemblies and 63,624 fuel rods in the Hoogenboom–Martin core. In the modeling, each rod contains 24 burnup regions (12 axially by 2 radially) to perform the depletion calculation. Table 2 predicts the memory storage using the memory model.

## 3. Domain decomposition method

Different from the data decomposition method, spatial domain decomposition (SDD) divides spatial geometry into domains, which are simulated separately by parallel processors, and particles crossing domains are communicated for continuing tracking.

As indicated in Fig. 1, the main steps involved in implementing SDD in particle transport MC code include (1) dividing

**Table 1 – Memory storage evaluation of data types.**

Data types	Unit storage	Scale	Maximum storage
Geometry	100 bytes	$10^0$ – $10^6$	0.1 GB
Material	4.8 KB	$10^0$ – $10^7$	10 GB
Nuclear data <sup>a</sup>	2 MB	$10^0$ – $10^2$	1 GB
Particle	64 bytes	$10^4$ – $10^6$	0.1 GB
Tally	70 bytes	$0$ – $10^{10}$	100 GB
Burnup	36 KB	$0$ – $10^7$	100 GB

GB, gigabyte; KB, kilobyte; MB, megabyte.  
<sup>a</sup> Assuming nuclides are in single temperature.

**Table 2 – Memory storage of H–M whole-core burnup.**

Data types	Memory storage
Material	7.3 GB
Nuclear data	400 MB
Particle	64 MB
Tally	64.1 GB
Burnup	55.0 GB
Total	126.9 GB

GB, gigabyte; MB, megabyte.

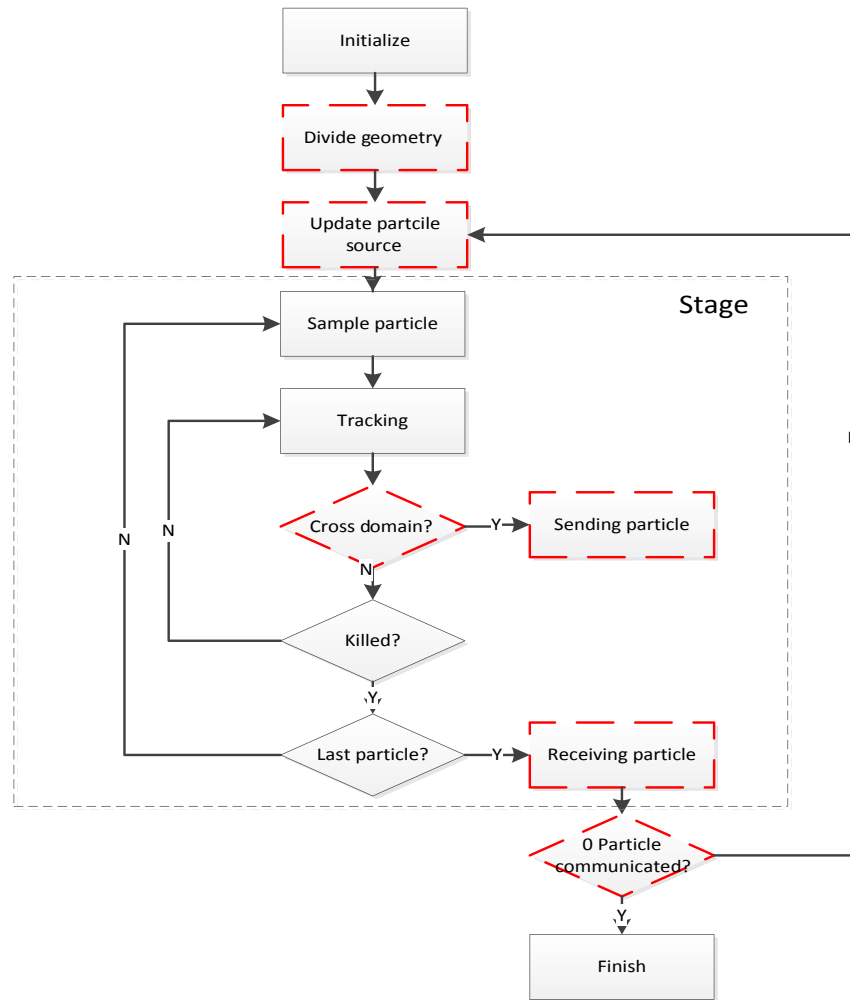


Fig. 1 – Flow chart of the domain decomposition method.

model geometry and decomposing data of tallies and materials simultaneously; (2) tracking particle histories by “stages,” which are bounded with particle communications; and (3) finishing simulation by checking all particle stages and communications.

It should be noted that except for the domain partition strategy, the particle communication algorithm is a key part of

SDD as well. The asynchronous particle communication algorithm is designed using the message passing interface [MPI] [11], as shown in Fig. 2. In the algorithm, particles flying out of current domains are buffered in a local processor, and a buffer will be sent to the remote processor when it is full, or when the current processor finishes tracking all particles. The MPI

```

1: decompose Geometry into Domains to all processes
2: for cycle = 1 to CycleNum
3:   do // stages
4:     | post non-blocking receives to all other processes(MPI_Irecv)
5:     | for neutron = 1 to NeutronNum
6:     |   | tracking history
7:     |   |   | test all receives(MPI_Test)
8:     |   |   |   | while(one receive success)
9:     |   |   |   |   | add received particle to local bank
10:    |   |   |   |   | if(not last rcv)post receive to the same target (MPI_Irecv)
11:    |   |   |   |   | end while loop
12:    |   |   |   | if(cross domain)
13:    |   |   |   |   | save it to particle buffer
14:    |   |   |   |   | if(buffer is full)send this buffer to target(MPI_Isend)
15:    |   |   |   |   | end one history
16:    |   |   |   | end for loop
17:    |   |   | send all buffer out(MPI_Isend)
18:    |   |   | wait all send and receive success(MPI_Wait)
19:    |   |   | reduce all received particles number(MPI_Allreduce)
20:    |   | while(received particle exist)
21:    |   | process cycle end
22:   end for loop
  
```

Fig. 2 – Asynchronous particle communication algorithm in the domain decomposition method.

nonblocking functions (MPI\_Isend/MPI\_Irecv and MPI\_Test/MPI\_Wait) are used for overlap of computing and passing messages, to shorten total communication time as much as possible.

Note that the particle buffer has a fixed size (Line 14 in the algorithm), which determines the period of communications. Theoretically, greater buffer size leads to less communication time, but the memory of buffers will increase as well. In practice, the utilization of nonblocking communication makes the parallel performance of the algorithm insensitive with the buffer size in a reasonable range. The buffer size is usually set as  $10^3$  in the SDD implementation in RMC.

#### 4. Coupling domain decomposition with MC burnup

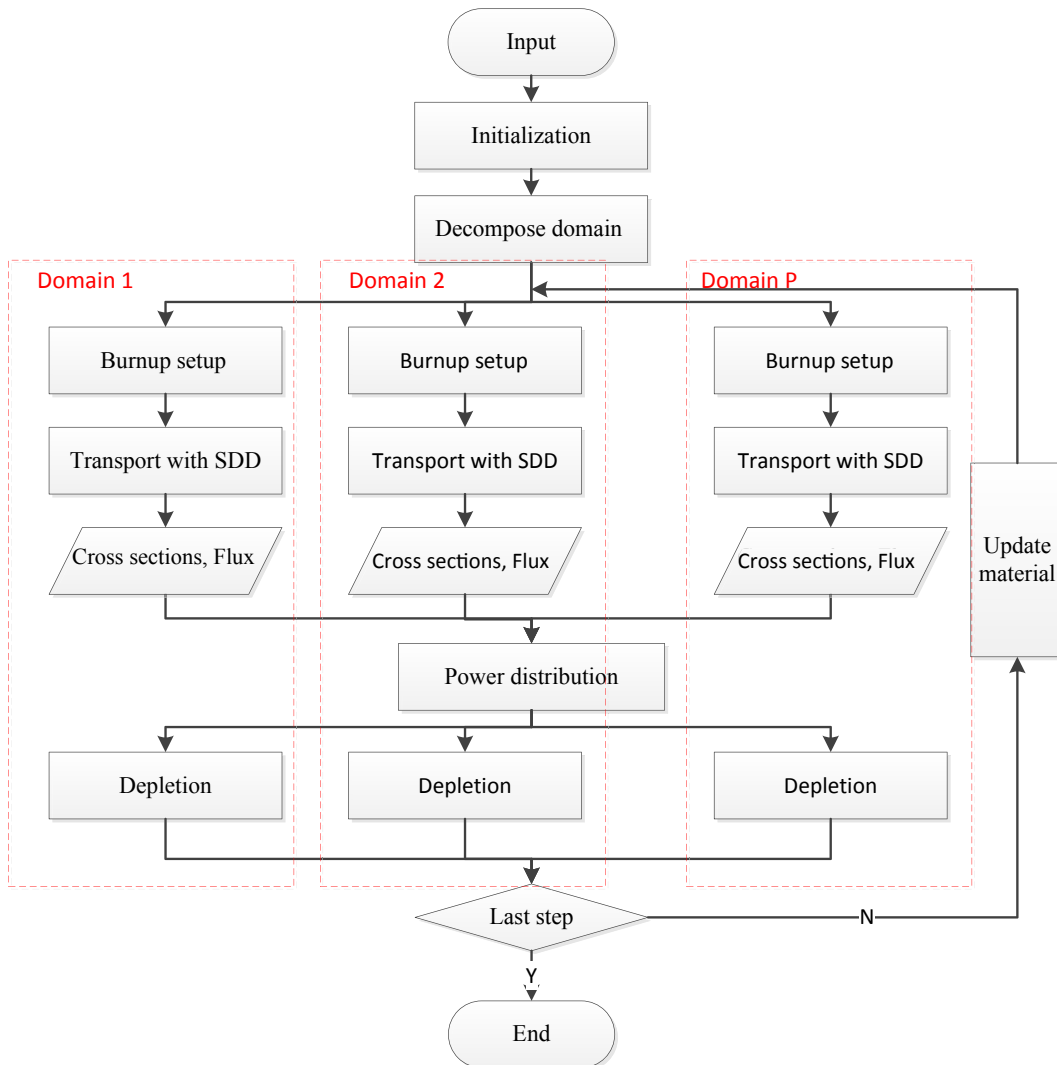
##### 4.1. Coupling strategy and memory estimation

It is known that the MC burnup calculation process is a combination of MC neutron transport simulation and depletion equation computation. The basic calculation unit in depletion, that is, a burnup cell, is a geometrical region

**Table 3 – Main calculation parameters used.**

Parameters	Data
Number of histories in Monte Carlo transport	1,000,000 particle/cycle 750 cycles (250 inactive)
Number of burnup regions	1,526,976
Number of isotopes/burnup region	1,487
Number of time steps	14
Step length ( $\Delta MWD/kgU$ )	0.1, 0.4, 0.5, $1 \times 11$
Number of inner steps in depletion	10
Burnup strategy	Predictor–corrector
Power density (W/gU)	30
Total burnup (MWD/kgU)	12
Number of parallel processors	96

as well. Therefore, it is possible to couple domain decomposition with MC burnup by utilizing consistent domain partition in the transport and depletion processes. In other words, burnup regions can be decomposed automatically according to geometry decomposition in transport.



**Fig. 3 – Frame of coupled domain decomposition. SDD, spatial domain decomposition.**

Through the coupling, geometry-based data including data on tally, material, and isotope densities (i.e., the main memory-consuming data types) are all decomposed.

From the memory model in the “Memory Evaluation of MC Codes” section, assuming the geometry is evenly decomposed into  $P$  domains, the memory size of one domain can be estimated as Eq. (4).

$$M_{\text{Domain}} \approx M_{\text{geo}} + M_{\text{cs}} + M_{\text{part\_Domain}} + \frac{M_{\text{mat}} + M_{\text{tally}} + M_{\text{burn}}}{P} \quad (4)$$

Three major data types are all proportional to the number of burnup cells, as shown in Eq. (5). Here, it is assumed that there are about 150 isotopes in every burnable cell or material for transport simulation. A total of four types

of one-group cross section for every isotope in every material should be enabled for the depletion calculation [12].

$$M_{\text{mat}} + M_{\text{tally}} + M_{\text{burn}} \approx N_{\text{burncell}} (4.8 \times 10^3 + 70 \times 600 + 3.6 \times 10^4) \times \text{bytes} \approx 0.08 N_{\text{burncell}} \text{MB} \quad (5)$$

Therefore, the necessary number of domain decomposition can be figured out according to the capacity of computer RAM. For example, to carry out MC burnup calculations on 4-GB RAM computers, assuming that memory size of data not decomposed is no more than 2 GB, the condition and result can be derived as follows [Eqs. (6) and (7)], indicating that no less than 42 domain partitions is adequate for a 1-million-region burnup calculation.

$$\frac{M_{\text{mat}} + M_{\text{tally}} + M_{\text{burn}}}{P} \leq 2\text{GB} \quad (6)$$

$$P \geq 4.14 N_{\text{burncell}} \times 10^{-5} \quad (7)$$

#### 4.2. Implementation of coupled domain decomposition

Coupled domain decomposition is implemented in RMC based on the aforementioned coupling strategy, as shown in Fig. 3. It can be seen that depletion regions are grouped according to domain partition in the transport process.

### 5. Numerical test of 3D full-core burnup calculation

The Hoogenboom–Martin core is used for 3D full-core burnup calculation tests. Table 3 shows the main calculation parameters used in the calculation; 1.5 million burnup regions are set up with 24 burnup regions (12 axially by 2 radially) for each rod, illustrated by Fig. 4. This full core is decamped into 96 domains with eight pieces radially and 12 slices axially. The partitions in radial planes are performed in a way to make every assembly fully belong to only one domain, while balancing the loads of domains as well as possible.

The “Inspur TS10000” cluster at Tsinghua University, Beijing, China was used to run the 96-domain parallel calculation.

(A)

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21
22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42
43	44	45	46	47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63
64	65	66	67	68	69	70	71	72	73	74	75	76	77	78	79	80	81	82	83	84
85	86	87	88	89	90	91	92	93	94	95	96	97	98	99	100	101	102	103	104	105
106	107	108	109	110	111	112	113	114	115	116	117	118	119	120	121	122	123	124	125	126
127	128	129	130	131	132	133	134	135	136	137	138	139	140	141	142	143	144	145	146	147
148	149	150	151	152	153	154	155	156	157	158	159	160	161	162	163	164	165	166	167	168
169	170	171	172	173	174	175	176	177	178	179	180	181	182	183	184	185	186	187	188	189
190	191	192	193	194	195	196	197	198	199	200	201	202	203	204	205	206	207	208	209	210
211	212	213	214	215	216	217	218	219	220	221	222	223	224	225	226	227	228	229	230	231
232	233	234	235	236	237	238	239	240	241	242	243	244	245	246	247	248	249	250	251	252
253	254	255	256	257	258	259	260	261	262	263	264	265	266	267	268	269	270	271	272	273
274	275	276	277	278	279	280	281	282	283	284	285	286	287	288	289	290	291	292	293	294
295	296	297	298	299	300	301	302	303	304	305	306	307	308	309	310	311	312	313	314	315
316	317	318	319	320	321	322	323	324	325	326	327	328	329	330	331	332	333	334	335	336
337	338	339	340	341	342	343	344	345	346	347	348	349	350	351	352	353	354	355	356	357
358	359	360	361	362	363	364	365	366	367	368	369	370	371	372	373	374	375	376	377	378
379	380	381	382	383	384	385	386	387	388	389	390	391	392	393	394	395	396	397	398	399
400	401	402	403	404	405	406	407	408	409	410	411	412	413	414	415	416	417	418	419	420
421	422	423	424	425	426	427	428	429	430	431	432	433	434	435	436	437	438	439	440	441

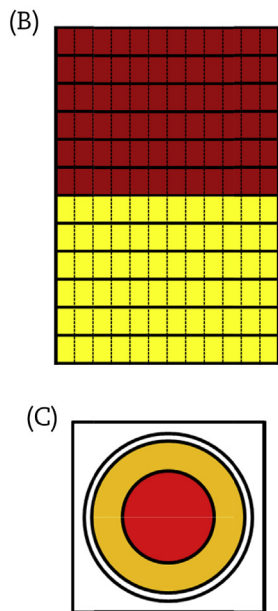


Fig. 4 – Domain and depletion region partition in three-dimensional full-core burnup. (A) Eight domains (8 colors) radially, without cutting each assembly; (B) 12 segments axially; and (C) two rings/rod.

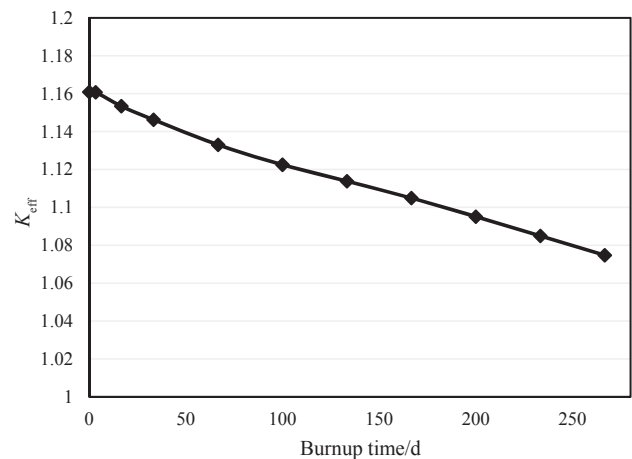


Fig. 5 – Variation of core K-effective ( $K_{\text{eff}}$ ) with burnup.



As estimated in Section 2, the total memory storage consumed in this calculation is up to 127 GB, which exceeds the memory capacity of a normal computer processor. Nevertheless, it is demonstrated that the RMC code integrated with the domain decomposition method is capable of completing the calculation, and memory storage/processor is about 3 GB. Figs. 5–10 show the results of the calculation.

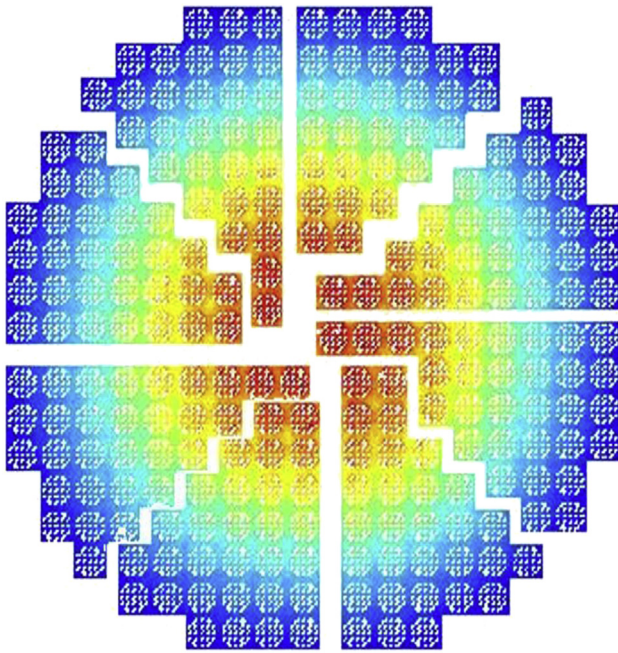


Fig. 6 – Core radial power distribution of Layer 0 (top) at 0 MWD/kgU, d, day.

Fig. 5 is the K-effective ( $K_{eff}$ ) variation along with burnup time. Figs. 6 and 7 are the radial pin-by-pin power distributions of the top layer of the core at 0 MWD/kgU and 5 MWD/kgU burnup points, respectively, whereas Figs. 8 and 9 are radial burnup distributions of the middle of the core, all plotted by jointing eight domains together. Fig. 10 shows the axial power distributions of 12 radially integrated segments at different burnup points.

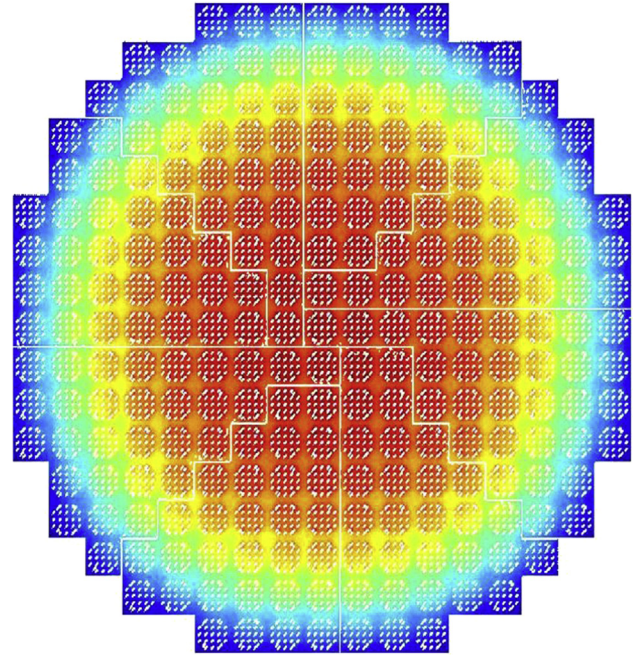


Fig. 8 – Core radial burnup distribution of Layer 6 (middle) at 5 MWD/kgU.

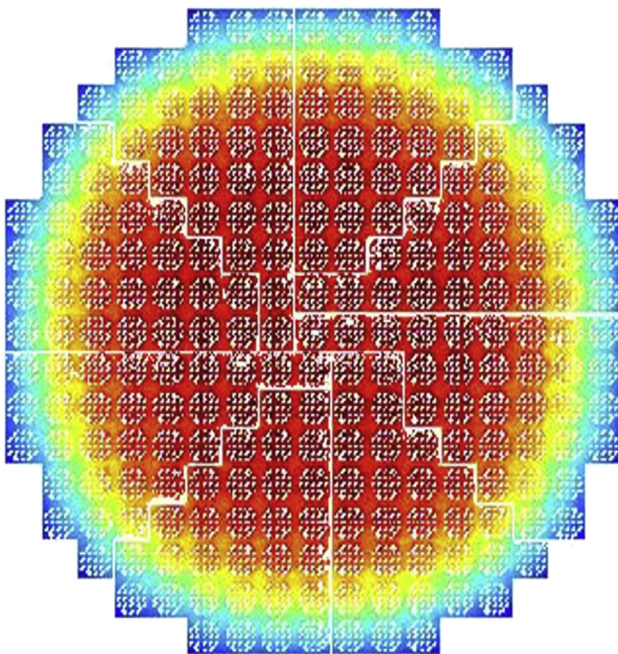


Fig. 7 – Core radial power distribution of Layer 0 (top) at 5 MWD/kgU.

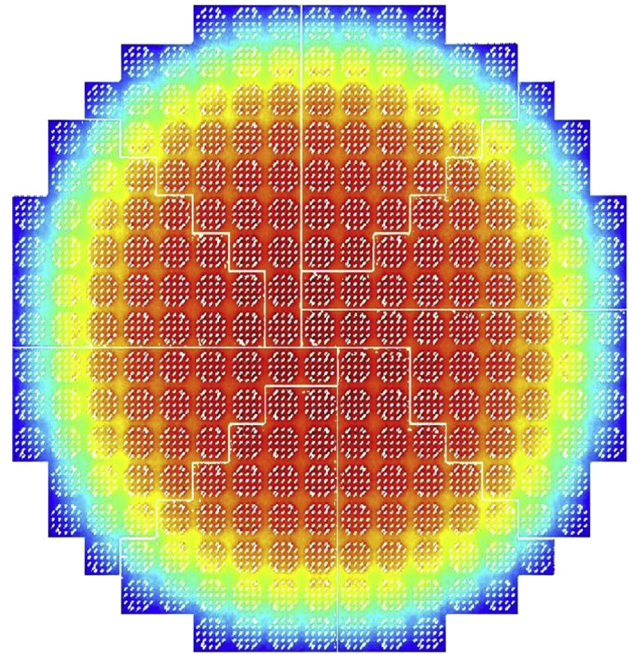


Fig. 9 – Core radial burnup distribution of Layer 6 (middle) at 8 MWD/kgU.

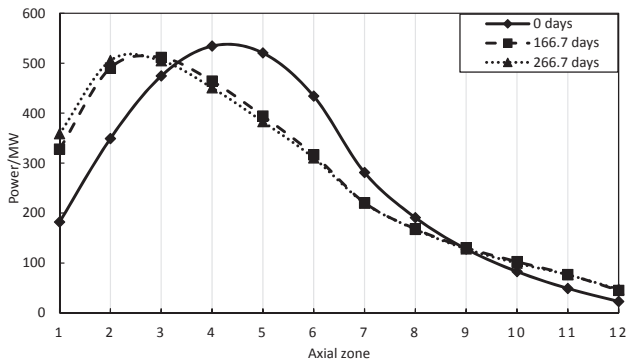


Fig. 10 – Axial power distributions.

## 6. Discussion

From a qualitative point of view, the results of the whole-core burnup calculation are reasonable. However, it should be noted that the numerical test is mainly performed to verify the capability of the domain decomposition method for extending the scale of MC burnup analysis. To achieve practical application in nuclear engineering, further improvements remain to be done. First, some more features need to be integrated, for example, criticality search function, to maintain the operation state of the core. Second, thermal-hydraulic coupling is required in reactor analysis to account for temperature feedback, which relies on efficient treatment of temperature-dependent nuclear cross sections, that is, on-the-fly Doppler broadening. Progress associated with these issues are under study [13,14].

In addition, the domain decomposition burnup strategy (e.g., the usage of a huge amount of data produced in depletion and the method of domain partition) is still not ideal for engineering use. Only one method of domain partition is used and performances are not tested in this study. Research shows that the pattern of decomposition, which determines source load imbalance between different domains, influences computing performance significantly [15].

In conclusion, the paper deals with the domain decomposition method and its coupling with burnup calculations in the RMC code. The successful running of the Hoogenboom–Martin benchmark with the assumed burnup cases demonstrates the effectiveness of domain decomposition methods for solving memory problems. It is indicated that domain-decomposed MC codes are capable of performing pin-wise full-core burnup calculations with millions of depletion regions. The next steps are to improve the practicability of the domain-decomposed burnup strategy and to integrate it with other latest features, such as thermal–hydraulics feedback, to eventually realize practical engineering application.

## Conflicts of interest

All contributing authors declare no conflicts of interest.

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