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TREE-AUTOMATIC WELL-FOUNDED TREES *

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ABSTRACT. We investigate tree-automatic well-founded trees. Using Delhomme’s decomposition technique for tree-automatic structures, we show that the (ordinal) rank of a tree-automatic well-founded tree is strictly below ω^ω . Moreover, we make a step towards proving that the ranks of tree-automatic well-founded partial orders are bounded by ω^{ω^ω} : we prove this bound for what we call upwards linear partial orders. As an application of our result, we show that the isomorphism problem for tree-automatic well-founded trees is complete for level $\Delta_{\omega^\omega}^0$ of the hyperarithmetical hierarchy with respect to Turing-reductions.

1. INTRODUCTION

Various classes of infinite but finitely presented structures received a lot of attention in algorithmic model theory [2]. Among the most important such classes of structures is the class of *string-automatic structures* [17]. A (relational) structure is string-automatic if its universe is a regular set of words and all relations can be recognized by synchronous multi-tape automata. During the past 15 years a theory of string-automatic structures has emerged. This theory was developed along two interrelated branches. The first is a structural branch, which leads to (partial) characterizations of particular classes of string-automatic structures [9, 16, 18, 19, 22]. The second is an algorithmic branch, which leads to numerous decidability and undecidability, as well as complexity results for important algorithmic problems for string-automatic structures [4, 18, 21]. One of the most fundamental

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results for string-automatic structures states that their first-order theories are uniformly decidable [17].

By replacing strings and string automata by trees and tree automata, respectively, Blumensath [3] generalized string-automatic structures to *tree-automatic structures* and proved that their first-order theories are still uniformly decidable. However compared to string-automatic structures, the theory of tree-automatic structures is less developed. The only non-trivial characterization of a class of tree-automatic structures we are aware of concerns ordinals. In [9] Delhommé proved that an ordinal is tree-automatic if and only if it is strictly below ω^{ω^ω} . Some complexity results for first-order theories of tree-automatic structures are shown in [21]. Recently, Huschenbett proved that it is decidable whether a given tree-automatic scattered linear order is string-automatic [12].

In this paper, we study tree-automatic well-founded trees. We stress here that in this paper the term *tree* always refers to order trees $\mathfrak{T} = (T, \leq)$, i.e., partial orders that have the shape of a tree, as opposed to successor trees where the edge relation is not transitive (except for trivial cases). Using Delhommé's decomposition technique for tree-automatic structures [9], we show that the rank of a tree-automatic well-founded tree is strictly below ω^ω . As a generalization of trees, we introduce the class of upwards linear partial orders. These are orders where the elements above any fixed element form a linear suborder. If such an order is well-founded and tree-automatic, then its rank is strictly below ω^{ω^ω} . Unfortunately, we do not know if this bound holds for all tree-automatic well-founded partial orders. We will show in Example 3.27 that there is no hope to extend Delhommé's technique in order to obtain bounds on the ranks of all tree-automatic well-founded partial orders. Thus, characterizing the ranks of well-founded partial orders requires the development of completely new techniques.

We apply this result to the isomorphism problem for tree-automatic well-founded trees. In [20], it was shown that the isomorphism problem for string-automatic well-founded trees is complete for level Δ_ω^0 of the hyperarithmetical hierarchy. In other words, the isomorphism problem for string-automatic well-founded trees is recursively equivalent to true arithmetic. We show that the rank of well-founded computable trees determines the complexity of the isomorphism problem in the following sense: The isomorphism problem for well-founded computable trees of rank at most λ (where λ is a computable limit ordinal) belongs to level Δ_λ^0 of the hyperarithmetical hierarchy. Since we know that the rank of a tree-automatic well-founded tree is strictly below ω^ω , this implies that the isomorphism problem for tree-automatic well-founded trees belongs to level $\Delta_{\omega^\omega}^0$ of the hyperarithmetical hierarchy. We also provide a corresponding lower bound. We prove that the isomorphism problem for tree-automatic well-founded trees is $\Delta_{\omega^\omega}^0$ -complete under Turing-reductions.

Let us remark that for non-well-founded order trees, the isomorphism problem is complete for Σ_1^1 (the first existential level of the analytical hierarchy) already in the string-automatic case [20], and this complexity is in a certain sense maximal, since the isomorphism problem for the class of all computable structures is Σ_1^1 -complete as well [6, 10]. Let us also emphasize that our proof techniques do not work for successor trees.

2. PRELIMINARIES

We write $\mathbb{N}_{>0}$ for $\mathbb{N} \setminus \{0\}$. For M, N sets, 2^M denotes the *powerset* of M and N^M denotes the *set of functions* from M to N . We denote by $M \sqcup N$ (by $\bigsqcup_{i \in N} M_i$, respectively) the

disjoint union of the sets M and N ($(M_i)_{i \in N}$, respectively). For any (partial) function f we write $\text{dom}(f)$ for the domain of f .

A (finite and relational) signature $\tau = (\mathcal{R}, \text{ar})$ is a finite set \mathcal{R} of relation symbols together with a map $\text{ar} : \mathcal{R} \rightarrow \mathbb{N}_{>0}$ assigning to each $R \in \mathcal{R}$ its arity $\text{ar}(R)$. A τ -structure $\mathfrak{S} = (S, (R^{\mathfrak{S}})_{R \in \mathcal{R}})$ consists of a domain S and an $\text{ar}(R)$ -ary relation $R^{\mathfrak{S}}$ on S for each $R \in \mathcal{R}$. For $T \subseteq S$, we denote the substructure of \mathfrak{S} induced by the restriction of \mathfrak{S} to the set T as $\mathfrak{S}|_T$. In this paper we consider only structures with countable domains.

A partial order $\mathfrak{A} = (A, \leq)$ is regarded as a structure over a signature consisting of a single binary relation symbol \leq . Every substructure of \mathfrak{A} is again a partial order. We call a subset $B \subseteq A$ an *antichain* if for all distinct $a, b \in B$, neither $a \leq b$ nor $b \leq a$.

Let A be a (not necessarily finite) set. We use \preceq to denote the prefix order on finite words in A^* , i.e., for $u, v \in A^*$, $u \preceq v$ if $v = uw$ for some $w \in A^*$. For a language $L \subseteq A^*$, let $\text{pref}(L) = \{w \in A^* \mid \exists u \in L : w \preceq u\}$ be its *prefix-closure*.

2.1. Trees, forests and upwards linear partial orders. Let $\mathfrak{P} = (P, \leq)$ be a partial order. We say that \mathfrak{P} is *connected* if the undirected graph $(P, \{(a, b) \mid a \leq b \text{ or } b \leq a\})$ is connected. For an element $p \in P$ let $\uparrow p = \{a \in P \mid p < a\}$ be its (strict) *upwards closure* and $\downarrow p = \{a \in P \mid a < p\}$ be its (strict) *downwards closure*. Two elements $a, b \in P$ are *comparable* if $a \leq b$ or $b \leq a$. We say that \mathfrak{P} is *upwards linear* if for every $p \in P$, $\uparrow p$ induces a linear order $\mathfrak{P}|_{\uparrow p}$. A *forest* is an upwards linear partial order $\mathfrak{F} = (F, \leq)$ such that for every $p \in P$ the set $\uparrow p$ induces a finite linear order. A *tree* is a forest which has a greatest element, which is called the *root* of the tree. Note that a forest is a disjoint union of (an arbitrary number of) trees and all its substructures are also forests. For a given forest \mathfrak{F} , we denote as $\langle \mathfrak{F} \rangle$ the tree that results from adding a new root (i.e., a new greatest element) to \mathfrak{F} . If F is the domain of \mathfrak{F} we write $\langle F \rangle$ for the domain of $\langle \mathfrak{F} \rangle$. For a node u in \mathfrak{F} , $\mathfrak{F}(u)$ denotes the subtree of \mathfrak{F} at u , i.e., $\mathfrak{F}(u) = \mathfrak{F}|_{\{v \in F \mid v \leq u\}}$. We define the *successor relation* of \mathfrak{F} as

$$E_{\mathfrak{F}} = \{(x, y) \in F^2 \mid x > y, \neg \exists z : x > z > y\}.$$

For $x \in F$ the set of *children* of x in \mathfrak{F} is $E_{\mathfrak{F}}(x) = \{y \in F \mid (x, y) \in E_{\mathfrak{F}}\}$. The set of leaves of \mathfrak{F} is $\text{leaves}(\mathfrak{F}) = \{x \in F \mid E_{\mathfrak{F}}(x) = \emptyset\}$. A partial order $\mathfrak{P} = (P, \leq)$ is *well-founded*, if it does not contain an infinite descending chain $a_1 > a_2 > a_3 > \dots$. The *sum* of two partial orders $\mathfrak{P}_1 = (P_1, \leq_1)$ and $\mathfrak{P}_2 = (P_2, \leq_2)$ with $P_1 \cap P_2 = \emptyset$ is the partial order $\mathfrak{P}_1 + \mathfrak{P}_2 = (P_1 \cup P_2, \leq)$, where $a \leq b$ if either $(a, b \in P_1 \text{ and } a \leq_1 b)$, or $(a, b \in P_2 \text{ and } a \leq_2 b)$, or $(a \in P_1 \text{ and } b \in P_2)$.

We write *wulpo* as an abbreviation for well-founded upwards linear partial order. Note that the class of wulpo's contains all ordinals and all well-founded forests and is closed under taking substructures.

Let $\mathfrak{P} = (P, \leq)$ be some partial order. In the following, we set $p < \infty$ for all $p \in P$. In particular, ∞ is comparable to every $p \in P \cup \{\infty\}$ and $\max(p, \infty) = \infty$ for every $p \in P \cup \{\infty\}$. For $p_1 \in P$ and $p_2 \in P \cup \{\infty\}$ with $p_1 \leq p_2$, we denote by $[p_1, p_2]$ the *interval* $\{p \in P \mid p_1 \leq p < p_2\}$. We call this interval *branching free* if for all $p \in [p_1, p_2]$ we have

$$\downarrow p = \downarrow p_1 + [p_1, p).$$

In other words: $[p_1, p_2]$ is branching free, if for every $p \in [p_1, p_2]$ and every $x < p$, the points x and p_1 are comparable.

Lemma 2.1. *Let $p_1 \in P$ and $p_2 \in P \cup \{\infty\}$ with $p_1 \leq p_2$. If $[p_1, p_2)$ is branching free, then for every $p \in [p_1, p_2)$ also $[p_1, p)$ is branching free.*

Proof. Assume that $x < p' \in [p_1, p)$. We have to show that x and p_1 are comparable. But this is clear, since also $x < p' \in [p_1, p_2)$ and the latter interval is branching free. \square

Lemma 2.2. *Let $\mathfrak{P} = (P, \leq)$ be an upwards linear partial order and let $a, c \in P$ and $b, d \in P \cup \{\infty\}$ with $a \leq b$ and $c \leq d$. If $[a, b)$ and $[c, d)$ are branching free and not disjoint, then a and c are comparable, and b and d are comparable.*

Proof. Assume that $[a, b)$ and $[c, d)$ are not disjoint and let $p \in [a, b) \cap [c, d)$. Since \mathfrak{P} is upwards linear, the assumptions $p < b$ and $p < d$ imply that b and d are comparable. Moreover, since $c \leq p \in [a, b)$ and $[a, b)$ is branching free, the elements c and a are comparable too. \square

We call a branching free interval *maximal branching free* if it is maximal with respect to set inclusion. We use this interval notation mostly for upwards linear partial orders. In this case every interval induces a linear suborder.

Lemma 2.3. *Let $\mathfrak{P} = (P, \leq)$ be an upwards linear partial order and let $a, c \in P$ and $b, d \in P \cup \{\infty\}$ with $a \leq b$ and $c \leq d$. If $[a, b)$ and $[c, d)$ are maximal branching free, then they are either disjoint or equal.*

Proof. If $[a, b)$ and $[c, d)$ are not disjoint, Lemma 2.2 implies that $\max(b, d)$ and $\min(a, c)$ are defined. We claim that $[\min(a, c), \max(b, d))$ is also branching free. Since we assumed $[a, b)$ and $[c, d)$ to be maximal branching free, this implies $[a, b) = [\min(a, c), \max(b, d)) = [c, d)$.

Let us now prove the claim. If $a = c$ or $b = d$ or $(a < c, d < b)$ or $(c < a, b < d)$ then $[\min(a, c), \max(b, d))$ is $[a, b)$ or $[c, d)$ and hence branching free. The remaining two cases are symmetric. Hence we can assume that $a < c$ and $b < d$ and we have to show that $[a, d)$ is branching free. Since $a \leq b$ and $a \leq c$, we have either $b \leq c$ or $c < b$. If $b \leq c$, then the intervals $[a, b)$ and $[c, d)$ would be disjoint. Hence we have

$$a < c < b < d.$$

To show that $[a, d)$ is branching free, let $x \in [a, d)$ and $y < x$. We have to show that a and y are comparable.

If $x \in [a, b)$ then a and y are comparable since $[a, b)$ is branching free. So, assume that $x < b$ does not hold. Since \mathfrak{P} is upwards linear and $a \leq b$, $a \leq x$, it follows $b \leq x$. In particular $x \in [c, d)$. Since $[c, d)$ is branching free, c and y must be comparable. If $c \leq y$ then also $a \leq y$. Hence $y < c$. Since $c \in [a, b)$ and $[a, b)$ is branching free, y and a are comparable. \square

Lemma 2.4. *Let \mathfrak{P} be a wulpo. Every node of \mathfrak{P} is contained in a unique maximal branching free interval.*

Proof. Fix a node $p \in P$. Firstly, we define a node $p_+ > p$, secondly we define a node $p_- \leq p$, and finally, we prove that $[p_-, p_+)$ is maximal branching free. If for all $p' > p$ and all $p'' < p'$, p'' and p are comparable, set $p_+ := \infty$. Note that $[p, p_+)$ is branching free. Otherwise, by well-foundedness and upwards linearity, there is a minimal element $p_+ > p$ for which there exists $p' < p_+$ incomparable to p . Thus, $[p, p_+)$ is branching free.

Lemma 2.2 implies that, for all $p_1 \leq p, p_2 \leq p$, if $[p_1, p_+)$ and $[p_2, p_+)$ are branching free, then p_1 and p_2 are comparable. By well-foundedness, there is a minimal $p_- \leq p$ such that $I := [p_-, p_+)$ is branching free.

In order to prove maximality, assume that for some $p_1 \leq p < p_2$, $[p_1, p_2)$ is branching free. First we show that $p_+ \notin [p_1, p_2)$: If $p_+ \in [p_1, p_2)$ then, since $[p_1, p_2)$ is branching free, for every $p' < p_+$, p' and p_1 are comparable. By upwards linearity, if p' and p_1 are comparable then also p' and p are comparable. Hence, every $p' < p_+$ is comparable to p . But this contradicts the choice of p_+ . Hence, we have $p_+ \notin [p_1, p_2)$.

Since $p < p_2$ and $p < p_+$, p_2 and p_+ are comparable. We cannot have $p_+ < p_2$. Hence, $\max(p_2, p_+) = p_+$. By Lemma 2.3, the interval

$$[\min(p_1, p_-), p_+) = [\min(p_1, p_-), \max(p_2, p_+))$$

is branching free. But by definition of p_- , I is the minimal branching free interval of the form $[p, p_+)$. Thus, $p_- \leq p_1$ and $[p_1, p_2)$ is contained in I . This shows that I is maximal branching free. \square

Lemma 2.3 and 2.4 imply that every wulpo is partitioned into its maximal branching free intervals.

Lemma 2.5. *Let \mathfrak{P} be a wulpo and $I \neq J$ two different maximal branching free intervals of \mathfrak{P} . For $p_1 \in I$ and $p_2 \in J$, if $p_1 < p_2$ then there is a $p \in \mathfrak{P}$ such that $p < p_2$ but p and p_1 are incomparable, i.e., $\downarrow p_2 \neq \downarrow p_1 \cup [p_1, p_2)$.*

Proof. Heading for a contradiction, assume that

$$\downarrow p_2 = \downarrow p_1 \cup [p_1, p_2). \quad (2.1)$$

Note that there is a $p_m \in \mathfrak{P} \cup \{\infty\}$ such that $K := [p_1, p_m) = \{p \in \mathfrak{P} \mid p_1 \leq p \leq p_2\}$. Due to (2.1), K is branching free. Using Lemma 2.3 we conclude that $I \cup K \cup J$ is branching free contradicting maximality of I and J . \square

We finally discuss the replacement of maximal branching free intervals in wulpo's. Let $\mathfrak{P} = (P, \leq)$ be a wulpo and $I = [p_0, p_1)$ a maximal branching free interval. For each $i \in I$ and each $p \in P \setminus I$ we have

$$\begin{aligned} p \leq i & \text{ if and only if } p \leq i' \text{ for all } i' \in I \text{ and} \\ p \geq i & \text{ if and only if } p \geq i' \text{ for all } i' \in I \end{aligned}$$

Let J be some ordinal. Then we can replace I by J and obtain the wulpo $\mathfrak{P}' = ((P \setminus I) \sqcup J, \leq')$ where $p \leq' p'$ if

- (1) $p, p' \in P \setminus I$ and $p \leq p'$,
- (2) $p, p' \in J$ and $p \leq p'$ in J ,
- (3) $p \in P \setminus I, p' \in J$ and $p \leq p_0$, or
- (4) $p \in J, p' \in P \setminus I$ and $p_0 \leq p'$.

Note that the maximal branching free intervals of \mathfrak{P}' are J and the maximal branching free intervals of \mathfrak{P} except for I . Analogously, one can also replace simultaneously all maximal branching free intervals of a wulpo by new ordinals. This kind of replacement will be used in Lemma 3.15.

2.2. Ordinal arithmetic and the ordinal rank. We use standard terminology concerning ordinals; see e.g. [25]. We briefly recall the operations of natural sum and natural product on ordinals. The natural sum of ordinals α and β , denoted by $\alpha \oplus \beta$, is the largest (well-founded) linearization of $\alpha \sqcup \beta$ (cf. [7], where the following statements can be found). It is commutative, associative and strictly monotone, i.e., $\alpha \oplus \beta < \alpha \oplus \gamma$ if and only if $\beta < \gamma$. Furthermore,

$$\alpha \oplus \beta = \omega^\gamma \text{ implies } \alpha = \omega^\gamma \text{ or } \beta = \omega^\gamma. \quad (2.2)$$

The natural product (cf. [7]) of α and β , denoted as $\alpha \otimes \beta$, is the largest (well-founded) linearization of the direct (Cartesian) product of α and β with the componentwise order. It is commutative, associative, and strictly monotone, i.e., for $\alpha \neq 0$ we have $\alpha \otimes \beta < \alpha \otimes \gamma$ if and only if $\beta < \gamma$. Furthermore,

$$\alpha \otimes \beta = \omega^{\omega^\gamma} \text{ implies } \alpha = \omega^{\omega^\gamma} \text{ or } \beta = \omega^{\omega^\gamma}. \quad (2.3)$$

Lemma 2.6. *Let α, β, γ be ordinals. Then*

$$(\alpha \otimes \beta) + (\alpha \otimes \gamma) \leq \alpha \otimes (\beta + \gamma)$$

where $+$ is the usual sum of ordinals.

Proof. Observe that $(\alpha \otimes \beta) + (\alpha \otimes \gamma)$ is a well-order with domain $(\alpha \times \beta) \sqcup (\alpha \times \gamma) = \alpha \times (\beta \sqcup \gamma)$ such that for all $a < a' \in \alpha$, $b < b' \in \beta$ and $c < c' \in \gamma$ we have $(a, b) < (a', c)$, $(a, b) < (a', b')$, $(a, c) < (a', c)$, $(a, b) < (a, b')$, and $(a, c) < (a, c')$. Thus, we conclude that it is a linearization of the product of the orders α and $\beta + \gamma$. Since $\alpha \otimes (\beta + \gamma)$ is the largest linearization of these two orders, the claim follows immediately. \square

Let us recall the (ordinal) rank of a well-founded partial order. For a set of ordinals M , we denote by $\sup(M)$ the supremum of M , where $\sup(\emptyset) = 0$.

Definition 2.7. Let \mathfrak{P} be a well-founded partial order. The *rank* of an element $p \in \mathfrak{P}$ is inductively defined by

$$\text{rank}(p, \mathfrak{P}) = \sup\{\text{rank}(p', \mathfrak{P}) + 1 \mid p' < p \in P\}.$$

The rank of \mathfrak{P} (also called the *ordinal height* of \mathfrak{P}) is

$$\text{rank}(\mathfrak{P}) = \sup\{\text{rank}(p, \mathfrak{P}) + 1 \mid p \in P\}.$$

If the partial order \mathfrak{P} is clear from the context, we will write $\text{rank}(p)$ instead of $\text{rank}(p, \mathfrak{P})$. Note that $\text{rank}(p, \mathfrak{P}) = \text{rank}(\mathfrak{P}|_{\downarrow p})$. Also note that the rank of the ordinal α is exactly α . Next, we collect some useful facts about the rank.

Lemma 2.8. *Let $\mathfrak{P} = (P, \leq)$ be a well-founded partial order and $S \subseteq P$ a subset. The induced suborder $\mathfrak{P}|_S$ has rank at most $\text{rank}(\mathfrak{P})$.*

Proof. By induction on the rank of an element $s \in S$ (with respect to \mathfrak{P}) one easily shows that the rank of s in $\mathfrak{P}|_S$ is bounded by the rank of s in \mathfrak{P} . \square

We will also use a result proved by Khoussainov and Minnes on decomposition of partial orders into suborders.

Lemma 2.9 ([16], Lemma 3.3). *Let $\mathfrak{P} = (P, \leq)$ be a well-founded partial order, and let $P = P_1 \sqcup P_2$ be a partition of the domain of \mathfrak{P} . Setting $\mathfrak{P}_i = \mathfrak{P} \upharpoonright_{P_i}$ for $i \in \{1, 2\}$, we have $\text{rank}(\mathfrak{P}) \leq \text{rank}(\mathfrak{P}_1) \oplus \text{rank}(\mathfrak{P}_2)$.*

Lemma 2.10. *Let $\mathfrak{P} = (P, \leq)$ be a well-founded partial order and $p_1 \in P$, $p_2 \in P \cup \{\infty\}$ with $p_1 \leq p_2$. If $[p_1, p_2]$ is branching free, then $\text{rank}(p) = \text{rank}(p_1) + \text{rank}([p_1, p])$ for each $p \in [p_1, p_2]$. Moreover, if $p_2 \in P$, then $\text{rank}(p_2) \geq \text{rank}(p_1) + \text{rank}([p_1, p_2])$.*

Proof. Let $p \in [p_1, p_2]$. We prove the first statement of the lemma by induction on $\alpha := \text{rank}([p_1, p])$. We distinguish three cases:

Case 1. $\alpha = 0$. We must have $[p_1, p] = \emptyset$, i.e., $p_1 = p_2$. Hence $\text{rank}(p_2) = \text{rank}(p_1) + 0$.

Case 2. α is a successor ordinal, i.e., $\alpha = \beta + 1$. Then by definition of the rank there is an element $p' \in [p_1, p]$ such that

$$\text{rank}(p', [p_1, p]) = \text{rank}([p_1, p']) = \beta.$$

Moreover, $\text{rank}(p) \geq \text{rank}(p') + 1$ since $p' < p$. By induction hypothesis, we conclude that

$$\text{rank}(p) \geq \text{rank}(p') + 1 = \text{rank}(p_1) + \beta + 1 = \text{rank}(p_1) + \alpha.$$

Moreover, for every $p' \in [p_1, p]$ we must have $\text{rank}([p_1, p']) \leq \beta$. Since $[p_1, p_2]$ is branching free, we have $\downarrow p' = \downarrow p_1 + [p_1, p']$ for every $p' < p$. Therefore

$$\text{rank}(p) = \sup\{\text{rank}(p') + 1 \mid p' \in [p_1, p]\}.$$

By induction,

$$\text{rank}(p') = \text{rank}(p_1) + \text{rank}([p_1, p']) \leq \text{rank}(p_1) + \beta.$$

Hence, $\text{rank}(p) \leq \text{rank}(p_1) + \beta + 1 = \text{rank}(p_1) + \alpha$.

Case 3. α is a limit ordinal. Since $[p_1, p_2]$ is branching free, for each $p' < p$ one of the following holds.

- (1) $p' \in [p_1, p]$ and there is some $\beta < \alpha$ such that $\text{rank}([p_1, p']) = \beta$. By induction hypothesis, $\text{rank}(p') = \text{rank}(p_1) + \beta$.
- (2) $p' < p_1$ whence $\text{rank}(p') \leq \text{rank}(p_1)$.

Thus, we have

$$\begin{aligned} \text{rank}(p) &\leq \sup\{\text{rank}(p_1) + \beta + 1 \mid \beta < \alpha\} \\ &= \sup\{\text{rank}(p_1) + \beta \mid \beta < \alpha\} \\ &= \text{rank}(p_1) + \alpha. \end{aligned}$$

On the other hand, $\text{rank}([p_1, p]) = \alpha$ implies that for each $\beta < \alpha$ there is a node $p_\beta \in [p_1, p]$ such that $\beta \leq \text{rank}([p_1, p_\beta]) < \alpha$. By induction, we get $\text{rank}(p_\beta) \geq \text{rank}(p_1) + \beta$. Thus,

$$\begin{aligned} \text{rank}(p) &\geq \sup\{\text{rank}(p_\beta) + 1 \mid \beta < \alpha\} \\ &\geq \sup\{\text{rank}(p_1) + \beta + 1 \mid \beta < \alpha\} \\ &= \sup\{\text{rank}(p_1) + \beta \mid \beta < \alpha\} \\ &= \text{rank}(p_1) + \alpha. \end{aligned}$$

Finally, the second statement of the lemma follows easily:

$$\begin{aligned} \text{rank}(p_2) &\geq \sup\{\text{rank}(p) + 1 \mid p \in [p_1, p_2)\} \\ &= \sup\{\text{rank}(p_1) + \text{rank}([p_1, p)) + 1 \mid p \in [p_1, p_2)\} \\ &= \text{rank}(p_1) + \text{rank}([p_1, p_2)) \end{aligned}$$

This concludes the proof of the lemma. \square

Corollary 2.11. *Let \mathfrak{P} be a partial order and $p_1 \leq p_2 \in \mathfrak{P}$ such that $[p_1, p_2)$ is branching free. For every $p \in [p_1, p_2)$, $\text{rank}(p) < \text{rank}(p_1) + \text{rank}([p_1, p_2))$.*

2.3. Finitely labeled trees. We identify a non-empty, finite, prefix-closed subset $T \subseteq \{0, 1\}^*$ with the tree (T, \succeq) and call T a *finite binary tree*. Note that ε is the largest element of the inverse prefix relation \succeq and hence the root of the tree. We denote the set of all finite binary trees as $\mathcal{T}_2^{\text{fin}}$. For $T \in \mathcal{T}_2^{\text{fin}}$ let

$$\text{fr}(T) = T \setminus \{0, 1\} \quad \text{and} \quad \text{cl}(T) = T \cup \text{fr}(T)$$

be the *frontier* and *closure* of T . Notice that $\text{cl}(T) \in \mathcal{T}_2^{\text{fin}}$.

A finite Σ -labeled binary tree is a pair (T, λ) , where $T \in \mathcal{T}_2^{\text{fin}}$ and $\lambda : T \rightarrow \Sigma$ is a labeling function. We denote the set of all finite Σ -labeled binary trees by $\mathcal{T}_{2, \Sigma}^{\text{fin}}$; elements of $\mathcal{T}_{2, \Sigma}^{\text{fin}}$ will be denoted by lower case letters (s, t, \dots) . When Σ is the singleton set $\{\#\}$, we will simply consider a tree $t \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ as unlabeled, i.e., $t \in \mathcal{T}_2^{\text{fin}}$. The set of leaves of $t = (T, \lambda)$ is $\text{leaves}(t) = \text{leaves}(T, \succeq)$. We use the following operations on labeled trees $t = (T, \lambda) \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$. For $d \in T$, the *subtree rooted at d* is the tree $t(d) = (U, \mu) \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ with $U = \{u \in \{0, 1\}^* \mid du \in T\}$ and $\mu(u) = \lambda(du)$. For trees $t_1, \dots, t_n \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ and nodes $d_1, \dots, d_n \in \text{cl}(T)$ forming an antichain (i.e., d_i is not a prefix of d_j for $i \neq j$) we consider the tree $t[d_1/t_1, \dots, d_n/t_n] \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ which is obtained from t by simultaneously replacing for each i the subtree rooted at d_i by t_i . Formally, $t[d_1/t_1, \dots, d_n/t_n] = (V, \nu)$ is defined by

$$\begin{aligned} V &= (T \setminus (\{d_1, \dots, d_n\}\{0, 1\}^*)) \cup \bigcup_{1 \leq i \leq n} d_i T_i \text{ and} \\ \nu(v) &= \begin{cases} \lambda_i(u) & \text{if } v = d_i u \text{ for some (unique) } i \in \{1, \dots, n\}, \\ \lambda(v) & \text{otherwise,} \end{cases} \end{aligned}$$

where $t_i = (T_i, \lambda_i)$.

We fix a new symbol $\diamond \notin \Sigma$. Let Σ_\diamond denote the set $\Sigma \sqcup \{\diamond\}$ and Σ_\diamond^n the n -fold product $\Sigma_\diamond \times \dots \times \Sigma_\diamond$. Let $\bar{t} = (t_1, \dots, t_n) \in (\mathcal{T}_{2, \Sigma}^{\text{fin}})^n$ be a tuple of trees with $t_i = (T_i, \lambda_i)$. The *convolution* of \bar{t} is the Σ_\diamond^n -labeled binary tree $\otimes \bar{t} = (T, \lambda) \in \mathcal{T}_{2, \Sigma_\diamond^n}^{\text{fin}}$ defined by

$$\begin{aligned} T &= \bigcup_{1 \leq i \leq n} T_i \text{ and} \\ \lambda(u) &= (\lambda'_1(u), \dots, \lambda'_n(u)), \text{ where} \\ \lambda'_i(u) &= \begin{cases} \lambda_i(u) & \text{if } u \in T_i, \\ \diamond & \text{if } u \in T \setminus T_i. \end{cases} \end{aligned}$$

Instead of $\otimes(t_1, t_2, \dots, t_n)$ we sometimes also write $t_1 \otimes \dots \otimes t_n$. Finally, for any n -ary relation $R \subseteq (\mathcal{T}_{2,\Sigma}^{\text{fin}})^n$ we define

$$\otimes R = \{\otimes \bar{t} \mid \bar{t} \in R\} \subseteq \mathcal{T}_{2,\Sigma_\diamond^n}^{\text{fin}}.$$

2.4. Tree automata and tree-automatic structures. Let Σ be a finite alphabet. A (top-down) *tree automaton* over Σ is a tuple $\mathcal{A} = (Q, \Delta, I, F)$, where Q is the finite set of states, $I \subseteq Q$ is the set of initial states, $F \subseteq Q$ is the set of final states, and $\Delta \subseteq (Q \setminus F) \times \Sigma \times Q \times Q$ is the transition relation.¹ Given a finite Σ -labeled binary tree $t = (T, \lambda) \in \mathcal{T}_{2,\Sigma}^{\text{fin}}$, a *successful run* of \mathcal{A} on t is a mapping $\rho : \text{cl}(T) \rightarrow Q$ such that

- (1) $\rho(\varepsilon) \in I$,
- (2) $\rho(\text{fr}(T)) \subseteq F$, and
- (3) $(\rho(d), \lambda(d), \rho(d0), \rho(d1)) \in \Delta$ for every $d \in T$.

Let $L(\mathcal{A})$ be the set of all $t \in \mathcal{T}_{2,\Sigma}^{\text{fin}}$ which admit a successful run of \mathcal{A} . A set $L \subseteq \mathcal{T}_{2,\Sigma}^{\text{fin}}$ is called *regular* if there is a tree automaton \mathcal{A} over Σ with $L = L(\mathcal{A})$.

A successful run ρ of \mathcal{A} on $t = (T, \lambda) \in \mathcal{T}_{2,\Sigma}^{\text{fin}}$ naturally defines a $((\Sigma \times (Q \setminus F)) \cup F)$ -labeled binary tree $\text{tree}(\rho) = (\text{cl}(T), \mu)$ such that $\mu(d) = (\lambda(d), \rho(d))$ for every $d \in T$ and $\mu(d) = \rho(d)$ for every $d \in \text{fr}(T)$. $\text{Run}(\mathcal{A}, t)$ denotes the set of all trees $\text{tree}(\rho)$ where ρ is a successful run of \mathcal{A} on t . Let $\text{Run}(\mathcal{A}) = \bigcup_{t \in L(\mathcal{A})} \text{Run}(\mathcal{A}, t)$. This is also a regular set: a tree automaton for $\text{Run}(\mathcal{A})$ can be obtained by replacing every transition $(p, a, p_1, p_2) \in \Delta$ by $(p, (a, p), p_1, p_2)$ and adding transitions (f, f, \perp, \perp) for each final state $f \in F$ where \perp is a new state which is the only final state of the new automaton. For notational simplicity we refer to $\text{tree}(\rho)$ simply as ρ .

A tree automaton $\mathcal{A} = (Q, \Delta, I, F)$ is called *bottom-up deterministic* if

- (1) for all $p, q \in Q$ and $a \in \Sigma$ there is a unique $r \in Q$ with $(r, a, p, q) \in \Delta$ and
- (2) F is a singleton set.

In this situation, for $F = \{q_0\}$, a tree $t = (T, \lambda) \in \mathcal{T}_{2,\Sigma}^{\text{fin}}$, a node $u \in \text{cl}(T)$, a set $U \subseteq \text{fr}(T)$ and a map $\zeta : U \rightarrow Q$ we recursively define a state $\mathcal{A}(t, u, \zeta) \in Q$ by

$$\mathcal{A}(t, u, \zeta) = \begin{cases} r & \text{if } u \in T \text{ and } (r, \lambda(u), \mathcal{A}(t, u0, \zeta), \mathcal{A}(t, u1, \zeta)) \in \Delta, \\ \zeta(u) & \text{if } u \in U, \\ q_0 & \text{if } u \in \text{fr}(T) \setminus U. \end{cases}$$

We omit the second parameter if $u = \varepsilon$ and the third one if $U = \emptyset$. Notice that $\mathcal{A}(t, u) = \mathcal{A}(t, u)$ for all $u \in T$. The tree t admits a successful run ρ of \mathcal{A} if and only if $\mathcal{A}(t) \in I$. It is well known that for each tree automaton \mathcal{A} one can compute a bottom-up deterministic tree automaton \mathcal{A}' with $L(\mathcal{A}) = L(\mathcal{A}')$.

An n -ary relation $R \subseteq (\mathcal{T}_{2,\Sigma}^{\text{fin}})^n$ is called *tree-automatic* if there is a tree automaton \mathcal{A}_R over Σ_\diamond^n such that $L(\mathcal{A}_R) = \otimes R$. A relational structure \mathfrak{S} is called *tree-automatic* over Σ if its domain S is a regular subset of $\mathcal{T}_{2,\Sigma}^{\text{fin}}$ and each of its atomic relations $R^{\mathfrak{S}}$ is tree-automatic; any tuple \mathbb{P} of automata that accept the domain and the relations of \mathfrak{S} is called a *tree-automatic presentation* of \mathfrak{S} . In this case, we write $\mathfrak{S}(\mathbb{P})$ for \mathfrak{S} . If a tree-automatic structure \mathfrak{S} is isomorphic to a structure \mathfrak{S}' , then \mathfrak{S} is called a *tree-automatic copy* of \mathfrak{S}' and \mathfrak{S}' is

¹In contrast to the usual definition, we disallow transitions starting in final states. Obviously this restricted version of tree-automata is equivalent to the usual one. Our definition simplifies some constructions because runs on a tree t only assume final states on the frontier of t .

tree-automatically presentable. In this paper we sometimes abuse the terminology referring to \mathfrak{S}' as simply tree-automatic and calling a tree-automatic presentation of \mathfrak{S} also a tree-automatic presentations of \mathfrak{S}' . We also simplify our statements by saying “given/compute a tree-automatic structure \mathfrak{S} ” for “given/compute a tree-automatic presentation \mathbb{P} of a structure $\mathfrak{S}(\mathbb{P})$ ”. The structures $(\mathbb{N}, +)$ and (\mathbb{N}, \cdot) are examples of tree-automatic structures.

A tree-automatic structure over a singleton alphabet (i.e., the domain of the structure is a subset of $\mathcal{T}_2^{\text{fin}}$) is called *unary tree-automatic*. Moreover, let

$$\mathcal{T}_{\text{bin}} = \{t \in \mathcal{T}_2^{\text{fin}} \mid \forall u \in t : u0 \in t \Leftrightarrow u1 \in t\}$$

be the set of all finite (unlabeled) full binary trees. We will make use of the following simple lemma.

Lemma 2.12. *For every tree-automatic structure \mathfrak{S} there is an isomorphic unary tree-automatic structure \mathfrak{S}' whose domain is a subset of \mathcal{T}_{bin} . Moreover, there is a computable isomorphism from \mathfrak{S} to \mathfrak{S}' .*

Proof. Let Σ be some finite alphabet; w.l.o.g. assume that $\Sigma = \{1, 2, \dots, n\}$. For $1 \leq i \leq n$ let $a_i = \text{pref}(\{0, 10, 110, \dots, 1^{i-1}0, 1^i\}) \in \mathcal{T}_{\text{bin}}$. We inductively define an injective mapping $\text{unlabel} : \mathcal{T}_{2, \Sigma}^{\text{fin}} \rightarrow \mathcal{T}_{\text{bin}}$ as follows: Let $t = (T, \lambda) \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ and for $i \in \{0, 1\}$ let $t_i = t(i)$ be the subtree of t rooted at node i , where we set $t_i = \emptyset$ if $i \notin T$. Then

$$\text{unlabel}(t) = \{\varepsilon, 0\} \cup 00\text{unlabel}(t_0) \cup 01\text{unlabel}(t_1) \cup 1a_{\lambda(\varepsilon)},$$

where we set $\text{unlabel}(\emptyset) = \emptyset$. By induction over the size of t it easily follows that the mapping unlabel is indeed injective. We show that for every tree-automatic relation $R \subseteq (\mathcal{T}_{2, \Sigma}^{\text{fin}})^k$, the relation

$$\text{unlabel}(R) = \{(\text{unlabel}(t_1), \dots, \text{unlabel}(t_k)) \mid (t_1, \dots, t_k) \in R\}$$

is also tree-automatic. Suppose $\mathcal{A} = (Q, \Delta, I, F)$ is a tree automaton recognizing the relation R . We construct a (top-down) tree automaton \mathcal{A}' as follows: The state set of \mathcal{A}' contains the set

$$Q \cup (Q \times Q) \cup \{\diamond, 1, 2, \dots, n\}^n \cup \{\perp\}$$

(in addition \mathcal{A}' contains some auxiliary states that we do not specify), where \perp is final (states from F are no longer final). For a state $q \in Q$, \mathcal{A}' contains the following transitions (we omit here the unique node label, which formally should be the second component of every transition):

$$(q, (p, r), (x_1, \dots, x_k)) \text{ if } (q, (x_1, \dots, x_k), p, r) \in \Delta \quad \text{and } (q, \perp, \perp) \text{ for } q \in F.$$

For a state $(p, r) \in Q \times Q$, Δ' contains the following transitions:

$$((p, r), p, r).$$

Finally, the tree automaton \mathcal{A}' contains additional states and transitions such that from a state $(x_1, \dots, x_k) \in \{\diamond, 1, 2, \dots, n\}^n$ only the tree $\text{pref}(s_1 \otimes s_2 \otimes \dots \otimes s_n)$ with

$$s_i = \begin{cases} a_{x_i} & \text{if } x_i \in \{1, \dots, n\}, \\ \{\varepsilon\} & \text{if } x_i = \diamond. \end{cases}$$

is accepted. Now let $t_1, \dots, t_k \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ and $q \in Q$. A straightforward induction on the size of trees shows that \mathcal{A} accepts the convolution $t_1 \otimes \dots \otimes t_k$ via a run ρ with $\rho(\varepsilon) = q$ if and only if \mathcal{A}' accepts $\text{unlabel}(t_1) \otimes \dots \otimes \text{unlabel}(t_k)$ via a run ρ' with $\rho'(\varepsilon) = q$.

The above considerations shows that for every tree-automatic structure \mathfrak{S} there exists an isomorphic unary tree-automatic structure \mathfrak{S}' , whose domain is a subset of \mathcal{T}_{bin} . An isomorphism between \mathfrak{S} and \mathfrak{S}' is given by the computable mapping unlabel . This proves the lemma. \square

Consider $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$, the first-order logic over the signature τ extended by the quantifiers \exists^∞ (“there are infinitely many”) and the chain-quantifier \exists^{chain} (if $\varphi(x, y)$ is some formula, then $\exists^{\text{chain}}\varphi(x, y)$ asserts that φ is a partial order and there is an infinite increasing φ -chain). Let $\varphi(x_1, \dots, x_m, y_1, \dots, y_n)$ be a formula of this logic in the signature of some structure \mathfrak{S} . For all $\bar{b} \in S^n$ we define the m -ary relation

$$\varphi^{\mathfrak{S}}(\cdot, \bar{b}) = \{\bar{a} \in S^m \mid \mathfrak{S} \models \varphi(\bar{a}, \bar{b})\}$$

on \mathfrak{S} . In case of $n = 0$ we simply write $\varphi^{\mathfrak{S}}$ instead of $\varphi^{\mathfrak{S}}(\cdot)$. The following theorem from [3, 13, 26] lays out the main motivation for investigating tree-automatic structures.

Theorem 2.13. *From a tree-automatic presentation \mathbb{P} of a structure $\mathfrak{S}(\mathbb{P})$ over the signature τ and an $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$ -formula $\varphi(\bar{x})$ one can compute a tree automaton accepting $\otimes \varphi^{\mathfrak{S}(\mathbb{P})}$. In particular, the $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$ -theory of any tree-automatic structure \mathfrak{S} is (uniformly) decidable.*

Note that the property of being a tree is expressible in $\text{FO}[\leq] + \exists^\infty$. Moreover, the chain-quantifier allows to define well-foundedness of a tree. Hence, we obtain:

Corollary 2.14. *The class of tree-automatic well-founded trees is decidable.*

Let \mathcal{K} be a class of tree-automatic presentations. The *isomorphism problem* $\text{Iso}(\mathcal{K})$ is the set of pairs $(\mathbb{P}_1, \mathbb{P}_2) \in \mathcal{K} \times \mathcal{K}$ of tree-automatic presentations with $\mathfrak{S}(\mathbb{P}_1) \cong \mathfrak{S}(\mathbb{P}_2)$. If \mathcal{K} is the class of tree-automatic presentations for a class \mathcal{C} of relational structures (e.g. trees), then we briefly speak of the isomorphism problem for (tree-automatic members of) \mathcal{C} . The isomorphism problem for the class of all tree-automatic structures is complete for Σ_1^1 , the first level of the analytical hierarchy; this holds already for string-automatic trees [18, 20].

3. BOUNDING THE RANK OF TREE-AUTOMATIC WULPO’S AND WELL-FOUNDED TREES

In this section we show that the rank of a tree-automatic wulpo is strictly below ω^{ω^ω} . As a corollary, we obtain that the rank of every tree-automatic well-founded tree is strictly below ω^ω . The first part of our proof is a refinement of Delhommé’s decomposition theorem for tree-automatic structures [9]. This result shows that for a given tree-automatic structure \mathfrak{A} and a first-order formula $\varphi(x, \bar{y})$ all substructures induced by φ (via different tuples \bar{s} of parameters) are composed from a finite set $\mathcal{S}_\varphi^{\mathfrak{A}}$ of structures using the operations of box-augmentation and sum-augmentation. Roughly speaking, a structure \mathfrak{A} is a sum-augmentation of structures $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ if there is a finite partition of the domain of \mathfrak{A} which induces the substructures $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ (see Definition 3.1). The structure \mathfrak{A} is a box-augmentation of $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ if its domain is a finite product of copies of the domains of $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ such that fixing all but the i^{th} component of this product arbitrarily results in a structure isomorphic to \mathfrak{B}_i (see Definition 3.1).

Let ν be a function assigning ordinals to structures. An ordinal α is ν -indecomposable if for all structures $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ and all sum- or box-augmentations \mathfrak{A} of them, $\nu(\mathfrak{A}) = \alpha$ implies $\nu(\mathfrak{B}_i) = \alpha$ for some $i \in \{1, \dots, n\}$ (see Definition 3.7). Delhommé’s result implies the following. The substructures of a tree-automatic structure \mathfrak{A} (induced by a fixed first-order

formula φ) only take finitely many ν -indecomposable values. Unfortunately, Delhommé never published a proof of this result. Since our further arguments rely on this proof, we reprove his result in Sections 3.1 and 3.2. More precisely, we strengthen his result in the following way.

- All box-augmentations occurring in the decomposition of tree-automatic substructures provided by Delhommé's result are tamely colorable. Roughly speaking, \mathfrak{A} is a tamely colorable box-augmentation of $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ if there are finite colorings of the \mathfrak{B}_i and a simple rule how to completely reconstruct \mathfrak{A} from the \mathfrak{B}_i and these colorings.
- We introduce a notion of ν -indecomposability for intervals of ordinals: we say $[\alpha_1, \alpha_2]$ is ν -indecomposable if
 - (1) for all structures $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ and all sum-augmentations \mathfrak{A} of them, $\nu(\mathfrak{A}) = \alpha_2$ implies $\nu(\mathfrak{B}_i) = \alpha_2$ for some $i \in \{1, \dots, n\}$ and
 - (2) for all structures $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ and all tamely colorable box-augmentations \mathfrak{A} of them, $\nu(\mathfrak{A}) = \alpha_2$ implies $\nu(\mathfrak{B}_i) \in [\alpha_1, \alpha_2]$ for some $i \in \{1, \dots, n\}$.

It follows directly that for a sequence $\alpha_0 < \alpha_1 < \alpha_2 < \dots$ of ordinals such that each interval $[\alpha_i, \alpha_{i+1}]$ is ν -indecomposable, the substructures of a tree-automatic structure \mathfrak{A} (induced by a first-order formula φ) take only finitely many values among the α_i under ν (Proposition 3.9).

In Section 3.3, we prove that the **rank** of any tree-automatic wulpo is bounded by ω^{ω^ω} and in the case of a tree-automatic well-founded forest its **rank** is bounded by ω^ω . We prove this result using **rank**-indecomposability: we can prove that all intervals of the form $[\omega^{\omega^\alpha}, \omega^{\omega^{\alpha+1}}]$ are **rank**-indecomposable (if we restrict the domain of **rank** to the class of wulpo's). Thus, there are infinitely many **rank**-indecomposable intervals below ω^{ω^ω} . A simple transfinite induction shows that for every well-founded partial order \mathfrak{P} of rank α and for all $\beta < \alpha$, there is a node $p \in \mathfrak{P}$ such that $\mathfrak{P}|_{\{p' \mid p' < p\}}$ is of rank β . Thus, for a fixed tree-automatic well-founded partial order \mathfrak{P} , the automaton corresponding to the formula $x < y$ induces substructures of all ranks $\beta < \text{rank}(\mathfrak{P})$. For a tree-automatic wulpo \mathfrak{P} we have seen that these substructures may only take finitely many values of the form ω^{ω^i} . Thus, one concludes immediately that $\text{rank}(\mathfrak{P}) < \omega^{\omega^\omega}$. The bound on the ranks of tree-automatic forests follows from an analogous proof: writing $\text{rank}_{\mathcal{F}}$ for the restriction of **rank** to forests, we can prove that the interval $[\omega^i, \omega^{i+1}]$ is $\text{rank}_{\mathcal{F}}$ -indecomposable.

Note that Delhommé's original result is too weak for our proof. Using the definition of **rank**-indecomposable values with respect to all box-augmentations, the ordinals 0 and 1 are the only **rank**-indecomposable values. The problem is that any forest containing an infinite antichain is the box-augmentation of two infinite antichains. Note that an infinite antichain has **rank** 1. Hence, ordinals above 1 are not $\text{rank}_{\mathcal{F}}$ -indecomposable with respect to Delhommé's original definition. The crucial point is that such a box-augmentation of two infinite antichains in general is not tamely colorable.

3.1. Delhommé's decomposition theorem for tree-automatic structures. In this section we reprove Delhommé's decomposition theorem [9]. Beforehand, we give the precise definitions of sum- and box-augmentations. After providing a proof of Delhommé's original result, we introduce tamely colorable box-augmentations and improve this result. *For the rest of this section, we fix a relational signature $\tau = (\mathcal{R}, \text{ar})$.*

Definition 3.1. Let \mathfrak{A} and $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ be τ -structures.

(1) We say that \mathfrak{A} is a *sum-augmentation* of $(\mathfrak{B}_1, \dots, \mathfrak{B}_n)$ if there is a partition

$$A = \bigsqcup_{1 \leq i \leq n} A_i$$

of the domain of \mathfrak{A} such that $\mathfrak{A}|_{A_i} \cong \mathfrak{B}_i$ for all $1 \leq i \leq n$.

(2) We say that \mathfrak{A} is a *box-augmentation* of $(\mathfrak{B}_1, \dots, \mathfrak{B}_n)$ if there is a bijection

$$\eta : \prod_{1 \leq i \leq n} B_i \rightarrow A$$

such that for each $1 \leq k \leq n$ and all $\bar{d} = (d_1, \dots, d_{k-1}, d_{k+1}, \dots, d_n) \in \prod_{1 \leq i \leq n, i \neq k} B_i$

$$\eta_k^{\bar{d}} : B_k \rightarrow A$$

$$\eta_k^{\bar{d}}(e) = \eta(d_1, \dots, d_{k-1}, e, d_{k+1}, \dots, d_n)$$

is an (isomorphic) embedding of \mathfrak{B}_k into \mathfrak{A} .

Remark 3.2. Whenever \mathfrak{A} is a sum- or box-augmentation of $(\mathfrak{B}_1, \dots, \mathfrak{B}_n)$, the \mathfrak{B}_i are substructures of \mathfrak{A} . In particular, if \mathfrak{A} is a forest or a wulpo, the \mathfrak{B}_i are also forests or wulpo's, respectively.

To simplify notation, for a τ -structure \mathfrak{A} , an $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$ -formula $\varphi(x, y_1, \dots, y_n)$, and a tuple $\bar{s} \in A^n$ we use $\mathfrak{A}|_{\varphi, \bar{s}}$ as an abbreviation for $\mathfrak{A}|_{\varphi^{\mathfrak{A}}(\cdot, \bar{s})}$. Delhommé's decomposition theorem is the following.

Theorem 3.3. *From a given tree-automatic τ -structure \mathfrak{A} and an $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$ -formula $\varphi(x, y_1, \dots, y_n)$, one can compute a finite set $\mathcal{S}_\varphi^{\mathfrak{A}}$ of tree-automatic structures such that for all $\bar{s} \in A^n$ the substructure $\mathfrak{A}|_{\varphi, \bar{s}}$ is a sum-augmentation of box-augmentations of elements from $\mathcal{S}_\varphi^{\mathfrak{A}}$.*

Proof. Let \mathcal{A}_R , where $R \in \mathcal{R}$, (resp. \mathcal{A}_φ) be a bottom-up deterministic tree automaton accepting $\otimes R^{\mathfrak{A}}$ (resp. $\otimes \varphi^{\mathfrak{A}}$) and Q_R (resp. Q_φ) be its state set. For each $t = (T, \lambda) \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ and all $R \in \mathcal{R}$ of arity r we define $\otimes_R t = \otimes(t, \dots, t) \in \mathcal{T}_{2, \Sigma^r}^{\text{fin}}$ where the convolution is made up of r many copies of t . We further consider the tree $t \otimes_n \emptyset = (T, \lambda') \in \mathcal{T}_{2, \Sigma_0^{1+n}}^{\text{fin}}$ with $\lambda'(u) = (\lambda(u), \diamond, \dots, \diamond)$, where the number of \diamond -symbols is n .

As a first step towards proving the claim, we construct the set $\mathcal{S}_\varphi^{\mathfrak{A}}$. Before we construct the set, let us give some intuition on the structures we are interested in. $\mathcal{S}_\varphi^{\mathfrak{A}}$ consists of structures \mathfrak{B}_γ that are substructures of the structures $\mathfrak{A}|_{\varphi, \bar{s}}$ which are obtained as follows: we fix an element t of $\mathfrak{A}|_{\varphi, \bar{s}}$ and we fix a node d in t that is outside of the domain of \bar{s} . Now the domain of \mathfrak{B}_γ consists of all elements of $\mathfrak{A}|_{\varphi, \bar{s}}$ that are obtained from t by replacing the subtree rooted at d by some other tree. Thus, the structures \mathfrak{B}_γ are the substructures of all $\mathfrak{A}|_{\varphi, \bar{s}}$ whose domain is obtained by only local changes to a fixed tree.

Let $\Gamma = \prod_{R \in \mathcal{R}} Q_R \times Q_\varphi \times \prod_{R \in \mathcal{R}} 2^{Q_R}$. For each $\gamma = ((q_R)_{R \in \mathcal{R}}, q_\varphi, (P_R)_{R \in \mathcal{R}}) \in \Gamma$ we define a structure \mathfrak{B}_γ with domain

$$B_\gamma = \{t \in \mathcal{T}_{2, \Sigma}^{\text{fin}} \mid \mathcal{A}_R(\otimes t) = q_R \text{ for all } R \in \mathcal{R} \text{ and } \mathcal{A}_\varphi(t \otimes_n \emptyset) = q_\varphi\} \quad (3.1)$$

and relation $R \in \mathcal{R}$ (with arity $r = \text{ar}(R)$) interpreted by

$$R^{\mathfrak{B}_\gamma} = \{\bar{t} \in (B_\gamma)^r \mid \mathcal{A}_R(\otimes \bar{t}) \in P_R\}. \quad (3.2)$$

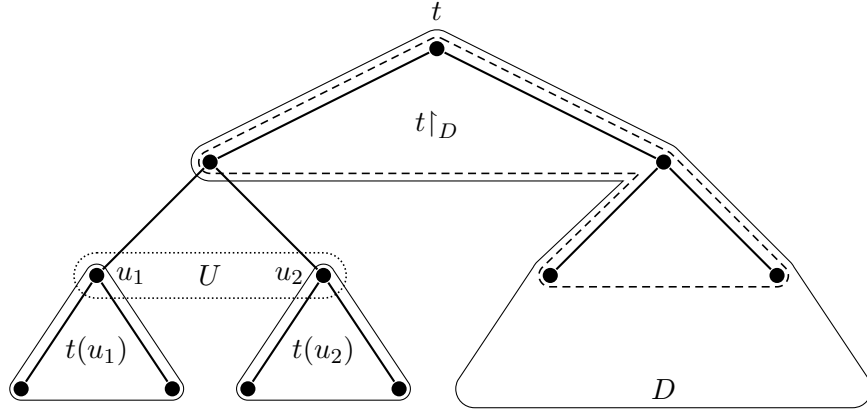


Figure 1: The \bar{s} -type of some example tree t . The dashed line surrounds $t|_D$ and U consists of all nodes within the dotted line.

Clearly, \mathfrak{B}_γ is tree-automatic. Finally, we let $\mathcal{S}_\varphi^\mathfrak{A}$ be the finite set

$$\mathcal{S}_\varphi^\mathfrak{A} = \{\mathfrak{B}_\gamma \mid \gamma \in \Gamma\}.$$

It remains to show that for each tuple $\bar{s} = (s_1, \dots, s_n) \in \mathfrak{A}^n$ the substructure $\mathfrak{A}|_{\varphi, \bar{s}}$ is a sum-augmentation of box-augmentations of structures from $\{\mathfrak{B}_\gamma \mid \gamma \in \Gamma\}$.

For this purpose, we fix such a tuple \bar{s} and consider the finite binary tree $D = \bigcup_{1 \leq i \leq n} S_i$, where $s_i = (S_i, \mu_i)$. The \bar{s} -type of a tree $t = (T, \lambda) \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ is the tuple

$$\text{tp}_{\bar{s}}(t) = (t|_D, U, (\zeta_R)_{R \in \mathcal{R}}, \zeta_\varphi),$$

where

- (1) $t|_D = (T \cap D, \lambda|_{(T \cap D)}) \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ is the restriction of t to the domain $T \cap D$,
- (2) $U = T \cap \text{fr}(D) \subseteq \text{fr}(T \cap D)$,
- (3) $\zeta_R : U \rightarrow Q_R$ is a map with $\zeta_R(u) = \mathcal{A}_R(\otimes_R t(u))$, and
- (4) $\zeta_\varphi : U \rightarrow Q_\varphi$ is a map with $\zeta_\varphi(u) = \mathcal{A}_\varphi(t(u) \otimes \emptyset)_n$.

Figure 1 illustrates the choice of $t|_D$ and U . Observe that for $t_u := t(u) \otimes \emptyset_n$ the equation

$$\otimes(t, \bar{s}) = \otimes(t|_D, \bar{s})[(u/t_u)_{u \in U}]$$

holds, whence

$$\mathcal{A}_\varphi(\otimes(t, \bar{s})) = \mathcal{A}_\varphi(\otimes(t|_D, \bar{s}), \zeta_\varphi). \quad (3.3)$$

Therefore, $\text{tp}_{\bar{s}}(t)$ determines whether $t \in \varphi^\mathfrak{A}(\cdot, \bar{s})$. Since D and \mathcal{R} are finite, there are only finitely many distinct \bar{s} -types. Consequently, the equivalence relation $\sim_{\bar{s}}$ on $\mathcal{T}_{2, \Sigma}^{\text{fin}}$ defined by $t \sim_{\bar{s}} t'$ if and only if $\text{tp}_{\bar{s}}(t) = \text{tp}_{\bar{s}}(t')$ has finite index. Due to (3.3), $\sim_{\bar{s}}$ saturates the set $\varphi^\mathfrak{A}(\cdot, \bar{s})$, i.e., $\varphi^\mathfrak{A}(\cdot, \bar{s})$ is a disjoint union of some $\sim_{\bar{s}}$ -classes. Assume that $Z_1, \dots, Z_\ell \subseteq \varphi^\mathfrak{A}(\cdot, \bar{s})$ are these $\sim_{\bar{s}}$ -classes. Then $\mathfrak{A}|_{\varphi, \bar{s}}$ is a sum-augmentation of $\mathfrak{A}|_{Z_1}, \dots, \mathfrak{A}|_{Z_\ell}$.

As the next step, we fix a single $\sim_{\bar{s}}$ -class $Z \subseteq \varphi^\mathfrak{A}(\cdot, \bar{s})$. We show that $\mathfrak{Z} = \mathfrak{A}|_Z$ is a box-augmentation of a tuple of structures from $\mathcal{S}_\varphi^\mathfrak{A}$. Let $(t, U, (\zeta_R)_{R \in \mathcal{R}}, \zeta_\varphi)$ be the \bar{s} -type corresponding to Z . Note that the domain of t is contained in the set D . For $u \in U$ we define

$$\gamma(u) = ((\zeta_R(u))_{R \in \mathcal{R}}, \zeta_\varphi(u), (P_R(u))_{R \in \mathcal{R}}) \in \Gamma,$$

where for each relation R

$$P_R(u) = \{q \in Q_R \mid \mathcal{A}_R(\otimes_R t, \zeta_R[u \mapsto q]) \in I_R\}. \quad (3.4)$$

Here $I_R \subseteq Q_R$ is the set of initial states of \mathcal{A}_R , and $\zeta_R[u \mapsto q]$ is the map ζ_R with the value at position u changed to q . Finally, we put $\mathfrak{X}_u = \mathfrak{B}_{\gamma(u)}$ and denote the domain of \mathfrak{X}_u by X_u . By definition of the domain $B_{\gamma(u)}$ of \mathfrak{X}_u (see (3.1)), $y \in X_u$ is equivalent to

$$\mathcal{A}_R(\otimes_R y) = \zeta_R(u) \text{ for all } R \in \mathcal{R} \quad (3.5)$$

$$\mathcal{A}_\varphi(y \otimes \emptyset) = \zeta_\varphi(u). \quad (3.6)$$

It remains to prove that \mathfrak{Z} is a box-augmentation of $(\mathfrak{X}_u)_{u \in U}$. First, observe that the map

$$\eta : \prod_{u \in U} X_u \rightarrow \mathcal{T}_{2, \Sigma}^{\text{fin}}$$

defined by

$$\eta((x_u)_{u \in U}) = t[(u/x_u)_{u \in U}]$$

is injective. Moreover, a tree $x = (X, \lambda) \in \mathcal{T}_{2, \Sigma}^{\text{fin}}$ is contained in the image of η if and only if

- (1) $x \upharpoonright_D = t$,
- (2) $X \cap \text{fr}(D) = U$, and
- (3) $x(u) \in X_u$ for each $u \in U$.

Note that by (3.5) and (3.6), condition (3) is equivalent to $\mathcal{A}_R(\otimes_R x(u)) = \zeta_R(u)$ for all $u \in U$ and $R \in \mathcal{R}$ and $\mathcal{A}_\varphi(x(u) \otimes \emptyset) = \zeta_\varphi(u)$ for all $u \in U$. Hence, the conjunction of (1), (2), and (3) is equivalent to $\text{tp}_{\bar{s}}(x) = (t, U, (\zeta_R)_{R \in \mathcal{R}}, \zeta_\varphi)$ which again is equivalent to $x \in Z$. Thus, η is a bijection onto its image Z .

Finally, we show that η is a componentwise embedding of the $(\mathfrak{X}_u)_{u \in U}$ into \mathfrak{Z} . For $u_0 \in U$ and $\bar{x} = (x_u)_{u \in U \setminus \{u_0\}} \in \prod_{u \in U \setminus \{u_0\}} X_u$ define the map $\eta_{u_0}^{\bar{x}} : X_{u_0} \rightarrow Z$ by

$$\eta_{u_0}^{\bar{x}}(x_{u_0}) = \eta((x_u)_{u \in U}).$$

For all $R \in \mathcal{R}$ of arity r and $(t_1, \dots, t_r) \in (X_{u_0})^r$ we have

$$\otimes_R(\eta_{u_0}^{\bar{x}}(t_1), \dots, \eta_{u_0}^{\bar{x}}(t_r)) = \left(\otimes_R t[(u/x_u)_{u \in U \setminus \{u_0\}}] \right) [u_0 / \otimes(t_1, \dots, t_r)].$$

Since $\mathcal{A}_R(\otimes_R x_u) = \zeta_R(u)$ for all $u \in U \setminus \{u_0\}$, this implies

$$\mathcal{A}_R(\otimes(\eta_{u_0}^{\bar{x}}(t_1), \dots, \eta_{u_0}^{\bar{x}}(t_r))) = \mathcal{A}_R(\otimes_R t, \zeta_R[u_0 \mapsto \mathcal{A}_R(\otimes(t_1, \dots, t_r))]). \quad (3.7)$$

We obtain the following chain of equivalences

$$\begin{aligned} (\eta_{u_0}^{\bar{x}}(t_1), \dots, \eta_{u_0}^{\bar{x}}(t_r)) \in R^{\mathfrak{Z}} &\iff (\eta_{u_0}^{\bar{x}}(t_1), \dots, \eta_{u_0}^{\bar{x}}(t_r)) \in R^{\mathfrak{A}} \\ &\iff \mathcal{A}_R(\otimes(\eta_{u_0}^{\bar{x}}(t_1), \dots, \eta_{u_0}^{\bar{x}}(t_r))) \in I_R \\ &\stackrel{(3.7)}{\iff} \mathcal{A}_R(\otimes_R t, \zeta_R[u_0 \mapsto \mathcal{A}_R(\otimes(t_1, \dots, t_r))]) \in I_R \\ &\stackrel{(3.4)}{\iff} \mathcal{A}_R(\otimes(t_1, \dots, t_r)) \in P_R(u_0) \\ &\stackrel{(3.2)}{\iff} (t_1, \dots, t_r) \in R^{\mathfrak{X}_{u_0}} \end{aligned}$$

showing that $\eta_{u_0}^{\bar{x}}$ is an embedding of \mathfrak{X}_{u_0} into \mathfrak{Z} . Consequently, \mathfrak{Z} is a box-augmentation of $(\mathfrak{X}_u)_{u \in U}$. \square

After all, let us take a closer look at the properties of the box-augmentation \mathfrak{Z} of $(\mathfrak{X}_u)_{u \in U}$ in the previous proof. Consider some $R \in \mathcal{R}$ and let $r = \text{ar}(R)$. For all $\bar{x}_1, \dots, \bar{x}_r \in \prod_{u \in U} X_u$ with $\bar{x}_i = (x_{i,u})_{u \in U}$ we have

$$\otimes(\eta(\bar{x}_1), \dots, \eta(\bar{x}_r)) = (\otimes_R t)[(u/ \otimes (x_{1,u}, \dots, x_{r,u}))_{u \in U}]$$

and hence

$$\mathcal{A}_R(\otimes(\eta(\bar{x}_1), \dots, \eta(\bar{x}_r))) = \mathcal{A}_R(\otimes_R t, \zeta), \quad (3.8)$$

where $\zeta : U \rightarrow Q_R$ is defined by $\zeta(u) = \mathcal{A}_R(\otimes(x_{1,u}, \dots, x_{r,u}))$. Thus, the tuple $(\zeta(u))_{u \in U}$ determines whether $(\eta(\bar{x}_1), \dots, \eta(\bar{x}_r)) \in R^{\mathfrak{A}}$.

This observation leads to an improved version of Theorem 3.3, namely Corollary 3.6 below. As an abstraction of our observation, we introduce the notion of tamely colorable box-augmentations in the following definition. Therein for $R \in \mathcal{R}$ a *finite R -coloring* of a τ -structure \mathfrak{A} is a map $\sigma : A^{\text{ar}(R)} \rightarrow C$ into a finite set C such that for all $\bar{x}, \bar{y} \in A^{\text{ar}(R)}$ with $\sigma(\bar{x}) = \sigma(\bar{y})$ we have $\bar{x} \in R^{\mathfrak{A}}$ if and only if $\bar{y} \in R^{\mathfrak{A}}$.

Definition 3.4. Let \mathcal{A} and $\mathfrak{B}_1, \dots, \mathfrak{B}_n$ be τ -structures. We say that \mathfrak{A} is a *tamely colorable box-augmentation* of $(\mathfrak{B}_1, \dots, \mathfrak{B}_n)$ if

- (1) there exists a bijection $\eta : \prod_{1 \leq i \leq n} B_i \rightarrow A$ witnessing that \mathfrak{A} is a box-augmentation of $(\mathfrak{B}_1, \dots, \mathfrak{B}_n)$ and
- (2) for every $R \in \mathcal{R}$ of arity $r = \text{ar}(R)$ there are finite R -colorings $\sigma_i : B_i^r \rightarrow C_i$ of each \mathfrak{B}_i such that the map $\sigma : A^r \rightarrow C_1 \times \dots \times C_n$ mapping $(a_1, \dots, a_r) \in A^r$ with $a_j = \eta(b_{j,1}, \dots, b_{j,n})$ for $1 \leq j \leq r$ to

$$\sigma(a_1, \dots, a_r) = (\sigma_1(b_{1,1}, \dots, b_{r,1}), \dots, \sigma_n(b_{1,n}, \dots, b_{r,n})),$$

is a finite R -coloring of \mathfrak{A} .

Roughly speaking, the last condition says that the colors assigned to the tuples of the “componentwise preimages” of the a_i under η already determine whether $(a_1, \dots, a_r) \in R^{\mathfrak{A}}$.

Remark 3.5. In the situation of Definition 3.4, assume that \mathfrak{A} and the \mathfrak{B}_i are directed graphs and all C_i are the same set, say $\{1, \dots, m\}$. For $1 \leq i \leq n$ the structure $\mathfrak{X}_i = (B_i; R_1^{\mathfrak{X}_i}, \dots, R_m^{\mathfrak{X}_i})$ with $R_c^{\mathfrak{X}_i} = \sigma_i^{-1}(c)$ can be regarded as a coloring of the edges of the complete directed graph on B_i with m colors. Since this coloring is compatible with the edge relation of \mathfrak{B}_i , the graph \mathfrak{A} is a generalized product—in the sense of Feferman and Vaught—of the structures $\mathfrak{X}_1, \dots, \mathfrak{X}_n$ using only atomic formulas.

We conclude by proving the “tamely-colorable” version of Theorem 3.3.

Corollary 3.6. *Given a tree-automatic τ -structure \mathfrak{A} and an $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$ -formula $\varphi(x, \bar{y})$, one can compute a finite set $\mathcal{S}_\varphi^{\mathfrak{A}}$ of tree-automatic τ -structures such that for all $\bar{s} \in A^n$ the substructure $\mathfrak{A}|_{\varphi, \bar{s}}$ is a sum-augmentation of tamely colorable box-augmentations of elements from $\mathcal{S}_\varphi^{\mathfrak{A}}$.*

Proof. We show that the box-augmentation \mathfrak{Z} of $(\mathfrak{X}_u)_{u \in U}$ constructed in the proof of Theorem 3.3 is tamely colorable. Therefore, consider the bijection $\eta : \prod_{u \in U} X_u \rightarrow Z$ witnessing this box-augmentation and fix some $R \in \mathcal{R}$ with arity $r = \text{ar}(R)$. Due to the definition of $R^{\mathfrak{X}_u}$ in (3.2), for each $u \in U$ the map $\sigma_u : X_u^r \rightarrow Q_R$ with $\sigma_u(\bar{x}) = \mathcal{A}_R(\otimes \bar{x})$ is a finite R -coloring of \mathfrak{X}_u . Finally, (3.8) shows that condition (2) from Definition 3.4 holds. \square

3.2. Sum- and box-indecomposability. Suppose that \mathcal{C} is a class of τ -structures and ν is a function assigning an ordinal $\nu(\mathfrak{A})$ to each $\mathfrak{A} \in \mathcal{C}$.² In this situation, we say that \mathcal{C} is *ranked* by ν .

Definition 3.7. Let \mathcal{C} be a class of τ -structures ranked by ν .

- (1) An ordinal α is called *ν -sum-indecomposable* if for every $\mathfrak{A} \in \mathcal{C}$ with $\nu(\mathfrak{A}) = \alpha$ and all τ -structures $\mathfrak{B}_1, \dots, \mathfrak{B}_m$ such that \mathfrak{A} is a sum-augmentation of $(\mathfrak{B}_1, \dots, \mathfrak{B}_m)$, there is $1 \leq i \leq m$ such that $\mathfrak{B}_i \in \mathcal{C}$ and $\nu(\mathfrak{B}_i) = \alpha$.
- (2) An interval $[\alpha_1, \alpha_2]$ of ordinals is called *ν -tamely-colorable-box-indecomposable* if for every $\mathfrak{A} \in \mathcal{C}$ with $\nu(\mathfrak{A}) = \alpha_2$ and all τ -structures $\mathfrak{B}_1, \dots, \mathfrak{B}_m$ such that \mathfrak{A} is a tamely colorable box-augmentation of $(\mathfrak{B}_1, \dots, \mathfrak{B}_m)$, there is $1 \leq i \leq m$ such that $\mathfrak{B}_i \in \mathcal{C}$ and $\nu(\mathfrak{B}_i) \in [\alpha_1, \alpha_2]$.

Remark 3.8. For classes \mathcal{C} which are closed under taking substructures, like the classes of forests and wulpo's, the requirement $\mathfrak{B}_i \in \mathcal{C}$ is always satisfied. Hence in this case explicitly requiring $\mathfrak{B}_i \in \mathcal{C}$ is not necessary.

The decomposition results from the previous section imply that ν may only take finitely many ν -sum-indecomposable and ν -tamely-colorable-box-indecomposable values among the substructures of the form $\mathfrak{A}|_{\varphi, \bar{s}}$ (defined just before Theorem 3.3) for a fixed $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$ -formula $\varphi(x, \bar{y})$.

Proposition 3.9. Let \mathcal{C} be a class of τ -structures ranked by ν and $\alpha_0 < \alpha_1 < \alpha_2 < \dots$ an infinite sequence of ν -sum-indecomposable ordinals such that $[\alpha_i, \alpha_{i+1}]$ is ν -tamely-colorable-box-indecomposable for all $i \in \mathbb{N}$. Moreover, let \mathfrak{A} be a tree-automatic τ -structure and $\varphi(x, y_1, \dots, y_r)$ an $\text{FO}[\tau] + \exists^\infty + \exists^{\text{chain}}$ -formula. Then there are only finitely many $i \in \mathbb{N}$ which admit a tuple $\bar{s} \in A^r$ with $\mathfrak{A}|_{\varphi, \bar{s}} \in \mathcal{C}$ and $\nu(\mathfrak{A}|_{\varphi, \bar{s}}) = \alpha_i$.

Proof. Let $\mathcal{S}_\varphi^\mathfrak{A}$ be the finite set of structures which exists by Corollary 3.6. Consider an $i \in \mathbb{N}_{>0}$ satisfying the condition above, witnessed by $\bar{s} \in A^r$, i.e., $\mathfrak{A}|_{\varphi, \bar{s}} \in \mathcal{C}$ and $\nu(\mathfrak{A}|_{\varphi, \bar{s}}) = \alpha_i$. There are structures $\mathfrak{X}_1, \dots, \mathfrak{X}_k$ such that each of them is a tamely colorable box-augmentation of elements from $\mathcal{S}_\varphi^\mathfrak{A}$ and $\mathfrak{A}|_{\varphi, \bar{s}}$ is a sum-augmentation of $(\mathfrak{X}_1, \dots, \mathfrak{X}_k)$. Due to the definition of ν -sum-indecomposability, there is a $1 \leq j \leq k$ such that $\mathfrak{X}_j \in \mathcal{C}$ and $\nu(\mathfrak{X}_j) = \alpha_i$. There are structures $\mathfrak{Y}_1, \dots, \mathfrak{Y}_\ell \in \mathcal{S}_\varphi^\mathfrak{A}$ such that \mathfrak{X}_j is a tamely colorable box-augmentation of $(\mathfrak{Y}_1, \dots, \mathfrak{Y}_\ell)$. By the definition of ν -tamely-colorable-box-indecomposability, there is an $1 \leq h \leq \ell$ such that $\mathfrak{Y}_h \in \mathcal{C}$ and $\nu(\mathfrak{Y}_h) \in [\alpha_{i-1}, \alpha_i]$. Thus, i belongs to the set

$$\{i \in \mathbb{N}_{>0} \mid \exists \mathfrak{B} \in \mathcal{S}_\varphi^\mathfrak{A} \cap \mathcal{C} : \nu(\mathfrak{B}) \in [\alpha_{i-1}, \alpha_i]\},$$

which is finite since each $\mathfrak{B} \in \mathcal{S}_\varphi^\mathfrak{A}$ satisfies the condition above for at most two distinct i 's. \square

²All our applications use isomorphism invariant functions, which means that $\nu(\mathfrak{A}) = \nu(\mathfrak{B})$ if $\mathfrak{A} \cong \mathfrak{B}$. In fact, the results of this section apply to arbitrary class functions but are only useful if the function is isomorphism invariant for the important part of the class \mathcal{C} because for other functions there are not enough indecomposable values. Thus, upon first reading of the following part, we recommend the reader to think of ν as an isomorphism invariant function.

3.3. Rank-tamely-colorable-box- and rank-sum-indecomposability. Let us briefly prove that the ordinals ω^α are rank-sum-indecomposable for the class of all well-founded partial orders. Afterwards, we characterize rank-tamely-colorable-box-indecomposable intervals for the classes of wulpo's and well-founded forests. For this purpose let rank_{ul} be the function obtained from rank by restriction of the domain to the class of wulpo's and analogously let rank_F be rank restricted to the class of well-founded forests.

Proposition 3.10. *The ordinals of the form ω^α are rank-sum-indecomposable for the class of well-founded partial orders.*

Proof. Let $\mathfrak{P} = (P, \leq)$ be a well-founded partial order and assume that \mathfrak{P} is a sum-augmentation of $(\mathfrak{P}_1, \dots, \mathfrak{P}_n)$. If $\text{rank}(\mathfrak{P} \upharpoonright_{P \setminus P_1}) < \omega^\alpha$ and $\text{rank}(\mathfrak{P}_1) < \omega^\alpha$, then

$$\text{rank}(\mathfrak{P}) \stackrel{\text{Lemma 2.9}}{\leq} \text{rank}(\mathfrak{P}_1) \oplus \text{rank}(\mathfrak{P} \upharpoonright_{P \setminus P_1}) < \omega^\alpha,$$

where the last inequality follows from Property (2.2) of \oplus (see page 6). Thus, $\text{rank}(\mathfrak{P}) = \omega^\alpha$ implies $\text{rank}(\mathfrak{P}_1) = \omega^\alpha$ or $\text{rank}(\mathfrak{P} \upharpoonright_{P \setminus P_1}) = \omega^\alpha$. The claim follows by induction on n . \square

It follows trivially that the ordinals of the form ω^α are rank_{ul} -sum-indecomposable and rank_F -sum-indecomposable.

We want to show that $[\omega^{\omega^\alpha}, \omega^{\omega^{\alpha+1}}]$ is a rank_{ul} -tamely-colorable-box-indecomposable interval for each ordinal α and $[\omega^\alpha, \omega^{\alpha+1}]$ is a rank_F -tamely-colorable-box-indecomposable interval for each ordinal α . We start with the observation that every box-decomposition of a wulpo only contains at most one proper wulpo in the sense that if a wulpo is a box-augmentation of $(\mathfrak{P}_1, \mathfrak{P}_2, \dots, \mathfrak{P}_n)$ then all but one of the \mathfrak{P}_i are disjoint unions of ordinals. In order to prove this fact, we introduce the following notation. Let $\mathfrak{P} = (P, \leq)$ be some partial order. We call $a \in P$ a *branching node* (of \mathfrak{P}), if there are $b, c \in P$ such that $b < a$, $c < a$ and neither $b \leq c$ nor $c \leq b$ (i.e., b and c are incomparable).

Lemma 3.11. *Let \mathfrak{P} , \mathfrak{P}_1 , and \mathfrak{P}_2 be wulpo's. If \mathfrak{P} is a box-augmentation of $(\mathfrak{P}_1, \mathfrak{P}_2)$ then \mathfrak{P}_1 or \mathfrak{P}_2 does not contain a branching node.*

Proof. Let $\mathfrak{P} = (P, \leq)$ and $\mathfrak{P}_i = (P_i, \leq_i)$. Heading for a contradiction assume that $a_i, b_i, c_i \in P_i$ for $i \in \{1, 2\}$ are nodes such that $b_i <_i a_i, c_i <_i a_i$ and neither $b_i \leq_i c_i$ nor $c_i \leq_i b_i$.

Let $\eta : P_1 \times P_2 \rightarrow P$ be the bijection that witnesses that \mathfrak{P} is a box-augmentation of $(\mathfrak{P}_1, \mathfrak{P}_2)$. Then \mathfrak{P} contains the chains

$$\begin{aligned} \eta(b_1, b_2) &< \eta(b_1, a_2) \text{ and} \\ \eta(b_1, b_2) &< \eta(a_1, b_2). \end{aligned}$$

Since \mathfrak{P} is a wulpo, the elements above $\eta(b_1, b_2)$ are linearly ordered and we may assume that $\eta(b_1, a_2) < \eta(a_1, b_2)$ without loss of generality. Thus, we obtain

$$\eta(b_1, c_2) < \eta(b_1, a_2) < \eta(a_1, b_2).$$

Furthermore, we have

$$\eta(b_1, c_2) < \eta(a_1, c_2).$$

Again, the elements above $\eta(b_1, c_2)$ are linearly ordered and we obtain that $\eta(a_1, c_2)$ and $\eta(a_1, b_2)$ are comparable in \mathfrak{P} . By definition of a box-augmentation, we obtain that c_2 and b_2 are comparable in \mathfrak{P}_2 as well, which contradicts our assumptions. Thus, \mathfrak{P}_1 or \mathfrak{P}_2 does not contain a branching node. \square

Corollary 3.12. *Let \mathfrak{P} and $\mathfrak{P}_1, \dots, \mathfrak{P}_n$ be wulpo's such that \mathfrak{P} is a box-augmentation of $(\mathfrak{P}_1, \dots, \mathfrak{P}_n)$. There is at most one $i \in \{1, \dots, n\}$ such that \mathfrak{P}_i contains a branching node, i.e., there is an $i \in \{1, \dots, n\}$ such that \mathfrak{P}_k is a disjoint union of well-orders for all $k \neq i$.*

Proof. Let $\eta : \prod_{i=1}^n \mathfrak{P}_i \rightarrow \mathfrak{P}$ be the bijection of the box-augmentation. Choose numbers $1 \leq j < k \leq n$ and a tuple $\bar{b} = (b_1, \dots, b_{j-1}, b_{j+1}, \dots, b_{k-1}, b_{k+1}, \dots, b_n) \in \prod_{i \in \{1, \dots, n\} \setminus \{j, k\}} P_i$ arbitrarily but fixed. Then $\eta_{j,k}^{\bar{b}} : P_j \times P_k \rightarrow P$ with $\eta_{j,k}^{\bar{b}}(b_j, b_k) = \eta(b_1, \dots, b_n)$ induces a box-augmentation of some subwulpo $\mathfrak{P}' \leq \mathfrak{P}$. Application of Lemma 3.11 yields the claim. \square

Remark 3.13. In the following, our proofs for the case of wulpo's and the case of well-founded forests proceed completely analogous. The difference in the results stems from an observation concerning Corollary 3.12: if a well-founded forest \mathfrak{F} is a box-augmentation of partial orders $\mathfrak{P}_1, \dots, \mathfrak{P}_n$ then the \mathfrak{P}_i occur as substructures of \mathfrak{F} . Hence, these are also well-founded forests. But if a well-founded forest is a disjoint union of ordinals, all these ordinals must be finite. Thus, each connected component of such a disjoint union has finite rank and the whole structure has rank at most ω . In contrast, if a disjoint union of ordinals is a tree-automatic wulpo each connected component is a tree-automatic ordinal whence its rank is bounded by ω^{ω^ω} and all smaller ordinals can be reached. This difference causes the different results with respect to box-indecomposability.

Lemma 3.14. *Let \mathfrak{P} and $\mathfrak{P}_1, \dots, \mathfrak{P}_m$ be wulpo's such that \mathfrak{P} is a box-augmentation of $(\mathfrak{P}_1, \dots, \mathfrak{P}_m)$ via $\eta : \prod_{i=1}^m \mathfrak{P}_i \rightarrow \mathfrak{P}$. Let $I_i \subseteq P_i$ be a well-order. The substructure $\mathfrak{P}' \subseteq \mathfrak{P}$ induced by $\eta(\prod_{i=1}^m I_i)$ is a well-order.*

Proof. Note that \mathfrak{P}' is a wulpo because it is a substructure of \mathfrak{P} . Thus, it suffices to show that \mathfrak{P}' is linear. Let $a_i^j \in I_i$ for each $1 \leq i \leq m$ and $j \in \{1, 2\}$. Set $m_i = \min(a_i^1, a_i^2)$. For $j \in \{1, 2\}$,

$$\eta(m_1, m_2, \dots, m_n) \leq \eta(a_1^j, m_2, \dots, m_n) \leq \dots \leq \eta(a_1^j, a_2^j, \dots, a_n^j).$$

Since the elements above $\eta(m_1, \dots, m_n)$ are pairwise comparable, the elements $\eta(a_1^1, \dots, a_n^1)$ and $\eta(a_1^2, \dots, a_n^2)$ are comparable. Since the a_i^j have been chosen arbitrarily any two elements of \mathfrak{P}' are comparable, i.e., \mathfrak{P}' is linear. \square

In the following lemma, the term “replacement” refers to the replacement operation introduced at the end of Section 2.1.

Lemma 3.15. *If a wulpo \mathfrak{P} is the box-augmentation of ordinals $\mathfrak{C}_1, \dots, \mathfrak{C}_n$ and a wulpo \mathfrak{Q} via the bijection η , then \mathfrak{P} results from \mathfrak{Q} by replacing every maximal branching free interval I of \mathfrak{Q} by $\eta((\prod_{i=1}^n \mathfrak{C}_i) \times I)$ (which is a well-order by Lemma 3.14).*

Proof. We first show the following two claims.

- (1) For all $q_0, q_1 \in \mathfrak{Q}$ and for all $b_i, c_i \in \mathfrak{C}_i$ the nodes $\eta(b_1, \dots, b_n, q_0)$ and $\eta(c_1, \dots, c_n, q_1)$ are incomparable if and only if q_0 and q_1 are incomparable.
- (2) For all $b_i, c_i \in \mathfrak{C}_i$ and $q_0, q_1, q \in \mathfrak{Q}$ such that $q_0 < q$ and $q_1 < q$ but q_0 and q_1 are incomparable, $\eta(b_1, \dots, b_n, q_1) < \eta(c_1, \dots, c_n, q)$.

The first statement is easy to show: If $\eta(b_1, \dots, b_n, q_0) \leq \eta(c_1, \dots, c_n, q_1)$ then the elements $\eta(\max(b_1, c_1), \dots, \max(b_n, c_n), q_0)$ and $\eta(\max(b_1, c_1), \dots, \max(b_n, c_n), q_1)$ are both above the

element $\eta(b_1, \dots, b_n, q_0)$ and therefore comparable. By the definition of a box-augmentation, also q_0 and q_1 are comparable. For the other direction, if $q_0 \leq q_1$, then

$$\eta(\min(b_1, c_1), \dots, \min(b_n, c_n), q_0) \leq \eta(\min(b_1, c_1), \dots, \min(b_n, c_n), q_1).$$

Thus, $\eta(b_1, \dots, b_n, q_0)$ and $\eta(c_1, \dots, c_n, q_1)$ are both above $\eta(\min(b_1, c_1), \dots, \min(b_n, c_n), q_0)$ and therefore comparable.

Let us now prove the second claim. Due to the first claim, $q_1 < q$ implies that $\eta(b_1, \dots, b_n, q_1)$ and $\eta(c_1, \dots, c_n, q)$ are comparable. Moreover,

$$\text{for all } a_i \in \mathfrak{C}_i, \eta(a_1, \dots, a_n, q_0) \text{ and } \eta(a_1, \dots, a_n, q_1) \text{ are incomparable,} \quad (3.9)$$

because q_0 and q_1 are incomparable and \mathfrak{P} is a box-augmentation. Similarly, it is clear that

$$\eta(\min(b_1, c_1), \dots, \min(b_n, c_n), q_i) \leq \eta(\min(b_1, c_1), \dots, \min(b_n, c_n), q) \leq \eta(c_1, \dots, c_n, q) \text{ and} \\ \eta(\min(b_1, c_1), \dots, \min(b_n, c_n), q_i) \leq \eta(b_1, \dots, b_n, q_i)$$

for $i \in \{0, 1\}$. Since \mathfrak{P} is a wulpo, the nodes above $\eta(\min(b_1, c_1), \dots, \min(b_n, c_n), q_i)$ are linearly ordered. Hence $\eta(c_1, \dots, c_n, q)$ and $\eta(b_1, \dots, b_n, q_i)$ are comparable for $i \in \{0, 1\}$. Note that

$$\eta(b_1, \dots, b_n, q_i) \leq \eta(c_1, \dots, c_n, q) \leq \eta(b_1, \dots, b_n, q_{1-i})$$

for $i \in \{0, 1\}$ would contradict (3.9). Analogously, if

$$\eta(c_1, \dots, c_n, q) \leq \eta(b_1, \dots, b_n, q_1) \text{ and} \\ \eta(c_1, \dots, c_n, q) \leq \eta(b_1, \dots, b_n, q_0),$$

then upwards linearity would lead to a contradiction with (3.9). Thus, we conclude that

$$\eta(b_1, \dots, b_n, q_1) \leq \eta(c_1, \dots, c_n, q),$$

which shows the second claim.

Using these two claims, we now show that for every pair I_1, I_2 of distinct maximal branching free intervals of \mathfrak{Q} and for all $q_1 \in I_1, q_2 \in I_2$, and for all $b_i, c_i \in \mathfrak{C}_i$ we have

$$q_1 < q_2 \iff \eta(b_1, \dots, b_n, q_1) < \eta(c_1, \dots, c_n, q_2). \quad (3.10)$$

First assume that $q_1 < q_2$. Since q_1 and q_2 come from different maximal branching free intervals, Lemma 2.5 implies that there is a $q_3 < q_2$ such that q_1 and q_3 are incomparable. Due to the second claim, this immediately shows that $\eta(b_1, \dots, b_n, q_1) < \eta(c_1, \dots, c_n, q_2)$. By symmetry, we conclude that $q_2 < q_1$ implies $\eta(b_1, \dots, b_n, q_2) < \eta(c_1, \dots, c_n, q_1)$. Moreover, because of the first claim, if q_1 and q_2 are incomparable then $\eta(b_1, \dots, b_n, q_1)$ and $\eta(c_1, \dots, c_n, q_2)$ are incomparable as well.

By (3.10), \mathfrak{P} is obtained from \mathfrak{Q} by replacing every maximal branching free interval I of \mathfrak{Q} by a box-augmentation of $(\mathfrak{C}_1, \dots, \mathfrak{C}_n, I)$, which is a well order by Lemma 3.14. This concludes the proof of the lemma. \square

Lemma 3.16. *Let \mathfrak{P} be a wulpo that is a box-augmentation of ordinals $\mathfrak{C}_1, \dots, \mathfrak{C}_n$ and a wulpo \mathfrak{Q} via the bijection η . Let $q \in \mathfrak{Q}$ and $q' \in \mathfrak{Q} \cup \{\infty\}$ such that $q < q'$ and $[q, q']$ is a maximal branching free interval in \mathfrak{Q} . Let $\alpha := \text{rank}([q, q']) > 0$ and 0_i the minimal element of \mathfrak{C}_i for each $1 \leq i \leq n$. Then*

$$\text{rank}(\eta(0_1, 0_2, \dots, 0_n, q), \mathfrak{P}) \leq \text{rank}(\mathfrak{C}_1) \otimes \text{rank}(\mathfrak{C}_2) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes \text{rank}(q, \mathfrak{Q}) \quad (3.11)$$

and for each $\hat{q} \in [q, q']$ and $c_i \in \mathfrak{C}_i$ we have

$$\text{rank}(\eta(c_1, c_2, \dots, c_n, \hat{q}), \mathfrak{P}) < \text{rank}(\mathfrak{C}_1) \otimes \text{rank}(\mathfrak{C}_2) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes (\text{rank}(q, \mathfrak{Q}) + \alpha). \quad (3.12)$$

Proof. We prove both claims simultaneously by induction. If $\text{rank}(q, \mathfrak{Q}) = 0$, then q is minimal in \mathfrak{Q} . By Lemma 3.15, $\eta(0_1, 0_2, \dots, 0_n, q)$ is minimal in \mathfrak{P} . Hence, its rank is also 0 as desired.

Now assume that (3.11) is true for some $q \in \mathfrak{Q}$ inducing a maximal branching free interval $[q, q'] \subseteq \mathfrak{Q}$ and let $c_i \in \mathfrak{C}_i$ and $\hat{q} \in [q, q']$. Due to Lemma 3.15, $I := \eta(\prod_{i=1}^n \mathfrak{C}_i \times [q, q'])$ is a maximal branching free interval of \mathfrak{P} with minimal element $\eta(0_1, 0_2, \dots, 0_n, q)$. The rank of I can be at most $\text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes \alpha$. The interval

$$J := [\eta(0_1, \dots, 0_n, q), \eta(c_1, \dots, c_n, \hat{q})) \subseteq \mathfrak{P}$$

is strictly contained in I whence

$$\text{rank}(J) < \text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes \alpha. \quad (3.13)$$

Thus, we get

$$\begin{aligned} & \text{rank}(\eta(c_1, \dots, c_n, \hat{q}), \mathfrak{P}) \\ & \stackrel{\text{Lem. 2.10}}{=} \text{rank}(\eta(0_1, \dots, 0_n, q), \mathfrak{P}) + \text{rank}(J) \\ & \stackrel{(3.13)}{<} \text{rank}(\eta(0_1, \dots, 0_n, q), \mathfrak{P}) + \text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes \alpha \\ & \stackrel{(3.11)}{\leq} (\text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes \text{rank}(q, \mathfrak{Q})) + (\text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes \alpha) \\ & \stackrel{\text{Lem. 2.6}}{\leq} \text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes (\text{rank}(q, \mathfrak{Q}) + \alpha). \end{aligned}$$

Finally, let $q, q' \in \mathfrak{Q}$ be nodes such that $[q, q']$ is maximal branching free and all maximal branching free intervals below q satisfy the claims. Lemma 2.4 shows that for each $q_1 < q$ there is a maximal branching free interval $I := [q_0, q_2)$ with $q_2 \leq q$ such that $q_1 \in I$. Then for all $c_i \in \mathfrak{C}_i$

$$\begin{aligned} \text{rank}(\eta(c_1, c_2, \dots, c_n, q_1), \mathfrak{P}) & \stackrel{\text{IH}}{<} \text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes (\text{rank}(q_0, \mathfrak{Q}) + \text{rank}(I)) \\ & \stackrel{\text{Lem. 2.10}}{\leq} \text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes (\text{rank}(q_2, \mathfrak{Q})) \\ & \leq \text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes (\text{rank}(q, \mathfrak{Q})). \end{aligned}$$

By Lemma 3.15 all elements below $\eta(0_1, 0_2, \dots, 0_n, q)$ are of the form $\eta(c_1, c_2, \dots, c_n, q_1)$ with $q_1 < q$. Thus,

$$\text{rank}(\eta(0_1, 0_2, \dots, 0_n, q) \leq \text{rank}(\mathfrak{C}_1) \otimes \dots \otimes \text{rank}(\mathfrak{C}_n) \otimes (\text{rank}(q, \mathfrak{Q})),$$

which proves the lemma. \square

Corollary 3.17. *Assume that \mathfrak{P} is a wulpo that is a box-augmentation of connected wulpo's $\mathfrak{Q}_1, \dots, \mathfrak{Q}_n$. Then*

$$\text{rank}(\mathfrak{P}) \leq \text{rank}(\mathfrak{Q}_1) \otimes \dots \otimes \text{rank}(\mathfrak{Q}_n).$$

Proof. By Corollary 3.12 we can assume that $\mathfrak{Q}_2, \dots, \mathfrak{Q}_n$ are ordinals with minimal elements $0_2, \dots, 0_n$. For each node $p \in \mathfrak{P}$ there is a maximal branching free interval $[p_0, p_1) \subseteq \mathfrak{P}$ such that $p \in [p_0, p_1)$ (where $p_1 = \infty$ is possible). There is a $q_0 \in \mathfrak{Q}_1$ such that $p_0 =$

$\eta(q_0, 0_2, \dots, 0_n)$ and there is a $q_1 \in \mathfrak{Q}_1 \cup \{\infty\}$ such that $I := [q_0, q_1)$ is maximal branching free. By Lemma 3.16, we have

$$\begin{aligned} \text{rank}(p, \mathfrak{P}) &< \text{rank}(\mathfrak{Q}_2) \otimes \dots \otimes \text{rank}(\mathfrak{Q}_n) \otimes (\text{rank}(q_0, \mathfrak{Q}_1) + \text{rank}(I)) \\ &\leq \text{rank}(\mathfrak{Q}_2) \otimes \dots \otimes \text{rank}(\mathfrak{Q}_n) \otimes \text{rank}(\mathfrak{Q}_1), \end{aligned}$$

where in case that $p = p_0$ the strict inequality follows from $\text{rank}(I) \geq 1$. \square

Having studied connected wulpo's, we now have to deal with disconnected ones. For this purpose, we have to restrict our attention to tamely colorable box-augmentations. As a first step, we analyze boxes of antichains. If a wulpo \mathfrak{P} is a tamely colorable box-augmentation of n antichains then the rank of \mathfrak{P} is bounded by some constant that only depends on the tame colorings of the antichains and on n . In order to prove this fact, we first introduce a notion of *same factor equivalence* on the elements of a box-augmentation. Elements are equivalent with respect to this equivalence if and only if their preimages in each of the factors of the box are contained in the same connected component.

Definition 3.18. Let \mathfrak{P} be a wulpo that is a box-augmentation of (wulpo's) $(\mathfrak{P}_1, \dots, \mathfrak{P}_n)$. Let $\mathfrak{P}_i = (P_i, \leq_i)$. For $a_i, b_i \in P_i$ we write $a_i \equiv_i b_i$ if a_i and b_i are in the same connected component of \mathfrak{P}_i . Due to upwards linearity, this is equivalent to saying that $a_i \equiv_i b_i$ if there exists some $c_i \in P_i$ such that $a_i \leq_i c_i$ and $b_i \leq_i c_i$.

Let η be the bijection witnessing that \mathfrak{P} is box-augmentation of $(\mathfrak{P}_1, \dots, \mathfrak{P}_n)$. For $a, b \in P$ we define the *same factors equivalence* by $a \equiv b$ if $a_i \equiv_i b_i$ for all $1 \leq i \leq n$ where $a = \eta(a_1, \dots, a_n)$ and $b = \eta(b_1, \dots, b_n)$. For $p \in P$ we write $[p]_{\equiv}$ for the equivalence class $\{p' \mid p' \equiv p\}$.

Lemma 3.19. Let $\mathfrak{P} = (P, \leq)$ be a wulpo that is a box-augmentation of $(\mathfrak{P}_1, \dots, \mathfrak{P}_n)$ via the bijection η . If $C \subseteq P$ is a chain that contains elements of k distinct \equiv -classes, then \mathfrak{P} contains a chain C' of length k such that $C' \subseteq \eta(A_1 \times \dots \times A_n)$ where A_i is an antichain in \mathfrak{P}_i .

Proof. Let $a^1 < a^2 < \dots < a^k$ be a chain of pairwise \equiv -inequivalent elements of \mathfrak{P} . There are $a_i^j \in \mathfrak{P}_i$ such that $a^j = \eta(a_1^j, \dots, a_n^j)$. Let

$$A_i^j := \{a_i^{j'} \mid a_i^j \equiv_i a_i^{j'}\}$$

be the connected component of a_i^j restricted to those elements that appear as factors of the $a^{j'}$ with $1 \leq j, j' \leq k$. Due to upwards linearity, for all $1 \leq i \leq n$ and $1 \leq j \leq k$ there is a minimal $m_i^j \in P_i$ such that for all $x \in A_i^j$, $x \leq_i m_i^j$. By definition, we have

$$m_i^j = m_i^{j'} \Leftrightarrow a_i^j \equiv_i a_i^{j'} \Leftrightarrow A_i^j = A_i^{j'}. \quad (3.14)$$

By definition of a box-augmentation we have

$$a^j = \eta(a_1^j, \dots, a_n^j) \leq \eta(m_1^j, a_2^j, \dots, a_n^j) \leq \dots \leq \eta(m_1^j, \dots, m_n^j).$$

Thus, $a^1 \leq a^j \leq \eta(m_1^j, \dots, m_n^j)$ for all $1 \leq j \leq k$. Due to upwards linearity, the set

$$C' := \{\eta(m_1^j, \dots, m_n^j) \mid 1 \leq j \leq k\}$$

forms a chain. Note that $\eta(m_1^j, \dots, m_n^j) = \eta(m_1^{j'}, \dots, m_n^{j'})$ would imply $m_i^j = m_i^{j'}$, i.e., $a_i^j \equiv_i a_i^{j'}$ (by (3.14)) for all $1 \leq i \leq n$. But this would lead to the contradiction $a^j \equiv a^{j'}$. Thus, the chain C' consists of k elements. We conclude by proving that $A_i := \{m_i^j \mid 1 \leq$

$j \leq k\}$ is an antichain in \mathfrak{P}_i . Heading for a contradiction assume that there are $j \neq j'$ such that $m_i^j < m_i^{j'}$. Then clearly $m_i^j \equiv_i m_i^{j'}$ holds and hence also $a_i^j \equiv_i a_i^{j'}$ holds. With (3.14) we conclude that $m_i^j = m_i^{j'}$ contradicting our assumption $m_i^j < m_i^{j'}$. \square

The previous result can be used to bound the length of ordinals occurring in the image of antichains in a tamely colorable box-augmentation.

Lemma 3.20. *Let $\mathfrak{P} = (P, \leq)$ be a wulpo which is a tamely colorable box-augmentation of $(\mathfrak{P}_1, \mathfrak{P}_2, \dots, \mathfrak{P}_n)$ via η . There is a constant $c \in \mathbb{N}$ such that the following holds: For all choices of antichains $A_i \subseteq P_i$ (for each $1 \leq i \leq n$), the substructure of \mathfrak{P} induced by $\eta(\prod_{i=1}^n A_i)$ does not contain a chain of length c .*

Proof. For each $1 \leq i \leq n$ we fix a finite coloring $\sigma_i : P_i \times P_i \rightarrow C_i$ of \mathfrak{P}_i such that the map $\sigma : P \times P \rightarrow C$ with $C = \prod_{1 \leq i \leq n} C_i$ and

$$\sigma(\eta(p_1, \dots, p_n), \eta(q_1, \dots, q_n)) = (\sigma_1(p_1, q_1), \dots, \sigma_n(p_n, q_n)).$$

is a finite coloring of \mathfrak{P} .

By Ramsey's theorem [23] there exists a constant $c \in \mathbb{N}$ such that every complete simple graph with at least c nodes whose edges are colored by $|C|$ colors contains a monochromatic triangle.

For the sake of a contradiction, assume that for each $1 \leq i \leq n$ there exists an antichain A_i in \mathfrak{P}_i such that $\eta(\prod_{i=1}^n A_i)$ contains a chain of length c . Due to the choice of c there exist three elements $p^1 > p^2 > p^3$ in this chain and a color $\bar{c} = (c_1, \dots, c_n) \in C$ such that

$$\sigma(p^1, p^2) = \sigma(p^1, p^3) = \sigma(p^2, p^3) = \bar{c}.$$

Let $p^j = \eta(p_1^j, \dots, p_n^j)$, where $p_i^j \in A_i$. For each $1 \leq i \leq n$ we obtain

$$\sigma_i(p_i^1, p_i^2) = \sigma_i(p_i^1, p_i^3) = \sigma_i(p_i^2, p_i^3) = c_i.$$

For $p' = \eta(p_1^2, p_2^1, \dots, p_n^1)$ we get $\sigma(p^1, p^3) = \sigma(p', p^3)$. Since $p^1 > p^3$, this implies $p' > p^3$. As the elements above p^3 are linearly ordered, we conclude that p' and p^1 are comparable. Recall that $\eta(\cdot, p_2^1, \dots, p_n^1)$ embeds \mathfrak{P}_1 into \mathfrak{P} . We conclude that p_1^1 and p_1^2 are comparable in \mathfrak{P}_1 . Since A_1 is an antichain and $p_1^1, p_1^2 \in A_1$, this implies $p_1^1 = p_1^2$.

Analogous arguments for the other coordinates show that $p_i^1 = p_i^2$ for each $1 \leq i \leq n$, i.e., $p^1 = p^2$. However, this contradicts $p^1 > p^2$. \square

We now head for the following result. Given a wulpo that is a tamely colorable box-augmentation we can write it as a finite sum-augmentation of wulpo's whose rank is bounded in terms of the ranks of the connected components of the factors of the box.

Lemma 3.21. *Let $\mathfrak{P} = (P, \leq)$ be a countable wulpo which is a tamely colorable box augmentation of $(\mathfrak{P}_1, \dots, \mathfrak{P}_n)$ via the bijection η . Let $c \in \mathbb{N}$ such that the image of antichains under η does not contain a chain of length c (exists by Lemma 3.20). Then there is a map $\mu : P \rightarrow \{1, \dots, c\}$ such that for all chains $L \subseteq P$, and all $p, p' \in L$ we have $\mu(p) = \mu(p')$ if and only if $p' \equiv p$.*

Proof. Fix an enumeration m_1, m_2, \dots of the minimal elements of \mathfrak{P} . Let μ_0 be the partial function with empty domain. We define μ as the limit of partial functions μ_i satisfying the lemma (restricted to their domain). This limit μ is a total function because we guarantee

that

$$\text{dom}(\mu_i) = \bigcup_{j=1}^i [m_j, \infty).$$

Assume that μ_j has already been defined.

First we set $\mu_{j+1}(x) = \mu_j(x)$ for all $x \in \text{dom}(\mu_j)$. By transfinite induction we extend μ_{j+1} to all nodes $y \in [m_{j+1}, \infty) \setminus \bigcup_{i=1}^j [m_i, \infty)$. For this purpose assume that $y \in [m_{j+1}, \infty) \setminus \bigcup_{i=1}^j [m_i, \infty)$ and assume that we have defined $\mu_{j+1}(x)$ for all $m_{j+1} \leq x < y$. Hence, the domain of the current μ_{j+1} is

$$D = \bigcup_{i=1}^j [m_i, \infty) \cup [m_{j+1}, y).$$

We assume that for all chains $L \subseteq D$, and all $p, p' \in L$ we have $\mu_{j+1}(p) = \mu_{j+1}(p')$ if and only if $p' \equiv p$. Let

$$J = \{1, \dots, c\} \setminus \mu_{j+1}(D \cap [m_{j+1}, \infty)).$$

Then $J \neq \emptyset$: $D \cap [m_{j+1}, \infty)$ is a chain. By Lemma 3.19 and the choice of c , each chain in \mathfrak{P} contains elements of at most $c-1$ many \equiv -classes. Hence, the elements from $D \cap [m_{j+1}, \infty)$ fall into at most $c-1$ many \equiv -classes. By our assumption on the current μ_{j+1} , this mapping takes at most $c-1$ different values among the elements from $D \cap [m_{j+1}, \infty)$. Hence, J is not empty.

Now we extend μ_{j+1} to y as follows:

$$\mu_{j+1}(y) := \begin{cases} \mu_{j+1}(p) & \text{if } p \in D \cap [m_{j+1}, \infty) \text{ and } p \equiv y, \\ \min(J) & \text{otherwise.} \end{cases}$$

Each of the μ_j is a well-defined partial function and setting $\mu = \bigcup_{j \in \mathbb{N}} \mu_j$ settles the claim. \square

Corollary 3.22. *Let \mathfrak{P} , c , and μ be defined as in Lemma 3.21. Let $1 \leq i \leq c$ and $a, b \in \mu^{-1}(i)$ such that a and b are in a connected component of the suborder induced by $\mu^{-1}(i) \subseteq \mathfrak{P}$. Then we have $a \equiv b$.*

Proof. Let $\mu(a) = i = \mu(b)$ and assume that a and b are in a connected component of the suborder induced by $\mu^{-1}(i) \subseteq \mathfrak{P}$. Since this suborder is again a wulpo, there exists $p \in \mu^{-1}(i)$ with $a \leq p$ and $b \leq p$. Lemma 3.21 together with $\mu(a) = \mu(p) = \mu(b)$ implies $a \equiv p \equiv b$, i.e., $a \equiv b$ by transitivity of \equiv . \square

Corollary 3.23. *Let \mathfrak{P} be a wulpo that is box-augmentation of (wulpo's) $(\mathfrak{P}_1, \mathfrak{P}_2, \dots, \mathfrak{P}_m)$ such that there is some $n \in \mathbb{N}$ with $\text{rank}(\mathfrak{P}_j) < \omega^{\omega^n}$ for all $1 \leq j \leq m$.*

- (1) *For μ as in Lemma 3.21, $\text{rank}(\mathfrak{P} \upharpoonright_{\mu^{-1}(i)}) \leq \omega^{\omega^n}$ for all i in the range of μ .*
- (2) *Moreover, if \mathfrak{P} is a forest and there is $n \in \mathbb{N}$ with $\text{rank}(\mathfrak{P}_j) < \omega^n$ for all $1 \leq j \leq m$, then $\text{rank}(\mathfrak{P} \upharpoonright_{\mu^{-1}(i)}) \leq \omega^n$ for all i in the range of μ .*

Proof. Let \mathfrak{C}_j be the set of connected components of \mathfrak{P}_j for all $1 \leq j \leq m$. Due to Corollary 3.22, $\mathfrak{P} \upharpoonright_{\mu^{-1}(i)}$ is isomorphic to some suborder of the disjoint union

$$\bigsqcup_{(D_1, \dots, D_m) \in \mathfrak{C}_1 \times \dots \times \mathfrak{C}_m} \mathfrak{P} \upharpoonright_{\eta(D_1 \times \dots \times D_m)}.$$

Note that every $D_j \in \mathfrak{C}_j$ is a substructure of \mathfrak{P}_j , hence $\text{rank}(D_j) < \omega^{\omega^n}$. Due to Lemma 3.17,

$$\text{rank}(\mathfrak{P}|_{\eta(D_1 \times \dots \times D_m)}) \leq \text{rank}(D_1) \otimes \dots \otimes \text{rank}(D_m) \stackrel{(2.3)}{<} \omega^{\omega^n}.$$

Thus,

$$\text{rank} \left(\bigsqcup_{(D_1, \dots, D_m) \in \mathfrak{C}_1 \times \dots \times \mathfrak{C}_m} \mathfrak{P}|_{\eta(D_1 \times \dots \times D_m)} \right) \leq \omega^{\omega^n}.$$

Hence, the same holds for the substructure $\mathfrak{P}|_{\mu^{-1}(i)}$.

If \mathfrak{P} is a forest, then without loss of generality $\mathfrak{P}_2, \dots, \mathfrak{P}_m$ are disjoint unions of finite ordinals (cf. Remark 3.13). Thus, $\text{rank}(\mathfrak{P}|_{\eta(D_1 \times \dots \times D_m)}) \leq \text{rank}(D_1) \otimes c = \alpha \otimes c$ for some finite ordinal $c < \omega$ and some $\alpha < \omega^n$. We get

$$\alpha \otimes c \leq \underbrace{\alpha \oplus \alpha \oplus \dots \oplus \alpha}_{c \text{ many}} \stackrel{(2.2)}{<} \omega^n,$$

where the first inequality follows from the fact that every linearization of the direct product $\alpha \otimes c$ can be viewed as a linearization of the disjoint union of c copies of α (one can actually show that equality holds at the place of the first inequality). We can now conclude completely analogously to the wulpo case. \square

We are now prepared to establish box-indecomposable intervals with respect to the rank function on the domain of wulpo's. Recall that we denote by rank_{ul} the function rank restricted to wulpo's and by rank_F its restriction to well-founded forests.

Proposition 3.24. *Let $i \in \mathbb{N}$.*

- (1) *The interval $[\omega^{\omega^i}, \omega^{\omega^{i+1}}]$ is rank_{ul} -tamely-colorable-box-indecomposable.*
- (2) *The interval $[\omega^i, \omega^{i+1}]$ is rank_F -tamely-colorable-box-indecomposable.*

Proof. For the first claim let $\mathfrak{P} = (P, \leq)$ be a wulpo of rank $\omega^{\omega^{i+1}}$. Heading for a contradiction, assume that it is a tamely colorable box-augmentation of wulpo's $(\mathfrak{P}_1, \dots, \mathfrak{P}_n)$ such that $\text{rank}(\mathfrak{P}_j) < \omega^{\omega^i}$ for all $1 \leq j \leq n$.

Take $c \in \mathbb{N}$ and the mapping $\mu : P \rightarrow \{1, \dots, c\}$ from Lemma 3.21. The preimages $\mu^{-1}(i)$ of μ induce a sum-decomposition of \mathfrak{P} into c parts $\mathfrak{P}|_{\mu^{-1}(1)}, \dots, \mathfrak{P}|_{\mu^{-1}(c)}$. Since $\omega^{\omega^{i+1}}$ is rank-sum-indecomposable by Proposition 3.10, we can assume (perhaps after renaming of colors) that $\text{rank}(\mathfrak{P}|_{\mu^{-1}(1)}) = \omega^{\omega^{i+1}}$. But Corollary 3.23 implies that $\text{rank}(\mathfrak{P}|_{\mu^{-1}(1)}) \leq \omega^{\omega^i}$ which is clearly a contradiction.

For the case of forests, we use exactly the same arguments, and the second part of Corollary 3.23 yields a similar contradiction. \square

Using Propositions 3.9, 3.10 and 3.24, we obtain directly the desired bounds on the ranks.

Theorem 3.25.

- (1) *Every tree-automatic wulpo has rank strictly below ω^{ω^ω} .*
- (2) *Every tree-automatic well-founded forest has rank strictly below ω^ω .*

Proof. Concerning the first claim, let $\alpha_i = \omega^{\omega^i}$ for each $i \in \mathbb{N}$. Due to Propositions 3.10 and 3.24 (part (1)), the sequence $\alpha_0 < \alpha_1 < \alpha_2 < \dots$ satisfies the conditions of Proposition 3.9. For the sake of a contradiction to this latter proposition, assume that \mathfrak{P} is a tree-automatic wulpo of rank at least ω^{ω^ω} . Let φ be the formula $\varphi(x, y) = x < y$. By well-foundedness, for each $i \in \mathbb{N}$ there is a $p \in P$ such that $\mathfrak{P}|_{\varphi, p}$ has rank α_i , contradicting Proposition 3.9.

For (2), we proceed analogously with $\alpha_i = \omega^i$. This time we use the second part of Proposition 3.24. \square

Note that this result has an analogous counterpart in the world of string-automatic structures: Khoussainov and Minnes [16] and Delhommé [9] independently proved that the ordinal ranks of string-automatic well-founded partial orders are the ordinals strictly below ω^ω and this bound is optimal because for each $\alpha < \omega^\omega$ the ordinal α is string-automatic. In contrast, we [14, 15] proved that every string-automatic forest has rank strictly below ω^2 and this bound is also optimal. In the next section, we construct tree-automatic well-founded trees of all ranks strictly below ω^ω . Moreover, since ordinals are wulpo's, our bound on the ranks of wulpo's is also optimal. Thus, tree-automatic forests realize much smaller ranks than tree-automatic partial orders (as it is also the case in the string-automatic setting). Note that there is no non-trivial bound on the ranks of arbitrary well-founded tree-automatic partial orders. But our new result on the subclass of all wulpo's and the analogy to the string-automatic case support the following conjecture.

Conjecture 3.26. Every tree-automatic well-founded partial order has rank below ω^{ω^ω} .

Beside upwards linearity, well-founded forests also have the property that one cannot embed an infinite linear order. One might wonder whether the class of tree-automatic partial orders without infinite linear suborder already satisfies that the ranks are bounded by ω^ω . The example below answers this question negatively by showing that for each $\alpha < \omega^{\omega^\omega}$ there exists such a partial order of rank α .

Example 3.27. Let $\alpha < \omega^{\omega^\omega}$ be an ordinal. We show that the direct product of α and ω^* (the reverse of ω) as *strict* partial orders is a tree-automatic partial order which has rank α and contains no infinite linear suborder.

More formally, we consider the partial order $\mathfrak{P} = (P, \preceq)$ on

$$P = \{(\beta, k) \mid \beta < \alpha, k \in \mathbb{N}\}$$

whose induced strict partial order \prec is given by

$$(\beta, k) \prec (\beta', k') \iff \beta < \beta' \text{ and } k > k'.$$

As a direct product of the tree-automatic structures α and ω^* , \mathfrak{P} is also tree-automatic.

Note that \mathfrak{P} does not contain an infinite linear suborder because any such suborder could be projected to an infinite strictly descending sequence in α or to an infinite strictly ascending sequence in ω^* .

We prove that $\text{rank}(\mathfrak{P}) = \text{rank}(\alpha)$ by transfinite induction. For $\beta = 0$ and $k \in \omega^*$ we have $\text{rank}((\beta, k), \mathfrak{P}) = 0$ because (β, k) is a minimal element. For $\beta > 0$, $k \in \omega^*$,

$$\begin{aligned} \text{rank}((\beta, k), \mathfrak{P}) &= \sup\{\text{rank}((\beta', k'), \mathfrak{P}) + 1 \mid (\beta', k') \prec (\beta, k) \in P\} \\ &= \sup\{\beta' + 1 \mid \beta' < \beta\} = \beta. \end{aligned}$$

Thus, we conclude that

$$\text{rank}(\mathfrak{P}) = \sup\{\text{rank}((\beta, k), \mathfrak{P}) + 1 \mid (\beta, k) \in \mathfrak{P}\} = \sup\{\beta + 1 \mid \beta < \alpha\} = \alpha.$$

The structure \mathfrak{P} also shows that Delhommé's decomposition technique cannot be used to prove our conjecture that all tree-automatic partial orders have ranks below ω^{ω^ω} . Note that for all $\alpha' \in \alpha$ the substructure $\mathfrak{P} \upharpoonright_{\{(\alpha', n) \mid n \in \omega^*\}}$ is isomorphic to the countably infinite antichain $(\omega^*, =)$. Similarly, for all $n \in \omega^*$ the substructure $\mathfrak{P} \upharpoonright_{\{(\alpha', n) \mid \alpha' \in \alpha\}}$ is isomorphic to the countably infinite antichain $(\alpha, =)$. Thus, \mathfrak{P} is a partial order of rank α that is a tamely colorable box-augmentation of two antichains. Since infinite antichains have rank 1, we conclude that applications of box-decomposition to well-founded partial orders may destroy all information about the rank. Thus, there is no hope that one could provide bounds on the ranks of tree-automatic well-founded partial orders using Delhommé's decomposition technique.

4. UPPER BOUND FOR THE ISOMORPHISM PROBLEM FOR WELL-FOUNDED TREES

We first introduce hyperarithmetical sets. Afterwards we provide the upper bound for the isomorphism problem of well-founded tree-automatic trees.

4.1. Hyperarithmetical sets. We use standard terminology concerning recursion theory; see e.g. [24]. We use the definition of the hyperarithmetical hierarchy from Ash and Knight [1], see also [11]. We first define inductively a set of *ordinal notations* $O \subseteq \mathbb{N}_{>0}$. Simultaneously we define a mapping $a \mapsto |a|_O$ from O into ordinals and a strict partial order $<_O$ on O . The set O is the smallest subset of $\mathbb{N}_{>0}$ satisfying the following conditions:

- $1 \in O$ and $|1|_O = 0$, i.e., 1 is a notation for the ordinal 0.
- If $a \in O$, then also $2^a \in O$. We set $|2^a|_O = |a|_O + 1$ and let $b <_O 2^a$ if and only if $b = a$ or $b <_O a$.
- If $e \in \mathbb{N}$ is such that Φ_e (the e^{th} partial computable function) is total, $\Phi_e(n) \in O$ for all $n \in \mathbb{N}$, and $\Phi_e(0) <_O \Phi_e(1) <_O \Phi_e(2) <_O \dots$, then also $3 \cdot 5^e \in O$. We set $|3 \cdot 5^e|_O = \sup\{|\Phi_e(n)|_O \mid n \in \mathbb{N}\}$ and let $b <_O 3 \cdot 5^e$ if and only if there exists $n \in \mathbb{N}$ with $b <_O \Phi_e(n)$.

An ordinal α is *computable* if there exists $a \in O$ with $|a|_O = \alpha$. The smallest non-computable ordinal is the Church-Kleene ordinal ω_1^{ck} . If $a \in O$ then the restriction of the partial order $(O, <_O)$ to $B = \{b \in O \mid b <_O a\}$ is isomorphic to the ordinal $|a|_O$ [1, Proposition 4.9]. Moreover, the set B is computably enumerable and an index for B can be computed from a [1, Proposition 4.10].

Next, we define the *hyperarithmetical hierarchy* based on ordinal notations. For this we define sets $H(a)$ for each $a \in O$ as follows:

- $H(1) = \emptyset$,
- $H(2^b) = H(b)'$ (the Turing jump of $H(b)$; see e.g. [24]),
- $H(3 \cdot 5^e) = \{\langle b, n \rangle \mid b <_O 3 \cdot 5^e, n \in H(b)\}$; here $\langle \cdot, \cdot \rangle$ denotes a computable pairing function.

Spector has shown that $|a|_O = |b|_O$ implies that $H(a)$ and $H(b)$ are Turing equivalent. The levels of the hyperarithmetical hierarchy can be defined as follows, where α is a computable ordinal.

- If α is infinite, then Σ_α^0 is the set of all subsets $A \subseteq \mathbb{N}$ that are recursively enumerable in some $H(a)$ with $|a|_O = \alpha$ (by Spector's theorem, the concrete choice of a is irrelevant). For $n > 0$ finite, one defines Σ_n^0 as the set of all subsets $A \subseteq \mathbb{N}$ that are recursively enumerable in $H(a)$ with $|a|_O = n - 1$.³
- Π_α^0 is the set of all complements of Σ_α^0 -sets.
- $\Delta_\alpha^0 = \Sigma_\alpha^0 \cap \Pi_\alpha^0$

A relation $R \subseteq \mathbb{N}^k$ is X_α^0 (with $X \in \{\Sigma, \Pi, \Delta\}$) if the set $\{\langle x_1, \dots, x_k \rangle \mid (x_1, \dots, x_k) \in R\}$ is X_α^0 , where $\langle \cdot, \dots, \cdot \rangle$ denotes a computable encoding of k -tuples.

For any two computable ordinals α, β , $\alpha < \beta$ implies $\Sigma_\alpha \cup \Pi_\alpha \subsetneq \Delta_\beta$. The union of all classes Σ_α^0 where $\alpha < \omega_1^{\text{ck}}$ yields the class of all *hyperarithmetical sets*. By a classical result of Kleene, the hyperarithmetical sets are exactly the sets in $\Delta_1^1 = \Sigma_1^1 \cap \Pi_1^1$, where Σ_1^1 is the first existential level of the analytical hierarchy, and Π_1^1 is the set of all complements of Σ_1^1 -sets.

For our purposes, it is convenient to present an alternative definition of the hyperarithmetical hierarchy using computable infinitary formulas, see [1, Chapter 7]. Fix a predicate $R(\bar{x}) \subseteq \mathbb{N}^k$ where $k \geq 1$. If R is computable, then a Σ_0^0 (resp. Π_0^0) *index* for R is a triple $(\Sigma, 0, e)$ (resp. $(\Pi, 0, e)$) where e is an index for R . Next, let $0 < \alpha < \omega_1^{\text{ck}}$ be a computable ordinal. *Case 1.* $\alpha = \beta + 1$ is a successor ordinal. Then, a Σ_α^0 (resp. Π_α^0) *index* for R is a triple (Σ, a, e) (resp. (Π, a, e)) where $|a|_O = \alpha$ and e is a Π_β^0 (resp. Σ_β^0) *index* for a predicate $P(\bar{x}, y) \subseteq \mathbb{N}^{k+1}$ such that for all $\bar{x} \in \mathbb{N}^k$:

$$R(\bar{x}) \Leftrightarrow \exists y : P(\bar{x}, y) \quad \left(\text{resp. } R(\bar{x}) \Leftrightarrow \forall y : P(\bar{x}, y) \right).$$

Case 2. α is a limit ordinal. Then, a Σ_α^0 (resp. Π_α^0) *index* for R is a triple (Σ, a, e) (resp. (Π, a, e)), where $|a|_O = \alpha$ and Φ_e is a total computable function such that the following holds: For all $n \in \mathbb{N}$, $\Phi_e(n)$ is a $\Pi_{\beta_n}^0$ (resp. $\Sigma_{\beta_n}^0$) *index* for some predicate $P_n(\bar{x}) \subseteq \mathbb{N}^k$, $\beta_0 < \beta_1 < \dots < \alpha$ with $\sup\{\beta_n \mid n \in \mathbb{N}\} = \alpha$, and for all $\bar{x} \in \mathbb{N}^k$:

$$R(\bar{x}) \Leftrightarrow \bigvee_{i \in \mathbb{N}} P_i(\bar{x}) \quad \left(\text{resp. } R(\bar{x}) \Leftrightarrow \bigwedge_{i \in \mathbb{N}} P_i(\bar{x}) \right).$$

Essentially, we can view a Σ_α^0 (resp. Π_α^0) *index* as a finite representation of a *computable infinitary formula* that defines the corresponding Σ_α^0 (resp. Π_α^0) predicate. For instance, the Σ_α^0 *index* (Σ, a, e) , where $|a|_O$ is a limit ordinal, represents the computable infinitary formula $\bigvee_{i \in \mathbb{N}} \varphi_i$, where φ_i is the computable infinitary formula represented by the index $\Phi_e(i)$. In this paper we use the notions of Σ_α^0 (resp. Π_α^0) predicates and indices interchangeably. Formally, we do not allow negation in computable infinitary formulas. But if $\varphi(\bar{x})$ defines the Σ_α^0 (resp. Π_α^0) set $A \subseteq \mathbb{N}^k$, then one can construct effectively a Π_α^0 (resp. Σ_α^0) formula for $\mathbb{N}^k \setminus A$; and we therefore may define this formula as $\neg\varphi(\bar{x})$, see [1, Theorem 7.1].

4.2. Hyperarithmetical Upper Bound. It turns out that the rank for well-founded computable trees yields an upper bound on the recursion-theoretic complexity of the isomorphism problem. Recall that we defined forests as particular partial orders. For the isomorphism problem, it is useful to assume that also the direct successor relation is computable. When speaking of a computable forest in the following theorem, we mean a forest $\mathfrak{F} = (F, \leq)$

³The distinction between finite and infinite ordinals is made in order to have a correspondence between the arithmetical hierarchy and the finite part of the hyperarithmetical hierarchy, see also [1, pp. 74 and 75].

such that F , \leq , and the direct successor relation $E_{\mathfrak{F}}$ are all computable.⁴ Note that the direct successor relation of a tree-automatic forest is even tree-automatic because it is first-order definable.

Theorem 4.1. *Let α be a computable ordinal and assume that $\alpha = \lambda + k$, where $k \in \mathbb{N}$ and λ is not a successor ordinal. The isomorphism problem for well-founded computable trees of rank at most α belongs to level $\Pi_{\lambda+2k+1}^0$ of the hyperarithmetical hierarchy.*

Proof. Let us fix a well-founded forest $\mathfrak{F} = (F, \leq)$. We define a computable infinitary $\Pi_{\lambda+2k+1}^0$ formula expressing $\mathfrak{F}(x) \cong \mathfrak{F}(y)$ for nodes x and y of \mathfrak{F} of rank at most $\lambda + k$, where $k \in \mathbb{N}$ and $\lambda = 0$ or λ is a limit ordinal. This suffices because the disjoint union of two computable trees is a computable forest.

Let $E = E_{\mathfrak{F}}$ be the direct successor relation of \mathfrak{F} . For every ordinal α we define a computable infinitary formula $\text{iso}_{\alpha}(x, y)$ over \mathfrak{F} as follows: Let

$$\text{iso}_0(x, y) = \forall z (\neg(x > z) \wedge \neg(y > z)).$$

This is a Π_1^0 formula. For a successor ordinal $\alpha + 1$ let⁵ $\text{iso}_{\alpha+1}(x, y)$ be

$$\forall u \in E(x) \cup E(y) \forall \ell \geq 1 \left(\begin{array}{l} \exists^{\geq \ell} v \in E(x) \text{ iso}_{\alpha}(u, v) \iff \exists^{\geq \ell} v \in E(y) \text{ iso}_{\alpha}(u, v) \\ \wedge \text{ iso}_{\alpha}(u, u). \end{array} \right)$$

Finally, for a limit ordinal α , we define $\text{iso}_{\alpha}(x, y)$ as

$$\bigwedge_{\beta < \alpha} \forall u \in E(x) \cup E(y) \forall \ell \geq 1 \left(\begin{array}{l} \exists^{\geq \ell} v \in E(x) \text{ iso}_{\beta}(u, v) \iff \exists^{\geq \ell} v \in E(y) \text{ iso}_{\beta}(u, v) \\ \wedge \bigvee_{\beta < \alpha} \text{ iso}_{\beta}(u, u). \end{array} \right)$$

Let us argue by induction that $\text{iso}_{\lambda+k}$ (where $k \in \mathbb{N}$ and $\lambda = 0$ or λ a limit ordinal) is a $\Pi_{\lambda+2k+1}^0$ -formula. The case $\lambda = k = 0$ is clear. For the case $k = 0$ and $\lambda > 0$ note that the subformula $\forall u \in E(x) \cup E(y) : \bigvee_{\beta < \lambda} \text{ iso}_{\beta}(u, u)$ is $\Pi_{\lambda+1}^0$. Moreover, the subformula

$$\forall u \in E(x) \cup E(y) \forall \ell \geq 1 : (\exists^{\geq \ell} v \in E(x) : \text{ iso}_{\beta}(u, v) \iff \exists^{\geq \ell} v \in E(y) : \text{ iso}_{\beta}(u, v))$$

is $\Pi_{\lambda'+2k'+3}^0$, where $\beta = \lambda' + k' < \alpha$ (with $k' \in \mathbb{N}$ and $\lambda' = 0$ or λ' a limit ordinal). Since λ is a limit ordinal, we have $\lambda' + 2k' + 3 < \lambda$. This shows that $\text{iso}_{\lambda}(x, y)$ is a conjunction of a Π_{λ}^0 -formula and a $\Pi_{\lambda+1}^0$ -formula, and hence also $\Pi_{\lambda+1}^0$. Finally, for the case that $k > 0$ and $\lambda > 0$ note that if $\text{iso}_{\lambda+k}(x, y)$ is a $\Pi_{\lambda+2k+1}^0$ -formula, then $\text{iso}_{\lambda+k+1}(x, y)$ is a $\Pi_{\lambda+2k+3}^0$ -formula, since two alternating quantifier blocks are added for successor ordinals.

It remains to show that if $\text{rank}(x) \leq \alpha$ and $\text{rank}(y) \leq \alpha$, then $\mathfrak{F} \models \text{iso}_{\alpha}(x, y)$ if and only if $\mathfrak{F}(x) \cong \mathfrak{F}(y)$. This is the content of Claim 2 below. As an auxiliary step we show the following claim.

Claim 1. If $\mathfrak{F} \models \text{iso}_{\alpha}(x, y)$ then $\text{rank}(x) \leq \alpha$ and $\text{rank}(y) \leq \alpha$.

We prove Claim 1 by induction on α . The case $\alpha = 0$ is clear because trees of rank 0 do only consist of one node. Next, consider an ordinal $\alpha > 0$ such that the claim holds for all $\beta < \alpha$. Assume that $\mathfrak{F} \models \text{iso}_{\alpha}(x, y)$. Note that the second conjunct of

⁴On the other hand, if we would omit the requirement of a computable direct successor relation in Theorem 4.1, then we would only have to replace the constants in the theorem by larger values.

⁵We use $\exists^{\geq \ell} x : \varphi(x)$ as an abbreviation for $\exists x_1 \cdots \exists x_{\ell} : \bigwedge_{1 \leq i < j \leq n} x_i \neq x_j \wedge \bigwedge_{i=1}^n \varphi(x_i)$. The quantifier $\exists^{\geq \ell} x$ can be encoded by an ordinary single existential quantifier.

$\text{iso}_\alpha(x, y)$ and the induction hypothesis imply that all children of x have $\text{rank} < \alpha$. Thus, $\text{rank}(x) \leq \sup\{\beta + 1 \mid \beta < \alpha\} = \alpha$. For y we conclude analogously.

Claim 2. If $\text{rank}(x) \leq \alpha$ and $\text{rank}(y) \leq \alpha$, then $\mathfrak{F} \models \text{iso}_\alpha(x, y)$ if and only if $\mathfrak{F}(x) \cong \mathfrak{F}(y)$.

We prove Claim 2 by induction on α . Again, the case $\alpha = 0$ is clear. Next, consider some ordinal $\alpha > 0$ and assume that $\text{rank}(x) \leq \alpha$ and $\text{rank}(y) \leq \alpha$.

First, assume that $\mathfrak{F}(x) \cong \mathfrak{F}(y)$. Fix $\beta < \alpha$, $u \in E(x) \cup E(y)$, and $\ell \geq 1$. By assumption, we have $\text{rank}(u) < \alpha$. Furthermore, due to Claim 1, if $\beta < \text{rank}(u)$, then there is no v such that $\text{iso}_\beta(u, v)$ holds. Thus, in this case

$$\left(\exists^{\geq \ell} v \in E(x) : \text{iso}_\beta(u, v) \iff \exists^{\geq \ell} v \in E(y) : \text{iso}_\beta(u, v) \right).$$

holds. Moreover, if $\text{rank}(u) \leq \beta < \alpha$, then Claim 1 and the induction hypothesis imply that

$$\exists^{\geq \ell} v \in E(x) : \text{iso}_\beta(u, v) \iff \exists^{\geq \ell} v \in E(y) : \text{iso}_\beta(u, v)$$

is equivalent to

$$\exists^{\geq \ell} v \in E(x) : \mathfrak{F}(u) \cong \mathfrak{F}(v) \iff \exists^{\geq \ell} v \in E(y) : \mathfrak{F}(u) \cong \mathfrak{F}(v).$$

Since $\mathfrak{F}(x) \cong \mathfrak{F}(y)$ the latter is clearly satisfied. Moreover, by induction hypothesis for each $u \in E(x) \cup E(y)$ we have $\text{iso}_{\text{rank}(u)}(u, u)$ (note that $\text{rank}(u) < \alpha$). Hence $\bigvee_{\beta < \alpha} \text{iso}_\beta(u, u)$ holds as well. Thus $\text{iso}_\alpha(x, y)$ is satisfied.

For the other direction, assume that $\mathfrak{F} \models \text{iso}_\alpha(x, y)$. Then for each $u \in E(x) \cup E(y)$, there is some $\beta < \alpha$ such that $\text{iso}_\beta(u, u)$ holds. Due to Claim 1 this implies that $\text{rank}(u) < \alpha$ for each child u of x or y . Using the induction hypothesis and Claim 1, we conclude that for all $u \in E(x) \cup E(y)$ and all ordinals $\text{rank}(u) \leq \beta < \alpha$,

$$\exists^{\geq \ell} v \in E(x) : \text{iso}_\beta(u, v) \iff \exists^{\geq \ell} v \in E(y) : \text{iso}_\beta(u, v) \quad (4.1)$$

is equivalent to

$$\exists^{\geq \ell} v \in E(x) : \mathfrak{F}(u) \cong \mathfrak{F}(v) \iff \exists^{\geq \ell} v \in E(y) : \mathfrak{F}(u) \cong \mathfrak{F}(v). \quad (4.2)$$

Since $\text{iso}_\alpha(x, y)$ holds, for each $u \in E(x) \cup E(y)$ there is a $\text{rank}(u) \leq \beta < \alpha$ such that (4.1) holds. Hence, (4.2) holds for all children u of x or y , which implies that $\mathfrak{F}(x) \cong \mathfrak{F}(y)$. \square

Remark 4.2. If we are only interested in the isomorphism problem for computable trees of rank at most λ for a limit ordinal λ it suffices to consider the Π_λ^0 -formula

$$\bigwedge_{\beta < \alpha} \forall u \in E(x) \cup E(y) \forall \ell \geq 1 : (\exists^{\geq \ell} v \in E(x) : \text{iso}_\beta(u, v) \iff \exists^{\geq \ell} v \in E(y) : \text{iso}_\beta(u, v))$$

(where the formulas iso_β are constructed as above). This is because if the rank of the root is bounded by λ , any child of the root has rank strictly below λ and satisfies $\text{iso}_\beta(u, u)$ for some $\beta < \lambda$.

Corollary 4.3. *The isomorphism problem for well-founded tree-automatic trees belongs to $\Delta_{\omega^\omega}^0$.*

Proof. The $\Sigma_{\omega^\omega}^0$ formula $\bigvee_{\alpha < \omega^\omega} \text{iso}_\alpha(x, y)$ expresses $\mathfrak{F}(x) \cong \mathfrak{F}(y)$ for all x, y with $\text{rank}(x), \text{rank}(y) < \omega^\omega$.

With Remark 4.2, we also obtain a $\Pi_{\omega^\omega}^0$ formula expressing $\mathfrak{F}(x) \cong \mathfrak{F}(y)$ for all x, y with $\text{rank}(x), \text{rank}(y) < \omega^\omega$.

Due to Theorem 3.25, the rank of every tree-automatic well-founded tree is strictly below ω^ω . Thus, the isomorphism problem for these trees belongs to $\Sigma_{\omega^\omega}^0 \cap \Pi_{\omega^\omega}^0 = \Delta_{\omega^\omega}^0$. \square

5. LOWER BOUND FOR THE ISOMORPHISM PROBLEM FOR WELL-FOUNDED TREES

In this section, we prove hyperarithmetical lower bounds for the isomorphism problem for well-founded tree-automatic trees. More precisely, we show that for every ordinal ω^i , there exists a well-founded tree \mathfrak{T} such that the set of all tree-automatic copies of \mathfrak{T} is hard for the class $\Pi_{\omega^i}^0$. Moreover, for the class of all well-founded tree-automatic trees, we prove that the isomorphism problem is $\Delta_{\omega^\omega}^0$ -hard under Turing-reductions.

5.1. Isomorphism for computable trees of rank $< \omega^\omega$. Basically, our hardness proof is a reduction from computable well-founded trees to tree-automatic well-founded trees. For this, we make use of a construction from [11], which works for all computable ordinals. We use this construction only for ordinals strictly below ω^ω . In this section, a computable tree is a computable prefix-closed subset $\mathfrak{S} \subseteq \mathbb{N}_{>0}^*$.⁶ We identify \mathfrak{S} with the tree (\mathfrak{S}, \succeq) . Recall that \succeq is the inverse prefix relation. Hence, the empty word ε is the root of (\mathfrak{S}, \succeq) .

First, we have to fix the so called fundamental sequence for every limit ordinal $< \omega^\omega$. Each ordinal $\alpha < \omega^\omega$ can be written in its Cantor normal form as

$$\alpha = \omega^{e_i} \cdot n_i + \omega^{e_{i-1}} \cdot n_{i-1} + \cdots + \omega^{e_1} \cdot n_1,$$

where $e_i > e_{i-1} > \cdots > e_1 \geq 0$ and $n_j > 0$ are natural numbers for $1 \leq j \leq i$. Assume that $e_1 > 0$ so that α is a limit ordinal. Then

$$\alpha = \sup\{\alpha_k \mid k \geq 1\},$$

where we define α_k as

$$\alpha_k = \omega^{e_i} \cdot n_i + \omega^{e_{i-1}} \cdot n_{i-1} + \cdots + \omega^{e_2} \cdot n_2 + \omega^{e_1} \cdot (n_1 - 1) + \omega^{e_1-1} \cdot k + 1, \quad (5.1)$$

for $k \geq 1$. Note that α_k is a successor ordinal and we call $(\alpha_k)_{k \geq 1}$ the fundamental sequence.

Next, we define for every ordinal $\alpha < \omega^\omega$ a computable well-founded tree $\mathfrak{S}_\alpha \subseteq \mathbb{N}_{>0}^*$ by induction on α . Let $\mathfrak{S}_0 = \{\varepsilon\}$ be the tree consisting of a single node. If $\alpha = \beta + 1 < \omega^\omega$ is a successor ordinal then

$$\mathfrak{S}_{\beta+1} = \{nu \mid n \in \mathbb{N}_{>0}, u \in \mathfrak{S}_\beta\} \cup \{\varepsilon\}.$$

Hence, $\mathfrak{S}_{\beta+1}$ consists of \aleph_0 many copies of \mathfrak{S}_β together with a new root. Finally let $\alpha < \omega^\omega$ be a limit ordinal with the fundamental sequence $(\alpha_k)_{k \geq 1}$ defined in (5.1). Then

$$\mathfrak{S}_\alpha = \{ku \mid k \in \mathbb{N}_{>0}, u \in \mathfrak{S}_{\alpha_k}\} \cup \{\varepsilon\}.$$

Thus, \mathfrak{S}_α consists of all the trees \mathfrak{S}_{α_k} ($k \in \mathbb{N}_{>0}$) together with a new root. By induction on $\alpha < \omega^\omega$, it is straightforward to show that \mathfrak{S}_α is well-founded and computable. Moreover, also the set of leaves $\text{leaves}(\mathfrak{S}_\alpha) \subseteq \mathbb{N}_{>0}^*$ is computable. Let $A \subseteq \text{leaves}(\mathfrak{S}_\alpha)$. Then we denote the structure consisting of the tree \mathfrak{S}_α together with the additional unary predicate A by (\mathfrak{S}_α, A) .

The following result is implicitly shown in [11, Proposition 3.2], where it is stated for all computable ordinals. But since we only defined fixed fundamental sequences for limit ordinals below ω^ω , we restrict to ordinals below ω^ω .

⁶For technical reasons, it is useful to exclude 0. This makes the construction of tree-automatic copies easier.

Theorem 5.1. *Given a limit ordinal $\alpha < \omega^\omega$ and $k \in \mathbb{N}_{>0} \cup \{\infty\}$ (resp., a successor ordinal $\alpha < \omega^\omega$), one can compute indices of a computable subset $L_k^\alpha \subseteq \text{leaves}(\mathfrak{S}_\alpha)$ (resp., computable subsets $A_\alpha, E_\alpha \subseteq \text{leaves}(\mathfrak{S}_\alpha)$) such that the following holds:*

From a natural number $n \in \mathbb{N}$ and a Π_α^0 index (Π, a, e) for a set $P \subseteq \mathbb{N}$ one can compute an index for a computable subset $T_{P,n} \subseteq \text{leaves}(\mathfrak{S}_\alpha)$ such that the following holds:

- $(\mathfrak{S}_\alpha, A_\alpha) \not\cong (\mathfrak{S}_\alpha, E_\alpha)$ and $(\mathfrak{S}_\alpha, L_\infty^\alpha) \not\cong (\mathfrak{S}_\alpha, L_k^\alpha)$ for all $k \in \mathbb{N}_{>0}$, respectively.
- If α is a successor ordinal, then

$$(\mathfrak{S}_\alpha, T_{P,n}) \cong \begin{cases} (\mathfrak{S}_\alpha, A_\alpha) & \text{if } n \in P \\ (\mathfrak{S}_\alpha, E_\alpha) & \text{if } n \notin P. \end{cases}$$

- If α is a limit ordinal, then

$$(\mathfrak{S}_\alpha, T_{P,n}) \cong \begin{cases} (\mathfrak{S}_\alpha, L_\infty^\alpha) & \text{if } n \in P \\ (\mathfrak{S}_\alpha, L_k^\alpha) \text{ for some } k \in \mathbb{N}_{>0} & \text{if } n \notin P. \end{cases}$$

Proof. In Proposition 3.2 from [11], the computable subset $A_\alpha \subseteq \text{leaves}(\mathfrak{S}_\alpha)$ is replaced by a computable tree $\mathcal{A}_\alpha \subseteq \mathfrak{S}_\alpha$ (and similarly for $E_\alpha, L_k^\alpha, T_{P,n}$). For our purpose it is more convenient to work with computable subsets of the leaves of the fixed tree \mathfrak{S}_α . Nevertheless, the construction works analogously to the proof of Proposition 3.2 from [11], except for the induction base $\alpha = 1$. In [11], the tree \mathcal{A}_1 consists of a single node and the tree \mathcal{E}_1 consists of a root with infinitely many children (i.e., $\mathcal{E}_1 \cong \mathfrak{S}_1$). For the construction of $T_{P,n}$, one chooses a computable set $Q \subseteq \mathbb{N} \times \mathbb{N}$ such that $P' = \{n \mid \exists x Q(x, n)\}$ is the complement of the Π_1^0 -set P . Then the tree $\mathcal{T}_{P,n}$ consists of the root and $x \in \mathbb{N}_{>0}$ is a child of the root if and only if there exists $y \leq x$ satisfying $Q(y, n)$. Hence, $\mathcal{T}_{P,n} \cong \mathcal{E}_1$ if there exists x with $Q(x, n)$ and $\mathcal{T}_{P,n} \cong \mathcal{A}_1$ if $\neg Q(x, n)$ for all x .

In our context, we define the subsets $A_1, E_1, T_{P,n} \subseteq \mathbb{N}_{>0} = \text{leaves}(\mathfrak{S}_1)$ as follows: let $A_1 = \emptyset$, E_1 be the set of all non-zero even numbers, and $T_{P,n} = \{2x \mid \exists y \leq x : Q(y, n)\}$. Then, we have $(\mathfrak{S}_1, T_{P,n}) \cong (\mathfrak{S}_1, E_1)$ if there exists x with $Q(x, n)$ and $(\mathfrak{S}_1, T_{P,n}) \cong (\mathfrak{S}_1, A_1)$ if $\neg Q(x, n)$ for all x .

Let us now discuss the induction step. As remarked above, it works analogously to the proof of Proposition 3.2 from [11], but we present the details for the sake of completeness. Let $\alpha < \omega^\omega$, $n \in \mathbb{N}$, and (Π, a, e) be a Π_α^0 index for a set $P \subseteq \mathbb{N}$. Hence, $|a|_O = \alpha$. We distinguish three cases:

Case 1. $\alpha = \beta + 1$ for a successor ordinal $\beta < \omega^\omega$. Note that we can compute indices for the sets $A_\beta, E_\beta \subseteq \text{leaves}(\mathfrak{S}_\beta)$ from β . The following constructions are illustrated in Figure 2. We define the sets $A_\alpha, E_\alpha \subseteq \text{leaves}(\mathfrak{S}_\alpha) = \{mu \mid m \in \mathbb{N}_{>0}, u \in \text{leaves}(\mathfrak{S}_\beta)\}$ as follows (here and in the following, we use square brackets to enclose single numbers in finite words over \mathbb{N}):

$$\begin{aligned} A_\alpha &= \{mu \mid n \in \mathbb{N}_{>0}, u \in E_\beta\}, \\ E_\alpha &= \{[2m]u \mid n \in \mathbb{N}_{>0}, u \in E_\beta\} \cup \{[2m-1]u \mid m \in \mathbb{N}_{>0}, u \in A_\beta\}. \end{aligned}$$

Clearly, indices for these computable sets can be computed from the indices for A_β and E_β . Since $(\mathfrak{S}_\beta, E_\beta) \not\cong (\mathfrak{S}_\beta, A_\beta)$, we also have $(\mathfrak{S}_\alpha, E_\alpha) \not\cong (\mathfrak{S}_\alpha, A_\alpha)$.

The index e from (Π, a, e) is a Σ_β^0 index for a set $Q \subseteq \mathbb{N} \times \mathbb{N}$ such that $n \in P$ if and only if $(n, m) \in Q$ for all $m \in \mathbb{N}$. From e one can compute a Π_β^0 index f for the complement of Q [1, Theorem 7.1]. By induction, one can compute from f and $(n, m) \in \mathbb{N} \times \mathbb{N}$ an index

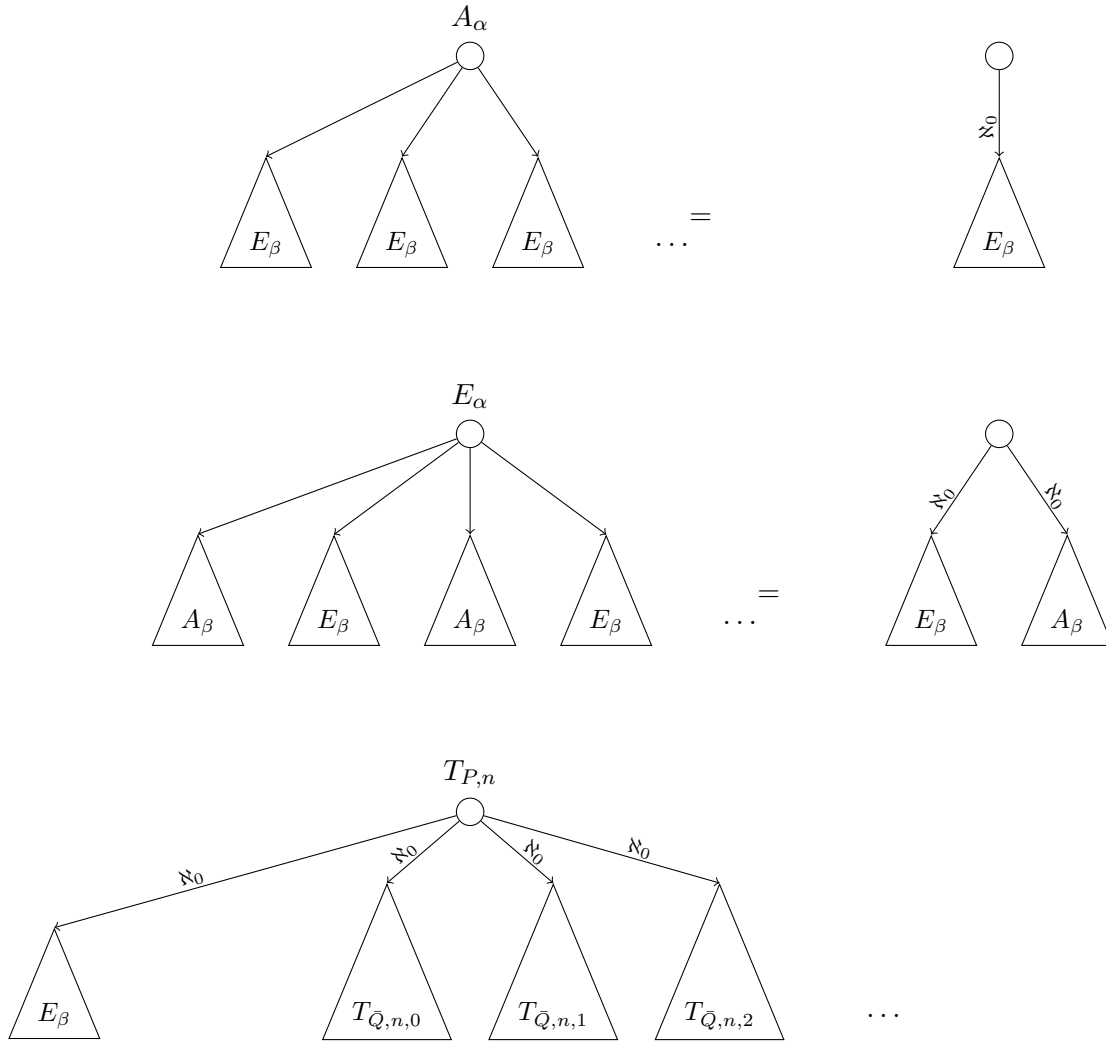


Figure 2: Trees for Case 1 from the proof of Theorem 5.1 (here and in the following pictures we write an \aleph_0 labeled edge from a node d to a subtree T if d has edges to countably infinitely many subtrees isomorphic to T).

of a computable set $T_{\overline{Q},n,m} \subseteq \text{leaves}(\mathfrak{S}_\beta)$ such that:⁷

$$(\mathfrak{S}_\beta, T_{\overline{Q},n,m}) \cong \begin{cases} (\mathfrak{S}_\beta, A_\beta) & \text{if } (n, m) \notin Q, \\ (\mathfrak{S}_\beta, E_\beta) & \text{if } (n, m) \in Q. \end{cases}$$

We define the computable set $T_{P,n} \subseteq \text{leaves}(\mathfrak{S}_\alpha)$ as follows (where we write $2^{\mathbb{N}}3^{\mathbb{N}}$ for $\{n \in \mathbb{N} \mid \exists k, l \in \mathbb{N} \ n = 2^k \cdot 3^l\}$):

$$T_{P,n} = \{[2^m 3^x]u \mid x, m \in \mathbb{N}, u \in T_{\overline{Q},n,m}\} \cup \{xu \mid x \in \mathbb{N}_{>0} \setminus 2^{\mathbb{N}}3^{\mathbb{N}}, u \in E_\beta\}.$$

⁷Formally, we have to encode the pair (n, m) into a single number in order to apply the induction hypothesis.

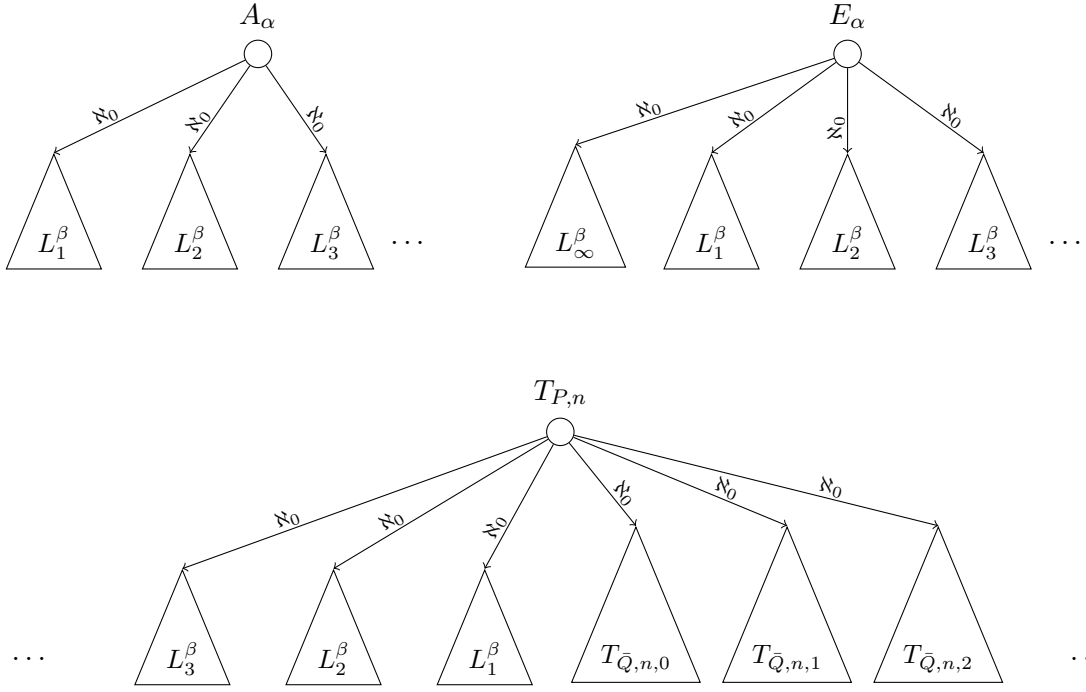


Figure 3: Trees for Case 2 from the proof of Theorem 5.1.

Note that the leaf-labeled tree $(\mathfrak{S}_\alpha, T_{P,n})$ consists of a root together with infinitely many copies of each of the trees $(\mathfrak{S}_\beta, T_{\bar{Q},n,m})$ for $m \in \mathbb{N}$ together with infinitely many copies of the tree $(\mathfrak{S}_\beta, E_\beta)$. Also note that each of the subtrees $(\mathfrak{S}_\beta, T_{\bar{Q},n,m})$ is either isomorphic to $(\mathfrak{S}_\beta, A_\beta)$ or $(\mathfrak{S}_\beta, E_\beta)$.

If there is an m with $(n, m) \notin Q$, then $(\mathfrak{S}_\alpha, T_{P,n})$ consists of infinitely many copies of $(\mathfrak{S}_\beta, A_\beta)$ and infinitely many copies of $(\mathfrak{S}_\beta, E_\beta)$. Hence, we have $(\mathfrak{S}_\alpha, T_{P,n}) \cong (\mathfrak{S}_\alpha, E_\alpha)$. On the other hand, if $(n, m) \in Q$ for all $m \in \mathbb{N}$ then $(\mathfrak{S}_\alpha, T_{P,n}) \cong (\mathfrak{S}_\alpha, A_\alpha)$.

Finally, an index for the computable set $T_{P,n}$ can be computed from (Π, a, e) and n .

Case 2. $\alpha = \beta + 1$ for a limit ordinal $\beta < \omega^\omega$. Hence, from a given $k \in \mathbb{N}_{>0} \cup \{\infty\}$ one can compute and index of a computable subset $L_k^\alpha \subseteq \text{leaves}(\mathfrak{S}_\beta)$. We define the sets $A_\alpha, E_\alpha \subseteq \text{leaves}(\mathfrak{S}_\alpha) = \{mu \mid m \in \mathbb{N}_{>0}, u \in \text{leaves}(\mathfrak{S}_\beta)\}$ as follows (cf. Figure 3):

$$A_\alpha = \{[2^{k-1}(2x-1)]u \mid k, x \in \mathbb{N}_{>0}, u \in L_k^\beta\},$$

$$E_\alpha = \{[2^k(2x-1)]u \mid k, x \in \mathbb{N}_{>0}, u \in L_k^\beta\} \cup \{[2x-1]u \mid x \in \mathbb{N}_{>0}, u \in L_\infty^\beta\}.$$

Again, indices for these computable sets can be computed from α . Since $(\mathfrak{S}_\beta, L_k^\beta) \not\cong (\mathfrak{S}_\beta, L_\infty^\beta)$ for all $k \in \mathbb{N}_{>0}$, we have $(\mathfrak{S}_\alpha, E_\alpha) \not\cong (\mathfrak{S}_\alpha, A_\alpha)$.

As in Case 1, the index e from (Π, a, e) is a Σ_β^0 index for a set $Q \subseteq \mathbb{N} \times \mathbb{N}$ such that $n \in P$ if and only if $(n, m) \in Q$ for all $m \in \mathbb{N}$. From e one can compute a Π_β^0 index f for the complement of Q . By induction, one can compute from f and $(n, m) \in \mathbb{N} \times \mathbb{N}$ an index

of a computable set $T_{\overline{Q},n,m} \subseteq \text{leaves}(\mathfrak{S}_\beta)$ such that

$$(\mathfrak{S}_\beta, T_{\overline{Q},n,m}) \cong \begin{cases} (\mathfrak{S}_\beta, L_\infty^\beta) & \text{if } (n, m) \notin Q, \\ (\mathfrak{S}_\beta, L_k^\beta) \text{ for some } k \in \mathbb{N}_{>0} & \text{if } (n, m) \in Q. \end{cases}$$

We define the computable set $T_{P,n} \subseteq \text{leaves}(\mathfrak{S}_\alpha)$ as the set

$$\{[2^{2m}(2x+1)]u \mid m, x \in \mathbb{N}, u \in T_{\overline{Q},n,m}\} \cup \{[2^{2k-1}(2x+1)]u \mid k \in \mathbb{N}_{>0}, x \in \mathbb{N}, u \in L_k^\beta\}.$$

The leaf-labeled tree $(\mathfrak{S}_\alpha, T_{P,n})$ consists of a root together with infinitely many copies of each of the trees $(\mathfrak{S}_\beta, T_{\overline{Q},n,m})$ for $m \in \mathbb{N}$ together with infinitely many copies of each of the trees $(\mathfrak{S}_\beta, L_k^\beta)$ for $k \in \mathbb{N}_{>0}$. Also note that each of the subtrees $(\mathfrak{S}_\beta, T_{\overline{Q},n,m})$ is isomorphic to one of the trees $(\mathfrak{S}_\beta, L_k^\beta)$ ($k \in \mathbb{N}_{>0} \cup \{\infty\}$).

If there is an m with $(n, m) \notin Q$, then $(\mathfrak{S}_\alpha, T_{P,n})$ consists of infinitely many copies of each of the trees $(\mathfrak{S}_\beta, L_k^\beta)$ for all $k \in \mathbb{N}_{>0} \cup \{\infty\}$. Hence, we have $(\mathfrak{S}_\alpha, T_{P,n}) \cong (\mathfrak{S}_\alpha, E_\alpha)$. On the other hand, if $(n, m) \in Q$ for all $m \in \mathbb{N}$ then $(\mathfrak{S}_\alpha, T_{P,n}) \cong (\mathfrak{S}_\alpha, A_\alpha)$.

Case 3. α is a limit ordinal. Let $(\alpha_k)_{k \geq 1}$ be the fundamental sequence of α . Recall that α_k is a successor ordinal. We define a computable subset L_∞^α of the leaves $\text{leaves}(\mathfrak{S}_\alpha) = \{ku \mid k \in \mathbb{N}_{>0}, u \in \text{leaves}(\mathfrak{S}_{\alpha_k})\}$ by

$$L_\infty^\alpha = \{ku \mid k \in \mathbb{N}_{>0}, u \in A_{\alpha_k}\}.$$

For $k \in \mathbb{N}_{>0}$ we define the computable set $L_k^\alpha \subseteq \text{leaves}(\mathfrak{S}_\alpha)$ by

$$L_k^\alpha = \{xu \mid 0 < x < k, u \in A_{\alpha_x}\} \cup \{xu \mid x \geq k, u \in E_{\alpha_x}\}.$$

Recall that $(\mathfrak{S}_{\alpha_k}, A_{\alpha_k}) \not\cong (\mathfrak{S}_{\alpha_k}, E_{\alpha_k})$ for all $k \geq 1$, whence $(\mathfrak{S}_\alpha, L_k^\alpha) \not\cong (\mathfrak{S}_\alpha, L_\infty^\alpha)$ for all $k \geq 1$.

For the Π_α^0 index (Π, a, e) we can assume (using a padding argument) that Φ_e is a total computable function such that for all $k \geq 1$, $\Phi_e(k)$ is a $\Pi_{\alpha_k}^0$ index for some set $Q_k \subseteq \mathbb{N}$ and $P = \bigcap_{k \geq 1} Q_k$. Let us define for $k \in \mathbb{N}_{>0}$ the set $P_k = \bigcap_{i=1}^k Q_k$. Since $\Pi_{\alpha_k}^0$ sets are closed under finite intersections, P_k is a $\Pi_{\alpha_k}^0$ set. Moreover, a $\Pi_{\alpha_k}^0$ index for P_k can be computed from k . By induction, from $k \in \mathbb{N}_{>0}$ and n one can compute an index for a computable set $T_{P_k,n} \subseteq \text{leaves}(\mathfrak{S}_{\alpha_k})$ such that

$$(\mathfrak{S}_{\alpha_k}, T_{P_k,n}) \cong \begin{cases} (\mathfrak{S}_{\alpha_k}, A_{\alpha_k}) & \text{if } n \in P_k, \\ (\mathfrak{S}_{\alpha_k}, E_{\alpha_k}) & \text{if } n \notin P_k. \end{cases}$$

Define the computable set

$$T_{P,n} = \{ku \mid k \in \mathbb{N}_{>0}, u \in T_{P_k,n}\}.$$

If $n \in P$, then $n \in P_k$ for all $k \in \mathbb{N}_{>0}$ and we get $(\mathfrak{S}_\alpha, T_{P,n}) \cong (\mathfrak{S}_\alpha, L_\infty^\alpha)$. On the other hand, if $n \notin P$ then there is some $k \in \mathbb{N}_{>0}$ such that $n \in P_i$ if and only if $i < k$. In this case we get $(\mathfrak{S}_\alpha, T_{P,n}) \cong (\mathfrak{S}_\alpha, L_k^\alpha)$. \square

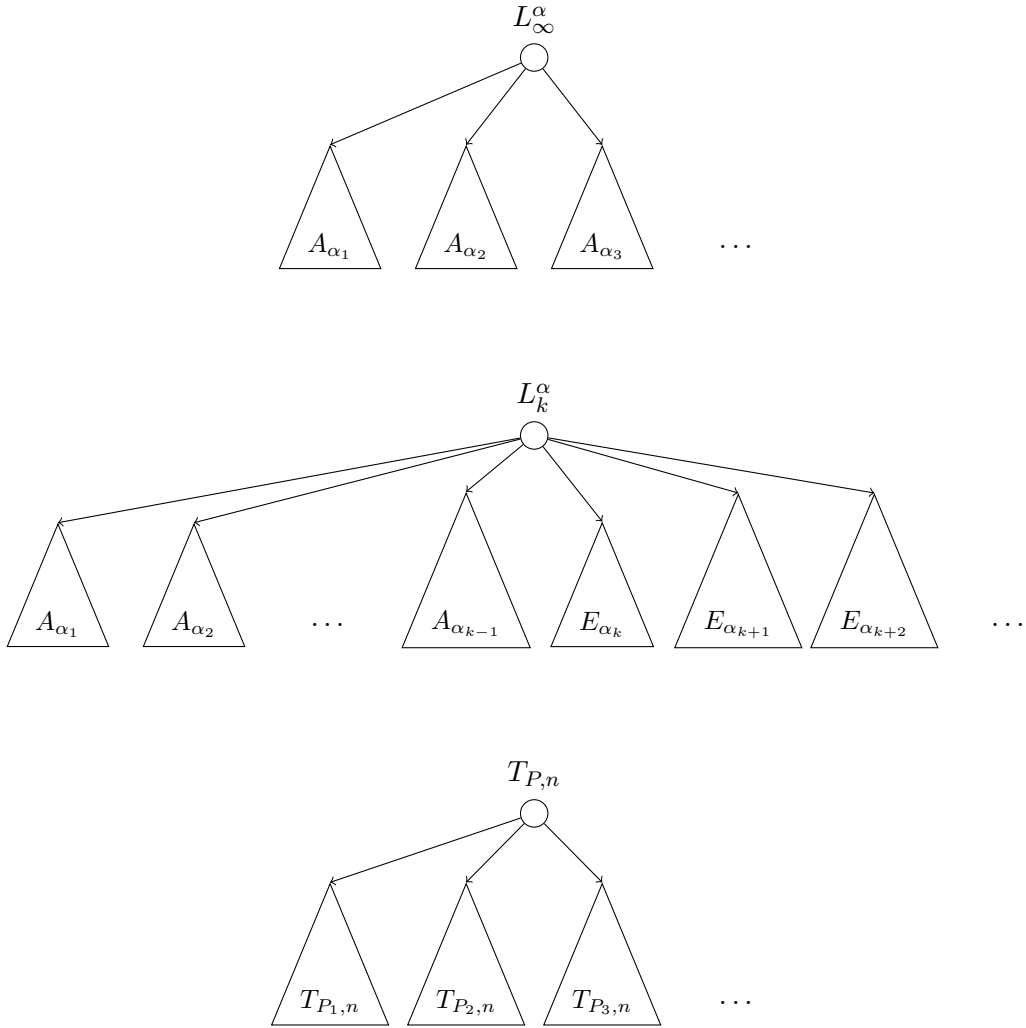


Figure 4: Trees for Case 3 from the proof of Theorem 5.1.

5.2. Tree-automaticity of the trees \mathfrak{S}_{ω^i} . In this section, we show that all trees \mathfrak{S}_{ω^i} ($i \geq 1$) from the previous section are tree-automatic. For this, we need the following lemma.

Lemma 5.2. *Let $i \geq 1$, $n \geq 1$, $\alpha \leq \omega^i$. Then we have*

$$\mathfrak{S}_\alpha \cup \{uv \mid u \in \text{leaves}(\mathfrak{S}_\alpha), v \in \mathfrak{S}_{\omega^i \cdot n}\} = \mathfrak{S}_{\omega^i \cdot n + \alpha}. \quad (5.2)$$

Note that the tree on the left-hand side of (5.2) is the tree that results from the tree \mathfrak{S}_α by replacing every leaf by a copy of the tree $\mathfrak{S}_{\omega^i \cdot n}$.

Proof. We prove the lemma by induction on $\alpha \leq \omega^i$. The case $\alpha = 0$ is clear. Next, assume that $\alpha = \gamma + 1$ is a successor ordinal. Then

$$\mathfrak{S}_\alpha = \{\varepsilon\} \cup \{nu \mid n \in \mathbb{N}_{>0}, u \in \mathfrak{S}_\gamma\}. \quad (5.3)$$

By induction hypothesis, we have

$$\mathfrak{S}_\gamma \cup \{uv \mid u \in \text{leaves}(\mathfrak{S}_\gamma), v \in \mathfrak{S}_{\omega^i \cdot n}\} = \mathfrak{S}_{\omega^i \cdot n + \gamma}. \quad (5.4)$$

Hence, we get

$$\begin{aligned} & \mathfrak{S}_\alpha \cup \{uv \mid u \in \text{leaves}(\mathfrak{S}_\alpha), v \in \mathfrak{S}_{\omega^i \cdot n}\} \\ & \stackrel{(5.3)}{=} \{\varepsilon\} \cup \{nu \mid n \in \mathbb{N}_{>0}, u \in \mathfrak{S}_\gamma\} \cup \{nu'v \mid n \in \mathbb{N}_{>0}, u' \in \text{leaves}(\mathfrak{S}_\gamma), v \in \mathfrak{S}_{\omega^i \cdot n}\} \\ & \stackrel{(5.4)}{=} \{\varepsilon\} \cup \{nw \mid n \in \mathbb{N}_{>0}, w \in \mathfrak{S}_{\omega^i \cdot n + \gamma}\} = \mathfrak{S}_{\omega^i \cdot n + \gamma + 1} = \mathfrak{S}_{\omega^i \cdot n + \alpha}. \end{aligned}$$

Finally, assume that $\alpha \leq \omega^i$ is a limit ordinal with the fundamental sequence $(\alpha_k)_{k \geq 1}$. Then $(\omega^i \cdot n + \alpha_k)_{k \geq 1}$ is our fundamental sequence for the ordinal $\omega^i \cdot n + \alpha$. We have

$$\mathfrak{S}_\alpha = \{\varepsilon\} \cup \{ku \mid k \in \mathbb{N}_{>0}, u \in \mathfrak{S}_{\alpha_k}\}. \quad (5.5)$$

By induction hypothesis, for every $k \geq 1$ we have

$$\mathfrak{S}_{\alpha_k} \cup \{uv \mid u \in \text{leaves}(\mathfrak{S}_{\alpha_k}), v \in \mathfrak{S}_{\omega^i \cdot n}\} = \mathfrak{S}_{\omega^i \cdot n + \alpha_k}. \quad (5.6)$$

Hence, we get

$$\begin{aligned} & \mathfrak{S}_\alpha \cup \{uv \mid u \in \text{leaves}(\mathfrak{S}_\alpha), v \in \mathfrak{S}_{\omega^i \cdot n}\} \\ & \stackrel{(5.5)}{=} \{\varepsilon\} \cup \{ku \mid k \in \mathbb{N}_{>0}, u \in \mathfrak{S}_{\alpha_k}\} \cup \{ku'v \mid k \in \mathbb{N}_{>0}, u' \in \text{leaves}(\mathfrak{S}_{\alpha_k}), v \in \mathfrak{S}_{\omega^i \cdot n}\} \\ & \stackrel{(5.6)}{=} \{\varepsilon\} \cup \{kw \mid k \in \mathbb{N}_{>0}, w \in \mathfrak{S}_{\omega^i \cdot n + \alpha_k}\} = \mathfrak{S}_{\omega^i \cdot n + \alpha}. \end{aligned}$$

□

Now we can prove tree-automaticity of \mathfrak{S}_{ω^i} .

Lemma 5.3. *For every $i \geq 1$, the tree $\mathfrak{S}_{\omega^i} \subseteq \mathbb{N}_{>0}^*$ is tree-automatic. Moreover, there is a unary tree-automatic copy (L, \leq) of \mathfrak{S}_{ω^i} together with a computable isomorphism $f : \mathfrak{S}_{\omega^i} \rightarrow (L, \leq)$.*

Proof. We prove the lemma by induction on i . Assume that we have already constructed a tree-automatic copy (L, \leq) of the tree \mathfrak{S}_{ω^i} over a unary alphabet (i.e., $L \subseteq \mathcal{T}_2^{\text{fin}}$ is regular) together with the computable isomorphism f . In addition, we assume that the root of (L, \leq) is the one-node tree $\{\varepsilon\}$; this property is preserved by the construction. We aim at constructing a unary tree-automatic copy of $\mathfrak{S}_{\omega^{i+1}}$ with root $\{\varepsilon\}$. Let us first construct a tree-automatic copy of the computable forest $\bigsqcup_{n \geq 1} \mathfrak{S}_{\omega^i \cdot n}$. This forest is isomorphic to the inverse prefix relation on the domain

$$\{nu \mid n \geq 1, u \in \mathfrak{S}_{\omega^i \cdot n}\} \subseteq \mathbb{N}_{>0}^*.$$

Define well-founded trees \mathfrak{T}_n for $n \geq 1$ inductively as follows. Let $\mathfrak{T}_1 = \mathfrak{S}_{\omega^i}$ and let \mathfrak{T}_{n+1} result from the tree \mathfrak{S}_{ω^i} by replacing every leaf by a copy of the tree \mathfrak{T}_n . Formally, we define

$$\mathfrak{T}_n = \{u_1 \cdots u_j v \mid 0 \leq j < n, u_1, \dots, u_j \in \text{leaves}(\mathfrak{S}_{\omega^i}), v \in \mathfrak{S}_{\omega^i}\}.$$

Lemma 5.2 implies $\mathfrak{T}_n = \mathfrak{S}_{\omega^i \cdot n}$ for $n \geq 1$. We construct a tree-automatic copy of $\bigsqcup_{n \geq 1} \mathfrak{T}_n$ using (L, \leq) . The universe of this copy is the set L' of all trees of the form $\text{pref}(\bigcup_{i=1}^n 0^i 1 t_i)$, where $n \geq 1$, $t_1, \dots, t_n \in L$ and there exists $1 \leq j \leq n$ such that t_j is a leaf of (L, \leq) for all $j < i$ and $t_j = \{\varepsilon\}$ for all $j > i$. Since the set of leaves of (L, \leq) is regular, the set L' is clearly regular. We define a tree-automatic partial order \leq' on the set L' by comparing the t_i componentwise. Let $t = \text{pref}(\bigcup_{i=1}^m 0^i 1 t_i) \in L'$ and $t' = \text{pref}(\bigcup_{i=1}^n 0^i 1 t'_i) \in L'$. Then $t \leq' t'$

if and only if $n = m$ and $t_i \leq t'_i$ for all $1 \leq i \leq n$. From this construction, it follows easily that

$$(L', \leq') \cong \bigsqcup_{n \geq 1} \mathfrak{T}_n = \bigsqcup_{n \geq 1} \mathfrak{S}_{\omega^i \cdot n}.$$

The set of roots of this forest is $\{\text{pref}(\bigcup_{i=1}^n 0^i 1) \mid n \geq 1\}$. Let us also define a computable isomorphism $f' : \bigsqcup_{n \geq 1} \mathfrak{T}_n \rightarrow (L', \leq')$. Take an element nw from $\bigsqcup_{n \geq 1} \mathfrak{T}_n$, where $n \in \mathbb{N}_{>0}$ and $w \in \mathfrak{T}_n$. There is a unique factorization $w = u_1 \cdots u_j v$ with $0 \leq j < n$, $u_1, \dots, u_j \in \text{leaves}(\mathfrak{S}_{\omega^i})$, and $v \in \mathfrak{S}_{\omega^i}$. Since the set $\text{leaves}(\mathfrak{S}_{\omega^i})$ is computable, we can compute this factorization. Next, using the computable isomorphism $f : \mathfrak{S}_{\omega^i} \rightarrow (L, \leq)$ define the trees $t_k = f(u_k)$ ($1 \leq k \leq j$), $t_{j+1} = f(v)$, and $t_{j+2}, \dots, t_n = \{\varepsilon\}$. Then set $f'(nw) = \text{pref}(\bigcup_{i=1}^n 0^i 1 t_i)$. It is straightforward to verify that this defines indeed an isomorphism from $\bigsqcup_{n \geq 1} \mathfrak{T}_n$ to (L', \leq') .

We derive from (L', \leq') a tree-automatic copy of the computable forest $\bigsqcup_{n \geq 1} \mathfrak{S}_{\omega^i \cdot n+1}$. This forest is isomorphic to the inverse prefix relation on

$$\{nu \mid n \in \mathbb{N}_{>0}, u \in \mathfrak{S}_{\omega^i \cdot n+1}\} = \mathbb{N}_{>0} \cup \{nmw \mid n, m \in \mathbb{N}_{>0}, w \in \mathfrak{S}_{\omega^i \cdot n}\}.$$

Note that in every node $u \in \mathcal{T}_2^{\text{fin}}$ of the forest (L', \leq') , the root of u has no right child (i.e., $1 \notin u$). Define the regular set of trees

$$L'' = \{\text{pref}(0^n) \mid n \geq 1\} \cup \{\text{pref}(\{1^m\} \cup t) \mid t \in L', m \geq 1\}.$$

On the set L'' we define the partial order \leq'' as follows. For $u, v \in L''$, let $u \leq'' v$ if and only if either

- $u = \text{pref}(\{1^m\} \cup s)$, $v = \text{pref}(\{1^m\} \cup t)$ with $s, t \in L'$ and $s \leq' t$ or
- $u = \text{pref}(\{1^m\} \cup t)$, $v = \text{pref}(0^n)$ with $t \in L'$ of the form $\text{pref}(\bigcup_{i=1}^n 0^i 1 t_i)$.

This order relation is clearly tree-automatic. Moreover, the construction implies that

$$(L'', \leq'') \cong (\mathbb{N}_{>0} \cup \{nmw \mid n, m \in \mathbb{N}_{>0}, w \in \mathfrak{S}_{\omega^i \cdot n}\}, \succeq) = \bigsqcup_{n \geq 1} \mathfrak{S}_{\omega^i \cdot n+1}.$$

A computable isomorphism $f'' : \bigsqcup_{n \geq 1} \mathfrak{S}_{\omega^i \cdot n+1} \rightarrow (L'', \leq'')$ can be defined as follows. For a root $n \in \mathbb{N}_{>0}$ let $f''(n) = \text{pref}(0^n)$. A node nmw with $n, m \in \mathbb{N}_{>0}$ and $w \in \mathfrak{S}_{\omega^i \cdot n}$ (hence, $mw \in \mathfrak{S}_{\omega^i \cdot n+1}$) is mapped to $f''(nmw) = \text{pref}(\{1^m\} \cup t)$ with $f'(nw) = t$.

Finally, we add to the forest (L'', \leq'') the root $\{\varepsilon\}$; this gives us a tree-automatic copy of the tree $\mathfrak{S}_{\omega^{i+1}}$ with root $\{\varepsilon\}$. A computable isomorphism is obtained by extending f'' by $f''(\varepsilon) = \{\varepsilon\}$. \square

5.3. Encoding Σ_2^0 -sets of binary trees. Theorem 5.1 and Lemma 5.3 show that the isomorphism problem for the following class of computable structures is $\Pi_{\omega^i}^0$ -hard for every $i \in \mathbb{N}$: the class contains all structures of the form (V, \sqsubseteq, X) where (V, \sqsubseteq) is the unary tree-automatic copy of \mathfrak{S}_{ω^i} from Lemma 5.3 and X is a computable unary predicate, which is moreover a subset of $\text{leaves}(V, \sqsubseteq)$. By Lemma 2.12 we can moreover assume that $V \subseteq \mathcal{T}_{\text{bin}}$, i.e., V consists of unlabeled full binary trees.

Let us define the set

$$\mathcal{T}_{\text{lef}} = \{\{\varepsilon\} \cup 0u \mid u \in \mathcal{T}_{\text{bin}}\}.$$

Thus, \mathcal{T}_{lef} contains all trees $t \in \mathcal{T}_2^{\text{fin}}$, where the root of t has no right child ($1 \notin t$), the root has a left child ($0 \in t$), and the subtree rooted at 0 belongs to \mathcal{T}_{bin} , i.e., is a full binary tree.

In this section, we describe an encoding of Σ_2^0 -subsets of \mathcal{T}_{lef} by sets of tree-automatic trees of height 3. Actually, we need this encoding only for computable subsets of \mathcal{T}_{lef} (instead of Σ_2^0 -sets), but the proof of Lemma 5.4 is not simpler for a computable set B .

Lemma 5.4. *There are two trees \mathfrak{U}_0 and \mathfrak{U}_1 of height 3 ($\mathfrak{U}_0 \not\cong \mathfrak{U}_1$) with the following property: From a given index of a Σ_2^0 -set $B \subseteq \mathcal{T}_{\text{lef}}$ one can effectively construct a tree-automatic forest \mathfrak{F}_B of height 3 such that:*

- The set of roots of \mathfrak{F}_B is \mathcal{T}_{lef} .
- For every $t \in \mathcal{T}_{\text{lef}}$, $\mathfrak{F}_B(t) \cong \mathfrak{U}_0$ if $t \in B$ and $\mathfrak{F}_B(t) \cong \mathfrak{U}_1$ if $t \notin B$.

Restricting to trees from \mathcal{T}_{lef} makes our encoding technically a bit simpler, and this restriction can be easily enforced later when we apply Lemma 5.4.

We prove Lemma 5.4 using a similar statement for words from [20]. First, we have to introduce a notation from [20]. Let $\mathcal{A} = (Q, \Sigma, \Delta, I, F)$ be a finite nondeterministic *string* automaton, where Q is the set of states, Σ is the input alphabet, $\Delta \subseteq Q \times \Sigma \times Q$ is the set of transition triples, I is the set of initial states, and F is the set of final states. A *successful run* of \mathcal{A} on a non-empty word $w \in \Sigma^+$ is a word $(q_0, a_1, q_1)(q_1, a_2, q_2) \cdots (q_{n-1}, a_n, q_n)$ over Δ such that $q_0 \in I$, $q_n \in F$, and $w = a_1 a_2 \cdots a_n$. The language $L(\mathcal{A})$ accepted by \mathcal{A} consists of all non-empty words (for technical reasons, the empty word was excluded in [20]) that have a successful run. We define a forest $\text{forest}(\mathcal{A})$ in the following. Clearly, $(L(\mathcal{A}), \preceq)$ (recall that \preceq is the prefix relation on words) is a (string-automatic) forest. The set of all leaves of the forest $(L(\mathcal{A}), \preceq)$ is first-order definable whence it is a regular language. It is the set of all words in $L(\mathcal{A})$ that are not a proper prefix of another word in $L(\mathcal{A})$. Let us denote the set of leaves of $(L(\mathcal{A}), \preceq)$ with $\text{leaves}(\mathcal{A})$. We define the forest $\text{forest}(\mathcal{A})$ as follows:

- The domain of $\text{forest}(\mathcal{A})$ is the regular set

$$L(\mathcal{A}) \sqcup \{r \in \Delta^+ \mid r \text{ is a successful run of } \mathcal{A} \text{ on some } v \in \text{leaves}(\mathcal{A})\}.$$

- The order relation \leq of $\text{forest}(\mathcal{A})$ is defined by $u \leq v$ if
 - (1) either $u, v \in L(\mathcal{A})$ and $v \preceq u$ or
 - (2) $v \in L(\mathcal{A})$ and $u \in \Delta^+$ is a successful run of \mathcal{A} on some $w \in \text{leaves}(\mathcal{A})$ with $v \preceq w$.

Clearly, $\text{forest}(\mathcal{A})$ is string-automatic. Intuitively, we take the forest resulting from the inverse prefix order on the regular language $L(\mathcal{A})$ and append to each leaf v of $(L(\mathcal{A}), \preceq)$ all successful runs of \mathcal{A} on v as children. All these children are leaves in $\text{forest}(\mathcal{A})$. In [20], the following lemma was proved.

Lemma 5.5. *There exist two trees \mathfrak{U}_0 and \mathfrak{U}_1 of height 3 ($\mathfrak{U}_0 \not\cong \mathfrak{U}_1$) with the following property: From a given index of a Σ_2^0 -set $A \subseteq \{0, 1\}^*1$ one can effectively construct a finite string automaton \mathcal{A} (over an alphabet Σ with $0, 1, \# \in \Sigma$) such that $\text{forest}(\mathcal{A})$ is a forest of height 3 with the following properties.*

- The set of roots of $\text{forest}(\mathcal{A})$ is $\{0, 1\}^*1\#$.
- For every $w \in \{0, 1\}^*1$, $\text{forest}(\mathcal{A})(w\#) \cong \mathfrak{U}_0$ if $w \in A$, and $\text{forest}(\mathcal{A})(w\#) \cong \mathfrak{U}_1$ if $w \notin A$.

In order to prove Lemma 5.4 using Lemma 5.5, we have to encode trees from \mathcal{T}_{lef} by words. A tree $t \in \mathcal{T}_{\text{bin}}$ can be encoded by a non-empty bracket expression, i.e., a word over the alphabet $\{(\cdot), \cdot\}$. Here, we view such a bracket expression as a binary string by identifying $($ with 0 and $)$ with 1. Thus, we define a nonempty word $\text{word}(t) \in \{0, 1\}^+$ as follows: Consider a depth-first left-to-right traversal of t . Each time, we move from a node v to one of its children, we write down 0. Each time, we move from a node vi to its parent

node v , we write down 1. The resulting word is $\text{word}(t)$. Formally, let $\text{word}(\{\varepsilon\}) = \varepsilon$ and for $t \in \mathcal{T}_{\text{bin}} \setminus \{\{\varepsilon\}\}$ such that $t = \{\varepsilon\} \cup 0t_1 \cup 1t_2$ let

$$\text{word}(t) = 0\text{word}(t_1)10\text{word}(t_2)1.$$

This mapping word is clearly injective. Finally, for $t = (\{\varepsilon\} \cup 0u) \in \mathcal{T}_{\text{lef}}$ with $u \in \mathcal{T}_{\text{bin}}$ let $\text{word}(t) = 0\text{word}(u)1$. Also the mapping $\text{word} : \mathcal{T}_{\text{lef}} \rightarrow 0\{0,1\}^*1$ is injective.

Let us now fix an alphabet Σ such that $0, 1, \# \in \Sigma$ (as in Lemma 5.5). Take a word $w = u\#v$ with $u \in \{0,1\}^*$, $v \in \Sigma^*$ and $u = \text{word}(t)$ for some tree $t \in \mathcal{T}_{\text{lef}}$. Note that the root of t has no right child in t . For $1 \leq i \leq |v|$ let a_i be the i^{th} symbol of v . We encode w by the Σ -labeled tree $\text{tree}(w) = (T, \lambda) \in \mathcal{T}_{2,\Sigma}^{\text{fin}}$, where

$$T = t \cup \{1^i \mid 1 \leq i \leq |v|\} \text{ and } \lambda(x) = \begin{cases} a_i & \text{if } x = 1^i \text{ for some } 1 \leq i \leq |v|, \\ \# & \text{else.} \end{cases}$$

Note that $\text{tree}(\text{word}(t)\#) = t$, since we identify an unlabeled tree with a tree where all nodes are labeled with $\#$.

Lemma 5.6. *From a given string automaton \mathcal{A} over Σ one can construct effectively a tree automaton \mathcal{B} over Σ such that for every tree $t \in \mathcal{T}_{\text{lef}}$ and every word $v \in \Sigma^*$ the following holds: The number of successful runs of \mathcal{A} on the string $\text{word}(t)\#v$ equals the number of successful runs of \mathcal{B} on the tree $\text{tree}(\text{word}(t)\#v)$.*

Proof. Let $\mathcal{A} = (Q, \Sigma, \Delta, I, F)$. Note that $\text{tree}(\text{word}(t)\#v)$ consists of the tree t to which we add at the root a right branch of length $|v|$. The i^{th} node on this branch is labeled with the i^{th} symbol of $\#v$ (we count from the root to the leaf). Essentially, the tree automaton \mathcal{B} simulates a tree-walking automaton \mathcal{W} (see [5] for a survey on tree walking automata, but we do not need a formal definition of tree walking automata) that walks over the tree $\text{tree}(\text{word}(t)\#v)$ in depth-first left-to-right order. Thereby, \mathcal{W} simulates the string automaton \mathcal{A} . The automaton \mathcal{W} starts in the root of the tree. In a first phase (which is finished if \mathcal{W} returns to the root for the first time), \mathcal{W} behaves as follows: each time, \mathcal{W} moves down in the tree (towards the leaves), it simulates a 0-labeled transition of \mathcal{A} , and each time, \mathcal{W} moves up in the tree (towards the root), it simulates a 1-labeled transition of \mathcal{A} . After the first phase, the tree t is fully traversed and \mathcal{W} goes into the right branch of $\text{tree}(\text{word}(t)\#v)$ (which is labeled with the word $\#v$) and continues the simulation of \mathcal{A} .

Here is a formal definition of the tree automaton $\mathcal{B} = (Q', \Delta', I', F')$, which simulates a tree-walking automaton \mathcal{W} with the above behavior. Fix an arbitrary final state $q_f \in F$. The state set of \mathcal{B} is

$$Q' = (Q \times Q \times Q) \cup (I \times Q) \cup Q,$$

the set of initial states is $I' = I \times Q$, and the set of final states is F . The set of transitions is $\Delta' = \Delta_1 \cup \Delta_2 \cup \Delta_3$, where

$$\begin{aligned} \Delta_1 &= \{((p_1, p_2), \#, (q_1, q_2, q_3), r) \mid p_1 \in I, (p_1, 0, q_1) \in \Delta, q_2 \in Q, (q_3, 1, p_2), (p_2, \#, r) \in \Delta\}, \\ \Delta_2 &= \{((p_1, p_2, p_3), \#, (q_1, q_2, q_3), (r_1, r_2, r_3)) \mid \\ &\quad (p_1, 0, q_1), q_2 \in Q, (q_3, 1, p_2), (p_2, 0, r_1), r_2 \in Q, (r_3, 1, p_3) \in \Delta\} \\ &\quad \cup \{((p, p, p), \#, q_f, q_f) \mid p \in Q\}, \\ \Delta_3 &= \{(p, a, q_f, q) \mid (p, a, q) \in \Delta\}. \end{aligned}$$

With the transitions in Δ_1 we split the simulation of the tree-walking automaton into its first and second phase, i.e., p_2 in Δ_1 is the state reached by the tree-walking automaton after traversing the tree t . The transitions in Δ_2 simulate the traversal of t , whereas the transitions in Δ_3 simulate the the string automaton \mathcal{A} on the right $\sharp v$ -labeled branch. \square

Let us now prove Lemma 5.4.

Proof of Lemma 5.4. Fix a Σ_2^0 -set $B \subseteq \mathcal{T}_{\text{lef}}$. Then the set $A = \text{word}(B) \subseteq 0\{0,1\}^*1$ belongs to Σ_2^0 as well (the range of the **word**-mapping is computable and on its range, the inverse of **word** is also computable). Therefore, we can apply Lemma 5.5 to the set $\text{word}(B)$. We obtain (effectively) a finite string automaton \mathcal{A} (over an alphabet Σ with $0, 1, \sharp \in \Sigma$) such that $\text{forest}(\mathcal{A})$ is a forest of height 3 with the following properties.

- The set of roots of $\text{forest}(\mathcal{A})$ is $\{0, 1\}^*1\sharp$.
- For every $w \in \{0, 1\}^*1$, $\text{forest}(\mathcal{A})(w\sharp) \cong \mathfrak{U}_0$ if $w \in A$, and $\text{forest}(\mathcal{A})(w\sharp) \cong \mathfrak{U}_1$ if $w \notin A$ for two nonisomorphic trees \mathfrak{U}_0 and \mathfrak{U}_1 .

To the string automaton \mathcal{A} we next apply Lemma 5.6. We obtain (effectively) a tree automaton \mathcal{B} over Σ such that for every tree $t \in \mathcal{T}_{\text{lef}}$ and every word $v \in \Sigma^*$ the following holds: The number of successful runs of \mathcal{A} on the string $\text{word}(t)\sharp v$ equals the number of successful runs of \mathcal{B} on the tree $\text{tree}(\text{word}(t)\sharp v)$. Since \mathcal{A} accepts every word from $\{0, 1\}^*1\sharp$ (this set is the set of roots of $\text{forest}(\mathcal{A})$), \mathcal{B} accepts every tree $t \in \mathcal{T}_{\text{lef}}$.

By taking the product with a deterministic tree automaton that accepts the set of trees

$$\{\text{tree}(\text{word}(t)\sharp v) \mid t \in \mathcal{T}_{\text{lef}}, v \in \Sigma^*\}$$

(which is regular), we can assume that

$$L(\mathcal{B}) \subseteq \{\text{tree}(\text{word}(t)\sharp v) \mid t \in \mathcal{T}_{\text{lef}}, v \in \Sigma^*\}.$$

For trees $t_1 = \text{tree}(\text{word}(t)\sharp v_1)$, $t_2 = \text{tree}(\text{word}(t)\sharp v_2)$ let us write $t_1 \sqsubseteq t_2$ if v_1 is a prefix of v_2 . Clearly, this is a tree-automatic relation. Let

$$\text{max}(\mathcal{B}) = \{t \in L(\mathcal{B}) \mid \text{there does not exist } t' \in L(\mathcal{B}) \text{ with } t \sqsubset t'\};$$

this set is regular as well. We can now construct a tree-automatic forest \mathfrak{F}_B of height 3 as follows. The set of nodes of \mathfrak{F}_B is

$$L(\mathcal{B}) \cup \bigcup_{t \in \text{max}(\mathcal{B})} \text{Run}(\mathcal{B}, t).$$

Since $\text{max}(\mathcal{B})$ is regular, this set is also regular. The order relation of the forest \mathfrak{F}_B is the tree-automatic relation

$$\sqsubseteq \cup \{(\rho, t) \mid t \in L(\mathcal{B}), \exists u \in \text{max}(\mathcal{B}) : t \sqsubseteq u, \rho \in \text{Run}(\mathcal{B}, u)\}.$$

The set of roots of \mathfrak{F}_B is (as required) \mathcal{T}_{lef} . Moreover, for every tree $t \in \mathcal{T}_{\text{lef}}$, the construction directly implies that $\mathfrak{F}_B(t) \cong \text{forest}(\mathcal{A})(\text{word}(t)\sharp)$. Hence, for every tree $t \in \mathcal{T}_{\text{lef}}$ we have $\mathfrak{F}_B(t) \cong \mathfrak{U}_0$ if and only if $\text{forest}(\mathcal{A})(\text{word}(t)\sharp) \cong \mathfrak{U}_0$ if and only if $\text{word}(t) \in A = \text{word}(B)$ if and only if $t \in B$, and similarly, $\mathfrak{F}_B(t) \cong \mathfrak{U}_1$ if and only if $t \notin B$. \square

5.4. Hardness for the isomorphism problem. Hardness of the isomorphism problem for well-founded tree-automatic trees is established through the following theorem.

Theorem 5.7. *From a given $i \in \mathbb{N}_{>0}$, one can compute a well-founded tree-automatic tree \mathfrak{V}_i such that the following holds: From a given $\Pi_{\omega^i}^0$ -set $P \subseteq \mathbb{N}$ (represented by a $\Pi_{\omega^i}^0$ index) and $n \in \mathbb{N}$ one can compute a well-founded tree-automatic tree $\mathfrak{W}_{P,n}$ such that $n \in P$ if and only if $\mathfrak{V}_i \cong \mathfrak{W}_{P,n}$.*

Proof. Fix $i \geq 1$, an arbitrary $\Pi_{\omega^i}^0$ -set $P \subseteq \mathbb{N}$, and $n \in \mathbb{N}$. According to Theorem 5.1 there exists effectively a computable subset $L_i = L_{\infty}^{\omega^i} \subseteq \text{leaves}(\mathfrak{S}_{\omega^i})$ such that the following holds. From an index for P and n one can compute an index for a computable subset $T_{P,n} \subseteq \text{leaves}(\mathfrak{S}_{\omega^i})$ such that $(\mathfrak{S}_{\omega^i}, L_i) \cong (\mathfrak{S}_{\omega^i}, T_{P,n})$ if and only if $n \in P$. By Lemma 5.3, the tree \mathfrak{S}_{ω^i} is tree-automatic and there exists a unary tree-automatic copy (S, \leq) of \mathfrak{S}_{ω^i} together with a computable isomorphism $f : \mathfrak{S}_{\omega^i} \rightarrow (S, \leq)$. Moreover, by Lemma 2.12 we can assume that $S \subseteq \mathcal{T}_{\text{bin}}$. Finally, by applying the computable bijection $t \mapsto \{\varepsilon\} \cup 0t$ from \mathcal{T}_{bin} to \mathcal{T}_{lef} , we can even assume that $S \subseteq \mathcal{T}_{\text{lef}}$. Since f is computable and bijective and L_i and $T_{P,n}$ are computable, the sets $L'_i = f(L_i) \subseteq S$ and $T'_{P,n} = f(T_{P,n}) \subseteq S$ are computable as well and in particular Σ_2^0 . We have $n \in P$ if and only if $(S, \leq, L'_i) \cong (S, \leq, T'_{P,n})$.

Next, using Lemma 5.4 we obtain two trees \mathfrak{U}_0 and \mathfrak{U}_1 of height 3 and from the indices of the computable sets L'_i and $T'_{P,n}$, we can compute two tree-automatic forests \mathfrak{G}_i and $\mathfrak{H}_{P,n}$ of height 3 such that the following holds.

- The set of roots of \mathfrak{G}_i (resp. $\mathfrak{H}_{P,n}$) is \mathcal{T}_{lef} .
- For every $t \in \mathcal{T}_{\text{lef}}$, $\mathfrak{G}_i(t) \cong \mathfrak{U}_0$ if $t \in L'_i$ and $\mathfrak{G}_i(t) \cong \mathfrak{U}_1$ otherwise.
- For every $t \in \mathcal{T}_{\text{lef}}$, $\mathfrak{H}_{P,n}(t) \cong \mathfrak{U}_0$ if $t \in T'_{P,n}$ and $\mathfrak{H}_{P,n}(t) \cong \mathfrak{U}_1$ otherwise.

Note that by Theorem 2.13, $\text{leaves}(S, \leq)$ is a regular set of trees since it is first-order definable in the tree-automatic tree (S, \leq) . Let \mathfrak{G}'_i (resp. $\mathfrak{H}'_{P,n}$) be the restriction of the forest \mathfrak{G}_i (resp. $\mathfrak{H}_{P,n}$) to those trees with a root from $\text{leaves}(S, \leq)$. It follows that \mathfrak{G}'_i and $\mathfrak{H}'_{P,n}$ are again tree-automatic forests of height 3. Moreover, we can assume that the intersection of the domains of \mathfrak{G}'_i (resp. $\mathfrak{H}'_{P,n}$) and S equals $\text{leaves}(S, \leq)$. Finally, let \mathfrak{V}_i (resp. $\mathfrak{W}_{P,n}$) be the well-founded tree-automatic tree obtained from the union of \mathfrak{G}'_i (resp. $\mathfrak{H}'_{P,n}$) and (S, \leq) . Hence, \mathfrak{V}_i (resp. $\mathfrak{W}_{P,n}$) results from the tree (S, \leq) by (i) replacing every leaf which belongs to L'_i (resp. $T'_{P,n}$) by the tree-automatic height-3 tree \mathfrak{U}_0 and by (ii) replacing every leaf which does not belong to L'_i (resp. $T'_{P,n}$) by the tree-automatic height-3 tree \mathfrak{U}_1 . Since $\mathfrak{U}_0 \not\cong \mathfrak{U}_1$, we have $n \in P$ if and only if $(S, \leq, L'_i) \cong (S, \leq, T'_{P,n})$ if and only if $\mathfrak{V}_i \cong \mathfrak{W}_{P,n}$. \square

Theorem 5.8. *The isomorphism problem for well-founded tree-automatic trees is $\Delta_{\omega^\omega}^0$ -hard under Turing-reductions.*

Proof. Let Φ_e be a total computable function that maps $i \in \mathbb{N}$ to an ordinal notation $a_i \in O$ with $|a_i|_O = \omega^i$. Hence $3 \cdot 5^e$ is an ordinal notation for ω^ω . Recall that $\Delta_{\omega^\omega}^0$ consists of all sets that are Turing-reducible to

$$H(3 \cdot 5^e) = \{\langle a, n \rangle \mid a <_O 3 \cdot 5^e, n \in H(a)\}.$$

This set is Turing-reducible to

$$A = \{\langle i, n \rangle \mid i \geq 1, n \in H(a_i)\}.$$

To see this, take a pair $\langle a, n \rangle$. First, check whether $a <_O 3 \cdot 5^e$. Since the set $\{b \in O \mid b <_O 3 \cdot 5^e\}$ is computably enumerable [1, Prop. 4.10], this is effectively possible using the

halting problem as an oracle. Clearly, the halting problem is computable in $H(a_1) = H(\omega)$ and hence in A . If $a <_O 3 \cdot 5^e$, we can compute effectively $i \in \mathbb{N}$ with $a <_O a_i$: Simply enumerate all sets $B_i = \{b \in O \mid b <_O a_i\}$ until a is found. Having $i \in \mathbb{N}$ with $a <_O a_i$ we can finally compute m such that $m \in H(a_i)$ if and only if $n \in H(a)$ (more precisely, from $a, a_i \in O$ one can compute an index for a many-one reduction of $H(a)$ to $H(a_i)$ [24, p. 437]).

Finally, we reduce the set A to the isomorphism problem for well-founded tree-automatic trees. Take a pair $\langle i, n \rangle$. From i, n (and a $\Pi_{\omega_i}^0$ index for the $\Delta_{\omega_i}^0$ -set $H(a_i)$, which can be computed from i) we can compute by Theorem 5.7 two well-founded tree-automatic trees \mathfrak{W} and \mathfrak{W} such that $n \in H(a_i)$ if and only if $\mathfrak{W} \cong \mathfrak{W}$. Hence, $\langle i, n \rangle \in A$ if and only if $\mathfrak{W} \cong \mathfrak{W}$. This proves the theorem. \square

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