

Appl Categor Struct (2015) 23:215–262
DOI 10.1007/s10485-013-9331-x

Monads of Regular Theories

Stanisław Szawiel · Marek Zawadowski

Received: 10 August 2012 / Accepted: 4 September 2013 / Published online: 29 September 2013
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Abstract We characterize the categories of semi-analytic monads, regular Lawvere theories, and regular operads that are equivalent to the category of regular equational theories. We also show that the category of all finitary monads on *Set* is monadic over the category of semi-analytic functors.

Keywords Operad · Monad · Lawvere theory · Equational theory · Factorization system · Distributive law

Mathematics Subject Classifications (2010) Primary 18C05 · 18C10 · 18C15 · 18D50 · 03F25 · 08B05 · 03G30; Secondary 18A32 · 18D10 · 18C20 · 08B20

1 Introduction

The category of algebras of a (finitary) equational theory can be equivalently described as a category of models of a Lawvere theory or as a category of algebras of a finitary monad on the category *Set* or a category of algebras of a (generalized) operad. In fact, the four categories of (finitary) equational theories, Lawvere theories, finitary monads on *Set* and (generalized) operads are equivalent. These equivalences induce a correspondence between various subcategories. In [14] we have given an intrinsic characterizations of equational theories and Lawvere theories that correspond to the analytic and polynomial monads on *Set*. This was achieved via correspondence with symmetric and rigid operads.

An equational theory is regular (cf. [12]) if it has a set of axioms where each equation contains the same variables on both sides. Thus the theories of monoids and

S. Szawiel · M. Zawadowski (✉)
Instytut Matematyki, Uniwersytet Warszawski, ul. S. Banacha 2, 00-913 Warszawa, Poland
e-mail: zawado@mimuw.edu.pl

S. Szawiel
e-mail: szawiel@mimuw.edu.pl

join-semilattices are regular but the theory of groups is not. The main objective of this paper is to describe the categories of regular Lawvere theories **RegLT**, semi-analytic monads **SanMnd**, and regular operads **RegOp** that correspond to the category of regular equational theories with regular morphisms.

In [14] symmetric operads proved to be very useful in describing correspondences of this kind. Full and regular operads (cf. [6, 13], Section 2) play the same role here. The category of full operads is the category of monoids in the monoidal category $Set^{\mathbb{F}}$, where \mathbb{F} is the skeleton of the category of finite sets, with the substitution tensor. This monoidal category was first considered in [6] and already there it was noticed that the category of monoids in $Set^{\mathbb{F}}$ is equivalent to the category of finitary monads on Set . The substitution tensor on $Set^{\mathbb{F}}$ can be described as follows. We can think of a functor $A : \mathbb{F} \rightarrow Set$ as of a signature on which we have an action of finite functions from \mathbb{F} . $A(n)$ is the set of n -ary function symbols. Then, if $a \in A(n)$ and $f : \underline{n} \rightarrow \underline{m}$, we can think of $f \cdot a = A(f)(a) \in A(m)$ as an m -ary operation obtained by substitution of variables along f so that the following equation

$$a(v_{f(1)}, \dots, v_{f(n)}) = (f \cdot a)(v_1, \dots, v_m) \tag{1}$$

holds. If we interpret objects A, B of $Set^{\mathbb{F}}$ in this way, the tensor $A \otimes B$ is the signature that arises by plugging tuples operations, say $a_1, \dots, a_n \in A$, into single operations, say $b \in B(n)$. Because of the Eq. 1 we need to identify some such composed terms. For example, if $f : \underline{2} \rightarrow \underline{2}$ and $b \in B(2)$, $a_1, a_2 \in A$, then we want to identify the terms

$$(a_1, a_2; (f \cdot b)) \quad \text{and} \quad (a_{f(2)}, a_{f(1)}; b)$$

in some way.

Still, there is a certain ambiguity concerning the interpretation of the expression

$$(a_1, \dots, a_n; b)$$

The ‘clonic’ (universal algebraic) interpretation will insist that all the a_i ’s are members of a single set $A(k)$ and the composed operation should be of arity k . In categorical notation, it would be a composed operation

$$X^k \xrightarrow{\langle a_1, \dots, a_n \rangle} X^n \xrightarrow{b} X$$

The ‘operadic’ (geometric) interpretation will take an arbitrary n -tuple of operations $a_i \in A(k_i)$ and the composed operation will have as arity the sum $\sum_{i=1}^n k_i$. In categorical notation, it would be a composed operation

$$X^{\sum_i k_i} \xrightarrow{a_1 \times \dots \times a_n} X^n \xrightarrow{b} X$$

In case when all functions from \mathbb{F} act on operations, both compositions (when appropriate identifications are made) lead to equivalent descriptions of the substitution tensor on $Set^{\mathbb{F}}$. This was shown in [13]. The relation between these two compositions can be explained as follows. Let \otimes temporarily denote the ‘clonic’ composition and \oplus , the ‘operadic’ composition. If $a_i \in A(k_i)$, $b \in B(n)$, $k = \sum_{i=1}^n k_i$ and $\iota_i : \underline{k_i} \rightarrow \underline{k}$ are the obvious inclusions, then

$$(a_1, \dots, a_n) \oplus b = ((\iota_1 \cdot a_1, \dots, \iota_n \cdot a_n) \otimes b$$

On the other hand, if $a_i \in A(k)$ and $\pi : k \cdot n \rightarrow k$ is the obvious projection, then

$$(a_1, \dots, a_n) \otimes b = \pi \cdot ((a_1, \dots, a_n) \oplus b)$$

Thus a monoid in $Set^{\mathbb{F}}$ consists of a set of operations that are equipped with not only a multiplication operator but also an action of the whole category \mathbb{F} . The left Kan extension $Set^{\mathbb{F}} \rightarrow \mathbf{End}$ along the inclusion $\mathbb{F} \rightarrow Set$ establishes the equivalence of those monoidal categories (\mathbf{End} is the category of finitary endofunctors on Set) and induces the equivalence of the category of full operads \mathbf{FOp} with the category \mathbf{Mnd} of finitary monads on Set .

We shall describe this tensor in $Set^{\mathbb{F}}$ in the ‘operadic’ style in Section 2. The correspondence between the category of full operads and the categories of equational theories, Lawvere theories, and finitary monads will be presented in Section 3.

Regular operads are monoids in the monoidal category $Set^{\mathbb{S}}$ with substitution tensor, where \mathbb{S} is the skeleton of the category of finite sets and surjections. Thus in regular operads only surjections act on operations.

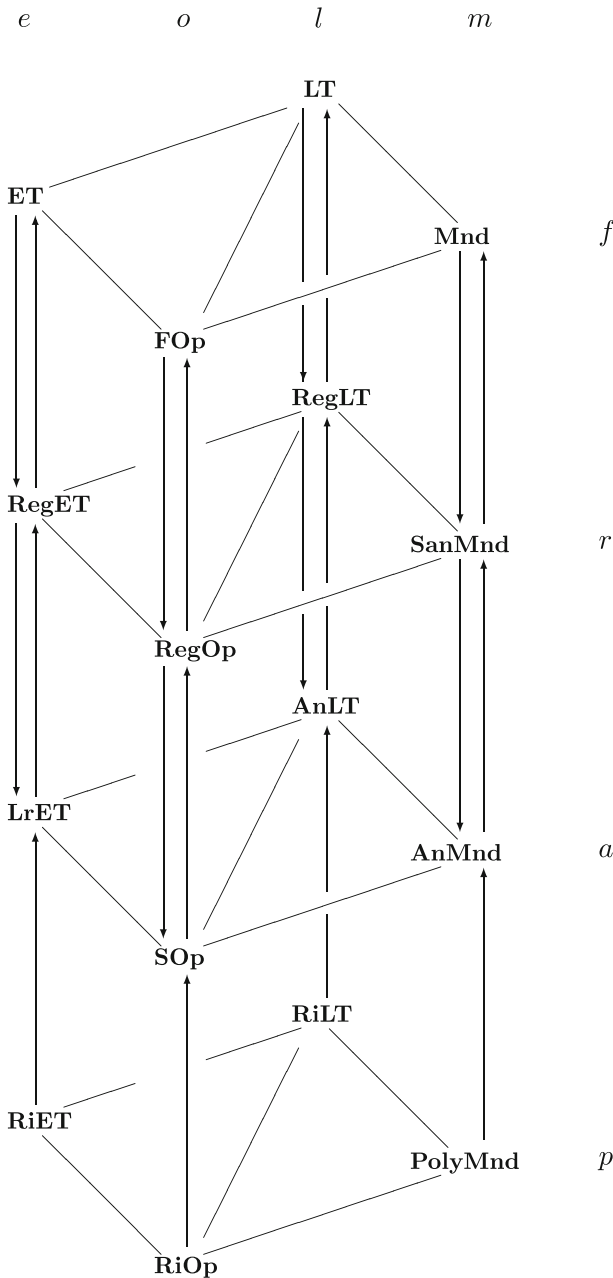
We identify the essential image of the left Kan extension $Set^{\mathbb{S}} \rightarrow \mathbf{End}$ along inclusion $\mathbb{S} \rightarrow Set$ as the category of semi-analytic functors and semi-cartesian natural transformations \mathbf{San} . We call these functors semi-analytic as they are similar to analytic ones. For example, semi-analytic functors also have presentations via series similar to the series that represent analytic functors; see Section 2. The category \mathbf{San} is a monoidal subcategory of \mathbf{End} and the monoids there form the category of semi-analytic monads equivalent to the category of regular theories; see Section 5.

As we shall show in Section 2, a finitary endofunctor on Set is semi-analytic iff it preserves pullbacks along monomorphisms and a natural transformation is semi-cartesian iff the naturality squares for monomorphism are pullbacks. Thus the notion of a semi-analytic monad is equivalent to the notion of a collection monads introduced in [10] and to the notion of a finitary taut monad introduced in [11]. The category of semi-analytic functors with semi-cartesian natural transformations \mathbf{San} was already considered in [11] as the category of finitary taut functors on Set . In [10] it was also shown that regular theories (called balanced in [10]) correspond to semi-analytic monads. This is the object part of the correspondence that we study in Section 6; cf. Theorem 6.2.

The category of regular Lawvere theories is defined very much in the spirit of the definition of the category of analytic Lawvere theories, cf. Section 2. A regular Lawvere theory is a Lawvere theory with nicely behaving isomorphisms and a factorization system consisting of the class of projections and the class of regular morphisms. Regular morphisms in a Lawvere theory are by definitions those morphisms that are right orthogonal to all projections. A regular interpretation of Lawvere theories is an interpretation of Lawvere theories that preserves regular morphisms. We show that the category of regular Lawvere theories is equivalent to the category of regular operads and hence also to the category of semi-analytic monads and to the category of regular equational theories; cf. Section 6.

The category of semi-analytic monads \mathbf{SanMnd} contains subcategories of cartesian $\mathbf{CartMnd}$ and weakly cartesian $\mathbf{wCartMnd}$ monads (i.e. those that preserve and weakly preserve pullbacks, respectively). In Section 7 we characterize the subcategories of $Set^{\mathbb{S}}$ that have as their essential images $\mathbf{wCartMnd}$ and $\mathbf{CartMnd}$. As these characterizations are a bit technical, we do not try to rephrase those conditions in terms of either Lawvere theories or equational theories.

In this way we shall describe the second level (level r) of the following diagram



The vertical lines denote adjoint equivalences. Thus up to equivalence there are only four categories in it, one on each level. One equivalent to the category of all finitary monads on Set , second equivalent to the category of all semi-analytic monads on Set , third equivalent to the category of all analytic monads on Set , and fourth equivalent

to the category of all polynomial monads on *Set*. These levels are denoted by letters *f*, *r*, *a*, and *p*, respectively. Thus all four columns of equational theories, Lawvere theories, monads and operads are equivalent. These columns are denoted by letters *e*, *l*, *m* and *o*, respectively. The vertical functors heading up, in all columns but the column of operads, are inclusions of subcategories. In the column of operads the functors heading up are more like free extensions of the actions. The lowest functors are full embeddings and the upper are embeddings that are full on isomorphisms. The vertical functors heading down, the right adjoints to those heading up, are monadic. All the squares in the diagram commute up to canonical isomorphism. The notation concerning categories involved is displayed in the above diagram.

The notation concerning functors is not displayed on the diagram but it is meant to be systematically referring to levels and columns they ‘connect’. The horizontal functors are denoted using letters from both columns they connect; the codomain by the script letter, the domain by its subscript, and the level is denoted by superscript. Thus the functor $\mathbf{AnMnd} \rightarrow \mathbf{AnLT}$ will be denoted by \mathcal{L}_m^a . We usually drop superscripts when it does not lead to confusion. Thus we can write, for example, $\mathcal{E}_o = \mathcal{E}_o^p : \mathbf{RiOp} \rightarrow \mathbf{RiET}$. The vertical functors heading up are denoted by the script letter \mathcal{P} with superscript indicating the column and subscript indicating the level of the codomain. The vertical functors heading down are denoted by the script letter \mathcal{Q} with subscript and superscript as with those heading up. Thus we have, for example, functors $\mathcal{P} = \mathcal{P}^o = \mathcal{P}_a^o : \mathbf{RiOp} \rightarrow \mathbf{SOp}$ and $\mathcal{Q} = \mathcal{Q}_a = \mathcal{Q}_a^m : \mathbf{Mnd} \rightarrow \mathbf{AnMnd}$. We will also refer to various diagonal morphisms and then we need to extend the notation concerning vertical functors by specifying both the columns of the domain and the codomain. For example, we write $\mathcal{P}_f^{ol} : \mathbf{SOp} \rightarrow \mathbf{LT}$ to denote one such functor and its right adjoint will be denoted by $\mathcal{Q}_a^{lo} : \mathbf{LT} \rightarrow \mathbf{SOp}$. In principle, this notation will leave the codomain not always uniquely specified but in practice it is sufficient and in fact usually much less is needed.

The paper is organized as follows. In Section 2 we introduce all the main categories that we shall study in the paper, i.e. the category of equational regular theories, full and regular operads, regular Lawvere theories, and semi-analytic functors and monads. We prove the equivalence of the three descriptions given. In Section 3 we briefly recall the equivalence between the category of full operads **FOP** and the three other descriptions of the category of equational theories: **ET**, **LT**, **Mnd**. In the following three Sections 4–6 we show that the category of regular operads **RegOp** is equivalent to the category of regular Lawvere theories **RegLT**, the category of semi-analytic monads **SanMnd**, and the category of regular equational theories **RegET**, respectively. In Section 4 we also show that **LT** is monadic over **RegLT**. In Section 5 we additionally show that **Mnd** is monadic over **San** and to this end we explain the distributive law that comes from the monoidal monad **W** on **San** for finitary functors. In the final Section 7 we characterize the subcategories of $\mathbf{Set}^{\mathcal{S}}$ that have as their essential images in **End** the categories of functors (weakly) preserving pullbacks.

1.1 Notation

$[n] = \{0, \dots, n\}$, $(n) = \{1, \dots, n\}$, ω - denotes the set of natural numbers. The set X^n is interpreted as $X^{(n)}$ when convenient. The skeletal category equivalent to the category of finite sets \mathbf{Set}_{fin} will be denoted by \mathbb{F} . We assume that the objects of \mathbb{F} are sets (n) for $n \in \omega$. The subcategories of \mathbb{F} with the same objects as \mathbb{F} but having as morphisms bijections, surjections and injections will be denoted by \mathbb{B} , \mathbb{S} ,

\mathbb{I} , respectively. S_n is the group of permutations of $[n]$. When S_n acts on a set A_n on the right and on the set B_n on the left, the set $A \otimes_n B$ is the usual tensor product of S_n -sets.

2 Presentations of Categories of Algebras

2.1 Equational Theories

By an *equational theory* we mean a pair of sets $T = (L, A)$, $L = \bigcup_{n \in \omega} L_n$ and L_n is the set of n -ary operations of T . The sets of operations of different arities are disjoint. The set $Tr(L, \vec{x}^n)$ of terms of L in context $\vec{x}^n = \langle x_1, \dots, x_n \rangle$ is the usual set of terms over L built with the help of variables from \vec{x}^n . We write $t : \vec{x}^n$ for the term t in context \vec{x}^n . Thus all the variables occurring in t are among those in \vec{x}^n . The set A is a set of equations in context $t = s : \vec{x}^n$, i.e. both $t : \vec{x}^n$ and $s : \vec{x}^n$ are terms in context.

A morphism of equational theories, an *interpretation*, $I : (L, A) \rightarrow (L', A')$ is given by a set of functions $I_n : L_n \rightarrow Tr(L', \vec{x}^n)$, for $n \in \omega$. The I_n s extend to functions $\bar{I}_n : Tr(L, \vec{x}^n) \rightarrow Tr(L', \vec{x}^n)$ in an obvious way. We require that for any $t = s : \vec{x}^n$ in A

$$A' \vdash \bar{I}(t) = \bar{I}(s) : \vec{x}^n$$

where $A' \vdash$ is the provability in the equational logic from axioms in the set A' . We identify two such interpretations if they are provably equal. In this way we have defined the category of equational theories **ET**.

A term in context $t : \vec{x}^n$ is *regular* if every variable in \vec{x}^n occurs in t at least once. A term in context $t : \vec{x}^n$ is *linear* if every variable in \vec{x}^n occurs in t at most once. A term in context $t : \vec{x}^n$ is *linear-regular* if it is both linear and regular. An equation $s = t : \vec{x}^n$ is *linear-regular* iff both $s : \vec{x}^n$ and $t : \vec{x}^n$ are linear-regular terms in contexts.

A *simple ϕ -substitution of a term in context $t : \vec{x}^n$* along a function $\phi : [n] \rightarrow [k]$ is a term in context denoted $\phi \cdot t : \vec{x}^k$ such that every occurrence of the variable x_i is replaced by the occurrence of $x_{\phi(i)}$. An *α -conversion of a term in context $t : \vec{x}^n$* is a simple ϕ -substitution of a term in context along a monomorphism $\phi : [n] \rightarrow [k]$.

An equational theory $T = (L, A)$ is a *linear-regular theory* iff every equation $s = t : \vec{x}^n$ that is a consequence of the theory T is a consequence of the set of linear-regular consequences of T . An interpretation is *regular (linear-regular)* iff it interprets function symbols as regular (linear-regular) terms.

A theory $T = (L, A)$ is a *rigid theory* iff it is linear-regular and for any linear-regular term in context $t : \vec{x}^n$ whenever $A \vdash t = \tau \cdot t : \vec{x}^n$, then τ is the identity permutation. $\tau \cdot t$ is the simple τ -substitution of a term in context $t : \vec{x}^n$ along a permutation $\tau \in S_n$.

We denote by **LrET** the subcategory of **ET** consisting of linear-regular theories and linear-regular interpretations. **RiET** denotes the full subcategory of **LrET** whose objects are rigid theories. **RegET** is a category of regular theories and regular interpretations. We have three inclusion functors

$$\mathbf{RiET} \longrightarrow \mathbf{LrET} \longrightarrow \mathbf{RegET} \longrightarrow \mathbf{ET}$$

with the first inclusion being full and the other two being full on isomorphisms (cf. [16]).

2.2 Operads

The symmetric operads provide another way of presenting models of an equational theory. This kind of presentation is usually very convenient, however the models defined by such operads are more specific than models of arbitrary equational theories. For example, if \mathcal{O} is a symmetric operad, then the free algebra functor $Set \rightarrow Alg(\mathcal{O})$ preserves weak wide pullbacks. Below we extend the definition of an operad so that it captures all the equational theories but still keeps the operadic flavor. The main difference is that instead of having just symmetric groups acting on sets of operations we have actions of the morphisms of the whole skeleton of the category of finite sets \mathbb{F} . Symmetric operads can be thought of as monoids for the substitution tensor on the category $Set^{\mathbb{B}}$. Similarly, \mathbb{F} -operads can be thought of as monoids for the substitution tensor on the category $Set^{\mathbb{F}}$. By **End** we denote the category of finitary endofunctors of Set . It is a strict monoidal category with the tensor being composition. The substitution tensor on $Set^{\mathbb{F}}$ makes the equivalence of categories $Set^{\mathbb{F}} \rightarrow \mathbf{End}$ given by the left Kan extension, a strong monoidal equivalence. This immediately shows that the category of \mathbb{F} -operads is equivalent to the category of finitary monads on Set . The definition below was first explicitly spelled out in [13].

A *full operad* (or \mathbb{F} -operad) \mathcal{O} consists of a family of sets \mathcal{O}_n , for $n \in \omega$, a unit element $\iota \in \mathcal{O}_1$, for any $k, n, n_1, \dots, n_k \in \omega$ with $n = \sum_{i=1}^k n_i$, a multiplication operation

$$\begin{aligned} * : \mathcal{O}_{n_1} \times \dots \times \mathcal{O}_{n_k} \times \mathcal{O}_k &\longrightarrow \mathcal{O}_n \\ (b_1, \dots, b_k, a) &\mapsto \langle b_1, \dots, b_k \rangle * a \end{aligned}$$

a left action of the morphisms on operations

$$\cdot : \mathbb{F}(n, m) \times \mathcal{O}_n \longrightarrow \mathcal{O}_m$$

for $n, m \in \omega$. The multiplication is associative with unit ι and compatible with the category action, i.e. for $a \in \mathcal{O}_n$

$$\langle a \rangle * \iota = a = \langle \iota, \dots, \iota \rangle * a; \quad 1_n \cdot a = a$$

for $a \in \mathcal{O}_n, b_i \in \mathcal{O}_{k_i}, c_{ij} \in \mathcal{O}_{m_{ij}}$, for $i \in (n), j \in (k_i]$, we have

$$\begin{aligned} \langle c_{1,1}, \dots, c_{1,k_1}, \dots, c_{n,1}, \dots, c_{n,k_n} \rangle * (\langle b_1, \dots, b_n \rangle * a) &= \\ = \langle \langle c_{1,1}, \dots, c_{1,k_1} \rangle * b_1, \dots, \langle c_{n,1}, \dots, c_{n,k_n} \rangle * b_n \rangle * a \end{aligned}$$

and for $\phi \in \mathbb{F}(n, m), a \in \mathcal{O}_n, \psi_i \in \mathbb{F}(k_i, l_i), b_i \in \mathcal{O}_{k_i}$ for $i \in (m)$, we have

$$\langle \psi_1 \cdot b_1, \dots, \psi_m \cdot b_m \rangle * (\phi \cdot a) = (\langle \psi_1 \dots, \psi_m \rangle \star \phi) \cdot (\langle b_{\phi(1)}, \dots, b_{\phi(m)} \rangle * a)$$

in \mathcal{O}_l , where

$$\langle \psi_1 \dots, \psi_m \rangle \star \phi : (k) = \sum_{j=1}^n (k_{\phi(j)}) \longrightarrow \sum_{i=1}^m (l_i) = (l)$$

is a function such that j -th summand of the domain $(k_{\phi(j)})$ is sent to the $\phi(j)$ -th summand $(l_{\phi(j)})$ of the codomain by the function $\psi_{\phi(j)}$, i.e.

$$\langle j, r \rangle \mapsto \langle \phi(j), \psi_{\phi(j)}(r) \rangle$$

for $j \in [n]$ and $r \in (k_{\phi(j)})$. This definition refers to the obvious lexicographic order on both $(k) = \sum_{j=1}^n (k_{\phi(j)})$ and $(l) = \sum_{i=1}^m (l_i)$. It is an extension of the multiplication operation in the operad of symmetries.

$$\begin{array}{ccc} \sum_{j=1}^n (k_{\phi(j)}) & \xrightarrow{\langle \psi_1 \dots, \psi_m \rangle \star \phi} & \sum_{i=1}^m (l_i) \\ \kappa_j \uparrow & & \uparrow \lambda_{\phi(j)} \\ (k_{\phi(j)}) & \xrightarrow{\psi_{\phi(j)}} & (l_{\phi(j)}) \end{array}$$

Remark The operation \star is defined on a family of functions indexed by $k, l \in \omega$ between sets:

$$\coprod_{n,m \in \omega, \phi \in \mathbb{F}(n,m)} \coprod_{k_i, l_i \in \omega, k = \sum_{j=1}^n k_{\phi(j)}, l = \sum_{i=1}^m l_i} \prod_{i=1}^m \mathbb{F}(k_i, l_i) \xrightarrow{\star} \mathbb{F}(k, l)$$

A morphism of full operads $f : \mathcal{O} \rightarrow \mathcal{O}'$ is a function that respects arities of operations, unit, compositions, and the actions of functions from \mathbb{F} .

The operation \star , when applied to morphisms in \mathbb{S} , returns a morphism in \mathbb{S} . Therefore these definitions make sense if we restrict morphisms to surjections, i.e. morphisms in \mathbb{S}^1 . In this way, we obtain the notion of a *regular operad*, a *morphism of regular operads*, and the whole category of regular operads denoted **RegOp**.

We have functors

$$\mathbf{FOp} \xrightarrow{\mathcal{Q}_r^o} \mathbf{RegOp} \xrightarrow{\mathcal{Q}_a^o} \mathbf{SOp}$$

‘restricting actions’ along the inclusions $\mathbb{B} \rightarrow \mathbb{S} \rightarrow \mathbb{F}$. They have left adjoints:

$$\mathbf{FOp} \xleftarrow{\mathcal{P}_f^o} \mathbf{RegOp} \xleftarrow{\mathcal{P}_r^o} \mathbf{SOp}$$

We sketch the definitions of those functors below.

For a symmetric operad \mathcal{O} the regular operad $\mathcal{P}_r^o(\mathcal{O})$ has

$$\coprod_{k \in \omega} \mathbb{S}(k, n) \otimes_k \mathcal{O}_k$$

as the set of n -ary operations. The unit of $\mathcal{P}_r^o(\mathcal{O})$ is

$$[1, \iota] \in \mathbb{S}(1, 1) \otimes_1 \mathcal{O}_1 \subset \coprod_{k \in \omega} \mathbb{S}(k, 1) \otimes_k \mathcal{O}_k$$

¹Clearly, these definitions make also sense if we restrict morphisms in \mathbb{F} to injections, i.e. morphisms in \mathbb{I} . But we will study such operads elsewhere.

The action of the category \mathbb{S} in $\mathcal{P}_r^o(\mathcal{O})$

$$\cdot : \mathbb{S}(n, m) \times \coprod_{k \in \omega} \mathbb{S}(k, n) \otimes_k \mathcal{O}_k \longrightarrow \coprod_{k \in \omega} \mathbb{S}(k, m) \otimes_k \mathcal{O}_k$$

is

$$\phi \cdot [f, a] = [\phi \circ f, a]$$

where $\phi : [n] \rightarrow [m]$, $f : [k] \rightarrow [n]$ are surjection and $a \in \mathcal{O}_k$. The composition in $\mathcal{P}_r^o(\mathcal{O})$

$$\begin{aligned} & \coprod_{l_1 \in \omega} \mathbb{S}(l_1, n_1) \otimes_{l_1} \mathcal{O}_{l_1} \times \dots \times \coprod_{l_m \in \omega} \mathbb{S}(l_m, n_m) \otimes_{l_m} \mathcal{O}_{l_m} \\ & \times \coprod_{k \in \omega} \mathbb{S}(k, m) \otimes_k \mathcal{O}_k \longrightarrow \coprod_{k \in \omega} \mathbb{S}(k, n) \otimes_k \mathcal{O}_k \end{aligned}$$

is given by

$$\langle [\psi_1, b_1], \dots, [\psi_m, b_m], [\phi, b] \rangle \mapsto \langle [\psi_1, \dots, \psi_m] \star \phi, \langle b_{\phi(1)}, \dots, b_{\phi(k)} \rangle \star a \rangle$$

where $\varphi \in \mathbb{S}(k, m)$, $a \in \mathcal{O}_k$, $\sum_{i=1}^m n_i = n$, $\psi_i \in \mathbb{S}(l_i, n_i)$, $b_i \in \mathcal{O}_{l_i}$, for $i \in [m]$. The definition of the functor \mathcal{P}_a^o on morphisms is left for the reader.

For a regular operad \mathcal{O} the full operad $\mathcal{P}_f^o(\mathcal{O})$ has

$$\coprod_{k \in \omega} \mathbb{I}(k, n) \otimes_k \mathcal{O}_k$$

as the set of n -ary operations. The action of the category \mathbb{F} in $\mathcal{P}_f^o(\mathcal{O})$

$$\cdot : \mathbb{F}(n, m) \times \coprod_{k \in \omega} \mathbb{I}(k, n) \otimes_k \mathcal{O}_k \longrightarrow \coprod_{k \in \omega} \mathbb{I}(k, n) \otimes_k \mathcal{O}_k$$

is

$$\phi \cdot [f, a] = [f', \phi' \cdot a]$$

where $\phi : [n] \rightarrow [m]$ is a function, $f : [k] \rightarrow [n]$ is an injection, $a \in \mathcal{O}_k$ and ϕ', f' is the epi-mono factorization of $\phi \circ f$:

$$\begin{array}{ccc} [k] & \xrightarrow{f} & [n] \\ \phi' \downarrow & & \downarrow \phi \\ [n'] & \xrightarrow{f'} & [m] \end{array}$$

The rest of the definition of $\mathcal{P}_f^o(\mathcal{O})$ is similar to the remaining part of the definition of $\mathcal{P}_r^o(\mathcal{O})$ given above.

2.3 Lawvere Theories

The category of Lawvere theories will be denoted by **LT**, see [7, 8, 14] for details. \mathbb{F}^{op} is the initial Lawvere theory with the obvious projections. Let **T** be any Lawvere theory. By $\pi : \mathbb{F}^{op} \rightarrow \mathbf{T}$ we denote the unique morphism from \mathbb{F}^{op} to **T**. The class of *projections* in **T** is the closure under isomorphisms of the image of the injections in \mathbb{F} . A morphism r in **T** is *regular* iff r is right orthogonal to all projection morphisms in **T**. By a factorization system we mean a factorization system in the sense of [4], see [3] sec 2.8.

$Aut(n)$ is the set of automorphisms of n in **T**. Recall from [14] that a Lawvere theory has *simple automorphisms* if the canonical function

$$\rho_n : S_n \times Aut(1)^n \longrightarrow Aut(n)$$

such that

$$(\sigma, a_1, \dots, a_n) \mapsto a_1 \times \dots \times a_n \circ \pi_\sigma$$

is a bijection, for $n \in \omega$.

A Lawvere theory **T** is a *regular Lawvere theory* iff the projection morphisms and the regular morphisms form a factorization system and **T** has simple automorphisms. A *regular interpretation* of Lawvere theories is a morphism of Lawvere theories that preserves regular morphisms. Thus we have a non-full subcategory of **LT** of regular Lawvere theories with regular interpretations **RegLT**. The theory \mathbb{F}^{op} is regular.

We have inclusion functors

$$\mathbf{AnLT} \xrightarrow{\mathcal{P}_r^l} \mathbf{RegLT} \xrightarrow{\mathcal{P}_f^l} \mathbf{LT}$$

2.4 Monads

We introduce here a notion of a semi-analytic monad that is broader than that of an analytic monad but it still retains some combinatorial flavor.

Recall that an analytic (endo)functor on *Set* can be defined by any of the following conditions

1. finitary functor preserving weak wide pullbacks;
2. Kan extension of a functor from \mathbb{B} to *Set* along the inclusion functor $\mathbb{B} \rightarrow \mathbf{Set}$;
3. functor (isomorphic to one) having an analytic presentation $\sum_{n \in \omega} X^n \otimes_n A_n$, where the n -coefficient A_n is a (left) S_n -set for $n \in \omega$.

Similarly, we shall define a *semi-analytic* (endo)functor on *Set* as a functor satisfying any of three equivalent conditions (see Proposition 2.1, Theorem 2.2)

1. finitary functor preserving pullbacks along monomorphisms;
2. Kan extension of a functor from \mathbb{S} to *Set* along the inclusion functor $\mathbf{i}_\mathbb{S} : \mathbb{S} \rightarrow \mathbf{Set}$;
3. functor (isomorphic to one) having a semi-analytic presentation $\sum_{n \in \omega} \left[\begin{smallmatrix} X \\ n \end{smallmatrix} \right] \otimes_n A_n$, where the category \mathbb{S} acts on coefficients A_n on the left (see below).

The meaning of the first two characterizations of the semi-analytic functors is clear. We shall describe the third one and show that it is equivalent to the other two.

For the time being, the first definition of a semi-analytic functor is the official one. A natural transformation $\phi : F \rightarrow G$ is *semi-cartesian* iff the naturality squares for monomorphisms are pullbacks. The category of semi-analytic functors with the semi-cartesian natural transformations will be denoted by **San**.

Examples

1. The functor $\mathcal{P}_{\leq n} : Set \rightarrow Set$ associating to a set X the set of subsets of X with at most n -elements is not analytic if $n > 2$, as it can be easily seen that it does not preserve weak pullbacks. However, it preserves pullbacks along monos and hence it is semi-analytic.
2. If U is a set, $n \in \omega$, then the functor $(-)^U_{\leq n} : Set \rightarrow Set$, associating to a set X the set of functions from U to X with an at most n -element image, is not analytic, if $|U| > n > 2$. Again, it can be easily seen that it does not preserve weak pullbacks. However, it is semi-analytic.
3. We will see later that the functor part of any monad on Set that comes from a regular equational theory is semi-analytic.

For a set X and $n \in \omega$ the set

$$\left[\begin{matrix} X \\ n \end{matrix} \right]$$

denotes the set of monomorphisms from (n) to X . If X has less than n elements, this set is empty and if X is a finite set, then it has $\binom{|X|}{n} \cdot n!$ elements. The notation is in analogy with the notation

$$\binom{X}{n}$$

denoting the set of n -element subsets of X .

Clearly, $\left[\begin{matrix} X \\ n \end{matrix} \right]$ is not functorial in X on its own but if we build a series with such sets and coefficients that are related by surjections, we do get a functor. To see this let

$$A : \mathbb{S} \rightarrow Set$$

be a functor or, equivalently, a sequence of coefficient sets $\{A_n\}_{n \in \omega}$ on which the category \mathbb{S} acts on the left. As S_n acts on $\left[\begin{matrix} X \\ n \end{matrix} \right]$ on the right, by composition it makes sense to form a set

$$\left[\begin{matrix} X \\ n \end{matrix} \right] \otimes_n A_n$$

and a whole coproduct

$$\hat{A}(X) = \sum_{n \in \omega} \left[\begin{matrix} X \\ n \end{matrix} \right] \otimes_n A_n$$

The above formula is functorial in X , i.e. \hat{A} can be defined on morphisms as follows. Let $f : X \rightarrow Y$ be a function and $[\vec{x}, a]$ an element of $\left[\begin{smallmatrix} X \\ n \end{smallmatrix} \right] \otimes_n A_n$. We take the epi-mono factorization α, \vec{y} of $f \circ \vec{x}$

$$\begin{array}{ccc} (n) & \xrightarrow{\vec{x}} & X \\ \alpha \downarrow & & \downarrow f \\ (m) & \xrightarrow{\vec{y}} & Y \end{array}$$

and we put

$$\hat{A}(f)([\vec{x}, a]) = [\vec{y}, A(\alpha)(a)]$$

which is an element in $\left[\begin{smallmatrix} Y \\ m \end{smallmatrix} \right] \otimes_m A_m$, i.e. in $\hat{A}(Y)$. As the factorization is unique up to isomorphism, $\hat{A}(f)$ is well defined.

If $\tau : A \rightarrow B$ is a natural transformation of functors, it induces a natural transformation of functors

$$\hat{\tau} : \hat{A} \longrightarrow \hat{B}$$

as follows. For $[\vec{x}, a]$ in $\left[\begin{smallmatrix} X \\ n \end{smallmatrix} \right] \otimes_n A_n$ we put

$$\hat{\tau}([\vec{x}, a]) = [\vec{x}, \tau(a)]$$

Proposition 2.1 *The functor $(\hat{-}) : Set^{\mathbb{S}} \rightarrow \mathbf{End}$ is well defined and it is isomorphic to the left Kan extension along the inclusion functor $i_{\mathbb{S}} : \mathbb{S} \rightarrow Set$.*

Proof It is well known (see [9]) that Kan extensions can be calculated using coends. Thus, for a functor $A : \mathbb{S} \rightarrow Set$ and a function $f : X \rightarrow Y$, we have

$$\begin{aligned} Lan_{i_{\mathbb{S}}}(A)(X) &= \int^{n \in \mathbb{S}} X^n \times A_n \xrightarrow{Lan_{i_{\mathbb{S}}}(A)(f)} \int^{n \in \mathbb{S}} Y^n \times A_n = Lan_{i_{\mathbb{S}}}(A)(Y) \\ &\| \vec{x}, a \| \mapsto \| f \circ \vec{x}, a \| \end{aligned}$$

where $\| \vec{x}, a \|$ is the equivalence class of the equivalence relation \approx on $\sum_{n \in \omega} X^n \times A_n$ generated by the relation \sim such that

$$\langle \vec{y} \circ \phi, a \rangle \sim \langle \vec{y}, \phi \cdot a \rangle$$

for $\vec{y} : (m) \rightarrow X, \phi : (n) \rightarrow (m) \in \mathbb{S}$ and $a \in A_n$. Let us call a representant $\langle \vec{x}, a \rangle$ of a class $\| \vec{x}, a \|$ *minimal* iff \vec{x} is injective. Any element $\langle \vec{x}, a \rangle$ is \sim -related to one of form $\langle \vec{y}, f \cdot a \rangle$ where f, \vec{y} is the surjective-injective factorization of \vec{x} , i.e. any element is \sim -related to a minimal one. It is easy to see that any two minimal representatives of \approx -equivalence class are \sim -related.

We define an isomorphism

$$\kappa_A : \hat{A} \longrightarrow Lan_{is}(A)$$

so that on a set X

$$(\kappa_A)_X : \hat{A}(X) \longrightarrow Lan_{is}(A)(X)$$

we put

$$[\vec{x}, a] \mapsto \|\vec{x}, a\|$$

where $\vec{x} : (n) \rightarrow X$ and $a \in A_n$. Taking the minimal elements in the equivalence class defines the inverse function. Thus $(\kappa_A)_X$ is a bijection. It is easy to see that κ_A defined this way is natural in X , i.e. it is a natural isomorphism.

For a natural transformation $\tau : A \rightarrow B$ in Set^S the square

$$\begin{CD} \hat{A} @<\tau_A<< Lan_{is}(A) \\ @V\hat{\tau}VV @VV Lan_{is}(\tau)V \\ \hat{B} @<\tau_B<< Lan_{is}(B) \end{CD}$$

obviously commutes, as both compositions of an element $[\vec{x}, a] \in \hat{A}(X)$ are equal to $\|\vec{x}, \tau(a)\|$. Thus

$$\kappa : (\hat{-}) \longrightarrow Lan_{is}$$

is also a natural isomorphism, as required. □

Theorem 2.2 *The functor $(\hat{-}) : Set^S \rightarrow \mathbf{End}$ is faithful, full on isomorphisms, and its essential image is the category of semi-analytic functors **San**.*

The theorem will be proved via a series of lemmas.

Lemma 2.3 *The essential image of the functor $(\hat{-}) : Set^S \rightarrow \mathbf{End}$ is contained in the category **San**.*

Proof First, we check that for $A \in Set^S$, \hat{A} preserves pullbacks along monos. Let

$$\begin{CD} P @>f'>> Y \\ @V\beta'VV @VV\beta V \\ Z @>f>> X \end{CD}$$

be a pullback in *Set* with β mono. We need to show that the square

$$\begin{array}{ccc}
 \sum_{n \in \omega} \begin{bmatrix} P \\ n \end{bmatrix} \otimes_n A_n & \xrightarrow{\hat{A}(f')} & \sum_{n \in \omega} \begin{bmatrix} Y \\ n \end{bmatrix} \otimes_n A_n \\
 \hat{A}(\beta') \downarrow & & \downarrow \hat{A}(\beta) \\
 \sum_{n \in \omega} \begin{bmatrix} Z \\ n \end{bmatrix} \otimes_n A_n & \xrightarrow{\hat{A}(f)} & \sum_{n \in \omega} \begin{bmatrix} X \\ n \end{bmatrix} \otimes_n A_n
 \end{array}$$

is also a pullback in *Set*. Let $[\bar{y}, a] \in \begin{bmatrix} Y \\ n \end{bmatrix} \otimes_n A_n$ and $[\bar{z}, b] \in \begin{bmatrix} Z \\ m \end{bmatrix} \otimes_m A_m$ be such that

$$\hat{A}(f)([\bar{z}, b]) = \hat{A}(\beta)([\bar{y}, a])$$

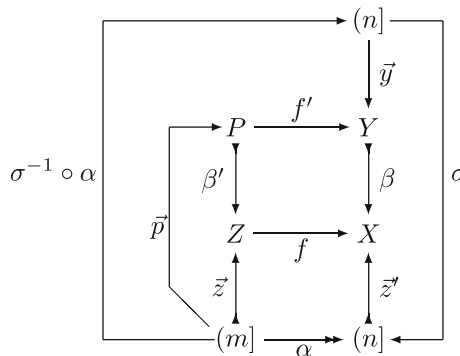
Let $\alpha : [m] \rightarrow [n']$ and $\bar{z}' : [n'] \rightarrow X$ be a surjection-injection factorization of the function $f \circ \bar{z} : [m] \rightarrow X$. From the above equation follows that $n' = n$ and there is a $\sigma \in S_n$ such that

$$\bar{z}' \circ \sigma = \beta \circ \bar{y}, \quad A(\alpha)(b) = A(\sigma)(a)$$

We have

$$f \circ \bar{z} = \bar{z}' \circ \alpha = \beta \circ \bar{y} \circ \sigma^{-1} \circ \alpha$$

Hence there is a function $\vec{p} : [m] \rightarrow P$ as in the following diagram



such that

$$\beta' \circ \vec{p} = \bar{z}, \quad f' \circ \vec{p} = \bar{y} \circ \sigma^{-1} \circ \alpha$$

The function \vec{p} is an injection since \bar{z} is. Thus $[\vec{p}, b]$ is an element of $\begin{bmatrix} P \\ m \end{bmatrix} \otimes_m A_m$.

We have

$$\hat{A}(\beta')([\vec{p}, b]) = [\beta' \circ \vec{p}, b] = [\bar{z}, b]$$

Moreover, on representatives we have

$$\begin{aligned} \langle f' \circ \vec{p}, b \rangle &= \langle \vec{y} \circ \sigma^{-1} \circ \alpha, b \rangle \sim \\ &\sim \langle \vec{y}, A(\sigma^{-1} \circ \alpha)(b) \rangle = \langle \vec{y}, A(\sigma^{-1})(A(\alpha)(b)) \rangle = \\ &= \langle \vec{y}, A(\sigma^{-1})(A(\sigma)(a)) \rangle = \langle \vec{y}, a \rangle \end{aligned}$$

and hence

$$\hat{A}(f')([\vec{p}, b]) = [f' \circ \vec{p}, b] = [\vec{y}, a]$$

As \hat{A} preserves monos, such an element $[\vec{p}, b]$ is unique and hence \hat{A} preserves pullbacks along monos.

Now let $\tau : A \rightarrow B$ be a morphism in Set^S . We shall show that $\hat{\tau} : \hat{A} \rightarrow \hat{B}$ is semi-cartesian. Let $f : X \rightarrow Y$ be an injection in Set . We need to show that the square

$$\begin{array}{ccc} \hat{A}(X) & \xrightarrow{\hat{\tau}_X} & \hat{B}(X) \\ \hat{A}(f) \downarrow & & \downarrow \hat{B}(f) \\ \hat{A}(Y) & \xrightarrow{\hat{\tau}_Y} & \hat{B}(Y) \end{array}$$

is a pullback. Let $[\vec{y}, a] \in \hat{A}(Y)$ and $[\vec{x}, b] \in \hat{B}(X)$ be such that

$$[\vec{y}, \tau_n(a)] = \hat{\tau}_Y([\vec{y}, a]) = \hat{B}(f)([\vec{x}, b]) = [f \circ \vec{x}, b]$$

where $a \in A_n$. Note that as f is an injection, so is $f \circ \vec{x}$ and hence $\langle f \circ \vec{x}, b \rangle$ represents an element in $\hat{B}(Y)$. The above equality means that $b \in B_n$ and there is a $\sigma \in S_n$ such that

$$\vec{y} \circ \sigma = f \circ \vec{x}, \quad \tau_n(a) = B(\sigma)(b)$$

We have an element $[\vec{x} \circ \sigma^{-1}, a]$ such that

$$\hat{A}(f)([\vec{x} \circ \sigma^{-1}, a]) = [f \circ \vec{x} \circ \sigma^{-1}, a] = [\vec{y}, a]$$

On representatives we have

$$\langle \vec{x} \circ \sigma^{-1}, \tau_n(a) \rangle \sim \langle \vec{x}, \hat{B}(\sigma^{-1}(\tau_n(a))) \rangle = \langle \vec{x}, b \rangle$$

and hence we have

$$\hat{\tau}_X([\vec{x} \circ \sigma^{-1}, a]) = [\vec{x} \circ \sigma^{-1}, \tau_n(a)] = [\vec{x}, b]$$

as well. This means that $\hat{\tau}$ is semi-cartesian indeed. □

Lemma 2.4 *The functor $(\hat{-}) : Set^S \rightarrow \mathbf{End}$ is faithful and full on semi-cartesian morphisms. In particular, it is full on isomorphisms.*

Proof One can easily verify that if two natural transformations $\hat{\tau}, \hat{\tau}' : \hat{A} \rightarrow \hat{B}$ agree on elements of the form $[1_{(m)}, a]$ for $m \in \omega$ and $a \in A_m$, then the natural transformations $\tau, \tau' : A \rightarrow B$ are equal. Thus $(\hat{-})$ is faithful.

To show that $(\hat{-})$ is full, let us fix two functors $A, B \in \text{Set}^{\mathcal{S}}$ and a semi-cartesian natural transformation $\psi : \hat{A} \rightarrow \hat{B}$. Let $m \in \omega$ and $a \in A_m$. We shall define $\tau_m(a)$. We have $\psi_{(m)}([1_{(m)}, a]) = [f, b] \in \hat{B}(m)$ for some injection $f : (k) \rightarrow (m)$ and $b \in B_k$. We claim that $k = m$ and f is a bijection. Suppose to the contrary that $k < m$. We have that $\hat{B}(f)([1_{(k)}, b]) = [f, b]$. Since f is an injection and ψ is semi-cartesian, the square

$$\begin{array}{ccc} \hat{A}(m) & \xrightarrow{\psi_{(m)}} & \hat{B}(m) \\ \hat{A}(f) \uparrow & & \uparrow \hat{B}(f) \\ \hat{A}(k) & \xrightarrow{\psi_{(k)}} & \hat{B}(k) \end{array}$$

is a pullback and hence there is an element $[g, c] \in \hat{A}(k)$ such that

$$\hat{A}(f)([g, c]) = [1_{(m)}, a], \quad \psi_{(m)}([g, c]) = [1_{(k)}, b]$$

The first equality implies that the proper mono $f \circ g$ is epi (as its codomain is (m) and its mono part is $1_{(m)}$). This is a contradiction. Hence f is a bijection and we can apply the functor B to it. We put

$$\tau_m(a) = B(f)(b)$$

Thus we have defined a transformation $\tau : A \rightarrow B$.

To show that τ is natural, let us fix a surjection $\beta : (m) \rightarrow (n)$ and $a \in A_m$. Then using the definitions of \hat{A}, \hat{B} and the naturality of ψ we have

$$\begin{aligned} [1_{(n)}, B(f)(\tau_m(a))] &= \hat{B}(f)([1_{(m)}, \tau_m(a)]) = \\ &= \hat{B}(f) \circ \psi_{(m)}([1_{(m)}, a]) = \psi_{(n)} \circ \hat{A}(f)([1_{(m)}, a]) = \\ &= \psi_{(n)}([1_{(n)}, \hat{A}(f)(a)]) = [1_{(n)}, \tau_n(\hat{A}(f)(a))] \end{aligned}$$

But this means that

$$B(f)(\tau_m(a)) = \tau_n(\hat{A}(f)(a))$$

Since a and f were arbitrary, τ is natural.

Finally, we show that $\hat{\tau} = \psi$. Let us fix a set X and $[\vec{x}, a] \in \hat{A}(X)$, where $\vec{x} : (k) \rightarrow X$ and $a \in A_k$. Note that $\hat{A}(\vec{x})([1_{(k)}, a]) = [\vec{x}, a]$. Using the naturality of ψ and $\hat{\tau}$ on \vec{x} and the definition of τ we have

$$\begin{aligned} \psi_X([\vec{x}, a]) &= \psi_X \circ \hat{A}(\vec{x})([1_{(k)}, a]) = \\ &= \hat{B}(\vec{x}) \circ \psi_{(k)}([1_{(k)}, a]) = \hat{B}(\vec{x})([1_{(k)}, \tau_k(a)]) = \\ &= \hat{B}(\vec{x}) \circ \hat{\tau}_{(k)}([1_{(k)}, a]) = \hat{\tau}_X \circ \hat{A}(\vec{x})([1_{(k)}, a]) = \\ &= \hat{\tau}_X([\vec{x}, a]) \end{aligned}$$

Since X and $[\vec{x}, a]$ were arbitrary, $\hat{\tau} = \psi$. □

Lemma 2.5 *Each semi-cartesian functor is in the essential image of $(\hat{\quad}) : \text{Set}^{\mathbb{S}} \rightarrow \mathbf{End}$.*

Proof Let us fix a semi-cartesian functor F . We will define a functor $A : \mathbb{S} \rightarrow \text{Set}$ and a natural isomorphism $\tau : \hat{A} \rightarrow F$. Put

$$A(n) = F(n) - \bigcup_{(m) \xrightarrow{f} (n) \in \mathbb{I}, m < n} F(f)(A(m))$$

Note that the sum over the empty index set is empty. For $\alpha : [n] \rightarrow [m]$ in \mathbb{S} we put

$$A(\alpha) = F(\alpha) \upharpoonright_{A[n]} : A[n] \rightarrow A[m]$$

We need to show that $A(\alpha)$ is well defined, i.e. for $a \in A(n)$ we have $F(\alpha)(a) \in A(m)$. Clearly, $F(\alpha)(a) \in F(m)$. Suppose to the contrary that $F(\alpha)(a) \notin A(m)$. Then there is a proper mono $f : (k) \rightarrow (m)$ and $b \in A(k)$ such that $F(\alpha)(a) = F(f)(b)$. We form a pullback of a surjection α and a proper mono f

$$\begin{array}{ccc} (p) & \xrightarrow{\alpha'} & (k) \\ f' \downarrow & & \downarrow f \\ (n) & \xrightarrow{\alpha} & (m) \end{array}$$

Thus f' is again a proper mono. F preserves this pullback, i.e. the square

$$\begin{array}{ccc} F(p) & \xrightarrow{F(\alpha')} & F(k) \\ F(f') \downarrow & & \downarrow F(f) \\ F(n) & \xrightarrow{F(\alpha)} & F(m) \end{array}$$

is again a pullback. Thus there is an element $c \in F(p)$ such that

$$F(\alpha')(c) = b, \quad F(f')(c) = a$$

The latter equality means that $a \notin A(n)$ contrary to the supposition. Thus $A : \mathbb{S} \rightarrow \text{Set}$ is a well defined functor.

For a set X we define the function

$$\sum_{n \in \omega} \left[\begin{array}{c} X \\ n \end{array} \right] \otimes_n A_n \xrightarrow{\tau_X} F(X)$$

by

$$\tau_X([\vec{x}, a]) = F(\vec{x})(a)$$

where $n \in \omega$, $\vec{x} : (n) \rightarrow X$ is an injection and $a \in A(n)$.

First, we show that the transformation $\tau : \hat{A} \rightarrow F$ is natural. Let $f : X \rightarrow Y$ be a function. Then we have an epi/mono factorization of f', \vec{y} of the function $f \circ \vec{x}$

$$\begin{array}{ccc} (n) & \xleftarrow{\vec{x}} & X \\ f' \downarrow & & \downarrow f \\ (m) & \xleftarrow{\vec{y}} & Y \end{array}$$

Thus we have

$$\begin{aligned} F(f)\tau_X([\vec{x}, a]) &= F(f \circ \vec{x})(a) = \\ &= F(\vec{y} \circ f')(a) = F(\vec{y})(F(f')(a)) = \\ &= F(\vec{y})(A(f')(a)) = \tau_Y([\vec{y}, A(f')(a)]) = \\ &= \tau_Y \hat{A}(f)([\vec{x}, a]) \end{aligned}$$

i.e. τ is natural.

It remains to show that τ is an isomorphism. Fix a set X . Let $x \in F(X)$. As F is finitary, there is an $n \in \omega$, $f : (n) \rightarrow X$ and $y \in F(n)$ such that $F(f)(y) = x$. Let α, g be an epi/mono factorization of f

$$\begin{array}{ccc} (n) & \xrightarrow{f} & X \\ \alpha \searrow & & \nearrow g \\ & (m) & \end{array}$$

Then $[g, A(\alpha)(y)] \in \begin{bmatrix} X \\ m \end{bmatrix} \otimes_m A_m$ and

$$\begin{aligned} \tau_X([g, A(\alpha)(y)]) &= F(g)(A(\alpha)(y)) = \\ &= F(g \circ \alpha)(y) = F(f)(y) = x \end{aligned}$$

Thus τ_X is onto.

Now suppose that $[\vec{x}, a] \in \begin{bmatrix} X \\ n \end{bmatrix} \otimes_n A_n$ and $[\vec{y}, b] \in \begin{bmatrix} X \\ m \end{bmatrix} \otimes_m A_m$ and

$$\tau_X([\vec{x}, a]) = \tau_X([\vec{y}, b])$$

We form a pullback of monos

$$\begin{array}{ccc} (n) & \xrightarrow{\vec{x}} & X \\ f \uparrow & & \uparrow \vec{y} \\ (k) & \xrightarrow{g} & (m) \end{array}$$

that F preserves, i.e. we have a pullback

$$\begin{array}{ccccc}
 A(n) & \longrightarrow & F(n) & \xrightarrow{F(\vec{x})} & F(X) \\
 & & \uparrow F(f) & & \uparrow F(\vec{y}) \\
 & & F(k) & \xrightarrow{F(g)} & F(m) \\
 & & & & \uparrow \\
 & & & & A(m)
 \end{array}$$

Since

$$F(\vec{x})(a) = \tau_X([\vec{x}, a]) = \tau_X([\vec{y}, b]) = F(\vec{y})(b)$$

there is a $c \in F(k)$ such that

$$F(f)(c) = a, \quad F(g)(c) = b$$

As $a \in A(n)$ and $b \in B(m)$, we must have that $n = k = m$ and both f and g are bijections. Put $\sigma = f \circ g^{-1}$. Then

$$\vec{x} \circ \sigma = \vec{y}, \quad a = A(\sigma)(b)$$

which means that $[\vec{x}, a] = [\vec{y}, b]$. Thus τ_X is an injection, as required. □

A monad (M, η, μ) on *Set* is a *semi-analytic monad* iff M is a semi-analytic functor and both η and μ are semi-cartesian natural transformations. The category of semi-analytic monads with the semi-cartesian morphisms will be denoted by **SanMnd**.

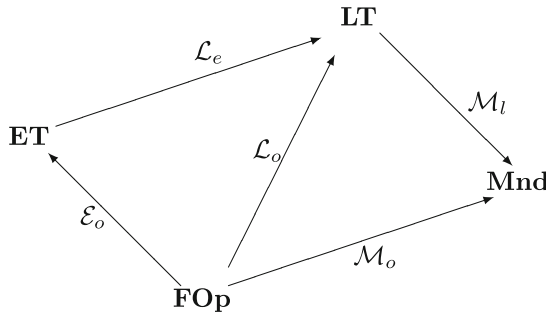
We have inclusion functors

$$\mathbf{AnMnd} \xrightarrow{\mathcal{P}_r^m} \mathbf{SanMnd} \xrightarrow{\mathcal{P}_f^m} \mathbf{Mnd}$$

3 Full Operads

In this section we describe the equivalence of the category of full operads **FOp** to categories of equational theories **ET**, Lawvere theories **LT** and monads **Mnd**. As we said, the categories **FOp** and **Mnd** are categories of monoids of two equivalent

monoidal categories, thus they are obviously equivalent. Therefore, we will just define the equivalences of categories



with the domain **FOp** and their essential inverses leaving all the verifications to the reader. In [13] it was shown that the category of full operads is equivalent to the category of abstract clones. The latter category is known to be equivalent to **ET**.

Let us fix a morphism $h : \mathcal{O} \rightarrow \mathcal{O}'$ of **F**-operads. We denote by $\kappa_i^n : \{1\} \rightarrow \{n\}$ the function that sends 1 to i , where $n \in \omega$ and we denote π_i^n the action of κ_i^n on the unit $\iota \in \mathcal{O}_1$, i.e. $\pi_i^n = \kappa_i^n \cdot \iota$.

3.1 The Functor \mathcal{E}_o

The equational theory $\mathcal{E}_o(\mathcal{O}) = (L, A)$ has as the set of operations $L = \coprod_{n \in \omega} \mathcal{O}_n$. For $f \in \mathcal{O}_k$, $f_i \in \mathcal{O}_{n_i}$, $i = 1, \dots, k$, $n = \sum_{i=1}^k n_i$, and $\psi : \{k\} \rightarrow \{n\}$ we have the following axioms in A :

$$f(f_1(x_1, \dots, x_{n_1}), \dots, f_k(x_{n-n_k+1}, \dots, x_n)) = ((f_1, \dots, f_k) \star f)(x_1, \dots, x_n) : \vec{x}^n$$

$$f(x_{\psi(1)}, \dots, x_{\psi(k)}) = (\psi \cdot f)(x_1, \dots, x_n) : \vec{x}^n$$

and

$$\iota(x_1) = x_1 : \vec{x}^1$$

The interpretation $\mathcal{E}_o(h) : \mathcal{E}_o(\mathcal{O}) \rightarrow \mathcal{E}_o(\mathcal{O}')$ sends a function symbol $f \in \mathcal{O}_n$ to the term $h(f)(x_1, \dots, x_n) : \vec{x}^n$. This ends the definition of the functor \mathcal{E}_o .

The essential inverse of the functor \mathcal{E}_o

$$\mathcal{O}p_e : \mathbf{ET} \longrightarrow \mathbf{FOp}$$

is defined as follows. For an equational theory $T = (L, A)$, the **F**-operad $\mathcal{Q} = \mathcal{O}p_e(T)$ has as n -ary operations \mathcal{Q}_n the terms in context \vec{x}^n modulo provability in T . The action

$$\cdot : \mathbb{F}(n, m) \times \mathcal{Q}_n \longrightarrow \mathcal{Q}_m$$

is defined as

$$\phi \cdot [t(x_1, \dots, x_n) : \vec{x}^n] = [t(x_{\phi(1)}, \dots, x_{\phi(n)}) : \vec{x}^m]$$

where $[t(x_1, \dots, x_n) : \vec{x}^n]$ is an operation in a \mathcal{Q}_n and $\phi : (n) \rightarrow (m)$ is a function. The composition in \mathcal{Q}

$$\star : \mathcal{Q}_{n_1} \times \dots \times \mathcal{Q}_{n_k} \times \mathcal{Q}_k \longrightarrow \mathcal{Q}_n$$

where $n = \sum_{i=1}^k n_i$ is defined by substitution with α -conversion, i.e. for $[t_i(x_1, \dots, x_{n_i}) : \vec{x}^{n_i}] \in \mathcal{Q}_{n_i}$ and $[s(x_1, \dots, x_k) : \vec{x}^k] \in \mathcal{Q}_k$ we have

$$\begin{aligned} &([t_1(x_1, \dots, x_{n_1}) : \vec{x}^{n_1}], \dots, [t_k(x_1, \dots, x_{n_k}) : \vec{x}^{n_k}]) \star [s(x_1, \dots, x_k) : \vec{x}^k] = \\ &= [s(t_1(x_{\alpha_1(1)}), \dots, x_{\alpha_1(n_1)}), \dots, t_k(x_{\alpha_k(1)}, \dots, x_{\alpha_k(n_k)})) : \vec{x}^k] \end{aligned}$$

where $\alpha_i : (n_i) \rightarrow (n)$, for $i = 1, \dots, k$ is the obvious embedding.

3.2 The Functor \mathcal{L}_o

A morphism

$$\langle f_1, \dots, f_k \rangle : n \longrightarrow k$$

in the Lawvere theory $\mathcal{L}_o(\mathcal{O})$ is a k -tuple such that $f_i \in \mathcal{O}$, for $i = 1, \dots, k$. The identity on n is

$$\langle \pi_1^n, \dots, \pi_n^n \rangle : n \longrightarrow n$$

and the i -th projection from n is

$$\pi_i^n : n \longrightarrow 1$$

Recall that π_i^n is the value of the action of the function $\kappa_i^n : (1) \rightarrow (n)$ that picks i on the unit t . The composition of morphisms

$$\langle g_1, \dots, g_n \rangle : m \longrightarrow n, \quad \langle f_1, \dots, f_k \rangle : n \longrightarrow k$$

is

$$\langle t_1, \dots, t_k \rangle : m \longrightarrow k$$

where

$$t_i = \psi \cdot ((g_1, \dots, g_n) \star f_i)$$

for $i = 1, \dots, k$, where $\psi : (n \cdot m) \rightarrow (m)$ is given by

$$\psi(r) = ((r - 1) \bmod m) + 1$$

for $r \in (n \cdot m)$. The interpretation $\mathcal{L}_o(h) : \mathcal{L}_o(\mathcal{O}) \rightarrow \mathcal{L}_o(\mathcal{O}')$ sends morphism $\langle f_1, \dots, f_k \rangle : n \longrightarrow k$ to morphism $\langle h(f_1), \dots, h(f_k) \rangle : n \longrightarrow k$. This ends the definition of the functor \mathcal{L}_o .

The essential inverse of the functor \mathcal{L}_o

$$\mathcal{O}_{p_I} : \mathbf{LT} \longrightarrow \mathbf{FOP}$$

is defined as follows. For a Lawvere theory \mathbf{T} the set of n -ary operation of the operad $\mathcal{Q} = \mathcal{O}_{p_I}(\mathbf{T})$ is $\mathcal{Q}_n = \mathbf{T}(n, 1)$. The action in \mathcal{Q} is given by

$$\phi \cdot f = f \circ \pi_\phi$$

where $f \in \mathcal{Q}_n, \phi$ as above, and $\pi : \mathbb{F}^{op} \rightarrow \mathbf{T}$ is the unique morphism from the initial Lawvere theory. The composition in \mathcal{Q} is given by products and composition in \mathbf{T}

$$(f_1, \dots, f_k) \star g = g \circ (f_1 \times \dots, \times f_k)$$

3.3 The Functor \mathcal{M}_o

Let $f : X \rightarrow Y$ be a function. The functor part of the monad $\mathcal{M}_o(\mathcal{O})$ is defined via the coend formula

$$\mathcal{M}_o(\mathcal{O})(X) = \int^{n \in \mathbb{F}} X^n \times \mathcal{O}_n$$

Thus an element of this set is an equivalence class of pairs $\langle \tilde{x}, t \rangle \in X^n \times \mathcal{O}_n$. We put

$$\mathcal{M}_o(\mathcal{O})(f)([\tilde{x}, t]) = [f \circ \tilde{x}, t]$$

The unit of the monad

$$\eta_X^{\mathcal{O}} : X \rightarrow \mathcal{M}(\mathcal{O})(X)$$

is given by

$$x \mapsto [\tilde{x}, t]$$

where $\tilde{x} : (1) \rightarrow X$ sends 1 to x . The multiplication

$$\mu_X^{\mathcal{O}} : \mathcal{M}_o^2(\mathcal{O})(X) = \int^{k, n_1, \dots, n_k \in \omega} X^n \times \mathcal{O}_{n_1} \times \dots \times \mathcal{O}_{n_k} \times \mathcal{O}_k \longrightarrow \mathcal{M}_o(\mathcal{O})(X)$$

where $n = \sum_{i=1}^k n_i$, is given by

$$[\tilde{x}, s_1, \dots, s_k, t] \mapsto [\tilde{x}, (s_1, \dots, s_k) * t]$$

The essential inverse of the functor \mathcal{M}_o

$$\mathcal{O}p_m : \mathbf{Mnd} \longrightarrow \mathbf{FOp}$$

is defined as follows. For a monad $T = (T, \eta, \mu)$ in \mathbf{Mnd} the set of n -ary operations in the operad $\mathcal{Q} = \mathcal{O}p_m(T)$ is $\mathcal{Q}_n = T(n)$. The action on $t \in \mathcal{Q}_n$ is

$$\phi \cdot t = T(\phi)(t)$$

Let $t_i \in T(n_i)$ and $s \in T(k)$. Then we have corresponding morphisms

$$[s] : (1) \longrightarrow T(k), \quad [\vec{t}] : (k) \longrightarrow T(n_1) + \dots + T(n_k)$$

‘picking’ s and the k terms t_1, \dots, t_k . The composition

$$(t_1, \dots, t_k) \star s$$

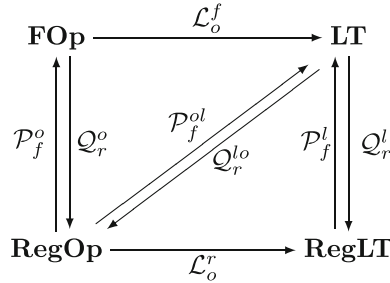
is the element ‘picked’ by the morphism

$$(1) \xrightarrow{[s]} T(k) \xrightarrow{T([\vec{t}]}) T(T(n_1) + \dots + T(n_k)) \xrightarrow{T([T(\alpha_1), \dots, T(\alpha_k)])} T^2(n) \xrightarrow{\mu} T(n)$$

in $T(n)$, where functions α_i are as above.

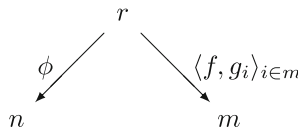
4 Regular Lawvere Theories vs Regular Operads

In this section we study the relations between Lawvere theories and regular operads. We shall describe the adjunction $\mathcal{Q}_r^{lo} \dashv \mathcal{P}_f^{lo}$ and the properties of the embedding \mathcal{P}_f^{lo} .

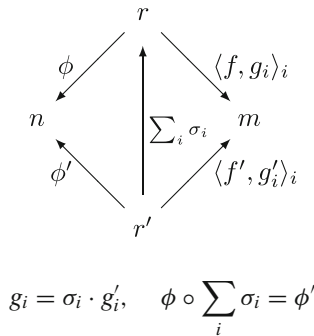


4.1 The Functor $\mathcal{P}_f = \mathcal{P}_f^{ol} : \mathbf{RegOp} \rightarrow \mathbf{LT}$

Let \mathcal{O} be a regular operad; $\iota, \cdot, *$ denote the unit, action (of \mathbb{S}), and composition in \mathcal{O} , respectively. We define a Lawvere theory $\mathcal{P}_f(\mathcal{O})$ as follows. The set of objects of $\mathcal{P}_f(\mathcal{O})$ is the set of natural numbers ω . A morphism from n to m is an equivalence class of spans



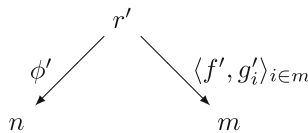
such that $\phi : (r] \rightarrow (n]$ is a function called *amalgamation*, $f : (r] \rightarrow (m]$ is a monotone function called the *arity function* (as it determines the arities of the operations g_i), $r_i = |f^{-1}(i)|$ and we have $g_i \in \mathcal{O}_{r_i}$ for $i \in (m]$ and $r = \sum_{i=1}^m r_i$, moreover ϕ and f are jointly mono (equivalently, ϕ is mono on the fibers of f). Two spans $\langle \phi, f, g_i \rangle_{i \in m}$ and $\langle \phi', f', g'_i \rangle_{j \in m'}$ are equivalent whenever $f = f'$ and there are permutations $\sigma_i : (r_i] \rightarrow (r'_i]$ for $i \in (m]$ such that



In particular, $|f^{-1}(i)| = r_i = |f'^{-1}(i)|$, for $i \in (m]$. By $\sum_i \sigma_i : r \rightarrow r$ we mean the permutation that is formed by placing permutations σ_i ‘one after another’. Thus, it

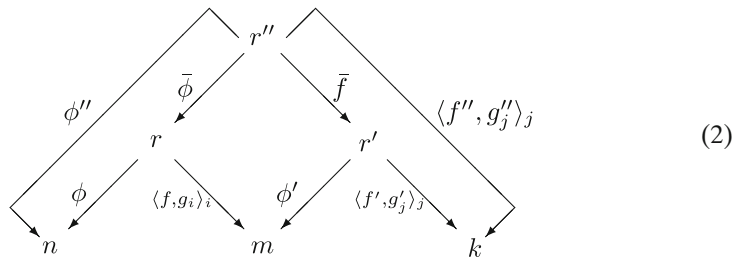
respects the fibers of f , i.e. $f \circ \sum_i \sigma_i = f'$. A morphism in the category $\mathcal{P}_f(\mathcal{O})$ is a class of spans modulo the above equivalence relation.

We could represent morphisms in $\mathcal{P}_f(\mathcal{O})$ as spans without the requirement that ϕ and f are jointly mono. But then the relation identifying the spans would be more complicated. Instead of permutations σ_i we would need to consider surjections. But this relation is not an equivalence relation and we would need to work with the equivalence relation generated by such a relation. However, it might happen, as with the compositions defined below, that an operation on spans naturally leads to a span whose amalgamation ϕ is not injective on fibers of f . In such a case we can *regularize* the span as follows. Let $\langle \phi, f, g_i \rangle_{i \in m}$ be a span as above but with ϕ not necessarily injective on the fibers of f . Let ϕ_i be the restriction of ϕ to the fiber $f^{-1}(i)$ for $i \in (m)$. Let $\phi_i = \phi'_i \circ \psi_i$ be an epi-mono factorization of ϕ_i and $g'_i = \psi_i \cdot g_i$, for $i \in (m)$. Then the *regularization* of the span $\langle \phi, f, g_i \rangle_{i \in m}$ is the span



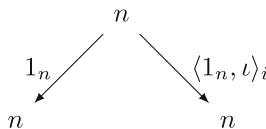
where $r' = \sum_{i \in (m)} r'_i$, $r'_i = |\text{dom}(\phi'_i)|$, for $i \in (m)$, $\phi' = (\phi'_1 + \dots + \phi'_m) : r' \rightarrow n$. f' is the monotone map sending the elements in the domain of ϕ'_i to i , for $i \in (m)$.

The composition of morphisms $\langle \phi, f, g_i \rangle_{i \in (m)} : n \rightarrow m$ and $\langle \phi', f', g'_j \rangle_{j \in (k)} : m \rightarrow k$ to $\langle \phi'', f'', g''_j \rangle_{j \in (k)} : n \rightarrow k$ is defined in two steps as follows. In the diagram

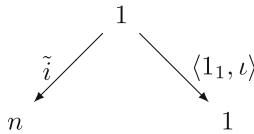


the square is a pullback of f along ϕ' . The function \bar{f} is so chosen that it is monotone. We define a span by $f'' = f' \circ \bar{f}$, $\phi'' = \phi' \circ \bar{\phi}$, and $g''_j = g'_j * \langle g_{\phi'(l)} \rangle_{l \in f'^{-1}(j)}$. Finally, we take a regularization of this span to get a regular span that represents the composition. We leave for the reader the verification that the composition is a congruence with respect to the equivalence relation on spans.

The identity on n is the span



The i -projection $\pi_i^n : n \rightarrow 1$ on i -th coordinate is the span



where $i \in (n]$ and $\tilde{i}(1) = i$.

For a morphism of regular operads $h : \mathcal{O} \rightarrow \mathcal{O}'$, we define a functor

$$\mathcal{P}_f(h) : \mathcal{P}_f(\mathcal{O}) \longrightarrow \mathcal{P}_f(\mathcal{O}')$$

so that for a morphism $\langle \phi, f, g_i \rangle_{i \in (m)} : n \rightarrow m$ in $\mathcal{P}_f(\mathcal{O})$, we have

$$\mathcal{P}_f(h)(\langle \phi, f, g_i \rangle_{i \in (m)}) = \langle \phi, f, h(g_i) \rangle_{i \in (m)} : n \rightarrow m$$

in $\mathcal{P}_f(\mathcal{O}')$. This ends the definition of the functor \mathcal{P}_f .

4.2 The Functor $\mathcal{Q}_r = \mathcal{Q}_r^{\text{lo}} : \mathbf{LT} \longrightarrow \mathbf{RegOp}$

Let \mathbf{T} be a Lawvere theory. Recall that $\pi : \mathbb{F}^{\text{op}} \rightarrow \mathbf{T}$ is the morphism from the initial Lawvere theory. The operad $\mathcal{Q}_r(\mathbf{T})$ consists of the operations of \mathbf{T} , i.e. morphisms to 1. In detail, it can be described as follows. The set of n -ary operations $\mathcal{Q}_r(\mathbf{T})_n$ is the set of n -ary operations $\mathbf{T}(n, 1)$ of \mathbf{T} , for $n \in \omega$. The action

$$\cdot : \mathbb{S}(n, m) \times \mathcal{Q}_r(\mathbf{T})_n \longrightarrow \mathcal{Q}_r(\mathbf{T})_m$$

is given, for $f \in \mathbf{T}(n, 1)$ and $\phi \in \mathbb{S}(n, m)$, by

$$\phi \cdot f = f \circ \pi_\phi$$

The identity of $\mathcal{Q}_r(\mathbf{T})$ is $\iota = 1_1 \in \mathbf{T}(1, 1)$. The composition

$$* : \mathcal{Q}_r(\mathbf{T})_{n_1} \times \mathcal{Q}_r(\mathbf{T})_{n_k} \times \mathcal{Q}_r(\mathbf{T})_k \longrightarrow \mathcal{Q}_r(\mathbf{T})_n$$

is given, for $f \in \mathcal{Q}_r(\mathbf{T})_k$ and $f_i \in \mathcal{Q}_r(\mathbf{T})_{n_i}$, where $i \in (k]$, $n = \sum_{i \in (k]} n_i$, by

$$\langle f_1, \dots, f_k \rangle * f = f \circ (f_1 \times \dots, \times f_k)$$

where $f_1 \times \dots, \times f_k$ is defined using the chosen projections in \mathbf{T} and \circ is the composition in \mathbf{T} .

If $F : \mathbf{T} \rightarrow \mathbf{T}'$ is a morphism of Lawvere theories, then the map of regular operads

$$\mathcal{Q}_r(F) : \mathcal{Q}_r(\mathbf{T}) \rightarrow \mathcal{Q}_r(\mathbf{T}')$$

is defined, for $f \in \mathcal{Q}_r(\mathbf{T})_n$, by

$$\mathcal{Q}_r(F)(f) = F(f)$$

This ends the definition of the functor \mathcal{Q}_r .

4.3 The Adjunction $\mathcal{P}_f \dashv \mathcal{Q}_r$ and the Properties of the Functor \mathcal{P}_f

We note for the record

Proposition 4.1 *The functors $\mathcal{P}_f : \mathbf{RegOp} \rightarrow \mathbf{LT}$ and $\mathcal{Q}_r : \mathbf{LT} \rightarrow \mathbf{RegOp}$ are well defined.*

We have an easy

Lemma 4.2 *Let \mathcal{O} be a regular operad and $n \in \omega$. An isomorphism on n in $P_r(\mathcal{O})$ has a representation by a span of the following form*

$$\begin{array}{ccc} & n & \\ \phi \swarrow & & \searrow \langle 1_n, a_i \rangle_i \\ n & & n \end{array}$$

where $\phi : (n) \rightarrow (n)$ is a bijection, $a_i \in \mathcal{O}_1$ is an invertible operation, i.e. there is $b_i \in \mathcal{O}_1$ such that $a_i * b_i = \iota = b_i * a_i$ for $i \in (n)$. It is the unique span in its equivalence class.

Proof Suppose that we have two spans

$$\begin{array}{ccccc} & r & & r' & \\ \phi \swarrow & & \searrow \langle f, g_i \rangle_i & \swarrow \phi' & \searrow \langle f', h_j \rangle_j \\ n & & n & & n \end{array} \tag{3}$$

representing two morphisms that compose to identity 1_n both ways. Then, since the displayed composition is 1_n , the functions ϕ, f' are surjections. Hence the functions ϕ', f are surjections, as well. Having this, it is easy to notice that ϕ sends elements in different fibers of f to different elements. Thus ϕ (and ϕ') must be in fact a bijection.

If we take a composition $\langle g_i \rangle_{i \in f^{-1}(j)} * h_j$ and regularize it multiplying by a surjection, we will get ι . This implies that h_j cannot be a nullary operation and that the arity of g_i 's is at most one. Now, as the above spans represent morphisms that compose both ways to identity, it is easy to see that we get the required description. \square

Proposition 4.3 *We have an adjunction $\mathcal{P}_f \dashv \mathcal{Q}_r$. The functor \mathcal{P}_f is faithful.*

Proof We shall show that $\mathcal{P}_f \dashv \mathcal{Q}_r$. For a regular operad \mathcal{O} the unit is

$$\begin{aligned} \eta_{\mathcal{O}} : \mathcal{O} &\longrightarrow \mathcal{Q}_r(\mathcal{P}_f(\mathcal{O})) \\ \mathcal{O}_n \ni g &\mapsto \langle 1_n, !, g \rangle \end{aligned}$$

For Lawvere theory \mathbf{T} the counit is

$$\begin{aligned} \varepsilon_{\mathbf{T}} : \mathcal{P}_f \mathcal{Q}_r(\mathbf{T}) &\longrightarrow \mathbf{T} \\ \langle \phi, f, g_i \rangle_{i \in \{m\}} &\mapsto (g_1 \times \dots \times g_m) \circ \pi_{\phi} \end{aligned}$$

We verify the triangular equalities. For $g \in \mathcal{Q}_r(\mathbf{T})_n = \mathbf{T}(n, 1)$ we have

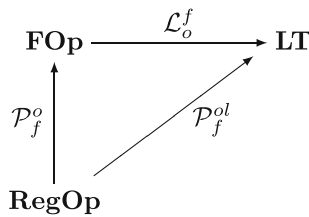
$$\begin{aligned} \mathcal{Q}_r(\varepsilon_{\mathbf{T}}) \circ \eta_{\mathcal{Q}_r(\mathbf{T})}(g) &= \\ &= \mathcal{Q}_r(\varepsilon_{\mathbf{T}})(\langle 1_n, !, g \rangle) = \\ &= g \circ \pi_{1_n} = g \end{aligned}$$

For $\langle \phi, f, g_i \rangle_{i \in \{m\}} \in \mathcal{P}_f(\mathcal{O})$ we have

$$\begin{aligned} \varepsilon_{\mathcal{P}_f(\mathcal{O})} \circ \mathcal{P}_f(\eta_{\mathcal{O}})(\langle \phi, f, g_i \rangle_{i \in \{m\}}) &= \\ &= \varepsilon_{\mathcal{P}_f(\mathcal{O})}(\langle \phi, f, \langle 1_r, !, g_i \rangle \rangle_{i \in \{m\}}) = \\ &= (\langle 1_{r_1}, !, g_1 \rangle \times \dots \times \langle 1_{r_m}, !, g_m \rangle) \circ \pi_{\phi} = \\ &= \langle \phi, f, g_i \rangle_{i \in \{m\}} \end{aligned}$$

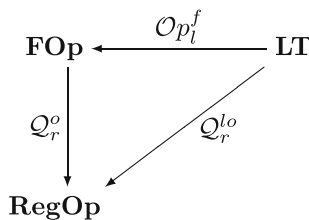
As the unit $\eta_{\mathcal{O}}$ is mono, \mathcal{P}_f is faithful. □

Corollary 4.4 *The triangle*



commutes up to isomorphism.

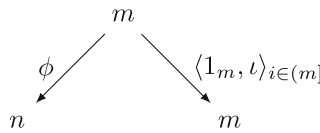
Proof By Proposition 4.3 it is enough to show that triangle



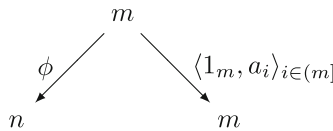
commutes up to isomorphism, where the functor $\mathcal{O}p_1^f$ is the left adjoint to \mathcal{L}_o^f and together they form an adjoint equivalence. The functor $\mathcal{O}p_1^f$ is defined as the functor $\mathcal{Q}_r = \mathcal{Q}_r^{l_o}$ except the action involved is the action of the whole category \mathbb{F} rather than its subcategory \mathbb{S} . As the functor \mathcal{Q}_r^o forgets this additional part of the action, the above diagram clearly commutes. \square

Proposition 4.5 *The functor $\mathcal{P}_f : \mathbf{RegOp} \rightarrow \mathbf{LT}$ is full on isomorphisms and its essential image is \mathbf{RegLT} . In particular, \mathbf{RegOp} is equivalent to \mathbf{RegLT} .*

Proof Recall that we have a unique morphism of Lawvere theories from the initial theory $\pi : \mathbb{F}^{op} \rightarrow \mathcal{P}_f(\mathcal{O})$. For a function $\phi : m \rightarrow n$, π_ϕ the morphism π_ϕ is represented by the span of form

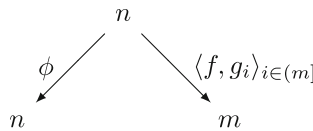


The class of projection morphisms is the closure under isomorphisms of the class of morphisms $\{\pi_\phi : \phi \in \mathbb{I}\}$ in $\mathcal{P}_f(\mathcal{O})$. Using Lemma 4.2, it is easy to see that the projection morphisms in $\mathcal{P}_f(\mathcal{O})$ are (represented by) the spans of the form



where ϕ is an injection and a_i is an invertible unary operation.

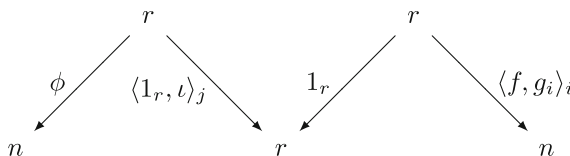
The regular morphisms in $\mathcal{P}_f(\mathcal{O})$ are (represented by) the spans of the form



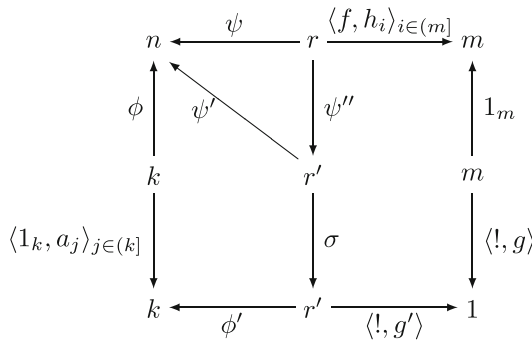
where ϕ is a bijection.

Clearly, both classes contain isomorphisms and are closed under composition.

Any morphism $\langle \phi, f, g_i \rangle_{i \in (m)} : n \rightarrow m$ in $\mathcal{P}_f(\mathcal{O})$ has a projection-regular factorization as follows



Thus to prove that projections and (what we have described as) regular morphisms form a factorization system, it remains to show that projection morphisms are left orthogonal to the regular morphisms. Let



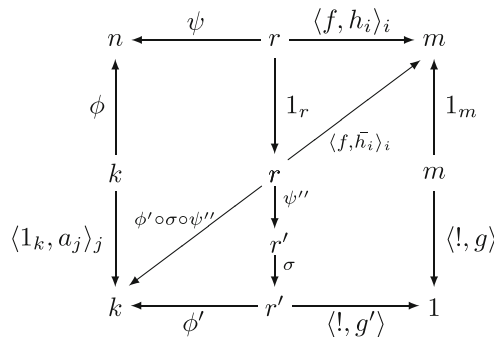
be a commutative square in $\mathcal{P}_f(\mathcal{O})$ with the left vertical morphism $\langle \phi, 1_r, a_i \rangle_{j \in \{k\}}$ a projection map (i.e. ϕ mono) and right vertical morphism $\langle 1_m, !, g \rangle$ a regular map. We have chosen the right bottom to be 1 to simplify notation but the general case is similar. The commutativity means that, with $\psi = \psi' \circ \psi''$ being an epi/mono factorization of ψ , there is a $\sigma \in \mathcal{S}_r$ such that

$$\psi' = \phi \circ \phi' \circ \sigma$$

and

$$\langle a_{\phi'(1)}, \dots, a_{\phi'(r)} \rangle * g' = \sigma \cdot \psi'' \cdot \langle (h_1, \dots, h_m) * g \rangle$$

Note that as $\phi \circ \phi'$ is mono, the representation of the composition $\langle \phi \circ \phi', a_{\phi'(1)}, \dots, a_{\phi'(r)} \rangle * g'$ is regular whereas the representation $\langle \psi, (h_1, \dots, h_m) * g \rangle$ is not in general. This is why we need to apply ψ'' to regularize it. Putting into the square a diagonal morphism $\langle \phi' \circ \sigma \circ \psi'', f, h_i \rangle_{i \in \{m\}}$



where

$$\bar{h}_i = \langle a_{\phi' \circ \sigma \circ \psi''(l)}^{-1} \rangle_{l \in f^{-1}(i)} * h_i$$

we see that the permutations 1_r and σ show that both triangles commute. An easy but tedious argument shows that such a diagonal morphism is in fact unique. It is

enough to consider the left triangle only and use the fact that $a_j s$ are invertible and ϕ is mono. Thus regular morphisms are indeed right orthogonal to the projections and $\mathcal{P}_f(\mathcal{O})$ is a regular Lawvere theory.

From the description of the functor $\mathcal{P}_f(h) : \mathcal{P}_f(\mathcal{O}) \rightarrow \mathcal{P}_f(\mathcal{O}')$ and the description of the structure of $\mathcal{P}_f(\mathcal{O})$ it is clear that $\mathcal{P}_f(h)$ sends the regular (projection) morphisms to the regular (projection) morphisms. Thus $\mathcal{P}_f(h)$ is a regular morphism of Lawvere theories.

Now let \mathbf{T} be any Lawvere theory. The class of regular morphisms in \mathbf{T} is right orthogonal to a class of morphisms and hence it is closed under composition, finite products and isomorphisms. Moreover, for any surjection in $\psi : [n] \rightarrow [m]$ in \mathbb{F} the image $\pi_\psi : m \rightarrow n$ in \mathbf{T} is regular as it is orthogonal to all projection morphisms in \mathbf{T} . Thus surjections in \mathbb{S} act on all regular morphisms $f : n \rightarrow 1$ in \mathbf{T} on the left

$$\cdot : \mathbb{S}(n, m) \times \mathbf{T}^r(n, 1) \longrightarrow \mathbf{T}^r(m, 1)$$

by

$$(\psi, f) \mapsto f \circ \pi_\psi$$

Hence the regular operations of any Lawvere theory \mathbf{T} form a regular operad. The unit is the identity morphism on 1. The composition $\langle f_1, \dots, f_n \rangle * f$ is defined to be $f \circ (f_1 \times \dots \times f_n)$. The action of \mathbb{S} is defined as above. Let us denote this operad by \mathbf{T}^r . We have an inclusion morphism of regular operads $\mathbf{T}^r \rightarrow \mathcal{Q}_r(\mathbf{T})$. By adjunction we get a morphism of Lawvere theories

$$\xi_{\mathbf{T}} : \mathcal{P}_f(\mathbf{T}^r) \longrightarrow \mathbf{T}$$

Clearly, $\xi_{\mathbf{T}}$ is bijective on objects. If \mathbf{T} is regular, then $\xi_{\mathbf{T}}$ is full (faithful) since the projection-regular factorization exists (is unique and $\pi : \mathbb{F}^{op} \rightarrow \mathbf{T}$ is faithful).

If $I : \mathbf{T} \rightarrow \mathbf{T}^r$ is a regular interpretation between any Lawvere theories, then the diagram

$$\begin{array}{ccc} \mathcal{P}_f(\mathbf{T}^r) & \xrightarrow{\mathcal{P}_f(I^r)} & \mathcal{P}_f(\mathbf{T}^{rr}) \\ \xi_{\mathbf{T}} \downarrow & & \downarrow \xi_{\mathbf{T}^r} \\ \mathbf{T} & \xrightarrow{I} & \mathbf{T} \end{array}$$

commutes, where I^r is the obvious restriction of I to \mathbf{T}^r . Thus the essential image of \mathcal{P}_f is indeed the category of regular Lawvere theories and regular interpretations. An isomorphic interpretation of Lawvere theories is always regular. Therefore \mathcal{P}_f is full on isomorphisms. □

We have

Proposition 4.6 *The functor $\mathcal{Q}_r : \mathbf{LT} \rightarrow \mathbf{RegOp}$ is monadic.*

Proof We shall verify the assumption of the Beck monadicity theorem. By Proposition 4.3 \mathcal{Q}_r has a left adjoint. It is easy to see that \mathcal{Q}_r reflects isomorphisms. We shall verify that \mathbf{LT} has and \mathcal{Q}_r preserves \mathcal{Q}_r -contractible coequalizers.

Let $I, I' : \mathbf{T}' \rightarrow \mathbf{T}$ be a pair of interpretations between Lawvere theories so that

$$\begin{array}{ccc} \mathcal{Q}_r(\mathbf{T}') & \begin{array}{c} \xrightarrow{\mathcal{Q}_r(I)} \\ \xrightarrow{\mathcal{Q}_r(I')} \\ \xleftarrow{r} \end{array} & \mathcal{Q}_r(\mathbf{T}) \begin{array}{c} \xrightarrow{q} \\ \xleftarrow{s} \end{array} \mathcal{O} \end{array}$$

is a split coequalizer in **RegOp**. We define a Lawvere theory $\mathbf{T}_{\mathcal{O}}$ so that a morphism from n to m in $\mathbf{T}_{\mathcal{O}}$ is an m -tuple $\langle g_1, \dots, g_m \rangle$ with $g_i \in \mathcal{O}_n$, for $i = 1, \dots, m$. The compositions and the identities in $\mathbf{T}_{\mathcal{O}}$ are defined in the obvious way from the compositions and the unit in \mathcal{O} . The projections $\bar{\pi}_i^n$ in $\mathbf{T}_{\mathcal{O}}$ are the images of the projections π_i^n in \mathbf{T} , i.e. $\bar{\pi}_i^n = q(\pi_i^n)$.

The functor $\tilde{q} : \mathbf{T} \rightarrow \mathbf{T}_{\mathcal{O}}$ is defined, for $f : n \rightarrow m$ in \mathbf{T} , as

$$\tilde{q}(f) = \langle q(\pi_1^m \circ f), \dots, q(\pi_m^m \circ f) \rangle$$

First, we verify that $\mathbf{T}_{\mathcal{O}}$ has finite products. For this, it is enough to verify that $\langle f_1, \dots, f_n \rangle * \bar{\pi}_i^n = f_i$, where $*$ is the multiplication in the operad \mathcal{O} . The uniqueness of the morphism into the product is obvious from the construction. We have routine calculations

$$\begin{aligned} & \langle f_1, \dots, f_n \rangle * \bar{\pi}_i^n \\ &= q \circ s(\langle f_1, \dots, f_n \rangle * q(\pi_i^n)) \\ &= \langle q \circ s(f_1), \dots, q \circ s(f_n) \rangle * (q \circ s \circ q(\pi_i^n)) \\ &= \langle q \circ s(f_1), \dots, q \circ s(f_n) \rangle * (q \circ \mathcal{Q}_r(I) \circ t(\pi_i^n)) \\ &= \langle q \circ s(f_1), \dots, q \circ s(f_n) \rangle * (q \circ \mathcal{Q}_r(I') \circ t(\pi_i^n)) \\ &= \langle q \circ s(f_1), \dots, q \circ s(f_n) \rangle * (q(\pi_i^n)) \\ &= q(\langle s(f_1), \dots, s(f_n) \rangle * \pi_i^n) \\ &= q(s(f_i)) = f_i \end{aligned}$$

It is obvious that \tilde{q} is a morphism of Lawvere theories and that $\mathcal{Q}_r(\tilde{q}) = q$. It remains to verify that \tilde{q} is a coequalizer in **LT**. Let $p : \mathbf{T} \rightarrow \mathbf{S}$ be a morphism in **LT** coequalizing I and I' .

$$\begin{array}{ccccc} \mathbf{T}' & \begin{array}{c} \xrightarrow{I} \\ \xrightarrow{I'} \end{array} & \mathbf{T} & \begin{array}{c} \xrightarrow{\tilde{q}} \\ \searrow p \end{array} & \mathbf{T}_{\mathcal{O}} \\ & & & & \downarrow \tilde{k} \\ & & & & \mathbf{S} \end{array}$$

The morphism $\mathcal{Q}_r(p)$ coequalizes $\mathcal{Q}_r(I)$ and $\mathcal{Q}_r(I')$ in **RegOp**. Thus there is a unique morphism k in **RegOp** making the triangle on the right

$$\begin{array}{ccccc}
 \mathcal{Q}_r(\mathbf{T}') & \xrightarrow{\mathcal{Q}_r(I)} & \mathcal{Q}_r(\mathbf{T}') & \xrightarrow{q} & \mathcal{O} \\
 & \xrightarrow{\mathcal{Q}_r(I')} & & & \downarrow k \\
 & & & \searrow \mathcal{Q}_r(p) & \mathcal{Q}_r(\mathbf{S})
 \end{array}$$

commute. We define the functor \tilde{k} so that

$$\tilde{k}(\langle f_1, \dots, f_n \rangle) = \langle k(f_1), \dots, k(f_n) \rangle$$

for any morphism $\langle f_1, \dots, f_n \rangle$ in $\mathbf{T}_{\mathcal{O}}$. The verification that \tilde{k} is the required unique functor is left for the reader. □

5 Semi-analytic Monads vs Regular Operads

The main objective of this section is to show that the square

$$\begin{array}{ccc}
 \mathbf{FOp} & \xrightarrow{\mathcal{M}_o^f} & \mathbf{Mnd} \\
 \mathcal{P}_f^o \uparrow & & \uparrow \mathcal{P}_f^m \\
 \mathbf{RegOp} & \xrightarrow{\mathcal{M}_o^r} & \mathbf{SanMnd}
 \end{array} \tag{4}$$

commutes up to isomorphism, with horizontal functors being equivalences of categories. The functor \mathcal{M}_o^f was defined in Section 3. \mathcal{P}_f^m is an inclusion.

To this end, we shall use Theorem 2.2 and some abstract considerations. For a regular operad \mathcal{O} , the monad $\mathcal{M}_o^r(\mathcal{O})$ on a set X is defined as follows. The functor $\mathcal{M}_o^r(\mathcal{O})$ is the application of the functor $(\hat{\cdot}) : \mathbf{Set}^{\mathbf{S}} \rightarrow \mathbf{End}$ to \mathcal{O} considered as a functor $\mathcal{O} : \mathbf{S} \rightarrow \mathbf{Set}$, see Section 2 for details.

In particular, for a set X we have

$$\mathcal{M}_o^r(\mathcal{O})(X) = \sum_{n \in \omega} \begin{bmatrix} X \\ n \end{bmatrix} \otimes_n \mathcal{O}_n$$

In the set $\begin{bmatrix} X \\ n \end{bmatrix} \otimes_n \mathcal{O}_n$ we identify $\langle \vec{x} \circ \sigma, a \rangle$ with $\langle \vec{x}, \sigma \cdot a \rangle$ for $a \in \mathcal{O}_n$, $\vec{x} : (n) \rightarrow X$ and $\sigma \in S_n$.

Let $\gamma : \mathbb{S} \rightarrow \mathbb{F}$ be the inclusion functor. It induces the following diagram of categories and functors that we describe below

$$\begin{array}{ccccccc}
 \mathbb{F} & & \mathbf{Set}^{\mathbb{F}} & \xrightarrow{\mathbf{i}_{\mathbb{F}}} & \mathbf{End} & \xleftarrow{\widehat{U}} & \mathbf{mon}(\mathbf{End}) = \mathbf{Mnd} \\
 \uparrow \gamma & & \uparrow \text{Lan}_{\gamma} & & \uparrow i^{sa} & & \uparrow \text{mon}(i^{sa}) \\
 & & \uparrow \downarrow \gamma^* & & \uparrow \downarrow (-)^{sa} & & \uparrow \downarrow \mathcal{P}_f^m \\
 \mathbb{S} & & \mathbf{Set}^{\mathbb{S}} & \xrightarrow{\mathbf{i}_{\mathbb{S}}} & \mathbf{San} & \xleftarrow{U} & \mathbf{mon}(\mathbf{San}) = \mathbf{SanMnd} \\
 & & & & \uparrow \downarrow & & \uparrow \downarrow \mathcal{Q}_r^m \\
 & & & & (\mathbb{W}, \eta, \mu) & & \mathbf{mon}(\mathbb{W}, \eta, \mu) = (\bar{\mathbb{W}}, \bar{\eta}, \bar{\mu})
 \end{array}
 \tag{5}$$

γ^* is the functor of ‘composing with γ ’. It has a left adjoint Lan_{γ} , the left Kan extension along γ . For $C \in \text{Set}^{\mathbb{S}}$ it is given by the coend formula

$$\text{Lan}_{\gamma}(C)(X) = \int^{n \in \mathbb{S}} X^n \times C(n)$$

The functor $i^{sa} : \mathbf{San} \rightarrow \mathbf{End}$ is just an inclusion. The equivalence

$$\mathbf{i}_{\mathbb{S}} : \text{Set}^{\mathbb{S}} \longrightarrow \mathbf{San}$$

is defined by left Kan extension that may be given by both the coend and coproduct formulas

$$\mathbf{i}_{\mathbb{S}}(C)(X) = \int^{n \in \mathbb{S}} X^n \times C(n) = \sum_{n \in \omega} \left[\begin{matrix} X \\ n \end{matrix} \right] \otimes_n C(n)$$

where $C \in \text{Set}^{\mathbb{S}}$. Similarly, the equivalence

$$\mathbf{i}_{\mathbb{F}} : \text{Set}^{\mathbb{F}} \longrightarrow \mathbf{End}$$

is defined by left Kan extension that is given by the coend formula

$$\mathbf{i}_{\mathbb{F}}(B)(X) = \int^{n \in \mathbb{F}} X^n \times C(n)$$

where $B \in \text{Set}^{\mathbb{F}}$. The following calculation shows that the left hand square in the above diagram commutes:

$$\begin{aligned}
 \mathbf{i}_{\mathbb{F}}(\text{Lan}_{\gamma}(C))(X) &\cong \int^{m \in \mathbb{F}} X^m \times \text{Lan}_{\gamma}(C)(m) \cong \\
 &\cong \int^{m \in \mathbb{F}} X^m \times \sum_{n \in \omega} (m)^n \otimes_n C(n) \cong \\
 &\cong \sum_{n \in \omega} X^n \otimes_n C(n)
 \end{aligned}$$

The functor $(-)^{sa}$, right adjoint to i^{sa} , is given by the formula

$$F^{sa}(X) = \sum_{n \in \omega} \begin{bmatrix} X \\ n \end{bmatrix} \otimes_n F(n)$$

for $F \in \mathbf{End}$. $(-)^{sa}$ associates to functors and natural transformations their ‘semi-analytic parts’.

Note that both \mathbf{San} and \mathbf{End} are strict monoidal categories with tensor given by composition, and i^{sa} is a strict monoidal functor. Thus its right adjoint $(-)^{sa}$ has a unique lax monoidal structure making the adjunction $i^{sa} \dashv (-)^{sa}$ a monoidal adjunction. This in turn gives us a monoidal monad (\mathbb{W}, η, μ) on \mathbf{San} .

We have a 2-natural transformation \mathcal{U}

$$\begin{array}{ccc} & \xrightarrow{\mathbf{mon}} & \\ \mathbf{MonCat} & \Downarrow \mathcal{U} & \mathbf{Cat} \\ & \xrightarrow{|\ - |} & \end{array}$$

where \mathbf{MonCat} is the 2-category of monoidal categories, lax monoidal functors, and monoidal transformations, \mathbf{mon} is the 2-functor associating the objects of monoids to monoidal categories, $|\ - |$ is the forgetful functor forgetting the monoidal structure, and \mathcal{U} is a 2-natural transformation whose component at a monoidal category M is the forgetful functor from monoids in M to the underlying category of M : $\mathcal{U}_M : \mathbf{mon}(M) \rightarrow |M|$ (cf. [15, 17]).

Applying \mathcal{U} to the monoidal adjunction $i^{sa} \dashv (-)^{sa}$ and monoidal monad \mathbb{W} we get an adjunction between categories of monoids and a monad on $\mathbf{mon}(\mathbf{San})$. The unnamed arrow in the above diagram is $\mathbf{mon}((-)^{sa})$. But the monoids in \mathbf{End} and \mathbf{San} are exactly monads and hence we get the right most adjunction $\mathcal{P}_f^m \dashv \mathcal{Q}_r^m$ together with the monad $(\bar{\mathbb{W}}, \bar{\eta}, \bar{\mu})$ on the category of semi-analytic monads.

On the other hand, on the categories $\mathbf{Set}^{\mathbb{F}}$ and $\mathbf{Set}^{\mathbb{S}}$ there are substitution tensors making $\mathbf{i}_{\mathbb{F}}$ and $\mathbf{i}_{\mathbb{S}}$ monoidal equivalences and $\mathbf{Lan}_{\gamma} \dashv \gamma^*$ a monoidal adjunction. Thus we can apply the 2-functor \mathbf{mon} to this adjunction and obtain an adjunction $\mathbf{mon}(\mathbf{Lan}_{\gamma}) \dashv \mathbf{mon}(\gamma^*)$ as in the diagram

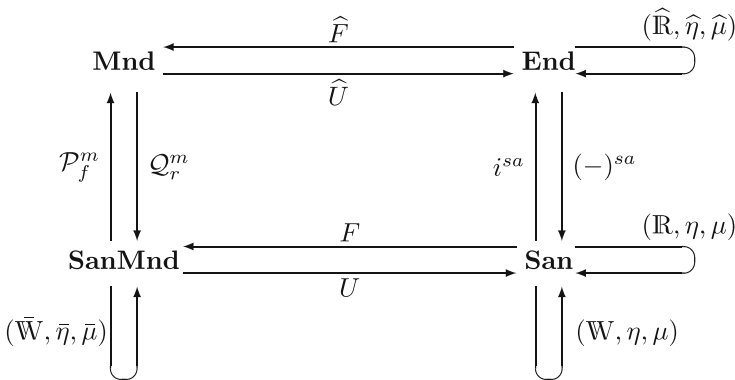
$$\begin{array}{ccc} \mathbf{FOp} = \mathbf{mon}(\mathbf{Set}^{\mathbb{F}}) & \xrightarrow{\mathcal{U}_{\mathbf{Set}^{\mathbb{F}}}} & \mathbf{Set}^{\mathbb{F}} \\ \uparrow \mathcal{P}_f^o \quad \downarrow \mathcal{Q}_r^o & & \uparrow \mathbf{Lan}_{\gamma} \quad \downarrow \gamma^* \\ \mathbf{RegOp} = \mathbf{mon}(\mathbf{Set}^{\mathbb{S}}) & \xrightarrow{\mathcal{U}_{\mathbf{Set}^{\mathbb{S}}}} & \mathbf{Set}^{\mathbb{S}} \\ & & \uparrow \mathbf{mon}(\gamma^*) \quad \downarrow \end{array} \tag{6}$$

The unnamed functor is $\mathbf{mon}(\mathbf{Lan}_{\gamma})$. But monoids in $\mathbf{Set}^{\mathbb{F}}$ and $\mathbf{Set}^{\mathbb{S}}$ are (equivalent to) the categories of full and regular operads, respectively. The verification that the left most square commutes serially is left for the reader. We obtain

Proposition 5.1 *The square (4) of categories and functors commutes up to isomorphism. The functor \mathcal{M}_o^r is an equivalence of categories.*

Proof We need to place the diagram (5) after the diagram (6). Both horizontal adjunctions in the square (4) are obtained from equivalent monoidal adjunctions $Lan_\gamma \dashv \gamma^*$ and $\mathbf{mon}(Lan_\gamma) \dashv \mathbf{mon}(\gamma^*)$. It remains to show that the identifications we obtained above are isomorphic to the functors \mathcal{M}_o^f and \mathcal{M}_o^r , respectively. This is left for the reader. \mathcal{M}_o^r is an equivalence of categories as a lift to the categories of monoids of a monoidal equivalence of monoidal categories. \square

There are free monads on finitary functors (cf. [1]) and free semi-analytic monads on semi-analytic functors. The adjunctions $F \dashv U$ and $\widehat{F} \dashv \widehat{U}$ induce monads \mathbb{R} and $\widehat{\mathbb{R}}$, respectively. $\widehat{\mathbb{R}}$ is the finitary version of what is called ‘the monad for all monads’ in [1]. Putting this additional data to the above diagram and simplifying it at the same time, we get a diagram



In the above diagram the square of left adjoints commutes. Thus, the square of right adjoints commutes as well. This shows in particular that the free monad on a semi-analytic functor is semi-analytic.

The monad \bar{W} is a lift of a monad W to the category of \mathbb{R} -algebras **SanMnd** and, by [2] we obtain

Theorem 5.2 *The monad \mathbb{R} for regular monads distributes over the monad W for finitary functors, i.e. we have a distributive law*

$$\lambda : \mathbb{R}W \longrightarrow W\mathbb{R}$$

*The category of algebras of the composed monad $W\mathbb{R}$ on **SanMnd** is equivalent to the category **Mnd** of all finitary monads on **San**.*

Remark We arrived at the above theorem with essentially no calculations. In [14], where we studied analytic functors, we have given explicit formulas for a similar distributive law. In the semi-analytic case such formulas would be much more involved and will not be given here.

6 Equational Theories vs Regular Operads

In this section we study the relations between regular equational theories and regular operads. We shall show that the square

$$\begin{array}{ccc}
 \mathbf{FOp} & \xrightarrow{\mathcal{E}_o^f} & \mathbf{ET} \\
 \mathcal{P}_f^o \uparrow & & \uparrow \mathcal{P}_f^e \\
 \mathbf{RegOp} & \xrightarrow{\mathcal{E}_o^r} & \mathbf{RegET}
 \end{array} \tag{7}$$

commutes up to isomorphism, with \mathcal{P}_f^e being an inclusion and both horizontal functors being equivalences of categories. \mathcal{P}_f^o was defined in Section 2 and \mathcal{E}_o^f was defined in Section 3. We shall define \mathcal{E}_o^r .

6.1 The Functor $\mathcal{E}_o^r : \mathbf{RegOp} \rightarrow \mathbf{ET}$

Let \mathcal{O} be a regular operad. We define an equational theory $\mathcal{E}_o^r(\mathcal{O}) = (L, A)$. As the set of n -ary function symbols we put $L_n = \mathcal{O}_n$ for $n \in \omega$. The set of axioms A contains the following three kinds of equations in context:

1. *unit*: $\iota(x_1) = x_1 : \vec{x}^1$ where $\iota \in \mathcal{O}_1$ is the unit of the operad \mathcal{O} ;
2. *action*: $a(x_{\tau(1)}, \dots, x_{\tau(m)}) = (\tau \cdot a)(x_1, \dots, x_n) : \vec{x}^n$ for all $a \in \mathcal{O}_m$ and surjections $\tau : (m) \rightarrow (n)$;
3. *multiplication*: $a(a_1(x_1, \dots, x_{k_1}), \dots, a_m(x_{k-k_m+1}, \dots, x_k)) = ((a_1, \dots, a_m) * a)(x_1, \dots, x_k) : \vec{x}^k$ where $a \in \mathcal{O}_m, a_i \in \mathcal{O}_{k_i}$ for $i \in 1, \dots, m, k = \sum_{i=1}^m k_i$;

Clearly, all equations are regular and hence the theory $\mathcal{E}_o^r(\mathcal{O})$ is regular.

Suppose that $h : \mathcal{O} \rightarrow \mathcal{O}'$ is a morphism of regular operads. We define the interpretation

$$\mathcal{E}_o^r(h) : \mathcal{E}_o^r(\mathcal{O}) \longrightarrow \mathcal{E}_o^r(\mathcal{O}')$$

For $a \in \mathcal{O}_n$ we put

$$\mathcal{E}_o^r(h)(a) = h(a)(x_1, \dots, x_n) : \vec{x}^n$$

for $n \in \omega$. Clearly, $\mathcal{E}_o^r(h)$ is a regular interpretation.

Proposition 6.1 *The square (7) commutes up to a natural isomorphism.*

Proof Let \mathcal{O} be a regular operad. We define an interpretation of equational theories

$$I_{\mathcal{O}} : \mathcal{E}_o^r(\mathcal{O}) \longrightarrow \mathcal{E}_o^f \mathcal{P}_f^o(\mathcal{O})$$

by

$$\mathcal{O}_n \ni a \mapsto [[1_n, a](x_1, \dots, x_n) : \vec{x}^n]$$

where $a \in \mathcal{O}_n$ and $[1_n, a]$ is an n -ary operation symbol of the theory $\mathcal{E}_o^f \mathcal{P}_f^o(\mathcal{O})$.

We need to verify that I is a well defined natural isomorphism. First, we need to verify that $I_{\mathcal{O}}$ preserves axioms. The unit axiom is obvious. To prove the action axioms, we fix $a \in \mathcal{O}_m$ and $\tau : [m] \rightarrow [n]$ and we calculate using the theory $\mathcal{E}_o^f \mathcal{P}_f^o(\mathcal{O})$

$$\begin{aligned} I_{\mathcal{O}}(a(x_{\tau(1)}, \dots, x_{\tau(m)})) &= [1_m, a](x_{\tau(1)}, \dots, x_{\tau(m)}) = \\ &= \tau \cdot [1_m, a](x_1, \dots, x_n) = [1_m, \tau \cdot a](x_1, \dots, x_n) = \\ &= I_{\mathcal{O}}(\tau \cdot a(x_1, \dots, x_n)) \end{aligned}$$

The calculations for the composition axioms are similar. The naturality of $I_{\mathcal{O}}$ is left for the reader.

We shall show that $I_{\mathcal{O}}$ is an isomorphism of theories. To this end, we define an inverse interpretation

$$J_{\mathcal{O}} : \mathcal{E}_o^f \mathcal{P}_f^o(\mathcal{O}) \longrightarrow \mathcal{E}_o^r(\mathcal{O})$$

given by

$$\mathbb{I}(k, n) \otimes_k \mathcal{O}_k \ni [f, a] \mapsto a(x_{f(1)}, \dots, x_{f(k)}) : \vec{x}^n$$

Again, we need to verify that $J_{\mathcal{O}}$ preserves the axioms and again we shall verify the action axiom only. Fix $\phi : [n] \rightarrow [m]$ and $[f, a] \in \mathbb{I}(k, n) \otimes \mathcal{O}_k$. Let $f' \circ \phi'$ be an epi-mono factorization in \mathbb{F} with $f' : [n'] \rightarrow [m]$. Using the theory $\mathcal{E}_o^r(\mathcal{O})$, we have

$$\begin{aligned} J_{\mathcal{O}}((\phi \cdot [f, a])(x_1, \dots, x_n)) &= [f', \phi' \cdot a](x_1, \dots, x_k) = \\ &= (\phi' \cdot a)(x_{f'(1)}, \dots, x_{f'(n')}) = a(x_{f' \circ \phi'(1)}, \dots, x_{f' \circ \phi'(n')}) = \\ &= a(x_{\phi \circ f(1)}, \dots, x_{\phi \circ f(n')}) = J_{\mathcal{O}}([f, a](x_{\phi(1)}, \dots, x_{\phi(n)})) \end{aligned}$$

Finally, we need to verify that $I_{\mathcal{O}}$ and $J_{\mathcal{O}}$ are mutually inverse one to the other. The composition $J_{\mathcal{O}}I_{\mathcal{O}}$ sends an operation $a \in \mathcal{O}_n$ to the term $a(x_1, \dots, x_n) : \vec{x}^n$, so it is the identity. For an operation $[f, a] \in \mathbb{I}(k, n) \otimes_k \mathcal{O}_k$ in theory $\mathcal{E}_o^f \mathcal{P}_f^o(\mathcal{O})$ we have

$$\begin{aligned} I_{\mathcal{O}}J_{\mathcal{O}}([f, a]) &= [1_n, a](x_{f(1)}, \dots, x_{f(n)}) : \vec{x}^n = \\ &= f \cdot [1_n, a](x_1, \dots, x_k) : \vec{x}^k = [f, a](x_1, \dots, x_k) : \vec{x}^k \end{aligned}$$

i.e. the composition $I_{\mathcal{O}}J_{\mathcal{O}}$ is the identity as well. □

Finally, we have

Theorem 6.2 *The functor $\mathcal{E}_o^r : \mathbf{RegOp} \longrightarrow \mathbf{RegET}$ is an equivalence of categories.*

Proof From Proposition 6.1 it follows that \mathcal{E}_o^r is faithful. To see that it is full, consider a regular interpretation $I : \mathcal{E}_o^r(\mathcal{O}) \rightarrow \mathcal{E}_o^r(\mathcal{O}')$. We shall define a morphism of regular operads $h : \mathcal{O} \rightarrow \mathcal{O}'$ such that $\mathcal{E}_o^r(h) = I$. For $f \in \mathcal{O}_n$, $I(f)(\vec{x}^n) : \vec{x}^n$ is a regular term. But in theories in the image of \mathcal{E}_o^r every regular term is equal to a simple term, i.e. there is an $f' \in \mathcal{O}'_n$ such that $f'(\vec{x}^n) = I(f)(\vec{x}^n) : \vec{x}^n$ is a theorem in the theory $\mathcal{E}_o^r(\mathcal{O}')$. Thus we put $h(f) = f'$. It is left for the reader to verify that h has the required property.

To see that \mathcal{E}_o^r is essentially surjective, let us fix a regular theory $T = (L, A)$. Then the regular terms in T form a regular operad called T^{ro} . The unit is the term $x_1 : x_1$. The composition is defined by the substitution (making sure that we make the variables disjoint in different substituted terms, via α -conversion). The action of a surjection $\phi : (n] \rightarrow (m]$ on a regular term $t(x_1, \dots, x_n) : \vec{x}^n$ is again a regular term $t(x_{\phi(1)}, \dots, x_{\phi(n)}) : \vec{x}^n$. Again, it is a matter of a routine verification to show that $\mathcal{E}_o^r(T^{ro}) \cong T$. \square

6.2 Examples

1. The terminal equational theory $\mathbb{1}$ has one constant symbol, say e , and can be axiomatized by a single axiom: $v_1 = e : \vec{v}^1$. As a Lawvere theory it is the category that has exactly one morphism between any two objects.
2. The embedding of the regular theories into all equational theories $P_f^e : \mathbf{RegET} \rightarrow \mathbf{ET}$ has a right adjoint $Q_f^e : \mathbf{ET} \rightarrow \mathbf{RegET}$ that we denote here $(-)^r$. For an equational theory $T = (L, A)$ the theory $T^r = (L^r, A^r)$ can be described as follows. The n -ary function symbols of L^r are all (not only regular) terms in context $t : \vec{v}^n$ over the language L . The equations in A^r are all the regular equations over L^r such that when interpreted as terms over L are consequences of the set of axioms A .
3. The value $\mathbb{1}^r$ of the functor $(-)^r$ on the terminal equational theory $\mathbb{1}$ is the terminal regular theory. It is the theory of join-semilattices SL :

$$(x \vee y) \vee z = x \vee (y \vee z), \quad x \vee y = y \vee x, \quad x \vee \perp = x \quad (x \vee x) = x$$

4. If $R = (L, A)$ is any regular equational theory, then the unique regular interpretation $! : R \rightarrow SL$ interprets every constant c from L as the constant term

$$\perp : \vec{v}^0$$

and for $n > 0$, every n -ary function symbol f from L as the term

$$v_1 \vee, \dots, \vee v_n : \vec{v}^n$$

5. The above interpretation $!$ induces a functor between categories of algebras

$$Alg(!) : Alg(SL) \longrightarrow Alg(R)$$

that ‘reinterprets’ sup-semilattices as algebras of any regular theory in the way described above. This functor has a left adjoint providing a join-semilattice reflections of models of R .

6. The theory $\mathbb{1}$ has a unique proper subtheory $\mathbb{1}^-$. It has no function symbols and can be axiomatized by a single axiom: $v_1 = v_2 : \vec{v}^2$. The regular part of this theory $(\mathbb{1}^-)^r$, denoted SL^- , is the theory of join-semilattices without a bottom element \perp . This theory is the terminal theory in the full subcategory of the regular theories without constant symbols. In particular, there is a unique regular interpretation $! : SG \longrightarrow SL^-$ of the theory of semigroups SG in SL^- . As before, it induces a functor between categories of algebras that has a left adjoint

$$Alg(SG) \longrightarrow Alg(SL^-)$$

This functor is called ‘greatest semilattice image functor’; cf. [5].

7 Cartesian and Weakly Cartesian Monads

In this section we shall investigate two (strict monoidal) subcategories of **San** and their categories of monoids. The category of (weakly) cartesian functors and (weakly) cartesian natural transformations will be denoted by **Cart** (**wCart**). The corresponding categories of monoids: the category of (weakly) cartesian monads will be denoted by **CartMnd** (**wCartMnd**). Thus we have embeddings full on isomorphisms

$$\mathbf{Cart} \longrightarrow \mathbf{wCart} \longrightarrow \mathbf{San}$$

which are strict monoidal and induce embeddings of categories of monoids

$$\mathbf{CartMnd} \longrightarrow \mathbf{wCartMnd} \longrightarrow \mathbf{SanMnd}$$

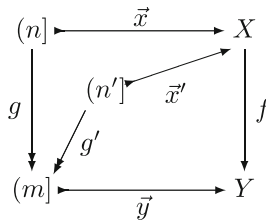
The characterizations of the subcategories of equational theories **ET** and Lawvere theories **LT** corresponding to **CartMnd** and **wCartMnd** are a bit technical and we are not going to describe them in detail here. Clearly, the objects are some regular theories satisfying additional conditions and similarly for morphisms. We shall content ourselves with a description of subcategories of Set^S whose essential images are **wCart** and **Cart**, respectively. Note, however, that if (T, η, μ) is a semi-analytic monad such that the functor part T is the left Kan extension of a functor $R : S \rightarrow Set$, then R is the functor of all regular operations in the equational theory corresponding to the monad T . Thus our description will in fact provide a description of the equational theories corresponding to monads in **wCart** and **Cart**.

Recall the description of functor $(-)^S : Set^S \rightarrow \mathbf{End}$ from Section 2, the Kan extension along $i_S : S \rightarrow Set$. We begin with the following observation

Lemma 7.1 *Let $A : S \rightarrow Set$ be a functor, $f : X \rightarrow Y$ a function, $[\vec{x} : (n) \rightarrow X, a \in A_n], [\vec{x}' : (n') \rightarrow X, a' \in A_{n'}] \in \hat{A}(X)$. If*

$$\hat{A}(f)([\vec{x}, a]) = \hat{A}(f)([\vec{x}', a'])$$

then there are $m \in \omega$, surjections $g : n \rightarrow m, g' : n' \rightarrow m$, and an injection $\vec{y} : m \rightarrow Y$ as in the diagram



such that $\vec{y} \circ g = f \circ \vec{x}, \vec{y} \circ g' = f \circ \vec{x}'$, and $A(g)(a) = A(g')(a')$. In particular

$$\hat{A}(f)([\vec{x}, a]) = [\vec{y}, A(g)(a)] = [\vec{y}, A(g')(a')] = \hat{A}(f)([\vec{x}', a'])$$

Proof Exercise. □

The following two Propositions identify the subcategory of Set^S whose essential image in **End** is **wCart**.

Proposition 7.2 *Let $A : \mathbb{S} \rightarrow \text{Set}$ be a functor. The functor $\hat{A} : \text{Set} \rightarrow \text{Set}$ weakly preserves pullbacks iff the functor A satisfies the following condition (WPB): for any pullback of surjections in \mathbb{F}*

$$\begin{array}{ccc} (p) & \xrightarrow{g_2} & (m) \\ f_2 \downarrow & & \downarrow f_1 \\ (n) & \xrightarrow{g_1} & (k) \end{array}$$

and a pair of elements $a \in A(m)$, $b \in A(n)$ such that $A(f_1)(a) = A(g_1)(b)$ there is an injection $h : (q) \rightarrow (p)$ and an element $c \in A(q)$ such that $f_2 \circ h$ and $g_2 \circ h$ are surjections and

$$A(g_2 \circ h)(c) = a, \quad \text{and} \quad A(f_2 \circ h)(c) = b$$

Proof

\Rightarrow First we verify that the condition (WPB) is necessary. Let us fix the square and elements a and b as in the above condition. Then we have elements

$$[1_{(m)}, a] \in \left[\begin{matrix} (m) \\ m \end{matrix} \right] \otimes_m A(m) \subseteq \sum_{k \in \omega} \left[\begin{matrix} (m) \\ k \end{matrix} \right] \otimes_k A(k) = \hat{A}(X)$$

and

$$[1_{(n)}, b] \in \left[\begin{matrix} (n) \\ n \end{matrix} \right] \otimes_n A(n) \subseteq \hat{A}(X)$$

such that

$$\hat{A}(f_1)([1_{(m)}, a]) = \hat{A}(g_1)([1_{(n)}, b])$$

As f_1, g_1 are surjections, this means that

$$A(f_1)(a) = A(g_1)(b)$$

As \hat{A} weakly preserves pullbacks, there is an $[h, c] \in \hat{A}(p)$ such that

$$\hat{A}(f_2)([h, c]) = [1_{(n)}, b], \quad \text{and} \quad \hat{A}(g_2)([h, c]) = [1_{(m)}, a]$$

Thus for some $h : (q) \rightarrow (p), c \in A(q)$ we have $[h, c] \in \left[\begin{matrix} (p) \\ q \end{matrix} \right] \otimes_q A(q)$ and using the diagram

$$\begin{array}{ccccc} (q) & \xrightarrow{g_3} & (m') & & \\ \downarrow f_3 & \searrow h & \downarrow \vec{x} & & \\ & & (p) & \xrightarrow{g_2} & (m) \\ & & \downarrow f_2 & & \downarrow f_1 \\ (n') & \xrightarrow{\vec{y}} & (n) & \xrightarrow{g_1} & (k) \end{array}$$

we have

$$[\vec{y}, A(f_3)(c)] = [1_{(n)}, b] \quad \text{and} \quad [\vec{x}, A(g_3)(c)] = [1_{(m)}, a]$$

This means that $n = n'$, \vec{y} is a bijection, $\vec{y} \circ f_3 = f_2 \circ h$ is a surjection and $b = A(\vec{y} \circ f_3)(c) = A(f_2 \circ h)(c)$. Similarly, we get that $m = m'$, $g_2 \circ h$ is a surjection and $a = A(g_2 \circ h)(c)$. Thus the condition (WPB) is necessary.

⇐ To show that the condition (WPB) is sufficient, we suppose that a functor $A : \mathbb{S} \rightarrow \text{Set}$ satisfies (WPB) and we shall show that $\hat{A} : \text{Set} \rightarrow \text{Set}$ weakly preserves pullbacks.

Let us fix a pullback in Set and

$$\begin{array}{ccc} R & \xrightarrow{g_2} & Y \\ f_2 \downarrow & & \downarrow f_1 \\ Z & \xrightarrow{g_1} & X \end{array} \tag{8}$$

and elements $[\vec{z}, b] \in \hat{A}(Z)$ and $[\vec{y}, a] \in \hat{A}(Y)$ such that

$$\hat{A}(f_1)([\vec{y}, a]) = \hat{A}(g_1)([\vec{z}, b])$$

i.e. in the diagram

$$\begin{array}{ccccc} (r] & \xrightarrow{h} & (q] & \xrightarrow{g_4} & (m] \\ & & \searrow v & & \swarrow \vec{y} \\ & & R & \xrightarrow{g_2} & Y \\ f_4 \downarrow & & \downarrow f_2 & & \downarrow f_1 \\ & & Z & \xrightarrow{g_1} & X \\ & & \swarrow \vec{z} & & \swarrow \vec{z}' \\ (n] & \xrightarrow{g_3} & (k] & \xleftarrow{\vec{y}'} & (k] \\ & & & & \uparrow \sigma \end{array} \tag{9}$$

we have

$$[\vec{y}', A(f_3)(a)] = [\vec{z}', A(g_3)(b)]$$

i.e. there is a permutation $\sigma \in S_k$ such that

$$\vec{y}' \circ \sigma = \vec{z}' \quad \text{and} \quad A(f_3)(a) = A(\sigma \circ g_3)(b)$$

Still, in the above diagram we can form a pullback in \mathbb{F}

$$\begin{array}{ccc} (q] & \xrightarrow{g_4} & (m] \\ f_4 \downarrow & & \downarrow f_3 \\ (n] & \xrightarrow{\sigma \circ g_3} & (k] \end{array}$$

and we have a morphism $v : (q) \rightarrow R$ into the pullback (8) such that

$$f_2 \circ v = \bar{z} \circ f_4, \quad g_2 \circ v = \bar{y} \circ g_4$$

Thus by (WPB) there is an injection $h : (r) \rightarrow (q)$ and a $c \in A(q)$ such that both $f_4 \circ h$ and $g_4 \circ h$ are surjections and

$$A(g_4 \circ h)(c) = a, \quad A(f_4 \circ h)(c) = b$$

As $f_4 \circ h$ is a surjection, \bar{z} is an injection and $f_2 \circ v \circ h = \bar{z} \circ f_4 \circ h$ we have

$$\hat{A}(f_2)(\hat{A}(v)([h, c])) = [\bar{z}, A(f_4 \circ h)(c)] = [\bar{z}, b]$$

Similarly

$$\hat{A}(g_2)(\hat{A}(v)([h, c])) = [\bar{y}, a]$$

Thus $\hat{A}(v)([h, c]) \in \hat{A}(P)$ is the sought element in the weak pullback

$$\begin{array}{ccccc} [h, c] \in \hat{A}(q) & \xrightarrow{\hat{A}(v)} & \hat{A}(P) & \xrightarrow{\hat{A}(g_2)} & \hat{A}(Y) \\ & & \hat{A}(f_2) \downarrow & & \downarrow \hat{A}(f_1) \\ & & \hat{A}(Z) & \xrightarrow{\hat{A}(g_1)} & \hat{A}(X) \end{array}$$

□

Proposition 7.3 *Let $\tau : A \rightarrow B$ be a natural transformation in Set^S . Then τ is weakly cartesian iff $\hat{\tau} : \hat{A} \rightarrow \hat{B}$ is weakly cartesian.*

Proof Assume that $\tau : A \rightarrow B$ is a weakly cartesian natural transformation, $f : X \rightarrow Y$ is a function. We shall show that the square

$$\begin{array}{ccc} \hat{A}(X) & \xrightarrow{\hat{\tau}_X} & \hat{B}(X) \\ \hat{A}(f) \downarrow & & \downarrow \hat{B}(f) \\ \hat{A}(Y) & \xrightarrow{\hat{\tau}_Y} & \hat{B}(Y) \end{array}$$

is a weak pullback. Let us fix elements

$$[\bar{y}, a] \in \begin{bmatrix} Y \\ n \end{bmatrix} \otimes_n A(n) \subseteq \hat{A}(Y), \quad \text{and} \quad [\bar{x}, b] \in \begin{bmatrix} X \\ m \end{bmatrix} \otimes_m B(m) \subseteq \hat{B}(X)$$

such that

$$\hat{B}(f)([\bar{x}, b]) = \hat{\tau}([\bar{y}, a]) \quad (= [\bar{y}, \tau_n(a)])$$

By Lemma 7.1 we have a surjection f'

$$\begin{array}{ccc} (m) & \xrightarrow{\vec{x}} & X \\ f' \downarrow & & \downarrow f \\ (n) & \xrightarrow{\vec{y}} & Y \end{array}$$

such that $B(f')(b) = \tau_n(a)$. Since τ is weakly cartesian, there is $c \in A(m)$ such that

$$\tau_m(c) = b, \quad \text{and} \quad A(f')(c) = a$$

Then

$$[\vec{x}, c] \in \left[\begin{array}{c} X \\ m \end{array} \right] \otimes_m A(m) \subseteq \hat{A}(X)$$

and moreover

$$\hat{A}(f)([\vec{x}, c]) = [\vec{y}, A(f')(c)] = [\vec{y}, a]$$

and

$$\hat{\tau}_X([\vec{x}, c]) = [\vec{x}, \tau_m(c)] = [\vec{x}, b]$$

To prove the converse, let us assume that $\hat{\tau} : \hat{A} \rightarrow \hat{B}$ is a weakly cartesian natural transformation and that $f : (m) \rightarrow (n)$ is a surjection in \mathcal{S} . We need to show that the square

$$\begin{array}{ccc} A(m) & \xrightarrow{\tau_m} & B(m) \\ A(f) \downarrow & & \downarrow B(f) \\ A(n) & \xrightarrow{\tau_n} & B(n) \end{array}$$

is weakly cartesian. Fix $a \in A(n)$ and $b \in B(m)$ such that $\tau_n(a) = B(f)(b)$. Consider a weak pullback in Set

$$\begin{array}{ccc} \hat{A}(m) & \xrightarrow{\hat{\tau}_m} & \hat{B}(m) \\ \hat{A}(f) \downarrow & & \downarrow \hat{B}(f) \\ \hat{A}(n) & \xrightarrow{\hat{\tau}_n} & \hat{B}(n) \end{array}$$

We have elements

$$[1_{(n)}, a] \in \left[\begin{array}{c} (n) \\ n \end{array} \right] \otimes_n A(n) \subseteq \hat{A}(n), \quad \text{and} \quad [1_{(m)}, b] \in \left[\begin{array}{c} (m) \\ m \end{array} \right] \otimes_m B(m) \subseteq \hat{B}(m)$$

such that

$$\hat{\tau}_n([1_{(n)}, a]) = [1_{(n)}, \tau_n(a)] = [1_{(n)}, B(f)(b)] = \hat{B}(f)([1_{(m)}, b])$$

Thus, by our assumption, for some $k \in \omega$ there is an $[\vec{x}, c] \in \left[\begin{smallmatrix} (m) \\ k \end{smallmatrix} \right] \otimes_k A(k)$ such that

$$\hat{A}(f)([\vec{x}, c]) = [1_{(n)}, a], \quad \text{and} \quad \hat{\tau}_m([\vec{x}, c]) = [1_{(m)}, b]$$

Thus $(\vec{x}, \tau_m(c)) \sim (1_{(m)}, b)$ and we have that $k = m$, \vec{x} is a bijection, $(\vec{x}, c) \sim (1_{(m)}, A(\vec{x})(c))$. Hence we also have $\tau_m(A(\vec{x})(c)) = b$. Moreover, as the square

$$\begin{array}{ccc} (m) & \xrightarrow{\vec{x}} & (m) \\ f \circ \vec{x} \downarrow & & \downarrow f \\ (n) & \xrightarrow{1_{(n)}} & (n) \end{array}$$

commutes, we have

$$A(f)(A(\vec{x})(c)) = A(f \circ \vec{x})(c) = a$$

Thus $A(\vec{x})(c)$ is the element sought for a and b and hence τ is weakly cartesian. \square

The final two Propositions identify the subcategory of Set^S whose essential image in **End** is **Cart**.

Proposition 7.4 *Let $A : S \rightarrow Set$ be a functor. The functor $\hat{A} : Set \rightarrow Set$ preserves pullbacks iff the functor A satisfies the condition (WPB) from the Proposition 7.2 and additionally satisfies the following condition (PB): suppose that the square*

$$\begin{array}{ccc} (p) & \xrightarrow{g_2} & (m) \\ f_2 \downarrow & & \downarrow f_1 \\ (n) & \xrightarrow{g_1} & (k) \end{array}$$

is a pullback of surjections in \mathbb{F} , $\vec{x} : (q) \rightarrow (p)$, $\vec{x}' : (q') \rightarrow (p)$ are two injections, $c \in A(q)$, $c' \in A(q')$ are two elements so that the functions

$$f_2 \circ \vec{x}, \quad f_2 \circ \vec{x}', \quad g_2 \circ \vec{x}, \quad g_2 \circ \vec{x}'$$

are surjections and

$$A(f_2 \circ \vec{x})(c) = A(f_2 \circ \vec{x}')(c'), \quad A(g_2 \circ \vec{x})(c) = A(g_2 \circ \vec{x}')(c').$$

Then $q = q'$ and there is $\sigma \in S_q$ such that

$$\vec{x}' \circ \sigma = \vec{c}, \quad A(\sigma)(c) = c'.$$

Proof Assume that A satisfies (WPB) and (PB). Thus, by Proposition 7.3, \hat{A} weakly preserves pullbacks. Let

$$\begin{array}{ccc} P & \xrightarrow{g_2} & Y \\ f_2 \downarrow & & \downarrow f_1 \\ Z & \xrightarrow{g_1} & X \end{array}$$

be a pullback in *Set*. We shall show that the square

$$\begin{array}{ccc}
 \hat{A}(P) & \xrightarrow{\hat{A}(g_2)} & \hat{A}(Y) \\
 \hat{A}(f_2) \downarrow & & \downarrow \hat{A}(f_1) \\
 \hat{A}(Z) & \xrightarrow{\hat{A}(g_1)} & \hat{A}(X)
 \end{array}$$

is a pullback, i.e. it satisfies the uniqueness condition. So let

$$[h, a] \in \left[\begin{array}{c} P \\ q \end{array} \right] \otimes_q A(q), \quad \text{and} \quad [h', a'] \in \left[\begin{array}{c} P \\ q' \end{array} \right] \otimes_{q'} A(q')$$

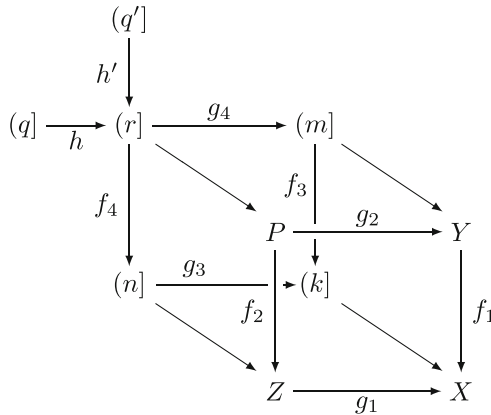
be such that

$$\hat{A}(f_2)([h, a]) = \hat{A}(f_2)([h', a']), \quad \text{and} \quad \hat{A}(g_2)([h, a]) = \hat{A}(g_2)([h', a'])$$

This implies the equalities of images of functions

$$im(f_2 \circ h) = im(f_2 \circ h'), \quad im(g_2 \circ h) = im(g_2 \circ h'), \quad im(g_1 \circ f_2 \circ h) = im(f_2 \circ h')$$

All these sets are finite, say, having n , m , and k elements, respectively. Thus we can form a commuting diagram



where the back square is a pullback of surjections (in \mathbb{F}) and all arrows from the back to the front are injections (for simplicity: inclusions). By definition the functions

$$f_4 \circ h, \quad f_4 \circ h', \quad g_4 \circ h, \quad f_4 \circ h'$$

are surjections. Since \hat{A} preserves monos, we also have

$$A(f_4 \circ h)(a) = A(f_4 \circ h')(a'), \quad A(g_4 \circ h)(a) = A(g_4 \circ h')(a')$$

Thus from the condition (PB) it follows that $q = q'$ and there is a $\sigma \in S_q$ such that $h' \circ \sigma = h$ and $A(\sigma)(a) = a'$, i.e. $[h, a] = [h', a']$, as required.

To prove the converse, we assume now that \hat{A} preserves pullbacks and we fix a pullback

$$\begin{array}{ccc} (r) & \xrightarrow{g_2} & (m) \\ f_2 \downarrow & & \downarrow f_1 \\ (n) & \xrightarrow{g_1} & (k) \end{array}$$

in \mathbb{F} of surjections. As \hat{A} weakly preserves pullbacks, the above pullback is sent by \hat{A} to a weak pullback in Set . Then, using Lemma 7.1 it is easy to see that the condition (PB) expresses the uniqueness part of the fact that \hat{A} sends the above square to a pullback in Set . \square

Proposition 7.5 *Let $\tau : A \rightarrow B$ be a natural transformation. Then τ is cartesian iff $\hat{\tau} : \hat{A} \rightarrow \hat{B}$ is cartesian.*

Proof First assume that $\tau : A \rightarrow B$ is a cartesian natural transformation. Fix a function $f : X \rightarrow Y$. By Proposition 7.3 the square

$$\begin{array}{ccc} \hat{A}(X) & \xrightarrow{\hat{\tau}_X} & \hat{B}(X) \\ \hat{A}(f) \downarrow & & \downarrow \hat{B}(f) \\ \hat{A}(Y) & \xrightarrow{\hat{\tau}_Y} & \hat{B}(Y) \end{array}$$

is a weak pullback. We shall show that it also satisfies the uniqueness property. Let us fix elements

$$[\vec{x}, a] \in \left[\begin{array}{c} X \\ n \end{array} \right] \otimes_n A(n) \subseteq \hat{A}(X), \quad \text{and} \quad [\vec{x}', a'] \in \left[\begin{array}{c} X \\ n' \end{array} \right] \otimes_{n'} A(n') \subseteq \hat{A}(X)$$

such that

$$\hat{\tau}([\vec{x}, a]) = \hat{\tau}([\vec{x}', a']), \quad \text{and} \quad \hat{A}(f)([\vec{x}, a]) = \hat{A}(f)([\vec{x}', a']) \tag{10}$$

The first equality means that $[\vec{x}, \tau_n(a)] = [\vec{x}', \tau_n(a')]$. Hence $n = n'$ and there is a $\sigma \in S_n$ such that

$$\vec{x}' \circ \sigma = \vec{x}, \quad \text{and} \quad \tau_n(a') = B(\sigma)(\tau_n(a)) = \tau_n(A(\sigma)(a))$$

By Lemma 7.1 and the second equality in Eq. 10 there are surjections g and g' and an injection \vec{y} as in the diagram

$$\begin{array}{ccc} (n) & \xrightarrow{\vec{x}} & X \\ & \searrow \sigma & \nearrow \vec{x}' \\ & (n') & \\ g \downarrow & & \downarrow f \\ (m) & \xrightarrow{\vec{y}} & Y \\ & \nearrow g' & \end{array}$$

such that

$$\bar{y} \circ g = f \circ \bar{x}, \quad \bar{y} \circ g' = f \circ \bar{x}', \quad \text{and} \quad A(g')(a') = A(g)(a)$$

Then we have

$$\bar{y} \circ g = f \circ \bar{x}' \circ \sigma = \bar{y} \circ g' \circ \sigma$$

As \bar{y} is mono $g = g' \circ \sigma$. Thus

$$A(g')(a') = A(g)(a) = A(g')(A(\sigma)(a))$$

As τ is a cartesian natural transformation, we get, from the fact that the naturality square for g' is a pullback, that $a' = A(\sigma)(a)$. But this means that

$$(\bar{x}', a') \sim (\bar{x}', A(\sigma)(a)) \sim (\bar{x}' \circ \sigma, a) \sim (\bar{x}, a)$$

i.e. $[\bar{x}', a'] = [\bar{x}, a]$ and $\hat{\tau}$ is cartesian.

To show the converse, assume that $\hat{\tau}$ is a cartesian natural transformation and $f : (n) \rightarrow (m)$ is a surjection. We need to show that the weak pullback

$$\begin{array}{ccc} A(n) & \xrightarrow{\tau_n} & B(n) \\ A(f) \downarrow & & \downarrow B(f) \\ A(m) & \xrightarrow{\tau_m} & B(m) \end{array}$$

satisfies the uniqueness condition, as well. Fix $a, a' \in A(n)$ such that

$$A(f)(a) = A(f)(a'), \quad \text{and} \quad \tau_n(a) = \tau_n(a')$$

By assumption the square

$$\begin{array}{ccc} \hat{A}(n) & \xrightarrow{\hat{\tau}(n)} & \hat{B}(n) \\ \hat{A}(f) \downarrow & & \downarrow \hat{B}(f) \\ \hat{A}(m) & \xrightarrow{\hat{\tau}(m)} & \hat{B}(m) \end{array}$$

is a pullback. We have elements $[1_{(n)}, a], [1_{(n)}, a'] \in \hat{A}(n)$ such that

$$\tau_{(n)}([1_{(n)}, a]) = [1_{(n)}, \tau_n(a)] = [1_{(n)}, \tau_n(a')] = \tau_{(n)}([1_{(n)}, a'])$$

and

$$\hat{A}(f)([1_{(n)}, a]) = [1_{(n)}, A(f)(a)] = [1_{(n)}, A(f)(a')] = \hat{A}([1_{(n)}, a'])$$

Thus $[1_{(n)}, a] = [1_{(n)}, a']$ and $a = a'$, i.e. τ is a cartesian natural transformation. □

Acknowledgements M. Zawadowski wants to thank M.M. Clementino for bringing to his attention the work of E. G. Manes in general and his notion of a taut monad in particular and P.T. Johnstone for referring him to [6].

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