

# Mixed Bundling Auctions

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## Abstract

We study multi-object auctions where agents have private and additive valuations for heterogeneous objects. We focus on the revenue properties of a class of dominant strategy mechanisms where a weight is assigned to each partition of objects. The weights influence the probability with which partitions are chosen in the mechanism. This class contains efficient auctions, pure bundling auctions, mixed bundling auctions, auctions with reserve prices and auctions with pre-packaged bundles. For any number of objects and bidders, both the pure bundling auction and separate, efficient auctions for the single objects are revenue-inferior to an auction that involves mixed bundling.

## 1 Introduction

Very little is known about the revenue maximizing auction for multiple, heterogeneous objects, and we doubt that the problem is analytically tractable. It is well known that reserve prices may increase expected revenue by inefficiently withholding the objects. An important and different insight is that an optimal auction will not necessarily allocate the objects to the buyers that value them most<sup>1</sup>. Palfrey (1983) introduced the study of *bundling* to the auction literature, and showed that a *pure bundling auction* is revenue

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<sup>1</sup>This, in principle, leaves scope for a resale market. As we want to focus on the role of bundling, we assume here that the resale-market is absent.

superior to separate (and efficient) auctions if there are only two bidders<sup>2</sup>. Under pure bundling, a single bidder gets all objects if she has the highest value for the entire bundle, although other bidders attach higher values to parts of the bundle. Under some assumptions on the distribution of valuations, Chakraborty (1999) showed that there is a critical number of bidders, below which bundling outperforms separate auctions, and above which separate auctions outperform the bundling auction. This analysis suggests that pure bundling may increase revenue by increasing competition (i.e., by "importing" demand from one object to another), but that it may also decrease revenue by creating inefficiencies. The logic behind Chakraborty's result is that, with many bidders, the inefficiency effect outweighs the competition effect.

The main purpose of this paper is to construct and analyze the properties of a class of *mixed bundling auctions* that improve upon both pure bundling and upon separate, efficient auctions. Moreover, we offer a methodology that can be used to identify the most promising bundling strategy.

In our model, valuations for bundles are equal to the sum of the valuations for the included objects, so that bundling is solely driven by strategic considerations of the auctioneer rather than by technological features of the demand functions (such as complementarities).

The constructed mechanisms assign to each partition of objects among the agents a partition-specific weight. For each realization of values, the implemented partition is the one that maximizes the sum of the bidders' values and the assigned weight. Thus, our mechanisms can be seen as weighted Vickrey-Clarke-Groves mechanisms, and it is a dominant strategy for bidders to reveal their true values. For example, assigning a positive weight to all partitions where some bundle of objects is not sold corresponds to introducing a reserve price on this bundle. Assigning a large positive weight to all partitions where the objects are sold to a single buyer corresponds to a pure bundling auction, while a smaller weight on such partitions will sometimes allow for separate sales, thus yielding mixed bundling.

The advantages of mixed bundling have been previously pointed out by the literature on monopolistic pricing. One might think that the optimal pricing strategy for a monopolist selling two goods (for which customers' valuations are known to be additive) is to price each good at the price a single-product monopolist would optimally charge. Adams and Yellen (1976) show (via examples) that this is not necessarily the case. McAfee et al. (1989) generalize the non-optimality of additive pricing to arbitrarily dis-

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<sup>2</sup>Jehiel and Moldovanu (2001) exploit Palfrey's results from the perspective of the revenue equivalence theorem.

tributed valuations, and suggest a class of situations in which the optimal prices include a discount for bundles - this is referred to as mixed bundling by Adams and Yellen<sup>3</sup>. Manelli and Vincent (2001) determine necessary conditions for the optimal monopolistic prices for each bundle of objects, and study how non-deterministic sale mechanisms may improve upon deterministic ones. Armstrong (1996) considers a more general model where objects are divisible, and he looks for optimal tariffs. His main result is that the seller will always want to withhold sales from some low-type buyers, no matter how profitable selling some quantity to them might be. We can apply Armstrong's argument to our auction model in order to show that the revenue of both pure and mixed bundling auctions can be increased by introducing reserve prices. Rochet and Choné (1998) further analyze the multi-product monopoly problem (by relaxing some of the assumptions used by Armstrong), and show how bunching (i.e. different types get the same allocation) is a robust feature of the optimal mechanism.

Since monopoly pricing can be viewed as an auction design with a unique potential bidder, one contribution of this paper is to generalize the idea of mixed bundling from monopoly pricing to the more involved theory of auctions.

Most of the auction literature (and, in particular, that literature searching for optimal mechanisms) has studied models where each bidder's information can be represented by a one-dimensional signal. Most prominently, Myerson (1981) and Riley and Samuelson (1981) find the revenue-maximizing auction<sup>4</sup> (in the class of Bayes-Nash implementable allocation rules) for the private values, one-object case. If bidders are ex-ante symmetric, the optimal auction is a second-price auction with a reserve price. Hence this mechanism is implementable in dominant strategies, and, as mentioned above, belongs to the class of mechanisms discussed here.

Models with one-dimensional signals are of limited use for studies of auctions for multiple, heterogeneous objects. Jehiel, Moldovanu and Stacchetti (1999) have shown that the main difficulty with multi-dimensional signals is that incentive compatibility constraints<sup>5</sup> yield a complex integrability condition on the vector of interim expected probabilities with which buyers get

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<sup>3</sup>In other cases the optimal prices contain a surcharge for the bundle instead of a discount. This raises the problem of deterring customers from buying the objects separately in order to evade the surcharge.

<sup>4</sup>Bulow and Roberts (1989) illustrate the parallels between optimal auction design and monopolistic pricing for one object.

<sup>5</sup>These constraints reduce to a monotonicity condition when private information is one-dimensional.

the various bundles<sup>6</sup>. Thus, under this constraint the maximization problem is very hard - this is the main obstacle in front of a characterization of the revenue-maximizing multi-object auction. Instead of trying to characterize such an auction, we focus here on a large class of mechanisms that includes pure and mixed bundling, as well as separate, efficient auctions. A major practical advantage is that revenue maximization in this class reduces to an optimization problem over vectors of real numbers (the weights attached to the partitions) that can be numerically solved by any personal computer.

Armstrong (2000) and Avery and Hendershott (2000) study 2-object auctions in discrete settings. In particular, Armstrong finds that the revenue maximizing auction is efficient. The discrepancy to our results stems from Armstrong's assumption that valuations are binary, i.e. they can either be "high" or "low", but not "intermediate". This leaves no scope for mixed bundling, as analyzed in this paper<sup>7</sup>.

As the focus of this paper is to identify and describe a tractable way of improving auction revenues rather than to characterize the optimal auction, we do not consider here how Bayes-Nash implementation might improve revenues over dominant strategy implementation. Mookerjee and Reichelstein (1992) describe a class of mechanism design problems (including the quest for a revenue-maximizing one object auction) where using the stronger requirement of dominant strategies is without loss of generality. Their result does not hold for general models with multidimensional signals, and very little is known about this issue in such contexts.

The paper is organized as follows: Section 2 lays out the model, and introduces the weighted Clarke-Groves-Vickrey mechanisms. In Section 3 we first calculate derivatives of revenue and expected revenue with respect to the weights attached to each partition of objects (Theorem 1). In Subsection 3.1 we use this Theorem to show how to improve upon pure bundling and pre-packaged auctions (Proposition 1). In Subsection 3.2 we use the Theorem in order to show how to improve upon separate, efficient auctions (Proposition 2). This results also identifies the most promising directions for increasing the weights on partitions around the efficient auction (where all weights are zero). The results in this Subsection depend on a condition on the distribution of values that requires symmetry among agents and independence across objects. Proposition 3 in Subsection 3.3 combines the above insights, and

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<sup>6</sup>These authors also translate the integrability constraint into a geometric condition about the boundaries of the regions in the type spaces where various allocations are chosen. The condition says that the vector representing the "jump" in the allocation must be perpendicular to the boundary where the jump occurs.

<sup>7</sup>Armstrong notes that when the type spaces increases in size, the revenue-maximizing auction should not be expected to remain efficient.

shows that a simple mixed bundling auction improves both upon the pure bundling auction and upon the efficient auction. This Subsection ends with a numerical example where we solve for the optimal level of mixed bundling. Proposition 4 in Subsection 3.4 shows that the revenue enhancing effect of mixed bundling is robust to the introduction of (optimal) reserve prices. Section 3.5 contains an example showing that the optimal weighted Vickrey-Clarke-Groves mechanism may be asymmetric even in symmetric situations. Section 4 concludes.

## 2 The Model

A seller  $s$  owns a set  $\mathcal{K}$  of  $K$  heterogeneous objects. There is a set  $\mathcal{N} = \{1, \dots, i, j, \dots, N\}$  of potential buyers (whom we also call bidders, or agents). Each bidder  $i$  privately knows her valuation  $v_k^i$  for each object  $k$ . For the seller, we assume  $v_k^s = 0$ . We assume that agent  $i$ 's valuation for a subset  $P_i \subseteq \mathcal{K}$  of objects is given by<sup>8</sup>  $v_{P_i}^i := \sum_{k \in P_i} v_k^i$ , and that her utility function is quasi-linear in money. Thus, a bidder receiving bundle  $P_i$  for a payment of  $T$  has a utility of  $\sum_{k \in P_i} v_k^i - T$ .

We assume that  $v_k^i$  can take values on an interval  $[\underline{v}_k, \bar{v}^k] \subseteq \mathbb{R}_0^+$ , and that  $v = (v_k^i)_{i,k}$  is distributed according to a continuous, strictly positive density function  $f : \Omega \rightarrow \mathbb{R}_0^+$  where  $\Omega = \prod_{k \in \mathcal{K}} [\underline{v}_k, \bar{v}^k]^N$ . The  $K$ -dimensional type of agent  $i$  is denoted by  $v^i \in \prod_{k \in \mathcal{K}} [\underline{v}_k, \bar{v}^k]$ , and  $v^{-i} \in \prod_{k \in \mathcal{K}} [\underline{v}_k, \bar{v}^k]^{N-1}$  denotes the  $K(N-1)$ -dimensional vector of all but bidder  $i$ 's valuations.

The outcome of an auction is a labeled partition  $P = (P_s, P_1, \dots, P_N) \in \mathcal{P}$  of objects and a vector of transfers  $(t^i)_{i \in \{1, \dots, N\}} \in \mathbb{R}^N$ . The interpretation is that bidder  $i \in \{1, \dots, N\}$  receives bundle  $P_i$ , the seller keeps the bundle  $P_s$ , and bidder  $i$  makes a payment of  $t^i$  to the seller.  $\mathcal{P}_{bun} \subset \mathcal{P}$  denotes the subset of pure bundling allocations, i.e.  $P \in \mathcal{P}_{bun} \Rightarrow \exists i \in \mathcal{N}$ , such that  $P_i = \mathcal{K}$ .

We study deterministic, dominant strategy, direct revelation mechanisms  $(g, t) : \Omega \rightarrow \mathcal{P} \times \mathbb{R}^N$  that specify for each vector of valuations  $v$  an allocation  $g(v)$  and payments from each agent  $t^i(v)$ . By payoff equivalence (see e.g. Green and Laffont (1979), and Holmstrom (1979)), the payments in any dominant strategy mechanism are, up to a constant, pinned down by the allocation rule  $g$ .

Two such mechanisms immediately come to mind: separate Vickrey auctions (one for each object), and a Vickrey auction for the bundle of all the

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<sup>8</sup>This model immediately extends to any kind of complementarities or substitutabilities between objects. We focus on additive valuations in order to single out the strategic effect of bundling.

objects:

1) In the separate auctions every object goes to the bidder who values it most:  $g(v) = \arg \max_{P \in \mathcal{P}} \left\{ \sum_{j \in \mathcal{N}} v_{P_j}^j \right\}$ . This allocation rule is implemented in dominant strategies by the well-known second-price transfer rule which charges bidder  $i$  the highest rejected bid on each object she receives.

2) In the (pure) bundling auction the set of all objects goes to the bidder who values highest the entire lot:  $g(v) = \arg \max_{P \in \mathcal{P}_{bun}} \left\{ \sum_{j \in \mathcal{N}} v_{P_j}^j \right\}$ . Charging the successful bidder the highest rejected bid on the entire lot makes this mechanism strategy-proof.

## 2.1 General $\lambda$ -auctions

We now define a large class of auctions that encompass many well-known formats, including the two mentioned above:

**Definition 1** For a given real vector  $\lambda = \{\lambda_P\}_{P \in \mathcal{P}}$  let

$$g_\lambda(v) := \arg \max_{P \in \mathcal{P}} \left\{ \sum_{j \in \mathcal{N}} v_{P_j}^j + \lambda_P \right\} \quad (1)$$

and

$$g_\lambda^{-i}(v) := \arg \max_{P \in \mathcal{P}} \left\{ \sum_{j \neq i} v_{P_j}^j + \underline{v}_{P^i}^i + \lambda_P \right\}$$

The  $\lambda$ -**auction**  $(g_\lambda, t_\lambda)$  is defined by the allocation rule<sup>9</sup>  $g = g_\lambda(v)$ , and by the transfer rule

$$t_\lambda^i(v) = - \left( \sum_{j \neq i} v_{g_\lambda(v)_j}^j + \lambda_{g_\lambda(v)} \right) + \tau_\lambda^i(v) \quad (2)$$

where  $\tau_\lambda^i(v) = \left( \sum_{j \neq i} v_{g_\lambda^{-i}(v)_j}^j \right) + \underline{v}_{g_\lambda^{-i}(v)_i}^i + \lambda_{g_\lambda^{-i}(v)}$ .

To interpret the above definition, we suggest a conceptual "trick": Consider an additional, fictitious agent who does not hold private information and cannot receive objects, but who attaches a value of  $\lambda_P$  to allocation

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<sup>9</sup>To circumvent the problem of tie-breaks in cases where argmax sets are not singletons, we assume throughout the paper that ties are broken according to a fixed order on  $\mathcal{P}$ . The tie-breaking rule does not matter given our assumption of an atomless distribution.

$P$ . The  $\lambda$ -auction is then the Vickrey-Clarke-Groves mechanism for the enlarged set of agents. The precise value of  $\tau_\lambda^i(v^{-i})$  - which is irrelevant for the dominance argument - is determined so that  $\underline{v}^i$ , the lowest type of bidder  $i$ , always receives zero utility in the truth-telling equilibrium of the  $\lambda$ -auction, in analogy to the Clarke mechanism. Note also that a bidder who receives no objects makes no payments<sup>10</sup>.

Together with payoff-equivalence, the above observations immediately yield:

**Lemma 1** *1) Truth-telling is a dominant strategy in an  $\lambda$ -auction. 2) Ex-post participation constraints are satisfied, and the lowest type of a bidder has a utility of zero in the truth-telling equilibrium. 3) Given the allocation rule  $g_\lambda$ , the transfer rule  $t_\lambda$  is the unique one with the above properties. Thus,  $t_\lambda$  is the highest transfer rule implementing  $g_\lambda$ .*

We now show how several well-known auction formats fit into the class of  $\lambda$ -auctions.

**Example 1** • *Separate second-price auctions for the  $K$  objects with reserve price  $r_k$  for object  $k$  are represented by the  $\lambda$ -auction where  $\lambda_P = \sum_{k \in P_s} r_k$ , i.e. the weight on an allocation is the sum of the reservation prices over all the objects that are not sold.*

- *The pure bundling auction is represented by the  $\lambda$ -auction where  $\lambda_P = \sum_k (\bar{v}^k - \underline{v}_k)$  for pure bundling allocations  $P \subset \mathcal{P}_{bun}$ , and  $\lambda_P = 0$  else. This ensures that  $\arg \max_{P \in \mathcal{P}} \left\{ \sum_{j \in \mathcal{N}} v_{P_j}^j + \lambda_P \right\} \in \mathcal{P}_{bun}$  for all possible realizations of  $v$ .*
- *More generally, consider any partition  $Q = \{Q_1, \dots, Q_\tau\}$  of  $\mathcal{K}$ . A pre-packaged auction with partition  $Q$  is represented by the  $\lambda$ -auction where  $\lambda_P = \sum_k (\bar{v}^k - \underline{v}_k)$  if the allocation  $P$  is compatible with the partition  $Q$  (i.e. each sub-bundle  $Q_\mu$  is allocated to one bidder<sup>11</sup>), and  $\lambda_P = 0$  else.*
- *Pure bundling with reserve price  $r$  on the bundle is represented by the  $\lambda$ -auction where  $\lambda_P = \sum_k (\bar{v}^k - \underline{v}_k)$  for  $P \in \mathcal{P}_{bun}$ ,  $\lambda_P = \sum_k (\bar{v}^k - \underline{v}_k) + r$  for  $P$  with  $P_s = \mathcal{K}$ , and  $\lambda_P = 0$  else.*

<sup>10</sup>This is because in this case  $g_\lambda^{-i}(v) = g_\lambda(v)$  and  $\underline{v}_{g_\lambda^{-i}(v)_i}^i = 0$ , as  $i$  receives no object.

<sup>11</sup>This means:  $\forall \mu \in \{1, \dots, \tau\}, \exists i \in \mathcal{N} : Q_\mu \subseteq P_i$ .

Weighted VCG mechanisms have been used by Roberts (1979) in order to characterize all dominant-strategy (not necessarily efficient) implementable choice rules in a general social choice setting. Note that, however large, the class of choice rules implemented by  $\lambda$ -auctions is, it is a strict subclass of dominant-strategy implementable choice rules<sup>12</sup>. Contrary to a conjecture in an earlier version of this paper, it is not even true that all anonymous allocation mechanisms are  $\lambda$ -auctions<sup>13</sup>. We focus on  $\lambda$ -auctions since this class is analytically tractable, and since, as illustrated above, it includes well known auction formats and their "convex combinations".

### 3 Optimal Bundling

In this Section we study how expected revenue depends on the choice of the  $\lambda$  parameters. The main working horse for the subsequent results is Theorem 1 which computes partial derivatives of revenue with respect to  $\lambda$ . We then apply this result in order to show how to improve the revenue from the pure bundling auction by making it slightly less inefficient, and how to improve upon the revenue of efficient auctions by rendering them slightly less efficient. After looking at these extreme cases, we focus on mixed-bundling auctions, a one-dimensional subclass of  $\lambda$ -auctions, which generalizes the idea of a discount on a bundle of objects in monopolistic mixed bundling.

Let us first introduce some notation. Given fixed, truthfully reported

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<sup>12</sup>In contrast to Roberts (1979), there is no simple characterization of all dominant strategy implementable choice rules in our framework. For example allocating goods only to bidder  $i$ , according to an arbitrary function of bidder  $j$ 's signal is implementable. The reason is, basically, the absence of allocative externalities. If bidders' preferences were allowed to depend on the entire profile of allocations (as opposed to their allocation only), Roberts' result would apply and the set of allocation rules implementable in dominant strategy would be the set of affine maximizers (i.e., the same as  $\lambda$ -auctions except that the weight given to the various bidders may be different). See Bickchandani et al. (2004), Gui et al. (2004) and Lavi et al. (2004) for recent research on the study of dominant strategy implementation in externality-free environments.

<sup>13</sup>Consider the following 1-object allocation mechanism: bidder  $i$  receives the good for a payment of  $t^i = 2v^j$  if  $v^i > 2v^j$ ; bidder  $j$  receives the good for a payment of  $t^j = 2v^i$  if  $v^j > 2v^i$ ; the seller keeps the object in all other cases. This mechanism is easily seen to be strategy-proof but it is not a  $\lambda$ -auction. We are grateful to a referee that offered this counter-example.



valuations  $v$  in the  $\lambda$ -auction, denote by

$$\begin{aligned} R_\lambda(v) &= \sum_i t_\lambda^i(v), \\ S_\lambda(v) &= \sum_i v_{g_\lambda(v)_i}^i, \\ BS_\lambda(v) &= S_\lambda(v) - R_\lambda(v) = \sum_i \left( v_{g_\lambda(v)_i}^i - t_\lambda^i(v) \right) \end{aligned}$$

realized revenue, social surplus and bidder surplus, respectively. Also, denote by  $R(\lambda) = \mathbb{E}_v [R_\lambda(v)]$ ,  $S(\lambda) = \mathbb{E}_v [S_\lambda(v)]$ , and  $BS(\lambda) = \mathbb{E}_v [BS_\lambda(v)]$  expected revenue, expected social surplus, and expected bidder surplus, respectively.

**Theorem 1** 1. Fix  $v$  and  $P \in \mathcal{P}$ . Let  $\mathbb{I}$  be the indicator function, and let  $P'' = \arg \max_{P' \neq P} \left\{ \sum_{j \in \mathcal{N}} v_{P'_j}^j + \lambda_{P'} \right\}$ , function. The revenue  $R_\lambda(v)$  is a piecewise linear function of  $\lambda_P$ , and has a single discontinuity of  $(\sum_j v_{P_j}^j - \sum_j v_{P''_j}^j)$  at  $\lambda_P^* := \left( \sum_{j \in \mathcal{N}} v_{P''_j}^j \right) + \lambda_{P''} - \sum_{j \in \mathcal{N}} v_{P_j}^j$ . Whenever differentiable, its derivative is given by:

$$\frac{\partial}{\partial \lambda_P} R_\lambda(v) = \sum_i \left( -\mathbb{I}_{\{(\lambda, v): g_\lambda(v) = P\}} + \mathbb{I}_{\{(\lambda, v): g_\lambda^{-i}(v) = P\}} \right) \quad (3)$$

2. Expected revenue is a differentiable function of  $\lambda_P$  with partial derivative given by:

$$\frac{\partial R}{\partial \lambda_P}(\lambda) = \frac{\partial S}{\partial \lambda_P}(\lambda) + \sum_{i \in \mathcal{N}} \left( -\Pr(g_\lambda(v) = P) + \Pr(g_\lambda^{-i}(v) = P) \right) \quad (4)$$

**Proof.** 1) The revenue effect of changing  $\lambda$  can be decomposed into a surplus generation effect, and a surplus extraction effect :

$$R_\lambda(v) = S_\lambda(v) - BS_\lambda(v).$$

Social surplus depends on  $\lambda$  only via the chosen allocation  $g_\lambda(v)$ . Therefore  $S_\lambda(v)$  is a step function of  $\lambda_P$  with a discontinuity of  $(\sum_{j \in \mathcal{N}} v_{P_j}^j - \sum_{j \in \mathcal{N}} v_{P''_j}^j)$  at the point  $\lambda_P^*$  where the allocation switches from  $P''$  to  $P$ , i.e, at

$$\lambda_P^* := \sum_{j \in \mathcal{N}} v_{P''_j}^j + \lambda_{P''} - \sum_{j \in \mathcal{N}} v_{P_j}^j.$$

Thus, its derivative is zero whenever it is well-defined.

Bidder  $i$ 's surplus is given by:

$$\begin{aligned} v_{g_\lambda(v)_i}^i - t_\lambda^i(v) &= \sum v_{g_\lambda(v)_j}^j + \lambda_{g_\lambda(v)} - \tau_\lambda^i(v) \\ &= \sum v_{g_\lambda(v)_j}^j + \lambda_{g_\lambda(v)} - \sum_{j \neq i} v_{g_\lambda^{-i}(v)_j}^j - \frac{v_{g_\lambda^{-i}(v)_i}^i}{\lambda_{g_\lambda^{-i}(v)}} - \lambda_{g_\lambda^{-i}(v)}. \end{aligned}$$

The bidders' surplus is given by:

$$BS_\lambda(v) = \sum_i \left[ \max_{P \in \mathcal{P}} \left\{ \sum_{j \in \mathcal{N}} v_{P_j}^j + \lambda_P \right\} - \max_P \left\{ \sum_{j \neq i} v_{P_j}^j + \frac{v_{P_i}^i}{\lambda_P} + \lambda_P \right\} \right]$$

Thus,  $BS_\lambda(v)$  is the difference of maxima of continuous functions. As such, it is continuous and piecewise linear in  $\lambda_P$ . Its derivative is given by:

$$\frac{\partial}{\partial \lambda_P} BS_\lambda(v) = \sum_i \left( \mathbb{I}_{\{(\lambda, v): g_\lambda(v) = P\}} - \mathbb{I}_{\{(\lambda, v): g_\lambda^{-i}(v) = P\}} \right)$$

The claim follows by combining the two insights above.

2) Again, we compute the revenue effect as the difference of the effects on total surplus, and on bidder surplus. The latter is computed by integration of equation 3 with respect to  $v$ :

$$\begin{aligned} \frac{\partial}{\partial \lambda_P} BS(\lambda) &= \frac{\partial}{\partial \lambda_P} \mathbb{E}_v [BS_\lambda(v)] \\ &= \mathbb{E}_v \left[ \frac{\partial}{\partial \lambda_P} BS_\lambda(v) \right] \\ &= \mathbb{E}_v \left[ \sum_i \mathbb{I}_{\{(\lambda, v): g_\lambda(v) = P\}} - \mathbb{I}_{\{(\lambda, v): g_\lambda^{-i}(v) = P\}} \right] \\ &= \sum_i (\Pr(g_\lambda(v) = P) - \Pr(g_\lambda^{-i}(v) = P)) \end{aligned}$$

For the second equation, we used Lebesgue's theorem and the Lipschitz-continuity of  $BS_\lambda(v)$  in  $\lambda$  in order to exchange the order of differentiation and integration.

It remains to show that, unlike realized social surplus, expected social surplus is differentiable as a function of  $\lambda_P$ , i.e. that  $\frac{\partial S}{\partial \lambda_P}(\lambda)$  exists. To do so, denote by

$$\partial_\lambda(P, P') := \left\{ v \in \Omega : P, P' \in \arg \max_{P \in \mathcal{P}} \left( \sum_{j \in \mathcal{N}} v_{P_j}^j + \lambda_P \right) \right\}$$

the boundary in the type space between areas where  $g_\lambda$  chooses allocations  $P$  and  $P'$ , respectively, and by  $d^{NK-1}v$  the  $NK - 1$  dimensional Lebesgue measure. With this notation, we compute  $\frac{\partial S}{\partial \lambda_P}(\lambda)$  by applying Leibniz's rule to integrate the discontinuity of realized social surplus  $S_\lambda(v)$  over all valuations at which the discontinuity occurs:

$$\frac{\partial S}{\partial \lambda_P}(\lambda) = \sum_{P' \in \mathcal{P}} \int_{\partial_\lambda(P, P')} \left( \sum_{j \in \mathcal{N}} v_{P_j}^j - \sum_{j \in \mathcal{N}} v_{P'_j}^j \right) f(v) d^{NK-1}v$$

Given the continuity of  $f$ , this expression is well defined for all  $\lambda$ . ■

Recall the interpretation of the  $\lambda$ -auction as the efficient auction for society including a virtual agent whose preferences are determined by  $\lambda$ . Changing these preferences affects the externality that agent  $i$  imposes on the virtual agent, and this translates into a change of agent  $i$ 's payment. For a fixed allocation  $g_\lambda(v)$ , the marginal externality is given by the term under the summation in equation (3). When the allocation  $g_\lambda(v)$  is affected by the change in  $\lambda$ , the aggregated externality on the virtual agent that results from changing the allocation from  $P''$  to  $P$  equals the discontinuity in aggregated payments of the real agents:

$$\lambda_{P''} - \lambda_P^* := \sum_j v_{P_j}^j - \sum_j v_{P''_j}^j.$$

For an illustration, consider a single-good auctioneer setting a reserve price, and recall that a reserve price corresponds to a weight on the allocation where the auctioneer keeps the object. To determine the optimal reserve price, the auctioneer is trading off the marginal loss of efficiency due to missing the opportunity to sell, against the marginal gain of additional revenue. For any realized  $v$ , the efficiency loss is captured by the discontinuity of  $R_\lambda(v)$ . The marginal revenue from agent  $i$  is captured by the term under the summation in equation (3): it equals 1 for realizations of  $v$  in which only the valuation of agent  $i$  exceeds the reserve price. In expectation, the efficiency loss is captured by the first term of equation (4), while the expected marginal revenue from  $i$  equals agent  $i$ 's expected externality on the virtual agent, represented by the term under summation in equation (4).

### 3.1 Improving upon pure bundling and pre-packaged auctions

We now apply Theorem 1 to the pure bundling auction, and we show how to increase revenue by allowing separate sales for extreme realizations. Recall

that the pure bundling auction is represented by the  $\lambda$ -auction where  $\lambda_P = \sum_k (\bar{v}^k - \underline{v}_k)$  for  $P \subset \mathcal{P}_{bun}$ , and  $\lambda_P = 0$  else. Decreasing the parameter  $\lambda_P$  for  $P \in \mathcal{P}_{bun}$  will allow allocations other than pure bundling to be chosen.

**Definition 2** *The  $c$ -bundling auction  $(g_c, t_c)$  is the  $\lambda$ -auction with:*

$$\lambda_P := \begin{cases} c & \text{if } P \in \mathcal{P}_{bun} \\ 0 & \text{else.} \end{cases}$$

We denote by  $R_c(v)$  the revenue of a  $c$ -bundling auction for a given realization of  $v$ , and by  $R(c)$  the expected revenue. Let  $c^* := \inf \{c : g_c(\Omega) \subseteq \mathcal{P}_{bun}\}$  be the smallest value of  $c$  for which the  $c$ -bundling auction is the pure bundling auction<sup>14</sup>.

**Proposition 1** 1. *For sufficiently small  $\varepsilon > 0$ , and for any realization of  $v$ , the revenue of the  $(c^* - \varepsilon)$ -bundling auction is higher than the revenue of the pure bundling auction:  $R_{c^* - \varepsilon}(v) \geq R_{c^*}(v)$ , with strict inequality for  $v$  such that  $g_{c^* - \varepsilon}(v) \notin \mathcal{P}_{bun}$ . In particular,  $R(c^* - \varepsilon) > R(c^*)$  for generic distribution functions.*

2. *The revenue of a pre-packaged auction with partition  $\mathcal{Q}$  is increased by uniformly lowering the parameter  $\lambda_P$  on all allocations  $P$  compatible with  $\mathcal{Q}$  below the corresponding threshold value  $c_{\mathcal{Q}}^*$ .*

**Proof.** 1) Consider first  $v$  with  $g_{c^* - \varepsilon}(v) \in \mathcal{P}_{bun}$ . By part 1 of Theorem 1 we get that

$$\begin{aligned} R_{c^*}(v) - R_{c^* - \varepsilon}(v) &= \int_{c^* - \varepsilon}^{c^*} \frac{\partial}{\partial c} R_c(v) dc \\ &= \int_{c^* - \varepsilon}^{c^*} \sum_{P \in \mathcal{P}_{bun}} \sum_i \left( -\mathbb{I}_{\{(c,v):g_c(v)=P\}} + \mathbb{I}_{\{(\lambda,v):g_c^{-i}(v)=P\}} \right) dc \\ &= \sum_i \int_{c^* - \varepsilon}^{c^*} \left( -\mathbb{I}_{\{(c,v):g_c(v) \in \mathcal{P}_{bun}\}} + \mathbb{I}_{\{(\lambda,v):g_c^{-i}(v) \in \mathcal{P}_{bun}\}} \right) dc \end{aligned}$$

As  $g_c(v) \in \mathcal{P}_{bun}$  for  $c \in [c^* - \varepsilon; c^*]$ , the integrand is non-positive and the result follows.

Consider now  $v$  with  $g_{c^* - \varepsilon}(v) \notin \mathcal{P}_{bun}$ . Denote by

$$\tilde{c}(v) := \inf \{c \in [c^* - \varepsilon; c^*] : g_c(v) \in \mathcal{P}_{bun}\}$$

---

<sup>14</sup>We assume that ties in the  $c^*$ -bundling auction are broken in favor of pure bundling allocations.

the lowest value of  $c$  for which a pure bundling allocation is chosen at realization  $v$ . By Theorem 1 the discontinuity of  $R(c)$  at  $c = \tilde{c}(v)$  equals exactly  $-\tilde{c}(v)$ . We get:

$$\begin{aligned} & R_{c^*}(v) - R_{c^*-\varepsilon}(v) \\ &= \sum_i \int_{c^*-\varepsilon}^{c^*} \left( -\mathbb{I}_{\{(c,v):g_c(v) \in \mathcal{P}_{bun}\}} + \mathbb{I}_{\{(\lambda,v):g_c^{-i}(v) \in \mathcal{P}_{bun}\}} \right) dc - \tilde{c}(v) \end{aligned}$$

As the first term is bounded above by  $N\varepsilon$ , this expression is strictly negative for small enough  $\varepsilon$ .

2) The pre-packaged auction with partition  $\mathcal{Q}$  can be interpreted as a set of separate pure-bundling auctions for each of the packages  $Q_\mu$ . Thus, the result follows by applying part 1 to each of these separate pure-bundling auctions. ■

### 3.2 Improving upon the efficient auction

By definition, the social surplus  $S(\lambda)$  is maximized at the efficient auction where  $\lambda = 0$ . Therefore,  $\nabla S(0) = 0$ . Together with Theorem 1, this observation yields:

#### Corollary 2

$$\frac{\partial R}{\partial \lambda_P}(0) = \sum_{i \in \mathcal{N}} \left( -\Pr(g_0(v) = P) + \Pr(g_0^{-i}(v) = P) \right)$$

Thus, a necessary condition for the efficient auction to maximize revenue in the class of  $(\lambda_P)_{P \in \mathcal{P}}$ -auctions is :

$$\forall P \in \mathcal{P}, \sum_{i \in \mathcal{N}} \left( -\Pr(g_0(v) = P) + \Pr(g_0^{-i}(v) = P) \right) = 0 \quad (5)$$

This is a highly non-generic condition on the density function governing the distribution of the signals  $v$ . In other words, by introducing a small  $\lambda_P$  for some  $P$  one can almost always improve upon the efficient auction<sup>15</sup>.

We now use Corollary 2 in order to identify the most promising directions in which to improve on the efficient auction. The results in the rest of this subsection use the following condition:

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<sup>15</sup> The FOC need not be sufficient for a global maximum, as  $R(\lambda)$  need not be a concave function of  $\lambda$ . Subsection 3.5 contains an example showing that the best  $\lambda$ -auction may be asymmetric even in symmetric situations, thereby illustrating the potential non-concavity of  $R(\lambda)$ .

**Condition 3 (SAIO)** Valuations  $v = (v_k^i)_{i,k}$  are distributed symmetrically across agents and independently across objects. Formally, the density function  $f$  satisfies: 1)  $f(\dots, v^i, \dots, v^j, \dots) = f(\dots, v^j, \dots, v^i, \dots)$ , and 2)  $f(v) = \prod_{k \in \mathcal{K}} f_k(v_k)$ , where  $f_k : [\underline{v}_k, \bar{v}_k]^N \rightarrow R_0^+$  is the marginal density with respect to object  $k$ .

**Proposition 2** Assume that condition SAIO is met, and denote by  $\#P := \#\{i \in \mathcal{N} : P_i \neq \emptyset\}$ , the number of bidders receiving some objects in allocation  $P$ . We have

$$\begin{aligned} \frac{\partial R}{\partial \lambda_P}(0) &= \frac{N^{K-1}(N - \#P) - (N - 1)^K}{N^{K-1}(N - 1)^K}, \text{ for } P \text{ with } P_s = \emptyset \quad (6) \\ \frac{\partial R}{\partial \lambda_P}(0) &= 0, \text{ for } P \text{ with } P_s \neq \emptyset \end{aligned}$$

In particular, introducing a bundling parameter  $\lambda_P = c$  for all  $P \in \mathcal{P}_{\text{bun}}$  leads to an increase in expected revenue versus the expected revenue in the efficient auction.

**Proof.** Consider first  $P$  with  $P_s = \emptyset$ . As the density function  $f$  is assumed to be symmetrical across bidders, the probability  $\Pr(k \in g_0(v)_i)$  of bidder  $i$  having the highest valuation for some object  $k$  equals  $\frac{1}{N}$ . As valuations are assumed to be independent across objects, the probability of an allocation  $P$  being efficient is obtained by multiplying this constant probability over all objects:  $\Pr(g_0(v) = P) = \frac{1}{N^K}$ . To determine  $\Pr(g_0^{-i}(v) = P)$  we apply the same logic, and we obtain  $\Pr(g_0^{-i}(v) = P) = \frac{1}{(N-1)^K}$  if  $P_i = \emptyset$ , and  $\Pr(g_0^{-i}(v) = P) = 0$  else.

Using Corollary 2, we obtain:

$$\begin{aligned} \frac{\partial R}{\partial \lambda_P}(0) &= \sum_{i \in \mathcal{N}} [-\Pr(g_0(v) = P) + \Pr(g_0^{-i}(v) = P)] \\ &= -\frac{N}{N^K} + \frac{N - \#P}{(N - 1)^K} = \frac{N^{K-1}(N - \#P) - (N - 1)^K}{N^{K-1}(N - 1)^K} \end{aligned}$$

For  $P$  with  $P_s \neq \emptyset$  the result follows from noticing that  $\Pr(g_0(v) = P) = \Pr(g_0^{-i}(v) = P) = 0$ . Introducing a small reserve price has an expected marginal effect of zero as both the size of the effect, and the probability of the realization of the case are marginal<sup>16</sup>.

<sup>16</sup>It is worth bearing in mind, that the partial derivatives at the efficient auction need not be a good indicator for the impact of significantly increasing  $\lambda_P$ . It is well known that

The second part of the Proposition follows by observing that  $\#P = 1$  for  $P \in \mathcal{P}_{\text{bun}}$ . Thus, we obtain

$$\begin{aligned} \frac{\partial R}{\partial c}(0) &= \sum_{P \in \mathcal{P}_{\text{bun}}} \frac{\partial R}{\partial \lambda_P}(0) = N \frac{N^{K-1}(N-1) - (N-1)^K}{N^{K-1}(N-1)^K} \\ &= N(N-1) \left( \frac{N^{K-1} - (N-1)^{K-1}}{N^{K-1}(N-1)^K} \right) > 0. \end{aligned}$$

■

**Remark 1** Note that  $\frac{\partial R}{\partial \lambda_P}(0)$  is decreasing in  $\#P$ , positive for  $\#P = 1$ , negative for  $\#P = \min\{N, K\}$  and positive for every fixed  $\#P \leq \min\{N, K\} - 1$  as either  $N$  or  $K$  tends towards infinity. Thus, in the space of parameter combinations  $\lambda = (\lambda_P)_{P \in \mathcal{P}}$  increasing  $\lambda_P$  locally around the efficient auction has the highest positive impact for  $P \in \mathcal{P}_{\text{bun}}$ . Raising  $\lambda_P$  for partitions  $P$  that allocate objects to all bidders,  $\#P = N$ , or that sell the objects as singletons,  $\#P = K$ , decreases expected revenue.

**Remark 2** To broaden our focus let us consider raising  $\lambda_P$  for all allocations  $P$  that are derived from a partition  $\mathcal{Q}: Q_1 \uplus \dots \uplus Q_\tau$ , i.e. for all  $P$  such that  $\forall \mu \in \{1, \dots, \tau\}, \exists i: Q_\mu = P_i$ . By definition, precisely  $\tau$  agents receive objects in each of these allocations. To compute the number of such  $P$ , observe that there are  $N$  possibilities to allocate bundle  $Q_1$ ,  $N-1$  possibilities to allocate bundle  $Q_2$ , and so on. Thus, we get

$$\sum_{P \text{ derived from } \mathcal{Q}} \frac{\partial R}{\partial \lambda_P}(0) = \frac{N!}{(N-\tau)!} \frac{N^{K-1}(N-\tau) - (N-1)^K}{N^{K-1}(N-1)^K}$$

The revenue impact only depends on  $\tau$ , the number of sub-bundles in the partition, but not on whether  $\mathcal{Q}$  divides up the objects evenly or not. Similarly to equation 6, the obtained expression is positive for small  $\tau$  and becomes negative for  $\tau$  close to  $\min\{N, K\}$ . But now it is not strictly decreasing in  $\tau$ . Therefore the largest impact comes from raising  $\lambda_P$  for all  $P$  with an intermediate  $\tau$ . For example, in a setting with 5 bidders and 3 objects, we obtain that, optimally,  $\tau = 2$ . Thus, the partition  $\mathcal{Q}$  should contain two sub-bundles of one and two objects, respectively.

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the introduction of a "large" reserve price can have a large impact on expected revenue.

Note also that the optimal one-object auction for heterogeneous bidders as in Myerson (1981) need not be an  $\lambda$ -auction. This is because, in general, it sets different reserve prices for different agents, which a  $\lambda$ -auction can not do.

The next example shows that, without condition SAIO,  $\frac{\partial R}{\partial \lambda_P}(0)$  may be negative even for  $P$  with  $\#P = 1$ .

**Example 2** Let  $[v_k, \bar{v}^k] = [0, 1]$  and assume that  $v^i$  are i.i.d. but that there is correlation between objects. Specifically let  $v^i$  be distributed on  $[0, 1]^K$  in a way that puts probability  $\frac{1}{2}$  uniformly on  $[0, \frac{1}{2}]^K$  and probability  $\frac{1}{2}$  on the "upper diagonal"  $\left\{ (x, \dots, x) \in [0, 1]^K : x \geq \frac{1}{2} \right\}$ . Here  $\frac{\partial R}{\partial \lambda_P}(0) < 0$  for  $P \in \mathcal{P}_{bun}$ .

**Proof.** Let  $P \in \mathcal{P}_{bun}$ . Due to symmetry between agents we have that

$$\frac{\partial R}{\partial \lambda_P}(0) = \frac{1}{N} \sum_{i \in \mathcal{N}} (\Pr(g_0^{-i}(v) \in \mathcal{P}_{bun}) - \Pr(g_0(v) \in \mathcal{P}_{bun}))$$

Let us first evaluate  $\Pr(g_0(v) \in \mathcal{P}_{bun})$ : With probability  $1 - \left(\frac{1}{2}\right)^N$  there is at least one agent  $i$  with signal  $v^i$  on the upper diagonal, in which case there is a bidder who has the highest valuation for all objects. With probability  $\left(\frac{1}{2}\right)^N$  this is not the case, and the probability of there being an overall winner is  $\left(\frac{1}{N}\right)^{K-1}$ . Thus

$$\Pr(g_0(v) \in \mathcal{P}_{bun}) = 1 - \left(\frac{1}{2}\right)^N + \left(\frac{1}{2}\right)^N \left(\frac{1}{N}\right)^{K-1} \quad (7)$$

Analogously, we get that

$$\Pr(g_0^{-i}(v) \in \mathcal{P}_{bun}) = 1 - \left(\frac{1}{2}\right)^{N-1} + \left(\frac{1}{2}\right)^{N-1} \left(\frac{1}{N-1}\right)^{K-1}$$

As  $\left[1 - \left(\frac{1}{2}\right)^N + \left(\frac{1}{2}\right)^N \left(\frac{1}{N}\right)^{K-1}\right]$  is increasing in  $N$ , we obtain that

$$\frac{\partial R}{\partial \lambda_P}(0) = \frac{1}{N} \sum_{i \in \mathcal{N}} (\Pr(g_0^{-i}(v) \in \mathcal{P}_{bun}) - \Pr(g_0(v) \in \mathcal{P}_{bun})) < 0$$

■

### 3.3 Mixed Bundling Auctions

So far we have made use of local arguments to show that neither pure bundling nor separate auctions can be optimal (under SAIO). In this subsection, we focus on the mixed bundling auctions as introduced in Definition 2. This class is parametrized by a parameter  $c$  that acts as a surcharge on



the payment of a bidder who receives some, but not all, of the objects (or, equivalently, as a discount on the payment of a bidder who receives the entire lot). By definition, this class encompasses the efficient auction, represented by the  $c$ -bundling auction with  $c = 0$ , and the pure bundling auction, represented by the  $c$ -bundling auction with  $c = c^* = \inf \{c : g_c(\Omega) \subseteq \mathcal{P}_{bun}\}$ . As a corollary to Propositions 1 and 2, we get:

**Proposition 3** *Assume that condition SAIO is met. Then there exists  $c \in (0, c^*)$  such that the  $c$ -bundling auction achieves a higher expected revenue than both the pure bundling auction and the efficient auction.*

**Proof.** Expected revenue is a continuous function of  $c$ . Thus, it must have a maximum on the closed interval  $[0, c^*]$ . Proposition 1 showed the maximum cannot be at  $c = c^*$ , while Proposition 2 showed that the maximum cannot be at  $c = 0$ . ■

To understand the comparative statics for  $c$  in the interior of  $[0, c^*]$ , consider the case of two bidders  $i, j$ . Then the formula for payments reduces to:

$$t_c^i(v) = \begin{cases} v_{\mathcal{K}}^j & \text{if } P^i = \mathcal{K}, \\ v_{P^i}^j + c & \text{if } \emptyset \neq P^i \subsetneq \mathcal{K}, \\ 0 & \text{if } P^i = \emptyset. \end{cases} \quad (8)$$

Increasing the bundling parameter by  $\varepsilon$  increases revenue by  $2\varepsilon$  for valuations such that the objects are sold separately. It decreases revenue discontinuously by  $c$  for valuations such that the  $\varepsilon$ -increase in  $c$  leads to a bundling allocation being chosen instead of the efficient one. Figure 1 shows the realized revenue  $R_c(v)$  for some fixed values of  $v$  as a function of  $c$ .

Proposition 2 stated that introducing mixed bundling to the efficient auction increases revenue (since the externality on the virtual bidder is added as a surcharge on bidders' payments). This is reflected in the positive slope of  $R_c(v)$  at  $c = 0$ . Proposition 1 stated that making the pure bundling auction slightly less inefficient by lowering  $c$  also increases revenue. This is reflected in the upward jump in  $R_c(v)$  when  $c$  moves from right to left.

We now analytically compute the optimal bundling parameter in a simple example.

**Example 3** *There are two bidders  $i, j$  competing for two objects  $A, B$ . Valuations are uniformly and independently distributed on  $[0, 1]$ . The revenue-maximizing bundling parameter is  $c = \frac{1}{3}$ .*

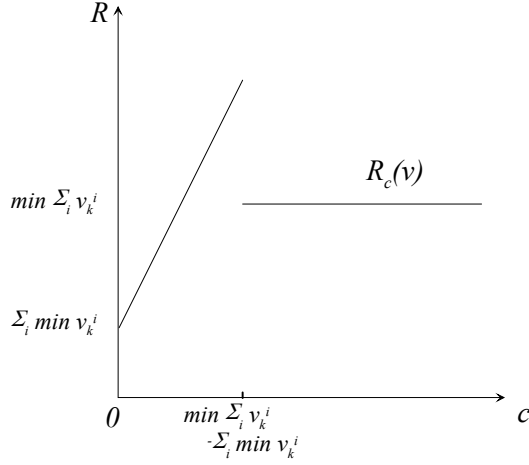


Figure 1: Realized Revenue  $R(v)$  as a function of  $c$  (fixing  $v$  such that bundling is inefficient).

**Proof.** By Theorem 1 we know that  $\frac{dR(c)}{dc} = \frac{dS(c)}{dc} + 2 \Pr(g_c(v) \notin \mathcal{P}_{bun})$ . As the marginal loss in expected social welfare is the inefficiency  $c$  times the marginal probability that this inefficiency is incurred,  $\frac{d}{dc} \Pr(g_c(v) \in \mathcal{P}_{bun})$ , we get that:

$$\frac{dR(c)}{dc} = -c \frac{d}{dc} \Pr(g_c(v) \in \mathcal{P}_{bun}) + 2 \Pr(g_c(v) \notin \mathcal{P}_{bun})$$

Note that  $\Pr(g_c(v) \notin \mathcal{P}_{bun}) = 2 \Pr(g_c(v) = [\emptyset, \{A\}, \{B\}])$ . Only for  $\{v_A^j < 1 - c, v_B^j > c\}$  there are  $v^i$  such that  $g_c(v) = [\emptyset, \{A\}, \{B\}]$ . The area of this region of  $i$ 's valuations is given by  $(1 - (v_A^j + c))(v_B^j - c)$ , as is illustrated in figure 2. This yields:

$$\Pr(g_c(v) \notin \mathcal{P}_{bun}) = 2 \int_c^1 \int_0^{1-c} (1 - (v_A^j + c))(v_B^j - c) dv_A^j dv_B^j = \frac{1}{2} (1 - c)^4$$

Thus, we get:

$$\frac{dR(c)}{dc} = (1 - c)^4 - 2c(1 - c)^3$$

Note that  $\frac{dR(c)}{dc} = 0$  for  $c = \frac{1}{3}$ . To calculate expected revenue, we use the fact that  $R(0) = \frac{2}{3}$  (twice the expectation of the second-order statistic of two random variables distributed uniformly on  $[0, 1]$ ), and we get

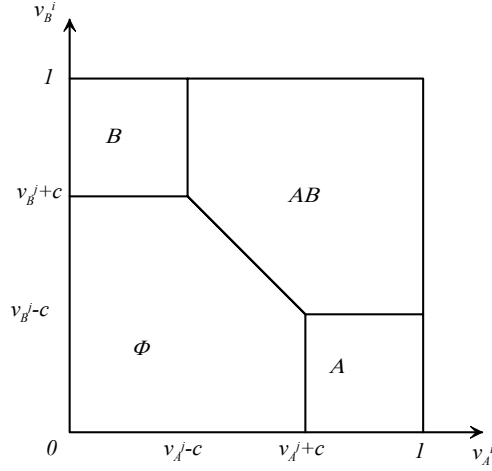


Figure 2:  $i$ 's allocation ( $\emptyset$ ,  $A$ ,  $B$  or  $AB$ ) as a function of  $v^i$  on a slice through the 4-dimensional type space for fixed  $v_A^j, v_B^j$ .

$$R(c) = \frac{2}{3} + c - 3c^2 + 4c^3 - \frac{5}{2}c^4 + \frac{3}{5}c^5.$$

Thus, we obtain  $R(0) = \frac{2}{3}$ ,  $R(c^*) = R(1) = \frac{23}{30} \approx 0.766$ , and  $R(\frac{1}{3}) = \frac{637}{810} \approx 0.786$ . The improvement caused by mixed bundling seems small in this particular example. But, of course, it can get quite large if valuations are increased. A meaningful calibration can only be done within the context of specific applications. ■

### 3.4 Reserve prices

Reserve prices are known to increase revenue (see Myerson. 1981). Propositions 1 and 2 show that when the seller cannot retain the objects, neither separate nor pure bundling auctions can be optimal. Could it be that separate auctions with well adjusted reserve prices are optimal? The following result shows this is not so, at least when there are sufficiently many bidders (and condition SAIO is met).

We denote by  $(g_{c,r}, t_{c,r})$  the  $c$ -bundling auction with a vector of reserve

prices  $r = (r_1, \dots, r_K)$ , i.e, the  $\lambda$ -auction<sup>17</sup> with:

$$\lambda_P := \begin{cases} c & \text{if } P \in \mathcal{P}_{bun}, \\ \sum_{k \in P_s} r_k & \text{else.} \end{cases}$$

Denote the expected revenue (resp. surplus, resp. bidders surplus) of the above auction by  $R(c, r)$  (resp.  $S(c, r)$ ,  $BS(c, r)$ ).

**Proposition 4** *Assume that condition SAIO is met. For sufficiently large  $N$ , it holds that  $\frac{\partial R}{\partial c}(0, r) > 0$ .*

**Proof.** Let  $F_k: [\underline{v}_k, \bar{v}^k] \rightarrow [0; 1]$  denote the cumulative distribution function for each  $v_k^i$ . We decompose the effect on  $R$  into the effect on bidder surplus

$$\frac{\partial BS}{\partial c}(0, r) = \sum_{i \in \mathcal{N}} (-\Pr(g_{0,r}(v) \in \mathcal{P}_{bun}) + \Pr(g_{0,r}^{-i}(v) \in \mathcal{P}_{bun}))$$

and into the effect on social surplus  $\frac{\partial S}{\partial c}(0, r)$ .

Consider first the effect on social surplus. By the formula from the proof of Theorem 1 we obtain:

$$\frac{\partial S}{\partial c}(0, r) = \sum_{P \in \mathcal{P}_{bun}} \sum_{P' \in \mathcal{P}} \int_{\partial_\lambda(P, P')} \left( \sum_{j \in \mathcal{N}} v_{P_j}^j - \sum_{j \in \mathcal{N}} v_{P'_j}^j \right) f(v) d^{NK-1}v$$

Since the reserve prices acts as positive weights on allocations where the seller keeps some object, we obtain for  $c = 0$  that  $\lambda_P \leq \lambda_{P'}$  for  $P \in \mathcal{P}_{bun}$  and  $P' \in \mathcal{P}$ . By definition,

$$\partial_\lambda(P, P') := \left\{ v \in \Omega : P, P' \in \arg \max_{P \in \mathcal{P}} \left( \sum_{j \in \mathcal{N}} v_{P_j}^j + \lambda_P \right) \right\}$$

Thus, we obtain that  $\sum_{j \in \mathcal{N}} v_{P_j}^j - \sum_{j \in \mathcal{N}} v_{P'_j}^j$  is non-negative on  $\partial_\lambda(P, P')$ , and therefore that  $\frac{\partial S}{\partial c}(0, r) \geq 0$ .

---

<sup>17</sup>Note that  $c$  acts as a discount on the reserve prices if a bidder purchases the full bundle. This is somewhat arbitrary, as we could also define a  $(c, r)$ -auction by setting  $\lambda_P = \sum_{k \in P_s} r_k + c$  for  $P = (K, \emptyset, \dots, \emptyset)$ .

Consider now the effect on bidders' surplus. Denoting by  $F_k$  the one-dimensional distribution function of  $v_k^i$ , we get that:

$$\begin{aligned} \Pr(g_{0,r}(v) \in \mathcal{P}_{bun}) &= \sum_i \prod_{k \in \mathcal{K}} \Pr(v_k^{(1)} > r_k) \Pr(v_k^i > v_k^j, \forall j \neq i | v_k^{(1)} > r_k) \\ &= N \prod_{k \in \mathcal{K}} (1 - F_k(r_k)^N) \Pr(v_k^i > v_k^j, \forall j \neq i) \\ &= \frac{1}{N^{K-1}} \prod_{k \in \mathcal{K}} (1 - F_k(r_k)^N) \end{aligned}$$

and similarly:

$$\Pr(g_{0,r}^{-i}(v) \in \mathcal{P}_{bun}) = \frac{1}{(N-1)^{K-1}} \prod_{k \in \mathcal{K}} (1 - F_k(r_k)^{N-1})$$

Note that the function  $g(N) := \frac{1}{N^{K-1}} \prod_{k \in \mathcal{K}} (1 - F_k(r_k)^N)$  strictly increases in  $N$  for large  $N$ . Thus, the effect on bidders' surplus is strictly positive, and the claim follows. ■

If the distributions of valuations are independent across agents, we know from Myerson (1981) that, under a regularity condition, the optimal reserve price in a one-object auction is independent of the number of bidders. Under these assumptions, Proposition 4 shows that, when the number of bidders is large enough, one can always improve upon the separate auctions with optimally set reserve prices by introducing a mixed bundling parameter  $c > 0$ . The Proposition can also be extended to show that no pure bundling auction with reserve price  $r$  on the pure bundle can be optimal.

### 3.5 Asymmetric $\lambda$ -auctions

An interesting question is whether in a symmetric setting the optimal  $\lambda$ -auction is symmetric in the sense that it treats the bidders anonymously. The following example<sup>18</sup> illustrates that this need not be the case.

**Example 4** *There are two bidders competing for one object. Valuations are independently distributed across bidders, and take two values<sup>19</sup>:  $\bar{v}$  with probability  $p$  and  $\underline{v}$  with probability  $1 - p$ , where  $\bar{v} > \underline{v}$ . The only parameter in a symmetric single-object  $\lambda$ -auction is the reserve price: setting a low*

<sup>18</sup>The same example may be used to illustrate that adding asymmetric bidder-specific weights (see Roberts, 1979) may also increase revenue.

<sup>19</sup>The same conclusions will hold for distributions with continuous density approaching the discrete one considered here.

reserve price yields expected revenue of  $p^2\bar{v} + (1 - p^2)\underline{v}$ , while setting a high reserve price yields  $p(2 - p)\bar{v}$ . We now construct an asymmetric  $\lambda$ -auction that outperforms both of these, and thus all, symmetric  $\lambda$ -auctions.

Let  $\lambda_i^\varepsilon = \varepsilon$  and  $\lambda_j^\varepsilon = \bar{v} - \underline{v}$  where  $i$  ( $j$ ) refers here to the allocation where  $i$  ( $j$ ) receives the object. For small values of  $\varepsilon$ , the allocation and the payments in the  $\lambda^\varepsilon$ -auction are given by the following table:

	$v^j = \underline{v}$	$v^j = \bar{v}$
$v^i = \bar{v}$	$i$ wins; $t^i = -\varepsilon + \underline{v} + (\bar{v} - \underline{v})$	$j$ wins; $t^j = -(\bar{v} - \underline{v}) + \bar{v} + \varepsilon$
$v^i = \underline{v}$	$j$ wins; $t^j = \underline{v}$	$j$ wins; $t^j = \underline{v}$

Thus, expected revenue is given by

$$\begin{aligned} & p(1 - p)(\bar{v} - \varepsilon) + [1 - p(1 - p)]\underline{v} + p^2\varepsilon \\ & = p(1 - p)\bar{v} + [1 - p(1 - p)]\underline{v} + p(2p - 1)\varepsilon \end{aligned}$$

This expression exceeds  $p^2\bar{v} + (1 - p^2)\underline{v}$  for  $p > \frac{1}{2}$  and exceeds  $p(2 - p)\bar{v}$  for large values of  $\underline{v}$ . Thus, whenever  $p > \frac{1}{2}$  and  $\underline{v}$  is sufficiently large, the asymmetric  $\lambda^\varepsilon$ -auction generates more revenue than the optimal symmetric  $\lambda$ -auction.

## 4 Conclusion

We have identified the role of mixed bundling for raising revenue in a multi-object auction without a resale market. We have shown that it is neither in the self-interest of the auctioneer to sell the objects in separate, efficient auctions, nor to force the sale of the entire bundle of objects to a single bidder.

The direct mechanisms used in this paper (which are theoretically easy to define and work with) can be seen as Vickrey-Clarke-Groves mechanisms for a society with an additional "virtual" agent, who receives no objects, but gets utility when bundling allocations are being chosen. Thus, even though the real agents' valuations are assumed to be additive, the virtual agent has "preferences" that imply complementarity among objects. An interesting question is whether one can implement mixed bundling via more "intuitive" bidding mechanisms (e.g., an ascending auction). The answer to this question is not entirely clear since complementarities are known to distort other agents' incentives in ascending-bid auctions (see for example Ausubel and Milgrom, 2002).

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