

COMMON PROTOCOL FOR DISTRIBUTED NETWORK FILE SYSTEM

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Summary Paper deals with common protocol for distributed network file system. Focus is on CIFS protocol from Microsoft, the enhanced version of Microsoft Server Message Block (SMB), that is proposed as possible common solution for file sharing among distributed systems. There are new requirements included as well, that are to be implemented due to recent changes in applications and devices and has been addressed in new generation of distributed file system protocols such as AFS2, CODA and WebDAV.

1. INTRODUCTION

A network file system is defined as any computer file system that supports sharing of files, printers over a computer network. Network file services enable computers attached to the network to access files stored on other computers in the same way that they access files on their own disks. That is, the client interface used by programs should not distinguish between local and remote files. It is up to the network file system to locate the files and to arrange for the transport of the data [1], [2].

The computer which provides access to its files and directories is called file server, and the computer using these files and directories is the client. A computer can be client and server at the same time. For communication between servers and clients, the specific network protocols are used. The most popular are the following protocols, see Tab.1:

Tab.1 Network file system protocols

| | Novell NetWare | DoD TCP/IP | Microsoft Network |
|---|----------------|----------------|-------------------|
| 7 | NCP | FTP, HTTP, NFS | SMB, RPC |
| 6 | | | |
| 5 | | | NetBIOS |
| 4 | SPX | TCP/UDP | NetBEUI |
| 3 | IPX | IP | |

- NFS(Network File System) –SunSystems, IETF
- NCP (NetWare Core Protocol) -Novell
- Server Message Block (SMB) –IBM, Microsoft
- File Transfer Protocol (FTP) -IETF
- Hypertext Transfer Protocol (HTTP) -IETF

These protocols were implemented in the different operation systems such as Microsoft Windows (SMB) and UNIX/LINUX (NFS) [3]. Therefore, their clients cannot use the different network filesystems by default, as shown in Fig.1a). There are multi-client solutions applied, with several implemented network file system protocols within given operation system. Fig.1b) shows multiclient configuration, with the protocol stacks for NFS and

SMB (SAMBA implementation of SMB server in LINUX) so users who work with UNIX or Microsoft operation systems, are able to use the both network file systems. Another option is to develop and implement a common network file system that would be supported by all platforms and suppliers. Then, just one protocol stack could be used for access and sharing files via Internet. This option is shown in Fig.1c).

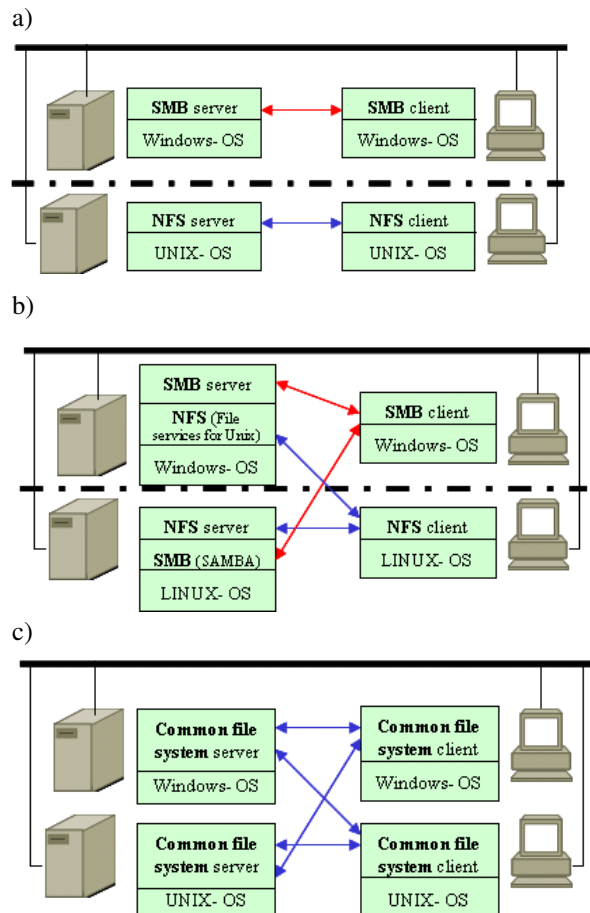


Fig.1 Possible access to network file systems

Due to Microsoft superiority in number of installed nodes in Internet, the introduction of common network file system protocol will require its strong support and acceptance. Microsoft is currently trying to establish its protocols as the official standard of the Internet. It could perspective replace the currently used FTP protocols and file transfers through the HTTP protocol, which represents substantially heterogeneous environment with problematic security. Microsoft has been trying to push Common Internet File System (CIFS), protocol through standardization process since 1996 [4]. CIFS drafts were even provided to IETF so that it could be approved among Internet standard protocols.

2. COMMON INTERNET FILE SYSTEMS

Common Internet File Systems (CIFS) is the enhanced version of Microsoft Server Message Block (SMB) protocol that defines a remote file-access to share data on local disks and network file servers, among distributed systems across intranets and the Internet. In addition to SMB, which was developed on the top of NETBIOS-API, CIFS is able to run over TCP/IP and utilizes the Domain Naming Service (DNS). It is optimized to support communication via Internet with various speed levels. CIFS complements Hypertext Transfer Protocol (HTTP) while providing more sophisticated file sharing than traditional protocols, such as FTP.

The CIFS protocol is not only available for servers and PCs, but also for embedded devices and distributed systems. It had to be scaled down in size and reworked for multitasking features to meet the needs of embedded devices such as real-time, deterministic response times, small memory, and a variety of microprocessor–file system combinations. The using of CIFS for servers and embedded devices is shown in Fig.2. The file sharing is independent on hosting platform and embedded operation system/file system.

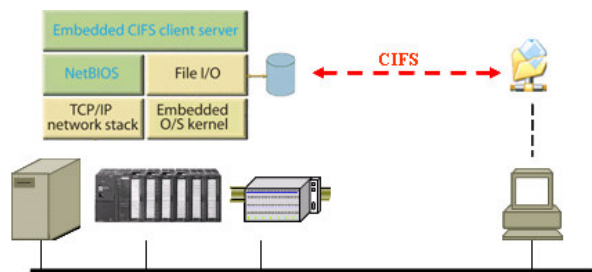


Fig. 2 CIFS protocol in servers and embedded devices

CIFS provides the following functions:

1 File Access

- Support of basic file operations (open, close, read, write and seek).
- Record lock and unlocking for multiple clients access and file sharing and locking.
- ② **Performance and portability**
 - Highly integrated within the operating system.
 - Support for all recent Microsoft platforms and other operation systems such as Unix, VMS, IBM, Macintosh
- ③ **Security**
 - Support of anonymous access(no authorization).
 - Enhanced security and authenticated access
- ④ **Optimization**
 - Improved protocol performance
 - Reduction of overhead in SMB protocol
- ⑤ **Global File Names**
 - No need to mount remote file systems,
 - Access based on globally significant names
- ⑥ **Unicode File Names**

Because CIFS is mainly based on SMB protocol we will explain below its core functions.

3. SMB PROTOCOL

The **SMB (Server Message Block)** protocol was originally developed by the company IBM in co-operation with the company Microsoft for one of the first network operating systems, the **PC Network Program v.1.0**. Today, this protocol is the most frequently used one in the field of peer-to-peer communication of LAN network file servers and clients. The SMB protocol is currently used by file and print servers IBM, Microsoft’s network operating systems (Lan Manager, Windows NT, XP, Vista). The protocol header and basic services are shown in Tab. 2 and Tab. 3.

Tab. 2 Header of SMB protocol

| | | | | | | | | | | | | | | |
|-----|---|---|---|---|----|----|----|----|-----|----|---|----|----|---|
| SID | S | C | R | A | RC | RS | NI | PI | UID | MI | P | PC | BL | B |
| D | | | | H | | | D | D | | D | | d | | |
| 1 | 3 | 1 | 1 | 1 | 2 | 15 | 2 | 2 | 2 | 2 | 1 | 2 | 2 | 1 |

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| MARKING | | MEANING |
|--------------|------------------------------|---------------------------------------------------|
| SID | SMB Identification | Identification of SMB protocol (0xFF) |
| S | Server | Identification of SMB server’s dialect |
| C | Command | Functional code of called SMB service |
| R | Return Class | SMB function return code class |
| AH | Processor AH register | Result of operation in register of AH processor |
| RC | Return Code | Return code of operation |
| RS | Reserve | Reserve for future extension |
| NID | Net Path ID | Identifier allocated to shared means |
| PID | Process ID | Identifier of client’s process |
| UID | User ID | Identifier of user |
| MID | Multiplex ID | Multiplex identifier of client process |
| Prmct | Parameter Count | Number of optional parameters for called function |

| | | |
|-----------|-----------------------|---------------------------------------|
| PC | <i>Parameter Code</i> | Called function parameter code |
| BL | <i>Buffer Length</i> | Length of data part of SMB block |
| B | <i>Buffer</i> | First octet of data part of SMB block |

Tab.3 Example of SMB services

| SMB BLOCK | CODE | SERVICE |
|----------------------------|------|---------------------------------|
| SMB_COM_CREATE_DIRECTORY | 0x00 | Creation of directory |
| SMB_COM_DELETE_DIRECTORY | 0x01 | Deletion of directory |
| SMB_COM_OPEN | 0x02 | Opening of file |
| SMB_COM_CREATE | 0x03 | Creation of directory |
| SMB_COM_CLOSE | 0x04 | Closing opened file |
| SMB_COM_DELETE | 0x06 | Deletion of file |
| SMB_COM_RENAME | 0x07 | Renaming of file |
| SMB_COM_READ | 0x0A | Reading from an opened file |
| SMB_COM_WRITE | 0x0B | Writing into opened file |
| SMB_COM_LOCK_BYTE_RANGE | 0x0C | Locking file byte range |
| SMB_COM_UNLOCK_BYTE_RANGE | 0x0D | Unlocking file byte range |
| SMB_COM_CHECK_DIRECTORY | 0x10 | Searching through directory |
| SMB_COM_SEEK | 0x12 | Transfer in open file |
| SMB_COM_TREE_CONNECT | 0x70 | Connection to shared means |
| SMB_COM_TREE_DISCONNECT | 0x71 | Disconnection from shared means |
| SMB_COM_NEGOTIATE | 0x72 | Negotiation of SMB parameters |
| SMB_COM_SESSION_SETUP_ANDX | 0x73 | Connection to SMB server |
| SMB_COM_LOGOFF_ANDX | 0x74 | Disconnection from SMB server |
| SMB_COM_OPEN_PRINT_FILE | 0xC0 | Opening print file |
| SMB_COM_WRITE_PRINT_FILE | 0xC1 | Writing to print file |
| SMB_COM_CLOSE_PRINT_FILE | 0xC2 | Closing print file |

The SMB block with a client's request and a server's answer are practically identical, they just differ in the data part and the length of the data buffer. An example of communication of a client with an SMB server can be seen in Fig. 3.

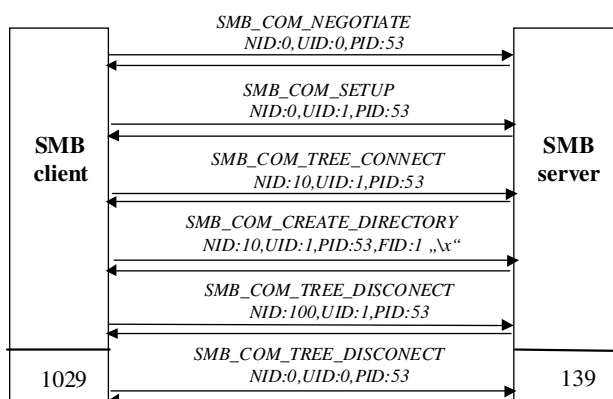


Fig.3 Example of client communication with SMB server

The protocol works on the basis of the client/server model. A server provides to network clients the so-called shared resources, usually shared disks, directories, print files and named pipes. Shared resources are identified through an universal network address called **UNC**

`\\server_name\resource_name.`

A client secures the formulation of requirements of shared server resources. The SMB server analyses requirements sent by a client in the form of an SMB block (packet), it checks an access permission, and based on that it carries out a required operation (file opening, writing into print file, creation of a directory, etc.) The reply, together with the result, is sent to the client in an identical SMB block. In the implementation of SMB servers, two different methods of securing access to shared means of the server are offered through:

- ① control of access at the share level
- ② control of access at the user level

In the first case the access to shared resources is allowed to all network nodes that have a valid password allocated to a resource of a server. During an attempt to access any server resource, a password is required from the client.

After its successful entering, an NID (Network ID) resource identifier is allocated to the client, through which it accesses the resource. Thus a number of passwords with different access levels can be allocated to each shared server resource. In the other case, the client is verified through a name and a password immediately after the connection to the server, and in the event of their correctness, the server allocates to the client its user identifier **UID** (User ID), through which the server derives access rights to individual shared resources. In a multi-user environment, when several processes can communicate simultaneously with the server, a couple of **PID** (*Process ID*) and **MID** (*Multiple ID*) identifiers can serve for their distinguishing.

The client sends the server an **SMB_COM_NEGOTIATE** request, through which the server and the client negotiate connection parameters and a supported protocol version. Subsequently, the client sends to the server a **SMB_COM_SETUP** request, in which there is the user's name and password. If the server works in the user-level mode, it allocates an UID identifier to the client. Through the **SMB_TREE_CONNECT** message the client connects to a shared directory with a TID identifier allocated by the server. Then the creation of a (x) FID subdirectory and the disconnection of the server follows.

```

Ethernet:  <006097ef0da6 -> WDgt1 23dedf>  type: IP<0x800>
Internet:  170.1.1.10 -> 170.1.1.3      hl: 5   ver: 4   tos: 0
  len: 106  id: 0x8502 fragoff: 0   flags: 0x2 ttl: 32  prot: TCP<6>
  xsum: 0x7f7c
TCP:      1029 -> netb-ses<139>  seq: 00ba651f  ack: 001351cf
  win: 7333  hl: 5   xsum: 0xd729  urg: 0   flags: <ACK><PUSH>
SMB: 'TreeConX'  retcode: 0x0  retclass: 0x0  AH: 0x0
  netpath id: 0x0000  pid: 0x871a  user id: 0x0000  mutiplex id: 26
  parmcnt<in word>: 4  smb_buflen: 4864

```

Fig.3: Example of SMB packet

4. CONCLUSION

Traditional network file systems were created in a world where the dominant network paradigm was the LAN/WAN. Lack of standardization caused that common standard has not been selected and approved yet, despite to ISO standard FTAM (File Transfer and Access Protocol) and recent Microsoft activities with its proposed solution -CIFS.

In addition the network file system challenge has become the distributed file system and connections via Internet/WAN, as there is move from self-contained LAN environments to a world of distributed system and even rarely connected devices. This kind of network environment introduces at least three main challenges:

- ❶ Authentication
- ❷ Data Transport
- ❸ Synchronization

Some of these issues could be solved at the application level rather than the file system level. For instance, Rsync, a powerful program that came out of the Samba project provides remote file synchronization over the network. Another alternative to traditional protocol is for example, **AFS** (Andrew File System) and **CODA**. Coda is a distributed file system developed from AFS2. Coda is designed for mobile computing in an occasionally connected environment. Coda clients maintain a persistent cache of copies of files from

the server. Coda checks the network both for the availability of connections between client and server, and for the approximate bandwidth of the connection.

As explained there are still the different solutions in place, such as NFS in Unix and CIFS/SMB in Microsoft world. The future development must consider not only requirement for common file system and protocol, but also to include the new requirements regarding distributed systems, replications, caching, which have not been taken into consideration yet.

Therefore the next development will have to solve the mentioned challenges and to come with a proposal for new solution, which will enable implementation of distributed file system for open and embedded systems as well.

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