Abstract

Temperley-Lieb algebras have been generalized to web spaces for rank 2 simple Lie algebras. Using these webs, we find a complete description of the Jones-Wenzl idempotents for the quantum \( \mathfrak{sl}(3) \) and \( \mathfrak{sp}(4) \) by single clasp expansions. We discuss applications of these expansions.

1. Introduction

After the discovery of the Jones polynomial [9, 10], its generalizations have been studied in many different ways. Using the quantum \( \mathfrak{sl}(2) \) representation theory, the Jones polynomial can be seen as a polynomial invariant of a colored link whose components are colored by the two dimensional vector representation of the quantum \( \mathfrak{sl}(2) \). By using all irreducible representations of the quantum \( \mathfrak{sl}(2) \), one can find the colored Jones polynomial and it has been extensively studied [5, 11, 19, 26, 35, 38].

The other direction is to use the representation theory of other complex simple Lie algebras from the original work of Reshetikhin and Turaev [30, 31]. These quantized simple Lie algebras invariants can be found by using the Jones-Wenzl idempotents and fundamental representations. In this philosophy, Kuperberg introduced web spaces of simple Lie algebras of rank 2, \( \mathfrak{sl}(3) \), \( \mathfrak{sp}(4) \) and \( G_2 \) as generalizations of Temperley-Lieb algebras corresponding to \( \mathfrak{sl}(2) \) [21]. Then he successively generalized the result for \( \mathfrak{sl}(2) \) [32] that the dimension of the invariant subspace of the tensor of irreducible representations of the quantum \( \mathfrak{sl}(3) \) and \( \mathfrak{sp}(4) \) is equal to the dimension of web spaces of the given boundary with respect to the relations in Fig. 5 and Fig. 12 respectively [21]. But there was no immediate generalization to other Lie algebras until new results for \( \mathfrak{so}(7) \) [37] and \( \mathfrak{sl}(4) \) [17]. The quantum \( \mathfrak{sl}(3) \) invariants have many interesting results [1, 2, 12, 13, 28, 34] also have been generalized to the quantum \( \mathfrak{sl}(n) \) [8, 14, 27, 33, 39]. An excellent review can be found in [6].

Ohtsuki and Yamada generalized Jones-Wenzl idempotents (these were called magic weaving elements) for the quantum \( \mathfrak{sl}(3) \) web spaces by taking the expansions in Proposition 3.1 and 3.4 as a definition of clasps [28]. On the other hand, Kuperberg abstractly proved the existence of generalized Jones-Wenzl idempotents for other simple
Lie algebras of rank 2, he called clasps [21]. In the recursive formula shown in Fig. 1, the resulting webs have two (one with one clasp) clasps, thus it is called a double clasps expansion of the clasp of weight \( n \). There is an expansion for which each resulting web has just one clasp as depicted in Fig. 3 [5]. We called it a single clasp expansion of the clasp of weight \( n \). These expansions are very powerful tools for graphical calculus [5, 15, 35]. We provide single clasp expansions of all quantum \( \mathfrak{sl}(3) \) clasps together with double, quadruple clasps expansions of all quantum \( \mathfrak{sl}(3) \) clasps. We also find single and double clasp expansions of some quantum \( \mathfrak{sp}(4) \) clasps.

Using expansions of clasps, Lickorish first found a quantum \( \mathfrak{sl}(2) \) invariants of 3-manifolds [23, 24]. Ohtsuki and Yamada did for the quantum \( \mathfrak{sl}(3) \) [28] and Yokota found for the quantum \( \mathfrak{sl}(n) \) [39]. For applications of single clasp expansions, first we provide a criterion which determines the periodicity of a link using colored quantum \( \mathfrak{sl}(3) \) and \( \mathfrak{sp}(4) \) link invariants. We discuss a generalization of 3\( j \), 6\( j \) symbols for the quantum \( \mathfrak{sl}(3) \) representation theory. At last, we review how \( \mathfrak{sl}(3) \) invariants extend for a special class of graphs.

The outline of this paper is as follows. In Section 2, we review the original Jones-Wenzl idempotents and provide some algebraic background of the representation theory of \( \mathfrak{sl}(3) \) and \( \mathfrak{sp}(4) \). We provide single clasp expansions of all clasps for the quantum \( \mathfrak{sl}(3) \) in Section 3. In Section 4 we study single clasp expansions of some clasps for the quantum \( \mathfrak{sp}(4) \). In Section 5, we will discuss some applications of the quantum \( \mathfrak{sl}(3) \) clasps and their single clasp expansions. In Section 6, we prove two key lemmas.

A part of the article is from the author’s Ph. D. thesis. Precisely, Section 3 and 6 are from [15, Section 2.3] and Section 4 is from [15, Section 2.4].

2. **Jones-Wenzl idempotents and algebraic background**

For standard terms and notations for representation theory, we refer to [4].

2.1. **Jones-Wenzl idempotents.** An explicit algebraic definition of Jones-Wenzl idempotents can be found in [5]. We will recall a definition of Jones-Wenzl idempotents which can be generalized for other simple Lie algebras. Let \( T_n \) be the \( n \)-th Temperley-Lieb algebra with generators, \( 1, e_1, e_2, \ldots, e_{n-1} \), and relations,

\[
e_i^2 = -(q^{1/2} + q^{-1/2})e_i, \]

\[
e_iei = e_ie_i \quad \text{if} \quad |i - j| \geq 2, \]

\[
e_i = e_ie_{i \pm 1}e_i. \]

For each \( n \), the algebra \( T_n \) has an idempotent \( f_n \) such that \( f_n x = xf_n = \epsilon(x)f_n \) for all \( x \in T_n \), where \( \epsilon \) is an augmentation. These idempotents were first discovered by Jones [9] and Wenzl [36]. They found a recursive formula:

\[
f_n = f_{n-1} + \frac{[n-1]}{[n]} f_{n-1}e_{n-1}f_{n-1},
\]
Fig. 1. A double clasps expansion of the clasp of weight $n$.

Fig. 2. Properties of the Jones-Wenzl idempotents.

Fig. 3. A single clasp expansion of the clasp of weight $n$.

as illustrated in Fig. 1 where we use a rectangular box to represent $f_n$ and the quantum integers are defined as

$$[n] = \frac{q^{n/2} - q^{-n/2}}{q^{1/2} - q^{-1/2}}.$$ 

Thus, they are named Jones-Wenzl idempotents (projectors). It has the following properties 1) it is an idempotent 2) $f_n e_i = 0 = e_i f_n$ where $e_i$ is a U-turn from the $i$-th to the $(i + 1)$-th string as shown in Fig. 2. The second property is called the annihilation axiom. We will discuss the importance of Jones-Wenzl idempotents in Section 2.4. In Fig. 3, $n$ stands for the number of strings and “$i$” stands for $i$-th string from the right. We will use this convention for the rest of the article.

2.2. The representation theory of $\mathfrak{sl}(3)$. The Lie algebra $\mathfrak{sl}(3)$ is the set of all $3 \times 3$ complex matrices with trace zero, which is an 8 dimensional vector space with the Lie bracket. Let $\lambda_i$ be a fundamental dominant weight of $\mathfrak{sl}(3)$, $i = 1, 2$. All finite dimensional irreducible representation of $\mathfrak{sl}(3)$ are determined by its highest weight $\lambda = a\lambda_1 + b\lambda_2$, denoted by $V_\lambda$ where $a$ and $b$ are all nonnegative integers. We will abbreviate $V_{a\lambda_1 + b\lambda_2}$ by $V(a, b)$. The dimension and the quantum dimension of the fun-
damental representation $V_{\lambda_1} \cong (V_{\lambda_2})^*$ of $\mathfrak{sl}(3)$ are 3, [3]. The weight space of a fundamental representation $V(1, 0)$ is $[1, 0]$, $[-1, 1]$ and $[0, -1]$. The weight space of a fundamental representation $V(0, 1)$ is $[0, 1]$, $[1, -1]$ and $[-1, 0]$. Thus, one can easily find the following decomposition formula for a tensor product of a fundamental representation and an irreducible representation,

$$V_{\lambda_1} \otimes V_{\lambda_1+b\lambda_2} \cong V_{(a+1)\lambda_1+b\lambda_2} \oplus V_{(a-1)\lambda_1+(b+1)\lambda_2} \oplus V_{a\lambda_1+(b-1)\lambda_2},$$

$$V_{\lambda_2} \otimes V_{\lambda_1+b\lambda_2} \cong V_{a\lambda_1+(b+1)\lambda_2} \oplus V_{a\lambda_1+(b-1)\lambda_2} \oplus V_{(a+1)\lambda_1+b\lambda_2},$$

with a standard reflection rule, a refined version of the Brauer’s theorem [7, pp.142]. Using these tensor rules, one can find the following lemma.

**Lemma 2.1.** For integers $a, b \geq 1$,

$$\dim(\text{Inv}(V_{\lambda_1}^{\otimes a} \otimes V_{\lambda_2}^{\otimes b} \otimes V_{(b-1)\lambda_1+\lambda_2})) = ab.$$

To compare the weight of cut paths and clasps, we recall the usual partial ordering of the weight lattice of $\mathfrak{sl}(3)$ as

$$a\lambda_1 + b\lambda_2 > (a + 1)\lambda_1 + (b - 2)\lambda_2,$$

$$a\lambda_1 + b\lambda_2 > (a - 2)\lambda_1 + (b + 1)\lambda_2.$$

### 2.3. The representation theory of $\mathfrak{sp}(4)$.

The Lie algebra $\mathfrak{sp}(4)$ is the set of all $4 \times 4$ complex matrices of the following form,

$$
\begin{bmatrix}
A & B \\
C & -'A
\end{bmatrix},
$$

where $'B = B$, $'C = C$

which is a 10 dimensional vector space with the Lie bracket, where $A$, $B$ and $C$ are $2 \times 2$ matrices. Let $\lambda_i$ be a fundamental dominant weight of $\mathfrak{sp}(4)$, $i = 1, 2$. All finite dimensional irreducible representation of $\mathfrak{sp}(4)$ are determined by its highest weight $\lambda = a\lambda_1 + b\lambda_2$, denoted by $V_\lambda$, where $a$ and $b$ are all nonnegative integers. We will abbreviate $V_{a\lambda_1+b\lambda_2}$ by $V(a, b)$. The dimension and the quantum dimension of the fundamental representation $V_{\lambda_i}(V_{\lambda_2})$ of $\mathfrak{sp}(4)$ are 4, [4] (5, [5], respectively). The weight space of a fundamental representation $V(1, 0)$ is $[1, 0]$, $[-1, 1]$, $[1, -1]$ and $[-1, 0]$. The weight space of a fundamental representation $V(0, 1)$ is $[0, 1]$, $[0, -1]$, $[2, -1]$, $[-2, 1]$ and $[0, 0]$. Thus, one can easily find the following decomposition formula for a tensor product of a fundamental representation and an irreducible representation,

$$V_{\lambda_1} \otimes V_{\lambda_1+b\lambda_2} \cong V_{(a+1)\lambda_1+b\lambda_2} \oplus V_{(a-1)\lambda_1+(b+1)\lambda_2} \oplus V_{(a+1)\lambda_1+(b-1)\lambda_2} \oplus V_{(a-1)\lambda_1+b\lambda_2},$$

$$V_{\lambda_2} \otimes V_{\lambda_1+b\lambda_2} \cong V_{a\lambda_1+(b+1)\lambda_2} \oplus V_{a\lambda_1+(b-1)\lambda_2} \oplus V_{(a+1)\lambda_1+b\lambda_2} \oplus V_{(a+2)\lambda_1+(b+1)\lambda_2} \oplus V_{a\lambda_1+b\lambda_2},$$

with a similar reflection rule. Using these tensor rules, one can find the following two lemmas.
Lemma 2.2. For a positive integer $n$,

$$\dim(\text{Inv}(V_{\lambda_1}^{\otimes n} \otimes V_{(n-1)\lambda_1})) = \frac{n(n+1)}{2}.$$ 

Lemma 2.3. For a positive integer $n$,

$$\dim(\text{Inv}(V_{\lambda_2}^{\otimes n} \otimes V_{(n-1)\lambda_2})) = \frac{n(n+1)}{2}.$$ 

There is a natural partial ordering of the $\mathfrak{sp}(4)$ weight lattice given by

$$a\lambda_1 + b\lambda_2 > (a - 2)\lambda_1 + (b + 1)\lambda_2,$$

$$a\lambda_1 + b\lambda_2 > (a + 2)\lambda_1 + (b - 2)\lambda_2.$$ 

2.4. Invariant vector spaces and web spaces. In this subsection, we briefly review the web spaces, full details can be found in [21]. Let $V_1$ be an irreducible representation of complex simple Lie algebras $\mathfrak{g}$. One of classical invariant problems is to characterize the vector space of invariant tensors

$$\text{Inv}(V_1 \otimes V_2 \otimes \cdots \otimes V_n),$$

together with algebraic structures such as tensor products, cyclic permutations and contractions. The dimension of such a vector space is given by Cartan-Weyl character theory; $\dim(\text{Inv}(V_1 \otimes V_2 \otimes \cdots \otimes V_n))$ is the number of copies of the trivial representation in the decomposition of $V_1 \otimes V_2 \otimes \cdots \otimes V_n$ into irreducible representations. For this algebraic space, we look for a geometric counterpart which can preserve the algebraic structure of the invariant spaces. The discovery of quantum groups opens the door for the link between invariant spaces and topological invariants of links and manifolds. For quantum $\mathfrak{sl}(2)$, the dimension of the invariant spaces of $V_1^{\otimes 2n}$ is the dimension of the $n$-th Temperley-Lieb algebra as a vector space which is generated by chord diagrams with $2n$ marked points on the boundary of the disk where $V_1$ is the vector representation of $\mathfrak{sl}(2)$. In particular, this space is free, i.e., there is no relation between chord diagrams. To represent any irreducible representations other than the vector representation, we use Jones-Wenzl idempotents as we described in Section 2.1. Then all webs in the web space of a tensor of irreducible representations $V_1 \otimes V_2 \otimes \cdots \otimes V_n$ can be obtained from webs in the web space of $V_1^{\otimes \sum_i i_k}$ and by attaching Jones-Wenzl idempotents of weight $i_k$ along the boundary (some webs become zero by the annihilation axiom, no longer a basis for web space and the other nonzero webs are called basis webs), where $V_1$ is the irreducible representation of the quantum $\mathfrak{sl}(2)$ of highest weight.
Fig. 4. Generators of the quantum $\mathfrak{sl}(3)$ web space.

$i$ and $k = 1, 2, \ldots, n$. For example [21], the web

is not a basis web of $V_2 \otimes V_3 \otimes V_4 \otimes V_5$, which instead has basis

where the Jones-Wenzl idempotents were presented by the thick gray lines instead of boxes.

A first generalization of Temperley-Lieb algebras was made for simple Lie algebras of rank 2, $\mathfrak{sl}(3)$, $\mathfrak{sp}(4)$ and $G_2$ [21]. Each diagram appears in a geometric counterpart of the invariant vectors is called a web, precisely a directed and weighted cubic planar graph. Unfortunately, some of webs are no longer linearly independent for simple Lie algebra other than $\mathfrak{sl}(2)$. For example, we look at the web space of $\mathfrak{sl}(3)$ representations. Let $V_{\lambda_1}$ be the vector representation of the quantum $\mathfrak{sl}(3)$ and $V_{\lambda_2}$ be the dual representation of $V_{\lambda_1}$. The web space of a fixed boundary (a sequence of $V_{\lambda_1}$ and $V_{\lambda_2}$) is a vector space spanned by the all webs of the given boundary which is generated by the webs in Fig. 4 (as inward and outward arrows) modulo by the subspace spanned by the equation of diagrams which are called a complete set of the relations, equations (1), (2) and (3) as illustrated in Fig. 5. We have drawn a web in Fig. 6. We might use the notation $+,-$ for $V_{\lambda_1}, V_{\lambda_2}$ but it should be clear. For several reasons, such as the positivity and the integrality [22], we use $-[2]$ in relation (2) but one can use a quantum integer [2] and get an independent result. If one uses [2], one can rewrite all results in here by multiplying each trivalent vertex by the complex number $i$. 
To define the generalization of Jones-Wenzl idempotents, *clasps*, we first generalize the annihilation axiom for other web spaces. We need to introduce new concepts: a *cut path* is a path which is transverse to strings of a web, and the *weight* of a cut path is the sum of weights of all decorated strings which intersect with the cut path. For example, the weight of the clasp as depicted in Fig. 7 is $2V_{\lambda_1}$, abbreviated by $(2, 0)$. Then we can generalize the annihilation axiom as follows: if we attach the clasp to a web which has a cut path of weight less than that of the clasp, then it is zero. Since the weight of the clasp shown in Fig. 7 is $(2, 1)$ and there is a cut path of weight $(2, 0)$, the web in Fig. 7 is zero by the annihilation axiom. For $\mathfrak{sl}(3)$, the clasp $\omega$ of weight $(a, b)$ is defined to be the web in the web space of $V_{\lambda_1}^{\otimes a} \otimes V_{\lambda_2}^{\otimes b} \otimes (V_{\lambda_1}^{*})^{\otimes a} \otimes (V_{\lambda_2}^{*})^{\otimes b}$, say $W$, which satisfies the annihilation axiom and the idempotent axiom ($\omega^2 = \omega$). One
can see the dimension of the web space of $W$ is one, i.e., all webs in the web space of $W$ are multiples of $\omega$. However, the clasp of weight $(a, b)$ is unique by the idempotent axiom (it is nonzero). An algebraic proof of the existence of clasps for $\mathfrak{sl}(3)$ and $\mathfrak{sp}(4)$ is given [21]. On the other hand, the double clasps expansion and the quadruple clasps expansion formulae [28] do concretely show the existence of the $\mathfrak{sl}(3)$ clasp. Using these expansions one can find Example 2.4 (we omit some of arrows on the edges of webs, but it should be clear).

Example 2.4. The complete expansions of the clasps of weight $(2, 0)$ and $(3, 0)$ are

$$
\begin{align*}
\begin{array}{c}
\begin{tikzpicture}
\draw (0,0) -- (1,0) -- (1,1) -- (0,1) -- (0,0);
\end{tikzpicture}
\end{array}
\end{align*}
+ \frac{1}{2}
\begin{align*}
\begin{array}{c}
\begin{tikzpicture}
\draw (0,0) -- (1,0) -- (1,1) -- (0,1) -- (0,0);
\end{tikzpicture}
\end{array}
\end{align*}
+ \frac{1}{3}
\begin{align*}
\begin{array}{c}
\begin{tikzpicture}
\draw (0,0) -- (1,0) -- (1,1) -- (0,1) -- (0,0);
\end{tikzpicture}
\end{array}
\end{align*}
+ \frac{1}{2}
\begin{align*}
\begin{array}{c}
\begin{tikzpicture}
\draw (0,0) -- (1,0) -- (1,1) -- (0,1) -- (0,0);
\end{tikzpicture}
\end{array}
\end{align*}
+ \frac{1}{2}
\begin{align*}
\begin{array}{c}
\begin{tikzpicture}
\draw (0,0) -- (1,0) -- (1,1) -- (0,1) -- (0,0);
\end{tikzpicture}
\end{array}
\end{align*}
.

3. Single clasp expansions for the quantum $\mathfrak{sl}(3)$ clasps

First we look at a single clasp expansion of the clasp of weight $(n, 0)$ where the weight $(a, b)$ stands for $a\lambda_1 + b\lambda_2$ in Section 3.1. We can easily find a single clasp expansion of the clasp of weight $(0, n)$ by reversing arrows in the equation presented in the formula of Proposition 3.1. In Section 3.2, we find a single clasp expansion of the clasp of weight $(a, b)$ and double clasps expansions. Kuperberg showed that for a fixed boundary, all webs of the given boundary are cut outs from the hexagonal tiling of the plane with the given boundary [21].

3.1. Single clasp expansions of a clasp of weight $(n, 0)$. First we find a single clasp expansions of a clasp of weight $(n, 0)$ in Proposition 3.1. It is worth to mention that i) this expansion can be obtained from a complete expansion (linear expansions of webs without any clasps) which can be found by using a double clasps expansion iteratively [28] and then attaching a clasp of weight $(n-1, 0)$ to each web in the expansion; ii) the single clasp expansion in Proposition 3.1 holds for any $\mathfrak{sl}(n)$ where $n \geq 4$ because $\mathfrak{sl}(3)$ is naturally embedded in $\mathfrak{sl}(n)$. By symmetries, there are four different single clasps expansions depending on where the clasp of weight $(n-1, 0)$ is located. For equation (4), the clasp is located at the southwest corner, which will be considered the standard expansion, otherwise, we will state the location of the clasp.

We demonstrate Proposition 3.1 for $n = 2, 3$ directly using the presentations of the clasps in Example 2.4. For $n = 2$, Proposition 3.1 is identical to the first formula of Example 2.4. For $n = 3$, by attaching the clasp of weight $(2, 0)$ to the southwest corner
of each web in the second formula of Example 2.4,

\[
\begin{align*}
\parbox{1cm}{\includegraphics{web1}} &= \parbox{1cm}{\includegraphics{web2}} + \frac{[2]}{[3]} \left( \parbox{1cm}{\includegraphics{web3}} + \parbox{1cm}{\includegraphics{web4}} \right) \\
&\quad + \frac{[1]}{[3]} \left( \parbox{1cm}{\includegraphics{web5}} + \parbox{1cm}{\includegraphics{web6}} \right) + \frac{1}{[2][3]} \parbox{1cm}{\includegraphics{web7}}.
\end{align*}
\]

Since \( \parbox{1cm}{\includegraphics{web8}} = 0 \), we find

\[
\begin{align*}
\parbox{1cm}{\includegraphics{web1}} &= \parbox{1cm}{\includegraphics{web2}} + \frac{[2]}{[3]} \left( \parbox{1cm}{\includegraphics{web3}} + \parbox{1cm}{\includegraphics{web4}} \right) \\
&\quad + \frac{[1]}{[3]} \parbox{1cm}{\includegraphics{web5}} + \frac{1}{[2][3]} \parbox{1cm}{\includegraphics{web7}}.
\end{align*}
\]

This verifies the \( n = 3 \) case of Proposition 3.1.

**Proposition 3.1.** For a positive integer \( n \),

\[
\begin{align*}
\parbox{1cm}{\includegraphics{web9}} &= \parbox{1cm}{\includegraphics{web10}} + \frac{n + 1 - i}{[n]} \sum_{i=2}^{n} \parbox{1cm}{\includegraphics{web11}} + \parbox{1cm}{\includegraphics{web12}}.
\end{align*}
\]

Proof. We prove the linear independence of the webs in the right-hand side of the equation (4). Suppose there exists a linear combination of webs which is zero, let us denote \( c_i \) be the coefficient of this linear combination corresponding to the \( i \)-th web.
in the right-hand side of the equation (4).

If we attach the clasp of weight \((n, 0)\) to the top of each web in the right-hand side of equation (4), the first web corresponding to the coefficient \(c_1\) is nonzero because it is a cut out from the hexagonal tiling of the plane. All other remaining webs corresponding to \(c_k\) where \(k \geq 2\) are zero because \(c_1 = 0\). Therefore, \(c_1 = 0\).

Inductively we assume all \(c_k = 0\) where \(k < i\). If we attach the clasp of weight \((n - i + 1, 0)\) to the left top of each web in the right side of equation (4), the \(i\)-th web corresponding to the coefficient \(c_i\) is nonzero because it is a cut out from the hexagonal tiling of the plane. All other remaining webs corresponding to \(c_k\) where \(k \geq i + 1\) are zero because the same reason. Therefore, \(c_i = 0\). This completes the proof of linear independency.

Now, we can show that the set of the webs in the right-hand side in equation (4) is a basis by counting the dimension of web spaces. If we set \(a = (n + 1), b = 1\), we find the dimension of the web space of \(V_{\lambda_1} \otimes V_{\lambda_2} \otimes V_{(n-1)\otimes_2}\) is \(n\) by Lemma 2.1. Therefore, these webs in right side of equation (4) form a basis for the single clasp expansion.
We put

\[
\begin{align*}
\begin{array}{c}
\cdots \\
\quad n \\
\quad n-1
\end{array}
\end{align*}
= a_1 \begin{array}{c}
\cdots \\
\quad n \\
\quad n-1
\end{array} + \sum_{i=2}^{n} a_i \begin{array}{c}
\cdots \\
\quad n \\
\quad n-1
\end{array}
\]

with some \(a_i\), since the webs in the right hand side span the web space which contains the web of the left hand side. If we attach a \(\bigwedge\) on the top of webs in equation (4), the left side of equation (4) becomes zero and all webs in the right-hand side of equation (4) become multiples of a web. Thus we get the following \(n - 1\) equations.

\[a_{n-1} - [2]a_n = 0.\]

For \(i = 1, 2, \ldots, n - 2\),

\[a_i - [2]a_{i+1} + a_{i+2} = 0.\]

From these equations, we are able to find the relations between the coefficients \(a_i\)’s. By a normalization, attaching the clasp of weight \((n, 0)\) to the top of each web in the equation, we find \(a_1 = 1\). Then other coefficients can be found subsequently.

3.2. Single clasp expansions of a non-segregated clasp of weight \((a, b)\). The most interesting case is a single clasp expansion of the clasp of weight \((a, b)\) where \(a \neq 0 \neq b\). By Lemma 2.1, we know the number of webs in a single clasp expansion of the clasp of weight \((a, b)\) is \((a + 1)b\). We need a set of basis webs with a nice rectangular order, but we can not find one in the general case. Even if one finds such a basis, each web in the basis would have many hexagonal faces which make it very difficult to get numerical relations. So we start from an alternative, non-segregated clasp. A non-segregated clasp is obtained from the segregated clasp by attaching a sequence of \(H\)’s until we get the desired shape of edge orientations. Fortunately, there is a canonical way to find them by putting \(H\)’s from the leftmost string of weight \(\lambda_2\) or \(-\) until it reach to the desired position. The left side of the equation in Fig. 8 is an example of a non-segregated clasp of weight \((2, 3)\). the right-hand side of the equation in Fig. 8 shows a sequence of \(H\)’s which illustrates how we obtain it from the segregated clasp of weight \((2, 3)\).

First of all, we can show that the non-segregated clasps are well-defined [15, Lemma 2.6]. One can prove that non-segregated clasps also satisfy two properties of segregated clasps: 1) two consecutive non-segregated clasps is equal to a non-segregated clasp, 2) if we attach a web to a non-segregated clasp and if it has a cut path whose
Fig. 8. A non-segregated clasp of weight (2, 3).

Fig. 9. Fillings for the boxes in equation (5).

weight is less than the weight of the clasp, then it is zero [15, Lemma 2.7]. We find a single clasp expansion of a non-segregated clasp of weight \((a, b)\) as shown in equation (5). Kuperberg showed that for a fixed boundary, the interior can be filled by a cut out from the hexagonal tiling of the plane with the given boundary [21]. For our cases, there are two possible fillings but we use the maximal cut out of the hexagonal tiling. We draw examples of the case \(i = 6, j = 5\) and the first one in Fig. 9 is not a maximal cut out and the second one is the maximal cut out which fits to the left rectangle and the last one is the maximal cut out which fits to the right rectangle as the number indicated in equation (5). An example of a single clasp expansion of a segregated clasp of weight (2, 2) can be found in [15, pp.18].
Theorem 3.2. For $a, b \geq 1$,

\[
\begin{array}{c}
\text{(5)} \\
= \sum_{i=0}^{b} \sum_{j=0}^{a} \frac{[b-i+1][b+j+1]}{[b][a+b+1]}
\end{array}
\]

Proof. Let us denote the web corresponding to the coefficient $[b-i+1][b+j+1]/[b][a+b+1] = a_{i,j}$ by $D_{i,j}$. First of all, all these webs in the equation (5) are nonzero because they are cut outs from the hexagonal tiling of the plane. These webs in the right hand side of the equation form a basis because their cardinality is the same as the dimension of the invariant space of $V_{\lambda_1} \otimes (b-1) \otimes V_{\lambda_2} \otimes V_{\lambda_3} \otimes V_{\lambda_4}$ and they are linearly independent. Suppose that a linear combination of webs in the right-hand side of the equation (5) is zero for some choice of $a_{i,j}$. By attaching the clasp of weight $(0, b)$ and $(0, a+1)$ one can see all webs but the webs $D_{i,j}, 1 \leq s \leq i, 0 \leq t \leq j$ vanish. It is clear that $a_{1,0} = 0$ by attaching the clasp of weight $(0, b)$ and $(0, a+1)$. Inductively, we can show $a_{i,j} = 0$ for all $i, j$.

To find $a_{i,j}$, we attach a $\bigvee$ or a $\bigcap$ to find one exceptional and three types of equations as follow.

\[
[3]a_{1,0} - [2]a_{1,1} - [2]a_{2,0} + a_{2,1} = 0.
\]

Type I: For $j = 0, 1, \ldots, a$,

\[
a_{b-1,j} - [2]a_{b,j} = 0.
\]

Type II: For $i = 1, 2, \ldots, b-2$ and $j = 0, 1, \ldots, a$.

\[
a_{i,j} - [2]a_{i+1,j} + a_{i+2,j} = 0.
\]

Type III: For $i = 1, 2, \ldots, b$ and $j = 0, 1, \ldots, a-2$.

\[
a_{i,j} - [2]a_{i,j+1} + a_{i,j+2} = 0.
\]
If we set $a_{1,0} = x$, then inductively one can see that the coefficient $a_{i,j}$ in the equation (5) is

$$\frac{[b-i+1][b+j+1]}{[b][b+1]} x.$$  

One might check that these are the right coefficients. Usually we normalize one basis web in the expansion to get a known value. But we cannot normalize for this expansion yet because it is not a segregated clasp. Thus we use a complicate procedure in Lemma 6.2 to find that the coefficient of $a_{1,a}$ is 1. Then, we find that $a_{1,0}$ is $[b+1]/[a+b+1]$ and it completes the proof of the theorem.

We find a double clasps expansion as shown in Theorem 3.3, the box between two clasps is filled by the unique maximal cut out from the hexagonal tiling with the given boundary as we have seen in Fig. 9.

**Theorem 3.3.** For $a, b \geq 1$,

Proof. It follows from Lemma 6.1 and Lemma 6.2.

The expansion in equation depicted in Proposition 3.4 was first used to define the segregated clasp of weight $(a, b)$ [28]. The clasps can be constructed from web spaces [21] and these two are known to be equal. We will apply Theorem 3.2 to demonstrate the effectiveness of single clasp expansions by deriving the coefficients in Proposition 3.4.
Proposition 3.4 ([28]). A quadruple clasps expansion of the segregated clasp of weight \((a, b)\) is

\[
\begin{align*}
\text{Fig. 10. Induction step for the proof of Proposition 3.4.}
\end{align*}
\]

Proof. Let us denote the \(k\)-th term in the right-hand side of equation by \(D(k)\). We induct on \(a + b\). It is clear for \(a = 0\) or \(b = 0\). If \(a \neq 0 \neq b\) then we use a segregated single clasp expansion of weight \((a, b)\) in the middle for the first equality. Even if we do not use the entire single clasp expansion of a segregated clasp, once we attach clasps of weight \((a, 0), (0, b)\) on the top, there are only two nonzero webs which are webs with just one \(U\)-turn. One of resulting webs has some \(H\)'s as in Fig. 10 but if we push them down to the clasp of weight \((a, b - 1)\) in the middle, it becomes a non-segregated clasp. For the second equality we use a non-segregated single clasp

\[
\begin{align*}
\sum_{k=0}^{\min(a,b)} (-1)^k \frac{a!b! [a+b-k+1]!}{(a-k)! (b-k)! [k]! [a+b+1]!} a - k
\end{align*}
\]
expansion at the clasp of weight \((a, b - 1)\) for which clasps of weight \((a - 1, b - 1)\)
are located at northeast corner. By the induction hypothesis, we have

\[
1 = \sum_{k=0}^{b-1} (1)^k \frac{[a]! [b - 1]! [a + b - k]!}{[a - k]! [b - 1 - k]! [k]! [a + b]!} D(k)
- \frac{[a + 1][a]}{[a + b + 1][a + b]} \sum_{k=0}^{b-1} (1)^k \frac{[a - 1]! [b - 1]! [a + b - 1 - k]!}{[a - 1 - k]! [b - 1 - k]! [k]! [a + b - a]!} D(k + 1)
= 1 \cdot D(0) + \sum_{k=1}^{b-1} (1)^k \left( \frac{[a + 1][a + b - 1]!}{[a - k]! [b - 1 - k]! [k]! [a + b]!} \right) D(k)
\]

\[
= D(0)
+ \sum_{k=1}^{b-1} (1)^k \frac{[a]! [b]! [a + b + 1 - k]!}{[a - k]! [b - k]! [k]! [a + b + 1]!} \left( \frac{[b - k][a + b + 1] + [k][a + 1]}{[b][a + b + k - 1]} \right) D(k)
+ (1)^b \frac{[a]! [b - 1]! [a + 1]!}{[a - b]! [0]! [b - 1]! [a + b + 1]!} D(b)
\]

\[
= \sum_{k=0}^{b} (1)^k \frac{[a]! [b]! [a + b + 1 - k]!}{[a - k]! [b - k]! [k]! [a + b + 1]!} D(k)
\]

4. Single clasp expansion for the quantum \(\mathfrak{sp}(4)\)

The quantum \(\mathfrak{sp}(4)\) webs are generated by a single web in Fig. 11 and a complete
set of relations is given in Fig. 12 [21]. Again, an algebraic proof of the existence
of the clasp of the weight \((a, b)\) using the annihilation axiom and the idempotent
axiom is given in [21]. On the other hand, one can use the double clasps expansions
in Corollary 4.3 and Corollary 4.5 to define the clasps of the weight \((n, 0)\) and \((0, n)\).
Unfortunately, we do not have any expansion formula for the clasp of the weight \((a, b)\)
where \(a \neq 0 \neq b\). Using these expansions one can find Example 4.1. We can define
tetravalent vertices to achieve the same end as in Fig. 13. We will use the these shapes
to find a single clasp expansion otherwise there is an ambiguity of a preferred direction
by the last relation presented in Fig. 12.

First we will find single clasp expansions of clasps of weight \((n, 0)\) and \((0, n)\) and
then use them to find coefficients of double clasps expansions of clasps of weight \((n, 0)\)
and \((0, n)\). But we are unable to find a single clasp expansion of the clasp of weight
\((a, b)\) where \(a \neq 0 \neq b\). Remark that the cut weight is defined slightly different way. A
cut path may cut diagonally through a tetravalent vertex, and its weight is defined as
Fig. 11. The generator of the quantum $sp(4)$ web space.

\[ = \frac{[6][2]}{[3]}, \quad = \frac{[6][5]}{[3][2]}, \quad = 0, \quad = -[2]^2 \]

Fig. 12. A complete set of relations of the quantum $sp(4)$ web space.

\[ = \quad = \quad = \]

Fig. 13. Tetravalent vertices.

\[ = \quad = \]

\[ n\lambda_1 + (k+k')\lambda_2, \text{ where } n \text{ is the number of type "1", single strands, that it cuts, } k \text{ is the number of type "2", double strands, that it cuts, and } k' \text{ is the number of tetravalent vertices that it bisects.} \]

\[ \text{EXAMPLE 4.1. The complete expansions of the clasps of weight (2, 0) and (3, 0) are} \]

\[ = + \frac{1}{[2]^2} \quad + \frac{[3][4]}{[2]^2[6]}, \]

\[ = + \frac{1}{[3]} \left( \quad + \quad \right) + \frac{1}{[2]^2[3]} \left( \quad + \quad \right) + \frac{[4]^2}{[2][3][8]} \left( \quad + \quad \right) + \frac{[2][4][6]}{[3]^2[8]} \left( \quad + \quad \right) + \frac{[4][6]}{[2]^2[3]^2[8]} \]

We demonstrate Theorem 4.2 for \( n = 2, 3 \) by using the presentations of clasps in Example 4.1. For \( n = 2 \), Theorem 4.2 is identical to the first formula of Example 4.1. For \( n = 3 \), we first attach the clasp of weight \((2, 0)\) to the southwest corner of each web in the second formula of Example 4.1. Since \( y = 0 \) and \( y = 0 \), we find

\[
\begin{align*}
\text{Fig. 5} & = \frac{1}{[3]} + \frac{1}{[2]^2[3]} + \frac{[2][4][6]}{[3]^2[8]} + \frac{[4][6]}{[2][3]^2[8]} + \frac{[4]^2}{[2][3][8]}.
\end{align*}
\]

We can confirm these coefficients are the same as given in Theorem 4.2.

Now, we state a single clasp expansion of the clasp of weight \((n, 0)\).

**Theorem 4.2.** For a positive integer \( n \),

\[
(6)
\]

Proof. By combining with the weight diagram of \( V_{\lambda_1}^{\otimes n} \) and minimal cut paths, we can find a set of nonzero webs for single clasp expansion of a clasp of weight \((n, 0)\) as in equation (6). Let us denote the web corresponding to the \( i \)-th in the first summation and \( j \)-th in the second summation by \( D_{i,j} \) and its coefficient by \( a_{i,j} \). First we will show that these webs are linearly independent. Suppose that a linear combination of the right-hand side of the equation in Fig. 5 is zero for some choice of \( a_{i,j} \). It is clear that \( a_{i,i+1} = 0 \) by attaching the clasp of weight \((n - i, 0)\) to left top of webs and the
clasp of weight \((i, 0)\) to the right top of each webs. By attaching the clasp of weight \((n - j + 1, 0)\) to left top of webs and the clasp of weight \((i, 0)\) to the right top of webs, inductively we can show \(a_{ij} = 0\) for all \(j \geq i + 1\). By Lemma 2.2, we know that the dimension of the web space of \(V_{k,1}^{\otimes n+1} \otimes V_{l,1}^{(n-1)\otimes 1}\) is \(n(n + 1)/2\). Thus, these webs in right hand side of the equation form a basis.

Now we are set to finds \(a_{i,j}\). For equations, we remark that the relations of webs shown in Fig. 14 can be easily obtained from the relations depicted in Fig. 12. Using these relations, we get the following \(n-1\) equations by attaching a \(\bigcap\). By attaching a \(\bigcup\), we have \((n-1)^2\) equations. There are two special equations and four different shapes of equation as follows.

\[
\begin{align*}
    a_{n-2,n-1} + \frac{[2][6]}{[3]} a_{n-2,n} - \frac{[2][6]}{[3]} a_{n-1,n} &= 0, \\
    -\frac{[2][6]}{[3]} a_{12} + \frac{[2][6]}{[3]} a_{13} + a_{23} + 1 + \frac{[2][6]}{[3]} b_2 - [2][4] b_3 &= 0.
\end{align*}
\]

Type I: For \(i = 1, 2, \ldots, n-3\),

\[
a_{i,i+1} + \frac{[2][6]}{[3]} a_{i,i+2} - [2][4] a_{i,i+3} - \frac{[2][6]}{[3]} a_{i+1,i+2} + \frac{[2][6]}{[3]} a_{i+1,i+3} + a_{i+2,i+3} = 0.
\]

Type II: For \(i = 0, 1, \ldots, n-2\),

\[a_{i,n-1} - [2]^2 a_{i,n} = 0.\]

Type III: For \(i = 0, 1, 2, \ldots, n-3, k = 2, 3, \ldots, n-i-1\),

\[a_{i,n-k} - [2]^2 a_{i,n-k+1} + [2]^2 a_{i,n-k+2} = 0.\]

Type IV: For \(i = 3, 4, \ldots, n, k = n - i + 3, n - i + 4, \ldots, n\),

\([2]^2 a_{n-k,i} - [2]^2 a_{n-k+1,i} + a_{n-k+2,i} = 0.\]
Then we check our answer satisfies the equations and it is clear that $a_{0,1} = 1$ by a normalization. Since these webs in the equation (6) form a basis, the coefficients are unique. Therefore, it completes the proof.

By attaching the clasp of weigh $(n - 1, 0)$ on the top of all webs in the equation presented in equation (6), we find the double clasp expansion of the clasp of weight $(n, 0)$.

**Corollary 4.3.** For a positive integer $n$,

\[
\begin{align*}
\begin{array}{c}
\text{n} \\
\text{n} - 1
\end{array}
&= \begin{array}{c}
\text{n - 1} \\
\text{n - 1}
\end{array} + \frac{[2n][n + 1][n - 1]}{[2n + 2][n][n]} + \frac{[n - 1]}{[n][2]} \\
\text{n - 1}
&= \text{n - 1} + \text{n - 1}.
\end{align*}
\]

Then we look at the clasp of weight $(0, n)$. The main idea for the clasp of weight $(n, 0)$ works exactly same except we replace the basis as shown in equation (7). For the linear independency, every idea of the proof of Theorem 4.2 works with the fact $a_{0,0} = 0$. As we did for the clasp of weight $(n, 0)$, we first find the equations as illustrated in Fig. 15 for the next step. The same as before, we set $a_{ij}$ be the coefficient of the web of $(i, j)$ in the summation. By attaching a $\cap$ and a $\cup$, we get the following equations and we can solve them successively as in Theorem 4.4.

\[
\begin{align*}
\alpha_{n-2,n-1} - [5][2]^2 \alpha_{n-2,n} + \frac{[6][5]}{[3][2]} \alpha_{n-1,n} &= 0, \\
-\frac{[3][2]^2 \alpha_{n-2,n} + [5] \alpha_{n-1,n}}{[3][2]} &= 0.
\end{align*}
\]
Type I: For $i = 0, 1, \ldots, n - 3$,

\[ a_{i,i+1} - [5][2]^2a_{i,i+2} + [3][2]^4a_{i,i+3} + \frac{[6][5]}{[3][2]}a_{i+1,i+2} - [5][2]^2a_{i+1,i+3} + a_{i+2,i+3} = 0. \]

Type II: For $i = 0, 1, \ldots, n - 2$,

\[ a_{i,n-1} - [4][2]a_{i,n} = 0. \]

Type III: For $i = 0, 1, \ldots, n - 3$ and $j = i + 1, i + 2, \ldots, n - 2$,

\[ a_{i,j} - [4][2]a_{i,j+1} + [2]^4a_{i,j+2} = 0. \]

Type IV: For $i = 0, 1, \ldots, n - 3$ and $j = i + 3, i + 4, \ldots, n$,

\[ [2]^4a_{i,j} - [4][2]a_{i+1,j} + a_{i+2,j} = 0. \]

Type V: For $i = 1, 2, \ldots, n - 2$

\[-[3][2]^2a_{i-1,i+1} + [2]^4a_{i-1,i+2} + [5]a_{i,i+1} - [3][2]^2a_{i,i+2} = 0. \]

**Theorem 4.4.** For a positive integer $n$,

\[ \sum_{i=0}^{n-1} \sum_{j=i+1}^{n} [2]^{2(i+j-i-j)} \frac{[2n + 1 - 2i][2n - 2j + 2]}{[2n][2n + 1]} \]

By attaching the clasp of weight $(0, n - 1)$ on the top of all webs shown in equation (7) we find the double clasps expansion of the clasp of weight $(0, n)$.

**Corollary 4.5.** For a positive integer $n$,
5. Applications of the quantum $\mathfrak{sl}(3)$ representation theory

In the section we will discuss some applications of the quantum $\mathfrak{sl}(3)$ representation theory.

5.1. Polynomial invariants of links. The HOMFLY polynomial $P_3(q)$ can be obtained by coloring all components by the vector representations of the quantum $\mathfrak{sl}(3)$ and the following skein relations

\[
P_3(\emptyset) = 1,
\]

\[
P_3(\bigcirc \cup D) = [3]P_3(D),
\]

\[
q^{3/2}P_3(L_+) - q^{-3/2}P_3(L_-) = (q^{1/2} - q^{-1/2})P_3(L_0),
\]

where $\emptyset$ is the empty diagram, $\bigcirc$ is the trivial knot and $L_+,$ $L_-$ and $L_0$ are three diagrams which are identical except at one crossing as illustrated in Fig. 16. On the other hand, the polynomial $P_3(q)$ can be computed by linearly expanding each crossing into a sum of webs as shown in Fig. 17 then by applying relations in Fig. 5 [1, 20, 27]. A benefit of using webs is that we can easily define the colored $\mathfrak{sl}(3)$ HOMFLY polynomial $G_3(L, \mu)$ of $L$ as follows. Let $L$ be a colored link of $l$ components say, $L_1, L_2, \ldots, L_l,$ where each component $L_i$ is colored by an irreducible representation $V_{a_i\lambda_1 + b_i \lambda_2}$ of the quantum $\mathfrak{sl}(3)$ and $\lambda_1, \lambda_2$ are the fundamental weights of $\mathfrak{sl}(3).$ The coloring is denoted by $\mu = (a_1\lambda_1 + b_1 \lambda_2, a_2\lambda_1 + b_2 \lambda_2, \ldots, a_l\lambda_1 + b_l \lambda_2).$ First we replace each component $L_i$ by $a_i + b_i$ copies of parallel lines and each of $a_i$ lines is colored by the weight $\lambda_1$ and each of $b_i$ lines is colored by the weight $\lambda_2.$ Then we put a clasp of the weight $(a_i\lambda_1 + b_i \lambda_2)$ for $L_i.$ If we assume the clasps are far away from crossings,
Fig. 18. Trihedron coefficients for $sl(2)$.

we expand each crossing as depicted in Fig. 17, then expand each clasp inductively by Theorem 3.3. The value we find after removing all faces by using the relations in Fig. 5 is the colored $sl(3)$ HOMFLY polynomial $G_3(L, \mu)$ of $L$. One can find the following theorem which is a generalization of a criterion to determine the periodicity of a link [1, 3].

**Theorem 5.1.** Let $p$ be a positive integer and $L$ be a $p$-periodic link in $S^3$ with the factor link $\overline{L}$. Let $\mu$ be a $p$-periodic coloring of $L$ and $\overline{\mu}$ be the induced coloring of $\overline{L}$. Then

$$G_3(L, \mu) = G_3(\overline{L}, \overline{\mu})^{p}, \text{ modulo } I_3,$$

where $\overline{L}$ is the factor link and $I_3$ is the ideal of $\mathbb{Z}[q^{\pm 1/2}]$ generated by $p$ and $[3]^p - [3]$.

**Proof.** Since the clasps are idempotents, for each component, we put $p - 1$ extra clasps for each copies of components by the rotation of order $p$. First we keep the clasps far away from the crossings. The key idea of the proof given in [1] is that if any expansion of crossings occurs in the link diagram, it must be used identically for all other $p - 1$ copies of the diagram. Otherwise there will be $p$ identical shapes by the rotation of order $p$, then it is congruent to zero modulo $p$. By the same philosophy, if any application of relations occurs, it must be used identically for all other $p - 1$ copies. Otherwise it is congruent to zero modulo $p$. Once there is an unknot in the fundamental domain of the action of order $p$, there are $p$ identical unknots by the rotation which occurs only once in the factor link. Therefore, we get the congruence $[3]^p - [3]$.

5.2. $3j$ and $6j$ symbols for the quantum $sl(3)$ representation theory. $3j$ symbols and $6j$ symbols for the quantum $sl(2)$ representation theory have many significant implications in mathematics and physics. $3j$ symbols are given in the equation shown in Fig. 18 [25]. Its natural generalization for the quantum $sl(3)$ representation theory was first suggested [21] and studied [16]. Let $\lambda_1, \lambda_2$ be the fundamental dominant weights of $sl(3, \mathbb{C})$. Let $V_{\lambda_1 + \lambda_2}$ be an irreducible representation of $sl(3, \mathbb{C})$ of highest
weight \(a\lambda_1 + b\lambda_2\). Now each edge of \(\Theta\) is decorated by an irreducible representation of \(\mathfrak{sl}(3)\), let say \(V_{a_1\lambda_1 + b_1\lambda_2}, V_{a_2\lambda_1 + b_2\lambda_2}\) and \(V_{a_3\lambda_1 + b_3\lambda_2}\) where \(a_i, b_j\) are nonnegative integers. Let \(d = \min\{a_1, a_2, a_3, b_1, b_2, b_3\}\). If \(\dim(\text{Inv}(V_{a_1\lambda_1 + b_1\lambda_2} \otimes V_{a_2\lambda_1 + b_2\lambda_2} \otimes V_{a_3\lambda_1 + b_3\lambda_2}))\) is nonzero, in fact \(d+1\), then we say a triple of ordered pairs \((a_1, b_1), (a_2, b_2), (a_3, b_3)\) is admissible. One can show \((a_1, b_1), (a_2, b_2), (a_3, b_3)\) is admissible if and only if there exist nonnegative integers \(k, l, m, n, o, p, q\) such that \(a_2 = d + l + p,\ a_3 = d + n + q,\ b_1 = d + k + p,\ b_2 = d + m + q,\ b_3 = d + o\) and \(k - n = o - l = m\). For an admissible triple, we can write its trihedron coefficients as a \((d + 1) \times (d + 1)\) matrix. Let us denote it by \(M_0\) \((a_1, b_1, a_2, b_2, a_3, b_3)\) or \(M_0(\lambda)\) where \(\lambda = a_1\lambda_1 + b_1\lambda_2 + a_2\lambda_1 + b_2\lambda_2 + a_3\lambda_1 + b_3\lambda_2\). Also we denotes its \((i, j)\) entry by \(\Theta(a_1, b_1, a_2, b_2, a_3, b_3; i, j)\) or \(\Theta(\lambda; i, j)\) where \(0 \leq i, j \leq d\). The trihedron shape of \(\Theta(a_1, b_1, a_2, b_2, a_3, b_3; i, j)\) is given in Fig. 19 where the triangles are filled by cut outs from the hexagonal tiling of the plane [21]. \(M_0(0, m + n, l, m + q, n + q, m + l), M_0(0, n + p, p + l, q, n + q, l)\) and \(M_0(i, j + k, k + l, m, j + m, j + l; 0, 0)\) were found in [16]. All other cases of \(3j\) symbols and \(6j\) symbols are left open.

5.3. \(\mathfrak{sl}(3)\) invariants of cubic planar bipartite graphs. The \(\mathfrak{sl}(3)\) webs are directed cubic bipartite planar graphs together circles (no vertices) where the direction of the edges is from one set to the other set in the bipartition. From a given directed cubic bipartite planar graph, we remove all circles by the relation (1) and then remove
the multiple edges by the relation (2) in Fig. 5. Using a simple application of the Euler characteristic number of a graphs in the unit disc, we can show the existence of a rectangular face [28]. By inducting on the number of faces, we provides the existence of the quantum $\mathfrak{sl}(3)$ invariants of directed cubic bipartite planar graphs. It is fairly easy to prove the quantum $\mathfrak{sl}(3)$ invariant does not depend on the choice of directions in the bipartition. Thus, the quantum $\mathfrak{sl}(3)$ invariant naturally extends to any cubic bipartite planar graph $G$, let us denote it by $P_G(q)$. By using a flavor of graph theory, we find a classification theorem and provide a method to find all 3-connected cubic bipartite planar graphs which is called prime webs [18]. As little as it is known about the properties of the quantum invariants of links, we know a very little how $P_G(q)$ tells us about the properties of graphs.

For symmetries of cubic bipartite planar graph, the idea of the Theorem 5.1 and 5.3 works for the $\mathfrak{sl}(3)$ graph invariants with one exception. There is a critical difference between these two invariants which is illustrated in Theorem 5.2.

**Theorem 5.2** ([18]). Let $G$ be a planar cubic bipartite graph with the group of symmetries $\Gamma$ of order $n$. Let $\Gamma_d$ be a subgroup of $\Gamma$ of order $d$ such that the fundamental domain of $G/\Gamma_d$ is not a basis web with the given boundary. Then

$$P_G(q) \equiv (P_{G/\Gamma_d}(q))^d \mod \mathcal{I}_d,$$

where $\mathcal{I}_d$ is the ideal of $\mathbb{Z}[q^{\pm 1/2}]$ generated by $d$ and $[3]^d - [3]$.

If the fundamental domain of $G/\Gamma$ is a basis web with the given boundary, then the main idea of the theorem no longer works and a counterexample was found as follows [18]. We look at an example $6_1$ as shown in Fig. 20. By a help of a machine, we can see that there does not exist an $\alpha \in \mathbb{Z}[q^{\pm 1/2}]$ such that


even though there do exist a symmetry of order 6 for $6_1$.

**5.4. Applications for the quantum $\mathfrak{sp}(4)$ representation theory.** A quantum $\mathfrak{sp}(4)$ polynomial invariant $G_{\mathfrak{sp}(4)}(L, \mu)$ can be defined [20, 21] where $\mu$ is a funda-
ment of the quantum \( \mathfrak{sp}(4) \). Since we have found single clasp expansion of the clasps of weight \((a, 0)\) and \((0, b)\), we can extend \( G_{\mathfrak{sp}(4)}(L, \mu) \) for \( \mu \) is an irreducible representations of weight either \((a, 0)\) and \((0, b)\). If we assume a coloring \( \mu = (a, 0) \) or \( \mu = (0, b) \), by the same idea of the proof of Theorem 5.1, we can find the following theorem from Corollary 4.3 and 4.5.

**Theorem 5.3.** Let \( p \) be a positive integer and \( L \) be a \( p \)-periodic link in \( S^3 \) with the factor link \( \overline{L} \). Let \( \mu \) be a \( p \)-periodic coloring of \( L \) and \( \overline{\mu} \) be the induced coloring of \( \overline{L} \). Then

\[
G_{\mathfrak{sp}(4)}(L, \mu) \equiv G_{\mathfrak{sp}(4)}(\overline{L}, \overline{\mu}) \pmod{\mathcal{I}_{\mathfrak{sp}(4)}},
\]

where \( \overline{L} \) is the factor link and \( \mathcal{I}_{\mathfrak{sp}(4)} \) is the ideal of \( \mathbb{Z}[q^{\pm 1/2}] \) generated by \( p, (-[6][2]/[3])^p + [6][2]/[3] \) and \( ([6][5]/[3][2])^p - [6][5]/[3][2] \).

In fact, Theorem 5.3 remains true even if \( \mu \) is any finite dimensional irreducible representation of \( \mathfrak{sp}(4) \), but we would not be able to obtain the actual polynomials because any expansion is not known for the clasp of the weight \((a, b)\) where \( a \neq 0 \neq b \).

6. The proof of lemmas

Let us recalled that the relation (3) in Fig. 5 is called a rectangular relation and the first (second) web in the right-hand side of the equality is called a horizontal (vertical, respectively) splitting. The web in the equation shown in Fig. 5 corresponding to the coefficient \( a_{i,j} \) is denoted by \( D_{i,j} \). After attaching \( H \)'s to \( D_{i,j} \) as illustrated in Fig. 21, the resulting web is denoted by \( \tilde{D}_{i,j} \). We find that \( \tilde{D}_{i,j} \) contains some elliptic faces. If we decompose each \( \tilde{D}_{i,j} \) into a linear combination of webs which have no elliptic faces, then the union of all these webs forms a basis. Let us prove that these webs actually form a basis which will be denoted by \( D_{i,j}^{BC} \). As vector spaces, this change, adding \( H \)'s as in Fig. 21, induces an isomorphism between two web spaces because its matrix representation with respect to these web bases \( \{D_{i,j}\} \) and \( \{D_{i,j}^{BC}\} \) is an \( (a+1)b \times (a+1)b \) matrix whose determinant is \( \pm [2]^{ab} \) because a single \( H \) contributes \( \pm [2] \) depends on the choice of the direction of \( H \).

To find a single clasp expansion of the segregated clasp of weight \((a, b)\), we have to find all linear expansions of \( \tilde{D}_{i,j} \) into a new web basis \( D_{i,j}^{BC} \). In general this is very complicated. Instead of using relations for linear expansions, we look for an alternative. From \( \tilde{D}_{i,j} \) we see that there are \( a+b+1 \) nodes on top and \( a+b-1 \) nodes right above the clasp. A \( Y \) shape in the web \( D_{i,j} \) forces \( \tilde{D}_{i,j} \) to have at least one rectangular face. Each splitting creates another rectangular face until it becomes a basis web (possibly using the relation (2) in Fig. 5 once). If we repeatedly use the rectangular relations as in equation (3) in Fig. 5, we can push up \( Y \)'s so that there are either two \( Y \)'s or one \( U \) shape at the top. A stem of a web is \( a+b-1 \) disjoint union of vertical lines which connect top \( a+b-1 \) nodes out of \( a+b+1 \) nodes to the clasp of weight \((a, b-1)\).
Fig. 21. A sequence of H’s which transforms $D_{ij}$ to a linear combinations of webs in the single clasp expansion of segregated clasp of weight $(a, b - 1)$.

together a $U$-turn or two $Y$’s on top. It is clear that these connecting lines should be mutually disjoint, otherwise, we will have a cut path with weight less than $(a, b - 1)$, i.e., the web is zero. Unfortunately some of stems do not arise naturally in the linear expansion of $\tilde{D}_{i,j}$ because it may not be obtained by removing elliptic faces. If a stem appears, we call it an admissible stem. For single clasp expansions, finding all these admissible stems will be more difficult than linear expansions by relations. But for double clasp expansions of segregated clasps of weight $(a, b)$ there are only few possible admissible stems whose coefficients are nonzero.

**Lemma 6.1.** After attaching the clasp of weight $(a, b - 1)$ to the top left side of webs $\tilde{D}_{i,j}$ from the equation in Fig. 21, the only non-vanishing webs are those three webs as depicted in Fig. 22.

Proof. It is possible to have two adjacent $Y$’s which appear in the second and third webs in Fig. 22 but a $U$-turn can appear in only two places because of the orientation of edges. If we attach the clasp of weight $(a, b - 1)$ to the northwest corner of the resulting web and if there is a $U$ or a $Y$ shape just below the clasp of weight $(a, b - 1)$, the web becomes zero. Therefore only these three webs do not vanish. 

In the following lemma, we find all $\tilde{D}_{i,j}$’s which can be transformed to each of the web in Fig. 22.

**Lemma 6.2.** Only $\tilde{D}_{1,a}(\tilde{D}_{2,a})$ can be transformed to the first (second, respectively) web in Fig. 22. Only the three webs, $\tilde{D}_{1,a-1}$, $\tilde{D}_{1,a}$ and $\tilde{D}_{2,a-1}$ can be transformed to the last web. Moreover, all of these transformations use only rectangular relations as in equation (3) except the transformation from $\tilde{D}_{1,a-1}$ to the third web uses the relation (2) in Fig. 5 exactly once.
Fig. 22. Three webs which do not vanish after attaching the clasp of weight \((a, b - 1)\) to the top left side of webs \(\tilde{D}_{i,j}\) from the equation in Fig. 21.

Fig. 23. \(D_{i,a}\) where \(i > 1\).

Proof. For the first web shown in Fig. 22, it is fairly easy to see that we need to look at \(\tilde{D}_{1,a}\), for \(i = 1, 2, \ldots, b\), otherwise the last two strings cannot be changed to the first web presented in Fig. 22 with a \(U\)-turn. Now we look at the \(D_{i,a}\) where \(i > 1\) as illustrated in Fig. 23. Since we picked where the \(U\) turn appears already, only possible disjoint lines are given as thick and lightly shaded lines but we cannot finish to have a stem because the darkly shaded string from the left top cannot be connected to the bottom clasp without being zero, i.e., if we connect the tick line to clasp, there will be either \(\bigwedge\) or a \(\bigcap\) right above of the clasp of weight \((a, b - 1)\).

So only nonzero admissible stems should be obtained from \(\tilde{D}_{1,a}\). As we explained before, one can see that there is a rectangular face in the web \(\tilde{D}_{1,a}\). Since the horizontal splitting makes it zero, we have to split vertically. For the resulting web, this process created one rectangular face at right topside of previous place. We have to split
vertically and the process are repeated until the last step, both splits do not vanish. The web in the last step is drawn in Fig. 24 with the rectangular face, darkly shaded. The vertical split gives us the first web in Fig. 22 and the horizontal split gives the third web in Fig. 22.

A similar argument works for the second web illustrated in Fig. 22. The third web depicted in Fig. 22 is a little subtle. First one can see that none of $\tilde{D}_{i,j}$ can be transformed if either $i > 2$ or $j < a - 1$. Thus, we only need to check $\tilde{D}_{1,a-1}$, $\tilde{D}_{1,a}$, $\tilde{D}_{2,a-1}$ and $\tilde{D}_{2,a}$ but we already know about $\tilde{D}_{1,a}$, $\tilde{D}_{2,a}$. Fig. 25 shows the nonzero admissible stem for $\tilde{D}_{1,a-1}$. As usual, we draw a stem as a union of thick and darkly shaded lines. Note that we have used relation (3) in Fig. 5 exactly once which contributes $-2$. The Fig. 26 shows the nonzero admissible stem for $\tilde{D}_{2,a-1}$. This completes the proof of the lemma.
Fig. 26. The nonzero admissible stem for $\tilde{D}_{2,a-1}$.

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