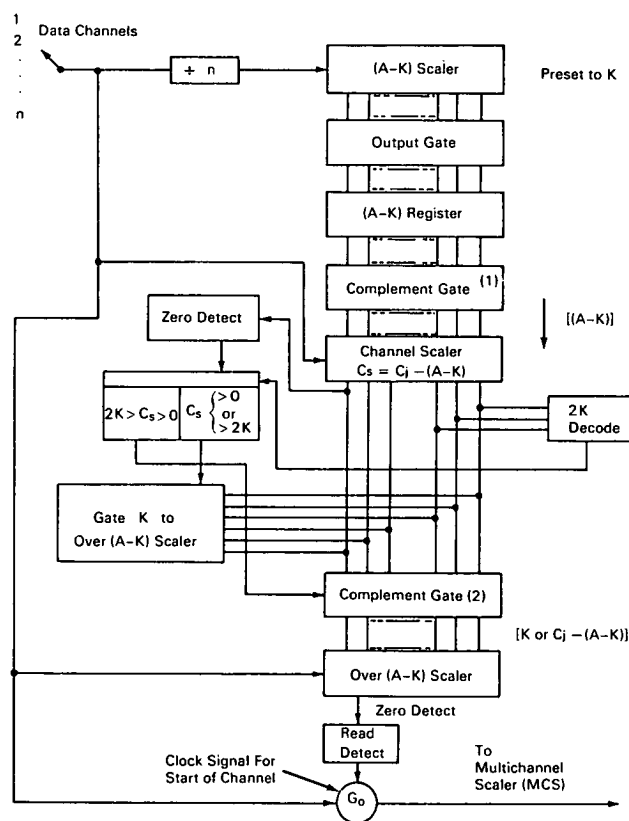


# AEC-NASA TECH BRIEF



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## Digital Filter Suppresses Effects of Nonstatistical Noise Bursts on Multichannel Scaler Digital Averaging Systems



### The problem:

To suppress the effects of large, nonstatistical noise bursts on data that have been averaged over many sweeps of a multichannel scaler. Digital data entering a multichannel scaler digital averaging system with a small signal-to-noise ratio will obscure or confuse data previously averaged by the system if the entering data pulse contains large, nonstatistical noise bursts. A

means of filtering out such distorted pulse data is required.

### The solution:

A digital noise filter, interposed between the sampled channels and the multichannel scaler digital averaging system, makes use of binary logic circuitry to compare the number of counts per channel ( $C_j$ ) with the average number of counts per channel ( $A_m$ ).

(continued overleaf)

If the channel count falls within  $A_m \pm K$  ( $\pm K$  = preset allowable count fluctuation), the channel count is accepted and passed to the multichannel scaler digital averaging system.

#### How it's done:

Before operating the digital filter (DF), a value  $K$  is chosen as the acceptable fluctuation limits about the average number of counts per channel  $A_m$ . The filter consists of four parallel paths: (1) a division scaler set to give one output count for every  $n$  input counts ( $n$  = number of channels being swept); (2) a channel scaler in which incoming counts for each channel  $j$  are accumulated; (3) the "OVER ( $A-K$ )" scaler which is used to output data to the multichannel scaler; and (4) a gate  $G_0$  to pass the output of the DF to the input of the multichannel scaler (MCS).

Before each sweep of the data channels, the ( $A-K$ ) scaler is preset to  $\bar{K}$ , the binary complement of  $K$ . At the end of the sweep, the content of this scaler is transferred by a parallel output gate to the ( $A-K$ ) register, and the ( $A-K$ ) scaler is cleared and reset to  $\bar{K}$ .

The DF reads each channel for a channel period of  $T_j$  during each sweep. At the start of each channel period  $T_j$ , the complement of ( $A-K$ ), ( $\overline{A-K}$ ), is non-destructively shifted by the complement gate (1) to the channel scaler. The channel scaler is thus preset to a value  $-(A-K)$ . The total counts in the channel scaler at the end of each period  $T_j$  will be  $C_s = C_j - (A-K)$ . The incoming counts for a data channel  $C_j$  fed to the channel scaler thus advance the scaler contents toward zero. If  $C_s$  reaches zero, then  $C_j$  exceeds the lower limit ( $A-K$ ). To determine whether or not  $C_j$  has exceeded the upper limit ( $A+K$ ), a  $2K$  coincidence circuit monitors the channel scaler. If the channel scaler count  $C_s$  passes zero but does not pass  $2K$ , then  $C_j$  lies between the limits ( $A+K$ ) and ( $A-K$ ) and the channel information is acceptable. If the count  $C_s$  is less than zero or greater than  $2K$ , the allowable limits have been passed and the data are not acceptable. For  $(A-K) < C_j < (A+K)$ , the binary complement of the contents of the channel scaler [effectively  $C_j - (A-K)$ ] are passed by complement gate (2) into the "OVER ( $A-K$ )" scaler. For  $(A+K) < C_j < (A-K)$ , the complement  $\bar{K}$  is shifted into the "OVER ( $A-K$ )" scaler.

At the start of the next channel period  $T_{j+1}$ , gate  $G_0$  is opened and remains so until the "OVER ( $A-K$ )" scaler is driven to zero. While  $G_0$  is open, either  $C_j$

– ( $A-K$ ) or  $K$  (depending on the contents of the scaler) is transferred into channel  $j+1$  of the MCS. Thus channel data accumulated in  $T_j$  are stored in channel  $j+1$ .

During  $T_1$ , no previous channel information is available, so channel 1 is used to accumulate the sum of all  $(A-K)_m$  for the whole run. Thus for  $T_1$  only, circuitry not shown closes  $G_0$ , permitting the value  $A-K$  to enter channel 1 of the MCS if  $C_1 > (A-K)$ . A small error occurs in the accumulated sum of the  $(A-K)_m$  whenever  $C_1 < (A-K)$ .

#### Notes:

1. A signal strength of  $S > K$  is recorded as  $K$  rather than  $S+K$ . In this case, either the DF is bypassed or  $K$  must be large relative to  $S$ .
2. For the DF described,  $C_j$  (counts accumulated for a given channel in each pass) is not signal plus background  $s+b$ , but the much smaller number  $s+K$ .
3. This DF has been constructed with 155 integrated circuit devices and has fluctuation limits from 8 to 768 about the average count (2 figure binary precision).
4. Additional details are contained in Review of Scientific Instruments, vol. 37, no. 6, June 1966, p. 769-771.
5. Inquiries concerning this innovation may be directed to:

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#### Patent status:

Inquiries about obtaining rights for commercial use of this innovation may be made to:

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