OBJECT-ORIENTED PROGRAMMING WITH MIXING IN AUA

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INTRODUCTION

My guess is that object-oriented programming will be in the 1980s what structured programming was in the 1970s. Everyone will be in favor of it. Every manufacturer will promote his products as supporting it. Every manager will pay lip service to it. Every programmer will practice it. And no one will know just what it is.

[Rentsch 82]

Recently, I wrote a paper discussing the lack of "true" object-oriented programming language features in Ada 83, why one might desire them in Ada and how they might be added in Ada 9X [Seidewitz 91]. The approach I took in this paper was to build the new object-oriented features of Ada 9X as much as possible on the basic constructs and philosophy of Ada 83. The object-oriented features proposed for Ada 9X [Ada9X 91b], while different in detail, are based on the same kind of approach.

Further consideration of this approach led me on a long reflection on the nature of object-oriented programming and its application to Ada. The results of this reflection, presented in this paper, show how a fairly natural objectoriented style can indeed be developed even in Ada 83. The exercise of developing this style is useful for at least three reasons:

- 1. It provides a useful style for programming object-oriented applications in Ada 83 until new features become available with Ada 9X;
- 2. It demystifies many of the mechanisms that seem to be "magic" in most object-oriented programming languages by making them explicit; and
- 3. It points out areas that are and are not in need of change in Ada 83 to make object-oriented programming more natural in Ada 9X.

In the next four sections I will address in turn the issues of object-oriented classes, mixins, self-reference and supertyping. The presentation is through a sequence of examples, similar to those in [Seidewitz 91]. This results in some overlap with that paper, but all the examples in the present paper are written entirely in Ada 83. I will return to considerations for Ada 9X in the last section of the paper.

CLASSES

An object represents a component of...[a] software system... An object consists of some private memory and a set of operations...A crucial property of an object is that its private memory can be manipulated only by its operations...A class describes the implementation of a set of objects that all represent the same kind of system component.

[Goldberg and Robson 83]

In Ada, an object is a variable or a constant that contains a value. The declared type of the object determines the set of possible values for the object and the set of operations that may be applied to the object. If this type is a private type, then the value of the object may only be changed through application of an operation. This corresponds to the object-oriented notion of a *class*.

Consider, for example, a simple financial account class implemented as a private type:

with Finance_Types; use Finance_Types; package Finance is

type ACCOUNT is limited private;

procedure Open	(The_Account With_Balance	:	in in	out	ACCOUNT; MONEY);
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procedure	Deposit		: in out ACCOUNT; : in MONEY);
procedure	Withdraw		: in out ACCOUNT; : in MONEY);
function	Balance_Of	(The_Account	: ACCOUNT) return MONEY;
private			

```
type ACCOUNT is
record
Balance : MONEY := 0.00;
```

end Finance;

end record;

The type Finance. ACCOUNT represents a class of account objects. The subprograms defined in package Finance are the allowable operations on objects of this class. The body of this package is straightforward. Note that for simplicity I will assume that a number of simple types (such as MONEY) are defined in a Finance_Types package.

The class defined by package Finance provides a simple but very general abstraction. In an object-oriented approach, such general classes are used as the basis for implementing more specialized classes. For example, a savings account is a specific kind of account that holds savings that earn interest. Other than some new operations associated with earning interest, a savings account is the same as the original general financial account. Thus we should be able to implement a savings account in terms of a general account:

```
with Finance_Types; use Finance_Types;
package Savings is
```

type ACCOUNT is limited private;

procedure	Open	(The_Account With_Balance	:	in out in	ACCOUNT; MONEY);
procedure	Set_Rate	(Of_Account To_Rate		in out in	ACCOUNT; RATE);
procedure	Deposit	(Into_Account The_Amount			
procedure	Withdraw	(From_Account The_Amount			
procedure	Earn_Interest	(On_Account Over_Time			ACCOUNT; INTERVAL);
function	Balance_Of	(The_Account	:	ACCOUN	() return MONEY;
function	Interest_On	(The_Account	:	ACCOUN	r) return MONEY;

private

```
type ACCOUNT is
  record
  Parent : Finance.ACCOUNT;
  Rate : RATE := 0.06;
  Interest : MONEY := 0.00;
  end record;
```

end Savings;

While this may not seem to gain us a lot in this simple example, such incremental extension of abstractions is fundamental to object-oriented techniques. The class of financial accounts is said to be the *superclass* of the class of savings accounts. Each savings account (of type Savings. ACCOUNT) has a unique *parent* financial account (of type Finance. ACCOUNT).

Now, three of the seven savings account operations (Open, Deposit and Withdraw) are syntactically and semantically the same as the corresponding financial account operations. Thus, we would like to *inherit* these financial account operations. Ada 83 has no direct way of doing this. Nevertheless, we can achieve the effect of inheritance for our present purposes by using *call-through* subprograms. For example, the Savings.Deposit

operation can easily be implemented as follows:

The expense of such call-throughs may be minimized through the use of pragma Inline.

Three other savings account operations (Set_Rate, Earn_Interest and Interest_On) provide the incremental new functionality of the savings account *subclass*. These operations are implemented in terms of the additional components of the representation of type Savings.ACCOUNT. For example:

procedure Earn_Interest (On_Account : in out ACCOUNT; Over_Time : in INTERVAL) is

Balance : constant MONEY := Balance_Of (On_Account);

begin

end Earn Interest;

Note that the Balance_Of operation used here is the subclass operation Savings.Balance_Of.

The remaining savings account operation, Balance_Of, is syntactically the same as the financial account operation, but it is semantically different. The balance of a savings account includes interest earned up to the present point in time:

```
function Balance_Of (The_Account : ACCOUNT) return MONEY is
begin
return Finance.Balance_Of (The_Account.Parent)
+ The Account.Interest;
```

-

end Balance_Of;

Note that while Balance_Of is not a call-through operation, the superclass operation Finance.Balance_Of is used in its implementation.

The usefulness of a superclass like the financial account class comes from the fact that it can provide a common basis for a *number* of subclasses. For example, a class of checking accounts may provide another subclass of financial accounts:

```
with Finance;
with Finance_Types; use Finance_Types;
package Checking is
```

type ACCOUNT is limited private;

procedure	Open	(The_Account With_Balance			ACCOUNT; MONEY);
procedure	Set_Fee	(Of_Account To_Fee		in out in	ACCOUNT; MONEY);
procedure	Deposit	(Into_Account The_Amount	::	in out in	ACCOUNT; MONEY);
procedure	Withdraw	(From_Account The_Amount			ACCOUNT; MONEY);
function	Balance_Of	(The_Account	:	ACCOUN	() return MONEY;

private

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```
type ACCOUNT is
  record
  Parent : Finance.ACCOUNT;
  Overdraft_Fee : MONEY := 10.00;
  end record;
```

end Checking;

In this simple example, the only difference between checking accounts and financial accounts is that overdrawing a checking account is not permitted. Further, each overdraft attempt (i.e., a returned check) is penalized by deducting a fee from the account. Thus, the implementation of Withdraw must be changed for checking accounts:

end if;

end Withdraw;

The savings account and checking account subclasses of the financial account class may themselves act as superclasses for even more specialized classes. Thus, a general class may be the root of a quite extended class hierarchy. Each subclass in the hierarchy incrementally extends the capabilities of its superclass, while inheriting common functionality.

MIXINS

A mixin is...a subclass definition that may be applied to different superclasses to create a related family of modified classes.

[Bracha and Cook 90]

A superclass may be used as the base for many subclasses. However, as described so far, a subclass is tied to one superclass. For instance, savings accounts are based specifically on the class defined by package Finance. There may be other types of accounts to which we want to added interest-bearing functionality such as that defined for savings accounts. For example, an interest-bearing checking account is basically a checking account with interest-bearing functionality added to it (or, alternatively, a savings account with checking functionality added).

Rather than recoding essentially the same interest-bearing functionality each time it is needed, we can capture this functionality in a generic package that takes a specific superclass as a parameter:

```
with Finance Types; use Finance_Types;
generic
  type SUPERCLASS is limited private;
  with function Balance_Of (The_Account : SUPERCLASS) return MONEY is <>;
package Interest is
  type MIXIN is limited private;
  type ACCOUNT is
    record
      Parent
                   : SUPERCLASS;
                   : MIXIN;
      Extension
    end record;
                          (Of Account : in out ACCOUNT;
 procedure Set Rate
                                       : in
                                                RATE);
                           To Rate
 procedure Earn_Interest (On_Account : in out ACCOUNT;
                           Over Time
                                       : in
                                                INTERVAL);
                          (The Account : ACCOUNT) return MONEY;
  function Balance Of
                                       6-6
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```

function Interest_On (The_Account : ACCOUNT) return MONEY;

```
private
```

```
type MIXIN is
  record
   Rate : RATE := 0.06;
   Interest : MONEY := 0.00;
  end record;
```

end Interest;

A generic package such as this is called a *mixin* because it provides an increment of functionality which may be "mixed-into" any superclass that has the operations required to fill in the generic parameters. Typically, mixins are used within a framework of *multiple inheritance*. For example, we can reconstruct the savings account class by inheriting from *both* the financial account class and an appropriate instantiation of the interest mixin:

```
with Finance, Interest;
with Finance_Types; use Finance_Types;
package Savings is
type ACCOUNT is limited private;
procedure Open (The_Account : in out ACCOUNT;
With_Balance : in MONEY);
...
function Interest_On (The_Account : ACCOUNT) return MONEY;
```

private

```
package Savings_Interest is
    new Interest (Finance.ACCOUNT, Finance.Balance_Of);
```

type ACCOUNT is new Savings_Interest.ACCOUNT;

end Savings;

The Parent component of the ACCOUNT type defined in mixin Interest is used to inherit from the parent superclass Finance via a call-through. For example:

The record type Savings_Interest.ACCOUNT is defined as a visible, rather than a private, type in the mixin to allow access to the Parent component. Note that it would not be possible to replace this use of a visible record component with a function that returns the parent object, because we need to use the parent as an in out parameter. The type MIXIN is never used itself outside of the mixin package.

Call-through subprograms are also needed to inherit from the mixin instantiation Savings_Interest. This is because the equivalent derived subprograms obtained from the derived type definition of Savings.ACCOUNT are hidden by the operations declared in the package specification, and in Ada 83 there must be a full subprogram body for each of these declarations. For example:

```
function Interest_On (The_Account : ACCOUNT) return MONEY is
begin
    return Savings_Interest.Interest_On
        (Savings_Interest.ACCOUNT(The_Account));
end Interest_On;
```

Having introduced the concept of mixins, we can, of course, also create a mixin that embodies the overdraft functionality of the checking account class:

```
with Finance Types; use Finance Types;
generic
  type SUPERCLASS is limited private;
  with procedure Withdraw
                            (From Account : in out SUPERCLASS;
                             The Amount : in MONEY) is <>;
  with function Balance_Of (The_Account : SUPERCLASS) return MONEY is <>;
package Draft is
  type MIXIN is limited private;
  type ACCOUNT is
    record
                    : SUPERCLASS;
      Parent
                    : MIXIN;
      Extension
    end record;
                     (Of Account : in out ACCOUNT;
  procedure Set Fee
                                   : in
                      To_Fee
                                            MONEY);
  procedure Withdraw (From Account : in out ACCOUNT;
                      The Amount
                                  : in
                                            MONEY);
private
  type MIXIN is
    record
      Overdraft_Fee : MONEY := 10.00;
    end record;
and Draft;
   Even our original financial account class can be converted to a mixin:
with Finance_Types; use Finance_Types;
generic
  type SUPERCLASS is limited private;
package Monetary is
  type MIXIN is limited private;
  type ACCOUNT is
    record
      Parent
                 : SUPERCLASS;
                  : MIXIN;
      Extension
    end record;
  procedure Open
                        (The Account : in out ACCOUNT;
                         With Balance : in MONEY);
  function Balance Of (The Account : ACCOUNT) return MONEY;
private
  type MIXIN is
    record
                  : MONEY := 0.00;
      Balance
    end record;
end Monetary;
```

Of course, this mixin does not require any superclass functionality to implement its operations. However, use of the mixin construct allows monetary account functionality to be mixed into any class.

The use of mixins causes traditional class hierarchies to collapse into pieces. Each piece is a mixin that provides a well-defined increment of functionality. We can then form specific classes from these pieces by instantiating a number of mixins and inheriting all necessary functionality from them. To provide a definite starting

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point for this process, we can define a root class that basically does nothing more than provide an empty record to which we can add mixins:

package Root is
type CLASS is limited private;

private

```
type CLASS is
record
null; _
end record;
```

end Root;

While this root class seems a bit pointless, the concept will prove useful in the next section.

At last we are ready to construct an interest-bearing checking account class without rewriting any savings account or checking account functionality. To do this, we simply mix together interest, draft and monetary account functionality. All Interest Bearing Checking.ACCOUNT operations are implemented as call-throughs to various mixin operations. Thus, from the three mixins Monetary, Interest and Draft, we can easily construct an interest-bearing checking account class, as well as reconstructing our original financial, savings and checking account classes.

Of course, in the actual Interest Bearing_Checking package, the three mixin generics must be instantiated in a specific sequential order. In the present case, we must first establish the basic monetary account functionality, then mix in interest and draft functionality. This results in the following implementation:

```
with Root, Monetary, Interest, Draft;
with Finance_Types; use Finance_Types;
package Interest_Bearing_Checking is
```

type ACCOUNT is limited private;

procedure Open (The Account : in out ACCOUNT; With Balance : in MONEY);

•••

function Interest_On (The_Account : ACCOUNT) return MONEY;

private

end Interest_Bearing_Checking;

Note that all the mixins are instantiated in the private part of the specification. Each instantiation uses the type and subprograms from the previous instantiation as arguments. The intermediate procedure Withdraw for type Checking_Interest.ACCOUNT is necessary because the instantiated mixin Checking_Interest only provides the interest-related operations on Checking_Interest.ACCOUNT. It is implemented as simply a call-through to Finance.Withdraw.

When an object of a given class is created its state components include those of the class and all its superclasses and it can perform operations of the class and its superclasses on the component state. References to "self" in operations of a superclass refer to the composite object on behalf of which the operation is to be performed.

[Wegner 87]

In the interest-bearing checking account package at the end of the last section, the Interest mixin was instantiated before the Draft mixin. It would seem that we could equally well have instantiated them in the opposite order:

with Root, Monetary, Interest, Draft; with Finance_Types; use Finance_Types; package Interest_Bearing_Checking is

type ACCOUNT is limited private;

procedure	Open	(The A	ccount :	in	out	ACCOUNT;
•	•	With_	Balance :	in		MONEY);

•••

function Interest_On (The_Account : ACCOUNT) return MONEY;

private

end Interest_Bearing_Checking;

Unfortunately, it turns out that this introduces a subtle error, as follows:

- In the new implementation, the Draft mixin is instantiated before the Interest mixin, using the Checking_Finance.Balance_Of operation.
- The implementation of the Withdraw operation in the Draft mixin uses the Balance_Of operation given as a generic formal superclass operation to determine if there is an overdraft. In this case, the actual subprogram used is Checking_Finance.Balance_Of, which does not add in any earned interest.
- The Interest Bearing Checking. Withdraw operation is inherited from the instantiation Checking Draft of the Draft mixin, so as to include the overdraft functionality. This means that accumulated interest is ignored when checking for an overdraft. This is clearly unfair to the customer!

The problem is that we do not really want to use the *superclass* Balance_Of operation in the Draft mixin instantiation. Rather, we need to use the Balance_Of operation from the composite *subclass* being constructed. However, we cannot use the subclass type Interest_Bearing_Checking.ACCOUNT in the instantiation of the Draft mixin, because that type cannot be fully defined yet. Thus, we must instead be sure to instantiate the Interest mixin first, so that the interest-bearing functionality is mixed into the Balance_Of operation before Draft is instantiated.

Such order dependencies are at best annoying sources of potential errors. At worst, they can introduce circular dependencies that make it impossible to mix together certain mixins. To avoid this, we need a mechanism that allows mixins to call subclass operations in addition to superclass operations. Following the parameterization approach that led us to mixins in the first place, we can include a second generic formal type parameter in mixins to

represent the subclass.

For example, we want the Draft mixin to use the subclass Balance_Of operation:

```
with Finance_Types; use Finance_Types;
generic
  type SUPERCLASS is limited private;
  with procedure Withdraw (From_Account : in out SUPERCLASS;
                             The Amount : in MONEY) is <>;
  type SUBCLASS is limited private;
  with function Balance_Of (The_Account : SUBCLASS) return MONEY is <>;
                                         : SUPERCLASS) return SUBCLASS is <>;
                            (Parent
  with function Self
package Draft is
  type MIXIN is limited private;
  type ACCOUNT is
    record
                     : SUPERCLASS;
      Parent
      Extension
                    : MIXIN;
    end record;
  procedure Set_Fee (Of_Account : in out ACCOUNT;
                                             MONEY);
                       To_Fee
                                    : in
  procedure Withdraw (From_Account : in out ACCOUNT;
                       The Amount : in
                                             MONEY);
                    (This_Account : ACCOUNT) return SUBCLASS;
  function Self
private
  type MIXIN is
    record
      Overdraft_Fee : MONEY := 10.00;
    end record;
end Draft;
   The Withdraw operation for this mixin is then implemented as follows:
                                                                         ~ ~
procedure Withdraw (From_Account : in out ACCOUNT;
                                            MONEY) is
                     The Amount : in
begin
  if The Amount <= Balance_Of(Self(From_Amount)) then
    Finance.Withdraw (From_Account.Parent, The_Amount);
  else
    Finance.Withdraw (From_Account.Parent,
                       From Account.Extension.Overdraft_Fee);
  end if;
end Withdraw;
Note the use of the function Self to convert an object of type Draft. ACCOUNT to the appropriate object of type
SUBCLASS. These odd little Self functions are the key to this approach. They allow us to use the subclass
```

operations as required. The question is, of course, how can we implement such a Self function? Strangely enough, we can implement

it in terms of the superclass Self function given as a generic formal parameter:

```
function Self (This_Account : ACCOUNT) return SUBCLASS is
begin
return Self (This_Account.Parent);
end Self;
```

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Obviously, this passing of the buck must end someplace. It ends with the root class, which we reimplement as follows:

generic

type SUBCLASS is limited private;

package Root is

type CLASS is limited private;

procedure Initialize (The Object : **in out** CLASS; To_Self : **in** SUBCLASS);

function Self (This_Object : CLASS) return SUBCLASS;

private

```
type CLASS is
  record
    Self : SUBCLASS;
  end record;
```

end Root;

Thus the mystery is resolved: the Self functions all ultimately access a Self component defined in the root class.

Now, the astute reader may have noticed that we have introduced a strange sort of circularity here. The representation of any class built on the root class will include a component of the subclass type. However, when we finish constructing a class from the root class and mixins, the result is the very subclass with which we need to instantiate the root class to begin with! To achieve this circularity, we must require that the subclass type be an access type. The Self component is then intended to be a *pointer* back to the complete, composite subclass object. (Actually, access types are also needed to allow the Self functions to work properly with subclass procedures that would otherwise have in out parameters.)

With inclusion of subclass parameters in mixins, we can now correctly implement the interest-bearing checking account class using either order of mixin instantiation:

with Finance_Types; use Finance_Types; package Interest_Bearing_Checking is

type ACCOUNT is limited private;

procedure	Open	(The_Account With_Balance				ACCOUNT; MONEY);
procedure	Close	(The_Account	:	in	out	ACCOUNT);
procedure	Set_Rate	(Of_Account To_Rate	-	in in		ACCOUNT; RATE);

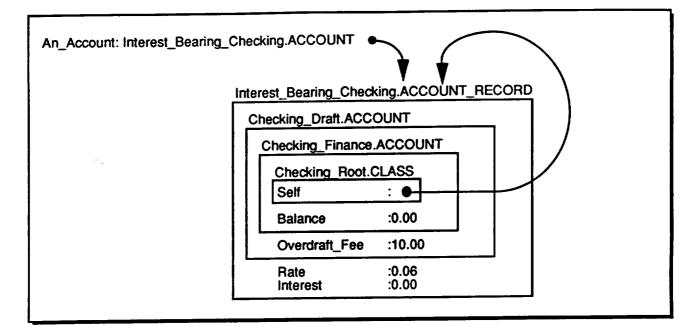
private

type ACCOUNT_RECORD; type ACCOUNT is access ACCOUNT_RECORD;

end Interest_Bearing_Checking;

An advantage of implementing a private type as an access type is that the details of the type representation can be deferred to the package body by using an incomplete type definition for ACCOUNT_RECORD in the private part of the specification. The use of an access type also allows the use of in rather than in out parameters in procedures such as Set_Rate, which is necessary for the use of Self functions.

Circular type definition is also achieved using the incomplete type definition for ACCOUNT_RECORD. The circle is closed by completing the definition of ACCOUNT_RECORD after all the mixin instantiations in the package body. The figure on the next page shows the structure of an Interest_Bearing_Checking.ACCOUNT object resulting from the following implementation:



```
with Root, Monetary, Interest, Draft;
package body Interest_Bearing_Checking is
  package Checking_Root is
    new Root (SUBCLASS => Interest_Bearing_Checking.ACCOUNT);
  use Checking Root;
  package Checking_Finance is
    new Monetary
                     => Checking_Root.CLASS,
      (SUPERCLASS
                    => Interest_Bearing_Checking.ACCOUNT);
       SUBCLASS
  use Checking_Finance;
  package Checking_Draft is
    new Draft
                    => Checking_Finance.ACCOUNT,
      (SUPERCLASS
                    => Interest_Bearing_Checking.ACCOUNT);
       SUBCLASS
  function Balance_Of (The_Account : in Checking_Draft.ACCOUNT)
    return MONEY;
  -- call-through to Finance.Balance_Of
  use Checking Draft;
  package Checking_Interest is
    new Interest
      (SUPERCLASS
                     => Checking Draft.ACCOUNT,
                     => Interest_Bearing_Checking.ACCOUNT);
       SUBCLASS
  type ACCOUNT_RECORD is new Checking_Interest.ACCOUNT;
```

end Interest Bearing Checking;

(Note that to simplify the instantiations, I have taken advantage of the box defaults on the generic formal subprogram parameters of the mixins.)

A disadvantage of using an access type is that interest-bearing checking accounts must be explicitly allocated. We can do this as part of the Open operation:

end Open;

Note the use of the root Initialize operation to set the Self component. The figure on the previous page shows the structure of nested records and self reference that results from the allocation and initialization of an Interest_Bearing_Checking.ACCOUNT object.

We also need to provide a way to deallocate interest-bearing checking accounts:

procedure Free is new Unchecked_Deallocation (ACCOUNT_RECORD, ACCOUNT);

procedure Close (The_Account : in out ACCOUNT) is begin Free(The_Account); end Close;

All the rest of the interest-bearing checking account operations are inherited from one or the other of the mixin instantiations.

SUPERTYPES

Subtyping is a substitutability relationship, i.e., an instance of a subtype can stand in for an instance of its supertype. How the subtype is implemented is totally irrelevant; all that matters is that it have the right behavior so that it can be substituted.

[Lalonde and Pugh 91]

Typically, the customer of a bank will have several accounts at that bank. Each bank account may be, say, a savings account, a checking account or an interest-bearing checking account. To manage all the bank accounts of one customer, we would like to create a bank account type that is the *supertype* of the types that represent the various classes of accounts. We could then create lists of bank accounts, define bank account operations, etc.

As discussed in the previous sections, each class is implemented in Ada by a private type that is distinct from all other class types. Nevertheless, we can still explicitly create a bank account supertype:

```
with Savings, Checking, Interest_Bearing_Checking;
with Finance_Types; use Finance_Types;
package Bank is
  type ACCOUNT_TYPE is (SAVINGS, CHECKING, INTEREST_CHECKING);
  type ACCOUNT (Kind : ACCOUNT TYPE := SAVINGS) is
    record
      case Kind is
                               => A_Savings_Account : Savings.ACCOUNT;
        when SAVINGS
                               => A_Checking_Account : Checking.ACCOUNT;
        when CHECKING
        when INTEREST_CHECKING => An_Interest_Checking_Account
                                     : Interest Bearing Checking.ACCOUNT;
      end case;
    end record;
                       (The Account : in out ACCOUNT;
 procedure Open
                       With Balance : in
                                             MONEY);
                       (The Account : in out ACCOUNT);
 procedure Close
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```

```
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```

procedure	Deposit	(Into_Account The_Amount					۰.
procedure	Withdraw	(From_Account The_Amount	:	in ou in	ut	ACCOUNT; MONEY);	
function	Balance_Of	(The_Account	:	ACCOU	UNT) return	MONEY;

end Bank;

The type Bank.ACCOUNT defines a supertype with subtypes Bank.ACCOUNT (SAVINGS), Bank.ACCOUNT (CHECKING) and Bank.ACCOUNT (INTEREST_CHECKING). Each subtype corresponds to one of the classes defined in previous sections. Note that a private type is unnecessary here, because we wish to be able to freely convert between the Bank.ACCOUNT subtypes and the class types.

The five operations defined in package Bank reflect the operations that are common to all the account types. Semantically, we wish each supertype operation to mirror the implementation of the appropriate subtype operation. For example, the statement:

Bank.Withdraw (From_Account => A, The_Amount => X);

should be equivalent to either Savings.Withdraw, Checking.Withdraw or Interest Bearing Checking.Withdraw, depending on the subtype of A. Since the subtype of A can, in general, only be determined at run-time, we are effectively asking that Bank.Withdraw be dynamically bound to the appropriate subtype operation.

We can achieve the effect of dynamic binding in Ada by implementing the bank account operations as dispatching or case-selection subprograms. For example:

procedure Withdraw (From Account : in out ACCOUNT; The Amount : in MONEY) is

begin

```
case Kind is
when SAVINGS =>
Savings.Withdraw (From_Account.A_Finance_Account, The_Amount);
when CKECKING =>
Checking.Withdraw (From_Account.A_Checking_Account, The_Amount);
when INTEREST_CHECKING =>
Interest_Bearing_Checking.Withdraw
(From_Account.An_Interest_Checking_Account, The_Amount);
```

end case;

end Withdraw;

Once we have the bank account supertype, we can create *polymorphic* data structures and operations that can handle all kinds of bank accounts. For example:

type CUSTOMER ACCOUNTS is array(POSITIVE range <>) of Bank.ACCOUNT;

function Total_Assets_Of (The_Accounts : CUSTOMER_ACCOUNTS) return MONEY is

Total : MONEY := 0.00;

begin

```
for I in The_Accounts'range loop
   Total := Total + Bank.Balance_Of (The_Accounts(I));
end loop;
return Total;
```

----·

end Total_Assets_Of;

The function defined above finds the total assets a customer has in his accounts, regardless of what kinds of accounts they are.

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It is important to note that to be included in a supertype, a class need only provide implementations for all the operations defined for the supertype. The ways in which various subtype classes implement these operations do not have to be related at all. For example, the Bank. ACCOUNT supertype is constructed from a number of classes implemented by various combinations of the mixins Monetary, Interest and Draft. These classes thus share some common implementation, but this is not at all important to the construction of the supertype.

Thus, supertypes and superclasses are really distinct concepts. Looking at it another way, the supertype provides a set of dispatching operations for those operations which are common to all its subtypes, regardless of how those operations may be implemented by the subtype classes or how the subtypes may be represented. A supertype that is constructed in this way from a given list of subtype classes is said to be the *union type* of those classes. Thus we have constructed a bank account supertype that is the union of the savings, checking and interest-bearing checking account classes.

It was noted earlier that the use of mixins causes a collapse of the original class hierarchy. Using union types, however, we can still form a type hierarchy by appropriately grouping classes. As well as the Bank. ACCOUNT union type, such a *type* hierarchy for account classes could include the union of the savings and interest-bearing checking account classes (an investment supertype treating interest-bearing checking account sa savings accounts) and the union of the checking and interest-bearing checking account sa schecking accounts. Note how it is possible for a class to be included in more than one union type.

CONSIDERATIONS FOR ADA 9X

There is a recognized need for improving Ada's support for data abstraction, and the construction of programs from preexisting components.

[Ada9X 91a]

The mixin-based style described in this paper combines the benefits of object-oriented mixins with the advantages of explicit parameterization through generics. With superclass and subclass parameterization, mixins are completely independent software components that can be mixed and matched in many combinations. This leads to a powerful paradigm known as parameterized programming that promotes highly reusable code (see, for example, [Goguen 84; Seidewitz and Stark 91]).

Unfortunately, as the reader can see from the examples in this paper, this style is awkward in places with Ada 83. In particular, the following areas especially need to be addressed in Ada 9X:

- 1. There needs to be a way to create a subclass type by simple extension of a class type and to parameterize this extension with a mixin. The proposed Ada 9X record extension mechanism [Ada9X 91b] fills this need admirably well.
- 2. There needs to be a simpler way to achieve self-reference during the combination of mixins. This need seems to be filled by the proposed mechanism in Ada 9X to allow type extensions as generic formal type parameters [Ada9X 91b]. This would probably necessitate the use of nested generics to allow the mixin type to be an extension of the SUPERCLASS type parameter and the SUBCLASS type parameter to be an extension of the mixin type. Such a construction would, however, eliminate the need for Self functions.
- 3. There needs to be a mechanism for constructing supertypes without having to explicitly code dispatch operations. Ada 9X does provide an automatic dispatching capability using "tagged records" [Ada9X 91b]. However, this capability can only be used if the subtypes are implemented as subclasses (type extensions) of the supertype. This perpetuates the confusion of superclass and supertype.

Thus the proposed object-oriented features for Ada 9X largely support the mixin style described in this paper. Unfortunately, the tagged record mechanism confuses type extension and dispatching. This is analogous to the equation of superclasses and supertypes in most typed object-oriented programming languages (such as C++ [Stroustrup 86]).

Requiring supertypes to be superclasses is inconvenient when we are using generic mixins to construct classes and wish to create a type hierarchy after the fact. Perhaps a better model for Ada 9X would be the "abstract type" mechanism of the languages Emerald [Black et al. 87] and POOL-I [America and van der Linden 90]. Even with the currently proposed Ada 9X features, however, a generics-based approach to mixins, such as that presented in this paper, could be an important contribution of Ada 9X back to the object-oriented programming community.

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