## POSSIBILITIES FOR IMPROVING THE EFFICIENCY OF LINEAR ERROR-CORRECTING CODES

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POSSIBILITIES FOR IMPROVING THE EFFICIENCY OF LINEAR ERROR-CORRECTING CODES<br>R. R. Varshamov<br>Computing Center of the Academy of Sciences of the Armenian SSR and Yerivan State University

The elaboration of design methods for the construction of noiseproof coding systems with a maximal data transmission rate (or equivalently with maximum efficiency of code subsets) when the main parameters of the system are specified is one of the main tasks of algebraic coding theory. Here $n$ is the length of a code block, $t$ is the number of detected errors, and $q$ is the base of the code. This note presents some results which are relevant to possibilities for constructing, within Hamming distance, noiseproof systems of signals in which the efficiency of codesubsets exceeds many times the efficiency of all similar known linear codes described in the literature.

The following theorem states our basic concept.
Theorem 1. Regardless of the magnitude of the number $\theta$ and method used to construct error-correcting codes correcting $t$ symmetric errors, values of the parameters $n$ and $q$ can always be found (on the basis of this method) for which the following is true:
a ) new cłasses of codes can be constructed in which the efficiency of code subsets is more than $\theta$ times greater than the efficiency of the corresponding codes that were constructed using the original method.
b ) the error-correcting capacity of the old codes is preserved.
We will assume that a method for constructing coding systems correcting a fixed number $t$ of errors is given if a completely defined method which allows the construction of such systems is specified for any arbitrary $n$ and $q$.
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The physical interpretation of Theorem 1 will be discussed using Hamming and Bose-Choudhuri-Hocquenghem codes as examples.

1. Hamming Codes. It is well known [1] that Hamming codes represent the class of linear codes correcting single symmetric errors. Although these codes are optimal in the class of linear codes, they belong to a class of codes for which a construction method has been developed. Therefore, according to the statement of theorem 1, new classes of codes allowing an improvement in the efficiency of Hamming codes (even without limit) can be constructed.

Henceforth we will denote by $g(n, q)$ the efficiency of the optimal Hamming code of length $n$ with base $q$. We will also use the notation $\varepsilon_{q}(n)=q^{\left[\log _{q} n\right]}$ where $[x]$ is the greatest integer not greater than $x$, and $\{\mathbf{x}\}=\mathbf{x}-[\mathrm{x}]$.

The case $q=3^{1}$
Theorem 2. For any integers $n$, satisfying the inequalities

$$
\log _{3} 3 / 2<\left\{\log _{3} n\right\}<2 \log _{3}(37 / 5)
$$

signal systems with base 3 correcting single symmetric errors exist in which the efficiency of code subsets $b(n, 3)$ is strictly greater than the efficiency of the corresponding Hamming codes and satisfies the inequality

$$
b(n, 3)>3^{1-\left(\log _{3}, 5_{n}\right\}} g(n, 3) .
$$

This result can be stated in somewhat stronger form, namely as the following theorem,

Theorem 3. Values of the parameter $n$ can always be found for which coding systems with base 3 correcting single errors can be constructed in whichthe efficiency of code subsets $b(n, 3)$ satisfies the inẹquality

$$
b(n, 3)>3^{1-\left(\log _{4}(\cdot n+o(n))\right]} g(n, 3),
$$

where $\delta(n)=1$ or 2 .
Thus, for example, the following theorem holds:

[^0]Theorem 4. Let $m=p_{1} \ldots p_{\sigma}, p_{s}, s=1, \ldots, \sigma$, be primes of the form $8 k+5$ and $g_{s}$ be a primitive root modulo $p_{s}$,

$$
\begin{gathered}
Q_{s}=p_{1} \ldots p_{s}\left(Q_{0}=1\right), \quad m_{s}=Q_{s}^{-1} m, \quad M_{s}=\frac{m-m_{s}}{4}, \\
p_{s}^{\prime}=\frac{p_{s}-1}{4}, \quad\left\{\log _{3} m\right\}>\log _{3} 2 .
\end{gathered}
$$

Then for any integers $\alpha$ and $n\left(\varepsilon_{3}(m)<2 n<m / 2\right)$, the set of a11 possible solutions of the congruence
were $x_{u}, u n^{2}$, taking on arbitrarily the values 0,1 and 2 , is a code of length $n$ with base 3 , correcting single symmetric errors. Among these codes the efficiency of the optimal code is

$$
b(n, 3)>3^{1-\left\{\log _{3} m\right)} g(n, 3), \quad \delta(n)=1 .
$$

The following theorem is also valid.
Theorem 5. Let $m=p_{1} \ldots p_{\sigma}$, where $p_{s}$ are primes of the form $8 k+3, m>3(2 m)$ and $q=3$.

Then for arbitrary integers $\alpha, \beta$ and $n \quad\left(\varepsilon_{3}(2 n)<2 n<m\right)$, the set of all possible solutions of the system of congruences

$$
\begin{gathered}
\sum_{\omega}^{\omega} \sum_{s=1}^{\sigma} \sum_{v=0}^{m_{s}-1} \sum_{u=1}^{2 p_{s}^{\prime}} Q_{s-1}\left(g_{s}^{2 u}+p_{s} v\right) x_{2 M}+2 p_{s}^{\prime} v+u=\alpha(\bmod m), \\
\sum_{u=1}^{n} x_{u}=\beta\left(\bmod _{2}\right) .
\end{gathered}
$$

where $x_{u}=0,1$ and $2(u \leq n)$ and $x_{u}=0(u>n)$, is a code of length $n$ correcting single symmetric errors. Among these codes, the efficiency of the optimal code is

$$
b(n, 3)>3^{1-(\log , 2 m)} g(n, 3), \quad \delta(n)=2
$$

We will now consider the general case of arbitrary $q=p_{0}^{\nu} 3$

[^1]Theorem $6 \quad(\nu>1)$. For any integer values $n$, satisfying the inequalities

$$
\log _{8}(q-1)^{-1} q \leqslant\left\{\log _{9} n\right\} \leqslant 1 / v
$$

signal systems with base $q$ correcting single symmetric errors exist in which the efficiency of code subsets $b(n, q)$ satisfies the relation

$$
b(n, q)=p^{v-1} g(n, q)
$$

Theorem $7(v=1)$. For any integer values $n$ satisfying the inequalities

$$
\log _{p}(p-1)^{-1} p \leqslant\left\{\log _{p} n\right\} \leqslant \log _{p} p / \omega(n), \quad 2 \leqslant \omega(n)<4
$$

signal systems with base $p$ exist in which the efficiency of code subsets $b(n, p)$ satisfies the inequality

$$
b(n, p)>p^{1-\| \ln _{p} g_{p} \omega(n) \prime} g(n, p)
$$

This result can be stated in somewhat stronger form.
Theorem 8. Values of $n$ exist such that codes with base $p$ correcting single symmetric errors can be constructed in which the efficiency of code subsets satisfies the inequality

$$
b(n, p)>p^{1-\left(\log _{p}{ }^{2 n)} g(n, p) .\right.}
$$

2. Bose-Chandhuri-Hocquenghem (B. C. H.) codes. B. C. H. codes are a generalization of Hamming codes to the multi-error correcting case, and like Hamming codes, they belong to a class of codes for which a construction method exists. Therefore, according to our concept, new more efficient codes correcting multiple errors can be constructed on the basis of these codes.

Let us denote by $G_{t}(n, q)$ the efficiency of optimal B. C. H. codes of length $n$ with base $q$ correcting $t$ symmetric errors.

The case $q=2, t=2$.

Theorem 9. For any integers $n\left(\left\{\log _{2} n\right\}<\log _{2}(2 / \sqrt{3})\right)$ and $q=2$, codes correcting binary symmetric errors can be constructed whose efficiency $B_{2}(n, 2)$ is strictly greater than the efficiency of the corresponding B. C. H. codes and satisfies the inequality

$$
\begin{equation*}
B_{2}(n, 2)>4^{1-\left(\log _{2} n \sqrt{3}\right)} G_{2}(n, 2) . \tag{1}
\end{equation*}
$$

Thus, for example, the following theorem holds.
Theorem 10. Let $p$ be a prime of the form $6 k-1$ satisfying the inequality $\left\{\log _{2} p\right\}<\log _{2}(2 / \sqrt{3})$.

Then for any integers $\alpha, \beta, \gamma u$ and $n\left(\varepsilon_{2}(p) \leq n \leq p\right)$, the set of all possible solutions of the system of congruences

$$
\begin{aligned}
& \sum_{u=1}^{n} u x_{u}=\alpha(\bmod p), \\
& \sum_{u=1}^{n} u^{2} x_{u}=\beta(\bmod p), \\
& \sum_{u=1}^{n} x_{u}=\gamma(\bmod 3),
\end{aligned}
$$

where $x_{u}=0$ or 1 , is a code of length $n$ with base 2 correcting two symmetric errors. Among these codes, the efficiency of the optimal codes is strictly greater than the efficiency of analogous B. C. H. codes, and it satisfies inequality 1 .

The case $q=3, t=2$.
Theorem 11. Let $n$ be any integer satisfying the inequality $\left\{\log _{3} n\right\}<\log _{3} 3 / 2$ and $g$ be a primitive element of the Galois field $\operatorname{GF}\left(\varepsilon_{3}(3 n)\right)$.

Then the set of all possible solutions of the equations

$$
\sum_{u=1}^{n} g^{-u} x_{u}=a, \quad \sum_{u=1}^{n} g^{u} x_{u}=b
$$

where $x_{u} \notin G F(3)$ and $a, b \notin G F\left(\varepsilon_{3}(3 n)\right)$, is a code of length $n$ with
base 3 , correcting two symmetric errors, whose efficiency $B_{2}(n, 3)$ is

$$
B_{2}(n, 3)=\left({\sqrt{3})^{1-(-1)}}^{(\log , n)} G_{2}(n, 3) .\right.
$$

The case $\mathrm{q}=4, \mathrm{t}=2$.
Theorem 12. Suppose an arbitrary natural number $n>5$ and a primitive element $g$ of the Galois field $\operatorname{GF}\left(\varepsilon_{4}(12 n)\right)$, are given.

Then the set of all possible solutions of the equations

$$
\sum_{u=1}^{n} g^{-u} x_{u}=a, \quad \sum_{u=1}^{n} g^{n} x_{u}=b
$$

where $x_{u} \notin \operatorname{GF}\binom{4}{x}$ and $a, b, \xi \in\left(\varepsilon_{4}(12 n)\right)$ is a code of length $n$ with base 4 correcting two symmetric errors, whose efficiency strictly exceeds the efficiency of the analogous B. C. H. code and satisfies the relation

$$
B_{2}(n, 4)=4^{n(n)} G_{2}(n, 4),
$$

where

$$
\left.h(n)=\left[\log _{4} n\right]-(-1)^{[4 \varepsilon} 4(n) / 3 n\right]
$$

However, on the basis of only the Bose-Chaudhuri-Hocquenghem results, one can also prove the following theorem.

Case of arbitrary p.
Theorem 13. For any arbitrary values $t \geq 2, p \geq 2 t$ and $n$ $\left(\left\{\log _{p} n\right\}<\log _{p}(p / \nu(n)), \quad 1 \leq \nu(n)<2\right)$, codes of length $n$ with base p, correcting $t$ symmetric errors can be constructed, whose efficiency $B_{t}(n, p)$ is strictly greater than the efficiency of the corresponding B. C. H. codes correcting the same number of errors, and whose efficiency satisfies the inequality

$$
B_{l}(n, p)>\left(p^{t-\left(\log _{p} v(n) n\right)}\right)^{2 t-1} G_{t}(n, p) .
$$

Note that when $n$ approaches infinity, the quantity $\nu(n)$ takes on the value 1 infinitely many times. Hence the following theorem holds.

Theorem 14. For fixed $t$ and $p \geq 2 t$, the inequality

$$
B_{l}(n, p)>\left(p^{1-\left(\log _{p} n^{\prime}\right)}\right)^{2 t-1} G_{l}(n, p)
$$

is satisfied infinitely many times as $n$ approaches infinity.

## REFERENCES

Peterson, W. Kody, ispravlyayushchiye oshibki (Error-correcting codes). Mir Press, Moscow, 1964.

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STANDARD TITLE PAGE



[^0]:    ${ }^{1}$ Our method is inefficient in the case $q=2$

[^1]:    $2 \quad x_{u}=0$ for all $u>n$
    3 p is an odd prime

