DESIGN AND IMPLEMENTATION OF NOVEL HIGH PERFORMANCE DOMINO LOGIC

A thesis submitted in partial fulfillment of the requirements for the award of the degree of

Doctor of Philosophy

in

VLSI Design and Embedded Systems

by

SRINIVASA V S SARMA D

Roll No: 510EC102

Under the Guidance of

Prof. KAMALAKANTA MAHAPATRA



Electronics and Communication Engineering Department

National Institute of Technology

Rourkela-769008

Odisha

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CERTIFICATE

This is to certify that the thesis report entitled "DESIGN AND

IMPLEMENTATION OF NOVEL HIGH PERFORMANCE DOMINO

LOGIC" submitted by Srinivasa V S Sarma D, Roll No: 510EC102, in

partial fulfillment of the requirements for the award of the degree of Doctor of

Philosophy with specialization in "VLSI Design and Embedded Systems"

in Electronics and Communication Engineering at the National Institute

of Technology, Rourkela is an authentic work under my supervision and

guidance.

To the best of my knowledge, the matter embodied in the thesis has not been

submitted to any other University / Institute for the award of any Degree or

Diploma.

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Dedicated to

My parents

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ABSTRACT

This dissertation presents design and implementation of novel high performance domino logic techniques with increased noise robustness and reduced leakages. The speed and overhead area became the primary parameters of choice for fabrication industry that led to invention of clocked logic styles named as Dynamic logic and Domino logic families. Most importantly, power consumption, noise immunity, speed of operation, area and cost are the predominant parameters for designing any kind of digital logic circuit technique with effective trade-off amongst these parameters depending on the situation and application of design.

Because of its high speed and low overhead area domino logic became process of choice for designing of high speed application circuits. The concerning issues are large power consumption and high sensitivity towards noise. Hence, there is a need for designing new domino methodology to meet the requirements by overcoming above mentioned drawbacks which led to ample opportunities for diversified research in this field. Therefore, the outcome of research must be able to handle the primary design parameters efficiently. Besides this, the designed circuit must exhibit high degree of robustness towards noise.

In this thesis, few domino logic circuit techniques are proposed to deal with noise and sub-threshold leakages. Effect of signal integrity issues on domino logic techniques is studied. Furthermore, having been subjected to process corner analysis and noise analysis, the overall performance of proposed domino techniques is found to be enhanced despite a few limitations that are mentioned in this work. Besides this, lector based domino and dynamic node stabilized techniques are also proposed and are investigated thoroughly. Simulations show that proposed circuits are showing superior performance. In addition to this, domino based Schmitt triggers with various hysteresis phenomena are designed and simulated. Pre-layout and post-layout simulation results are compared for proposed Schmitt trigger. Simulations reveal that proposed Schmitt trigger techniques are more noise tolerant than CMOS counterparts. Moreover, a test chip for domino based Schmitt trigger is done in UMC 180 nm technology for fabrication.

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ABBREVIATIONS

MOSFET Metal Oxide Semiconductor Field Effect Transistor

CMOS Complementary Metal Oxide Semiconductor

NMOS N-channel Metal Oxide Semiconductor

PMOS P-Channel Metal Oxide Semiconductor

GND Ground

VLSI Very Large Scale Integration

DCVSL Differential Cascode Voltage Swing Logic

PTL Pass Transistor Logic

DPTL Differential Pass Transistor Logic

TTL Transistor-Transistor Logic

CML Current Mode Logic

CLK Clock

PDP Power-Deley-Product

UNG Unity Noise Gain

ANTE Average Noise Threshold Energy

STHD Skew Tolerant High Speed Domino

FEL Front-End of Line

BEL Back End of Line

PUD Pull-Up Device

PDN Pull-Down Device

LCT Leakage Control Transistor

NN Normal-Normal

FF Fast-Fast

SS Slow-Slow

FS Fast-Slow

SF Slow-Fast

ST Schmitt Trigger

VTC Voltage Transfer Characteristics

NM Noise Margin

DRC Design Rule Check

LVS Layout Versus Schematic

RCX Parasitic Extraction

Op-Amp Operational Amplifier

CHAPTER 1 INTRODUCTION

1.1 Introduction

Complementary Metal Oxide Semiconductor for wide variety of applications in VLSI field became the logic style of choice for the design of digital semiconductor domain because of its low power dissipation and ease of design with increased robustness [1-6]. This became the major advantage of CMOS logic over the other available manufacturing processes then, which suffered from flow of leakage currents or constant dissipation of bias currents. The rapid development of VLSI technology made a remarkable shift in the fabrication industry with its emerging qualities like high speed, low power, increased robustness and low area overhead. Scaling brought impeccable change in the recent trends. The evolution of various logic families like pseudo NMOS, DCVSL, PTL, and DPTL changed the ongoing market trend in manufacturing field.

Then speed and overhead area became the primary parameters of choice for fabrication industry that led to invention of clocked logic styles named as Dynamic logic and Domino logic families. Power consumption, noise immunity, speed of operation, area and cost are the predominant parameters that have to be taken into consideration before designing any kind of digital logic circuit technique. There may be a requirement for the effective trade-off between any two parameters depending on the situation and application of design. Sometimes, the design techniques might not meet all the mentioned requirements in their application, but still an optimization may be followed in order to proceed further in research areas.

Because of its high speed and low overhead area domino logic became process of choice for many digital circuits. The concerning issues are large power consumption and high sensitivity towards noise. Hence, there is a need for designing new domino methodology or improving existing techniques to meet the requirements by overcoming the drawbacks which led to ample opportunities for diversified research in this field. Therefore, the outcome of research must be able to handle the primary design parameters efficiently. Besides this, the designed circuit must exhibit high degree of robustness towards noise.

In this thesis, few domino logic circuit techniques are proposed to deal with noise and sub-threshold leakages. Furthermore, few existing circuits have also been modified to improve response. Proposed logic techniques are effective in increasing the immunity of system towards noise and sub-threshold leakage issues. This logic is further modified using various types of conditional keepers to design an energy-efficient circuit. Schmitt trigger, using proposed technique, is designed and investigated for its operation. A test chip for domino based Schmitt trigger is done in UMC 180 nm technology.

1.2 History

The revolution in integration industry and IC design made an impeccable shift in VLSI industry in the 1960s. According to Moore's law, the number of transistors that can be accommodated or integrated on a single die would exponentially grow with time [1]. Figure 1.1 shows Moore's prediction. It is observed that the complexity of integration doubles approximately every year. In the early 1970s, the microprocessor has begun to grow up in integration complexity and high performance.

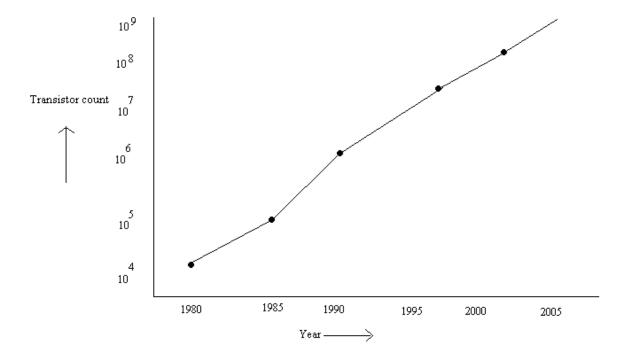


Fig. 1.1 Moore's law

1.3 Motivation

Besides its classical advantage of high speed operation, Domino logic family suffers from low noise sensitivity and large power consumption [5-10]. Significant research has been going in this field in order to stabilize this domino with reference to designing parameters. Several techniques have been proposed to overcome the mentioned drawbacks and most of them, however, partially improve the design parameters in various applications.

1.4 Objectives of the research work

The main aim is to design and implement domino logic circuit techniques to deal with noise issues and enhance the primary design parameters like power, speed, leakages, noise, area and cost. [1, 2]

The main objectives of this thesis are

- (1) Study of existing domino logic circuit techniques,
- (2) Simulating the benchmark circuits for analyzing the overall functionality,
- (3) Improving the existing methodologies by modifying the topologies or if possible introducing novel techniques,
- (4) Making the comparison of improved circuits with existing ones and
- (5) Designing of application based circuit (Schmitt Trigger) techniques based on improved methods.

1.5 Thesis structure and over all contribution

Chapter 1: Introduction

We present a generalized introduction about the broad area of research from the very basic level. In this chapter we also present the organization of the thesis and chapter wise contribution.

Chapter 2: Overview of logic styles and related work

Here the research area is primarily focused on present working environment from a broader angle to this field. This chapter gives overview of standard logic styles in brief and introduces the dynamic logic followed by domino logic circuits with description. The research area is primarily focused on present working environment-Domino logic from a broader angle. Description of various circuit styles along with their advantages and disadvantages is illustrated with corresponding figures. In addition to this, the functioning

of domino logic with the encroachment of down scaling of process technology is investigated with analysis. Technique, which uses a PMOS keeper at dynamic node, to alleviate inevitable charge lost is reviewed and corresponding simulation result is presented in Table 2.2. A brief review on domino logic is conducted and issues related to domino logic are brought out that facilitated us proceeding to the next chapter.

Chapter 3: Novel Domino logic topologies

This chapter gives general introduction to domino logic family with detailed literature survey. Standard benchmark domino logic circuit schemes followed by the analysis of their functionality with simulation results are thoroughly investigated. In addition to this, novel domino logic circuit techniques are proposed and are analyzed in detail with equivalent circuit diagrams in all operating phases along with simulation results. Moreover, analysis of benchmark circuits and proposed techniques includes variation on all the design parameters at different ambient conditions. Furthermore, noise analysis is carried out which includes the need for robustness, various noise metric parameters for measuring noise immunity or robustness of domino circuits such as UNG, ANTE along with the method of calculations, various sources of noise in domino logic circuits and their role on operating region. Besides this, description of process corner analysis and various corners involved in it along with their significant role on the overall functionality of the designed domino logic circuit is presented. Also the consequences of subjecting the device to the extreme corners with the boundary limitations are discussed. Result section shows the calculations and comparisons of all the parameters of standard benchmark circuits and proposed domino techniques. The primary design parameters such dynamic power, leakage or static power, total power, PDP (power-delay-product), UNG and ANTE for wide fan-in circuits of existing and proposed techniques are measured. The comparisons along with graphical analysis and tabulations are made and discussed the functionality with pros and cons.

Chapter 4: Signal integrity issues and modified circuit techniques

This chapter gives general introduction to need for power reduction and leakage minimization. It discusses signal integrity issues in detail with simulations. A review on prior works related to leakage power reduction schemes and the lector technique is

presented. Also, modified lector domino scheme and dynamic node stabilizing technique are proposed in this chapter. All the simulations are done at CMOS 90 nm process technology with 1 V power supply. Process corner and noise analyses are carried out for proposed schemes. Also, the primary design parameters such dynamic power, leakage or static power, total power, PDP (power-delay-product), UNG and ANTE for various fan-in circuits of existing and proposed techniques are measured and comparisons are made and conclusions are drawn in result section.

Chapter 5: Design of various domino based Schmitt trigger circuits

This chapter briefly explains conventional Schmitt triggers using op-amp and CMOS logic. Next, it demonstrates proposed domino Schmitt trigger circuits along with analysis. A novel domino logic based noise tolerant Schmitt trigger circuit is designed and simulated. The simulation results with discussion are presented from which concluding remarks are made.

Chapter 6: Chip tape-out

This chapter explains complete VLSI design flow which includes the process of chip tape-out. It compares overall functionality of domino based Schmitt trigger in 180 nm and 90 nm technologies. The difference between pre-layout and post-layout simulation is investigated. The chip tape-out of proposed domino Schmitt trigger circuit is done in CMOS 180 nm process technology.

Chapter 7: Conclusion and future work

The total work is concluded in this chapter and it also discussed future scope briefly for further improvements based on this work.

1.6 Conclusion

Therefore, the need for high speed, low power, less area and more noise tolerant qualities is briefly discussed. As technology is growing with rapid improvements in VLSI fields with diversified applications, there is need to achieve these target applications with up-coming research outcomes. For that purpose, updating these present technological library files is mandatory so as to coagulate the quality of research methodologies. This chapter is presenting overall idea behind this work from very basic level. It briefed up the generalized introduction, related history behind this work, motivation and main objectives of this work. Thesis structure along with overall chapter wise contribution is also given.

CHAPTER 2 OVERVIEW OF LOGIC STYLES AND RELATED WORK

2.1 CMOS AND NMOS

FrankWanlass was the first person who in the year 1963 proposed Complementary Metal Oxide Semiconductor which for wide variety of applications became the logic style of choice for the design of digital semiconductor domain because of its low power dissipation (almost no power) when the gate inputs are not altered [6]. This is implied from the fact that CMOS structure consists of both PMOS field effect transistors which can expeditiously drive a strong one and NMOS field effect transistors which can adequately drive a strong zero at the output node. Therefore this peculiar combination of complementary transistors allows CMOS logic gate circuits which are to be implemented in such a way that the output node voltage level is always connected to either supply voltage rail or ground rail but not both simultaneously, which there by implies that as long as the logic inputs of CMOS circuit are not altered, there is no power dissipation by the circuit. This became the major advantage of CMOS logic over the other available manufacturing processes then, which suffered from flow of leakage currents or constant dissipation of bias currents.

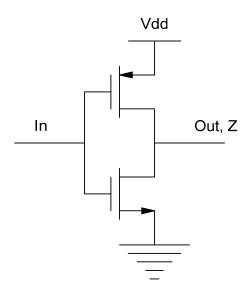


Fig. 2.1 Static CMOS logic Inverter

Fig. 2.1 shows the schematic representation of a static CMOS NAND gate. The logic gate has two inputs A and B with an output Z. A logic high voltage level at both the input nodes A and B turns on the NMOS transistors in the pull down device MN1 and MN2, while turning off the PMOS transistors in the pull up device MP1 and MP2 which results the output node Z to be a logic low voltage level. When either input A or B is off, however, the path for discharge of output node to the ground rail is broken, with an existing path to the power supply rail (Vdd) being established. This causes the output node Z to reach a logic high voltage level. While the NAND gate represents a simple logic function, it shows how contention current between the power supply rail and ground rail can effectively be avoided in static CMOS circuits. This lack of contention current implies that when the inputs to the CMOS circuit are not changing, often referred as a ideal state of operation or standby mode, almost no power dissipation occurs, except for a small but finite amount of leakage current which flows through the transistors present in the circuit due to the imperfect structural manner in which the MOSFET device itself acts as a switch due to the persistent scaling in the physical dimensions of CMOS processes, actively driven by the cost advantages of having a very smaller silicon wafer areas for digital logic functional circuits, MOS transistors had become imperfect switches, leading to the flow of greater leakage currents.

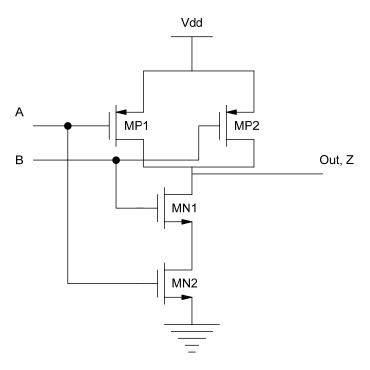


Fig. 2.2 Static CMOS 2-input NAND gate

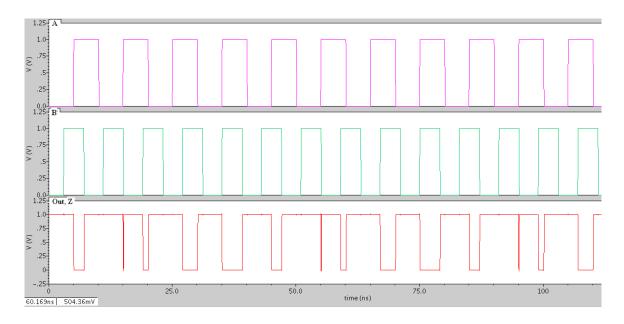


Fig. 2.3 Static CMOS 2-input NAND gate simulation

It is known fact that the CMOS logic family which resulted in significant power savings was apparent to Frank Wanlass who in the year 1963 proved the practicality of CMOS logic and other technical advantages of CMOS design with the technology of a massive monolithic implementation for the very first time. But when this idea of implementation proved unfeasible, Frank demonstrated the same concept with discrete transistors. So, here these particular CMOS implementations alleviated standby power effectively by six orders of magnitude over the other PMOS logic and bipolar logic implementations available. Besides giving impressive results, this advantage related to CMOS logic would not prove decisive for so many years. The former monolithic designs were very small, with the small amount of standby power. Due to the deficient maturity of MOS transistors, in the 1960s, the bipolar logic family raced ahead of MOS technology in various applications. Then the Transistor-transistor logic (TTL) developed in 1962, and Emitter-Coupled Logic (ECL) also referred as Current Mode Logic (CML) developed in 1966 provided effective techniques for the digital design of bipolar logic transistors in the semiconductor industry which is rapidly increasing. There was a point of time in the early years where prime user of CMOS logic was the watch industry. In that era of time battery life was given highest priority than speed. After that MOS technology had begun to mature in the 1970s rapidly, by the contribution of Intel Corporation limited with much of its early

and advanced industrial development. Then the first and foremost world's microprocessor 4004 was released by Intel.

The Intel microprocessor 4004 was invented with a 10 µm line width of PMOS transistor and it used around 2300-2400 transistors running at the chip speed of 108 kHz [8]. Again after having contributed a great research in this field Intel released an 8-bit 8080 microprocessor in the year 1974 which was manufactured with a 6 µm line width of NMOS transistor and used 6000 transistors at the chip speed of 2MHz. Due to the higher mobility of electrons over holes NMOS logic became faster than PMOS logic and hence NMOS obviously started to become the choice of selection.

Advantages of CMOS Logic:

- (1) Robustness(less sensitive to noise).
- (2) Simple approach for implementing logic gates.
- (3) Easy to translate logic to FETs.
- (4) Good noise margins since FETs are in cut off & sizing not critical
- (5) No static power dissipation.
- (6) Low power consumption.

Disadvantages of CMOS Logic:

- (1) Complexity of circuits increases with increased Fan-in.
- (2) For a M-input logic gate, 2M-transistors are required which resulted in relatively large implementation area.
- (3) Propagation delay of CMOS logic gates deteriorates rapidly as a function of the Fan-in.

Fig. 2.4 gives the schematic implementation of a 2 input NAND gate only with NMOS transistors. The PMOS transistors shown in Fig.2.2, MP1 and MP2 are now removed in this implementation of logic and replaced by a resistor, R as shown. In fact the conceptual resistor is implemented by a NMOS transistor with depletion mode. Now when either input A or B is low, the output Z is at Vdd. When both the inputs A and B are at logic high state then the output Z is completely discharged which resulted in logic low value. If one wants this logic to function properly then there is a condition to be satisfied which is related to the current-driving capability of NMOS transistors present in the pull down network which needs to be much greater than the current-driving capability of the resistor in the pull up network. Now this concept demonstrated that the output can be driven to a logic low

voltage level but at the cost of higher power dissipation due to inclusion of the resistor in the pull up network. Besides this standby power dissipation, NMOS logic is a bit slower than CMOS logic due to the presence of weak pull up resistor that functions very slowly when a low to high transition takes place. However the observed drawbacks might have made the NMOS logic appear to be unappealing, still NMOS logic designs are much more compact than CMOS logic circuits. The implementation in Fig. 2.4 has only two NMOS transistors and a resistor R, in comparison with CMOS logic design which has four transistors required to implement its logic. In this way the NMOS logic uses fewer transistors and it is simpler process than CMOS logic design. Then the need to move to CMOS technology had therefore arisen only when the advanced level of integration on integrated circuits (ICs) made the huge standby power dissipation on the NMOS logic circuit design unacceptable. When the 8086/8088 family of microprocessors was released in market in the year 1978 by Intel Corporation Limited, this period of transition occurred since all those implemented designs were almost identical to the family of 8088 microprocessors with an 8-bit bus while the family of 8086 microprocessors with a 16-bit bus only. The power dissipation of 8086 family of microprocessors was 1.5 W with 29,000 transistors at the clock rate of 5-10 MHz, which crossed the nominal 1 W per chip power limit for processing of plastic packaging. The increments in integration levels meant that a 32-bit microprocessor would nominally dissipate 5-6 W, leading to few severe reliability obstacles. Only 250 mW was the power consumption of the CMOS version of the 8086 and 80C86 [9].

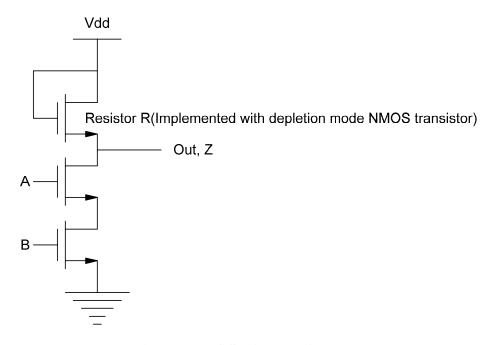


Fig. 2.4 NMOS 2-input NAND gate

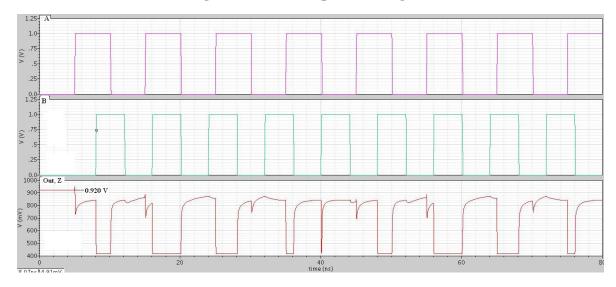


Fig. 2.5 NMOS 2-input NAND gate simulation

The capability of CMOS logic to minimize the total power consumption with increasing large scale integration made it to be best technology which could effectively utilize the advancements in the field of fabrication industry in the recent trends [10]. Also this peculiar CMOS logic maintained till today the advantage of manufacturing large number of digital IC designs using this technology.

With the progression of the manufacturing of semiconductor devices, the biggest challenge was the ability to design and verify all the circuit designs using the increased

number of transistors available then. This demand was achieved by Electronic Design Automation (EDA) software which was developed in the hardware industry in early years. In early days of industry, there might have been an assumption that all the existing techniques for the implementation of digital CMOS logic would be replaced by the ASIC design and its methodologies. That has not happened since many digital IC designs have got their own specific needs that cannot be achieved by standard ASIC techniques. There is a rapid and notable advancement regarding the capabilities of ASIC technology tools in the recent years.

The important common benefits of custom IC design [13-20]:

- (1) The ability to optimize across distinctly separate levels of abstractions in the ASIC design methodologies available which leads to the development of standard cell library used by ASIC design through sequential approach and the opportunity which is provided by custom IC design for using logic families other than standard static logic.
- (2) Also it can utilize certain type of logic families, specifically dynamic logic families (also referred as clocked logic families), which automated the specific design frameworks.

2.2 Different static logic styles

2.2.1 Pseudo N-MOS

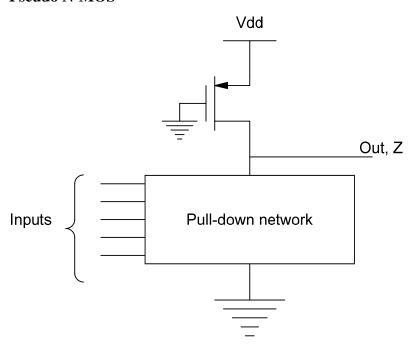


Fig. 2.6 Pseudo N-MOS Logic

Advantages:

For N-input logic gate implementation, only (N+1) transistors are needed. This count is less when compared with Static CMOS Logic.

Disadvantages:

- (1) Noise Margin reduces.
- (2) Static power dissipation increases.

2.2.2 Differential Cascode Voltage Swing Logic (DCVSL)

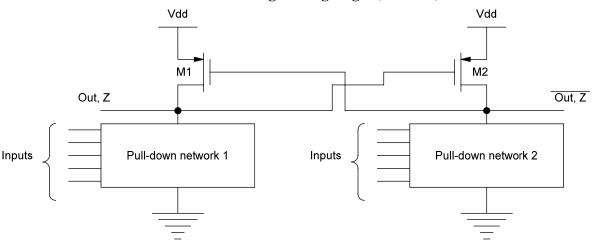


Fig. 2. 7 Differential Cascode Voltage Swing Logic

Pull-down network 1 (PDN1) and Pull-down network 2 (PDN2) are mutually exclusive, i.e. PDN2 = Complement of (PDN1)

Advantages:

- (1) Provides rail-to-rail swing.
- (2) Completely eliminates static currents, thus static power dissipation is eliminated.
- (3) A functional logic and its inverse can simultaneously be implemented.

Disadvantages:

- (1) Exhibits the problem of increased design complexity.
- (2) Power dissipation problem due to cross-over currents.
- (3) During the transition, there is a period of time when PMOS & PDN are turned 'ON' simultaneously, producing a short-circuit path.

2.2.3 Pass Transistor Logic

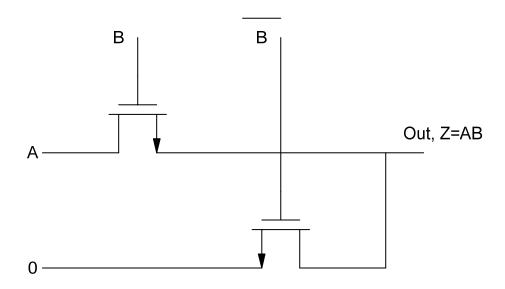


Fig. 2.8 Pass Transistor Logic implementation of 2-input AND gate

Advantages:

- (1) Effective reduction of number of transistors required for implementing a combinational functional logic is the main advantage of PTL family. This could be achieved by driving all the three terminals of MOSFET, Gate, Source and Drain, by active inputs.
- (2) The process of reducing the number of devices has the additional advantage of lowering capacitance.

Applications:

This logic is used in Multiplexers and Latches.

PTL gates cannot be cascaded by connecting the output of a gate to the input of another Pass Transistor.

2.2.4 Differential / Complementary Pass Transistor Logic

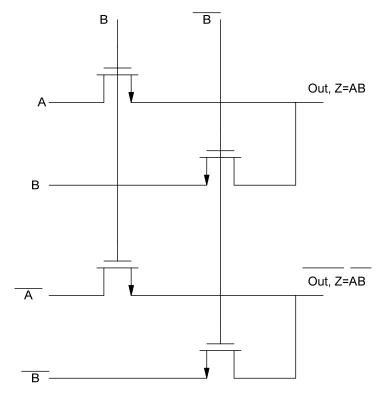


Fig. 2.9 Differential/Complementary Pass Transistor Logic implementation of 2-input AND gate

Advantages:

- (1) Adders, sub tractors and other complex gates such as XNOR, XOR can efficiently be realized using significantly small number of transistors in this DPT Logic.
- (2) This logic is belonging to the family of Static gates, since all the output nodes are always connected to either Vdd rail or ground rail through an existing low resistance path which is an advantage for Noise resiliency.

Disadvantages:

- (1) Static power dissipation is present.
- (2) Reduced Noise Margins.

Remedy for drawbacks:

- (1) Using Level restoration circuit techniques.
- (2) Multiple threshold transistor techniques.
- (3) Implementing Transmission-Gate logic.

Here is a particular type of logic technique called "Transmission Gate logic" which is most widely used technique to deal with problems like voltage drop.

• Transmission-Gate logic:

- ➤ It is built on the basis of the complementary properties of NMOS & PMOS transistors.
- ➤ It combines both device flavors by replacing a PMOS in parallel with a NMOS as shown in Fig.1.10.

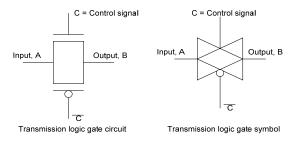


Fig. 2.10 Transmission Gate

Transmission Gate is acting as a Bi-directional switch which is controlled by a control signal-'C' through Gate terminal of MOSFET transistor. If C='1', then both the MOSFETs are on, so they allow signal to pass through the gate. Therefore A=B, while in other case if C='0', then both the MOSFETs are off, so they are in cut-off mode, Thus there is an open circuit between the nodes A and B.

Advantages:

- ➤ It enables rail-to-rail swing although it requires 2-transistors & more control signals.
- ➤ Using these Transmission gates, complex gates can efficiently be built.

Here is the implementation of 2:1 Multiplexer using Transmission Gate Logic.

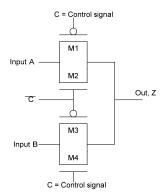


Fig. 2.11 2:1 Multiplexer using Transmission Gate

• When C='0', then transistors M1 and M2 will be ON, and there by output Z is taking signal A.

• When C='1', then transistors M3 and M4 will be ON, thus they are passing signal B to the output node Z.

2.3 Dynamic CMOS logic design

In the Integrated Circuit design industry, by late 1970s, Dynamic logic also referred as the "Clocked logic" was popular in the digital logic design. It can be distinguished from the static logic family with the usage of a driving signal called clock signal in the implementation of Combinational functional logic circuit designs. The use of clock signal in Dynamic logic is to evaluate the combinational function but a sequential logic circuit has also got its own clock signal where it is used to synchronize the transitions in sequential logic circuits.

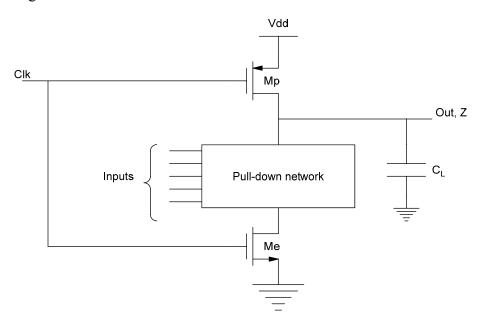


Fig. 2.12 Dynamic CMOS logic

As it has a clock signal which is of a pulse type with two levels '0' (logic low) and '1' (logic high), the basic operation is divided into 2 phases: Pre-charge and Evaluation.

When Clk goes low, pre-charge transistor Mp will be turned 'ON' and transistor Me will be 'OFF'. Thus irrespective of condition of inputs applied to Pull Down Network (PDN), node 'Output' gets charged to 'Vdd' and other nodes may pre-charge to (Vdd – Vth.n) depending on values of inputs.

When Clk is high, transistor Mp will be turned 'OFF' and transistor Me will be 'ON'. Since the actual combinational functional logic will be evaluated through Pull Down Network (PDN) in this phase, this is referred as Evaluation phase. If the input combination to PDN has configured a conducting path then Output nodal voltage may be discharged to 'Gnd'; else Output node stays at logic high. Importantly the inputs must be stable before Clk signal goes to logic high since once the output node has been discharged to gnd, it won't go to logic high again until the next cycle arrives. Thus glitches (dynamic hazards), noise pulses on input signals, cannot exceed the threshold voltage of Me transistor, which is a stringent condition to be highly required for domino logic gates than for static CMOS logic gates.

Advantages:

- (1) There is no static power consumption with an addition of a clock signal input which uses a sequence off phases called Pre-charge and Evaluation.
- (2) Increased speed and reduced implementation area.
- (3) This dynamic logic is twice as fast as the nominal static CMOS logic since it uses only fast NMOS transistors in its evaluation phase in Pull Down Network.
- (4) It is amenable to transistor sizing optimizations.
- (5) Glitches (Dynamic Hazards), due to gates which have non-zero propagation do not occur in dynamic logic.

Glitches (Dynamic Hazards):

The finite propagation delay from one logic block to next logic block when a signal is passing through a gate from input node to output node causes spurious or abrupt transitions at the output node, which are known as Glitches. All the gates have a non-zero propagation delay.

Disadvantages:

- (1) More power consumption because this dynamic logic significantly increases the number of transistors required which are switching at any given instant of time.
- (2) Problems will arise when cascading one gate to next gate. Thus the straightforward cascading of gates is not possible.

Signal Integrity Issues in Dynamic Design:

There are several important considerations that must be taken into account if one wants Dynamic circuits to function properly. They are

- (1) Charge leakage
- (2) Charge sharing
- (3) Capacitive coupling and
- (4) Clock feed through

Charge leakage and Charge sharing occur in Evaluation phase.

2.4 Domino logic circuits

Dynamic logic with an addition of a static inverter at the output node results in a complete domino logic block. This is nothing but a CMOS based exploitation of dynamic CMOS logic circuit techniques which are established on either NMOS or PMOS logic transistors. This particular logic was initially developed to achieve high speed operation in the logic circuits [25-27].

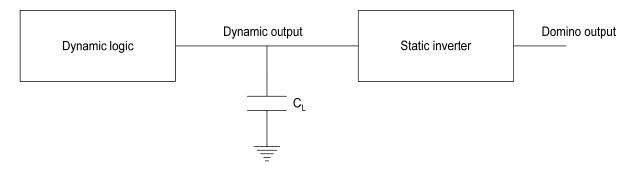


Fig. 2.13 Block diagram of Domino logic

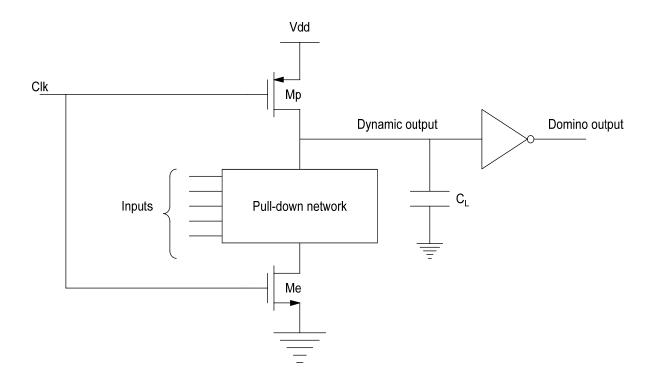


Fig. 2.14 Domino CMOS logic

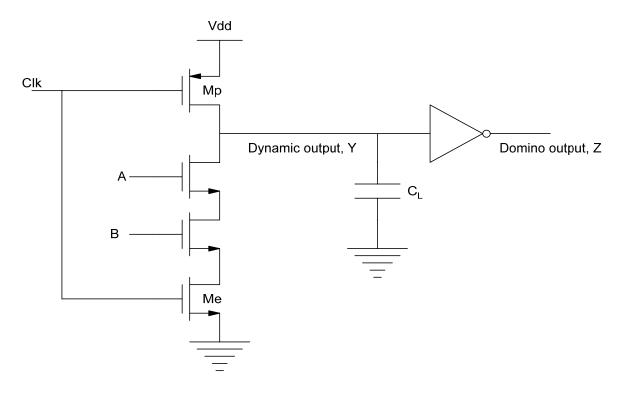


Fig. 2.15 Domino CMOS 2-input AND gate

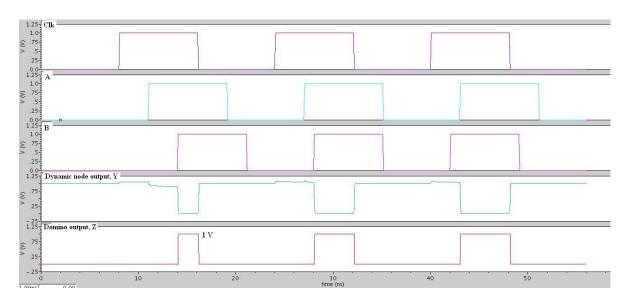


Fig. 2.16 Domino CMOS 2-input AND gate simulation

The 2-input domino AND gate shown in Fig. 2.15 is used to illustrate the logical functionality, the advantage of increased speed, and few challenges involved in using this Domino CMOS logic family. It is observed that there are two inputs, A and B, along with the driving element clock signal, Clk. Since it is an implementation of purely a combinational circuit-AND gate which does not require any clock signal, unlike sequential logic circuits, the presence of clock signal makes it strange. Domino CMOS logic like dynamic logic is also a clocked type of logic family in which every single logic gate has presence of clock signal. If the clock becomes low, then the dynamic node gets charged to supply voltage, causing the domino output, Z to go low due to the presence of inverter between these two nodes. Now this mechanism represents that the gate output to go low logic level once it has been driven high logic level. So this operating period of the block when the clock signal and output are low is known as the Pre-charge phase. Next phase is Evaluation phase which starts when clock signal is high. During this Evaluation phase the actual functional logic is evaluated through Pull Down Network. Therefore the output, Z may go high if both inputs A and B are high, which results the evaluation or dynamic node to be driven to a low logic value. The Evaluation phase as mentioned earlier is the actual functional operating cycle in domino logic circuits with the Pre-charge phase enabling the succeeding Evaluation phase to occur. The significant application of the presence of clock

signal in domino logic ensures that the critical path only traverses through NMOS transistors of Pull Down Network present in the Evaluation phase. One of the advantages of domino cell is that there is no need for the input signals to drive any PMOS transistor present in Pull Up Network as the domino cell only switches from low logic level to high logic level direction.

Now, for a particular instant of current drive, the lack of a PMOS transistor implies that the effective width of PMOS transistor which loads down its previous stage logic favors this domino logic over the static logic. Here it is very critical because the key to very high speed is ensuring that the advantage of high speed can be acquired without loading down the logic block greatly.

Advantages:

- (1) Domino logic circuits allow nominal rail-to-rail swing.
- (2) Domino logic circuits comparatively have smaller areas than static CMOS logic circuits.
- (3) Higher operating speeds are possible since parasitic capacitances in cells are reliably smaller.
- (4) Domino logic ensures free glitch operation as each gate can make only single transition.

Disadvantages:

- (1) There is relatively deterioration of noise immunity due to the presence of problems like unavoidable leakage currents and charge sharing issues.
- (2) Relatively large power consumption when compared to the static CMOS logic.
- (3) Only non inverting functional structures can be built since there is an inverting buffer at the dynamic node.
- (4) Charge distribution may also be taken care off.

List of Important Logic Families

- Resistor Transistor Logic (RTL)
- Resistor Capacitor Transistor Logic (RCTL)
- Diode Transistor Logic (DTL)
- Transistor Transistor Logic (TTL)
- Direct Coupled Transistor Logic (DCTL)
- Integrated Injection Logic (IIL)
- High Threshold Logic (HTL)
- Emitter Coupled Logic/Current Mode Logic (ECL/CML)
- Metal Oxide Semiconductor Logic (MOSL)
- Complementary Metal Oxide Semiconductor Logic (CMOSL)
- Pseudo NMOS Logic
- Enhancement NMOS Logic
- Differential Cascode Voltage Swing Logic (DCVSL)
- Pass Transistor Logic (PTL)
- Differential Pass Transistor Logic (DPTL)
- Dynamic CMOS Logic
- Domino CMOS Logic/Domino Logic

2.4.1 Impact on power consumption

Power consumption is the one of the predominant constraints which plays a vital role in the process of designing any digital dynamic logic circuit. Most of the versatile applications in the microprocessors, digital signal processors and dynamic RAM are based on the technology platform provided by domino CMOS logic family due to their advantage of achieving high speed operation with relatively low device count. Of course there might be inevitable problems like leakage currents and charge sharing issues which normally degrade the degree of robustness in terms of noise immunity or levels of noise margin of the domino CMOS logic family. Nevertheless there is significantly huge power dissipation, in comparison with the other existing static CMOS logic circuits. The continuous down scaling trend of CMOS technology is making the situation even bitterer from genesis to genesis. This chapter investigates the functioning of domino logic circuits with the

encroachment of scaling of process technology. A simple Domino AND gate is designed and simulated at different process technologies.

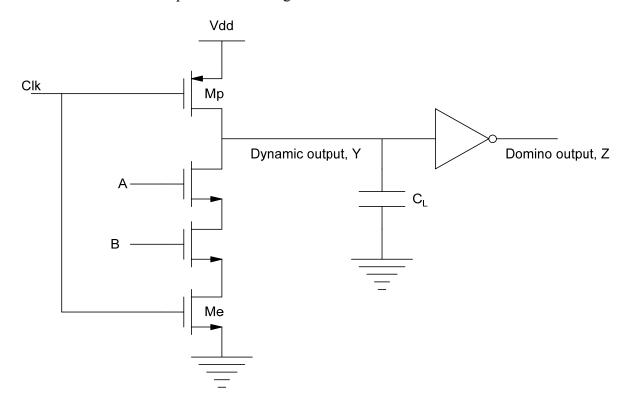


Fig. 2.17 Domino CMOS 2-input AND gate

Table 2.1 Comparison of parameters with technology scaling for Domino CMOS 2-input AND gate

Technology (in nm)	Dynamic power dissipation (in W)	Vth,n (in V)	Vth,p (in V)
1200	3.260E-12	0.607	-0.832
500	3.276E-12	0.708	-0.918
350	3.375E-12	0.549	-0.680
250	11.663E-12	0.365	-0.562
180	14.538E-12	0.372	-0.394

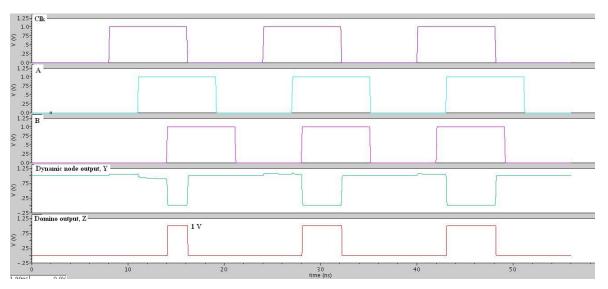


Fig. 2.18 Domino CMOS 2-input AND gate simulation

From the above tabulations, it is evident that the power consumption increases with the down scaling of process technology. Particularly, the logic designer is forced to go for a lower bias voltage or supply voltage so as to reduce the dynamic power consumption. This again demands the minimization of sub threshold voltage to achieve required functional execution of logic along with the consorted increment in the sub threshold leakage current. Therefore an adequately sized parallel PMOS device keeper needs to be inserted to deal with the leakage current issue. Secondly the most important factor of domino CMOS logic, for which it is being preferred mostly over the other logic styles, is speed which will also degrade with the down scaling trend of CMOS process technology due to the presence of inevitable keeper current.

Fig. 2.17 shows domino CMOS 2-input AND gate. In this implementation there is a pull down network (PDN) for realizing the actual functional logic along with clock signal that periodically operates the two phases called pre-charge and evaluation as shown in Fig. 2.19. The total parasitic capacitance at the dynamic node is represented by C_L. If clock is at logic low level, then pull up transistor Mp will turn on which will result the total circuit in pre-charge phase where the dynamic node will get charged to supply voltage through pull up transistor. And also in this phase, the input signals applied to the PDN may be allowed to switch and may get settled to their fixed values. As the evaluation transistor Me is in off state, the discharging path to ground is now disrupted. Now when the Clk is at logic high level, then Mp will turn off and immediately Me will turn on which is known as evaluation phase.

Let's look at the evaluation phase where in, two possible conditions exist for the dynamic nodal voltage to get stabilized. If the current combination of input signals corresponds to a low logic level domino output voltage, then the dynamic node voltage should maintain its nodal potential at supply voltage Vdd, in the form of charge stored by load capacitor. Similarly, if the same input combination leads to a high logic level domino output voltage, then the dynamic nodal voltage ought to be discharged to ground level through the conducting path provided by the NMOS transistors present inside the Pull Down Network..

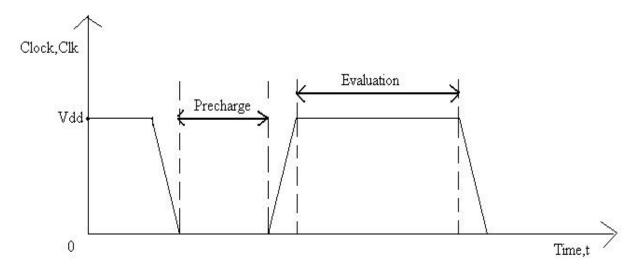


Fig. 2.19 Clock signal in Domino logic Circuit

2.4.2 Technique to compensate charge lost, through PMOS keeper

However, in the previous session there is an assumption that there exists no charge leakage from the dynamic nodal capacitor, C_L . But in real time conditions, however, leakages are found during the clock's evaluation phase through copious provisions, like the sub threshold leakage currents, the gate terminal tunneling currents..., etc despite the fact that the input combination of Pull Down Network is not allowing the dynamic load capacitor C_L , to discharge from Vdd to ground. Although the leakage current is an inevitable parameter in MOS devices, it is very small but finite. Due to concentration gradient between source and drain terminals of MOSFET, absolute temperature, and other conditions like inappropriate doping concentration it flows. It is an inevitable problem in dynamic circuits. Besides the above issues there may a problem with charge sharing because of which the charge of load capacitor, C_L might be shared with the consorted drain

capacitors of the NMOS devices in the Pull Down Network. Therefore, a PMOS keeper as given in Fig.2.4 has to be introduced in order to refill the unavoidable charge lost from the dynamic load capacitor C_L , so that the noise margin levels could be maintained at an operating controlled ambient. So a weak PMOS keeper device is always advisable in parallel with pre-charge transistor as shown in figure since during the clock's evaluation phase, eventually the load capacitor C_L , if the charge stored on it is to discharge through the conducting PDN, will slowly discharge due to the presence of contention keeper current.

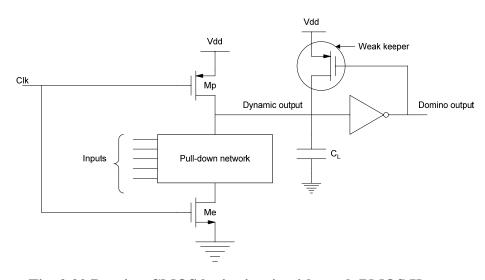


Fig. 2.20 Domino CMOS logic circuit with weak PMOS Keeper

$$P_d = SfC_L V dd^2 (2.1)$$

The above equation explains that the dynamic power consumption, P_d is varying proportionally with the square of the bias voltage, Vdd.

Here S denotes switching scale activity factor, f gives the rate at which the device switches called switching rate and C_L represents the dynamic load capacitor.

From this equation it can mathematically be concluded that the effective reduction in the parameter Vdd, results in prominent minimization of the dynamic power consumption P_d . Added to this, it must also be taken into account that the time delay t_d , will get affected with power supply voltage Vdd, and threshold voltage Vth from the following equation.

$$t_d = \frac{C_L V dd^2}{K (V dd - V_{TH})^{\alpha}}$$
 (2.2)

Here α gives the velocity saturation index whose value is 1.3(approximately 1) for a short channel device and 2 for a long channel device [11-20]. K denotes the parameter which is described by the process of CMOS technology. Now let's take an ideal case from the equation that in order make time delay t_d, independent of Vdd, the threshold voltage V_{TH} , must be set to zero for a short channel device whose approximated value of α is 1. Thus it gives a striking effect on minimizing the time delay t_d with the reduction of Vdd irrespective of V_{TH}. But this assumption is highly impossible as each device possesses nonzero threshold voltage. Therefore this case is to be optimized. The threshold voltage V_{TH}, is to be decreased in such a way that time delay t_d, must not get affected by any means. This operation again results in the subsequent exponential increment in the sub threshold leakage current I_{leak}. Hence in order to compensate this difficulty of getting higher leakages, the logic design engineer is any way forced to go for a PMOS keeper which is a bit larger than earlier minimum sized keeper device. Now with this, it gives two possible conditions for the process of discharging of dynamic load capacitor C_L during the evaluation phase of the clock signal. Firstly, if load capacitor C_L is to discharge, then the entire process will be decelerated due to the keeper device contention current. Secondly, as long as the contention keeper current is maintaining a greater value than the discharging current, the load capacitor C_L will never get discharged to ground level at all.

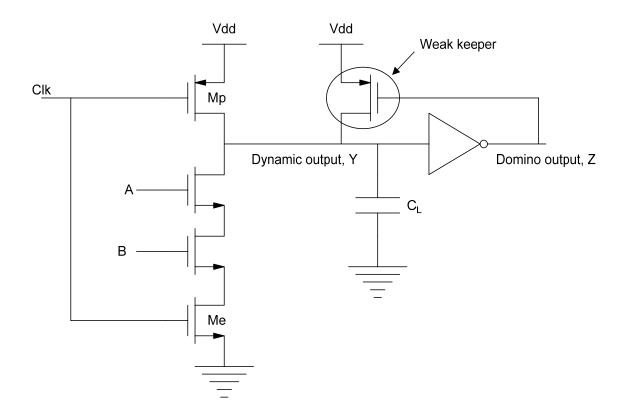


Fig. 2.21 Domino CMOS 2 Input AND gate with weak PMOS Keeper

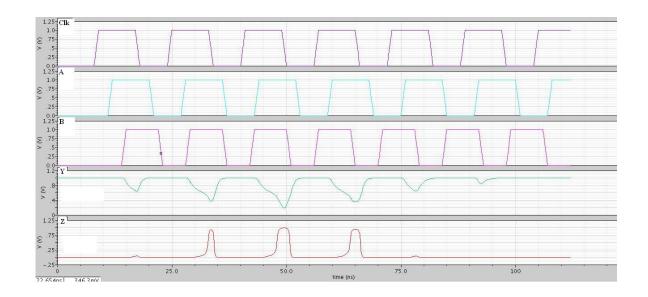


Fig. 2.22 A Domino CMOS 2 Input AND gate with weak PMOS Keeper simulation

Table 2.2 Comparison of parameters with technology scaling for Domino CMOS 2input AND gate with PMOS keeper

Technology (in nm)	Power dissipation (in W)	Vth,n (in V)	Vth,p (in V)
1200	21.896E-12	0.607	-0.832
500	21.880E-12	0.708	-0.918
350	22.086E-12	0.549	-0.680
250	38.682E-12	0.365	-0.562
180	60.856E-12	0.372	-0.394

Observations from PMOS keeper circuit:

- (1) The motivation for the reduction of the dynamic power consumption P_d , demands the logic designer to choose relatively small power supply voltage Vdd along with an effective lower threshold voltage (V_{TH}), so as to maintain the performance and reliability of logic circuits.
- (2) The process of reducing the threshold voltage (V_{TH}), results in the increment of the sub threshold leakage current, exponentially which demands the entrepreneur to select a greater than the small sized device keeper. This, during the clock's evaluation phase, will in turn increase the contention keeper current of the parallel PMOS transistor which will gradually decelerate the process of discharging of dynamic load capacitor.
- (3) Also it is concluded that in comparison with other existing logic styles, this particular domino CMOS logic technology loses its basic and fundamental advantage of high speed operation with the down scaling trend of the CMOS technology which became an essential requirement and imperative solution for all the applications where large number of faster NMOS devices are used in parallel inside the Pull Down Network (PDN), with the subsequent increment in the sub threshold leakage current of devices.

2.5 Conclusion

Therefore, this chapter has given overview of standard logic styles in brief and introduced the dynamic logic followed by domino logic circuits with description. The research area is primarily focused on present working environment-Domino logic from a broader angle. Description of various circuit styles along with their advantages and disadvantages is illustrated. In addition to this, the functioning of domino logic with the encroachment of down scaling of process technology is investigated and this analysis reveals that with deep sub-micron, more power is consumed. This is presented in Table 2.1. Technique, which uses a PMOS keeper at dynamic node, to alleviate inevitable charge-lost is reviewed and corresponding simulation result is presented in Table 2.2. A detailed review on domino logic is conducted and issues related to domino logic are brought out that facilitated us proceeding to the next chapter.

CHAPTER 3 NOVEL DOMINO LOGIC TOPOLOGIES

3.1Introduction

Domino CMOS logic circuit family explores ample applications in the process of designing high performance circuits because of its increased speed and reduced implementation area, despite its high noise sensitivity for which the sub-threshold leakage current through the evaluation network is the prime factor. The issue becomes much more complex with constant down scaling of process technology [30-35].

Due to the fast-growing trend of down scaling of the process technology especially in the field of deep sub-micron, noise-tolerance has become one of the primary concerning issues in the design of VLSI chip manufacturing field. This noise in the digital VLSI integrated circuit design normally relates to the fluctuation which may cause the variation of a nodal voltage from its base value. There have been distinctly various sources of noise in the deep-submicron regime which are typically associated with crosstalk, charge sharing, inevitable leakage currents and minute variations of the supply voltage from its base value. The leakage current is the most predominant factor which is found to be rising exponentially with the down scaling of process technology [31, 32]. However, the bias voltage is to be reduced to limit dynamic power consumption. Simultaneously, the down scaling of the threshold voltage (V_{TH}) of the MOS devices to ensure high performance causes the sub-threshold leakage current to increase exponentially, as it is dependent on -V_{TH}. Besides this, there is a definite increment in the gate leakage current with the subsequent reduction of the gate oxide thickness due to the impact of gate oxide carrier tunneling.

Therefore, to achieve high performance in various aspects, several design techniques are widely exploited. Domino logic design is one amongst them which is faster than static counterpart. Also these gates are much more compact, exclusively when dealt with circuits with wide fan-in gates such as multiplexers, which became the process of choice of high performance in microprocessors and digital signal processors. Added to this, domino logic gates exhibit high noise sensitivity when compared with static CMOS logic gates due to the low threshold voltage, which is almost equal to V_{TH} of NMOS transistors in the pull-

down network. Noise sensitivity is one of the major concerning issues in recent days, distinctly in the process of designing circuits with wide fan-in gates.

Versatile techniques have been illustrated to alleviate the noise leakage issues of several wide fan-in domino logic gates [33-40] and there is an improvement of the noise robustness. Furthermore, the reduction in speed and the increment in power dissipation became more tactile with continuous down scaling of process technology.

Novel techniques are proposed in this chapter to increase the noise robustness of domino logic gates with low power and reduced leakages. In order to demonstrate the proficiency of proposed scheme, comparisons are made with existing techniques and it is illustrated that from the perception of Power Delay Product (PDP), the proposed techniques exhibit PDP lower than that of the existing techniques. Simulation results reveal that the proposed technique can attain high degree of noise robustness with low power and low leakages in CMOS 90nm technology platform used for analyzing wide fan-in logic circuits. Also the predominant noise metric parameters such as UNG and ANTE are significantly improved in proposed techniques. To probe further, dealing with sub-threshold leakages, the proposed schemes effectively minimize the leakage power.

In section 2, several existing domino logic schemes dealing with noise tolerance associated with leakages and power dissipation are discussed. Section 3 demonstrates the proposed novel domino logic techniques. Analyzing the noise metric parameters like UNG and ANTE is executed in section 4. Section 5 describes the process corner analysis. Simulation results through tabulations and graphs along with discussion are given in section 6 and concluding remarks are made in section 7.

3.2 Different high-performance noise tolerant circuit techniques

3.2.1 Wide fan-in Domino OR gate-Footless and Footed schemes

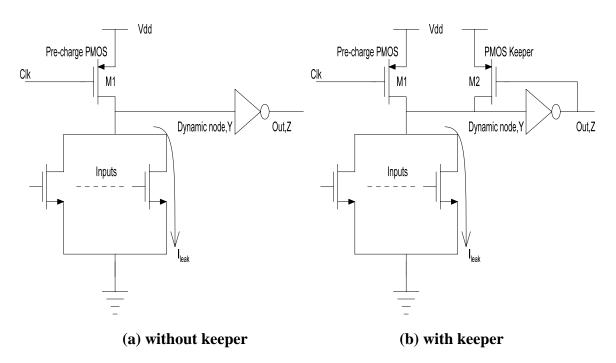


Fig. 3.1 Wide fan-in domino OR gate-footless

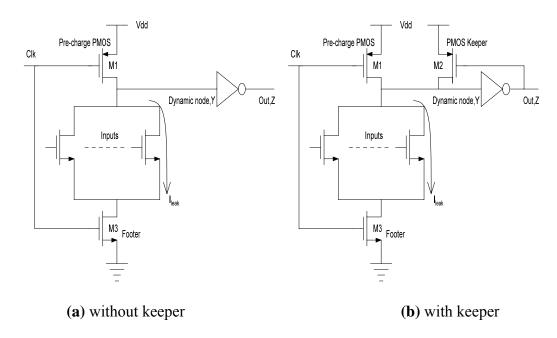


Fig. 3.2 Wide fan-in domino OR gate-footed

Fig. 3.1 and Fig. 3.2 show the implementation of wide fan-in domino OR gate footless scheme and footed scheme with and without keeper wherein the footless technique is

frequently exploited in high-performance logic circuits because of its quick discharging process of the dynamic node, with minimal capacitive load of the Clk signal that manages the entire domino logic operation [30-40].

During the pre-charge phase (when Clk is low) the dynamic node gets charged to supply voltage Vdd through pre-charge PMOS transistor. When the Clk becomes high the evaluation phase starts where the actual logic function is being evaluated by elaborating the input signals through gate terminal during which the dynamic node doesn't possess any connection to Vdd rail through pre-charge PMOS device.

Employing keeper technique is an advantage to avoid the charge lost from dynamic node when needed a strong one at this node [45]. There may be a chance of getting deterioration of strength of voltage at dynamic node due to the inevitable flow of sub-threshold leakage currents through pull-down network in Clk's evaluation mode and hence this voltage drop needs to be recompensed by other means in order to stabilize the node. This could be achieved by employing PMOS keeper device between the dynamic node and supply rail whose gate terminal is driven by domino output. As long as domino output Z, is low, keeper charges dynamic nodal voltage up to Vdd which may have been forced to go for wrong discharge by the flow of sub-threshold leakage current through pull-down network even if all the inputs of pull-down network are low. This effect becomes much more severe with the occurrence of noise pulse or glitch at any of the input nodes. However, the noise voltage impulse triggers notable effects such as the exponential rise of sub-threshold leakage current through the pull-down network with the abrupt fluctuation of gate-tosource voltage (V_{GS}) which leads to forceful discharge of the dynamic node. Also there is a reduction in gate leakage current, because of reduction of V_{DG} with V_{GS} increasing, which is almost irrelevant when compared to the significant increment of the sub-threshold leakage current. In reality, the impact of V_{DG} on gate leakage current is not as much significant as the effect of V_{GS} on the sub-threshold leakage current. Hence there is troublesome mechanism associated with discharge of dynamic node owing to the existence of noise glitch at the input nodes of pull-down network. Besides these issues, the secondary cause that forces dynamic node to discharge wrongly during the evaluation phase is an abrupt change at the ground level nodal voltage since if at all exists a negative pulse, at the ground level, it may increase the effective V_{GS} of NMOS devices of the pulldown network which in turn spurs greater flow of sub-threshold leakage current that again stimulates faulty process of discharging the dynamic node.

Thus existing circuit techniques however reduce the mentioned effects, either by minimizing the flow of sub-threshold leakage currents [33] or by using keeper transistor [37] that yields enhanced charge restitution at the dynamic node.

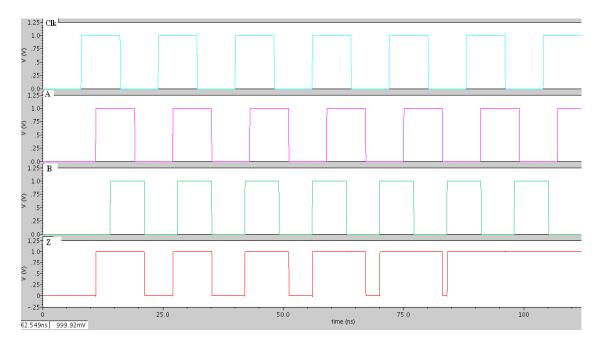


Fig. 3.3 Domino 2-input OR gate-footless simulation

The output Z is low as long as Clk is in pre-charge phase since the dynamic node is charged to Vdd. When evaluation starts, the actual logic function operation takes place. When the all the inputs are low then the dynamic node is at logic high and now if any one of inputs is high then that corresponding transistor in pull-down network turns on providing a discharging path for dynamic node and resulting in a logic high domino output. Since there is no footer between pull-down network and ground rail, the discharging mechanism is faster than that of the footed scheme which makes it favorable for widely being used in high performance circuit designs. In footed technique the footer transistor turns off in pre-charge mode and will turn on in evaluation phase. Thus if the dynamic node once discharged, cannot be charged immediately as it is driven by Clk line and therefore it again charges to supply voltage until the next pre-charge phase initiates. The simulation results for footless and footed schemes are shown in Fig. 3.3 and Fig. 3.4. A, B

represent applied inputs and Z shows the corresponding output that takes OR operation between A and B.

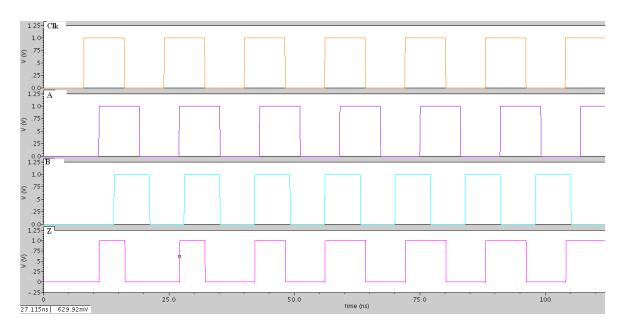


Fig. 3.4 Domino 2-input OR gate-footed simulation

3.2.2 Wide fan-in Domino OR gate Diode footed scheme

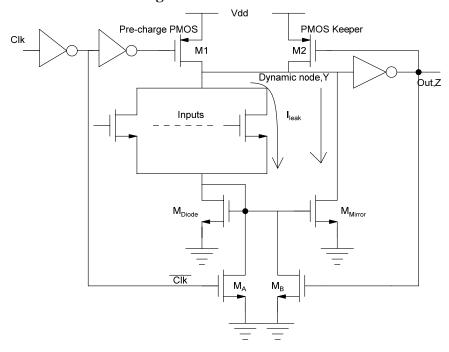


Fig. 3.5 Wide fan-in Domino OR gate-Diode footed scheme

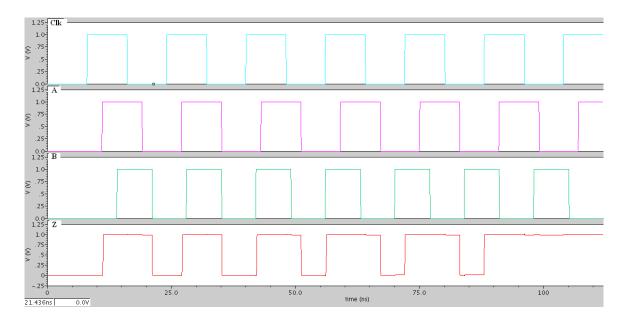


Fig. 3.6 Domino 2-input OR gate-Diode footed scheme simulation

A Wide fan-in Domino OR gate Diode footed scheme is shown in Fig. 3.5 [40]. The flow of sub-threshold leakage current in the evaluation phase through pull-down network is controlled by the diode footer transistor-M_{Diode}, through the phenomenon called stack effect [40-55]. The leakage current through pull-down network constitutes a nominal voltage drop across the diode footer which makes the V_{GS} of the 'off' NMOS devices inside the pull-down network negative and thereby reducing the flow of sub-threshold leakage current exponentially. Also the diode voltage drop increases the body effect of the pull-down devices, which helps in the subsequent reduction of sub-threshold leakage reduction [49]. The simulation of domino 2-input OR gate using diode footed scheme is given in Fig. 3.6. Besides this, the diode footer device doubles the threshold voltage of gate and hence the new switching threshold voltage becomes 2V_{th-n}. Higher the gate switching threshold voltage, greater the noise immunity. But this in turn, as a drawback of this technique, raises equivalent resistance of the evaluation path of pull-down network by the diode footer that eventually makes the gate slower. An alternative discharging path has been established and also assures high performance by the mirror network which consists of the NMOS devices MA, MB and MMirror. The cause behind the deterioration in the performance is the decrement of evaluation current by the diode footer. In order to stabilize this evaluation current, the mirror transistor is connected as shown in Fig. 3.5. M_{Mirror} is used for mirroring action of evaluation current and is drained from the dynamic node.

Thus, the total evaluation current is the sum of evaluation current through the pull-down network and the mirrored evaluation current.

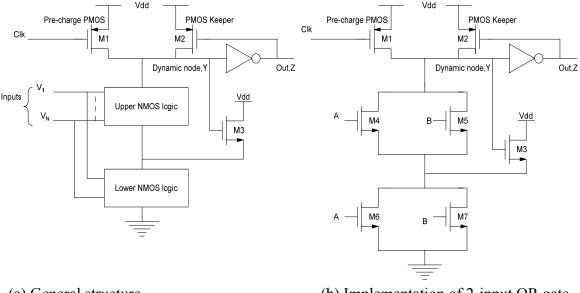
The mirror ratio is defined as the ratio of the current driving capability of the mirror transistor to that of the diode footer and the relation is given in following equation [59].

Mirror ratio =
$$\frac{\left(\frac{W}{L}\right)_{(Mirror)}}{\left(\frac{W}{L}\right)_{(Diode)}}$$
 (3.1)

During the pre-charge phase of Clk signal, transistor M_A is on which turns off the mirror transistor (M_{Mirror}) to limit the short-circuit current through M_{Mirror} in pre-charge phase. M_A turns ON during Pre-charge mode only and in evaluation it turns OFF. Thus the corresponding explanation is given based on this condition. Domino output node triggers the transistor M_B to pull down both footer node and dynamic node to a logic low level, if the output Z goes high in the evaluation phase. Thus any possible short circuit power dissipation is prevented in the static inverter during the evaluation phase. Owing to the nominal reduction in the sub-threshold leakage current of pull-down network, small PMOS keeper is enough to suffice the purpose. The overall performance can be enhanced by increasing the mirror ratio but at the cost of lower robustness, because the device M_{Mirror} also drains small amount of leakage current from the dynamic node. Therefore there is a trade-off between circuit robustness and its performance by the mirror ratio. The impact when upsizing the PMOS keeper in domino circuit resembles the same result when downsizing of M_{Mirror} is done in this technique.

Low area overhead is the main merit of this technique and the demerit is that high degree of circuit robustness could be achieved with gate delay penalty. Since the diode connected NMOS increases the equivalent resistance of pull-down path, the gate becomes slower. To avoid gate noise sensitivity, the following technique is preferred.

3.2.3 Wide fan-in Domino OR gate-Replicated evaluation scheme



(a) General structure

(b) Implementation of 2-input OR gate

Fig. 3.7 Wide fan-in Domino OR gate-Replicated evaluation scheme

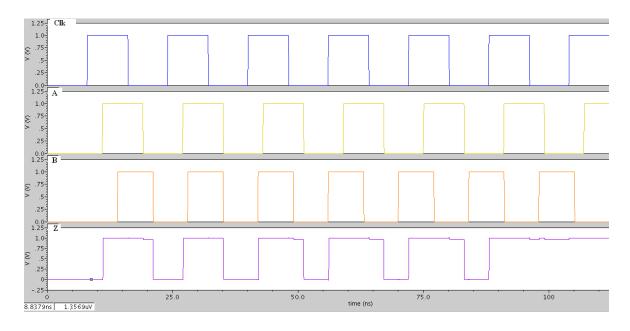


Fig. 3.8 Domino 2-input OR gate-Replicated evaluation scheme simulation

Fig. 3.7 depicts Wide fan-in Domino OR gate with Replicated evaluation scheme and implementation of 2-input OR gate which minimizes the gate noise sensitivity by retroflexing the evaluation network [40-50]. This circuit simulation is shown in Fig. 3.8. This technique effectively mitigates the flow of sub-threshold leakage current by means of the NMOS transistor M3 inserted between two evaluation networks that intends to increase the nodal voltage between them which in turn lessens the V_{GS} of the NMOS devices of the upper evaluation network. Nevertheless, the parameters gate delay and implemented area are substantially compromised. To probe further the significant increment in the circuit complexity of this technique due to the inclusion of two series connected evaluation networks makes the capacitive load at each input line double which countermands the basic advantages of a domino logic gate as compared to the static counterpart with minimal input node capacitive loads.

3.2.4 Wide fan-in Domino OR gate-Dynamic node footed scheme

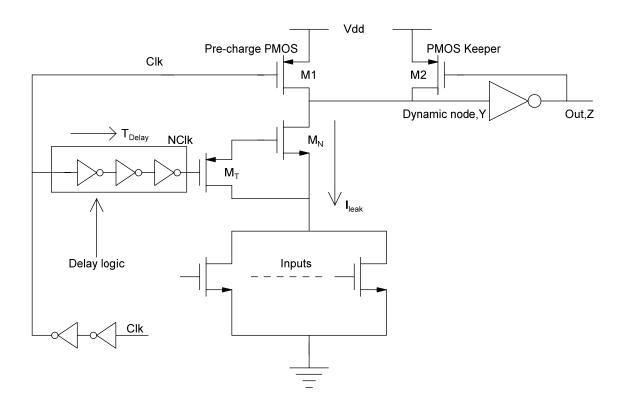


Fig. 3.9 Wide fan-in Domino OR gate-Dynamic node footed scheme

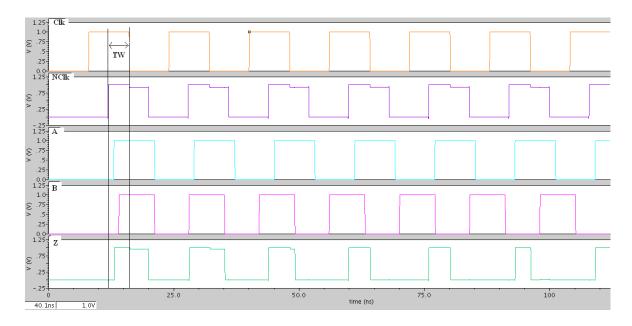


Fig. 3.10 Domino 2-input OR gate-Dynamic node footed scheme simulation

Fig. 3.9 gives Wide fan-in Domino OR gate Dynamic node footed scheme which has an NMOS device M_N, inserted between the dynamic node and the evaluation network. The series of static inverters, that constitutes some delay and the PMOS transistor M_T, turn M_N on properly when required. In the evaluation phase, the finite amount of delay of the series of inverters (T_{Delay}) makes both the Clk and NClk lines high. This is called "Transparency Window (TW)" during which the gate elaborates the signal inputs and the dynamic node can finally be discharged. Once the T_{Delay} is elapsed, M_N turns off due to NClk signal that becomes low. Therefore in this situation, the charge leakage from the dynamic node is minimized exponentially and also the noise robustness is increased with the stacking effect. As a merit, a very low area over head is also assured by this technique similar to the diode-footed technique. Coming to the drawbacks, due to extra added transistor M_N, there has been significant increment of capacitive load with Clk line as well as enhanced discharging path equivalent resistance. The parameters noise immunity and gate delay are dominantly affected by T_{Delay} through the series of PMOS transistors connected as shown. The simulation of domino 2-input OR gate with dynamic node footed scheme is shown in Fig. 3.10. Fig. 3.11 gives the description of transparency window along with other phases of operation.

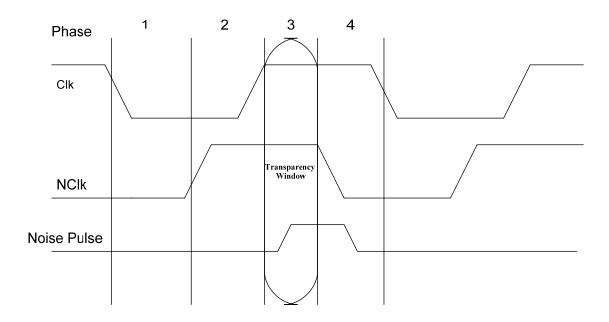


Fig. 3.11 Transparency Window-phase3 waveform

- (1) If the T_{Delay} is longer, then the gate transparency window becomes wider that results in lowering the immunity towards noise but with faster gates. On the other hand
- (2) If the T_{Delay} is smaller, then there will be thick transparency window which leads to lowering the gate speed with increased noise robustness.

The above analyzed circuit techniques, however, reduce the subsequent flow of inevitable leakage currents through the dynamic node by means of device staking effect. Nevertheless these schemes are not meant to reduce the overall leakage current. Upsizing the keeper PMOS transistor ensures faster replenishing of charge storage at the dynamic node since the electric current flow from supply rail to dynamic node becomes much faster with widening of keeper device. Despite this merit, there is troubling issue with DC contention current flow through PMOS keeper and pull-down network with the discharging of dynamic node when the evaluation network is turned on. Demarcated by this phenomenon the PMOS keeper drives the dynamic node to preserve its charge with its provision of alternative path. This peculiar action results in lowering the gate speed along with increased power dissipation owing to the existence of DC contention current flow through the static inverter at output, the PMOS keeper path and the pull-down network.

3.2.5 Wide fan-in Domino OR gate-Clock delayed single keeper scheme

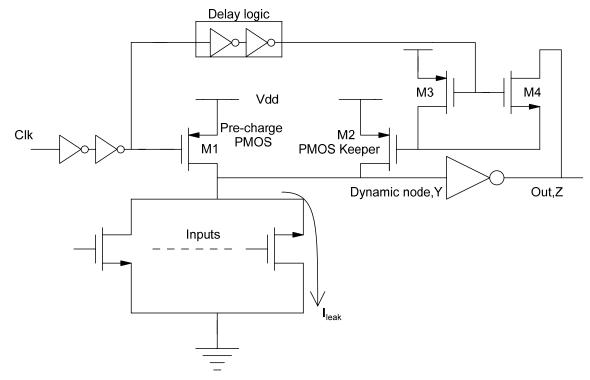


Fig. 3.12 Wide fan-in Domino OR gate-Clock delayed single keeper scheme

Fig. 3.12 presents a clock delay based circuit scheme that diminishes DC contention current flow due to PMOS keeper by means of an extra added circuit which limits the operation of keeper by making it on after a finite amount of delay and off when the evaluation phase is about to commence [50-60]. This implies the designing of a delay logic circuit (DLC) to deal with DC contention current related to keeper, dynamic node and evaluation network necessitates appropriate functioning. If the evaluation network is turned on, then the dynamic node normally discharges very fast. However, the delay logic circuit turns this keeper device off during this period. Another merit is when the dynamic node is forced to be wrongly discharged due to random noise impulse at the inputs of PDN then the delay circuit replenishes the charge lost from it by turning the wide PMOS keeper on. Fig. 3.12 represents Wide fan-in Domino OR gate-Clock delayed single keeper scheme designed for high speed domino proposed in [9]. Despite increased capacitive load of Clk signal due to extra delay logic circuit, this scheme is efficiently reducing the gate noise sensitivity. The simulation of domino 2-input OR gate with clock delayed single keeper scheme is shown in Fig. 3.13.

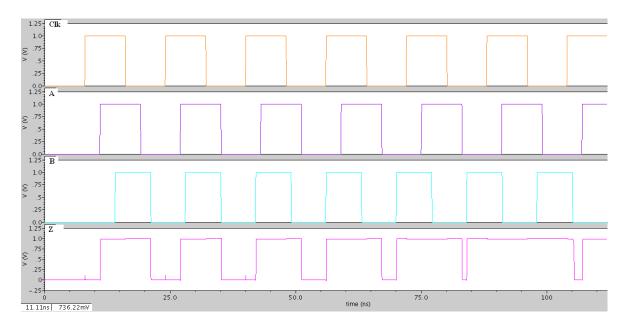


Fig. 3.13 Domino 2-input OR gate-Clock delayed single keeper scheme simulation

3.2.6 Wide fan-in Domino OR gate-Clock delayed dual keeper scheme

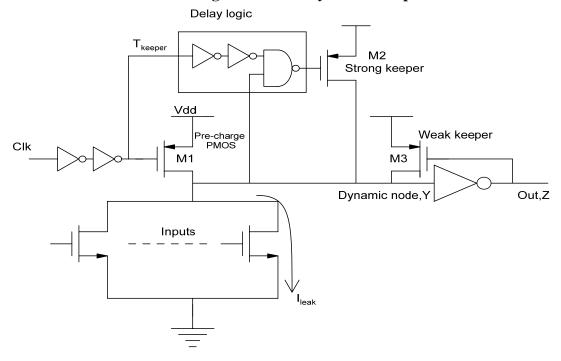


Fig. 3.14 Wide fan-in Domino OR gate-Clock delayed dual keeper scheme

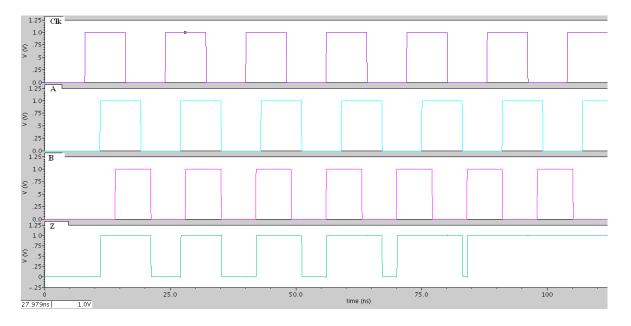


Fig. 3.15 Domino 2-input OR gate-Clock delayed dual keeper scheme simulation

Fig. 3.14 shows Clock delayed dual keeper scheme implementation as the name itself indicates that there are two keepers, a weak keeper and a strong keeper, used in the operation wherein the former one functions before the commencement of evaluation phase as usual to pre-charge the dynamic node and latter device will turn on whenever the noise glitch occurs [50-70]. This scheme has a delay logic circuit that is implemented using conditional keeper domino that uses a static NAND gate along with two series of cascaded inverters. Also the delay logic circuit replaces its NAND gate with a static inverter in order to implement another high speed domino scheme called Skew tolerant high speed (STHS) circuit. Simulation result reveals that above scheme increases capacitance of Clk line because of the extra delay logic circuit and causes power dissipation even if the output does not switch. Fig. 3.15 gives the simulation of domino 2-input OR gate implemented using clock delayed dual keeper scheme. A, B represent two inputs and Z shows the OR gate operation by using the dual keeper scheme.

3.2.7 Wide fan-in Domino OR gate-Skew tolerant high speed scheme

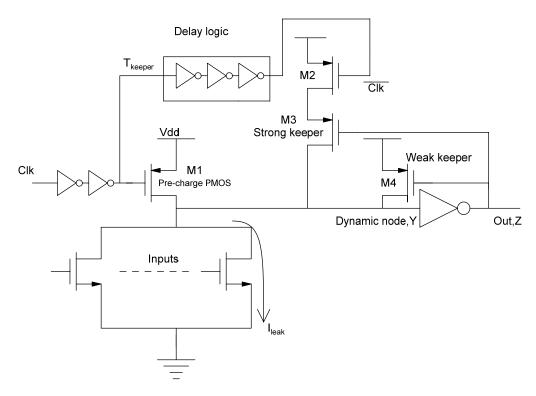


Fig. 3.16 Wide fan-in Domino OR gate-Skew tolerant high speed scheme

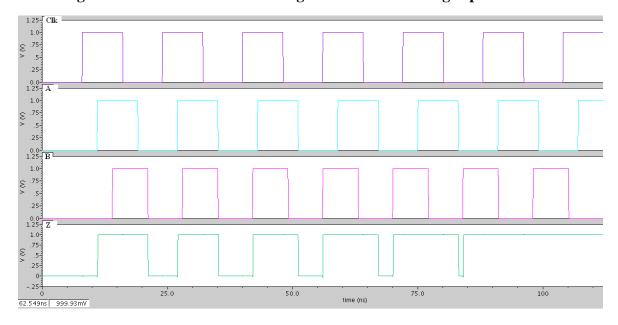


Fig. 3.17 Domino 2-input OR gate-Skew tolerant high speed scheme simulation

Fig. 3.16 depicts a new high speed clock delay based domino logic circuit called Skew tolerant high speed domino logic circuit (STHS) which as described earlier uses two keepers [50-80], a weak one and a strong keeper, where the former one works before commencement of evaluation phase as usual to pre-charge the dynamic node and latter device will turn on whenever the noise pulse or any glitch occurs. The difference between STHS scheme and dual keeper scheme is that the designing of delay logic circuit along with structure of conditional keeper circuit at dynamic node. The delay logic circuit replaces the NAND gate with a static inverter as shown above in order to implement this topology and also both the strong keeper and weak keeper are simultaneously driven by domino output. Simulation results reveal that the above clock delay based domino schemes increase capacitance of Clk line because of the extra added delay logic circuit and cause power dissipation even when the output is not in the switching mode. The corresponding simulation of domino 2-input OR gate using skew tolerant high speed scheme is shown in Fig. 3.17.

Therefore, with variation in the functioning of delay logic circuit, the trade-off amongst gate delay, noise immunity and power dissipation can strategically be achieved. If the delay of DLC is larger then, the gate becomes faster along with low power dissipation and reduced noise immunity. Otherwise, the small delay makes the gate robust but with increased gate delay and power dissipation penalties.

3.2.8 Wide fan-in Domino OR gate-Source following evaluation gate (SFEG) scheme

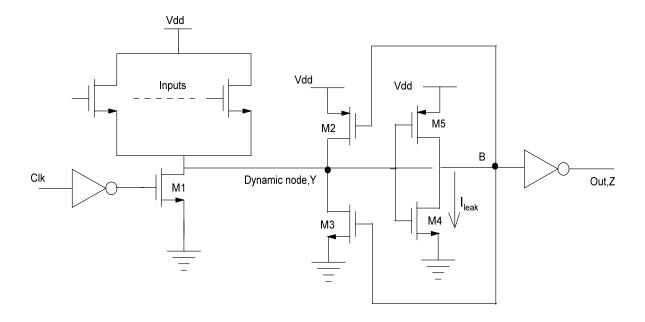


Fig. 3.18 Wide fan-in Domino OR gate-Source following evaluation gate (SFEG) scheme

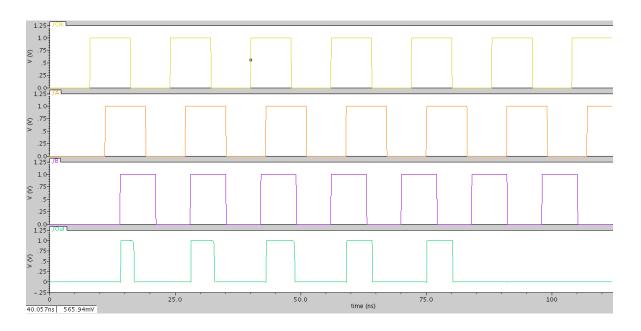


Fig. 3.19 Domino 2-input OR gate -Source following evaluation gate (SFEG) scheme simulation

The Wide fan-in Domino OR gate source following evaluation gate (SFEG) scheme, shown in Fig. 3.18 demonstrated improved noise robustness by replacing PMOS devices in the pull-up network with NMOS transistors [40-70]. Now the dynamic node gets charged through the leakage currents which flow through the evaluation network. Therefore, due to this peculiar phenomenon, there is subsequent reduction in the V_{GS} of the NMOS devices that leads to exponential decrement in sub-threshold leakage currents. Other merit of this scheme is the node triggering static inverter does not couple with the dynamic node so that if at all any leakage exists, that is only due to the NMOS device (M4) as shown in Fig. 3.18. The simulation of domino 2-input OR gate implemented using source following evaluation gate scheme is shown in Fig. 3.19.

The demerit is associated with producing strong one at dynamic node since pull-up network which is normally built up with PMOS transistors, specifically in this scheme, consists of NMOS devices which can charge the dynamic node only up to Vdd– V_{TH} during the gate switching period. Therefore, this threshold voltage drop needs to be recompensed and this can be established by the PMOS device M2. However, owing to the presence of finite delay of feedback connection that triggers the transistor M2, the same device is not instantly turning on which results in succeeding flow of short-circuit current, during the period of gate switching, through the devices M4–M5 as shown , leading to increased dynamic power dissipation.

3.3 UNG & ANTE

There are two kinds of metric parameters in digital logic circuits which are to be taken in to account for the measurements that are widely exploited to compare noise immunity for robustness of corresponding designed digital logic circuits. UNG and ANTE are the main parametric quantities that will evaluate how robust the circuit is towards noise and whether the circuit is able to sustain, in an ambient where there is much vulnerability of getting affected by random changes occurring in surrounding premises, without altering its nominal desired logic functional operation.

Unity Noise Gain (UNG):

It is the amplitude of the noise pulse at the input node which causes the glitch or dynamic hazard with the same amplitude at the output node. As the name implies that there is a gain of voltage amplitude of output glitch to the voltage amplitude of input noise pulse and that gain factor is required to be unity in order to possess a high degree of robustness. Closer the voltage gain value towards unity (one), higher the noise immunity or robustness of the circuit.

$$UNG = \{V_{noise}; when V_{noise} = V_{out}\}\$$

Average Noise Threshold Energy (ANTE):

It is the average input noise energy that the circuit can tolerate without altering its normal logical functioning. Also it is the UNG measurement but with an average calculation for few set of values within the conditional boundaries. Larger the ANTE, greater the circuit robustness.

$$ANTE = \left(\frac{1}{K}\right) \sum V_{noise}^2 T_{noise}$$

Where V_{noise} = Amplitude of input noise pulse,

 T_{noise} = Pulse width of input noise pulse and

K = Number of iterations or observations.

For the measurement of these quantities, noise pulses are applied at the input nodes of pull-down network in the evaluation period of clock signal, and the applied noise pulses possess certain amplitude with pulse width and time duration. There are two methods to take the measurement of output glitch by changing the shape of input noise pulse in different aspects. Varying the voltage amplitude of input noise pulse randomly by keeping its pulse width constant, to observe the corresponding changes occurring in the glitch is one method to calculate the UNG. Otherwise keeping the amplitude of input noise pulse voltage fixed at one level and computing the respective changes in the amplitude of output glitch for all

the iterative trails of changing the pulse width of input noise pulse is another method to figure out the UNG.

In this work the UNG computations are done with respect to the effective changes in the voltage level amplitude of input noise pulse for different iterative trails by keeping the pulse width constant. The simulations are done in CMOS 90nm process technology with 1 V power supply as bias voltage where the input noise pulse voltage amplitude is varied in iterative trails to observe the corresponding reflections at the output glitch by keeping the input noise pulse width constant at 7 pS. The same noise pulse is used for carrying out the complete analysis of the UNG measurement for different benchmark circuits and proposed schemes. The UNG is defined as per the equation given below and Fig. 3.20 shows the typical way of calculating the UNG with reference to clock signal and input noise pulse through waveforms. It is noticed that 0.5 V amplitude of output glitch is observed with the same amplitude of input noise pulse along with the driving element clock signal which is in the evaluation mode (when Clk=1) with an amplitude 1 V. Therefore the UNG is calculated to be 0.5 V.

$$UNG = \{V_{noise}; when V_{noise} = V_{out}\}$$

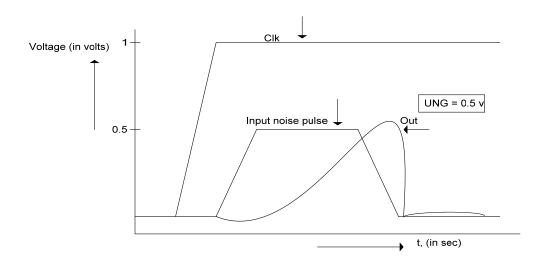


Fig. 3.20 Typical UNG measurement wave form

Noise and leakage issues in domino logic circuits

Noise is dominant parameter for the analysis of any digital logic circuit at different ambient conditions where the functionality may vary depending on the application and design requirement. Noise in digital circuits has got its own significant role to govern the functionality of the particular design altogether. There are different sources of noise in deep submicron regime. Dynamic logic assures high performance when compared to static counterpart circuits. But there are few significant parametric considerations that must be taken into account for proper functioning of clocked logic circuits. These issues are referred as signal integrity issues which are mentioned here but detailed explanation with complete analysis is given in the next chapter. They are

- (1) Crosstalk,
- (2) Charge leakage currents,
- (3) Charge sharing,
- (4) Capacitive coupling,
- (5) Clock feed through and
- (6) Small variations of nominal supply voltage.

3.4 Process Corner analysis

In the process of semiconductor manufacturing industry, the process corner analysis, which plays a vital role as an effective technique for Design of Experiments (DoE), refers to random variations of physical fabrication parameters exploited in employing the design of an Integrated Circuit (IC) to a silicon wafer. The extreme variations in the parameters which will show their impact on the overall functionality of the design are analyzed by this process corner analysis within which that particular circuit that is etched onto the wafer is expected to operate correctly without deviating from its normal functionality. Thus a fabricated design being tested at these various process corners may run slower or faster than specified speed of operation but if it does not operate at all at any of these conditional variations or corners, then the circuit design is considered to be an inadequate design.

Robustness of any designed digital IC is the primary designing metric parameter which must be scrutinized thoroughly by the manufacturers of semiconductor industry at different ambient conditions by subjecting the design to extreme conditional variations like threshold voltage, clock frequency, and temperature. This typical phenomenon is called design characterization that results in getting new observations and solutions which can be plotted with the help of a graphical technique known as 'shmoo' plot with the prior indication of boundary limitations of the circuit design beyond which the design functions to fail for a given set of combinational process corner variations at different environmental conditions.

This, also referred as corner-lot analysis, is the most effective method in digital electronics as it includes the process of subjecting the design to the extreme conditional variations directly on the speed of switching activity of transistors which is irrelevant for normal analog circuits.

In VLSI microprocessor design and semiconductor fabrication, this process corner analysis constitutes few variations typically five from the normal doping concentrations of transistors on a silicon wafer which can cause predominant alterations in the operation through the duty cycle and slew rate of signals of digital circuits and can sometimes lead to the catastrophic failure of the total system.

Types of corners:

There are two kinds of process corners typically FEL (Front End of Line) and BEL (Back End of Line) where the former one is normally employed at the schematic design level and the latter one includes the effect of PVT (Process-Voltage-Temperature) variations on on-chip interconnections.

FEL corners:

Nomenclature for FEL corners includes two-letter representation where the first letter refers to NMOS device corner and the second letter is related to PMOS device corner. So this convention has got 5 typical corners known as NN (Normal-Normal) or TT (Typical Typical), FF (Fast-Fast), SS (Slow-Slow), FS (Fast-Slow) and SF (Slow-Fast) that exhibit the mobility of carriers which is greater and lower than the normal specified conditions respectively. Therefore a corner represented as FS, for example, indicates a fast NMOS device and a slow PMOS device.

Amongst five corners, the first three corners (NN, FF and SS) are known as even corners since both the types of devices are evenly or simultaneously affected and do not adversely

affect the total functionality of the digital circuit design. As a result the device can operate at faster or slower clock frequencies. The remaining two corners, typically FS and SF, are called "skewed" corners which are the primary cause for the concerning issues in the operating zone because in case of FS corner, the switching activity of NMOS is much faster than that of PMOS device and in case of SF corner, the switching activity of NMOS is much slower than that of PMOS device. Therefore this kind of un-even variation may lead to considerable distortion in the normal operation of circuit design.

BEL corner:

This deals with the effect of PVT (Process-Voltage-Temperature) variations on onchip interconnections.

In this chapter, the process corner analysis is carried out for the proposed domino logic circuit techniques with increased fan-in by subjecting them to distinct ambient conditions and the corresponding observations are tabulated from which the conclusions are drawn.

3.5 Novel high-performance noise tolerant domino logic circuit techniques

3.5.1 Wide fan-in domino OR gate with proposed technique-1

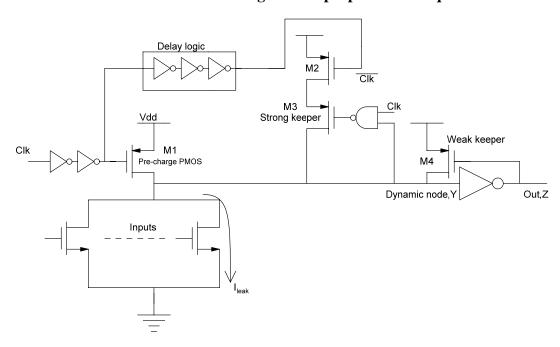


Fig. 3.21 Wide fan-in domino OR gate with proposed technique-1

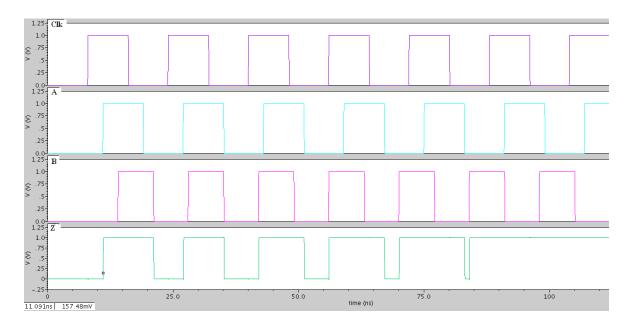


Fig. 3.22 Domino 2-input OR gate with proposed technique-1 simulation

The proposed technique-1 assures significantly high performance in terms of noise tolerance in domino logic gates, which takes the benefits of the utilization of static NAND gate along with the two PMOS series connected keepers as shown in Fig. 3.21 through the delayed logic network at the dynamic node in a feedback manner. The novel scheme is implemented in a foot-less domino technique as it possesses high performance than footed technique since the discharging mechanism of dynamic node in footless technique is faster than that of the footed technique. M1 is the pre-charge PMOS device to charge the dynamic node up to Vdd in clock's pre-charge phase. M4 is the weak keeper transistor connected in feedback manner through domino output node-Z. M2 and M3 are two series connected PMOS devices which are triggered by the status of clock signal and NAND gate response. A static NAND gate is selected in the process of designing the conditional keeper network at the dynamic node which plays a crucial role in stabilizing the dynamic node in terms of noise marginal levels. The analysis is carried out by implementing the proposed technique to a 2-input domino OR gate footless scheme and measured all the design parameters. Output 'Z' does not change in Pre-charge mode as long as all the inputs are not providing the discharge path. If any discharging path is available through the inputs when they are at logic high level, then only 'Z' discharges. And this status is continuing until the next pre-charge phase of clock signal arrives. Hence output 'Z' is changing when Clk='0' w.r.t A and B values. The simulation result of domino 2-input OR gate with proposed scheme-1 is given in Fig. 3.22.

During pre-charge phase:

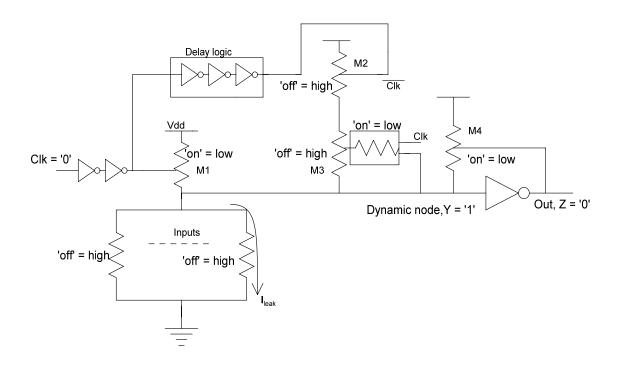


Fig. 3.23 Pre-charge operation of proposed technique-1

When clock becomes low, transistor M1 turns on providing a conducting path from supply rail to dynamic node and hence dynamic node Y gets charged to Vdd. The conditional keeper network as shown in Fig. 3.23 turns off and is not providing any conducting path from Vdd to dynamic node as M3 is turned off due to output response of NAND gate which is driven by clock line and dynamic node Y. As dynamic node Y takes a transition from low-to-high with clock signal at logic zero in pre-charge phase, the NAND gate produces logic 1 which makes M3 off and thereby making conditional keeper network completely turned off. M4 is the conventional PMOS keeper connected in feedback manner between nodes Y and Z respectively which is also providing an alternative conducting path from Vdd to dynamic node Y in pre-charge phase through which any leakages in pull-down network if exist can be managed. The flow of inevitable subthreshold leakage currents through pull-down network is represented in the Fig. 3.23. These leakages will cause a conducting path from dynamic node to ground through which

the dynamic node is forced to be wrongly discharged. This undesirable effect is pacified in pre-charge phase efficiently with the help of keeper device M4.

During evaluation phase:

Case (1): Pull-down network (PDN) is off

This analysis starts with an initial assumption that pull-down network is off when clock takes a transition from 0 to 1. Therefore in this case the node Y is supposed to maintain its pre-charged value since there is no discharging path available as it is assumed that pull-down network is completely off. M1 is now turned off disconnecting the conducting path from supply rail to dynamic node as clock takes a transition from low-to-high. Looking at the status of conditional keeper, from the Fig. 3.24 it is clear that M3 is turned on by NAND gate response which is producing logic zero output with its two active high inputs from node Y and clock line respectively. M2 is also turning on driven by inverted clock signal. Hence both M2 and M3 are providing a conducting path from supply rail to dynamic node which is required for managing any leakages in pull-down network. M4 is also providing a conducting path from bias rail to dynamic node as it is on since node Z is at logic 0. As long as the pull-down network is off during the clock's evaluation period, always strong one at dynamic node and strong zero at domino output are being produced which are highly desired phenomena by the conditional keeper network.

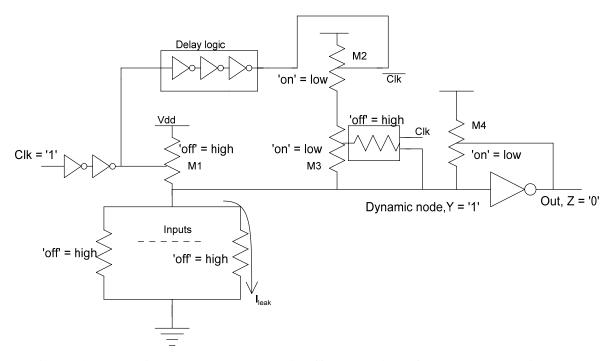


Fig. 3.24 Evaluation phase when PDN is off - operation of proposed technique-1

Case (2): Pull-down network (PDN) is on

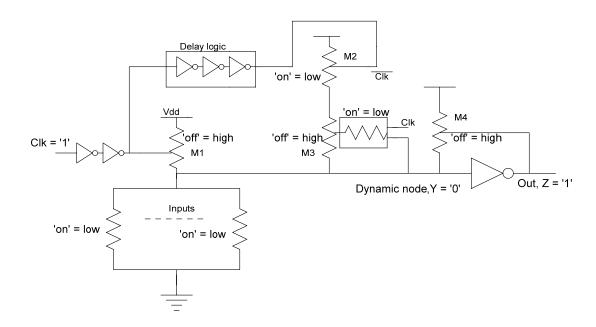


Fig. 3.25 Evaluation phase when PDN is on - operation of proposed technique-1

Now let us consider the transitions occurring at the input nodes of pull-down devices during evaluation phase. From the Fig. 3.25 it is evident that if one or more pull-down devices are turned on, then there is a discharging path provided by evaluation network that

discharges the dynamic node completely. Once the discharging process happens, then the NAND gate is turned to be on that makes M3 off and there by disconnecting the conducting path from supply rail to dynamic node through conditional keeper network despite the on status of PMOS transistor M2 which is of no use. M4 is also becoming off as the domino output switches from low-to-high state. Therefore during Pre-charge phase the dynamic node and domino node are producing strong logic levels without any degradation in the strength of voltage at the corresponding nodes. Similarly during the evaluation period, when the pull-down network is off then the dynamic node is efficiently maintaining its strong one logic level against the leakages in pull-down network and also when the pull-down network is on, then the complete discharging phenomenon takes place producing strong zero. In both the cases the proposed scheme is expeditiously giving the desired outcome by overcoming the sub-threshold leakages in pull-down network.

3.5.2 Wide fan-in domino OR gate with proposed technique-2

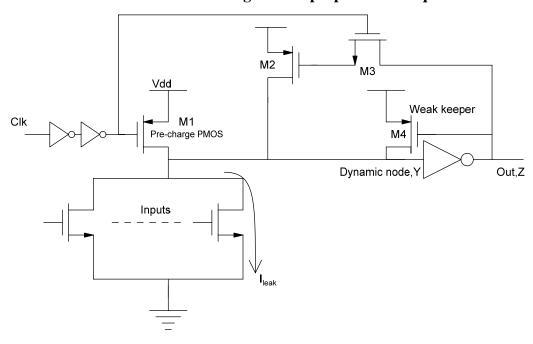


Fig. 3.26 Wide fan-in domino OR gate with proposed technique-2

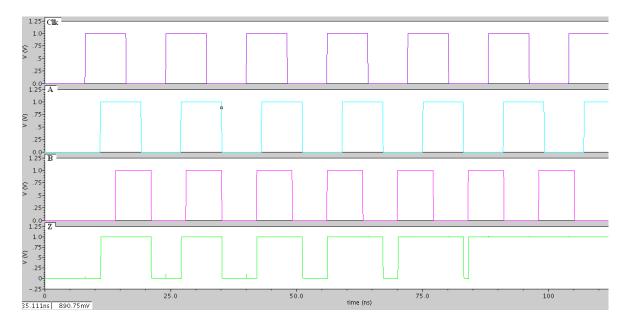


Fig. 3.27 Domino 2-input OR gate with proposed technique-2 simulation

The proposed technique-2 ensures significantly high performance in terms of noise tolerance in domino logic gates. Fig. 3.26 shows the implementation of 2-input domino OR gate with proposed technique-2. The novel scheme is implemented in a foot-less domino technique as it possesses high performance than footed technique since the discharging mechanism of dynamic node in footless technique is faster than that of the footed technique. M1 is the pre-charge PMOS device to charge the dynamic node up to Vdd in clock's pre-charge phase. M4 is the weak keeper transistor connected in feedback manner through domino output node-Z. M2 and M3 are two transistors, PMOS and NMOS respectively, which constitute the conditional keeper network at dynamic node. M2 is triggered by M3 which is drained through domino output node Z and driven by clock signal. This proposed circuit technique has conditional keeper network which comprises of two transistors, M2 and M3 only where as proposed-1 possesses two transistors along with static NAND gate. The process of designing the conditional keeper network at the dynamic node plays an important role in stabilizing the dynamic node in terms of noise marginal levels. The analysis is carried out by implementing the proposed technique to a 2-input domino OR gate footless scheme and measured all the design parameters. The simulation result of domino 2-input OR gate with proposed scheme-2 is given in Fig. 3.27.

During pre-charge phase:

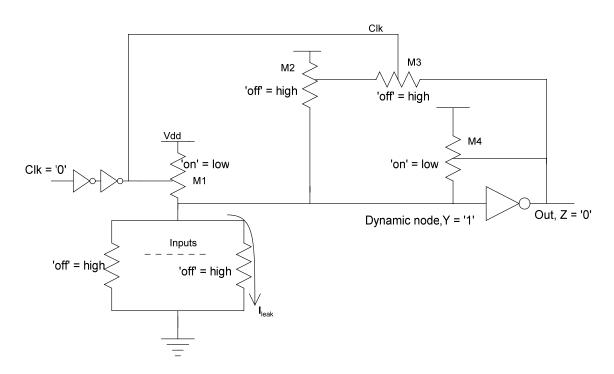


Fig. 3.28 Pre-charge operation of proposed technique-2

When clock becomes low, the pre-charge PMOS device M1 turns on providing a conducting path from supply rail to dynamic node and hence dynamic node Y gets charged to Vdd. The conditional keeper network as shown in Fig. 3.28 turns off and is not providing any conducting path from supply rail to dynamic node Y as transistor M3, driven by clock signal, is turned off which cannot turn M2 on. As dynamic node Y takes a transition from low-to-high with clock signal at logic zero in pre-charge phase, the dynamic node produces logic 1 which is inverter at domino output Z. M4 is the another PMOS keeper device which charges the dynamic node up to Vdd during pre-charge phase and also through this keeper transistor the leakage voltage drop if exists at dynamic node due to flow of inevitable leakage currents in pull-down network can be compensated. M4 is the conventional PMOS keeper connected in feedback manner between nodes Y and Z respectively. The flow of sub-threshold leakage currents through pull-down network is represented in the Fig. 3.28. These leakages will cause a conducting path from dynamic node to ground through which the dynamic node is forced to be wrongly discharged. This undesirable effect is alleviated in pre-charge phase effectively with the help of keeper device M4.

During evaluation phase:

Case (1): Pull-down network (PDN) is off

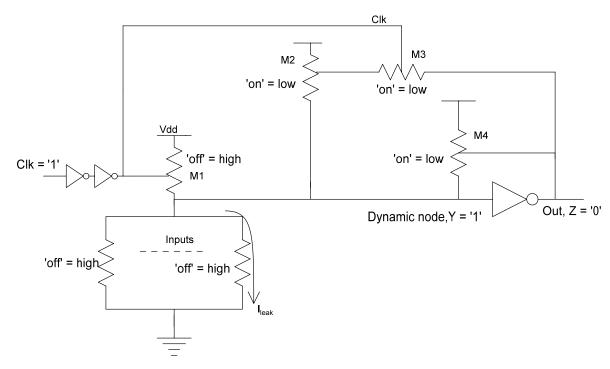


Fig. 3.29 Evaluation phase when PDN is off - operation of proposed technique-1

This analysis starts with an assumption that pull-down network is initially off when clock takes a transition from 0 to 1. Therefore in this case the node Y is supposed to maintain its pre-charged value since there is no discharging path available as it is assumed that pull-down network is completely off. M1 is now turned off disconnecting the conducting path from supply rail to dynamic node as clock takes a transition from low-to-high. Considering the status of conditional keeper, from the Fig. 3.29, it is clear that M3 is turned on by clock signal which in return will turn M2 on by passing the domino output because a logic low activates PMOS transistor. Also M4 is as usual charging the dynamic node since node Z is at logic low that makes M4 on. Hence both M2 and M4 are providing low resistance or conducting paths from supply rail to dynamic node which are required for managing any leakages in pull-down network. Therefore as long as the pull-down network is off during the clock's evaluation period, always strong one at dynamic node and strong zero at domino output are being produced which are highly desired phenomena by the conditional keeper network.

Case (2): Pull-down network (PDN) is on

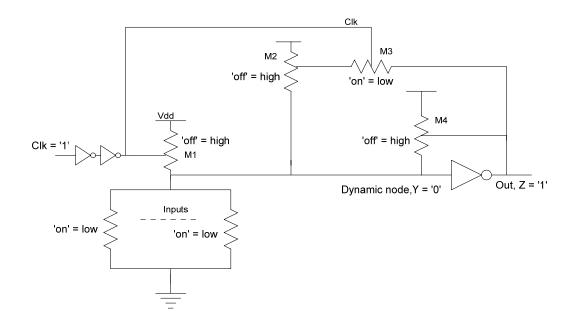


Fig. 3.30 Evaluation phase when PDN is on - operation of proposed technique-2

Now let us consider the transitions occurring at the input nodes of pull-down devices during evaluation phase. From the Fig. 3.30, it is clear that if one or more pull-down devices are turned on, then there is a discharging path provided by pull-down network that discharges the dynamic node completely. Once the discharging process takes place, then the dynamic node is turned to be low which in turn produces a strong one at domino output that will again result in turning the transistor M4 off. As domino output node Z takes a transition from low-to-high, this logic one will be passed by the transistor M3, which was already on, to M2 that results in turning it off. This clearly implies that there exists no low resistance path or conducting path from bias rail to dynamic node. Therefore during precharge phase of clock signal, the dynamic node and domino node are producing strong logic levels without any deterioration in the strength of voltage levels at the corresponding nodes. Similarly during the evaluation period, when the pull-down network is off then the dynamic node is efficiently maintaining its strong one logic level against the leakages existing in pull-down network and also when the pull-down network is on, then the complete discharging phenomenon takes place producing strong zero. In both the cases the proposed scheme is expeditiously giving the required result by mitigating the impact of the sub-threshold leakages in pull-down network.

3.5.3 Wide fan-in domino OR gate with proposed technique-3

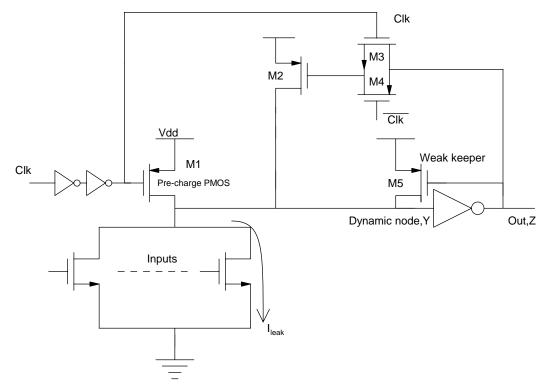


Fig. 3.31 Wide fan-in domino OR gate with proposed technique-3

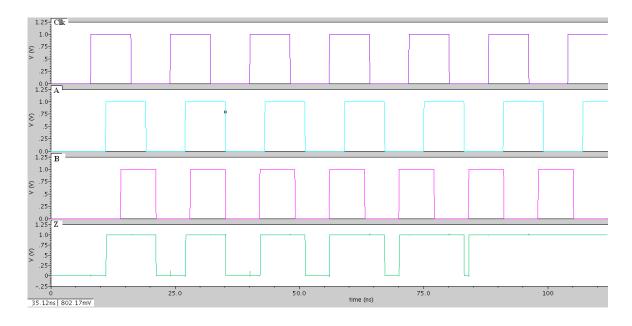


Fig. 3.32 Domino 2-input OR gate with proposed technique-3 simulation

The proposed technique-3 assures significantly very high performance in terms of noise robustness in domino logic gates which takes the benefits of the utilization of transmission gate logic circuit at conditional keeper network along with the PMOS connected keeper as shown in Fig. 3.31 at the dynamic node in a feedback manner. The novel technique is implemented in a foot-less domino technique as it possesses high performance than footed technique since the discharging mechanism of dynamic node in footless technique is faster than that of the footed technique. M1 is the pre-charge PMOS device that pre-charges the dynamic node up to Vdd in clock's pre-charge phase. M5 is the weak keeper transistor connected in feedback manner through domino output node-Z. M3 and M4 are two parallel connected NMOS and PMOS transistors forming the basic transmission logic gate circuit which is triggered by the status of a control signal clock at the two gate terminals of M3 and M4 respectively. Transistor M2 is driven by output of transmission gate. The concept of utilizing the benefits of selecting the transmission logic gate in the process of designing the conditional keeper network at the dynamic node plays a vital role in stabilizing the dynamic node in terms of noise marginal levels and robustness. The analysis is carried out by implementing the proposed technique-3 to a 2input domino OR gate footless scheme and measured all the design parameters. The simulation result of domino 2-input OR gate with proposed scheme-3 is given in Fig. 3.32. The transmission logic gate circuit used at conditional keeper network is shown in Fig. 3.33.

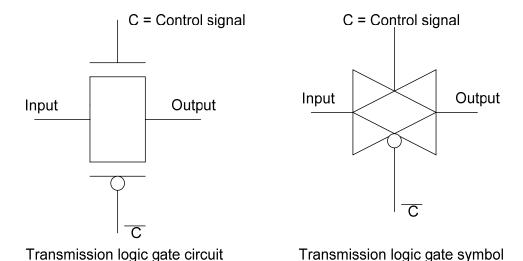


Fig. 3.33 A Transmission logic gate circuit

Transmission logic gate circuit:

- ➤ It is built on the basis of the complementary properties of both NMOS & PMOS transistors.
- ➤ It combines both device flavors by replacing a PMOS transistor in parallel with NMOS transistor as shown in Fig.3.41.

Transmission logic gate circuit is acting as a bi-directional switch which is driven by a control signal C, through the gate terminals of MOSFET devices. The functioning is running in two phases of control element which is of pulse signal.

• If C= '1', then both the transistors are on and as a result the input signal is passed to output node through them.

Therefore, Output=Input, if C= '1'.

• If C= '0', then both the devices are off which indicates that they are in cut-off mode that results in an open circuit between the Input and Output nodes.

Advantages of using transmission logic gate circuit:

- It enables rail-to-rail swing although it requires 2-transistors with control signal.
- > Complex gates can efficiently be built with this logic circuit.

The implementation of 2:1 Multiplexer using transmission logic gate is shown below.

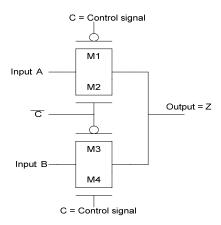


Fig. 3.34 A 2:1 Multiplexer using transmission logic gate circuit

When C='0', then transistors M1 and M2 will turn on and as a result output Z is taking input signal A while on the other hand if C='1', then transistors M3 and M4 will turn on and so they are passing input signal B to the output node Z.

The analysis of proposed technique-3 is carried out below.

During pre-charge phase:

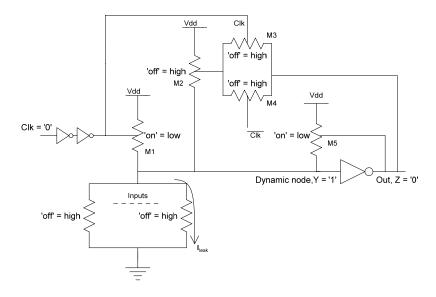


Fig. 3.35 Pre-charge operation of proposed technique-3

It starts when clock becomes low. The transistor M1 turns on providing a low resistance or conducting path from supply rail to dynamic node and hence dynamic node Y gets charged to Vdd. The conditional keeper network as shown in Fig. 3.35 turns off and is not providing any conducting path from Vdd to dynamic node as M2 is turned off due to output response of transmission logic gate circuit which is driven by clock line with its two inverted control signals at the corresponding gate terminals of M3 and M4 respectively. As dynamic node Y takes a transition from low-to-high with clock signal at logic zero in precharge phase, the transmission logic gate is still off since M3 and M4 are turned off by clock signal that results in turning off transistor M2. Therefore the conducting path through M2 is disrupted. M5 is the conventional PMOS keeper connected in feedback manner between nodes Y and Z respectively which is also providing an alternative conducting path from Vdd to dynamic node Y in pre-charge phase through which any leakages in pulldown network if exist can be managed. The flow of inevitable sub-threshold leakage currents through pull-down network is shown in the Fig. 3.35. These leakages will cause a conducting path from dynamic node to ground through which the dynamic node is forced to be wrongly discharged. This undesirable effect is pacified in pre-charge phase efficiently with the help of keeper device M5.

During evaluation phase:

Case (1): Pull-down network (PDN) is off

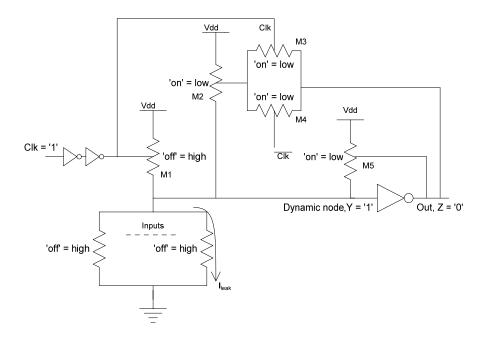


Fig. 3.36 Evaluation phase when PDN is off - operation of proposed technique-3

This analysis starts with an initial assumption that pull-down network is off when clock takes a transition from 0 to 1. Therefore in this phase the node Y is supposed to maintain its pre-charged value since there is no discharging path available as it is assumed that pull-down network is completely off. M1 is no longer connected to Vdd. Thus it is turned off disconnecting the conducting path from supply rail to dynamic node as clock takes a transition from low-to-high. Considering the status of conditional keeper, from the Fig. 3.36 it is clear that M3 and M4 are turned on by gate control signal clock which can make transmission gate circuit on that results in passing strong zero from domino node Z to transistor M2. Hence a low resistance or conducting path is provided by M2 as it is PMOS device that turns on with active low gate input. Also conventional keeper M5 is charging the dynamic node to supply voltage. Therefore both M2 and M5 are providing conducting paths from supply rail to dynamic node which are required for managing any leakages in pull-down network. As long as the pull-down network is off during the clock's evaluation period, always strong one at dynamic node and strong zero at domino output are being produced which are highly desired phenomena by the conditional keeper network.

Case (2): Pull-down network (PDN) is on

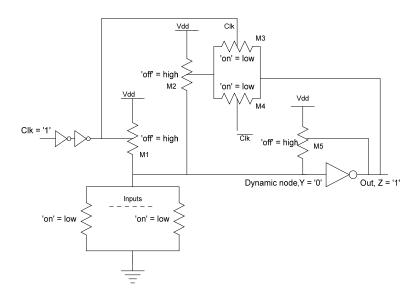


Fig. 3.37 Evaluation phase when PDN is on - operation of proposed technique-3

Now let us consider the transitions occurring at the input nodes of pull-down devices during evaluation phase. From the Fig. 3.37, it is evident that if one or more pull-down devices are turned on, then there exists a discharging path provided by pull-down network that discharges the dynamic node completely. Once the discharging process takes place, then the dynamic node is turned to be low which in turn produces a strong one at domino output that will again result in turning the transistor M5 off. As domino output node Z takes a transition from low-to-high, this logic one will be passed by transmission gate, which was already on, to M2 that results in turning it off. This clearly implies that neither M2 nor M5 is providing a low resistance path or conducting path from bias rail to dynamic node. Therefore during pre-charge phase of clock signal, the dynamic node and domino node are producing strong logic levels without any deterioration in the strength of voltage levels at the corresponding nodes. Similarly during the evaluation period, when the pulldown network is off then the dynamic node is efficiently maintaining its strong one logic level against the leakages existing in pull-down network and also when the pull-down network is on, then the complete discharging phenomenon takes place producing strong zero. In both the cases the proposed scheme is expeditiously giving the required result by mitigating the impact of the sub-threshold leakages in pull-down network.

3.6 Simulation results and discussion

Table 3.1 Comparison of typical power parameters and power-delay-product for standard and proposed domino logic techniques

without keeper Footless with keeper 9.373E-6 2.513E-8 9.398E-6 37.44E-12 350.92-18 Footed without keeper 0.203E-6 2.750E-8 0.230E-6 47.24E-12 9.58E-18 Footed with keeper 0.323E-6 2.754E-8 0.351E-6 72.13E-12 23.29E-18 Pooted with keeper 0.323E-6 2.739E-8 10.569E-6 113.99E-12 1201.68E-18 Replicated evaluation 11.461E-6 4.207E-8 11.503E-6 56.42E-12 646.62E-18 Pooted with keeper 16.598E-6 144.442E-8 18.042E-6 40.98E-12 680.18E-18 Replicated evaluation 16.598E-6 144.442E-8 18.042E-6 40.98E-12 680.18E-18 Clock delayed footed 9.118E-3 9.112E-3 18.230E-3 37.63E-12 343.11E-15 Single keeper Clock delayed gull keeper 9.828E-6 57.50E-12 560.74E-18 Skew tolerant high speed 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 Proposed - 1 9.910E-6 6.8E-8 9.978E-6	Domino logic technique	Dynamic power (in W)	Leakage or static power (in W)	Total power (in W)	Total propagation delay, T _p	Power- delay- product (in W-sec)
keeper Coted without keeper 0.203E-6 2.750E-8 0.230E-6 47.24E-12 9.58E-18 Footed with keeper 0.323E-6 2.754E-8 0.351E-6 72.13E-12 23.29E-18 Diode footed 10.542E-6 2.739E-8 10.569E-6 113.99E-12 1201.68E-13 Replicated evaluation 11.461E-6 4.207E-8 11.503E-6 56.42E-12 646.62E-18 Dynamic node footed 16.598E-6 144.442E-8 18.042E-6 40.98E-12 680.18E-18 Clock delayed single keeper 9.118E-3 9.112E-3 18.230E-3 37.63E-12 343.11E-15 Clock delayed dual keeper 9.752E-6 7.6E-8 9.828E-6 57.50E-12 560.74E-18 Skew tolerant high speed 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		9.292E-6	2.098E-8	9.312E-6	21.53E-12	200.05E-18
keeper 0.323E-6 2.754E-8 0.351E-6 72.13E-12 23.29E-18 Diode footed 10.542E-6 2.739E-8 10.569E-6 113.99E-12 1201.68E-13 Replicated evaluation 11.461E-6 4.207E-8 11.503E-6 56.42E-12 646.62E-18 Dynamic node footed 16.598E-6 144.442E-8 18.042E-6 40.98E-12 680.18E-18 Clock delayed single keeper 9.118E-3 9.112E-3 18.230E-3 37.63E-12 343.11E-15 Clock delayed dual keeper 9.752E-6 7.6E-8 9.828E-6 57.50E-12 560.74E-18 Skew tolerant high speed 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		9.373E-6	2.513E-8	9.398E-6	37.44E-12	350.92-18
keeper Diode footed 10.542E-6 2.739E-8 10.569E-6 113.99E-12 1201.68E-13 Replicated evaluation 11.461E-6 4.207E-8 11.503E-6 56.42E-12 646.62E-18 Dynamic node footed 16.598E-6 144.442E-8 18.042E-6 40.98E-12 680.18E-18 Clock delayed single keeper 9.118E-3 9.112E-3 18.230E-3 37.63E-12 343.11E-15 Clock delayed dual keeper 9.752E-6 7.6E-8 9.828E-6 57.50E-12 560.74E-18 Skew tolerant high speed 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		0.203E-6	2.750E-8	0.230E-6	47.24E-12	9.58E-18
Replicated evaluation 11.461E-6 4.207E-8 11.503E-6 56.42E-12 646.62E-18 Dynamic node footed 16.598E-6 144.442E-8 18.042E-6 40.98E-12 680.18E-18 Clock delayed single keeper 9.118E-3 9.112E-3 18.230E-3 37.63E-12 343.11E-15 Clock delayed dual keeper 9.752E-6 7.6E-8 9.828E-6 57.50E-12 560.74E-18 Skew tolerant high speed 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		0.323E-6	2.754E-8	0.351E-6	72.13E-12	23.29E-18
evaluation Dynamic node footed 16.598E-6 144.442E-8 18.042E-6 40.98E-12 680.18E-18 Clock delayed single keeper 9.118E-3 9.112E-3 18.230E-3 37.63E-12 343.11E-15 Clock delayed dual keeper 9.752E-6 7.6E-8 9.828E-6 57.50E-12 560.74E-18 Skew tolerant high speed 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 Proposed - 1 9.910E-6 4930.2E-8 80.004E-6 393.76E-12 12.09E-15 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18	Diode footed	10.542E-6	2.739E-8	10.569E-6	113.99E-12	1201.68E-18
footed 9.118E-3 9.112E-3 18.230E-3 37.63E-12 343.11E-15 Clock delayed single keeper 9.752E-6 7.6E-8 9.828E-6 57.50E-12 560.74E-18 Clock delayed dual keeper 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 Skew tolerant high speed 9.842E-6 4930.2E-8 80.004E-6 393.76E-12 12.09E-15 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18	1	11.461E-6	4.207E-8	11.503E-6	56.42E-12	646.62E-18
single keeper Clock delayed dual keeper 9.752E-6 7.6E-8 9.828E-6 57.50E-12 560.74E-18 Skew tolerant high speed 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 SFEG 30.704E-6 4930.2E-8 80.004E-6 393.76E-12 12.09E-15 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		16.598E-6	144.442E-8	18.042E-6	40.98E-12	680.18E-18
dual keeper 9.842E-6 5.1E-8 9.893E-6 61.45E-12 604.79E-18 SFEG 30.704E-6 4930.2E-8 80.004E-6 393.76E-12 12.09E-15 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		9.118E-3	9.112E-3	18.230E-3	37.63E-12	343.11E-15
high speed 4930.2E-8 80.004E-6 393.76E-12 12.09E-15 Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		9.752E-6	7.6E-8	9.828E-6	57.50E-12	560.74E-18
Proposed - 1 9.910E-6 6.8E-8 9.978E-6 59.69E-12 591.52E-18 Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18		9.842E-6	5.1E-8	9.893E-6	61.45E-12	604.79E-18
Proposed - 2 9.747E-6 5.1E-8 9.799E-6 55.58E-12 541.73E-18	SFEG	30.704E-6	4930.2E-8	80.004E-6	393.76E-12	12.09E-15
	Proposed - 1	9.910E-6	6.8E-8	9.978E-6	59.69E-12	591.52E-18
Proposed - 3 9.702E-6 5.0E-8 9.753E-6 58.58E-12 568.34E-18	Proposed - 2	9.747E-6	5.1E-8	9.799E-6	55.58E-12	541.73E-18
	Proposed - 3	9.702E-6	5.0E-8	9.753E-6	58.58E-12	568.34E-18

Proposed circuits exhibit high degree of noise robustness in terms of UNG and ANTE at discrete process corners. This was achieved because of the inclusion of novel conditional keeper circuits at dynamic node in order to stabilize it and also to control leakage current phenomenon. Thus, PDP is also quite high compared to existing ones since there was a trade-off between these constraint parameters as it was not possible to completely nullify

the leakage issue by maintaining high degree of noise robustness. One parameter will have to suffer while fabricating the other one, so we just optimized this trade-off by keeping the PDP at nominal and tolerable level while improving the noise robustness.

Table 3.2 UNG and ANTE comparison of standard and proposed domino logic techniques with Fan-in 2

	techniques with Fan-in 2			
Domino logic technique	UNG (in V)	ANTE (in V²*pico-sec)		
Foot-less without keeper	373.1E-3	3.34		
Foot-less with keeper	696.7E-3	6.53		
Footed without keeper	402.5E-3	3.58		
Footed with keeper	895E-3	8.11		
Diode footed	468.7E-3	4.32		
Replicated evaluation	432E-3	4.10		
Dynamic node footed	458.7E-3	4.22		
Clock delayed single keeper	604.2E-3	5.7		
Clock delayed dual keeper	752.2E-3	6.8		
Skew tolerant high speed	893E-3	7.88		
SFEG	457.2E-3	4.18		
Proposed - 1	776.48E-3	7.004		
Proposed - 2	968.55E-3	8.84		
Proposed - 3	984.443E-3	9.24		

The simulations are done for all the benchmark circuits and proposed techniques in CMOS 90nm process technology with 1V power supply as bias voltage. The primary design parameters like dynamic power dissipation, propagation delay, UNG, ANTE, leakages currents, static power consumption and power-delay-product are calculated for different ambient conditions and compared with standard results.

UNG against Process corner analysis for variable Fan-in:

Table 3.3 UNG comparison of proposed domino logic techniques with Fan-in=2 at different Process Corner analysis

Domino logic technique	UNG (in V)				
	NN	FF	SS	FS	SF
Proposed-1	776.48E-3	768.83E-3	842E-3	636.33E-3	
Proposed-2	968.55E-3	918.9E-3		728.64E-3	
Proposed-3	984.443E-3	928.55E-3		739.34E-3	

Table 3.4 UNG comparison of proposed domino logic techniques with Fan-in=4 at different Process Corner analysis

Domino logic technique	UNG (in V)				
	NN	FF	SS	FS	SF
Proposed-1	800E-3	788.592E-3	875E-3	657.51E-3	
Proposed-2	992.059E-3	934.75E-3		746.70E-3	
Proposed-3		946.2E-3		758E-3	

Table 3.5 UNG comparison of proposed domino logic techniques with Fan-in=8 at different Process Corner analysis

Domino logic technique	UNG (in V)				
	NN	FF	SS	FS	SF
Proposed-1	845E-3	828.5E-3	952E-3	697.53E-3	
Proposed-2		965.33E-3		782.555E-3	
Proposed-3		976E-3		794E-3	

Table 3.6 UNG comparison of proposed domino logic techniques with Fan-in=16 at different Process Corner analysis

Domino logic technique	UNG (in V)						
teeninque	NN	FF	SS	FS	SF		
Proposed-1	941E-3	910.3E-3		775E-3			
Proposed-2				849.82E- 3			
Proposed-3				863E-3			

Table 3.7 UNG comparison of proposed domino logic techniques with Fan-in=32 at different Process Corner analysis

Domino logic		UNG (in V)						
technique	NN	FF	SS	FS	SF			
Proposed-1				951E-3				
Proposed-2				973E-3				
Proposed-3				984E-3				

ANTE against Fan-in for various Process corners:

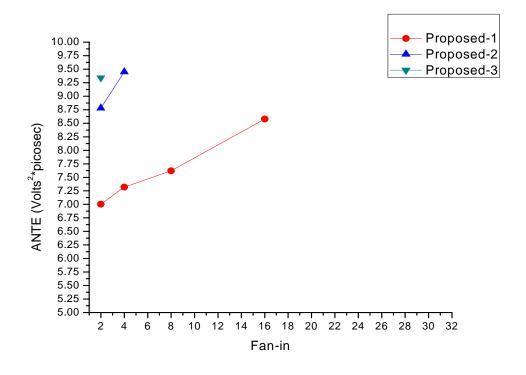


Fig. 3.38 ANTE Vs Fan-in for Process corner=NN

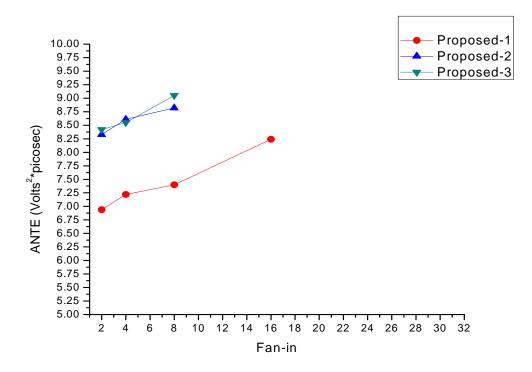


Fig. 3.39 ANTE Vs Fan-in for Process corner=FF

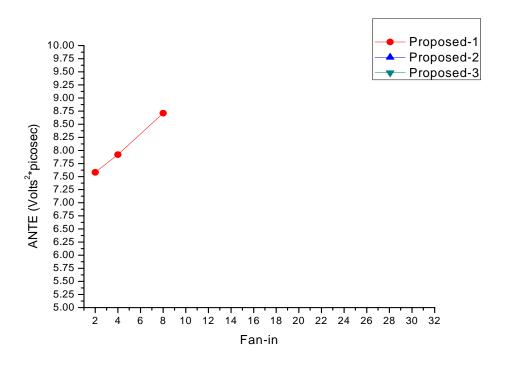


Fig. 3.40 ANTE Vs Fan-in for Process corner=SS

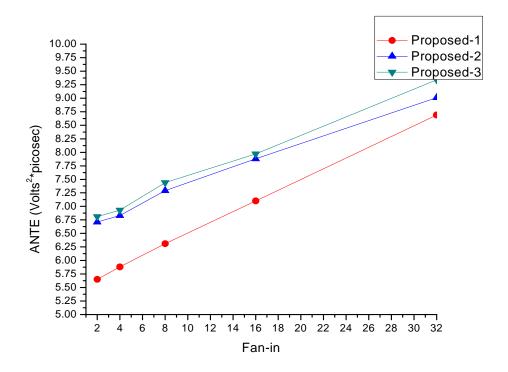


Fig. 3.41 ANTE Vs Fan-in for Process corner=FS

Discussion

Therefore the proposed domino logic circuit techniques are simulated at distinct ambient conditions by subjecting them to various process corners and thereby observations are tabulated. Comparisons are made with reference to benchmark circuit techniques and conclusions are drawn. From the tabulated results and graphical analysis, it is evident that the proposed domino logic circuit techniques are exhibiting high speed and high degree of robustness in terms of noise metric parameters like UNG, ANTE. Also in all the three proposed techniques, the UNG is almost closer to unity factor which is highly desired for greater noise tolerant circuits. ANTE is also following the same path in assuring high performance as UNG. Besides these noise metric parameters, PDP (power-delay-product) and total power consumption are considerably lower than those of few existing schemes.

In proposed technique-1, this has been achieved by employing the static NAND gate in the process of designing the conditional keeper network at dynamic node which possesses high noise margin. In proposed-2, a single NMOS is connected between domino node and keeper device in feed-back manner that passes the logic zero effectively. In

proposed-3 the transmission gate is adopted in conditional keeper network which stabilizes the dynamic node from leakages efficaciously, by passing both strong one and strong zero since it is a mixer of both NMOS and PMOS flavors.

Having analyzed the results from process corner analysis, it is clear that proposed circuit techniques assure high noise robustness with increased fan-in. There are limitations too for the functionality of proposed domino logic circuit techniques at various corners. Proposed techniques are functioning efficiently, at all the corners for lower fan-in circuits but with increased fan-in, the operating region of design has started becoming limited to few corners only which might be due to the reason that the circuits are being operated at lower bias supply voltage around 1 V. Thus increasing the bias voltage range would facilitate broadening the operating region of proposed domino techniques at all the process corners for wide fan-in circuits but while doing so the power consumption must be taken care off as it is directly proportional to square of supply voltage. Hence optimization through trade-off between supply voltage and process corners is necessary while improving the proposed domino logic circuit techniques.

Table 3.1 shows the comparison of typical power parameters such as dynamic power, static power, propagation delay, and power-delay-product for standard and proposed domino logic techniques. It is evident that the total power consumption and PDP of proposed circuits are considerably lowered when compared with benchmark circuits. Table 3.2 gives comparison of UNG and ANTE for standard and proposed domino logic techniques with fan-in 2 and it is clear that proposed circuits possess improved UNG and ANTE values.

Looking at the process-corner analysis, amongst five typical corners (NN, SS, FF, FS and SF), Normal-Normal corner is providing nominal switching threshold voltages for NMOS and PMOS devices as specified by EDA tool. In general, all the designed circuits will operate as per design specifications and according to user constraints at NN corner. Thus evaluation and estimation of overall performance of designed circuit by subjecting at NN process corner alone does not finish the task. In fact the circuit is required to be tested by subjecting at all the extreme corners. Having done that entire task, then only the performance of circuit can be judged. Always the SS corner is assuring increased noise

tolerance and which is also observed from tabulations. Highest UNG is recorded in SS corner since both NMOS and PMOS transistors are slow running devices which implied that their switching threshold voltages are high. Thus, the circuit is made less sensitive to noise glitches by these SS corner devices and hence the nature of being responsive to gate input noise glitches is gradually reducing, that in turn increases UNG. So, in comparison with other process corners, SS corner always exhibits highest UNG. Also in SF corner, as NMOS switching threshold voltage is higher and PMOS switching threshold voltage is lower, it also contributes increment in noise gain but not as efficient as SS corner because the PMOS threshold voltage is responding to noise impulses at gate inputs which makes the circuit more sensitive to noise glitches. Normally, in comparison with NN, FS and FF corners, circuit at SF corner exhibits better immunity towards noise. Having analyzed the proposed circuits at all the extreme corners, SF corner is not at all giving the response. This is because of the design of conditional keeper circuit which is suitable for high speed applications. Thus SF corner, which runs with slow NMOS and fast PMOS devices, is not suitable for proposed circuits. At FF process corner, both NMOS and PMOS devices are of very high speed with reduced switching threshold voltages and consume more power. Thus normal circuits usually become more sensitive to noise glitches at gate inputs due to their lower threshold voltages and as a result UNG is reduced. With increased fan-in, the UNG in other process corners is gradually lowering since more voltage is required to turn pulldown network on and discharging phenomenon is becoming slow. But proposed circuits are functioning at FF corner efficiently as they are designed for high speed applications. Normally, with increased fan-in, the UNG lowers but the proposed circuit techniques exhibit the greater noise robustness for wide fan-in also which made them suitable for high speed applications. The performance of circuit at FS process corner lies in between the corresponding performances at NN and FF corners. As NMOS possesses lower threshold voltage, it becomes more sensitive to noise. Despite slow PMOS device, the corner is lowering the noise immunity. From the tabulations, it is evident that proposed techniques exhibit high degree of robustness at FS corner too.

Having analyzed the results from process corner analysis, it is clear that proposed circuit techniques assure high noise robustness with increased fan-in despite few limitations at various corners. Proposed techniques are functioning efficiently, at all the

corners for lower fan-in circuits but with increased fan-in, the operating region of design has started becoming limited to few corners only which might be due to the reason that the circuits are being operated at lower bias voltage around 1 V. For example, from the simulations, it is clear that the UNG is getting limited at NN, FF and FS process corners for higher fan-in. Thus increasing the bias voltage range would facilitate broadening the operating region of proposed domino techniques at all the process corners for wide fan-in circuits but while doing so the power consumption must be taken care off as it is directly proportional to square of supply voltage. Hence optimization through trade-off between supply voltage and process corners is necessary while improving the proposed domino logic circuit techniques.

Therefore, the proposed circuits are designed for high speed applications with greater noise immunity. This is investigated from the tabulations of the analysis of UNG against process corner analysis for various fan-in circuits. Normally process corners SS and SF assure higher noise tolerance than FF and FS corners at which the circuits are becoming more sensitive and vulnerable to noise glitches. But in this chapter the proposed circuit techniques are exhibiting significantly improved noise tolerance even at FF and FS corners also. This is because of the design of conditional keeper network particularly the proposed-3 technique shows greater noise robustness than proposed-1 and proposed-2, since transmission gate is chosen at keeper network which passes both strong one and strong zero efficiently and stabilizes the leakages in pull-down network. Improved average noise threshold energy (ANTE) can also be observed from tabulations.

3.7 Conclusion

Therefore this chapter in section-1 gave general introduction to domino logic family.

Section-3.2 discussed standard benchmark domino logic circuit schemes followed by the analysis of their functionality with simulation results.

In section-3.3, the novel domino logic circuit techniques are proposed and analyzed their functionality in detail with equivalent circuit diagrams in all operating regions along with simulation results.

In section-3.4, the noise analysis is carried out which includes the need for robustness, various metric parameters for measuring noise immunity or robustness of domino circuits such as UNG, ANTE along with the method of calculations, various sources of noise in domino logic circuits and their role on operating region.

Section-3.5 gives the description of process corner analysis and various corners involved in it along with their significant role on the overall functionality of the designed domino logic circuit. Also the consequences of subjecting the device to the extreme corners with the boundary limitations are discussed.

Section-3.6 is the result section which showed the calculations and comparisons of all the parameters of standard benchmark circuits and proposed domino techniques. The primary design parameters such dynamic power, leakage or static power, total power, PDP (power-delay-product), UNG and ANTE for various fan-in circuits of existing and proposed techniques are measured. The comparisons along with graphical analysis are made and discussed the functionality with pros and cons.

The highest UNG of 984.4E-3 V and ANTE of 9.45E-12 V-Sec are exhibited by proposed-3 circuit with Fan-in=32 at FS process corner. Also, it is observed that the UNG of 992.05E-3 V and ANTE of 9.50E-12 V-Sec are showed by proposed-2 with Fan-in= 4 at NN process corner. Similarly, the proposed-1 circuit is exhibiting its highest UNG and ANTE of 952E-3 V and 8.75E-12 V-Sec respectively with Fan-in=8 at SS process corner. These are the conclusions drawn from tabulations and plots. Thus the proposed circuits show improved performance in terms of noise robustness over the existing bench mark circuits and the elaborated discussion on the comparison is also provided in discussion part.

CHAPTER 4 SIGNAL INTEGRITY ISSUES & MODIFIED CIRCUIT TECHNIQUES

4.1 Introduction

Power dissipation is one of the important parameters in the process of design of CMOS based VLSI circuits. Large power consumption affects battery life in durability and reliability. There are many sources for power dissipation amongst which main sources are load capacitor, short-circuit conducting path and leakage current. The load capacitor contributes definite amount of power dissipation at dynamic node while charging and discharging. When there exists a conducting path from supply rail to ground rail during the transition period of logic gates, then it leads to short-circuit power dissipation. Leakage power is due to the reverse-biased diode currents due to charge storage between drain and substrate or body terminals of MOSFET, and sub-threshold currents which flow when there is a phenomenon called carrier diffusion between source and drain terminals of 'off' state devices.

When a circuit is designed with equal rise time and fall times then the short-circuit power dissipation can effectively be minimized [10-20]. The predominant component of the power dissipation results from switching activity of the logic gate. Due to continuous down scaling of process technology especially in deep sub-micron regime, the feature size or overall dimension of the device is becoming smaller and thereby reducing load capacitances. Also scaling requires minimizing bias voltage and threshold voltage [120-122]. The voltage scaling is benefitted due to the quadratic relation between dynamic power consumption and supply voltage as the power varies linearly with square of the supply voltage but at the cost of drastic increase of gate delay in the operation of the circuit when bias voltage reaches the threshold voltage level in sub-threshold region [2], [116-122]. Thus to make the time delay parameter independent of supply voltage, the threshold voltage needs to be minimized. We know, that

Power = Voltage. Current

$$= \frac{Energy}{time} \tag{4.1}$$

Thus,
$$P_{Dynamic} = \frac{(Energy \ at \ dynamic \ node)}{time}$$

Energy at dynamic node corresponds to charge storage on load capacitor in the form of electric field which is governed by the following equation,

Energy =
$$(\frac{1}{2}) C_L V dd^2$$

= $SC_L V dd^2$

Where S = constant with value 0.5.

Therefore,
$$P_{Dynamic} = \frac{SC_LVdd^2}{tim e}$$

As frequency is rate of change of time, the factor $\frac{1}{time}$ converges to f.

Now,
$$P_{Dynamic} = SC_LVdd^2f$$

$$\Rightarrow P_{Dynamic} \propto Vdd^2$$
.

The relation between time delay and supply voltage is given by the following equation.

$$T_{delay} = \left(\frac{C_L V dd}{K (V dd - V_{th})^{\alpha}}\right) \tag{4.2}$$

Where $_{\alpha}$ = velocity saturation index

_atakes the value of 2 for long channel devices and 1.2 for a short channel device.

K = CMOS technology dependence parameter

Thus, lowering the threshold voltage makes time delay independent of supply voltage which is not possible practically as each device possesses certain threshold voltage. Furthermore, to stabilize the overall performance of CMOS logic circuits, the ratio of the supply voltage to threshold voltage must be at least 5 or above [3] which also assures increased noise margin and eliminates the so called hot-carrier effects in short-channel devices [4]. Reducing threshold voltage results in exponential rise of the sub-threshold leakage current [5]. The trade-off between the device threshold voltage and bias voltage for Intel microprocessor is discussed in [6]. In [9] it is demonstrated that the leakage power

is about 0.01% of the total switching power for 1- millimeter technology and increases to 10% for 0.1 millimeter technology. This implies a drastic increase in leakage power with the advancement of down scaling of process technology from genesis to genesis. Also it is estimated that within the few generations in future, the leakage power dissipation will become equal to the total switching or active power dissipation. Therefore, the efficient minimizing techniques for leakage power will become very crucial in the deep sub-micron regime. Despite going for new methodologies for reducing leakages, noise immunity of the circuit must also be considered as in deep sub-micron, the circuit becomes more prone to noise effects. The noise margin levels will get narrowed down with continuous down scaling which must be increased for better performance from the perspective of noise robustness. UNG and ANTE define the robustness of the clocked logic circuits and greater values of these parameters assure higher robustness of the circuits.

In this chapter, we propose a new leakage power reduction technique called domino lector technique. The rest of the chapter is organized as follows. Section 4.2 explains signal integrity issues. Section 4.3 describes the prior works related to leakage power reduction schemes and the lector technique. Modified lector domino scheme and dynamic node stabilizing technique are demonstrated in section 4.4. Simulation results along with discussion are presented in section 4.5 and in section 4.6 concluding remarks are made.

4.2 Signal integrity issues in clocked logic circuits

Noise is one of the important parameters that will be taken into consideration for the analysis of any digital logic circuit at different ambient conditions where the functionality may vary depending on the application and design requirement. Noise in digital circuits has got its own significant role to govern the functionality of the particular design altogether. There are different sources of noise in deep submicron regime [. Dynamic logic assures high performance when compared to static counterpart circuits. But there are few significant parametric considerations that must be taken into account for the proper functioning of dynamic logic circuits. These problems are referred as signal integrity issues in clocked logic circuits which include

- (1)Crosstalk,
- (2) Charge leakage currents,
- (3) Charge sharing,
- (4) Capacitive coupling,
- (5)Clock feed through and
- (6) Small variations of nominal supply voltage.

(1) Crosstalk Noise:

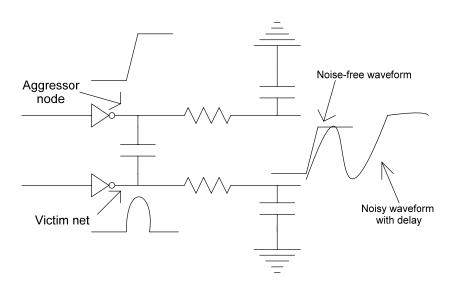


Fig. 4.1 Cross talk noise effect

Example of cross talk noise effect is shown in Fig. 4.1. It usually occurs on a wire which is associated with the switching action of neighboring wire. The switching wire is referred as the aggressor and the other wire is named as the victim. The reason for the occurrence is the phenomenon called capacitive coupling of the wires. Thus it is clear that it is not a random noise since it occurs only when aggressor wire's switching action happens. The switching action of output takes place in the evaluation phase wherein it is highly sensitive to the input. High input impedance is one of the most desired parameters in designing process of a logic circuit. But the circuit having been designed with high impedance of

output node is also associated with the drawback of becoming more susceptible to cross talk noise since the output node itself will make the circuit more prone to crosstalk effects. The relatively high output node impedance is the cause behind the cross-talk noise in the digital circuits. Hence the switching time of output and the phase of input wire at which it is sensitive to noise are noticed. Now this crosstalk noise can effectively be alleviated by properly laying out the aggressor and victim wires in such a way that there must not be any overlapping of evaluation phases. Sometimes CAD tools will solve this problem even after the layout is done by default. With the help of this technique the crosstalk noise, being the predominant noise source, can easily be eliminated for these kinds of circuits.

(2) Charge leakage currents:

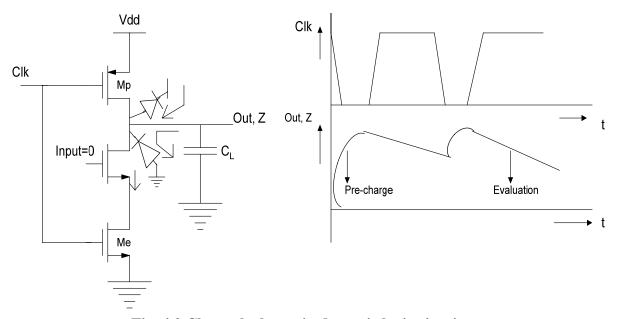


Fig. 4.2 Charge leakages in dynamic logic circuit

Domino logic circuits are highly sensitive to sub-threshold leakage currents. Despite their high speed of operation, they always suffer from high noise sensitivity through leakages which are inevitable in deep sub-micron regime. The dynamic gate operation primarily counts on the stored charge at the dynamic node by the load capacitor. In clock's pre-charge mode as the pre-charge PMOS provides conducting path from supply rail to dynamic node, the node will get charged to Vdd. Now once the evaluation phase commences, the output node should remain in its pre-charged value as long as the evaluation transistors are turned off. Yet, there is a voltage drop observed at this node

which is due to leakages eventually causing the malfunctioning of the operation of logic circuit. The main reason for this issue is the operation of devices in sub-threshold mode and continuous down scaling of process technology.

Charge leakage in dynamic logic circuits is given in Fig. 4.2. The two diodes shown in Fig. 4.2 are the reverse-biased. The charge stored on load capacitor, C_L will slowly discharge due to these diode leakage sources during evaluation period which deteriorates the strength of voltage level at dynamic node. Hence a minimal clock rate in the order of few kHz is highly desired for driving the dynamic circuits which in turn makes them un-favorable for the usage of these schemes for designs with low performance applications where there is no requirement of minimal clock rate. Also there is leakage current from pre-charge PMOS device due to the upper reverse bias diode and the subthreshold mode of operation. Like in cross-talk, the high impedance state of output also causes leakages during the clock's evaluate mode, when the evaluation network is turned off. Therefore, the leakage issue can be managed by minimizing the output impedance during the period of clock's evaluation. There are compensating techniques too, to overcome this phenomenon. Most commonly used scheme is employing a PMOS keeper at the dynamic node to replenish the charge lost from it which may, for better processing, need to be re-sized in order to alter its functional operation depending on the situational mode of operational requirement. This is usually done by employing a keeper transistor at the dynamic node whose purpose is to provide a conducting path from supply rail to dynamic node to compensate charge lost from it. Nevertheless there may be a chance of getting always a conducting or short circuit path as shown in Fig. 4.3 irrespective of status of pull-down network that results in flow of static currents and thereby causing increased static power dissipation in the circuit. Therefore to lessen this problem associated with keeper, the keeper, as shown in Fig. 4.4, is always connected in a feedback configuration.

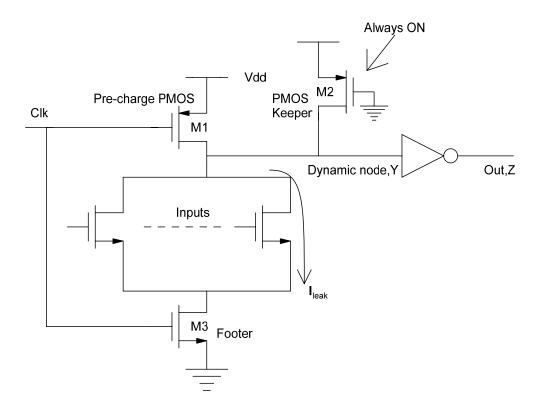


Fig. 4.3 Keeper with always ON configuration

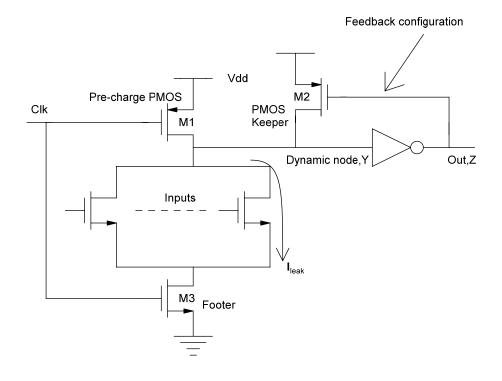


Fig. 4.4 Keeper with feedback configuration

Normally a small keeper is preferred as it refills the charge lost by providing the path from supply voltage rail, but as devices are continuously scaled down, the leakage current is increasing, and hence this small keeper, perhaps, might not be sufficient to compensate this drawback which in turn necessitates the usage of larger or a wide keeper. Therefore the usage of keeper along with proper sizing and optimization is necessary in designing the conditional keeper network at the dynamic node which is implemented in the proposed circuit schemes.

(3) Charge sharing:

This noise occurs because of sharing of stored charge at the dynamic node among the parasitic capacitances or junction capacitances of devices within the gate. Due to this, there is slight reduction in the strength of voltage at dynamic node that may result in erroneous output. Hence this problem needs to be eliminated in order to boost up circuit performance especially when it is used in a cascaded system of similar circuits on a giant network. The impact will be on the overall propagation delay and power dissipation. Domino logic finds it as one of the most common and inevitable signal integrity issues.

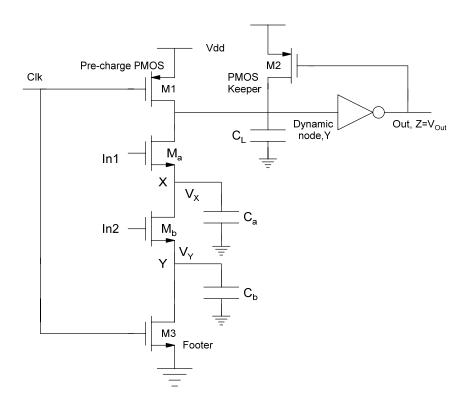


Fig. 4.5 Charge sharing analysis with 2-input domino AND gate

Let us consider a simple 2-input domino AND gate for the charge sharing analysis. Fig. 4.5 depicts the AND gate domino circuit with corresponding capacitances. During the clock's pre-charge period, the output node is charged to Vdd with the assumption that pull-down network is off and that the capacitances C_a and C_b are completely discharged. Now let us consider that input-In1 makes a transition from 0 to 1 while In2 remains at 0 during the evaluation period, turning the NMOS device M_a on. In this case there will be charge distribution of initially stored charge on load capacitor between the C_L and C_a which leads to a voltage drop in the output voltage that cannot be retrieved due to its dynamic nature. Similarly when In2 signal takes a transition from 0 to 1, again this will turn the transistor M_b on so that the total load capacitor will now get shared amongst all the three capacitors. This implies the dynamic node voltage is eventually shared amongst all the nodal capacitors when their corresponding devices are turned on along with load capacitor at dynamic node in the evaluation period of clock signal. Fig. 4.5 is used for the complete analysis of charge sharing phenomenon and also we extend this analysis to derive for a generalized case with 'p' number of nodal capacitors connected for a high fan-in AND gate logic circuit.

Charge sharing analysis is carried out here with the circuit shown in Fig. 4.5.

Here the analysis is carried out with the assumptions of defining the initial conditions as per following. $V_{Out}(t=0) = Vdd$ and $V_X(t=0) = 0$. Two possible cases must be taken in to account for the further analysis.

 $V_{th}(X)$ and $V_{th}(Y)$ are threshold voltages associated with nodes X and Y respectively.

 V_{Out} (final) is the output node voltage during the evaluation period of clock signal with all pull-down transistors set to 0.

Case (1):
$$\Delta V_{Out} < V_{th}$$

Charge at node X is given by $V_X=Vdd - V_{th}(X)$

$$=> C_L*(Vdd) = C_L*[V_{Out}(final)] + C_a*(Vdd - V_{th}(X))$$

$$=> Vdd = V_{Out}(final) + \frac{C_a}{C_L}(Vdd - V_{th}(X))$$

$$=> [V_{Out}(final) - Vdd] = \Delta V_{out}$$

$$= -\frac{C_a}{C_L}(Vdd - V_{th}(X))$$
(4.1)

Case (2): $\Delta V_{Out} > V_{th}$

Applying law of conservation of charge, we get the total charge distribution as follows.

$$[V_{Out}(final)]*(C_a+C_L) = Vdd*C_{L,}$$

$$=> V_{Out}(final)*C_a+V_{Out}(final)*C_L+Vdd*C_a-Vdd(C_a+C_L)=0$$

$$=> V_{Out}(final)*C_a+V_{Out}(final)*C_L+Vdd*C_a-Vdd*Ca-Vdd*C_L=0$$

$$=> V_{Out}(final)*[C_a+C_L] - Vdd*[C_a+C_L] = -Vdd*C_a$$

$$=> [V_{Out}(final) - Vdd]*[C_a+C_L] = -Vdd*C_a$$

$$=> [V_{Out}(final) - Vdd] = -Vdd*(\frac{C_a}{C_a + C_I})$$

$$\Rightarrow \Delta V_{Od} = - \text{Vdd}*(\frac{C_a}{C_a + C_I}). \tag{4.2}$$

Consider only C_L and C_a neglecting C_b .

Thus $Vx=[Vdd - V_{th}(X)]$, since the total supply voltage is now reduced by threshold voltage of M_a transistor at the node X.

Therefore, by applying the law of conservation of charge at the dynamic node we get,

Total charge at load capacitor = sum of distributed charges between C_L and C_a which is given by the following equation.

$$C_L*Vdd = C_L*V_{Out}(final) + C_a*V_X$$

Substituting the value of V_X in the above equation, we get

$$C_L*Vdd=C_L*V_{Out}(final) + C_a*(Vdd - V_{th}(X))$$

$$=> C_L*[Vdd - V_{Out}(final)] = C_a*[Vdd - V_{th}(X)]$$

$$\Rightarrow \frac{C_L}{C_a} = \frac{[Vdd - V_{th}(X)]}{[Vdd - V_{Out}(final)]}$$
(4.3)

Now let us consider the two capacitors C_a and C_b and analyze the charge sharing phenomenon.

$$V_X = Vdd - V_{th}(X)$$
 and

$$V_Y = V_X - V_{th}(Y)$$

Since the voltage at node Y is equal to supply voltage with the reduction factor by the threshold of $M_{\rm b}$

Applying the law of conservation of charge at the dynamic node we get,

Total charge stored at load capacitor = charge at load capacitor + charge at C_a + charge at C_b .

Mathematically it is governed by the following equation.

$$C_L*Vdd = C_L*V_{Out}(final) + C_a*V_X + C_b*V_Y$$

$$\Rightarrow C_L*Vdd = C_L*V_{Out}(final) + C_a*[Vdd - V_{th}(X)] + C_b*[V_X-V_{th}(Y)]$$

$$\Rightarrow C_L*[Vdd - V_{Out}(final)] = C_a*[Vdd - V_{th}(X)] + C_b*[V_X - V_{th}(Y)]$$

$$\Rightarrow C_L*[Vdd - V_{Out}(final)] = C_a*[Vdd - V_{th}(X)] + C_b[\{Vdd - V_{th}(X)\} - V_{th}(Y)]$$

$$\Rightarrow C_L*[Vdd - V_{Out}(final)] - C_a*[Vdd - V_{th}(X)] - C_b*[\{Vdd - V_{th}(X)\} - V_{th}(Y)] = 0$$

Charge sharing mechanism between two nodal capacitors C_a and C_b respectively is given by the following equation.

(4.4)

 C_L^* [Vdd - V_{Out} (final)] - C_a^* [Vdd - V_{th} (X)] - C_b^* [Vdd - V_{th} (X) - V_{th} (Y)]=0

$$C_a^*(Vx) = C_a^*(V_{XOut(final)}) + C_b^*(V_{Y)}$$

By substituting the value of V_Y in the above equation we get,

$$\Rightarrow C_a^*(V_X) = C_a^*(V_{XOut}(final)) + C_b^*[V_X - V_{th}(Y)]$$

$$\Rightarrow C_a^*[V_X - V_{XOut}(final)] = C_b^*[V_X - V_{th}(Y)]$$

$$\Rightarrow \frac{C_a}{C_b} = \frac{V_X - V_{th}(Y)}{V_Y - V_{YOut}(final)}$$

$$(4.5)$$

Therefore the total charge sharing phenomenon in the given circuit shown in Fig. 4.6 with load capacitor C_L and two nodal capacitors C_a and C_b is governed by

$$C_L *[Vdd - V_{Out}(final)] - C_a *[Vdd - V_{th}(X)] - C_b *[Vdd - V_{th}(X) - V_{th}(Y)] = 0$$
 (4.6)

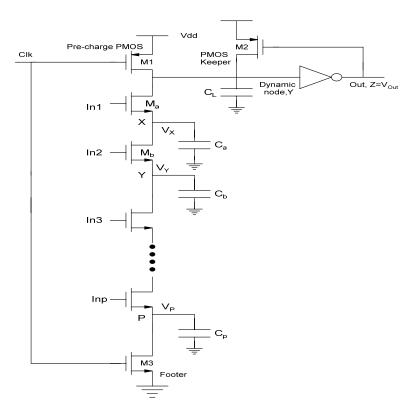


Fig. 4.6 Charge sharing analysis with wide fan-in (with fan-in=P) domino AND gate

This can be generalized to a wide fan-in domino AND gate circuit with 'P' number of nodal capacitors and the total charge sharing phenomenon is expressed by the following equation. The wide fan-in AND gate with fan-in 'P' is taken as shown in the Fig. 4.6. There are nodal capacitors C_a , C_b , C_c , C_dup to C_p along with a load capacitor C_L at the dynamic node.

Using the equation (4.6) we can derive the charge sharing mechanism for this wide fan-in AND gate circuit.

$$C_{L} *[Vdd - V_{Out}(final)] - C_{a} *[Vdd - V_{th}(X)] - C_{b} *[Vdd - V_{th}(X) - V_{th}(Y)] - C_{c} *[Vdd - V_{th}(X) - V_{th}(Y)] - V_{th}(Y) -$$

The solution for this charge sharing phenomenon is to charge the critical internal nodes up to Vdd but this in turn comes with increased over head area and capacitance penalties.

(4) Capacitive coupling:

The term 'coupling' implies the act of matching or pairing off two devices so that they move together by sharing energy. Thus the name capacitive coupling indicates that the matching or pairing off of two capacitances so that they initiate the combined mechanism that results in a discrete variation in the operation of the associated mechanical device. Therefore it is evident that it is also a kind of charge sharing phenomenon but when it comes to parasitic capacitances or charge sharing mechanism with respect to the inbuilt parasitic capacitive devices, it is often referred as 'capacitive coupling'. Coupling of energy or sharing of charge stored among the parasitic capacitances is capacitive coupling. Back-gate coupling is another form of capacitive coupling which is often called output-to-input coupling. Fig. 4.7 demonstrates the effect of back-gate coupling in the dynamic logic circuits. Additionally, this effect is causing erroneous functioning of the logic circuit when designed for an application at low supply voltages.

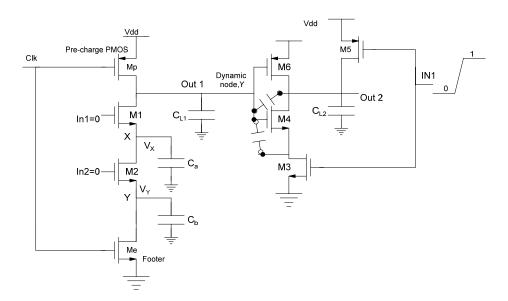


Fig. 4.7 Capacitive coupling phenomenon in domino logic circuits

The circuit shown in Fig. 4.7 has a dynamic 2-input NAND gate which drives the static NAND gate. Initially the transition in the IN1 makes the Out2 to discharge which now couples capacitively to the dynamic node or Out1 through the gate-to-source and gate-to-

drain capacitances of transistor M4, resulting in the deterioration of the strength of the nodal voltage at Out1. Therefore a special care is highly required while designing the dynamic logic circuits to pacify the intensity of capacitive coupling effect.

(5) Clock Feed-through:

Clock feed-through is a peculiar case of capacitive coupling where the coupling action takes place between the clock input and dynamic node output. As the clock is connected to pre-charge PMOS device which has got its gate-to-drain capacitance as one of the parasitic capacitances, it is involving in the coupling action. The result is an abrupt rise in the voltage above supply voltage level at the dynamic node specifically on the low-to-high transition of the clock line with the initial assumption that the evaluation network is turned off.

The troubling mechanism of this effect is that due to the rise of voltage level beyond Vdd, that makes the reverse biased junctions of pre-charge PMOS device forward biased leading to the flow of constant currents which will run uncontrollably even after the removal of supply voltage, causing static or short circuit currents to flow and there by resulting in static power dissipation. This particular phenomenon is called 'latch-up' problem. Thus once a device gets latched-up while it is in operating mode, then it results in increased static power dissipation even after the removal of supply voltage through the continuous flow of short circuit currents. Therefore while simulating the dynamic logic circuits especially designed for high speed applications, care must be taken about the boundaries of feed through voltage levels since they must not cross the bias supply voltage level. Also latch-up is encountered in CMOS digital logic circuits wherein it occurs when the output voltage level drops below the ground level voltage turning PMOS pull up network on and resulting in making all reverse biased junctions forward biased that will conduct effectively for the flow of static currents. There are few precautions to be taken in order to design latch-up free circuits at the system level as well as at the fabrication level by proper alignment of inner layers with appropriate doping concentrations of corresponding regions for fine response. The precautions are mentioned below. Fig. 4.8 shows the clock feed-through phenomenon in dynamic logic circuits.

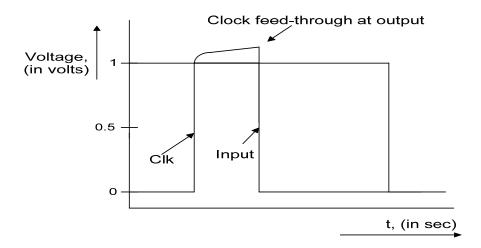


Fig. 4.8 Clock feed-through phenomenon dynamic logic circuits

Precautions:

- 1. Always make sure that the power supply rails must be off before plugging a board.
- 2. Usage of ESD (Electro Static Discharge) protection layers on all electronic appliances carefully to protect the designs from electrostatic currents which can trigger latch-up associated with input-output pads.
- 3. Designs must not get exposed to radiations including X-rays, cosmic rays which can produce electron-hole pairs as they hit the design chip through penetration since these carriers in turn generate substrate or well currents.
- **4.** Abrupt occurrence of transients on the supply or ground rail which may likely to happen if more numbers of transistors are switching simultaneously, can trigger the latch-up phenomenon in the designed circuit.

(6)Small variations from supply voltage:

This is associated with latch-up mechanism. Owing to the glitches in the circuit, there might be minute variations occurring in the nominal supply voltages which may sometimes cause severe iterative issues like latch-up problems that will repeatedly run for long time uncontrollably causing leakages and static currents to flow inside the device and there by resulting in static power dissipation. This may be due to the mismatching in the alignment of physical layers of devices while processing the lithography, etching and other fabrication steps involved. Therefore, proper care must be taken in order to avoid such

iterative cyclic problems in the circuits, while processing the device and arranging the inner layers through several physical fabrication steps. The amplitude of glitches will also show its considerable impact on the operation of the circuit that runs especially with low bias voltage. Also the voltage fluctuation at the ground level results in supply voltage variations in few cases. Hence the variations from supply voltages must be eliminated to get an error-free mechanism of the designed logic circuits.

4.3 Related work on lector power reduction technique

In static CMOS circuit designs, scaling is meant for reducing threshold voltage, device dimensions (typically channel length), bias voltage and increasing speed of operation. Due to down scaling of process technology the channel length is getting shortened which results in speeding up the operation. Besides this, there is rapid increment in the flow of inevitable sub-threshold leakage currents in the sub-threshold region in deep-sub-micron. Thus the continuous reduction of the threshold voltage causes significant increment in sub-threshold leakage current which results in static or leakage power dissipation. In this chapter, existing scheme known as "leakage control transistor (domino lector or domino LCT)" is described to minimize the leakages without affecting dynamic power consumption [60]. This lector scheme introduces two transistors called leakage control transistors of one PMOS and one NMOS each within the implementation of this technique to cut-down the leakage current through the operation of triggering each gate terminal of leakage control device by source terminal of other device.

This peculiar arrangement keeps one of the LCTs in cut-off region for any combination of inputs which increases the effective equivalent resistance of the conducting path from Vdd rail to ground rail, causing significant reduction in sub-threshold leakage currents. Before adding these LCTs to control leakage currents in the circuit, initially the gate-level net list of the circuit must be converted to a static CMOS implementation. The efficient functioning in both switching and non-switching modes of the circuit is the merit feature of lector technique which results in mitigating the leakages efficiently when compared to other schemes. Added to this, the proposed domino lector technique overcomes the limitations by other benchmark methods for the reduction of leakage. Copious schemes for reduction of leakage power have been presented in the literature. The

proposal in [67] demonstrates to make use of the dependent relation of the sub-threshold leakage current on the gate inputs. By adding an extra control network, the leakages can be controlled when it is in idle state and brought to its original operating state when activated again. This implies that there must be a sense of history to store the original data as the circuit is activated again from idle state which needs to include memory storage elements like latches or flip-flops and thereby increasing the overhead area of the circuit [8]. Power gating and clock gating are the other system level schemes used in some designs to reduce the power by turning off the supply voltage to unusable logic blocks where they make use of sleep transistors either NMOS or PMOS in the path from supply rail to ground rail. When the circuit is in active mode then sleep transistor will be turned on by a sleep signal while in other case, circuit's idle state, sleep transistor will be turned off and thereby creating virtual rails of power supply and ground nodes. Therefore, when the circuit is in active mode, the switching speed gets affected. Thus, the idle regions need to be identified along with the generation of sleep signals that trigger the sleep devices. This could be achieved by adding an extra hardware circuit which in turn consumes much power. MTCMOS (Multiple threshold voltage CMOS) scheme is other leakage reduction technique presented in [11, 12]. These devices are working at lower threshold voltages and footer transistor higher threshold NMOS transistor. The gating transistor operates as sleep transistor. Noise margins are reduced due to reverse conduction path [4]. Also, the drawback is that degradation in the performance due to series connected high threshold transistors in the switching paths. Modified version of MTCMOS technique is dual threshold technique that consists of transistors with two different threshold voltages. Critical path is evaluated by lower threshold transistors. Higher threshold voltage transistors are employed for gates which are connected in non-critical path [4, 13, 14]. The similarity between MTCMOS and Dual threshold voltage technique is that both require an additional mask for fabricating these devices on wafers according to their threshold voltages which makes it complex. Clock gating and power gating techniques for lowering power dissipation are most appropriate and compatible for system level design but not for circuit level. Another limitation is slow-latency as some time is required for returning to normal condition for the sub-circuit because when the activation is done, immediately they are not turned on.

With the help of series connected transistors, in path between supply rail and ground rail, called stack effect, this lector technique is proposed for reducing leakage power. It is given in [70-100] that the path wherein more than one transistor off between supply rail and ground rail is less leaky than that of only a single transistor off. Lector scheme introduces two leakage control transistors as described in such a way that one of them is at cut-off region. Fig. 4.9 shows the implementation of simple 2-input NAND gate using static CMOS logic and lector scheme.

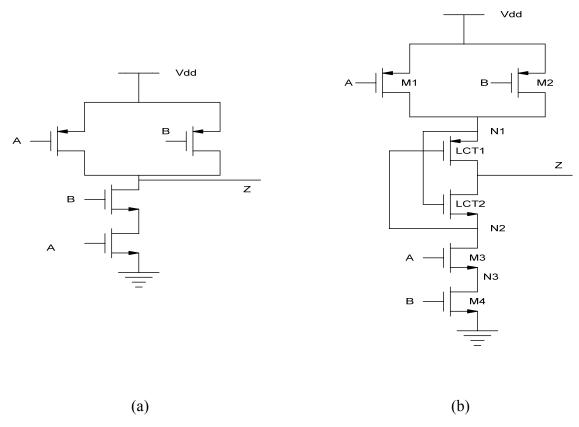


Fig. 4.9 Implementation of 2-input NAND gate using (a) static CMOS (b) lector scheme

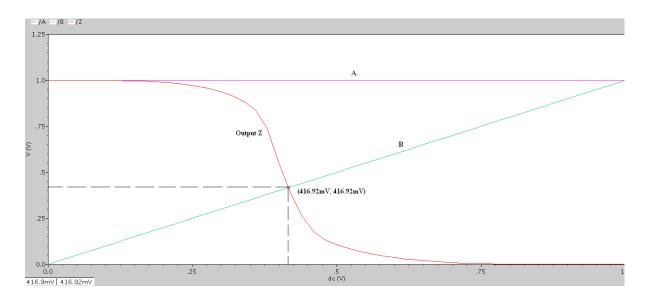


Fig. 4.10 DC characteristics of 2-input static CMOS NAND gate

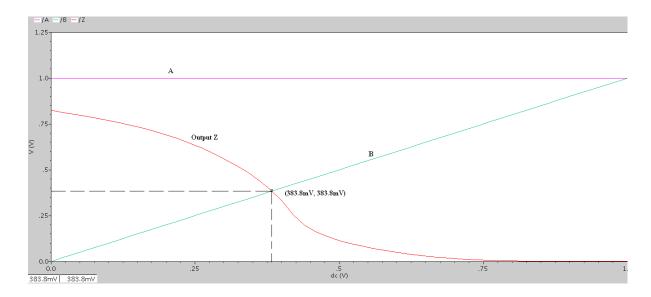


Fig. 4.11 DC characteristics of 2-input lector NAND gate

Fig. 4.10 and Fig. 4.11 show the DC characteristics of 2-input NAND gate using static CMOS and lector technique respectively where one of the inputs (input A) is fixed at 1 V and other input (input B) is varied from 0 V to 1 V.

Let's take the analysis of lector 2-input NAND gate circuit. When A is high (1 V) and B is low (0 V), then the node voltage at N2 is 800 mV which is not enough to turn LCT1 off completely. Thus, the equivalent resistance of device LCT1 is lesser than that of its off state condition which in turn allows conduction through low resistance path. Despite its lower equivalent resistance compared to its completely off state resistance, the flow of leakage current is controlled and there by resulting in lower leakage power by increasing the resistance of conducting path from Vdd rail to ground rail. In the similar manner, if both, A and B, are high then the node voltage at N1 is 200mV keeping LCT2 in cut-off mode. The condition of all transistors for various combinations of input signals is tabulated in Table 4.1. So, the inclusion of leakage control transistors increases the resistance of the conducting path from Vdd rail to ground rail. But, this also leads increased propagation delay of the circuit. To nullify this delay drawback, the leakage control transistors (LCT) are sized in such a way that the total propagation delay matches to that of conventional static CMOS NAND gate.

Table 4.1 Condition of all transistors of lector 2-input NAND gate for all possible combinations of inputs

	Input combination (A, B)					
Transistor, Q	(0, 0)	(0, 1)	(1, 0)	(1, 1)		
M1	On	On	Off	Off		
M2	On	Off	On	Off		
LCT1	Cut-off	Cut-off	Cut-off	On		
LCT2	On	On	On	Cut-off		
M3	Off	Off	On	On		
M4	Off	On	Off	On		

4.4 Proposed Domino lector technique and dynamic node stabilizing technique

Domino lector technique has the combination of LCT1 and LCT2 inserted between pre-charge PMOS transistor and pull-down network. The two transistors called leakage control transistors of PMOS and NMOS each are inserted between the nodes N1 and N2 between pull-down network and pre-charge PMOS as shown in Fig. 4.12 to obtain lector

domino AND gate. Both the drain terminals of LCT1 and LCT2 are shorted together across which dynamic output Y is observed and domino output Z takes the inversion of Y. The source terminal of each transistor is triggering the gate terminal other transistor and also the node potentials at N1 and N2 are controlling LCT1 and LCT2 respectively. This peculiar connection keeps one of the leakage control transistors in cut-off mode irrespective of current combination of inputs applied at gate terminals of evaluation network. In direct method, the lector technique is applied to conventional dynamic logic circuit which is shown in Fig. 4.12. In proposed scheme the wiring has been re-configured to improve the performance further. Comparatively the proposed scheme exhibits high degree of robustness along with leakage reduction. In direct method, the LCT combination is placed in between pre-charge PMOS transistor and pull-down network and output is observed across the mid-point between LCT1 and LCT2. In proposed technique, shown in Fig. 4.13, leakage control transistors are extended in such a way that pull-down network is inserted in between them. The performance can be improved with this kind of reconfigurable domino lector technique.

In direct method, the leakage control transistors increase the equivalent resistance so that any leakages or static currents from supply rail to ground rail can be controlled. But this technique may not have control over any leakages in pull-down network since LCTs are placed above it. The conducting path between supply rail and pull-down network is controlled by LCT network. Any gate leakages in evaluation network transistors can cause wrong discharge of the dynamic node and these inevitable leakage currents may not be controlled by LCT combination.

Therefore in order to increase the performance of direct method implementation, the proposed circuit technique is demonstrated in Fig. 4.13. This has LCT combination network extended between pre-charge PMOS device M1 and footer transistor M2 which can have a complete control over the full conducting path from supply rail to ground rail. Not only the static currents from Vdd rail to ground rail, but the gate leakage currents inside pull-down network, glitch stirred fluctuations at ground level can also be efficiently alleviated. When required, LCT combination provides low resistance path for discharging of dynamic node and maintains the same dynamic node charge as and when needed

depending upon the current combination of the inputs to the pull-down network transistors. Apart from LCT network, clock is also acting as driving element. The drain potential of footer transistor M2 is triggering the LCT1 and the drain potential of pre-charge PMOS device M1 is controlling LCT2.

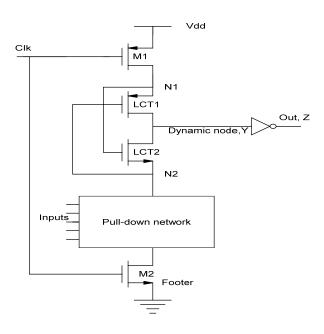


Fig. 4.12 Direct method of implementing lector domino logic circuit

Fig. 4.13 shows the implementation of proposed lector domino circuit technique and Fig. 4.14 represents the simulation result of proposed lector 2-input domino OR gate logic circuit. Potentials at nodes N1, N2 and the corresponding change in the voltage levels along with pre-charge and evaluation phases of clock signal can be observed from the simulation graph. Output-Z is following logic OR gate operation of two applied gate inputs of pull-down network. Transistors M1 and M2 are driven by clock signal and LCT1 and LCT2 are triggered by nodal potentials at N1 and N2 respectively.

In pre-charge phase, M1 is turned on and M2 is turned off. Node N1 is at Vdd which is sufficient to turn LCT2 on and node N2 is at 127.7mV which is enough to turn LCT1 on. Therefore there exists a conducting path from supply rail to dynamic node Y that charges node Y to Vdd. Since LCT2 is on and footer M2 is off, there is no discharging path during pre-charge phase irrespective of status of current combination of inputs to pull-down network. LCT1 and LCT2 are controlling the path between nodes N1 and N2 so that any

leakages in this path can effectively be managed and thereby minimizing leakage power. However the dynamic node is able to produce strong one during pre-charge phase.

During evaluation period, clock makes M1 off and M2 on. When pull-down network is completely off (all pull-down network transistors are off) then N1 is at 940 mV and N2 is at 0 V. This will turn both LCT1 and LCT2 on. When any one of pull-down transistors is on, then N1 comes around 233 mV-249 mV which may not be sufficient to turn LCT2 on and N2 will be at 0 V keeping LCT1 on. When all pull-down transistors are on then N1 will be at 220mV and N2 will be at 0 V keeping the same state of LCT2 off and LCT1 on. In all the possible combinations of inputs, the LCT network is controlling the equivalent resistance of path between nodes N1 and N2 associated with status of pull-down network transistors and thereby mitigating leakage/static power effectively. Therefore the dynamic node is charged to Vdd when strong one is needed and is discharged when a conducting path is provided by evaluation network by eliminating wrong discharge process. Furthermore, the noise parameters UNG and ANTE are improved compared to direct method scheme. The only drawback associated with this proposed technique is that its increased propagation delay with the inclusion of LCTs and this could be managed by replacing them with properly sized transistors in the conducting path so that the propagation delay becomes equal to that of conventional domino scheme.

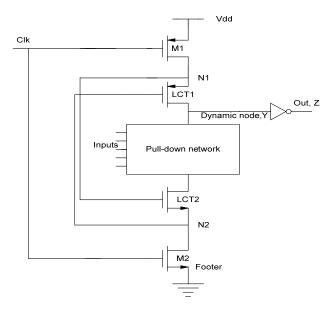


Fig. 4.13 Proposed lector domino logic circuit technique

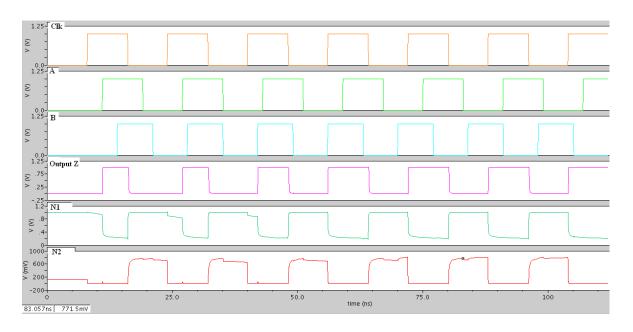


Fig. 4.14 Simulation of Proposed lector 2-input domino OR gate logic circuit

Dynamic node stabilizing technique (Power reduction)

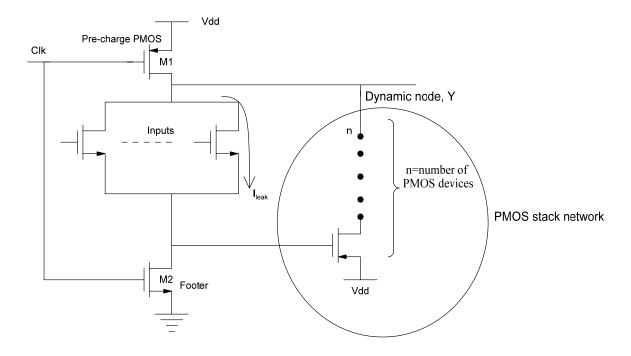


Fig. 4.15 Proposed dynamic node stabilizing technique

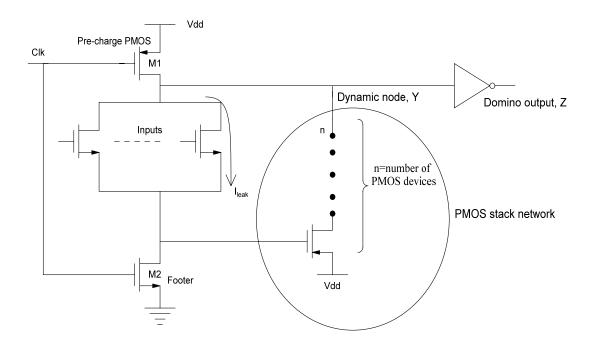


Fig. 4.16 Proposed dynamic node stabilizing technique applied to basic domino logic circuit

Fig. 4.15 and Fig. 4.16 show proposed dynamic node stabilizing technique applied to wide fan-in dynamic logic and domino logic circuits. Unlike domino scheme which includes static inverter at output, the basic pure clocked circuit (dynamic logic) is implemented with this scheme. The idea behind this technique is to produce strong logic values (both strong ones and strong zeros) at dynamic node without including static flavor at output. It consists of series of PMOS transistors, called PMOS stack network, connected in between the drain terminal of footer M2 and dynamic node Y. It also reduces the total power consumption by utilizing the flow of leakages through evaluation network with large number of PMOS devices. In spite of limiting the flow of sub-threshold leakage current through pull-down network in the evaluation phase, it is utilized to turn the stack network on which energizes the dynamic node by passing the strong one through supply voltage. Thus, the inevitable leakage current causes degradation in the strength of voltage at dynamic node which must be recompensed. This is achieved by the stack network of PMOS devices in the circuit. As the pull-down network is off, no current is supposed to flow through it except the small but finite sub-threshold leakage current which causes the deterioration in the strength of voltage stored at the dynamic node. This novel technique demonstrates the effective utilization of leakage currents. The leakage current through pull down network constitutes a nominal voltage drop across the footer M2 which turns the stack PMOS network on. This

technique effectively produces strong logic values (both strong one and strong zero) at output. The demerit is that high degree of circuit robustness could be achieved with gate delay penalty.

During pre-charge mode, M1 is turned on providing a conducting path from Vdd rail to dynamic node Y and there by charging it up to Vdd. Irrespective of status of pull-down network, dynamic node produces strong one as PMOS keeper is charging it to supply voltage against any inevitable leakages in pull-down network. Footer M2 is off which breaks the discharging path from dynamic node to ground. Therefore dynamic node is charges to Vdd in this phase.

During evaluation phase, M1 is turned off and M2 will be turned on. The actual logic function is evaluated by elaborating the gate terminals of pull-down network transistors. When all inputs are low then pull-down network is completely off and the dynamic node is supposed to maintain its pre-charged voltage. Thus, if any leakages or static currents are flowing through pull-down network, then they will constitute a small voltage drop across the footer M2 which is sufficient to turn PMOS stack network on and as a result all the stack devices will be turned on passing voltage Vdd to dynamic node Y. In this way the leakage is used to trigger the stack network which energizes the dynamic node by passing Vdd. This indicates that the very cause behind the voltage drop at the dynamic node is again compensating this voltage drop effect by charging it through stack PMOS network.

During evaluation mode, when pull-down network is on then it discharges the dynamic node completely by providing a low-resistance path. As footer is already on the nominal voltage drop across footer will turn off the PMOS stack network and there by invalidating the presence of this network.

Fig. 4.15 is indicating the proposed stabilizing technique applied for basic wide fan-in dynamic logic circuit. It has 'n'-number of PMOS transistors in stack network. Also power reduction mechanism is associated with value of 'n'. As the value of-n increases, comparatively less power is consumed.

Fig. 4.16 has shown the same technique applied to basic wide fan-in domino logic circuit design. The simulation is performed with 2-input OR gate circuit. As there is no requirement for producing strong logic levels at domino output which is done by default static inverter, this applied technique improves the UNG and ANTE noise metric parameters compared to conventional domino logic circuit. As 'n' increases, both static and dynamic power consumption are minimized. This could be achieved by analyzing the circuit in detail by drawing the equivalent circuit with corresponding (W/L) ratios of PMOS stack network devices. All the gate terminals of stack network devices are connected together and are fed to drain terminal of footer M3. When this is activated, supply voltage (Vdd) is passed to dynamic node through cascaded network.

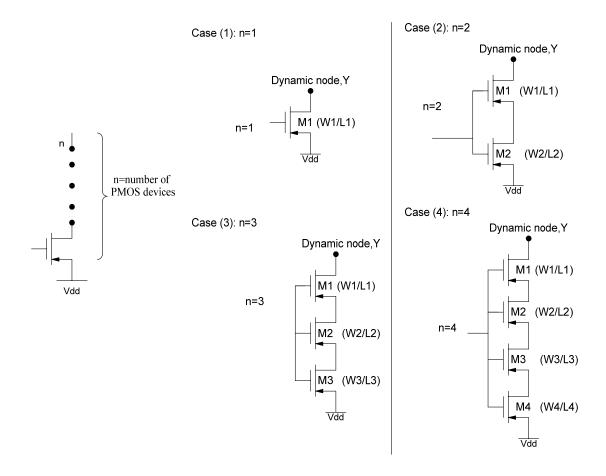


Fig. 4.17 Impact of PMOS stack network on power consumption

Let us assume, for the ease of analysis, that all the PMOS devices are of same dimensions. Thus $W_1=W_2=....=W$ and $L_1=L_2=....=L$. The equivalent resistance offered by stack

network is calculated individually in all the cases and corresponding power consumption is formulated.

Case (1): n=1

The equivalent $\left(\frac{W}{L}\right)$ ratio is $\left(\frac{W_{l}}{L_{l}}\right)$.

$$P_{Dynamic} = \frac{V_{dynamic}^2}{R_{equivalent}}$$
, where $R_{equivalent} \alpha \frac{1}{\left(\frac{W_1}{L_1}\right)}$

Since $(W_1=W_2=...)=W$ and $(L_1=L_2=...)=L$

$$\Rightarrow R_{equivalent} = K. \frac{1}{\left(\frac{W_1}{L_1}\right)} = K. \frac{1}{\left(\frac{W}{L}\right)} = R1 \text{ [where K is a constant]}$$

Therefore, $P_{\text{Dynamic}} = \frac{V^2}{R_1} = P1$ (Let us say)

Case (2): n=2

The equivalent $\left(\frac{W}{L}\right)$ ratio is effective of {series connected $\left(\frac{W_1}{L_1}\right)$ and $\left(\frac{W_2}{L_2}\right)$ }

$$\Rightarrow \left(\frac{W_1(or)W_2}{L_1 + L_2}\right)$$

$$P_{\text{Dynamic}} = \frac{V_{\text{dynamic}}^2}{R_{\text{equivalent}}}$$

Since $(W_1=W_2=....)=W$ and $(L_1=L_2=....)=L$

$$R_{equivalent} = K. \frac{1}{\left(\frac{W_1(or)W_2}{L_1 + L_2}\right)} = K. \frac{1}{\left(\frac{W}{2L}\right)} = R2$$

Therefore,
$$P_{\text{Dynamic}} = \frac{V^2}{R_2} = P2$$
.

The width of series connected transistors remains unchanged and channel length increases since channel length of each transistor is added to that of other transistor and as a result the total length will be summation of all individual channel lengths.

Thus, clearly R2 > R1 and P2 < P1.

Case (3): n=3

The equivalent $\left(\frac{W}{L}\right)$ ratio is effective of {series connected $\left(\frac{W_1}{L_1}\right)$, $\left(\frac{W_2}{L_2}\right)$ and $\left(\frac{W_3}{L_3}\right)$ }

$$\Rightarrow \left(\frac{W_1(or)W_2(or)W_3}{L_1 + L_2 + L_3}\right)$$

$$P_{Dynamic} = \frac{V_{dynamic}^2}{R_{equivalent}}$$

Since $(W_1=W_2=....)=W$ and $(L_1=L_2=....)=L$

$$R_{equivalent} = K. \frac{1}{\left(\frac{W_1(or)W_2(or)W_3}{L_1 + L_2 + L_3}\right)} = K. \frac{1}{\left(\frac{W}{3L}\right)} = R3$$

Therefore,
$$P_{\text{Dynamic}} = \frac{V^2}{R_3} = P3$$
.

Clearly R3 > R2 > R1 and thus P3 < P2 < P1.

Case (4): n=4

The equivalent $\left(\frac{W}{L}\right)$ ratio is effective of {series connected $\left(\frac{W_1}{L_1}\right)$, $\left(\frac{W_2}{L_2}\right)$, $\left(\frac{W_3}{L_3}\right)$ and $\left(\frac{W_4}{L_4}\right)$ }

$$\Rightarrow \left(\frac{W_1(or)W_2(or)W_3(or)W_4}{L_1 + L_2 + L_3 + L_4}\right)$$

$$P_{\text{Dynamic}} = \frac{V_{\text{dynamic}}^2}{R_{\text{equivalent}}}$$

Since $(W_1=W_2=...)=W$ and $(L_1=L_2=...)=L$

$$\underset{\Rightarrow}{R_{equivalent}} = K. \frac{1}{\left(\frac{W_{1}(or)W_{2}(or)W_{3}(or)W_{4}}{L_{1} + L_{2} + L_{3} + L_{4}}\right)} = K. \frac{1}{\left(\frac{W}{4L}\right)} = R4$$

Therefore,
$$P_{\text{Dynamic}} = \frac{V^2}{R_4} = P4$$
.

Clearly R4 > R3 > R2 > R1 and thus P4 < P3 < P2 < P1.

From the mathematical analysis, it is evident that as the number of PMOS stack network transistors-n increases, the corresponding effective (W/L) ratio of all devices connected in series is getting reduced which results in increasing the equivalent resistance at the dynamic node and thereby lowering the dynamic power. Therefore higher the 'n' value, lower the power dissipation. There are limitations too for this configuration. In few applied circuits this proposed technique results in lowering the power but at the cost of propagation delay. Thus, both static as well as dynamic power consumption parameters are strategically minimized with the increase of number of stack PMOS devices.

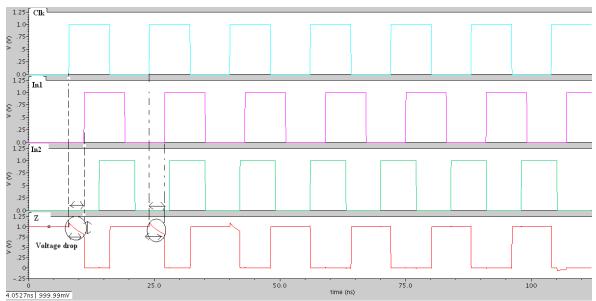


Fig. 4.18 Dynamic node voltage drop for 2-input dynamic NOR gate

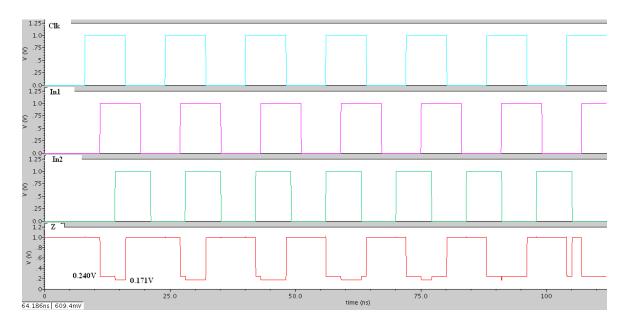


Fig. 4.19 Dynamic node voltage drop for 2-input dynamic NOR gate with n=1

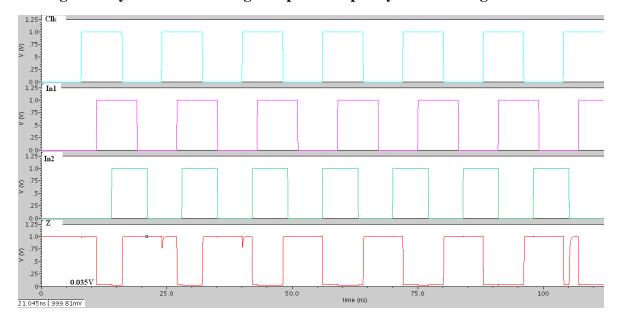


Fig. 4.20 Dynamic node voltage drop for 2-input dynamic NOR gate with n=8

Fig. 4.18 shows dynamic node voltage drop for 2-input dynamic NOR gate that could not produce strong one during clock's evaluation period when pull-down network is off. There is slight degradation in the voltage level at dynamic node due to inevitable leakages which therefore needs to be compensated. The drawback of dynamic logic over domino logic is that straightforward cascading of dynamic circuits is not possible due to this output voltage degradation phenomenon and this problem becomes more tactile in cascaded system. Since after passing few stages of cascaded network, there will surely be deterioration in the

strength of output voltage and this degraded signal might not be able to drive the next stage which leads to catastrophic failure of complete cascaded system.

In this proposed technique this drawback is considerably managed and could be overcome by charging the dynamic node during evaluation period of clock signal with the help of external added devices. When a single PMOS transistor is added in between dynamic node and footer, it boosts up this node by passing strong one. When there is no conventional discharging path for dynamic node, it retains its stored value. Due to leakages, there might be voltage drop that has to be charged again in order to produce a stronger level output. The PMOS gets turned on with the leakages through pull-down network during evaluation phase and does pass Vdd to dynamic node, replenishing the leakage drop. Furthermore power consumption is also reduced with increasing number of PMOS stack network. Both static power and dynamic power consumption are minimized with increase in number of PMOS stack transistors.

There are limitations too as it increases the propagation delay with more number of PMOS devices but produces strong zeros at dynamic node. Fig. 4.19 portrays that dynamic node which does not completely discharge dynamic node and as a result 0.241V is remaining at this node for dynamic NOR gate with n=1. Thus to improve this discharging phenomenon, more number of PMOS devices need to be added in the stack network. Fig. 4.20 presents dynamic node voltage drop for 2-input NOR gate with n=8 and it efficiently discharges dynamic node and thereby producing strong ones and zeros. Similarly, when the value of n increases to 16, 32 or any higher value, the discharging phenomenon take place effectively, resulting in generation of strong zeros and strong ones at dynamic node but with gate delay penalty. The simulation results along with tabulations and graphs are given in section 4.5.

4.5 Simulation results and discussion

Table 4.2 Comparison of various parameters for Static CMOS NAND and Lector NAND techniques

Technique	Dynamic power (in W)	Leakage or static power (in W)	Total power (in W)	Total propagation delay, T _p (in sec)	Power- delay- product (in W- sec)	$NM_{L} = V_{IL} - V_{OL} (in V)$	$NM_{H}=$ $V_{OH} V_{IH}$ (in V)
Static CMOS NAND	7.723E-8	2.501E- 6	2.57E- 6	19.451E-12	150.2E- 20	0.311	0.481
Lector NAND	7.932E-8	1.765E- 6	1.84E- 6	33.034E-12	262E-20	0.211	0.320

Table 4.3 Comparison of various parameters for domino lector-direct method and proposed domino lector technique

Domino logic technique	Dynamic power (in W)	Leakage or static power (in W)	Total power (in W)	Total propagation delay, T _p (in sec)	Power- delay- product (in W-sec)
Domino lector-direct method	0.658E-6	2.535E-8	0.683E-6	110.72E-12	72.853E- 18
Domino lector- Proposed technique	0.618E-6	2.512E-8	0.643E-6	111.164E-12	68.699E- 18

Table 4.4 UNG and ANTE comparison of domino lector-direct method and proposed domino lector technique with Fan-in=2

Domino logic technique	UNG (in V)	ANTE (in V²*picosec)
Domino lector-		
Direct method	468E-3	4.21
Domino lector-		
proposed technique	556E-3	5.14

UNG against Process corner analysis for variable Fan-in:

Table 4.5 UNG comparison of proposed domino logic techniques with Fan-in=2 at different Process Corner analysis

Domino logic	UNG (in V)					
technique	NN	FF	SS	FS	SF	
Domino lector- direct method	468E-3	377.33E-3	599.42E-3	398.56E- 3	585.33E- 3	
Domino lector- Proposed technique	556E-3	423.52E-3	825E-3	468.51E- 3	730E-3	

Table 4.6 UNG comparison of proposed domino logic techniques with Fan-in=4 at different Process Corner analysis

Domino logic technique	UNG (in V)						
100	NN	FF	SS	FS	SF		
Domino lector- direct method	450.78E-3	337.83E-3	600.24E-3	373.03E- 3	595.5E-3		
Domino lector- Proposed technique	550E-3	392E-3		445E-3	750E-3		

Table 4.7 UNG comparison of proposed domino logic techniques with Fan-in=8 at different Process Corner analysis

Domino logic technique	NN	UN FF	IG (in V)	FS	SF
Domino lector- direct method	425.48E-3	305.24E-3	611.01E-3	353.26E- 3	605.12E- 3
Domino lector- Proposed technique	543.12E-3	370.85E-3		428.44E- 3	782.24E- 3

Table 4.8 UNG comparison of proposed domino logic techniques with Fan-in=16 at different Process Corner analysis

Domino logic technique	NN	UNG (in V) FF SS FS SF			
Domino lector- direct method	403.99E-3	287.26E-3	641.23E-3	331.33E- 3	624.08E- 3
Domino lector- Proposed technique	537.55E-3	353.03.9E-3		405.64E- 3	810E-3

Table 4.9 UNG comparison of proposed domino logic techniques with Fan-in=32 at different Process Corner analysis

Domino logic	UNG (in V)					
technique	NN	FF	SS	FS	SF	
Domino lector- direct method	388.44E-3	268.52E-3	659.23E-3	322.20E-3	635.12E-3	
Domino lector- Proposed technique	522.22E-3	309.58E-3		350.87E-3	856.86E-3	

Analysis of dynamic node stabilizing technique (Power reduction):

Table 4.10 Comparison of power and delay parameters of proposed technique for 2-input dynamic NOR gate with variable n (number of PMOS stack devices)

Number of	Dynamic	Leakage or	Total power	Total	Power-
PMOS	power (in	static power	(in W)	propagation	delay-
stack	W)	(in W)		delay (in	product (in
devices (n)				sec)	W-sec)
0(Basic-	0.070E-6	5.585E-11	0.070E-6	18E-12	1.26E-18
NOR)					
1	10.756E-6	25.55E-6	36.306E-6	26.89E-12	289.22E-18
2	6.679E-6	15.82E-6	22.499E-6	26.48E-12	176.85E-18
4	3.81E-6	8.9E-6	12.7E-6	28.9E-12	110.109E-18
8	2.10E-6	4.712E-6	6.81E-6	33.53E-12	70.413E-18
16	1.20E-6	2.411E-6	3.61E-6	39.71E-12	47.652E-18
24	0.933E-6	1.617E-6	2.550E-6	45.15E-12	42.124E-18
32	0.814E-6	1.21E-6	2.02E-6	50.26E-12	40.911E-18

Table 4.11 Comparison of power and delay parameters of proposed technique for 2-input domino OR gate with variable n (number of PMOS stack devices)

Number of	Dynamic	Leakage or	Total power	Total	Power-delay-
PMOS stack	power (in W)	static power	(in W)	propagation	product (in
devices (n)		(in W)		delay (in sec)	W-sec)
0(Basic-OR)	0.203E-6	2.75E-8	0.23E-6	47.24E-12	9.589E-18
1	11.94E-6	28.96E-6	40.90E-6	51.21E-12	611.44E-18
2	6.984E-6	16.595E-6	23.579E-6	51.29E-12	358.20E-18
4	4.00E-6	9.518E-6	13.51E-6	54.19E-12	216.76E-18
8	2.291E-6	4.403E-6	6.694E-6	59E-12	135.16E-18
16	1.43E-6	2.311E-6	3.74E-6	66.75E-12	95.45E-18
24	1.214E-6	1.655E-6	2.869E-6	72.78E-12	88.35E-18
32	1.152E-6	1.245E-6	2.397E-6	78.27E-12	90.16E-18

Table 4.12 UNG and ANTE comparison of proposed technique applied for domino 2input OR gate for various stack devices

Number of stack PMOS	UNG (in V)	ANTE (in V ² *picoSec)
devices (n)		
0	402.5E-3	3.58
1	923E-3	8.91
2	770.3E-3	7.22
4	670.75E-3	6.38
8	638E-3	5.84
16	632.5E-3	5.63
24	629E-3	5.53
32	625E-3	5.41

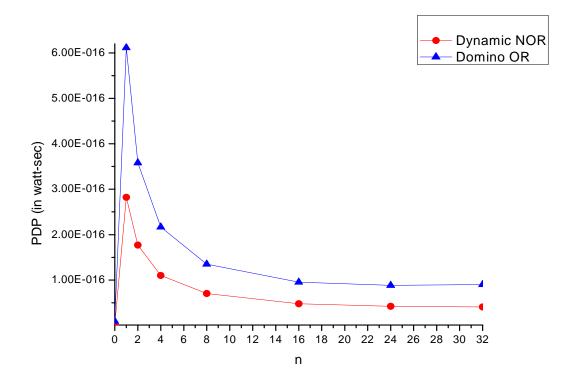


Fig. 4.21 Variation of Power-delay-product with number (n) of PMOS stack devices for proposed dynamic and domino 2-input OR gate

Discussion

From the tabulations, it is evident that the lector domino logic circuit implemented in direct method and proposed technique is exhibiting high degree of noise robustness in terms of leakages and noise metric parameters. Furthermore both the techniques are subjected to distinct process corners for wide fan-in circuits and investigated the overall functionality of designs. In direct method of implementing the lector domino circuit, the LCT combination is limited to the path between dynamic node and pull-down network whereas in the proposed scheme, the LCT combination is extended in such a way that it is controlling the resistance of total conducting path between supply rail and ground rail and thereby reducing static currents through this path. This is the major advantage of this scheme over direct method.

Table 4.3 shows reduced power-delay-product of proposed lector domino circuit technique over direct method because the proposed circuit is discharging without leakages in pull-down network. Table 4.4 reveals increased UNG because less voltage is sufficient

to turn a single NMOS in pull-down network. As fan-in increases, more voltage is required to make the evaluation network on which results in reduction of the noise gain.

Analyzing the process-corner analysis, amongst five typical corners (NN, SS, FF, FS and SF), Normal-Normal corner is providing nominal switching threshold voltages for NMOS and PMOS devices as specified by EDA tool. Generally, all the designed circuits will operate as per design specifications and according to user constraints at NN corner. Thus evaluation and estimation of overall performance of designed circuit by subjecting at NN process corner alone does not finish the task. In fact the circuit is required to be tested by subjecting at all the extreme corners. Having done that entire task, then only the performance of circuit can be judged. Always the SS corner is assuring increased noise tolerance and which is also observed from tabulations. Highest UNG is recorded in SS corner since both NMOS and PMOS transistors are slow running devices which implied that their switching threshold voltages are high. Thus, the circuit is made less sensitive to noise glitches and hence the nature of being responsive to gate input noise glitches is gradually reducing, that in turn increases UNG. So, in comparison with other process corners, SS corner always exhibits highest UNG. Also in SF corner, as NMOS switching threshold voltage is higher and PMOS switching threshold voltage is lower, it also contributes increment in noise gain but not as efficient as SS corner because the PMOS threshold voltage is responding to noise impulses at gate inputs which makes the circuit more sensitive to noise glitches. Normally, in comparison with NN, FS and FF corners, circuit at SF corner exhibits better immunity towards noise. At FF process corner, both NMOS and PMOS devices are of very high speed with reduced switching threshold voltages and consume more power. The least UNG is noted in FF case. This becomes more tactile with wide fan-in circuits. With increased fan-in, the UNG in other process corners is gradually lowering since more voltage is required to turn pull-down network on and discharging phenomenon is becoming slow. Thus normal circuits usually become more sensitive to noise glitches at gate inputs due to their lower threshold voltages and as a result UNG is reduced. But proposed circuits are functioning at FF corner efficiently. Normally, with increased fan-in, the UNG lowers. The performance of circuit at FS process corner lies in between the corresponding performances at NN and FF corners. As NMOS possesses lower threshold voltage, it becomes more sensitive to noise. Despite

slow PMOS device, the corner is lowering the noise immunity. From the tabulations, it is evident that proposed techniques exhibit high degree of robustness at FS corner too.

The topology of LCT combination in proposed lector domino logic circuit technique assures high noise tolerance than that of direct method. UNG measurements at various process corners for wide fan-in circuits reveal that proposed techniques are less sensitive to external and internal ambient noise glitches. This could be achieved by configuring the circuit with series connected transistors, in the path between supply rail and ground rail, called stack effect. This possesses the conducting path wherein more than one transistor is off between supply rail and ground rail which is less leaky than that of only a single transistor is off. Lector scheme introduced two leakage control transistors as described in such a way that one of them is at cut-off region.

Having analyzed the results from process corner analysis, it is clear that proposed circuit techniques assure high noise robustness with increased fan-in despite few limitations at various corners. Proposed techniques are functioning efficiently, at all the corners for lower fan-in circuits but with increased fan-in, the operating region of design has started becoming limited to few corners only which might be due to the reason that the circuits are being operated at lower bias voltage of 1V. For example, from the simulations, it is clear that the UNG is getting lowered at NN, FF and FS process corners for higher fan-in. Thus increasing the bias voltage range would facilitate broadening the operating region of proposed domino techniques at all the process corners for wide fan-in circuits but while doing so the power consumption must be taken care off as it is directly proportional to square of supply voltage. Hence optimization through trade-off between supply voltage and process corners is necessary while improving the proposed domino logic circuit techniques.

The proposed technique for stabilizing dynamic node also operates efficiently and is able to energize the dynamic node effectively. As the number of PMOS stack devices is increasing, the power consumption is less and more stronger values (strong one) are being generated by circuit at dynamic node. This can be observed from Fig. 4.21. It is also noted from the simulations that both static current and dynamic current are being limited and as a

result total power consumption is minimized with more number of PMOS stack devices. Despite these benefits, proposed technique suffers from increased propagation delay.

Table 4.9 and 4.10 reveal that proposed technique for stabilizing dynamic node is assuring high performance from the perspectives of dynamic power, leakage power, propagation delay and PDP. With increased stack PMOS devices, the equivalent resistance of stack network is increasing and as result the power consumption is lowered. But the main disadvantage is that, the complete discharging phenomenon is not occurring since small amount of voltage is remaining at dynamic node which is due to increased PMOS devices in stack network and as result discharging process is becoming slow and ineffective. It is evident from Table 4.11 that, UNG and ANTE metric parameters are reducing with increased stack PMOS devices and when compared with conventional domino OR gate circuit, proposed technique assures higher noise tolerance.

Fig. 4.16 has shown the same technique applied to basic domino logic design. The simulation is performed with 2-input OR gate circuit. As there is no requirement for producing strong logic levels at domino output which is done by default static inverter, this applied technique improves the UNG and ANTE noise metric parameters compared to conventional domino logic circuit. As 'n' increases, both static and dynamic power consumption are minimized. This could be achieved by analyzing the circuit in detail by drawing the equivalent circuit with corresponding (W/L) ratios of PMOS stack network devices.

4.6 Conclusion

Therefore this chapter in section-1 gives general introduction to need for power reduction and leakage minimization. Section 4.2 discussed signal integrity issues in detail with simulations. Section 4.3 described the prior works related to leakage power reduction schemes and the lector technique. Modified lector domino scheme and dynamic node stabilizing technique are proposed in section 4.4. Simulation results along with discussion are presented in section 4.5.

All the simulations are done at CMOS 90 nm process technology with 1 V power supply. The proposed domino logic circuit techniques are simulated at distinct ambient conditions by subjecting them to various process corners and thereby observations are

tabulated. Comparisons are made and conclusions are drawn. Noise analysis is carried out which includes the need for robustness, various metric parameters for measuring noise immunity or robustness of domino circuits such as UNG, ANTE along with the method of calculations, various sources of noise in domino logic circuits and their role on operating region.

Process corner analysis and various corners involved in it along with their significant role on the overall functionality of the designed lector domino logic circuit are described. Also the consequences of subjecting the device to the extreme corners with the boundary limitations are discussed.

The result section shows the calculations and comparisons of all the parameters of proposed lector domino techniques. The primary design parameters such dynamic power, leakage or static power, total power, PDP (power-delay-product), UNG and ANTE for various fan-in circuits of existing and proposed techniques are measured. The comparisons along with tabulations are made and discussed the functionality with pros and cons.

Thus the proposed circuits are exhibiting improved leakage reduction and greater noise immunity.

CHAPTER 5 DESIGN OF VARIOUS DOMINO BASED SCHMITT TRIGGER CIRCUITS

5.1 Introduction

Schmitt trigger (ST) is a comparator based application circuit that possesses hysteresis. This can be obtained by implementing positive feedback to the non-inverting differential amplifier. Also it is an active circuit which converts analog input data into digital output. The name "trigger" is given to this active circuit since output or the final response is triggered or actively driven by corresponding change at the input signal and hence it is referred as 'Schmitt trigger'. The output of Schmitt trigger remains unchanged or retains its previously stored value until the input changes through some threshold point. This Schmitt trigger circuit was first invented by the U.S scientist named Otto. H. Schmitt in the year 1934 [5-20]. The peculiar phenomenon of Schmitt trigger circuit is that it exhibits 'hysteresis' behavior which is bounded by two typical threshold values called upper threshold point and lower threshold point. Hysteresis also acts as memory state. Thus, if the input is crossing the upper threshold value, then it triggers the output and according to that output reaches logic high. On the other hand, if the input is below lower threshold value, then the corresponding change drives output and hence it becomes logic low. As it possesses two threshold points, there is a chance of getting a state wherein the input lies in between upper threshold value and lower threshold value. Therefore, if this case arises then the output retains its value and implies that the previous data is stored by the design when input lies in between two threshold values. This peculiar dual threshold phenomenon is called hysteresis. It also implies that the Schmitt trigger functions as a memory storage device as a bi-stable circuit (typical basic latch or flip-flop circuit) since two stable states are being stored by this design in this hysteresis mode. Schmitt trigger can be constructed by using latch and vice-versa.

In this chapter Schmitt trigger is designed using domino logic circuit techniques by applying the proposed circuit techniques and is simulated. It is observed that the Schmitt trigger possesses various hysteresis phenomena with different techniques applied.

In this chapter, we propose a novel leakage power reduction domino Schmitt trigger circuits. The rest of the chapter is organized as follows. Section 5.2 explains conventional

Schmitt triggers using op-amp and CMOS logic. Section 5.3 demonstrates proposed domino Schmitt trigger circuits along with analysis. In section 5.4, simulation results with discussion are presented and in section 5.5 concluding remarks are made.

5.2 Conventional Schmitt triggers

5.2.1 Op-amp based Schmitt trigger

Non-linear operational amplifier circuit

An op-amp circuit, connected without any negative feedback, constantly saturates at either its positive or negative saturation voltage point, is often referred as non-linear circuit because the circuit functions beyond its normal linear region except in transition state that occurs between positive and negative saturation states [5-10]. The very basic and simple non-linear circuit is open-loop polarity indicator as shown in Fig. 5.1.

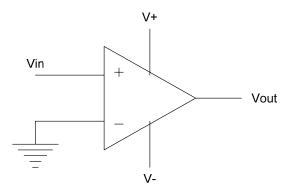


Fig. 5.1 Basic open-loop polarity indicator

The input voltage Vin, is fed directly to non-inverting terminal of op-amp and inverting terminal is grounded. As there is no feed-back connection, the range of input voltage Vin, across which the operation is completely linear, has been considerably small. Thus, the positive input, which is amplified by the open-loop gain of op-amp, drives the Vout to its upper saturation point and in similar manner the small negative input forces Vout to its lower saturation point. Therefore the circuit depending upon the polarity of the Vin, does shift Vout to either V+ or V- consequently. Fig. 5.2 shows the comparator circuit where a reference voltage signal Vr, is added to its inverting terminal that makes the circuit as open-loop comparator.

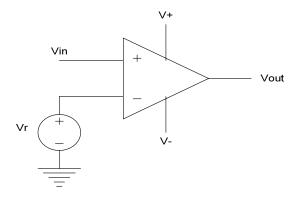


Fig. 5.2 Basic open-loop comparator

Therefore, the operation goes like this. When Vin crosses or becomes more positive than Vr, then Vout shifts to its positive saturation point V+, while in other case when Vin is below Vr, then Vout is forced to its negative saturation point V-. This indicates that the comparator circuit is more prone to noise glitches when input voltage Vin is closer to reference voltage Vr. During this period, the differential voltage, '(non-inverting terminal voltage) - (inverting terminal voltage)' comes close to zero and minute noise glitches at input node may cause Vout to swing between V+ and V- unpredictably. This problem, which needs to be alleviated, could be solved by using a positive feedback. Thus a positive feed-back Schmitt trigger lessens this drawback.

Schmitt trigger is basically a comparator circuit wherein the reference voltage is fraction of output voltage through feedback. The main difference between comparator and Schmitt trigger is that, in a comparator output voltage reaches to either positive threshold point or negative threshold point whenever the input voltage exceeds the reference voltage. Schmitt trigger, unlike comparator which does not possess memory, stores the previous or most recent data at output node and does hold it even if input voltage becomes zero. Schmitt trigger also acts as bi-stable multi-vibrator since it possesses two stable states when the input signal is zero: one stable state with positive output and other with negative output. The op-amp based Schmitt trigger is shown in Fig. 5.3. The output change with respect to various reference voltages is given in Fig. 5.4. The typical hysteresis voltage of op-amp based Schmitt trigger is presented in Fig. 5.5.

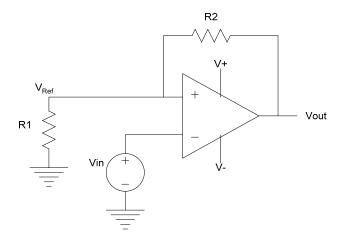


Fig. 5.3 Op-amp based Schmitt trigger configuration

Let us assume the voltage between inverting and non-inverting terminal is Vx. By applying KCL at input of op-amp we get,

$$\frac{Vx}{R1} + \frac{Vx - Vout}{R2} = 0 \cdot \text{Clearly, Vx=Vin.}$$

$$\text{Thus } \frac{Vin}{R1} + \frac{Vin - Vout}{R2} = 0$$

$$\Rightarrow Vin\left(\frac{1}{R1} + \frac{1}{R2}\right) = Vout\left(\frac{1}{R2}\right)$$

$$\Rightarrow Vout = Vin\left(\frac{R1 + R2}{R1}\right)$$

$$\Rightarrow \frac{Vout}{Vin} = \left(\frac{R1 + R2}{R1}\right)$$

If the Vout is at V+ then V_{th-upper} becomes positive and is given by following equation.

$$V_{th-upper} = \left(\frac{R1}{R1 + R2}\right)V + \tag{5.1}$$

If the Vout is at V- then V_{th-lower} becomes negative and is given by following equation.

$$V_{th-lower} = \left(\frac{R1}{R1 + R2}\right)V - \tag{5.2}$$

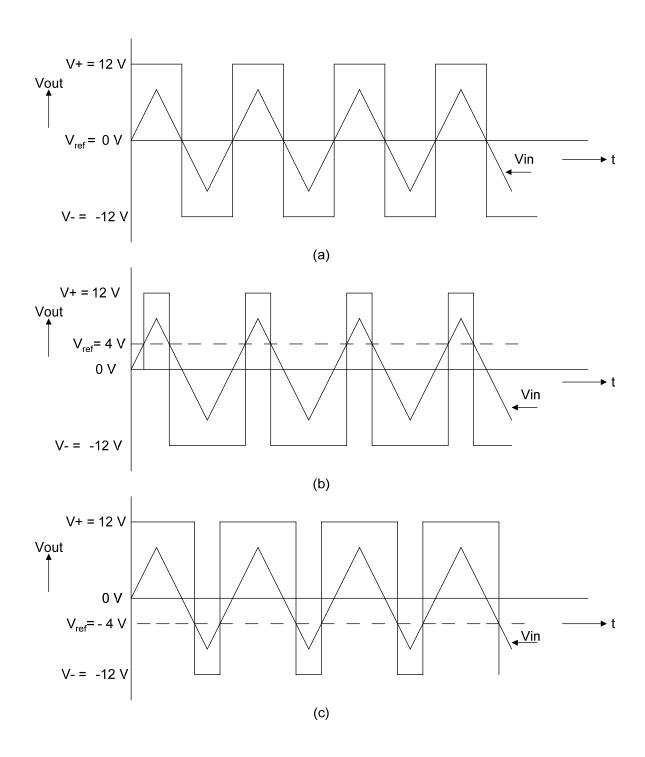


Fig. 5.4 Output of op-amp based Schmitt trigger with respect to various reference signal voltages When $V_{ref}=0$ V, (b) When $V_{ref}=4$ V and (c) When $V_{ref}=-4$ V [5]

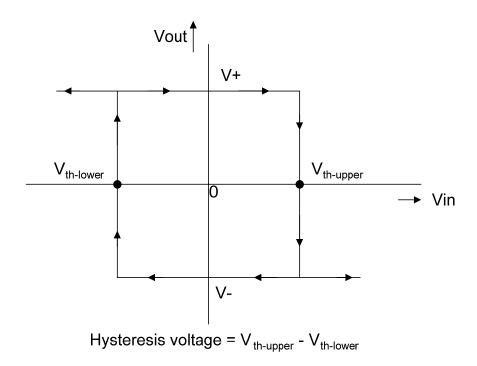


Fig. 5.5 Hysteresis of op-amp based Schmitt trigger

Hysteresis is exhibited by the transfer characteristic of Schmitt trigger where the state of output is followed in the path. From Fig. 5.5, it is clear that Schmitt trigger also performs inverting operation. Moreover, the large positive input voltage brings the output to negative value and the large negative input voltage shifts the output to positive value.

Therefore, the Vout of Schmitt trigger can also become zero when Vin is zero which also assures zero hysteresis because $V_{th\text{-upper}} - V_{th\text{-lower}} = 0$. This is highly unstable state that cannot be sustained indefinitely since a slight noise impulse or glitch can cause the output to fall in one of its states which are highly stable. Thus, for a Schmitt trigger with both input and output are at ground state, if noise glitch at input (non-inverting terminal) makes small positive voltage, then it is amplified at output and this rise will be attenuated by the voltage divider circuit that consists of R1 and R2. Despite this attenuation, the amplified voltage will appear as a positive voltage at non-inverting terminal of op-amp. The effective rise in the difference between non-inverting terminal voltage and inverting terminal voltage will further be amplified by op-amp and makes Vout to become even more positive. Again the amplified output voltage appears at non-inverting terminal which undergoes further amplification and as a result the output Vout will finally be forced to its upper saturation point. This cyclic action illustrates the effect of

positive feedback. Thus the upper saturation point is a stable state in which the positive output triggers the op-amp differential voltage in positive direction. In similar manner, there exists lower saturation point where, both Vout and differential voltage become negative and the negative Vout constantly drives differential voltage in negative direction. Thus a Schmitt trigger always remains in one of its two stable states unless and until an exceptionally large external impulse triggers the output to other un-known state which seldom happens. Hence this phenomenon, of being stable in both states, called 'bistability' makes Schmitt trigger favorable for design of electronic memory devices.

5.2.2 CMOS Schmitt trigger

The CMOS based Schmitt trigger circuit is widely used as a regenerative circuit whose voltage transfer characteristics (VTC) are similar to that of CMOS inverter but with two different threshold voltages called upper threshold voltage ($V_{th\text{-upper}}$) and lower threshold voltage ($V_{th\text{-lower}}$) which constitute typical hysteresis phenomenon [60]. Having possessed this hysteresis behavior, these Schmitt trigger circuits are used as detectors of high-to-low and low-to-high transitions in noisy ambient effectively. As it possesses bi-stable nature, at any input voltage, highly stabilized output is achieved without any undefined or indeterminate region. It has a sense of history since output remains unchanged and stays in the previous state, when input voltage lies in between upper threshold and lower threshold values.

This hysteresis is a highly required phenomenon in certain applications wherein greater noise margins and restoration of stable logic levels need to be established. It is the difference in the output response owing to the change in the direction of input signal. This indicates that, in case of comparator, when a noisy input crosses the threshold point of comparator, it leads to multiple iterative transitions at output node if the latency of comparator is less than the time between abrupt transitions. Therefore, this can efficiently be managed by possessing dual threshold values called upper threshold and lower threshold by Schmitt trigger circuit. At Schmitt trigger output node, in order to cause multiple transitions, the abrupt impulse voltage must be greater than the threshold difference, which makes it more robust towards noise sensitivity. Re-shaping of waveforms, filtering or cleaning up the noise components are typical covering procedures

that exploit Schmitt triggers widely. Schmitt triggers must be used when a square wave form is required to be generated from any kind of noisy input which also includes the conversion of sinusoidal signal to square wave. It also converts slow transition edges to fast transition edges. Therefore this hysteresis is established by incorporating essential positive feedback. Furthermore, V_H also varies according to transistor sizes.

Hysteresis implies that when the input voltage of Schmitt trigger is increased from 0 V to Vdd, it gives a response which differs from response obtained when input voltage is reduced from Vdd to 0 V. When Vin is increased from 0 V to Vdd, Vout stays at Vdd until Vin reaches above upper threshold value while in other case, Vout is at 0 V until Vin comes below lower threshold value. The CMOS Schmitt trigger is shown in Fig. 5.6 [5-10]. Fig. 5.7 gives the Voltage Transfer Characteristic curve of CMOS Schmitt trigger that establishes hysteresis. The upper threshold (V_{th-upper}) and lower threshold (V_{th-lower}) values are determined by pull-down network and pull-up network respectively and their corresponding formulae are given from Fig. 5.7.

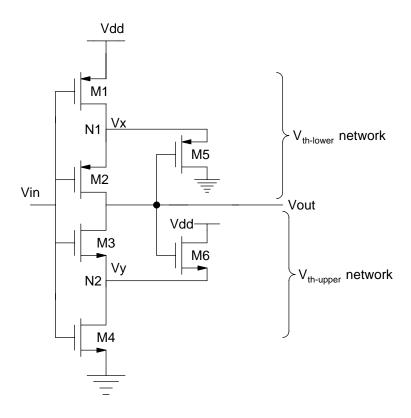


Fig. 5.6 CMOS Schmitt trigger (ST)-1

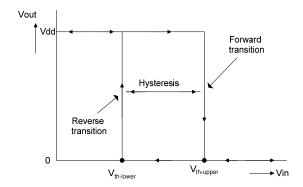


Fig. 5.7 Voltage Transfer Characteristic (VTC) curve of CMOS Schmitt trigger-1

From the VTC, it is clear that $V_{\text{th-lower}} < V_{\text{th-upper}}$ and Hysteresis voltage (VH) = $V_{\text{th-upper}}$ - $V_{\text{th-lower}}$.

Vout = Vdd when $Vin < V_{th-upper}$,

= 0 when $Vin > V_{th-upper}$.

Similarly, Vout = 0 when $Vin > V_{th-lower}$,

= Vdd when $Vin < V_{th-lower}$.

When $(V_{th\text{-lower}} < Vin < V_{th\text{-upper}})$, then Vout = Vout-previous (previous data/history).

Transistors M1, M2 and M5 determine lower threshold voltage and devices M3, M4 and M6 establish upper threshold voltage. M3 and M4 are in series combination and Vin drives both of them. When Vin = 0 then Vout=Vdd and M6 is turned on and it acts as feedback path from supply rail. As long as Vin is increasing, M6 keeps M3 off even after M4 turns on. The mathematical expression for $V_{th-upper}$ is given below.

$$V_{th-upper} = \frac{Vdd + \left(\sqrt{\frac{\beta 4}{\beta 6}}\right)Vtn}{1 + \left(\sqrt{\frac{\beta 4}{\beta 6}}\right)}, \text{ where } \sqrt{\frac{\beta 4}{\beta 6}} \text{ is the beta ratio of M4 and M6 devices}$$

respectively.

Beta calculation: The drain current equation of N-MOSFET is given below.

$$I_{d,linear} = \mu_n C_{ox} \left(\frac{W}{L} \right) \left[\left(V_{gs} - V_t \right) V_{ds} - \frac{{V_{ds}}^2}{2} \right]$$
 (5.3)

The above equation can also be written as

$$\Rightarrow I_{d,linear} = \beta_n \left[\left(V_{gs} - V_t \right) V_{ds} - \frac{V_{ds}^2}{2} \right]$$

Therefore, $\beta_n = \mu_n C_{ox} \left(\frac{W}{L} \right) = \frac{\mu_n \varepsilon}{t_{ox}} \left(\frac{W}{L} \right)$, where $C_{ox} = \frac{\varepsilon}{t_{ox}}$ and similarly,

$$\beta_p = \mu_p C_{ox} = \frac{\mu_p \mathcal{E}}{t_{ox}}$$

 β_n = design variable which is associated with $\left(\frac{W}{L}\right)$ of channel.

 μ_n =mobility of carriers (electrons in case of NMOS and holes in case of PMOS).

 $_{\varepsilon}$ =permittivity of silicon dioxide and t_{α} = gate oxide thickness.

Normally β_n is 2.5 to 3 times of β_p as mobility of NMOS is higher than that of PMOS.

Thus, the variation of beta ratio will affect the characteristics of output.

Therefore
$$\frac{\beta 4}{\beta 6} = \frac{\left(\frac{W}{L}\right)_4}{\left(\frac{W}{L}\right)_6}$$
 (5.4)

In the similar manner, while determining lower threshold voltage, M5 is acting as feedback PMOSFET. The formula for $V_{th-lower}$ is calculated from following equation.

$$V_{th-lower} = \frac{\left(Vdd - |Vtp|\right)\left(\sqrt{\frac{\beta 1}{\beta 5}}\right)}{1 + \left(\sqrt{\frac{\beta 1}{\beta 5}}\right)}, \text{ where } \frac{\beta 1}{\beta 5} = \frac{\left(\frac{W}{L}\right)}{\left(\frac{W}{L}\right)_{5}}$$
(5.5)

The beta ratio characteristics of Schmitt trigger results in the circuit design with large MOSFETs because the transistors connected in series should be made large in order to recompense for the resistance while the threshold voltages are determined by channel

dimensions, $\left(\frac{W}{L}\right)$ of M5 and M6. The simulation waveforms of CMOS Schmitt trigger are shown in Fig. 5.8 and Fig. 5.9.

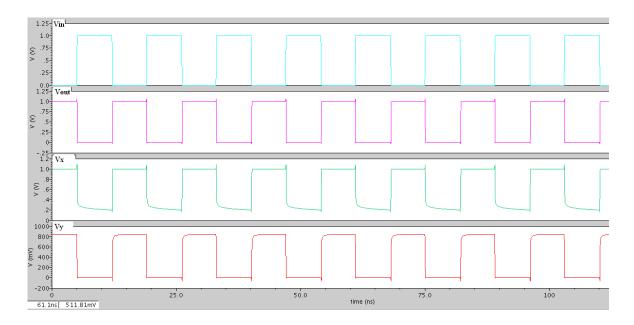


Fig. 5.8 Simulation of transient response of CMOS Schmitt trigger (ST)-1

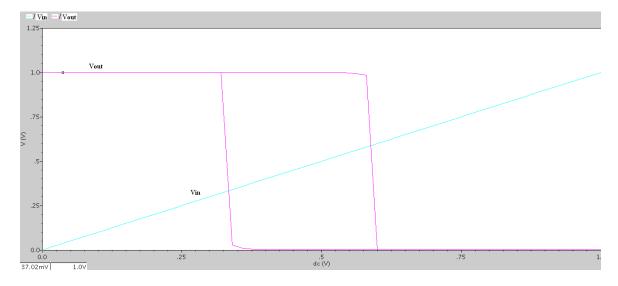


Fig. 5.9 Simulation of DC response (VTC) of CMOS Schmitt trigger (ST)-1

It is observed from the simulations that, in transient response, the inverting operation takes place between Vin and Vout. The voltages at nodes N1 and N2 are Vx and Vy respectively. When Vin is low, M1 and M2 are turned on while M3 and M4 are turned off. Thus M1 and M2 charge output node to Vdd. Vx at node N1 also reaches Vdd since M1 offers low resistance path so that N1 charges to Vdd. Moreover this Vx does not discharge through

M5 which is off as logic high is driving its gate terminal. In pull-down network, M6 is turned on due to its logic high driven gate voltage and thus it passes Vdd to node N2 but due to off status of M4, this Vy cannot be discharged to ground and as a result it stays at 844 mV as long as Vin is zero.

In VTC characteristic curve, balanced input and output characteristics are observed. For slow change in input, there is fast transition at output node which is desired for symmetric response. The power, noise margin calculations are listed in results section.

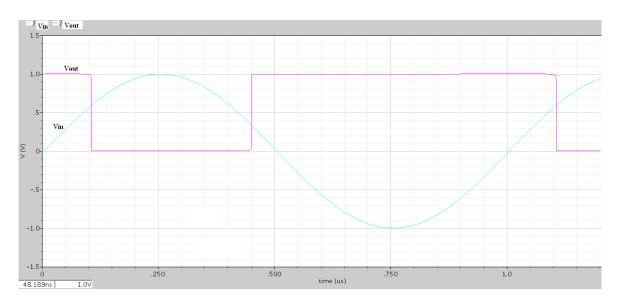


Fig. 5.10 Simulation of sinusoidal response of CMOS Schmitt trigger (ST)-1

Fig. 5.10 portrays the response of CMOS Schmitt trigger for sinusoidal input with 1 GHz frequency. It's evident that it is producing square wave with two stable states for input sinusoidal signal. All the corresponding voltage levels of varying input signal have been quantized to principle binary stable states: '0' and '1'.

The Schmitt trigger circuit-2 to improve the hysteresis further is shown in Fig. 5.11. It uses dual threshold action in determining $V_{\text{th-upper}}$ and $V_{\text{th-lower}}$. The main application behind increasing hysteresis width is to improve its bi-stability phenomenon. The corresponding simulation is plotted in Fig. 5.12.

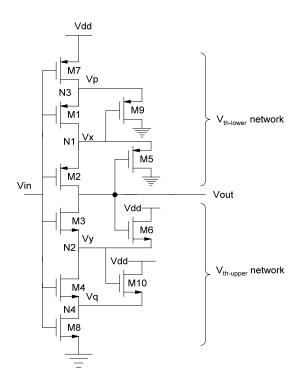


Fig. 5.11 CMOS Schmitt trigger (ST)-2

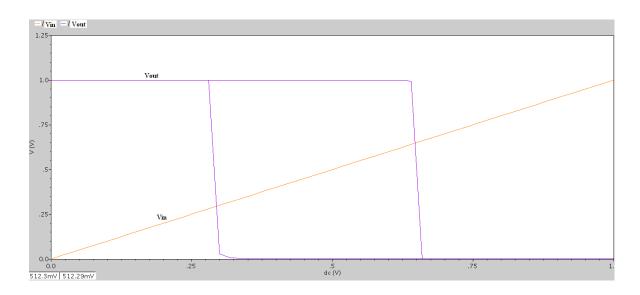


Fig. 5.12 Simulation of DC response (VTC) of CMOS Schmitt trigger (ST)-2

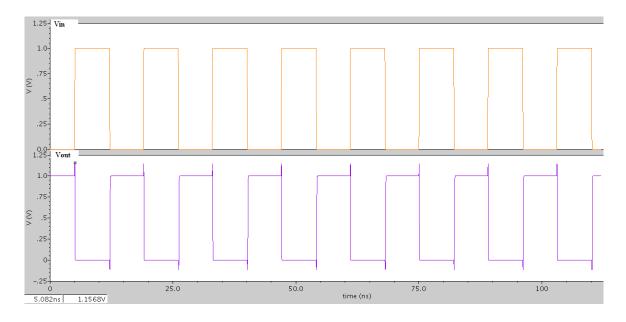


Fig. 5.13 Simulation of transient response of CMOS Schmitt trigger (ST)-2

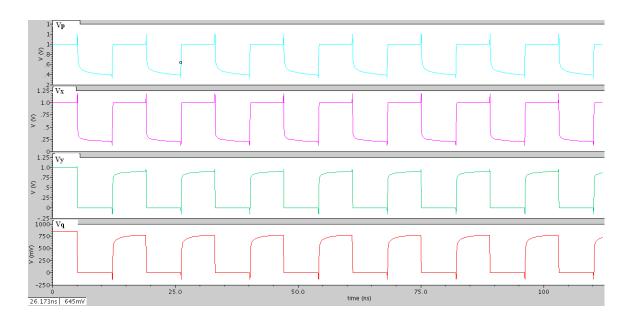


Fig. 5.14 Node voltages of CMOS Schmitt trigger (ST)-2

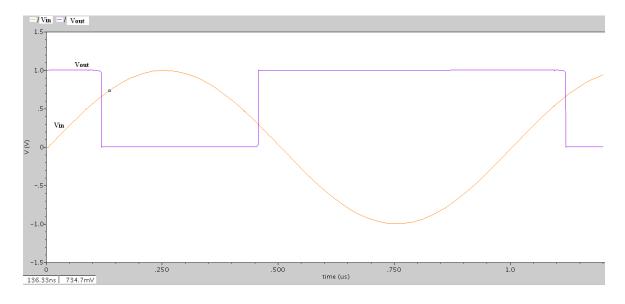


Fig. 5.15 Simulation of sinusoidal response of CMOS Schmitt trigger (ST)-2

It is observed from the simulations that, in transient response, the inverting operation takes place between Vin and Vout. The voltages at nodes N1, N2, N3 and N4 are Vx, Vy, Vp and Vq respectively. When Vin is low, M1, M2 and M7 are turned on while M3, M4 and M8 are turned off. Thus M1, M2 and M7 charge output node up to Vdd. Vp at node N3 and Vx at node N1 also charge to Vdd since M7 and M1 offer low resistance path so that N3 and N1 charge to Vdd. Moreover this Vx and Vp do not discharge through M5 and M9 which are off as logic high is driving their gate terminals. In pull-down network, M6 and M10 are turned on due to their logic high driven gate voltages and thus they pass Vdd to nodes N2 and N4 but due to off status of M8, this Vq cannot be discharged to ground and as a result it stays at 844 mV as long as Vin is zero. In VTC characteristic curve, more symmetric characteristic nature than that of Schmitt trigger-1 is observed with increased hysteresis. Also, there is fast transition at output node for slow change in input. The power, noise margin calculations are listed in results section.

Hysteresis improvement is the main constraint in the process of designing the proposed Schmitt trigger circuits. Thus, the Schmitt trigger circuit-2 to improve the hysteresis further is shown in Fig. 5.11. It uses dual threshold action in determining $V_{th-upper}$ and $V_{th-lower}$. The main application behind increasing hysteresis width is to improve its bi-stability phenomenon. The corresponding simulation is plotted in Fig. 5.12.

Various CMOS Schmitt trigger configurations

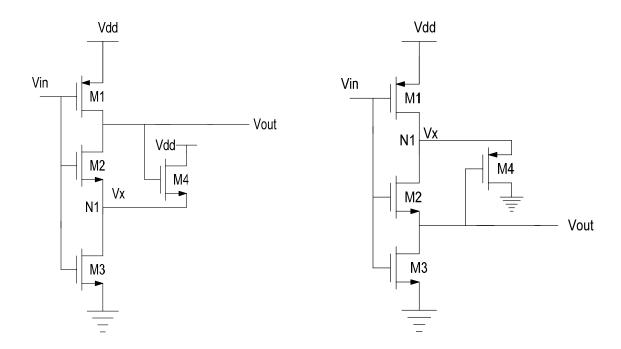


Fig. 5.16 Schmitt trigger (ST)-3

Fig. 5.17 Schmitt trigger (ST)-4

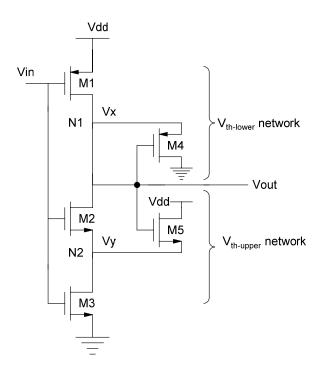


Fig. 5.18 Schmitt trigger (ST)-5

The simulations of various configurations of Schmitt trigger shown in Fig. 5.16, Fig. 5.17, and Fig. 5.18 are presented in Fig. 5.19, Fig. 5.20, and Fig. 5.21 respectively [60-100]. Variable hysteresis phenomena are observed from the simulations. Depending upon the application, these circuits are preferred. For example, Fig. 5.16 consumes less power than basic CMOS Schmitt trigger given in Fig. 5.6.

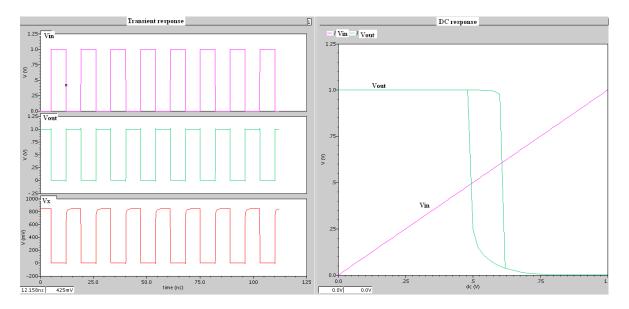


Fig. 5.19 Transient and DC response of Schmitt trigger (ST)-3

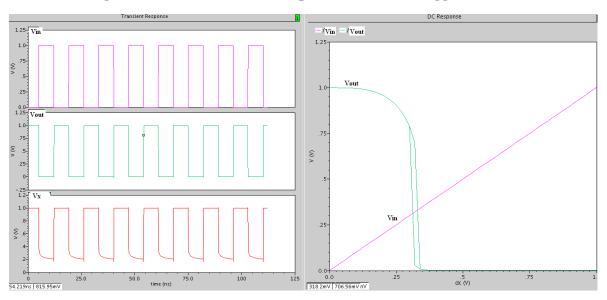


Fig. 5.20 Transient and DC response of Schmitt trigger (ST)-4

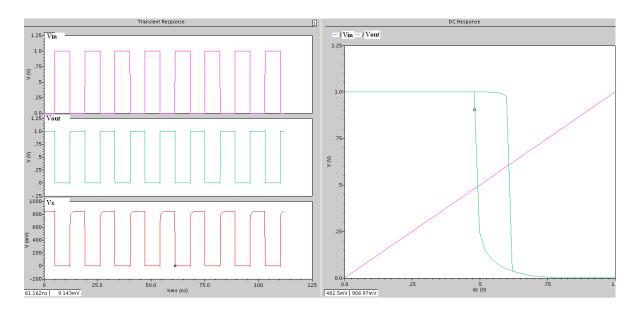


Fig. 5.21 Transient and DC response of Schmitt trigger (ST)-5

Parameters such as total power, leakage power, low noise margin, high noise margin, and hysteresis voltage are computed and corresponding comparison is made.

5.3 Proposed domino Schmitt trigger

In this section, Schmitt trigger circuit is designed using domino logic technique which can efficiently reduce leakages and thereby mitigating leakage power. This section has two proposed circuits with large and zero hysteresis phenomena. These domino logic configurations can explore wide variety of applications mostly in signal conditioners to nullify leakage noise from digital signals. Relaxation oscillators can be implemented using closed loop negative feedback configuration that uses the proposed circuits. Also, function generators and switching power supplies may find applications of domino Schmitt trigger circuits. The proposed domino Schmitt trigger-1 is shown in Fig. 5.22.

5.3.1 Proposed domino Schmitt trigger-1

The proposed technique-1 assures significantly high performance in terms of leakage noise tolerance in domino logic gates which takes the benefits of the utilization of PMOS keeper as shown in Fig. 5.22. The novel scheme is implemented in a footed domino technique as it possesses reduced leakages or static currents than foot-less technique. M1 is the pre-charge PMOS device to charge the dynamic node up to Vdd in clock's pre-charge phase. M5 is the weak keeper transistor connected in feedback manner through domino output node-Z. M2 and M3 are two series connected NMOS devices which are triggered

by applied input. The mechanism of the proposed circuit-1 in pre-charge and evaluation phases is described with equivalent circuit diagrams in Fig. 5.25, Fig. 5.26 and Fig. 5.27.

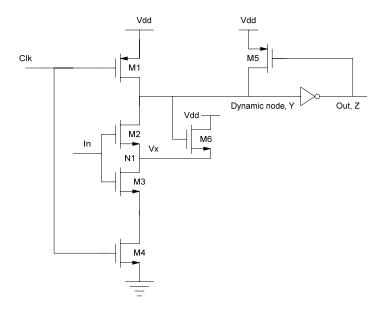


Fig. 5.22 Proposed domino Schmitt trigger-1

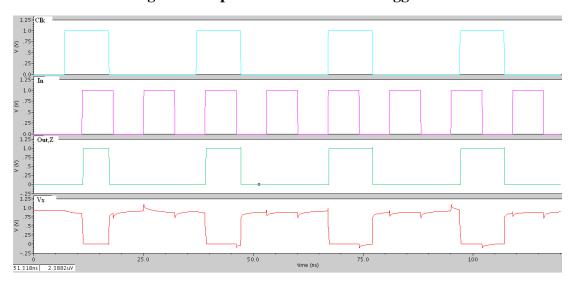


Fig. 5.23 Simulation of transient response of domino Schmitt trigger-1

Clock is the driving element in this domino based Schmitt trigger circuit. Hence the nominal functionality of existing Static CMOS Schmitt trigger could also be achieved by the domino based proposed circuit with increased width of hysteresis voltage. Also, Clock in fundamental sequential circuits is used to synchronize the transitions but whereas in these circuits, clock drives the functionality of Schmitt trigger in its evaluation phase with

increased performance by alleviating the adverse affects which occur in existing Schmitt trigger circuits.

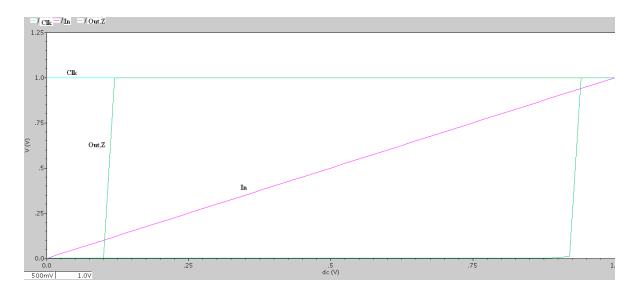


Fig. 5.24 Simulation of DC response (VTC) of domino Schmitt trigger-1

During pre-charge phase:

The pre-charge operation of proposed domino Schmitt trigger is explained in Fig. 5.25. It starts when clock becomes low. The transistor M1 turns on providing a low resistance or conducting path from supply rail to dynamic node and hence dynamic node Y gets charged to Vdd. The PMOS keeper M5, driven by output node Z as shown in Fig. 5.25, also turns on and is providing conducting path from Vdd rail to dynamic node Y, due to inversion operation between nodes Y and Z. As dynamic node Y takes a transition from low-to-high with clock signal at logic zero in pre-charge phase, it turns M6 on and thus Vdd is passed to node N1 and its corresponding nodal voltage Vx charges to 'Vdd-Vtn'. This threshold drop is due to NMOS device switching threshold since NMOS cannot pass stronger one compared to PMOS device. As M4 is off, no discharging is provided and as a result node Y is remaining at logic high state as long as clock is in pre-charge phase. M5 is the conventional PMOS keeper connected in feedback manner between nodes Y and Z respectively which is also providing an alternative conducting path from Vdd to dynamic node Y in pre-charge phase through which any leakages in pull-down network if exist can be managed. These leakages will cause a conducting path from dynamic node to ground

through which the dynamic node is forced to be wrongly discharged. This undesirable effect is pacified in pre-charge phase efficiently with the help of keeper device M5.

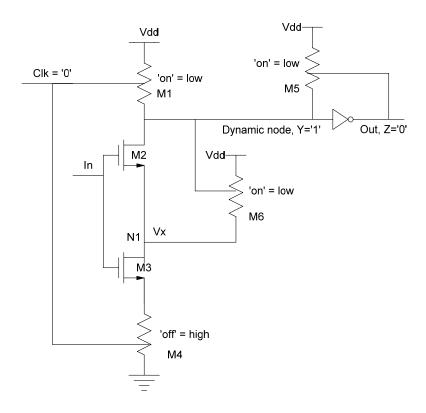


Fig. 5.25 Pre-charge operation of proposed domino Schmitt trigger-1 During evaluation phase:

Case (1): Pull-down network (PDN) is off

This analysis starts with an initial assumption that pull-down network is off (In='0') when clock takes a transition from 0 to 1. Therefore in this phase the node Y is supposed to maintain its pre-charged value since there is no discharging path available as it is assumed that pull-down network, consisting of transistors M2 and M3, is completely off. Node N1 is getting charged to Vdd due to path provided by transistor M6 and as a result Vx becomes 'Vdd-Vtn'. M1 is no longer connected to Vdd. Thus it is turned off disconnecting the conducting path from supply rail to dynamic node as clock takes a transition from low-to-high. Considering the status of PMOS keeper M5, from the Fig. 5.26, it keeps charging dynamic node Y to supply voltage which is required for managing any leakages in pull-down network. As long as the pull-down network is off during the clock's evaluation

period, always strong one at dynamic node and strong zero at domino output are being produced which are highly desired phenomena.

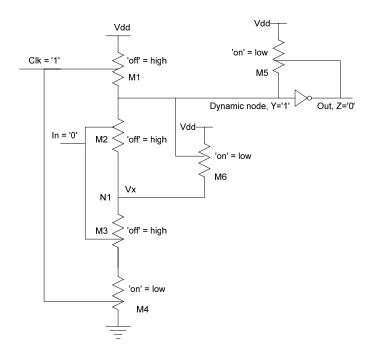


Fig. 5.26 Evaluation phase when PDN is off - operation of proposed domino Schmitt trigger-1

Case (2): Pull-down network (PDN) is on

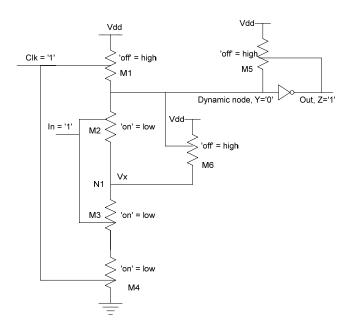


Fig. 5.27 Evaluation phase when PDN is on - operation of proposed domino Schmitt trigger-1

Now let us consider the transitions occurring at the input node of pull-down devices during evaluation phase. From the Fig. 5.27, it is evident that if M2 and M3 are turned on when In='1', then there exists a discharging path provided by pull-down network that discharges the dynamic node completely. Once the discharging process takes place, then the dynamic node is turned to be low which in turn produces a strong one at domino output that will again result in turning the transistor M5 off. This clearly implies that neither M1 nor M5 is providing a low resistance path or conducting path from bias rail to dynamic node.

Therefore, the proposed domino Schmitt trigger-1 is exhibiting improved robustness towards leakages by reducing leakage power effectively. Despite the increment in dynamic power compared to static CMOS Schmitt trigger, it is effective against leakage power along with large hysteresis. During pre-charge phase of clock signal, the dynamic node and domino node are producing strong logic levels without any deterioration in the strength of voltage levels at the corresponding nodes. Similarly during the evaluation period, when the pull-down network is off then the dynamic node is efficiently maintaining its strong one logic level against the leakages existing in pull-down network and also when the pull-down network is on, then the complete discharging phenomenon takes place producing strong zero at dynamic node. In both the cases the proposed scheme is expeditiously giving the required result by mitigating the impact of the sub-threshold leakages in pull-down network.

5.3.2 Proposed domino Schmitt trigger-2

The proposed technique-2, shown in Fig. 5.28, gives zero hysteresis with reduced leakage power. But, the drawback of this scheme is increased dynamic power when compared with proposed domino Schmitt trigger-1. M1 is the pre-charge PMOS device to charge the dynamic node up to Vdd in clock's pre-charge phase. M2 and M3 are two series connected NMOS devices which are triggered by applied input. The mechanism of the proposed circuit-2 in pre-charge and evaluation phases is described with equivalent circuit diagrams in Fig. 5.31, Fig. 5.32 and Fig. 5.33. The voltage Vx initially stays at 1 V in pre-charge phase and it slightly reduces in evaluation mode when pull-down network is off. Vy at node N2 charges up to 844 mV during pre-charge phase. When evaluation started, it drops to 225 mV with pull-down network off. When pull-down network is on, it

completely discharges and both Vx and Vy become zero. The simulations are given in Fig. 5.29 and Fig. 5.30.

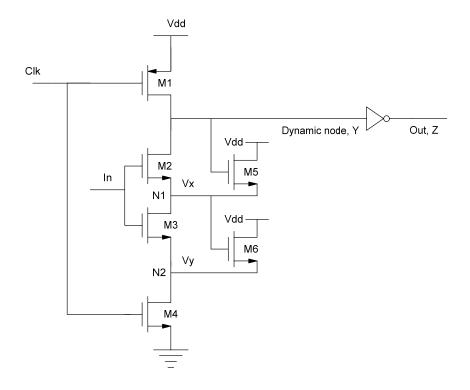


Fig. 5.28 Proposed domino Schmitt trigger-2

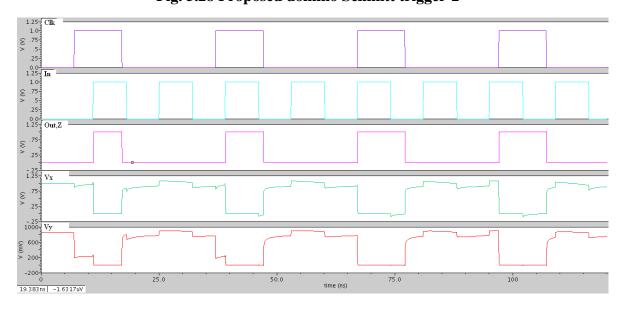


Fig. 5.29 Simulation of transient response of domino Schmitt trigger-2

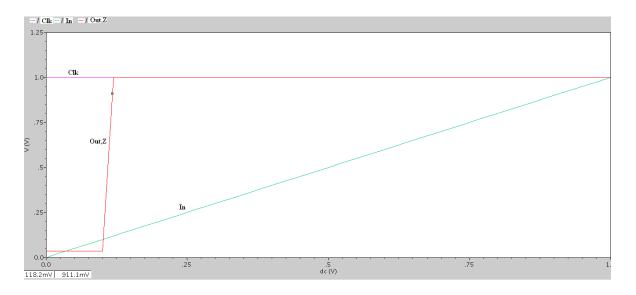


Fig. 5.30 Simulation of DC response (VTC) of domino Schmitt trigger-2 During pre-charge phase:

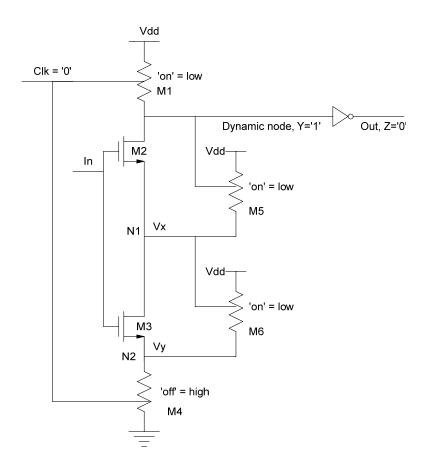


Fig. 5.31 Pre-charge operation of proposed domino Schmitt trigger-2

The pre-charge operation of proposed domino Schmitt trigger is explained in Fig. 5.31. It starts when clock becomes low. The transistor M1 turns on providing a low resistance or conducting path from supply rail to dynamic node so that dynamic node Y gets charged to Vdd. As dynamic node Y takes a transition from low-to-high with clock signal at logic zero in pre-charge phase, it turns M5 on and thus Vdd is passed to node N1 and its corresponding nodal voltage Vx charges to Vdd. This Vx is sufficient to turn M6 on and as a result Vy, nodal voltage at N2, charges to 'Vx-Vtn'This threshold drop is due to NMOS device switching threshold since NMOS cannot pass stronger one compared to PMOS device. As M4 is off, no discharging is provided and as a result node Y is remaining at logic high state as long as clock is in pre-charge phase.

During evaluation phase:

Case (1): Pull-down network (PDN) is off

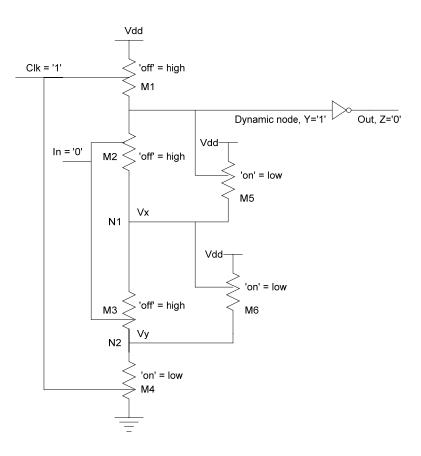


Fig. 5.32 Evaluation phase when PDN is off - operation of proposed domino Schmitt trigger-2

This analysis starts with an initial assumption that pull-down network is off (In='0') when clock takes a transition from 0 to 1. Therefore, in this phase the node Y is supposed to maintain its pre-charged value since there is no discharging path available as it is assumed that pull-down network, consisting of transistors M2 and M3, is completely off. Node N1 is getting charged due to path provided by transistor M5 and as a result Vx becomes 'Vdd-Vtn'. M1 is no longer connected to Vdd. Thus it is turned off disconnecting the conducting path from supply rail to dynamic node as clock takes a transition from low-to-high. As long as the pull-down network is off during the clock's evaluation period, always strong one at dynamic node and strong zero at domino output are being generated. The corresponding equivalent circuit for evaluation phase with pull-down network off is shown in Fig. 5.32.

Case (1): Pull-down network (PDN) is on

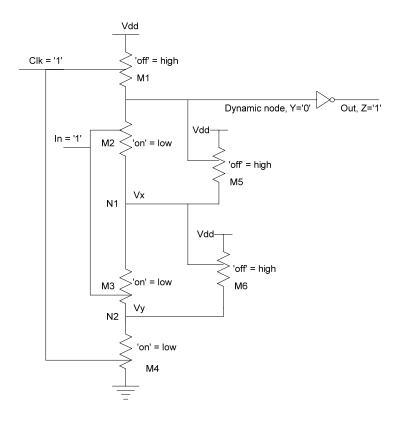


Fig. 5.33 Evaluation phase when PDN is on - operation of proposed domino Schmitt trigger-2

From the Fig. 5.33, it is evident that if M2 and M3 are turned on when In='1', then there exists a discharging path provided by pull-down network that discharges the dynamic node completely. Once the discharging process takes place, then the dynamic node is turned to be low which in turn produces a strong one at domino output. Vx and Vy completely discharge producing zero voltages at nodes N1 and N2 respectively.

Therefore, the proposed domino Schmitt trigger-2 is exhibiting improved robustness towards leakages by reducing leakage power effectively. Despite the increment in dynamic power compared to static proposed domino Schmitt trigger-1, it is effective against leakage power along with zero hysteresis. During pre-charge phase of clock signal, the dynamic node and domino node are producing strong logic levels without any deterioration in the strength of voltage levels at the corresponding nodes. Similarly during the evaluation period, when the pull-down network is off then the dynamic node efficiently maintains its strong one logic level against the leakages existing in pull-down network and also when the pull-down network is on, then the complete discharging phenomenon takes place producing strong zero at dynamic node. In both the cases the proposed scheme is giving the required result by minimizing sub-threshold leakages.

5.4 Results and discussion

All the simulations are done in CMOS 90 nm technology with power supply 1V. Typical parameters such as leakage power, dynamic power, propagation delay, power-delay-product, hysteresis voltage, noise margin level and undefined region calculations are done and comparison is made. Proposed domino Schmitt trigger circuits are exhibiting improved performance from perspective of leakage power, hysteresis voltage and undefined region of operation. The corresponding comparison tabulations are given below. Thus, proposed Domino Schmitt trigger circuit design techniques assure the following benefits.

- (a) Increased noise robustness,
- (b) Minimized undefined region,
- (c) Greater width of Hysteresis voltages and
- (d) Minimized un-wanted flow of leakages.

Table 5.1 Comparison of typical power parameters and power-delay-product of various CMOS and proposed domino Schmitt trigger circuits

Schmitt	Dynamic	Leakage	Total power Total		Power-	
trigger	power (in	power (in	(in W)	propagation	delay-	
	W)	W)		delay (in sec)	product (in	
					W-sec)	
ST-1	0.235E-6	0.069-6	0.304E-6	53.341E-12	12.535E-18	
ST-2	0.347E-6	0.047E-6	0.394E-6	81.674E-12	28.340E-18	
ST-3	0.158E-6	0.022E-6	0.180E-6	31.409E-12	4.962E-18	
ST-4	0.108E-6	0.069E-6	0.177E-6	28.92E-12	3.123E-18	
ST-5	0.177E-6	0.003E-6	0.180E-6	35.50E-12	6.283E-18	
Proposed						
domino ST-1	0.327E-6	0.004E-6	0.331E-6	110.92E-12	36.270E-18	
Proposed						
domino ST-2	2.678E-6	0.011E-6	2.689E-6	64.53E-12	172.811E-18	

Table 5.2 Comparison of noise margin, hysteresis voltage and undefined regions of various CMOS and proposed domino Schmitt trigger circuits

	NM _L (in V)		NM _H (in V)		Hysteresis	Undefined region (in V)	
Schmitt					voltage		
trigger	High-	Low-	High-	Low-to-	(in V)	High-	Low-to-
	to-low	to-high	to-low	high		to-low	high
ST-1	0.58	0.300	0.381	0.661	0.280	0.0396	0.0395
ST-2	0.640	0.260	0.321	0.698	0.378	0.0396	0.0424
ST-3	0.6001	0.4606	0.361	0.449	0.114	0.0396	0.0907
ST-4	0.322	0.2088	0.641	0.772	0.122	0.0372	0.0192
ST-5	0.600	0.460	0.361	0.446	0.112	0.0394	0.0941
Proposed							
dominoST-1	0.0804	0.9203	0.881	0.041	0.8399	0.0392	0.0393
Proposed							
dominoST-2		0.0804		0.8804	0	0	0.0392

Discussion

Table 5.1 shows that proposed domino Schmitt trigger circuits reduce leakage power efficiently when compared to CMOS Schmitt trigger circuits but at the cost of propagation delay. Due to inclusion of clock signal as the driving element and static CMOS at the output node, proposed methods are reducing flow of leakages during pre-charge and evaluation phases consequently. Also, with the increase of number of transistors, there is significant increment in total power and propagation delay.

Table 5.2 reveals that, the undefined region (UR) [5-10] is significantly minimized, allowing the circuit to operate at more noise voltage levels which is highly desired for a noise robust circuit. The least undefined region (UR), between noise margin voltage levels, is observed in proposed domino Schmitt triggers. Noise margin means the set of voltage values or the margin of voltage levels which are to be considered as either logic zero or logic one by the circuit. This implies, when input voltage in increasing from 0 V to Vdd, the circuit takes them as digital binary logic levels.

For example, a circuit being operated at 5 V power supply is defined with V_{IL} and V_{IH} as 2 V and 3 V respectively. Generally, V_{OL} and V_{OH} parameters take ground voltage (0V) and supply voltage (Vdd). The definition of noise margin levels is given in Fig. 5.34.

Low noise margin (NM_L) is the noise region over which all the applied input voltage levels are considered as logic zero. From Fig. 5.34, NM_L = V_{IL} - V_{OL} .

High noise margin (NM_H) is the noise region over which all the applied input voltage levels are considered as logic one. Thus, $NM_H = V_{OH} - V_{IH}$.

Therefore, Low noise margin, $NM_L = V_{IL} - V_{OL} = 2V - 0V = 2V$.

High noise margin, $NM_H = V_{OH}-V_{IH}= 5V-3V=2V$.

Undefined region = V_{IH} - V_{IL} = 3V-2V = 1V.

Therefore, all the input voltages from 0V to 2V are treated as logic zero and the circuit operates accordingly. Similarly, all the input voltages ranging from 3V to 5V (Vdd) are considered as logic one and the circuit functions according to logic one input.

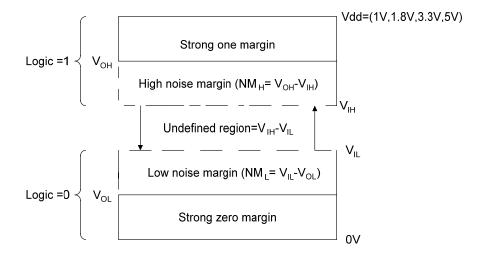


Fig. 5.34 Noise Margin levels

Here, while increasing the input voltage from 2V to 3V, the circuit cannot respond for these voltage levels. Since there may be a chance of getting a voltage level which is lying in the undefined region and to which circuit does not respond as it does not understand what kind of input has been applied at its input terminal. It leads to erroneous operation. Hence alleviating undefined region is highly required in digital circuits by enhancing noise margin levels. If the region between V_{IH} and V_{IL} is minimized, then circuit possesses greater noise immunity. Smaller the undefined region, greater the noise immunity. This is achieved in proposed domino Schmitt trigger circuits and the circuits are operating for all set of input values. Apart from these, the proposed circuits exhibit various hysteresis phenomena with improved noise immunity. The Undefined region existing between the Noise marginal levels, V_{IH} and V_{IL} is also minimized which is highly desired for proper operation of the circuit especially while working with digital logic states. This implies that the proposed circuits provide highly stable input and output logic levels for quantization. Therefore in these aspects the proposed circuits are better than existing ones.

5.5 Conclusion

Therefore, in this chapter, we proposed domino based Schmitt trigger circuits which are more noise tolerant than static CMOS counterpart. Leakage power is effectively minimized in proposed circuits. Furthermore, proposed circuits can operate at all set of input voltage levels as they possess significantly alleviated undefined regions than those of CMOS Schmitt trigger circuits. There is increment in propagation delay and power consumption also.

CHAPTER 6 CHIP TAPE-OUT READY

6.1 Introduction

Test chip is the pre-fabrication stage of VLSI design flow for bringing an electronic chip (e-chip) form to the designed integrated circuit, which will, in next stage referred as final fabrication stage, be sent to manufacturing section of fabrication lab in the form of the photo mask. The basic VLSI design flow is categorized into bottom-up and top-down approaches which are also known as back-end and front-end developments.

Bottom-up approach (back-end design):

It is a kind of development of the application which does not interact with the user directly. It is usually built on the server based platform. Unlike front-end which can directly be communicated by user, it interacts with server and then provides the final results or corresponding data to user. Mechanism of back-end is little bit complex than that of front-end.

Top-down approach (front-end design):

The name itself indicates it is on the front-line of application, which straightaway interacts with user through an interface. It bridges the gap between the back-end environment and front-end designer so that data is provided to user directly. Besides this, it contains many standard libraries for mapping synthesizable designs with concerned library files.

Both the approaches converge at top-level design of VLSI design flow and continue together towards the further steps involved in fabrication procedure. The primary steps involved in these methods are explained in detail section 6.2. Moreover, the test chip of proposed Schmitt trigger circuit is done in CMOS 180 nm process technology through bottom-up approach of basic VLSI design flow. The schematic level implementation, layout of circuit, extraction of parasitic components, transistor level simulations and other related results are presented in section 6.3. Results are given in section 6.4 and concluding remarks are made in section 6.5.

6.2 VLSI design flow

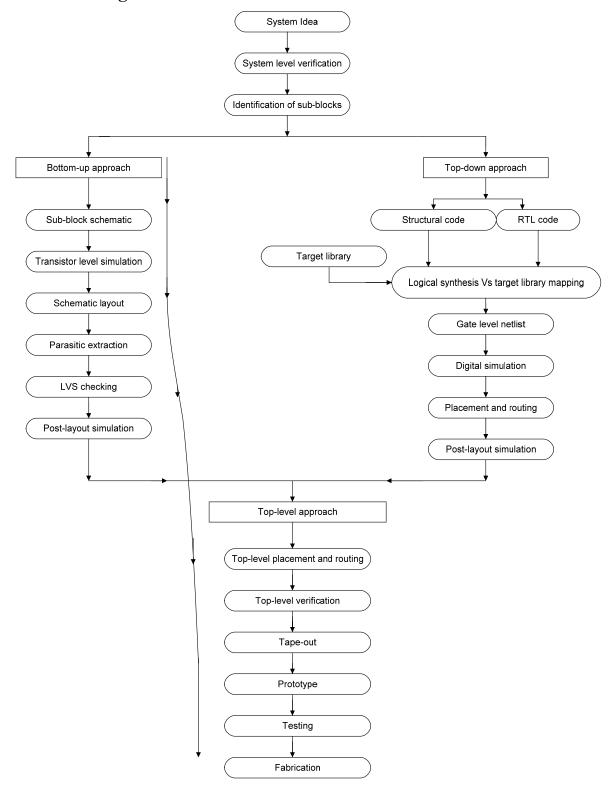


Fig. 6.1 Flow chart of basic VLSI design flow

6.2.1 Bottom-up (Back-end design) approach

Sub-block schematic:

This section gives the schematic implementation of logic circuit at transistor level. The passive and active devices can be taken from analog library of EDA tool and then will be connected to construct logic circuit. Symbols can also be included as connecting devices which can be created from circuit using EDA tool. After placing all the devices in appropriate places, routing is done through wiring. All the input and output nodes are connected to pins with appropriate direction. Every circuit must be biased with DC power supply that varies according to process technology for minimum supply voltage below which the circuit cannot operate. Thus proper bias voltage source needs to be connected and grounding of whole circuit must be ensured.

Transistor level simulation:

The functionality of sub-block schematic circuit needs to be verified which is done at transistor level simulation stage through waveforms of input and output pins. There are two types of verification: functional and formal verification. As this work is carried out in bottom-up approach, only functional verification is possible. Once output is following the functionality of the designed logic circuit, then it can be processed to layout level.

Schematic layout:

This is built on stick diagram of circuit which must be done manually before proceeding for layout. Once stick diagram is finished then standard layout is to be drawn for logic circuit accordingly. Before taking it to layout environment, all the input and bias power supply sources must be removed from schematic circuit and the resultant circuit will be left with input and output pins. Ground and power supply pins need to be replaced by input-output direction pins. The EDA tool will provide all the required layers for drawing layout with distinct colors. Apart from many layers, there are basic layers which are mostly used such as Metal1, Poly, N-Well and Diffusion (both p-type and n-type) layers. Also, for connecting these layers, contacts are used. All these components, including various layers and contacts, possess minimum size to be maintained while drawing in order to reduce the overhead area failing which leads to DRC (Design Rule Check) errors. Thus once layout is drawn then EDA tool will check for DRC errors. All the layers must satisfy, the design

rules which include, minimum width of layers, spacing between layers, type of contacts taken with appropriate size and color, labeling or naming the corresponding nodes and others.

Parasitic extraction:

After DRC is finished without errors, the layout will be taken to next stage called parasitic extraction for extracting parasitic components which include inbuilt resistors and capacitors at various nodes inside the circuit. In spite of being very small in magnitude, these will contribute notable impact on the performance of circuit through signal integrity issues which are discussed in chapter 4. For example, the major inevitable problems like charge sharing, leakage currents, capacitive coupling are few consequences due to parasitic components.

LVS check:

Layout versus schematic checking is required for checking the dissimilarities between the schematic of circuit and layout version. LVS should match so that post-layout simulation is possible. LVS checking needs matching of all the nodes, connections, appropriate layers, input-output pins and their corresponding names. When LVS checking is matched properly then post level simulation is done.

Post-layout simulation:

The name itself indicates that the simulation usually done after the layout is drawn. The significance of repeating simulation procedure again at this level is to investigate the impact of circuit parasitic components on the performance of designed logic circuit functionality. It is the verification of complete design with all constraints added after creation of layout. Therefore, if any error is encountered at this stage, then circuit needs to be amended to nullify the concerned error. Normally, owing to the presence of circuit parasites, more power is consumed when compared to simulation at schematic level. Moreover, propagation delay may also be increased. Thus, having finished this post-layout simulation, the bottom-up approach along with top-down approach converges to top-level approach and will be subjected to various top-level ambient conditions involved in final

fabrication process of manufacturing the electronic circuit chip. The top-down approach is briefed up in sub-section 6.2.2.

6.2.2 Top-down (Front-end design) approach

Firstly, coding is done using hardware description languages such as VHDL or Verilog by the designer out of which the gate level net-list will be generated. The coding can be of any type. The structural and RTL code will be synthesized with the help of target library standard files. The synthesis results in creation of gate level net-list which needs to undergo digital simulation. This is functional verification. Test bench can also be written for functional verification which is more effective than normal digital simulation. After digital simulation, placement and routing that includes positioning the components on chip and connecting all the I/O ports with pins using wiring, is done after which post-layout simulation is processed. This implies that post-layout simulation is common procedure before frond-end and back-end designs converge into top-level approach.

6.2.3 Top-level approach

Top-level approach is the final stage in the VLSI flow that includes making of test chip for circuit design. It starts with placement and routing to make sure that all the I/O pins and devices are properly placed and connected. Then verification of design takes place and is common phenomenon which, in fact, needs to be done at each and every level to make design more accurate and faultless. Next comes, chip tape-out which generates a photo mask of designed circuit that will be sent to final fabrication of chip. Chip tape-out results in prototype chip which is a sample model for testing the designed circuit to even enhance the accuracy and precision further. Once it gets tested and is found to be error-free then manufacturing engineers will go for final fabrication of test-chip.

The complete VLSI flow is shown in Fig. 6.1 where this test-chip of proposed circuit is done using back-end approach. All the related data is provided along with design specifications which include dimensions of devices taken (width and length of channel), supply voltage and applied input voltage.

6.3 Test chip of proposed domino Schmitt trigger circuit-1

Sub-block schematic:

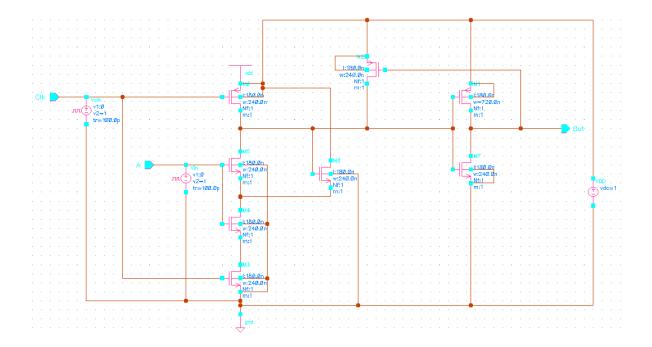


Fig. 6.2 Schematic implementation of proposed domino Schmitt trigger-1

Fig. 6.2 shows the schematic implementation of proposed Schmitt trigger-1 in EDA tool environment. All the components are taken and are connected through wiring. Input voltage sources are connected to clock (Clk) and input (A) nodes which act as input pins. DC power supply of 1V is connected to node Vdd. Output node is connected to 'Out' pin.

Transistor level simulation:

The functionality of circuit shown in Fig. 6.2 is verified at transistor level simulation stage through waveforms of input and output pins. The corresponding simulation wave form windows are given in Fig. 6.3 and Fig. 6.4 respectively. It is evident from the simulations that the output is error free and is following the functionality of the designed logic circuit properly. In transient analysis, the duration for running the simulation is 210 nanoseconds on time axis so that the operation includes 7 periodic cycles of clock input which makes use of verifying output in all possible combinations with respect to input signal which can be seen in Fig. 6.3 and in DC analysis the input (A) voltage is varied from 0V to 1V linearly, shown in Fig. 6.4, and plotted the corresponding hysteresis behavior of proposed circuit.

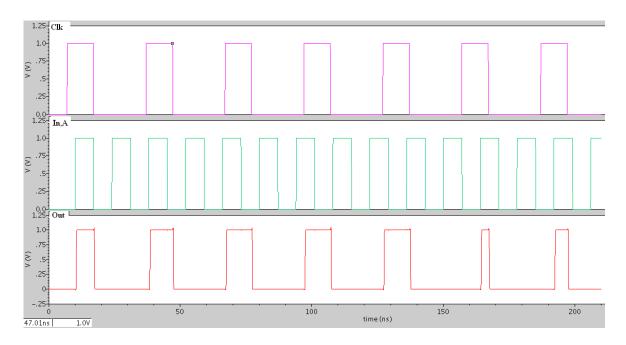


Fig. 6.3 Transistor level simulation of proposed domino Schmitt trigger-1-transient response

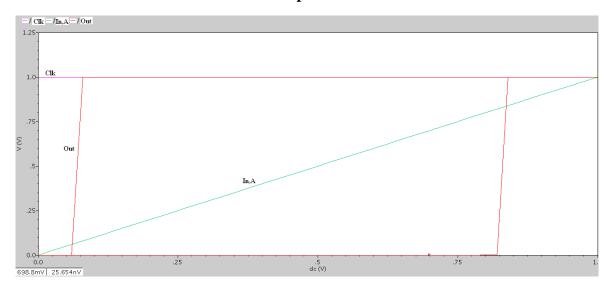


Fig. 6.4 Transistor level simulation of proposed domino Schmitt trigger-1-DC response

Schematic layout:

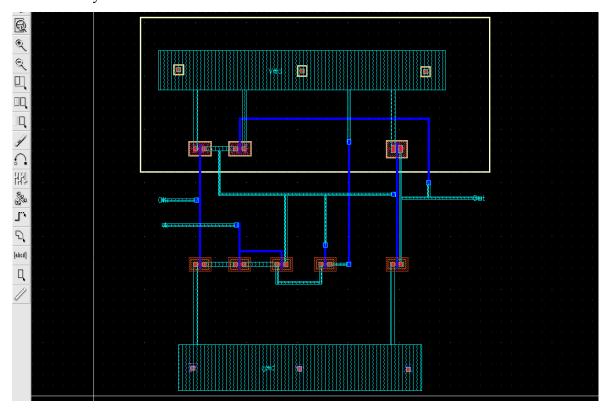


Fig. 6.5 Schematic layout of proposed domino Schmitt trigger-1

Checking for DRC:

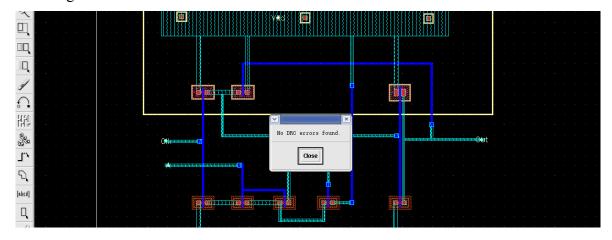


Fig. 6.6 DRC report of proposed domino Schmitt trigger-1

From Fig. 6.6, it is clear that, no DRC errors are found which indicates that the layout drawn for proposed circuit is error free.

Parasitic extraction:



Fig. 6.7 Parasitic extraction report of proposed domino Schmitt trigger-1

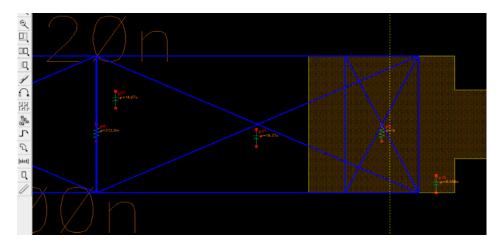


Fig. 6.8 Parasitic components (resistors and capacitors) of extracted layout of proposed domino Schmitt trigger-1

Fig. 6.7 and Fig. 6.8 show the parasitic extraction report, after DRC of layout is finished without errors. The parasitic components which include inbuilt resistors and capacitors can be seen at various nodes inside the circuit from Fig. 6.8. This will contribute significant role on signal integrity issues.

LVS (Layout Versus Schematic) check:

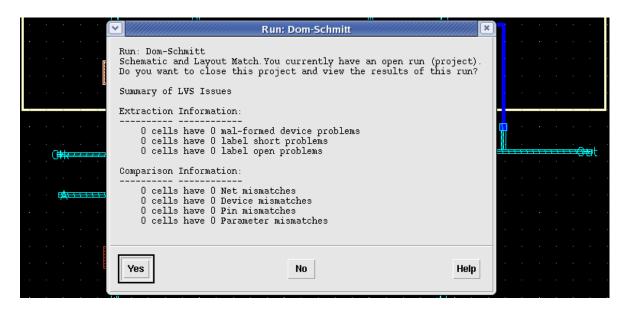


Fig. 6.9 Summary report after LVS check for proposed domino Schmitt trigger-1

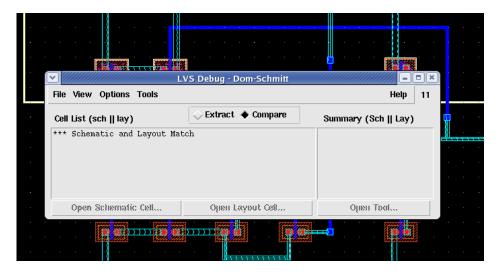


Fig. 6.10 Report of LVS check for proposed domino Schmitt trigger-1

LVS checking is done and reports are given in Fig. 6.9 and Fig. 6.10 which show that layout of Schmitt trigger-1 and its corresponding schematic are matched in all aspects which include layers, pins, labels and other circuit ports.

Post-layout simulation:

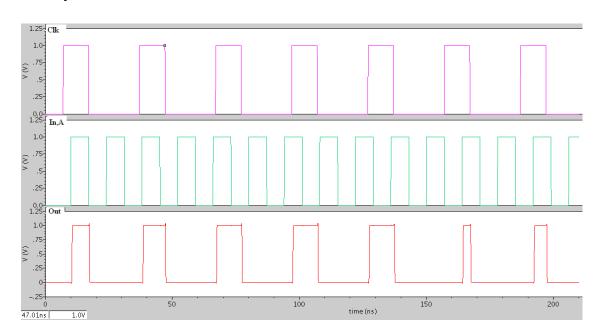


Fig. 6.11 Post-layout simulation of proposed domino Schmitt trigger-1-transient response

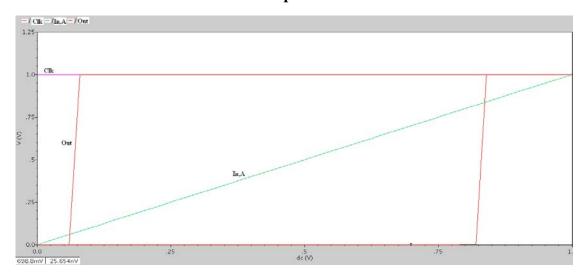


Fig. 6.12 Post-layout simulation of proposed domino Schmitt trigger-1-DC (VTC) response

Post-layout simulation is done after the layout is drawn. Fig. 6.11 and Fig. 6.12 represent the resultant response of post-layout simulation. The significance of repeating simulation procedure again at this level is investigated and concluded that there is impact of circuit parasitic components on the overall performance of designed Schmitt trigger functionality. It is done after adding all constraints once layout is created. Owing to the presence of

circuit parasitic components, it is observed that more power is consumed when compared to simulation at schematic level and the comparison of pre-layout and post-layout simulations is made. Moreover, circuit propagation delay in post-layout simulation is greater than that of schematic level. Therefore, having finished this post-layout simulation, the bottom-up approach is completely verified with proposed Schmitt trigger circuit and results are presented.

Chip tape-out ready of proposed domino Schmitt trigger-1:

The final photo mask of chip-tape out of proposed domino Schmitt trigger-1 circuit is presented in Fig. 6.13. All the inputs including driving element clock signal, outputs, bias supply, and ground ports are connected to corresponding pins of final chip.

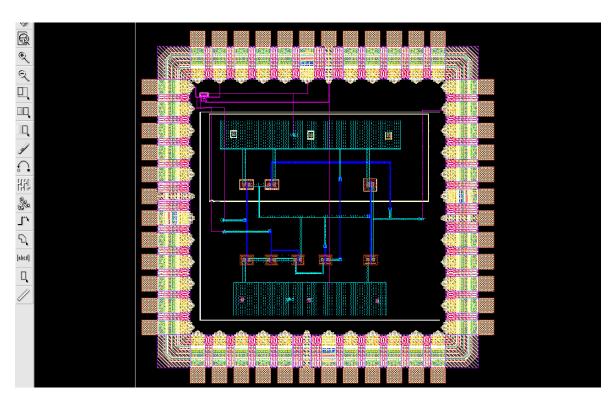


Fig. 6.13 Final chip tape-out ready of proposed domino Schmitt trigger-1 circuit

Top-level approach:

Placement and routing after post-layout simulation is done at top-level, for re-arranging components with wiring procedure and others connections, which is to be verified once again.

Chip tape-out:

Thus, final chip-tape out is made for proposed domino Schmitt trigger-1 and its photo mask will next be sent to fabrication lab for manufacturing IC chip. Tape-out is prefabrication stage where the photo mask of final result is given to the manufacturers. Then, physical fabrication will be done which includes steps such as wafer marching, packing, testing and finally followed by design of physical Integrated circuit chip.

Fabrication:

A fine response at post-layout simulation level does not always ensure a productive outcome since the verification of real-time performance of the chip could only be done by testing the fabricated prototype. The final step in IC design flow is adding pads to the circuit. Once it is done successfully, the chip will then be sent to fabrication lab for manufacturing. In spite of using the step called parasitic extraction for identifying the realistic performance of circuit conditions extended to a higher order degree from the actual photo mask layout level, most of the extraction procedures along with simulation models, employed especially in latest advanced design tools, possess inevitable boundary restrictions which must always be one of the main design considerations from very basic level.

Final testing:

Having finished the manufacturing of chip, final testing has to be carried out once again for checking its over functionality and if the response is matched with circuit desired functionality then that IC will be released in market and will be brought to customer.

6.4 Results and discussion

The designed domino Schmitt trigger-1 is biased with DC supply voltage of 1V in CMOS 180 nm process technology. PMOS devices are taken with increased width of channel than the width of NMOS channel about 3 times. Both pulse and DC voltages of 1 V are applied at input nodes of circuit. Table 6.1 shows comparison of design parameters of domino Schmitt trigger circuit-1 simulated at different process technologies.

Table 6.1 Comparison of design parameters of domino Schmitt trigger-1 at CMOS 90 nm and CMOS 180 nm process technology

Parameter	Domino Schmitt trigger-1		
	180 nm	90 nm	
Dynamic power consumption(in W)	0.444E-6	0.327E-6	
Static power consumption(in W)	4.042E-9	4.7E-9	
Total power consumption(in W)	0.448E-6	0.331E-6	
Propagation delay(in sec)	384.26E-12	110.92E-12	
Power-delay-product (in W-sec)	170.611E-18	36.270E-18	
Peak power(in W)	14.280E-6	31.924E-6	
Upper threshold voltage(in V)	0.840	0.939	
Lower threshold voltage(in V)	0.059	0.100	
Hysteresis width(in V)	0.781	0.839	

Table 6.2 Comparison of design parameters of domino Schmitt trigger-1 at prelayout and post-layout simulation stages

	Transistor level simulation		Post-layout simulation	
Parameter	180 nm	90 nm	180 nm	90 nm
Dynamic power (in W)	0.444E-6	0.327E-6	0.486E-6	0.335E-6
Static power (in W)	4.042E-9	4.7E-9	4.85E-9	5.4E-9
Total power (in W)	0.448E-6	0.331E-6	0.490E-6	0.340E-6
Propagation delay(in sec)	384.26E-12	110.92E-12	388.86E-12	113.65E-12
Power-delay-product(in W-	170.611E-	36.270E-18	188.985E-18	38.072E-18
sec)	18			

Discussion

From Table 6.1, it is clear that, the propagation delay and power-delay-product in 90 nm process technology are less than those of 180 nm process technology. This is due to reduced transit time for carriers across the length of channel which results in increased speed of operation of device. Besides transient analysis, peak power and hysteresis width are observed to be increased in DC analysis in 90 nm process technology.

Table 6.2 compares the computed design parameters at transistor level (pre-layout) and post-layout simulations. Compared to pre-layout simulation, post-layout simulation results are observed to be increased in all design parameters. This is the resultant of inclusion of extraction process of parasitic components from the layout of proposed circuit. Since, these parasites provide conducting path in off state of device, static power dissipation increases. Moreover, propagation delay and power-delay-product seem to be raised both in 90 nm and 180 nm process technologies in post-layout simulation.

6.5 Conclusion

The domino logic based Schmitt trigger-1 is designed and is simulated in CMOS 180 nm process technology platform and observed significant variation in the simulation results through comparison from the tabulations provided. The circuit analysis is carried out through the dominant performance parameters like power consumption, propagation delay and hysteresis width. Table 6.1 and Table 6.2 give the comparison of design parameters of domino Schmitt trigger-1. The chip tape-out of proposed domino Schmitt trigger circuit-1 is done in CMOS 180 nm process technology. As a whole, in this chapter, the bottom-up approach of the basic VLSI design flow is followed while designing chip tape out as this work is carried out in analog environment that starts the designing of circuits from subblock schematic level and converges to top-level approach after post-level simulation stage. The chip-tape out is being sent for fabrication.

Furthermore, this chip is not yet physically fabricated. We provide the prototype of test-chip ready to vendor so that the final process could be carried out in fabrication lab. In any case, this is not a part of the thesis. However, we will publish fabrication and test results on a later date as a part IP of our lab.

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CHAPTER 7 CONCLUSIONS & FUTURE WORK

7.1 Conclusions

The general conclusions drawn from this thesis and some suggested new directions of research are presented in this chapter. The rapid evolution of VLSI technology established a remarkable transformation in the fabrication industry with its emerging qualities like high speed and low area overhead. Scaling brought impeccable change in the recent trends. Power consumption, noise immunity, speed of operation, area and cost are the predominant parameters that have to be taken into consideration before designing any kind of digital logic circuit technique. There may be a trade-off between any two parameters depending on the situation and application of design. Sometimes, the design techniques might not meet all the mentioned requirements in their application, but still an optimization may be followed in order to proceed further in research areas. Because of high speed and low overhead area domino logic became process of choice for many digital designs. However, the real alarming issues are large power consumption and high noise sensitivity. Hence, there is a need for designing new domino methodology or improving existing techniques to meet the requirements by overcoming the above drawbacks. Therefore, the outcome of research must be able to handle the primary design parameters effectively. Besides this, the designed circuit must exhibit high degree of noise robustness.

> General conclusions

The work done in this thesis to achieve above mentioned objective is summarized below

➤ Novel domino topologies

The research area-Domino logic is primarily focused on present working environment from a broader angle with the study of standard logic styles and related work in brief. This thesis proposed novel domino logic circuit techniques to deal with noise and sub-threshold leakages. Furthermore, existing circuits are also modified to improve the response. When compared to existing techniques, proposed logic techniques are very effective in increasing the immunity of system towards noise and sub-threshold leakage issues. Energy-efficient circuits are proposed in chapter 3 which explore conditional

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keepers very efficaciously to sustain from noise glitches and sub-threshold leakage currents. Computations of noise metric parameters like UNG and ANTE are done for all proposed circuit techniques to examine how robust they are for noise affects. From the calculations of UNG, it is evident that the proposed domino circuit techniques are exhibiting superior noise metric parameter-UNG which is almost nearer to unity factor (1) that reveals its improved noise robustness. The proposed circuits score over existing circuits in terms of noise and leakage reduction and this is established through corresponding UNG and ANTE comparisons.

Process corner analysis

Moreover, all the existing and proposed circuits are subjected to distinct process corners like NN, FF, SS, FS and SF, to investigate the change in the performance. Having done the process corner analysis, it is observed that despite few limitations and boundary conditions, proposed circuit techniques are exhibiting high degree of noise robustness with reduced leakage power. Furthermore, the trade-off between power consumption and propagation delay is effectively optimized. All the computational results are tabulated and comparison is made followed by analysis through discussion. Results reveal that proposed techniques are very strong enough against noise glitches.

Signal integrity issues

Besides these techniques, there is other proposed circuit technique for stabilizing dynamic node which is presented in chapter 4 to avoid the leakage drop from dynamic node. It also acts as low power technique with increased number of PMOS stack transistors. Lector based domino technique also reduced leakage power and it increased UNG and ANTE noise parameters.

> Schmitt trigger using domino

Domino based Schmitt trigger circuits are also proposed with large hysteresis and zero hysteresis voltages which find variety of applications. Particularly large hysteresis domino Schmitt trigger circuit is effective in improving noisy immunity of the system while the other Schmitt trigger with zero hysteresis finds applications in I/O pads. Both are exhibiting superior performance in primary design parameters like dynamic power, leakage power, propagation delay, hysteresis width and noise metric parameters which include UNG and ANTE.

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> Test chip tape-out ready

A test chip of Domino Schmitt Trigger is also fabricated using UMC 180 nm technology.

The complete analysis of versatile domino logic circuits is carried out at different process ambient conditions. Existing Domino logic circuits are investigated thoroughly and issues related to this logic style are brought out. The significant variations in design parameters are also noted in all circuit techniques through simulation while improving the performance in various aspects. There are a few pros and cons with proposed techniques in certain parameters. In certain cases while improving a particular design parameter, there is notable deterioration in quality of other parameters which may be worsening the situation, causing drawbacks. However, we have come up with circuits that facilitate leakage reduction, increase the speed of operation, improved noise robustness and low power consumption further. Furthermore, novel Schmitt trigger circuits are designed indicates superior performance in terms of power consumption, propagation delay and hysteresis width.

7.2 Future work

There are ample opportunities to explore advanced research based on this work. This work can further be extended to improve the performance in all design parameters. Complete elimination of signal integrity issues like charge leakage, charge sharing and capacitive coupling effects is not possible as they are inevitable in all circuits. However, mitigating them in a comparative manner with existing schemes facilitates to proceed further to explore improved quality of research. Proposed circuit techniques too have few limitations at some process corners which can be investigated thoroughly and can be improved further striving for the best outcome.

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