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A Generalized Grothendieck Inequality and Nonlocal Correlations that Require High Entanglement

Jop Briët^{1,*}, Harry Buhrman^{1,*}, Ben Toner^{2,**}

¹ Centrum Wiskunde & Informatica (CWI), and University of Amsterdam, Science Park 123,

1098 SJ Amsterdam, The Netherlands. E-mail: j.briet@cwi.nl; jop.briet@cwi.nl; buhrman@cwi.nl

² School of Physics, The University of Melbourne, Victoria 3010, Australia. E-mail: bentoner@bentoner.com

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Abstract: Suppose that Alice and Bob make local two-outcome measurements on a shared entangled quantum state. We show that, for all positive integers d, there exist correlations that can only be reproduced if the local Hilbert-space dimension is at least d. This establishes that the amount of entanglement required to maximally violate a Bell inequality must depend on the number of measurement settings, not just the number of measurement outcomes. We prove this result by establishing a lower bound on a new generalization of Grothendieck's constant.

1. Introduction

Grothendieck's inequality first arose in the study of norms on tensor products of Banach spaces [Gro53]. It has since found many applications in mathematics and computer science, including approximation algorithms [AN04, CW04, BOFV10a] and communication complexity [LS07, LMSS08]. In quantum information, Grothendieck's inequality quantifies the difference between the classical and quantum values of certain simple Bell inequalities, as established by Tsirelson [Tsi87]. Tsirelson's work has been the starting point of considerable recent research into quantum nonlocality [CHTW04, AGT06, RT07, BPA⁺08, PGWP⁺08, BBLV09, JPPG⁺10].

We start by stating the inequality in its strongest form, in terms of the real Grothendieck constant K_G .

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Definition 1. The **real Grothendieck constant of order** n, is the smallest real number $K_G(n)$ such that: For all positive integers r and for all real $r \times r$ matrices $M = (M_{ij})$, the inequality

$$\max_{\substack{\boldsymbol{a}_1,\dots,\boldsymbol{a}_r\\\boldsymbol{b}_1,\dots,\boldsymbol{b}_r}} \sum_{i,j} M_{ij} \boldsymbol{a}_i \cdot \boldsymbol{b}_j \le K_G(n) \max_{\substack{\alpha_1,\dots,\alpha_r\\\beta_1,\dots,\beta_r}} \sum_{i,j} M_{ij} \alpha_i \beta_j \tag{1}$$

holds, where the maximum on the left-hand side is taken over all sequences $a_1, \ldots, a_r, b_1, \ldots, b_r$ of n-dimensional real unit vectors, $a_i \cdot b_j$ denotes the Euclidean inner product of a_i and b_j , and the maximum on the right-hand side is taken over all sequences $\alpha_1, \ldots, \alpha_r, \beta_1, \ldots, \beta_r$ of real numbers in the set $\{-1, +1\}$.

The real Grothendieck constant, *denoted* K_G , *is defined as* $\lim_{n\to\infty} K_G(n)$.

The exact value of K_G is unknown. The tightest version of the inequality known is due to Krivine [Kri79], who proved that $K_G \leq \pi/(2\ln(1 + \sqrt{2})) \approx 1.78$. Davie [Dav84] and, independently, Reeds [Ree91] are responsible for the best lower bounds: they showed that $K_G \gtrsim 1.68$. Raghavendra and Steurer have shown that K_G can be approximated within an error η in time exp(exp($O(1/\eta^3))$) [RS09].

In this paper, we give a new generalization of Grothendieck's inequality. We replace the maximization over scalars on the right-hand side of Eq. (1) with a maximization over real unit vectors of dimension m < n. More formally:

Definition 2. Let *m* and *n* be positive integers with m < n. Define $K_G(n \mapsto m)$ to be the smallest real number such that: For all positive integers *r* and for all real $r \times r$ matrices $M = (M_{ij})$, the inequality

$$\max_{\substack{a_1,\dots,a_r\\b_1,\dots,b_r}} \sum_{i,j} M_{ij} a_i \cdot b_j \le K_G(n \mapsto m) \max_{\substack{a_1',\dots,a_r'\\b_1',\dots,b_r'}} \sum_{i,j} M_{ij} a_i' \cdot b_j'$$
(2)

holds, where the maximum on the left-hand side is taken over all sequences $a_1, \ldots, a_r, b_1, \ldots, b_r$ of n-dimensional real unit vectors, and the maximum on the right-hand side is taken over all sequences $a'_1, \ldots, a'_r, b'_1, \ldots, b'_r$ of m-dimensional real unit vectors. This generalizes Definition 1 in the sense that $K_G(n) = K_G(n \mapsto 1)$.

Building on the techniques Grothendieck himself [Gro53] used to prove the original lower bound on K_G , we prove the following lower bound on $K_G(n \mapsto m)$.

Theorem 1. For all m < n,

$$K_G(n \mapsto m) \ge \frac{m}{n} \left(\frac{\Gamma(\frac{m}{2})}{\Gamma(\frac{m+1}{2})} \frac{\Gamma(\frac{n+1}{2})}{\Gamma(\frac{n}{2})} \right)^2 \tag{3}$$

$$= 1 + \frac{1}{2m} - \frac{1}{2n} - O(\frac{1}{m^2}).$$
(4)

After the appearance of this work in preprint form, it was shown in [BOFV10b] that for positive semidefinite matrices M, like the ones considered here, the lower bound on $K_G(\infty \mapsto m)$ given in Theorem 1 is in fact optimal. Note that for general matrices, trivially, we have $K_G(n \mapsto m) \leq K_G(n)$. Better upper bounds on $K_G(\infty \mapsto m)$ for general matrices were found in [BOFV10a]. *Related work.* Definition 2 is but the latest in a long history of generalizations of Grothendieck's inequality. Previously, Grothendieck's inequality has been generalized as follows:

- Replacing the real scalars, vectors and matrices with complex ones results in the definition of the *complex Grothendieck constant*.
- Restricting to positive definite matrices *M* results in a tighter inequality [Rie74, Nes98].
- Rather than proving inequalities that hold for all matrices, we can prove inequalities that only hold for all matrices M of some fixed size, say $r \times s$. This refinement has been studied by Fishburn and Reeds [FR94], and results in the definition of a constant which they denote $K_G(r, s)$, not to be confused with our $K_G(n \mapsto m)$.
- Observe that Eq. (1) has a bipartite structure, in the following sense: on the lefthand side, the sum is of inner products $a_i \cdot b_j$ of a vector from $\{a_1, \ldots, a_r\}$ with a vector from $\{b_1, \ldots, b_r\}$; there are no inner products $a_i \cdot a_j$ or $b_i \cdot b_j$. A similar observation applies to the right-hand side. So if we consider a graph with vertices labelled by the vectors a_i and b_j , and draw an edge between vertex a_i and b_j whenever $M_{ij} \neq 0$, then the resulting "interaction graph" is bipartite. Alon et al. have generalized Grothendieck's inequality to general graphs that are not necessarily bipartite [AMMN06].

Application to quantum correlations. Suppose that Alice and Bob share an entangled quantum state and that each performs a local two-outcome measurement on their part of that state. We are interested in the (classical) *correlations* between their measurement outcomes, i.e., the probability that they obtain the same measurement outcome, for various choices of the local measurements. It is a well-established fact (both theoretically [Bel64,CHSH69] and experimentally [AGR81,AGR82,ADR82,TBZG98,RKM⁺01]), that there are entangled states for which these correlations are *nonlocal*, meaning that they cannot be explained in a classical universe (more formally, they are inconsistent with all *local hidden variable models*).

As a corollary of Theorem 1, we show that there are nonlocal quantum correlations that require entangled states with local support on a Hilbert space of dimension at least d, for any d (even with arbitrary shared randomness). This confirms a conjecture of Brunner et al. [BPA⁺08], proving that what they term *dimension witnesses* exist with binary outcomes. In particular, Brunner et al. conjectured that two-outcome measurements may be sufficient to test the dimension of any bipartite quantum system. This work establishes that their conjecture is true. This strengthens the fundamental result described in the preceding paragraph: the fact that there exist nonlocal correlations is the d = 1 case of our result.

Brunner et al. pointed out that the same result would follow if one could prove that the Grothendieck constants $K_G(n)$ are strictly increasing in n. This is plausible but we do not know how to prove it. Our proof sidesteps this issue.

In addition to the work of Brunner et al., there is some other work on lower-bounding the amount of entanglement required to reproduce certain correlations. Pál and Vértesi construct correlations that cannot be reproduced if each party has only a smalldimensional quantum system [PV08, PV09]. Wehner, Christandl and Doherty show how to obtain lower bounds using information-theoretic arguments [WCD08].

Another lower bound follows from communication complexity, but for the more general problem of reproducing correlations of measurements with more than two outcomes. The Hidden Matching quantum communication complexity problem (HM(n))

[BYJK08] can be formulated as a nonlocal correlation, where a maximally entangled state of dimension *n* is used to reproduce the correlations perfectly. On the other hand, using the classical bounded error one-way communication complexity lower bound for HM(*n*), it follows that one needs $\omega(\sqrt{n})$ bits of one-way communication to approximately reproduce these correlations classically. This in turn yields a lower bound on the dimension of the entangled state of $\sqrt{n}/\log n$ for any quantum strategy that approximates these correlations. This follows because any smaller dimensional state can be used to establish a classical one-way protocol that approximates these correlations and uses less than $\omega(\sqrt{n})$ bits of communication, by simply communicating a classical description of an approximation of the state that Bob has after Alice did her measurement. (See also [BRSW10].)

Outline. The paper is structured as follows. We define notation in Sect. 2. In Sect. 3, we rework the definition of $K_G(n \mapsto m)$ in order to work in the limit $r \to \infty$, which makes things simpler. Then, in Sect. 4, we prove our main result, Theorem 1. In Sect. 5, we describe the consequences for quantum nonlocality. Readers wishing to skip the details of the proof can read Sect. 5 immediately after Sect. 3.

Note. After a preliminary version of this paper was submitted to the twelfth workshop on Quantum Information Processesing (QIP 2009) on 20 October, 2008, we learned of a paper by Vértesi and Pál [VP09], who obtain similar results independently. Without explicitly defining $K_G(n \mapsto m)$, they prove that this quantity is strictly increasing with m when $n \to \infty$ using essentially the same methods that we do, and use this result to confirm Brunner et al.'s conjecture, giving a Bell inequality that has an infinite number of measurement settings; they obtain dimension witnesses with finite measurement settings using different methods. In this paper, we obtain dimension witnesses with a finite number of measurement settings from our lower bounds on $K_G(n \mapsto m)$.

2. Notation

We write [n] for the set $\{1, ..., n\}$. The unit sphere in \mathbb{R}^n is denoted S_{n-1} . We write da for the Haar measure on S_{n-1} , normalized such that $\int da = 1$. The norm ||a|| of a vector a is always the Euclidean norm. In the Introduction and Sect. 6, subscripts label vectors; in the remainder of the paper, subscripts on a vector denote its components. Variables in lowercase roman type will typically be vectors on the unit sphere; variables in lowercase Greek type will typically be scalars.

3. An Equivalent Definition of $K_G(n \mapsto m)$

To establish a lower bound on $K_G(n \mapsto m)$ per Eq. (2), we need to exhibit an $r \times r$ matrix M and then calculate (or at least bound) both sides of Eq. (2). We will work in the limit $r \to \infty$ and so we start by showing that this indeed allows us to lower bound $K_G(n \mapsto m)$.

Lemma 1. The constant $K_G(n \mapsto m)$ is bounded by

$$K_G(n \mapsto m) \ge \sup_{M:S_{n-1} \times S_{n-1} \to [-1,1]} \left(\frac{1}{D(M)} \int dadb M(a,b) a \cdot b \right), \tag{5}$$

where the supremum is over measurable functions $M : S_{n-1} \times S_{n-1} \rightarrow \mathbb{R}$ with $\int dadb |M(a, b)| = 1$ and the denominator

$$D(M) = \max_{A,B:S_{n-1}\to S_{m-1}} \int dadb M(a,b) A(a) \cdot B(b),$$
(6)

with the maximum over functions $A, B : S_{n-1} \rightarrow S_{m-1}$.

Observe that we can rewrite the conventional definition of $K_G(n \mapsto m)$ as

$$K_G(n \mapsto m) = \lim_{r \to \infty} \sup_{M_{ij}} \left(\frac{\max_{(a_i, b_j)} \sum_{i, j} M_{ij} a_i \cdot b_j}{\max_{(a'_i, b'_j)} \sum_{i, j} M_{ij} a'_i \cdot b'_j} \right).$$

The following lemma implies Lemma 1.

Lemma 2. For any measurable function $M'(\boldsymbol{a}, \boldsymbol{b})$ with $\boldsymbol{a}, \boldsymbol{b} \in S_{n-1}$, and any $\eta > 0$, there exist an r and matrix $M_{ij} \in \mathbb{R}^r \times \mathbb{R}^r$, such that

$$\frac{\max_{(\boldsymbol{a}_{i},\boldsymbol{b}_{j})}\sum_{i,j}M_{ij}\boldsymbol{a}_{i}\cdot\boldsymbol{b}_{j}}{\max_{(\boldsymbol{a}_{i}',\boldsymbol{b}_{j}')}\sum_{i,j}M_{ij}\boldsymbol{a}_{i}'\cdot\boldsymbol{b}_{j}'} \geq \frac{\int d\boldsymbol{a}\,d\boldsymbol{b}M'(\boldsymbol{a},\boldsymbol{b})\boldsymbol{a}\cdot\boldsymbol{b}}{\max_{A,B:S_{n-1}\to S_{m-1}}\int d\boldsymbol{a}\,d\boldsymbol{b}M'(\boldsymbol{a},\boldsymbol{b})A(\boldsymbol{a})\cdot\boldsymbol{B}(\boldsymbol{b})} - \eta.$$
(7)

We prove Lemma 2 using an ε -net, which is defined as follows.

Definition 3 (ε -net). For fixed $\varepsilon > 0$, a set of vectors $E_n^{\varepsilon} = \{w_1, w_2, \dots \in S_{n-1}\}$ is an ε -net for S_{n-1} if for all $a \in S_{n-1}$, there exists a vector $u \in E$ that satisfies $||a - u||_2 \le \varepsilon$.

Lemma 3. For $0 < \varepsilon < 1$, there is an ε -net for S_{n-1} with $|E_n^{\varepsilon}| = (3/\varepsilon)^n$.

Proof. We follow a standard volume argument (see e.g. [Pis99, Lemma 4.10]). Let E_n^{ε} be a maximal set of vectors satisfying $\|\boldsymbol{u} - \boldsymbol{v}\|_2 \ge \varepsilon$ for all $\boldsymbol{u}, \boldsymbol{v} \in E_n^{\varepsilon}$, where the existence of such a set is guaranteed by Zorn's Lemma. Then E_n^{ε} is an ε -net for S_{n-1} . We bound $|E_n^{\varepsilon}|$ using a volume argument. The open balls of radius $\varepsilon/2$ around each point $\boldsymbol{u} \in E_n^{\varepsilon}$ are pairwise disjoint and all contained in the ball of radius $1 + \varepsilon/2$ about the origin. Hence

$$|E_n^{\varepsilon}| \le \frac{(1+\varepsilon/2)^n}{(\varepsilon/2)^n} = \left(\frac{2}{\varepsilon}+1\right)^n \le \left(\frac{3}{\varepsilon}\right)^n.$$
(8)

Proof (of Lemma 2). For some $\varepsilon > 0$ let $E_n^{\varepsilon} = \{w_1, w_2, ...\} \subseteq S_{n-1}$ be an ε -net for S_{n-1} . Suppose we use this net, to divide the unit *n*-sphere up into $|E_n^{\varepsilon}|$ disjoint regions: For $w \in E_n^{\varepsilon}$, let R_w be the set of points on S_{n-1} that are closer to w than to any other point in E_n^{ε} , and assign points equidistant to two or more points in the net in some arbitrary way. We set

$$M_{ij} = \int_{a \in R_{w_i}} \int_{b \in R_{w_j}} dadb M'(a, b).$$

Define $\delta := \max_{i,j} \{ |\boldsymbol{w}_i \cdot \boldsymbol{w}_j - \boldsymbol{a} \cdot \boldsymbol{b}| | \boldsymbol{a} \in R_{\boldsymbol{w}_i}, \boldsymbol{b} \in R_{\boldsymbol{w}_j} \}$. Then by the triangle inequality, we have

$$\left| \sum_{i,j} \int_{a \in R_{w_i}} da \int_{b \in R_{w_j}} db M(a, b) \left(w_i \cdot w_j - a \cdot b \right) \right|$$

$$\leq \sum_{i,j} \int_{a \in R_{w_i}} da \int_{b \in R_{w_j}} db |M'(a, b)| |w_i \cdot w_j - a \cdot b| \leq \delta$$

Hence, for the numerators we get:

$$\max_{(\boldsymbol{a}_i,\boldsymbol{b}_j)}\sum_{i,j}M_{ij}\boldsymbol{a}_i\cdot\boldsymbol{b}_j\geq\int d\boldsymbol{a}\,d\boldsymbol{b}M'(\boldsymbol{a},\boldsymbol{b})\boldsymbol{a}\cdot\boldsymbol{b}-\delta.$$

For the denominators, we have

$$\max_{A,B:S_{n-1}\to S_{m-1}}\int da\,dbM'(a,b)A(a)\cdot B(b)\geq \max_{(a'_i,b'_j)}\sum_{i,j}M_{ij}a'_i\cdot b'_j$$

since we can always pick $A(a) = a'_i$ and $B(b) = b'_j$ for all $a \in R_i$ and $b \in R_j$. The result follows from the fact that we can let δ become arbitrarily small by decreasing ε . To see this, note that by the Cauchy-Schwarz inequality, we have for every $i, j \in \{1, ..., |E_n^{\varepsilon}|\}$ and $a \in R_{w_i}$ and $b \in R_{w_j}$,

Hence, we have $\delta \leq 2\varepsilon$.

Although the inequality in Eq. (5) is sufficient for our need, it is not hard to see that in fact equality holds. We give a formal proof of this fact in Sect. 6.

4. Lower Bound on $K_G(n \mapsto m)$

We prove Theorem 1 by considering a specific example due to Grothendieck himself [Gro53]: For $a, b \in S_{n-1}$, take $M(a, b) = a \cdot b$.

We start by calculating the denominator D(M). To do this, we need to work out which embeddings $A, B : S_{n-1} \rightarrow S_{m-1}$ achieve the maximum in Eq. (6). It turns out that the maximum is achieved when A and B are equal. Informally, we should try to preserve as much of the structure of S_{n-1} as possible, and it is natural to conjecture that the best embedding is a projection onto an *m*-dimensional subspace. This is indeed the case. We prove this in the following lemma.

Lemma 4. For the function $M(a, b) = a \cdot b$, the optimal embedding $A : S_{n-1} \rightarrow S_{m-1}$ is a projection. In particular, the denominator D(M) is given by

$$D = \frac{1}{m} \left(\int d\boldsymbol{a} \left(\sum_{i=1}^{m} a_i^2 \right)^{1/2} \right)^2, \tag{9}$$

where a_1, \ldots, a_n are the components of **a**.

Proof. We prove this result in two steps. First, we show that the maximum is achieved by a *weighted* projection. Second, we show that the best projection is the one with uniform weights.

We need to calculate

$$D(M) = \max_{A,B:S_{n-1} \to S_{m-1}} \int dadb M(a, b) A(a) \cdot B(b),$$
(10)

with the maximum over functions $A, B : S_{n-1} \to S_{m-1}$. For $M(a, b) = a \cdot b$, we can write

$$(\boldsymbol{a} \cdot \boldsymbol{b}) (A(\boldsymbol{a}) \cdot B(\boldsymbol{b})) = (\boldsymbol{a} \otimes A(\boldsymbol{a})) \cdot (\boldsymbol{b} \otimes B(\boldsymbol{b})), \qquad (11)$$

(this trick is motivated by a similar one used by Krivine in proving his upper bound on K_G [Kri79]), which allows us to write D(M) as a maximization over the inner product of two vectors,

$$D = \max_{A,B:S_{n-1}\to S_{m-1}} \left(\int da \, a \otimes A(a) \right) \cdot \left(\int db \, b \otimes B(b) \right)$$
(12)

$$= \max_{A:S_{n-1}\to S_{m-1}} \left\| \int da \, a \otimes A(a) \right\|^2, \tag{13}$$

where the second equality follows from the fact that the inner product is maximized when vectors are parallel. Let $\int da \, a \otimes A(a) = \chi v$, where v is an (n+m)-dimensional unit vector and $\chi \ge 0$ is what we want to maximize. Applying the singular value decomposition—known in quantum information theory as the Schmidt decomposition (see, for example [NC00])—we can write

$$\boldsymbol{v} = \sum_{i=1}^{m} \sqrt{\gamma_i} \, \boldsymbol{x}_i \otimes \boldsymbol{y}_i, \tag{14}$$

where, for each $i \in [m]$, $\gamma_i \ge 0$, $\sum_i \gamma_i = 1$, and $\{x_1, \ldots, x_m\}$ and $\{y_1, \ldots, y_m\}$ are orthonormal sets in \mathbb{R}^n and \mathbb{R}^m respectively. Therefore, in order to maximize

$$\chi = \mathbf{v} \cdot \int d\mathbf{a} \, \mathbf{a} \otimes A(\mathbf{a})$$

= $\int d\mathbf{a} \sum_{i} \sqrt{\gamma_{i}} (\mathbf{a} \cdot \mathbf{x}_{i}) (A(\mathbf{a}) \cdot \mathbf{y}_{i})$
= $\int d\mathbf{a} A(\mathbf{a}) \cdot \left(\sum_{i} \sqrt{\gamma_{i}} (\mathbf{a} \cdot \mathbf{x}_{i}) \mathbf{y}_{i}\right),$ (15)

we should choose A(a) to be

$$\frac{\sum_{i} \sqrt{\gamma_{i}} (\boldsymbol{a} \cdot \boldsymbol{x}_{i}) \boldsymbol{y}_{i}}{\|\sum_{i} \sqrt{\gamma_{i}} (\boldsymbol{a} \cdot \boldsymbol{x}_{i}) \boldsymbol{y}_{i}\|} = \frac{\sum_{i} \sqrt{\gamma_{i}} (\boldsymbol{a} \cdot \boldsymbol{x}_{i}) \boldsymbol{y}_{i}}{(\sum_{i} \gamma_{i} (\boldsymbol{a} \cdot \boldsymbol{x}_{i})^{2})^{1/2}},$$
(16)

a weighted projection onto some *m*-dimensional subspace, the particular choice of which does not matter. Substituting this into Eq. (15) and then Eq. (13) and choosing a basis for \mathbb{R}^n by extending x_1, \ldots, x_m so that $a_i = a \cdot x_i$ establishes that

$$D = \max_{A,B:S_{n-1}\to S_{m-1}} \int dadb(a \cdot b) \left(A(a) \cdot B(b)\right) = \left(\chi(\gamma_1,\ldots,\gamma_m)\right)^2, \quad (17)$$

where

$$\chi(\gamma_1,\ldots,\gamma_m) = \int da \left(\sum_{i=1}^m \gamma_i a_i^2\right)^{1/2}.$$
 (18)

It remains to show that weights γ_i can be taken to be equal. To prove this, suppose that χ is maximized by $(\gamma_1^*, \gamma_2^*, \dots, \gamma_m^*)$. Then, by symmetry, the maximum is also achieved by $(\gamma_2^*, \gamma_1^*, \dots, \gamma_m^*)$, and indeed, by any other permutation σ of the γ_i^* . Hence

$$\chi(\gamma_1^*,\ldots,\gamma_m^*) = \frac{1}{m!} \sum_{\sigma} \chi(\gamma_{\sigma(1)}^*,\ldots,\gamma_{\sigma(m)}^*)$$
(19)

$$= \frac{1}{m!} \sum_{\sigma} \int da \left(\sum_{i=1}^{m} \gamma_{\sigma(i)} a_i^2 \right)^{1/2}$$
(20)

$$= \int da \frac{1}{m!} \sum_{\sigma} \left(\sum_{i=1}^{m} \gamma_{\sigma(i)}^* a_i^2 \right)^{1/2}$$
(21)

$$\leq \int da \left(\frac{1}{m!} \sum_{\sigma} \sum_{i=1}^{m} \gamma_{\sigma(i)}^* a_i^2 \right)^{1/2} \tag{22}$$

by Jensen's inequality and the concavity of $(\cdot)^{1/2}$. But the coefficient of a_i^2 in this expression is just

$$\frac{1}{m!} \sum_{\sigma} \gamma_{\sigma(i)}^* = \frac{1}{m} \sum_{i} \gamma_i^* = \frac{1}{m} \times 1 = \frac{1}{m}.$$
(23)

Thus the maximum is achieved by uniform weights.

With Lemma 4 in hand, the proof of Theorem 1 is straightforward.

Proof (of Theorem 1). Take $M(a, b) = a \cdot b$ in Lemma 5. It follows from Lemma 4 that

$$K_G(n \mapsto m) \ge \frac{m}{n} \left(\frac{Y_n}{Y_m}\right)^2,$$
 (24)

where

$$Y_k := \int_{\boldsymbol{a} \in S_{n-1}} \boldsymbol{d} \boldsymbol{a} \left(\sum_{i=1}^k a_i^2\right)^{1/2}$$

and we evaluated the numerator in Eq. (5) by observing that it is the same as the denominator when m = n, and so we already calculated it as a special case of Lemma 4. We can evaluate Y_k using a trick similar to that used to calculate the surface area of the *n*-sphere. Define

$$C_k := \int_{\boldsymbol{a} \in \mathbb{R}^n} d\boldsymbol{a} \left(\sum_{i=1}^k a_i^2 \right)^{1/2} e^{-\|\boldsymbol{a}\|_2^2}.$$

Introducing spherical coordinates, and writing $r = ||a||_2$, we have

$$C_{k} = Y_{k} \Omega_{n} \int_{0}^{\infty} dr r^{n-1} (r^{2})^{1/2} e^{-r^{2}} = Y_{k} \pi^{n/2} \frac{\Gamma\left(\frac{n+1}{2}\right)}{\Gamma\left(\frac{n}{2}\right)},$$
(25)

where Γ is the well-known gamma function, and we pick up a factor of $\Omega_n = 2\pi^{n/2}/\Gamma(n/2)$, the surface area of the unit sphere in *n* dimensions, because of our normalization convention.

On the other hand, we have

$$C_{k} = \int_{-\infty}^{\infty} da_{1} \cdots da_{k} \left(\sum_{i=1}^{k} a_{i}^{2} \right)^{1/2} e^{-(a_{1}^{2} + \dots + a_{k}^{2})} \int_{-\infty}^{\infty} da_{k+1} \cdots da_{n} e^{-(a_{k+1}^{2} + \dots + a_{n})^{2}}.$$
(26)

We can interpret $\left(\sum_{i=1}^{k} a_i^2\right)^{1/2}$ as the norm of a point in *k*-dimensional space, and write the first integral (over *k* variables) as

$$\Omega_k \int_0^\infty dr' r'^k e^{-r'^2} = \frac{2\pi^{k/2}}{\Gamma(\frac{k}{2})} \cdot \frac{\Gamma\left(\frac{k+1}{2}\right)}{2}.$$

The second integral of Eq. (26) is simply $(\sqrt{\pi})^{n-k}$. Comparing these two ways to evaluate C_k , we conclude that

$$Y_k = \frac{\Gamma(\frac{k+1}{2})}{\Gamma(\frac{k}{2})} \cdot \frac{\Gamma(\frac{n}{2})}{\Gamma(\frac{n+1}{2})}$$
(27)

and

$$K_G(n \mapsto m) \ge \frac{m}{n} \left(\frac{Y_n}{Y_m}\right)^2$$
 (28)

$$= \frac{m}{n} \left(\frac{\Gamma(\frac{m}{2})}{\Gamma(\frac{m+1}{2})} \frac{\Gamma(\frac{n+1}{2})}{\Gamma(\frac{n}{2})} \right)^2.$$
(29)

For all integers $1 \le m < n$, this bound is nontrivial, i.e., is strictly greater than 1. This is because the function

$$f(n) = \frac{1}{\sqrt{n}} \frac{\Gamma(\frac{n+1}{2})}{\Gamma(\frac{n}{2})}$$
(30)

is strictly increasing for n = 1, 2, ... (see Sect. 7 for a proof). Asymptotically, we have

$$K_G(n \mapsto m) \ge 1 + \frac{1}{2m} - \frac{1}{2n} - O(\frac{1}{m^2}),$$
(31)

.

where the approximation follows from the asymptotic series (see the answer to Exercise 9.60 in [GKP94])

$$\frac{\Gamma(k+\frac{1}{2})}{\Gamma(k)} = \sqrt{k}(1-\frac{1}{8k}+\frac{1}{128k^2}+\cdots).$$
(32)

5. Quantum Nonlocality

Here we describe the application to quantum nonlocality. Suppose that two parties, Alice and Bob, each have a *d*-dimensional quantum system, described by Hilbert spaces $\mathcal{H}_A \cong \mathbb{C}^d$ and $\mathcal{H}_B \cong \mathbb{C}^d$, respectively. Alice and Bob each make a two-outcome measurement on their own system, resulting in outcomes α , $\beta \in \{\pm 1\}$, respectively. Suppose the set of Alice's possible measurements is M_A , and the set of Bob's possible measurements is M_B . An observable is a Hermitian operator with eigenvalues in $\{\pm 1\}$. Alice's a^{th} possible measurement is specified by an observable A_a on \mathcal{H}_A ; Bob's b^{th} measurement by an observable B_b on \mathcal{H}_B (and all observables specify valid measurements). If the joint system of Alice and Bob is in pure state $|\psi\rangle \in \mathbb{C}^d \otimes \mathbb{C}^d$, then the *joint correlation*—the expectation of the product of Alice and Bob's outcomes, given that Alice performs measurement *a* and Bob measurement *b*—is

$$E[\alpha\beta|ab] = \langle \psi|A_a \otimes B_b|\psi\rangle. \tag{33}$$

In computer science, such correlations are studied in the context of *XOR nonlocal games* [CHTW04, CSUU08].

We say that a set of joint correlations, $\{E[\alpha\beta|ab]: a \in [a_{\max}], b \in [b_{\max}]\}$, is *pured-quantum-realizable* if there is a state $|\psi\rangle \in \mathbb{C}^d \otimes \mathbb{C}^d$ and for all $a \in [a_{\max}]$, there are observables A_a on \mathbb{C}^d and for all $b \in [b_{\max}]$, there are observables B_b on \mathbb{C}^d such that $E[\alpha\beta|ab] = \langle \psi | A_a \otimes B_b | \psi \rangle$. A set of joint correlations is *d-quantum-realizable* if it is a probabilistic mixture of pure-*d*-quantum-realizable correlations (this definition accounts for allowing Alice and Bob to share an arbitrary large amount of shared randomness, use POVMs, and share a mixed state). A set of joint correlations is *finitely quantum-realizable* if there is some *d* such that the correlations are *d*-quantum-realizable. Note that a set of correlations is *local* if they are 0-quantum-realizable.

We prove the following theorem.

Theorem 2. For any d, there are correlations that are finitely quantum-realizable, but which are not d-quantum-realizable.

We now describe the correlations that we use to prove Theorem 2. Fix some integer n. Alice and Bob's possible measurements are parametrized by unit vectors in \mathbb{R}^n , a and b, respectively. (Note that each party here has an infinite number of possible measurements; we'll reprove the theorem with finite sets of measurements in the next subsection.) The joint correlations are given by

$$E[\alpha\beta|ab] = a \cdot b, \tag{34}$$

where $a \cdot b$ is just the Euclidean inner product of a and b. For all n, these correlations are finitely quantum-realizable, as the following result shows.

Lemma 5 (Tsirelson [Tsi87]). Let $|\psi\rangle$ be a maximally entangled state on $\mathbb{C}^d \otimes \mathbb{C}^d$ where $d = 2^{\lfloor n/2 \rfloor}$. Then there are two mappings from unit vectors in \mathbb{R}^n to observables on \mathbb{C}^d , one taking **a** to A_a , the other taking **b** to B_b , such that

$$\langle \psi | A_{\boldsymbol{a}} \otimes B_{\boldsymbol{b}} | \psi \rangle = \boldsymbol{a} \cdot \boldsymbol{b},$$

for all unit vectors **a**, **b**.

To show they are not d-quantum-realizable, we will use the following characterization.

Lemma 6 ([Tsi87,AGT06]). Suppose Alice and Bob measure observables A_a and B_b on a pure quantum state $|\psi\rangle \in \mathbb{C}^d \otimes \mathbb{C}^d$. Then we can associate a real unit vector $A(a) \in \mathbb{R}^{2d^2}$ with A_a (independent of B_b), and a real unit vector $B(b) \in \mathbb{R}^{2d^2}$ with B_b (independent of A_a) such that

$$E[\alpha\beta|ab] = \langle \psi|A_a \otimes B_b|\psi\rangle = A(a) \cdot B(b).$$
(35)

Proof (of Theorem 2). Let $n = 2d^2 + 1$, and consider the joint correlations described in Eq. (34). By Lemma 5, these correlations are finitely quantum-realizable. To show they are not *d*-quantum-realizable, we will show that they lie outside the convex hull of the set of pure-*d*-quantum-realizable correlations.

We do this in the standard way using a Bell inequality, which is a linear function on the vector of correlations, i.e.,

$$B(E[\alpha\beta|ab]) = \int dadb M(a, b) E[\alpha\beta|ab], \qquad (36)$$

for some function M(a, b), where the integral is over all unit vectors a, b in \mathbb{R}^n . Here we take $M(a, b) = a \cdot b$ so as to apply the results of Sect. 4. Substituting for $E[\alpha\beta|ab]$ using Eq. (34), we have

$$B(E[\alpha\beta|ab]) = \int dadb(a \cdot b)^2.$$
(37)

For any pure-*d*-quantum-realizable correlations, by Lemma 6 there are vectors A(a) and B(b) in \mathbb{R}^{2d^2} , such that the resulting correlations are given by

$$E[\alpha\beta|ab]_d = A(a) \cdot B(b).$$
(38)

Evaluating *B* on these correlations, we must have

$$B(E[\alpha\beta|ab]_d) = \int dadb(a \cdot b)A(a) \cdot B(b)$$
⁽³⁹⁾

$$\leq \max_{A,B} \int dadb(a \cdot b) A(a) \cdot B(b)$$
(40)

$$\leq \frac{1}{K_G^{\leq}(n \mapsto 2d^2)} \int dadb (a \cdot b)^2 \tag{41}$$

$$\frac{1}{K_{G}^{\leq}(n\mapsto 2d^{2})}B(E[\alpha\beta|\boldsymbol{ab}]),\tag{42}$$

where $K_{\overline{G}}^{\leq}(n \mapsto 2d^2)$ is our lower bound on $K_G(n \mapsto 2d^2)$. Since we chose $n = 2d^2+1$, $K_{\overline{G}}^{\leq}(n \mapsto n-1) > 1$ by Theorem 1. Hence $B(E[\alpha\beta|ab]_d) < B(E[\alpha\beta|ab])$ for all pure-*d*-quantum-realizable correlations, which implies that $E[\alpha\beta|ab]$ lies outside the convex hull of vectors of the form $E[\alpha\beta|ab]_d$. We conclude that the correlations in Eq. (34) are not *d*-quantum-realizable.

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We conclude this section by noting a slight variant of our result. Suppose we only wish to *approximately* reproduce some set of quantum correlations, which were originally produced by performing two-outcome measurements on an entangled state with support $\mathbb{C}^D \otimes \mathbb{C}^D$. Then this can be done with a constant amount of entanglement, where the constant depends on the desired accuracy of the approximation, but not on the dimension D [CHTW09]. In other words, in order to strengthen Theorem 6 to construct correlations that are not approximately d-quantum-realizable, we would need to look beyond two-outcome measurements.

5.1. Reducing the number of questions. A possible objection to the example above is that the number of questions is taken to be infinite. Here we reduce to a finite number of questions in the most straightforward way possible, by discretizing the unit *n*-sphere using an ε -net.

Theorem 3. There are two-party two-outcome correlations with $\exp(poly(d))$ measurement settings that are finitely quantum-realizable but not *d*-quantum-realizable.

Proof. To convert the quantum correlations of Eq. (34) into correlations with only a finite number of settings, fix $0 < \varepsilon < 1$ (to be chosen later) and let E_n^{ε} be an ε -net for S_{n-1} with $(3/\varepsilon)^n$ settings, the existence of which is guaranteed by Lemma 3. Consider the following correlations: Alice's set of possible measurements is E_n^{ε} , and so is Bob's (note that we implicitly apply Lemma 5 here). If Alice performs a measurement $\mathbf{u} \in E_n^{\varepsilon}$ and Bob a measurement $\mathbf{v} \in E_n^{\varepsilon}$, the joint correlation should satisfy

$$E[\alpha\beta|\boldsymbol{u}\boldsymbol{v}] = \boldsymbol{u}\cdot\boldsymbol{v},\tag{43}$$

just as in our earlier example. These correlations, being a subset of those considered above, are finitely quantum-realizable.

As in the proof of Lemma 2, we use the ε -net to divide the unit sphere into $|E_n^{\varepsilon}|$ regions: For $u \in E_n^{\varepsilon}$, let R_u be the set of points on S_{n-1} that are closer to u than to any other point in E_n^{ε} , and assign points equidistant to two or more points in the net in some arbitrary way. Consider the Bell inequality

$$B_{\text{finite}}(E[\alpha\beta|uv]) = \sum_{u \in E_n^{\varepsilon}} \int_{a \in R_u} da \sum_{v \in E_n^{\varepsilon}} \int_{b \in R_v} db (a \cdot b) E[\alpha\beta|uv].$$
(44)

Evaluating this on the correlations $E[\alpha\beta|uv] = u \cdot v$, we obtain

$$B_{\text{finite}}(E[\alpha\beta|uv]) = \sum_{u \in E_n^{\varepsilon}} \int_{a \in R_u} da \sum_{v \in E_n^{\varepsilon}} \int_{b \in R_v} db \, (a \cdot b)(u \cdot v)$$
(45)

$$\geq -2\varepsilon + \int da \int db (a \cdot b)^2 \tag{46}$$

$$= B(E[\alpha\beta|ab]) - 2\varepsilon, \tag{47}$$

where we used

$$\boldsymbol{u} \cdot \boldsymbol{v} = \boldsymbol{a} \cdot \boldsymbol{b} + (\boldsymbol{u} - \boldsymbol{a}) \cdot \boldsymbol{b} + \boldsymbol{u} \cdot (\boldsymbol{v} - \boldsymbol{b})$$
(48)

$$\geq a \cdot b - \|u - a\|_2 - \|v - b\|_2 \tag{49}$$

$$\geq \boldsymbol{a} \cdot \boldsymbol{b} - 2\varepsilon,\tag{50}$$

and related the value of the Bell inequality with a finite number of settings to the value of our earlier Bell inequality with infinite settings using Eq. (37).

Now consider a pure *d*-dimensional quantum strategy. Let A(u) be the $2d^2$ dimensional real unit vector associated with Alice's measurement u and B(v) be the vector associated with Bob's measurement v by Lemma 6 and suppose that the resulting correlations are $E[\alpha\beta|uv]_d$. We need to show that $B_{\text{finite}}(E[\alpha\beta|uv]_d)$ is strictly smaller than $B_{\text{finite}}(E[\alpha\beta|uv])$. We'll do this by relating $B_{\text{finite}}(E[\alpha\beta|uv]_d)$ to $B(E[\alpha\beta|ab])$ and applying the inequality proved in Eq. (42). The finite strategy (i.e., the mappings A(u) and B(v)) induces a strategy for the correlations where we had an infinite number of questions. Recall that a strategy in the infinite case was equivalent to a mapping from S_{n-1} to unit vectors in \mathbb{R}^{2d^2} by Lemma 6. Consider the mapping defined as follows: First map a to the closest point u in the ε -net, then to the vector A(u), and similarly for Bob's strategy. For this strategy:

$$B(E[\alpha\beta|ab]_d) = \sum_{\boldsymbol{u}\in E_n^{\varepsilon}} \int_{\boldsymbol{a}\in R_u} d\boldsymbol{a} \sum_{\boldsymbol{v}\in E_n^{\varepsilon}} \int_{\boldsymbol{b}\in R_v} d\boldsymbol{b} (\boldsymbol{a}\cdot\boldsymbol{b})A(\boldsymbol{u})\cdot B(\boldsymbol{v})$$
(51)

$$=\sum_{\boldsymbol{u}\in E_n^{\varepsilon}}\int_{\boldsymbol{a}\in R_{\boldsymbol{u}}}d\boldsymbol{a}\sum_{\boldsymbol{v}\in E_n^{\varepsilon}}\int_{\boldsymbol{b}\in R_{\boldsymbol{v}}}d\boldsymbol{b}\,(\boldsymbol{a}\cdot\boldsymbol{b})E[\alpha\beta|\boldsymbol{u}\boldsymbol{v}]$$
(52)

$$= B_{\text{finite}}(E[\alpha\beta|\boldsymbol{u}\boldsymbol{v}]_d).$$
(53)

Combining the above calculations with Eq. (42), we calculate

$$B_{\text{finite}}(E[\alpha\beta|\boldsymbol{u}\boldsymbol{v}]_d) = B(E[\alpha\beta|\boldsymbol{a}\boldsymbol{b}]_d)$$
(54)

$$\leq \frac{1}{K_{\overline{G}}^{\leq}(n\mapsto 2d^2)}B(E[\alpha\beta|\boldsymbol{ab}]).$$
(55)

Take $n = 2d^2 + 2$ (for ease of calculation, so that the Γ -functions cancel), and we have

$$K_{\bar{G}}^{\leq}(2d^2 + 2 \mapsto 2d^2) = 1 + \Delta_d,$$
 (56)

where $\Delta_d = 1/(4d^2(d^2 + 1))$.

Now we can put everything together. Starting with Eq. (55), we have

$$B_{\text{finite}}(E[\alpha\beta|\boldsymbol{u}\boldsymbol{v}]_d) \le \frac{1}{1+\Delta_d} B(E[\alpha\beta|\boldsymbol{a}\boldsymbol{b}])$$
(57)

$$< \left(1 - \frac{\Delta_d}{2}\right) B(E[\alpha\beta|ab])$$
(58)

$$= B(E[\alpha\beta|ab]) - \frac{\Delta_d}{2n},$$
(59)

where (58) is valid because $0 < \Delta_d < 1$, and in Eq. (59) we substituted

$$B(E[\alpha\beta|ab]) = \int da \int db(a \cdot b)^2 = \frac{1}{n}Y_n^2 = \frac{1}{n},$$
(60)

as calculated in Sect. 4.

Thus, so long as we choose $2\varepsilon \leq \Delta_d/(2n)$, we can combine Eqs. (59) and (47) to obtain

$$B_{\text{finite}}(E[\alpha\beta|\boldsymbol{u}\boldsymbol{v}]_d) < B(E[\alpha\beta|\boldsymbol{a}\boldsymbol{b}]) - \frac{\Delta_d}{2n} \le B(E[\alpha\beta|\boldsymbol{a}\boldsymbol{b}]) - 2\varepsilon \le B_{\text{finite}}(E[\alpha\beta|\boldsymbol{u}\boldsymbol{v}]),$$
(61)

i.e., we have proved that $B_{\text{finite}}(E[\alpha\beta|uv]_d) < B_{\text{finite}}(E[\alpha\beta|uv])$ meaning that our correlations are not *d*-quantum-realizable. The number of measurement settings is

$$(3/\varepsilon)^n = (12n/\Delta_d)^n = \exp(O(d^2\log d)).$$
(62)

6. Proof of Equality in Lemma 1

The fact that equality holds in the inequality in Eq. (5) of Lemma 1 follows directly from the following proposition.

Proposition 1. For all positive integers r, for all $r \times r$ real-valued matrices M_{ij} , and any $\epsilon > 0$, there exists a measurable function $M' : S_{n-1} \times S_{n-1} \rightarrow \mathbb{R}$ with $\int dadb |M'(a, b)| = 1$, such that

$$\frac{\int da \, db M'(a, b) a \cdot b}{\max_{A, B: S_{n-1} \to S_{m-1}} \int da \, db M'(a, b) A(a) \cdot B(b)} \ge \frac{\max_{(a_i, b_j)} \sum_{i, j} M_{ij} a_i \cdot b_j}{\max_{(a'_i, b'_j)} \sum_{i, j} M_{ij} a'_i \cdot b'_j} - \epsilon.$$
(63)

Proof. Let $f, g: [r] \to S_{n-1}$ and $f', g': [r] \to S_{m-1}$ be vector valued functions, and let f^* and g^* be such that they give a sequence $(f^*(i), g^*(j))_{i,j=1}^r = (a_i^*, b_j^*)_{i=1}^r$ that maximizes $\sum_{i,j} M_{ij} a_i \cdot b_j$. Set

$$M'(\boldsymbol{a},\boldsymbol{b}) = \sum_{i,j} M_{ij} \delta(\boldsymbol{a} - \boldsymbol{a}_i^*) \delta(\boldsymbol{b} - \boldsymbol{b}_j^*),$$

where $\delta(\cdot)$ denotes the Dirac delta function. This causes the numerators of (63) to be equal. For the denominator of the left-hand side of (63), we have

$$\max_{A,B:S_{n-1}\to S_{m-1}} \int da db M'(a,b) A(a) \cdot B(b)
= \max_{A',B':(f^{*}(i),g^{*}(j))\to S_{m-1}} \sum_{i,j} M_{ij} A'(f^{*}(i)) \cdot B'(g^{*}(j))
\leq \max_{f',g'} \sum_{i,j} M_{ij} f'(i) \cdot g'(j),$$
(64)

where the inequality follows because the second maximization is over a subset of the set that the last equation is maximized over. The problem that the delta function is a distribution and not really a function may be remedied by considering an approximation to it sufficient to obtain the inequality to within ϵ . All that remains is to rescale M'(a, b) so that it satisfies $\int dadb |M'(a, b)| = 1$. This gives the result.

7. Proof that the Bound on $K_G(n \mapsto m)$ is Nontrivial

Here we establish the following lemma.

Lemma 7. The function

$$f(n) = \frac{1}{\sqrt{n}} \frac{\Gamma(\frac{n+1}{2})}{\Gamma(\frac{n}{2})}$$
(65)

is strictly increasing on integers n = 1, 2, ...

Proof. For $n \le 9$, just evaluate f(n). For n > 9, we use the following bound on log $\Gamma(x)$, first proved by Robbins [Rob55] for integer values of x, but which Matsunawa observed [Mat76, Remark 4.1] is also valid for real values of $x \ge 2$:

$$\sqrt{2\pi}x^{x+1/2}e^{-x+1/(12x+1)} < \Gamma(x+1) < \sqrt{2\pi}x^{x+1/2}e^{-x+1/(12x)}.$$
(66)

Using this bound, we obtain

$$\log \frac{f(n+1)}{f(n)} = -\frac{1}{2} \log \left(1 + \frac{1}{n}\right) + \log \frac{n}{2} + 2 \log \Gamma(\frac{n}{2}) - 2 \log \Gamma(\frac{n+1}{2})$$
$$\geq -\frac{1}{2} \log \left(1 + \frac{1}{n}\right) + \log \left(1 + \frac{1}{n/2 - 1}\right) - n \log \left(1 + \frac{1}{n-2}\right)$$
$$+ \frac{2}{6n - 11} - \frac{2}{6n - 6}.$$

Now use

$$\frac{1}{n} - \frac{1}{2n^2} + \frac{1}{3n^3} - \frac{1}{4n^4} \le \log\left(1 + \frac{1}{n}\right) \le \frac{1}{n} - \frac{1}{2n^2} + \frac{1}{3n^3},\tag{67}$$

(which is valid for all $n \ge 1$), and we obtain

$$\log \frac{f(m+10)}{f(m+9)} \ge \left(14m^7 + 679m^6 + 13923m^5 + 155346m^4 + 1005620m^3 + 3684139m^2 + 6679947m + 3828140\right) / \left(12(m+7)^4(m+8)(m+9)^3(6m+43)\right),$$

which is obviously positive when $m \ge 0$, i.e., when $n \ge 9$. Thus f(n) is strictly increasing.

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