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Social Choice among Complex Objects:

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Abstract

We present a geometric model of social choice when the latter takes place among bundles of interdependent elements, that we will call objects. We show that the outcome of the social choice process is highly dependent on the way these bundles are formed. By bundling and unbundling the same set of constituent elements an authority enjoys a vast power of determining the social outcome. We provide necessary and sufficient conditions under which a social outcome may be a local or global optimum for a set of objects, and we show that, by appropriately redefining the set of objects, intransitive cycles may be broken and the median voter may be turned into a loser.

1 Introduction

Social choice theory usually assumes that agents are faced with a set of exogenously given and mutually exclusive alternatives. These alternatives are given in the sense that the pre-choice process through which they are constructed is not analyzed. Moreover, these alternatives are "simple", in the sense that they are one-dimensional objects or, even when they are multidimensional, they are simply points in some portion of the homogeneous $\mathbb{R}^n$ space and they lack an internal structure that limits the set of possible alternatives.

Many choices in real life situations depart substantially from this simple setting. Choices are often made among bundles of interdependent elements. Those bundles may be formed in a variety of ways, which in turn affect the selection process of a social outcome. Let us consider, for instance, a typical textbook example of social choice, i.e. the case of a group of friends deciding how to spend the evening by democratic pairwise majority voting. The textbook would start from a given and predefined choice set as $X = \{A, B, C, D, \ldots\}$ where the set $\{A, B, C, D, \ldots\}$ could stand for *movie, concert, restaurant, dinner at home, ...*. At closer scrutiny, these alternatives are neither primitive nor exogenously given. Going to the movies or to a restaurant are labels for bundles of elements (e.g. with whom, where, when, movie genre, director, type of food, etc.) and everyone's preference is unlikely to be expressed before these labels are specified in their constituting elements. A skillful member of the group could easily obtain a social outcome close to the one he or she prefers by carefully crafting the objects $A, B, C, D, \ldots$ and possibly by designing a new set of objects.

Moving on, to more serious examples, candidates and parties in political elections stand for complex bundles of interdependent policies and personality traits. Committees and boards are called upon to decide upon packages of policies, e.g. a recruitment package that a university governing board has to approve. In principle, any combination of elements (subject to a budget or some other constraint) could be considered and compared (e.g. through majority voting) with any other, but in reality only a relatively small number of packages undergo examination. Typically, the bundling of elements serves the purpose of reducing the number of alternatives to be examined, by decomposing the whole space of alternatives into smaller subspaces.

In this paper we present a geometric model of social choice among bundles of elements, which we call *objects*. Geometric approaches have already been used in the literature on social choice. Donald Saari has greatly contributed to establishing general geometric representations of voting models and voting paradoxes (Saari 1994, Saari 2000a, Saari 2000b), and we

will argue later that our representation is more general in many respects. (Eckmann 1954), (Eckmann, Ganea, and Hilton 1962), (Chichilnisky 1980), (Chichilnisky 1983) study the problem of the existence of a social decision function from a topological point of view and show that the paradoxes of social choice are partly a consequence of the topological structure of the spaces of ordinal preferences. On the other hand, (Baryshnikov 1997) discusses the possibility of introducing topological methods in the combinatorial paradigm of social choice theory. (Weinberger 2004) and (Terao 2007) extend well-known results on social choice functions to, respectively, CW complexes and arrangements, thus obtaining new results for both mathematical objects. In this respect, our model is a novel contribution to the analysis of discrete problems of social choice.

We deal with the issue of object construction that has not been dealt with by economic models. The main idea is that, first, generally objects are not simply aggregations of primitive components but have an internal structure that is likely to determine interdependencies and non-separabilities in individual preferences. In the "what shall we do tonight" choice setting, my preferences on the *with whom* element is likely to be highly interdependent with the other elements, as I may well find a given person a perfect companion for an evening at the movies but relatively dislike her or his company if we finally decide to go to a restaurant. Second, objects provide structure to the choice problem. Consider again the "what shall we do tonight" case. A possible reply to our point on bundles would be that the choice set X is underspecified and that we should start from a choice set formed by all possible combinations of the elements. However, for obvious combinatorial arguments, this set, even in this simple example, would be so large that any exhaustive choice procedure, e.g. pairwise majority voting, could not be completed in a feasible time span. In our approach, objects decompose the search space into quasi-separable subspaces (Simon 1982) and simplify the computational task, making decisions possible.

There is also another way objects can contribute to making the determination of a social outcome easier. We will show that, by appropriate object construction, intransitive cycles that often characterize social decisions can almost always be eliminated. In general, coarse objects, i.e. those made of many elements, tend to produce many cycles, whereas fine objects do not. However finer objects do so by increasing the number of locally stable social optima and thereby making the social outcome more manipulable through the control of object construction, initial conditions and agendas. All in all, we will analyze how different sets of objects strike different balances in the trade-off between decidability and non-manipulability.

Our work is closely linked to the literature on agenda power (McKelvey

1976, Plott and Levine 1978), and we will show that we generalize some of its results in the sense that even agenda power is subject to manipulation through object design. It is also related to recent literature on multidimensional voting models (Kramer 1972, Shepsle 1979, Denzau and Mackay 1981, Enelow and Hinich 1983, Shepsle 1979) and the one that has begun to analyze decision-making when agents group states of the world into coarse categories (Mullainathan 2000, Fryer and Jackson 2008), but there are fundamental differences that make our model new.

In order to formally analyze the properties of a social choice model with object construction power we use tools of hyperplane arrangements theory and graph theory. Thanks to them, we show that the way objects are constructed by bundling (or unbundling) elements can strongly impact on the outcome of social choice. We show that, in general, by appropriately constructing objects, the outcome of a social choice process, e.g. pairwise majority voting, may be heavily manipulated. An authority who has the power to construct objects may obtain a desired outcome even when the latter is chosen democratically. We will prove necessary and sufficient conditions for any social outcome to be a locally or globally stable social optimum for a social decision rule under a specific set of objects. We will also show an algorithmic procedure to determine which set of objects, agendas and initial conditions can lead society, through a given social choice procedure, to select a given outcome. Moreover, the object construction power that we describe and analyze in this paper will be proven to be stronger and more general than the well-known agenda power in the sense that also the agenda power may be manipulated by appropriate object construction. Finally, we will discuss how another classic result, i.e. the so-called median voter theorem, may be reverted by appropriate modification of the set of objects, thus transforming the median voter into an outright loser.

This paper is structured as follows: in section 2 we outline our geometric model. A key ingredient of our analytic approach is the theory of hyperplane arrangements. Then, in section 3 we draw the main results concerning how, through object construction, it is possible to manipulate social outcomes, create or eliminate agenda power, create or eliminate intransitive cycles, and turn the median voter into a loser. We will also demonstrate that our object construction power is stronger and more general than the traditionally considered agenda power and that objects strike a balance in the trade-off between decidability and non-manipulability. In a companion paper (Marengo and Pasquali 2008) many analytical results obtained here are illustrated through examples and computational models. The interested reader can refer to it for algorithmic implementations of most of the arguments contained in the present paper.

2 Definitions and structure of the model

Features   We assume that choices are made over bundles of elements that we call *features*. The number of features is finite and each feature may take one value out of a finite set of alternatives. We call $F = \{f_1, \ldots, f_n\}$ the set of such features. In order to simplify the notation and without loss of generality, we assume that all features may take the same number $m + 1$ of values in the set $\{0, 1, 2, \ldots, m\}$ $\forall i = 1, \ldots, n$.

The space of social outcomes is given by the $(m+1)^n$ $n-$tuples specifying one value for every feature. We call X the set of all possible social outcomes and $x = (v_1.\cdots.v_n)$ a generic element thereof.

Example 1. Let us consider an example of 3 features $\{f_1, f_2, f_3\}$ taking a value out of the binary set $\{0, 1\}$. This is equivalent to, e.g., the case of presence or absence of three possible traits. Thus the space X of possible social outcomes is a set of 8 ordered 3-tuples of the form $x = (v_1.v_2.v_3)$ for $v_i \in \{0, 1\}$, $i = 1, 2, 3$. This can be regarded as the example in Introduction: a group of friends deciding how to spend the evening. Hence the 3 features could be {where to go, when to go, how to go} taking a value in $\{0, 1\}$, e.g. where to go {restaurant, movies}, when to go {8pm, 10pm}, how to go {by car, walking}. An element $x = (v_1.v_2.v_3)$ in X would be, for example, (restaurant.8pm.car).

The hyperplane arrangement   Given a set of features $F = \{f_1, \ldots, f_n\}$ taking values in the set $\{0, 1, 2, \ldots, m\}$, we can associate to it the hyperplane arrangement[1] in $\mathbb{R}^n$

$$\mathcal{A}_{n,m} = \left\{ H_{i,j} \mid \alpha_{H_{i,j}} = \lambda_i - j \right\}_{\substack{1 \leqslant i \leqslant n \\ 0 \leqslant j \leqslant m-1}}$$

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[1]For the reader who is unfamiliar with hyperplane arrangements, we provide some basic notions in (Settepanella 2010)

where indices $1 \leqslant i \leqslant n$ and $0 \leqslant j \leqslant m - 1$ vary, respectively, among the index of features and the index of values taken by features.

The social outcome $x = (v_1. \cdots .v_n)$ will correspond to the chamber C_x that contains the open set

$$\{(\lambda_1, \ldots, \lambda_n) \in \mathbb{R}^n \mid v_j - 1 < \lambda_j < v_j, \ j = 1, \ldots, n\}.$$

Let us remark that in our framework the number n of features gives the dimension of the space $\mathbb{R}^n$ and the values $v_i \in \{0, 1, \ldots, m - 1\}$ taken by the feature f_i provide a set of parallel hyperplanes.

Example 2. In example 1 the associated arrangement $\mathcal{A}_{3,2}$ is simply the arrangement given by the coordinate hyperplanes of an orthogonal Cartesian system in $\mathbb{R}^3$. Hence we get 8 chambers corresponding to the 8 different octants in which the orthogonal Cartesian system divides $\mathbb{R}^3$. Each social outcome $x \in X$ corresponds to an octant C_x.

Let us remark that, in general, there isn't any relation between the n-tuple $x = (v_1. \cdots .v_n)$ and the point $(v_1, \ldots, v_n) \in \mathbb{R}^n$. Nevertheless, depending on what we are choosing upon (e.g. if features take values on real numbers), we can slightly modify the arrangement of hyperplanes defining it in such a way that the point $(v_1, \ldots, v_n) \in \mathbb{R}^n$ is contained in the chamber C_x, $x = (v_1. \cdots .v_n)$. This is just one of the many way in which our framework can be adapted in order to include different informations.

Social decision rules Consider a population of ν agents. Each agent i is characterized by a system of *transitive* preferences $\succeq_i$, i.e. transitive relations, over the set of social outcomes X. The set of systems of transitive preferences $\succeq_i$ is denoted by $\mathcal{P}$. Let $\mathcal{P}^\nu$ be the Cartesian product of ν copies of $\mathcal{P}$, a social decision rule $\mathcal{R}$ is a function:

$$\mathcal{R} : \mathcal{P}^\nu = \prod_{i=1}^{\nu} \mathcal{P} \longrightarrow \overline{\mathcal{P}}$$

$$(\succeq_1, \ldots, \succeq_\nu) \longmapsto \succeq_{\mathcal{R}(\succeq_1, \ldots, \succeq_\nu)}$$

which determines a system of social preferences $\succeq_{\mathcal{R}(\succeq_1, \ldots, \succeq_\nu)}$ from the preferences of ν individual agents. With $\overline{\mathcal{P}}$ we denote the set of systems of (non-necessarily transitive[2]) social preferences, i.e. a set of relations. Any

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[2]It is well known that individual transitivity does not guarantee transitivity for a social rule (Condorcet de Caritat marquis de 1785).

type of social decision rule may fit into our framework, provided individual preferences are expressed sincerely. For the time being we rule out the possibility of strategic misrepresentation of a person's preferences.

In order to completely describe $\overline{\mathcal{P}}$, let us recall that if $\Delta = \{(x,x) \in X \times X\}$ is the *diagonal* of the Cartesian product $X \times X$, then an element $\succeq_{\mathcal{R}} \in \overline{\mathcal{P}}$ defines a subset:

$$Y_{1,\succeq_{\mathcal{R}}} \subset X \times X \setminus \Delta$$

as follows: a pair (x,y) is in $Y_{1,\succeq_{\mathcal{R}}}$ if and only if $x \succeq_{\mathcal{R}} y$.

Without loss of generality we will indicate by $\succeq_{\mathcal{R}} \in \overline{\mathcal{P}}$ a generic element in the image of $\mathcal{R}$.

We call $Y_{0,\succeq_{\mathcal{R}}}$ the set of *relevant* social outcomes, i.e. the set of all social outcomes on which social preferences are expressed:

$$Y_{0,\succeq_{\mathcal{R}}} = \{x \in X \mid \exists y : (x,y) \in Y_{1,\succeq_{\mathcal{R}}} \text{ or } (y,x) \in Y_{1,\succeq_{\mathcal{R}}}\}.$$

The subset $Y_{0,\succeq_{\mathcal{R}}} \subseteq X$ is in one to one correspondence with a subset $\mathcal{C}_{\succeq_{\mathcal{R}}} \subset \mathcal{C}$ of the set of chambers in $\mathcal{A}_{n,m}$:

$$C_x \in \mathcal{C}_{\succeq_{\mathcal{R}}} \Leftrightarrow x \in Y_{0,\succeq_{\mathcal{R}}}.$$

Example 3. Let us consider the space of social outcomes X as in example 1. Let us assume that the aggregated preferences of the group of friends upon the 3 features give rise to the following rule $\succ_{\mathcal{R}}$:

$(0.0.0)$ preferred to all, except $(1.1.0) \succ_{\mathcal{R}} (0.0.0), (0.0.1) \succ_{\mathcal{R}} (0.0.0)$;
$(0.1.0) \prec_{\mathcal{R}} (0.1.1), (0.1.0) \prec_{\mathcal{R}} (1.1.1), (0.1.0) \prec_{\mathcal{R}} (1.0.0),$
$(0.1.0) \succ_{\mathcal{R}} (1.0.1), (0.1.0) \succ_{\mathcal{R}} (1.1.0), (0.1.0) \prec_{\mathcal{R}} (0.0.1)$;
$(0.1.1) \succ_{\mathcal{R}} (1.1.1), (0.1.1) \succ_{\mathcal{R}} (1.0.0), (0.1.1) \succ_{\mathcal{R}} (1.0.1),$
$(0.1.1) \succ_{\mathcal{R}} (1.1.0), (0.1.1) \prec_{\mathcal{R}} (0.0.1)$;
$(1.1.1) \succ_{\mathcal{R}} (1.0.0), (1.1.1) \succ_{\mathcal{R}} (1.0.1), (1.1.1) \succ_{\mathcal{R}} (1.1.0), (1.1.1) \succ_{\mathcal{R}} (0.0.1)$;
$(1.0.0) \succ_{\mathcal{R}} (1.0.1), (1.0.0) \succ_{\mathcal{R}} (1.1.0), (1.0.0) \prec_{\mathcal{R}} (0.0.1)$;
$(1.0.1) \succ_{\mathcal{R}} (1.1.0), (1.0.1) \prec_{\mathcal{R}} (0.0.1)$;
$(1.1.0) \prec_{\mathcal{R}} (0.0.1)$.

For the sake of simplicity, when possible, we will use notation $v_1 v_2 v_3$ instead of $(v_1.v_2.v_3)$. Thus $Y_{0,\succeq_{\mathcal{R}}} = \{v_1 v_2 v_3 \mid v_j = 0 \text{ or } 1\} = X$.
$Y_{1,\succeq_{\mathcal{R}}}$ is given by all pairs $(v_1 v_2 v_3, g_1 g_2 g_3)$ such that $v_1 v_2 v_3 \succ_{\mathcal{R}} g_1 g_2 g_3$; i.e., for example, $(000, 101), (011, 000) \in Y_{1,\succeq_{\mathcal{R}}}$ while $(000, 011) \notin Y_{1,\succeq_{\mathcal{R}}}$.

Let us remark that when individual preferences are *strict*, as in the example above, then they are total orders on the space of relevant social outcomes $Y_{0,\succ_{\mathcal{R}}}$ (i.e. antisymmetric, transitive relations). While the aggregated preferences give rise, in general, to non transitive relations, i.e. the rules $\succ_{\mathcal{R}}$ are not, in general, orders.

If social preferences are complete over X, then $Y_{0,\succeq_{\mathcal{R}}} = X$. In the sequel we will consider only complete preferences, although our framework may easily accommodate the more general case where $Y_{0,\succeq_{\mathcal{R}}} \subseteq X$.

Graphs   We can represent the sets $Y_{0,\succeq_{\mathcal{R}}}$ and $Y_{1,\succeq_{\mathcal{R}}}$ respectively as the set of vertices and edges of an *oriented graph* $\mathcal{Y}_{\succeq_{\mathcal{R}}}$. Two vertices x and y in $Y_{0,\succeq_{\mathcal{R}}}$ are connected by an edge if and only if $(x,y) \in Y_{1,\succeq_{\mathcal{R}}}$ or $(y,x) \in Y_{1,\succeq_{\mathcal{R}}}$, while the orientation is from x to y in the former case and from y to x in the latter. In a natural way this construction applies to all subsets $Y \subset X \times X \setminus \Delta$.

Without loss of generality we will denote by x the vertices of $\mathcal{Y}_{\succeq_{\mathcal{R}}}$ and by (x,y) its edges.

Example 4. The social preferences of example 3 are fully described by the oriented graph in figure 1.
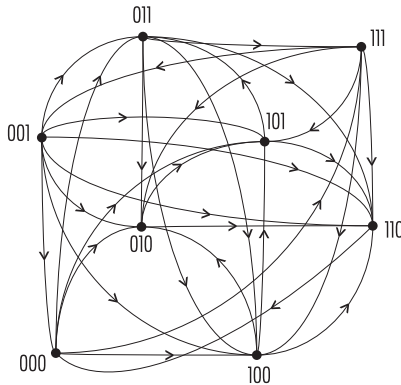

Figure 1: The oriented graph derived from example 3

Notice that the assumption of complete preferences guarantees that we will deal only with *connected* graphs.

A *cycle of length h* in the oriented graph $\mathcal{Y}_{\succeq_{\mathcal{R}}}$ is a subgraph γ_I with vertices $\{x_1, \ldots, x_h\}$ and edges $\{(x_1, x_2), (x_2, x_3), \ldots, (x_h, x_1)\}$. It corresponds to a Condorcet-Arrow cycle, i.e. to the sequence $x_1 \succeq_{\mathcal{R}} x_2 \succeq_{\mathcal{R}} \ldots \succeq_{\mathcal{R}} x_h \succeq_{\mathcal{R}} x_1$. In example 3 and in the corresponding graph of figure 1 we find a three outcomes cycle $000 \succ_{\mathcal{R}} 101 \succ_{\mathcal{R}} 011 \succ_{\mathcal{R}} 000$.

Salvetti's complex As shown by (Salvetti 1987), we can put together the hyperplane arrangements theory and the graph description in a more general context that allows us to use both geometric and algebraic tools. Indeed we can construct a CW-complex $\mathcal{S}(\mathcal{A})^3$ having the homotopy type of the complement $M(\mathcal{A})$ when the arrangement $\mathcal{A}$ is the complexification of a real one. In this complex each element $\succeq_\mathcal{R} \in \overline{\mathcal{P}}$ is equivalent to a subcomplex of the 1-skeleton of $\mathcal{S}(\mathcal{A}_{n,m})$.

Let us characterize the 0 and 1-dimensional cells (also called 0 and 1-skeletons or vertices and edges) of the Salvetti complex and show how they are related to the graph $\mathcal{Y}_{\succeq_\mathcal{R}}$ for a given social rule $\succeq_\mathcal{R}$.

Let $\mathcal{A}_{n,m}$ be the arrangement associated to the space of social outcomes and $\mathcal{S}(\mathcal{A}_{n,m})$ the correspondent Salvetti complex. The set of generators $\mathcal{S}_0(\mathcal{A}_{n,m})$ of the 0-skeleton of the complex $\mathcal{S}(\mathcal{A}_{n,m})$ is in a one-to-one correspondence with the set of chambers in $\mathcal{A}_{n,m}$, i.e. with the set of social outcomes X. Thus:

$$\mathcal{S}_0(\mathcal{A}_{n,m}) = Y_{0,\succeq_\mathcal{R}} = X$$

for any given social rule $\succeq_\mathcal{R}$.

Example 5. The 0-skeleton $\mathcal{S}_0(\mathcal{A}_{3,2})$ associated to example 1 is given by 8 points that can be chosen freely one in each chamber of the coordinate arrangement $\mathcal{A}_{3,2}$, i.e. one in each octant. Thus they can be represented as the vertices of a cube or, more generally, as the vertices of a 3-dimensional polyhedron.

As remarked before the vertex $x \in \mathbb{R}^3$ corresponding to the social outcome $x = v_1 v_2 v_3$ is not the point (v_1, v_2, v_3), e.g. 000 is not the point $(0, 0, 0)$ which is the origin of the orthogonal system and hence doesn't lie in any octant.

Let us recall that two chambers C_x and C_y are said to be *adjacent* if and only if they are separated by only one hyperplane H. If x and y are the 0-cells corresponding to C_x and C_y, then we can consider the edge (x, y) between them.

The generators of the 1-skeleton can be described as the elements in the set:

$$\mathcal{S}_1(\mathcal{A}_{n,m}) = \{(x, y) \in X \times X \setminus \Delta \mid x \text{ and } y \text{ are adjacent}\}.$$

Two elements in $\mathcal{S}_1(\mathcal{A}_{n,m})$ are *consecutive* if and only if the second entry of the first pair is equal to the first entry of the second pair: for example, (x, y) and (y, z). Given a subset of consecutive elements in $\mathcal{S}_1(\mathcal{A}_{n,m})$

$$\{(x_1, x_2), (x_2, x_3), \ldots, (x_{k-2}, x_{k-1}), (x_{k-1}, x_k)\} \in \mathcal{S}_1(\mathcal{A}_{n,m})$$

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[3]In the appendix we briefly recall the construction of the first boundary map of this complex.

we define their *formal sum* as:

$$(x_1, x_k) = \sum_{j=1}^{k-1} (x_j, x_{j+1}). \tag{1}$$

Let C_x and C_y be two adjacent chambers separated by the hyperplane H. We say that we move from C_x to C_y across H if we move along the edge (x, y). Thus, by moving across hyperplanes, we can reach any chamber of the arrangement starting from a given one. If we only cross each hyperplane once, we call this path *minimal*. Obviously, there are many minimal paths, depending on which order we cross hyperplanes, however, they are all homotopically equivalent (Salvetti 1987).

It follows that, given a rule $\succeq_{\mathcal{R}}$, any edge $(x, y) \in Y_{1, \succeq_{\mathcal{R}}}$ can be written as a formal sum of a minimal number of consecutive elements in $\mathcal{S}_1(\mathcal{A}_{n,m})$. The number of elements is exactly the number of hyperplanes that separate the two social outcomes $x, y \in X$.

Remark 2.1. Notice that Salvetti's complex is a CW-complex in $\mathbb{C}^n$, but it has a underlying real structure which is a purely simplicial complex. This structure can be used in order to recast and generalize some existing geometric models of voting such as those provided by (Saari 1994).

3   Object construction and social outcomes

In this section we develop our main results. We define objects as bundles of features, and show that, in general, by appropriate manipulations of the set of objects, almost every social outcome may be obtained from a given social rule (e.g. majority voting). Such "object construction power", that is, the power of determining social outcomes by appropriately bundling or unbundling features, is stronger than the agenda power traditionally studied in the literature (McKelvey 1976). We then show that this power of manipulation also includes the possibility of breaking or creating intransitive cycles *à* la Condorcet-Arrow and of overturning the median voter effect. Finally we discuss the emerging trade-off between non manipulability and decidability. Coarser objects make social decisions less manipulable but they also increase the likelihood of intransitive cycles and the time required to reach a socially optimal outcome (if any). On the other hand finer objects make cycles less likely and allow reaching a social outcome more quickly, but the number of locally optimal social outcomes greatly increases and, therefore, the scope for manipulability becomes broader.
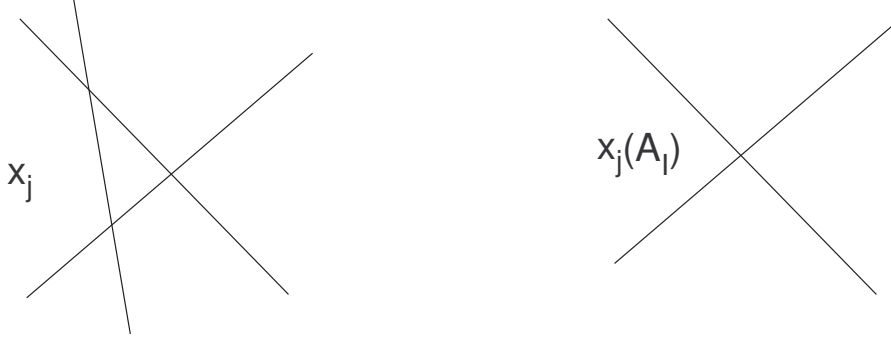
Figure 2: A graphical characterization of object instantiation

3.1 Objects in social choice

In this paragraph we give the main definitions and notations useful in the rest of the paper.

Objects schemes With the notation introduced above, let $\mathcal{A}_{n,m}$ be the hyperplane arrangement defined by n features $F = \{f_1, \ldots, f_n\}$ indexed by the set $\{1, \ldots, n\}$ and m possibilities for each of them. Let $\succeq_{\mathcal{R}}$ be the social preferences over the set of social outcomes X. Given a subset $I \subset \{1, \ldots, n\}$, an *object* $\mathcal{A}_I$ is a non empty subset of the arrangement $\mathcal{A}_{n,m}$ of the form:

$$\mathcal{A}_I = \big\{ H_{i,j} \big\}_{\substack{i \in I \\ 0 \leq j \leq m-1}}.$$

The cardinality of $\mathcal{A}_I$ is called the *size* of the object $\mathcal{A}_I$. We will also denote by $\mathcal{A}_{I^c} = \mathcal{A}_{n,m} \setminus \mathcal{A}_I$ the complement of the arrangement $\mathcal{A}_I$ in $\mathcal{A}_{n,m}$.

A set of objects $A = \{\mathcal{A}_{I_1}, \ldots, \mathcal{A}_{I_k}\}$ such that $\cup_{j=1}^k I_j = \{1, \ldots, n\}$ is an *objects scheme*. Notice that the subsets $I_1, \ldots, I_k$ don't have to partition the set I as some index h may belong to more than one subset. Since each subset I_j naturally corresponds to a subset $\{f_i\}_{i \in I_j}$ of features, this is equivalent to notice that an objects scheme does not have to partition the set of features, as some of them may belong to more than one object. However, we require that the union of all objects covers all the features, otherwise the remaining features would not be decided on and could be simply omitted from the model.

Let y be a social outcome in X, i.e. a chamber of $\mathcal{A}_{n,m}$, then the *object instantiation* $y(\mathcal{A}_I)$ is the chamber of the subarrangement $\mathcal{A}_I$ which contains the chamber corresponding to y, as shown, for instance, in figure 2.

Moreover, the *size* of an objects scheme is the size of its largest object:

$$| A |= \max\{| \mathcal{A}_{I_1} |, \ldots, | \mathcal{A}_{I_k} |\}.$$

Neighbors of a social outcome   Given an objects scheme $A = \{\mathcal{A}_{I_1}, \ldots, \mathcal{A}_{I_k}\}$, we say that a social outcome y is a *preferred neighbor* of a social outcome x with respect to an object $\mathcal{A}_{I_h} \in A$ if the following conditions hold:

1. $y \succeq_{\mathcal{R}} x$,

2. $y(\mathcal{A}_{I_h^c}) = x(\mathcal{A}_{I_h^c})$, i.e. x and y belong to the same chamber of the arrangement $\mathcal{A}_{I_h^c}$,

3. $y(\mathcal{A}_{I_h}) \neq x(\mathcal{A}_{I_h})$, i.e. x and y belong to different chambers of the arrangement $\mathcal{A}_{I_h}$.

In words, a preferred neighbor of a social outcome x with respect to an object is another social outcome y which is preferred to x and differs from x only in features that belong to that object. The set of all preferred neighbors of the social outcome x with respect to $\mathcal{A}_{I_h} \in A$ is denoted by $\Phi(x, \mathcal{A}_{I_h})$. The set of all preferred neighbors of the social outcome x is denoted by $\Phi(x, A) = \bigcup_{j=1}^{k} \Phi(x, \mathcal{A}_{I_j})$.

A social outcome $y \in \Phi(x, \mathcal{A}_{I_h})$ is said to be a *best neighbor* of a social outcome x with respect to an object $\mathcal{A}_{I_h} \in A$ if

$$y \succeq_{\mathcal{R}} w \quad \forall w \in \Phi(x, \mathcal{A}_{I_h}).$$

The set of all best neighbors of the social outcome x with respect to $\mathcal{A}_{I_h} \in A$ is denoted by $B(x, \mathcal{A}_{I_h})$. The set of all best neighbors of the social outcome x is denoted by $B(x, A) = \bigcup_{j=1}^{k} B(x, \mathcal{A}_{I_j})$.

Paths and reachable social outcome   A *path $P(x, y, A)$ through A, starting from x and ending in y*, is a sequence of best neighbors with respect to objects in A, i.e. a sequence

$$x = x_0 \preceq_{\mathcal{R}} x_1 \preceq_{\mathcal{R}} \ldots \preceq_{\mathcal{R}} x_{s+1} = y \tag{2}$$

such that there exist objects $\mathcal{A}_{I_{h_1}}, \ldots, \mathcal{A}_{I_{h_s}} \in A$ with $x_i \in B(x_{i-1}, \mathcal{A}_{I_{h_i}})$ for all $1 \leqslant i \leqslant s$.

A social outcome y is said to be *reachable from x with respect to an objects scheme A* if there exists a path $P(x, y, A)$.

Local optima A social outcome x is said to be a *local optimum* for A if and only if for all $y \in \Phi(x, A)$, $(x, y) \in Y_{1, \succeq_{\mathcal{R}}}$. In the case of strict preferences, this amounts to the condition $\Phi(x, A) = \emptyset$.

Agenda Let $A = \{\mathcal{A}_{I_1}, \ldots, \mathcal{A}_{I_k}\}$ be an objects scheme. An *agenda* α of A is an ordered tuple of indices $(h_1, \ldots, h_t)$ such that the set $\{h_1, \ldots, h_t\} = \{1, \ldots, k\}$. An agenda α states the order in which objects $\mathcal{A}_{I_i}$ are decided upon. In general all objects of the given scheme appear at least once in an agenda, but they may appear more than once. In fact, because of non-separabilities, it is possible that after modifications of other objects, an object previously set into what appeared to be a locally optimal social outcome, may undergo further improvements. In order to study only the stable limit properties of agendas, we suppose that the agenda is repeated over and over again until either a local optimum or a cycle is reached.

The ordered tuple of objects $(\mathcal{A}_{I_{h_1}}, \ldots, \mathcal{A}_{I_{h_t}})$ is denoted by A_α. The set of all possible agendas of A is denoted by $\Lambda(A)$. Let $\alpha = (h_1, \ldots, h_t)$ be an agenda. A path

$$x_0 \preceq_{\mathcal{R}} x_1 \preceq_{\mathcal{R}} \ldots \preceq_{\mathcal{R}} x_s \tag{3}$$

is said to be *ordered along* α if

$$x_i \in B(x_{i-1}, \mathcal{A}_{I_{h_q+1}})$$

where h_q is the remainder of the division of $i - 1$ by t. Such a path will be denoted by $P(x_0, x_s, A_\alpha)$.

Maximal paths A path is said to be *maximal* if it ends in either a local optimum or a limit cycle, i.e. a cycle such that there is no social outcome not belonging to the cycle itself which is a preferred neighbor to any social outcome in the cycle. Thus a limit a cycle, once entered, cannot be exited.

Basin of attraction The *basin of attraction* $\Psi(x, A)$ of a local optimum x *with respect to an objects scheme* A is the set of the social outcomes y such that there exists a maximal path $P(y, x, A)$ that ends up in x.

The *ordered basin of attraction* $\Psi(x, A_\alpha)$ of x *with respect to an agenda* α *of* A is the set of the social outcomes y such that there exists a maximal path $P(y, x, A_\alpha)$ that ends up in x. Clearly, we have

$$\Psi(x, A) = \bigcup_{\alpha \in \Lambda(A)} \Psi(x, A_\alpha).$$

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Global optima A social outcome $z \in X$ is said to be a *global optimum for an agenda* α if $\Psi(z, A_\alpha) = X$ holds. It is said to be a *global optimum for the objects scheme* A if and only if $\Psi(z, A_\alpha) = X$ holds for all agendas $\alpha \in \Lambda(A)$, i.e. it is a global optimum for all the agendas of A.

It is easy to verify that local and global optima strictly depend on the choice of the objects scheme A.

Example 6. Let us go back to example 3 and denote by H_i the hyperplane kernel of the polynomial $\alpha_{H_i} = \lambda_i$ for $i = 1, 2, 3$. If the objects scheme is given by $A = \{\{H_1, H_2, H_3\}\}$, i.e. agents have to decide where, when and how simultaneously or by $A = \{\{H_1, H_2\}, \{H_3\}\}$, i.e. agents decide where and when, given how, and, separately, they decide upon how, then the social rule always produces a cycle, that may be either $001 \succ_\mathcal{R} 011 \succ_\mathcal{R} 111 \succ_\mathcal{R} 001$ or $000 \succ_\mathcal{R} 010 \succ_\mathcal{R} 110 \succ_\mathcal{R} 000$ depending on the chosen agenda and initial condition. While if $A = \{\{H_1, H_3\}, \{H_2\}\}$ or $A = \{\{H_2, H_3\}, \{H_1\}\}$ then there exists a unique global optimum 001 for all agendas.

The dependence of the optimum on the objects scheme is very strong. Indeed, it is easy to provide examples in which two different social outcomes $z_1, z_2 \in X$ are global optima for two different choice of objects schemes. More precisely, there exist two objects schemes A_1 and A_2 such that $\Psi(z_i, A_i) = X$ and $\Psi(x, A_i) = \emptyset$ for all social outcomes $x \neq z_i$ and $i = 1, 2$.

Example 7. Let us consider the space of social outcomes X as in example 1 with the following preferences $\succ_\mathcal{R}$:

$(0.0.0)$ preferred to all except $(0.1.1) \succ_\mathcal{R} (0.0.0)$

$(0.1.1)$ preferred to all except $(1.0.1) \succ_\mathcal{R} (0.1.1)$

$(1.0.1) \prec_\mathcal{R} (1.0.0), (1.0.1) \succ_\mathcal{R} (1.1.1), (1.0.1) \succ_\mathcal{R} (0.0.1),$

$(1.0.1) \prec_\mathcal{R} (1.1.0), (1.0.1) \prec_\mathcal{R} (0.1.0);$

$(0.0.1) \succ_\mathcal{R} (1.1.1), (0, 0, 1) \prec_\mathcal{R} (0, 1, 0), (0.0.1) \prec_\mathcal{R} (1.1.0), (0.0.1) \prec_\mathcal{R} (1.0.0);$

$(1.1.0) \succ_\mathcal{R} (1.1.1), (1, 1, 0) \succ_\mathcal{R} (1, 0, 0), (1.1.0) \prec_\mathcal{R} (0.1.0);$

$(1.1.1) \prec_\mathcal{R} (0.1.0), (1.1.1) \prec_\mathcal{R} (1.0.0);$

$(0.1.0) \succ_\mathcal{R} (1.0.0).$

With the objects scheme $A_1 = \{\{H_1, H_2\}, \{H_3\}\}$ the social outcome $(0.0.0)$ is the unique global optimum, while with the objects scheme $A_2 = \{\{H_2, H_3\}, \{H_1\}\}$ the unique global optimum is $(0.1.1)$.

From now on A will be an objects scheme for a given social rule $\succ_\mathcal{R}$ on a social outcome space X that corresponds to the arrangement $\mathcal{A}_{n,m}$.

A fundamental question is whether and under which conditions, given a social outcome $z \in X$, it is possible to choose an objects scheme A such that the basin of attraction $\Psi(z, A)$ is not empty, i.e. such that z is a local optimum or a global one for some objects scheme and some agenda. The answers to these questions are given in the next two sections.

3.2   Objects and manipulability: local optimality

Suppose one of the agents has a preferred social outcome $z \in X$ and some form of power in the determination of the objects, e.g. because he or she is chairing a committee or can somehow influence the cognitive framing of the choice, for instance through persuasion (Mullainathan, Schwartzstein, and Shleifer 2008). An interesting question is whether there exists an objects scheme $A = \{\mathcal{A}_{I_1}, \ldots, \mathcal{A}_{I_k}\}$ such that z is a local optimum for A, and which combination of initial conditions and agenda will make social choice converge to z. The conditions under which this is possible will be provided in theorem 2 that we will prove after introducing notions of separation and distance in hyperplane arrangements.

Let x and y be two social outcomes. They are said to be *separated by an hyperplane* $H \in \mathcal{A}_{n,m}$ if H separates the chambers C_x and C_y. In this case, the notation $x \mid H \mid y$ will be used. Moreover, x and y are said to be *prominently separated* if there exist two hyperplanes $H_{i_1,j_1}, H_{i_2,j_2} \in \mathcal{A}_{n,m}$ with $i_1 \neq i_2$ (i.e. non-parallel) such that $x \mid H_{i_1,j_1} \mid y$ and $x \mid H_{i_2,j_2} \mid y$ hold. We will say that $x = (v_1. \cdots .v_n)$ and $y = (g_1. \cdots .g_n)$ are *separated by the feature* f_i if $v_i \neq g_i$, i.e. the value of the feature f_i of y differs from that of x. The set of the features that separate x and y is denoted by $\overline{\mathcal{H}}_{x,y}$.

Distance of social outcomes   The *distance* between x and y is the minimum number of hyperplanes that separate x from y. The *prominent distance* $d_p(x, y)$ is the number of features that separate x from y, i.e. $\#\overline{\mathcal{H}}_{x,y}$. Note that $d_p(x, y)$ equals the minimum number of hyperplanes that prominently separate x from y.

Recall that, by definition, if $H_{i,\bar{j}}$ belongs to the object $\mathcal{A}_I$ for some $\bar{j}$, then $H_{i,j}$ belongs to $\mathcal{A}_I$ for all $0 \leqslant j < m_i - 1$. Therefore, the subarrangement

$$\mathcal{H}_{x,y} = \left\{ H_{i,j} \in \mathcal{A}_{n,m} \mid i \in \overline{\mathcal{H}}_{x,y}, \ 0 \leqslant j < m_i - 1 \right\}$$

of $\mathcal{A}_{n,m}$ has been considered. Note that, if we have $d_p(x, y) = 1$ and $d(x, y) > 1$, all the hyperplanes in $\mathcal{H}_{x,y}$ are parallel.

Remark 1. The sets $\mathcal{H}_{x,y}$ and $\overline{\mathcal{H}}_{x,y}$ are strictly interconnected. For instance, we will use the fact that $\mathcal{H}_{x,y}$ is contained in $\mathcal{H}_{z,w}$ if and only if $\overline{\mathcal{H}}_{x,y}$ is contained in $\overline{\mathcal{H}}_{z,w}$.

From now on, for the sake of simplicity, we will consider only social rules with strict preferences. This restriction is almost always unnecessary, but it greatly simplifies the presentation. Therefore, from now on, $\succ_{\mathcal{R}}$ will always denote a social rule with complete strict preferences.

If z is a social outcome such that all social outcome $x_1, \ldots, x_k$ with $x_j \succ_{\mathcal{R}} z$ verify $d_p(x_j, z) > 1$, then x_j and z are prominently separated at least by two hyperplanes.

It follows that we can build an objects scheme A_z such that $\mathcal{H}_{x_j,z} \not\subseteq \mathcal{A}$ for all $\mathcal{A} \in A_z$ and all $1 \leq j \leq k$. For example, if $H_j^1, H_j^2 \in \mathcal{H}_{x_j,z}$ are two hyperplanes related to different features for $1 \leq j \leq k$, then let us consider an objects scheme A_z such that for any x_j there exist two objects $\mathcal{A}_j^1, \mathcal{A}_j^2$ in A_z with $H_j^1 \in \mathcal{A}_j^1$, $H_j^2 \in \mathcal{A}_j^2$ and $\{H_j^1, H_j^2\} \not\subseteq \mathcal{A}$ for all $\mathcal{A} \in A_z$.

It is obvious that such an objects scheme exists. Moreover z is a local optimum for A_z. Indeed for all $x_j \succ_{\mathcal{R}} z$ and for all $\mathcal{A} \in A_z$ the chambers $C_{x_j}(\mathcal{A}^c)$ and $C_z(\mathcal{A}^c)$ are always separated by H_j^1 or H_j^2. That is $x_j(\mathcal{A}^c) \neq z(\mathcal{A}^c)$ and then $\Phi(z, \mathcal{A}) = \emptyset$ for all $\mathcal{A} \in A_z$. It follows that $z \in \Psi(z, A_z)$ and then $z \in \Psi(z, A_{z,\alpha})$ for all agendas α, i.e. z is a local optimum.

On the other hand if x is a social outcome $x \succ_{\mathcal{R}} z$ such that $d_p(x, z) = 1$ then for any objects scheme A there is at least one object $\mathcal{A}$ such that all hyperplanes H separating x from z are in $\mathcal{A}$. Then, by definition, $x \in \Phi(z, \mathcal{A}) \neq \emptyset$. Hence we get the following easy fact.

Theorem 2. Let z be a social outcome. Then, there exists an objects scheme A_z for which z is a local optimum if and only if the inequality $d_p(w, z) > 1$ holds for any social outcome w with $w \succ z$.

Remark 3.1. Theorem 2 also tells us that the condition for a social outcome to be a local optimum does not depend on the agenda. If z is a local optimum for an agenda $\alpha \in \Lambda(A)$ it will be a local optimum also for any other agenda $\beta \in \Lambda(A)$, Thus object construction power is, in some sense, stronger than agenda power. Notice, however, that the basin of attraction of a local optimum is, in general, different for different agendas and, therefore, even when starting from the same initial social outcome the choice processes may end up in different local optima for different agendas.

Moreover in order to verify if a social outcome is a local optimum we need only to verify if it is greater than all the $n(m + 1) - 1$ social outcomes with prominent distance equal one from it, instead of checking all $(m+1)^n - 1$ social outcomes in X. The above construction implicitly gives also an algorithmic description of how to construct an objects scheme which makes social choice converge to a desired local optimum.

Finally, notice that the independence of our construction from the 1-dimensional distance, i.e. the distance along the 1-dimensional family of

hyperplanes $\{H_{i,j}\}_{0 \leq j \leq m-1}$ for a fixed i, is a consequence of the independence of the choice from the order in which a 1-dimensional list of objects is given.

Example 8. Let $\succ_{\mathcal{R}}$ be the social decision rule of example 7. Thanks to Theorem 2 it is very simple to verify that (0.0.0) and (0.1.1) are the only local optima for $\succ_{\mathcal{R}}$. Moreover, exploiting the above procedure, we can construct all objects schemes for which (0.0.0) and (0.1.1) are reachable. Hence we can notice that, if A_1 and A_2 are the objects schemes of example 7, then their basins of attraction coincide with X, $\Psi((0.0.0), A_1) = X$ and $\Psi((0.1.1), A_2) = X$, i.e. (0.0.0) and (0.1.1) are global optima for, respectively, object schemes A_1 and A_2. On the contrary, if we consider for instance the objects scheme $A_3 = \{\{H_1\}, \{H_2\}, \{H_3\}\}$ then (0.0.0) and (0.1.1) are both reachable local optima, each of them with its own basin of attraction such as $\Psi((0.0.0), A_3) \subset X$ and $\Psi((0.1.1), A_3) \subset X$.

If a social outcome z does not meet the necessary and sufficient condition of theorem 2 it cannot be a local optimum. However, an agent with object construction power may nevertheless make choices converge to another social outcome close enough to z.

Free social outcome We say that a social outcome z is *free* with respect to a social decision rule $\succ_{\mathcal{R}}$ if and only if for any social outcome x such that $x \succ_{\mathcal{R}} z$ then $d_p(x, z) > 1$. Thus, by the theorem 2, z is the local optimum for an objects scheme A_z if and only if z is free.

Moreover a social outcome $\bar{z}$ has *minimal distance* from z with respect to $\succ_{\mathcal{R}}$ if and only if $\bar{z}$ is free with respect to a decision rule $\mathcal{R}$ and

$$d(z, \bar{z}) = \mathtt{min}\{d(z, x) \mid x \text{ is free }\}$$

Thus it is always possible to build an objects scheme $A_{\bar{z}}$ such that the social outcome $\bar{z}$ of minimal distance from z with respect to $\succ_{\mathcal{R}}$ is a local optimum.

Remark 3.2. If we consider the classical 1-dimensional problem, then the prominent distance between two social outcomes x and z is always $d_p(x, z) = 1$. It follows:

- the social outcome z is free if and only if z is an optimum, i.e. for any social outcome x, $z \succ_{\mathcal{R}} x$;

- if the social outcome of minimal distance from z with respect to $\succ_{\mathcal{R}}$ exists then it is the only optimum;

- if the social outcome of minimal distance from z with respect to $\succ_{\mathcal{R}}$ does not exist, then our theorem simply recovers the usual intransitive cycles.

3.3   Objects and manipulability: global optimality

We now analyze under which conditions a social outcome z is a global optimum for an agenda α and under which conditions it is a global optimum for all agendas $\alpha \in \Lambda(A)$.

From now on, given a social outcome $z \in X$, we will denote by A_z an objects scheme such that $\Phi(z, A_z) = \emptyset$, i.e. it admits z as local optimum.

Let us remark that, given a social outcome $z \in X$, if z is a global optimum for an agenda $\alpha \in \Lambda(A_z)$ then the following two conditions hold:

1. if $x \in X$ is a social outcome such that $x \succ_{\mathcal{R}} z$ then there exist social outcomes $\{x_1, \ldots, x_k\}$ such that

$$z \prec_{\mathcal{R}} x \prec_{\mathcal{R}} x_1 \prec_{\mathcal{R}} \ldots \prec_{\mathcal{R}} x_k \prec_{\mathcal{R}} z$$

   Indeed if this is not the case, all paths starting from x would never end up in z for all possible objects schemes A; i.e. $x \notin \Psi(z, A)$ for any objects scheme A.

2. $\Phi(x, A_z) \neq \emptyset$ for all social outcomes $x \neq z$. Otherwise if $\Phi(x, A_z) = \emptyset$, then $x \in \Psi(x, A_{z,\alpha})$ for all agendas $\alpha \in \Lambda(A_z)$.

The first condition is simple to verify. The second one is true for all social outcomes $x \in X$ that are not free (see theorem 2), while free social outcomes must satisfy the following conditions.

Theorem 3. Let z be a free social outcome. Then, there exists an objects scheme A_z such that $\Phi(z, A_z) = \emptyset$ and $\Phi(x, A_z) \neq \emptyset$ for all free social outcomes x if and only if the condition

$$\exists y \succ x \text{ such that } \mathcal{H}_{w,z} \nsubseteq \mathcal{H}_{x,y} \quad \forall w \succ z \tag{4}$$

holds for all free social outcomes $x \neq z$.

Proof: *Sufficiency:* if $\Phi(x, A_z) \neq \emptyset$ then there exists a social outcome y in $\Phi(x, A_z)$. By construction y satisfies the following two conditions:

1. $y \succ_{\mathcal{R}} x$

2. there exists $\mathcal{A} \in A_z$ such that $\mathcal{H}_{x,y} \subset \mathcal{A}$ or, equivalently, $\mathcal{H}_{x,y}^c \supset \mathcal{A}^c$.

Moreover $\Phi(z, A_z) = \emptyset$ implies that for all objects $\mathcal{A} \in A_z$ and for all $w \succ_\mathcal{R} z$ the intersection $\mathcal{A}^c \cap \mathcal{H}_{w,z} \neq \emptyset$, then $\mathcal{H}_{x,y}^c \cap \mathcal{H}_{w,z} \neq \emptyset$, i.e. $\mathcal{H}_{w,z} \not\subseteq \mathcal{H}_{x,y}$ for all $w \succ_\mathcal{R} z$. Hence y is the social outcome verifying condition (4).

Necessity: let $G(z, A_z) \subset X$ be the set

$$G(z, A_z) = \{x \in X \mid \Phi(x, A_z) = \emptyset\}$$

i.e. if $x \in G(z, A_z)$ then it is a local optimum for A_z. If $G(z, A_z) = \emptyset$ we proved necessity, otherwise given $x \in G(z, A_z)$ we can build a new objects scheme A_z' as follows. Let $y \succ_\mathcal{R} x$ as in the hypothesis, then we can consider the new objects scheme

$$A_z' = A_z \cup \mathcal{H}_{x,y}$$

Clearly $y \in \Phi(x, \mathcal{H}_{x,y})$, i.e. x is not a local optimum for A_z' while z is still a local optimum for A_z'. Indeed $\Phi(z, \mathcal{A}) = \emptyset$ for all $\mathcal{A} \in A_z$ by construction while $\Phi(z, \mathcal{H}_{x,y}) = \emptyset$ by hypothesis:

let w be a social outcome in $\Phi(z, \mathcal{H}_{x,y}) \neq \emptyset$. Then $w \succ_\mathcal{R} z$ and $\mathcal{H}_{w,z} \subset \mathcal{H}_{x,y}$, but this is an absurd since $\mathcal{H}_{w,z} \not\subseteq \mathcal{H}_{x,y}$ for all $w \succ_\mathcal{R} z$ by condition (4).

We have a new objects scheme A_z' such that z is a local optimum for A_z' and $G(z, A_z') \subset G(z, A_z) \setminus \{x\}$. We can iterate this construction until we obtain an objects scheme A such that z is a local optimum and $G(z, A) = \emptyset$. $\square$

The above theorem gives us a necessary and sufficient condition for the existence of an objects scheme A_z such that for all possible agendas and all possible starting social outcomes $x \in X$ the voting process ends up in z or in a cycle. Following the proof of this theorem we can also provide an algorithm which computes such an objects scheme A_z (Amendola and Settepanella 2012). Moreover, given a starting social outcome $x \in X$, condition (4) allows us to construct an objects scheme A_z and an agenda α in A such that $x \in \Psi(z, A_{z,\alpha})$.

Example 9. Let us consider the social outcome $z = (0.0.0)$ and the social rule $\succ_\mathcal{R}$ as in example 7. Thanks to Theorem 3 it is easy to prove the existence of an objects scheme A such that $(0.0.0)$ is the only possible optimum.

Indeed the only social outcome x preferred to z, i.e. $x \succ_{\mathcal{R}} z$, is $x = (0.1.1)$ and the social outcome $(1.0.1) = y \succ_{\mathcal{R}} x$ verifies that

$$\mathcal{H}_{x,z} = \{H_2, H_3\} \nsubseteq \mathcal{H}_{x,y} = \{H_1, H_2\}.$$

Then, by theorem 3, there is an objects scheme A_z for which z is the only possible optimum. Hence or z is a global optimum for the objects scheme A_z, i.e. $\Psi(z, A) = X$, or there is a social outcome $x \in X$ such that any path through A_z starting from x ends up in a cycle.

The same argument applies for the social outcome $z = (0.1.1)$.

Example 10. Let us consider the space of social outcomes X as in example 1 with the following preferences $\succ_{\mathcal{R}}$:

$(0.0.0)$ preferred to all except $(0.1.1) \succ_{\mathcal{R}} (0.0.0)$

$(0.1.1)$ preferred to all except $(1.0.0) \succ_{\mathcal{R}} (0.1.1)$

$(1.0.1) \prec_{\mathcal{R}} (1.0.0), (1.0.1) \succ_{\mathcal{R}} (1.1.1), (1.0.1) \succ_{\mathcal{R}} (0.0.1),$

$(1.0.1) \prec_{\mathcal{R}} (1.1.0), (1.0.1) \prec_{\mathcal{R}} (0.1.0)$

$(0.0.1) \succ_{\mathcal{R}} (1.1.1), (0.0.1) \prec_{\mathcal{R}} (0.1.0), (0.0.1) \prec_{\mathcal{R}} (1.1.0), (0.0.1) \prec_{\mathcal{R}} (1.0.0)$

$(1.1.0) \succ_{\mathcal{R}} (1.1.1), (1.1.0) \succ_{\mathcal{R}} (1.0.0), (1.1.0) \prec_{\mathcal{R}} (0.1.0)$

$(1.1.1) \prec_{\mathcal{R}} (0.1.0), (1.1.1) \prec_{\mathcal{R}} (1.0.0)$

$(0.1.0) \succ_{\mathcal{R}} (1.0.0)$.

The social outcome $z = (0.0.0)$ is a local optimum by Theorem 2, but it is not a global optimum for any objects scheme. Indeed the social outcome $x = (0.1.1)$ is a free outcome and it verifies $x \succ_{\mathcal{R}} z$. Moreover $y = (1.0.0)$ is the only social outcome preferred to x, i.e. $y \succ_{\mathcal{R}} x$, and it verifies

$$\mathcal{H}_{x,z} = \{H_2, H_3\} \subset \mathcal{H}_{x,y} = \{H_1, H_2, H_3\}.$$

Then, by Theorem 3, it follows that if $(0.0.0)$ is a local optimum for an objects scheme A, then also $(0.1.1)$ is a local optimum for A and hence $(0.1.1) \notin \Psi((0.0.0), A)$.

3.4   Objects and manipulability: obtaining a social outcome from a given *status quo*

So far we have analyzed cases in which objects, agendas and initial *status quo* may all be manipulated in order to obtain a desired social outcome. In this section instead we assume that the initial *status quo* is exogenously given and only objects and agenda are subject to manipulation. In theorem 4 we

give a necessary and sufficient condition that a *status quo* x needs to satisfy in order to belong to the basin of attraction of a chosen social outcome z with respect to some objects scheme A. In other words, given an agent with preferred free social outcome z, we find the conditions under which he or she may obtain z from the given initial condition.

Universal basin of attraction Let $\Pi(\mathcal{A}_{n,m})$ be the set of all possible objects schemes in $\mathcal{A}_{n,m}$. The *universal basin of attraction* of a social outcome $z \in X$ is the set
$$\Psi(z) = \bigcup_{A \in \Pi(\mathcal{A}_{n,m})} \Psi(z, A),$$
i.e. the set of all the social outcomes x such that there exists an objects scheme for which there is a path starting from x and ending up in z.

By Theorem 2, the universal basin of attraction of the social outcome z is non-empty if and only if z is *free*.

As above, let $x \in X$ be a starting social outcome. Let us define the following set:

$$G_x^z = \{y \succ_{\mathcal{R}} x \mid \mathcal{H}_{w,z} \nsubseteq \mathcal{H}_{x,y} \quad \forall w \succ_{\mathcal{R}} z \text{ and } B(x, \mathcal{H}_{x,y}) \neq \emptyset\}. \tag{5}$$

Proposition 3.3. Given two social outcomes $x, z \in X$, if x is in the universal basin of attraction $\Psi(z)$ then $G_x^z \neq \emptyset$.

Proof. Let us assume that G_x^z is empty then either:

i) condition 4 is not met by any of the social outcomes $y \in X$. Therefore, by theorem 3, it follows that $\Phi(x, A_z) = \emptyset$ for all A_z such that $\Phi(z, A_z) = \emptyset$, i.e. if x is the starting social outcome then all maximal paths in A_z end up in x for any agenda α;

or:

ii) $\Phi(x, A_z) \neq \emptyset$ but $B(x, A_z) = \emptyset$ for all objects schemes A_z such that $\Phi(z, A_z) = \emptyset$, i.e. if x is the starting social outcome then all voting processes end up in a cycle.

In both cases $x \notin \Psi(z)$ and this concludes the proof $\qquad \square$

Suppose now that x is a social outcome such that G_x^z is non-empty (therefore z is free [4]). If $B(x, \mathcal{H}_{x,y})$ is non-empty, its cardinality is one. The only element of $B(x, \mathcal{H}_{x,y})$ will be denoted by $b_{x,y}$ and we can define the set

$$BG_x^z = \{b_{x,y} \mid y \in G_x^z\}.$$

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[4]For the sake of completeness, we assume G_x^z to be $\emptyset$ if z is not free.

Then we can consider the following finite subsets[5] in X:

$$E_0^z = \{z\};$$
$$E_1^z = \{x \in X \setminus \{z\} \mid z \in BG_x^z\};$$
$$E_2^z = \{x \in X \setminus \cup_{i=0}^1 E_i^z \mid E_1^z \cap BG_x^z \neq \emptyset\};$$
$$\vdots$$
$$E_h^z = \{x \in X \setminus \cup_{i=0}^{h-1} E_i^z \mid E_{h-1}^z \cap BG_x^z \neq \emptyset\};$$
$$E_{h+1}^z = \{x \in X \setminus \cup_{i=0}^h E_i^z \mid E_h^z \cap BG_x^z \neq \emptyset\} = \emptyset$$

where h is the smallest integer such that E_{h+1}^z is empty, and:

$$E^z = \bigcup_{i=1}^h E_i^z. \tag{6}$$

We have the following:

Theorem 4. Let x and z be two social outcomes. Then x is in the universal basin of attraction $\Psi(z)$ if and only if x belongs to E^z, i.e.

$$\Psi(z) = E^z.$$

Proof. *Sufficiency:* the proof is by induction. If $x \in \Psi(z)$, then there exists an objects scheme $A = \{\mathcal{A}_{I_1}, \ldots, \mathcal{A}_{I_t}\}$ such that $x \in \Psi(z, A)$ and an agenda $\alpha = (h_k, \ldots, h_1)$ for A such that the ordered path along α

$$x = x_k \prec_{\mathcal{R}} x_{k-1} \prec_{\mathcal{R}} \ldots \prec_{\mathcal{R}} x_1 \prec_{\mathcal{R}} x_0 = z$$

is maximal.
The social outcome $x_0 = z$ is in E^z and, by induction, we assume that there is a $0 \le j \le k$ such that $x_j \in E^z$, i.e. $x_j \in E_i^z$ for an integer $0 \le i \le h$. By construction $x_j \in B(x_{j+1}, \mathcal{A}_{h_{j+1}})$ then $x_j \in B(x_{j+1}, \mathcal{H}_{x_{j+1}, x_j}) \neq \emptyset$. Moreover $\mathcal{H}_{w,z} \nsubseteq \mathcal{H}_{x_{j+1}, x_j}$ $\forall w \succ_{\mathcal{R}} z$, indeed if $w \succ_{\mathcal{R}} z$ is a social outcome such that $\mathcal{H}_{w,z} \subseteq \mathcal{H}_{x_{j+1}, x_j} \subseteq \mathcal{A}_{h_j}$ then $z \notin \Psi(z, A)$ which is absurd. It follows that $x_j \in BG_{x_{j+1}}^z$.

Since $x_j \in E_i^z \cap BG_{x_{j+1}}^z \neq \emptyset$ then either $x_{j+1} \in E_{i+1}^z$ or $x_{j+1} \in \cup_{s=0}^{i-1} E_s^z$, i.e. $x_{j+1} \in E^z$.

Necessity: viceversa, if $x \in E^z$ then there is an integer $0 \le i \le h$ such that $x \in E_i^z$ and we can construct a path

$$x = x_i \prec_{\mathcal{R}} x_{i-1} \prec_{\mathcal{R}} \ldots \prec_{\mathcal{R}} x_1 \prec_{\mathcal{R}} x_0 = z \tag{7}$$

_____

[5]Once again, for the sake of completeness, we assume all these sets to be empty if z is not free.

such that $x_j \in E_j^z \cap BG_{x_{j+1}}$ for $0 \le j \le i-1$.

Let us consider an objects scheme A_z such that $z \in \Psi(z, A_z)$, then

$$A = \{\mathcal{H}_{x_j, x_{j+1}}\}_{0 \le j \le i-1} \cup A_z$$

still satisfies $z \in \Psi(z, A)$.

Moreover let $\alpha = (h_1, \ldots, h_k)$ be an agenda for A such that the h_t-object in the ordered scheme A_α is $\mathcal{H}_{x_{i-t-1}, x_{i-t}}$ for $0 \le t \le i-1$. Then, by construction, the path (7) is exactly the maximal path $P(x, z, A_\alpha)$ and $x \in \Psi(z)$. $\qquad \square$

Example 11. Let us consider again the social outcome $z = (0.0.0)$ and the social decision rule $\succ_{\mathcal{R}}$ as in example 7. The social outcome $w = (0.1.1)$ is the only outcome preferred to z and $\mathcal{H}_{w,z} = \{H_2, H_3\}$. It is easy to verify that:

$$E_0^z = \{(0.0.0)\};$$
$$E_1^z = X \setminus \{(0.0.0), (0.1.1), (1.1.1)\};$$
$$E_2^z = \{(0.1.1), (1.1.1)\}.$$

Then the universal basin of attraction $\Psi(z)$ is the whole set X.

With the above notations, if we consider the social rule $\succ_{\mathcal{R}}$ defined in example 10, then $G_w^z = \emptyset$ and hence $w = (0.1.1) \notin E^z$.

3.5   Breaking intransitive cycles

In this section we show that also intransitive cycles which, as well known, can be generated by any social aggregation rule are subject to manipulation through object construction. The main result of this section, theorem 5, demonstrates that almost any intransitive cycle can be broken by appropriate modification of the objects scheme. It actually turns out that cycles can be broken by introducing new local optima. Thus, we observe a trade-off between decidability, i.e. the possibility of reaching some social optimum in a feasible time, and non-manipulability, that is, the convergence of the social decision process to a unique global outcome that does not depend on initial condition and agenda. The balance in this trade-off is struck by the objects scheme. In particular, we will show that coarse objects containing many features tend to produce many cycles and few local optima, whereas fine objects containing only one or a few features are much less likely to induce cycles but tend to generate many local optima and, therefore, greatly increase the opportunity for manipulation.

Suppose that for some objects scheme A' social preferences may encounter a set of cycles Γ. We will say that an objects scheme $A \ne A'$ *breaks* the cycles

Γ if social choice does not enter into any cycle belonging to Γ with the new objects scheme.

Given a cycle $\gamma : x_1 \prec_{\mathcal{R}} \ldots \prec_{\mathcal{R}} x_k \prec_{\mathcal{R}} x_{k+1} = x_1$, let us define the following subset of hyperplanes in $\mathcal{A}_{n,m}$:

$$\mathcal{H}_\gamma = \bigcup_{1 \leq j \leq k+1} \mathcal{H}_{x_j, x_{j+1}}. \tag{8}$$

We will say that a cycle $\gamma : x_1 \prec_{\mathcal{R}} \ldots \prec_{\mathcal{R}} x_k \prec_{\mathcal{R}} x_{k+1} = x_1$ is *free* if there are at least two consecutive indices $j, j+1$ such that $d_p(x_j, x_{j+1}) > 1$.

Then we have the following:

Theorem 5. Let $\mathcal{Y}_{\succ_{\mathcal{R}}}$ be the oriented graph for a given social rule $\succ_{\mathcal{R}}$ and $\mathcal{A}_{n,m}$ the arrangement associated to the space of social outcomes. Then for any given set $\Gamma = \{\gamma_1, \ldots, \gamma_h\}$ of free cycles in $\mathcal{Y}_{\succ_{\mathcal{R}}}$ there exists an objects scheme A such that all voting process will never end up in any cycle of Γ. Moreover this scheme can be chosen with maximum size.

In order to prove this theorem we remark that if $\gamma \in \mathcal{Y}_{\succ_{\mathcal{R}}}$ is a cycle and A an objects scheme such that:

1. $\mathcal{H}_\gamma \nsubseteq \mathcal{A}$ for any object $\mathcal{A} \in A$,

2. any path ordered along α and starting in x does not end up in γ for any agenda α,

then γ disappears in A.

Lemma 3.4. Let $\gamma \in \mathcal{Y}_{\succ_{\mathcal{R}}}$ be a cycle and A an objects scheme such that at least two consecutive outcomes $x_j \prec_{\mathcal{R}} x_{j+1}$ in the cycle satisfy $\mathcal{H}_{x_j, x_{j+1}} \nsubseteq \mathcal{A}$ for any object $\mathcal{A} \in A$. Then γ disappears in A.

Proof. Condition 1 above is trivially satisfied. Moreover, by hypothesis, let $x_j \prec_{\mathcal{R}} x_{j+1}$ be two consecutive outcomes in γ such that $\mathcal{H}_{x_j, x_{j+1}} \nsubseteq \mathcal{A}$ for any object $\mathcal{A} \in A$. Then, by definition, $x_{j+1} \notin B(x_j, \mathcal{A})$ for all objects $\mathcal{A} \in A$, i.e. it is not possible to move from x_j to x_{j+1} directly and then condition 2 holds. This conclude the proof. $\square$

As a direct consequence of the above Lemma any free cycle $\gamma \in \mathcal{Y}_{\succ_{\mathcal{R}}}$ disappears for the trivial object scheme

$$A = \{\{H_{i,j}\}_{\substack{1 \leq i \leq n \\ 0 \leq j \leq m-1}}\} \quad .$$

Proof of Theorem 5. Let us consider the finite set $\mathbb{A}_\Gamma$ of all possible objects schemes A of $\mathcal{A}_{n,m}$ such that all voting processes in A will never end up in any cycle in Γ. This set is not empty, indeed the trivial objects scheme

$$A = \left\{ \{H_{i,j}\}_{\substack{1 \leq i \leq n \\ 0 \leq j \leq m-1}} \right\}$$

is such that any social choice process will never end up in any cycle in Γ. Moreover the set of the sizes of elements $A \in \mathbb{A}_\Gamma$ is a finite non-empty set of natural numbers $\mathbb{N}$ and then it admits a maximum.
This conclude the proof. $\qquad \square$

Example 12. In the example 6 the objects schemes $A = \{\{H_1, H_3\}, \{H_2\}\}$ and $A = \{\{H_2, H_3\}, \{H_1\}\}$ break at once all cycles in the graph $\mathcal{Y}_{\succ_\mathcal{R}}$ of example 3.

Remark 3.5. A cycle γ which is not free can also be broken in most cases, but not always. An example is given by the cycle $\gamma : 00 \prec_\mathcal{R} 01 \prec_\mathcal{R} 11 \prec_\mathcal{R} 10 \prec_\mathcal{R} 00$ in the case $\mathcal{A}_{2,1}$. Indeed, in this case, the only two objects schemes are $A = \{\{H_1, H_2\}\}$ and $A = \{\{H_1\}, \{H_2\}\}$ and none of them breaks γ.

The above theorem proves the existence of a non trivial objects scheme A such that all the cycles in a given set $\Gamma = \{\gamma_1, \ldots, \gamma_h\}$ are *broken*, while Lemma 3.4 shows how to construct it.

All in all, a cycle γ may be *broken*, i.e. it disappears for a given objects scheme A, if there exist consecutive social outcomes x, w involved in the cycle γ such that $\mathcal{H}_{x,w} \nsubseteq \mathcal{A}$ for all $\mathcal{A}$ belonging to the scheme A. Thus, the number of cycles broken by a given objects scheme A increases when the size of A, i.e. the cardinality of its largest object, decreases.

Symmetrically, the number of local optima increases when the size of the objects scheme A decreases. Indeed by theorem 2, a free social outcome z is a local optimum for a scheme A if and only if for all social outcomes $x \succ_\mathcal{R} z$ $\mathcal{H}_{x,z} \nsubseteq \mathcal{A}$, for all $\mathcal{A}$ belonging to the scheme A. It follows that if the objects scheme A is composed by k objects, then we have at least $(k-1)!$ different local optima. Thus we obtain a decidability vs. non-manipulability trade-off. Schemes composed of few large objects tend to produce cycles, while schemes of many small objects tend to produce considerably fewer cycles but increasingly many local optima.

Remark 3.6. From what we have seen until now, it is obvious that, in order to work with the object construction power, it is helpful to find the hyperplanes involved in the fundamental cycles that generate all the others.
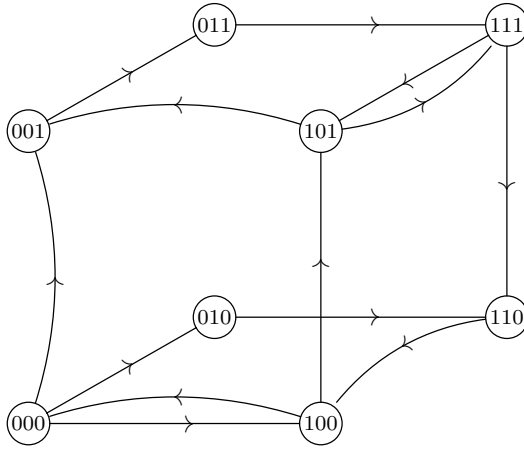
Figure 3: Reduced graph from example 3

In this respect, given an oriented graph $\mathcal{Y}_{\succeq_\mathcal{R}}$, it is helpful to consider a reduced homotopic one. Hence, for example, if $(x, y) \in Y_{1, \succeq_\mathcal{R}}$ is an edge given by a formal sum with coefficient 1 of edges which are in $Y_{1, \succeq_\mathcal{R}}$ then it can be deleted. So the graph in figure 1 strongly simplifies in that in figure 3[6].

To be precise, the graph in figure 3 does not exactly represent Salvetti's one, as the latter is in the complexification of $\mathcal{A}_{3,2}$, i.e. in $\mathbb{C}^3 \simeq \mathbb{R}^6$, which allows paths to "go around" hyperplanes (e.g. the line and the arc joining 000 and 100 form, in the complex space, a loop around the hyperplane of equation $\lambda_1 = 0$). However figure 3 gives a useful, though not totally precise, visual representation.

Looking at this figure we remark that it is not possible to retrieve the social rule from the graphic representation of the Salvetti's graph: we don't know if agents prefer the social outcome 101 to 111 or viceversa; we only know that there is a cycle involving these two preferences. In order to get the precise information about two social outcomes we have to look at the rule.

Nevertheless the reduced graph carries very useful information. For example it is relatively simple to find the fundamental cycles that generate all the others. Looking at the graph in fig. 3 it is clear that, for this base, the

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[6]For simplicity reason we omitted, in the figure, the hyperplanes, i.e. the coordinate hyperplanes of equations $\lambda_i = 0$, $i = 1, 2, 3$.

fundamental cycles are:

$$001 \succ_{\mathcal{R}} 011 \succ_{\mathcal{R}} 111 \succ_{\mathcal{R}} 001,$$
$$000 \succ_{\mathcal{R}} 010 \succ_{\mathcal{R}} 110 \succ_{\mathcal{R}} 000,$$
$$000 \succ_{\mathcal{R}} 011 \succ_{\mathcal{R}} 111 \succ_{\mathcal{R}} 110 \succ_{\mathcal{R}} 000,$$
$$000 \succ_{\mathcal{R}} 100 \succ_{\mathcal{R}} 101 \succ_{\mathcal{R}} 110 \succ_{\mathcal{R}} 000. \tag{9}$$

By construction, the first one depends on the existence of the edge $(111, 001)$, while the other three depend on $(110, 000)$. This information allows us to determine which set of objects can break the cycles in the given social rule.

Clearly, there are many ways we can reduce the set $Y_{1, \succeq_{\mathcal{R}}}$ to a basic number of elements, and this number is not unique. However, all the reduced graphs represent exactly the same pairwise social preferences.

3.6 On the median voter theorem

In addition to intransitive cycles and agenda power, also another classic result of social choice theory, the median voter theorem, is subject to manipulability through object construction. By appropriately modifying the objects scheme, the outcome of a social decision may be made as distant as one wishes from the median voter's preferred social outcome.

Recall that the median voter theorem in its strong version (Black 1958, Downs 1957) states that if a voter with median preferences exists, her/his most preferred outcome will beat any other alternative in pairwise majority voting.

Indeed given a rule $\succ_{\mathcal{R}}$ with a median voter m, if z is the opposite social outcome with respect to the median voter's preferred one and z is free, then it is possible to find a scheme with preferred outcome z.

This is equivalent to saying that in our construction it is always possible to manipulate the objects scheme in such a way that the median voter theorem does not apply and social choice may converge to social outcomes very distant from the median voter's preferred one.

A   Appendix: Salvetti's complex

In Section 2 we described the cells of a complex. Obviously there is a boundary map on Salvetti's complex which is also very informative [7]. We will

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[7]The boundary map is a well know map in mathematics and it is a basic ingredient of Homology Theory

introduce the boundary only for the 1-skeleton of the complex $\mathcal{S}(\mathcal{A}_{n,m})$.
Let us consider the free abelian groups $S_1(\mathcal{A}_{n,m})$ and $S_0(\mathcal{A}_{n,m})$ generated by $\mathcal{S}_1(\mathcal{A}_{n,m})$ and $\mathcal{S}_0(\mathcal{A}_{n,m})$ considering all the formal sums with integer coefficients. The boundary map $\partial : S_1(\mathcal{A}_{n,m}) \longrightarrow S_0(\mathcal{A}_{n,m})$ is defined as follows:

$$\partial(x_i, x_j) = x_j - x_i$$

for any generator $(x_i, x_j) \in \mathcal{S}_1(\mathcal{A}_{n,m})$ and the map extends to $S_1(\mathcal{A}_{n,m})$ by linearity, i.e. if $c = \partial \sum_{j=1}^{k-1} a_j(x_{i_j}, x_{i_{j+1}})$, $a_j \in \mathbb{Z}$, is an element in $S_1(\mathcal{A}_{n,m})$, then

$$\partial c = \partial \sum_{j=1}^{k-1} a_j(x_{i_j}, x_{i_{j+1}}) = \sum_{j=1}^{k-1} a_j \partial(x_{i_j}, x_{i_{j+1}}).$$

Given a social rule $\succeq_\mathcal{R}$ we obtain a graph $\mathcal{Y}_{\succeq_\mathcal{R}}$ which is, as we have seen in Section 2, a subcomplex of Salvetti's complex in 0 and 1-dimensions. Then we can compute its boundary map. We may begin noticing that pairwise preferences correspond to formal sums with coefficients 1 in $Y_{1,\succ_\mathcal{R}}$ and that if two such formal sums are equivalent from a social point of view, i.e. they correspond to the same pairwise preferences, then they have the same boundary. Indeed, by linearity, the boundary of the sum in (1) is given by:

$$\partial \sum_{j=1}^{k-1}(x_j, x_{j+1}) = \sum_{j=1}^{k-1} \partial(x_j, x_{j+1}) =$$
$$\sum_{j=1}^{k-1} x_{j+1} - \sum_{j=1}^{k-1} x_j = x_k - x_1 = \partial(x_1, x_k),$$

namely it depends only by its initial and final vertices.
Recall that a cycle $\gamma_I \in \mathcal{Y}$ is a subgraph with vertices $\{x_1, \ldots, x_h\}$ and edges $\{(x_1, x_2), (x_2, x_3), \ldots, (x_h, x_1)\}$. Hence the boundary $\partial \gamma_I$ of a cycle is equal to 0. Indeed it is the formal sum with coefficient 1 of the edges in the subgraph γ_I with same initial and final vertex x_1, i.e.

$$\partial \gamma_I = \partial(x_1, x_2) + \partial(x_2, x_3) + \ldots + \partial(x_h, x_1) =$$
$$x_2 - x_1 + x_3 - x_2 + \ldots + x_1 - x_h = 0.$$

Then to study the boundary map of a graph $\mathcal{Y}_{\succeq_\mathcal{R}}$ is equivalent to study the Condorcet-Arrow's cycles of the corresponding social rule. The boundary operator transforms a geometric problem into an algebraic one and makes it possible to develop an algorithm that computationally performs the reduction of a given directed graph and allows to find cycles.
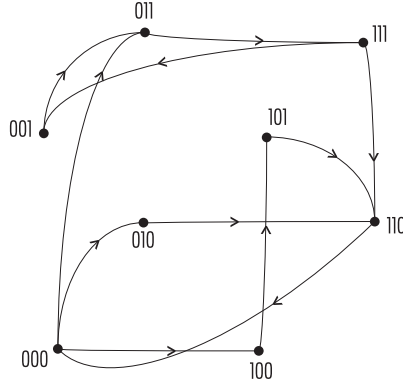
Figure 4: Reduced graph with a base of 11 elements

Example 13. In example 3 the boundary of each edge in $Y_{1,\succ_{\mathcal{R}}}$ is

$$\partial(v_1 v_2 v_3, g_1 g_2 g_3) = g_1 g_2 g_3 - v_1 v_2 v_3.$$

We get a boundary matrix of 28 columns, one for each edge, and 8 rows, one for each vertex, with two entries equal 1 and -1 (those corresponding to the vertices of the edge in question) in each column and all the other entries equal to 0. Now our algorithmic implementation must simply delete any column which is the sum of two or more other columns. In our example, we can get the following reduced matrix:

$$\begin{bmatrix}
-1 & -1 & -1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & -1 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & -1 & 0 & 0 \\
1 & 0 & 0 & 0 & 1 & -1 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & -1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & -1 \\
0 & 0 & 0 & -1 & 0 & 0 & 1 & 0 & 1 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & -1 & -1 & 0 & 0 & 0
\end{bmatrix}$$

where the columns are ordered as follows:

$\{(000, 011), (000, 010), (000, 100), (110, 000), (001, 011), (011, 111),$
$(111, 110), (111, 001), (010, 110), (100, 101), (101, 110)\}$

and the rows are ordered lexicographically. If the sum, with coefficient 1, of three or more columns is the zero column, then we have a cycle. Indeed, for

example, the sum of the 5-th 6-th and 8-th columns is 0 and it corresponds to the first cycle described in (9). Hence from the above matrix it is computational easy to retrieve all cycles. Moreover we get that the graph in figure 1 can be reduced to the one in figure 4 where cycles are easier to be noticed.

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