

THREE METHODS TO CALCULATE THE PROBABILITY OF RUIN

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ABSTRACT

The first method, essentially due to GOOVAERTS and DE VYLDER, uses the connection between the probability of ruin and the maximal aggregate loss random variable, and the fact that the latter has a compound geometric distribution. For the second method, the claim amount distribution is supposed to be a combination of exponential or translated exponential distributions. Then the probability of ruin can be calculated in a transparent fashion; the main problem is to determine the nontrivial roots of the equation that defines the adjustment coefficient. For the third method one observes that the probability of ruin is related to the stationary distribution of a certain associated process. Thus it can be determined by a single simulation of the latter. For the second and third methods the assumption of only proper (positive) claims is not needed.

KEYWORDS

Probability of ruin; discretization; combination of exponentials; simulation.

1. INTRODUCTION

Traditionally, practitioners have approximated the probability of ruin by the expression e^{-Ru} , where R is the adjustment coefficient (by some authors called insolvency coefficient or Lundberg's constant) and u the initial surplus. From a technical point of view, the need for such an approximation has become less important; thanks to the arrival of efficient computers and even personal computers, the exact probability of ruin can be calculated. This has been demonstrated by several authors, i.a. THORIN and WIKSTAD (1976), SHIU (1988), MEYERS and BEEKMAN (1987), PANJER (1986), and indirectly by STRÖTER (1985).

In this paper we shall present three methods; they have the merit that they can be explained in elementary terms and they can be implemented numerically without any difficulty.

The method of upper and lower bounds (section 2) is a method of numerical analysis and is essentially due to GOOVAERTS and DE VYLDER (1984). The main drawback of this method is that it is limited to the situation where negative claims are excluded.

Method 2 (section 3) is analytical in nature (but it can be understood without extensive knowledge of complex analysis); it generalizes a method that has been proposed by BOHMAN (1971). If the claim amount distribution is a combination

of exponential or translated exponential distributions, the probability of ruin will be of the form as shown in formula (42). The main task is the numerical determination of the nontrivial (possibly) complex roots of the equation that defines the adjustment coefficient.

The probability of ultimate ruin can also be obtained by simulation (section 4), although this seems to be a paradoxical idea at first sight. If the claims are reinsured by (for example) an excess of loss contract, the distribution of retained claims cannot be approximated by a combination of exponential distributions, and Method 2 cannot be applied. Method 3 is generally applicable and does not have the drawbacks of Method 1 (no negative claims) or Method 2 (no reinsurance).

2. METHOD OF UPPER AND LOWER BOUNDS

2.1. Introduction

In this section we shall present a method that leads to the bounds that are due to GOOVAERTS and DE VYLDER (1984); our derivation will be very similar to PANJER's presentation (1986) and along the ideas of TAYLOR (1985).

2.2. The Model

In the following we shall use the model and the notation of continuous time ruin theory as it is explained in the text by BOWERS et al. (1986, sections 12.2, 12.5, 12.6). Thus

$$(1) \quad U(t) = u + ct - S(t)$$

is the insurer's surplus at time $t \geq 0$. Here $u \geq 0$ is the initial surplus, c the rate at which the premiums are received, and $S(t)$ the aggregate claims between 0 and t . It is assumed that $S(t)$ is a compound Poisson process, given by the Poisson parameter λ (claim frequency) and the distribution function $P(x)$ of the individual claim amounts. In this section we assume that $P(0) = 0$ (no negative "claims"); afterwards this assumption will be dropped.

The mean claim size is denoted by p_1 . Of course we assume that c exceeds λp_1 , the expected payment per unit time. The relative security loading θ is defined by the condition that $c = (1 + \theta)\lambda p_1$.

We denote by $\psi(u)$ the probability of "ruin", i.e. that $U(t)$ is negative for some $t > 0$. It is well known that $\psi(0) = 1/(1 + \theta)$. For notational convenience we denote this quantity by q .

The maximal aggregate loss,

$$(2) \quad L = \max_{t \geq 0} \{S(t) - ct\},$$

is a random variable of great interest, since

$$(3) \quad 1 - \psi(u) = \Pr(L \leq u), \quad u \geq 0,$$

i.e. the probability of survival is the distribution function of L , see BOWERS et al. [1986, formula (12.6.2.)]. We can write L as a random sum,

$$(4) \quad L = L_1 + L_2 + \dots + L_N,$$

see BOWERS et al. [1986, formula (12.6.5)]. Here N is the number of record highs of the process $S(t)$ -ct and has a geometric distribution:

$$(5) \quad \Pr(N = n) = (1 - q) q^n, \quad n = 0, 1, 2, \dots$$

The common distribution function of the L_i 's is

$$(6) \quad H(x) = \frac{1}{p_1} \int_0^x [1 - P(y)] dy.$$

Furthermore, the random variables N, L_1, L_2, \dots are independent. Thus L has a compound geometric distribution,

$$(7) \quad \Pr(L \leq u) = \sum_{n=0}^{\infty} (1 - q) q^n H^{*n}(u).$$

Together with (3) this yields the convolution formula for the probability of ruin, which is often attributed to BEEKMAN (1974, section 13.4), but can also be found in DUBOURDIEU (1952, p. 246).

2.3. Derivation of the bounds

Since $H(x)$ is a continuous distribution function, the expression of the right hand side of (7) cannot be evaluated directly. According to PANJER (1986), the idea is to replace $H(x)$ by one or several discrete distributions. Here we prefer to go one step back and use (4) as a starting point.

For the ease of presentation and notation we assume that the interval of discretisation is the unit interval (in fact this means that the monetary unit is identical to the length of the interval of discretisation). Then we introduce two new random variables that are closely related to L :

$$(8) \quad L^l = [L_1] + [L_2] + \dots + [L_N],$$

and

$$(9) \quad L^u = [L_1 + 1] + [L_2 + 1] + \dots + [L_N + 1].$$

Thus the idea is to round the summands in (4) to the next lower integer, which gives (8), or to the next higher integer, which gives (9). Clearly

$$(10) \quad L^l \leq L \leq L^u,$$

which implies that

$$(11) \quad \Pr(L^l > u) \leq \Pr(L > u) \leq \Pr(L^u > u)$$

for all u . Since $\psi(u) = \Pr(L > u)$ is a continuous function for $u > 0$, it follows that

$$(12) \quad \Pr(L^l \geq u) \leq \psi(u) \leq \Pr(L^u > u), \quad u > 0.$$

Since the distributions of L^l and L^u can be calculated recursively, these bounds are of a practical interest.

2.4. Numerical evaluation of the bounds

Let h_k^l denote the probability that a given summand in (8) is equal to k , i.e., that a given summand in (4) is between k and $k+1$. Thus

$$(13) \quad h_k^l = H(k+1) - H(k), \quad k = 0, 1, 2, \dots$$

Let h_k^u denote the corresponding probability for the summands in (9). Thus

$$(14) \quad h_k^u = H(k) - H(k-1), \quad k = 1, 2, 3, \dots$$

Here $H(x)$ is given by formula (6). We want to calculate

$$(15) \quad f_i^l = \Pr(L^l = i), \quad i = 0, 1, 2, \dots$$

and

$$(16) \quad f_i^u = \Pr(L^u = i), \quad i = 0, 1, 2, \dots$$

Then

$$(17) \quad 1 - \sum_{i=0}^{u-1} f_i^l \leq \psi(u) \leq 1 - \sum_{i=0}^u f_i^u, \quad u = 0, 1, \dots,$$

according to (12).

The probabilities (15) and (16) can be calculated recursively by the formulas

$$(18) \quad f_0^l = \frac{1-q}{1-qh_0^l},$$

$$(19) \quad f_i^l = \frac{q}{1-qh_0^l} \sum_{k=1}^i h_k^l f_{i-k}^l, \quad i = 1, 2, \dots$$

and

$$(20) \quad f_0^u = 1-q,$$

$$(21) \quad f_i^u = q \sum_{k=1}^i h_k^u f_{i-k}^u, \quad i = 1, 2, \dots$$

These formulas can be derived as follows. By the law of total probability

$$(22) \Pr(L^l = i) = \Pr(L^l = i | N = 0) \cdot \Pr(N = 0)$$

$$+ \sum_{k=0}^i \Pr(L^l = i | N \geq 1, [L_1] = k) \cdot \Pr(N \geq 1, [L_1] = k)$$

for $i = 0, 1, 2, \dots$. Thus for $i = 0$

$$(23) f_0^l = 1 - q + q f_0^l h_0^l,$$

which gives (18), and for $i = 1, 2, \dots$

$$(24) f_i^l = q \sum_{k=0}^i f_{i-k}^l h_k^l,$$

which yields (19). For the derivation of (20) and (21) we simply replace the superscript “ l ” by “ u ” and observe that $h_0^u = 0$.

FIGURE 1. Combination of exponential densities.

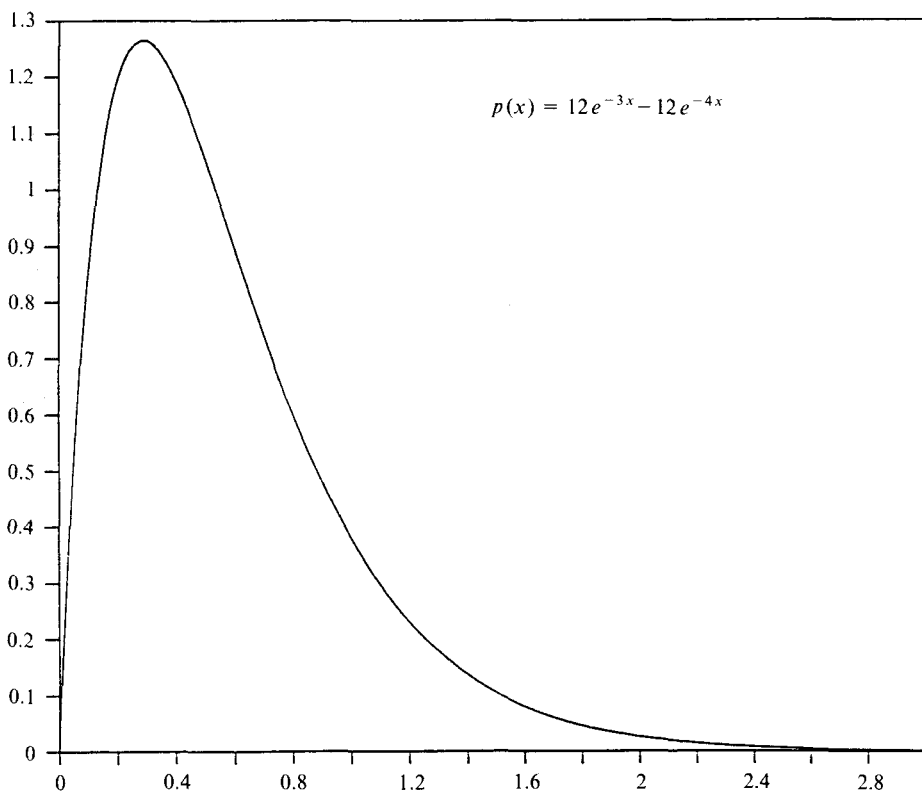


TABLE I
UPPER AND LOWER BOUNDS FOR THE PROBABILITY OF RUIN

u	lower bounds			upper bounds		
.0	.583333	.583333	.583333	.583333	.583333	.583333
.5	.373585	.374626	.375144	.376177	.376692	.377718
1.0	.226752	.228198	.228921	.230367	.231089	.232535
1.5	.136653	.138040	.138736	.140132	.140831	.142234
2.0	.082274	.083424	.084002	.085165	.085750	.086926
2.5	.049528	.050410	.050855	.051753	.052206	.053119
3.0	.029814	.030461	.030788	.031449	.031783	.032459
3.5	.017947	.018406	.018639	.019110	.019350	.019835
4.0	.010804	.011122	.011284	.011613	.011780	.012120
4.5	.006504	.006720	.006831	.007057	.007172	.007406
5.0	.003915	.004061	.004135	.004288	.004366	.004526
5.5	.002357	.002454	.002504	.002606	.002658	.002766
6.0	.001419	.001483	.001516	.001583	.001618	.001690
6.5	.000854	.000896	.000918	.000962	.000985	.001033
7.0	.000514	.000541	.000555	.000585	.000600	.000631
7.5	.000309	.000327	.000336	.000355	.000365	.000386
8.0	.000186	.000198	.000204	.000216	.000222	.000236
8.5	.000112	.000119	.000123	.000131	.000135	.000144
9.0	.000068	.000072	.000075	.000080	.000082	.000088
9.5	.000041	.000044	.000045	.000048	.000050	.000054
10.0	.000024	.000026	.000027	.000029	.000031	.000033

.005
.01
.02

length of interval of discretisation

REMARK: In order to keep the paper self-contained, we gave an elementary proof of the recursive formulas (19) and (21). They could also have been obtained from the observation that the compound geometric distribution is a special case of the family of the compound distributions considered by PANJER (1981). For the reader who is familiar with renewal theory formulas (19) and (21) are easy to understand. The solution of such a renewal equation is a compound geometric distribution. Thus in this particular context we use this relationship backwards, i.e., that a discrete compound geometric distribution can be interpreted as the solution of a renewal equation.

2.5. Illustration

For a numerical illustration we assume a claim amount distribution with a probability density function

$$(25) \quad p(x) = 12(e^{-3x} - e^{-4x}), \quad x > 0,$$

which is shown in figure 1. Written as

$$(26) \quad p(x) = 4(3e^{-3x}) - 3(4e^{-4x}), \quad x > 0,$$

it can be interpreted as a combination of exponential densities, where the coefficients are 4 and -3 . The mean claim size is

$$(27) \quad p_1 = 4 \frac{1}{3} - 3 \frac{1}{4} = \frac{7}{12}.$$

We assume further that $\lambda = c = 1$; thus $\theta = 5/7$.

From (26) we get

$$(28) \quad P(x) = 1 - 4e^{-3x} + 3e^{-4x}, \quad x > 0.$$

Using (6) we obtain

$$(29) \quad H(x) = 1 - \frac{16}{7}e^{-3x} + \frac{9}{7}e^{-4x}, \quad x > 0.$$

The method of upper and lower bounds has been used for discretisation intervals with length .02, .01 and .005. The resulting bounds are displayed in table 1. (The first line, $u = 0$, shows the known exact value, $q = 1/(1+\theta) = 7/12$).

3. COMBINATIONS OF EXPONENTIAL DISTRIBUTIONS AND THEIR TRANSLATIONS

3.1. Improper claims

In section 2 negative claims were excluded, $P(0) = 0$. If the claims can be negative, $0 < P(0) < 1$, the method of section 2 cannot be applied. (The basic formulas (3) and (4) still hold; but the parameter q of the geometric distribution of N is unknown, as well as the common distribution of the L_i 's). We shall present two methods for this more general situation. The first is for a particular family of claim amount distributions; the second will be discussed in section 4.

3.2. A special family of claim amount distributions

We assume that the claim amount distribution is either a *combination* of exponentials, with probability density function of the form

$$(30) \quad p(x) = \sum_{i=1}^n A_i \beta_i e^{-\beta_i x}, \quad x > 0,$$

or else a distribution that is obtained if a combination of exponentials is translated by $\tau > 0$ to the left; then the probability density function is

$$(31) \quad p(x) = \sum_{i=1}^n A_i \beta_i e^{-\beta_i(x+\tau)}, \quad x > -\tau.$$

The β_i 's are positive parameters; for simplicity we assume $0 < \beta_1 < \beta_2 < \dots < \beta_n$. Some of the A_i 's may be negative, but of course

$$(32) \quad A_1 + A_2 + \dots + A_n = 1.$$

If all the A_i 's are positive, (30) is a *mixture* of exponential densities, and (31) is a translation of a *mixture* of exponential densities.

The family of combinations of exponentials is much richer (and therefore more useful) than the family of mixtures; note that for the *latter* the mode is necessarily at 0. In the following we shall treat (30) as a special case ($\tau = 0$) of (31). We shall show how the probability of ruin can be calculated, if the claim amount distribution has such a density.

3.3. A functional equation

By distinguishing according to time (say t) and amount (say x) of the first claim we see that the probability of ruin satisfies the following functional equation:

$$(33) \quad \begin{aligned} \psi(u) = & \int_0^\infty \lambda e^{-\lambda t} \int_{-\tau}^{u+ct} \psi(u+ct-x) dP(x) dt \\ & + \int_0^\infty \lambda e^{-\lambda t} [1 - P(u+ct)] dt \end{aligned}$$

In a sense, this equation has a unique solution:

LEMMA: The functional equation

$$(34) \quad \begin{aligned} g(u) = & \int_0^\infty \lambda e^{-\lambda t} \int_{-\tau}^{u+ct} g(u+ct-x) dP(x) dt \\ & + \int_0^\infty \lambda e^{-\lambda t} [1 - P(u+ct)] dt \end{aligned}$$

has exactly one solution $g(u)$, $u \geq 0$, with the property that $g(\infty) = 0$.

PROOF: Since $\psi(u)$ is a solution, such a solution exists. To show the uniqueness, we assume that there are two different solutions $g_1(u)$ and $g_2(u)$ with $g_1(\infty) = g_2(\infty) = 0$. Then their difference $\delta(u) = g_1(u) - g_2(u)$ satisfies the equation

$$(35) \quad \delta(u) = \int_0^\infty \lambda e^{-\lambda t} \int_{-\tau}^{u+ct} \delta(u+ct-x) dP(x) dt.$$

Now let $m = \max_{u \geq 0} |\delta(u)|$ and let v be a point at which the maximum is

attained: $m = |\delta(v)|$. From (35) it follows that

$$\begin{aligned}
 (36) \quad m &= |\delta(v)| \leq m \int_0^\infty \lambda e^{-\lambda t} \int_{-\tau}^{v+ct} dP(x) dt \\
 &= m \int_0^\infty \lambda e^{-\lambda t} P(v+ct) dt < m,
 \end{aligned}$$

which is impossible. Thus it is not possible that there are two different solutions that vanish at ∞ . Q.E.D.

The lemma provides us with a simple tool to determine the probability of ruin. If we can construct a function $g(u)$ that satisfies (34) and vanishes at ∞ , we know that it is the function $\psi(u)$.

3.4. Construction of a solution

We try to solve (34) by an expression of the form

$$(37) \quad g(u) = \sum_{k=1}^n C_k e^{-r_k u}, \quad u \geq 0.$$

We substitute this and (31) into (34). After some calculations we obtain the condition that

$$\begin{aligned}
 (38) \quad \sum_{k=1}^n C_k e^{-r_k u} &= \sum_{i=1}^n \sum_{k=1}^n \frac{A_i \beta_i C_k \lambda}{(\beta_i - r_k)(\lambda + r_k c)} e^{-r_k(u+\tau)} \\
 &\quad - \sum_{i=1}^n \sum_{k=1}^n \frac{A_i \beta_i C_k \lambda}{(\beta_i - r_k)(\lambda + \beta_i c)} e^{-\beta_i(u+\tau)} \\
 &\quad + \sum_{i=1}^n \frac{A_i \lambda}{\lambda + \beta_i c} e^{-\beta_i(u+\tau)}.
 \end{aligned}$$

Comparison of the coefficients of $C_k e^{-r_k u}$ yields the condition that

$$(39) \quad 1 = \sum_{i=1}^n \frac{A_i \beta_i \lambda}{(\beta_i - r_k)(\lambda + r_k c)} e^{-r_k \tau}.$$

But this means that r_1, r_2, \dots, r_n must be roots of the equation

$$(40) \quad \lambda + cr = \lambda \sum_{i=1}^n A_i \frac{\beta_i}{(\beta_i - r)} e^{-r\tau},$$

which is the equation that defines the adjustment coefficient.

Now we compare the coefficients of $\frac{\lambda A_i}{\lambda + \beta_i c} e^{-\beta_i(u+\tau)}$ in (38). This gives

$$(41) \quad \sum_{k=1}^n \frac{\beta_i}{\beta_i - r_k} C_k = 1.$$

Our preliminary result is the following: if r_1, r_2, \dots, r_n are solutions of (40), and if C_1, C_2, \dots, C_n satisfy (41) for $i = 1, 2, \dots, n$, expression (37) is a solution of (34). But how can we be sure that $g(\infty) = 0$?

3.5. The probability of ruin

The function (37) vanishes at ∞ , if all the r_k 's are positive or (if some are complex) have a positive real part. Thus the question arises, if equation (40) has n such solutions. Two cases have to be distinguished.

- 1) If all the A_i 's are positive (i.e. the claim amount distribution is a mixture or a translated mixture of exponentials) the answer is easy to obtain. The geometric argument of BOWERS et al. (1986, figure 12.7) can also be applied if $\tau > 0$. Thus in this case equation (40) has n positive solutions r_1, r_2, \dots, r_n with $0 < r_1 = R < \beta_1 < r_2 < \beta_2 < \dots < r_n < \beta_n$.
- 2) If some of the A_i 's are negative, the situation is more complicated, since some of the solutions of (40) may be complex. But even here one can show that equation (40) has exactly n solutions that are positive or have a positive real part; the interested reader is referred to DUFRESNE and GERBER (1988, Appendix). It is possible that some of the r_k 's coincide; in the following we exclude this unlikely situation from our discussion.

Thus we have found the following result: The probability of ruin is

$$(42) \quad \psi(u) = \sum_{k=1}^n C_k e^{-r_k u};$$

here r_1, r_2, \dots, r_n are the n roots of equation (40) that are positive or have a positive real part, and the coefficients C_1, C_2, \dots, C_n are the solutions of the system of linear equations

$$(43) \quad \sum_{k=1}^n \frac{\beta_i}{\beta_i - r_k} C_k = 1, \quad i = 1, 2, \dots, n.$$

(Such a system has a unique solution, which we shall determine in the following paragraph).

REMARK: CRAMÉR (1955, section 5.14) derived (42) for the case 1) above; however he did not give any explicit formulas for the C_k 's.

3.6. The coefficients

To determine the coefficients we consider the rational function

$$(44) \quad Q(x) = \frac{\sum_{j=1}^n \frac{1}{\beta_j} \prod_{k=1}^n (\beta_j - r_k) \prod_{\substack{i=1 \\ i \neq j}}^n \left(\frac{x - \beta_i}{\beta_j - \beta_i} \right)}{\prod_{k=1}^n (x - r_k)} .$$

We note that $Q(\beta_j) = 1/\beta_j, j = 1, 2, \dots, n$. By the principle of partial fractions there are unique coefficients D_1, D_2, \dots, D_n such that

$$(45) \quad \sum_{k=1}^n \frac{D_k}{x - r_k} = Q(x) .$$

They can be determined by the condition that the expression on the two sides must be equal for n different values of x , for example for $x = \beta_j (j = 1, 2, \dots, n)$:

$$(46) \quad \sum_{k=1}^n \frac{D_k}{\beta_j - r_k} = Q(\beta_j) = \frac{1}{\beta_j} .$$

But this system is equivalent to (43), and we find as a first result that

$$(47) \quad C_k = D_k, \quad k = 1, 2, \dots, n .$$

Finally we determine these coefficients as follows. First we multiply (45) by $(x - r_h)$ to get

$$(48) \quad \sum_{k=1}^n \frac{x - r_h}{x - r_k} D_k = \frac{\sum_{j=1}^n \frac{1}{\beta_j} \prod_{k=1}^n (\beta_j - r_k) \prod_{\substack{i=1 \\ i \neq j}}^n \left(\frac{x - \beta_i}{\beta_j - \beta_i} \right)}{\prod_{\substack{k=1 \\ k \neq h}}^n (x - r_k)}$$

For $x = r_h$ this gives

$$(49) \quad D_h = \frac{\sum_{j=1}^n \frac{1}{\beta_j} \prod_{k=1}^n (\beta_j - r_k) \prod_{\substack{i=1 \\ i \neq j}}^n \left(\frac{r_h - \beta_i}{\beta_j - \beta_i} \right)}{\prod_{\substack{k=1 \\ k \neq h}}^n (r_h - r_k)} .$$

In view of (47) this is the desired result.

REMARK: The case where $c = 0$ is of some unexpected interest: Then $\psi(u)$ can be interpreted as the probability of ruin in a discrete time model in which the annual premium is τ and the probability density function of the annual aggregate claims is given by (30).

3.7. The case $\tau = 0$

In the special case where the claim amount distribution is a combination of exponentials (without translation), there is an alternative and somewhat simpler expression for the coefficients.

First of all, there is an alternative way to get the r_k 's. We replace

$$\frac{\beta_i}{\beta_i - r} \quad \text{by} \quad 1 + \frac{r}{\beta_i - r}$$

in (40), subtract λ on both sides, and divide the resulting equation by r to get the condition that

$$(50) \quad c = \lambda \sum_{i=1}^n \frac{A_i}{\beta_i - r}.$$

Then r_1, r_2, \dots, r_n are the roots of this equation.

Now consider the rational function

$$(51) \quad \tilde{Q}(x) = \frac{1}{x} \frac{\lambda \sum_{i=1}^n \frac{A_i}{\beta_i - x} - \lambda p_1}{\lambda \sum_{i=1}^n \frac{A_i}{\beta_i - x} - c},$$

where $p_1 = \sum_{i=1}^n A_i/\beta_i$. Since $\tilde{Q}(x)$ has the same poles as $Q(x)$ and since

$$\tilde{Q}(\beta_j) = \frac{1}{\beta_j} = Q(\beta_j), \quad j = 1, 2, \dots, n,$$

we conclude that $\tilde{Q}(x)$ is in fact identical to the function $Q(x)$. Thus we gather from (45) that

$$(52) \quad \sum_{k=1}^n \frac{D_k}{x - r_k} = \tilde{Q}(x).$$

Multiplication by $x - r_h$ gives

$$(53) \quad \sum_{k=1}^n \frac{x-r_h}{x-r_k} D_k = \frac{1}{x} \frac{\lambda \sum_{i=1}^n \frac{A_i}{\beta_i-x} - \lambda p_1}{\left(\lambda \sum_{i=1}^n \frac{A_i}{\beta_i-x} - c \right) / (x-r_h)}.$$

Now we let $x \rightarrow r_h$. The the denominator of the expression on the right will tend to the derivative of the function

$$\lambda \sum_{i=1}^n \frac{A_i}{\beta_i-x} - c$$

at $x = r_h$. Thus in the limit we obtain from (53) the formula

$$(54) \quad D_h = \frac{1}{r_h} \frac{\sum_{i=1}^n \frac{A_i}{\beta_i-r_h} - p_1}{\sum_{i=1}^n \frac{A_i}{(\beta_i-r_h)^2}},$$

which (if $\tau = 0$) can be used instead of (49).

REMARK: A formula that is essentially identical to (54) has been given by BOHMAN (1971) for the case of mixtures of exponentials. A result similar to (54) can be found in CRAMÉR (1955, section 5.14), see also SHIU (1984). The importance of combinations of exponentials has been recognized by THORIN and WIKSTAD (1977) and GERBER, GOOVAERTS and KAAS (1987).

3.8. Illustration

We consider two examples:

a. For the first example we use the combination of exponential densities of section 2.5, where $\tau = 0$ and

$$(55) \quad n = 2, \quad \beta_1 = 3, \quad \beta_2 = 4, \quad A_1 = 4, \quad A_2 = -3,$$

and assume, as before, $\lambda = c = 1$. The solutions of (50) are

$$r_1 = 1, \quad r_2 = 5.$$

Then we get from (54) and (47)

$$C_1 = 5/8, \quad C_2 = -1/24.$$

Thus the probability of ruin is given by the expression

$$\psi(u) = \frac{5}{8} e^{-u} - \frac{1}{24} e^{-5u}.$$

TABLE 2
THE PROBABILITY OF RUIN FOR COMBINATIONS OF EXPONENTIAL DISTRIBUTIONS AND THEIR
TRANSLATIONS

u	a.	b.
	$\tau = 0$	$\tau = 0.1$
.0	0.583333	0.584204
.5	.375661	.365203
1.0	.229644	.219122
1.5	.139433	.130687
2.0	.084583	.077873
2.5	.051303	.046396
3.0	.031117	.027642
3.5	.018873	.016468
4.0	.011447	.009812
4.5	.006943	.005845
5.0	.004211	.003483
5.5	.002554	.002075
6.0	.001549	.001236
6.5	.000940	.000736
7.0	.000570	.000439
7.5	.000346	.000261
8.0	.000210	.000156
8.5	.000127	.000093
9.0	.000077	.000055
9.5	.000047	.000033
10.0	.000028	.000020

The numerical values are shown in table 2, column *a.*, and confirm our findings of table 1.

b. For the second example we translate the probability density (26) by 0.1 to the left. Thus the claim amount distribution is now given by $\tau = 0.1$ and (55). The mean size is now

$$p_1 = 7/12 - 0.1 = 29/60.$$

We assume the same relative security loading as in the first example, $\theta = 5/7$. Thus, if for example $c = 1$, we assume $\lambda = 35/29$. From (40) we find

$$r_1 = 1.035774, \quad r_2 = 4.817225.$$

From (49) and (47) we obtain

$$C_1 = 0.618102, \quad C_2 = -0.033898.$$

Then the probability of ruin is given by the expression

$$\psi(u) = C_1 e^{-r_1 u} + C_2 e^{-r_2 u}.$$

The numerical values are shown in table 2, column *b.*

4. SIMULATION

4.1. Introduction

If the probability of a certain event is to be found by simulation, the usual proceeding is to repeat the stochastic experiment a number of times and to see each time, if the event does or does not take place. Then the observed empirical frequency is used to estimate the probability of the event.

Obviously, this method of brute force is not very practical, if we are to find the probability of *ultimate* ruin. Nevertheless the probability of ultimate ruin can be obtained by simulation, if the following two facts are kept in mind: Firstly, the probability of ruin is equivalent to the stationary distribution of a certain associated process, and secondly, this stationary distribution can be obtained by pathwise simulation.

4.2. Duality

Let us introduce

$$(56) \quad L(t) = S(t) - ct, \quad t \geq 0,$$

the aggregate loss at time t , and

$$(57) \quad M(t) = \max_{0 \leq z \leq t} L(z), \quad t \geq 0,$$

the maximal aggregate loss in the interval from 0 to t . Then

$$(58) \quad 1 - \psi(u, t) = \Pr(M(t) \leq u),$$

i.e., the probability of survival to time t (a function of the initial surplus) is the distribution function of $M(t)$. Note that (58) generalizes (3).

We shall also consider

$$(59) \quad W(t) = L(t) - \min_{0 \leq z \leq t} L(z).$$

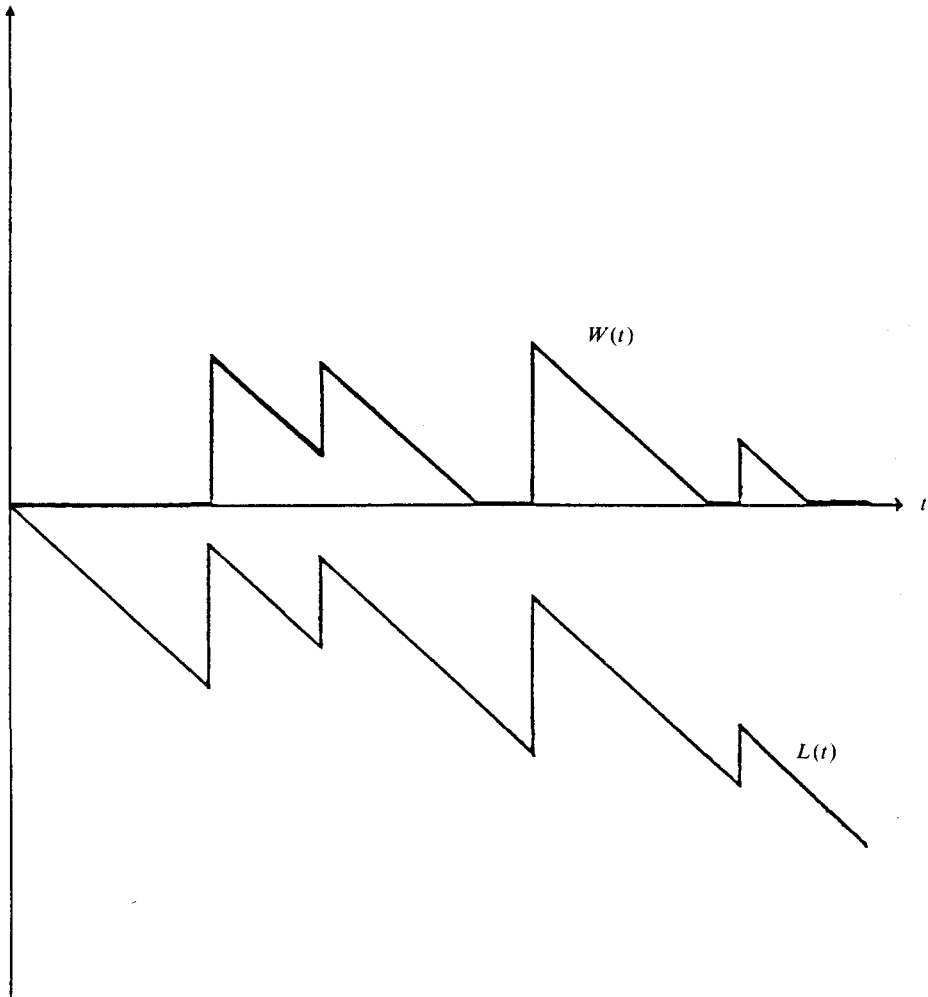
The process $\{W(t)\}$ is obtained from the process $\{L(t)\}$ by introduction of a retaining barrier at 0. This is illustrated in figure 2.

Let

$$(60) \quad F(x, t) = \Pr(W(t) \leq x)$$

denote the distribution function of $W(t)$. Now we rewrite (59) as follows

$$(61) \quad W(t) = \max_{0 \leq z \leq t} \{L(t) - L(z)\}.$$

FIGURE 2. Construction of the process $\{W(t)\}$.

Since the process $\{L(t)\}$ has stationary and independent increments, it follows that $W(t)$ and $M(t)$ have the same distribution. Therefore

$$(62) \quad 1 - \psi(u, t) = F(u, t).$$

Let

$$(63) \quad F(x) = \lim_{t \rightarrow \infty} F(x, t)$$

denote the stationary distribution of the process $\{W(t)\}$. Then it follows from (62) that

$$(64) \quad 1 - \psi(u) = F(u).$$

It remains to show how $F(u)$ can be obtained efficiently.

REMARKS: The fact that $W(t)$ and $M(t)$ are identically distributed is well known, see for example FELLER (1966, VI.9) or SEAL (1972).

4.3. Determination of the stationary distribution

The distribution $F(x)$ can be obtained by pathwise simulation in the following fashion: For a particular value of x let $D(x, t)$ denote the total duration of time that the process $\{W(z)\}$ is below the level x before time t . Then one can show that

$$(65) \quad \frac{D(x, t)}{t} \rightarrow F(x) \quad \text{for } t \rightarrow \infty.$$

This is essentially an application of the Strong Law of Large Numbers and can be found in HOEL, PORT and STONE (1972, section 2.3). From (64) and (65) we see that the probability of ruin can indeed be obtained by simulation, where the process $\{W(t)\}$ has to be simulated only once.

4.4. Practical implementation

We simulate T_1, T_2, \dots , the times when the claims occur, and X_1, X_2, \dots , the corresponding claim amounts. Instead of keeping track of $D(x, t)$, it is easier to keep track of $D_n(x)$, the duration of the time that the process $\{W(t)\}$ is below the level x before the time of the n -th claim. Then it follows from (65) and the fact that $T_n \rightarrow \infty$ for $n \rightarrow \infty$ that

$$(66) \quad \frac{D_n(x)}{T_n} \rightarrow F(x) \quad \text{for } n \rightarrow \infty.$$

Thus if n is sufficiently large, $1 - \psi(u) = F(u)$ is estimated by the value of $D_n(u)/T_n$.

For a given value of x , $D_n(x)$ can be computed recursively as follows. First $D_1(x) = T_1$. Then for $n = 1, 2, \dots$

$$(67) \quad D_{n+1}(x) = \begin{cases} D_n(x) + T_{n+1} - T_n & \text{if } W_n + X_n \leq x \\ D_n(x) + \left(T_{n+1} - T_n - \frac{W_n + X_n - x}{c} \right)_+ & \text{if } W_n + X_n > x. \end{cases}$$

Here W_n denotes the value of $W(t)$ immediately before the n -th claim. Thus $(W_n + X_n)_+$ is the value of $W(t)$ just after the n -th claim. Therefore we can calculate the W_n 's recursively according to the formula

$$(68) \quad W_{n+1} = [(W_n + X_n)_+ - c(T_{n+1} - T_n)]_+,$$

with starting value $W_1 = 0$.

4.5. Illustration

We consider the two examples of section 3.8. Each time the simulation has been carried out through 1 million claims, and the results are shown in table 3a and table 3b. We note that in each column the estimators for the probability of ruin approach the exact value (taken from table 2) as simulation progresses. We note that the convergence is quite satisfactory, perhaps with the exception of the "large" values of u , where the probability of ruin is small. In any case, if one is not satisfied with the convergence, the simulation can be continued to obtain more precise results. Of course for this it would be advisable to do the simulation on a mainframe computer! (The simulations above have been carried out on a PC).

TABLE 3a
PROBABILITY OF RUIN BY SIMULATION

	u = 0	1	2	3	4	5	6	7
n								
500	0.588746	0.254429	0.127094	0.076168	0.028443	0.011798	0.002875	0.000119
100000	0.581629	0.221634	0.077008	0.027749	0.009953	0.003666	0.001398	0.000596
150000	0.581775	0.222313	0.077642	0.027530	0.009794	0.003529	0.001290	0.000486
200000	0.582200	0.223009	0.078331	0.027859	0.009804	0.003509	0.001262	0.000454
250000	0.583733	0.224995	0.079864	0.028884	0.010299	0.003720	0.001387	0.000503
300000	0.583160	0.224523	0.079211	0.028330	0.009968	0.003550	0.001305	0.000468
350000	0.582982	0.224037	0.079002	0.028149	0.009757	0.003403	0.001233	0.000450
400000	0.582879	0.223622	0.078979	0.028174	0.009790	0.003423	0.001225	0.000443
450000	0.582922	0.223279	0.078748	0.028017	0.009683	0.003379	0.001186	0.000421
500000	0.583052	0.223807	0.079288	0.028304	0.009855	0.003432	0.001200	0.000429
550000	0.583466	0.224405	0.079577	0.028410	0.010002	0.003567	0.001309	0.000494
600000	0.583861	0.224748	0.079699	0.028368	0.009971	0.003535	0.001293	0.000479
650000	0.583791	0.224570	0.079593	0.028316	0.009929	0.003488	0.001261	0.000457
700000	0.583743	0.224869	0.079917	0.028544	0.010073	0.003550	0.001261	0.000445
750000	0.583600	0.224708	0.079819	0.028493	0.010052	0.003563	0.001284	0.000465
800000	0.583540	0.224802	0.079999	0.028697	0.010231	0.003646	0.001303	0.000460
850000	0.583609	0.224789	0.079971	0.028726	0.010262	0.003675	0.001334	0.000486
900000	0.583758	0.225065	0.080087	0.028787	0.010314	0.003707	0.001349	0.000488
950000	0.583820	0.225186	0.080183	0.028808	0.010274	0.003668	0.001321	0.000481
1000000	0.583951	0.225311	0.080232	0.028808	0.010221	0.003613	0.001301	0.000478
exact	0.583333	0.229644	0.084583	0.031117	0.011447	0.004211	0.001549	0.000570

TABLE 3b
PROBABILITY OF RUIN BY SIMULATION

	u = 0	1	2	3	4	5	6	7
n								
500	0.593045	0.243940	0.091826	0.034948	0.008195	0.000000	0.000000	0.000000
100000	0.584593	0.214373	0.074654	0.027335	0.010103	0.003662	0.001404	0.000710
150000	0.583346	0.212020	0.072741	0.026267	0.009432	0.003226	0.001140	0.000527
200000	0.582293	0.211215	0.072141	0.025765	0.009309	0.003361	0.001177	0.000509
250000	0.582775	0.212572	0.073010	0.026148	0.009469	0.003316	0.001122	0.000453
300000	0.583533	0.213759	0.073920	0.026537	0.009531	0.003335	0.001153	0.000439
350000	0.584417	0.214489	0.074245	0.026771	0.009687	0.003398	0.001142	0.000403
400000	0.584564	0.214564	0.074286	0.026818	0.009652	0.003421	0.001131	0.000376
450000	0.584437	0.214486	0.074457	0.027060	0.009865	0.003507	0.001167	0.000395
500000	0.584456	0.214311	0.074225	0.026841	0.009749	0.003449	0.001122	0.000361
550000	0.584328	0.213919	0.073795	0.026510	0.009533	0.003350	0.001072	0.000337
600000	0.584500	0.214244	0.074049	0.026507	0.009519	0.003382	0.001111	0.000372
650000	0.584395	0.214239	0.074237	0.026734	0.009640	0.003467	0.001162	0.000393
700000	0.584397	0.214095	0.074108	0.026711	0.009612	0.003425	0.001144	0.000386
750000	0.584731	0.214306	0.074305	0.026732	0.009607	0.003393	0.001118	0.000373
800000	0.584664	0.214142	0.074188	0.026633	0.009533	0.003340	0.001084	0.000355
850000	0.584454	0.213887	0.074059	0.026543	0.009460	0.003308	0.001068	0.000353
900000	0.584524	0.213940	0.074002	0.026448	0.009437	0.003318	0.001089	0.000375
950000	0.584671	0.214196	0.074177	0.026499	0.009438	0.003325	0.001124	0.000405
1000000	0.584531	0.214211	0.074280	0.026556	0.009482	0.003337	0.001123	0.000400
exact	0.584204	0.219122	0.077873	0.027642	0.009812	0.003483	0.001236	0.000439

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