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Gröbner–Shirshov bases for metabelian Lie algebras[☆]

Yongshan Chen, Yuqun Chen^{*}

School of Mathematical Sciences, South China Normal University, Guangzhou 510631, PR China

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ABSTRACT

In this paper, we establish the Gröbner–Shirshov bases theory for metabelian Lie algebras. As applications, we find the Gröbner–Shirshov bases for partial commutative metabelian Lie algebras related to circuits, trees and some cubes.

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1. Introduction

The class of metabelian Lie algebras is an important class of Lie algebras and attracts much attention. Let us mention the recent papers by E. Daniyarova, I. Kazatchkov, and V. Remeslennikov [4–6] on algebraic geometry of free metabelian Lie algebra, S. Findik and V. Drensky [7,8] on automorphisms of free metabelian Lie algebras, and V. Kurlin [9] on the Backer–Campbell–Hausdorff formula for free metabelian Lie algebras. Gröbner–Shirshov bases theory would be useful on this class of algebras. This theory was first considered by V.V. Talapov [10] in 1982. However, there are serious gaps in his paper. He missed several cases when he defined compositions. This means the theory was not established correctly. We refine his idea and complete the results.

It is well known that for many kinds of algebras, if $A_i = (X_i | S_i)$, $i = 1, 2$, are defined by generators and defining relations, where S_1 and S_2 are Gröbner–Shirshov bases respectively, then $S_1 \cup S_2$ is a

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^{*} Corresponding author.

E-mail addresses: jackalshan@126.com (Y. Chen), yqchen@scnu.edu.cn (Y. Chen).

Gröbner–Shirshov basis for the free product $A_1 * A_2 = (X_1 \cup X_2 | S_1 \cup S_2)$ of A_1 and A_2 , for example, associative algebras, Lie algebras and for all classes with compositions of inclusion and intersection only (cf. [2,3]). We prove that it is not the case for metabelian Lie algebras, see Theorem 3.1, even in the case of $S_2 = \emptyset$. On the other hand, if $S_i \subset A_i^{(2)}$, then $S_1 \cup S_2$ is a Gröbner–Shirshov basis for the free metabelian Lie product $A_1 * A_2$, see Proposition 3.2.

Throughout this paper, all algebras will be considered over a field \mathbf{k} of arbitrary characteristic. Suppose that \mathcal{L} is a Lie algebra. Then \mathcal{L} is called a metabelian Lie algebra if $\mathcal{L}^{(2)} = 0$, where $\mathcal{L}^{(0)} = \mathcal{L}$, $\mathcal{L}^{(n+1)} = [\mathcal{L}^{(n)}, \mathcal{L}^{(n)}]$. More precisely, the variety of metabelian Lie algebras is given by the identity

$$(x_1x_2)(x_3x_4) = 0.$$

2. Composition-Diamond lemma for metabelian Lie algebras

Let us begin with the construction of a free metabelian Lie algebra. Let X be a set and $Lie(X)$ be the free Lie algebra generated by X . Then $\mathcal{L}_{(2)}(X) = Lie(X)/Lie(X)^{(2)}$ is the free metabelian Lie algebra generated by X . Any metabelian Lie algebra \mathcal{ML} is a homomorphic image of a free metabelian Lie algebra generated by some X , that is, \mathcal{ML} can be presented by generators X and defining relations S : $\mathcal{ML} = \mathcal{L}_{(2)}(X|S)$.

We call a non-associative monomial on X left-normed if it is of the form $(\dots((ab)c)\dots)d$. In the sequel, the brackets in the expression of left-normed monomials are omitted.

Let X be well ordered. For an arbitrary set of indices j_1, j_2, \dots, j_m , define an associative word

$$\langle a_{j_1} \cdots a_{j_m} \rangle = a_{i_1} \cdots a_{i_m},$$

where $a_{i_1} \leq \dots \leq a_{i_m}$ and i_1, i_2, \dots, i_m is a permutation of the indices j_1, j_2, \dots, j_m .

Let

$$R = \{u = a_0a_1a_2 \cdots a_n \mid a_i \in X (0 \leq i \leq n), a_0 > a_1 \leq \dots \leq a_n, n \geq 1\}$$

and $N = X \cup R$, where $u = a_0a_1a_2 \cdots a_n$ is left-normed.

Then N forms a linear basis of the free metabelian Lie algebra $\mathcal{L}_{(2)}(X)$, i.e., $\mathcal{L}_{(2)}(X) = \mathbf{k}N$, see [1].

We call elements of N regular words on X and those of R regular R -words. Therefore, for any $f \in \mathcal{L}_{(2)}(X)$, f has a unique presentation $f = f^{(1)} + f^{(0)}$, where $f^{(1)} \in \mathbf{k}R$ and $f^{(0)} \in \mathbf{k}X$. Moreover, the multiplication table of regular words is the following, $u \cdot v = 0$ if both $u, v \in R$, and

$$a_0a_1a_2 \cdots a_n \cdot b = \begin{cases} a_0 \langle a_1a_2 \cdots a_nb \rangle & \text{if } a_1 \leq b, \\ a_0ba_1a_2 \cdots a_n - a_1b \langle a_0a_2 \cdots a_n \rangle & \text{if } a_1 > b. \end{cases}$$

If $u = a_0a_1 \cdots a_n \in R$, then the regular words a_i ($0 \leq i \leq n$) and $a_0 \langle a_{i_1} \cdots a_{i_l} \rangle$ ($l \leq n, a_{i_1}, \dots, a_{i_l}$ is a subsequence of the sequence a_1, \dots, a_n) are called subwords of u . The words a_i ($2 \leq i \leq n$), and also a_1 if $a_0 > a_2$ are called strict subwords of u .

Define the length of regular words:

$$|a| = 1, \quad |a_0a_1a_2 \cdots a_n| = n + 1,$$

where $a, a_0, \dots, a_n \in X$. Now we order the set N degree-lexicographically, i.e., for any $u, v \in N$,

$$u > v \quad \text{if } |u| > |v| \quad \text{or} \quad |u| = |v|, \quad u >_{lex} v.$$

Through out this paper, we will use this ordering.

The largest monomial occurring in $f \in \mathcal{L}_{(2)}(X)$ with nonzero coefficient is called the leading word of f and is denoted by \bar{f} . Then we have $\overline{a_0a_1a_2 \cdots a_n \cdot b} = a_0 \langle a_1a_2 \cdots a_nb \rangle$ and $|\overline{u \cdot b}| = |u| + 1$. For any $f \in \mathcal{L}_{(2)}(X)$, we call f monic, (1)-monic and (0)-monic if the coefficients of \bar{f} , $\overline{f^{(1)}}$ and $\overline{f^{(0)}}$ are 1 respectively.

Lemma 2.1. For any $u, v \in N$, if $u > v$ then

$$(\forall b \in N) \quad u \cdot b \neq 0 \quad \Rightarrow \quad \overline{u \cdot b} > \overline{v \cdot b}.$$

Proof. The result is obvious if either $u, v \in X$ or $|u| > |v|$. Suppose that $u = a_0 a_1 a_2 \cdots a_n, v = a'_0 a'_1 a'_2 \cdots a'_n \in R$ and $b \in X$. If $a_0 > a'_0$ then we are done. If $a_0 = a'_0$, then $\langle a_1 a_2 \cdots a_n b \rangle > \langle a'_1 a'_2 \cdots a'_n b \rangle$ in $[X]$ since the deg-lex ordering on $[X]$ is monomial, where $[X]$ is the free commutative monoid generated by X . Now, the result follows. \square

Let $S \subset \mathcal{L}_{(2)}(X)$. We denote $u_s = s v_1 v_2 \cdots v_n$, where $v_i \in N, s \in S$ and $n \geq 0$. We call u_s an s -word (or S -word). It is clear that each element of the ideal $Id(S)$ of $\mathcal{L}_{(2)}(X)$ generated by S is a linear combination of S -words.

Definition 2.2. Let $S \subset \mathcal{L}_{(2)}(X)$. Then the following two kinds of polynomials are called normal S -words:

- (i) $sa_1 a_2 \cdots a_n$, where $a_i \in X (1 \leq i \leq n), a_1 \leq a_2 \leq \cdots \leq a_n, s \in S, \bar{s} \neq a_1$ and $n \geq 0$;
- (ii) us , where $u \in R, s \in S$ and $\bar{s} \neq u$.

By a simple observation, we have

$$\overline{sa_1 a_2 \cdots a_n} = \begin{cases} c_0 \langle c_1 \cdots c_k a_1 a_2 \cdots a_n \rangle & \text{if } \bar{s} = c_0 c_1 \cdots c_k, \\ c_0 a_1 a_2 \cdots a_n & \text{if } \bar{s} = c_0 > a_1, \\ a_1 c_0 a_2 \cdots a_n & \text{if } \bar{s} = c_0 < a_1, \end{cases}$$

and $\overline{us} = a_0 \langle a_1 \cdots a_k s^{(0)} \rangle$, where $u = a_0 \langle a_1 \cdots a_k \rangle$. That is to say, if u_s is a normal s -word, then $\overline{u_s}$ contains either \bar{s} as a subword or $s^{(0)}$ as a strict subword.

A regular word u is called S -irreducible if for any $s \in S, u$ contains neither \bar{s} as a subword nor $s^{(0)}$ as a strict subword. Denote by $Irr(S)$ the set of all S -irreducible words. This means

$$Irr(S) = \{u \mid u \in N, u \neq \overline{v_s} \text{ for any normal } S\text{-word } v_s\}.$$

Remark. For any $s \in \mathcal{L}_{(2)}(X)$,

$$sa_1 a_2 \cdots a_n = sa_1 a_{j_2} \cdots a_{j_n},$$

where $\langle a_{j_2} \cdots a_{j_n} \rangle = a_2 \cdots a_n$.

Lemma 2.3. Let $S \subset \mathcal{L}_{(2)}(X)$ and $Id(S)$ be the ideal of $\mathcal{L}_{(2)}(X)$ generated by S . Then for any $f \in Id(S), f$ can be written as a linear combination of normal S -words.

Proof. It is suffice to show that any S -word $u_s = s u_1 u_2 \cdots u_n$ is a linear combination of normal S -words, where $u_i \in N, 1 \leq i \leq n$. We may assume that s is monic. The proof will be proceeded by induction on n .

There is nothing to prove if $n = 0$.

Assume that $n = 1$. If $\bar{s} \neq u_1$, then either su_1 or u_1s is normal. If $\bar{s} = u_1$, then $s = u_1 + \sum_{\bar{s} > v_j \in N} \alpha_j v_j, \alpha_j \in \mathbf{k}$ and

$$su_1 = s \left(s - \sum_{v_j < \bar{s}} \alpha_j v_j \right) = -s \sum_{v_j < \bar{s}} \alpha_j v_j = - \sum_{v_j < \bar{s}} \alpha_j s v_j,$$

where for each j , either $v_j s$ or sv_j is normal.

For $n \geq 2$, if $\exists u_i \in R$ ($i \geq 2$), then $su_1u_2 \cdots u_n = 0$; if $u_1 \in R$, then $(su_1)a_2 \cdots a_n = s(u_1a_2 \cdots a_n)$ which is the above case. So we may assume that $u_s = sa_1a_2 \cdots a_n$ is normal and $u_{n+1} = a \in X$. Then

$$u_s \cdot u_{n+1} = sa_1a_2 \cdots a_n \cdot a = sa_1 \langle a_2 \cdots a_n a \rangle.$$

If $a \geq a_1$, then $sa_1 \langle a_2 \cdots a_n a \rangle$ is normal. If $a < a_1$, then

$$\begin{aligned} u_s \cdot u_{n+1} &= sa_1aa_2 \cdots a_n \\ &= saa_1a_2 \cdots a_n - ((a_1a)s)a_2 \cdots a_n \\ &= saa_1a_2 \cdots a_n - a_1aa_2 \cdots a_n \cdot s. \end{aligned}$$

Clearly, by the previous proof, $a_1aa_2 \cdots a_n \cdot s$ is normal. Now $saa_1a_2 \cdots a_n$ is already normal provided that $\bar{s} \neq a$. If $\bar{s} = a$, then we substitute a by $-\sum_{\bar{s} > v_j \in N} \alpha_j v_j$ where $s = a + \sum_{\bar{s} > v_j \in N} \alpha_j v_j$, and the result follows now. \square

Lemma 2.4. Let u_s be a normal S -word and $w \in N$. If $\bar{u}_s < w$, then

$$(\forall a \in X) \quad w \cdot a \neq 0 \quad \Rightarrow \quad \overline{u_s \cdot a} < \overline{w \cdot a}.$$

Proof. Suppose that $w = b_0b_1 \cdots b_m$ where $m \geq 0$. Then

$$\overline{w \cdot a} = \begin{cases} b_0 \langle b_1 \cdots b_m a \rangle & \text{if } m > 0, \\ b_0 a & \text{if } m = 0 \text{ and } b_0 > a, \\ ab_0 & \text{if } m = 0 \text{ and } b_0 < a. \end{cases}$$

If $u_s = sa_1a_2 \cdots a_n$, then

$$\bar{u}_s = \begin{cases} c_0 \langle c_1 \cdots c_k a_1 a_2 \cdots a_n \rangle & \text{if } \bar{s} = c_0 c_1 \cdots c_k, \\ c_0 a_1 a_2 \cdots a_n & \text{if } \bar{s} = c_0 > a_1, \\ a_1 c_0 a_2 \cdots a_n & \text{if } \bar{s} = c_0 < a_1 \end{cases}$$

and

$$u_s \cdot a = \begin{cases} sa_1 \langle a_2 \cdots a_n a \rangle & \text{if } a \geq a_1, \\ saa_1a_2 \cdots a_n - a_1aa_2 \cdots a_n \cdot s & \text{if } a < a_1. \end{cases}$$

Therefore,

$$\overline{u_s \cdot a} = \begin{cases} c_0 \langle c_1 \cdots c_k a_1 a_2 \cdots a_n a \rangle & \text{if } \bar{s} = c_0 c_1 \cdots c_k, \\ c_0 \langle a_1 a_2 \cdots a_n a \rangle & \text{if } \bar{s} = c_0 > a_1, \\ a_1 \langle c_0 a_2 \cdots a_n a \rangle & \text{if } \bar{s} = c_0 < a_1. \end{cases}$$

If $u_s = a_0a_1 \cdots a_n \cdot s$, then $\bar{u}_s = a_0 \langle a_1 \cdots a_n s^{(0)} \rangle$ and $\overline{u_s \cdot a} = a_0 \langle a_1 \cdots a_n s^{(0)} a \rangle$. Since $\bar{u}_s < w$, in both cases we have $\overline{u_s \cdot a} < \overline{w \cdot a}$. \square

Definition 2.5. Let f and g be momic polynomials of $\mathcal{L}_{(2)}(X)$ and α and β be the coefficients of $\overline{f^{(0)}}$ and $\overline{g^{(0)}}$ respectively. We define seven different types of compositions as follows:

1. If $\bar{f} = a_0 a_1 \cdots a_n$, $\bar{g} = a_0 b_1 \cdots b_m$ ($n, m \geq 0$) and $\text{lcm}(AB) \neq \langle a_1 \cdots a_n b_1 \cdots b_m \rangle$, where $\text{lcm}(AB)$ denotes the least common multiple in $[X]$ of associative words $a_1 \cdots a_n$ and $b_1 \cdots b_m$, then let $w = a_0 \langle \text{lcm}(AB) \rangle$. The composition of type I of f and g relative to w is defined by

$$C_I(f, g)_w = f \left\langle \frac{\text{lcm}(AB)}{a_1 \cdots a_n} \right\rangle - g \left\langle \frac{\text{lcm}(AB)}{b_1 \cdots b_m} \right\rangle.$$

2. If $\bar{f} = \overline{f^{(1)}} = a_0 a_1 \cdots a_n$, $\overline{g^{(0)}} = a_i$ for some $i \geq 2$ or $\overline{g^{(0)}} = a_1$ and $a_0 > a_2$, then let $w = \bar{f}$. The composition of type II of f and g relative to w is defined by

$$C_{II}(f, g)_w = f - \beta^{-1} a_0 a_1 \cdots \hat{a}_i \cdots a_n \cdot g,$$

where $a_0 a_1 \cdots \hat{a}_i \cdots a_n = a_0 a_1 \cdots a_{i-1} a_{i+1} \cdots a_n$.

3. If $\bar{f} = \overline{f^{(1)}} = a_0 a_1 \cdots a_n$, $\bar{g} = \overline{g^{(0)}} = a_1$ and $a_0 \leq a_2$ or $n = 1$, then let $w = \bar{f}$. The composition of type III of f and g relative to w is defined by

$$C_{III}(f, g)_{\bar{f}} = f + g a_0 a_2 \cdots a_n.$$

4. If $\bar{f} = \overline{f^{(1)}} = a_0 a_1 \cdots a_n$, $g^{(1)} \neq 0$, $\overline{g^{(0)}} = a_1$ and $a_0 \leq a_2$ or $n = 1$, then for any $a < a_0$ and $w = a_0 \langle a_1 \cdots a_n a \rangle$, the composition of type IV of f and g relative to w is defined by

$$C_{IV}(f, g)_w = f a - \beta^{-1} a_0 a a_2 \cdots a_n \cdot g.$$

5. If $\bar{f} = \overline{f^{(1)}} = a_0 a_1 \cdots a_n$, $f^{(0)} \neq 0$, $g^{(1)} \neq 0$ and $\overline{g^{(0)}} = b \notin \{a_i\}_{i=1}^n$, then let $w = a_0 \langle a_1 \cdots a_n b \rangle$. The composition of type V of f and g relative to w is defined by

$$C_V(f, g)_w = f b - \beta^{-1} a_0 a_1 \cdots a_n \cdot g.$$

6. If $\overline{f^{(0)}} = \overline{g^{(0)}} = a$ and $f^{(1)} \neq 0$, then for any $a_0 a_1 \in R$ and $w = a_0 \langle a_1 a \rangle$, the composition of type VI of f and g relative to w is defined by

$$C_{VI}(f, g)_w = (a_0 a_1) (\alpha^{-1} f - \beta^{-1} g).$$

7. If $f^{(1)} \neq 0$, $g^{(1)} \neq 0$ and $\overline{f^{(0)}} = a > \overline{g^{(0)}} = b$, then for any $a_0 > a$ and $w = a_0 b a$, the composition of type VII of f and g relative to w is defined by

$$C_{VII}(f, g)_w = \alpha^{-1} (a_0 b) f - \beta^{-1} (a_0 a) g.$$

Immediately, we have $\overline{C_\lambda(f, g)_w} < w$.

Remark. In the paper of V.V. Talapov [10], only the compositions of types I, II and III are defined.

Definition 2.6. Given a set S of monic polynomials of $\mathcal{L}_{(2)}(X)$ and $w \in N$, a polynomial $f \in \mathcal{L}_{(2)}(X)$ is called trivial modulo S and w , denoted by $f \equiv 0 \pmod{S, w}$, if f is a linear combination of normal S -words whose leading words are less than w , i.e., $f = \sum_i \alpha_i u_{s_i}$, where $\alpha_i \in \mathbf{k}$, u_{s_i} are normal S -words and $\overline{u_{s_i}} < w$. For any $f, g \in \mathcal{L}_{(2)}(X)$, we say $f \equiv g \pmod{S, w}$ if $f - g \equiv 0 \pmod{S, w}$.

The set S is a Gröbner–Shirshov basis in $\mathcal{L}_{(2)}(X)$ if S is closed under compositions, which means every composition of any two elements of S is trivial modulo S and corresponding w , i.e., $(\forall f, g \in S) C_\lambda(f, g)_w \equiv 0 \pmod{S, w}$.

Lemma 2.7. *If $sa_1a_2 \cdots a_n$ is a normal s -word with leading word w , then for any $a_{i_1} < \bar{s}$,*

$$sa_1a_2 \cdots a_n \equiv sa_{i_1}a_{i_2} \cdots a_{i_n} \pmod{(s, w)},$$

where $\langle a_{i_1}a_{i_2} \cdots a_{i_n} \rangle = a_1a_2 \cdots a_n$.

Proof. There is nothing to prove if $a_{i_1} = a_1$. Suppose that $a_{i_1} = a_j > a_1$ for some $j \geq 2$. Then we have

$$\begin{aligned} sa_1a_2 \cdots a_n &= sa_1a_ja_2 \cdots \hat{a}_j \cdots a_n \\ &= sa_ja_1a_2 \cdots \hat{a}_j \cdots a_n + (a_1a_j)a_2 \cdots \hat{a}_j \cdots a_n \cdot s. \end{aligned}$$

Since $a_{i_1} < \bar{s}$, it is easy to see that $\overline{(a_1a_j)a_2 \cdots \hat{a}_j \cdots a_n \cdot s} < \overline{sa_1a_2 \cdots a_n} = w$. The result follows. \square

The following lemma plays a key role in this paper.

Lemma 2.8. *Let S be a Gröbner–Shirshov basis in $\mathcal{L}_{(2)}(X)$. If $w = \overline{u_{s_1}} = \overline{u_{s_2}}$, where $s_1, s_2 \in S$ and u_{s_1}, u_{s_2} are normal S -words, then for some $0 \neq \alpha \in \mathbf{k}$,*

$$u_{s_1} \equiv \alpha u_{s_2} \pmod{(S, w)}.$$

Proof. There are three main cases to consider.

Case 1. $u_{s_1} = s_1a_1a_2 \cdots a_n, u_{s_2} = s_2b_1b_2 \cdots b_m$.

(1.1) If $\bar{s}_1 = s_1^{(1)} = c_0c_1 \cdots c_k$ and $\bar{s}_2 = s_2^{(1)} = d_0d_1 \cdots d_l$, then $c_0 = d_0$ and

$$w = c_0 \langle c_1 \cdots c_k a_1 a_2 \cdots a_n \rangle = d_0 \langle d_1 \cdots d_l b_1 b_2 \cdots b_m \rangle = c_0 \langle \text{lcm}(CD)T \rangle,$$

where $T \in [X]$ such that $\langle c_1 \cdots c_k a_1 a_2 \cdots a_n \rangle = \langle d_1 \cdots d_l b_1 b_2 \cdots b_m \rangle = \langle \text{lcm}(CD)T \rangle$. Thus, By Lemmas 2.7 and 2.4 we have

$$\begin{aligned} s_1a_1a_2 \cdots a_n - s_2b_1b_2 \cdots b_m &= s_1 \left\langle \frac{\text{lcm}(CD)}{c_1 \cdots c_k} T \right\rangle - s_2 \left\langle \frac{\text{lcm}(CD)}{d_1 \cdots d_l} T \right\rangle \\ &\equiv \left(s_1 \left\langle \frac{\text{lcm}(CD)}{c_1 \cdots c_k} \right\rangle - s_2 \left\langle \frac{\text{lcm}(CD)}{d_1 \cdots d_l} \right\rangle \right) \langle T \rangle \\ &\equiv C_1(s_1, s_2)_{w'} \langle T \rangle \\ &\equiv 0 \pmod{(S, w)}, \end{aligned}$$

where $w' = c_0 \langle \text{lcm}(CD) \rangle$ and $w = \overline{w' \langle T \rangle}$.

(1.2) If $\bar{s}_1 = s_1^{(1)} = c_0c_1 \cdots c_k$ and $\bar{s}_2 = s_2^{(0)} = d$, then there are two subcases to be discussed.

(1.21) If $d > b_1$ then

$$w = c_0 \langle c_1 \cdots c_k a_1 a_2 \cdots a_n \rangle = db_1b_2 \cdots b_m,$$

which implies $c_0 = d$ and $\langle c_1 \cdots c_k a_1 a_2 \cdots a_n \rangle = b_1b_2 \cdots b_m$.

Hence,

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n - s_2 b_1 b_2 \cdots b_m &\equiv s_1 a_1 a_2 \cdots a_n - (s_2 c_1 \cdots c_k) a_1 a_2 \cdots a_n \\ &\equiv (s_1 - s_2 c_1 \cdots c_k) a_1 a_2 \cdots a_n \\ &\equiv C_I(s_1, s_2)_{\bar{s}_1} a_1 a_2 \cdots a_n \\ &\equiv 0 \pmod{(S, w)}. \end{aligned}$$

(1.22) If $d < b_1$ then $a_1 \geq c_1$. In fact, if $a_1 < c_1$ ($< c_0$), then $w = c_0 a_1 \langle c_1 \cdots c_k a_2 \cdots a_n \rangle = b_1 d b_2 \cdots b_m$, which implies $c_0 = b_1$, $a_1 = d$ and $\langle c_1 \cdots c_k a_2 \cdots a_n \rangle = b_2 \cdots b_m$. This is impossible because $c_1 < c_0 = b_1 \leq b_i$ ($2 \leq i \leq m$). Thus we have $a_1 \geq c_1$ and

$$w = c_0 c_1 \langle c_2 \cdots c_k a_1 a_2 \cdots a_n \rangle = b_1 d b_2 \cdots b_m,$$

which implies $c_0 = b_1$, $c_1 = d$ and $\langle c_2 \cdots c_k a_1 a_2 \cdots a_n \rangle = b_2 \cdots b_m$.

By noting that $c_0 = b_1 \leq b_i = c_2$ for some $2 \leq i \leq m$, we have

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n + s_2 b_1 b_2 \cdots b_m &= s_1 a_1 a_2 \cdots a_n + (s_2 c_0 c_2 \cdots c_k) a_1 a_2 \cdots a_n \\ &= (s_1 + s_2 c_0 c_2 \cdots c_k) a_1 a_2 \cdots a_n \\ &\equiv C_{III}(s_1, s_2)_{\bar{s}_1} a_1 a_2 \cdots a_n \\ &\equiv 0 \pmod{(S, w)}. \end{aligned}$$

(1.3) If $\bar{s}_1 = \overline{s_1^{(0)}} = c$ and $\bar{s}_2 = \overline{s_2^{(0)}} = d$, then we have $n = m$. Thus, we may assume that $n = m \geq 1$. There are two subcases to consider.

(1.31) If either $c > a_1$, $d > b_1$ or $c < a_1$, $d < b_1$, then

$$w = c a_1 \cdots a_n = d b_1 \cdots b_n$$

or

$$w = a_1 c a_2 \cdots a_n = b_1 d b_2 \cdots b_n,$$

which implies $c = d$, $a_i = b_i$ for any i .

It is easy to see that

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n - s_2 b_1 b_2 \cdots b_n &= (s_1 - s_2) a_1 \cdots a_n \\ &= C_I(s_1, s_2)_{\bar{s}_1} a_1 \cdots a_n \\ &\equiv 0 \pmod{(S, w)}. \end{aligned}$$

(1.32) If $c > a_1$ but $d < b_1$, then

$$w = c a_1 \cdots a_n = b_1 d b_2 \cdots b_n,$$

which implies $c = b_1$, $d = a_1$, $a_i = b_i$ for any $i > 1$.

Obviously,

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n + s_2 b_1 b_2 \cdots b_n &= (s_1 \bar{s}_2 - \bar{s}_1 s_2) a_2 \cdots a_n \\ &= (s_1 (\bar{s}_2 - s_2) - (\bar{s}_1 - s_1) s_2) a_2 \cdots a_n \\ &\equiv 0 \pmod{(S, w)}. \end{aligned}$$

Case 2. $u_{s_1} = s_1 a_1 a_2 \cdots a_n$, $u_{s_2} = b_0 b_1 b_2 \cdots b_m \cdot s_2$. We may assume that s_2 is (0)-monic and $\overline{s_2^{(0)}} = d$. Then $w = b_0 \langle b_1 \cdots b_m d \rangle$.

(2.1) If $\bar{s}_1 = \overline{s_1^{(1)}} = c_0 c_1 \cdots c_k$, then $c_0 = b_0$ and

$$w = c_0 \langle c_1 \cdots c_k a_1 a_2 \cdots a_n \rangle = b_0 \langle b_1 \cdots b_m d \rangle.$$

(2.11) If $d \notin \{c_i\}_{i=1}^k$, then there exists an a_i ($1 \leq i \leq n$) such that $d = a_i$. Thus,

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n - b_0 b_1 b_2 \cdots b_m \cdot s_2 &\equiv (s_1 a_i) a_1 a_2 \cdots \hat{a}_i \cdots a_n - (c_0 c_1 \cdots c_k \cdot s_2) a_1 a_2 \cdots \hat{a}_i \cdots a_n \\ &\equiv (s_1 \overline{s_2^{(0)}} - \bar{s}_1 s_2) a_1 a_2 \cdots \hat{a}_i \cdots a_n \pmod{(S, w)}. \end{aligned}$$

If $s_2^{(1)} = 0$, then

$$\begin{aligned} (s_1 \overline{s_2^{(0)}} - \bar{s}_1 s_2) a_1 a_2 \cdots \hat{a}_i \cdots a_n &= (s_1 \bar{s}_2 - \bar{s}_1 s_2) a_1 a_2 \cdots \hat{a}_i \cdots a_n \\ &\equiv 0 \pmod{(S, w)}. \end{aligned}$$

If $s_1^{(0)} = 0$, i.e., $s_1 = s_1^{(1)} = \bar{s}_1 + r_1^{(1)}$, then let $s_2^{(0)} = \overline{s_2^{(0)}} + r_2^{(0)}$ and we have

$$\begin{aligned} s_1 \overline{s_2^{(0)}} - \bar{s}_1 s_2 &= (\bar{s}_1 + r_1^{(1)}) \overline{s_2^{(0)}} - \bar{s}_1 s_2^{(0)} \\ &= r_1^{(1)} \overline{s_2^{(0)}} - \bar{s}_1 r_2^{(0)} \\ &= r_1^{(1)} \overline{s_2^{(0)}} - \bar{s}_1 r_2^{(0)} + r_1^{(1)} r_2^{(0)} - r_1^{(1)} r_2^{(0)} \\ &= r_1^{(1)} \overline{s_2^{(0)}} - s_1 r_2^{(0)} \\ &= r_1^{(1)} s_2 - s_1 r_2^{(0)}, \end{aligned}$$

which implies $(s_1 \overline{s_2^{(0)}} - \bar{s}_1 s_2) a_1 a_2 \cdots \hat{a}_i \cdots a_n \equiv 0 \pmod{(S, w)}$ immediately.

If $s_2^{(1)} \neq 0$ and $s_1^{(0)} \neq 0$, then

$$\begin{aligned} (s_1 \overline{s_2^{(0)}} - \bar{s}_1 s_2) a_1 a_2 \cdots \hat{a}_i \cdots a_n &\equiv C_V(s_1, s_2)_{w'} a_1 a_2 \cdots \hat{a}_i \cdots a_n \\ &\equiv 0 \pmod{(S, w)}, \end{aligned}$$

where $w' = c_0 \langle c_1 \cdots c_k d \rangle$ and $w = \overline{w' a_1 a_2 \cdots \hat{a}_i \cdots a_n}$.

(2.12) If $d = c_i$ for some $i \geq 2$, or $d = c_1$ and $c_0 > c_2$, then

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n - b_0 b_1 b_2 \cdots b_m \cdot s_2 &\equiv s_1 a_1 a_2 \cdots a_n - (c_0 c_1 \cdots \hat{c}_i \cdots c_k \cdot s_2) a_1 a_2 \cdots a_n \\ &\equiv (s_1 - c_0 c_1 \cdots \hat{c}_i \cdots c_k \cdot s_2) a_1 a_2 \cdots a_n \\ &\equiv C_{II}(s_1, s_2)_{\bar{s}_1} a_1 a_2 \cdots a_n \\ &\equiv 0 \pmod{(S, w)}, \end{aligned}$$

where $c_i = d$.

(2.13) If $d = c_1$ and $c_0 \leq c_2$, then by the form of w , we have $b_0 b_1 \cdots b_m = c_0 \langle c_2 \cdots c_k a_1 \cdots a_n \rangle \in R$, which implies $c_2 \geq c_0 > a_1$. Thus,

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n - b_0 b_1 b_2 \cdots b_m \cdot s_2 &= s_1 a_1 a_2 \cdots a_n - c_0 a_1 \langle c_2 \cdots c_k a_2 \cdots a_n \rangle \cdot s_2 \\ &= (s_1 a_1 - c_0 a_1 c_2 \cdots c_k \cdot s_2) a_2 \cdots a_n \\ &= C_{IV}(s_1, s_2)_{w'} a_2 \cdots a_n \\ &\equiv 0 \pmod{(S, w)}, \end{aligned}$$

where $w' = c_0 \langle c_1 \cdots c_k a_1 \rangle$ and $w = w' a_2 \cdots a_n$.

(2.2) If $\bar{s}_1 = \overline{s_1^{(0)}} = c$ and $\overline{s_2^{(0)}} = d$, then $n = m + 1 \geq 2$ since $w = b_0 \langle b_1 \cdots b_m d \rangle$ and $m \geq 1$.

(2.21) If $c > a_1$, then $w = c a_1 \cdots a_n = b_0 \langle b_1 \cdots b_m d \rangle$, which implies $b_0 = c$.

(2.211) If $d \geq b_1$, then $a_1 = b_1, a_2 \cdots a_n = \langle b_2 \cdots b_m d \rangle$ and

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n - b_0 b_1 b_2 \cdots b_m \cdot s_2 &= (s_1 b_1 d) b_2 \cdots b_m - ((b_0 b_1) \cdot s_2) b_2 \cdots b_m \\ &= (s_1 b_1 d - (b_0 b_1) \cdot s_2) b_2 \cdots b_m \\ &= (s_1 b_1 \overline{s_2^{(0)}} - (\bar{s}_1 b_1) \cdot s_2^{(0)}) b_2 \cdots b_m \\ &= (s_1 b_1 (\overline{s_2^{(0)}} - s_2^{(0)}) - ((\bar{s}_1 - s_1) b_1) \cdot s_2^{(0)}) b_2 \cdots b_m \\ &= (s_1 b_1) (r_2^{(0)} b_2 \cdots b_m) - (r_1 b_1) b_2 \cdots b_m \cdot s_2 \\ &\equiv 0 \pmod{(S, w)}, \end{aligned}$$

where $s_2^{(0)} = \overline{s_2^{(0)}} + r_2^{(0)}$ and $s_1 = \bar{s}_1 + r_1$.

(2.212) If $d < b_1$, then $w = c d b_1 \cdots b_m$. Suppose that $s_1 = c + \sum_{c_i < c} \alpha_i c_i$, $s_2^{(0)} = d + \sum_{d_j < d} \beta_j d_j$. Thus,

$$\begin{aligned} s_1 a_1 a_2 \cdots a_n - b_0 b_1 b_2 \cdots b_m \cdot s_2 &= s_1 d b_1 b_2 \cdots b_m - c b_1 b_2 \cdots b_m \cdot s_2 \\ &= (s_1 d b_1 - (c b_1) \cdot s_2) b_2 \cdots b_m \\ &= \left(s_1 d b_1 - (s_1 b_1) s_2 + \sum_{c_i < c} \alpha_i (c_i b_1) s_2 \right) b_2 \cdots b_m \\ &= \left(s_1 b_1 d + (b_1 d) s_1 - (s_1 b_1) s_2 + \sum_{c_i < c} \alpha_i (c_i b_1) s_2 \right) b_2 \cdots b_m \end{aligned}$$

$$\begin{aligned}
 &= \left(s_1 b_1 (d - s_2) + (b_1 d) s_1 + \sum_{c_i < c} \alpha_i (c_i b_1) s_2 \right) b_2 \cdots b_m \\
 &= \left(- \sum_{d_j < d} \beta_j s_1 b_1 d_j + (b_1 d) s_1 + \left(\sum_{c_i < c} \alpha_i c_i b_1 \right) \cdot s_2 \right) b_2 \cdots b_m \\
 &= \left(- \sum_{d_j < d} \beta_j s_1 d_j b_1 + \sum_{d_j < d} \beta_j (b_1 d_j) s_1 + (b_1 d) s_1 + \left(\sum_{c_i < c} \alpha_i c_i b_1 \right) \cdot s_2 \right) b_2 \cdots b_m \\
 &\equiv 0 \pmod{(S, w)}.
 \end{aligned}$$

(2.22) If $c < a_1$, then $w = a_1 c a_2 \cdots a_n = b_0 (b_1 \cdots b_m d)$ and $a_1 = b_0$. In this case, $d \geq b_1$, and then $b_1 = c$, $d = a_i$ for some $i \geq 2$. Otherwise, if $d < b_1$, then $d = c$. This implies $a_i = b_{i-1}$ for any $i \geq 1$ and $b_0 = a_1 \leq a_2 = b_1$, which is a contradiction. Therefore,

$$\begin{aligned}
 s_1 a_1 a_2 \cdots a_n + b_0 b_1 b_2 \cdots b_m \cdot s_2 &= -(a_1 s_1) a_2 \cdots a_n + a_1 b_1 b_2 \cdots b_m \cdot s_2 \\
 &= -((a_1 s_1) d) a_2 \cdots \hat{a}_i \cdots a_n + (a_1 c) a_2 \cdots \hat{a}_i \cdots a_n \cdot s_2 \\
 &= ((s_1 a_1) \overline{s_2^{(0)}} + (a_1 \bar{s}_1) \cdot s_2) a_2 \cdots \hat{a}_i \cdots a_n \\
 &= ((s_1 a_1) (\overline{s_2^{(0)}} - s_2^{(0)}) + (a_1 (\bar{s}_1 - s_1)) \cdot s_2) a_2 \cdots \hat{a}_i \cdots a_n \\
 &\equiv 0 \pmod{(S, w)}.
 \end{aligned}$$

Case 3. $u_{s_1} = a_0 a_1 a_2 \cdots a_n \cdot s_1$, $u_{s_2} = b_0 b_1 b_2 \cdots b_n \cdot s_2$. We may assume that both s_1 and s_2 are (0)-monic. Suppose that $\overline{s_1^{(0)}} = c$ and $\overline{s_2^{(0)}} = d$. Then $w = a_0 \langle a_1 a_2 \cdots a_n c \rangle = b_0 \langle b_1 b_2 \cdots b_n d \rangle$ and $a_0 = b_0$.

(3.1) If $c = d$, then $a_i = b_i$ for all i and

$$a_0 a_1 a_2 \cdots a_n \cdot s_1 - b_0 b_1 b_2 \cdots b_n \cdot s_2 = a_0 a_1 a_2 \cdots a_n \cdot (s_1 - s_2).$$

If $s_1^{(1)} = s_2^{(1)} = 0$, i.e., $\bar{s}_1 = \overline{s_1^{(0)}} = \overline{s_2^{(0)}} = \bar{s}_2 = c$, then

$$a_0 a_1 a_2 \cdots a_n \cdot (s_1 - s_2) = a_0 a_1 a_2 \cdots a_n \cdot C_I(s_1, s_2) \equiv 0 \pmod{(S, w)}.$$

If $s_1^{(1)} \neq 0$, then

$$\begin{aligned}
 a_0 a_1 a_2 \cdots a_n \cdot (s_1 - s_2) &= ((a_0 a_1) (s_1 - s_2)) a_2 \cdots a_n \\
 &= C_{VI}(s_1, s_2)_{w'} a_2 \cdots a_n \\
 &\equiv 0 \pmod{(S, w)},
 \end{aligned}$$

where $w' = a_0 \langle a_1 c \rangle$.

(3.2) If $c \neq d$, say, $c > d$, then $w = a_0 \langle c d a_1 \cdots \hat{a}_i \cdots a_n \rangle = a_0 \langle c d b_1 \cdots \hat{b}_j \cdots b_n \rangle$ for some a_i and b_j .

(3.21) If $d \geq b_1$, then $w = a_0 b_1 \langle c d b_2 \cdots \hat{b}_j \cdots b_n \rangle = a_0 a_1 \langle c d a_2 \cdots \hat{a}_i \cdots a_n \rangle$, which implies $a_1 = b_1$, $a_2 \cdots \hat{a}_i \cdots a_n = b_2 \cdots \hat{b}_j \cdots b_n$. Thus,

$$\begin{aligned}
 a_0 a_1 a_2 \cdots a_n \cdot s_1 - b_0 b_1 b_2 \cdots b_n \cdot s_2 &= ((a_0 b_1 d) \cdot s_1) a_2 \cdots \hat{a}_i \cdots a_n - ((a_0 b_1 c) \cdot s_2) b_2 \cdots \hat{b}_j \cdots b_n \\
 &= (a_0 b_1 d \cdot s_1 - a_0 b_1 c \cdot s_2) b_2 \cdots \hat{b}_j \cdots b_n
 \end{aligned}$$

$$\begin{aligned}
 &= (a_0b_1(d - s_2) \cdot s_1 - a_0b_1(c - s_1) \cdot s_2)b_2 \cdots \hat{b}_j \cdots b_n \\
 &\equiv 0 \pmod{(S, w)}.
 \end{aligned}$$

(3.22) If $d < b_1$, then $w = a_0db_1 \cdots b_n = a_0a_1(a_2 \cdots a_nc)$, which implies $a_1 = d$ and $c = b_i$ for some i .

(3.221) If $c = b_1 < a_0$, then $a_i = b_i$ ($i \geq 2$) and $w = a_0dcb_2 \cdots b_n$. We have

$$\begin{aligned}
 a_0a_1a_2 \cdots a_n \cdot s_1 - b_0b_1b_2 \cdots b_n \cdot s_2 &= ((a_0d) \cdot s_1)a_2 \cdots a_n - ((a_0c) \cdot s_2)a_2 \cdots a_n \\
 &= (a_0d \cdot s_1 - a_0c \cdot s_2)a_2 \cdots a_n.
 \end{aligned}$$

If $s_1^{(1)} = 0$, then we may suppose that $s_1 = c + \sum_{c_i < c} \alpha_i c_i$ and $s_2^{(0)} = d + \sum_{d_j < d} \beta_j d_j$. We have

$$\begin{aligned}
 (a_0ds_1 - a_0cs_2)a_2 \cdots a_n &= ((a_0s_1)d + s_1da_0 - a_0c \cdot s_2)a_2 \cdots a_n \\
 &= \left((a_0s_1)s_2 - a_0c \cdot s_2 + s_1da_0 + \sum_{d_j < d} \beta_j s_1 a_0 d_j \right) a_2 \cdots a_n \\
 &= \left(a_0(s_1 - c)s_2 + s_1da_0 + \sum_{d_j < d} \beta_j s_1 d_j a_0 d_j - \sum_{d_j < d} \beta_j (a_0 d_j) s_1 \right) a_2 \cdots a_n \\
 &= \left(\sum_{c_i < c} \alpha_i (a_0 c_i) s_2 + s_1 da_0 + \sum_{d_j < d} \beta_j s_1 d_j a_0 d_j - \sum_{d_j < d} \beta_j (a_0 d_j) s_1 \right) a_2 \cdots a_n \\
 &\equiv 0 \pmod{(S, w)}.
 \end{aligned}$$

If $s_2^{(1)} = 0$, then we have

$$\begin{aligned}
 (a_0ds_1 - a_0c \cdot s_2)a_2 \cdots a_n &= (a_0ds_1 - a_0s_2c - s_2ca_0)a_2 \cdots a_n \\
 &= (a_0(d - s_2)s_1 - a_0s_2(c - s_1) - s_2ca_0)a_2 \cdots a_n \\
 &\equiv 0 \pmod{(S, w)}.
 \end{aligned}$$

If $s_i^{(1)} \neq 0$ ($i = 1, 2$), then let $w' = a_0dc$. We have $w = w'a_2 \cdots a_n$ and

$$\begin{aligned}
 (a_0ds_1 - a_0c \cdot s_2)a_2 \cdots a_n &= C_{VII}(s_2, s_1)_{w'} a_2 \cdots a_n \\
 &\equiv 0 \pmod{(S, w)}.
 \end{aligned}$$

(3.222) If $c = b_i > b_1$ for some $i \geq 2$, then

$$\begin{aligned}
 a_0a_1a_2 \cdots a_n \cdot s_1 - b_0b_1b_2 \cdots b_n \cdot s_2 &= ((a_0db_1) \cdot s_1 - (a_0b_1c) \cdot s_2)b_2 \cdots \hat{b}_i \cdots b_n \\
 &= (a_0b_1d \cdot s_1 + b_1da_0 \cdot s_1 - a_0b_1c \cdot s_2)b_2 \cdots \hat{b}_i \cdots b_n \\
 &\equiv (a_0b_1\overline{s_2^{(0)}} \cdot s_1 - a_0b_1\overline{s_1^{(0)}} \cdot s_2)b_2 \cdots \hat{b}_i \cdots b_n \\
 &\equiv ((a_0b_1) \cdot s_2) \cdot s_1 - ((a_0b_1) \cdot s_1) \cdot s_2)b_2 \cdots \hat{b}_i \cdots b_n \\
 &\equiv 0 \pmod{(S, w)}.
 \end{aligned}$$

The proof is complete. \square

Theorem 2.9 (Composition-Diamond lemma for metabelian Lie algebras). *Let $S \subset \mathcal{L}_{(2)}(X)$ be a nonempty set of monic polynomials and $Id(S)$ be the ideal of $\mathcal{L}_{(2)}(X)$ generated by S . Then the following statements are equivalent.*

- (i) S is a Gröbner–Shirshov basis.
- (ii) $f \in Id(S) \Rightarrow \bar{f} = \bar{u}_s$ for some normal S -word u_s .
- (iii) $Irr(S) = \{u \mid u \in N, u \neq \bar{v}_s \text{ for any normal } S\text{-word } v_s\}$ is a \mathbf{k} -basis for $\mathcal{L}_{(2)}(X)/Id(S) = \mathcal{L}_{(2)}(X)/Id(S)$.

Proof. (i) \Rightarrow (ii). Let S be a Gröbner–Shirshov basis and $0 \neq f \in Id(S)$. Then by Lemma 2.3 f has an expression $f = \sum \alpha_i u_{s_i}$, where $0 \neq \alpha_i \in \mathbf{k}$, u_{s_i} are normal S -words. Denote $w_i = \bar{u}_{s_i}$, $i = 1, 2, \dots$. We may assume without loss of generality that

$$w_1 = w_2 = \dots = w_l > w_{l+1} \geq w_{l+2} \geq \dots$$

for some $l \geq 1$.

The claim of the theorem is obvious if $l = 1$.

Now suppose that $l > 1$. Then $\bar{u}_{s_1} = w_1 = w_2 = \bar{u}_{s_2}$. By Lemma 2.8, for some $\alpha \in \mathbf{k}$,

$$u_{s_2} \equiv \alpha u_{s_1} \pmod{(S, w_1)}.$$

Thus,

$$\begin{aligned} \alpha_1 u_{s_1} + \alpha_2 u_{s_2} &= (\alpha_1 + \alpha \alpha_2) u_{s_1} + \alpha_2 (u_{s_2} - \alpha u_{s_1}) \\ &\equiv (\alpha_1 + \alpha \alpha_2) u_{s_1} \pmod{(S, w_1)}. \end{aligned}$$

Therefore, if $\alpha_1 + \alpha \alpha_2 \neq 0$ or $l > 2$, then the result follows from the induction on l . For the case $\alpha_1 + \alpha \alpha_2 = 0$ and $l = 2$, we use the induction on w_1 . Now the result follows.

(ii) \Rightarrow (iii). For any $f \in \mathcal{L}_{(2)}(X)$, we have

$$f = \sum_{\bar{u}_{s_i} \leq \bar{f}} \alpha_i u_{s_i} + \sum_{\bar{v}_j \leq \bar{f}} \beta_j v_j,$$

where $\alpha_i, \beta_j \in \mathbf{k}$, $v_j \in Irr(S)$ and u_{s_i} are normal S -words. Therefore, the set $Irr(S)$ generates the algebra $\mathcal{L}_{(2)}(X)/Id(S)$.

On the other hand, suppose that $h = \sum \alpha_i v_i = 0$ in $\mathcal{L}_{(2)}(X)/Id(S)$, where $\alpha_i \in \mathbf{k}$, $v_i \in Irr(S)$. This means that $h \in Id(S)$. Then all α_i must be equal to zero. Otherwise, $\bar{h} = \bar{v}_j$ for some j which contradicts (ii).

(iii) \Rightarrow (i). For any $f, g \in S$, we have

$$C_\lambda(f, g)_w = \sum_{\bar{u}_{s_i} < w} \alpha_i u_{s_i} + \sum_{\bar{v}_j < w} \beta_j v_j.$$

Since $C_\lambda(f, g)_w \in Id(S)$ and by (iii), we have

$$C_\lambda(f, g)_w = \sum_{\bar{u}_{s_i} < w} \alpha_i u_{s_i}.$$

Therefore, S is a Gröbner–Shirshov basis. \square

Lemma 2.10. (See [10].) *Suppose that $f \in \mathcal{L}_{(2)}(X)$. Then there exists an element $f' \in \mathcal{L}_{(2)}(X)$ such that $Id(f) = Id(f')$, $\bar{f}' \leq \bar{f}$, $f'^{(0)} = f^{(0)}$ and no word occurring in $f'^{(1)}$ contains $\bar{f}^{(0)}$ as a strict subword.*

Proof. If no word occurring in $f^{(1)}$ contains $\overline{f^{(0)}}$ as a strict subword, then we are done. If $\overline{f^{(1)}}$ contains $\overline{f^{(0)}}$ as a strict subword, say $\bar{f} = \overline{f^{(1)}} = a_0 a_1 \cdots a_n$, $\overline{f^{(0)}} = a_i$ for some $i \geq 2$ or $\overline{f^{(0)}} = a_1$ and $a_0 > a_2$, then let f_1 be the composition of type II of f and itself:

$$f_1 = C_{II}(f, f)_{\bar{f}} = f - \beta^{-1} a_0 a_1 \cdots \hat{a}_i \cdots a_n \cdot f,$$

where $a_i = \overline{f^{(0)}}$. It is obvious that $Id(f) = Id(f_1)$, and $\bar{f}_1 < \bar{f}$, $f_1^{(0)} = f^{(0)}$. If $\overline{f_1^{(1)}}$ contains $\overline{f^{(0)}}$ as a strict subword, we again consider the composition $f_2 = C_{II}(f_1, f_1)_{\bar{f}_1}$, and so on. By induction on the leading word, we obtain an element f' such that $Id(f) = Id(f')$, $\bar{f}' \leq \bar{f}$, $f'^{(0)} = f^{(0)}$, and either $f' = f'^{(0)}$ or $\overline{f'^{(1)}}$ does not contain $\overline{f^{(0)}}$ as a strict subword.

Arguments analogous to the one given above for the leading word also apply to other regular R -words occurring in the expansion of f and containing $\overline{f^{(0)}}$ as a strict subword. Finally, we have the one we want. \square

Lemma 2.11. Suppose that $\bar{f} = \overline{f^{(1)}} = a_0 a_1 \cdots a_n$, $g^{(1)} \neq 0$, $\overline{g^{(0)}} = a_1$ and $a_0 \leq a_2$ or $n = 1$. If $f^{(0)} = 0$, then for $a = a_1 < a_0$ and $w = a_0(a_1 \cdots a_n a)$, the composition of type IV of f and g is trivial.

Proof. We may suppose that g is (0)-monic. Then

$$\begin{aligned} C_{IV}(f, g)_w &= f a_1 - \bar{f} \cdot g \\ &= r_f^{(1)} \cdot \overline{g^{(0)}} - \bar{f} \cdot r_g^{(0)} \\ &= r_f^{(1)} \cdot \overline{g^{(0)}} - \bar{f} \cdot r_g^{(0)} + r_f^{(1)} \cdot r_g^{(0)} - r_f^{(1)} \cdot r_g^{(0)} \\ &= r_f^{(1)} (\overline{g^{(0)}} + r_g^{(0)}) - (\bar{f} + r_f^{(1)}) \cdot r_g^{(0)} \\ &= r_f^{(1)} \cdot g - f \cdot r_g^{(0)} \\ &\equiv 0 \pmod{\{f, g\}, w}, \end{aligned}$$

where $f = f^{(1)} = \bar{f} + r_f^{(1)}$ and $g^{(0)} = \overline{g^{(0)}} + r_g^{(0)}$. \square

Lemma 2.12. The compositions of type I, V and VI formed by f itself are always trivial.

Proof. For type I and VI, the result is obvious. We only check type V. Suppose that $\bar{f} = \overline{f^{(1)}} = a_0 a_1 \cdots a_n$, $\overline{f^{(0)}} = b \notin \{a_i\}_{i=1}^n$, and $w = a_0(a_1 \cdots a_n b)$. We have

$$\begin{aligned} C_V(f, f)_w &= f b - \beta^{-1} a_0 a_1 \cdots a_n \cdot f \\ &= f \cdot \overline{f^{(0)}} - \beta^{-1} \bar{f} \cdot f \\ &= f \cdot \overline{f^{(0)}} - f \cdot \beta^{-1} (r^{(1)} + \beta \overline{f^{(0)}} + r^{(0)}) \\ &= -\beta^{-1} f \cdot (r^{(1)} + r^{(0)}) \\ &= \beta^{-1} r^{(1)} \cdot f - \beta^{-1} f \cdot r^{(0)} \\ &\equiv 0 \pmod{f, w}, \end{aligned}$$

where $f^{(1)} = \bar{f} + r^{(1)}$ and $f^{(0)} = \beta \overline{f^{(0)}} + r^{(0)}$, $\beta \in \mathbf{k}$. \square

Remark. If a subset S of $\mathcal{L}_{(2)}(X)$ is not a Gröbner–Shirshov basis, then one can add all nontrivial compositions of polynomials of S to S . Continuing this process repeatedly, we finally obtain a Gröbner–Shirshov basis S^C that generates the same ideal as S . Such a process is called Shirshov’s algorithm and S^C is called a Gröbner–Shirshov complement of S . By Lemma 2.10, we may assume that any element of the original relation set S has no composition of type II formed by itself and Shirshov’s algorithm does not involve compositions discussed in Lemmas 2.11 and 2.12.

3. Applications

Suppose that \mathcal{A} is a metabelian Lie algebra and $Y = \{a_i, i \in I\} \cup \{b_j, j \in J\}$ is a \mathbf{k} -basis of \mathcal{A} , where $\{a_i\}$ is a basis of $\mathcal{A}^{(1)}$ and $\{b_j, j \in J\}$ is linear independent modulo $\mathcal{A}^{(1)}$. Suppose that I and J are well-ordered sets. The set of multiplications of Y , say M , consists of the following:

$$\begin{aligned} m_{1ij}: a_i b_j - \sum \gamma_{ij}^k a_k, \\ m_{2ij}: b_i b_j - \sum \delta_{ij}^k a_k \quad (i > j), \\ m_{3ij}: a_i a_j \quad (i > j), \end{aligned}$$

where $\gamma_{ij}^k, \delta_{ij}^k \in \mathbf{k}$. Then we have $\mathcal{A} = \mathcal{L}_{(2)}(Y|M)$ and since $\text{Irr}(M) = Y$, by Theorem 2.9, M is a Gröbner–Shirshov basis for \mathcal{A} with respect to $a_i > b_j$.

Let S denote the free metabelian Lie product of \mathcal{A} and a free metabelian Lie algebra generated by a well-ordered set $X = \{x_h \mid h \in H\}$, i.e.,

$$S = \mathcal{A} * \mathcal{L}_{(2)}(X) = \mathcal{L}_{(2)}(X \cup Y|M).$$

Theorem 3.1. *Let the notion be as above. Then with respect to $x_h > a_i > b_j$, a Gröbner–Shirshov complement M^C of M in $\mathcal{L}_{(2)}(X \cup Y)$ consists of M and some X -homogeneous polynomials without (0)-part, whose leading words are of the form $xy \cdots$ with an a_i as a strict subword, $x \in X, a_i, y \in Y$.*

Proof. For convenience, we call the X -homogeneous polynomials described in the theorem to satisfy property P_X .

Since M is a Gröbner–Shirshov basis in $\mathcal{L}_{(2)}(Y)$, we need to check the compositions which are formed by M itself and involve some elements in X . The possible types are VI and VII.

First, we check type VI. Suppose that $\overline{m_{1ij}^{(0)}} = \overline{m_{1st}^{(0)}} = a_l$ and the corresponding w is of the forms $xx'a_l, xba_l$ and $x(aa_l)$ for some $x, x' \in X, b \in \{b_j\}$ and $a \in \{a_i\}$.

If $w = xx'a_l$, then

$$\begin{aligned} C_{VI}(m_{1ij}, m_{1st})_w &= (xx')((\gamma_{ij}^l)^{-1} m_{1ij} - (\gamma_{st}^l)^{-1} m_{1st}) \\ &= - \sum_{k < l} (\gamma_{ij}^l)^{-1} \gamma_{ij}^k xx' a_k + \sum_{k < l} (\gamma_{st}^l)^{-1} \gamma_{st}^k xx' a_k \\ &= - \sum_{k < l} (\gamma_{ij}^l)^{-1} \gamma_{ij}^k xa_k x' + \sum_{k < l} (\gamma_{st}^l)^{-1} \gamma_{st}^k xa_k x' \\ &\quad + \sum_{k < l} (\gamma_{ij}^l)^{-1} \gamma_{ij}^k x' a_k x - \sum_{k < l} (\gamma_{st}^l)^{-1} \gamma_{st}^k x' a_k x \end{aligned}$$

and obviously it satisfies P_X .

If $w = xba_l$, then

$$\begin{aligned} C_{VI}(m_{1ij}, m_{1st})_w &= (xb)((\gamma_{ij}^l)^{-1}m_{1ij} - (\gamma_{st}^l)^{-1}m_{1st}) \\ &= -\sum_{k<l} (\gamma_{ij}^l)^{-1}\gamma_{ij}^k xba_k + \sum_{k<l} (\gamma_{st}^l)^{-1}\gamma_{st}^k xba_k \end{aligned}$$

and it still satisfies P_X .

If $w = xaa_l$, then

$$\begin{aligned} C_{VI}(m_{1ij}, m_{1st})_w &= (xa)((\gamma_{ij}^l)^{-1}m_{1ij} - (\gamma_{st}^l)^{-1}m_{1st}) \\ &= -\sum_{k<l} (\gamma_{ij}^l)^{-1}\gamma_{ij}^k xaa_k + \sum_{k<l} (\gamma_{st}^l)^{-1}\gamma_{st}^k xaa_k \\ &\equiv -\sum_{a_k < a} (\gamma_{ij}^l)^{-1}\gamma_{ij}^k xa_k a + \sum_{a_k < a} (\gamma_{st}^l)^{-1}\gamma_{st}^k xa_k a \\ &\quad - \sum_{a_k \geq a} (\gamma_{ij}^l)^{-1}\gamma_{ij}^k xaa_k + \sum_{a_k \geq a} (\gamma_{st}^l)^{-1}\gamma_{st}^k xaa_k \pmod{(M, w)}, \end{aligned}$$

and again the remainder satisfies P_X .

$C_{VI}(m_{1ij}, m_{2st})_w, C_{VI}(m_{2ij}, m_{2st})_w$ are similar to $C_{VI}(m_{1ij}, m_{1st})_w$.

Second, we check type VII. Suppose that $m_{1ij}^{(0)} = a_p > a_q = m_{1st}^{(0)}$ and $w = xaqa_p$. Then

$$\begin{aligned} C_{VII}(m_{1ij}, m_{1st})_w &= (\gamma_{ij}^p)^{-1}(xa_q)m_{1ij} - (\gamma_{st}^q)^{-1}(xa_p)m_{1st} \\ &= -\sum_{k<p} (\gamma_{ij}^p)^{-1}\gamma_{ij}^k xa_q a_k + \sum_{k<q} (\gamma_{st}^q)^{-1}\gamma_{st}^k xa_p a_k - x(a_p a_q) \\ &= -\sum_{q \leq k < l} (\gamma_{ij}^p)^{-1}\gamma_{ij}^k xa_q a_k - \sum_{k<q} (\gamma_{ij}^p)^{-1}\gamma_{ij}^k xa_k a_q - \sum_{q \leq k < l} (\gamma_{ij}^p)^{-1}\gamma_{ij}^k x(a_q a_k) \\ &\quad + \sum_{k<q} (\gamma_{st}^q)^{-1}\gamma_{st}^k xa_k a_p + \sum_{k<q} (\gamma_{st}^q)^{-1}\gamma_{st}^k x(a_p a_k) - x(a_p a_q) \\ &\equiv -\sum_{q \leq k < l} (\gamma_{ij}^p)^{-1}\gamma_{ij}^k xa_q a_k - \sum_{k<q} (\gamma_{ij}^p)^{-1}\gamma_{ij}^k \\ &\quad + xa_k a_q \sum_{k<q} (\gamma_{st}^q)^{-1}\gamma_{st}^k xa_k a_p \pmod{(M, w)}, \end{aligned}$$

and the remainder has property P_X . One may check that $C_{VII}(m_{1ij}, m_{2st})_w$ and $C_{VII}(m_{2ij}, m_{2st})_w$ are the same as $C_{VII}(m_{1ij}, m_{1st})_w$, which have property P_X .

Observing the above and the definition of compositions, we know that the nontrivial compositions of polynomials satisfying P_X themselves are only of type I and the results again satisfy P_X . Also by the definition of compositions and property P_X , the compositions of M and polynomials satisfying P_X are only of type II and the results still satisfy P_X . The theorem is proved. \square

Observing the proof of the above theorem, we have the following proposition.

Proposition 3.2. Let $\mathcal{A}_i = \mathcal{L}_{(2)}(X_i | S_i)$, where $S_i \subset \mathcal{L}_{(2)}(X_i)^{(1)}$, $i = 1, 2$. Then $S_1^C \cup S_2^C$ is a Gröbner–Shirshov basis for the free metabelian Lie product $\mathcal{A}_1 * \mathcal{A}_2$, where S_i^C is a Gröbner–Shirshov complement of S_i in $\mathcal{L}_{(2)}(X_i)$, $i = 1, 2$.

Now, we consider partial commutative metabelian Lie algebras related to some graphs.

Let $\Gamma = (V, E)$ be a graph, where V is the set of vertices and E the set of edges. For $e \in E$ we call $o(e)$ the origin of e and $t(e)$ the terminus. We say a metabelian Lie algebra is partial commutative related to a graph $\Gamma = (V, E)$, denoted by \mathcal{ML}_{Γ} , if

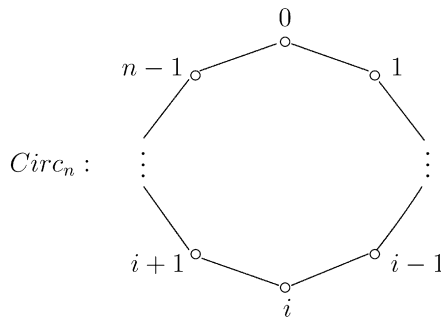
$$\mathcal{ML}_{\Gamma} = \mathcal{L}_{(2)}(V \mid [o(e), t(e)] = 0, e \in E).$$

In this section, we find Gröbner–Shirshov bases for partial commutative metabelian Lie algebras related to circuits, trees and 3-cube.

The following algorithm gives a Gröbner–Shirshov basis for partial commutative metabelian Lie algebras with a finite relation set.

Algorithm 3.3. Input: relations f_1, \dots, f_s of $\mathcal{L}_{(2)}(X)$, $f_i = xx'$, $F = \{f_1, \dots, f_s\}$.
 Output: a Gröbner–Shirshov basis $H = \{h_1, \dots, h_t\}$ for $\mathcal{L}_{(2)}(X|F)$.
 Initialization: $H := F$
 While: $f_i = x_{i_0}x_{i_1} \cdots x_{i_n}$, $f_j = x_{j_0}x_{j_1} \cdots x_{j_m}$, and $x_{i_0} = x_{j_0}$, $x_{i_1} \neq x_{j_1}$
 Then Do: $h := \max\{x_{i_1}, x_{j_1}\} \min\{x_{i_1}, x_{j_1}\} \langle x_{t_1}, x_{t_2} \cdots x_{t_l} \rangle$
 where $\{x_{t_1}, x_{t_2}, \dots, x_{t_l}\} = \{x_{i_0}, x_{i_2}, \dots, x_{i_n}\} \cup \{x_{j_2}, \dots, x_{j_m}\}$
 If: there is no $f_j \in H$ such that f_j is a subword of h
 Do: $H := H \cup \{h\}$
 End

Definition 3.4. Let n be a positive integer. A circuit (of length n), denoted by $Circ_n$, is a graph for which the set of vertices is $\mathbf{Z}/n\mathbf{Z}$ and the orientation is given by n edges $e_{i,i+1}$, $i \in \mathbf{Z}/n\mathbf{Z}$, with $o(e_{i,i+1}) = i$ and $t(e_{i,i+1}) = i + 1$.



Theorem 3.5. For the partial commutative metabelian Lie algebra related to $Circ_n$

$$\mathcal{ML}_{Circ_n} = \mathcal{L}_{(2)}(\mathbf{Z}/n\mathbf{Z} \mid [i + 1, i] = 0, i \in \mathbf{Z}/n\mathbf{Z}),$$

with the usual ordering on natural numbers, a Gröbner–Shirshov basis for \mathcal{ML}_{Circ_n} consists of the following relations:

$$f_0: [n - 1, 0] = 0,$$

$$f_i: [i, i - 1] = 0, \quad 1 \leq i \leq n - 1,$$

$$g_j: [j, 0, j + 1, j + 2, \dots, n - 1] = 0, \quad 2 \leq j \leq n - 2,$$

where the brackets $[\dots]$ are the left-normed brackets.

Proof. The only possible compositions are of type I by f_{n-1}, f_0 and g_j, f_j , where the corresponding w 's are $[n - 1, 0, n - 2]$ and $[j, 0, j - 1, j + 1, j + 2, \dots, n - 1]$ respectively.

For the first one, $w = [n - 1, 0, n - 2]$ and

$$\begin{aligned} C_I(f_{n-1}, f_0)_w &= [n - 1, n - 2] \cdot 0 - [n - 1, 0, n - 2] \\ &= [n - 2, 0, n - 1] \\ &\equiv 0 \pmod{(g_{n-2}, w)}. \end{aligned}$$

For the second one, $w = [j, 0, j - 1, j + 1, j + 2, \dots, n - 1]$ and

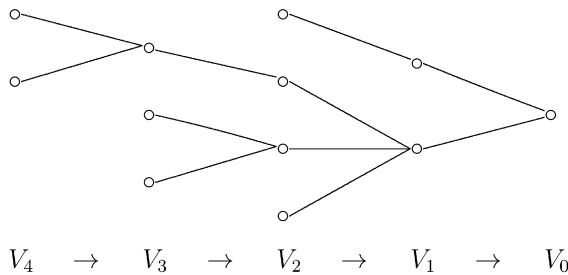
$$\begin{aligned} C_I(g_j, f_j)_w &= [j, 0, j + 1, j + 2, \dots, n - 1] \cdot (j - 1) - [j, j - 1, 0, j + 2, \dots, n - 1] \\ &= [j - 1, 0, j, j + 1, j + 2, \dots, n - 1]. \end{aligned}$$

Then it is trivial modulo f_2 if $j = 2$ and modulo g_{j-1} if $j \geq 3$. \square

Definition 3.6. A tree is a connected nonempty graph without circuits.

A geodesic in a tree is a path without backtracking. The length of the geodesic from v to v' is called the distance from v to v' , and is denoted by $l(v, v')$.

Fix a vertex v_0 of a tree Γ . For each integer $n \geq 0$, let V_n be the set of vertices v of Γ such that $l(v_0, v) = n$. Then the set of vertices of Γ is the union of V_n and $V_i \cap V_j = \emptyset, i \neq j$. If $v \in V_n$ with $n \geq 1$, there is a single vertex $v' \in V_{n-1}$ from v_0 to which v is adjacent.



We linearly order the set of vertices $V = \bigcup_{n \geq 0} V_n$ such that v_0 is the smallest element and for any $v \in V_i, v' \in V_j, v < v'$ if $i < j$. Then the partial commutative metabelian Lie algebra related to the tree Γ is defined by:

$$\mathcal{ML}_\Gamma = \mathcal{L}_{(2)}(V|R),$$

where

$$R = \{[v', v] = 0 \mid v' \in V_{n+1}, v \in V_n, v' \text{ and } v \text{ are adjacent}, n \geq 0\}.$$

Theorem 3.7. The relation set R forms a Gröbner–Shirshov basis for the partial commutative metabelian Lie algebra \mathcal{ML}_Γ related to the tree Γ .

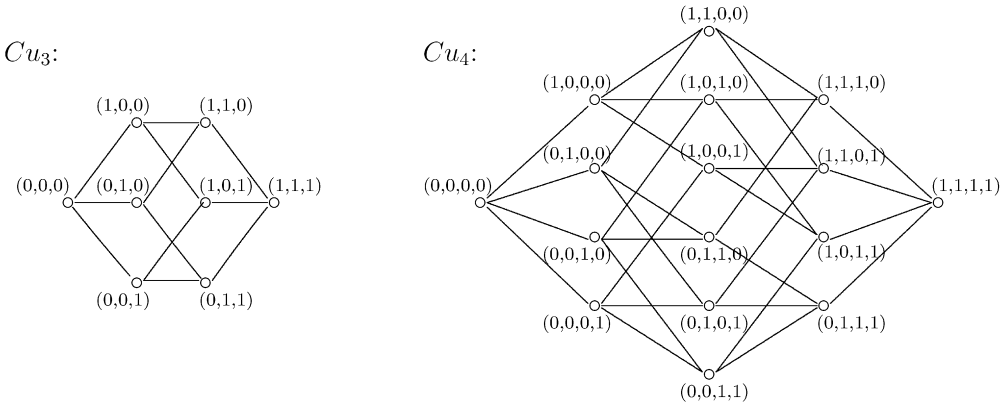
Proof. It is obvious that for any $v' \in V_{n+1}$, there is only one element $v \in V_n$ such that the relation $[v', v] = 0$ lies in R , which means there is no composition in R at all. Thus, R is a Gröbner–Shirshov basis automatically. \square

By Theorems 2.9 and 3.7, we have the following corollary.

Corollary 3.8. A linear basis of \mathcal{ML}_Γ consists of regular words $v_0 v_1 \cdots v_n$ ($n \geq 0$) on V satisfying the following condition: if $v_0 > v_i$ ($i \geq 1$), then $l(v_0, v_i) \neq 1$.

Definition 3.9. Let n be a positive integer. An n -cube, denoted by Cu_n , is a graph for which the set of vertices $V_n = \{(\varepsilon_1, \varepsilon_2, \dots, \varepsilon_n) \in \mathbb{R}^n \mid \varepsilon_i = 0 \text{ or } 1\}$ and two vertices $\varepsilon = (\varepsilon_1, \varepsilon_2, \dots, \varepsilon_n)$, $\delta = (\delta_1, \delta_2, \dots, \delta_n)$ are adjacent if $\exists i$, such that $\varepsilon_i = \delta_i + 1 \pmod 2$ and $\varepsilon_j = \delta_j$ for any $j \neq i$.

For example, 3-cube and 4-cube are the followings:



We order all vertices lexicographically. The distance of ε and δ is $d(\varepsilon, \delta) = \sum_{i=1}^n |\varepsilon_i - \delta_i|$. Then the partial commutative metabelian Lie algebra related to the n -cube Cu_n is defined by:

$$\mathcal{ML}_\Gamma = \mathcal{L}_{(2)}(V_n \mid \varepsilon\delta = 0, d(\varepsilon, \delta) = 1).$$

Theorem 3.10. A Gröbner–Shirshov basis S for the partial commutative metabelian Lie algebra related to 3-cube

$$\mathcal{ML}_{Cu_3} = \mathcal{L}_{(2)}(V_3 \mid \varepsilon\delta, d(\varepsilon, \delta) = 1, \varepsilon > \delta)$$

is the union of the following:

- $R_2 = \{[\varepsilon\delta] \mid d(\varepsilon, \delta) = 1\},$
- $R_3 = \{[\varepsilon\delta]\mu \mid d(\varepsilon, \delta) = 2, \mu\varepsilon, \mu\delta \in R_1\},$
- $R_4 = \{[\varepsilon\delta]\mu\gamma \mid d(\varepsilon, \delta) = 3, \mu\varepsilon \in R_2, \mu\delta\gamma \in R_3\},$
- $R_5 = \{[\delta_1\delta_2]\gamma\langle\mu_1\mu_2\rangle \mid d(\delta_1, \delta_2) = 2, \gamma\delta_i\mu_i \in R_3, i = 1, 2\},$
- $R'_5 = \{[\delta_1\delta_2]\gamma\mu\mu' \mid d(\delta_1, \delta_2) = 2, \gamma\delta_1 \in R_2, \gamma_2\mu\mu' \in R_4, d(\mu, \delta_1) \neq 1\},$

where $[\varepsilon\delta] = \max\{\varepsilon, \delta\} \min\{\varepsilon, \delta\}$.

By Algorithm 3.3, we have that a reduced Gröbner–Shirshov basis (it means there is no composition of type I, II, III) for the partial commutative metabelian Lie algebra related to 4-cube \mathcal{ML}_{Cu_4} consists of 268 relations.

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