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Von Neumann-Morgenstern Clutters

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**Abstract** - A clutter on a set  $X$  is a simple hypergraph with pairwise not-comparable hyperedges, hence in particular any set of Von Neumann-Morgenstern (VNM) -stable sets of an irreflexive simple digraph is a clutter. A clutter  $(X, E)$  is representable by VNM-stable sets or VNM if there exists an irreflexive simple digraph  $(X, \Delta)$  such that  $E$  is a set of VNM-stable sets of  $(X, \Delta)$ . The class of VNM clutters on a set  $X$  is characterized.

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Let  $H = (X, \mathbf{E})$  be a *simple hypergraph* i.e.  $X$  is a set, and  $\mathbf{E} \subseteq \mathcal{P}(X)$ , where  $\mathcal{P}(X)$  is the power set of  $X$ :  $X$  will also be referred to as the set of vertices of  $H$ , and  $\mathbf{E}$  as the set of hyperedges of  $H$ . Then,  $H$  is a *clutter* iff  $\mathbf{E}$  is a  $\subseteq$ -antichain, namely  $E \not\subseteq E'$  for any two *distinct*  $E, E' \in \mathbf{E}$ . In particular the ‘boundary’ cases with  $\mathbf{E} = \{\emptyset\}$ , or  $\mathbf{E} = \emptyset$  are allowed. [Notice that such a terminology is slightly at variance with that of other authors, who sometimes denote as ‘simple hypergraphs’, ‘Sperner systems’ or ‘Sperner families’ those clutters  $(X, \mathbf{E})$  -as defined above- such that  $\emptyset \notin \mathbf{E}$  and  $\bigcup \mathbf{E} = X$ : see e.g. Berge (1989)]. A *simple digraph* is a pair  $D = (X, \Delta)$  where  $X$  is a set and  $\Delta \subseteq X \times X$ ;  $D$  is *irreflexive* -or loopless- iff  $(x, x) \notin \Delta$  for any  $x \in X$ . A *Von Neumann-Morgenstern stable set* (henceforth *VNM-stable set*) of an irreflexive simple digraph  $(X, \Delta)$  is a set  $S \subseteq X$  such that (i) internal stability i.e.  $[(x, y) \notin \Delta \text{ for any } x, y \in S]$ , and (ii) external stability i.e. [for any  $z \in X \setminus S$  there exists  $x \in S$  such that  $(x, z) \in \Delta$ ] hold (see e.g. Von Neumann and Morgenstern (1953), Schmidt and Ströhlein (1985), and Ghoshal, Laskar and Pillone (1998)). Let  $\mathcal{S}(X, \Delta)$  denote the set of all VNM-stable sets of  $(X, \Delta)$ . Clearly,  $S_1 \not\subseteq S_2$  for any two *distinct*  $S_1, S_2 \in \mathcal{S}(X, \Delta)$  (otherwise, properties (i) and (ii) of  $S_1$  turn out to be mutually inconsistent). Thus, for any irreflexive simple digraph  $(X, \Delta)$ ,  $(X, \mathcal{S}(X, \Delta))$  is a clutter. This observation motivates the following

**Definition 1** (*VNM clutters*) *A clutter  $H = (X, \mathbf{E})$  is representable by VNM-stable sets or VNM iff there exists an irreflexive simple digraph  $(X, \Delta)$  such that  $\mathbf{E} \subseteq \mathcal{S}(X, \Delta)$ .*

In general, a clutter may or may not be VNM, as made clear by the following examples

**Example 2** *Consider  $X = \{x, y, z\}$  with  $x \neq y \neq z \neq x$ , and*

$$\mathbf{E} = \{\{x, y\}, \{x, z\}, \{y, z\}\}.$$

*It is easily checked that  $(X, \mathbf{E})$  is a clutter, but not a VNM one. Indeed, suppose to the contrary that there exists an irreflexive simple digraph  $(X, \Delta)$  such that  $\mathbf{E} \subseteq \mathcal{S}(X, \Delta)$ . Then, by property (ii) of VNM-stable sets as applied to  $\{x, y\}$ ,  $\Delta \cap \{(x, z), (y, z)\} \neq \emptyset$ , which contradicts the assumption that both  $\{x, z\}$  and  $\{y, z\}$  satisfy property (i) of VNM-stable sets.*

**Example 3** *Consider now  $X$  as defined above in the previous example,  $\mathbf{E}' = \{\{x, y\}, \{x, z\}\}$ , and  $\Delta' = \{(y, z), (z, y)\}$ . Clearly,  $\mathbf{E}' = \mathcal{S}(X, \Delta')$  i.e.  $(X, \mathbf{E}')$  is indeed a VNM clutter.*

**Example 4** Let  $(X, \mathbf{E})$  be a clutter such that  $E \cap E' = \emptyset$  for any pair of distinct  $E, E' \in \mathbf{E}$ .

Then, define  $\Delta = \{(x, y) \in X \times X : \{x, y\} \not\subseteq E \text{ for all } E \in \mathbf{E}\}$ .

Clearly, for any  $E \in \mathbf{E}$  and any  $x, y \in E$ ,  $(x, y) \notin \Delta$  by definition. Moreover, suppose there exists  $z \in X \setminus E$  such that  $(x, z) \notin \Delta$  for all  $x \in E$ . Then, for any  $x \in E$  there exists  $E' \in \mathbf{E}$  such that  $\{x, z\} \subseteq E'$ . Since by construction  $E \neq E'$  and  $E \cap E' \neq \emptyset$ , the existence of such an  $E'$  contradicts our starting hypothesis. It follows that  $E$  is indeed a VNM-stable set of  $(X, \Delta)$ . Therefore,  $(X, \mathbf{E})$  is a VNM clutter.

**Remark 5** (Kernel-representable clutters) Let  $(X, \Delta)$  be an irreflexive simple digraph, and  $\Delta^{-1} \subseteq X \times X$  the inverse of  $\Delta$ , namely, for any  $x, y \in X$ ,  $(x, y) \in \Delta^{-1}$  iff  $(y, x) \in \Delta$ . Of course,  $(X, \Delta^{-1})$  is also an irreflexive simple digraph. A subset of vertices  $K \subseteq X$  is a kernel of  $(X, \Delta)$  - written  $K \in \mathcal{K}(X, \Delta)$  - iff  $K \in \mathcal{S}(X, \Delta^{-1})$ . Now, consider any clutter  $(X, \mathbf{E})$  and declare it kernel-representable (KR) iff there exists an irreflexive simple digraph such that  $\mathbf{E} \subseteq \mathcal{K}(X, \Delta)$ . It is immediately checked that, in view of the foregoing observations, a clutter is KR iff it is VNM.

The foregoing observations and examples raise the issue of identifying VNM clutters ( in the same vein, a representation problem concerning finite lattices and stable matchings is addressed by Blair (1984)).

The aim of this note is in fact to provide a simple characterization of VNM clutters. In order to state our result in a most concise manner, let us first introduce some further auxiliary notions.

**Definition 6** (Conjugation relation of a simple hypergraph) Let  $H = (X, \mathbf{E})$  be a simple hypergraph. Then the conjugation relation  $I_H \subseteq X \times X$  of  $H$  is defined as follows: for any  $x, y \in X$ ,  $(x, y) \in I_H$  iff there exists an  $E \in \mathbf{E}$  such that  $\{x, y\} \subseteq E$ .

**Definition 7** (Complete bigraph) Let  $Y, Z$  be two nonempty sets. A (simple, nontrivial) bigraph on  $(Y, Z)$  is a triple  $(Y, Z, R)$  where  $R \subseteq Y \times Z$ . A bigraph  $(Y, Z, R)$  is complete if  $R = Y \times Z$ .

**Definition 8** Let  $(X, \Delta)$  be a simple digraph. A  $\Delta$ -basis of a complete bigraph is a  $B \subseteq X$  such that  $(B, X \setminus B, \Delta \cup \Delta^{-1})$  is a complete bigraph.

**Definition 9** (Complete bigraph generator (CBG)) *An hyperedge  $E \in \mathbf{E}$  of a clutter  $H = (X, \mathbf{E})$  is a complete bigraph generator iff it is an  $I_H$ -basis of a complete bigraph.*

A clutter  $H = (X, \mathbf{E})$  is *CBG-free* iff no  $E \in \mathbf{E}$  is a complete bigraph generator. Moreover, clutter  $H$  is *conjugation-saturated* iff it is CBG-free and  $I_H \neq I_{H'}$  for any CBG-free clutter  $H' = (X, \mathbf{E}')$  such that  $\mathbf{E} \subset \mathbf{E}'$ .

Then, we have the following

**Theorem 10** *Let  $H = (X, \mathbf{E})$  be a clutter. Then,  $H$  is VNM iff it is CBG-free.*

**Proof.** Observe that, by definition,  $H$  is not CBG iff for any  $E \in \mathbf{E}$  and any  $z \in X \setminus E$  there exists  $x_z \in E$  such that  $(z, x_z) \notin I_H$ .

Then, suppose  $H = (X, \mathbf{E})$  is VNM, and let  $(X, \Delta)$  be an irreflexive simple digraph such that  $\mathbf{E} \subseteq \mathcal{S}(X, \Delta)$ . Now, assume that there exist an hyperedge  $E \in \mathbf{E}$  and a vertex  $z \in X \setminus E$  such that  $(z, x) \in I_H$  for all  $x \in E$ . Thus, for any  $x \in E$  there exists an  $E_{zx} \in \mathbf{E}$  such that  $\{x, z\} \subseteq E_{zx}$  whence, by internal stability of  $E_{zx}$ ,  $(x, z) \notin \Delta$ : but then, external stability of  $E$  is violated. It follows that  $H$  is not VNM, a contradiction.

Conversely, suppose that  $H = (X, \mathbf{E})$  is such that for any  $E \in \mathbf{E}$  and any  $z \in X \setminus E$  there exists  $x_z \in E$  with  $(z, x_z) \notin I_H$ . Then, define  $\Delta^H \subseteq X \times X$  by the following rule: for any  $x, y \in X$ ,  $\{(x, y), (y, x)\} \subseteq \Delta^H$  if  $(x, y) \notin I_H$ , and  $\{(x, y), (y, x)\} \cap \Delta^H = \emptyset$  if  $(x, y) \in I_H$ : notice that  $(X, \Delta^H)$  is irreflexive. Therefore, for any  $E \in \mathbf{E}$  and any  $x, y \in E$ ,  $(x, y) \notin I_H$  hence  $(x, y) \notin \Delta^H$ , and  $E$  satisfies internal stability with respect to  $(X, \Delta^H)$ . Moreover, by hypothesis, for any  $z \in X \setminus E$  there exists  $x_z \in E$  with  $(z, x_z) \notin I_H$  i.e. in particular  $(x_z, z) \in \Delta^H$ , by definition of  $\Delta$ , hence  $E$  also satisfies external stability with respect to  $(X, \Delta^H)$ . Thus,  $\mathbf{E} \subseteq \mathcal{S}(X, \Delta^H)$  as required. ■

**Remark 11** *Notice that, as it is easily checked, the trivial clutters  $(\emptyset, \{\emptyset\})$  and  $(X, \{\emptyset\})$  with a nonempty  $X$  are both VNM. Moreover, it is worth emphasizing that the foregoing Theorem also implies that clutter  $(X, \mathbf{E})$  of Example 2 is not VNM because it is clearly not CBG-free (indeed, each of its hyperedges is a complete bigraph generator). On the contrary, its odd-cyclicity is not key to its being not VNM. To see this, consider clutter  $(\{x, y, z, u, v\}, \mathbf{E})$  with  $\mathbf{E} = \{\{x, y\}, \{y, z\}, \{z, u\}, \{u, v\}, \{v, x\}\}$ ,*

*which is also odd-cyclic (namely, there exist a positive integer  $k$ , and  $2k + 1$  distinct hyperedges  $E_i \in \mathbf{E}$  and vertices  $x_i$ ,  $i = 1, \dots, 2k + 1$  such*

that  $\{x_i, x_{i+1}\} \subseteq E_i$ ,  $i = 1, \dots, 2k$  and  $\{x_1, x_{2k+1}\} \subseteq E_{2k+1}$ . However,  $(\{x, y, z, u, v\}, \mathbf{E})$  is CBG-free hence by Theorem 10 is indeed a VNM clutter: to confirm the latter statement it is only to be checked that there is no  $E_i \in \mathbf{E}$  comprising one of the following pairs:  $\{x, z\}$ ,  $\{x, u\}$ ,  $\{y, v\}$ ,  $\{y, u\}$ ,  $\{z, v\}$ .

**Remark 12** In view of Theorem 10 it is easily checked that a remarkable class of clutters which are not VNM is provided by (nontrivial) Steiner triple systems i.e. clutters  $H = (X, \mathbf{E})$  such that  $\#X \geq 4$ ,  $\mathbf{E} \subseteq \{Y \subseteq X : \#Y = 3\}$  and for any two distinct  $x, y \in X$  there exists precisely one  $E \in \mathbf{E}$  with  $\{x, y\} \subseteq E$ . Indeed, any hyperedge  $E$  of such a clutter is a CBG: to check the latter statement, take any  $E \in \mathbf{E}$  and observe that there exist  $x \in X \setminus E$  and  $y \in E$ , and for any such  $x, y$  there exists by assumption an  $E' \in \mathbf{E}$  with  $\{x, y\} \subseteq E'$ .

Let us denote a clutter  $H = (X, \mathbf{E})$  as (strictly) VNM-complete if it is representable as the set of all VNM-stable sets of  $(X, \Delta^H)$  i.e.  $\mathbf{E} = \mathcal{S}(X, \Delta^H)$ , where  $\Delta^H = (\bigcup \mathbf{E} \times \bigcup \mathbf{E}) \setminus I_H$  as defined above in the proof of Theorem 10. Then, we have the following straightforward corollary of the previous theorem, namely

**Corollary 13** A clutter  $H = (X, \mathbf{E})$  is VNM-complete iff it is conjugation-saturated. In particular, that requires  $\bigcup \mathbf{E} = X$ .

**Proof.** Assume that  $H = (X, \mathbf{E})$  is VNM-complete and there exists  $\mathbf{E}' \supset \mathbf{E}$  such that  $H' = (X, \mathbf{E}')$  is CBG-free and  $I_H = I_{H'}$ . But then, by definition  $\mathcal{S}(X, \Delta^{H'}) = \mathcal{S}(X, \Delta^H)$ , and by Theorem 10,  $\mathbf{E}' \subseteq \mathcal{S}(X, \Delta^{H'}) = \mathcal{S}(X, \Delta^H)$ , a contradiction since by hypothesis  $\mathbf{E} = \mathcal{S}(X, \Delta^H)$ .

Conversely, let  $H = (X, \mathbf{E})$  be conjugation-saturated, and suppose that it is not VNM-complete, namely  $\mathbf{E} \subset \mathcal{S}(X, \Delta^H)$ . Then, by Theorem 10 there exists a CBG-free  $H' = (X, \mathbf{E}')$  such that  $\mathbf{E} \subset \mathbf{E}' = \mathcal{S}(X, \Delta^H)$ . Since  $H$  is conjugation-saturated,  $I_H \neq I_{H'}$  hence by definition  $I_H \subset I_{H'}$ . But then, there exist  $E' \in \mathcal{S}(X, \Delta^H) \setminus \mathbf{E}$  and  $x, y \in E' \setminus \bigcup \mathbf{E}$  whence -by definition of  $\Delta^H$ -  $x\Delta^H y$  which contradicts internal stability of  $E'$  with respect to  $(X, \Delta^H)$ .

Finally, observe that if  $X \setminus \bigcup \mathbf{E} \neq \emptyset$  then  $(x, y) \notin I_H$  for any  $x \in X \setminus \bigcup \mathbf{E}$  and  $y \in X \setminus \{x\}$  whence by definition  $\{x\} \in \mathcal{S}(X, \Delta^H) \setminus \mathbf{E}$  and thus  $(X, \mathbf{E})$  is not VNM-complete. ■

As an example, observe that clutter  $H = (\{x, y, z\}, \{\{x, y\}\})$  where  $x \neq y \neq z \neq x$  is clearly CBG-free but *not* conjugation-saturated since clutter  $H' = (\{x, y, z\}, \{\{x, y\}, \{z\}\})$  is such that  $I_{H'} = \{(x, y), (y, x)\} = I_H$ . Thus, the foregoing Corollary entails that  $H'$  is not VNM-complete.

On the contrary, for any set  $X$  with  $\#X \geq 3$  and any  $x \in X$  and positive integer  $1 < r < \#X$  take a ‘minimal’  $r$ -uniform star-clutter with centre  $x$ , i.e. a clutter  $H_x^* = (X, \mathbf{E})$  such that  $\bigcup \mathbf{E} = X$ ,  $\#E = r$  for every  $E \in \mathbf{E}$ , and for any two distinct  $E, E' \in \mathbf{E}$ ,  $E \cap E' = \{x\}$ : by construction,  $H_x^*$  is CBG-free since for any two distinct  $E, E' \in \mathbf{E}$  and any  $y \in E \setminus \{x\}$ ,  $z \in E' \setminus \{x\}$  it turns out that  $\{(y, z), (z, y)\} \cap I_{H_x^*} = \emptyset$ . Moreover,  $H_x^*$  is also a conjugation-saturated clutter since for any CBG-free clutter  $H' = (X, \mathbf{E}')$  with  $\mathbf{E}' \supset \mathbf{E}$ , and any  $E' \in \mathbf{E}' \setminus \mathbf{E}$  there exist two distinct  $E_1, E_2 \in \mathbf{E}$  and  $y \in E_1 \setminus E_2$ ,  $z \in E_2 \setminus E_1$  such that  $\{y, z\} \subseteq E'$ , hence  $\{(y, z), (z, y)\} \subseteq I_{H'}$  while  $\{(y, z), (z, y)\} \cap I_{H_x^*} = \emptyset$  by construction of  $H_x^*$ . Therefore, Corollary 13 implies that  $H_x^*$  is indeed VNM-complete.

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