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NEW AND IMPROVED SPANNING RATIOS FOR YAO GRAPHS

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ABSTRACT. For a set of points in the plane and a fixed integer $k > 0$, the Yao graph Y_k partitions the space around each point into k equiangular cones of angle $\theta = 2\pi/k$, and connects each point to a nearest neighbor in each cone. It is known for all Yao graphs, with the sole exception of Y_5 , whether or not they are geometric spanners. In this paper we close this gap by showing that for odd $k \geq 5$, the spanning ratio of Y_k is at most $1/(1 - 2\sin(3\theta/8))$, which gives the first constant upper bound for Y_5 , and is an improvement over the previous bound of $1/(1 - 2\sin(\theta/2))$ for odd $k \geq 7$.

We further reduce the upper bound on the spanning ratio for Y_5 from 10.9 to $2 + \sqrt{3} \approx 3.74$, which falls slightly below the lower bound of 3.79 established for the spanning ratio of Θ_5 (Θ -graphs differ from Yao graphs only in the way they select the closest neighbor in each cone). This is the first such separation between a Yao and Θ -graph with the same number of cones. We also give a lower bound of 2.87 on the spanning ratio of Y_5 .

In addition, we revisit the Y_6 graph, which plays a particularly important role as the transition between the graphs ($k > 6$) for which simple inductive proofs are known, and the graphs ($k \leq 6$) whose best spanning ratios have been established by complex arguments. Here we reduce the known spanning ratio of Y_6 from 17.6 to 5.8, getting closer to the spanning ratio of 2 established for Θ_6 .

Finally, we present the first lower bounds on the spanning ratio of Yao graphs with more than six cones, and a construction that shows that the Yao-Yao graph (a bounded-degree variant of the Yao graph) with five cones is not a spanner.

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1 Introduction

The complete Euclidean graph defined on a point set S in the plane is the graph with vertex set S and edges connecting each pair of points in S , where each edge xy has as weight the Euclidean distance $|xy|$ between its endpoints x and y . Although this graph is useful in many different contexts, its main disadvantage is that it has a quadratic number of edges. As such, much effort has gone into the development of various methods for constructing graphs that *approximate* the complete Euclidean graph. What does it mean to approximate this graph? One standard approach is to construct a spanning subgraph with fewer edges (typically linear) with the additional property that every edge e of the complete Euclidean graph is approximated by a path in the subgraph whose weight is not much more than the weight of e . This gives rise to the notion of a t -spanner. A t -spanner of the complete Euclidean graph is a spanning subgraph with the property that, for each pair of vertices x and y , the weight of a shortest path in the subgraph between x and y is at most $t \geq 1$ times $|xy|$. The *spanning ratio* is the smallest t for which the subgraph is a t -spanner. Spanners find many applications, such as approximating shortest paths or minimum spanning trees. For a comprehensive overview of geometric spanners and their applications, we refer the reader to the book by Narasimhan and Smid [15].

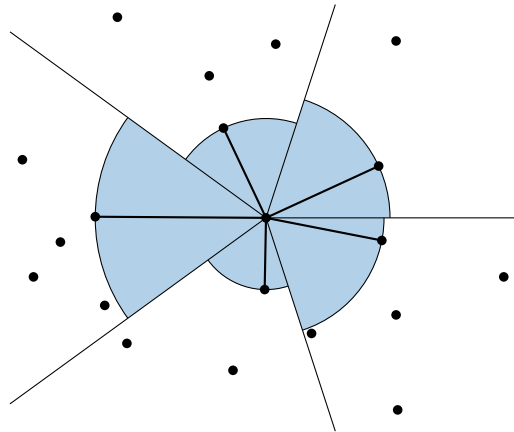


Figure 1: An example of the construction of the Yao graph with 5 cones.

One of the simplest ways of constructing a t -spanner is to first partition the plane around each vertex x into a fixed number of cones¹ and then add edges connecting x to a closest vertex in each cone (see Figure 1). Intuition suggests that this would yield a graph whose spanning ratio depends on the number of cones. Indeed, this is one of the first approximations of the complete Euclidean graph, referred to as *Yao graphs* in the literature, introduced independently by Flinchbaugh and Jones [13] and Yao [16]. We denote the Yao graph by Y_k where k is the number of cones, each having angle $\theta = 2\pi/k$. Yao used these graphs to simplify computation of the Euclidean minimum spanning tree. Flinchbaugh and Jones studied their graph-theoretic properties. Neither of them actually proved that they are t -spanners.

¹The orientation of the cones is the same for all vertices.

To the best of our knowledge, the first proof that Yao graphs are spanners was given by Althöfer *et al.* [1]. They showed that for every $t > 1$, there exists a k such that Y_k is a t -spanner. It appears that some form of this result was known earlier, as Clarkson [9] already remarked in 1987 that Y_{12} is a $1 + \sqrt{3}$ -spanner, albeit without providing a proof or a reference. Bose *et al.* [6] provided a more specific bound on the spanning ratio, by showing that for $k > 8$, Y_k is a spanner with spanning ratio at most $1/(\cos \theta - \sin \theta)$. This was later strengthened to show that for $k > 6$, Y_k is a $1/(1 - 2\sin(\theta/2))$ -spanner [4]. Damian and Raudonis [11] showed that Y_6 is a 17.64-spanner, and Bose *et al.* [5] showed that Y_4 is a 663-spanner. For $k < 4$, El Molla [12] showed that there is no constant t such that Y_k is a t -spanner. This leaves open only the question of whether Y_5 is a constant spanner.

In this paper we close this gap by showing that for odd $k \geq 5$, the spanning ratio of Y_k is at most $1/(1 - 2\sin(3\theta/8))$. This gives the first constant upper bound for Y_5 and implies that Y_k is a constant spanner if and only if $k \geq 4$. For odd $k \geq 7$, our result also improves on the previous bound of $1/(1 - 2\sin(\theta/2))$. A more careful analysis allows us to reduce the upper bound on the spanning ratio of Y_5 from 10.9 to $2 + \sqrt{3} \approx 3.74$. We also give a lower bound of 2.87 on the spanning ratio of Y_5 . This complements a recent result on the spanning ratio of Θ_5 , which differs from Y_5 only in the distance measure it uses to select the closest neighbor in each cone: instead of Euclidean distance, it projects each vertex on the bisector of the cone and selects the vertex with the closest projection. Bose *et al.* [7] showed that Θ_5 has a spanning ratio in the interval [3.79, 9.96]. Because our upper bound of 3.74 on the spanning ratio of Y_5 is slightly lower than their lower bound for Θ_5 , this result establishes the first separation between the spanning ratio of Yao and Θ -graphs. For all other $k \geq 4$, it is unclear which of Θ_k or Y_k has a better spanning ratio.

In addition, we revisit the Y_6 graph, which plays a particularly important role as the transition between the graphs ($k > 6$) for which simple inductive proofs are known, and the graphs ($k \leq 6$) whose best spanning ratios are established by complex arguments. Here we reduce the known spanning ratio of Y_6 from 17.64 to 5.8, thus moving toward the spanning ratio of 2 established for Θ_6 [3]. In contrast to Y_5 , we present a lower bound of 2 on the spanning ratio of Y_6 , showing that it can never improve upon Θ_6 in this regard.

Finally, we present the first lower bounds on the spanning ratio of Yao graphs with more than six cones, and a construction that shows that the Yao-Yao graph with five cones is not a spanner. Yao-Yao graphs are closely related to Yao graphs; a precise definition can be found in Section 4.5.

Before delving into these problems, we introduce a few definitions common to all sections of this paper. In particular, we start with a more precise definition of the Yao graph Y_k . For a fixed k , let $Q_i(a)$ be the half-open cone of angle $2\pi/k$ with apex a , including the angle range $[i, i+1) \cdot 2\pi/k$, for $i = 0, \dots, k-1$, where angles are measured counterclockwise from the positive x -axis. The directed graph \vec{Y}_k includes exactly one directed edge from a to a closest point in $Q_i(a)$, for each $i = 0, \dots, k-1$. If there are several equally-closest points within $Q_i(a)$, then ties are broken arbitrarily. The graph Y_k is the undirected version of \vec{Y}_k . We use D_a^b to denote the disk sector that is the intersection of the disk with center a and radius $|ab|$ and the cone with apex a containing point b . For any two points $a, b \in S$, we denote the length of a shortest path in Y_k from a to b by $p(a, b)$.

2 Spanning ratio of Y_k , for odd k

In this section we study the spanning properties of the Yao graphs Y_k defined on a plane point set S by an odd number of cones $k \geq 5$, each of angle $\theta = 2\pi/k$. For $k = 5$ in particular, this is the first result showing that Y_5 is a constant spanner. For odd values $k > 5$, we improve the currently known bound on the spanning ratio of Y_k .

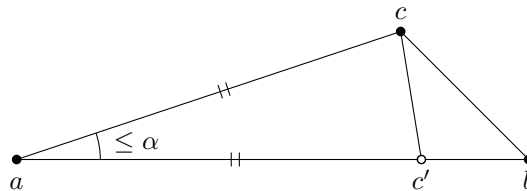


Figure 2: If α is small, there is a close relation between $|ac|$ and $|bc|$.

Lemma 1. *Given three points a , b , and c , such that $|ac| \leq |ab|$ and $\angle bac \leq \alpha < \pi$, then*

$$|bc| \leq |ab| - (1 - 2 \sin(\alpha/2)) |ac|.$$

Proof. Let c' be the point on ab such that $|ac| = |ac'|$ (see Figure 2). Since acc' forms an isosceles triangle,

$$|cc'| = 2 \sin(\angle bac/2) |ac| \leq 2 \sin(\alpha/2) |ac|.$$

Now, by the triangle inequality,

$$\begin{aligned} |bc| &\leq |cc'| + |c'b| \\ &\leq 2 \sin(\alpha/2) |ac| + |ab| - |ac'| \\ &= |ab| - (1 - 2 \sin(\alpha/2)) |ac|. \end{aligned} \quad \square$$

Theorem 2. *For any odd integer $k \geq 5$, the graph Y_k has spanning ratio at most $t = 1/(1 - 2 \sin(3\theta/8))$, where $\theta = 2\pi/k$.*

Proof. Let $a, b \in S$ be an arbitrary pair of points. We show that there is a path in Y_k from a to b no longer than $t|ab|$. For simplicity, let $Q(a)$ denote the cone with apex a that contains b , and let $Q(b)$ denote the cone with apex b that contains a . Rotate the point set S such that $Q(a)$ coincides with $Q_0(a)$, as depicted in Figure 3. We assume without loss of generality that b lies below the bisector of $Q(a)$; the case when b lies above this bisector is symmetric.

Let α be the angle formed by the segment ab with the bisector of $Q(a)$, and let β be the angle formed by ab with the bisector of $Q(b)$. Since k is odd, the bisector of $Q(a)$ is parallel to the bottom boundary of $Q(b)$. Hence, we have that $\alpha = \theta/2 - \beta$. Assume without loss of generality that α is the smaller of these two angles (if not, we exchange the roles of a and b). It follows that $\alpha \leq \theta/4$.

Our proof is by induction on the distance $|ab|$ (more formally, on the *rank* of $\{a, b\}$ among all pairs of points when ordered by distance). In the base case $|ab|$ is minimal, which

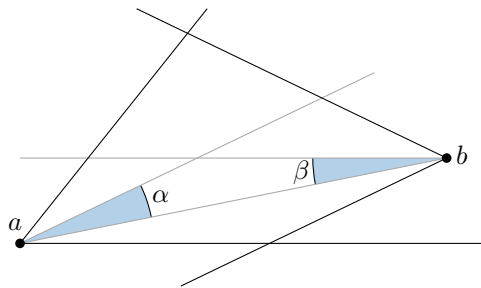


Figure 3: Since opposite cones are not symmetric, either α or β is small.

means that there is no point $c \in Q(a)$ that is strictly closer to a than b . Therefore either $ab \in Y_k$, in which case $p(a, b) = |ab|$ and our proof for the base case is finished, or there is a point $c \in Q(a)$ such that $|ab| = |ac|$ and $ac \in Y_k$. In this latter case, since $\alpha \leq \theta/4$ and $k \geq 5$, the angle between ab and ac is at most $\theta/2 + \alpha \leq 3\theta/4 \leq 3/4 \cdot (2\pi/5) = 3\pi/10$. This is less than $\pi/3$, which implies that $|bc| < |ab|$. This contradicts our assumption that $|ab|$ is minimal. It follows that $ab \in Y_k$ and the base case holds.

For the inductive step, let $c \in Q(a)$ be such that $\vec{ac} \in \vec{Y}_k$. If c coincides with b , then $p(a, b) = |ab|$ and the proof is finished. So assume that $c \neq b$. Because c is the closest vertex to a in this cone, and because $\angle cab \leq \theta/2 + \alpha \leq 3\theta/4$, we can apply Lemma 1 to derive $|cb| \leq |ab| - (1 - 2\sin(3\theta/8))|ac| = |ab| - |ac|/t$, which is strictly less than $|ab|$. Thus we can use the inductive hypothesis on cb to determine a path between a and b of length

$$p(a, b) \leq |ac| + t|cb| \leq |ac| + t \left(|ab| - \frac{|ac|}{t} \right) = t|ab|. \quad \square$$

Applying this result to Y_5 yields a spanning ratio of $1/(1 - 2\sin(3\pi/20)) \approx 10.868$. This is the first known upper bound on the spanning ratio of Y_5 and fully settles the question of which Yao graphs are spanners.

Corollary 3. *The graph Y_k is a spanner if and only if $k \geq 4$.*

Next we lower the upper bound on the spanning ratio of Y_5 by taking a closer look at all feasible configurations.

Theorem 4. *The graph Y_5 has spanning ratio at most $2 + \sqrt{3} \approx 3.74$.*

Here we also use induction on the pairwise distances between pairs of points in S . Consider the same configuration used in the proof of Theorem 2: $a \in Q(b)$ and $b \in Q(a)$ are points in S , and we seek a short path from a and b . Without loss of generality, we assume that $Q(a)$ coincides with $Q_0(a)$, and b lies below its bisector. If $|ab|$ is minimal (the base case) or $ab \in Y_5$, arguments similar to the ones used in the proof of Theorem 2 show that $p(a, b) = |ab|$ and our proof is finished. So let $c \in Q(a)$ and $d \in Q(b)$ be the vertices in S such that $\vec{ac} \in \vec{Y}_5$ and $\vec{bd} \in \vec{Y}_5$, and let $\phi = \angle cab$, and $\psi = \angle dba$ (see Figure 4a).

Now, instead of applying Lemma 1 for the maximum value of ϕ (as in the proof of Theorem 2), we apply Lemma 1 only for values $\phi \leq \bar{\theta}$ or $\psi \leq \bar{\theta}$, for some threshold angle $\bar{\theta}$

(to be determined later). These cases yield a spanning ratio of $t \geq 1/(1 - 2\sin(\bar{\theta}/2))$. We handle the remaining cases differently, so for the remainder of the proof, we assume that $\phi > \bar{\theta}$ and $\psi > \bar{\theta}$. We compute an exact value of $\bar{\theta}$ shortly, but for now we only need that $\theta/2 < \bar{\theta} < 3\theta/4$. This implies that neither c nor d can lie below ab , as this would make the corresponding angle smaller than $\theta/2$.

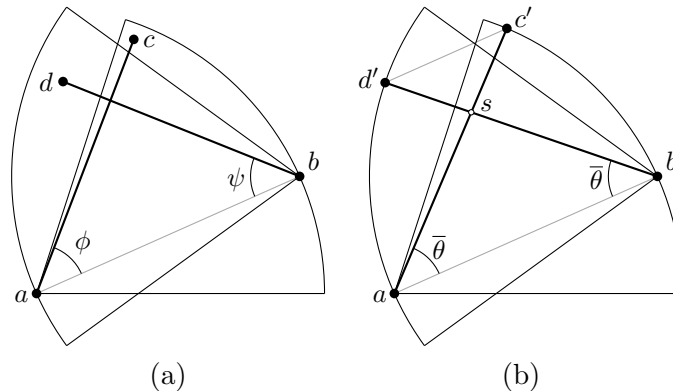


Figure 4: (a) Two vertices of Y_5 with their closest vertices. (b) The worst-case situation when ac and bd cross.

First consider the case where ac and bd intersect. In this case, instead of directly applying an inductive argument to either cb or da , we bound the distance cd and use induction to show that $|ac| + t|cd| + |db| \leq t|ab|$. To derive this bound, consider the point c' such that $\angle c'ab = \bar{\theta}$ and $|ac'| = |ab|$ and the analogously defined point d' (see Figure 4b). Let s be the intersection point between ac' and bd' . When ac and bd intersect, the distance $|cd|$ can be increased by rotating c towards b and d towards a . Since both ϕ and ψ must be larger than $\bar{\theta}$, the worst case occurs when $\phi = \psi = \bar{\theta}$, leaving c and d on the boundary of $\triangle c'd's$. As $c'd'$ is the longest side of this triangle, it follows that $|cd| < |c'd'|$. Using the fact that the triangles $\triangle c'd's$ and $\triangle abs$ are similar and isosceles, we can compute $|c'd'|$:

$$\begin{aligned}
 |c'd'| &= 2|c's| \cos \bar{\theta} \\
 &= 2(|ac'| - |as|) \cos \bar{\theta} \\
 &= 2 \left(|ab| - \frac{|ab|}{2 \cos \bar{\theta}} \right) \cos \bar{\theta} \\
 &= (2 \cos \bar{\theta} - 1)|ab|
 \end{aligned}$$

Recall that our aim is to use induction on cd to obtain a short path from a to b . We now compute the spanning ratio t required for the inequality $|ac| + t|cd| + |db| \leq t|ab|$ to hold. By the inequality above, we have that $|ac| + t|cd| + |db| \leq |ab| + t(2 \cos \bar{\theta} - 1)|ab| + |ab|$. This latter term is bounded above by $t|ab|$ for any $t \geq 1/(1 - \cos \bar{\theta})$.

So far we derived two constraints on t and $\bar{\theta}$: $t \geq 1/(1 - 2\sin(\bar{\theta}/2))$ and $t \geq 1/(1 - \cos \bar{\theta})$. Because $\sin \bar{\theta}$ is increasing and $\cos \bar{\theta}$ is decreasing for all values of $\bar{\theta}$ under

consideration, we minimize t by choosing $\bar{\theta}$ such that $1/(1 - 2\sin(\bar{\theta}/2)) = 1/(1 - \cos\bar{\theta})$. This yields $\bar{\theta} = \arccos(\sqrt{3} - 1) \approx 0.75$ and $t = 2 + \sqrt{3} \approx 3.74$.

Now consider what happens when one of ac or bd is “short”, under some notion of short captured by the following lemma.

Lemma 5. *Let $\triangle abc$ be a triangle with angle $\alpha = \angle cab$ and longest side ab . Let $\lambda > 1$ be a real constant. Then*

$$|ac| \leq \frac{2\lambda^2 \cos \alpha - 2\lambda}{\lambda^2 - 1} |ab| \quad \text{implies that} \quad |ac| + \lambda|bc| \leq \lambda|ab|.$$

Proof. First, note that the first inequality above implies $\lambda > 1/\cos \alpha$, as $|ac|$ would be non-positive otherwise. By the law of cosines, $|bc| = \sqrt{|ac|^2 + |ab|^2 - 2|ab||ac| \cos \alpha}$. By substituting this in the inequality $|ac| + \lambda|bc| \leq \lambda|ab|$, we see that it only holds if $|ac| \leq \frac{2\lambda^2 \cos \alpha - 2\lambda}{\lambda^2 - 1} |ab|$, as stated by the lemma. \square

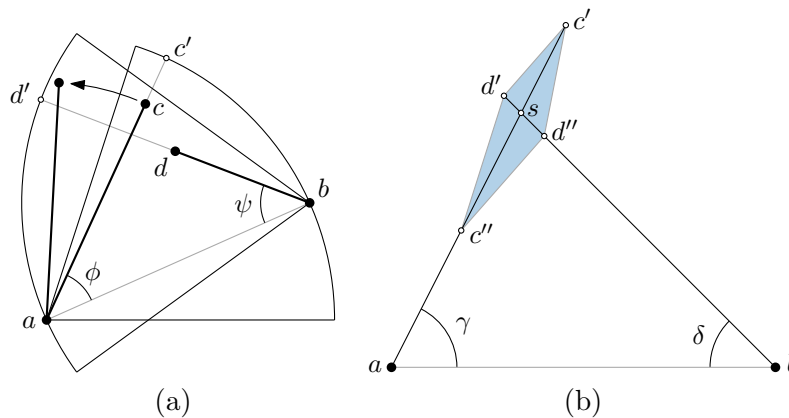


Figure 5: (a) The rotation to maximize $|cd|$ when ac and bd do not cross. (b) Illustration of Lemma 6.

The only case left to consider is when ac and bd are both long, but they do not intersect. In this case, we again seek to bound the distance $|cd|$. If we can show that $|cd| \leq (2\cos\bar{\theta} - 1)|ab|$, we can apply the same argument as for the intersecting case and we are done. Let c' be the point on the extension of ac with $|ac'| = |ab|$, and let d' be the analogous point on the extension of bd (see Figure 5a). If ac does not intersect bd' , we can rotate d away from c by increasing ψ . Similarly, if bd does not intersect ac' , we can rotate c away from d by increasing ϕ . Thus, the distance $|cd|$ is maximized when $\phi + \psi$ is maximal, which in our context happens when $\phi + \psi = 3\theta/2 = 3\pi/5$. Note that in most cases, rotating this far moves the corresponding vertex past the boundary of the cone. But since we are only trying to find an upper bound, this is not a problem.

Now let c'' be the point on the line through ac with $|ac''| = \frac{2t^2 \cos \phi - 2t}{t^2 - 1} |ab|$, and let d'' be the point on the line through bd with $|bd''| = \frac{2t^2 \cos \psi - 2t}{t^2 - 1} |ab|$. If c lies on ac'' , Lemma 5 tells us that $|ac| + t|bc| \leq t|ab|$, which is exactly what we need. The only difficulty is that

the location of c changed during the rotation. But since the rotation preserved $|ac|$ and only increased $|bc|$, the inequality must hold for the configuration before the rotation as well. The same argument applies for the case when d lies on bd'' . The situation where c and d lie on $c''c'$ and $d''d'$, respectively, is handled by the following lemma.

Lemma 6. *Let $a, b, c, d \in S$. Let $\gamma = \angle cab$ and $\delta = \angle dba$ such that $\gamma > \bar{\theta}$, $\delta > \bar{\theta}$, and $\gamma + \delta = 3\pi/5$. Let $t = 2 + \sqrt{3}$. If*

$$\frac{2t^2 \cos \gamma - 2t}{t^2 - 1} |ab| \leq |ac| \leq |ab| \quad \text{and} \quad \frac{2t^2 \cos \delta - 2t}{t^2 - 1} |ab| \leq |bd| \leq |ab|,$$

then $|cd| \leq (2 \cos \bar{\theta} - 1) |ab|$.

Proof. Assume without loss of generality that $|ab| = 1$ and $\gamma \geq \delta$, so that $3\pi/10 \leq \gamma \leq 3\pi/5 - \bar{\theta}$ and $\bar{\theta} \leq \delta \leq 3\pi/10$. We first define the relevant points, depicted in Figure 5b. Let c' be the point on the extension of ac with $|ac'| = |ab|$, and let d' be the analogous point on the extension of bd . Let s be the intersection of ac' and bd' . Let c'' be the point on the line through ac with $|ac''| = \frac{2t^2 \cos \gamma - 2t}{t^2 - 1} |ab|$, and let d'' be the point on the line through bd with $|bd''| = \frac{2t^2 \cos \delta - 2t}{t^2 - 1} |ab|$.

Since c and d are restricted to the segments $c'c''$ and $d'd''$, respectively, their maximum distance is one of $|c'd'|$, $|c''d'|$, $|c'd''|$, and $|c''d''|$. Instead of computing these distances directly as a function of γ , we use the triangle inequality to bound them in terms of $|c's|$, $|sc''|$, $|d's|$, and $|sd''|$. We first establish $|sc''| > |c's| > 0$, by examining the derivatives. Let $c_1 = 2t^2/(t^2 - 1)$ and $c_2 = 1/\sin(3\pi/5)$. We derive

$$\frac{d|ac''|}{d\gamma} = \frac{d\left(\frac{2t^2 \cos \gamma - 2t}{t^2 - 1}\right)}{d\gamma} = \frac{-2t^2 \sin \gamma}{t^2 - 1} = -c_1 \sin \gamma, \tag{1}$$

$$\frac{d|bd''|}{d\gamma} = \frac{d\left(\frac{2t^2 \cos \delta - 2t}{t^2 - 1}\right)}{d\gamma} = \frac{d\left(\frac{2t^2 \cos(3\pi/5 - \gamma) - 2t}{t^2 - 1}\right)}{d\gamma} = \frac{2t^2 \sin(3\pi/5 - \gamma)}{t^2 - 1} = c_1 \sin(3\pi/5 - \gamma), \tag{2}$$

$$\frac{d|as|}{d\gamma} = \frac{d\left(\frac{\sin \delta}{\sin(\gamma + \delta)}\right)}{d\gamma} = \frac{d\left(\frac{\sin(3\pi/5 - \gamma)}{\sin(3\pi/5)}\right)}{d\gamma} = \frac{-\cos(3\pi/5 - \gamma)}{\sin(3\pi/5)} = -c_2 \cos(3\pi/5 - \gamma), \tag{3}$$

$$\frac{d|bs|}{d\gamma} = \frac{d\left(\frac{\sin \gamma}{\sin(\gamma + \delta)}\right)}{d\gamma} = \frac{d\left(\frac{\sin \gamma}{\sin(3\pi/5)}\right)}{d\gamma} = \frac{\cos \gamma}{\sin(3\pi/5)} = c_2 \cos \gamma. \tag{4}$$

Let

$$x_1 = |as| - |ac''| = \frac{\sin \delta}{\sin(\gamma + \delta)} - \frac{2t^2 \cos \gamma - 2t}{t^2 - 1}, \tag{5}$$

$$x_2 = |ac'| - |as| = 1 - \frac{\sin \delta}{\sin(\gamma + \delta)}, \tag{6}$$

$$y_1 = |bs| - |bd''| = \frac{\sin \gamma}{\sin(\gamma + \delta)} - \frac{2t^2 \cos \delta - 2t}{t^2 - 1}, \tag{7}$$

$$y_2 = |bd'| - |bs| = 1 - \frac{\sin \gamma}{\sin(\gamma + \delta)}. \tag{8}$$

Note that the values of x_1 and y_1 could be negative if c'' or d'' lie past s . Substituting c_1 , c_2 , and (1) - (4) in the equalities above yields

$$\frac{dx_1}{d\gamma} = \frac{d(|as| - |ac''|)}{d\gamma} = -c_2 \cos(3\pi/5 - \gamma) + c_1 \sin \gamma, \quad (9)$$

$$\frac{dx_2}{d\gamma} = \frac{d(|ac'| - |as|)}{d\gamma} = \frac{d(1 - |as|)}{d\gamma} = c_2 \cos(3\pi/5 - \gamma), \quad (10)$$

$$\frac{dy_1}{d\gamma} = \frac{d(|bs| - |bd''|)}{d\gamma} = c_2 \cos \gamma - c_1 \sin(3\pi/5 - \gamma), \quad (11)$$

$$\frac{dy_2}{d\gamma} = \frac{d(|bd'| - |bs|)}{d\gamma} = \frac{d(1 - |bs|)}{d\gamma} = -c_2 \cos \gamma. \quad (12)$$

Recall that $c_1 = 2t^2/(t^2 - 1)$, $c_2 = 1/\sin(3\pi/5)$, and $3\pi/10 \leq \gamma \leq 3\pi/5 - \bar{\theta}$. We verify the following:

$$\begin{aligned} \frac{d^2x_1}{d\gamma^2} &= -c_2 \sin(3\pi/5 - \gamma) + c_1 \cos \gamma \\ &> -1.1 \sin(3\pi/10) + 2.1 \cos(3\pi/5 - \bar{\theta}) > 0, \\ \frac{d^2x_2}{d\gamma^2} &= c_2 \sin(3\pi/5 - \gamma) > 0, \\ \frac{d^2y_1}{d\gamma^2} &= -c_2 \sin \gamma + c_1 \cos(3\pi/5 - \gamma) \\ &> -1.1 \sin(3\pi/5 - \bar{\theta}) + 2.1 \cos(3\pi/10) > 0, \\ \frac{d^2y_2}{d\gamma^2} &= c_2 \sin \gamma > 0. \end{aligned}$$

Therefore, by substituting $\gamma = 3\pi/10$ or $\gamma = 3\pi/5 - \bar{\theta}$ as the lower- or upper-bound of γ into (9) - (12), we can verify the following ranges:

$$\begin{aligned} -c_2 \cos(3\pi/10) + c_1 \sin(3\pi/10) &\leq \frac{dx_1}{d\gamma} \leq -c_2 \cos \bar{\theta} + c_1 \sin(3\pi/5 - \bar{\theta}), \\ c_2 \cos(3\pi/10) &\leq \frac{dx_2}{d\gamma} \leq c_2 \cos \bar{\theta}, \\ c_2 \cos(3\pi/10) - c_1 \sin(3\pi/10) &\leq \frac{dy_1}{d\gamma} \leq c_2 \cos(3\pi/5 - \bar{\theta}) - c_1 \sin \bar{\theta}, \\ -c_2 \cos(3\pi/10) &\leq \frac{dy_2}{d\gamma} \leq -c_2 \cos(3\pi/5 - \bar{\theta}). \end{aligned} \quad (13)$$

Specifically, we can verify that

$$\frac{dx_1}{d\gamma} \geq \max \left(\frac{dx_2}{d\gamma}, \left| \frac{dy_1}{d\gamma} \right|, \left| \frac{dy_2}{d\gamma} \right| \right), \quad (14)$$

which implies

$$\frac{d(x_1 - x_2)}{d\gamma} = \frac{dx_1}{d\gamma} - \frac{dx_2}{d\gamma} > 0.$$

By simply plugging in $\gamma = 3\pi/10$ into (5) and (6), we verify that $(x_1 - x_2) > 0$ when $\gamma = 3\pi/10$ and hence $x_1 > x_2$ for all $\gamma \in [3\pi/10, 3\pi/5 - \bar{\theta}]$. Similarly, we have $x_2 > 0$ when $\gamma = 3\pi/10$, and hence by (13), $x_2 > 0$ for all $\gamma \in [3\pi/10, 3\pi/5 - \bar{\theta}]$. These together yield $x_1 > x_2 > 0$. By the triangle inequality,

$$\begin{aligned} |c''d''| &\leq |sc''| + |sd''| = |x_1| + |y_1| = x_1 + |y_1|, \\ |c''d'| &\leq |sc''| + |sd'| = |x_1| + |y_2| = x_1 + |y_2|, \\ |c'd''| &\leq |sc'| + |sd''| = |x_2| + |y_1| \leq x_1 + |y_1|, \\ |c'd'| &\leq |sc'| + |sd'| = |x_2| + |y_2| \leq x_1 + |y_2|. \end{aligned}$$

By (14),

$$\begin{aligned} \frac{d(x_1 + |y_1|)}{d\gamma} &\geq \frac{d(x_1)}{d\gamma} - \left| \frac{d(y_1)}{d\gamma} \right| \geq 0, \\ \frac{d(x_1 + |y_2|)}{d\gamma} &\geq \frac{d(x_1)}{d\gamma} - \left| \frac{d(y_2)}{d\gamma} \right| \geq 0. \end{aligned}$$

By substituting $\gamma = 3\pi/5 - \bar{\theta}$ into (5), (7), and (8), one can easily verify that $x_1 + |y_1| \leq 2 \cos \bar{\theta} - 1$ and $x_1 + |y_2| \leq 2 \cos \bar{\theta} - 1$ when γ is maximized. Therefore $\max(x_1 + |y_1|, x_1 + |y_2|) \leq 2 \cos \bar{\theta} - 1$ for all $\gamma \in [3\pi/10, 3\pi/5 - \bar{\theta}]$, and hence $|cd| \leq \max(|c''d''|, |c''d'|, |c'd''|, |c'd'|) \leq 2 \cos \bar{\theta} - 1$ as required. \square

This completes the proof for the upper bound. Next, we prove a lower bound on the spanning ratio.

Theorem 7. Y_5 has spanning ratio at least 2.87.

Proof. The inductive proof of the upper bound on the spanning ratio of Y_5 suggests a construction for a lower bound. It is based on recursively attaching the “lattice cross” shown in Figure 4b to pairs of non-adjacent points (e.g., pairs $\{a, d'\}$, $\{b, c'\}$, $\{c', d'\}$ in Figure 4b). This recursive construction results in a fractal-like shape, starting from the pair $\{a, b\}$ (see Figure 6). However, the growth of the fractal is limited by collisions of neighboring fractal branches that create shortcuts to the paths, as shown in the circled area of Figure 6. This construction yields a spanning ratio of 2.66.

We adjust the shape of the fractal to increase the spanning ratio. In Figure 7, we obtain a spanning ratio of more than 2.87 by equalizing the length of all shortest paths between a and b . The coordinates for the points in Figure 7 can be found in Appendix A. \square

3 Spanning ratio of Y_6

In this section we fix $k = 6$ and show that, for any pair of points $a, b \in S$, $p(a, b) \leq 5.8|ab|$. We also establish a lower bound of 2 for the spanning ratio of Y_6 . Our proof is inductive and it relies on two simple lemmas, which we introduce next.

Let $a, b \in S$ and let $\vec{ac} \in \vec{Y}_6$ be the edge from a within the cone that includes b . The next two lemmas will be relevant in the context where we seek to bound $p(a, b)$ by applying the induction hypothesis to $p(c, b)$. The basic geometry is illustrated in Figure 8.

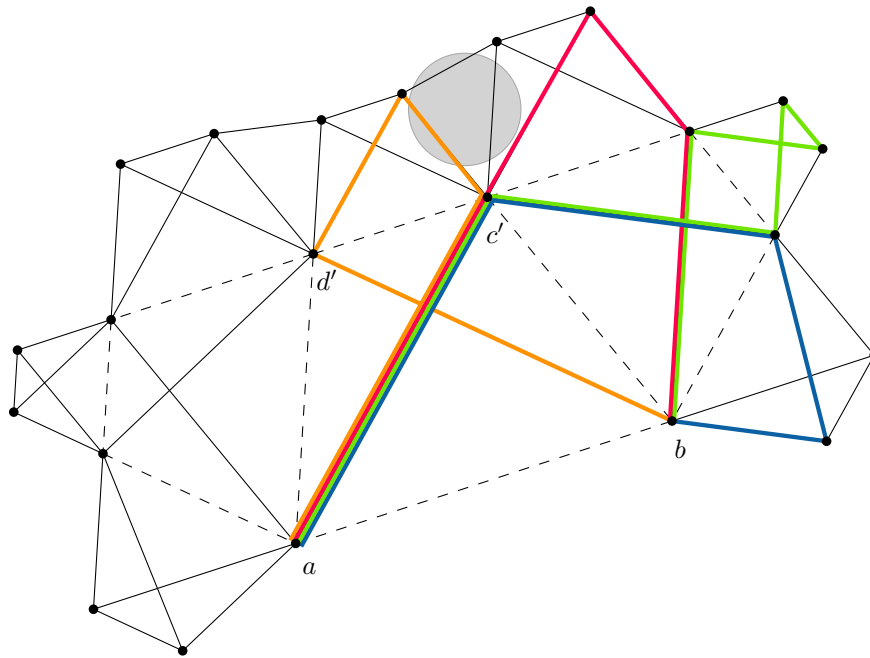


Figure 6: Spanning ratio 2.66. The fractal growth is limited by collision of branches in the circled area. The shortest paths between a and b are colored.

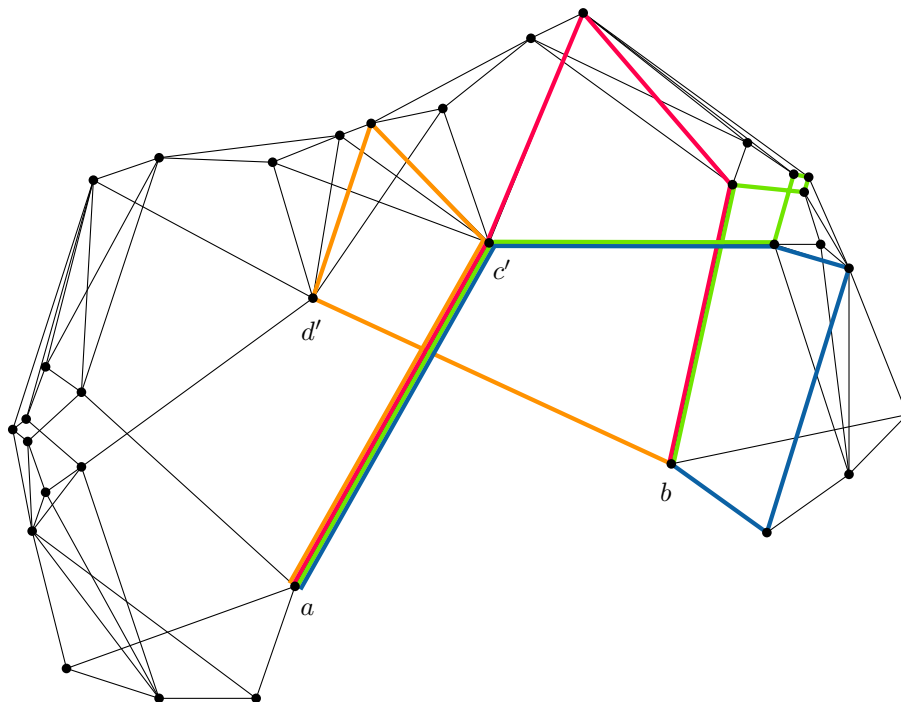


Figure 7: The spanning ratio is increased to 2.87 by adjusting the shape to equalize the lengths of the shortest paths between a and b .

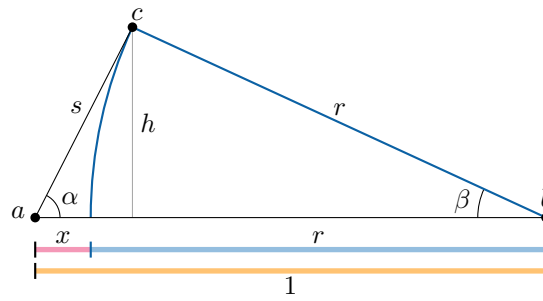


Figure 8: Notation for triangle $\triangle abc$. The dimensions have been normalized so that $|ab| = 1$.

Lemma 8 (Triangle). *Let $\triangle abc$ be labeled as in Figure 8, with $|ac| \leq |ab|$, $|bc| < |ab|$, $x = |ab| - |bc|$ and $s = |ac|$. The ratio s/x is equal to some function t that depends on α and β :*

$$\frac{s}{x} = t(\alpha, \beta) = \frac{\cos(\beta/2)}{\cos(\alpha + \beta/2)}. \quad (15)$$

Proof. Normalize the triangle so that $|ab| = 1$; this does not alter the quantity we seek to compute, s/x . Let $|bc| = r$ to simplify notation. Then $x = 1 - r$ and $x \geq 0$ because $r = |bc| \leq |ab| = 1$. Note that each of the angles $\angle cab$ and $\angle cba$ is strictly less than $\pi/2$, because $|ac| \leq |ab|$ and $|bc| \leq |ab|$. Thus the projection of c onto ab is interior to the segment ab . Computing the altitude h of $\triangle abc$ in two ways yields

$$s \sin \alpha = r \sin \beta.$$

Also projections onto ab yield

$$s \cos \alpha + r \cos \beta = 1.$$

Solving these two equations simultaneously yields expressions for r and s as functions of α and β :

$$r = \frac{\sin \alpha}{\sin \alpha \cos \beta + \cos \alpha \sin \beta}, \quad s = \frac{\sin \beta}{\sin \alpha \cos \beta + \cos \alpha \sin \beta}$$

Now we can compute $s/x = s/(1 - r)$ as a function of α and β . This simplifies to

$$\frac{s}{x} = \frac{\cos(\beta/2)}{\cos(\alpha + \beta/2)}$$

as claimed. \square

The following lemma derives an upper bound on the function $t(\alpha, \beta)$ from Lemma 8, which will be used in Theorem 10 to derive an optimal value for δ .

Lemma 9. *Let $a, b, c \in S$ satisfy the conditions of Lemma 8, and let $t(\alpha, \beta)$ be as defined in (15). Let $\delta \in (0, \pi/3)$ be a fixed positive angle. If $\alpha \leq \pi/3 - \delta$, or $\beta \leq \pi/3 - \delta$, then*

$$t(\alpha, \beta) \leq t(\pi/3, \pi/3 - \delta) = \frac{\cos(\pi/6 - \delta/2)}{\sin(\delta/2)}.$$

Proof. The derivative of $t(\alpha, \beta)$ with respect to α is

$$\frac{\partial t}{\partial \alpha} = \frac{\sin \alpha + \sin(\alpha + \beta)}{1 + \cos(2\alpha + \beta)} > 0.$$

This means that, for a fixed β value, $t(\alpha, \beta)$ reaches its maximum when α is maximum. Similarly, the derivative of $t(\alpha, \beta)$ with respect to β is

$$\frac{\partial t}{\partial \beta} = \frac{\sin \alpha}{2 \cos(\alpha + \beta/2)^2} > 0.$$

So for a fixed value α value, $t(\alpha, \beta)$ reaches its maximum when β is maximum. Because $|ac| \leq |ab|$, $\beta \leq \angle acb$. The sum of these two angles is $\pi - \alpha$, therefore $\beta \leq \pi/2 - \alpha/2$. This along with the derivations above implies that, for a fixed value $\alpha \leq \pi/3 - \delta$, $t(\alpha, \beta) \leq t(\alpha, \pi/2 - \alpha/2) \leq t(\pi/3 - \delta, \pi/3 + \delta/2)$ (we substituted $\alpha = \pi/3 - \delta$ in this latter inequality). Next we evaluate

$$\frac{t(\pi/3, \pi/3 - \delta)}{t(\pi/3 - \delta, \pi/3 + \delta/2)} = \frac{\cos(\pi/6 - \delta/2) \sin(3\delta/2)}{\cos(\pi/6 + \delta/2) \sin(\delta/2)} > 1.$$

It follows that $t(\pi/3, \pi/3 - \delta)$ is maximal. \square

We are now ready to prove the main result of this section.

Theorem 10. *The graph Y_6 has spanning ratio at most 5.8.*

This result follows from the following lemma, with the variable δ substituted by the quantity $\delta_0 = 0.324$ that minimizes $t(\delta)$. (It can be easily verified that $t(\delta) \geq t(0.324)$ and $t(0.324) < 5.8$.)

Lemma 11. *Let $\delta \in (0, \pi/9)$ be a strictly positive real value. The graph Y_6 has spanning ratio bounded above by*

$$t = t(\delta) = \max \left\{ \frac{\cos(\pi/6 - \delta/2)}{\sin(\delta/2)}, \frac{2}{1 - \frac{\sin(2\delta)}{\sin(\pi/6 + 2\delta)}} \right\}. \quad (16)$$

Proof. The proof is by induction on the pairwise distance between pairs of points $a, b \in S$. Without loss of generality let $b \in Q_0(a)$.

Base case We show that, if $|ab|$ is minimal, then $\vec{ab} \in \vec{Y}_6$ and so $p(a, b) = |ab|$. If $\vec{ab} \in \vec{Y}_6$, then the lemma holds. So assume that $\vec{ab} \notin \vec{Y}_6$; we will derive a contradiction. Because $\vec{ab} \notin \vec{Y}_6$, there must be another point $c \in Q_0(a)$ such that $\vec{ac} \in \vec{Y}_6$ and $|ac| = |ab|$. Let α_1 and α_2 be the angles that ab and ac make with the horizontal respectively. Because both $\alpha_1, \alpha_2 \in [0, \pi/3)$, necessarily $|\alpha_1 - \alpha_2| < \pi/3$. Thus $|bc| < |ab| = |ac|$, contradicting the assumption that $|ab|$ is minimal. So in fact it must be that $\vec{ab} \in \vec{Y}_6$, and the lemma is established.

Main idea of the inductive step It has already been established that Y_7 is a spanner [4]; the sector angles for Y_7 are $2\pi/7$. The main idea of our inductive proof is to partition the $\pi/3$ -sectors of Y_6 into peripheral cones of angle δ , for some fixed $\delta \in (0, \pi/9)$, leaving a central sector of angle $\pi/3 - 2\delta$. (The δ -cones are the shaded regions in Figure 9.)

When an edge of Y_6 falls inside the central sector, induction will apply, because an edge within the central sector makes definite progress toward the goal in that sector (as it does in Y_7), ensuring that the remaining distance to be covered is strictly smaller than the original. This idea is captured by the following lemma.

Lemma 12 (Induction Step). *Let $a, b, c \in S$ such that b and c lie in the same cone with apex a , and $\vec{ac} \in \vec{Y}_6$. Let $\alpha = \angle cab$ and $\beta = \angle cba$. If either $\alpha < \pi/3 - \delta$ or $\beta < \pi/3 - \delta$, then we may use induction on $p(c, b)$ to conclude that $p(a, b) \leq t|ab|$.*

Proof. This configuration is depicted in Figure 8. Because $\vec{ac} \in \vec{Y}_6$ and b and c lie in the same cone with apex a , we have that $|ac| \leq |ab|$. Because at least one of α or β is strictly smaller than $\pi/3$, we have that $|cb| < |ab|$. Thus the conditions of Lemma 8 are satisfied, so we can use Lemma 8 to bound $|ac|$ in terms of $x = |ab| - |bc|$: since $|ac|/x < t$, $|ac| < tx$. Because $|cb| < |ab|$, we may apply induction to bound $p(c, b)$: $p(c, b) \leq t|cb|$. Hence

$$p(a, b) \leq |ac| + p(c, b) \leq tx + t|cb| = t(x + |cb|) = t|ab|. \quad \square$$

We will henceforth use the symbol $\boxed{\text{Induct}}$ as shorthand for applying Lemma 12 to a triangle equivalent to that in Figure 8.

Lemma 12 leaves out Y_6 edges falling within the δ -cones, that could conceivably *not* make progress toward the goal. For example, following one edge of an equilateral triangle leaves one exactly as far away from the other corner as at the start. However, we will see that when all relevant edges of Y_6 fall within the δ -cones, the restricted geometric structure ensures that progress toward the goal is indeed made, and again induction applies.

Inductive step The inductive step proof first handles the cases where edges of Y_6 directed from a or from b fall in the central portion of the relevant sectors, and so satisfy Lemma 9, and so Lemma 12 applies.

Recall that $b \in Q_0(a)$ by our assumption. If $\vec{ab} \in \vec{Y}_6$, then $p(a, b) = |ab|$ and we are finished. Assuming otherwise, there must be a point $c \in Q_0(a)$ such that $\vec{ac} \in \vec{Y}_6$ and $|ac| \leq |ab|$. For the remainder of the proof, we are in this situation. The proof now partitions into three parts: (1) when only $Q_0(a)$ is relevant and leads to $\boxed{\text{Induct}}$; (2) when $Q_2(b)$ leads to $\boxed{\text{Induct}}$; (3) when we fall into a special situation, for which induction also applies, but for different reasons.

(1) The $Q_0(a)$ sector Consider $\triangle abc$ as previously illustrated in Figure 8. If either b or c is not in one of the δ -cones of $Q_0(a)$, then $\alpha = \angle bac < \pi/3 - \delta$: $\boxed{\text{Induct}}$.

Now assume that both b and c lie in δ -cones of $Q_0(a)$. If they both lie within the same δ -cone (Figure 9a), then again α is small: $\boxed{\text{Induct}}$. So without loss of generality let b

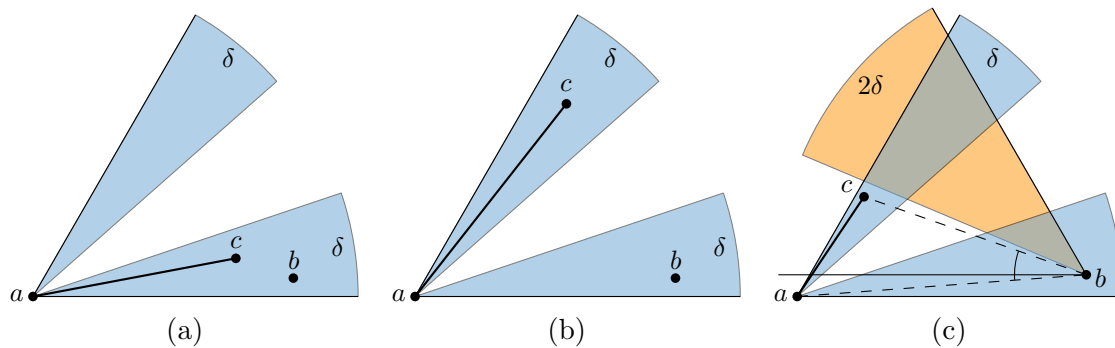


Figure 9: Different cases for the location of c ; the δ -cones are shaded. (a) b and c in the same δ -cone. (b) b and c in different δ -cones. (c) c in the upper δ -cone of $Q_0(a)$, but not in the upper 2δ -cone of $Q_2(b)$; here $\angle abc$ is “small”.

lie in the lower δ -cone, and c in the upper δ -cone of $Q_0(a)$; see Figure 9b. We cannot apply induction in this situation because the ratio s/x in Lemma 8 has no upper bound.

(2) The $Q_2(b)$ sector Now we consider $Q_2(b)$, the sector with apex at b aiming to the left of b , and assume that $c \in Q_2(b)$. Refer to Figure 9c. The case $c \notin Q_2(b)$ will be discussed later (special situation).

Because b may subtend an angle as large as δ at a with the horizontal, the “upper 2δ -cone” of $Q_2(b)$ becomes the relevant region. If c is not in the upper 2δ -cone of $Q_2(b)$ (as depicted in Figure 9c), then $\triangle abc$ satisfies Lemma 9 with $\angle abc < \pi/3 - \delta$: Induct. Note that this conclusion follows even if c is in the small region outside of and below $Q_2(b)$: the angle $\angle abc$ at b is then very small.

Assume now that c is in the upper 2δ -cone of $Q_2(b)$. Let $d \in Q_2(b)$ be the point such that $\vec{bd} \in \vec{Y}_6$. We now consider possible locations for d . If $d = c$, then $p(a, b) \leq |ac| + |cb| \leq 2|ab|$, and we are finished. So assume henceforth that d is distinct from c .

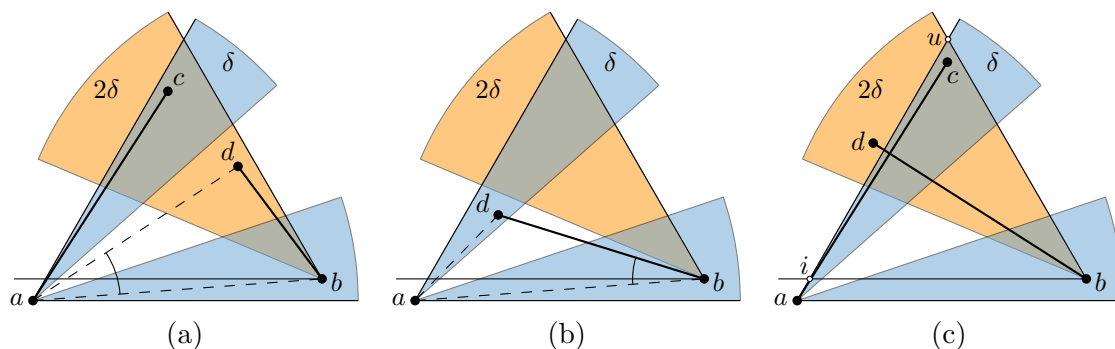


Figure 10: Different cases for the location of d . (a) d not in the upper δ -cone of $Q_0(a)$: $\angle bad$ is small. (b) d not in the upper 2δ -cone of $Q_2(b)$: $\angle abd$ is small. (c) d in the upper 2δ -cone of $Q_2(b)$, but not necessarily in the upper δ -cone of $Q_0(a)$: $|cd| < |ab|$ (Lemma 13).

If d is not in the upper δ -cone of $Q_0(a)$ (Figure 10a), then $\triangle abd$ satisfies Lemma 9 with the roles of a and b reversed: bd takes a step toward a , with the angle at a satisfying $\angle bad < \pi/3 - \delta$: Induct.

If d is not in the upper 2δ -cone of $Q_2(b)$ (Figure 10b), then $\triangle abd$ satisfies Lemma 9 again with the roles of a and b reversed and this time the angle at b bounded away from $\pi/3$, $\angle abd < \pi/3 - \delta$: Induct.

Assume now that d is in the upper 2δ -cone of $Q_2(b)$. Recall that we are in the situation where c also lies in the upper 2δ -cone of $Q_2(b)$, so it is close to d . (Note however that c and d may lie on either the same side, or on opposite sides of the upper ray of $Q_0(a)$.) See Figure 10c. This suggests the strategy of following ac and db , connected by $p(c, d)$. We show that in fact $|cd| < |ab|$, so the inductive hypothesis can be applied to $p(c, d)$. More precisely, we show the following result.

Lemma 13. *Let $a, b, c, d \in S$ be as in Figure 10c, with $\vec{bd} \in \vec{Y}_6$, $b, c \in Q_0(a)$ and $c, d \in Q_2(b)$. If both c and d lie above the lower rays bounding the upper 2δ -cones of $Q_0(a)$ and $Q_2(b)$, then for any $0 \leq \delta \leq \pi/9$,*

$$|cd| \leq \frac{\sin(2\delta)}{\sin(\pi/6 + 2\delta)} |ab|. \quad (17)$$

Note that c lies in the intersection region between the upper 2δ -cones of $Q_0(a)$ and $Q_2(b)$, because $c \in Q_0(a) \cap Q_2(b)$ (by the statement of the lemma). However, Lemma 13 does not restrict the location of d to the same region. Indeed, d may lie either below or above the upper ray bounding $Q_0(a)$, as long as it satisfies the condition $|\vec{bd}| \leq |bc|$. (This condition must hold because c, d are in the same sector $Q_2(b)$, and $\vec{bd} \in \vec{Y}_6$.) To keep the flow of our main proof uninterrupted, we defer a proof of Lemma 13 to Section 5.1.

By Lemma 13 we have $|cd| < |ab|$. Thus we can use the induction hypothesis to show that $p(c, d) \leq t|cd|$. We know that $|ac| \leq |ab|$, because both b and c are in $Q_0(a)$ and $\vec{ac} \in \vec{Y}_6$. We also know that $|bd| \leq |bc|$ because both c and d are in $Q_2(b)$ and $\vec{bd} \in \vec{Y}_6$. Let u and i be the upper and lower intersection points between the rays bounding $Q_2(b)$ and the upper ray of $Q_0(a)$, as in Figure 10c. Note that $\triangle bui$ is equilateral, and because c lies in this triangle, we have $|bc| \leq |bu| = |bi| \leq |ab|$. It follows that $|bd| \leq |ab|$. So in this situation (illustrated in Figure 10c), we have:

$$\begin{aligned} p(a, b) &\leq |ac| + p(c, d) + |bd| \\ &\leq 2|ab| + p(c, d) \\ &\leq 2|ab| + t|cd| \\ &\leq 2|ab| + t \frac{\sin(2\delta)}{\sin(\pi/6 + 2\delta)} |ab| \\ &\leq t|ab|. \end{aligned}$$

Here we have applied Lemma 13 to bound $|cd|$. Note that the latter inequality above is true for the value of t from (16).

(3) Special situation The only case left to discuss is the one in which c lies in the upper δ -cone of $Q_0(a)$ and to the right of the upper ray of $Q_2(b)$. This situation is depicted in Figure 11. Next consider $Q_4(c)$. Because $b \in Q_4(c)$, there exists $\vec{c}z \in \vec{Y}_6$, with $z \in Q_4(c)$ and $|cz| \leq |cb|$. Clearly $z \in Q_0(a) \cup Q_5(a)$. Note that the disk sector $D_a^c \subset Q_0(a)$ with center a and radius $|ac|$ must be empty, because $\vec{a}c \in \vec{Y}_6$.

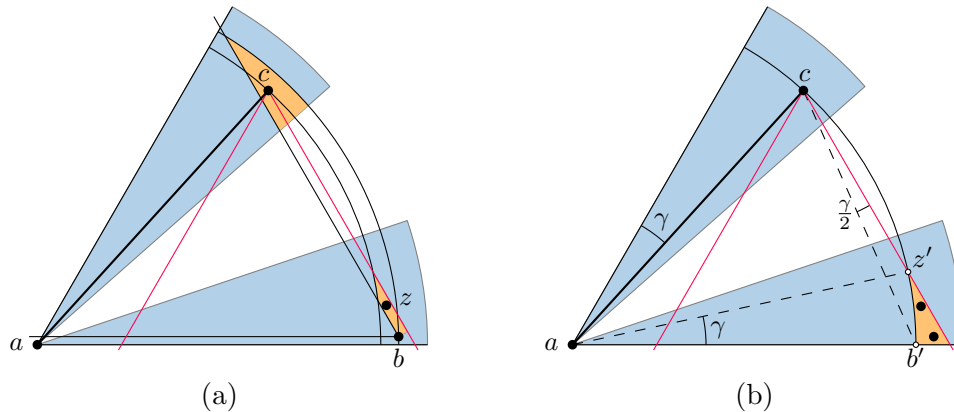


Figure 11: (a) The special situation, with $c \notin Q_2(b)$ and $z \in Q_0(a)$. (b) Both b and z must lie in the highlighted region, so we can apply Lemma 13.

Case 3(a) If $z \in Q_0(a)$, then z lies in the lower δ -cone of $Q_0(a)$ and to the right of D_a^c , close to b . See Figure 11a. In this case we show that the quantity on the right side of inequality (17) is a loose upper bound on $|bz|$, and that similar inductive arguments hold here as well. Let the circumference of D_a^c intersect the right ray of $Q_4(c)$ and the lower ray of $Q_0(a)$ at points $z' \neq c$ and b' , respectively. Refer to Figure 11b. Let $\gamma \leq \delta$ be the angle formed by ac with the upper ray of $Q_0(a)$. Then $\angle z'ab' = \gamma$ and $\angle z'cb' = \gamma/2$ by the inscribed angle theorem. This implies that both b' and z' lie in the intersection region between the lower δ -cone of $Q_0(a)$ and the right $\delta/2$ -cone of $Q_4(c)$. Thus $a, b, c, z \in S$ satisfy the conditions of Lemma 13, with the roles of b and c reversed: $|bz| \leq \sin(2\delta)/\sin(\pi/6 + 2\delta) \cdot |ac|$.

Arguments similar to the ones used in the proof of Lemma 13 show that $|cz| \leq |ac|$. This along with $|ac| \leq |ab|$ (because $\vec{a}c \in \vec{Y}_6$) and the above inequality imply

$$p(a, b) \leq |ac| + |cz| + p(z, b) \leq 2|ab| + t|bz| \leq t|ab|$$

for any t satisfying the conditions stated by this lemma.

Case 3(b) Assume now that $z \notin Q_0(a)$. Then $z \in Q_5(a)$, as depicted in Figure 12. In this case z lies in the disk sector D_c^b (because $|cz| \leq |cb|$) and below the horizontal line through a (because D_a^c is empty). This implies that there exists $\vec{a}e \in \vec{Y}_6$, with $e \in Q_5(a)$ and $|ae| \leq |az|$. Similarly, there exists $\vec{b}f \in \vec{Y}_6$, with $f \in Q_3(b)$ and $|bf| \leq |bz|$. If e lies above the lower 2δ -cone of $Q_5(a)$, then $\angle bae \leq \pi/3 - \delta$, which leads to **Induct** and settles this case. Similarly, if f lies above the lower δ -cone of $Q_3(b)$, then $\angle abf \leq \pi/3 - \delta$, which again leads to **Induct**. Otherwise, we show that the following lemma holds.

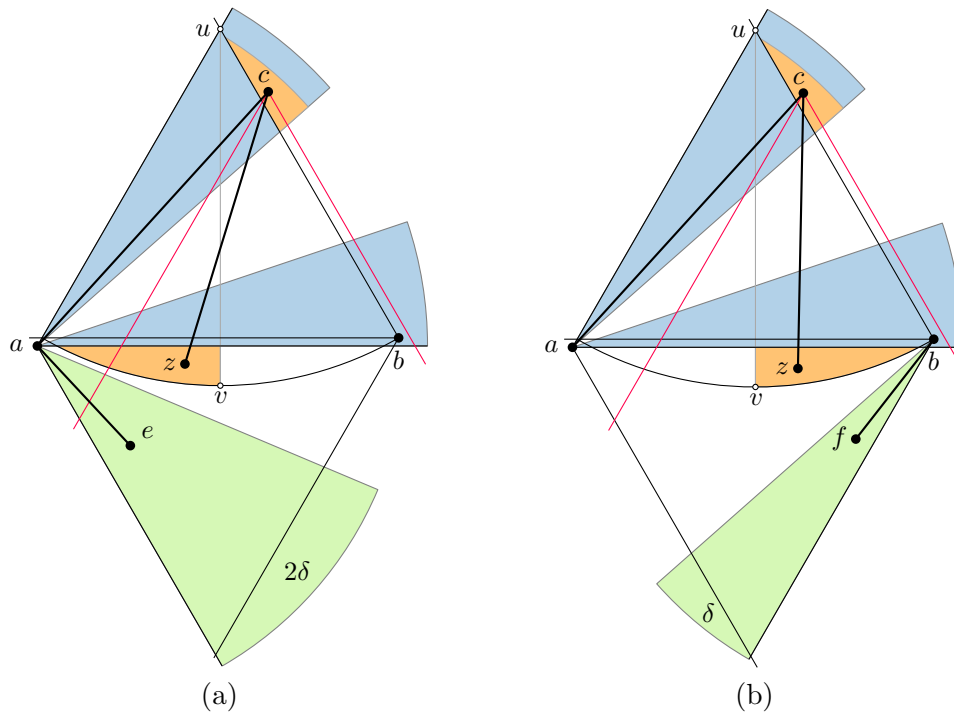


Figure 12: The special situation, with $c \notin Q_2(b)$ and $z \in Q_5(a)$. (a) z lies left of uv : $e \in Q_3(b)$. (b) z lies right of uv : $f \in Q_5(a)$.

Lemma 14. *Let $a, b, c, z \in S$ be in the configuration depicted in Figure 12, with $\vec{ac}, \vec{cz} \in \vec{Y}_6$. Let $\vec{ae}, \vec{bf} \in \vec{Y}_6$, with e in the lower 2δ -cone of $Q_5(a)$ and f in the lower δ -cone of $Q_3(b)$. Then at least one of the following is true: (a) $e \in Q_3(b)$, or (b) $f \in Q_5(a)$.*

We defer a proof of Lemma 14 to Section 5.2.

Lemma 14 guarantees that, if condition (a) holds, then ae may not cross the lower ray bounding $Q_3(b)$. This case reduces to one of the cases depicted in Figure 10, with e playing the role of c and the path passing under ab rather than above. Because ae does not cross the lower ray bounding $Q_3(b)$, the special situation depicted in Figure 11 (with e playing the role of c) may not occur in this case. Similarly, condition (b) from Lemma 14 reduces to one of the cases depicted in Figure 10, with the roles of a and b reversed and with f playing the role of c ; the special situation depicted in Figure 11 (with bf playing the role of ac) may not occur in this case. Having exhausted all cases, we conclude the proof. \square

Next we establish a lower bound on the spanning ratio of Y_6 .

Theorem 15. *The graph Y_6 has spanning ratio at least 2.*

Proof. We construct a lower bound example by extending the shortest path between two points a and b . Let $b \in Q_1(a)$ lie arbitrarily close to the cone boundary separating $Q_0(a)$ and $Q_1(a)$ (see Figure 13a). Let $\varepsilon > 0$ be a small constant, and let $c \in Q_1(a)$ such that

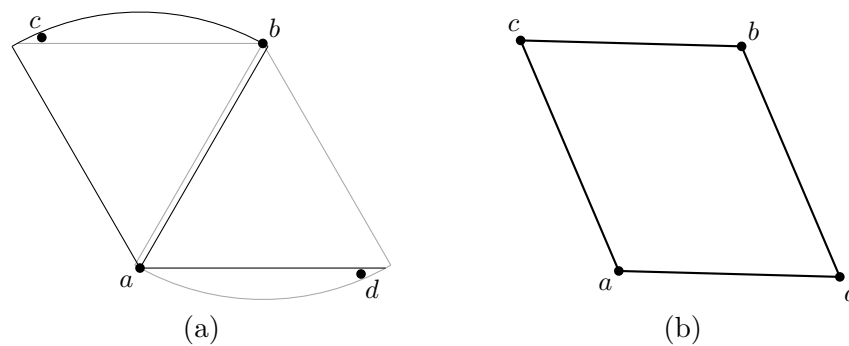


Figure 13: Y_6 has spanning ratio at least 2. (a) The lower bound construction. (b) The resulting Y_6 graph.

$|ac| = |ab| - \varepsilon$ and $|bc| \simeq |ab|$. Similarly, let $d \in Q_4(b)$ such that $|bd| = |ab| - \varepsilon$ and $|ad| \simeq |ab|$. Then the corresponding Y_6 graph is as depicted in Figure 13b. Note that there are two shortest paths between a and b , both of length $\simeq 2|ab|$. \square

4 Other lower bounds

In this section we provide lower bounds for all Yao graphs with at least six cones. We do this by dividing this group of graphs in four families, depending on their number of cones. We distinguish Yao graphs with $4x + 2$ cones ($x \geq 1$), $4x + 3$ cones, $4x + 4$ cones, and $4x + 5$ cones. This division was introduced by Bose *et al.* [8] to improve the analysis of Θ -graphs. It is based on the relation between the line perpendicular to the bisector of a cone and the cone boundaries. For example, in a Yao graph with $4x + 2$ cones, this line coincides with a cone boundary, whereas in a graph with $4x + 4$ cones it coincides with the bisector of another cone. For an overview of all bounds derived in this section, see Figure 25 and Table 1 in Section 6.

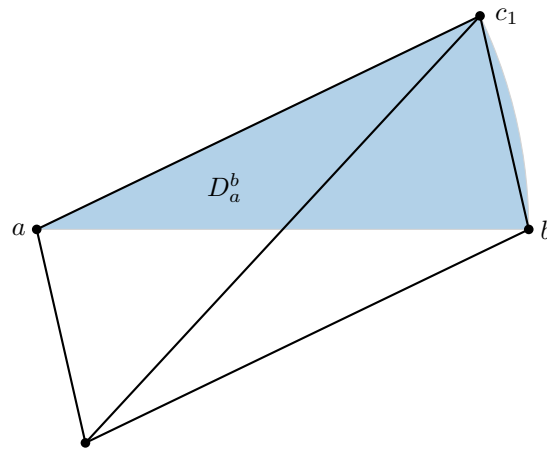
4.1 Lower bound for Y_{4x+2}

In this section we provide a lower bound for Yao graphs with $4x + 2$ cones ($x \geq 1$).

Theorem 16. *For all $x \geq 1$, the graph Y_{4x+2} has spanning ratio at least $1 + 2 \sin(\theta/2)$, where $\theta = 2\pi/(4x + 2)$.*

Proof. We construct the lower bound example by extending the shortest path between two vertices a and b . We describe only how to extend one of the shortest paths between these vertices. To extend all shortest paths, the same modification is performed in each of the analogous cases (see Figure 14).

First, we place b arbitrarily close to a cone boundary of a . Next, we ensure that there is no edge between a and b by placing a vertex c_1 in the corner of D_a^b that is furthest from b (see Figure 14). One of the shortest paths in the resulting graph visits a , c_1 and b . Thus, to obtain a lower bound for Y_{4x+2} , we compute the length of this path.

Figure 14: The construction of the lower bound for Y_{4x+2}

By construction, we have that $\angle c_1ab = \theta$, hence we can express the various line segments as follows:

$$|ac_1| = |ab|$$

$$|c_1b| = 2 \sin\left(\frac{\theta}{2}\right) \cdot |ab|$$

Hence, the total length of the shortest path is $|ac_1| + |c_1b|$, which can be rewritten to

$$\left(1 + 2 \sin\left(\frac{\theta}{2}\right)\right) \cdot |ab|,$$

proving the theorem. □

4.2 Lower bound for Y_{4x+3}

In this section we provide a lower bound for Yao graphs with $4x + 3$ cones ($x \geq 1$).

Theorem 17. *For all $x \geq 1$, the graph Y_{4x+3} has spanning ratio at least*

$$1 + 2 \sin\left(\frac{3\theta}{8}\right) + 4 \frac{(\sin(\frac{13\theta}{16}) + \sin(\frac{19\theta}{16})) \sin(\frac{\theta}{16}) \sin(\frac{3\theta}{8})}{\sin(2\theta)},$$

where $\theta = 2\pi/(4x + 3)$.

Proof. We construct the lower bound example by extending the shortest path between two vertices a and b in three steps. We describe only how to extend one of the shortest paths between these vertices. To extend all shortest paths, the same modification is performed in each of the analogous cases (see Figure 15).

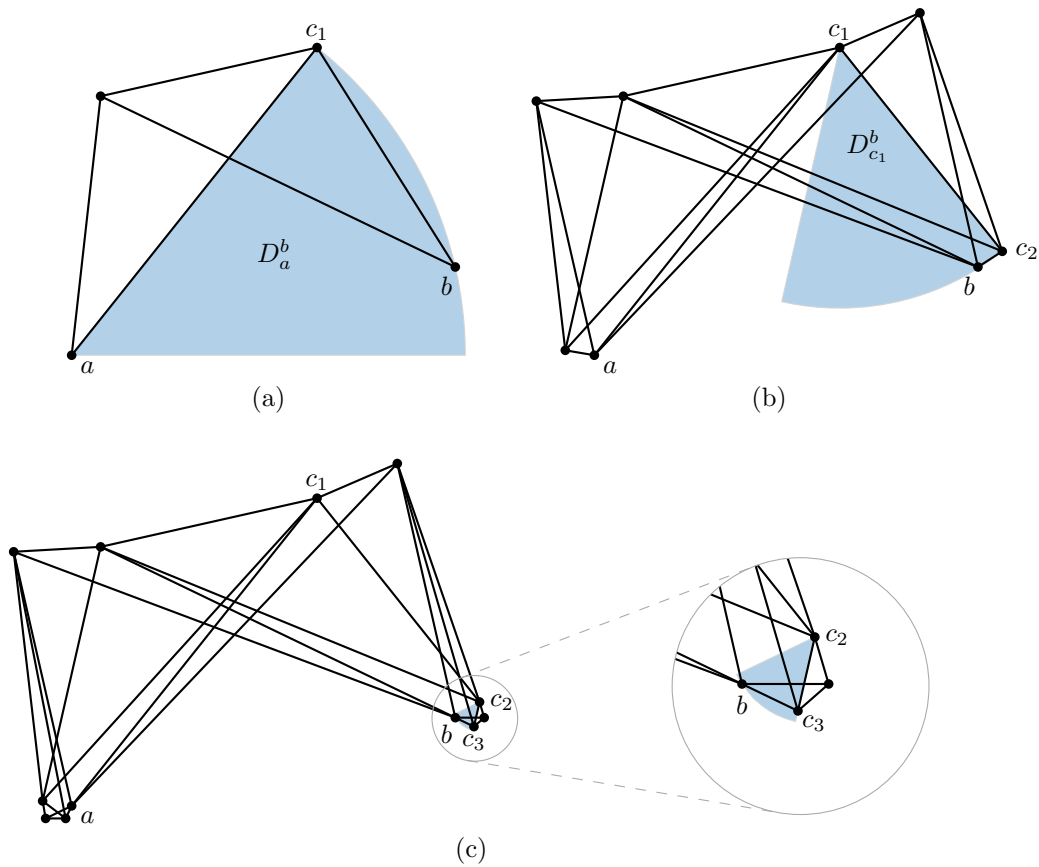


Figure 15: The construction of the lower bound for Y_{4x+3}

First, we place b such that the angle between ab and the bisector of the cone of a that contains b is $\theta/4$. Next, we ensure that there is no edge between a and b by placing a vertex c_1 in the corner of D_a^b that is furthest from b (see Figure 15a). Next, we place a vertex c_2 in the corner of $D_{c_1}^b$ that is closest to b , since we need to ensure that there is no edge between a and c_2 (see Figure 15b). Finally, we place a vertex c_3 in $D_{c_2}^b$ such that $D_{c_3}^a$ contains b (see Figure 15c). This ensures that no shortcut to a is created by c_3 . One of the shortest paths in the resulting graph visits a, c_1, c_2, c_3 and b . Thus, to obtain a lower bound for Y_{4x+3} , we compute the length of this path.

By construction, we have that $\angle c_1ab = 3\theta/4$. To compute the total length of the shortest path between a and b , we also need $\angle c_2c_1b = \angle c_2c_1a - \angle bc_1a$ (see Figure 16). Since $\angle c_2c_1a = \pi/2 - \theta/4$ and $\angle bc_1a = (\pi - 3\theta/4)/2$, it follows that $\angle c_2c_1b = \theta/8$. Finally, we need angles $\angle bc_3c_2 = \pi - 2\theta$, $\angle c_3c_2b = \angle c_3c_2c_1 - \angle bc_2c_1 = (\pi/2 + 3\theta/4) - (\pi/2 - \theta/16) = 13\theta/16$, and $\angle c_2bc_3 = \pi - \angle bc_3c_2 - \angle c_3c_2b = 19\theta/16$.

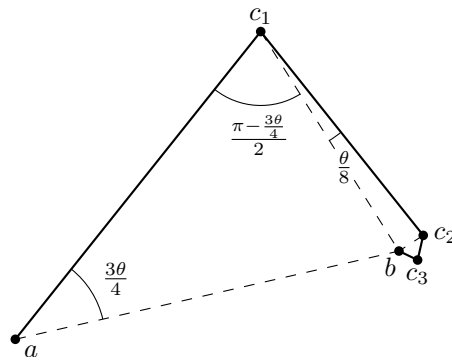


Figure 16: The lower bound for Y_{4x+3}

We can now express the various line segments as follows:

$$|ac_1| = |ab|$$

$$|c_1c_2| = |c_1b| = 2 \sin\left(\frac{3\theta}{8}\right) \cdot |ab|$$

$$|c_2b| = 2 \sin\left(\frac{\theta}{16}\right) \cdot |c_1b| = 4 \sin\left(\frac{\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right) \cdot |ab|$$

$$|c_2c_3| = \frac{\sin\left(\frac{19\theta}{16}\right)}{\sin(\pi - 2\theta)} \cdot |c_2b| = \frac{\sin\left(\frac{19\theta}{16}\right)}{\sin(2\theta)} \cdot |c_2b| = 4 \frac{\sin\left(\frac{19\theta}{16}\right) \sin\left(\frac{\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right)}{\sin(2\theta)} \cdot |ab|$$

$$|c_3b| = \frac{\sin\left(\frac{13\theta}{16}\right)}{\sin(\pi - 2\theta)} \cdot |c_2b| = \frac{\sin\left(\frac{13\theta}{16}\right)}{\sin(2\theta)} \cdot |c_2b| = 4 \frac{\sin\left(\frac{13\theta}{16}\right) \sin\left(\frac{\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right)}{\sin(2\theta)} \cdot |ab|$$

Hence, the total length of the shortest path is $|ac_1| + |c_1c_2| + |c_2c_3| + |c_3b|$, which can be rewritten as

$$\left(1 + 2 \sin\left(\frac{3\theta}{8}\right) + 4 \frac{(\sin\left(\frac{13\theta}{16}\right) + \sin\left(\frac{19\theta}{16}\right)) \sin\left(\frac{\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right)}{\sin(2\theta)}\right) \cdot |ab|,$$

proving the theorem. □

4.3 Lower bound for Y_{4x+4}

In this section we provide a lower bound for Yao graphs with $4x + 4$ cones ($x \geq 1$).

Theorem 18. *For all $x \geq 1$, the graph Y_{4x+4} has spanning ratio at least*

$$1 + 2 \sin\left(\frac{\theta}{2}\right) \left(1 + \tan\left(\frac{\theta}{2}\right)\right),$$

where $\theta = 2\pi/(4x + 4)$.

Proof. We construct the lower bound example by extending the shortest path between two vertices a and b in three steps. We describe only how to extend one of the shortest paths between these vertices. To extend all shortest paths, the same modification is performed in each of the analogous cases (see Figure 17).

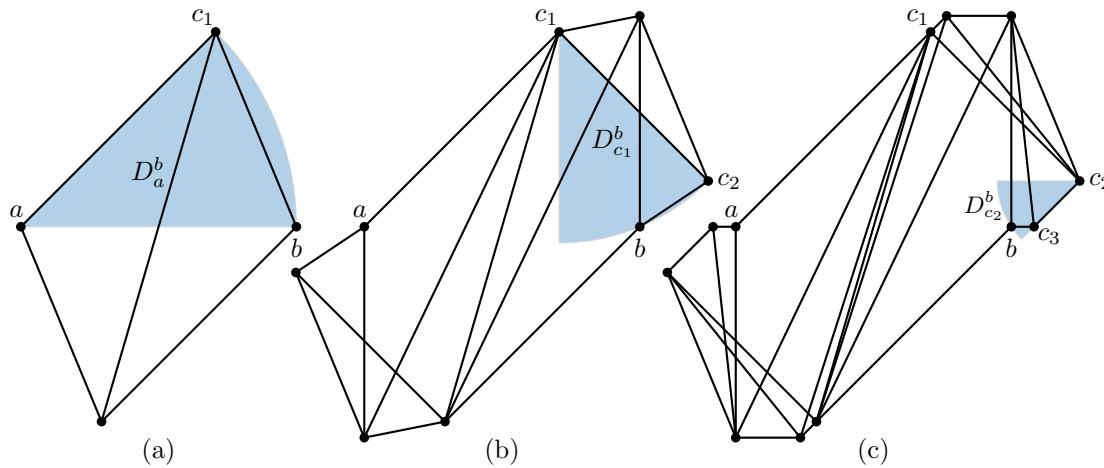


Figure 17: The construction of the lower bound for Y_{4x+4}

First, we place b arbitrarily close to a cone boundary of a . Next, we ensure that there is no edge between a and b by placing a vertex c_1 in the corner of D_a^b that is furthest from b (see Figure 17a). Next, we place a vertex c_2 in the corner of $D_{c_1}^b$ that is furthest from a (see Figure 17b). Finally, we place a vertex c_3 on the intersection of $D_{c_2}^b$ and $D_b^{c_2}$ (see Figure 17c). This ensures that no shortcut to a is created by c_3 . One of the shortest paths in the resulting graph visits a, c_1, c_2, c_3 and b . Thus, to obtain a lower bound for Y_{4x+4} , we compute the length of this path.

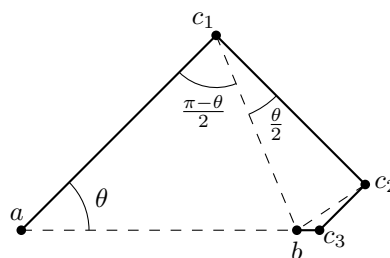


Figure 18: The lower bound for Y_{4x+4}

By construction, we have that $\angle c_1ab = \theta$. To compute the total length of the shortest path between a and b , we also need $\angle c_2c_1b = \angle c_2c_1a - \angle bc_1a$ (see Figure 18). Since $\angle c_2c_1a = \pi/2$ and $\angle bc_1a = (\pi - \theta)/2$, it follows that $\angle c_2c_1b = \theta/2$. Finally, we need angles $\angle bc_3c_2 = \pi - \theta$, $\angle c_3c_2b = \angle c_3c_2c_1 - \angle bc_2c_1 = \pi/2 - (\pi/2 - \theta/4) = \theta/4$, and $\angle c_2bc_3 = \pi - \angle bc_3c_2 - \angle c_3c_2b = 3\theta/4$.

Hence, we can express the various line segments as follows:

$$\begin{aligned}
 |ac_1| &= |ab| \\
 |c_1c_2| &= |c_1b| = 2 \sin\left(\frac{\theta}{2}\right) \cdot |ab| \\
 |c_2b| &= 2 \sin\left(\frac{\theta}{4}\right) \cdot |c_1b| = 4 \sin\left(\frac{\theta}{4}\right) \sin\left(\frac{\theta}{2}\right) \cdot |ab| \\
 |c_2c_3| &= \frac{\sin\left(\frac{3\theta}{4}\right)}{\sin(\pi - \theta)} \cdot |c_2b| = \frac{\sin\left(\frac{3\theta}{4}\right)}{\sin(\theta)} \cdot |c_2b| = 4 \frac{\sin\left(\frac{3\theta}{4}\right) \sin\left(\frac{\theta}{4}\right) \sin\left(\frac{\theta}{2}\right)}{\sin(\theta)} \cdot |ab| \\
 |c_3b| &= \frac{\sin\left(\frac{\theta}{4}\right)}{\sin(\pi - \theta)} \cdot |c_2b| = \frac{\sin\left(\frac{\theta}{4}\right)}{\sin(\theta)} \cdot |c_2b| = 4 \frac{\sin\left(\frac{\theta}{4}\right) \sin\left(\frac{\theta}{4}\right) \sin\left(\frac{\theta}{2}\right)}{\sin(\theta)} \cdot |ab|
 \end{aligned}$$

Hence, the total length of the shortest path is $|ac_1| + |c_1c_2| + |c_2c_3| + |c_3b|$, which can be rewritten to

$$\left(1 + 2 \sin\left(\frac{\theta}{2}\right) + 4 \frac{(\sin\left(\frac{3\theta}{4}\right) + \sin\left(\frac{\theta}{4}\right)) \sin\left(\frac{\theta}{4}\right) \sin\left(\frac{\theta}{2}\right)}{\sin(\theta)}\right) \cdot |ab|,$$

which is equal to

$$\left(1 + 2 \sin\left(\frac{\theta}{2}\right) \left(1 + \tan\left(\frac{\theta}{2}\right)\right)\right) \cdot |ab|,$$

proving the theorem. □

4.4 Lower bound for Y_{4x+5}

In this section we provide a lower bound for Yao graphs with $4x + 5$ cones ($x \geq 1$).

Theorem 19. *For all $x \geq 1$, the graph Y_{4x+5} has spanning ratio at least*

$$1 + 2 \sin\left(\frac{3\theta}{8}\right) + 4 \sin\left(\frac{5\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right),$$

where $\theta = 2\pi/(4x + 5)$.

Proof. We construct the lower bound example by extending the shortest path between two vertices a and b in two steps. We describe only how to extend one of the shortest paths between these vertices. To extend all shortest paths, the same modification is performed in each of the analogous cases (see Figure 19).

First, we place b such that the angle between ab and the bisector of the cone of a that contains b is $\theta/4$. Next, we ensure that there is no edge between a and b by placing a vertex c_1 in the corner of D_a^b that is furthest from b (see Figure 19a). Next, we place a vertex c_2 in the corner of $D_{c_1}^b$ that is furthest from b (see Figure 19b). One of the shortest

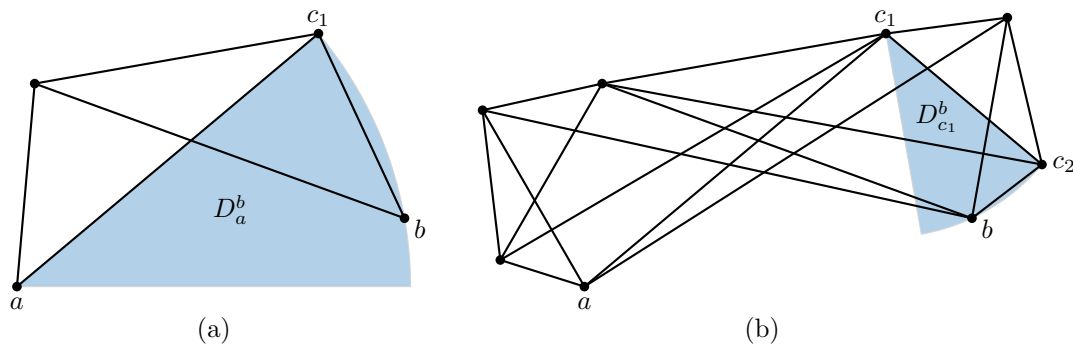


Figure 19: The construction of the lower bound for Y_{4x+5}

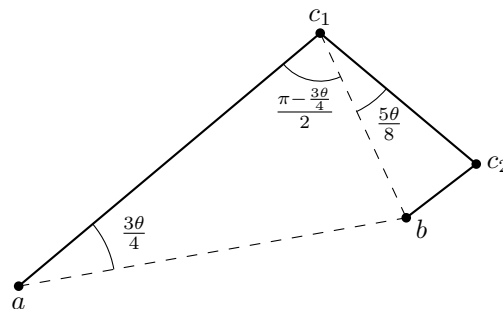


Figure 20: The lower bound for Y_{4x+5}

paths in the resulting graph visits a, c_1, c_2 and b . Thus, to obtain a lower bound for Y_{4x+5} , we compute the length of this path.

By construction, we have that $\angle c_1ab = 3\theta/4$. To compute the total length of the shortest path between a and b , we also need $\angle c_2c_1b = \angle c_2c_1a - \angle bc_1a$ (see Figure 20). Since $\angle c_2c_1a = \pi/2 + \theta/4$ and $\angle bc_1a = (\pi - 3\theta/4)/2$, it follows that $\angle c_2c_1b = 5\theta/8$. Hence, we can express the various line segments as follows:

$$\begin{aligned}
 |ac_1| &= |ab| \\
 |c_1c_2| &= |c_1b| = 2 \sin\left(\frac{3\theta}{8}\right) \cdot |ab| \\
 |c_2b| &= 2 \sin\left(\frac{5\theta}{16}\right) \cdot |c_1b| = 4 \sin\left(\frac{5\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right) \cdot |ab|
 \end{aligned}$$

Hence, the total length of the shortest path is $|ac_1| + |c_1c_2| + |c_2b|$, which can be rewritten to

$$\left(1 + 2 \sin\left(\frac{3\theta}{8}\right) + 4 \sin\left(\frac{5\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right)\right) \cdot |ab|,$$

proving the theorem. □

4.5 The Yao-Yao graph with 5 cones is not a spanner

One disadvantage of Yao graphs is that the maximum degree of a vertex might be $n - 1$. For example, this happens when $n - 1$ points are spread evenly on a circle centered on the last point. The *Yao-Yao graph*, introduced by Li *et al.* [14], solves this problem by first constructing the directed Yao graph, and then discarding all but the shortest incoming edge in each cone (ties are broken arbitrarily). As a result, each vertex in the Yao-Yao graph with k cones, denoted by YY_k , has maximum degree $2k$: one incoming and one outgoing edge per cone. In the resulting graph, the directions of the edges are typically ignored.

Of course, discarding all these edges has a cost: the spanning ratio increases. For a long time, it was unknown whether Yao-Yao graphs were even spanners. The first answers to this question were negative. Damian, El Molla, and Pinciu [10, 12] showed that, even though Y_4 and Y_6 are spanners, YY_k is not a spanner for $k \leq 4$ and $k = 6$. The first positive results followed soon afterwards, when Bauer and Damian [2] proved that YY_k is a spanner for all $k = 6x$, with $x \geq 6$. The spanner status of all other Yao-Yao graphs is still open.

We close the gap among Yao-Yao graphs with six or fewer cones, by presenting a construction of a YY_5 graph whose stretch factor is unbounded. Figure 21 shows the initial steps of constructing such a YY_5 graph, where the path between a and b can grow horizontally to the right by adding more points following the pattern, exceeding any bound on the stretch factor.

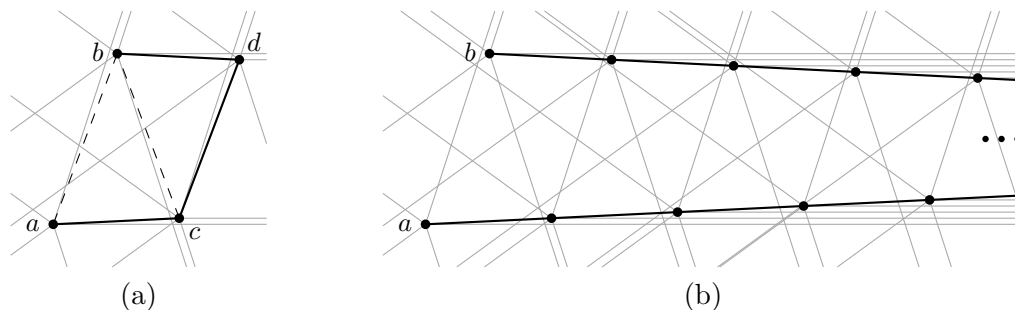


Figure 21: The construction of a YY_5 graph with unbounded stretch factor. The gray lines are the boundaries of the cones. (a) The first four vertices. The dashed edges are in Y_5 , but are discarded in favor of the bold edges. (b) The pattern continued to the right.

The pattern begins with four vertices, positioned as in Figure 21a. Vertex b is placed close to the upper boundary of $Q_0(a)$, and c is placed in $Q_0(a) \cap Q_4(b)$, near the intersection of their boundaries. Similarly, d is placed in $Q_0(c) \cap Q_4(b)$ near the intersection of their boundaries. In the construction of the Yao-Yao graph on these points, the Yao edge \vec{ba} is discarded since \vec{ca} is shorter, and \vec{cb} is discarded because \vec{db} is shorter. By recursively applying this pattern with c and d in the roles of a and b , we can eliminate the edge cd and push the connecting edge arbitrarily far to the right. Since the distance between a and b remains the same, we can construct a YY_5 graph with arbitrarily large spanning ratio.

Theorem 20. *The graph YY_5 is not a constant spanner.*

5 Deferred proofs

In this section, we prove the remaining lemmas required for the proof of Theorem 10.

5.1 Proof of Lemma 13

Lemma 13. *Let $a, b, c, d \in S$ be as in Figure 22a, with $\vec{bd} \in \vec{Y}_6$, $b, c \in Q_0(a)$ and $c, d \in Q_2(b)$. If both c and d lie above the lower rays bounding the upper 2δ -cones of $Q_0(a)$ and $Q_2(b)$, then for any $0 \leq \delta \leq \pi/9$,*

$$|cd| \leq \frac{\sin(2\delta)}{\sin(\pi/6 + 2\delta)} |ab|$$

Proof. Let u and v be the top and bottom points of the intersection quadrilateral R between the upper 2δ -cones of $Q_0(a)$ and $Q_2(b)$. See Figure 22. Then $c \in R$. For any $\delta \leq \pi/9$, the angles opposite to the diagonal uv of R are bounded below by $5\pi/9$, therefore uv is the diameter of R .

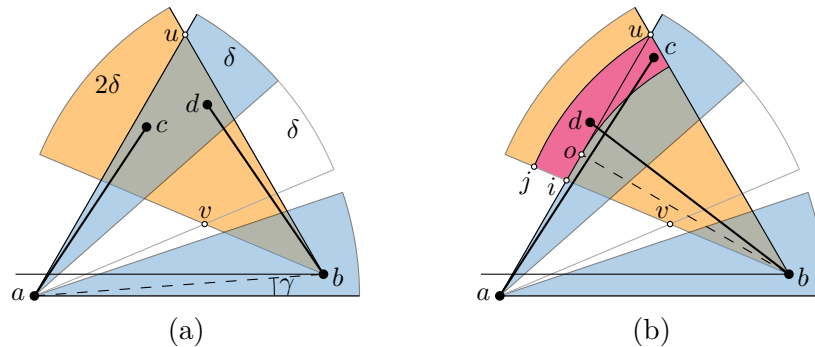


Figure 22: (a) The configuration for Lemma 13. (b) The distance $|cd|$ is still bounded by $|uv|$ when d lies outside of $Q_0(a)$.

Assume first that $d \in R$ as well. In this case, the quantity $|cd|$ is bounded above by the length $|uv|$ of the diameter of R . Let γ be the angle formed by ab with the horizontal. We show that $|uv|$ is maximized when $\gamma = 0$. Set a coordinate system with the origin at a . Scale the point set S so that $|ab| = 1$. Then the coordinates of b are $(\cos \gamma, \sin \gamma)$. The point u is at the intersection of the two lines passing through a and b with slopes $\tan \pi/3$ and $-\tan \pi/3$ respectively, given by $y = \sqrt{3}x$ and $y = -\sqrt{3}(x - \cos \gamma) + \sin \gamma$. Solving for x and y gives the coordinates of u

$$x_u = \frac{\sqrt{3} \cos \gamma + \sin \gamma}{2\sqrt{3}}, \quad y_u = \frac{\sqrt{3} \cos \gamma + \sin \gamma}{2}.$$

Similarly, the point v is at the intersection of two lines given by $y = \tan(\pi/3 - 2\delta)x$ and $y = -\tan(\pi/3 - 2\delta)(x - \cos \gamma) + \sin \gamma$. Solving for x and y gives the coordinates of v

$$x_v = \frac{\tan(\pi/3 - 2\delta) \cos \gamma + \sin \gamma}{2 \tan(\pi/3 - 2\delta)},$$

$$y_v = \frac{\tan(\pi/3 - 2\delta) \cos \gamma + \sin \gamma}{2}.$$

We can now compute $|uv| = \sqrt{(x_u - x_v)^2 + (y_u - y_v)^2}$ as a function of γ and δ . The

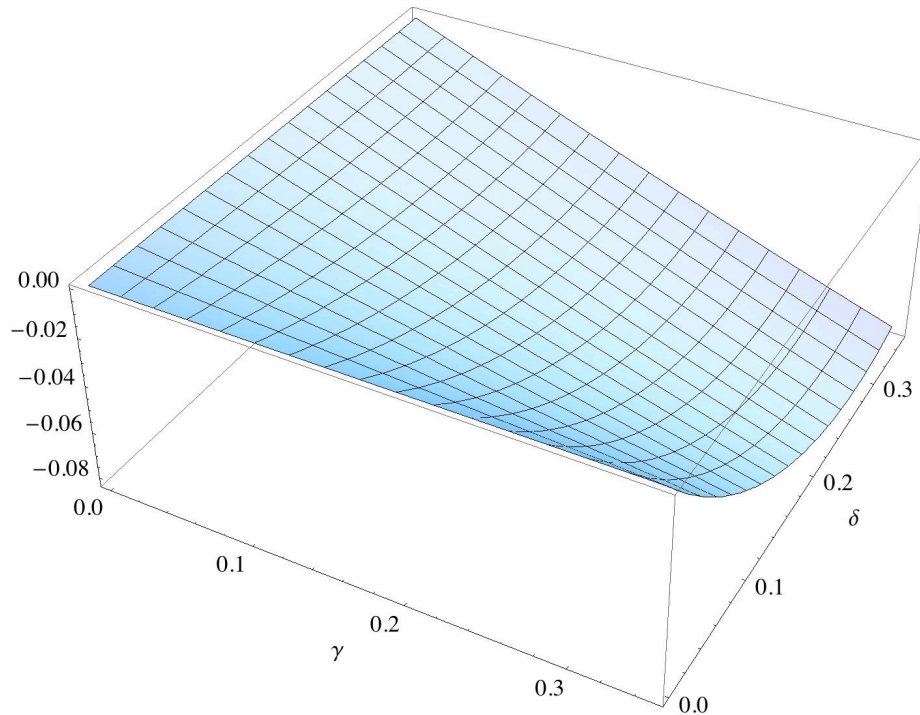


Figure 23: The derivative of $|uv|$ with respect to γ , for $\gamma, \delta \in [0, \pi/9]$.

derivative of this function with respect to γ is represented as a graph in Figure 23 for $\gamma, \delta \in [0, \pi/9]$. Note that this function is negative on the given interval, therefore $|uv|$ increases as γ decreases. Thus $|uv|$ is maximum when $\gamma = 0$. We now set $\gamma = 0$ and compute $|uv| \leq \sqrt{3}/2 - \cot(2\delta + \pi/6)/2 = \sin(2\delta)/\sin(2\delta + \pi/6)$ as claimed.

Assume now that $d \notin R$, so d lies above the upper ray bounding $Q_0(a)$. Let i and o be the points where the upper ray bounding $Q_0(a)$ intersects the lower ray bounding the upper 2δ -cone of $Q_2(b)$ and the bisector of $Q_2(b)$, respectively. Let j be the intersection point between the lower ray bounding the upper 2δ -cone of $Q_2(b)$ and the circumference of D_b^u . If $\delta > \pi/12$ (and so $2\delta > \pi/6$), then i lies below o , otherwise i coincides with, or lies above o . We define the path $p(i)$ to be the line segment io concatenated with the arc of the disk sector D_b^o , if $\delta > \pi/12$, or simply the arc of the disk sector D_b^i , if $\delta \leq \pi/12$. Since d is to the left of the line supporting io , we have that $|bo| < |bd| < |bc|$. (This latter inequality follows from the fact that $\vec{bd} \in \vec{Y}_6$.) This implies that c also lies left of $p(i)$, so both c and d lie in the strip delimited by D_b^u , $p(i)$ and the two rays bounding the upper 2δ -cone of $Q_2(b)$.

Thus $|cd|$ is no greater than the diameter of this strip, which we show to be no greater than the diameter of R . For this, it suffices to show that $\max\{|ui|, |uj|, |ij|\} \leq |uv|$.

Because ui is an edge of R , $|ui|$ is clearly no greater than the diameter $|uv|$ of R . Next we show that $|uj| \leq |uv|$. From the isosceles triangle Δbuj we derive $\angle ujb = \pi/2 - \delta$. Angle $\angle uvj$ is exterior to Δuvb , therefore $\angle uvj = \angle vub + 2\delta \leq \pi/6 + 2\delta$ (note that $\angle vub = \pi/6$ when ab is horizontal, otherwise $\angle vub < \pi/6$). It follows that $\angle ujb \leq \angle uvj$ for any $\delta \leq \pi/9$. This along with the law of sines applied to Δujb yields $|uj| \leq |uv|$.

It remains to show that $|ij| < |uv|$. We will in fact show that $|ij| < |uj|$, which along with the conclusion above that $|uj| \leq |uv|$, yields $|ij| < |uv|$. Angle $\angle uij$ is exterior to Δuib , therefore $\pi/3 \leq \angle uij \leq \pi/3 + 2\delta$. Earlier we showed that $\angle ujb = \pi/2 - \delta \geq 7\pi/18$, for any $\delta \leq \pi/9$. It follows that $\angle iuj \leq \pi - (7\pi/18 + \pi/3) = 5\pi/18$ is the smallest angle of Δuij , therefore $|ij| < |uv|$. This completes the proof. \square

5.2 Proof of Lemma 14

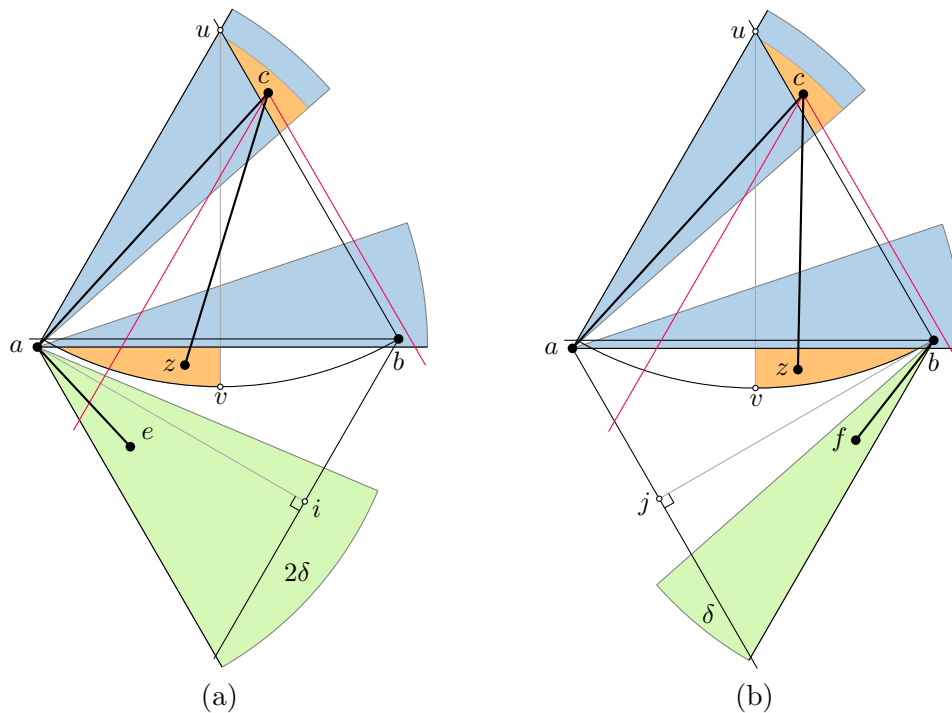


Figure 24: The configuration for Lemma 14, with $c \notin Q_2(b)$ and $z \in Q_5(a)$. (a) z lies left of uv : $e \in Q_3(b)$. (b) z lies right of uv : $f \in Q_5(a)$.

Lemma 14. *Let $a, b, c, z \in S$ be in the configuration depicted in Figure 24, with $\vec{ac}, \vec{cz} \in \vec{Y}_6$. Let $\vec{ae}, \vec{bf} \in \vec{Y}_6$, with e in the lower 2δ -cone of $Q_5(a)$ and f in the lower δ -cone of $Q_3(b)$. Then at least one of the following is true:*

- (a) $e \in Q_3(b)$
- (b) $f \in Q_5(a)$

Proof. We define four intersection points u , v , i and j as follows: u is at the intersection between the top rays of $Q_0(a)$ and $Q_2(b)$; v is at the intersection between the bisector of $\angle aub$ and the boundary of the disk sector D_u^b ; i is the foot of the perpendicular from a on the lower ray of $Q_3(b)$; and j is the foot of the perpendicular from b on the lower ray of $Q_5(a)$. Refer to Figure 24.

Note that $|ae| \leq |ai|$ implies condition (a), and $|bf| \leq |bj|$ implies condition (b). We show that the first holds if z lies to the left of or on uv , and the latter holds if z lies to the right of or on uv (and so at least one of the two conditions holds). We first show that $z \in D_u^b$. This follows immediately from the inequality $|uz| + |cb| < |cz| + |ub|$ (which can be derived using the triangle inequality twice on the triangles induced by the diagonals of $ucbz$), and the fact that $|cz| \leq |cb|$ (because $\vec{cz} \in \vec{Y}_6$). It follows that $|uz| < |ub|$, therefore $z \in D_u^b$.

Condition (a) Assume that z lies to the left of uv (as in Figure 24a). Because $z \in D_u^b$ is below the horizontal through a , $\angle azv$ is obtuse and therefore $|az| \leq |av|$ (equality holds when z coincides with v). Also $|ae| \leq |az|$, because z and e are in the same sector $Q_5(a)$ and $\vec{ae} \in \vec{Y}_6$. It follows that $|ae| \leq |av|$. We now show that $|av| \leq |ai|$, which implies $|ae| \leq |ai|$, thus settling this case.

Let $\gamma \in [0, \delta]$ be the angle formed by ab with the horizontal through a . Then $\angle abi = \pi/3 - \gamma$ and $|ai| = |ab| \sin(\pi/3 - \gamma)$. The law of sines applied to $\triangle uav$ tells us that

$$\frac{|av|}{\sin \pi/6} = \frac{|ua|}{\sin \angle uva} = \frac{|uv|}{\sin \angle uav}.$$

Note that $|uv| = |ub| \leq |ua|$, because v lies on the circumference of D_u^b and a lies outside of this disk. This along with the latter equality above yields $\angle uav \leq \angle uva$. The sum of these two angles is $5\pi/6$ (recall that uv is the bisector of $\angle aub$), therefore $\angle uva \geq 5\pi/12$. Also note that $\angle uva < \pi/2$, because v lies strictly below the horizontal through a (otherwise d may not exist). It follows that $\sin \angle uva \geq \sin 5\pi/12$. Substituting this in the equality above yields $|av| \leq |ua| \sin \pi/6 / \sin 5\pi/12$. The law of sines applied to triangle $\triangle abu$ yields $|au| = |ab| \sin(\pi/3 + \gamma) / \sin \pi/3$, which substituted in the previous equality yields

$$|av| \leq |ab| \frac{\sin(\pi/3 + \gamma) \sin \pi/6}{\sin \pi/3 \sin 5\pi/12}.$$

Thus the inequality $|av| \leq |ai|$ holds for any γ satisfying

$$\frac{\sin(\pi/3 + \gamma) \sin \pi/6}{\sin \pi/3 \sin 5\pi/12} \leq \sin(\pi/3 - \gamma).$$

It can be easily verified that this inequality holds for any $\gamma \leq \delta \leq 23\pi/180$, and in particular for the δ values restricted by Lemma 11.

Condition (b) Assume now that z lies to the right of uv (as in Figure 24b). In this case $|bf| \leq |bz| \leq |bv|$. We now show that $|bv| \leq |bj|$, which implies $|bf| \leq |bj|$, thus

settling this case. From the right triangle $\triangle baj$ with angle $\angle baj = \pi/3 + \gamma$, we derive $|bj| = |ab| \sin(\pi/3 + \gamma)$. Next we derive an upper bound on $|bv|$. From the isosceles triangle $\triangle vub$, having angle $\angle vub = \pi/6$, we derive $|bv| = 2|bu| \sin \pi/12$. The law of sines applied to triangle $\triangle uab$ gives us $|ub| = |ab| \sin(\pi/3 - \gamma) / \sin \pi/3$, which substituted in the previous equality yields $|bv| = 2|ab| \sin(\pi/3 - \gamma) \sin \pi/12 / \sin \pi/3$. Thus the inequality $|bv| \leq |bj|$ holds for any γ value satisfying

$$\frac{2 \sin(\pi/3 - \gamma) \sin \pi/12}{\sin \pi/3} \leq \sin(\pi/3 + \gamma).$$

It can be verified that this inequality holds for any $\gamma \leq \delta \leq \pi/3$, and in particular for the δ values restricted by Lemma 11. \square

6 Conclusion

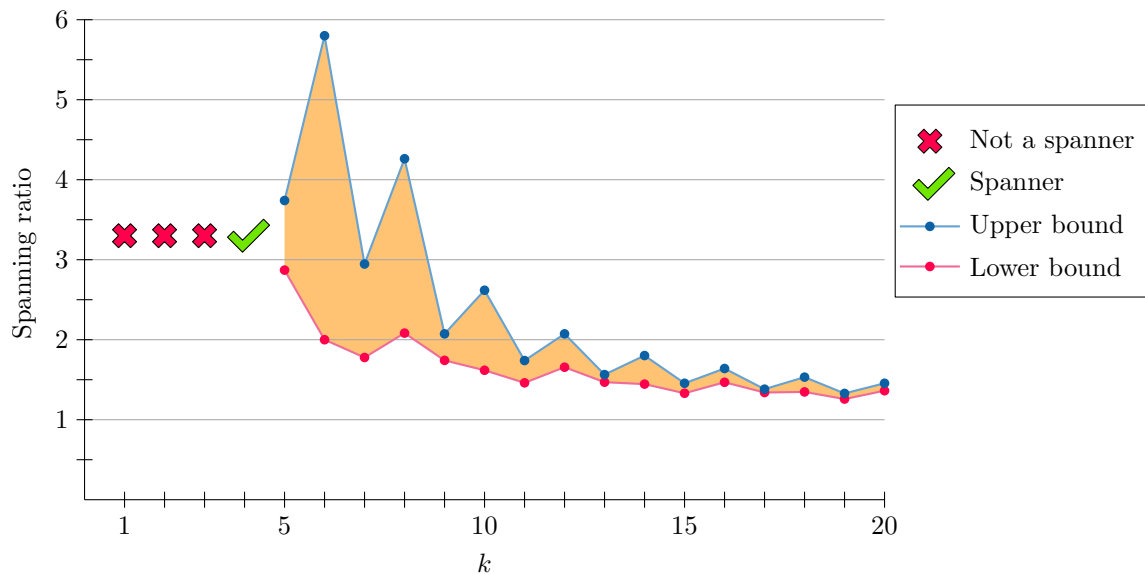


Figure 25: A plot showing the best known bounds on the spanning ratio of Yao graphs.

The main contributions of this paper are threefold. First we prove an upper bound of 3.74 on the spanning ratio of Y_5 . This settles the question of which Yao graphs are spanners, by establishing that Y_k is a spanner if and only if $k \geq 4$. In addition, the new upper bound falls just below the lower bound of 3.79 established for Θ_5 [7], making it the first separation between the spanning ratio of a Yao graph and its peer Θ -graph. For all other k it is still unknown which of Y_k or Θ_k has a better spanning ratio. By exploiting the asymmetry of Yao graphs with an odd number of cones, we also improve the upper bound on the spanning ratio of all graphs Y_k , for odd $k \geq 7$. The improvement is particularly significant for low values of k – for instance, we reduce the the upper bound on the spanning ratio of Y_7 from 7.562 to 2.946. Figure 25 and Table 1 summarize the currently known bounds on the spanning ratio of Yao graphs, including the results of this paper.

Second, we reduce the upper bound on the spanning ratio of Y_6 from 17.64 [11] to 5.8. We attribute this significant improvement to our direct approach of proving that Y_6 is a spanner, instead of proving that Y_6 spans the edges of another spanner. We believe that this upper bound can be reduced further. In fact, we conjecture that the spanning ratio of Y_6 matches the spanning ratio of 2 established for Θ_6 .

Finally, we present a collection of lower bounds, including the first lower bounds on the spanning ratio of Y_k , for $k \geq 6$. The gaps left between the lower and upper bounds are fairly small, especially for odd k . This is mostly due to our improved upper bound for odd Yao graphs, leading us to believe that there is still room to improve the upper bounds on the spanning ratios of Y_k for even k .

Our last lower bound also brings some light into the obscure world of Yao-Yao graphs, by proving that YY_5 is not a spanner. In this area, many interesting questions remain unanswered. For instance, even though Bauer and Damian [2] showed that there exist an infinite number of Yao-Yao graphs that are spanners, the first member of this class is YY_{36} , and nothing is known about the spanning properties of YY_7 through YY_{35} . In particular, it would be interesting to know whether, as with Yao and Θ -graphs, there is a constant k such that all Yao-Yao graphs with more than k cones are spanners, or whether there exist an infinite number of Yao-Yao graphs that are not spanners.

Table 1: Lower and upper bounds on the spanning ratio of Yao graphs.

k	Lower bound	Upper bound
$k = 2, 3$	∞ [12]	Not a spanner
$k = 4$	Open	$16(23 + 13\sqrt{2}) \approx 663$ [5]
$k = 5$	2.87	$2 + \sqrt{3} \approx 3.74$
$k = 6$	2	5.8
$k = 4x+2$	$1 + 2 \sin\left(\frac{\theta}{2}\right)$	$\frac{1}{1 - 2 \sin\left(\frac{\theta}{2}\right)}$ [4]
$k = 4x+3$	$1 + 2 \sin\left(\frac{3\theta}{8}\right) + 4 \frac{(\sin\left(\frac{13\theta}{16}\right) + \sin\left(\frac{19\theta}{16}\right)) \sin\left(\frac{\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right)}{\sin(2\theta)}$	$\frac{1}{1 - 2 \sin\left(\frac{3\theta}{8}\right)}$
$k = 4x+4$	$1 + 2 \sin\left(\frac{\theta}{2}\right) \left(1 + \tan\left(\frac{\theta}{2}\right)\right)$	$\frac{1}{1 - 2 \sin\left(\frac{\theta}{2}\right)}$ [4]
$k = 4x+5$	$1 + 2 \sin\left(\frac{3\theta}{8}\right) + 4 \sin\left(\frac{5\theta}{16}\right) \sin\left(\frac{3\theta}{8}\right)$	$\frac{1}{1 - 2 \sin\left(\frac{3\theta}{8}\right)}$

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A Lower bound coordinates

The following table lists the coordinates of the points in the Y_5 graph shown in Figure 7, whose spanning ratio is more than 2.87.

(0, 0)	<i>a</i>	(341, 264)
(252, 82)	<i>b</i>	(-179, 97)
(130, 230)	<i>c'</i>	(-180, 112)
(12, 193)	<i>d'</i>	(-91, -75)
(30, 302)		(316, 36)
(293, 269)		(352, 229)
(321, 229)		(303, 297)
(-143, 130)		(-167, 63)
(-143, 80)		(-167, 147)
(193, 384)		(-26, -75)
(158, 367)		(371, 213)
(-135, 272)		(51, 310)
(-91, 287)		(-176, 37)
(-153, -55)		(344, 274)
(371, 75)		(-189, 105)
(410, 115)		(99, 320)
(334, 276)		(-15, 284)

Table 2: Coordinates of the points in Figure 7