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The Development of a Knowledge-Based System for Semantic Data Modeling

Ву

William H. Fenton

A Thesis

Presented to the Graduate Committee

of Lehigh University

in Candidacy for the Degree of

Master of Science

in

Industrial Engineering

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This thesis is accepted and approved in partial fulfillment of the requirements for the degree of Master of Science.

date)

Professor in Charge

Chairman of Department

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Abstract

The basis for a sound and logical organization of a corporation's data is the building of logical data structures. Logical data structures support the data definitions and data relationships as related to an organization's business processes and functions.

Organizations implement business processes and functions using database management systems (DBMS). DBMS's are designed and engineered to manipulate data from a physical aspect and do not represent the meaning of the data. The semantics of business processes and data definitions are embedded in program code and replicated across systems. This creates inconsistency and unnecessary expenditures in data processing resources.

Limited research has been devoted to the development of a Semantic Data Model (SDM) to support business processes and data definitions.

Information about the business and data is left to individual interpretation by systems analysts.

This thesis deals with the development of a knowledge-based system to assist in recording the definitions, relationships, and semantics of data in an organization. The benefits of such a system are: (a) user involvement in data definition; (b) the reduction of ambiguous data; and (c) the basis for the implementation of a physical database, regardless of the DBMS.

Chapter I. Semantic Data Structure Background

Today, many organizations and companies develop information systems to establish a position within their market place or to provide a service that is new. Either way, the organization has an opportunity to gain an edge over its competition. In doing so, information systems departments are receiving requests to analyze and develop systems or change the current data processing services their respective departments receive.

The operating department requests focus on the individual needs of each department. In each case, the request satisfies a department requirement, contributing to achieving its business goal. Typically, how a request is satisfied and at what cost are overlooked.

When analysts review all department requests, they discover there is duplication of effort and redundancy of information within the organization's information systems. With duplication of effort, inconsistencies arise between data and the rules to process and analyze that data. These information systems interface and exchange data only when required, and usually only after a political struggle occurs over data ownership.

Following is an actual case of how a company attempted to develop a corporate information system.

Company X developed a new customer database. The system is intended to support and monitor the development of, customer contracts from a customer's initial request through completion. The customer system contains a price module. The price module records the price for computer hardware repair services. The system was developed and implemented under the assumption that prices will not change for each service charged, and that eighty percent of the contracts will use a standard price.

The new customer database system was implemented but never used because of two design oversights. The first was due to an original assumption. The company found that in practice only twenty percent of the contracts were written with standard pricing, while the remaining eighty percent were variations of the standard price. Secondly, the price module was only one of ten modules to be developed. When work began on the other modules, it was discovered that the price data did not meet the requirements to support implementation of the system. A project reassessment estimated it will cost the company \$1.5 million over their budget to redesign the system.

The problem company X faces is that the time and effort needed to develop a corporate data model consistent with the business rules (i.e., the way business is actually conducted), is reflected in the projected expenditure of \$1.5 millon needed to revise the system. What may have been a small cost to change in the early part of the system's life cycle will now cost three to five times as much. Information is

segregated among various files; hence, the high cost in time and dollars to evaluate the impact of the scope change. Note that this is the fifth time in six years the system has been designed and implemented.

Many systems fail because the company organizations do not expend the necessary time and effort to design and develop a corporate data model. Without a data model, departments cannot take advantage of the corporation's data. With continued problems of inconsistent and mistimed information, corporations will continue expending millions of dollars to develop <u>independent</u> information systems, instead of <u>integrated</u> information systems.

Chapter II. Problem Statement

A data model defines the entities, attributes, and data relationships of a business. The data meaning, usage, and relationships are left to individual interpretation. This opening allows multiple meanings of the data to exist. As a result, analysts can design and implement independent applications; but, in an effort to reduce costs, the meaning of the data must be extracted and documented so that it will be given the same interpretation in each application.

For example, consider a car sale. To a salesman, the sale of the car helps to meet a sales quota and to obtain a commission. To the bank, the car is collateral for a loan. To a mechanic, it is future work in the form of repairs. In each case the car is viewed in a slightly different manner.

In recognition of the critical need for a semantic database model, a number of researchers have recently considered the design of such a model [Abrial 1974, Bachman 1977, Biller+Neuhold 1978, Chen 1976, Codd 1979, El-Masri+Wiederhold 1979, Hammer+McLeod...] Although there are substantial and important differences in the detail of these various semantic database models, they have all attempted to increase the degree to which a schema captures the meaning of data in a form intelligible to its designers and users [18, pp. 194].

To assist in the derivation and to capture the meaning of business data, Hammer & McLeod have developed a formalism called 'Semantic Data Model' (SDM) [13]. They declare that the semantics of an organization's business data can be represented by a data model's

entities, attributes, and data relationships; and SDM will provide more of the meaning of business data than other data models of the current day.

Hammer and McLeod declare relationships between data extend farther than data model entities. In contrast, Appleton defines a relationship between entities as Business Rules [2]. The Business Rules explain the constraints which exist among zero, one, or more entities. The Business Rule controls the information contained within the business model, and what data will be created from the business model. Similarities exist between SDM and Business Rule methods, as shown in Figure 1.

Hammer/McLeod	<u>Appleton</u>
Class Attribute(s) Interclass Connections	Entity Attribute(s) Business Rules

Figure 1

Each defines an entity or class (i.e., Purchase Order, Customer, Part Number). Attributes describe the characteristics of the entity or class (i.e., Name, Address, City, State). In addition, SDM extends attributes to the entire entity called Class Attributes and occurrences of each entity called Member Attributes. There is a difference between Interclass Connections and Business Rules.

Hammer and McLeod extend SDM relationships and constraints to the entity and attribute level. They emphasize that attributes belong to a class which is bound by constraints. These constraints exist not just between entities but, as stated by Appleton attributes as well. This SDM recursive format documents the data interrelations throughout the data model, thereby extending the business rules to a class and an attribute.

Chapter III. SDM Rule Implementation

The state of the s

Hammer and McLeod developed a detailed methodology. To understand and document the methodology requires reviewing the entire publication and the SDM syntax. The SDM syntax alone will not provide the entire set of constraints for a class or attribute to exist. The major task in the development of the SDM/Prolog system is to translate the SDM syntax (Appendix A) and validation criteria into a set of Programming Logic (Prolog) rules.

Prolog is an interpretive programming language developed at the University of Marseille by Alan Colmerauer in the early 1970s [8,9]. Since 1970 a group at the University of Edinburgh has made several improvements to Prolog, which is considered to be the defacto standard [16].

Prolog is a declarative language, where the programmer specifies what the program is supposed to achieve by using a number of facts and rules to find a solution to the problem. Compared to procedural languages — such as Pascal, Basic, or Cobol — Prolog instructs the computer what to do, not how to do it.

The intent of Prolog, used as a knowledge-based system for business systems, is to capture as much knowledge of business rules as possible from an expert, and translate that knowledge into rules and facts the

computer understands in order to assist or simulate the decision-making process.

The particular version of Prolog used for this system is Turbo Prolog [5,6] for personal computers. Turbo Prolog departs from the Edinburgh standard, particularly because it requires strong data typing. This feature is required in order to compile code.

To manage SDM/Prolog development, each SDM rule is divided into small manageable sections and placed into a hierarchical structure according to function (see Figures 2 and 3 at end of the chapter). To translate the rule hierarchy into Prolog syntax requires that an analyst first translate the rule into an English sentence. Prolog lends itself more readily to implementation of a rule if that rule is first translated into an English-like structure.

Rule translation uses a numbering system similar to the standard paragraph numbering system. When a new subsection is created, a letter or number is appended to the rule. For example, Rule 1A is the second rule within Rule 1. This technique groups rules within the same category and provides a cross reference between the rule's English translation and the Prolog implementation of the rule. When a rule is used more than once, it is generalized and documented by a comment. Each rule implemented in the SDM/Prolog system is translated into in English-like sentence structure (see Appendix C). For example, the

English translation of a valid SDM class name whether it is a Base Class or Nonbase Class is:

interplated between the property of the contract of the property of the contract of the contra

Each CLASS has a CLASS NAME. A CLASS NAME is composed of uppercase letters and underscores. A CLASS NAME must be unique with respect to all CLASS NAMEs used in the business model. We will call this Rule 1.

The implementation of Rule_1 in Prolog will be a series of predicates.

For the rule to be true in Prolog, <u>all</u> of the predicates which comprise it must be true. One possible translation is:

```
rule_1:-
    lineinput(20,20,45,7,7,"Class name: ","",INPUT),
    upper_lower(INPUT,CIASS),
    isname(CIASS),
    class_list(LIST),
    not(member(CIASS,LIST)),
    retract(class_list(LIST)),
    assertz(class_list([CIASS|LIST]),
    write("Class added"),
    !.

rule_1:-
    write("Class name is invalid or already exists").
```

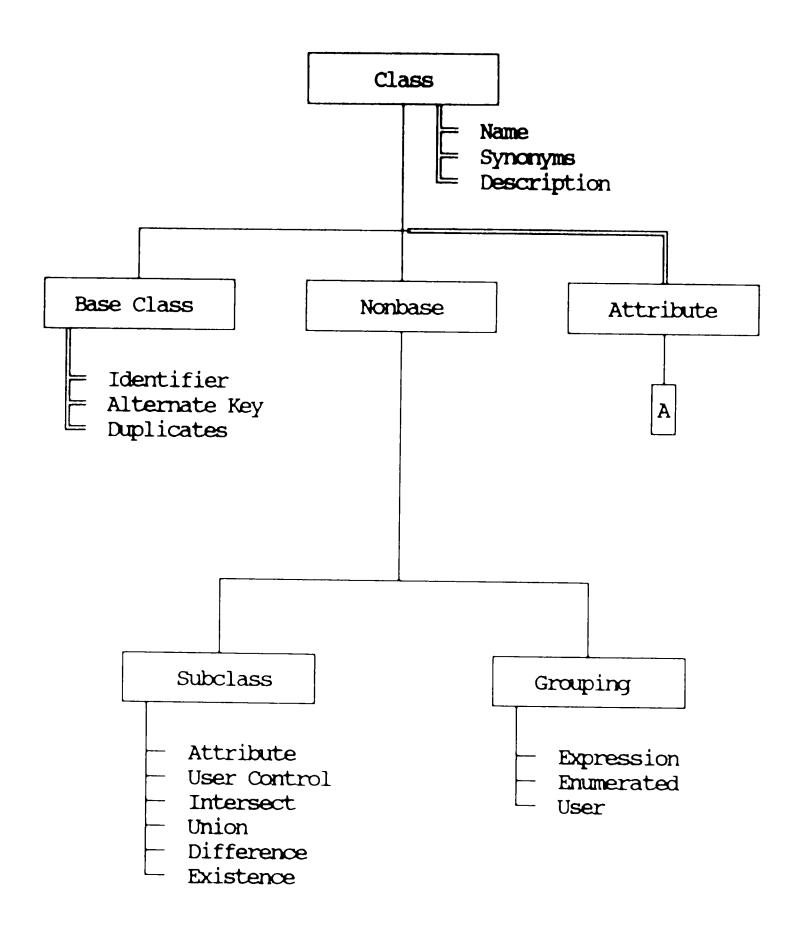
Lineinput is a screen input predicate documented in Turbo Prolog Toolbox [6, pp. 256], Class_list is a Prolog database predicate, and Member is predicate to check whether X is a member of a Prolog list, discussed in Clocksin and Mellish [10, pp. 49]. All other predicates are documented in Turbo Prolog [5].

The Prolog rule can be read as follows:

"Prompt the user for the class name, convert the class name to lower case for validating and processing, check that the class name is composed only of letters and underscores, retrieve the existing list of classes from the database, insure the class name is unique, delete the old class list, create the new class list by concatenating the new class and old class list, write a message, and force Prolog to stop backtracking.

If the first occurrence of the rule_1 predicate fails, write an error message."

SDM Features Hierarchy Diagram



Legend: Functions listed by a double line (=) are common SDM features. A single line (-) indicates one of the following must be selected, e.g. Class has the common features Name, Synonyms, and Description and is either a Base Class or Nonbase Class).

Figure 2

SDM Features Hierarchy Diagram (cont.)

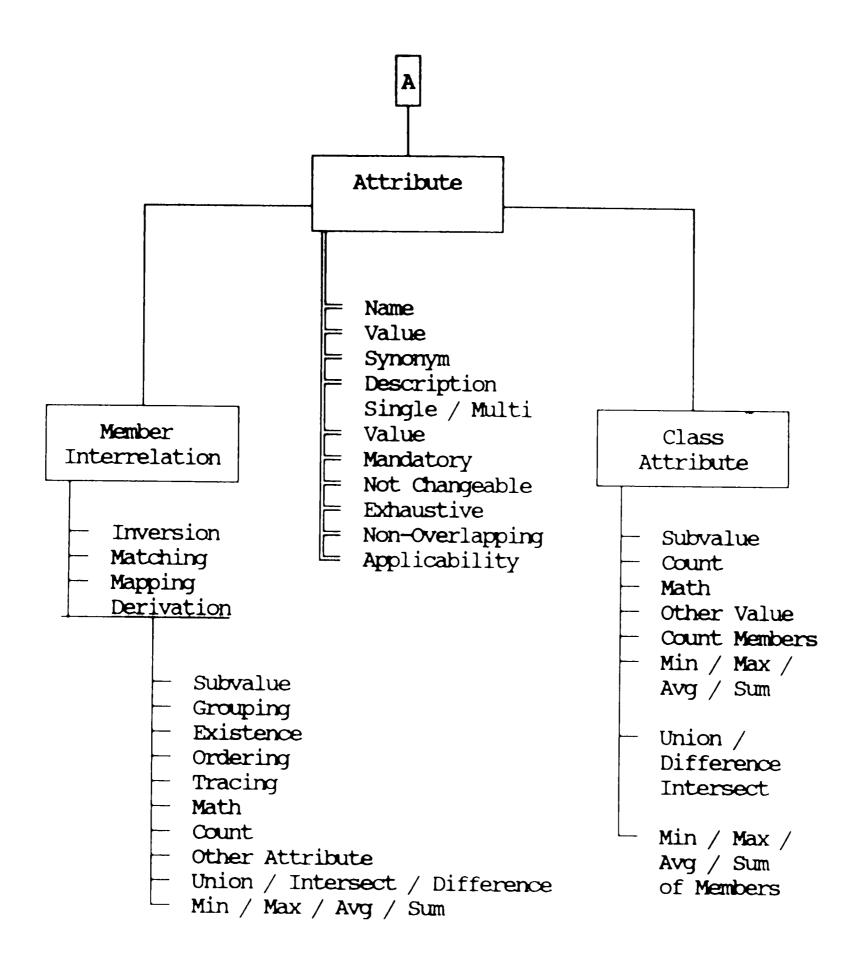


Figure 3

Chapter IV. System Overview

To test the SDM/Prolog system, this thesis will use the Tanker Application example of Hammer and McLeod [13]. This application is designed to monitor and control ships that enter or leave the United States coastal waters. A copy of the Tanker System example can be found in Appendix B.

To validate the Tanker Application, SDM/Prolog is designed to be an interactive system using question and answer menus to define an SDM class or attribute. A selection is processed by placing the cursor at the function or data value. This menu selection method provides a simple way to reduce rule coding, validate data, and control flow. If a selection is aborted, the system ignores any changes and returns to the main system menu. SDM/Prolog supports on-line reporting. It allows a user to list data from the database as the system is built. The limited reporting capabilities route the information to the screen.

Chapter V. Observations

An alternate validation technique would be to design and implement the system by translating the SDM Tanker Application into Prolog clauses and have the SDM/Prolog rules validate the clauses. This technique was tested, but proved to be cumbersome and produced complicated rule definition.

SDM rules are flexible, with complexity based on limited data requirements. A simple form of a Base Class requires a name, member attribute, and logical key. Defining a Nonbase Class requires a name and interclass connection. A more complicated format is a member attribute, with an attribute predicate defined with a mapping clause. The logic rules coded into the SDM/Prolog system were input and/or output driven, rather than deriving knowledge from the knowledge expert to build a knowledge base similar in functionality to the Berkshire MECHANIC automobile knowledge-base system [3].

The degree of interpretation of a class, attribute, or data relation still remains with an individual's business knowledge, experience, and expertise. In the SDM Tanker Application, the Nonbase Class HULL_NUMBERS is defined as the:

interclass connection: subclass of STRINGS where format is number where integer.

An alternative would be:

interclass connection: subclass of INTEGER.

SDM/Prolog's goal is to assist in managing these various interpretations and to document the meaning of data so that it will be given the same interpretation in every application.

Chapter VI. Future Enhancements

This thesis lays the foundation work for future enhancements and provides further functionality not addressed here. Due to time constraints, the scope of the system was defined that not all functions of SDM are included. Areas of SDM enhancements and directions for further research include:

o Help System

Without a detailed background of SDM and knowledge of Prolog, the user is limited to understanding the system functions. A Help system that encompasses the overall system, each functional area, and field level help would enhance the system. This function could be developed using an in-house help system, or the help system documented in the Turbo Prolog Toolbox [6].

o Screen Design

Design the system to display a full screen of data instead of responding to one question at a time, as shown in Figure 4. This will reduce the number of user responses and speed data entry. The suggested method is to use the screen definition and manipulation facilities documented in Turbo Prolog Toolbox [6].

Class Name: Base Class (Y/N):	Nonbase	(Y/N): _
Member Attributes:		_
Class Attributes:		_
Logical Key(s): (1) (2)		
Create Class defin	ition: (Y/N)	

Figure 4

o Reports

Reports are generated on-line only. Additional reports may be on-line and printed:

Cross Reference
Value class to member attribute
Member attributes to Class
Base Class
Nonbase Class

o System

Isolate the input/output activity from the knowledge base for queries.

One Prolog system may be written to manage the creation of the knowledge base and another to manage the queries.

Chapter VII. Conclusion

The goal of this thesis was to demonstrate that Hammer and McLeod's Semantic Data Modeling methodology, based on a structured syntax and set of constraints, can be implemented using Prolog, in a system called SDM/Prolog.

The results of documenting the SDM Tanker Monitoring Application as a skeletal semantic database model, using the SDM/Prolog system, supports business data and rules can be extracted from a business knowledge expert and recorded in a knowledge based system such that they will be given the same interpretation in each business application.

Chapter VIII. Location of SDM/Prolog System

In compliance with thesis preparation guidelines, the following materials are on file in the office of thesis advisor Professor John C. Wiginton.

- a. two SDM/PROLOG system object code diskettes,
- b. two SDM/PROLOG source code diskettes,
- c. SDM/PROLOG program source listings,
- d. Thesis

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APPENDIX A SDM Data Definition Language Syntax

(Source: Hammer and McLeod, [12, pp. 381-384])

Legend

- (1) The left side of a production is separated from the right by a "<---".
- (2) The first level of indentation in the syntax description is used to help separate the left and right sides of a production; all other indentation is in the SDM data definition language.
- (3) Syntactic categories are capitalized while all literals are in lowercase.
- (4) {} means optional.
- (5) [] means one of the enclosed choices must appear; choices are separated by a ';' (when used with '{}' one of the choices may optionally appear).
- (6) <> means one or more of the enclosed can appear, separated by spaces with optional comma and an optional 'and' at the end.
- (7) <<>> means one or more of the enclosed can appear, vertically appended.
- (8) * * encloses a 'meta' description of a syntactic category (to informally explain it).

```
SCHEMA <=
  << CLASS >>
CIASS <=
  < CLASS NAME >
    { description: CLASS DESCRIPTION }
    { [ BASE CLASS FEATURES; INTERCLASS CONNECTION ] }
    { MEMBER ATTRIBUTES }
    { CLASS ATTRIBUTES }
CLASS NAME <=
  * string of capitals possibly including special characters *
CLASS DESCRIPTION <=
  * string *
BASE CLASS FEATURES <=
  {[ duplicates allowed; duplicates not allowed ]}
  {<< IDENTIFIERS >>}
IDENTIFIERS <=
  [ ATTRIBUTE NAME; ATTRIBUTE NAME + IDENTIFIERS ]
```

```
MEMBER ATTRIBUTES <-
  member attributes:
    << MEMBER ATTRIBUTE >>
CLASS ATTRIBUTES <=
  class attributes:
    << CLASS ATTRIBUTE >>
 INTERCLASS CONNECTION <=
     [ SUBCLASS; GROUPING CLASS ]
SUBCLASS <=
  subclass of CLASS NAME where SUBCLASS PREDICATE
GROUPING <=
  [ grouping of CLASS NAME on common value of < ATTRIBUTE NAME >
    { groups defined as classes are < CLASS NAME > };
    grouping of CLASS NAME consisting of classes < CLASS NAME >;
    grouping of CLASS NAME as specified ]
SUBCLASS PREDICATE <=
  [ ATTRIBUTE PREDICATE;
    specified;
    is in CLASS NAME and is in CLASS NAME;
    is not in CLASS NAME;
    is in CLASS NAME or is in CLASS NAME;
    is a value of ATTRIBUTE NAME of CLASS NAME;
    format is FORMAT ]
ATTRIBUTE PREDICATE <=
  [ SIMPLE PREDICATE; ( ATTRIBUTE PREDICATE );
    not ATTRIBUTE PREDICATE;
    ATTRIBUTE PREDICATE and ATTRIBUTE PREDICATE;
    ATTRIBUTE PREDICATE or ATTRIBUTE PREDICATE ]
SIMPLE PREDICATE <=
  [ MAPPING SCALER COMPARATOR [ CONSTANT; MAPPING ];
    MAPPING SET COMPARATOR [ CONSTANT; CLASS NAME; MAPPING ]]
MAPPING <=
  [ ATTRIBUTE NAME; MAPPING.ATTRIBUTE NAME ]
SCALER COMPARATOR <
  [ EQUAL COMPARATOR; >; >=; <; =<; ]
EQUAL COMPARATOR <=
  [ =; >< ]
SET COMPARATOR <=
  [is { properly } contained in; { properly } contains ]
```

```
CONSTANT <-
  * a string or number constant *
FORMAT <==
  * a name class definition pattern *
   (see [16])
MEMBER ATTRIBUTE <==
  < ATTRIBUTE NAME >
    { ATTRIBUTE DESCRIPTION }
    value class: CLASS NAME
    { inverse: ATTRIBUTE NAME }
    {[ match: ATTRIBUTE NAME of CLASS NAME on ATTRIBUTE NAME;
       derivation: MEMBER ATTRIBUTE DERIVATION ]}
    { single valued; multivalued { with size between CONSTANT and
      CONSTANT }}
    { may not be null }
    { not changeable }
    { exhausts value class }
    { no overlap in values }
CLASS ATTRIBUTE <=
  < ATTRIBUTE NAME >
    { ATTRIBUTE DESCRIPTION }
    value class: CLASS NAME
    { derivation: CLASS ATTRIBUTE DERIVATION }
    { single valued; multivalued { with size between CONSTANT and
      CONSTANT }}
    { may not be null }
    { not changeable }
ATTRIBUTE NAME <=
  * string of lowercase letters beginning with a capital and possibly
    including special characters *
ATTRIBUTE DESCRIPTION <=
  * string *
MEMBER ATTRIBUTE DERIVATION <=
  [ INTER ATTRIBUTE DERIVATION; MEMBER SPECIFIC DERIVATION ]
CLASS ATTRIBUTE DERIVATION <=
  [ INTER_ATTRIBUTE_DERIVATION; CLASS_SPECIFIC_DERIVATION ]
```

```
INTER ATTRIBUTE DERIVATION <=
  same as MAPPING;
    subvalue of MAPPING where [ is in CLASS_NAME;
      ATTRIBUTE PREDICATE 1;
    where [ is In MAPPING and is in MAPPING; is in MAPPING or is in
      MAPPING; is in MAPPING and is not in MAPPING ];
    = MAPPING EXPRESSION;
    [ maximum; minimum; average; sum ] of MAPPING;
      number of { unique } members in MAPPING ]
MEMBER SPECIFIC DERIVATION <=
  [ ordering by [ increasing; decreasing ] < MAPPING >
    { within < MAPPING >};
    if in CLASS NAME;
    [ up to CONSTANT; all ] levels of values of ATTRIBUTE NAME;
        contents ]
CLASS SPECIFIC DERIVATION <=
  [ number of { unique } members in this class;
    [ maximum; minimum; average; sum ] of ATTRIBUTE NAME over members
        of this class ]
MAPPING EXPRESSION <=
  [ MAPPING; < MAPPING >; MAPPING NUMBER OPERATOR MAPPING ]
NUMBER OPERATOR <=
  [ +; -; *; /; ! ]
```

APPENDIX B SDM Tanker Monitoring Application

(Source: Hammer and McLeod, [12, pp. 376-381])

SHIPS

Name

Hull number

description: all ships with potentially hazardous cargoes that may enter U.S. coastal waters

and the production of the form of the state of the state

```
member attributes:
  Name
   value class: SHIP NAMES
  Hull number
   value class: HULL NUMBERS
   may not be null
   not changeable
  Type
   description: the kind of ship, for example, merchant or fishing
   value class: SHIP TYPE NAMES
  Country of registry
   value class: COUNTRIES
    inverse : Ships registered here
  Name of home port
   value class: PORT NAMES
  Cargo types
   description: the type(s) of cargo the ship can carry
   value class: CARGO TYPE NAMES
   multivalued
 Captain
   description: the current captain of the ship
   value class: OFFICERS
               : Officer of ASSIGNMENTS on Ship
   match
 Engines
   value class: ENGINES
   multivalued with size between 0 and 10
   exhausts value class
   no overlap in values
 Incidents involved in
   value class: INCIDENTS
   inverse
            : Involved ship
   multivalued
identifiers:
```

INSPECTIONS

description: inspections of oil tankers

member attributes:

Tanker

description: the tanker inspected

value class: OIL TANKERS inverse : Inspections

Date

value class: DATES Order for tanker

description: the ordering of the inspections for a tanker with

the most recent inspection having value 1

value class: INTEGERS

derivation: order by decreasing Date within Tanker

class attributes:

Number

description: the number of inspections in the database

value class: INTEGERS

derivation: number of members in this class

identifiers:

Tanker + Date

COUNTRIES

description: countries of registry for ships

member attributes:

Name

value class: COUNTRY NAMES

Ships_registered_here value class: SHIPS

inverse : Country of registry

multivalued

identifiers

Name

OFFICERS

description: all certified officers of ships

member attributes:

Name

value class: PERSON NAMES

Country of license

value class: COUNTRIES

Date_commissioned value class: DATES

Seniority

value class: INTEGERS

derivation: order by Date commissioned

Commander

description: the officer in direct command of this officer

value class: OFFICERS

Superiors

value class: OFFICERS

derivation: all level of values of Commander

inverse : Subordinates

multivalued Subordinates

value class: OFFICERS inverse : Superiors

multivalued

Contacts

value class: OFFICERS

derivation: where is in Superiors or is in Subordinates

identifiers:

Name

ENGINES

description: ship engines

member attributes: Serial number

value class: ENGINE SERIAL NUMBERS

Kind of engine

value class: ENGINE TYPE NAMES

identifiers:

Serial number

INCIDENTS.

description: accidents involving ships

member attributes:

Involved_ship

value class: SHIPS

inverse : Incidents involved in

Date

value class: DATES

Description

description: textual explanation of the accident

value class: INCIDENT DESCRIPTIONS

Involved captain

value class: OFFICERS

identifiers:

Involved ship + Date + Description

ASSIGNMENTS

description: assignments of captains to ships

member attributes:

Officer

value class: OFFICERS

Ship

value class: SHIPS

identifiers:

Officer + Ship

OIL TANKERS

description: oil-carrying ships

interclass connection: subclass of SHIPS where Cargo types

contains 'oil'

member attributes:

Hull_type

description: specification of single or double hull

value class: HULL TYPE NAMES

Is tanker banned

value class: YES/NO

derivation: if in BANNED SHIPS

Inspections

value class: INSPECTIONS

inverse : Tanker

multivalued

Number_of_times_inspected value class: INTEGERS

derivation: number of unique members in Inspections

Last inspection

value class: MOST RECENT INSPECTIONS

inverse : Tanker
Last two inspections

value class: INSPECTIONS

derivation: subvalue of Inspections where Order for tanker =<2

multivalued
Date_last_examined
value class: DATES

derivation: same as Last inspection. Date

Oil_spills_involved_in value class: INCIDENTS

derivation: subvalue of Incidents involved in where is in

OIL SPILLS

multivalued

class attributes:

Absolute_top_legal_speed value class: KNOTS

Top_legal_speed_in_miles_per_hour value class: MILES PER HOUR

derivation : = Absolute_top_legal_speed/1.1

RURITANIAN SHIPS

interclass connection: subclass of SHIPS where Country.Name = 'Ruritania'

RURITANIAN OIL TANKERS

interclass connection: subclass of OIL TANKERS where

Country.Name = 'Ruritania'

MERCHANT SHIPS

interclass connection: subclass of SHIPS where Type = 'merchant'

employed a light for the contract of the contr

member attributes:

Cargo types

value class: MERCHANT CARGO TYPE NAMES

OIL SPILLS

interclass connection: subclass of INCIDENTS where Description = 'oil spill'

member attributes:

Amount spilled

value class: GALIONS

Severity

derivation : = Amount_spilled/100,000

class attributes:

Total spilled

value class: GALLONS

derivation: sum of Amount spilled over members of this class

MOST RECENT INSPECTIONS

interclass connection: subclass of INSPECTIONS where

Order for tanker = 1

DANGEROUS CAPTAINS

description: captains who have been involved in an accident interclass connection: subclass of OFFICERS where IS A VALUE OF

Involved captain of INCIDENTS

BANNED SHIPS

description: ships banned from U.S. coastal waters

interclass connection: subclass of SHIPS where specified

member attributes:

Date banned

value class: DATES

OIL TANKERS REQUIRING INSPECTION

interclass connection: subclass of OIL TANKERS where specified

BANNED OIL TANKERS

interclass connection: subclass of SHIPS where is in BANNED SHIPS

and is in OIL TANKERS

SAFE SHIPS

description: ships that are considered good risks

interclass connection: subclass of SHIPS where is not in

BANNED SHIPS

SHIPS TO BE MONITORED

description: ships that are considered bad risks

interclass connection: subclass of SHIPS where is in BANNED SHIPS or

is in OIL TANKERS REQUIRING INSPECTION

SHIP TYPES

description: types of ships

interclass connection: grouping of SHIPS on common value of Type

groups defined as classes are MERCHANT SHIPS

member attributes:

Instances

description: the instances of the type of ship

value class: SHIPS

derivation : same as Contents

multivalued

Number_of_ships_of_this type

value class: INTEGERS

derivation: number of members in Contents

CARGO TYPE GROUPS

interclass connection: grouping of SHIPS on common value of

Cargo types

TYPES OF HAZARDOUS SHIPS

interclass connection: grouping of SHIPS consisting of classes BANNED SHIPS, BANNED OIL TANKERS, SHIPS TO BE MONITORED

CONVOYS

interclass connection: grouping of SHIPS as specified

member attributes:

Oil tanker constituents

description: the oil tankers that are in the convoy (if any)

value class: SHIPS

derivation: subvalue of Contents where is in OIL TANKERS

multivalued

CARGO TYPE NAMES

description: the types of cargo

interclass connection: subclass of STRINGS

MERCHANT CARGO TYPE NAMES

interclass connection: subclass of CARGO TYPE NAMES where specified

COUNTRY NAMES

interclass connection: subclass of STRINGS where specified

ENGINE SERIAL NUMBERS

interclass connection: subclass of STRINGS where format is

'H'

number where integer and >= 1 and =< 999

1-1

number where integer and >= 0 and =< 999999

DATES

description: calendar dates in the range '1/1/85' to '12/31/89' interclass connection: subclass of STRINGS where format is

month:number where >= 1 and =< 12

1/1

day :number where integer and >= 1 and =< 31 year :number where integer and >= 1970 and =< 2000 where if (month = 4 or = 5 or = 9 or = 11) then day

>= 30) and (if month = 2 then day =< 29)

ordering by year, month, day

ENGINE TYPE NAMES

interclass connection: subclass of STRINGS where specified

GALLONS

interclass connection: subclass of STRINGS where format is number

where integer

HULL NUMBERS

interclass connection: subclass of STRINGS where format is number

where integer

HULL TYPE NAMES

description: single or double

interclass connection: subclass of STRINGS where specified

INCIDENT DESCRIPTIONS

description: textual description of an accident

interclass connection: subclass of STRINGS

ROIS

interclass connection: subclass of STRINGS where format is number where integer

MILES PER HOUR

interclass connection: subclass of STRINGS where format is number where integer

PORT NAMES

interclass connection: subclass of STRINGS

PERSON NAMES

interclass connection: subclass of STRINGS

SHIP NAMES

interclass connection: subclass of STRINGS

SHIP TYPE NAMES

description: the names of the ship types, for example, merchant interclass connection: subclass of STRINGS where specified

APPENDIX C English Rule Interpretation

Class Generic Rules

- Rule_1 The class name must be unique and composed of uppercase letters and underscores. The class name will be stored as lower case and converted to uppercase for display.
- A class name can have synonyms (optional). Each synonym must be unique, composed of uppercase letters and underscores. The synonym will be stored as lower case and converted to uppercase for display.
- Rule_3 A class name can have descriptive text (optional) describing the meaning and contents of the class. The text is a string of characters.

Base Class Rules

- Rule_4 A class name must be a base class or nonbase class.
- A base class must have at least one (1) logical key. A logical key is composed of member attributes, and each member attribute must exist within the base class. A logical key can have duplicate or unique values. Its default is duplicate values. A member attribute cannot be repeated within a logical key.

Nonbase Class Rules

Rule_6 A nonbase class has one interclass connection, either a sub-class or grouping class connection. A subclass name must exist and cannot be equal to the nonbase class name.

> attribute-defined difference existence format intersection none union where-specified

Rule_6Al An attribute-defined predicate defines a subclass S,

composed of members of C, based on a member attribute equation:

equality contains

Rule_6A2 A difference predicate defines a subclass S composed of members from class C that do not belong to class D. Class D must exist and must be a subclass of C. Class

name S cannot be equal to class name C, D, or E.

Rule_6A3 An existence predicate defines a subclass S composed of members from class C, based on a member attribute of C. The member attribute must exist in class C. Class name

S cannot be equal to class name C.

Rule 6A4 A format predicate is free form textual description.

An intersection predicate defines a subclass S composed of members from class C that belong to class D <u>and</u> E. Class D and E must exist and both must be a subclass of C. Class name S cannot be equal to class name C, D, or

E.

Rule_6A6 No predicate will be defined.

A union predicate defines a subclass S composed of members from class C that belong to class D or E. Class D and E must exist and both must be a subclass of C. Class name S cannot be equal to class name C, D, or E.

Rule 6A8

A where-specified predicate defines a subclass S, composed of members from class C, manually controlled.

Equality defines an equation based on equal (=), not equal (>), greater than or equal (>), greater than (>), less than or equal (=<). The member attribute must exist, the comparator must be one of the above, and either a constant or another member attribute. The member attributes cannot be equal, and the second member attribute must exist. For example cargo_type = 'oil'. (Member attribute was not implemented).

Contains defines an equation based on properly contained, contained, properly contains, contains. The member attribute must exist, the set comparator must be one of the above, and either a constant, member attribute, or class name. The member attributes cannot be equal, and the second member attribute must exist. For example cargo type = 'oil'. The class name must exist, and cannot be equal to the class being created. (Member attribute and class name were not implemented).

A member attribute name must be unique within the class and all subclasses, composed of letters and underscores. The member attribute name will be stored as lower case and the first letter of the name will be converted to uppercase for display.

Rule_7A A member attribute can exhaust its value class and valid values are 'Yes' or 'No'. The default is 'No'.

A member attribute cannot be changed once a value is assigned. The value can be changed only to correct an error. Valid values are 'Yes' or 'No'. The default is 'No'.

Rule_7C A member attribute can be mandatory. Valid values are 'Yes' or 'No'. The default is 'No'. If a member attribute has no value class (Rule_8), the member attribute value cannot be mandatory.

Rule_7D A member attribute can be single valued or multivalued.

Valid values are 'S', single valued; or 'M',
multivalued. The default is single valued.

Rule_7D1 A multivalued attribute can have one of the following (optional):

nonoverlapping value restriction

Nonoverlapping means each member of the value class of the attribute is used only once. Valid values are 'Yes' or 'No'. Default is 'No'.

Rule_7D1B Value restriction is a range of values a member attribute can contain. The lower range is an integer. The upper range is an integer. The upper range integer must be greater than the lower range integer.

A member attribute can have descriptive text (optional) describing the meaning and contents of the member attribute. The text is a string of characters.

Rule_8 A member attribute can have a value class (optional). The value class name must exist. If there is no value class, the member attribute cannot be mandatory (Rule_7C). (The no value class option was not implemented).

APPENDIX D Prolog Implementation of SDM

code=4000

include "tdoms.pro"

```
database - sdmdb
  base class
                       (string, stringlist)
  class def
                       (string, string)
  class list
                       (stringlist)
  class members
                       (string, stringlist)
  class syn
                       (string, stringlist)
  member attr
                      (string, string, string, stringlist)
  non base
                       (string, string, string, string)
  stnd class
                      (stringlist)
include "tpreds.pro"
include "lineinp.pro"
include "menu2.pro"
include "status.pro"
predicates
 asc order
                     (string, string)
 append
                     (stringlist, stringlist, stringlist)
 base list
                     (string, string)
 create class
                     (string, string)
 create def
                     (string, string)
 delete
                     (string, stringlist, stringlist)
 get class
                     (symbol, string)
 get class members (string, stringlist)
 ф
 init
 init stnd class
                     (stringlist)
 insert
                     (string, stringlist, stringlist)
 insert sort
                     (stringlist, stringlist)
 makenblist
                     (string, stringlist)
 member
                     (string, stringlist)
 menuchoice
                    (integer)
 pause
 print class members
                           (stringlist)
 print list
                           (stringlist)
 print member
                           (string)
 print member attr list
                           (stringlist)
 print member attr
                           (string)
 print member def
                          (string)
print non base
                          (string, string, string)
rule1
                    (string)
rule2
                    (string, string)
rule3
                    (string)
rule6
                    (string, string, string, string)
rule6a
                    (integer, string, string, string, string)
rule6a1
                    (string, string)
rule7
                    (string)
rule7a
                    (string, string)
```

```
rule7b
                      (stringlist)
  rule7bl
                      (integer, stringlist)
  rule7b2
                      (integer, stringlist)
                      (integer, stringlist, string)
  selection
  update class members (string, string)
  valid name
                     (string)
qoal
  retractall(, sdmdb),
  makewindow(\overline{1},23,14," Semantic Data Model ",0,0,24,80),
  go.
clauses
go:-
  init,
  repeat,
  menu (8,27,31,11,
       [" Class Add", " Synonym Add", " Class Definition",
        " Member Attributes"," Non-base Add"," List Classes",
        " List Synonyms"," List Class Definition",
        " List Member Attributes "," List Nonbase Classes",
        " Exit"],
        " Main Menu ",1,Choice),
  menuchoice (Choice),
  clearwindow,
  Choice=11,
  exit.
init:-
  assertz(class list([])),
  assertz(stnd class(["integers",
    "numbers", "reals", "strings", "yes/no"])),
  stnd class(SLIST),
  init stnd class(SLIST),
 makestatus (112,"").
init stnd class([]).
init_stnd class([X|TAIL]):-
  init stnd class(TAIL),
  assertz(base class(X,[])).
```

```
menuchoice(1):- /* create class*/
  changestatus (""),
  lineinput(10,10,45,7,7,"Class name: ","", INPUT),
  upper lower (INPUT, CLASS),
  valid name (CLASS),
  rule1 (CLASS),
  meru(15,10,31,11,["No","Yes"]," Create definition? ",1,S),
  selection(S,["n","y"],DEFOPT),
  create def(DEFOPT, DEF),
  menu(15,30,31,11,["Base Class","Non Base Class"],
               " Select Class Type ",1,S2),
  selection(S2,["b","n"],CLASSOPT),
  create class(CLASSOPT,CLASS),
  class list(LIST),
  retract(class list(LIST)),
  insert sort([CLASS|LIST], SORTED LIST),
  assertz(class list(SORTED LIST)),
  assertz(class def(CLASS, DEF)),
  changestatus ("Class added").
menuchoice(2):- /* create synonyms*/
  get class(all,CIASS),
  upper lower (UPPER, CLASS),
  concat (UPPER, " synonym: ", MENUTEXT),
  lineinput(10,10,45,7,7,MENUTEXT,"",INPUT),
  upper lower(INPUT,SYN),
  valid name (SYN),
  rule2(CIASS, SYN),
  class syn(CLASS, SYNLIST),
  retract(class syn(CLASS,SYNLIST)),
  insert sort([SYN|SYNLIST],SORTLIST),
  assertz(class syn(CLASS, SORTLIST)),
  changestatus ("Synonym added"),
  !.
menuchoice(3):- /* class definition*/
  get class(all,CLASS),
  rule3(CIASS).
menuchoice(4):- /* member attributes*/
  get class(all,CLASS),
  rule7(CLASS).
/*
menuchoice(5):- /* non-base classes*/
  get class(all,NBCLASS),
  rule6.
*/
menuchoice(6):- /* no classes exist*/
  class list([]),
  beep,
  changestatus ("No Classes have been defined"),!.
```

```
menuchoice(6):- /* list all classes*/
  class list(LIST),
  print list(LIST),
  pause,
  changestatus(""),!.
menuchoice(7):- /* list synonyms*/
  class_syn(CLASS,LIST),
  upper lower (UPPER, CLASS),
  write("Class ",UPPER," synonym(s): "),
  print list(LIST),nl,
  pause,
  fail.
menuchoice(7):- /* no synonyms exist*/
  not(class_syn(_,_)),
  beep,
  changestatus ("No Synonyms have been defined").
menuchoice(8):- /* list class definition */
  class def(Class, Text),
  Text >< "",
  upper lower(Upper, Class),
  write("Class ", Upper, " definition: "), nl,
  write (Text), nl,
  pause,
  fail.
menuchoice(8):- /* no definitions exist*/
  not(class def(_,_)),
  beep,
  changestatus ("No Class definitions have been defined").
menuchoice(9):- /* list member attributes*/
  class list(LIST),
  print class members (LIST).
menuchoice(10):-/* list non-base class
                                                 */
  non_base(CLASS, SUBCLASS, TYPE, FORM, FORM2),
  upper lower (UPPER, CLASS),
  write(UPPER), nl,
  upper lower(UPPER2, SUBCLASS),
  write(" interclass connection: subclass of ", UPPER2),
  print non base(TYPE, FORM, FORM2), nl,
  pause,
  fail.
menuchoice(). /* menuchoice will always succeed */
```

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```
valid_name(X):- /* valid name*/
  not(isname(X)),
  changestatus ("Invalid name"),
  !,fail.
valid_name(_).
rule1(CLASS):- /* class cannot exist*/
  class list(CLIST),
  stnd class (STCLASS),
  append(CLIST,STCLASS,LIST),
  member (CLASS, LIST),
  changestatus ("Class name already exists"),
  !,fail.
rule1(_).
rule2( ,SYN):- /* syn cannot exist */
  class syn(CLASS,CLIST),
  stnd class(STCLASS),
  append(CLIST, STCLASS, LIST),
  member (SYN, LIST),
  upper lower (UPPER, CLASS),
  concat ("Synonym exists for class", UPPER, TEXT2),
  changestatus (TEXT2),
  !,fail.
rule2( ,SYN):- /* syn cannot exist as a class*/
  class list(LIST),
  member (SYN, LIST),
  changestatus ("Synonym exists as a class"),
  !,fail.
rule2(CIASS, ):-/* create dummy synonym*/
  not(class syn(CLASS, )),
  assertz(class_syn(CLASS,[])),
  ! .
rule2(_,_).
rule3(Class):- /* create class definition*/
  class_def(Class, Textin),
  edit(Textin, Textout) ,
  Textin > Textout,
 retract(class_def(Class, )),
  !,
 Textout > "",
 assertz(class def(Class, Textout)).
```

```
rule3 (Class):-
  not(class_def(Class, )),
  edit("", Textout),
  Textout × "",
  assertz(class def(Class, Textout)).
rule6 (NBCLASS, SUBCLASS, TYPE, FORM, FORM2):-/* non-base class*/
  get class(sub, SUBCLASS),
/* delete(NBCLASS, SUBCLASSL, SUBCLASSL2), */
  menu(13,10,31,11,["Attribute defined ","Difference",
                   "Existence", "Format", "Intersection",
                   "None", "Union", "Where Specified"],
                   " Select Subclass Type ",1,S),
  rule6a(S,NBCLASS,SUBCLASS,TYPE,FORM,FORM2).
rule6a(1, ,SUBCL,NBTYPE,FORM,FORM2):- /* non-base type = attribute */
  get_class_members(SUBCL,MLIST),
  menu(13,10,31,11,MLIST," Select Attribute ",1,S),
  selection(S,MLIST,FORM),
  menu(13,10,31,11,["equal","not equal","less than",
                   "less than or equal", "greater than",
                   "greate than or equal", "contains",
                  "properly contains", "contained in",
                   "is properly contained in "],
                   " Select Comparator ",1,S2),
   selection(S2, ["aeq", "ane", "alt", "ale", "agt", "age", "aco", "apc",
          "acd", "apd"], NBTYPE),
  menu(13,10,31,11,["Constant","Attribute"]," Select 2nd type ",1,
  selection(S3,["C","A"],MAPTYPE2),
  rule6a1(MAPTYPE2, FORM2),
  !..
rule6a(2, ,SUBCL,NBTYPE,FORM,FORM2):-/* non-base type = difference*/
  NBTYPE = "d",
  FORM2 = ""
  makenblist(SUBCL, NBLIST),
/* delete(NBCL, NBLIST, NBLIST2),/* change nblist = nblist3*/
  delete(SUBCL, NBLIST2, NBLIST3),
*/ menu(15,20,31,11,NBLIST," Select Difference Sub class ",1,S),
  selection(S,NBLIST,FORM),
rule6a(3,_,_,NBTYPE,FORM,FORM2):-/* non-base type = existence*/
 NBTYPE = "e"
  FORM
       = "existence",
 FORM2 = "",
 !.
```

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```
rule6a(4,_,_,NBTYPE,FORM,FORM2):-/* non-base type = format*/
  NBIYPE = "f",
  FORM2 = min
  edit("", FORM),
   !.
rule6a(5,_,SUBCL,NBTYPE,FORM,FORM2):-/* non-base type = intersection */
  NBIYPE = "i",
  makenblist(SUBCL, NBLIST),
/* delete(NBCL,NBLIST,NBLIST2),*/
  menu (15, 20, 31, 11, NBLIST,
    "Intersection: select 1st Sub class ",1,S1),
  selection(S1, NBLIST, FORM),
  delete(FORM, NBLIST, NBLIST2),
  menu(15,20,31,11,NBLIST2," Select 2nd Sub class ",1,S2),
  selection(S2, NBLIST2, FORM2),
  ! .
rule6a(6,_,_,NBTYPE,FORM,FORM2):-/* non-base type = none*/
  NBTYPE = "n"
  FORM = ""
  FORM2 = ""
  !.
rule6a(7,_,SUBCL,NBTYPE,FORM,FORM2):-/* non-base type = union */
  NBIYPE = "u",
  makenblist(SUBCL, NBLIST),
  menu(15,20,31,11,NBLIST," Union: select 1st Sub class ",1,S1),
  selection(S1, NBLIST, FORM),
  delete(FORM, NBLIST, NBLIST2),
  menu(15,20,31,11,NBLIST2," Select 2nd Sub class ",1,S2),
  selection(S2, NBLIST2, FORM2),
  !.
rule6a(8,_,_,NBTYPE,FORM,FORM2):-/* non-base type = specified*/
  NBTYPE = "s"
  FORM = ""
  FORM2 = ""
  ! .
```

```
rule7(CIAMS):- /* create member attributes*/
  upper lower (UPPER, CLASS),
  correct (CFFER, " member rame: ", MIEXT),
  lineingut(10,10,45,7,7,MTEXT,"",Input),
  upper lower (INFUT, MEMBER),
  valid name (MEMBER),
  rule7a(CLASS,MEMBER),
  get class (value, VCIASS),
  rule7b(MTRLIST),
  menu(8,27,31,11,["No","Yes"]," Create definition? ",1,S),
  selection(S,["n","y"],DEFOPT),
  create def(DEFOPT, DEF),
  asserta (member attr(CLASS, MEMBER, VCLASS, DEF, ATTRLIST)),
  update class members (CLASS, MEMBER),
  changestatus ("Member attribute added"),
  ! .
rule6a1("C", CONSTANT):-
  lineinput(10,10,45,7,7,"Contant is: ","", CONSTANT), !.
rule6al("A", ):-
  changestatus ("Function not available"),
  fail.
rule7a(C,M):-
               /* member cannot exist*/
  member attr(C,M,_,_,),
  changestatus ("Member attribute already exists"),
  !,fail.
rule7a(_,_).
rule7b(ATTR CHARACTERISTICS):-
  menu repeat (10, 10, 31, 11, ["Mandatory", "Optional"],
                  " Attribute is ",2,S1),
  selection(S1,["M","O"],MANDOPT),
  append([],[MANDOPT],OPTLIST1),
  menu_repeat(10,40,31,11,["Changeable","Not Changeable"],
                  " Attribute is ",1,S2),
  selection(S2,["C","NC"],CHGOPT),
  append(OPTLIST1,[CHGOPT],OPTLIST2),
  menu_repeat(15,10,31,11,["No","Yes"],
                 " Exhausts Value Class ",1,S3),
  selection(S3,["N","E"],EXHOPT),
  append(OPTLIST2,[EXHOPT],OPTLIST3),
 menu_repeat(15,40,31,11,["Single Value","Mult-Value"],
                 " Attribute is ",1,S4),
  rule7b1(S4,SMLIST),
  append (OPILIST3, SMLIST, ATTR CHARACTERISTICS),
 removewindow, removewindow, removewindow, removewindow,
  !.
```

```
rule7b1(1,["S"]).
rule7b1(2,MLIST):-
  menu repeat(12,10,31,11,["No","Yes"]," Overlap in Values ",1,S1),
  selection(S1,["N","O"],OVEROPT),
  append([],[OVEROPT],MLIST1),
  menu_repeat(12,40,31,11,["No","Yes"]," Range Limitation ",1,S2),
  rule7b2 (S2,RANGE),
  append (MLIST1, RANGE, MLIST),
  removewindow.
  removewindow.
rule7b2(1,[]).
rule7b2(2,RANGE):-
  lineinput(16,40,15,7,7,"Low Range: ","",NUMSTART),
  lineinput(16,60,15,7,7,"High Range: ","",NUMEND),
  append([NUMSTART, NUMEND],[],RANGE).
create def("y", Textout):-/* invoke Prolog editor */
  edit("", Textout),
  clearwindow,
  !.
create def( ,Text):-/* if editor aborted set null */
  Text="".
create class("b", BCIASS):-/* create base class*/
  assertz(base class(BCLASS,[])), !.
create class("n",NBCLASS):-/* create non base class*/
  rule6 (NBCLASS, SUBCLASS, TYPE, FORM, FORM2),
  assertz(non base(NBCLASS, SUBCLASS, TYPE, FORM, FORM2)),
 base list (NBCLASS, SUBCLASS).
```

```
print non base ("d", TEXT, ):-/* print difference sub class */
   upper lower (UPPER, TEXT),
  write(" where is not in ", UPPER),!.
print_non_base("f",TEXT,_):-/* print format sub class*/
  write(" where format is ", TEXT),!.
print_non_base("n",_,_):-!./* print none sub class*/
print_non_base("s",_,_):-/* print user controlled sub*/
  write(" where specfied"),!.
print non base ("u", TEXT, TEXT2):-/* print intersection sub class */
  upper lower (UPPER, TEXT),
  write(" where is in ",UPPER),nl,
  upper lower(UPPER2, TEXT2),
  write(" and is in ", UPPER2),!.
print non base("aeq", TEXT, TEXT2):-
  frontstr(1,TEXT,FIRST,REST),
  upper lower(FUPPER, FIRST),
  concat (FUPPER, REST, MATTR),
  concat("'", TEXT2, CONSTANT),
  concat(CONSTANT,"", CONSTANT2),
  write(" where ", MATTR, " = ", CONSTANT2),!.
print_class_members([])./* print member attributes*/
print class members([Class|Tail]):-
  print class members (Tail),
  upper lower (Upper, Class),
  write(Upper),
  write("\n member attributes: "),
  print member(Class),
  changestatus ("Press ENTER to return"),
  pause.
print member(C):-/* print member details
  member attr(C,M,V,D,A),
  frontstr(1,M,FIRST,REST),
  upper_lower(FUPPER, FIRST),
  concat (FUPPER, REST, MATTR),
  upper lower(VCLASS,V),
             ",MATTR),
  write("\n
  print member def(D),
  write("\n
                value class: ", VCLASS),
  print member attr list(A),
  pause,
  fail.
/*print member( ).*/
```

```
print number attr list([]).
print_member_attr_list([H|TAIL]):-
  print member attr(H),!,
  print member attr list(TAIL).
print_member_attr("M"):-
  write("\n may not be null").
print_member_attr("NC"):-
  write("\n
               not changeable").
print_member_attr("E"):-
  write("\n exhausts value class").
print_member_attr(_):- !.
print_member def(D):-
  D \times ""
  write("\n
            description: ",D),!.
print_member_def().
```

```
get class(all,C):-/* display memu of classes*/
  changestatus (""),
  class list(CHOICELIST),
  meru(8,27,31,11,CHOICELIST," Select Class from menu ",1,S),
  selection(S,CHOICELIST,C),!.
get_class(all,_):-
  class list([]),
  beep,
  changestatus ("No Classes have been defined"),
  !,fail.
get class(sub,C):-/* display menu of sub-classes*/
  changestatus (""),
  stnd class(STCLASS),
  class list(CLIST),
  append(CLIST, STCLASS, CLASSLIST),
  insert sort(CLASSLIST, SORTLIST),
  menu(8,27,31,11,SORTLIST," Select Subclass ",1,S),
  selection(S,SORTLIST,C),
  ! .
get class(value,C):-/* display memu of value classes*/
  changestatus(""),
  class list(CLIST),
  stnd class(SLIST),
  append(CLIST, SLIST, CHOICELIST),
  insert sort(CHOICELIST, SORTLIST) ,
  menu(8,27,31,11,SORTLIST," Select Value Class ",1,S),
  selection(S,SORTLIST,C),
  ! .
get class members(C,ML):-
  class members(C,ML), !.
get class members( , ):-
  changestatus ("No members exist for non base class"),
  pause,
 fail.
```

```
append([],LIST,LIST)./* append two lists into one*/
append([X|L1],LIST2,[X|L3]):-
  append(L1,LIST2,L3).
base list(NB,SC):-/* add subclass to base class*/
  base class(SC, LIST),
  not(member(NB, LIST)),
  retract(base class(C, LIST)),
  insert sort([NB|LIST],SORTLIST),
  assertz(base_class(C,SORTLIST)), !.
base list(NB,SC):-/* locate base class of sub class */
  base class(C,LIST),
  member (SC, LIST),
  retract(base class(C,LIST)),
  insert_sort([NB|LIST],SORTLIST),
  assertz(base class(C, SORTLIST)), !.
delete(,[],[])./* delete X from list*/
delete(X,[X|L1],L1):-!.
delete(X,[Y|L1],[Y|L2]):-!,
  delete(X, L1, L2).
insert_sort([],[])./* sort routine*/
insert sort([X|TAIL],SORTED LIST):-
  insert sort (TAIL, SORTED TAIL),
  insert(X,SORTED_LIST).
insert(X,[Y|SORTED LIST],[Y|SORTED LIST1]):-
  asc order(X,Y), !,
  insert (X, SORTED LIST, SORTED LIST1).
insert(X, SORTED LIST, [X | SORTED LIST]).
asc order(X,Y):- X>Y.
makenblist(SC, NBLIST):-/* make non-base selection list */
  base class(SC, NBLIST), !.
makenblist(SC,NBLIST):-
  base_class( ,LIST),
  member(SC, LIST),
  delete(SC, LIST, NBLIST), !.
member(X,[X|_])./* is X a found in LIST*/
member(X, [ | \overline{Y}]) :- member(X, Y).
```

```
/* pause*/
  changestatus ("Press ENTER to continue"),
  readin(_).
print_list([]):- !./* print a member list*/
print list([X|Y]):-
  upper lower (UPPER, X),
  write("\n
                 ",UPPER),
  print list(Y).
selection(0,[],_):- fail./* convert menu select to value */
selection(1,[H ],H):-!.
selection(N,[ T],Item):-
  N1=N-1,
  selection(N1,T,Item).
update class members(C,M):-
  class members (C, ML),
  retract(class members(C,ML)),
  insert_sort([M ML],SML),
  assertz(class members(C,SML)),
  !.
update class members(C,M):-
  assertz(class members(C,[M])),
  ! .
```

Vita

The author was born in Abington, Pennsylvania, on February 11, 1957, and is the son of William H. Fenton and E. Charlotte Fenton. Upon graduating from Northeast High School, Philadelphia, Pennsylvania, in June of 1975, he attended Clarion State University (now Clarion University), Clarion, Pennsylvania. He graduated in May of 1979 with a Bachelors of Science in Business Administration. Upon graduation from Clarion he was employed at Air Products and Chemicals Inc., Allentown, Pennsylvania, where he was a commercial business systems programmer. In January of 1985 he took a Project Manager position with Systems and Computer Technology, Malvern, Pennsylvania, where he managed the development of a student information system for Colorado State University, University of California at San Diego, Michigan State University, and University of Southwest Louisiana. In March of 1987 he took a Database Analyst position with Sorbus, Frazer, Pennsylvania. In January of 1988 he was transferred to Bell Atlantic Enterprises, Princeton, New Jersey, where he administers database management systems to support application systems development within Bell Atlantic Enterprise subsidiaries. He currently holds the title of Senior Database Analyst.