

Order

Linear Extensions and Comparable Pairs in Partial Orders

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Linear Extensions and Comparable Pairs in Partial Orders

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Abstract We study the number of linear extensions of a partial order with a given proportion of comparable pairs of elements, and estimate the maximum and minimum possible numbers. We also show that a random interval partial order on n elements has close to a third of the pairs comparable with high probability, and the number of linear extensions is $n! 2^{-\Theta(n)}$ with high probability.

Keywords Partial orders · Linear extensions · Comparable pairs · Concentration inequalities

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1 Introduction

Initially our interest was in the random interval order P_n , where we pick n intervals independently and uniformly at random in $(0, 1)$ (see the start of Section 6 for a precise definition). It turns out that with high probability about a third of the $\binom{n}{2}$ possible pairs are comparable, and the number $e(P_n)$ of linear extensions is $n! 2^{-\Theta(n)}$. But, given the proportion of comparable pairs, is this a large number of linear extensions? How few or many could there be? The main part of the paper is devoted to answering such questions.

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We investigate how few or many linear extensions a general partially ordered set Q or Q_n of n points may have, when it has a certain proportion of the possible edges in its comparability graph (we recall basic definitions in the next section). It is well known – see for example the Proposition at the end of [12] – that the comparability graph determines the number of linear extensions $e(Q)$, in the sense that posets with isomorphic (undirected) comparability graphs have the same number of linear extensions.

This number $e(Q)$ is an important invariant of a poset which is related to, for example, the worst case number $c(Q)$ of pairwise comparisons required to determine a (hidden) total order when we are initially given a partial order compatible with it. It is easy to see by a bisection argument that $c(Q) \geq \log_2 e(Q)$; and there has been much work over the years on how close to the truth this is, with one well known result being the bound of Fredman [15] that $c(Q) \leq \log_2 e(Q) + 2n$, see for example [8] for a recent algorithmic result in this direction and references to earlier literature.

We focus primarily on posets Q_n of n points where the comparability graph is dense, with about $\delta \binom{n}{2}$ edges for some constant $\delta \in (0, 1)$ – we refer to such a poset as a dense poset – but we are interested also in what happens when $\delta = \delta(n)$ tends to 0 or 1 as $n \rightarrow \infty$. Given a positive integer n and $0 < \delta < 1$, we let

$$f^+(n, \delta) = \max \left\{ e(Q_n) : \text{comp}(Q_n) \geq \delta \binom{n}{2} \right\}$$

and

$$f^-(n, \delta) = \min \left\{ e(Q_n) : \text{comp}(Q_n) \leq \delta \binom{n}{2} \right\},$$

where $\text{comp}(Q_n)$ denotes the number of edges in the comparability graph of the poset Q_n . Of course $1 \leq f^-(n, \delta), f^+(n, \delta) \leq n!$.

We are interested to study the asymptotic values of functions $f^+(n, \delta)$ and $f^-(n, \delta)$. Some steps in the investigation of the precise maximum values $f^+(n, \delta)$ were taken in [14], which concentrated on small numbers of edges rather than the comparatively large values we shall mainly examine. It is known from [14] that a partial order with the maximum number of linear extensions for a given number of vertices and comparable pairs is a semiorder. We are not aware of any previous work on the minimum values $f^-(n, \delta)$.

The main results of our paper are outlined in Theorem 1 below. The main thrust is that, for each fixed $0 < \delta < 1$, we have $f^+(n, \delta) = n! 2^{-\Theta(n)}$ and $f^-(n, \delta) = 2^{\Theta(n)}$. Let us spell this out more fully.

Theorem 1. *The maximum values $f^+(n, \delta)$ and minimum values $f^-(n, \delta)$ satisfy the following, as $n \rightarrow \infty$.*

(a) *For each $0 < \delta < 1$ there are constants $c_1(\delta)$ and $c_2(\delta)$ such that*

$$0 < c_1(\delta) < \left(\frac{f^+(n, \delta)}{n!} \right)^{1/n} < c_2(\delta) < 1 \quad \text{for } n \text{ sufficiently large;}$$

and further, $c_1(\delta) \rightarrow 1$ as $\delta \rightarrow 0$, and $c_2(\delta) \rightarrow 0$ as $\delta \rightarrow 1$.

(b) For each $0 < \delta < 1$ there are constants $c_3(\delta)$ and $c_4(\delta)$ such that

$$1 < c_3(\delta) < f^-(n, \delta)^{1/n} < c_4(\delta) < \infty \quad \text{for } n \text{ sufficiently large;}$$

and further, $c_3(\delta) \rightarrow \infty$ as $\delta \rightarrow 0$, and $c_4(\delta) \rightarrow 1$ as $\delta \rightarrow 1$.

Plan of the paper

After some preliminaries in the next section, in Section 3 we study posets consisting of either disjoint chains or disjoint antichains, and derive bounds on the numbers of linear extensions and comparable pairs.

In Section 4 we consider the maximum values $f^+(n, \delta)$. We prove part (a) of Theorem 1, and give values for $c_1(\delta)$ and $c_2(\delta)$, see (15) and (18). Also, in (20) we give a formula for $f^+(n, \delta)$ in the very dense case, and in Proposition 5 we describe the rate of convergence of $(f^+(n, \delta)/n!)^{1/n}$ to 0 as $\delta = \delta(n) \rightarrow 1$.

In Section 5 we similarly consider the minimum values $f^-(n, \delta)$. We prove part (b) of Theorem 1, and give values for $c_3(\delta)$ and $c_4(\delta)$, see (24) and (22). Also, in (26) we give a formula for $f^-(n, \delta)$ in the very sparse case, and in Proposition 10 we describe the rate of growth of $f^-(n, \delta)^{1/n}$ as $\delta = \delta(n) \rightarrow 0$.

In Section 6 we investigate the random interval order P_n (getting the historical sequence of events out of order): we show that whp $\text{comp}(P_n)$ is about $\frac{1}{3}\binom{n}{2}$ and $e(P_n)$ is $n!2^{-\Omega(n)}$, and thus whp $e(P_n) = n!2^{-\Theta(n)}$ by part (a) of Theorem 1. Finally, in Section 7, we make a few concluding remarks and conjectures; and in particular we conjecture that $(f^+(n, \frac{1}{2})/n!)^{1/n} \rightarrow \frac{1}{2}$ and $f^-(n, \frac{1}{2})^{1/n} \rightarrow 2$, as $n \rightarrow \infty$.

2 Preliminaries

A *poset* (partially ordered set) P is a set of points equipped with an irreflexive, antisymmetric transitive relation \prec , see e.g. [4], [27]. All posets in this paper are finite. Typical notations for a poset will be P or Q_n , where n indicates the number of points. A *linear extension* of a poset is a total order $<$ on the ground set of points such that whenever $x \prec y$ in the partial order, then we have $x < y$ in the total order too. The number of linear extensions of a poset P is denoted by $e(P)$.

We say that two points $x \neq y$ are *comparable* in a poset if $x \prec y$ or $y \prec x$; otherwise they are *incomparable*. The *comparability graph* G is the undirected graph with vertex set the set of points and an edge between x and y if and only if they are comparable: the *incomparability graph* is the complement \overline{G} of the comparability graph G .

A *chain* in P is a set T of points, any two of which are comparable: such a set T can be totally ordered, i.e. enumerated as $T = \{t_1, t_2, \dots, t_r\}$ with $t_1 \prec t_2 \prec \dots \prec t_r$. The maximum number of points in a chain is called the *height* of P and is denoted by $h(P)$. By contrast, an *antichain* in P is a set of points, no two of which are comparable. The maximum number of points in an antichain is called the *width* of P and is denoted by $w(P)$.

Recall that the comparability graph of P determines $e(P)$ [28]. If P and Q are posets on the same set of points, and the comparability graph of P is a strict subgraph of that of Q , then clearly $e(P) > e(Q)$ – see e.g. [25]. However, the naïve intuition that the more edges in the comparability graph the fewer linear extensions

is inaccurate: for example the six-vertex poset consisting of two three-element chains with no comparabilities between them has 6 edges in the comparability graph and $\binom{6}{3} = 20$ linear extensions (as we shall see), but the six-element poset comprising of one maximum element, one minimum element and an antichain of four elements all between the maximum and minimum element has 9 edges in the comparability graph and $4! = 24$ linear extensions. We thus word our results in terms of posets which have at most, or at least, a certain proportion of edges in the comparability graph to mitigate these lack of monotonicity issues.

We shall sometimes use the *level structure* of the poset P on a set V of points. The first level L_1 consists of all minimal elements, that is, those points $x \in V$ for which $y \preceq x \Rightarrow y = x$. The second level consists of the minimal elements of the induced subposet on $V \setminus L_1$ and generally the i th level L_i consists of the minimal elements of the induced subposet on $V \setminus \cup_{j=1}^{i-1} L_j$. Note that the number of non-empty levels is $h(P) \leq |V|$ and that each level is an antichain.

For positive functions $f(n)$ and $g(n)$ we as usual write $f(n) = O(g(n))$ if there is a constant K such that $f(n) \leq K \cdot g(n)$ for all large enough n , and similarly $f(n) = \Omega(g(n))$ if $f(n) \geq K \cdot g(n)$ for all large enough n . If $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$, then we write $f(n) = \Theta(g(n))$. We also use the standard notations $f(n) = o(g(n))$ if $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$, and $f(n) \sim g(n)$ if $f(n) = (1 + o(1))g(n)$ as $n \rightarrow \infty$. It is a little less standard to use the notation $f(n) = \omega(g(n))$ when $g(n) = o(f(n))$, that is, when $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$.

When we talk later about properties of a random poset Q_n with n points, we shall say that Q_n has a given property \wp *with high probability* (whp) if

$$\mathbb{P}(Q_n \text{ has } \wp) \rightarrow 1 \quad \text{as } n \rightarrow \infty.$$

We use Stirling's formula $n! \sim \sqrt{2\pi n}(n/e)^n$, and make repeated use of the related inequality $n! \geq (n/e)^n$ which follows from the series expansion of e^n . Also, we write $[n]$ to denote the set of positive integers $\{1, \dots, n\}$.

Graphs and entropy

A subset of the vertices of a graph G , such that every two are connected by an edge, is called a *clique*. The *clique number* of G is the cardinality of the largest clique and the *chromatic number* $\chi(G)$ of G is the minimum number of colours needed to colour its vertices in such a way that any two vertices joined by an edge receive different colours. G is *perfect* if for every induced subgraph of G the chromatic number equals the clique number, see for example [22]. It is well known that the comparability graph of a poset is perfect, and the complement of a perfect graph is perfect.

Given a graph G on vertex set $[n]$, the *clique polytope* $\mathcal{C}(G)$ of G is defined by

$$\mathcal{C}(G) = \left\{ \mathbf{x} \in [0, 1]^n : \sum_{i \in C} x_i \leq 1 \text{ for each clique } C \text{ of } G \right\},$$

see for example [9]. Stanley [26] proved that, if G is the comparability graph of a poset Q_n , then the number of linear extensions $e(Q_n)$ is equal to the volume of $\mathcal{C}(G)$ multiplied by $n!$.

The *entropy* $H(G)$ of a graph G is defined by

$$H(G) = \min_{\mathbf{x} \in \mathcal{C}(G)} -\frac{1}{n} \sum_{i=1}^n \log_2 x_i.$$

It satisfies $0 \leq H(G) \leq \log_2 n$, and indeed $H(G) \leq \log_2 \chi(G)$ (see e.g. [24]). Let Q_n be a partial order on $[n]$, with comparability graph G and incomparability graph \overline{G} . Then, by [10, Theorem 2],

$$H(G) + H(\overline{G}) = \log_2 n.$$

We are interested here in graph entropy since there are bounds on the number of linear extensions of Q_n in terms of the entropy of G or \overline{G} . Kahn and Kim [18] proved

$$n! 2^{-nH(G)} \leq e(Q_n) \leq 2^{nH(\overline{G})}, \tag{1}$$

and Cardinal et al. [8] proved

$$e(Q_n) \geq 2^{\frac{1}{2}nH(\overline{G})}. \tag{2}$$

3 The chain and antichain examples

We introduce two important standard examples of partial orders where we can write down the number of linear extensions. The chain example $C = C(n_1, \dots, n_k)$ is the poset consisting of disjoint chains of sizes n_1, \dots, n_k with no comparabilities between the chains. We have

$$e(C) = \binom{n}{n_1, \dots, n_k} \leq k^n. \tag{3}$$

To see the equality here, when for example $k = 2$, observe that the only choices to make are which positions in the total order are occupied by the n_1 points from the first chain. It follows that for any poset Q_n with width at most k we have

$$e(Q_n) \leq k^n, \tag{4}$$

since by Dilworth's Theorem [11] a poset Q_n with $w(Q_n) \leq k$ can be partitioned into k or fewer chains.

The antichain example is the poset $A = A(n_1, \dots, n_k)$ consisting of disjoint antichains A_1, \dots, A_k of sizes n_1, \dots, n_k such that if $x \in A_i, y \in A_j$ and $i < j$ then $x \prec y$. Observe that the comparability graphs for the posets A and C are complementary (with the natural choice of sets of points). We have

$$e(A) = \prod_{i=1}^k n_i! = n! / \binom{n}{n_1, \dots, n_k} \geq n! / k^n. \tag{5}$$

It follows by considering the level structure that for any poset Q_n with height at most k we have

$$e(Q_n) \geq n! / k^n. \tag{6}$$

As an aside, recall that by the Kleitman-Rothschild Theorem [19], almost all posets Q_n on $[n]$ have height 3, and thus by (6) $e(Q_n) \geq n! 3^{-n}$ for almost all such posets. In fact, by [5] almost all posets Q_n have

$$(e(Q_n)/n!)^{1/n} \sim 2^{-\frac{3}{2}} \approx 0.35.$$

See [6] and the references there for much more precise results on this and on the average number (which is a factor of order \sqrt{n} larger than where $e(Q_n)$ is concentrated).

We shall be most interested in the special cases when the n_i are as balanced as possible. The balanced chain example $\tilde{C}(n, k)$ is defined, for all integers $1 \leq k \leq n$, as $C(n_1, \dots, n_k)$ where $\lfloor n/k \rfloor \leq n_1 \leq \dots \leq n_k \leq \lceil n/k \rceil$, and the n_i sum to n . Then $e(\tilde{C}(n, k)) \leq k^n$ by (3) or (4). Also, for each given point there are at most $\lceil n/k \rceil - 1 \leq \frac{n-1}{k}$ comparable points, so

$$\text{comp}(\tilde{C}(n, k)) \leq \frac{1}{2} n \frac{n-1}{k} = \frac{1}{k} \binom{n}{2}. \quad (7)$$

The balanced antichain example $\tilde{A}(n, k)$ is defined similarly, for all integers $1 \leq k \leq n$, as $\tilde{A}(n_1, \dots, n_k)$, where the n_i are as above. Then $e(\tilde{A}(n, k)) \geq n!/k^n$ by (5) or (6). Also, by (7) and taking complements

$$\text{comp}(\tilde{A}(n, k)) \geq \left(1 - \frac{1}{k}\right) \binom{n}{2}. \quad (8)$$

For $1 \leq t \leq (n+1)/2$, we shall also use the examples $C(t, n-t)$ and $A(t, n-t)$ where the poset consists of two disjoint chains in the former case and two disjoint antichains in the latter, of cardinalities t and $n-t$. We set $\tau := t/n$. Then

$$e(C(t, n-t)) = \binom{n}{t} \leq \left(\frac{e}{\tau}\right)^{\tau n}, \quad (9)$$

and

$$\text{comp}(C(t, n-t)) = \binom{n}{2} - \tau(1-\tau)n^2 \leq (1-\tau) \binom{n}{2}, \quad (10)$$

since $(1-\tau)n^2 \geq \frac{n-1}{2n}n^2 = \binom{n}{2}$. Also

$$e(A(t, n-t)) = t!(n-t)! = n! / \binom{n}{t} \geq n! \left(\frac{\tau}{e}\right)^{\tau n}, \quad (11)$$

and

$$\text{comp}(A(t, n-t)) = \tau(1-\tau)n^2 \geq \tau \binom{n}{2}. \quad (12)$$

Note that $f(\tau) = \left(\frac{\tau}{e}\right)^{\tau}$ is continuous and decreasing on $[0, 1]$, and $f(0) = 1$. Thus $g(\tau) = \left(\frac{e}{\tau}\right)^{\tau} = 1/f(\tau)$ is continuous and increasing on $[0, 1]$, with $g(0) = 1$.

4 Maximum numbers of extensions: $f^+(n, \delta)$

In this section we prove part (a) of Theorem 1, and give values for $c_1(\delta)$ and $c_2(\delta)$, see (15) and (18). Also, in (20) we give a formula for $f^+(n, \delta)$ in the very dense case, and in Proposition 5, we describe the rate of convergence of $(f^+(n, \delta)/n!)^{1/n}$ to 0 as $\delta = \delta(n) \rightarrow 1$.

Lower bounds

Let $0 \leq \delta < 1$. Let $k = \lceil (1 - \delta)^{-1} \rceil$, and let Q_n be the balanced antichain example $\tilde{A}(n, k)$. Then by (8), since $1 - \frac{1}{k} \geq \delta$, we have $\text{comp}(Q_n) \geq \delta \binom{n}{2}$; and since $k \leq \frac{2-\delta}{1-\delta}$ we have

$$e(Q_n) \geq n! k^{-n} \geq n! \left(\frac{1-\delta}{2-\delta} \right)^n.$$

Thus

$$\left(\frac{f^+(n, \delta)}{n!} \right)^{1/n} \geq \left(\frac{e(Q_n)}{n!} \right)^{1/n} \geq \frac{1-\delta}{2-\delta} > \frac{1}{2}(1-\delta). \quad (13)$$

This result covers all δ with $0 \leq \delta < 1$, but even for $\delta = 0$ the lower bound is only $\frac{1}{2}$, and we want a lower bound which approaches 1 as $\delta \rightarrow 0$.

Assume that $0 < \delta \leq \frac{1}{2}$. Let $t = \lceil \delta n \rceil$, let $\tau = t/n$, and let Q_n be the two-antichain example $A(t, n-t)$. By (12) since $\delta \leq \tau < \delta + \frac{1}{n}$ we have $\text{comp}(Q_n) \geq \delta \binom{n}{2}$; and then by (11), since $(x/e)^x$ is decreasing on $[0, 1]$ we have

$$\left(\frac{f^+(n, \delta)}{n!} \right)^{1/n} \geq \left(\frac{e(Q_n)}{n!} \right)^{1/n} \geq \left(\frac{\delta + \frac{1}{n}}{e} \right)^{\delta + \frac{1}{n}} \geq \left(\frac{2\delta}{e} \right)^{2\delta} \quad (14)$$

for $\frac{1}{n} \leq \delta \leq \frac{1}{2}$. By (13) and (14), we may set

$$c_1(\delta) = \begin{cases} \max \left\{ \frac{1-\delta}{2-\delta}, \left(\frac{2\delta}{e} \right)^{2\delta} \right\} & \text{if } 0 < \delta \leq \frac{1}{2} \\ \frac{1-\delta}{2-\delta} & \text{if } \frac{1}{2} < \delta < 1. \end{cases} \quad (15)$$

Note that the lower bound $c_1(\delta)$ tends to 1 as $\delta \rightarrow 0$, as desired.

Upper bounds

We first use a martingale concentration inequality to prove the upper bound $e(Q_n) \leq n! 2^{-\Omega(n)}$ on the number of linear extensions of a dense poset.

Lemma 2. *There is an absolute positive constant c , which can be taken to be $\frac{\log_2 e}{32} \approx 0.045$, such that the following holds. Let $0 < \delta \leq 1$, and let Q_n be a poset on n points with $\text{comp}(Q_n) \geq \delta \binom{n}{2}$. Then*

$$e(Q_n) \leq n! 2^{-c\delta^2 n}.$$

Proof. Let S_n be the set of all $n!$ permutations of $[n]$. For $\tau \in S_n$ let $g(\tau)$ be the number of conflicts of τ with the poset Q_n : that is, the number of ordered pairs (i, j) , such that $i < j$ in τ and $j \prec i$ in the partial order. Obviously, τ is a linear extension of Q_n if and only if $g(\tau) = 0$. If τ and τ' differ by one transposition σ ,

i.e. $\tau = \tau' \circ \sigma$ in the symmetric group (σ acting first say), then if σ interchanges x and y , it changes the number of conflicts by at most $n - 1$ for each of x and y , so

$$|g(\tau') - g(\tau)| \leq 2(n - 1). \quad (16)$$

When τ is chosen uniformly at random from S_n , we have that

$$\mathbb{E}(g(\tau)) = \text{comp}(Q_n)/2$$

since for each $x \neq y$, $\mathbb{P}(x < y) = \mathbb{P}(y < x) = \frac{1}{2}$ (where $<$ refers to the ordering τ). Put $s = \text{comp}(Q_n)/2 \geq \delta n(n - 1)/4$. We may now use Theorem 7.4 of [21], together with the inequality (16), to obtain

$$\begin{aligned} \mathbb{P}(g(\tau) = 0) &= \mathbb{P}(g(\tau) - s \leq -s) \\ &\leq \exp\left(\frac{-2s^2}{4n(n - 1)^2}\right) \leq e^{-\delta^2 n/32}. \end{aligned}$$

Hence

$$e(Q_n) = n! \mathbb{P}(g(\tau) = 0) \leq n! e^{-\delta^2 n/32}$$

which completes the proof. \square

The last lemma covers all δ with $0 < \delta \leq 1$, but even when $\delta = 1$ we just find $e(Q_n) \leq n! 2^{-cn}$. We need another result to show that when $\delta \rightarrow 1$ we have that $(f^+(n, \delta)/n!)^{1/n} = o(1)$. This follows from Lemma 4 below (which says nothing asymptotically if $\delta < 1 - 2/e \approx 0.264$). In fact we shall see that if $\delta \rightarrow 1$, but not too quickly, then $(f^+(n, \delta)/n!)^{1/n} = \Theta(1 - \delta)$. First, here is a preliminary observation: we omit the easy proof.

Observation 3. *Let Q be a partial order on a set S of at least two elements, and let $v \in S$. Suppose that $S' = S \setminus \{v\}$ is a chain, and v is incomparable to exactly x points in S' . Then $e(Q) = x + 1$.*

Lemma 4. *Let $0 < \delta < 1$ and suppose that $\text{comp}(Q_n) \geq \delta \binom{n}{2}$. Then*

$$e(Q_n) \leq n! e^{\frac{2}{1-\delta}} \left(\frac{e(1-\delta)}{2}\right)^n.$$

Proof. Let \overline{G} be the complement of the comparability graph G of the poset Q_n and let $E(\overline{G})$ be the set of edges of \overline{G} . Consider the vertices in the order $1, 2, \dots, n$ and let x_v be the ‘back-degree’ of vertex v in \overline{G} (that is, the number of edges uv in \overline{G} with $u < v$). Then, by considering building up the linear extension step-by-step and using Observation 3, we may see that

$$e(Q_n) \leq \prod_{v=1}^n (x_v + 1) \leq \left(\frac{1}{n} \sum_{v=1}^n (x_v + 1)\right)^n,$$

by the arithmetic-geometric means inequality. But

$$\sum_{v=1}^n (x_v + 1) = |E(\overline{G})| + n \leq (1 - \delta) \binom{n}{2} + n = n \left(\frac{1 - \delta}{2}(n - 1) + 1\right).$$

Hence

$$e(Q_n) \leq \left(\frac{1 - \delta}{2}n + 1\right)^n = \left(\frac{1 - \delta}{2}n\right)^n \left(1 + \frac{2}{(1 - \delta)n}\right)^n.$$

But the last factor is at most $e^{\frac{2}{1-\delta}}$, and using $n^n \leq e^n n!$ concludes the proof. \square

It follows from the last lemma that if $\delta \geq 0$ and $1 - \delta \geq d/n$ for some $d > 0$ then

$$(f^+(n, \delta)/n!)^{1/n} \leq \frac{1}{2}e^{\frac{2}{d}+1}(1 - \delta). \quad (17)$$

Given $0 \leq \delta < 1$ we may define $c_2(\delta)$ as follows. Let $n_0 = n_0(\delta) = 6/(1 - \delta)$; and let

$$c_2(\delta) = \min \left\{ 2^{-c\delta^2}, 2(1 - \delta) \right\}, \quad (18)$$

where $c > 0$ is the constant from Lemma 2. Then $c_2(\delta) < 1$, $c_2(\delta) \rightarrow 0$ as $\delta \rightarrow 1$; and by Lemma 4, and inequality (17) with $d = 6$ (noting that $e^{4/3} < 4$),

$$(f^+(n, \delta)/n!)^{1/n} \leq c_2(\delta) \quad \text{for all } n \geq n_0,$$

as required in part (a) of Theorem 1.

Now we consider the very dense case, when we may determine $f^+(n, \delta)$ exactly. For, if Q_n has i incomparable pairs, then clearly $e(Q_n) \leq 2^i$. Thus we always have

$$f^+(n, \delta) \leq 2^{\lfloor (1-\delta) \binom{n}{2} \rfloor}. \quad (19)$$

But, if $i \leq n/2$, then $\tilde{A}(n, n - i)$ has i incomparable pairs and $e(\tilde{A}(n, n - i)) = 2^i$. Hence

$$f^+(n, \delta) = 2^{\lfloor (1-\delta) \binom{n}{2} \rfloor} \quad \text{when } (1 - \delta) \binom{n}{2} \leq n/2. \quad (20)$$

Next we determine the rate of convergence of $(f^+(n, \delta)/n!)^{1/n}$ to 0 as $\delta = \delta(n) \rightarrow 1$.

Proposition 5. *Let $\delta = \delta(n) \rightarrow 1$ as $n \rightarrow \infty$. Then*

$$(f^+(n, \delta)/n!)^{1/n} = \Theta(\max\{1 - \delta, 1/n\}).$$

Proof. It suffices to consider the cases (a) $1 - \delta \geq 1/n$ and (b) $1 - \delta \leq 1/n$. In case (a), the lower bound from (13) and the upper bound from (17) show that $(f^+(n, \delta)/n!)^{1/n} = \Theta(1 - \delta)$. In case (b), $(1 - \delta) \binom{n}{2} \leq n/2$, so by (19) we have $f^+(n, \delta) \leq 2^{n/2}$: thus $1 \leq (f^+(n, \delta))^{1/n} \leq \sqrt{2}$, and so

$$1/n \leq (f^+(n, \delta)/n!)^{1/n} \leq e\sqrt{2}/n.$$

This completes the proof. \square

We proved two upper bounds on $e(Q_n)$ above, in Lemmas 2 and 4. To close this section, we consider briefly whether the upper bound in (1) based on entropy, namely $e(Q_n) \leq 2^{nH(\overline{G})}$, could have been useful here.

Let k be an integer at least 2, let $\delta = \frac{1}{k}$, and let Q_n be the balanced antichain example $\tilde{A}(n, k)$, with comparability graph G . We have $\text{comp}(Q_n) \geq \delta \binom{n}{2}$ by (8). But $H(G) \leq \log_2 \chi(G) \leq \log_2 k$, so $H(\overline{G}) = \log_2 n - H(G) \geq \log_2(n/k)$. Hence the upper bound from (1) is

$$2^{nH(\overline{G})} \geq (n/k)^n.$$

In the case $\delta = \frac{1}{2}$, the upper bound is at least $(n/2)^n \gg n!$, so of course this tells us nothing. For smaller δ (when $k \geq 3$) we see that the upper bound from (1) is at least

$$(n/k)^n = n!((1 + o(1))e\delta)^n,$$

which is at least a factor of about 2^n greater than the upper bound in Lemma 4.

5 Minimum numbers of extensions: $f^-(n, \delta)$

In this section we prove part (b) of Theorem 1, and give values for $c_3(\delta)$ and $c_4(\delta)$, see (24) and (22). Also, in (26) we give a formula for $f^-(n, \delta)$ in the very sparse case, and in Proposition 10, we describe the rate of growth of $f^-(n, \delta)^{1/n}$ as $\delta = \delta(n) \rightarrow 0$.

Upper bounds

Let $0 < \delta \leq 1$. Set $k = \lceil \delta^{-1} \rceil$, and let Q_n be the balanced chain example $\tilde{C}(n, k)$. We have that $\text{comp}(Q_n) \leq \frac{1}{k} \binom{n}{2} \leq \delta \binom{n}{2}$ by (7), and $e(Q_n) \leq k^n$ by (3). Thus

$$f^-(n, \delta)^{1/n} \leq \lceil \delta^{-1} \rceil \leq 2\delta^{-1}. \quad (21)$$

We wish also to show that $f^-(n, \delta)^{1/n} \rightarrow 1$ when $\delta = \delta(n) \rightarrow 1$ (which the last result does not give). Assume that $\frac{1}{2} \leq \delta < 1$. Let $t = \lceil (1 - \delta)n \rceil$, let $\tau = t/n$ and note that $1 - \delta \leq \tau < 1 - \delta + 1/n$. Let Q_n be the two-chain example $C(t, n - t)$. By (10) (since $t \leq (n + 1)/2$) we have $\text{comp}(Q_n) \leq (1 - \tau) \binom{n}{2} \leq \delta \binom{n}{2}$. Hence by (9)

$$f^-(n, \delta)^{1/n} \leq e(Q_n)^{1/n} \leq \left(\frac{e}{\tau}\right)^\tau < \left(\frac{e}{1 - \delta + 1/n}\right)^{1 - \delta + 1/n},$$

since $g(x) = (e/x)^x$ is increasing on $(0, 1)$. Thus if $\frac{1}{2} \leq \delta < 1$ and $1 - \delta \geq 1/n$ then

$$f^-(n, \delta)^{1/n} < \left(\frac{e}{2(1 - \delta)}\right)^{2(1 - \delta)},$$

again using the fact that $g(x)$ is increasing. We now see that we can set

$$c_4(\delta) = \begin{cases} \lceil \delta^{-1} \rceil & \text{if } 0 < \delta < \frac{1}{2} \\ \min \left\{ \lceil \delta^{-1} \rceil, \left(\frac{e}{2(1 - \delta)}\right)^{2(1 - \delta)} \right\} & \text{if } \frac{1}{2} \leq \delta < 1. \end{cases} \quad (22)$$

For, given $0 < \delta < 1$, by (21) and the last inequality we have

$$f^-(n, \delta)^{1/n} \leq c_4(\delta) \quad \text{once } n \geq 1/(1 - \delta).$$

Note that the upper bound here tends to 1 as $\delta \rightarrow 1$.

Lower bounds

Let us state a preliminary inequality corresponding to inequality (6). For any poset Q_n with height at most k we have

$$e(Q_n) \geq 2^{n-k}. \quad (23)$$

To see this, let Q_n have $h \leq k$ levels, and let level i contain $r_i \geq 1$ elements for $i = 1, \dots, h$. Then, using the inequality $r_i! \geq 2^{r_i-1}$, we have

$$e(Q_n) \geq \prod_{i=1}^h r_i! \geq \prod_{i=1}^h 2^{r_i-1} = 2^{n-h} \geq 2^{n-k}.$$

Our first lower bound covers the whole range $0 < \delta < 1$, and in particular when δ is near to 1.

Lemma 6. *Let $0 < \delta < 1$, and let Q_n be a poset on n vertices with $\text{comp}(Q_n) \leq \delta \binom{n}{2}$. Then*

$$e(Q_n) \geq \max\{2^{(1-\sqrt{\delta})(n-1)}, 2^{\frac{1}{2}(1-\delta)n}\}.$$

Note that $1 - \sqrt{\delta} > \frac{1}{2}(1 - \delta)$, and $1 - \sqrt{\delta} = \frac{1}{2}(1 - \delta) + O((1 - \delta)^2)$ as $\delta \rightarrow 1$. Thus the first term in the maximum above is slightly better in terms of δ . The second term is the one we use to define c_3 in (24) below, where we need the factor n not $(n - 1)$ in the exponent.

Proof. For both of the terms in the maximum, we will use inequality (23), together with an upper bound on the height $h = h(Q_n)$ following from $\binom{h}{2} \leq \delta \binom{n}{2}$. Note first that $h < \sqrt{\delta}(n - 1) + 1$; for otherwise

$$\binom{h}{2} \geq \frac{1}{2}(\sqrt{\delta}(n - 1) + 1)\sqrt{\delta}(n - 1) > \frac{1}{2}(\sqrt{\delta}n)\sqrt{\delta}(n - 1) = \delta \binom{n}{2}.$$

Thus by (23) we have $e(Q_n) > 2^{(1-\sqrt{\delta})(n-1)}$.

It remains to show that $e(Q_n) \geq 2^{\frac{1}{2}(1-\delta)n}$. Since $\delta < 1$ we have $e(Q_n) \geq 2$, so we may assume that $(1 - \delta)n > 2$. We **claim** that $h < (1 + \delta)n/2$: once we have proved this, the required lower bound on $e(Q_n)$ will follow directly from (23). But if $h \geq (1 + \delta)n/2$, then $\delta \binom{n}{2} \geq \binom{h}{2} \geq \frac{1}{8}(1 + \delta)n((1 + \delta)n - 2)$. This may be rewritten as

$$n\delta^2 - 2(n - 1)\delta + n - 2 \leq 0;$$

and this quadratic inequality gives $(1 - \delta)n \leq 2$, which contradicts $(1 - \delta)n > 2$. This completes the proof of the claim, and thus of the lemma. \square

We may see quickly that $f^-(n, \delta)^{1/n} \rightarrow \infty$ when $n \rightarrow \infty$ and $\delta \rightarrow 0$, as required in part (b) of Theorem 1. We noted in the proof of the last lemma that the height $h = h(Q_n)$ satisfies $h < \sqrt{\delta}(n - 1) + 1$, so $h < \sqrt{\delta}n + 1$. Hence, by inequality (6) (and using $n! \geq (n/e)^n$),

$$e(Q_n)^{1/n} \geq \frac{n}{e(\delta^{\frac{1}{2}}n + 1)} = \frac{1}{e} \delta^{-\frac{1}{2}} \left(1 + \frac{1}{\delta^{\frac{1}{2}}n}\right)^{-1}.$$

But, if $n \geq 11\delta^{-\frac{1}{2}}$, then

$$\left(1 + \frac{1}{\delta^{\frac{1}{2}}n}\right)^{-1} \geq 1 - \frac{1}{\delta^{\frac{1}{2}}n} \geq \frac{10}{11} > \frac{\epsilon}{3}.$$

Hence, for each $0 < \delta \leq 1$,

$$f^-(n, \delta)^{1/n} \geq \frac{1}{3} \delta^{-\frac{1}{2}} \quad \text{for all } n \geq 11\delta^{-\frac{1}{2}}.$$

This gives a lower bound as required on $f^-(n, \delta)^{1/n}$; but we can obtain better lower bounds for this case in terms of the dependence on δ . In the bound in Lemma 8 below we essentially replace $\delta^{-\frac{1}{2}}$ by δ^{-1} .

We first recall a preparatory lemma. A proof of this can be found in for example [1] Theorem 6.3 – the result is also stated as exercise 57 to Chapter 3 of [27] and the equality part to exercise 20 in [20, subsection 5.1.4].

Lemma 7. Let Q_n be a poset on n vertices, and for each $t \in Q_n$ let $\lambda_t := |\{s \in Q_n : s \preceq t\}|$ be the size of the principal downset for t (often referred to as the hook length). Then

$$e(Q_n) \geq \frac{n!}{\prod_{t \in Q_n} \lambda_t},$$

with equality if and only if Q_n is a downward-branching forest.

Lemma 8. Let $0 < \delta \leq 1$ and let Q_n be a poset on n vertices with $\text{comp}(Q_n) \leq \delta \binom{n}{2}$. Then

$$e(Q_n) \geq n! \left(\frac{2}{\delta(n-1) + 2} \right)^n \geq e^{1-2/\delta} \left(\frac{2}{e\delta} \right)^n.$$

Proof. With notation as in the last lemma,

$$\sum_{t \in Q_n} \lambda_t = \text{comp}(Q_n) + n,$$

and by the arithmetic-geometric means inequality

$$\left(\frac{\sum_{t \in Q_n} \lambda_t}{n} \right)^n \geq \prod_{t \in Q_n} \lambda_t.$$

Hence, by Lemma 7

$$\begin{aligned} e(Q_n) &\geq \frac{n!}{\prod_{t \in Q_n} \lambda_t} \geq n! \left(\frac{n}{\sum_{t \in Q_n} \lambda_t} \right)^n \\ &= n! \left(\frac{n}{n + \text{comp}(Q_n)} \right)^n \geq n! \left(\frac{n}{n + \delta \binom{n}{2}} \right)^n \\ &= n! \left(\frac{2}{\delta(n-1) + 2} \right)^n, \end{aligned}$$

which gives the first inequality required. Using $n! \geq \left(\frac{n}{e}\right)^n$ and $1 + x \leq e^x$, we see that the last term is at least

$$\left(\frac{n}{e}\right)^n \left(\frac{2}{\delta(n-1) + 2} \right)^n = \left(\frac{2}{e\delta}\right)^n \left(1 + \frac{\frac{2}{\delta} - 1}{n} \right)^{-n} \geq \left(\frac{2}{e\delta}\right)^n e^{1-2/\delta},$$

which completes the proof. \square

By the last lemma, if $n \geq 6/\delta$ then

$$f^-(n, \delta)^{1/n} \geq e^{-\frac{2}{\delta n}} \frac{2}{e\delta} \geq e^{-\frac{1}{3}} \frac{2}{e\delta} \geq \frac{1}{2\delta}.$$

Thus, using also Lemma 6, we see that we may set

$$c_3(\delta) = \max \left\{ 2^{\frac{1}{2}(1-\delta)}, \frac{1}{2\delta} \right\}. \quad (24)$$

Next we consider the very sparse case. By Lemma 8, if $\text{comp}(Q_n) = j$ (so the density $\delta = \delta(Q_n)$ satisfies $\delta(n-1)/2 = j/n$) then

$$e(Q_n) \geq n! (1 + \delta(n-1)/2)^{-n} = n! (1 + j/n)^{-n} \geq n! e^{-j}. \quad (25)$$

If $\text{comp}(Q_n) = j \leq n/2$, then we can have $e(Q_n) = n! 2^{-j}$ (when the comparability graph is a matching). Inequality (25) is sufficient to prove Proposition 10 below, but we can improve it, and obtain the equality (26) which matches (20) nicely. We need a preliminary result, perhaps of interest in its own right.

Proposition 9. *For each partial order Q on $[n]$ we have*

$$e(Q) \geq n! 2^{-\text{comp}(Q)},$$

and equality holds if and only if the comparability graph is a (partial) matching.

Proof. Let us think of a partial order Q as the set of ordered pairs (u, v) of distinct elements such that $u < v$ in Q . Observe that $\text{comp}(Q) = |Q|$. Suppose that $|Q| = k$. There is an increasing sequence Q_0, Q_1, \dots, Q_k of partial orders on $[n]$ such that $|Q_i| = i$ for $i = 0, \dots, k$ and $Q_k = Q$. Thus the inequality in the proposition will follow from the next claim.

Claim. Let Q and Q' be partial orders on $[n]$ such that $Q' = Q \cup \{(u, v)\}$ where $(u, v) \notin Q$. Then $e(Q) \leq 2e(Q')$.

To prove the claim, note first that $(x, u) \in Q'$ implies $(x, v) \in Q$ (since $(x, v) \in Q'$ by transitivity, and $(x, v) \neq (u, v)$), and similarly $(v, y) \in Q'$ implies $(u, y) \in Q$. Let $LE(Q)$ denote the set of linear extensions of Q . Let $\pi \in LE(Q)$ with $\pi(u) > \pi(v)$, and let π' be formed from π by swapping u and v : it suffices to show that $\pi' \in LE(Q')$. Consider u first. If $(x, u) \in Q'$ then, since $(x, v) \in Q$ (as we noted), we have $\pi'(x) = \pi(x) < \pi(v) = \pi'(u)$; $\pi'(u) < \pi'(v)$; and if $(u, y) \in Q'$ for some $y \neq v$ then $(u, y) \in Q$ and so $\pi'(u) < \pi(u) < \pi(y) = \pi'(y)$. Considering v similarly, we see that $(x, v) \in Q'$ implies $\pi'(x) < \pi'(v)$ and $(v, y) \in Q'$ implies $\pi'(v) < \pi'(y)$. It follows that $\pi' \in LE(Q')$, as required.

Finally, consider when equality holds. By symmetry, we must have equality when the comparability graph is a matching. Now suppose it is not a matching, and note that we may choose Q_2 to consist of two comparable pairs $(u, x), (v, x)$ or $(x, u), (x, v)$. But now $e(Q_2) = n!/3$; and so, arguing as before,

$$e(Q) \geq (4/3)n! 2^{-\text{comp}(Q)} > n! 2^{-\text{comp}(Q)},$$

which completes the proof. \square

By the last result and the discussion preceding it we have

$$f^-(n, \delta) = n! 2^{-\lceil \delta \binom{n}{2} \rceil} \quad \text{when } \delta(n-1) \leq 1. \quad (26)$$

Now we determine the rate of growth of $f^-(n, \delta)^{1/n}$ as $\delta = \delta(n) \rightarrow 0$.

Proposition 10. *Let $\delta = \delta(n) \rightarrow 0$ as $n \rightarrow \infty$. Then*

$$f^-(n, \delta)^{1/n} = \Theta(\min\{\delta^{-1}, n\}).$$

Proof. It suffices to consider the cases (a) $\delta \geq 1/n$ and (b) $0 < \delta \leq 1/n$. In case (a), the upper bound from (21) and the lower bound from Lemma 8 show that $f^-(n, \delta)^{1/n} = \Theta(\delta^{-1})$. In case (b), let $j = \lceil \delta \binom{n}{2} \rceil$, and note that $j \leq n/2$. Hence if $\text{comp}(Q_n) \leq j$ then by (25) or Proposition 9

$$e(Q_n) \geq n! e^{-j} \geq n! e^{-n/2}.$$

Thus

$$n \geq (n!)^{1/n} \geq f^-(n, \delta)^{1/n} \geq (n!)^{1/n} e^{-\frac{1}{2}} \geq n e^{-\frac{3}{2}},$$

and so $f^-(n, \delta)^{1/n} = \Theta(n)$, as required. \square

Example: the lattice of subsets of $[t]$

An interesting example of a poset is the lattice of $n = 2^t$ subsets of $[t]$, ordered by inclusion. Let us call this poset L_n . It is easy to see that there are 3^t pairs A, B with $A \subseteq B \subseteq [t]$, and it follows that $\text{comp}(L_n) = 3^t - 2^t = n^2 (\frac{3}{4})^t - n$: thus

$$\delta(n) := \frac{\text{comp}(L_n)}{\binom{n}{2}} \sim 2 \cdot \left(\frac{3}{4}\right)^{\log_2 n} = 2n^{-\log_2(4/3)} = o(1).$$

(Note that $\log_2(4/3) \approx 0.415$.) Brightwell and Tetali [7], improving on [23], give a very good estimate of $e(L_n)$, which by Stirling's formula implies that

$$(e(L_n)/n!)^{1/n} \sim ((\pi e/2) \log_2 n)^{-1/2}.$$

Thus we are not close to having as many linear extensions as possible (given $\delta = \delta(n)$), but closer than to the opposite case of having as few as possible. Indeed, as $n \rightarrow \infty$, by part (a) of Theorem 1, $(f^+(n, \delta)/n!)^{1/n} \rightarrow 1$, and by Proposition 10, $f^-(n, \delta)^{1/n} = \Theta(\delta^{-1}) = \Theta(n^{\log_2(4/3)})$.

Let us close this section by considering briefly whether the lower bounds (1) and (2) on $e(Q_n)$ based on entropy can yield good lower bounds on $f^-(n, \delta)$. The first lower bound $n! 2^{-nH(G)}$ tells us nothing when δ is at least about $\frac{1}{2}$. For let Q_n be the balanced chain example $\tilde{C}(n, 2)$, with comparability graph G . Then $\text{comp}(Q_n) \sim \frac{1}{2} \binom{n}{2}$. Also $\chi(\overline{G}) = 2$, so $H(\overline{G}) \leq 1$ and thus $H(G) \geq \log_2(n/2)$. Hence

$$n! 2^{-nH(G)} \leq n! 2^{-n \log_2(n/2)} = n! \left(\frac{2}{n}\right)^n = o(1).$$

Now consider smaller δ , say $\delta = 1/k$ where $k \geq 3$. Consider the balanced chain example $\tilde{C}(n, k)$. Arguing as above we have $\chi(\overline{G}) = k$, so $H(\overline{G}) \geq \log_2(n/k)$. Hence

$$n! 2^{-nH(G)} \leq n! 2^{-n \log_2(n/k)} = n! \left(\frac{k}{n}\right)^n = \left(\frac{1+o(1)}{e\delta}\right)^n.$$

Thus we do not quite match the bound $(\frac{1}{2\delta})^n$ above.

The second lower bound (2), namely $2^{\frac{1}{2}nH(\overline{G})}$, does give a lower bound of the form $2^{\Omega(n)}$ for each $0 < \delta < 1$. For if $\text{comp}(Q_n) \geq \delta \binom{n}{2}$, then as we have seen $e(Q_n) \geq c_3(\delta)^n$ for n sufficiently large, and so for the comparability graph G we have by the upper bound in (1) that

$$2^{\frac{1}{2}nH(\overline{G})} \geq e(Q_n)^{\frac{1}{2}} \geq (c_3(\delta)^{\frac{1}{2}})^n.$$

In particular this gives

$$f^-(n, \delta)^{1/n} \geq 2^{\frac{1}{2}H(\overline{G})} = \Omega(\delta^{-\frac{1}{2}}) \text{ as } \delta \rightarrow 0.$$

However, using (2) we do not obtain $f^-(n, \delta)^{1/n} = \Omega(\delta^{-1})$ as $\delta \rightarrow 0$. For, let $k = \lceil \delta^{-1} \rceil$, and let Q_n be the balanced chain example $\tilde{C}(n, k)$ (as at the start of this section), with comparability graph G . Recall that $\text{comp}(Q_n) \leq \delta \binom{n}{2}$, and $H(\overline{G}) \leq \log_2 k$. Thus,

$$2^{\frac{1}{2}H(\overline{G})} \leq k^{\frac{1}{2}} = \lceil \delta^{-1} \rceil^{\frac{1}{2}} < 2\delta^{-\frac{1}{2}}.$$

6 Random interval partial orders

Here we examine random interval posets, see [13]. Given a family of n closed intervals $I_j = [a_j, b_j]$ in the real line for $j = 1, \dots, n$, we form an *interval partial order* on $\{1, \dots, n\}$ by setting $i < j$ if and only if $b_i < a_j$. Thus two elements of the poset are comparable when their corresponding intervals do not intersect, and otherwise they are incomparable. A chain of the poset is a set of pairwise non-intersecting intervals; and an antichain is a set of pairwise intersecting intervals (that is, by Helly's Theorem, a set of intervals containing a common point).

Suppose that we start with $2n$ independent random variables X_1, \dots, X_n and Y_1, \dots, Y_n , each from the uniform distribution on $[0, 1]$; and form n closed intervals I_j , for $j = 1, \dots, n$, where $I_j = [X_j, Y_j]$ if $X_j < Y_j$ and $[Y_j, X_j]$ otherwise. (The event that $X_j = Y_j$ has probability zero so can be ignored.) These intervals I_j yield a *random interval poset* P_n . It would not matter if we replaced the uniform distribution on $[0, 1]$ by any continuous distribution. Also equivalently, we could generate the intervals from a uniform random perfect matching on $\{1, \dots, 2n\}$.

It is easy to see that the probability that two random intervals are comparable is $1/3$ – see e.g. [17]. Indeed, for $i \neq j$, the probability that the four end points of the intervals I_i and I_j come in any given order is $1/4! = 1/24$; and there are exactly 8 orders such that the intervals are disjoint. It follows that the number Z of edges in the comparability graph satisfies $\mathbb{E}[Z] = \frac{1}{3} \binom{n}{2}$.

Also, given any values $x_1, \dots, x_n, y_1, \dots, y_n$, as end points of n intervals, changing the values of x_k and y_k can affect only the edges incident to node k in the comparability graph, so the number of edges can vary by at most $n - 1$. Hence by the bounded differences inequality (Lemma 1.2 of [21], or see for example Theorem 6.3 of [2]), for each $t > 0$,

$$\mathbb{P}(|Z - \mathbb{E}[Z]| \geq t) \leq 2 \exp\left(-\frac{2t^2}{n(n-1)^2}\right).$$

Thus, for $0 < a < 1$

$$\mathbb{P}\left(\left|Z - \frac{1}{3} \binom{n}{2}\right| \geq a \binom{n}{2}\right) \leq 2 \exp\left(-\frac{1}{2} a^2 n\right)$$

and we see that a random interval order is dense whp.

Again, we want to estimate how many linear extensions P_n has. Let A_n denote the size of a largest set of intersecting intervals in our family of n random intervals

– in the poset language, this is a maximum antichain. It was shown in [17] that whp $A_n = \frac{n}{2} + o(n)$. But $e(P_n) \geq A_n!$ so by Stirling's formula, we have that whp

$$\log_2 e(P_n) \geq \log_2 \left(\frac{n}{2} + o(n) \right)! = \left(\frac{1}{2} + o(1) \right) n \log_2 n.$$

We shall refine the idea above to show that a random interval partial order has many more linear extensions.

Theorem 11. *The number $e(P_n)$ of linear extensions of a random interval partial order P_n satisfies*

$$e(P_n) = n! 2^{-\Theta(n)} \text{ whp.}$$

Proof. By Lemma 2 it will suffice for us to show that $e(P_n) \geq n! 2^{-O(n)}$ whp; and in fact we shall show that

$$e(P_n) \geq n! 2^{-(6+o(1))n} \text{ whp.} \quad (27)$$

For positive integers i and j with i odd and $i < 2^j$, let $I(i, j)$ denote the interval $(\frac{i-1}{2^j}, \frac{i+1}{2^j})$. Thus for each $j \geq 1$ there are 2^{j-1} intervals $I(i, j)$; and these intervals partition $[0, 1]$, except that they omit the $2^{j-1} + 1$ end points $\frac{i}{2^{j-1}}$ for $0 \leq i \leq 2^{j-1}$.

Let $0 < a < b < 1$ and consider the interval (a, b) . Let $j(a, b)$ be the least integer $j \geq 1$ such that $\frac{i}{2^j} \in (a, b)$ for some positive integer i . There is a unique such i : for if there were at least two such i , then there would be at least one even i , and we could replace $i/2^j$ by $(i/2)/2^{j-1}$, giving a smaller value of j and contradicting the definition of j . Since i is unique, we may call it $i(a, b)$.

We denote the interval $I(i(a, b), j(a, b))$ by $J(a, b)$. For positive integers i and j with i odd and $i < 2^j$, let $\mathcal{A}(i, j)$ be the random set of all intervals (a, b) amongst I_1, \dots, I_n such that $J(a, b) = I(i, j)$. Each interval (a, b) in $\mathcal{A}(i, j)$ contains the point $i/2^j$, so the sets $\mathcal{A}(i, j)$ are antichains. It follows also that, if the midpoint of $I(i, j)$ (i.e. $i/2^j$) is less than the midpoint of $I(i', j')$ (i.e. $i'/2^{j'}$) then no interval in $\mathcal{A}(i', j')$ can precede any interval in $\mathcal{A}(i, j)$ in the interval partial order.

Let $N(i, j) = |\mathcal{A}(i, j)|$, so $\sum_{i,j} N(i, j) = n$. Then from the above

$$e(P_n) \geq \prod_{i,j} N(i, j)! \geq \prod_{i,j} (N(i, j)/e)^{N(i,j)}$$

and so

$$\log_2 e(P_n) \geq \sum_{i,j} N(i, j) \log_2 N(i, j) - n \log_2 e.$$

Here we sum only over i, j such that $N(i, j) \neq 0$. Fix j , and let $N_j = \sum_i N(i, j)$. Then, since we sum over at most 2^{j-1} values of i , by convexity of $x \log_2 x$

$$\sum_i N(i, j) \log_2 N(i, j) \geq N_j \log_2 (N_j / 2^{j-1}).$$

Hence

$$\log_2 e(P_n) \geq \sum_j N_j \log_2 (N_j / 2^{j-1}) - n \log_2 e. \quad (28)$$

Let the random variable J_k be $j(a, b)$ for the k th random interval. Then $\mathbb{P}(J_k = j) = 2^{-j}$ for each $j = 1, 2, \dots$. To see this, condition on the value x of the first end

point of the interval (which could be the left or the right end point). If $x \in (\frac{i-1}{2^j}, \frac{i}{2^j})$ for some odd i , then $J_k = j$ if and only if the second end point falls in $(\frac{i}{2^j}, \frac{i+1}{2^j})$, which has probability 2^{-j} . On the other hand, if $x \in (\frac{i}{2^j}, \frac{i+1}{2^j})$ for some odd i , then $J_k = j$ if and only if the second end point falls in $(\frac{i-1}{2^j}, \frac{i}{2^j})$, which also has probability 2^{-j} .

By the definition of N_j , it is the number of intervals I_k such that $J_k = j$. Then N_j has the binomial distribution with parameters n and 2^{-j} , which has mean $\mathbb{E}(N_j) = n2^{-j}$ and variance $n2^{-j}(1 - 2^{-j}) \leq n2^{-j}$. Hence, by Chebyshev's inequality,

$$\mathbb{P}(|N_j - n2^{-j}| \geq n2^{-j-1}) \leq \frac{n2^{-j}}{(n2^{-j-1})^2} = \frac{4}{n} 2^j.$$

Let $j_0 := j_0(n) = \log_2 n - h(n)$, where $h(n) \rightarrow \infty$ slowly as $n \rightarrow \infty$: we set $h(n) \sim \log_2 \log_2 n$. Then, by the union bound,

$$\begin{aligned} & \mathbb{P}(|N_j - n2^{-j}| \geq n2^{-j-1} \text{ for some } j = 1, \dots, j_0) \\ & \leq \frac{4}{n} \sum_{j=1}^{j_0} 2^j \leq \frac{8}{n} 2^{j_0} = 2^{-h(n)+3} = o(1). \end{aligned}$$

Also,

$$\mathbb{E}(\sum_{j>j_0} N_j) = n \sum_{j>j_0} 2^{-j} \leq n2^{-j_0} = 2^{h(n)}.$$

Hence by Markov's inequality, whp

$$\sum_{j>j_0} N_j \leq h(n)2^{h(n)} \quad \text{and so} \quad \sum_{j=1}^{j_0} N_j \geq n - h(n)2^{h(n)}.$$

Let B_n be the event that $|N_j - n2^{-j}| \leq n2^{-j-1}$ for each $j = 1, \dots, j_0$ and $\sum_{j=1}^{j_0} N_j \geq n - h(n)2^{h(n)}$. By the above $\mathbb{P}(B_n) = 1 - o(1)$. Condition on B_n . Then by (28)

$$\begin{aligned} \log_2 e(P_n) + n \log_2 e & \geq \sum_{j=1}^{j_0} N_j \log_2(N_j 2^{-j+1}) \\ & \geq \sum_{j=1}^{j_0} N_j \log_2(n2^{-j-1} 2^{-j+1}) \\ & = \sum_{j=1}^{j_0} N_j \log_2 n - \sum_{j=1}^{j_0} N_j \cdot 2j. \end{aligned}$$

But (still conditioned on the event B_n)

$$\sum_{j=1}^{j_0} N_j \cdot 2j \leq (3/2)n \sum_{j=1}^{j_0} 2^{-j} \cdot 2j \leq cn,$$

where $c = (3/2) \sum_{j \geq 1} 2^{-j} \cdot 2j = 6$. Hence

$$\begin{aligned} \log_2 e(P_n) &\geq \left(\sum_{j=1}^{j_0} N_j \right) \log_2 n - (6 + \log_2 e)n \\ &\geq (n - h(n)2^{h(n)}) \log_2 n - (6 + \log_2 e)n \\ &= n \log_2(n/e) - (6 + o(1))n. \end{aligned}$$

We have now shown that the event B_n happens whp, and when it happens we have $e(P_n) \geq n! 2^{-(6+o(1))n}$. Hence (27) holds, and we are done. \square

From Theorem 11, together with the upper bound in part (a) of Theorem 1, we see that the random interval order has about ‘as many linear extensions as possible’. Another model of random partial order is the random k -dimensional partial order R_n . To form this we take k independent random linear orders \leq_i each uniformly distributed on $[n]$, and set $x \preceq y$ in R_n when $x \leq_i y$ for each $i = 1, \dots, k$. Brightwell [3] has given a precise estimate for $e(R_n)$, which shows that $e(R_n)^{1/n} = \Theta(n^{1-\frac{1}{k}})$ whp; thus we are meeting numbers of linear extensions spread more widely through the range of possible values given in Theorem 1.

It would be interesting to compute estimates on the number of linear extensions of other, random or deterministic, models of partial orders; see for example [16] and [29] for two such models.

7 Concluding remarks

We have proved Theorem 1, concerning maximum and minimum possible numbers of linear extensions, and in the process we have given more detailed information. In particular, we saw that if $\delta \rightarrow 1$ but not too quickly then $(f^+(n, \delta)/n!)^{1/n} = \Theta(1 - \delta)$; and if $\delta \rightarrow 0$ but not too quickly then $f^-(n, \delta)^{1/n} = \Theta(\delta^{-1})$. What can we say about other rates of convergence? Can we improve the values $c_i(\delta)$?

Let us focus on δ of the form $1/k$ for an integer $k \geq 2$. From our earlier results, $\text{comp}(\tilde{A}(n, k)) \sim \delta \binom{n}{2}$ and $e(\tilde{A}(n, k)) = n! (1 - \delta)^n 2^{O(\log_2 n)}$. Is $\tilde{A}(n, k)$ extremal for having many linear extensions? Similarly, $\text{comp}(\tilde{C}(n, k)) \sim \delta \binom{n}{2}$ and $e(\tilde{C}(n, k)) = \delta^{-n} 2^{O(\log_2 n)}$; and we may ask if $\tilde{C}(n, k)$ is extremal for having few linear extensions.

Conjecture 12. *Let $\delta = 1/k$ for an integer $k \geq 2$. Then (a)*

$$(f^+(n, \delta)/n!)^{1/n} \rightarrow 1 - \delta \quad \text{as } n \rightarrow \infty;$$

and (b), more boldly, each partial order Q_n on $[n]$ with $\text{comp}(Q_n) \geq \text{comp}(\tilde{A}(n, k))$ satisfies $e(Q_n) \leq e(\tilde{A}(n, k))$.

Conjecture 13. *Let $\delta = 1/k$ for an integer $k \geq 2$. Then (a)*

$$f^-(n, \delta)^{1/n} \rightarrow 1/\delta \quad \text{as } n \rightarrow \infty;$$

and (b), more boldly, each partial order Q_n on $[n]$ with $\text{comp}(Q_n) \leq \text{comp}(\tilde{C}(n, k))$ satisfies $e(Q_n) \geq e(\tilde{C}(n, k))$.

For partial orders Q_n on n points with a given density of edges in the comparability graph G , we have bounded the possible number of linear extensions, which equals $n!$ times the volume of the clique polytope $\mathcal{C}(G)$. It could be interesting to investigate bounds on the volume of the clique polytope $\mathcal{C}(G)$ for graphs G from other classes of perfect graphs (or indeed from more general classes).

We have also shown in Theorem 11 that a random interval partial order P_n whp has $e(P_n) = n! 2^{-\Theta(n)}$. Can this be pinned down more precisely? Is there a constant $c > 0$ such that whp $(e(P_n)/n!)^{1/n} \rightarrow c$? (If there is such a constant c , then $c \geq 2^{-6}$ by the inequality (27).)

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List of responses to reviewers' comments

Colin McDiarmid David Penman Vasileios Iliopoulos

September 27, 2017

Here is a list of our responses to reviewer's comments in the file with name "MPI2.ref.pdf".

- Page 4, line 2 after "Graphs and Entropy" header. Comma added after "edge" in the definition of the clique number of a graph.
- Page 4, lines 12-13 after "Graphs and Entropy" header. "by [27]" is reworded.
- Proof of Lemma 6 is now amended, where we show that

$$e(Q_n) \geq \max\{2^{(1-\sqrt{\delta})(n-1)}, 2^{\frac{1}{2}(1-\delta)n}\}.$$

The second term is the one we use to define $c_3(\delta)$, where we need the factor n not $(n-1)$ in the exponent.

- Statement of Lemma 7 is corrected accordingly.
- Proof of Proposition 9 is rewritten according to referee's comments. In particular, for a poset Q on ground set $[n]$,

$$e(Q) \geq n! 2^{-\text{comp}(Q)},$$

and equality holds if and only if the comparability graph of Q is a (partial) matching.

- Page 17. The paragraph regarding the comparison of the numbers of linear extensions of a random interval order and of a random 2-dimensional order is now removed.