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# On the sensitivity matrix of the Nash bargaining solution* 

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#### Abstract

In this note we provide a characterization of a subclass of bargaining problems for which the Nash solution has the property of disagreement point monotonicity. While the original $d$-monotonicity axiom and its stronger notion, strong $d$-monotonicity, were introduced and discussed by Thomson [15], this paper introduces local strong $d$-monotonicity and derives a necessary and sufficient condition for the Nash solution to be locally strong $d$-monotonic. This characterization is given by using the sensitivity matrix of the Nash bargaining solution w.r.t. the disagreement point $d$. Moverover, we present a sufficient condition for the Nash solution to be strong $d$-monotonic.


Keywords: Nash bargaining solution, $d$-monotonicity, diagonally dominant Stieltjes matrix. Jel-codes: C61, C62, C71, C78.

## 1 Introduction

In this note we introduce the notion of local strong $d$-monotonicity for solutions of bargaining problems. Thomson introduced and discussed in ([14]) the disagreement point monotononicity property (d-monotonicity) for solutions of bargaining problems. This property states that, if some agent increases his threatpoint while the threatpoint of the other players remains constant then this agent's payoff increases (or at least not decreases). He also considered the stronger, strong d-monotonicity requirement, which states that not only this agent's payoff does not decrease but also the payoffs of none of the other agents increases. Thomson shows by means of a counterexample that the Nashsolution ( $N$-solution) does not satisfy this notion of strong $d$-monotonicity.
This notion of $d$-monotonicity is a global property in the sense that this property should hold for every positive increment of the threatpoint at every threatpoint $d$.
We will consider here the local version of this property. That is, we are interested in the effect of

[^0]changes on the point of agreement for a fixed feasible set if one (arbitrarily chosen) player unilaterally changes his disagreement point something. If this player is the only one who gains from such a small (positive) deviation and this property holds irrespective of which player alters his threatpoint we call the bargaining solution local strongly d-monotonic at the threatpoint $d$.
Given some threatpoint and the corresponding bargaining point, this notion tells us something about the stability of the realized bargaining point. This, in the following sense. Assume that the threatpoint can be controlled to some extent by an exogenous authority (e.g a European commission who might consider to change some directives which might favor some outside options of participating countries). If the bargaining point is local strongly $d$-monotonic at $d$ then whenever this threatpoint is changed at one entry only, this action will be disapproved by all other players. This, in contrast to the case that such a change in the threatpoint is benefitial for some other player(s) too. In that case it is rational for that (those) other player(s), at least, to be not against such a change in the threatpoint. So, a less number of players will be against a reopening of the bargaining process in such a case. In this sense, the threshold to reopen the bargaining process will be lower, and the bargaining point might be called less stable.
So, this notion of local strong $d$-monotonicity can be viewed as a new independent axiom for a bargaining solution which implies stability. We give in section 3 below a necessary and sufficient condition of domain restriction over which the $N$-solution has the property of local strong $d$-monotonicity. Furthermore, we present in this section a sufficient condition for strong $d$-monotonicity.
Section 2 introduces some notation and preliminary results, whereas section 4 considers some examples. Finally section 5 concludes.

## 2 Preliminaries

Following Thomson [15], we define an $n$-person bargaining problem to be a pair ( $S, d$ ), where $S \subset \mathbb{R}^{n}$ is called the feasible set, $\mathbb{R}^{n}$ the utility space and $d$ the disagreement point.
Thomson considers two classes of bargaining problems: 1) $\bar{\Sigma}^{n}$, where the feasible set $S$ is assumed to be convex, compact and such that there exists a $x \in S$ with $x>d$ (here we use the vector inequality notation); and 2) $\Sigma^{n}$, which is a subclass of $\bar{\Sigma}^{n}$, the so-called class of comprehensive bargaining problems. This subclass is obtained by considering just those elements in $S$ satisfying the additional property that whenever $x \in S$ and $d<\bar{x} \leq x$, then $\bar{x} \in S$.
We will consider in this paper a subclass $\Sigma_{P}^{n}$ of $\Sigma^{n}$. We assume that the (fixed) feasible set in this subclass $\Sigma_{P}^{n}$ satisfies the additional requirement that the set P of (weak) Pareto optimal solutions can be described by a smooth strictly concave function $\varphi$, that is $\sum_{P}^{n}=\left\{\left(x_{1}, \ldots, x_{n}\right)^{T} \in S \mid x_{i} \geq d_{i}, x_{n} \leq\right.$ $\varphi\left(x_{1}, \ldots, x_{n-1}\right)$, and whenever $x \in S$ and $d \leq y \leq x$, then $\left.y \in S\right\}$. This class of problems (for larger classes of bargaining problems, see e.g., [11] or [15]) is particular popular in applied economic sciences (see e.g. the literature on policy coordination [12], [16], [6], [3], [13]).
Given this class of $n$-person bargaining problems, a solution is a function $F$ associating with every $(S, d)$ in this class the point of agreement $F(S, d) \in S$. Since we consider here a fixed feasible set, the dependence of $F$ on $S$ will be omitted. $F$ is called the Nash solution, $N$, if for every fixed pair $(S, d), F(S, d)$ is assigned the point where the product $\Pi\left(x_{i}-d_{i}\right)$ is maximized for $x \in S$ with $x \geq d$.

For notational convenience $\mathbf{n}$ denotes the set $\{1, \cdots, n\}$. Furthermore, $I$ is the identity matrix, $e_{i}$ the $i^{\text {th }}$ standard basis vector in $\mathbb{R}^{n}, v^{T}$ the transpose of a vector/matrix $v, e$ the vector $(1, \cdots, 1)^{T}$
and $\underline{0}$ the zero vector $(0, \cdots, 0)^{T}$. The dimension of these vectors will be clear from the context. Furthermore, $\operatorname{diag}\left(a_{i}\right)$ denotes a diagonal matrix with as its $i^{t h}$ diagonal entry $a_{i} ;(A \mid B)$ the extended matrix of $A$ and $B$; and $\operatorname{sgn}(a)$ the sign of the number $a$. If $x:=\left(x_{1}, \cdots, x_{n}\right)$ is a vector, $x_{-}$is the truncated vector $\left(x_{1}, \cdots, x_{n-1}\right) . \varphi_{i}^{\prime}$ denotes the $i$-th partial derivative of $\varphi$.

The property of local strong $d$-monotonicity with respect to the disagreement point $d$ is now formalized as follows:

Definition 1: A bargaining solution $F$ on $\Sigma_{P}^{N}$ is called local strongly d-monotonic at a problem $(S, d) \in \Sigma_{P}^{N}$, if $F$ is differentiable in $d$, and for all $i$ and $j \neq i, \frac{\partial F_{j}(S, d)}{\partial d_{i}} \leq 0$ and $\frac{\partial F_{i}(S, d)}{\partial d_{i}} \geq 0$. []

In the ensueing analysis the set of so-called $M$-matrices arise in a natural way. An $M$-matrix is an $n \times n$ matrix with nonpositive off-diagonal entries whose inverse exists and is entry-wise nonnegative. Symmetric $M$-matrices are called Stieltjes matrices. From Berman et al. [1, pp.141] we recall the following result.

## Lemma 1:

1) Symmetric $M$-matrices are positive definite.
2) Symmetric positive definite matrices with nonpositive off-diagonal entries are $M$-matrices.

Unfortunately, the inverse of a nonsingular nonnegative matrix is not in general an $M$-matrix. In literature the problem has been addressed to characterize all matrices which do have this property. This turns out to be a difficult problem. A class of matrices that satisfy this property are e.g. the so-called strictly ultrametric matrices (see Nabben et al. [8] and [9]).
Finally, we call a symmetric square matrix $A=\left(a_{i j}\right)$ diagonally dominant if $\left|a_{i i}\right| \geq \sum_{j \neq i}\left|a_{i j}\right|$, for all $i$.

## 3 Theoretical Results

By assumption, the Nash bargaining solution $x^{N}:=\left(x_{1}^{N}, \cdots, x_{n}^{N}\right)$ is determined by the argument that solves the maximization problem

$$
\max _{x_{-}} f\left(x_{-}\right):=\max _{x_{-}} \Pi_{i \in \mathbf{n}-\mathbf{1}}\left(x_{i}-d_{i}\right)\left(\varphi\left(x_{-}\right)-d_{n}\right)
$$

where $\varphi^{\prime}{ }_{i}<0$ and $\varphi^{\prime \prime}$ is negative definite.
This maximization problem has, according to Nash [10], exactly one solution. Obviously, this solution lies not on the edge of the Paretofrontier $P$ of $\Sigma_{P}^{N}$, i.e., it is an interior point of $P$. Thus, the first order conditions yield that the Nash bargaining solution is uniquely determined by:

$$
\begin{equation*}
g_{i}\left(x_{-}^{N}, d\right)=0, \forall i \in \mathbf{n}-\mathbf{1}, \tag{1}
\end{equation*}
$$

where $g_{i}\left(x_{-}, d\right):=\varphi\left(x_{-}\right)-d_{n}+\left(x_{i}-d_{i}\right) \varphi_{i}^{\prime}\left(x_{-}\right), i \in \mathbf{n}-\mathbf{1}$.
Note that all derivatives in these $\mathrm{n}-1$ equations are evaluated at the Nash solution. To simplify notation we will drop this argument whenever it is the Nash solution. So, unless stated differently,
we assume from now on that the argument in the derivatives will always be the Nash solution.
Remark 1 Recall from the two-player case that the Nash solution can geometrically also be characterized as that point $x^{N}$ on the curve $\varphi$ that has the property that the line tangent to $\varphi$ at $x^{N}$ intersects the $d_{1}$-axis at the point $d_{1}+2\left(x_{1}^{N}-d_{1}\right)$. For the multi-player case this generalizes as follows. Consider the plane tangent to the graph of $\varphi$ at $x$, i.e.,

$$
\begin{equation*}
y_{n}\left(y_{-}\right)=\varphi\left(x_{-}\right)+\varphi^{\prime}\left(x_{-}\right)\left(y_{-}-x_{-}\right) \tag{2}
\end{equation*}
$$

Then $x^{N}$ is the point that satisfies the property that this plane intersects the vertical plane through the Nash point parallel to the $d_{i}$-axis in the $d_{n}$-plane at the point $y_{-}=\left(x_{1}^{N}, \cdots, x_{i-1}^{N}, d_{i}+2\left(x_{i}^{N}-\right.\right.$ $\left.\left.d_{i}\right), x_{i+1}^{N}, \cdots, x_{n-1}^{N}\right), i \in \mathbf{n}-\mathbf{1}$. Substitution of this into (2) yields the equations (1).

Since the solution of the above optimization problem is a maximum location we know that the second order derivative $H$ of $f$ evaluated at the Nash solution will be semi-negative definite. Simple calculations show that

$$
\begin{equation*}
H=D g^{\prime} \tag{3}
\end{equation*}
$$

where the $i^{t h}$ entry, $d_{i i}$, of the diagonal matrix $D$ is $\Pi_{j \neq i \in \mathbf{n}-\mathbf{1}}\left(x_{j}^{N}-d_{j}\right)$ and

$$
\begin{equation*}
g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right)=\frac{\partial g}{\partial x_{-}}\left(x_{-}^{N}, d\right)=\left(e e^{T}+I\right) \operatorname{diag}\left(\varphi_{i}^{\prime}\right)+\operatorname{diag}\left(x_{i}^{N}-d_{i}\right) \varphi^{\prime \prime} \tag{4}
\end{equation*}
$$

We will assume throughout this note additionally that $H$ is invertible. In particular it follows then from (3) that the inverse of $g^{\prime}$ exists and $g^{\prime-1}=H^{-1} D$. According the implicit function theorem

$$
\frac{\partial x_{-}^{N}}{\partial d}=-\left\{\frac{\partial g}{\partial x_{-}}\left(x_{-}^{N}, d\right)\right\}^{-1} \frac{\partial g}{\partial d}
$$

It is easily verified that

$$
\begin{equation*}
\frac{\partial g}{\partial d}=-\left(\operatorname{diag}\left(\varphi_{i}^{\prime}\right) \mid e\right) \tag{5}
\end{equation*}
$$

To complete the picture of $\frac{\partial x_{i}^{N}}{\partial d_{j}}$ we still have to consider $\frac{\partial x_{n}^{N}}{\partial d_{j}}$. To that end we recall that $x_{n}^{N}=\varphi\left(x_{-}^{N}\right)$. Consequently,

$$
\frac{\partial x_{n}^{N}}{\partial d_{j}}=\varphi^{\prime}\left(\begin{array}{c}
\frac{\partial x_{1}^{N}}{\partial d_{j}} \\
\vdots \\
\frac{\partial x_{n-1}^{N}}{\partial d_{j}}
\end{array}\right), \text { where } \varphi^{\prime}:=\left(\varphi_{1}^{\prime}, \cdots, \varphi_{n-1}^{\prime}\right)
$$

So, with $L:=\binom{I}{\varphi^{\prime}}$, we have that $\frac{\partial x^{N}}{\partial d}=-L\left\{\frac{\partial g}{\partial x_{-}}\left(x_{-}^{N}, d\right)\right\}^{-1} \frac{\partial g}{\partial d}$.
Before we present the sensitivity matrix we introduce for notational convenience

$$
v_{i}^{N}:=\frac{x_{i}^{N}-d_{i}}{\sqrt{\varphi\left(x_{-}^{N}\right)-d_{n}}} \text { and } G:=\left(-\left(e e^{T}+I\right)+\left(\varphi-d_{n}\right) \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right) \varphi^{\prime \prime} \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right)\right)^{-1}
$$

Theorem 1: Under the assumption that the hamiltonian of the Pareto frontier evaluated at the $N$ solution is invertible, the sensitivity matrix of the $N$-solution is given by

$$
\begin{equation*}
\frac{\partial x^{N}}{\partial d}=-\binom{\operatorname{diag}\left(v_{i}^{N}\right)}{-v_{n}^{N} e^{T}} G\left(\left.\operatorname{diag}\left(\frac{1}{v_{i}^{N}}\right) \right\rvert\, \frac{-1}{v_{n}^{N}} e\right) . \tag{6}
\end{equation*}
$$

Proof: Using $(4,5)$ it is clear that

$$
\begin{equation*}
\frac{\partial x^{N}}{\partial d}=\binom{I}{\varphi^{\prime}}\left(\left(e e^{T}+I\right) \operatorname{diag}\left(\varphi_{i}^{\prime}\right)+\operatorname{diag}\left(x_{i}^{N}-d_{i}\right) \varphi^{\prime \prime}\right)^{-1}\left(\operatorname{diag}\left(\varphi_{i}^{\prime}\right) \mid e\right) \tag{7}
\end{equation*}
$$

Some elementary rewriting of this equation (7) gives:

$$
\begin{array}{r}
\frac{\partial x^{N}}{\partial d}=\binom{I}{\varphi^{\prime}}\left(\left(e e^{T}+I\right) \operatorname{diag}\left(\frac{\varphi_{i}^{\prime}\left(x_{i}^{N}-d_{i}\right)}{x_{i}^{N}-d_{i}}\right)+\operatorname{diag}\left(x_{i}^{N}-d_{i}\right) \varphi^{\prime \prime} \operatorname{diag}\left(x_{i}^{N}-d_{i}\right) \operatorname{diag}\left(\frac{1}{x_{i}^{N}-d_{i}}\right)\right)^{-1} \\
\left(\left.\operatorname{diag}\left(\frac{\varphi_{i}^{\prime}\left(x_{i}^{N}-d_{i}\right)}{x_{i}^{N}-d_{i}}\right) \right\rvert\, e\right)
\end{array}
$$

From (1) we have that at the N -solution

$$
\begin{equation*}
\varphi_{i}^{\prime}\left(x_{i}^{N}-d_{i}\right)=\varphi_{j}^{\prime}\left(x_{j}^{N}-d_{j}\right)=-\left(\varphi-d_{n}\right) \tag{8}
\end{equation*}
$$

Using this, we can rewrite the above equation as follows

$$
\begin{aligned}
\frac{\partial x^{N}}{\partial d}= & \binom{I}{\varphi^{\prime}} \operatorname{diag}\left(x_{i}^{N}-d_{i}\right)\left(-\left(\varphi-d_{n}\right)\left(e e^{T}+I\right)+\operatorname{diag}\left(x_{i}^{N}-d_{i}\right) \varphi^{\prime \prime} \operatorname{diag}\left(x_{i}^{N}-d_{i}\right)\right)^{-1} \\
& \left(\left.-\left(\varphi-d_{n}\right) \operatorname{diag}\left(\frac{1}{x_{i}^{N}-d_{i}}\right) \right\rvert\, e\right) \\
= & \binom{\operatorname{diag}\left(v_{i}^{N}\right)}{-v_{n}^{N} e^{T}}\left(-\left(e e^{T}+I\right)+\operatorname{diag}\left(v_{i}^{N}\right) \varphi^{\prime \prime} \operatorname{diag}\left(v_{i}^{N}\right)\right)^{-1}\left(\left.-\operatorname{diag}\left(\frac{1}{v_{i}^{N}}\right) \right\rvert\, \frac{1}{v_{n}^{N}} e\right) \\
= & \binom{\operatorname{diag}\left(v_{i}^{N}\right)}{-v_{n}^{N} e^{T}}\left(-\left(e e^{T}+I\right)+\operatorname{diag}\left(\frac{\varphi_{i}^{\prime} v_{i}^{N}}{\varphi_{i}^{\prime}}\right) \varphi^{\prime \prime} \operatorname{diag}\left(\frac{\varphi_{i}^{\prime} v_{i}^{N}}{\varphi_{i}^{\prime}}\right)\right)^{-1}\left(\left.-\operatorname{diag}\left(\frac{1}{v_{i}^{N}}\right) \right\rvert\, \frac{1}{v_{n}^{N}} e\right) .
\end{aligned}
$$

From this, using (8) and the above introduced notation, (6) is obtained.
Elementary spelling out (6) shows that the sensitivity matrix can also be written as

$$
\frac{\partial x^{N}}{\partial d}=-\left(\begin{array}{cc}
\operatorname{diag}\left(v_{i}^{N}\right) G \operatorname{diag}\left(\frac{1}{v_{i}^{N}}\right) & \frac{-1}{v_{n}^{N}} \operatorname{diag}\left(v_{i}^{N}\right) G e  \tag{9}\\
-v_{n}^{N} e^{T} G \operatorname{diag}\left(\frac{1}{v_{i}^{N}}\right) & e^{T} G e
\end{array}\right) .
$$

Since, by assumption, $\varphi^{\prime \prime}$ is negative definite $G$ is negative definite too. Using this, it follows immediately from (9) that all diagonal entries of the sensitivity matrix are always positive. Or stated differently,

Corollary 1: The N-solution is d-monotonic.

Next, we address the question under which conditions on $\varphi$ the N -solution is local strongly $d$ monotonic. We have the following result:

Theorem 2: The N -solution is local strongly $d$-monotonic if and only if $-G$ is a diagonally dominant Stieltjes matrix.

Proof: Consider (9). Since $v_{i}^{N}>0$ it follows that $\operatorname{sgn}\left(\left(\frac{\partial x^{N}}{\partial d}\right)_{i j}\right)=\operatorname{sgn}\left(-G_{i j}\right), i, j \in \mathbf{n}-\mathbf{1}$. As already noted before, $-G$ is positive definite. So, by Lemma 1.2), $-G$ is a Stieltjes matrix. Moreover it follows from (9) that $\operatorname{sgn}\left(\left(\frac{\partial x^{N}}{\partial d}\right)_{i n}\right)=\operatorname{sgn}\left(\frac{v_{i}^{N}}{v_{n}^{N}} e_{i}^{T} G e\right)$. So, $\left(\frac{\partial x^{N}}{\partial d}\right)_{i n} \leq 0$ if and only if entry $i$ of $G e \leq 0$, $i \in \mathbf{n}-\mathbf{1}$. Or, stated differently, $-G$ is diagonally dominant.

## Remark 2

1) Since $G=-\operatorname{diag}\left(\varphi_{i}^{\prime}\right)\left(g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right)\right)^{-1},-G^{-1}$ is nonnegative if and only if $g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right)$ is nonpositive. Furthermore, as already noticed before, $-G$ is a positive definite matrix. Therefore, an equivalent statement of Theorem 2 is: the $N$-solution is local strongly d-monotonic if and only if $g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right)$ is nonpositive and $\operatorname{diag}\left(\varphi_{i}^{\prime}\right)\left(g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right)\right)^{-1}$ is a diagonally dominant matrix with nonpositive offdiagonal entries.
This clarifies the statement about ultrametric matrices we made at the end of the preliminaries.
Notice that $g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right)$ is nonpositive if, e.g., $\varphi^{\prime \prime}$ is a nonpositive matrix.
2) Since $-G$ is an $M$-matrix it follows that $e e^{T}-\left(\varphi-d_{n}\right) \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right) \varphi^{" \prime} \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right)$ is nonnegative. Or, equivalently, $\varphi^{\prime T} \varphi^{\prime}-\left(\varphi-d_{n}\right) \varphi^{\prime \prime}$ is a nonnegative matrix.
3) In the two-player case the $N$-solution is always strongly d-monotonic.

In the three-player case the $N$-solution is local strongly d-monotonic if and only if $g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right)$ is nonpositive and $g_{x_{-}}^{\prime}\left(x_{-}^{N}, d\right) \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right)$ is a diagonally dominant matrix. This follows by a simple elaboration of $-G^{-1}$.
4) In case the Pareto frontier has no extreme bendings, i.e. $\varphi^{\prime}$ is almost constant, $\varphi^{\prime \prime}$ is almost zero. In that case $G$ almost equals $-\left(e e^{T}+I\right)^{-1}$ which is a diagonally dominant Stieltjes matrix. Since this is the case independent of the choice of the threatpoint, under these conditions the $N$-solution will be strongly $d$-monotonic.

Next we derive a sufficient condition on the Pareto frontier under which the $N$-solution will be (global) strongly $d$-monotonic. Inspired by items 1 ) and 2 ) of the above remark we will consider the case that $\varphi$ " is a nonpositive matrix. The result is stated in Theorem 3. Its proof uses the next lemma.

Lemma 2: Assume $S$ is an invertible matrix and $D$ is a positive diagonal matrix. Consider $P:=$ $(S+D)^{-1}$.

1) If $S^{-1}$ is diagonally dominant, then $P$ is diagonally dominant.
2) If $S^{-1}$ is a Stieltjes matrix, then $P$ is a Stieltjes matrix.

Proof: 1) First notice that

$$
\begin{equation*}
(S+D)^{-1}=D^{-1}-D^{-1}\left(D^{-1}+S^{-1}\right)^{-1} D^{-1} \tag{10}
\end{equation*}
$$

Next consider

$$
H:=\left(\begin{array}{cc}
D^{-1}+S^{-1} & D^{-1} \\
D^{-1} & D^{-1}
\end{array}\right)
$$

Due to our assumptions, it is easily verified that $H$ is diagonally dominant. From e.g. Lei et al. [7] (see also Carlson et al. [2]) we conclude then that the Schur complement of $H$, which equals (10), is also diagonally dominant.
2) Since by assumption $S^{-1}$ is a Stieltjes matrix, by Lemma 1.1), $S^{-1}$ is a positive definite matrix. From this it is obvious that $P$ will be positive definite too. So, the diagonal entries of $P$ are positive. Furthermore since, by assumption, both $S$ and $D$ are a nonnegative matrix also $S+D$ is a nonnegative matrix. Next we consider the off-diagonal entries of $P$. Since both $D^{-1}$ and $S^{-1}$ are Stieltjes matrices, also $D^{-1}+S^{-1}$ is a Stieltjes matrix. So, in particular, all entries of $\left(D^{-1}+S^{-1}\right)^{-1}$ are nonnegative. From (10) it is obvious then that all off-diagonal entries of $P$ are nonpositive. Since we already argued above that $P$ is positive definite, Lemma 1.2) shows that $P$ is a Stieltjes matrix. []

Theorem 3: Assume that at any point of the Pareto frontier

$$
\begin{equation*}
\Phi^{-1}:=-\left[\operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right) \varphi^{\prime \prime} \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right)\right]^{-1} \tag{11}
\end{equation*}
$$

is a diagonally dominant Stieltjes matrix. Then, the N -solution is strongly d-monotonic.
Proof: What has to be shown is that irrespective of the choice of the threatpoint $d$ the N -solution will be local strongly d-monotonic. Or, equivalently (see Theorem 2), that matrix $-G$ is irrespective of the choice of the threatpoint $d$ a diagonally dominant Stieltjes matrix.
To that end first note that, since $\Phi^{-1}$ is a diagonally dominant Stieltjes matrix, also $\left(\varphi-d_{n}\right) \Phi^{-1}$ is a diagonally dominant Stieltjes matrix. So, by Lemma 2,

$$
\begin{equation*}
P:=\left(I+\left(\varphi-d_{n}\right) \Phi\right)^{-1} \tag{12}
\end{equation*}
$$

is a diagonally dominant Stieltjes matrix. Next consider $-G$. We have

$$
-G=\left(\left(e e^{T}+I\right)+\left(\varphi-d_{n}\right) \Phi\right)^{-1}=\left(e e^{T}+P^{-1}\right)^{-1}=P-P e\left(e^{T} P e+1\right)^{-1} e^{T} P
$$

Since $P$ is diagonally dominant $P e \geq 0$. Consequently, $P e\left(e^{T} P e+1\right)^{-1} e^{T} P \geq 0$. So, all off-diagonal entries of $-G$ are nonpositive. Obviously, $-G$ is a positive definite matrix and all entries of $-G^{-1}$ are nonnegative. So, by Lemma 1.2), $-G$ is a Stieltjes matrix.
Furthermore it follows from (13) that

$$
-G e=\left(P-P e\left(e^{T} P e+1\right)^{-1} e^{T} P\right) e=\left(1-\frac{e^{T} P e}{1+e^{T} P e}\right) P e \geq 0
$$

That is, $-G$ is diagonally dominant.
Remark 3:

1) Clearly, (11) is only satisfied if $\varphi$ " is nonpositive. Furthermore, it is easily verified that for the scalar case $\Phi=\left[\frac{1}{\varphi^{\prime}}\right]^{\prime}$, whereas for the multivariable case, with $S:=\operatorname{diag}\left(\varphi_{i}^{\prime}\right), \Phi=S\left[\frac{1}{\varphi_{i}^{\prime}} \cdots \frac{1}{\varphi_{n-1}^{\prime}}\right]^{\prime} S^{-1}$. This relationship might be helpful in getting a better intuition about the conditions under which a bargaining solution satisfies the strong $d$-monotonicity property.
2) Consider the next statements:
i) $\Phi^{-1}$ is a diagonally dominant Stieltjes matrix.
ii) $\left(e e^{T}+\left(\varphi-d_{n}\right) \Phi\right)^{-1}$ is a diagonally dominant Stieltjes matrix.
iii) $-G$ is a diagonally dominant Stieltjes matrix.

Then, i) $\Rightarrow$ ii) $\Rightarrow$ iii). The first implication can be shown by using the fact that $\left(e e^{T}+\Phi\right)^{-1}=$ $\frac{1}{\alpha} \Phi^{-1}\left(\alpha I-e e^{T} \Phi^{-1}\right)$, where $\alpha=1+\sum_{i, j=1}^{n} \Phi_{i, j}^{-1}>0$, whereas the second implication follows using similar arguments as in the proof of Theorem 3. Unfortunately none of the reverse implications holds.
3) By a local interpretation of (11) one immediately obtains, using (1), that the $N$-solution is local strongly $d$-monotonic if $-\left[\operatorname{diag}\left(x_{i}^{N}-d_{i}\right) \varphi^{\prime \prime}\left(x_{-}^{N}\right) \operatorname{diag}\left(x_{i}^{N}-d_{i}\right)\right]^{-1}$ is a diagonally dominant Stieltjes matrix.

## Remark 4:

The above analysis can also be used to study the case of weighted Nash solutions. That is the solution that solves the maximization problem

$$
\max _{x_{-}} \tilde{f}\left(x_{-}\right):=\max _{x_{-}} \Pi_{i \in \mathbf{n}-\mathbf{1}}\left(x_{i}-d_{i}\right)^{\alpha_{i}}\left(\varphi\left(x_{-}\right)-d_{n}\right)^{\alpha_{n}}, \alpha_{i}>0 .
$$

Introducing the weight matrix $W_{k} \in \mathbb{R}^{k \times k}$ as $W_{k}:=\operatorname{diag}\left(\frac{\alpha_{n}}{\alpha_{i}}\right)$ we have that with

$$
\begin{gathered}
\tilde{G}:=\left(-\left(e e^{T}+W_{n-1}\right)+\left(\varphi-d_{n}\right) \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right) \varphi^{\prime \prime} \operatorname{diag}\left(\frac{1}{\varphi_{i}^{\prime}}\right)\right)^{-1}, \\
\frac{\partial x^{N}}{\partial d}=-W_{n}\binom{\operatorname{diag}\left(v_{i}^{N}\right)}{-v_{n}^{N} e^{T}} \tilde{G}\left(\left.\operatorname{diag}\left(\frac{1}{v_{i}^{N}}\right) \right\rvert\, \frac{-1}{v_{n}^{N}} e\right) .
\end{gathered}
$$

Using this it follows then that the results of Corollary 1 and Theorems 2 and 3 apply for this case as well, with $G$ replaced by $\tilde{G}$.

## 4 Examples

In this section we provide two examples. The first example provides a number of Pareto frontiers for which the $N$-solution is strongly $d$-monotonic. Intuitively it demonstrates that if the frontier is not too extremely bending one may expect that this property holds.
The second example may be interpreted as a cartel-formation game. For different sets of parameters we present numerically the set of threatpoints where the $N$-solution is local strongly $d$-monotonic.

Example 1:

1) Assume that $\varphi^{\prime \prime}$ is a nonpositive diagonal matrix (so, $\varphi\left(x_{-}\right)$is e.g. a plane or $\varphi\left(x_{-}\right)=r+b^{T} x_{-}+$ $\frac{1}{2} x_{-}^{T} A x_{-}$, where $b, x_{-}$are $n$ - 1 -dimensional vectors with $b \leq 0$ and $d \geq 0$ and $A$ a nonpositive diagonal matrix). Then, for every choice of $d,-G$ is a diagonally dominant Stieltjes matrix. So, see

Theorem 2, the $N$-solution is strongly $d$-monotonic.
2) Assume that the Pareto frontier has a constant curvature, that is

$$
\varphi\left(x_{-}\right)=\sqrt{r^{2}-x_{1}^{2}-\cdots-x_{n-1}^{2}} .
$$

Then $\varphi_{i}^{\prime}=\frac{-x_{i}}{\varphi\left(x_{-}\right)}, \varphi_{i j}^{\prime \prime}=\frac{-x_{i} x_{j}}{\varphi^{3}\left(x_{-}\right)}$if $i \neq j$ and $\varphi_{i i}^{\prime \prime}=\frac{-\left(\varphi^{2}\left(x_{-}\right)+x_{i}^{2}\right)}{\varphi^{3}\left(x_{-}\right)}$. Consequently,

$$
\Phi=\frac{1}{\varphi\left(x_{-}\right)}\left(\begin{array}{ccccc}
\frac{\varphi^{2}\left(x_{-}\right)+x_{1}^{2}}{x_{1}^{2}} & 1 & \cdots & \cdots & 1 \\
1 & \frac{\varphi^{2}\left(x_{-}\right)+x_{2}^{2}}{x_{2}^{2}} & 1 & \cdots & 1 \\
\vdots & \ddots & & & \vdots \\
\vdots & & & \ddots & 1 \\
1 & \cdots & \cdots & 1 & \frac{\varphi^{2}\left(x_{-}\right)+x_{n-1}^{2}}{x_{n-1}^{2}}
\end{array}\right)
$$

and

$$
\Phi^{-1}=\frac{1}{r^{2} \varphi\left(x_{-}\right)}\left(\begin{array}{ccccc}
\left(r^{2}-x_{1}^{2}\right) x_{1}^{2} & -x_{1}^{2} x_{2}^{2} & \cdots & \cdots & -x_{1}^{2} x_{n-1}^{2} \\
-x_{2}^{2} x_{1}^{2} & \left(r^{2}-x_{2}^{2}\right) x_{2}^{2} & -x_{2}^{2} x_{3}^{2} & \cdots & x_{2}^{2} x_{n-1}^{2} \\
\vdots & \ddots & & & \vdots \\
\vdots & & & \ddots & -x_{n-2}^{2} x_{n-1}^{2} \\
-x_{n-1}^{2} x_{1}^{2} & \cdots & \cdots & -x_{n-1}^{2} x_{n-2}^{2} & \left(r^{2}-x_{n-1}^{2}\right) x_{n-1}^{2}
\end{array}\right)
$$

Obviously $\Phi^{-1}$ is a Stieltjes matrix. Furthermore $\Phi^{-1} e=\frac{\varphi}{r^{2}}\left[x_{1}^{2}, \cdots, x_{n-1}^{2}\right]^{T}$. So $\Phi^{-1}$ is diagonally dominant too. Therefore (see Theorem 3) the $N$-solution is strongly $d$-monotonic. It is easily verified that this result also holds if $\varphi\left(x_{-}\right)$is replaced by $\varphi\left(x_{-}\right)=\sqrt{r^{2}-\alpha_{1} x_{1}^{2}-\cdots-\alpha_{n-1} x_{n-1}^{2}}, \alpha_{i}>0$.
3) Assume that the Pareto frontier is described by $\varphi\left(x_{-}\right)=\prod_{i=1}^{n-1}\left(b_{i}-x_{i}\right)^{\alpha_{i}}$, where $0<\alpha_{i}<1$ and $x_{i} \leq b_{i}, i \in \mathbf{n - 1}$. Note that this type of functions includes e.g. the Cobb-Douglas function which often occurs in economics. Then,

$$
\varphi_{i}^{\prime}=\frac{-\alpha_{i}}{b_{i}-x_{i}} \varphi ; \varphi_{i j}^{\prime \prime}=\frac{\alpha_{i} \alpha_{j}}{\left(b_{i}-x_{i}\right)\left(b_{j}-x_{j}\right)} \varphi, i \neq j ; \text { and } \varphi_{i i}^{\prime \prime}=\frac{\alpha_{i}\left(-1+\alpha_{i}\right)}{\left(b_{i}-x_{i}\right)^{2}} \varphi .
$$

Elementary calculations show that then $-\Phi=\frac{1}{\varphi}\left(e e^{T}+\operatorname{diag}\left(\frac{-1}{\alpha_{i}}\right)\right)$. Consequently,

$$
\begin{align*}
-G & =\left(e e^{T}+I+\left(\varphi-d_{n}\right) \Phi\right)^{-1}=\left(e e^{T}+I-\left(\varphi-d_{n}\right) \frac{1}{\varphi}\left(e e^{T}+\operatorname{diag}\left(\frac{-1}{\alpha_{i}}\right)\right)\right)^{-1} \\
& =\varphi\left(d_{n} e e^{T}+\operatorname{diag}\left(\varphi+\frac{\varphi-d_{n}}{\alpha_{i}}\right)\right)^{-1}=\varphi\left(D^{-1}-D^{-1} e\left(e^{T} D^{-1} e+\frac{1}{d_{n}}\right)^{-1} e^{T} D^{-1}\right) \tag{13}
\end{align*}
$$

where $D:=\operatorname{diag}\left(\varphi+\frac{\varphi-d_{n}}{\alpha_{i}}\right)$. From (13) it is easily verified that $-G$ is a Stieltjes matrix. Furthermore $-G e=\left(1-\frac{e^{T} D^{-1} e}{e^{T} D^{-1} e+\frac{1}{d_{n}}}\right) D^{-1} e$. Clearly this vector is positive, so $-G$ is diagonally dominant too. Since, irrespective of the location of the threatpoint $d,-G$ is a diagonally dominant Stieltjes matrix we conclude that the $N$-solution is strongly $d$-monotonic.

## Example 2:

Consider 3 firms who sell an amount $x_{i}$ of a product on a market. The price, $p$, they get on the market depends on the quantity sold by all firms. That is, $p=c-\alpha_{1} x_{1}-\alpha_{2} x_{2}-\alpha_{3} x_{3}$, where $\alpha_{i}$ and $c$ are some positive constants. The costs for producing $x_{i}$ are $C_{i}\left(x_{i}\right)$. So the profits for firm $i$ are $\pi_{i}=p x_{i}-C_{i}\left(x_{i}\right)$. Next consider the parameterized joint profit function

$$
\pi=\sum_{i=1}^{3} \lambda_{i} \pi_{i}, \text { where } \sum_{i=1}^{3} \lambda_{i}=1 \text { and } \lambda_{i} \geq 0
$$

By maximizing for all possible parametercombinations, $\lambda:=\left(\lambda_{1}, \lambda_{2}, \lambda_{3}\right), \pi$ one obtains the Pareto frontier characterizing all possible joint maximal profits the firms can obtain by cooperation.
We will consider two different specifications for the cost functions $C_{i}$.
Case 1: $C_{i}\left(x_{i}\right)=\beta x_{i}$.
In this case straightforward (though lengthy) calculations show that the Pareto frontier is given by the plane:

$$
\alpha_{1} \pi_{i}+\alpha_{2} \pi_{2}+\alpha_{3} \pi_{3}=\left(\frac{c-\beta}{2}\right)^{2} .
$$

So, whatever the threatpoint $d$ is, the $N$-solution is strongly $d$-monotonic in this case (see Example 1.1)).

Case 2: $C_{i}\left(x_{i}\right)=\beta_{i} x_{i}^{2}$.
In this case it is not possible to derive an analytic expression for the Pareto frontier. We will briefly indicate how one may pursue in this case numerically to verify the local strong $d$-monotonicity of the $N$-solution. Differentiation of $\pi$ w.r.t. $x_{i}$ yields 3 first order conditions in $x_{i}$ for every $\lambda$ (in this case this is just a set of linear equations). From this one can solve $x_{i}$ in terms of $\lambda$. Using this one can determine then for an arbitrary threatpoint, with e.g. the numerical algorithm outlined in Douven [4, Section 3.3.2] (see also Engwerda [5, Section 6.4]), the with this threatpoint corresponding $N$-solution.
From the seven equations $\pi_{i}-p x_{i}-C_{i}\left(x_{i}\right)=0, \frac{\partial \pi}{\partial x_{i}}=0, i=1,2,3$ and $\sum_{i=1}^{3} \lambda_{i}=1$, one can then implicitly solve $\left(\pi_{3}, x_{1}, x_{2}, x_{3}, \lambda_{1}, \lambda_{2}, \lambda_{3}\right)$ as a function of $\pi_{1}$ and $\pi_{2}$. In particular the implicit function theorem can be used to find analytic expressions for the derivative and hamiltonian of $\pi_{3}=\varphi\left(\pi_{1}, \pi_{2}\right)$ at the $N$-solution. From this then straightforwardly the monotonicity property can be verified.
Figures 1-3 present some results for this example for different parameters and threatpoints. In all these cases the parameters $\alpha_{1}=1, \beta_{i}=1$ and $c=5$ remained unchanged.
Figure 1 reports the local strong $d$-monotonicity property of the $N$-solution, if the threatpoint $d=0$, for different values of the parameters $\alpha_{2}$ and $\alpha_{3}$ (both ranging between 0.1 and 0.5 ). A dot (empty space) indicates that with that choice of parameters the zero-threatpoint is (not) local strongly $d$-monotonic. Not shown here is that for large values of these parameters the zero-threatpoint is also local strongly $d$-monotonic. So a situation with two firms having a small impact on the price compared to the third firm seems to be not stable (in the sense discussed in the introduction). Here the notion "small" should however be interpreted in the light of the other modelparameters that were kept constant. Some additional experiments suggest that the level of the $\beta_{i}$ parameters are more important than that of the $c$ parameter for this comparison. In case all firms have a substantial effect on the price, the $N$-solution is always local strongly $d$-monotonic.
Figure 2 reports the local strong $d$-monotonicity property for different threatpoints in case $\alpha_{2}=$


Figure 1: Case $\alpha_{1}=1, \beta_{i}=1, c=5$, $d=0$. Dot=local strongly $d$-monotonic.


Figure 2: Case $\alpha_{1}=1, \alpha_{2}=\alpha_{3}=\frac{1}{4}$, $\beta_{i}=1, c=5$.
$\alpha_{3}=\frac{1}{4}$. The threatpoint $d_{1}$ ranges here from 0 to 1.9 and $d_{2}=d_{3}$ ranges from 0 to 3.1. To complete the three dimensional picture we plotted in Figure 3 for three different values of $d_{1}(0,0.5$ and 0.7 , respectively) using the same modelparameters the monotonicity result if the other two threatpoints $d_{2}$ and $d_{3}$ range between 0 and 3.1. Notice that the more closer the threatpoint is to the Pareto frontier, the more this frontier resembles a plane. For that reason in fact the dots in the right and upper part of these graphs extend until the Pareto frontier. For numerical simplicity we did not plot this extension.
Concluding, this example demonstrates that in case firms have a substantial effect on the price, the $N$-solution is strong $d$-monotonic. In case at least one firm has a "small" (see above discussion) impact on the price there exist area's of threatpoints where the $N$-solution is not local strongly $d$ monotonic. Furthermore we observe the phenomenon that if at a certain threatpoint the $N$-solution is local strongly $d$-monotonic this does not imply that at every larger threatpoint the $N$-solution will have this property too.

## 5 Concluding remarks

In this note we derived, under some technical conditions, the sensitivity matrix of the Nash bargaining solution w.r.t. the disagreement point $d$. In particular, this makes it possible to analyze the local strong $d$-monotonicity of the $N$-solution. We showed that the $N$-solution satisfies this property if and only if a certain matrix, $-G$, evaluated at the Nash bargaining solution is a diagonally dominant Stieltjes matrix. Using this result, a class of bargaining problems was characterized for which the $N$-solution satisfies the strong d-monotonicity property. The results were illustrated in a number of examples.
The condition under which the Nash solution is local strongly $d$-monotonic is phrased in terms of the (second) order derivative of the Pareto frontier. Unfortunately at this moment a clear intuition


Figure 3: Case $\alpha_{1}=1, \alpha_{2}=\alpha_{3}=\frac{1}{4}, \beta_{i}=1, c=5$.
about the set of problems for which the $N$-solution is (local) strongly d-monotonic is lacking. From the condition it is clear that in case the Pareto frontier has no extreme bendings, the property will hold. This implies that if in the bargaining problem the interests of the players are similar the $N$ solution will be (local) strongly $d$-monotonic. Finding a geometric interpretation of the conditions and, from that, more intuition about the set of bargaining problems for which the $N$-solution satisfies the monotonicity properties remains, however, an open problem.

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