## СвМ <br> Discussion paper


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Tilburg University

Center<br>for<br>Economic Research

# No. 9846 <br> A NOTE ON GAMES CORRESPONDING TO SEQUENCING SITUATIONS WITH DUE DATES 

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May 1998

# A NOTE ON GAMES CORRESPONDING TO SEQUENCING SITUATIONS WITH DUE DATES 

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#### Abstract

It is shown that sequencing situations in which all jobs have equal processing times, the due date date of each job is a multiple of its processing time and the cost of each job is linear in its completion time, yield the same class of convex games as the sequencing situations in in which all jobs have equal processing times, the ready time of each job is a multiple of its processing time and the cost of each job is linear in its completion time.


Journal of Economic Literature Classification Number: C71

## Keywords:

Convex cooperative games, one-machine sequencing situations, due dates, ready times.

In a one-machine sequencing situation there is a queue of agents, each with one job, before a machine. Each job has to be processed on the machine. The finite set of agents is denoted by $N$ and $|N|=n$. By a bijection $\sigma: N \rightarrow\{1, \ldots, n\}$ we can describe the position of the agents in the queue. Specifically, $\sigma(i)=j$ means that player $i$ is in position $j$. The due date $d_{i}$ of the job of agent $i$ is the latest time the processing of this job has to be completed. The processing time $p_{i}$ of the job of agent $i$ is the time the machine takes to handle this job. We assume that every agent has a linear cost function $c_{i}:[0, \infty) \rightarrow I R$ defined by $c_{i}(t)=\alpha_{i} t$ with $\alpha_{i}>0$ the cost coefficient of player $i$. The completion time $C(\sigma, i)$ of the job of agent $i$ if processed according to a bijection $\sigma$ (in a semi-active way) is the sum of the earliest time the job can start w.r.t. $\sigma$ and its processing time. In this note we concentrate on sequencing situations that satisfy
(A1) $\quad d_{i} \in\{1, \ldots, n\}$ and $p_{i}=1$ for all $i \in N$
Further, it is assumed that there is an initial bijection $\sigma_{0}: N \rightarrow\{1, \ldots, n\}$ on the jobs of the players before the processing of the machine starts with the properties
(A2) $\quad d_{i} \leq d_{j}$ for all $i, j \in N$ with $\sigma_{0}(i)<\sigma_{0}(j)$, and $C\left(\sigma_{0}, i\right) \leq d_{i}$ for all $i \in N$ and
(A3) $\quad \sigma_{0}(i)=C\left(\sigma_{0}, i\right)$ for all $i \in N$.
Note that the assumptions $(A 1)-(A 2)$ imply that in the initial bijection there is no time gap in the job processing and that in particular the last job that is processed

[^0]according to $\sigma_{0}$ is completed at time $n$. In spite of the conclusion that assumption (A3) is superfluous, we have added it here for the sake of convenience and symmetry with ready time sequencing situations discussed later on. A sequencing situation as described above is denoted by $\left(N, \sigma_{0}, d, p, \alpha\right)$ and will be refered to as a $d$-sequencing situation.

The total costs $c_{\sigma}(S)$ of a coalition $S \subseteq N$ w.r.t. a bijection $\sigma$ are given by

$$
c_{\sigma}(S):=\sum_{i \in S} \alpha_{i}(C(\sigma, i))
$$

The (maximal) cost savings of a coalition $S$ depend on the set of admissible rearrangements of this coalition. Since each job has to be completed before its due date, we will consider only those $\sigma: N \rightarrow\{1, \ldots, n\}$ that satisfy $C(\sigma, i) \leq d_{i}$. Such a bijection $\sigma: N \rightarrow\{1, \ldots, n\}$ will be called admissible for $S$ if it satisfies $P(\sigma, i)=P\left(\sigma_{0}, i\right)$ for all $i \in N \backslash S$, where $P(\sigma, i)=\{j \in N \mid \sigma(j)<\sigma(i)\}$. Hence, we consider an order to be admissible for $S$ if each agent outside $S$ has the same starting time as in the initial order. Moreover, the agents of $S$ are not allowed to jump over players outside $S$. The set of all admissible rearrangements for a coalition $S$ is denoted by $\Sigma_{S}$. Note that by the assumptions on the initial and admissible bijections we have for any $\sigma \in \Sigma_{S}$ that $\sigma(i)=C(\sigma, i)$ for all $i \in N$.

Given a sequencing situation ( $N, \sigma_{0}, d, p, \alpha$ ) the corresponding sequencing game is defined in such a way that the the worth of a coalition $S$ is equal to the maximal cost savings the coalition can achieve by means of admissible rearrangements. Formally we have

$$
\begin{equation*}
v(S)=\max _{\sigma \in \Sigma_{S}}\left\{\sum_{i \in S} \alpha_{i} C\left(\sigma_{0}, i\right)-\sum_{i \in S} \alpha_{i} C(\sigma, i)\right\} \tag{1}
\end{equation*}
$$

From the definition of admissible rearrangements it follows that the essential coalitions for sequencing games are the connected ones. A coalition $S$ is called connected with respect to $\sigma_{0}$ if for all $i, j \in S$ and $k \in N, \sigma_{0}(i)<\sigma_{0}(k)<\sigma_{0}(j)$ implies $k \in S$.

Next, we describe the special class of one-machine sequencing situations, in which all jobs have equal processing times and the ready time of each job is a multiple of the processing time and the corresponding class of games. The description of these sequencing games is identical to the sequencing situations corresponding to due dates. The only difference is that there is no due date imposed on a player but a ready time. The ready time $r_{i}$ of the job of agent $i$ is the earliest time that the job can be processed on the machine. We will concentrate on sequencing situations that satisfy
(B1) $\quad r_{i} \in\{0, \ldots, n-1\}$ and $p_{i}=1$ for all $i \in N$.
The initial order $\sigma_{0}$ has the properties
(B2) $\quad r_{i} \leq r_{j}$ for all $i, j \in N$ with $\sigma_{0}(i)<\sigma_{0}(j)$ and $C\left(\sigma_{0}, i\right) \geq r_{i}+1$ for all $i \in N$ and

$$
\text { (B3) } \quad \sigma_{0}(i)=C\left(\sigma_{0}, i\right) \text { for all } i \in N \text {. }
$$

Note that the assumptions $(B 1)-(B 3)$ imply that in the initial bijection $\sigma_{0}$ there are no time gaps in the job processing and that the job that is processed last is completed at time $n$. A sequencing situation as described above is denoted by ( $N, \sigma_{0}, r, p, \alpha$ ) and will be refered to as an $r$-sequencing situation.
In $r$-sequencing situations we will only consider those bijections $\sigma: N \rightarrow\{1, \ldots, n\}$ that satisfy $C(\sigma, i) \geq r_{i}+1$ for all $i \in N$. The set of admissible rearrangements, denoted by
$\mathcal{A}_{S}$, has the same restrictions with respect to interchanging positions between players of a coalition $S$ as before. Hence, we may again conclude that for any $\sigma \in \mathcal{A}_{S}$ we have that $\sigma(i)=C(\sigma, i)$. The corresponding sequencing game is defined by

$$
\begin{equation*}
v(S)=\max _{\sigma \in \mathcal{A}_{S}}\left\{\sum_{i \in S} \alpha_{i} C\left(\sigma_{0}, i\right)-\sum_{i \in S} \alpha_{i} C(\sigma, i)\right\} \tag{2}
\end{equation*}
$$

Hamers, Borm and Tijs (1995) show that sequencing games arising from $r$-sequencing situations are convex by establishing relations between optimal orders of subcoalitions. These relations are obtained by analysing the procedure described in Rinnooy Kan (1976) that provides an optimal order. For the optimal order in d-sequencing situations we can use the procedure of Smith (1956), which operates similar to the procedure of Rinnooy Kan (1976). Both procedures aim for having the jobs with the largest cost coefficient $\alpha_{i}$ as far as possible at the front of the queue. The Smith-procedure has to take into account the due dates, whereas the Rinnooy Kan-procedure has to take into account the ready times. For this reason the Smith-procedure starts at the end of the queue, whereas the the Rinnooy Kan-procedure starts at the front of the queue. In spite of this difference it is possible for $d$-sequencing situations to establish similar relations between optimal orders of various subcoalitions as for $r$-sequencing situations. However, where in the Rinnooy Kan-procedure these relations are established if a player is added at the end of a (sub)queue, in the Smith-procedure these relations can be established if a player is added at the front of a (sub)queue. Following exactly the same line of argument it can be infered that sequencing games arising from $d$-sequencing situations are convex games.

In fact, we will show even a stronger result: both classes of sequencing situations generate the same class of sequencing games.
Theorem 1 Let $R(N)$ and $D(N)$ be the class of sequencing games that arise from $r$ sequencing situations and d-sequencing situations, respectively. Then $R(N)=D(N)$.
Proof: We show that $R(N) \subseteq D(N)$. Let $(N, v) \in R(N)$. Let $\left(N, \sigma_{0}, r, p, \alpha\right)$ be an $r$-sequencing situation that generates the game $(N, v)$. W.l.o.g. we can take $\sigma_{0}(i)=i$ for all $i \in N$. Let $S=\{i, i+1, \ldots, j\}$, be a connected set w.r.t. $\sigma_{0}$. Then

$$
\begin{equation*}
v(S)=\max \left\{\sum_{k=i}^{j} \alpha_{k} k-\sum_{k=i}^{j} \alpha_{k} x_{k} \mid x_{k} \geq r_{k}+1 \forall k \in S,\left\{x_{i}, \ldots, x_{j}\right\}=\{i, \ldots, j\}\right\} \tag{3}
\end{equation*}
$$

Consider the $d$-sequencing situation $\left(N, \tau_{0}, d, p, \beta\right)$ in which for all $i \in N$ we define $\tau_{0}(i)=n+1-i, d_{i}=n-r_{i}$ and $\beta_{i}=c+\left(\alpha_{n}-\alpha_{i}\right)$ with $c=\max _{i \in N} \alpha_{i}$.
We first show that $\left(N, \tau_{0}, d, p, \beta\right)$ satisfies the assumptions $(A 1)-(A 3)$. Obviously, (A3) is a consequence of (B1), while (A1) follows immediately from the definition of $d$ and (B1). If $\tau_{0}(l)<\tau_{0}(m)$ then $m<l$ which implies that $r_{m} \leq r_{l}$. The definition of $d$ yields immediately that $d_{l} \leq d_{m}$. Further, we have for any $l \in N$ that $\sigma_{0}(l)=l \geq r_{l}+1=$ $n+1-d_{l}$. This implies that $d_{l} \geq n+1-l=\tau_{0}(l)=C\left(\tau_{0}, l\right)$. Hence (A2) is satisfied. Note that from the definition of $\tau_{0}$ it follows that $S$ is also connected w.r.t. $\tau_{0}$. Then for the game $(N, w)$ corresponding to ( $N, \tau_{0}, d, p, \beta$ ) it holds that

$$
\begin{align*}
& w(S)=\max \left\{\sum_{k=i}^{j} \beta_{k}(n+1-k)-\sum_{k=i}^{j} \beta_{k} y_{k} \mid y_{k} \leq d_{k} \forall k \in S\right. \\
&\left\{y_{i}, \ldots, y_{j}\right\}=\{n+1-j, \ldots, n+1-i\}\} \tag{4}
\end{align*}
$$

Let $\hat{y}$ be an optimal solution of (4). By defining $\hat{x}$ by $\hat{x}_{k}=n+1-\hat{y}_{k}$ for all $k \in\{i, \ldots j\}$ we have

$$
\begin{aligned}
w(S) & =\sum_{k=i}^{j} \beta_{k}(n+1-k)-\sum_{k=i}^{j} \beta_{k} \hat{y}_{k} \\
& =\sum_{k=i}^{j}\left(c+\alpha_{n}-\alpha_{k}\right)(n+1-k)-\sum_{k=i}^{j}\left(c+\alpha_{n}-\alpha_{k}\right)\left(n+1-\hat{x}_{k}\right) \\
& =\left(c+\alpha_{n}\right) \sum_{k=i}^{j}\left(\hat{x}_{k}-k\right)+\sum_{k=i}^{j} \alpha_{k}\left(k-\hat{x}_{k}\right) \\
& =\sum_{k=i}^{j} \alpha_{k}\left(k-\hat{x}_{k}\right) \\
& \leq v(S)
\end{aligned}
$$

where the first equality holds since $\hat{y}$ is optimal, the second equality by the definition of $\tau_{0}, \beta$ and $\hat{x}$, the third equality and fourth equality by straightforward calculations. The inequality holds by (3) since $\hat{x}_{k}=n+1-\hat{y}_{k} \geq n+1-d_{k}=n+1-\left(n-r_{k}\right)=r_{k}+1$ and $\left\{\hat{x}_{i}, \ldots, \hat{x}_{j}\right\}=\{i, \ldots, j\}$.
Let $\hat{x}$ be an optimal solution of (3). By defining $\hat{y}$ by $\hat{y}_{k}=n+1-\hat{x}_{k}$ for all $k \in S$ we can show in the same way as above that $v(S) \leq w(S)$, which completes the first part of this proof.
Obviously, the second part, $D(N) \subseteq R(N)$, can be dealt with in an analogous way.

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