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Toughness and the existence of fractional *k*-factors of graphs $\stackrel{\text{transform}}{\to}$

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Abstract

The toughness of a graph G, t(G), is defined as $t(G) = \min\{|S|/\omega(G-S)|S \subseteq V(G), \omega(G-S) > 1\}$ where $\omega(G-S)$ denotes the number of components of G - S or $t(G) = +\infty$ if G is a complete graph. Much work has been contributed to the relations between toughness and the existence of factors of a graph. In this paper, we consider the relationship between the toughness and the existence of fractional k-factors. It is proved that a graph G has a fractional 1-factor if $t(G) \ge 1$ and has a fractional k-factor if $t(G) \ge k - 1/k$ where $k \ge 2$. Furthermore, we show that both results are best possible in some sense. © 2007 Published by Elsevier B.V.

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1. Introduction

The graphs considered here will be finite undirected graphs which may have multiple edges but no loops. We refer the readers to [2] for the terminologies not defined here. Let G be a graph. We use V(G) and E(G) to denote its vertex set and edge set, respectively. We use G[S] and G - S to denote the subgraph of G induced by S and V(G) - S, respectively, for $S \subseteq V(G)$ and $N_G(x)$ to denote the set of vertices adjacent to x in G. A subset S of V(G) is called a covering set (an independent set) of G if every edge of G is incident with at least (at most) one vertex of S. Let S and T be two disjoint subsets of V(G), we use E(S, T) to denote the set of edges with one end in S and the other end in T and set $S - S' = S \setminus S'$.

Let g and f be two integer-valued functions defined on V(G) with $g(x) \leq f(x)$ for any $x \in V(G)$. A subgraph F of G is called a (g, f)-factor if $g(x) \leq d_F(x) \leq f(x)$ holds for any vertex $x \in V(G)$. A (g, f)-factor is called an [a, b]-factor if $g(x) \equiv a$ and $f(x) \equiv b$. An [a, b]-factor is called a k-factor if a = b = k. Let $h : E(G) \to [0, 1]$ be a function. Let $k \geq 1$ be an integer. If $\sum_{e \geq x} h(e) = k$ holds for any vertex $x \in V(G)$, we call $G[F_h]$ a fractional k-factor of G with indicator function h where $F_h = \{e \in E(G) | h(e) > 0\}$. A fractional 1-factor is also called a fractional perfect matching [7].

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A graph is *t*-tough if for any $S \subseteq V(G)$ and $\omega(G - S) > 1$, we have

 $|S| \ge t\omega(G-S)$

holds where $\omega(G - S)$ denotes the number of components of (G - S). A complete graph is *t*-tough for any positive real number *t*. If *G* is not complete, there exists the largest *t* such that *G* is *t*-tough. This number is denoted by t(G) and is called the toughness of *G*. We define $t(K_n) = +\infty$. If *G* is not complete,

$$t(G) = \min\left\{\frac{|S|}{\omega(G-S)} \middle| S \subseteq V(G) \text{ and } \omega(G-S) \ge 2\right\}$$

The toughness of a graph was first introduced by Chvátal in [3]. Since then, much work has been contributed to the relations between toughness and the existence of factors of a graph. The most famous result is that of [4] which confirms a conjecture stated by Chvátal. Its main result is the following Lemma.

Lemma 1.1. Let G be a graph. If G is k-tough, $|V(G)| \ge k + 1$ and k|V(G)| is even, then G has a k-factor.

The result is sharp since for any positive real number ε , there exists a graph G that has no k-factor with $t(G) \ge k - \varepsilon$ [4]. Katerinis considered toughness and the existence of [a, b]-factors in [5]. In this paper we discuss the relationship between toughness and the existence of fractional k-factors. In [1] Anstee gave a necessary and sufficient condition for a graph to have a fractional (g, f)-factor for which we gave a new proof. The following result can be found in [6].

Lemma 1.2. Let $k \ge 1$ be an integer. A graph G has a fractional k-factor if and only if for any subset S of V(G),

$$k|T| - d_{G-S}(T) \leqslant k|S|, \tag{1}$$

where $T = \{x \in V(G) - S | d_{G-S}(x) \leq k-1\}$ and $d_{G-S}(T) = \sum_{x \in T} d_{G-S}(x)$.

In particular, for k = 1, we have the following result.

Lemma 1.3 (Scheinerman and Ullman [7]). A graph G has a fractional perfect matching if and only if for any $S \subseteq V(G)$,

$$i(G-S) \leqslant |S|,\tag{2}$$

where $i(G - S) = |\{x \in V(G) - S | d_{G-S}(x) = 0\}|.$

Our main results are the following two theorems.

Theorem 1.1. Let G be a connected graph with $|V(G)| \ge 2$. Then G has a fractional perfect matching if $t(G) \ge 1$.

Theorem 1.2. Let $k \ge 2$ be an integer. A graph G with $|V(G)| \ge (k+1)$ has a fractional k-factor if $t(G) \ge k - 1/k$.

The result in Theorem 1.1 is sharp. To see this, consider the graph $G_1 = K_m \vee (m+1)K_1$ where \vee means "join" and *m* is an arbitrary positive integer. It is easy to find out that $t(G_1) = m/(m+1) < 1$ and (1) does not hold if we let $S = V(K_m)$. By Lemma 1.3 G_1 has no fractional perfect matching. But $t(G_1) \to 1$ when $m \to +\infty$.

To see Theorem 1.2 is also sharp, we construct the following graph G_k : $V(G_k) = A \cup B \cup C$ where A, B and C are disjoint with |A| = |B| = (nk + 1)(k - 1), and |C| = n(k - 1). Both A and C are cliques in G_k , while B is isomorphic to $(nk + 1)K_{k-1}$. Other edges in G_k are a perfect matching between A and B and all the pairs between B and C. If k = 2, let $S = (A - \{u\}) \cup C$ where $u \in A$, then |S| = 3n and $\omega(G - S) = 2n + 1$; if $k \ge 3$, let $S = (A - \{u\}) \cup \{v\} \cup C$ where $u \in A$ and $v \in B$ is matched to u in G_k . Then |S| = (nk + n + 1)(k - 1) and $\omega(G - S) = nk + 2$. This follows that

$$t(G_k) = \begin{cases} \frac{3n}{2n+1} & \text{if } k = 2, \\ \frac{(nk+n+1)(k-1)}{nk+2} & \text{if } k \ge 3. \end{cases}$$

But (1) does not hold if we let S = C. Thus by Lemma 1.2 G_k has no fractional k-factor. It is easy to see that $t(G_k)$ can be made arbitrarily close to k - 1/k when n is large enough. In this sense, the result in Theorem 1.2 is also sharp.

Remark. A graph *G* that satisfies the condition of Theorem 1.1 has 1-factors if |V(G)| is even by Lemma 1.1. However, a graph *G* that satisfies the conditions of Theorem 1.2 does not necessarily have a *k*-factor even if k|V(G)| is even [4].

2. Proofs of theorems

At first let us prove Theorem 1.1.

Proof of Theorem 1.1. If *G* is complete, obviously *G* has a fractional perfect matching as $|V(G)| \ge 2$. In the following we assume that *G* is not complete. Suppose that *G* satisfies the conditions in Theorem 1.1, but *G* has no fractional perfect matching. From Lemma 1.3, there exists a subset *S* of V(G) such that

$$i(G-S) > |S|.$$

Since G is connected, $S \neq \emptyset$. Thus $i(G - S) \ge 2$. Then

$$t(G) \leqslant \frac{|S|}{\omega(G-S)} \leqslant \frac{|S|}{i(G-S)} < 1$$

contradicting to $t(G) \ge 1$. \Box

To prove Theorem 1.2, we need the following Lemmas.

Lemma 2.1 (*Chvátal* [3]). If a graph G is not complete, then $t(G) \leq \frac{1}{2}\delta(G)$.

Lemma 2.2. Let G be a graph and let H = G[T] such that $d_G(x) = k - 1$ for every $x \in V(H)$ and no component of H is isomorphic to K_k where $T \subseteq V(G)$ and $k \ge 2$. Then there exists an independent set I and the covering set C = V(H) - I of H satisfying

$$|V(H)| \leq \left(k - \frac{1}{k+1}\right)|I| \tag{3}$$

and

$$|C| \leqslant \left(k - 1 - \frac{1}{k+1}\right)|I|. \tag{4}$$

Proof. Suppose that *H* has *m* components. For each component H_n , let I_n be a maximum independent set of H_n . First we claim that for each vertex $x \in I_n$ and $d_{H_n}(x) = k - 1$, there exists a vertex $y \in I_n - \{x\}$ such that $N_{H_n}(x) \cap N_{H_n}(y) \neq \emptyset$. For this, we show that $H_n[N_{H_n}(x)]$ is not complete. Otherwise, $H'_n = H_n[\{x\} \cup N_{H_n}(x)]$ is isomorphic to K_k . Since H_n is connected and for every vertex $x \in V(H_n)$, $d_{H_n}(x) \leqslant k - 1$, it follows that $H_n = H'_n$, which contradicts to that H_n is not isomorphic to K_k . Now if for any $y \in I_n - \{x\}$, $N_{H_n}(x) \cap N_{H_n}(y) = \emptyset$, then $E(\{x\} \cup N_{H_n}(x), I_n - \{x\}) = \emptyset$. Let x' and y' be two vertices in $H_n[N_{H_n}(x)]$ that are not adjacent. Then $(I_n - \{x\}) \cup \{x', y'\}$ will be an independent set of H_n , contradicting to that I_n is a maximum independent set of H_n . So what we desire follows. Let $I'_n = \{x | x \in I_n \text{ and } d_{H_n}(x) = k - 1\}$ and $I''_n = I_n - I'_n$. Then for every $x \in I''_n$, $d_{H_n}(x) \leqslant k - 2$. Note that both of I'_n and I''_n are independent sets of H_n . Since for every vertex $x \in I'_n$, $d_{H_n}(x) = k - 1$, and for every $x \in I''_n$, $d_{H_n}(x) \leqslant k - 2$, where $k \ge 2$, by the above claim we have the following inequality:

$$|V(H_n)| \leq k|I'_n| - \left\lceil \frac{|I'_n|}{2} \right\rceil + (k-1)|I''_n| \leq k|I_n| - \left\lceil \frac{|I_n|}{2} \right\rceil \leq \left(k - \frac{1}{k+1}\right)|I_n|$$

for each n = 1, ..., m. Let $I = \bigcup_{n=1}^{m} I_n$, Then $|I| = \sum_{n=1}^{m} |I_n|$ and I is a maximum independent set of H. Thus

$$|V(H)| = \sum_{n=1}^{m} |V(H_n)| \leq \sum_{n=1}^{m} \left(k - \frac{1}{k+1}\right) |I_n| = \left(k - \frac{1}{k+1}\right) |I|$$

which is inequality (3). Let C = V(H) - I. Then |C| = |V(H)| - |I| and the result (4) follows easily from (3). The proof is completed. \Box

The following Lemma 2.3 is similar to Lemma 5 of [5]. However, it has been strengthened not only in its conditions but also in its result.

Lemma 2.3. Let G be a graph and let H = G[T] such that $\delta(H) \ge 1$ and $1 \le d_G(x) \le k - 1$ for every $x \in V(H)$ where $T \subseteq V(G)$ and $k \ge 2$. Let T_1, \ldots, T_{k-1} be a partition of the vertices of H satisfying $d_G(x) = j$ for each $x \in T_j$ where we allow some T_j to be empty. If each component of H has a vertex of degree at most k - 2 in G, then H has a maximal independent set I and a covering set C = V(H) - I such that

$$\sum_{j=1}^{k-1} (k-j)c_j \leqslant \sum_{j=1}^{k-1} (k-2)(k-j)i_j,$$

where $c_j = |C \cap T_j|$ and $i_j = |I \cap T_j|$ for every $j = 1, \ldots, k-1$.

Proof. Since $\delta(H) \ge 1$ and each component of *H* has a vertex of degree at most k - 2 in *G*, we have $k \ge 3$. We prove the lemma by induction on |V(H)|. If |V(H)| = 2, then *H* is isomorphic to K_2 . Let $V(H) = \{x, y\}$ and suppose $x \in T_{i_0}$ and $y \in T_{i_0}$. We may assume $i_0 \le j_0$. Let $I = \{x\}$ and $C = V(H) - \{x\} = \{y\}$. Then

$$\sum_{j=1}^{k-1} (k-j)c_j = k - j_0 \leqslant k - i_0 \leqslant (k-2)(k-i_0) = \sum_{j=1}^{k-1} (k-2)(k-j)i_j$$

and the result follows. Now we assume that the result holds when |V(H)| < L. Now we consider $|V(H)| = L \ge 3$. Let $m = \min\{j | T_j \neq \emptyset\}$. Then $1 \le m \le k - 2$. Take any $y \in T_m$. Then $H - (\{y\} \cup N_H(y))$ may have some isolated vertices in *H*. Let I'' be the set of *y* and these isolated vertices. Now let $H' = H - (I'' \cup N_H(y))$. If $x \in I'' - \{y\}$, then we can see that $d_H(x) \le d_H(y)$ and $d_G(x) \ge d_G(y)$ by the definition of I'' and *m*.

If |V(H')| = 0, put I = I'' and $C = V(H) - I = N_H(y)$. Note that $T_j = \emptyset$ and $i_j = 0$ for j < m. Since |V(H')| = 0, we have

$$\sum_{j=m}^{k-1} c_j \leqslant m.$$

Thus

$$\sum_{j=1}^{k-1} (k-j)c_j \leq \sum_{j=m}^{k-1} (k-m)c_j$$

= $(k-m) \sum_{j=m}^{k-1} c_j \leq (k-m)m$
 $\leq (k-2)(k-m) \leq \sum_{j=1}^{k-1} (k-2)(k-j)i_j$

Obviously *I* is a maximal independent set of *H*.

So we suppose that $|V(H')| \neq 0$ or $V(H) \neq I'' \cup N_H(y)$. Note that if a vertex v is only adjacent to $N_H(y)$, then v is in $(I'' - \{y\})$. If vertex v is adjacent to a vertex $u \in (I'' - \{y\})$, then u is not an isolated vertex of $H - (\{y\} \cup N_H(y))$ in H. This contradicts to that $u \in (I'' - \{y\})$. Thus it follows that $\delta(H') \ge 1$. Clearly $\Delta(H') \le k - 1$. It is obvious that $|V(H')| \ge 2$. From the definition of H' and $\Delta(H) \le k - 1$ we can also see that each component of H' has a vertex of degree at most k - 2 in G as follows. If a component H_0 of H' is also a component of H, clearly, H_0 has a vertex of degree at most k - 2 by the hypothesis. Otherwise, a component H_0 of H' is not a component of H. Then H_0 is a component of $H_1 - (I'' \cup N_H(y))$ where H_1 is a component of H. Note that there are at least one edge e = uv

joining H_0 and $I'' \cup N_H(y)$ in H_1 . We may assume that vertex v is in H_0 . Since $d_G(x) \leq k - 1$ for every vertex x of H_0 , $d_G(v) \leq k - 1$. Thus vertex v is adjacent to at most k - 2 vertices in H_0 . It is easy to see that H_0 must have a vertex of degree at most k - 2. Let $T'_j = T_j \cap V(H')$. Since |V(H')| < L, by induction hypothesis, there exists a maximal independent set I' and a covering set C' = V(H') - I' of H' such that

$$\sum_{j=1}^{k-1} (k-j)c'_j \leqslant \sum_{j=1}^{k-1} (k-2)(k-j)i'_j,$$

where $i'_j = |I' \cap T'_j|$ and $c'_j = |C' \cap T'_j|$. Now let $I = I' \cup I''$ and $C = V(H) - I = C' \cup N_H(y)$. Obviously, *I* is a maximal independent set of *H*. Then

$$\sum_{j=1}^{k-1} (k-2)(k-j)i_j \ge \sum_{j=1}^{k-1} (k-2)(k-j)i'_j + m(k-m)$$
$$\ge \sum_{j=1}^{k-1} (k-j)c'_j + m(k-m).$$

Since $d_G(y) \leq m$ and $m = \min\{j | T_j \neq \phi\}$, we have

$$\sum_{j=1}^{k-1} (k-j)c_j \leqslant \sum_{j=1}^{k-1} (k-j)c'_j + m(k-m)$$

Thus

$$\sum_{j=1}^{k-1} (k-j)c_j \leqslant \sum_{j=1}^{k-1} (k-2)(k-j)i_j$$

completing the proof. \Box

Proof of Theorem 1.2. If *G* is complete, since $|V(G)| \ge k + 1$, obviously, *G* has a fractional *k*-factor. In the following we assume that *G* is not complete. Suppose that *G* satisfies the conditions of Theorem 1.2, but has no fractional *k*-factors. From Lemma 1.2 there exists a subset *S* of V(G) such that

$$k|T| - d_{G-S}(T) > k|S|,$$
(5)

where $T = \{x \in V(G) - S | d_{G-S}(x) \le k - 1\}$. By Lemma 2.1, we have $\delta(G) \ge 2t(G) \ge 2k - 2/k \ge k + 1$. Therefore $S \ne \emptyset$ by (5). Let *l* be the number of the components of H' = G[T] which are isomorphic to K_k and let $T_0 = \{x \in V(H') | d_{G-S}(x) = 0\}$. Let *H* be the subgraph obtained from $H' - T_0$ by deleting those *l* components isomorphic to K_k . If |V(H)| = 0, then from (5) we obtain

$$k|T_0| + lk > k|S|$$

or

$$1 \leq |S| < |T_0| + l.$$

Hence $\omega(G-S) \ge l + |T_0| > 1$ and

$$t(G) \leq \frac{|S|}{\omega(G-S)} = \frac{|S|}{l+|T_0|} < 1.$$

This contradicts that $t(G) \ge k - 1/k \ge \frac{3}{2}$.

Now we consider that |V(H)| > 0 and $\delta(H) \ge 1$. Let $H = H_1 \cup H_2$ where H_1 is the union of components of H which satisfies that $d_{G-S}(x) = k - 1$ for every vertex $x \in V(H_1)$ and $H_2 = H - H_1$. By Lemma 2.2, H_1 has a maximum independent set I_1 and the covering set $C_1 = V(H_1) - I_1$ such that

$$|V(H_1)| \leqslant \left(k - \frac{1}{k+1}\right)|I_1| \tag{6}$$

and

$$|C_1| \leq \left(k - 1 - \frac{1}{k+1}\right) |I_1|.$$
 (7)

On the other hand, it is obvious that $\delta(H_2) \ge 1$ and $\Delta(H_2) \le k-1$. Let $T_j = \{x \in V(H_2) | d_{G-S}(x) = j\}$ for $1 \le j \le k-1$. By the definition of *H* and H_2 we can also see that each component of H_2 has a vertex of degree at most k-2 in G-S. According to Lemma 2.3, H_2 has a maximal independent set I_2 and the covering set $C_2 = V(H_2) - I_2$ such that

$$\sum_{j=1}^{k-1} (k-j)c_j \leqslant \sum_{j=1}^{k-1} (k-2)(k-j)i_j,$$
(8)

where $c_j = |C_2 \cap T_j|$ and $i_j = |I_2 \cap T_j|$ for every j = 1, ..., k-1. Set W = V(G) - S - T and $U = S \cup C_1 \cup C_2 \cup (N_G(I_2) \cap W)$. Then since $|C_2| + |(N_G(I_2) \cap W) \leq \sum_{j=1}^{k-1} ji_j$ we obtain

$$|U| \leq |S| + |C_1| + \sum_{j=1}^{k-1} j i_j \tag{9}$$

and

$$\omega(G-U) \ge t_0 + l + |I_1| + \sum_{j=1}^{k-1} i_j, \tag{10}$$

where $t_0 = |T_0|$. Let t(G) = t. Then when $\omega(G - U) > 1$, we have

$$|U| \ge t\omega(G-U). \tag{11}$$

In addition, (11) also holds when $\omega(G - U) = 1$ as by Lemma 2.1 for any $x \in T$,

$$|U| \ge d_{G-S}(x) + |S| \ge d(x) \ge 2t.$$

By (9)–(11),

$$|S| + |C_1| + \sum_{j=1}^{k-1} ji_j \ge t(t_0 + l) + t|I_1| + t\sum_{j=1}^{k-1} i_j$$

or

$$|S| + |C_1| \ge \sum_{j=1}^{k-1} (t-j)i_j + t(t_0+l) + t|I_1|.$$
(12)

From (5) we have

$$kt_0 + kl + |V(H_1)| + \sum_{j=1}^{k-1} (k-j)i_j + \sum_{j=1}^{k-1} (k-j)c_j > k|S|.$$

Combining with (12) we have

$$kt_0 + kl + |V(H_1)| + \sum_{j=1}^{k-1} (k-j)i_j + \sum_{j=1}^{k-1} (k-j)c_j + k|C_1| > \sum_{j=1}^{k-1} (kt-kj)i_j + kt(t_0+l) + kt|I_1|.$$

Thus

$$\sum_{j=1}^{k-1} (k-j)c_j + |V(H_1)| + k|C_1| > \sum_{j=1}^{k-1} (kt - kj - k + j)i_j + k(t-1)(t_0 + l) + kt|I_1|$$

$$\geq \sum_{j=1}^{k-1} (kt - kj - k + j)i_j + kt|I_1|.$$
(13)

By (6) and (7),

$$|V(H_1)| + k|C_1| \leq \left[k - \frac{1}{k+1} + k\left(k - 1 - \frac{1}{k+1}\right)\right] |I_1|$$

= $(k^2 - 1)|I_1|.$ (14)

Using (8), (13) and (14), we have

$$\sum_{j=1}^{k-1} (k-2)(k-j)i_j + (k^2-1)|I_1| > \sum_{j=1}^{k-1} (kt-kj-k+j)i_j + kt|I_1|.$$

Thus at least one of the following two cases must hold. *Case* 1: There is at least one *j* such that

$$(k-2)(k-j) > kt - kj - k + j.$$

It follows that

$$t < \frac{k^2 - k + j}{k}.$$

But $j \leq (k-1)$, we have

$$t < k - \frac{1}{k},$$

contradicting to the toughness condition of Theorem 1.2. Case 2: $k^2 - 1 > kt$. In this case we have

$$t < k - \frac{1}{k}.$$

This also contradicts to the toughness condition of Theorem 1, completing the proof of the theorem. \Box

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