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Ultra-Low Power Circuits Using Graphene P-N Junctions and Adiabatic Computing

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Abstract

Recent works have proven the functionality of electrostatically controlled graphene p-n junctions that can serve as basic primitive for the implementation of a new class of compact graphene-based reconfigurable multiplexer logic gates. Those gates, referred as RG-MUXes, while having higher expressive power and better performance w.r.t. standard CMOS gates, they also have the drawback of being intrinsically less power/energy efficient.

In this work we address this problem from a circuit perspective, namely, we revisit RG-MUXes as devices that can operate *adiabatically* and hence with ultra-low (ideally, almost zero) power consumption. More specifically, we show how to build basic logic gates and, eventually, more complex logic functions, by appropriately interconnecting graphene-based p-n junctions as to implement the adiabatic charging principle.

We provide a comparison in terms of power and performance against both adiabatic CMOS and their non-adiabatic graphene-based counterparts; characterization results collected from SPICE simulations on a set of representative functions show that the proposed ultra-low power graphene circuits can operate with 1.5 to 4 orders of magnitude less average power w.r.t. adiabatic CMOS and non-adiabatic graphene counterparts respectively. When it comes to performance, adiabatic graphene shows 1.3 (w.r.t. adiabatic CMOS) to 4.5 orders of magnitude (w.r.t. non-adiabatic technologies) better power-delay product.

Keywords: Adiabatic Circuits, Graphene p-n junction, Low Power Circuits, Graphene Nanoelectronics

1. Introduction

In the last decade, graphene rapidly emerged as a potential candidate to replace silicon in the next generation of electronic circuits due to its astounding electro-mechanical properties [1]. Graphene is a stretchable and transparent electrical conductor with carrier mobility and saturation velocity far larger than standard silicon-based semiconductors. Those features, combined with the possibility of arranging graphene with other materials to form new composites, represent a perfect match for the growing market of flexible and wearable mobile applications.

Besides these superlatives, however, graphene also shows an indisputable limit, that is, the lack of an energy band gap; conduction and valence band touch each other at zero-energy, where the Fermi Energy E_F passes, thereby preventing the material to implement the OFF state. In other words, the ON/OFF current ratio in graphene is quite below the value reached in silicon, resulting in a weak separation between logic 0's and 1's. Needless to say, this characteristic has been initially used as an argument to support the inadequacy of graphene in the implementation of electronic devices. However, recent studies have proven possible ways to overcome this drawback.

On one hand people have addressed the problem by following a “semiconductor-like” strategy focusing on possible fabrication techniques that physically open an energy gap in the material. Most of them belong to industries that do not want to waste the huge investments done for silicon and would like to replicate as much as possible the successful story of silicon semiconduc-

tors with minimal efforts. Graphene Nanoribbons [2] (GNRs) are the most popular embodiment of this class of approaches. GNRs consist of narrow stripes of graphene that show an energy band gap inversely proportional to their width; like standard semiconductors, GNRs can be used to implement Field Effects Transistors (GNR-FETs), e.g., those presented in [3] and [4]. Although even very narrow GNRs exhibit energy gaps sufficiently large for use as a semiconductor to implement graphene-FETs [5], it is quite difficult to fabricate samples of graphene with perfect edges. Edge roughness alters the level of disorder in the material and results in significant degradation of device characteristics and its electrical properties [6].

Another approach exploits the semi-metallic nature of graphene and tries to accommodate it by means of alternative solutions that do not require any physical distortion of the lattice structure. *Electrostatic doping* [7] is the most representative strategy belonging to this class. It allows a fine-tuning of the Fermi Energy E_F that can be shifted down into the valence band (to obtain p-type graphene) or up into the conduction band (to obtain n-type graphene) using external electrical field applied through metal gates. Face-to-face regions with opposite doping profiles form an equivalent p-n junction [8], the key component behind any electronic circuit. It is worth emphasizing that since p-n junctions are built on a pristine sheet of graphene, they preserve the main characteristics of the material.

At this preliminary stage it is hard to predict which of these strategies will prevail in the electronics market, and in how

much time; we embrace the basic principle that an efficient use of graphene should inevitably exploit its intrinsic properties rather than trying to change them. That brought our attention to the second class of methods, i.e., the implementation of digital circuits based on electrostatically controlled graphene p-n junctions.

Graphene p-n junctions can serve as basic switch for a complex logic gate, called RG-MUX (Reconfigurable Gate MultipleXer) because it implements the functionality of a multiplexer. RG-MUXes use a wide graphene sheet (around 190nm) due to which the material defects are within the limits [9, 10]. The authors of [11] propose the realization of various logic gates using RG-MUXes and show that these gates have superior performance and smaller area than traditional CMOS-based ones. RG-MUXes have been characterized for timing and power [12, 13], and various synthesis and design styles that exploit graphene p-n junction [14] and MUX-based [15] have also been investigated [16, 17, 18]. In spite of its great advantage in terms of speed and area, however, RG-MUXes have the drawback of being less energy-efficient than equivalent CMOS gates [19]. This is due to a larger gate capacitance, a consequence of the larger gate area. Energy benefits can be obtained only indirectly thanks to the smaller size of such devices, which allow shorter interconnects [19].

In this work, we revisit RG-MUXes with the objective of making them also energy-efficient, while preserving the highly desirable characteristics of graphene. We do this by recognizing that RG-MUXes, for both structural and functional reasons, naturally lend themselves to implement logic elements that can operate adiabatically. Adiabatic logic [20], aims at mimicking an adiabatic, i.e., without energy exchange, computation process in digital logic. The basic idea of reducing energy dissipation during the switching process relies on the use of a variable (ramp) power supply to recycle a portion of the energy from the load capacitance. Although regarded as a mostly theoretical and somehow exotic computational style, research on the topic has been constantly active over the years, providing several demonstrations of working implementations [21, 22, 23]. The basic building block of any adiabatic circuit is the adiabatic amplifier, a buffer that uses adiabatic charging to drive a capacitive load. In its traditional embodiment, the adiabatic amplifier is implemented using transmission gates (TGs) [20].

Extending a previous work [24], this paper provides three main new contributions: (i) we show that an RG-MUX, by properly assigning voltage values at its terminals, can naturally operate as an adiabatic amplifier, and that each of its two p-n junctions can be regarded as TGs; (ii) we also characterize the figures of merit of these gates and compare them against both adiabatic and non-adiabatic CMOS implementations; (iii) we show how to build fast and extremely low-power adiabatic logic gates based on graphene p-n junctions. Simulation results obtained from SPICE simulations using a dedicated Verilog-A model for the p-n junction, show that the proposed adiabatic gates are 1.5X more power efficient than equivalent adiabatic CMOS-based implementations, still showing more than one order of magnitude better power-delay product. We also demonstrate that such devices are able to operate with about 4X less power

and about 4.5X improvement in power-delay product with respect to non-adiabatic counterparts.

2. Background

2.1. Graphene physics

Theorized since 1947 [25], but isolated for the first time only sixty years later [26, 27, 28], graphene is the most surprising allotrope of carbon. It consists of a one-atom-thick sheet of graphite where all the carbon atoms form covalent bonds in a single plane. The resulting two-dimensional (2D) structure is made up of atoms packed in a hexagonal crystal lattice with a carbon-carbon bond length a_{C-C} of 1.42Å. This gives graphene special electrical properties [29] that are a direct expression of its unique energy band structure.

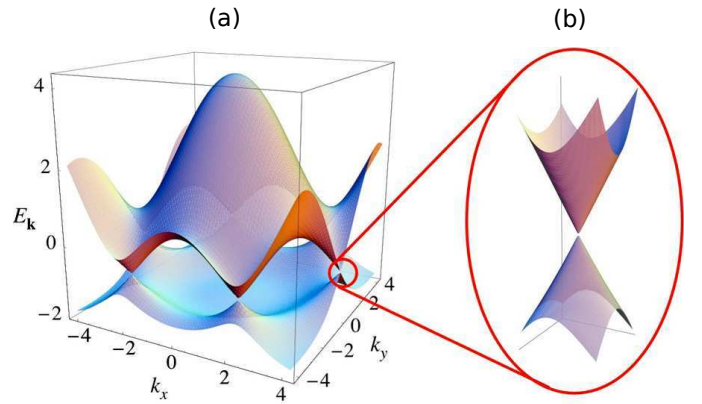


Figure 1: Graphene energy band structure (a) and the linear energy dispersion around the Fermi Energy (b). Image credits to [29].

As for any solid material, it is possible to derive an analytical expression of the electron dispersion by solving the time-independent Schrödinger's equation through the periodic potential of the lattice [30, 31]. An exact solution is computationally challenging, but reasonable and accurate approximations are possible. We refer here to the solution proposed by [32], in which the equation describing the energy dispersion E of graphene is given in (1) and plotted in Figure 1-a. Interested readers can refer to [32, 33] for a formal derivation, which is out of the scope of this work.

$$E^\pm(\mathbf{k}) = \pm \gamma \sqrt{1 + 4 \cos\left(\frac{\sqrt{3}a}{2}k_x\right) \cos\left(\frac{a}{2}k_y\right) + 4 \cos^2\left(\frac{a}{2}k_y\right)} \quad (1)$$

The positive and negative energy branches are the conduction bands (E^+) and valence bands (E^-), respectively; the vector $\mathbf{k} = \{k_x, k_y\}$ represents the 2D wavevector; γ is a fitting parameter whose value can range from 2.7eV to 3.3eV; $a = \sqrt{3}a_{C-C}$ is the side length of the parallelogram representing the primitive cell in the Bravais lattice [32].

From the plot, we notice that the conduction and valence curves touch each other near the edges of the Brillouin Zone, i.e., at

zero energy, where the Fermi energy E_F passes (Figure 1-b). This gapless spectrum provides graphene with semi-metallic properties, different from metals (where E_F is in the conduction band) and semiconductors (where E_F falls in the bandgap). As a consequence, graphene can only implement a weak “OFF” state. Recent works have shown, however, the possibility to implement equivalent p-n junctions by means of electrostatic doping. Those p-n junctions can be thereby used to implement logic gates.

2.2. Graphene p-n junction

Figure 2 shows the basic structure of a graphene p-n junction that uses electrostatic doping. Such device is composed of four layers, namely: (i) the bottom layer, that includes two split gates (referred as back gates in the figure) made of conductive material separated by a distance D ; (ii) an insulating layer of oxide, which is placed on top of the split gates; (iii) a wide graphene sheet; and, (iv) two electrodes (front metal contacts in the figure) placed on top of the graphene sheet, which serve to supply a reference current to the device.

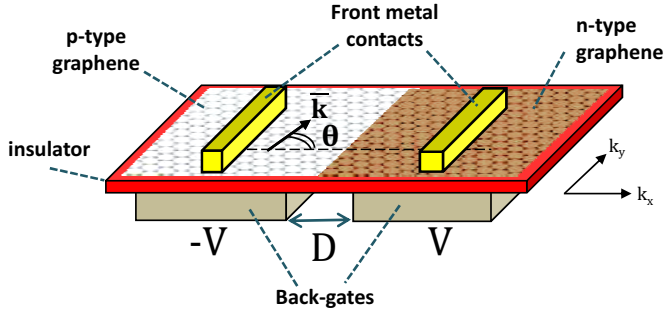


Figure 2: Graphene p-n Junction.

The application of a negative voltage on a back-gate shifts E_F towards the valence band resulting in p-type graphene in the above region. On the other hand, a positive voltage shifts E_F towards the conduction band leading to n-type graphene [34]. In this way, by applying asymmetric voltages to the two back gates a p-n junction is formed [35]. The front metal contacts represent the conceptual source (left) and the probe (right) that are emitting and receiving carriers.

As demonstrated in [36], carriers injected in the p-region through the left front contact cross the potential barrier at the p-n junction with a transmission probability $T(\theta)$ which depends on two parameters: the angle θ between the electron's wave vector \bar{k} and the normal of the junction, and the width D of the p-n transition region. The analytical expression of $T(\theta)$ is:

$$T(\theta) = \cos^2(\theta)e^{-\pi k D \sin^2 \theta} \quad (2)$$

The transmission probability $T(\theta)$ is thus 1 for carriers that travel orthogonally with respect to the junction ($\theta = 0$), regardless of D , and decreases exponentially for larger values of θ , e.g., $T(\theta) = 0$ for $\theta = \pi/2$. Notice that when a symmetric control voltage is simultaneously applied at two adjacent back-gates ($(+V, +V)$ or $(-V, -V)$), the graphene layer is entirely of p- or n-type, respectively; the device is thus transparent to the charge flow.

2.3. Electrical model of graphene p-n junction

Fig. 3 shows a schematic of the electrical model implemented borrowing the works of [11, 19]. Terminals A and Z denote the left and right front contacts respectively, whereas S and U denote the back gate potentials on the two gates. The resistor R_{AZ} represents the resistive equivalent of a graphene p-n junction between input A and output Z. Its value ranges from $R_{ON} = 300\Omega$, under p-p or n-n configurations, to $R_{OFF} = 10^7\Omega$, under p-n or n-p configurations. The analytical expression for the junction resistance is expressed as:

$$R_{AZ} = \frac{R_0}{N_{ch}T(\theta)} \quad (3)$$

In (3), the transmission probability of the carriers across the junction is given by (2), $R_0 = \frac{h}{4q^2}$ is the quantum resistance per propagation mode, and N_{ch} is the number of excited propagation modes [37]. The electrical model also includes parasitics of the metal contacts. The resistors R_c at the front pins A and Z model the resistance of the metal-to-graphene contacts [38]. The lumped capacitance C_g at the back-gates, i.e., C_{gs} at S and C_{gu} at U, consists of the series of the oxide capacitance C_{ox} and the quantum capacitance of the graphene sheet C_q , namely $C_g = 1/(C_{ox}^{-1} + C_q^{-1})$. C_{gU} and C_{gS} denote the gate capacitances at the back gates S and U respectively. This electrical model was implemented in Verilog-A and included in our SPICE simulations.

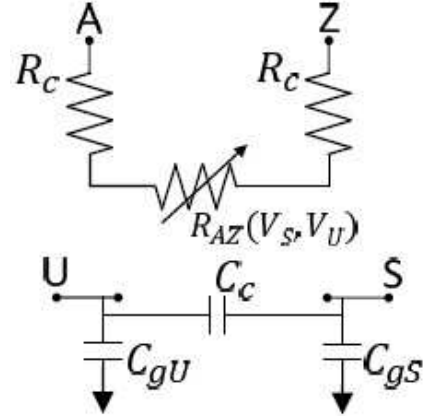


Figure 3: Electrical model of graphene p-n junction.

Fig 4 shows the variation of the junction resistance obtained through SPICE simulations; there are two curves: one with the potential of back-gate U fixed at either a positive potential ($+V_{dd}/2$, cross mark) and at a negative potential ($-V_{dd}/2$, plus mark). Let us consider the curve with the potential of U fixed at $+V_{dd}/2$ and let us vary the potential of the other back-gate S from $-V_{dd}/2$ to $+V_{dd}/2$. As it can be seen, at $-V_{dd}/2$, the junction resistance is very large, and as it approaches $+V_{dd}/2$ the resistance settles at few hundred Ω s. In practice when the two back-gates have the same polarity the junction resistance is small.

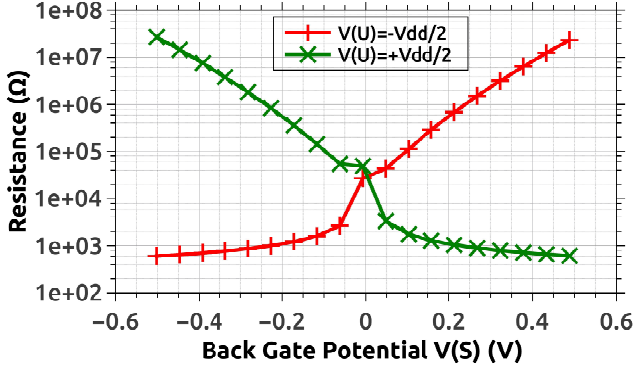


Figure 4: Junction resistance variation with the back gate potential

2.4. RG-MUX: structure and functionality

By using the transmission properties of the p-n junction previously described, it is possible to connect two junctions together as to implement a voltage controlled selector, which is called RG-MUX (Re-configurable Graphene MUX). Figure 5 shows the 3D (a), bottom (b) and top (c) views of the RG-MUX as introduced in [11]. It consists of two back-to-back p-n junctions: j_1 (left), controlled by back-gate pair $S-\bar{U}$, and j_2 (right) controlled by back-gate pair $S-U$.

The back-gates \bar{U} and U are driven with fixed and opposite voltages, i.e., $\bar{U} = -V_{dd}/2$ and $U = +V_{dd}/2$. This forces the graphene sheet to be p-type on the left side and n-type on the right side¹. The shared back-gate S in the center serves as logic input.

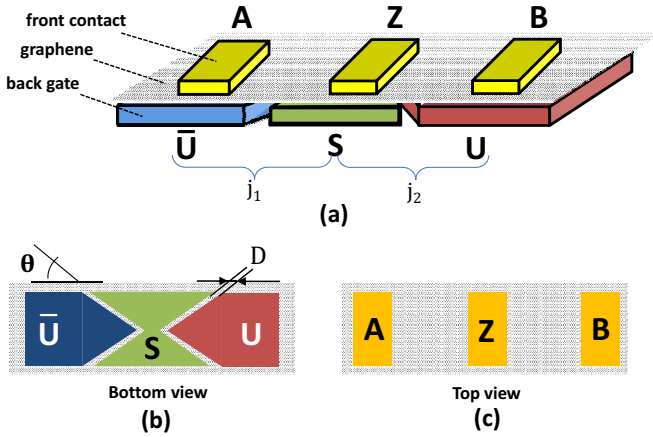


Figure 5: RG-MUX Structure.

Figure 6 shows a conceptual model of the RG-MUX depending on the voltage assignments at its back-gate inputs.

When $S = "0"$, the central graphene becomes p-type. This forms a p-p "junction" on the left, i.e., a zero barrier potential, and a p-n junction on the right, i.e., a barrier potential through which carriers are transmitted with a transmission probability $T(\theta)$. It is worth noticing that θ is physically imposed at 45° by construction (back-view in Figure 5). The resulting p-p-n doping

¹In this implementation, $+V_{dd}/2$ and $-V_{dd}/2$ correspond to logic "1" and logic "0" respectively

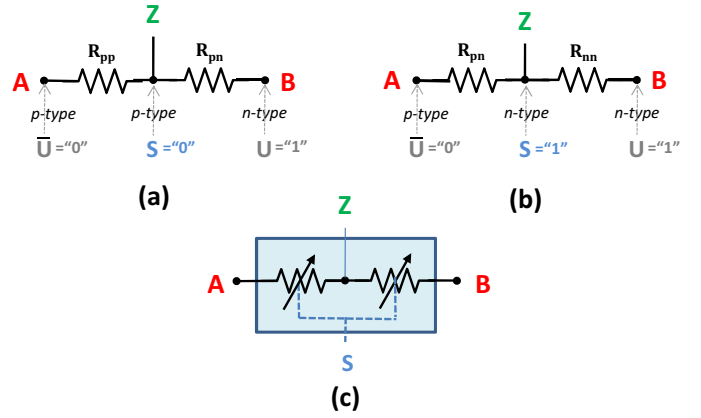


Figure 6: Operations of the RG-MUX: The Two Possible Scenarios Depending on the Voltage at the Back-Gates (a) and (b), and a Conceptual Model (c).

profile of the graphene sheet creates a low-resistive path R_{pp} between the front contacts $A-Z$ and a high-resistive R_{pn} path between the front contacts $B-Z$ (Figure 6-a). This forces the output Z to follow the input signal associated with the lowest resistance, i.e., $Z = A$.

Similarly, when $S = "1"$ the central graphene region is n-doped, leading to a p-n-n doping profile of the graphene layer. Therefore a low-resistive path between the contacts $B-Z$ makes the output to follow B , i.e., $Z = B$ (Figure 6-b).

From the above analysis, the RG-MUX can be seen as a 3-input, 1-output device, where A and B are the data input, S the control input, and Z the output. Depending on the polarity of S the resistances on the two input-to-output paths are properly set (refer to Figure 6-c). From a functional point of view, the RG-MUX implements a multiplexer, where S is the selection input, i.e., $Z = S \cdot B + \bar{S} \cdot A$.

3. Adiabatic Computing

The development of adiabatic logic in CMOS technology has been widely investigated in the past decades [39]. However, several of the proposed techniques show some critical downsides, mainly regarding: (i) increase in logic design complexity; (ii) the need of multiple supplies for proper interfacing between stages; and (iii) a self-charging effect at the output nodes, which is due to leakage currents. In the proposed graphene-based implementation the first two concerns are removed by construction, whereas only the third issue remains open. In this section, we employ a single-supply adiabatic strategy that allows the realization of compact logic circuits operating at low power.

3.1. Adiabatic switching

Consider the application of a step voltage source (from 0 to V_{DD}) to an RC circuit (Figure 7-a). The total energy supplied by the voltage source is given by $E_{supplied} = Q \cdot V_{DD} = CV_{DD}^2$. The energy stored in a capacitor when the output has reached the final value is given by $E_C = \frac{1}{2}CV_{DD}^2$, i.e., only half of the supplied energy. The other half is dissipated across the resistor in the form of heat (E_R). This scenario represents the traditional "computation" model, in which energy waste is maximum.

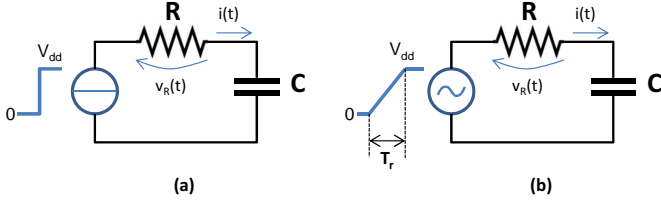


Figure 7: Traditional (a) and Adiabatic (b) Charging.

In contrast, adiabatic logic uses a slow-rising ramp supply signal with transition time T_r to change from 0 to V_{DD} (Figure 7-b). Assuming that T_r is such that the capacitor is able to charge instantaneously to the input supply voltage, the current is given by:

$$i(t) = C \frac{dv(t)}{dt} = \frac{CV_{DD}}{T_r} \quad (4)$$

The energy dissipated across the resistor is obtained by integrating the power across the resistor over time T_r :

$$E = \int_0^{T_r} v_R(t)i(t)dt = \int_0^{T_r} R \frac{C^2 V_{DD}^2}{T_r^2} dt = \frac{RC}{T_r} CV_{DD}^2 \quad (5)$$

where $v_R(t)$ is the voltage drop across the resistor.

The minimum transition time for which the step and ramp source produce the same energy dissipation can be computed from (5) and the expression of $E_{supplied}$. For $T_r > 2RC$, adiabatic circuits are more energy efficient than regular circuits with step input [40].

3.2. The adiabatic amplifier

The adiabatic amplifier [20] is the basic building block used in adiabatic circuits; from a functional point of view it is a simple buffer that uses adiabatic charging to drive a capacitive load. In traditional CMOS implementations, the typical realization uses transmission gates (TGs) as shown in Figure 8. The choice of TGs is because of their fully-restoring feature, i.e., the NMOS transistor is used to pass logic “0” whereas the PMOS is used to pass logic “1”, so the output is always strongly driven and the levels are never degraded.

The basic operations of the adiabatic amplifier is quite intuitive. Depending on the values of X and \bar{X} , one of the output capacitances will be adiabatically (through V_A) charged. Clearly, the two sides are mutually exclusive and only one side drives the supply voltage to the output².

The circuit uses dual-rail encoding for both inputs (because driving the TGs require the double polarity of X) and outputs; in the latter case this is required for interconnecting the amplifier to other adiabatic circuit elements. Operations occur in three phases, corresponding to the three “regions” of the ramp input. The input is first set to a stable value; then the “charge” phase starts, in which the supply V_A is ramped and the load capacitor is adiabatically charged. In the second “evaluate” phase, when

²The original implementation of the adiabatic amplifier includes two NMOS clamp at the two outputs driven by X and \bar{X} , which are not shown here.

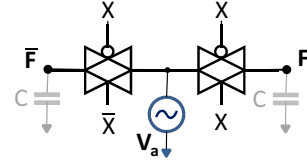


Figure 8: Adiabatic Amplifier Implemented using TGs

V_A has reached the final value, output voltages are stable and can be used by next logic stages. In the third “recovery” phase, the load capacitor discharges back into V_A as it ramps down to 0V.

4. Adiabatic through RG-MUXes

4.1. RG-MUX as an adiabatic amplifier

If we compare the RG-MUX of Figure 5 and its model in Figure 6 with the adiabatic amplifier of Figure 8 we can immediately spot many similarities. Both devices operate as selectors with variable resistance from inputs to outputs: in the RG-MUX, signal S determines the low- and high-resistance paths, whereas in the adiabatic amplifier it is represented by the signal X . There is, however, a significant difference between the two devices, namely the flow of “computation”, which occurs in opposite directions. In the RG-MUX the central node (Z) is an output, whereas in the adiabatic amplifier the central node is connected to the ramped supply voltage. The opposite occurs for the two side nodes (A and B) which are inputs in the RG-MUX and outputs in the adiabatic amplifier.

Figure 9 shows the input and output signals assignment of an RG-MUX so that it can be used to implement the operations of an adiabatic amplifier.

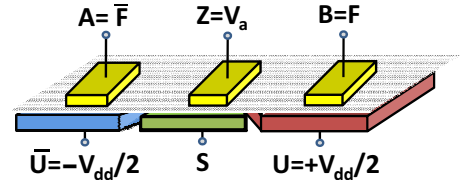


Figure 9: Adiabatic Amplifier using the RG-MUX.

Specifically, we notice how signal Z (an output in a RG-MUX) is now connected to the variable supply voltage V_a , whereas pins A and B are now the dual-rail encoded values of the output. The role of the back-gate input S (together with the specific encoding of U and \bar{U} , which are normally fixed) is similar to the role of X and \bar{X} signals in the amplifier: $S = “0”$ (“1”) is equivalent to $X = 0$ ($X = 1$), that is, output \bar{F} (F) is charged by V_a .

4.2. Implementing adiabatic logic gates

The adiabatic amplifier is the simplest of many possible adiabatic logic gates and functionally corresponds to a buffer. In order to implement complex logic functions, we adopt the architectural template of a generic logic function of [20], in which a

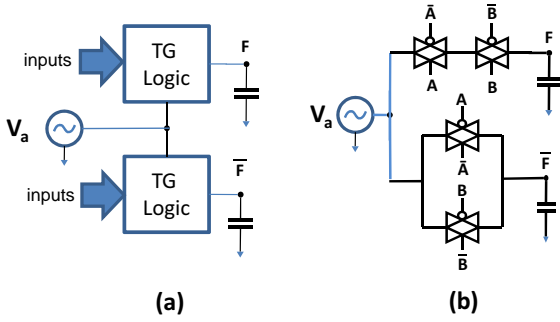


Figure 10: General Architecture of an Adiabatic Logic Function (a) and an AND/NAND gate (b).

parallel of the structure of classical CMOS gates is maintained. In practice, the complementary pull-up and pull-down networks of CMOS are replaced by two equivalent networks of TGs. The new pull-up network is used to charge the output F , whereas the pull-down one will charge the complemented output \bar{F} (Figure 10-a). Figure 10-b shows an example of a logic gate which realizes an AND/NAND function.

The implementation of arbitrarily complex logic functions requires the availability of the graphene counterpart of a TG. Since the RG-MUX is essentially a side-by-side combination of two junctions (in the p-n/n-n or p-p/p-n pattern), we can use a single graphene p-n junction (“half” RG-MUX) to implement a TG. Figure 11 illustrates the basic concept of this equivalence.

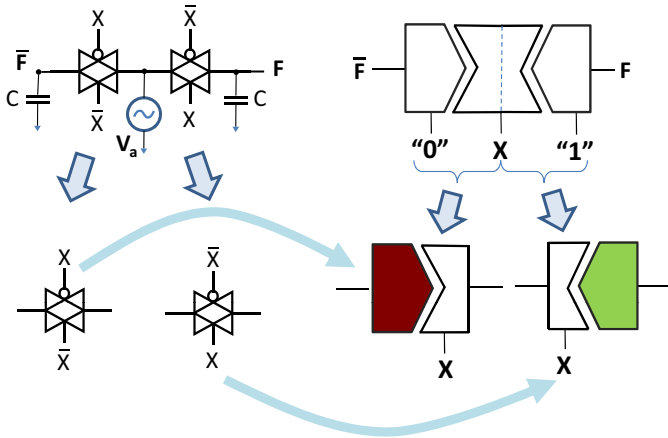


Figure 11: P-N Junction based TG Realization

The symbol on the top right corner of the figure denotes the schematic of an RG-MUX used as an adiabatic amplifier (as in Figure 9), and mimics the structure of the diagonally-cut back-gates. Values “0” and “1” represent the encoding of \bar{U} and U . We conceptually “cut” this device (and its corresponding symbol) into two one-sided graphene p-n junctions, denoted by the double-wedged symbols in the bottom left corner. The horizontal pins represent the I/O signals, whereas the bottom pin represent the back-gate signal equivalent to pin S of the RG-MUX. For ease of readability we do not explicitly represent the

\bar{U}/U pins, which are assumed to be fixed at “0” and “1” respectively. The symbol on the left (dark red back-gate) denotes $\bar{U} = “0”$, i.e., $-V_{dd}/2$, and is therefore equivalent of the TG with signal transmitted when $X = 0$; the symbol on the right (light green back-gate) denotes $\bar{U} = “1”$, i.e., $V_{dd}/2$, and corresponds to the TG transmitting signal with $X = 1$. In other words, by splitting the RG-MUX, one can obtain two TGs with positive and negative polarity, with no need of complementing the input signal X .

Using the general architecture of Figure 10, the basic logic gates (INV/BUF, AND/NAND and OR/NOR) based on these TGs can be implemented as shown in Figure 12. V_a denotes the ramped-supply voltage that enables adiabatic operations. Two p-n junctions are needed to realize an INV/BUF, whereas four junctions are needed for AND/NAND and OR/NOR.

For the sake of clarity, we depict in Figure 13 the signal waveforms of a 2-input AND (left plot) and of a 2-input OR (right plot). Signal $V(a)$ represents the power supply, $V(X)$ and $V(Y)$ represent the two input signals, and signals $V(F)$ and $V(Fbar)$ represent the true and complemented outputs of the circuits. Let us consider the plot on the left. When the two input signals are both at logic 1, i.e., the AND operator evaluates to 1, the supply signal $V(a)$ is propagated to the output; whereas, in the other three cases, namely when the AND operator evaluates to logic 0, $V(a)$ is not allowed to propagate. In other words, a graphene p-n junction based circuit propagates the supply signal to the output node if, and only if, the logic function evaluates to logic 1. This is an example of the intrinsic *pass-through* characteristics of graphene p-n junctions. As the plot suggests, an important aspect to be considered is the leakage-induced charging of the output nodes, referred as the *Self-Charging* effect. Consider the situation in which the input X/\bar{X} is kept constant for several cycles. The I_{off} current drained from V_a charges the output capacitance. The effect is plotted in Figure 13, where \bar{F} (plot $v(Fbar)$ in the picture) slowly charges to intermediate voltage levels. After some cycles, this will cause a bit-flip at the output node. A quantitative analysis of the output voltage refresh rate is presented in the Section 5.

5. Simulation results

5.1. Power dissipation across RC-circuits

The first experiment compares the energy benefits obtained by adiabatic operations on graphene-based devices against those achieved with CMOS. More precisely, we measured the power dissipation for the RC circuit shown in Figure 7 for both technologies. The CMOS configuration consists of a TG designed using minimum size PMOS and NMOS available in a 40nm technology library provided by STMicroelectronics, driving a capacitance corresponding to the input pin/gate capacitance of a minimum sized inverter, i.e., $0.6fF$. The equivalent resistance is $2.56k\Omega$, obtained by the parallel of the resistances of the p- and n- networks in the linear region of operation. The graphene configuration consists of a graphene p-n junction driving a capacitance equal to the back-gate capacitance of an RG-MUX, i.e., $0.62fF$. The resistance is that of the junction (including

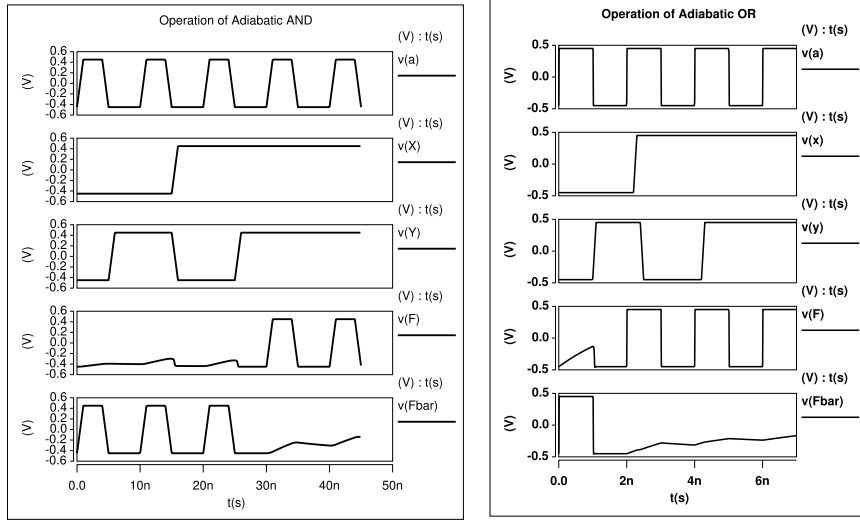


Figure 13: Operation of the Adiabatic AND/NAND & OR/NOR

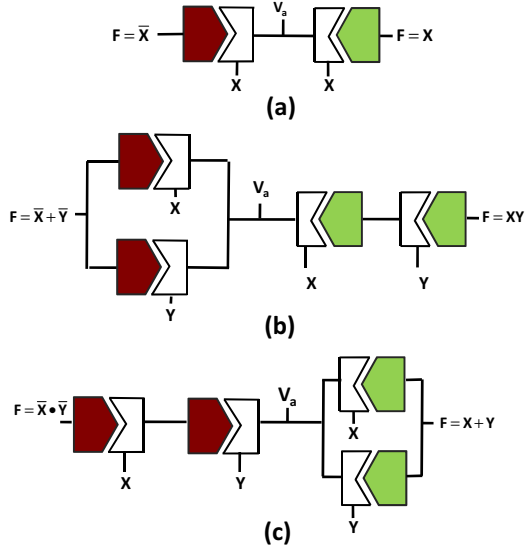


Figure 12: Adiabatic Logic Gates: INV/BUF (a), AND/NAND (b), OR/NOR (c).

the contact resistance), i.e., $1.38k\Omega$. In both cases, power dissipation after application of the ramp input signal is estimated through the power command available in HSPICE over a signal transition at the output node. Experimental results are plotted in Fig 14.

We notice that for very small values of T_r (< 10 ps) the equivalent CMOS network is more energy-efficient than the graphene one. However, such small values of T_r actually correspond to non-adiabatic operations. As a matter of fact, even the steepest transition in a conventional non-adiabatic circuit takes a few picoseconds of rise/fall time. Such superiority corresponds to the intrinsic energy efficiency of CMOS over graphene-based devices mentioned in Section 1. For $T_r > 10ps$, graphene and CMOS initially exhibit similar power consumption, before graphene starts consuming less power. Notice that, although the

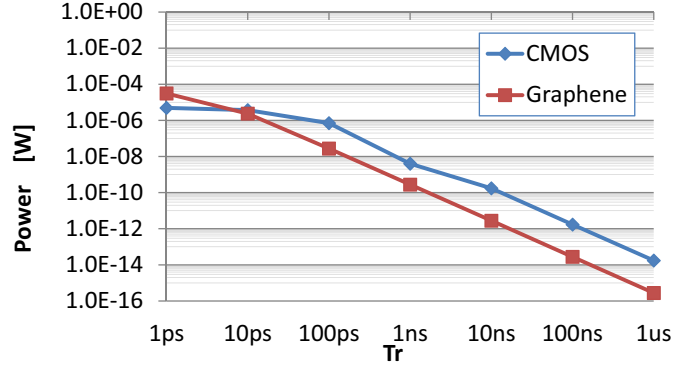


Figure 14: Power Dissipation of Adiabatic RC Circuit vs. T_r : CMOS vs. Graphene-Based.

truly adiabatic region for graphene starts at $T_r = 2 \cdot R \cdot C \approx 1.7ps$, power consumption is already lower for smaller T_r values. On average, there are about two orders of magnitude consistently over the entire range of T_r values greater than $10ps$.

Please note that we reported a large range of T_r values although some of the larger values in the μs range are clearly impractical but for some specific low-throughput applications.

5.2. Power characterization of the adiabatic amplifier

We characterized the power consumption of the adiabatic amplifier using HSPICE, by assessing in particular the breakdown of power between ON and OFF states in the two junctions. Figure 15 shows these two components as a function of T_r . The plot reports the average power P_{on} and P_{off} of the two p-n junctions of the amplifier. Subscripts “on” and “off” denote the fact that the two junctions have complementary states.

The figure clearly shows that both components of power dissipation decreases as T_r increases, yet with different scales. While total power (sum of P_{on} and P_{off}) exhibits about seven orders of magnitude power reduction over the range of T_r values

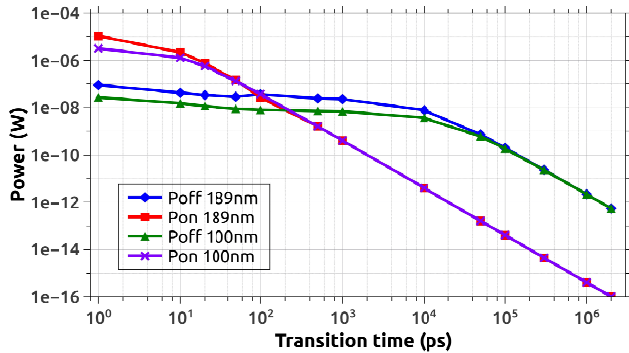


Figure 15: On vs. Off Power in Adiabatic Amplifier.

reported, for smaller T_r , power is determined by active power P_{on} , whereas around $T_r = 100ps$ the two components have similar weight, and for larger T_r values the off current (junction leakage) becomes dominant. Experiments were replicated for different device width (W), i.e., the size of the p-n junction along the k_y axis depicted in Figure 2. In particular, we characterized devices assuming $W = 189nm$ and $W = 100nm$. As can be seen from the plot, for lower values of T_r , increasing the device width results in increased off-state power. Also, increasing device width decreases the junction resistance, as per (3), and increases the average current across the branch. However, for large values of T_r , the signal at \bar{F} charges towards logic “1” irrespective of the junction resistance making constant off-state current. Similarly, at low T_r the power in the on-state increases with device width due to the increase of the average current across the junction and remains constant for larger values of T_r .

5.3. Relation with operating frequency

An important concern to be considered is the relation of T_r with the operating frequency of the circuit. Assuming a 50% duty cycle for the variable power supply. The load capacitor has to completely charge to its final value in the given duty cycle and remain in that state for a significant time. Therefore, although power savings are maximum for large T_r 's, there exists an optimal T_r for each frequency, calculated as the value that allows the capacitor to charge to its final value in the middle of the duty cycle. As an intuitive rule of thumb, higher frequencies imply smaller T_r values. Calculation of such optimal value yields the following expression of T_r (in seconds) vs. frequency (in Hz): $T_{r,opt}(f) = \frac{1}{2 \cdot f}$.

5.4. Leakage currents and self-charging effect

To avoid the self-charging effect (see Section 4), adiabatic circuits must periodically refresh the output voltages. Since I_{off} increases for increasing T_r , the lower the operating frequency the faster the charging of the capacitor and thus the more frequent the need of refreshing the voltage. The actual refresh frequency depends on what level of the output voltage should be considered as *critical*, e.g., some percentage of V_{dd} .

The time taken to reach that level, expressed in cycles at the desired frequency, will determine the feasibility of the operations. Since we can refresh at most once per cycle, this time should be

Table 1: Refresh Time in Cycles.

f [MHz]	% of V_{dd}				
	10%	20%	30%	40%	50%
1000	1.30	2.55	4.65	9.54	∞
500	0.52	1.42	2.48	4.54	∞
100	0.24	0.36	0.45	0.54	1.52
10	0.09	0.15	0.21	0.26	0.32
1	0.06	0.11	0.17	0.22	0.28

greater than 1 cycle. Tab 1 reports the results of an exploration of feasibility by spanning different frequency values ($1 \div 1000$ MHz), and different voltage thresholds (up to 50% of V_{dd} in steps of 10%). Entries in bold are the feasible ones. We can observe that very low frequencies are not feasible even for large voltage degradation at the output, whereas frequencies in the order of a few hundred of MHz are mostly feasible. Using the relation between T_r and f of Section 5.3, this corresponds to T_r 's in the order of $1ns$, in the region in which there is an adiabatic benefit over CMOS (please refer to Figure 14).

5.5. Characterization of simple logic functions

We realized 46 simple logic functions using the adiabatic approach presented in this paper and a classical standard cell approach which is non-adiabatic. The characterization was done for a range of transition time starting from $1ps$ to $1ns$. We compare the characterization results of graphene technology with that of conventional CMOS technology. Adopted logic functions are summarized in Table 2. We limited our experiments to functions having no more than six devices in series. We then compared the following four implementations:

Adiabatic graphene: logic functions are realized according to the methodology described in this paper, i.e., each p-n junction acting as a transmission gate. A ramp signal is fed to the common intersection point of the dual logic design style.

Adiabatic CMOS: logic functions are implemented using MOSFET based transmission gates, i.e., parallel connection of NMOS/PMOS. Also in this case a ramp signal is used to exploit the adiabatic charging principle.

Non-adiabatic graphene: logic functions are first synthesized using the ABC synthesis tool using a subset of all possible gates, namely INV, BUF, AND, OR. Then, each gate is replaced with the corresponding RG-MUX appropriately configured.

Non-adiabatic CMOS: logic functions are created using standard complementary MOS architectures, namely, by means of pull-up/down networks connected to V_{dd}/GND ; the pull-up network is implemented using PMOS whereas the pull-down network is realized using NMOS with proper size.

The simulations for estimating power/performance was carried out using Synopsys HSPICE. For graphene technology we use the electrical model presented in Section 2.3 written in Verilog-A, whereas for CMOS technology we use the models from the ST-Microelectronics at 40nm technology node. The load capacitance was fixed at $0.6fF$, which is the minimum gate capacitance in both technologies.

5.5.1. Power characterization

Power dissipation results for the logic functions in Table 2 is plotted in Figure 16. We report the average power dissipation

Table 2: Some Boolean functions with no more than six series P-N junctions.

Function	Function	Function	Function
F00	\bar{a}	F23	$a + (b \odot d) \cdot c$
F01	$a \odot b$	F24	$(a \odot d) + (b \odot d) \cdot c$
F02	$a + b$	F25	$a + (b \odot d) \cdot (c \odot d)$
F03	$a \cdot b$	F26	$(a \odot d) + ((b \odot d) \cdot (c \odot d))$
F04	$(a \odot b) + c$	F27	$(a \odot d) \cdot b \cdot c$
F05	$(a \odot b) \cdot c$	F28	$(a \odot d) \cdot (b \odot d) \cdot c$
F06	$(a \odot b) + (a \odot c)$	F29	$(a \odot d) \cdot (b \odot d) \cdot (c \odot d)$
F07	$(a \odot b) \cdot (a \odot c)$	F30	$(a \odot d) + (b \odot e) + c$
F08	$(a \odot b) + (c \odot d)$	F31	$(a \odot d) + (b \odot d) + (c \odot e)$
F09	$(a \odot b) \cdot (c \odot d)$	F32	$((a \odot d) + (b \odot e)) \cdot c$
F10	$a + b + c$	F33	$((a \odot d) + b) \cdot (c \odot e)$
F11	$(a + b) \cdot c$	F34	$((a \odot d) + (b \odot d)) \cdot (c \odot e)$
F12	$a + (b \cdot c)$	F35	$((a \odot d) + (b \odot e)) \cdot (c \odot d)$
F13	$a \cdot b \cdot c$	F36	$(a \odot d) + ((b \odot e) \cdot c)$
F14	$(a \odot d) + b + c$	F37	$a + ((b \odot d) \cdot (c \odot e))$
F15	$(a \odot d) + (b \odot d) + c$	F38	$(a \odot d) + ((b \odot e) \cdot (c \odot e))$
F16	$(a \odot d) + (b \odot d) + (c \odot d)$	F39	$(a \odot d) + ((b \odot e) \cdot (c \odot d))$
F17	$((a \odot d) + b) \cdot c$	F40	$(a \odot d) \cdot (b \odot e) \cdot c$
F18	$((a \odot d) + (b \odot d)) \cdot c$	F41	$(a \odot d) \cdot (b \odot d) \cdot (c \odot e)$
F19	$((a \odot d) + b) \cdot (c \odot d)$	F42	$(a \odot d) + (b \odot e) + (c \odot f)$
F20	$((a \odot d) + (b \odot d)) \cdot (c \odot d)$	F43	$((a \odot d) + (b \odot e)) \cdot (c \odot f)$
F21	$(a + b) \cdot (c \odot d)$	F44	$(a \odot d) + ((b \odot e) \cdot (c \odot f))$
F22	$(a \odot d) + (b \cdot c)$	F45	$(a \odot d) \cdot (b \odot e) \cdot (c \odot f)$

over the 46 logic functions versus the input signal transition time. First key observation is that, irrespective of the technology, non-adiabatic implementations exhibit higher power dissipation compared to adiabatic ones. However, non-adiabatic graphene technology has the least power dissipation due to the smaller junction resistance compared to that of a MOSFET. For the adiabatic implementations, graphene technology has smaller power dissipation over CMOS technology over the entire range of T_r , with a maximum gain of about 1.5 orders of magnitude at $T_r = 1ns$.

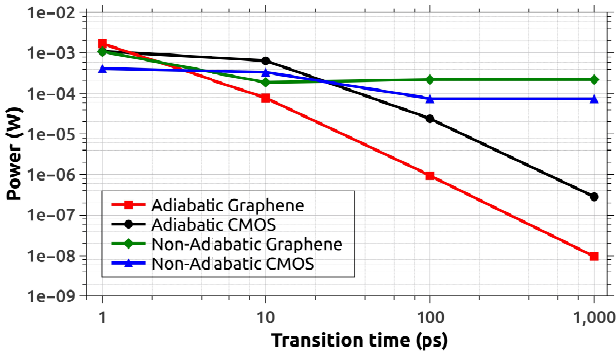


Figure 16: Average Power of logic functions in Table 2.

5.5.2. Power-delay product characterization

Figure 17 depicts the power delay product (PDP) of adiabatic and non-adiabatic implementations in graphene and CMOS technologies. The PDP plotted is averaged over the 46 logic functions for a given transition time. The objective of this plot is to show that power savings in adiabatic implementations do not come at the expense of performance loss. As can be seen

from the plot, the PDP of graphene-based adiabatic implementations is more than an order of magnitude smaller than that of adiabatic CMOS. However, in the non-adiabatic approach it is the CMOS technology that has better PDP than a graphene-based implementation.

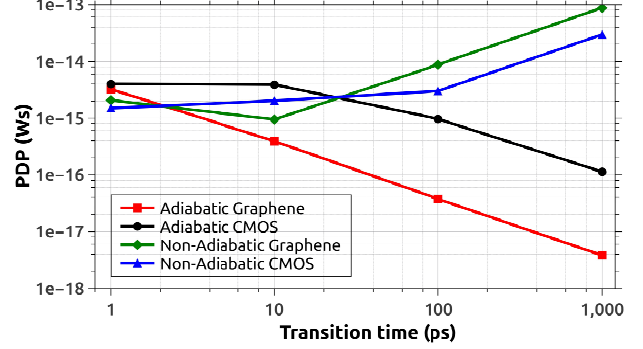


Figure 17: Average Power Delay Product (PDP) of logic functions in Table 2.

5.6. A CAD tool for adiabatic graphene circuits

The logic functions described in the previous section were built manually. In order to allow an automated synthesis usable for realistically-sized designs, we implemented a CAD tool that maps a generic Boolean function onto a graphene-based p-n junction technology. We adopted Binary Decision Diagrams (BDDs) [41, 42] as a common data structure to represent the functions of the circuits. A generic node of a BDD (Figure 18-a) has two in-going edges f_x and f'_x , representing the positive and negative co-factors of the function, respectively. They propagate the input ramp-signals generated by the two terminal nodes 0 and 1 to the out-going edge, labeled f . Such evaluation is based on the value assumed by a generic primary input x associated to the decision node.

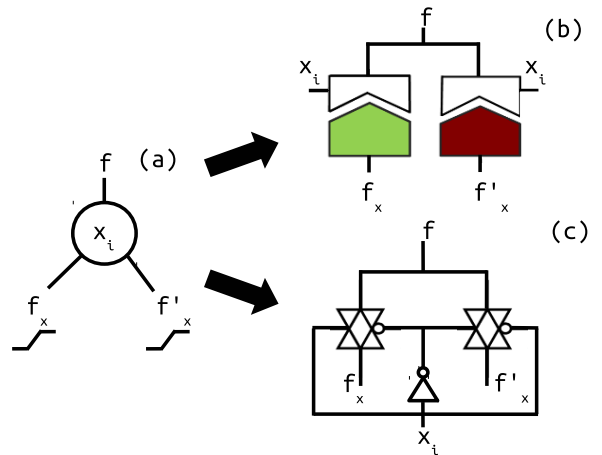


Figure 18: Decision node in BDD. (a) Logic structure, (b) PN-BDD realization, and (c) TG-BDD realization.

For a fair comparison, we compared the two following implementations:

- **PN-BDD:** each internal BDD node is mapped with graphene p-n junctions transmission gates, as showed in

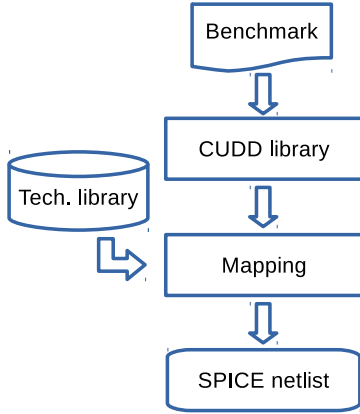


Figure 19: Design flow.

Figure 18-b.

- **TG-BDD**: each internal BDD node is mapped with MOSFET transmission gates, as depicted in Figure 18-c.

In both cases we exploit the adiabatic charging principle to achieve the minimum power consumption. The adopted design flow is illustrated in Figure 19. Each benchmark is first processed with the open-source CUDD library [43] in order to generate the corresponding BDD structure; then, a TCL script maps each decision node in the corresponding technology, i.e., PN-BDD or TG-BDD. Node descriptions are stored in a technology library which is fed to the TCL mapping script. The output is a SPICE-compliant netlist.

Table 3 summarizes the results we obtained on a set of open-source benchmarks. Columns **PI** and **PO** show the total number of primary inputs and outputs of the design and the total number of nodes in the BDD structure. Columns **PN-BDD** and **TG-BDD** report, for the two implementation styles, the number of p-n junctions (for the PN-BDD) and transistors (for the TG-BDD) as well as the resulting area.

For the PN-BDD structure, the equivalent area of a single BDD node is given by the sum of the two p-n junctions (refer to Figure 18), which corresponds to $0.382\mu\text{m}^2$ [11]. The total area occupation for a given circuit is simply obtained by multiplying this number by the number of BDD nodes. The transmission gate implementation conversely consists of six MOSFET transistors, i.e., an PMOS/NMOS for each transmission gate, and one PMOS/NMOS for the inverter stage. In our case, we adopted minimum-sized transistors defined in a commercial technology library from STMicroelectronics at 40nm node, where each PMOS has an equivalent area of $0.205\mu\text{m}^2$, and each NMOS has an equivalent area of $0.147\mu\text{m}^2$. Therefore, the total area of a single TG-BDD node is equal to $1.056\mu\text{m}^2$. It is therefore clear that the first advantage of the PN-BDD structure is in the number of devices needed to implement a BDD node. Infact, a TG-BDD node requires an additional inverter stage that is not needed in a PN-BDD node. Another key-point is signal degradation: a stack of MOSFET transmission gates connected in series shows a higher signal degradation

w.r.t. a series of p-n junctions. For this reason, in our mapping script, we added a signal-restoration stage for each stack of ten MOSFET transmission gates, by employing a minimum-sized buffer from the same library with an equivalent area of $1.038\mu\text{m}^2$. Results reported in Table 3 shows that the proposed PN-BDD mapping strategy requires, on average, almost 70% less devices w.r.t. the TG-BDD implementation, which translates to a roughly 39% area saving.

Concerning the efficiency of the PN-BDD over the TG-BDD strategy, we can refer to Figure 20, where we depict the averaged power consumption of the two techniques in function of the transition time (T_r) of the input signal. Both structures are considered in adiabatic configuration. As can be seen from the plot, in non-adiabatic region, i.e., $T_r = 1\text{ps}$, TG-BDD (circle mark) is slightly more efficient (about 3.5%), whereas, as T_r increases PN-BDD (square mark) results in less power consumption. The best case is recorded at $T_r = 1\text{ns}$ with a difference of 1.8 orders of magnitude, which is mainly due to the lower power requirements of an adiabatic p-n junction w.r.t. a MOSFET transmission gate.

In Figure 21, we also address the performances of the proposed mapping solutions by means of the power-delay product (PDP). As can be seen from the plot, the PN-BDD structure shows lower PDP w.r.t. the TG-BDD counterpart over the entire T_r range. The motivation lies in two fundamental differences of the structures, namely the presence of an inverter stage in each TG-BDD node, which increases the configuration time, and the additional signal-restoring stages that increase the total propagation delay.

Table 3: Benchmark characteristics with area occupation for PN-BDD and TG-BDD mapping techniques.

	PI	PO	Nodes	PN-BDD		TG-BDD	
				P-N junctions	Area [μm^2]	Transistors	Area [μm^2]
9sym	9	1	25	50	19.1	150	26.4
5xp1	7	10	42	84	32.08	252	44.35
risc	8	31	79	158	60.35	474	83.42
max1024	10	6	250	500	191	1508	266.11
ex5	8	63	250	500	191	1500	264
apex2	38	3	329	658	251.35	2182	402.33
alu4	14	8	351	702	268.16	2338	431.90
c432	36	7	1166	2332	890.82	7900	1469.95
apex1	45	45	1246	2492	951.94	7804	1402.36
cps	24	109	1693	3386	1293.45	11310	2091.93
xparc	41	73	2268	4536	1732.75	15544	2906.11
c880	60	26	9010	18020	6883.64	61320	11431.2
seq	41	42	9159	18318	6997.47	58250	10542.04
c1908	33	25	9731	19462	7434.48	68698	12998.30
avg.				5085.57	1942.68	17087.85	3168.60

6. Conclusions

This paper explores the potential of graphene p-n junctions for implementing logic devices that operate according to the principle of adiabatic computation.

Results show that these gates can operate with about 2 orders of magnitude lower power than their CMOS counterpart, overcoming a limitation of these graphene-based elements observed in previous works. Moreover, these power benefits do not require very low frequencies as in typical adiabatic operations,

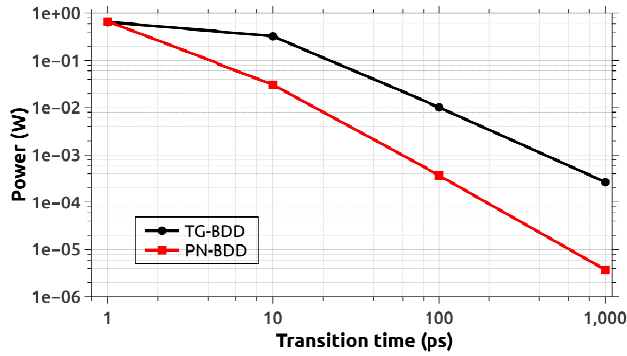


Figure 20: Average power consumption of benchmarks in Table 3.

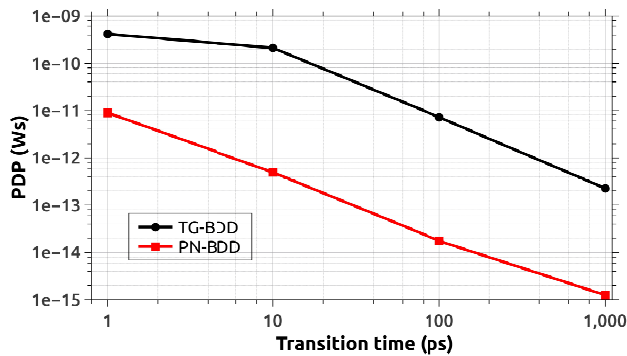


Figure 21: Average Power Delay Product (PDP) of logic functions in Table 3.

and can be achieved with operating frequencies in the order of the hundred of MHz.

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