

The Socially Stable Core in Structured Transferable Utility Games

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Abstract

We consider cooperative games with transferable utility (TU-games), in which we allow for a social structure on any coalition, for instance a network, a hierarchical ordering or a dominance relation. For every coalition the relative strength of a player within that coalition is induced by its social structure and is measured by a power function. We call a payoff vector socially stable if there is a collection of coalitions that can sustain it and at which all players have the same power. A payoff vector is called economically stable if it belongs to the core of the superadditive cover of the game. The socially stable core consists of the socially and economically stable payoff vectors.

We show that the socially stable core is non-empty if the game itself is socially stable. The socially stable core consists of a finite number of faces of the set of economically stable payoff vectors. Generically, it consists of a finite number of payoff vectors. Convex TU-games have a non-empty socially stable core, irrespective of the underlying social structure. When the game is permutationally convex, the socially stable core is shown to be non-empty if the power vectors are permutationally consistent and is shown to contain a unique element if the power vectors are permutationally compatible. We demonstrate the usefulness of the concept of the socially stable core by applying it to structured hierarchy games. We also present applications concerning sequencing games and the distribution of water.

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1 Introduction

A situation in which a finite set of players can obtain certain payoffs by cooperation can be described by a *cooperative game with transferable utility*, or simply a TU-game. In a TU-game players only differ with respect to their position in the game. Examples of models in which players not only differ with respect to their position in the game, but also are part of some relational structure (which possibly affects the cooperation possibilities or payoff distributions) are *games in coalition structure* and *games with limited communication structure*. In games with coalition structure it is assumed that the set of players is partitioned into disjoint sets which represent social groups. For a particular player it is more easy to cooperate with players in his own group than to cooperate with players in other groups (see, e.g., Aumann and Drèze (1974), Owen (1977), Hart and Kurz (1983) and Winter (1989)). In games with limited communication structure the edges of an undirected graph on the set of players represent binary communication links between the players such that players can cooperate only if they are connected (see, e.g. Myerson (1977), Kalai, Postlewaite and Roberts (1978), Owen (1986) and Borm, Owen and Tijs (1992)).

Another line of research in the field of cooperative games are situations in which the players are part of some hierarchical structure such as *games with a permission structure*. In such games it is assumed that players in a TU-game are part of a hierarchical organization in which there are players that need permission from other players before they are allowed to cooperate within a coalition, see for instance Gilles, Owen and van den Brink (1992) and van den Brink and Gilles (1996). Related is also the model of Faigle and Kern (1992) who consider feasible rankings of the players. Demange (2004) provides a rationale for the fact that a hierarchical structure is a widespread organizational form in many areas.

In this paper we consider TU-games with an exogenously given social structure on any subset of the players, for instance a network, a hierarchical ordering or some dominance relation. We assume that for every coalition the underlying social structure is represented by a *power vector*, whose components reflect the strengths of the individual members of the coalition within the social structure. From the literature several methods are known to measure the power of the individual members. The power vector of a coalition is completely determined by the social structure along which the coalition is organized.

The concept of socially stable core has been introduced in Herings, van der Laan and Talman (2003) within the more general framework of structured games with non-transferable utility. For a payoff vector to be in the socially stable core, there should be neither incentives to deviate from an economic point of view, nor from a social one. A payoff vector is *economically stable* if it is feasible and undominated, i.e. when the payoff vector is in the superadditive cover core of the game, and therefore in the core whenever the game is superadditive. No player has an incentive to deviate from a core payoff vector from

an economic point of view. Socially motivated deviations do not occur when all individuals are equally powerful at the proposed payoff vector. This is formalized by considering the power vectors of all coalitions that could realize the proposed payoff vector for its members. If there is a weighted sum of these power vectors that gives all individuals the same power, then individuals are said to be equally powerful at the proposed payoff and the payoff vector is called *socially stable*. The socially stable core consists of all payoff vectors that are both economically and socially stable. In Herings, van der Laan and Talman (2003), the property of social stability for a socially structured non-transferable utility game has been defined, and games satisfying this property are referred to as socially stable NTU-games. It has been shown that a socially stable game always has a non-empty socially stable core.

In this paper we consider the class of socially structured TU-games. We show that any convex game has a non-empty socially stable core. When the game is permutationally convex for some permutation, the socially stable core is shown to contain the corresponding marginal vector if the power vectors are permutationally consistent with respect to this permutation. When the game is permutationally convex and the power vectors are permutationally compatible for some permutation, the corresponding marginal vector is the unique element of the socially stable core. In general, the socially stable core consists of a finite union of (disjoint) faces of the set of economically stable payoff vectors. Moreover, it is proved that, generically, the socially stable core consists of a finite number of elements only. We demonstrate the usefulness of the concept of the socially stable core by applying it to structured hierarchy games, in which the underlying social structure of the whole set of players is a hierarchy and the social structure of any subset of players is induced by it. For power vectors representing this hierarchy, we prove that the socially stable core consists of only one element, if the game is superadditive. Finally, we consider some applications in which the hierarchy on the players is a linear order that arises from the characteristics of the economic situation. These applications concern sequencing games and the water distribution problem.

The structure of the paper is as follows. Section 2 gives some preliminaries. In Section 3 socially structured transferable utility games are introduced as well as the associated solution concept of the socially stable core. In Section 4 the main theorems are presented. Section 5 studies structured hierarchy games, Section 6 sequencing games, and Section 7 the water distribution problem. Section 8 concludes.

2 Preliminaries

A situation in which a finite set of players can obtain certain payoffs by cooperation can be described by a *cooperative game with transferable utility*, or simply a TU-game, being a pair (N, v) , with $N = \{1, 2, \dots, n\}$ a finite set of n players and $v: 2^N \rightarrow \mathbb{R}$ a characteristic function assigning to any coalition $S \subseteq N$ of players a real number $v(S)$ as the *worth* of coalition S with $v(\emptyset) = 0$, i.e. the members of coalition S can obtain a total payoff of $v(S)$ by agreeing to cooperate. In this paper we assume that N is a fixed set of players. We therefore denote a game (N, v) shortly by its characteristic function v . A TU-game v is *superadditive* if $v(S \cup T) \geq v(S) + v(T)$ for any pair of subsets $S, T \subseteq N$ such that $S \cap T = \emptyset$. Further, a TU-game v is *convex* if $v(S \cup T) + v(S \cap T) \geq v(S) + v(T)$ for all $S, T \subseteq N$.

A payoff vector is a vector $x \in \mathbb{R}^n$ assigning payoff x_i to player $i \in N$. For a vector $y \in \mathbb{R}^n$, $y(S) = \sum_{i \in S} y_i$ denotes the sum over S of the components of y . A solution F assigns a set $F(v) \subset \mathbb{R}^n$ of payoff vectors to every TU-game v . A well-known set-valued solution is the *core*, assigning to every game v the (possibly empty) set

$$C(v) = \{x \in \mathbb{R}^n \mid x(N) = v(N) \text{ and } x(S) \geq v(S), \text{ for all } S \subseteq N\}.$$

It is well-known that the core of a TU-game v is non-empty if and only if v is balanced (see Bondareva (1963)).

A more general notion of core is obtained by considering payoff configurations. The concept of payoff configuration was used by Aumann and Maschler (1964) when introducing the notion of bargaining set, see also e.g. Owen (1982) or Friedman (1989). A payoff configuration of a TU-game v is a pair (x, Π) , where x is a payoff vector and Π a partition $\{S_1, \dots, S_m\}$ of N , satisfying $x(S_j) = v(S_j)$, for all $j = 1, \dots, m$. Let PC be the set of all payoff configurations. A payoff configuration (x, Π) is undominated if $x(S) \geq v(S)$ for all $S \subseteq N$. We now may define the *superadditive cover* as the set of all payoff vectors of undominated payoff configurations. Denoting this set as $\tilde{C}(v)$, we have that

$$\tilde{C}(v) = \{x \in \mathbb{R}^n \mid \exists \Pi \text{ s.t. } (x, \Pi) \in PC \text{ and } x(S) \geq v(S), \text{ for all } S \subseteq N\}.$$

Notice that $\tilde{C}(v) = C(\tilde{v})$, where \tilde{v} is the characteristic function given by the superadditive cover of v ,

$$\tilde{v}(S) = \max_{\Pi \in \Pi|_S} \sum_{T \in \Pi} v(T),$$

with $\Pi|_S$ the collection of all partitions of S , i.e. the set of payoff vectors of undominated payoff configurations is equal to the core of its superadditive cover. When the game v is superadditive and thus $\tilde{v} = v$, it holds that $\tilde{C}(v) = C(\tilde{v}) = C(v)$. It is for this reason that the results carry over to this more general core notion.

Aumann and Drèze (1974) take a fixed partition structure, say $B \in \Pi|_N$, and define the core under partition structure B , denoted by $C(v, B)$, as the set of undominated payoff vectors satisfying $x(S) = v(S)$ for all $S \in B$. When $\tilde{v}(N) = \sum_{S \in B} v(S)$, it holds that $C(v, B) = \tilde{C}(v) = C(\tilde{v})$. For a payoff vector x in $\tilde{C}(v)$ and for any partition $B \in \Pi|_N$ satisfying $\sum_{S \in B} v(S) = \tilde{v}(N)$ we have that $x(S) = v(S)$ for all $S \in B$.

An important point-valued solution is the *Shapley value*. This value can be defined in several ways, for instance as a weighted sum of the so-called marginal contributions (see Shapley (1953)) or as an equal distribution of the so-called Harsanyi dividends of coalitions (see Harsanyi (1959)) among the players in the coalitions. Because of reasons later on in this paper, we use here the concept of marginal vector to define the Shapley value. For a permutation $\pi: N \rightarrow N$, assigning rank number $\pi(i) \in N$ to any player $i \in N$, define $\pi^i = \{j \in N | \pi(j) \leq \pi(i)\}$, i.e. π^i is the set of all players with rank number at most equal to the rank number of i , including i himself. Then the *marginal value vector* $m^\pi(v) \in \mathbb{R}^n$ of game v and permutation π is given by

$$m_i^\pi(v) = v(\pi^i) - v(\pi^i \setminus \{i\}), \quad i \in N,$$

and thus assigns to player i his marginal contribution to the worth of the coalition consisting of all his predecessors in π . The Shapley value is equal to the average of the marginal value vectors over all permutations. When v is convex, the core of v is equal to the convex hull of all marginal value vectors and thus the Shapley value is in the core.

A game v is *permutationally convex*, see Granot and Huberman (1982), if there exists a permutation π such that for all $1 \leq j < k < n$ it holds that $\max[v(S), v(\pi^j \cup S) - v(\pi^j)] \leq v(\pi^k \cup S) - v(\pi^k)$ for all $S \subset N \setminus \pi^k$. When a game v is permutationally convex with respect to the permutation π , it holds that the corresponding marginal vector $m^\pi(v)$ is in the core and, hence, the core is non-empty.

3 Structured TU-Games

Following Herings, van der Laan and Talman (2003) for games with non-transferable utility, we assume in this paper that in a transferable utility game any subset of agents, including the grand coalition, is internally organized according to some social structure. A social structure could be a communication network (e.g., see Myerson, 1977), a hierarchy with a principal at the top and a function allocating to any of the other players a unique superior such that from each such player there is a unique path to the principal, or some other more general dominance relation. On any proper coalition, the social structure could be completely induced by the social structure on the grand coalition, as is often the case in models of network formation (e.g., see Jackson and Wolinsky, 1996), but it is also allowed

that the social structure on a coalition is not induced by the social structure on the grand coalition but instead for instance is the most efficient internal organization for the coalition itself. In this paper we assume that the social structures on all the coalitions are given exogeneously and that the characteristic function v assigns to any coalition S the value $v(S)$ of the coalition S when organized according to its given social structure.

We assume that for every coalition S there is a power vector $p(S)$ that represents the strength of the position of every player $i \in S$ within the underlying social structure of coalition S . In case the social structure is given by a directed or undirected graph on the set of players, we may use one of the power measures known from the literature. A well-known power measure is the degree or score measure, see e.g. Rubinstein (1980). According to this measure the power of a node in a given graph is equal to its (out)degree, being the number of edges in the graph to which the node is incident or in case of a directed graph the number of edges leaving the node. Other power measures for directed graphs have been proposed in for instance van den Brink and Gilles (2000) and Herings, van der Laan and Talman (2005). In the latter paper, the positional power of a player in a directed graph takes into account the power of its successors and in a hierarchy the resulting power vector has the property that every player has more power than any of his successors. In case of undirected graphs, any of the so-called centrality measures may be used as a way to measure the power of a player.

A (socially) structured TU-game can therefore be described by a finite set of players, $N = \{1, \dots, n\}$, a characteristic function $v : 2^N \rightarrow \mathbb{R}$, assigning to any coalition S its value $v(S)$, and a social structure on every coalition, represented by a power function $p : \mathcal{N} \rightarrow \mathbb{R}^n$, assigning to every coalition S the power vector $p(S)$ of its players within the underlying social structure of S , where $\mathcal{N} = 2^N \setminus \{\emptyset\}$ is the collection of all non-empty subsets of N . It is assumed that both the characteristic function and the social structure (or equivalently its corresponding power vector) are exogeneously given. For mathematical reasons we take the power vector $p(S)$ in \mathbb{R}^n and not in $\mathbb{R}^{|S|}$. For a power vector $p(S)$ of coalition S it holds that $p(S)$ is a nonnegative vector in \mathbb{R}^n , $p_i(S) = 0$ for any i not in S , and $p_i(S) > 0$ for at least one player i in S . We now have the following definition of a structured TU-game (STG).

Definition 3.1 (Structured TU-Game)

A *structured TU-game* is given by the triple $\Gamma = (N, v, p)$ with N a finite set of players, v a characteristic function, and p a power function.

In a structured TU-game we are interested in payoff vectors that are *socially* and *economically stable*. If for some payoff vector $x \in \mathbb{R}^n$ and coalition $S \in \mathcal{N}$ it holds that $x(S) \leq v(S)$, then coalition S can obtain value $x(S)$ without cooperating with players

outside S and we say that coalition S sustains x . If within coalition S an individual at payoff vector x , sustained by S , has more power than some of the other individuals in S and x cannot be sustained by any other coalition, then this individual is able to increase his payoff at the expense of these other individuals. The payoff vector x is in that case not socially stable. This process can only be stopped by a credible threat of some of the other individuals to form another coalition. Such a threat is only credible if the deviating individuals can guarantee their coalition members at least the same payoff as before, and have the power to enforce the outcome that leads to these payoffs.

A payoff vector x is called socially stable if nonnegative real numbers or weights can be assigned to the coalitions S sustaining x in such a way that the weighted total power of every agent is equal to 1 and thus each agent has the same power. In the sequel, for $S \in \mathcal{N}$, let $e^S \in \mathbb{R}^n$ denote the vector given by $e_i^S = 1$ when $i \in S$ and $e_i^S = 0$ otherwise.

Definition 3.2 (Socially Stable Payoff)

For a structured TU-game $\Gamma = (N, v, p)$, a payoff vector $x \in \mathbb{R}^n$ is *socially stable* if the system of equations

$$\sum_{\{S|x(S) \leq v(S)\}} \lambda_S p(S) = e^N$$

has a nonnegative solution.

Notice that a payoff vector x satisfying $x(N) \leq v(N)$ is socially stable whenever $p(N) = e^N$. The next example shows that this may not hold when $p(N) \neq e^N$. In this example and also in later examples we often use, for ease of notation, as argument of v and p the elements of S instead of the set S itself.

Example 3.3 Take $N = \{1, 2, 3\}$, $v(1) = v(2) = v(3) = 0$, $v(1, 2) = 4$, $v(1, 3) = v(2, 3) = 2$ and $v(1, 2, 3) = 5$, and $p(i) = e^{\{i\}}$ for $i = 1, 2, 3$, $p(1, 2) = (2, 1, 0)^\top$, $p(1, 3) = (2, 0, 1)^\top$, $p(2, 3) = (0, 1, 2)^\top$ and $p(N) = (3, 2, 1)^\top$. Consider the payoff vector $x = (2, 2, 1)^\top$. Then $x(S) \leq v(S)$ for $S = \{1, 2\}$ and $S = N$. Clearly, e^N is not a nonnegative combination of $p(1, 2)$ and $p(N)$, so x is not socially stable. Observe that for any S sustaining this x it holds that $p_1(S) > p_2(S) > p_3(S)$. When we increase at x the payoff to player 1 at the expense of for instance player 2 until $y = (3, 1, 1)^\top$, then $y(S) \leq v(S)$ for $S = \{1, 2\}$, $S = \{2, 3\}$ and $S = N$. Now e^N is a nonnegative linear combination of $p(1, 2)$, $p(2, 3)$ and $p(N)$ and thus y is socially stable. In fact e^N is a nonnegative linear combination of only $p(1, 2)$ and $p(2, 3)$. As a result any payoff vector x satisfying $x(S) \leq v(S)$ for $S = \{1, 2\}$ and $S = \{2, 3\}$ is socially stable, for instance $x = (4, 0, 2)^\top$. \square

In the sequel it will be useful to define stability of a collection of coalitions without reference to a particular payoff vector.

Definition 3.4 (Stable Collection of Coalitions)

A collection of coalitions in \mathcal{N} , $\{S_1, \dots, S_k\}$, is *stable* if the system of equations

$$\sum_{j=1}^k \lambda_j p(S_j) = e^N$$

has a nonnegative solution. A stable collection of coalitions in \mathcal{N} is *minimal* if no proper subset of it is stable.

A socially stable payoff vector is therefore a payoff vector whose components can be achieved by every element of some stable collection of coalitions for its members. Observe that stability of a collection of coalitions reduces to balancedness of the collection when for every $S \in \mathcal{N}$ it holds that $p(S) = e^S$. In this case every member of a coalition has the same power. This might happen when for example the social structure of the grand coalition is a network connecting each pair of agents. In this way stability can be seen as a generalization of balancedness. In the remaining of this paper we denote the particular case of balancedness by $p = e$.

A socially stable payoff vector may not be feasible. A payoff vector x is said to be feasible if $x(N) \leq \tilde{v}(N)$, i.e. when the total payoff can be attained by cooperating according to some partition of the grand coalition. In Example 3.3 the socially stable vector $(4, 0, 2)^\top$ is not feasible, because $x(N) = 6 > 5 = \tilde{v}(N) = v(N)$ (because the game is superadditive). Furthermore, social stability of a payoff vector x does not imply that x is undominated, i.e. there may exist an $S \in \mathcal{N}$ and $y \in \mathbb{R}^n$ satisfying $y(S) \leq v(S)$ and $y_i > x_i$ for all $i \in S$. A payoff vector that is both feasible and undominated is called economically stable.

Definition 3.5 (Economically Stable Payoff)

For a structured TU-game $\Gamma = (N, v, p)$, a payoff vector $x \in \mathbb{R}^n$ is *economically stable* if $x(N) = \tilde{v}(N)$ and $x(S) \geq v(S)$ for all $S \in \mathcal{N}$.

The definition says that a payoff vector x is economically stable if and only if x is in the superadditive cover core $\tilde{C}(v)$ of v . The set of all socially and economically stable payoff vectors is called the socially stable core of the game.

Definition 3.6 (Socially Stable Core)

The *socially stable core* of a structured TU-game $\Gamma = (N, v, p)$, to be denoted $SC(v, p)$ for N , consists of the set of socially and economically stable payoff vectors of Γ .

A payoff vector x therefore lies in the socially stable core if and only if x is feasible and undominated (economic stability) and can be sustained by a socially stable collection of

coalitions (social stability). We want to stress that feasibility requires that $x(N) \leq \tilde{v}(N)$, but that sustainability is defined with respect to v . We conclude this section with two examples. The first example shows that a core element does not need to be socially stable.

Example 3.7 Consider the game (N, v, p) with $N = \{1, 2, 3\}$, $v(1) = v(2) = v(3) = 0$, $v(1, 2) = 1$, $v(1, 3) = 2$, $v(2, 3) = 3$ and $v(N) = 6$. The social structure on N is given by the hierarchy such that agent 1 dominates agent 3 and agent 3 dominates agent 2, e.g. $p_1(S) = 3$ when $1 \in S$, $p_2(S) = 1$ when $2 \in S$ and $p_3(S) = 2$ when $3 \in S$. Clearly, v is convex and therefore superadditive and thus $\tilde{v}(N) = v(N)$. The core is given by

$$C(v) = \{x \in \mathbb{R}_+^3 \mid x_1 + x_2 + x_3 = 6, 0 \leq x_1 \leq 3, 0 \leq x_2 \leq 4, 0 \leq x_3 \leq 5\}.$$

Take any payoff vector x in the relative interior of $C(v)$. Such a vector can only be sustained by the grand coalition N . Since $p^N \neq e^N$, the collection $\{N\}$ is not stable and therefore x is not an element of the socially stable core. As can easily be seen, the only socially stable element in the core is $(3, 0, 3)^\top$, sustained by the stable collection $\{\{2\}, \{2, 3\}, N\}$. In this socially stable core element the weakest agent, agent 2, gets his own value, the middleman, agent 3, gets his marginal contribution when he joins the weakest agent, and the strongest agent, agent 1, gets his marginal contribution when he joins the coalition of the weakest agent and the middleman. \square

The next example considers a game with four players, which is not superadditive. In particular it stresses the difference between feasibility and sustainability by the grand coalition N .

Example 3.8 Take (N, v, p) with $N = \{1, 2, 3, 4\}$, $v(1, 2) = v(3, 4) = 2$, $v(S) = 2$ when $|S| = 3$, $v(N) = 3$ and $v(S) = 0$ otherwise; and $p(1, 2) = (2, 1, 0, 0)^\top$, $p(3, 4) = (0, 0, 1, 2)^\top$ and $p(S) = e^S$ for all other S . Clearly $\tilde{v}(N) = 4$ and the set of undominated feasible payoff vectors is given by

$$\tilde{C}(v) = C(\tilde{v}) = \{x \in \mathbb{R}_+^4 \mid x_1 + x_2 = 2, x_3 + x_4 = 2\}.$$

Now take any strictly positive vector $x \in \tilde{C}(v)$. Then $x(S) > v(S)$ for all $S \neq \{1, 2\}, \{3, 4\}$. Since e^N cannot be written as a nonnegative linear combination of $p(1, 2)$ and $p(3, 4)$, x is not socially stable. In this case, the requirement of social stability removes all strictly positive vectors from $\tilde{C}(v)$. It then easily follows that the unique element of $SC(v, p)$ is the payoff vector $x = (2, 0, 0, 2)^\top$. For this payoff vector it holds that $x(S) \leq v(S)$ when S is $\{2\}, \{3\}, \{1, 2\}$ and $\{3, 4\}$. Since e^N is a nonnegative linear combination of $p(1, 2)$, $p(3, 4)$, $p(2)$ and $p(3)$, it holds that x is sustained by the stable collection $\{\{2\}, \{3\}, \{1, 2\}, \{3, 4\}\}$ and thus $x \in SC(v, p)$. Observe that for instance $x = (0, 2, 0, 2)^\top$ is not socially stable, because $p_1(1, 2) > p_2(1, 2)$.

Notice that not any $x \in \tilde{C}(v)$ is socially stable although $p(N) = e^N$. The reason is that for each $x \in \tilde{C}(v)$ it holds that $x(N) = \tilde{v}(N) > v(N)$ and thus x is not sustained by N . To obtain total payoff equal to 4, N has to split itself in $\{1, 2\}$ and $\{3, 4\}$, so x is sustained by these two coalitions. However, as shown above, the collection of these two coalitions is not socially stable. \square

4 Non-emptiness of the Socially Stable Core

In this section we give sufficient conditions for the non-emptiness of the socially stable core of a structured TU-game (N, v, p) . When $p = e$, we always have that $SC(v, p) = \tilde{C}(v)$. In case v is also superadditive and thus $SC(v, p) = \tilde{C}(v) = C(v)$, we know from Bondareva (1963) that the socially stable core is non-empty if and only if the game is balanced. However, for an arbitrary power function p and arbitrary game v , the socially stable core is a subset of the superadditive cover core and might be empty even if the game is superadditive and balanced. The balancedness condition for superadditive TU-games is not sufficient for the non-emptiness of the socially stable core when $p \neq e$. The next definition of social stability of the game has been given in the NTU-context in Herings, van der Laan and Talman (2003).

Definition 4.1 (Socially Stable Game)

A structured TU-game $\Gamma = (N, v, p)$ is *socially stable* if every socially stable payoff x of Γ is feasible.

It should be observed that this social stability condition reduces to the usual balancedness condition of a TU-game when $p = e$ and the game is superadditive. We now have the following theorem.

Theorem 4.2

A structured TU-game $\Gamma = (N, v, p)$ has a non-empty socially stable core if Γ is socially stable.

Proof. The proof follows from a more general theorem for socially structured NTU-games given in Herings, van der Laan and Talman (2003). \square

The theorem above requires feasibility to be shown for any socially stable payoff vector. This may be a demanding task. On the other hand, social stability is sufficient but not necessary. The next theorem says that for every given power function a convex game has a non-empty socially stable core. Observe that the proof holds for any game whose core contains all marginal vectors. However, this does not generalize the result, because a

game is convex iff the core contains all marginal vectors, see e.g. Ichiishi (1981). Observe that convexity implies superadditivity, so $x(N) = v(N)$ for any undominated payoff vector.

Theorem 4.3

If v is a convex game, then for every power function p the structured TU-game $\Gamma = (N, v, p)$ has a non-empty socially stable core.

Proof. To prove non-emptiness, we first construct for a given power function p a stable collection of coalitions.

Step 1. Set $k = 1$, $S_k = N$, $q^k = e^N$, and $r_k = n$. Goto Step 2.

Step 2. Define $T_k = \{j \in S_k \mid p_j(S_k)/q_j^k = \max_{h \in S_k} p_h(S_k)/q_h^k\}$ and $t_k = |T_k|$. Define $\lambda_k = (\max_{h \in S_k} p_h(S_k)/q_h^k)^{-1}$. Goto Step 3.

Step 3. For $j \in T_k$, define $\pi(j) \in N$ such that $\{\pi(j) \mid j \in T_k\} = \{r_k, r_k - 1, \dots, r_k - t_k + 1\}$. If $r_k = t_k$, define $k^* = k$ and stop the procedure; otherwise set $k = k + 1$ and goto Step 4.

Step 4. Set $S_k = S_{k-1} \setminus T_{k-1}$, $q^k = q^{k-1} - \lambda_k p(S_{k-1})$ and $r_k = r_{k-1} - t_{k-1} = |S_k| > 0$. Return to step 2.

By construction we have that the collection $\{S_1, S_2, \dots, S_{k^*}\}$ is a stable collection of coalitions and that $\pi = (\pi(1), \pi(2), \dots, \pi(n))$ is a permutation of the elements of N such that for any $k = 1, \dots, k^*$ it holds that

$$S_k = \{\pi(1), \pi(2), \dots, \pi(\ell_k)\}, \text{ with } \ell_k = \sum_{h=k}^{k^*} t_h.$$

Next take the payoff vector x equal to the marginal vector $m^\pi(v)$. Then it follows for any $k = 1, \dots, k^*$ that

$$\sum_{j \in S_k} x_j = \sum_{j=1}^{\ell_k} (v(\pi^j) - v(\pi^j \setminus \{j\})) = v(\pi^{\ell_k}) = v(S_k).$$

By construction of the sets S_k it follows that x is socially stable. Moreover, since v is convex we also have that $x \in C(v)$. Hence $x \in SC(v, p)$. \square

Observe that the marginal vector constructed in the proof is unique if and only if $k^* = n$ and thus $|T_k| = 1$ for all k . When for some k , T_k contains multiple players, we can take any order of the players within T_k in Step 3. So, in general there are $\prod_k |T_k|$ different

permutations satisfying the conditions and inducing a marginal vector in $SC(v, p)$. We, however, remark that different permutations may induce the same marginal vector.

The assumptions in Theorems 4.2 and 4.3 are independent, i.e. socially stability of Γ does not imply convexity of v and vice versa. Clearly, when $p = e$, Γ is socially stable if and only if v satisfies balancedness. However, balancedness of v does not imply convexity. The following example shows that convexity of v does not imply social stability of Γ .

Example 4.4 Take $N = \{1, 2, 3\}$, $v(1) = v(3) = v(1, 3) = 0$, $v(2) = v(1, 2) = v(2, 3) = v(1, 2, 3) = 1$. Take any power vector function p such that $p(1, 2) = (2, 1, 0)^\top$ and $p(2, 3) = (0, 1, 2)^\top$. This game is convex, so the socially stable core of $\Gamma = (N, v, p)$ is non-empty. The payoff vector $x = (1, 0, 1)^\top$ is socially stable, being sustained by the stable collection $\{\{1, 2\}, \{2, 3\}\}$. However, x is not feasible and thus $\Gamma = (N, v, p)$ is not socially stable. \square

Theorem 4.3 states that for convex games v the socially stable core is non-empty for any power function p . When we make joint assumptions on p and v , it is possible to weaken the assumptions on the characteristic function v . We first define the notion of π -consistency.

Definition 4.5 A power function $p : \mathcal{N} \rightarrow \mathbb{R}^n$ is π -consistent for a permutation π on N , when for all coalitions S and for all players i and j in S it holds that $\pi(i) < \pi(j)$ implies $p_i(S) \leq p_j(S)$.

When the power function p is π -consistent, the ranking of the players determined by permutation π is consistent with the ranking of the power of players in any coalition. We will show that the marginal vector $m^\pi(v)$ belongs to the socially stable core of $\Gamma = (N, v, p)$, when v is permutationally convex for the permutation π and p is π -consistent.

Theorem 4.6 Consider a structured TU-game $\Gamma = (N, v, p)$, where, for some permutation π on N , v is permutationally convex and p is π -consistent. Then the socially stable core of Γ contains the vector $m^\pi(v)$ as an element.

Proof. As in the proof of Theorem 4.3 it can be constructively shown that the collection $\{\pi^i \mid i \in N\}$ is socially stable. Notice that some of the weights might be zero, so the collection is not necessarily minimal. Take the payoff vector x equal to the marginal vector $m^\pi(v)$. Then it follows for any $i \in N$ that

$$\sum_{j \in \pi^i} x_j = \sum_{j=1}^i (v(\pi^j) - v(\pi^j \setminus \{j\})) = v(\pi^i).$$

It follows that x is socially stable. Moreover, since v is permutationally convex for the permutation π , it follows from Granot and Huberman (1982) that $x \in C(v)$. Hence $x \in$

$SC(v, p)$. □

Of course, also under the conditions of Theorem 4.6 the socially stable core may contain multiple elements. Clearly, when $p = e$ the power function is π -consistent for any π and $SC(v, p) = \tilde{C}(v)$. However, when the power function p is such that for some permutation π a player i has little power in any coalition involving players from $N \setminus \pi^i$, the socially stable core can be shown to consist of a unique element given by the marginal vector $m^\pi(v)$. To make this statement more precise, we introduce the notion of π -compatibility.

Definition 4.7 A power function $p : \mathcal{N} \rightarrow \mathbb{R}^n$ is π -compatible for a permutation π of N , when for all players $i \in N$ and for all coalitions S containing i such that $S \setminus \pi^i \neq \emptyset$, it holds that $p_i(S) < \sum_{j=1}^n p_j(S)/n$.

When a power function is π -compatible, the power of a player i in any coalition that involves another player that is ranked higher according to π , is less than the average power $\sum_{j=1}^n p_j(S)/n$. Notice that neither π -compatibility implies π -consistency nor π -consistency implies π -compatibility.

Theorem 4.8 Consider a structured TU-game $\Gamma = (N, v, p)$, where for some permutation π of N , v is permutationally convex and p is π -compatible. Then the socially stable core of Γ contains the marginal vector $m^\pi(v)$ as its unique element.

Proof. Without loss of generality, we may normalize power functions such that $\sum_{j=1}^n p_j(S) = n$ and we may assume that the permutation π corresponds to the ordering $\pi(i) = n + 1 - i$, $i = 1, \dots, n$. Let the payoff vector x belong to $SC(v, p)$ and let $\{S_1, \dots, S_m\}$ be a stable collection of coalitions with $(\lambda_1, \dots, \lambda_m)$ a vector of nonnegative weights such that $\sum_{j=1}^m \lambda_j p(S_j) = e^N$ and $x(S_j) \leq v(S_j)$ for $j = 1, \dots, m$. Notice that $\sum_{j=1}^m \lambda_j = 1$.

We define the ordering \prec^ℓ on \mathcal{N} by $S \prec^\ell T$ if and only if the lowest ranked individual in $S \cup T$ not in $S \cap T$ belongs to S . Without loss of generality, we may choose $\{S_1, \dots, S_m\}$ to be minimal, and we can order the coalitions such that $S_j \prec^\ell S_{j+1}$.

We claim that, for $i = 1, \dots, m$, this stable collection satisfies $i \in S_i \subset \{i, \dots, n\}$. Suppose that $\{S_1, \dots, S_m\}$ does not contain the singleton coalition $\{n\}$. Since p is π -compatible, it follows that $p_n(S_j) < 1$, $j = 1, \dots, m$, so

$$1 = \sum_{j=1}^m \lambda_j p_n(S_j) < 1,$$

a contradiction. Consequently, $\{S_1, \dots, S_m\}$ does contain the singleton coalition $\{n\}$, and by the properties of \prec^ℓ it follows that $S_m = \{n\}$.

We now use an induction argument to proceed. Assume it is true that, for some $k' \leq m$, for $k = 1, \dots, k'$, $n - k \in S_{m-k} \subset \{n - k, \dots, n\}$. We show that then $n - k - 1 \in$

$S_{m-k-1} \subset \{n-k-1, \dots, n\}$. Obviously, it is not the case that $S_{m-k-1} \subset \{n-k, \dots, n\}$ as this would violate the minimality of $\{S_1, \dots, S_m\}$. Suppose it is not true that $S_{m-k-1} \subset \{n-k-1, \dots, n\}$, so the lowest ranked player in S_{m-k-1} is $i' < n-k-1$. For all S_j , it holds that $p_{n-k-1}(S_j) < 1$, so $1 = \sum_{j=1}^m \lambda_j p_{n-k-1}(S_j) < 1$, a contradiction. It follows that the lowest ranked player in S_{m-k-1} is $n-k-1$. This completes the proof of the induction step, and it follows as a corollary that $m = n$.

It remains to be shown that $x = m^\pi(v)$. We denote $m^\pi(v)$ by y .

Obviously, it holds that $x_n = y_n = v(\{n\})$. We proceed with an induction argument. Assume it is true that, for some k' , $x_i = y_i$ for $i = n-k', \dots, n$, and $\sum_{i=n-k'}^n x_i = v(\{n-k', \dots, n\})$. We will show that $x_{n-k'-1} = y_{n-k'-1}$.

Since x is economically stable, it follows that $\sum_{i=n-k'-1}^n x_i \geq v(\pi^{k'+1})$. From the induction hypothesis that $y_i = x_i$ for $i = n-k', \dots, n$, we then obtain $x_{n-k'-1} \geq y_{n-k'-1}$. Since y is economically stable, it follows that $\sum_{i \in S_{n-k'-1}} y_i \geq v(S_{n-k'-1}) = \sum_{i \in S_{n-k'-1}} x_i$. Since $n-k'-1 \in S_{n-k'-1} \subset \pi^{k'+1}$, it follows that $x_{n-k'-1} \leq y_{n-k'-1}$. We have shown that $x_{n-k'-1} = y_{n-k'-1}$. \square

When the power function p is π -compatible, any player has so much power compared to his lower-ranked players that he is able to extract all payoffs from them, up to the point where the lower-ranked players could form a deviating coalition. The socially stable core consists of a unique element, corresponding to the marginal vector $m^\pi(v)$. We come back to this property in the next sections.

In general the socially stable core consists of a finite number of faces of $\tilde{C}(v)$. To show this, for a collection of coalitions $\mathcal{F} \subset \mathcal{N}$, define

$$C^{\mathcal{F}}(v) = \{x \in \tilde{C}(v) \mid x(S) = v(S) \text{ for all } S \in \mathcal{F}\},$$

i.e., $C^{\mathcal{F}}(v)$ is a (possibly empty) face of $\tilde{C}(v)$. Clearly, when $x \in SC(v, p)$ lies in the (relative) interior of $C^{\mathcal{F}}(v)$, then due to the linearity of the constraints every point of the face $C^{\mathcal{F}}(v)$ belongs to $SC(v, p)$. The next example shows a socially stable core consisting of two disjoint faces.

Example 4.9 Take $N = \{1, 2, 3\}$, $v(1) = v(2) = v(3) = 0$, $v(1, 2) = 4$, $v(1, 3) = 1$, $v(2, 3) = 3$, $v(1, 2, 3) = 6$. Since the game is superadditive, $\tilde{C}(v) = C(v) = \{x \in \mathbb{R}^3 \mid x_1 + x_2 + x_3 = 6, 0 \leq x_1 \leq 3, 0 \leq x_2 \leq 5, 0 \leq x_3 \leq 2\}$. Take $p(1, 2) = (3, 1, 0)^\top$, $p(1, 3) = (2, 0, 1)^\top$, $p(2, 3) = (0, 1, 3)^\top$, $p(1, 2, 3) = (2, 3, 1)^\top$. For this power function the socially stable core $SC(v, p)$ consists of two zero-dimensional faces of the core, one being the vertex $(3, 1, 2)^\top$ and the other one being the marginal vector $(1, 5, 0)^\top$. The vector $(3, 1, 2)^\top$ is sustained by the stable collection $B_1 = \{\{1, 2, 3\}, \{1, 2\}, \{2, 3\}\}$ and the vector $(1, 5, 0)^\top$ by the stable collection $B_2 = \{\{1, 2, 3\}, \{1, 3\}, \{3\}\}$. \square

The example shows the generic case that the socially stable core consists of zero-dimensional faces of $\tilde{C}(v)$ and thus consists of a finite number of payoff vectors. To prove this, observe that once the number of players is fixed, a structured TU-game is completely determined by the tuple of payoffs v , which can be represented by a vector in \mathbb{R}^{2^n-1} , and the tuple of power functions, which can be represented by a vector in $S^{(2^n-1)^n}$, the unit simplex in $\mathbb{R}^{(2^n-1)^n}$. The Euclidean topology and Lebesgue measure on $\mathbb{R}^{2^n-1} \times S^{(2^n-1)^n}$ therefore induce a topology and a measure on structured TU-games.

Theorem 4.10 *Let N be the set of players. Then there is an open set of payoffs and power functions with full Lebesgue measure $V \times P$ such that for any $(v, p) \in V \times P$, the socially stable core of the structured TU-game $\Gamma = (N, v, p)$ is either empty or consists of a finite number of elements.*

Proof. Define the closed subset W of \mathbb{R}^n with measure zero by

$$W = \cup_{(S_1, \dots, S_{n-1}) \in \mathcal{N}^{n-1}} \text{span} [e^{S_1}, \dots, e^{S_{n-1}}].$$

Next we define the open subset V of \mathbb{R}^{2^n-1} with full measure by

$$V = \{v \in \mathbb{R}^{2^n-1} \mid \forall (S_1, \dots, S_n) \in \mathcal{N}^n \text{ with } S_j \neq S_{j'} \text{ when } j \neq j', (v_{S_1}, \dots, v_{S_n}) \notin W\}.$$

Finally, we define the open subset P of $S^{(2^n-1)^n}$ with full measure as the set of vectors $p = (p(S))_{S \in \mathcal{N}}$ for which it holds that any selection of n vectors from the vectors $p(S)$, $S \in \mathcal{N}$, and e^N yields an independent set of vectors.

We now examine the socially stable core for the structured TU-game $\Gamma = (N, v, p)$, where $(v, p) \in V \times P$. All socially stable core elements are obtained by considering, for all minimal stable collections of coalitions $\{S_1, \dots, S_m\}$, the solutions to the system of equations

$$x(S_j) = v(S_j), \quad j = 1, \dots, m.$$

In fact, the union over all stable collections of solutions to the corresponding system is a superset of the socially stable core. Since $\{S_1, \dots, S_m\}$ is a stable collection, there exists a vector of nonnegative weights λ such that

$$\sum_{j=1}^m \lambda_j p(S_j) = e^N.$$

Moreover, $\{S_1, \dots, S_m\}$ is minimal, so that the vectors $p(S_j)$ are independent, and in particular $m \leq n$. Since $p \in P$, it holds that $m = n$.

Consider the system of equations,

$$x(S_j) = v(S_j), \quad j = 1, \dots, n.$$

If the vectors $e^{S_j}, j \in N$, are all independent, it follows that this system has exactly one solution, and therefore we obtain at most one socially stable core element. When the vectors $e^{S_j}, j \in N$, are not all independent, it follows from the definition of V that this system of equations has no solution. Hence, there is at most one solution for each minimal stable collection. Since the number of minimal stable collections is finite, this proves the theorem. \square

Theorem 4.10 shows that in general the socially stable core refines the superadditive cover core to a great extent. There is typically only a finite number of payoff vectors in the socially stable core.

5 Structured hierarchy games

A hierarchical structure is a widespread organizational form in many areas. A hierarchy on N is represented by a tree on N . A tree is a directed graph on N such that there is one top-player, say player 1, and for each player $i \neq 1$ there is a unique path of directed edges $(i_1, i_2), \dots, (i_{j-1}, i_j)$ such that $i_1 = 1$ and $i_j = i$. Following Myerson (1977) we assume that players can't cooperate when they are not connected. This yields the so-called restricted game v^r assigning to any coalition S the sum of the worths of the components of S in the graph, i.e.

$$v^r(S) = \sum_{T \in C(S)} v(T),$$

where $C(S)$ is the collection of maximally connected subsets of S in the graph. Clearly in a hierarchy graph we have that a subset S of nodes is connected if and only if there is a unique player $j \in S$ such that for any other player $i \in S$ there is a unique path of directed edges in the graph from j to i and all players on each of these paths are also in S . From Kaneko and Wooders (1982) and Le Breton, Owen and Weber (1992) it follows that $C(\tilde{v}^r)$ is not empty, i.e., the core of the restricted game of the superadditive cover of v is not empty, see also Demange (1994), who shows that $C(v^r)$ is not empty if v is superadditive. From this Demange (2004) argues that hierarchies yield stability, providing a rationale for the fact that a group organizes itself in hierarchies so as to achieve coordination.

In this section we consider structured hierarchy games. In such games the social structure on the set N of agents is given by a hierarchy on N and the social structure on any subset S is induced by this hierarchy. The characteristic function reflects the hierarchy, in the sense that $v = v^r$. The power function reflects the hierarchy in the sense that a player has a higher power than any of its subordinates. To be more precise, for $i \in N$, let $D(i)$ denote the set of subordinates of player i , i.e. for each $j \in D(i)$ there is a unique directed path from i to j . Then the power vector for the whole set of players is a vector

$p(N)$ such that for each i and j it holds that $p_i(N) > p_j(N)$ whenever $j \in D(i)$, i.e., the top player 1 has the highest power and the power is strictly decreasing when descending in the hierarchy. A natural example is $p_i(N) = |D(i)| + 1$ for all $i \in N$, i.e., the power of a player is equal to the number of its subordinates including itself. Further, for any $S \subset N$ and $i \in S$ we take $p_i(S) = p_i(N)$, although it is sufficient assume that $p_i(S) > p_j(S)$ whenever $j \in D(i)$ and $i \in S$. We now have the following theorem.

Theorem 5.1 *Consider a structured hierarchy TU-game $\Gamma = (N, v, p)$ where v is super-additive. Then the socially stable core of Γ contains the vector x^* , with*

$$x_j^* = v(D(j) \cup \{j\}) - v(D(j)), \text{ for all } j \in N, \quad (1)$$

as its unique element.

Proof. Let π be any permutation on N such that for all i and j it holds that $\pi(j) < \pi(i)$ when $p_j(N) < p_i(N)$. Clearly, the collection $\{\pi^1, \dots, \pi^n\}$ is stable. In Demange (2004) it has been shown that $m^\pi(v) \in C(v)$. Since $v = v^r$, we have that $x^* = m^\pi(v)$, being sustained by the stable collection $\{\pi^1, \dots, \pi^n\}$. Therefore x^* is an element of the socially stable core. So, it remains to be proven that x^* is the unique socially stable element of the core for any power function p reflecting the hierarchy. Let x be an arbitrary element of the socially stable core and let $B = \{S_1, \dots, S_m\}$ be a stable collection sustaining x . For $i \neq 1$, let $b(i)$ be the unique predecessor of i in the hierarchy and define $b(1) = 0$. Recall that player 1 is the topman of the hierarchy. Since $p_{b(i)}(S) > p_i(S)$ for any S with $b(i) \in S$, for every player $i \in N$ it must hold that there exists an index $h(i) \in \{1, \dots, m\}$ such that $i \in S_{h(i)}$ and $b(i) \notin S_{h(i)}$. Since $D(i)$ is the number of subordinates of player $i \in N$, we have that $|D(i)| = 0$ when i is a leave (being a node without subordinates), $|D(1)| = n - 1$, and $|D(j)| > |D(i)|$ for all j and $i \in D(j)$. We proceed now by induction on $|D(i)|$. Let node i be a leave, i.e., $D(i) = \emptyset$. Notice that in a hierarchy such a node always exists. Then $\{i\}$ is a component of $S_{h(i)}$ in the graph, since $b(i) \notin S_{h(i)}$. From (i) x is sustained by $S_{h(i)}$, (ii) x is in the core and (iii) $v = v^r$ it follows that $x_i = v(i)$, and thus $x_i = x_i^*$ when $|D(i)| = 0$.

Now, let $x_i = x_i^*$ for any player $i \in N$ with $|D(i)| \leq k$ for some k , $0 \leq k < n - 1$. When there is no player $j \in N$ with $|D(j)| = k + 1$, then $x_i = x_i^*$ also holds for any player $i \in N$ with $|D(i)| \leq k + 1$. Otherwise, let $j \in N$ be any player with $|D(j)| = k + 1$. Since $S_{h(j)}$ sustains x and x is in the core, we have that $x(S_{h(j)}) = v(S_{h(j)})$. Since $v = v^r$, it holds that $v(S_{h(j)})$ is equal to the sum of the values of all the components of $S_{h(j)}$ in the graph. Together with the fact that x is in the core this implies that $x(S) = v(S)$ for any component S in the graph of $S_{h(j)}$. Let $C(j)$ be the component of $S_{h(j)}$ in the graph containing j . For any $i \in N$, define $D^0(i) = D(i) \cup \{i\}$. Since $b(j) \notin S_{h(j)}$ we have

that $C(j) \subset D^0(j)$. Moreover, because of the tree structure and the fact that $j \in C(j)$, we have that $S_{h(j)} \setminus C(j)$ is partitioned in a collection of components $\{D^0(i_1), \dots, D^0(i_l)\}$ with $i_s \in D(j)$ for all $s = 1, \dots, l$. Because of the induction hypothesis we have that $x(D^0(i_s)) = x^*(D^0(i_s)) = v(D^0(i_s))$ for all s . Hence,

$$\begin{aligned} x(D^0(j)) &= x(C(j)) + \sum_{s=1}^l x(D^0(i_s)) = v(C(j)) + \sum_{s=1}^l v(D^0(i_s)) \\ &\leq v(D^0(j)) \leq x(D^0(j)), \end{aligned}$$

where the first inequality follows from the superadditivity of v and the latter inequality from the fact that x is in the core. Hence $x(D^0(j)) = v(D^0(j))$.

To complete the proof, observe that $D(j)$ is partitioned in the graph in a collection of components $\{D^0(j_1), \dots, D^0(j_k)\}$, where $\{j_1, \dots, j_k\}$ is the set of followers of j . From the induction hypothesis we have that $x(D^0(j_s)) = x^*(D^0(j_s)) = v(D^0(j_s))$ for all s . It then follows that

$$\begin{aligned} x_j &= x(D^0(j)) - \sum_{s=1}^k x(D^0(j_s)) = v(D^0(j)) - \sum_{s=1}^k v(D^0(j_s)) \\ &= v(D(j) \cup \{j\}) - v(D(j)) = x_j^*, \end{aligned}$$

where the second last equality comes from the fact that $v = v^r$. This proves the induction step. Since the number of players is finite this completes the proof. \square

We remark that the permutation π for which $x^* = m^\pi(v)$ is in general not unique. Every permutation π that satisfies the condition in the proof yields x^* as marginal vector. The same holds for the choice of the power vectors. Every power vector $p(N)$ satisfying $p_i(N) > p_j(N)$ whenever $j \in D(i)$, yields the same socially stable core, consisting of the unique element x^* . The outcome x^* is very natural. Every individual receives a payoff equal to what he is contributing when he joins his subordinates. The core typically also contains other elements than the vector $x^* = m^\pi(v)$. In particular, following Demange (2004), we can replace the hierarchy graph by a communication graph by replacing all directed edges by undirected edges. Then we can choose arbitrarily some other player and reconstruct a new directed graph with this player as the top-player. Clearly, also the marginal vector with respect to a permutation reflecting this new directed graph is in the core of the game. So, for a given topman, the socially stable core with respect to a power function reflecting the corresponding hierarchy reduces the core containing several marginal vectors to one marginal vector. The next sections discuss some examples when the hierarchy is a linear order.

6 Sequencing games

A one-machine sequencing situation, see e.g. Curiel (1988) or Hamers (1995) is described as a triple (N, q, c) , where $N = \{1, \dots, n\}$ is the set of jobs in a queue to be processed, $q \in \mathbb{R}_+^n$ is an n -vector with q_i the processing time of job i and $c = (c_i)_{i \in N}$ is a collection of cost functions $c_i: \mathbb{R}_+ \rightarrow \mathbb{R}_+$, specifying the costs $c_i(t)$ when t is the total time needed to complete job i . For an ordering ρ on N describing the positions of the jobs in the queue, the completion time of job i is given by $T_i(\rho) = \sum_{\{j | \rho(j) \leq \rho(i)\}} q_j$, i.e. the completion time is the sum of its waiting time and its own processing time, and the costs of processing i are given by $C_i(\rho) = c_i(T_i(\rho))$. The total costs of a coalition $S \subseteq N$ given an ordering ρ are equal to $C_S(\rho) = \sum_{i \in S} C_i(\rho)$. In the sequel we assume without loss of generality that the initial positions of the jobs in the queue are given by the ordering ρ^0 with $\rho^0(i) = i$ for all $i \in N$, so that the costs of a coalition S of jobs according to ρ^0 are given by

$$C_S(\rho^0) = \sum_{i \in S} C_i(\rho^0) = \sum_{i \in S} c_i\left(\sum_{\{j | j \leq i\}} q_j\right), \quad S \subseteq N.$$

Now, each coalition S of jobs can obtain cost savings by rearranging the jobs among the members of S . Then the minimal cost of the grand coalition is given by

$$C_N = \min_{\rho} C_N(\rho).$$

However, members of any other coalition S can only rearrange their positions under the condition that the members of S are not allowed to ‘jump’ over jobs outside S . So, an ordering ρ is admissible for S if for any $j \notin S$ the set of its predecessors does not change with respect to the initial situation, i.e. if for any $j \notin S$ it holds that $\{k \in N \mid \rho(k) < \rho(j)\} = \{k \in N \mid \rho^0(k) < \rho^0(j)\} = \{k \in N \mid k < j\}$. Let $\mathcal{A}(S)$ be the set of admissible orderings for S . Then the minimal cost of S is given by

$$C_S = \min_{\rho \in \mathcal{A}(S)} C_S(\rho).$$

This gives the cost savings *sequencing* game (N, v) with N the set of jobs as the set of players and characteristic function v given by

$$v(S) = C_S(\rho^0) - C_S, \quad S \subseteq N.$$

In the following, we use the terminology players instead of jobs. Obviously, since for any S only orderings in $\mathcal{A}(S)$ are admissible, only connected coalitions (i.e. coalitions of consecutive players) can realise cost reductions. For $i < j$, denote the set $\{i, i+1, \dots, j\}$ of consecutive players by $[i, j]$ and let \mathcal{L} denote the set of all coalitions of consecutive players, i.e.

$$\mathcal{L} = \{T \in \mathcal{N} \mid T = [i, j], 1 \leq i \leq j \leq n\}.$$

For some $S \in \mathcal{N}$, let $P(S)$ be the unique minimal partition of S in coalitions of consecutive players, i.e. $T \in \mathcal{L}$ if $T \in P(S)$ and $T_1 \cup T_2 \notin \mathcal{L}$ for any pair $T_1, T_2 \in P(S)$. Then it holds that

$$v(S) = \sum_{T \in P(S)} v(T), \quad S \in \mathcal{N},$$

i.e. the value of a coalition S is equal to the sum of the values of the coalitions of consecutive players in its unique minimal partition. From this property it follows immediately that the characteristic function v is superadditive and that the game v coincides with its restricted game v^r , i.e. $v = v^r$. Moreover, v satisfies permutational convexity, implying that v has a non-empty core. More precisely, let u and ℓ be the two permutations on N defined by $u(i) = i, i = 1, 2, \dots, n$, and $\ell(i) = n + 1 - i, i = 1, 2, \dots, n$. Further, denote $\mu(v) = m^u(v)$ and $\lambda(v) = m^\ell(v)$ as the corresponding marginal value vectors. Then it immediately follows that v satisfies the permutational convexity conditions for the two permutations u and ℓ , implying that the two marginal vectors $\mu(v)$ and $\lambda(v)$ are in the core $C(v)$.

We now want to apply the results of the previous section. Therefore, observe that the set \mathcal{L} of connected coalitions equals the set of connected sets of nodes in the undirected line graph on N with the set of edges given by $\mathcal{E} = \{(i, i + 1) | i = 1, \dots, n - 1\}$. We can convert this undirected graph into a linearly ordered hierarchy by taking either player 1 or player n as topman. In the former case the hierarchy is decreasing from 1 to n , in the latter case from n to 1. Now, let p be a power function representing the hierarchy. In case player 1 is the topman we have for every $S \in \mathcal{N}$ that

$$p_i(S) < p_j(S) \quad \text{when } i > j, \quad i, j \in S \quad (2)$$

and in case player n is the topman we have that

$$p_i(S) < p_j(S) \quad \text{when } i < j, \quad i, j \in S. \quad (3)$$

The following result follows now immediately from Theorem 5.1.

Corollary 6.1

For any sequencing game it holds:

- (i) $\lambda(v)$ is the unique element of $SC(v, p)$ when p satisfies condition (2).
- (ii) $\mu(v)$ is the unique element of $SC(v, p)$ when p satisfies condition (3).

In case (i) player n gets its own value and player 1 is able to extract all gains from his cooperation with his subordinates, in case (ii) player 1 gets its own value and player n is the topman who is able to extract all gains from cooperation. To discuss these results, we now consider the special case of linear costs, i.e., $c_i(t) = \alpha_i t$ for all $t \geq 0$ with $\alpha_i > 0$. It is well-known that in this case the characteristic function v is convex, see e.g. Curiel (1988)

or Hamers (1995). Moreover, in these references it has been shown that under linear costs for each coalition $[i, j] \in \mathcal{L}$ it holds that

$$v([i, j]) = \sum_{\{h, k \in [i, j] | h < k\}} g_{hk},$$

where $g_{hk} = \max(0, \alpha_k q_h - \alpha_h q_k)$ is the gain of a switch between player h and k in any ordering such that player h is directly in front of k . Now, the net-costs of player h resulting from marginal vector $\mu(v)$ are given by the costs of the waiting time in the initial order minus the savings obtained from cooperation, i.e., the net-costs $c_h^u(N, q, c)$ of player $h \in N$ in the linear cost sequencing situation (N, q, c) induced by $m^u(v)$ is given by

$$\begin{aligned} c_h^u(N, q, c) &= C_h(\rho^0) - \mu_h(v) = C_h(\rho^0) - (v([1, h]) - v([1, h-1])) \\ &= C_h(\rho^0) - \left(\sum_{\{i, j \in [1, h] | i < j\}} g_{ij} - \sum_{\{i, j \in [1, h-1] | i < j\}} g_{ij} \right) \\ &= C_h(\rho^0) - \sum_{i=1}^{h-1} g_{ih}. \end{aligned}$$

So, according to this solution all the savings obtained from a switch of player h with any of its predecessors goes to player h . Fernández, Borm, Hendrickx and Tijs (2005) show that this cost-assignment rule is the unique solution being stable (i.e. $c^u(N, q, c)$ is in the core of the cost-game for any linear cost sequencing situation (N, q, c)) and satisfying the so-called property of *Drop Out Monotonicity* (DOM). Clearly, the stableness property follows from the fact that $\mu(v)$ is in the core of the cost-savings game. To state DOM, let (N_{-k}, q_{-k}, c_{-k}) with player set $N_{-k} = N \setminus \{k\}$ be the $(n-1)$ -player sequencing situation obtained when player k leaves the queue (i.e., job k is cancelled) and let v^{-k} be the corresponding characteristic function. Then a cost assignment rule r assigning costs $r_h(N, q, c)$ for all $h \in N$ satisfies DOM if for any linear cost situation (N, q, c) it holds that

$$r_h(N_{-k}, q_{-k}, c_{-k}) \leq r_h(N, q, c), \quad h \in N_{-k},$$

i.e. if one of the players leaves the queue, for each of the remaining players the costs are non-increasing.

The cost rule c^u induced by $\mu(v)$ indeed satisfies DOM. In fact, it follows by straightforward calculations that

$$c_h^u(N_{-k}, q_{-k}, c_{-k}) = c_h(N, q, c), \quad h = 1, \dots, k-1,$$

and

$$c_h^u(N_{-k}, q_{-k}, c_{-k}) = c_h^u(N, q, c) - \min(\alpha_h q_k, \alpha_k q_h) < c_h^u(N, q, c), \quad h = k+1, \dots, n.$$

So, for the players in front of k there is no change in the net-costs, whereas for any player h after k the decrease $\alpha_h q_k$ in initial costs when k leaves the queue is bigger than the loss

of the cost-savings g_{kh} (if positive) from a switch between k and h . The DOM property advocated in Fernández et al. (2002) seems to be very appealing and reasonable: when player k drops out, the players in front of k are not affected, while for the players after k the costs are decreasing.

On the other hand, let (N^j, q, c) denote the adjusted sequencing situation in which some player j refuses to cooperate with any player $k > j$. As a consequence of this refusal of j , any coalition $[i, h]$, $i \leq j < h$, cannot form and hence all the gains g_{ih} of a switch between i and h , $i \leq j < h$, cannot be realized anymore. Let v^j be the corresponding characteristic function of the cost-savings game. For $S \in \mathcal{N}$, let $P^j(S)$ be the unique minimal partition of S in coalitions of consecutive players not containing both j and $j+1$, i.e. $T \in \mathcal{L}$ if $T \in P^j(S)$ and if $T_1 \cup T_2 \in \mathcal{L}$ for some pair $T_1, T_2 \in P^j(S)$, then $j \in T_1$ and $j+1 \in T_2$ (or reversely). Then it holds that

$$v^j(S) = \sum_{T \in P^j(S)} v(T), \quad S \in \mathcal{N},$$

i.e. the value of a coalition S is equal to the sum of the values of the coalitions in its unique minimal partition $P^j(S)$. From this property it follows immediately that the net-cost according to $\mu(v^j)$ becomes equal to

$$c_h^u(N^j, q, c) = \begin{cases} C_h(\rho^0) - \sum_{i=1}^{h-1} g_{ih} = c_h^u(N, q, c), & h = 1, \dots, j, \\ C_h(\rho^0) - \sum_{i=j+1}^{h-1} g_{ih}, & h = j+1, \dots, n. \end{cases}$$

Comparing this with the costs $c_h^u(N, q, c)$ of the original situation, it follows that the costs do not change for the players $1, \dots, j$, whereas a player h after j loses all the gains g_{ih} , $i \leq j$, and therefore suffers from an increase in the costs with $\sum_{i \leq j} g_{ih}$. So, the unique core outcome satisfying DOM has the serious drawback that it does not give any incentive to a player to cooperate with its successors in the queue: not cooperating does not hurt her. To make the point more clear, consider a two player sequencing situation. Of course, nothing happens if the initial order of 1 before 2 is optimal already. So, suppose it is optimal to reverse the initial order and to place 2 in front of 1 generating a decrease g_{12} of the costs. According to the payoff vector $\mu(v)$, this decrease is fully assigned to player 2. Why should player 1 be willing to cooperate by agreeing to take the second position? On the contrary, player 1 has the power to play the noncooperative ultimatum game and to offer the first place in the queue to player 2 if player 2 is willing to give all the gains of this change to player 1, i.e. player 1 is willing to sell his place against a price equal to $\alpha_1 q_2$ (the additional costs of waiting for player 1) plus the gains g_{12} of this trade.

Extending this reasoning we obtain that any player $j < n$ can decide upon whether or not to cooperate with his successors. This is modelled by the hierarchy with player one as the topman and a corresponding power function satisfying condition (2). According to

Corollary 6.1 then $\lambda(v) = m^\ell(v)$ is the unique element in $SC(v, p)$. So, the marginal vector $\lambda(v)$ is economically stable and socially stable with respect to a social structure reflecting the dominance of any player j over its successors in the sequence. The resulting costs induced by $\lambda(v)$ are given by

$$c_h^\ell(N, q, c) = C_h(\rho^0) - \lambda_h(v) = C_h(\rho^0) - \sum_{h < j} g_{hj}, \quad h \in N.$$

Of course, this cost rule does not satisfy DOM, but any player i gets the gains of switching with her successors and therefore is willing to cooperate with her successors.

7 The water distribution problem

In their paper ‘Sharing a river’ Ambec and Sprumont (2002) consider the problem of the optimal distribution of water to agents located along a river from upstream to downstream. Let $N = \{1, \dots, n\}$ be the set of agents, numbered successively from upstream to downstream and let $f_i \geq 0$ be the flow of water entering the river between agent $i - 1$ and i , $i = 1, \dots, n$, with f_1 the inflow before the most upstream agent 1. Agent i , $i = 1, \dots, n$, has a quasi-linear utility function given by $u^i(x_i, t_i) = b^i(x_i) + t_i$, where t_i is a monetary compensation to agent i , x_i is the amount of water allocated to agent i , and $b^i: \mathbb{R}_+ \rightarrow \mathbb{R}$ a continuous non-decreasing function yielding the benefit $b^i(x_i)$ to agent i of the consumption x_i of water. An allocation is a pair $(x, t) \in \mathbb{R}_+^n \times \mathbb{R}^n$ of water distribution and compensation scheme, satisfying

$$\sum_{i=1}^n t_i \leq 0 \quad \text{and} \quad \sum_{i=1}^j x_i \leq \sum_{i=1}^j f_i, \quad j = 1, \dots, n.$$

The first condition is a budget condition and says that the total amount of compensations is non-positive, i.e. the compensations only redistribute the total welfare. The second condition reflects that any agent can use the water that entered upstream, but that the water inflow downstream of some agent can not be allocated to this agent. Because of the quasi-linearity and the possibility of making money transfers, an allocation is Pareto optimal if and only if the distribution of the water streams maximizes the total benefits, i.e. the water distribution $x^* \in \mathbb{R}_+^n$ solves the following maximization problem:

$$\max_x \sum_{i=1}^n b^i(x_i) \quad \text{s.t.} \quad \sum_{i=1}^j x_i \leq \sum_{i=1}^j f_i, \quad j = 1, \dots, n. \quad (4)$$

A welfare distribution allocates the total benefits of an optimal water distribution x^* over the agents, i.e. it is a vector $z \in \mathbb{R}^n$ assigning utility z_i to agent i and satisfying $\sum_{i=1}^n z_i = \sum_{i=1}^n b^i(x_i^*)$. Clearly, any welfare distribution can be implemented by the allocation (x, t) with $x_i = x_i^*$ and $t_i = z_i - b^i(x_i^*)$, $i = 1, \dots, n$.

The problem to find a reasonable welfare distribution can be modelled as a TU-game. Obviously, for any pair of players i, j with $j > i$ it holds that the water inflow entering the river before the upstream agent i can only be allocated to the downstream agent j if all agents between i and j cooperate, otherwise any agent between i and j can take the flow from i to j for its own use. Hence, only coalitions of consecutive agents are admissible. Clearly, for $S = N$, $v(N) = \sum_{i=1}^n b^i(x_i^*)$ with $x^* \in \mathbb{R}^n$ the solution of maximization problem (4). For any connected coalition $S = [i, j] \in \mathcal{L}$, its worth is given by

$$v([i, j]) = \sum_{h=i}^j b^h(x_h^{*S}),$$

where $x^{*S} = (x_h^{*S})_{h=i}^j$ solves

$$\max_{x_i, \dots, x_j} \sum_{h=i}^j b^h(x_h) \quad \text{s.t.} \quad \sum_{k=i}^h x_k \leq \sum_{k=i}^h f_k, \quad h = i, \dots, j. \quad (5)$$

Without loss of generality we normalize the benefit functions by taking $b^i(f_i) = 0$, implying that $v(\{i\}) = b^i(f_i) = 0$, $i = 1, \dots, n$, so that the values $v(S)$ for $|S| \geq 2$ represent the net-gains of cooperating. Again, for an arbitrary coalition $S \in \mathcal{N}$, the value $v(S)$ is equal to the sum of its consecutive parts, i.e.,

$$v(S) = \sum_{T \in P(S)} v(T), \quad S \in \mathcal{N},$$

with $P(S)$ the minimal partition as defined in the previous section. Clearly, the game v is superadditive, $v = v^r$, and hence it follows from Granot and Huberman (1982) that v is permutationally convex for the permutations u and ℓ . Consequently, both marginal vectors $\mu(v)$ and $\lambda(v)$ are core solutions. In case all functions b^i are differentiable with derivative going to infinity as x_i tends to zero, strictly increasing and strictly concave, Ambec and Sprumont (2002) have even shown that the game is convex and hence that the core contains all marginal vectors.

Under the conditions for convexity, Ambec and Sprumont (2002) have shown that the marginal vector corresponding to the permutation u is the unique element in the core of the game satisfying a so-called fairness condition. This condition (quite different from the fairness condition of Myerson to characterize the Shapley value) says that any coalition S gets at most its aspiration level, being the highest utility it can obtain when it may use all the water of all the agents $1, \dots, \hat{s}$, where $\hat{s} = \max\{s \mid s \in S\}$. Clearly, this implies that any coalition $[1, j]$ can get at most $v([1, j])$, $j = 1, \dots, n$, so that it trivially follows that indeed the marginal vector $m^u(v)$ assigning $\mu_i(v) = v([1, i]) - v([1, i-1])$, $i = 1, \dots, n$, is the unique candidate in the core satisfying the aspiration requirements. For the proof that it indeed satisfies the requirements we refer to Ambec and Sprumont (2002).

As in the sequencing situation again we have that the payoff vector $\mu(v)$ has the property that when a player j does not want to cooperate, the players in front of j , including j itself, are not hurt. Like in the sequencing game, this is a very counterintuitive outcome. Although any upstream coalition $[1, j]$ can prevent that coalition $[j + 1, n]$ gets more than $v([j + 1, n])$ by using all flows f_1, \dots, f_j by itself, all benefits from cooperating go to the coalition $[j + 1, n]$. Again the outcome $\mu(v)$ has the serious drawback that it does not give any incentive to a player j to cooperate with its successors in the queue. Repeating the reasoning once more, again consider a two agent situation. In this case there is no gain of cooperation when in the optimal solution player 1 fully consumes its upstream inflow f_1 . However, suppose it is optimal to allocate a part of f_1 to the second agent. According to the outcome $\mu(v)$, agent 1 is just compensated by agent 2 for its loss of utility, i.e. player 1 receives a compensation $t_1 = b^1(f_1) - b^1(x_1^*)$, giving her utility $b^1(f_1) = v(\{1\}) = 0$. So, like in the sequencing game, there is no reason for player 1 to cooperate. However, player 1 has the power to play the noncooperative ultimatum game and to pass the stream $f_1 - x_1^*$ to player 2 if this player is willing to give up all the gains of cooperation, i.e. player 1 is willing to sell this stream against a price (or compensation) t_1 equal to $v(\{1, 2\}) - v(2) = v(\{1, 2\})$. Player 2 is indifferent to accepting this offer or not and therefore is willing to accept the offer (or any slightly lower price). Also in this river game we may argue that player 1 is in control of whether or not to cooperate by letting through a part of its water inflow f_1 . In general, any player is in control of cooperation with its successors. This can be modelled by a hierarchy consisting on the line graph with player 1 as the topman and corresponding power function satisfying condition (2). According to Corollary 6.1, $\lambda(v)$ is the unique element in the socially stable core, yielding an outcome reflecting to a social structure in which any player j dominates its successors in the sense that player j controls the water inflow up to j .

8 Conclusions

In this paper we consider structured transferable utility games. In such games each coalition is organized according to some social structure, for instance a network, a hierarchy or some other dominance relation. There exist several measures in the literature to measure the power of the players within a coalition having some social structure. In this way the social structure of a coalition can be represented by a power vector. A collection of coalitions is stable if weights can be assigned to these coalitions such that the total weighted power of every agent is the same. Stability generalizes the well-known concept of balancedness, where all members of a coalition have the same power. The socially stable core consists of the economically stable payoff vectors that are sustained by a stable collection of coalitions.

The socially stable core is a subset of the superadditive cover core, and therefore of the core whenever the game is superadditive. We provide several results on non-emptiness and uniqueness of the socially stable core. If the game is convex, then the socially stable core is non-empty for any social structure. Also, the socially stable core is non-empty if for some permutation the game is permutationally convex and the power function is permutationally consistent. If the game is permutationally convex and the power function is permutationally compatible, then the socially stable core consists of only one element, being the marginal vector corresponding to the permutation. Generically, the socially stable core consists of a finite number of elements only. Therefore, the socially stable core is a forceful selection device for the superadditive cover core.

The socially stable core concept is applied to structured hierarchy games. Under rather mild conditions the socially stable core of a structured hierarchy game consists only of a specific marginal vector. This vector is completely determined by the hierarchical structure on the players. The sequencing problem and the water distribution problem are modeled as two examples of structured hierarchy games. As a consequence, we are able to advocate a unique solution for these two problems.

In this paper the social structure has been given exogenously. In the network literature, the social structure is the result of a link formation process. It would be interesting to combine insights provided by this paper with those coming from the network formation literature. In this way one can both endogenize the social structure and the distribution of network payoffs, which opens up an interesting line of further research.

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