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Data parallel patterns in Erlang/OpenCL

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Alla mia famiglia

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Chapter 1

Introduction

This thesis presents the Skel OCL library for the Erlang programming language [3][13] implementing data parallel patterns using OpenCL [6][17], to tap into Graphic Processing Units (GPU) tremendous parallel processing abilities.

In the following sections we introduce the motivation (1.1) for such a library, then we will continue describing the goals of the project (1.2) and stating the achievements of this thesis (1.3). Finally, the last section (1.4) outlines the rest of the thesis, summarizing the others chapters.

1.1 Background and Motivation

Over the past decade, we witnessed the rise of ubiquitous parallel hardware. Not only CPUs have nowadays multiple cores, but also GPUs, which have stopped being just specialized graphics processors and have also become very capable computing engines sporting a remarkable number of cores (currently in the range of hundreds to thousands).

Since GPUs must, owing to its execution model, execute multi-threaded code with threads operating on different data with the same code (this paradigm is called STMD, Single Thread Multiple Data), they are used as CPU co-processors to exploit data parallelism (definition in 2.1.4) in the computation at hand.

The industry standard programming framework to operate on GPU is OpenCL (see chapter 2). Its main competitor is NVIDIA’s proprietary solution called CUDA [19]; in contrast to OpenCL, CUDA can only work in conjunction with NVIDIA hardware.

In recent years Erlang has seen a rapid increase in adoption thank to its built-in support for concurrency and distribution (among other very desirable features, see chapter 2). Declaredly, Erlang is ill suited for that class of problems for which iterative performance and high arithmetic intensity are prime requirements. This limitation can be overcome harnessing the processing power of GPUs using the OpenCL framework.

Erlang is also the one of the languages on which of the European Paraphrase project [10] is focusing its work. Citing from the project’s website:

The ParaPhrase project aims to produce a new structured design and implementation process for heterogeneous parallel architectures, where developers exploit a variety of parallel patterns to develop component based applications that can be mapped to the available hardware resources, and which may then be dynamically re-mapped to meet application needs and hardware availability.

Therefore extending Erlang so to support GPU programming falls entirely within the scope of the project.

One of OpenCL original goals is to allow developers to write a single program that can run on a wide range of systems, from cell phones, to nodes in massive supercomputers. It manages to achieve this goal exposing the hardware instead of hiding it behind some user-friendly abstraction. Erlang programmers, being Erlang a functional language, are instead accustomed to very high level programming style. Dealing with all the intricacies of the specific hardware platform would be a very rude awakening from the elegance of functional programming.

It is then an imperative to hide most of the tedious aspects of OpenCL programming behind some façade that would make feel the Erlang programmer at home. An answer to this requirement can be found in the methodologies developed by the research on algorithmic skeletons.

The algorithmic skeleton concept has been introduced by Murray Cole with his PhD thesis in 1988 [14]. Algorithmic skeletons were described as pre-defined patterns encapsulating the structure of a parallel computation that are provided to user as building blocks to be used to write applications. Each skeleton corresponded to a single parallelism exploitation pattern. Skeletons have been categorized in several classes [22]; we focus on the ones belonging to the class of data parallelism exploitation patterns. Many skeletons programming frameworks, developed by several academic groups, are already available targeting a range of diverse systems and programming languages (e.g. FastFlow [9], ScepPU [11][16], SkelCL [7], SkeTo [8])

1.2 Goals

The goal of the project developed in this thesis is to provide Erlang with an easy to use, high performance, data parallel skeleton library exploiting GPUs' processing power. The implementation is required to be vendor neutral, so OpenCL is the only possible choice. The skeletons must be easy to use, possibly be a nearly perfect drop-in replacement to standard Erlang functions with similar semantic (minimal disruption principle [15]).

1.3 Achievements

We developed a prototype Erlang library, using OpenCL, implementing the data parallel skeletons to accelerate computations on lists of Erlang floats. The user code for such computations is specified writing OpenCL kernels using the OpenCL C (see 2.1.5) programming language. The library has a very user-friendly API since it is modeled after Erlang's list module, with which every Erlang programmer is

familiar. The outcome of the tests (see chapter 4) shows, on one hand, a significant performance improvement over a pure Erlang implementation and, on the other hand, an acceptable performance loss over a C++/OpenCL reference implementation.

1.4 Thesis Outline

The rest of the thesis has the following structure:

- *Chapter 2* introduces and then give some details of, the technologies used in the project: OpenCL and Erlang. In particular it discuss the mechanism provided by Erlang to implement native functions in C that we used to implement the library.
- *Chapter 3* presents the Skel OCL library, starting from the architectural design, continuing with the implementation and ending with an example of its use.
- *Chapter 4* is concerned with the analysis of the experiments carried out to benchmark the library in order to validate its performance. After a description of the test platform, three different applications that use the skeletons provided by the library will be shown. The final paragraph presents and evaluates the result of the tests.
- *Chapter 5* presents the conclusions, recapping the achievements of the thesis and suggesting some directions for future work.

Chapter 2

Background Technologies

This chapter introduces the two technologies used in the project: OpenCL and Erlang. In first part it covers OpenCL, presenting the conceptual foundations of the standard and how NVIDIA implemented it. The second part concerns with Erlang. There is firstly a brief introduction of the platform and the language, then an introduction of Erlang’s mechanism to implement native functions: the NIF API. We concentrate on this API as it is one of the main building blocks of the projects; NIFs are, in fact, the mechanism chosen to use OpenCL from Erlang.

2.1 OpenCL

OpenCL [6] is an industry standard framework for programming computers composed of a combination of CPUs, GPUs, and other processors. These so-called heterogeneous systems have become an important class of platforms, and OpenCL is the first industry standard that directly addresses these architectures. It is a framework for parallel programming and it includes a language, API, libraries and a runtime system to support software development.

OpenCL delivers high level of portability by exposing the underlying hardware structure: the programmer must explicitly define the platform, its context, and how work is scheduled onto different devices.

OpenCL goal is to provide a low-level hardware abstraction layer for portability so that high-level frameworks can be built on.

First released in 2008, the OpenCL specification version 2.0 is available since November 2013; to date, very few vendors released 2.0 implementation [12], in fact the most supported version of the standard is still version 1.1.

Let's now take a look at OpenCL conceptual foundations, namely the different models defined by the standard:

- *Platform model* defines a high-level description of the heterogeneous system.
- *Execution model* defines an abstract representation of how streams of instructions execute on the heterogeneous platform.
- *Memory model* defines the collection of memory regions used by OpenCL and how they interact during an OpenCL computation.
- *Programming models* defines the high-level abstractions a programmer uses when designing algorithms to implement an application.

2.1.1 Platform Model

The OpenCL *platform model* defines a high-level representation of any heterogeneous platform used with OpenCL.

An OpenCL platform always includes a single host; it interacts with the environment external to the OpenCL program, including I/O or interaction with the user.

The host is connected to one or more OpenCL *devices*. The device is where the kernels (a stream of instructions) execute, therefore an OpenCL device is referred to as a *compute device*: it can be a CPU, a GPU, a DSP, or any other processor provided by the hardware and supported by the vendor.

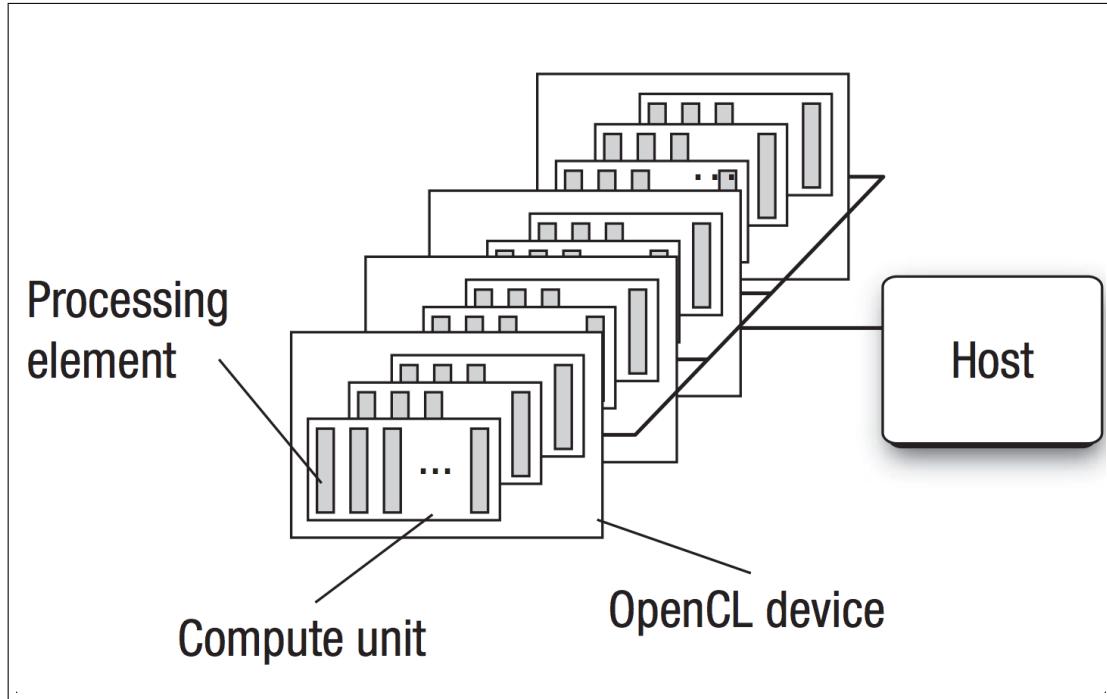


Figure 2.1: The OpenCL platform model with one host and one or more OpenCL devices. Each OpenCL device has one or more compute units, each of which has one or more processing elements.

The OpenCL devices are further divided into *compute units* and these are further divided into one or more *processing elements* (PEs). On-device computations occur within the PEs.

2.1.2 Execution Model

An OpenCL application is comprised of two distinct parts: the *host program* and a collection of one or more *kernels*. The host program runs, unsurprisingly, on the host. OpenCL does not define the details of how such program works, only how it interacts with objects defined within OpenCL.

The kernels execute on the OpenCL devices. Coded in the OpenCL C programming language and compiled with the OpenCL compiler, they do the actual work of an OpenCL application. Kernels are typically simple functions that transform input memory objects into output ones.

The OpenCL execution model defines how the kernels execute.

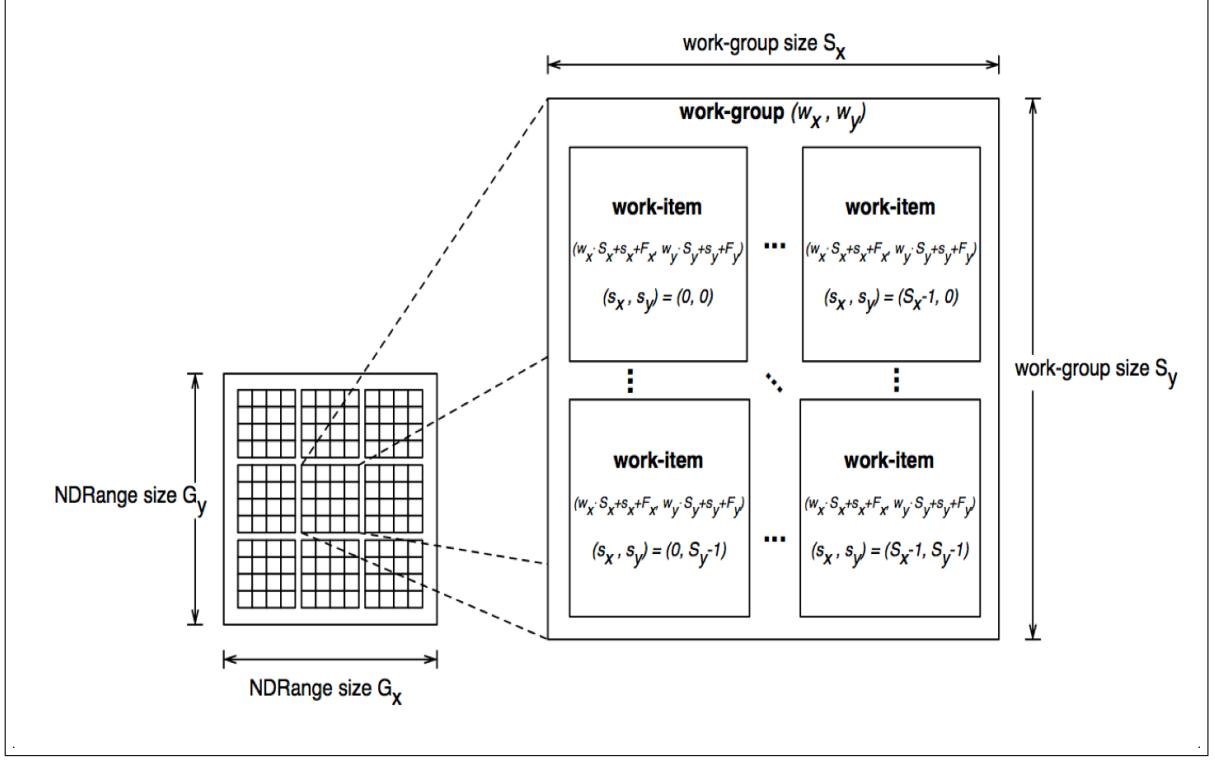


Figure 2.2: An example of an NDRange index space showing work-items, their global IDs and their mapping onto the pair of work-group and local IDs.

The host program issues a command that submits the kernel, which is defined on the host, for execution on an OpenCL device. When this command is issued by the host, the OpenCL runtime system creates an integer index space; an instance of the kernel executes for each point in this index space. Each instance of an executing kernel is called *work-item*, which is identified by its coordinates in the index space. These coordinates are the global ID of the work-item.

Although the sequence of instructions defined by a single kernel is the same for all work-items, the behavior of each one can vary because of branch statements within the code or data selected through the global ID.

Work-items are organized into *work-groups*. The work-groups provide a more coarse-

grained decomposition of the index space and exactly span the global index space. That is, work-groups are the same size in corresponding dimensions, and this size evenly divides the global size in each dimension. Work-groups are assigned a unique ID with the same dimensionality as the index space used for the work-items. Work-items are assigned also a local ID; therefore they can be uniquely identified either by its global ID or by a combination of its local ID and work-group ID.

The index space spans an N-dimensioned range of values and thus is called an `NDRange`. The N in this N-dimensional index space can, currently, be 1, 2, or 3. A `NDRange` is defined by an integer array of length N specifying the size of the index space in each dimension. Each work-item's global and local ID is an N-dimensional tuple.

Work-groups are dispatched to compute units so to be executed by its processing elements.

Context and command queues.

The host defines a context for the execution of the kernels that includes the following resources:

- *Devices*: The collection of OpenCL devices to be used by the host.
- *Kernels*: The OpenCL functions that run on OpenCL devices.
- *Program Objects*: The program source and executable that implement the kernels.
- *Memory Objects*: A set of memory objects visible to the host and the OpenCL devices. Memory objects contain values that can be operated on by instances of a kernel.

The context is created and manipulated by the host using functions from the OpenCL API. The host creates a data structure called *command-queue* to coordinate execution

of the kernels on the devices. The host places commands into the command-queue, which are then scheduled onto the devices within the context. The kinds of commands available are:

- *Kernel execution commands*: Execute a kernel on the processing elements of a device.
- *Memory commands*: Transfer data to, from, or between memory objects, or map and unmap memory objects from the host address space.
- *Synchronization commands*: Constrain the order of execution of commands.

The command-queue schedules commands for execution on a device. These execute asynchronously between the host and the device. The relative order in which commands execute depends on what mode are used for the queue: beginitemize

- *In-order Execution*: Commands are launched in the order they appear in the command-queue and complete in order. In other words, a preceding command on the queue completes before the following command begins. That establishes a serial execution order of commands in a queue.
- *Out-of-order Execution*: Commands are issued in order, but do not wait to complete before following commands execute. The programmer enforces any order constraints through explicit synchronization commands.

Kernel execution and memory commands submitted to a queue generate event objects. These are used to control execution between commands and to coordinate execution between the host and devices.

It is possible to associate multiple queues with a single context that will run concurrently and independently. We used this feature to optimize the performance of the library. In particular, we used two independent queues to concurrently submit commands to the GPU (see 3.2 for details).

2.1.3 Memory model

The execution model tells us how the kernels execute, how they interact with the host, and how they interact with other kernels. The Memory model, instead, define the details of memory objects and the rules to safely use them.

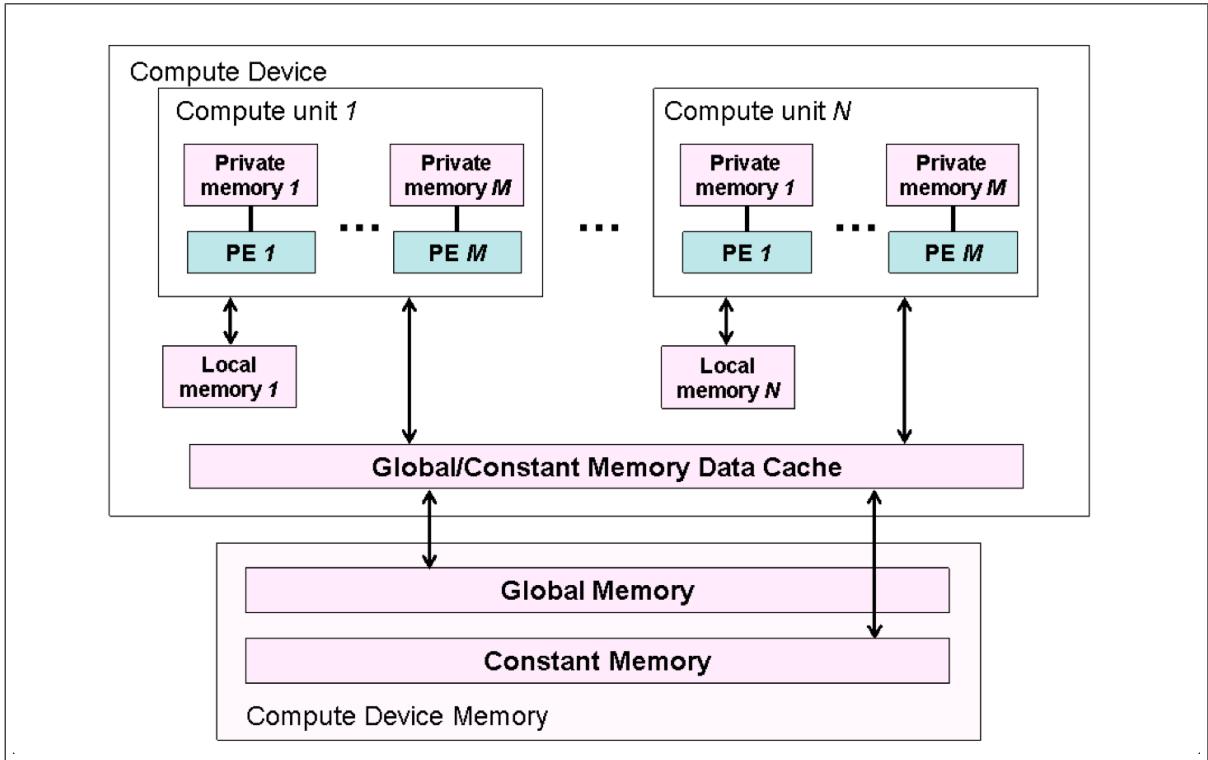


Figure 2.3: The OpenCL conceptual device architecture with processing elements (PE), compute units and devices. The host is not shown.

The OpenCL memory model defines five distinct memory regions:

- *Host memory*: It's visible only to the host. OpenCL defines only how the host memory interacts with OpenCL objects and constructs.
- *Global memory*: It permits read/write access to all work-items in all work-groups. Work-items can read from or write to any element of a memory object in global memory. Reads and writes to global memory may be cached depending on the capabilities of the device.
- *Constant memory*: This memory region of global memory remains constant during the execution of a kernel. The host allocates and initializes memory

objects placed into constant memory. Work-items have read-only access to these objects.

- *Local memory*: This memory region is local to a work-group. This memory region can be used to allocate variables that are shared by all work-items in that work-group. It may be implemented as dedicated regions of memory on the OpenCL device. Alternatively, the local memory region may be mapped onto sections of the global memory.
- *Private memory*: This region of memory is private to a work-item. Variables defined in one work-item's private memory are not visible to other work-items.

These memory regions form a hierarchy as shown in figure 2.3. Memory accesses of a work-item are increasingly slower as it request data from the upper levels, with the global one being the slowest.

OpenCL uses a relaxed consistency memory model: the state of memory visible to a work-item is not guaranteed to be consistent across the collection of work-items at all times.

Within a work-item memory has load/store consistency. Local memory is consistent across work-items in a single work-group at a work-group barrier. Global memory is consistent across work-items in a single work-group at a work-group barrier, but there are no guarantees of memory consistency between different work-groups executing a kernel.

Memory consistency for memory objects shared between enqueued commands is enforced at a synchronization point.

2.1.4 Programming model

The OpenCL execution model supports data parallel and task parallel programming models, as well as supporting hybrids of these two models. Ultimately, what really drives the design of OpenCL is the data parallel one.

In *data-parallel* programming the problem is described in terms of collections of data elements that can be updated in parallel. The parallelism is expressed by concurrently applying the same stream of instructions to each data element: the parallelism is in the data. The index space associated with the OpenCL execution model defines the work-items and how the data maps onto the work-items. In a strictly data parallel model, there is a one-to-one mapping between the work-item and the element in a memory object over which a kernel can be executed in parallel. OpenCL implements a relaxed version of the data parallel programming model where a strict one-to-one mapping is not a requirement.

In particular, OpenCL provides a hierarchical data parallel programming model: there are two ways to specify the hierarchical subdivision. In the *explicit* model a programmer defines the total number of work-items to execute in parallel and also how the work-items are divided among work-groups. In the *implicit* model, a programmer specifies only the total number of work-items to execute in parallel, and the division into work-groups is managed by the OpenCL implementation.

In a *task-parallel* programming model, programmers directly define and manipulate parallel tasks. Problems are decomposed into tasks that can run concurrently, which are then mapped onto processing elements (PEs) of a parallel computer for execution. OpenCL supports task parallel programming defining a model in which a single instance of a kernel is executed independent of any index space. It is logically equivalent to executing a kernel on a compute unit with a work-group containing a single work-item. Under this model, users express parallelism by using vector data types implemented by the device or enqueueing multiple tasks.

2.1.5 The OpenCL C language

OpenCL kernels are written using the OpenCL C language. It is based on the ISO/IEC 9899:1999 C language specification (referred to in short as C99) with some restrictions and specific extensions to the language for parallelism. standard, with some extensions.

The following C99 features (among others) are not supported:

- Function pointers
- Recursion
- Bit field struct members are not supported
- Variable-length arrays and structures with flexible (or unsized) arrays are not supported

Instead, OpenCL adds the following features to C99:

- *Vector data types.* In OpenCL C, vector data types can be used in the same way scalar types are used in C. This makes it much easier for developers to write vector code because similar operators can be used for both vector and scalar data types. It also makes it easy to write portable vector code because the OpenCL compiler is now responsible for mapping the vector operations in OpenCL C to the appropriate vector ISA for a device. Vectorizing code also helps improve memory bandwidth because of regular memory accesses and better coalescing of these memory accesses.
- *Address space qualifiers.* OpenCL devices such as GPUs implement a memory hierarchy. The address space qualifiers are used to identify a specific memory region in the hierarchy.
- *Additions for parallelism:* These include support for work-items, work-groups, and synchronization between work-items in a work-group.
- *An extensive set of functions* such as math, integer, geometric, functions.

2.1.6 NVIDIA OpenCL implementation

Being the platform on which the project has been developed and tested, equipped with a NVIDIA GPU, let's have a look at NIVIDIA OpenCL implementation.

To date, NVIDIA supports OpenCL up to version 1.1, lagging behind others vendors such as AMD.

NVIDIA's unenthusiastic support to OpenCL is due to the fact that they market their own proprietary standard for GPU computing, named CUDA [19], which is OpenCL's main competitor.

NVIDIA's OpenCL implementation is, therefore, based on CUDA: e.g. Kernels written in OpenCL C are compiled into PTX, which is CUDA's instruction set.

Let's see in more detail how CUDA architecture maps onto OpenCL one.

A CUDA device is built around a scalable array of multithreaded Streaming Multiprocessors (SMs). A multiprocessor corresponds to an OpenCL compute unit.

An SM executes a CUDA thread for each OpenCL work-item and a thread block for each OpenCL work-group. A kernel is executed over an OpenCL NDRange by a grid of thread blocks.

When an OpenCL program on the host invokes a kernel, the work-groups are enumerated and distributed as thread blocks to the multiprocessors with available execution capacity. The threads of a thread block execute concurrently on one multiprocessor. As thread blocks terminate, new blocks are launched on the vacated multiprocessors.

A multiprocessor is designed to execute hundreds of threads concurrently. To manage such a large amount of threads, it employs an architecture called SIMT (Single-Instruction, Multiple-Thread).

To maximize utilization of its functional units, it leverages thread-level parallelism by using hardware multithreading, more so than instruction-level parallelism within a single thread (instructions are pipelined, but unlike CPU cores they are executed in order and there is no branch prediction and no speculative execution).

The SM creates, manages, schedules, and executes threads in groups of 32 parallel threads called *warps*. Individual threads composing a warp start together at the same program address, but they have their own instruction address counter and register state and are therefore free to branch and execute independently.

When a multiprocessor is given one or more thread blocks to execute, it partitions them into warps that get scheduled by a warp scheduler for execution.

A warp executes one common instruction at a time, so full efficiency is realized when all 32 threads of a warp agree on their execution path. If threads of a warp diverge via a data-dependent conditional branch, the warp serially executes each branch path taken, disabling threads that are not on that path, and when all paths complete, the threads converge back to the same execution path. Branch divergence occurs only within a warp; different warps execute independently regardless of whether they are executing common or disjoint code paths.

The execution context (program counters, registers, etc.) for each warp processed by an SM is maintained on-chip during the entire lifetime of the warp. Switching from one execution context to another has therefore no cost, and at every instruction issue time, the warp scheduler selects a warp that has threads ready to execute (active threads) and issues the next instruction to those threads.

2.2 Erlang

2.2.1 Introduction

Erlang [3] is a functional general-purpose programming language and runtime environment developed by Ericsson since the 1980s; it has built-in support for concurrency, distribution and fault tolerance.

Since its open source release, Erlang has been adopted by many leading telecom and IT companies. Nowadays it is successfully being used in other industries including banking, finance and e-commerce.

Erlang is distributed with a large collection of libraries called OTP (Open Telecom platform) [4] allowing the user to easily develop applications using anything from telecommunication protocols to HTTP servers. It also provides implementations of several patterns that have proven useful in massively concurrent development over the years. Most production Erlang applications are actually Erlang/OTP applications. OTP is also open source and distributed with Erlang.

Although Erlang is a general-purpose language, it is well suited especially for the application domain for which it was designed. This is a niche mostly consisting of distributed, reliable, soft real-time concurrent systems. These types of applications are telecommunication systems, Servers and Database applications, which require soft real-time behavior.

However, the most common class of problems where Erlang is not indicated is characterized by iterative performance being a prime requirement. To such class belong applications like image processing, signal processing and sorting large volumes of data.

It is evident how much Erlang could benefit from GPU raw processing power in solving the very same class of problems for which it is, currently, ill suited.

2.2.2 Erlang Features

Let's now have a look at Erlang's core features:

- *Concurrency* - A process in Erlang is the entity that carries out work. They are fast to create, suspend or terminate being much more lightweight than OS ones. A running Erlang system may even have millions of concurrent processes, each of them having its own memory area, which is not shared with the other processes. The memory each process allocates can change dynamically during its execution upon request. Processes communicate only by asynchronous message passing. The sending of messages is non-blocking; the receiving, however, suspends the process until a matching message is delivered to its mailbox, i.e. a message queue.
- *Distribution* - Erlang is designed, from the ground up, to run in a distributed environment. An Erlang VM is called a node and a distributed Erlang system is just a set of networked Erlang nodes. Every node can create parallel processes running on other nodes in other machines independently from their operating system. Communication between processes residing on different nodes uses the same mechanism used by processes on the same node to communicate.
- *Robustness* - *Erlang* supports an exception detection and recovery mechanism based on the catch/throw-style. It also offers a supervision infrastructure: a process can register to be the supervisor of another, and receive a notification message if the supervised process terminates. The node under supervision can even reside in a different machine; once notified the supervisor can implement its preferred recovery policy.
- *Hot code replacement* - The high availability needed by Erlang's original domain of application, telecom systems, demanded the system not to be halted even during software updates. Thus, Erlang provides a way of replacing running code without stopping the system: the runtime system maintains a global table containing the addresses for all the loaded modules. These addresses are updated whenever new modules replace old ones. Future calls invoke functions

in the new modules. It is also possible for two versions of a module to run simultaneously in a system.

- *Soft real-time* - Erlang supports developing soft real-time applications with response time demands in the order of milliseconds.
- *Memory management* - Memory is managed by the VM automatically using garbage collection techniques. When a process terminates, its memory is simply reclaimed.

Moving to the more pragmatic aspects of the language, Erlang is a dynamically typed language. Variables don't need to be declared before their use, they are single assigned: once bound to a value they cannot be reassigned. Like in the vast majority of functional languages, values are immutable.

Erlang's basic data types are number, atom, function type, binary, reference, process identifier, and port identifier.

- *Atoms* are constant literals, and are similar to the enumerations used in other programming languages.
- *Functions* are first class objects of the language: they are a data type, can be passed as an argument to other functions, or can be a returning result of a function.
- *Binaries* are a reference to an area of raw memory so they are used as an efficient way for storing and transferring large amounts of data.
- *References* are unique values generated on a node for the purpose of identifying entities.
- *Process and port identifiers* are used to uniquely identify different processes and ports. Ports are used to pass binary messages between Erlang nodes and external programs possibly written in other programming languages and residing in another OS process. A port is created and owned by an Erlang process, which is responsible for coordinating all the messages passing through that port.

The native (implemented in C) counterpart of an Erlang function is the so-called NIF (*Natively implemented Function*). The NIF mechanism will be analyzed in the next section, being one of the cornerstones of the project.

Besides its basic data types, Erlang provides the classical functional data structures such as tuples and lists, adding to that also records.

Tuples and lists are collections of any valid Erlang data-type. From a tuple we can only extract a particular element, but a list can be split and combined. Records in Erlang are akin to C structures, having a fixed number of named fields.

The basic units of code in Erlang are *Modules*. Every module has a number of functions that can be called from other modules if the programmer exports them.

Functions are, instead, the basic unit of abstraction; they consist of several clauses. The right clause to be evaluated is chosen at runtime by pattern matching the calling arguments against the specified pattern.

Erlang doesn't provide constructs for looping; recursion is used instead. To reduce stack consumption, tail call optimization is implemented.

2.2.3 Native Implemented Functions (NIF)

Erlang, from version R14B onwards, provides an API to implement native functions, in C/C++, that are called in the same way as the Erlang implemented ones.

Such functions are called NIF (Native Implemented Functions) and, to interact with Erlang's VM, they use the API provided in the `erl_nif` module [2].

A NIF library contains native implementation of some functions of an Erlang module. Each NIF must also have an implementation in Erlang that will be invoked if the function is called before the NIF library has been successfully loaded, so to throw

an exception or, perhaps, provide a fallback implementation if the NIF library is not implemented for some architecture.

Since a native function is executed as a direct extension of the native code of the VM, its execution is not made in a safe environment. That is, if the native function misbehaves, so the whole VM will. In particular, quoting the module's documentation [2]:

- *A native function that crashes will crash the whole VM*
- *An erroneously implemented native function might cause a VM internal state inconsistency, which may cause a crash of the VM, or miscellaneous misbehaviors of the VM at any point after the call to the native function.*
- *A native function that does lengthy work (more than 1ms long) before returning will degrade responsiveness of the VM, and may cause miscellaneous strange behaviors. Such strange behaviors include, but are not limited to, extreme memory usage, and bad load balancing between schedulers. Strange behaviors that might occur due to lengthy work may also vary between OTP releases.*

As workarounds for the last limitation, several methods are suggested in the documentation, depending on the ability to fully control the code to execute in the native function. If that is the case, the best approach is to divide the work into multiple chunks of work and call the native function multiple times; a function is provided for helping with such work division.

If full control of the code is not possible (e.g. calling third-party libraries) then is recommended to dispatch the work to another thread, return from the native function, and wait for the result. The thread can send the result back to the calling thread using message passing.

As another way to solve the problem, in version R17B a new kind of schedulers have been introduced for dealing with long running functions.

A "dirty NIF" is a NIF that cannot be split and cannot execute in a millisecond or less so it performs work that the Erlang runtime cannot handle cleanly. Building the system with the currently experimental support for dirty schedulers, and dispatching the work to a dirty NIF, the user can ignore the problem altogether since such schedulers will take care of everything.

Although the usage of dirty NIFs is generally discouraged, for some applications it may be considered acceptable. For example, if the only work done by an Erlang program is carried out by the NIF, scheduling interferences are not a problem.

NIF API Functionalities

Every interaction between a NIF library and Erlang, is performed through the NIF API functions.

Several functions are provided for the following functionalities:

Handling of Erlang terms

Any Erlang terms can be passed to a NIF as function arguments and be returned as function return values. The terms can only be read, written and queried using API functions.

All terms belong to an environment. The lifetime of a term is controlled by the lifetime of its environment object. All API functions that read or write terms have the environment that the term belongs to, as the first function argument. There are two types of environments:

- *A process bound environment* is passed as the first argument to all NIFs. All function arguments passed to a NIF and its return value will belong to that environment. The environment is only valid in the thread where it was supplied as argument until the NIF returns. It is not only useless but also dangerous to store pointers to process bound environments between NIF calls.
- *A process independent environment* is user created and can be used to store

terms between NIF calls and to send term. A process independent environment with all its terms is valid the user explicitly invalidates it.

Terms can be copied between environments with `enif_make_copy`.

Binaries

Terms of type binary, like every other Erlang term, must be handled only through calls to API functions. Instances of them, however, are always allocated by the user.

A mutable binary must in the end either be freed with `enif_release_binary` or made read-only by transferring it to an Erlang term with `enif_make_binary`. Read-only binaries do not have to be released.

Resource objects

Extensively used in Skel OCL to allow OpenCL objects handling in Erlang programs (see Chapter 3), Resource objects (ROs) allows for a safe way to return pointers to native data structures from a NIF. A resource object is just a block of memory allocated by the user. A handle ("safe pointer") to this memory block can then be returned to Erlang, but it is totally opaque in nature. It can be stored and passed between processes on the same node, but the only real end usage is to pass it back as an argument to a NIF. The NIF can then get back a pointer to the memory block that is guaranteed to still be valid. A resource object will not be deallocated until the VM has garbage collected the last handle term and the user has released the resource (in any order). All resource objects are created as instances of some resource type: this makes resources from different modules to be distinguishable.

Threads and concurrency

A NIF is thread-safe without any explicit synchronization as long as it acts as a pure function and only reads the supplied arguments. As soon as it writes towards a shared state, it must supply its own explicit synchronization. This includes terms in process independent environments that are shared between. Resource objects will also require synchronization if treated as mutable.

Chapter 3

The Skel OCL library

Skel OCL is a prototype for an Erlang data parallel skeleton library that exploits GPUs' processing power using the OpenCL programming framework.

Skel OCL allows the Erlang programmer to use map and reduce data parallel skeletons on lists of Erlang floats. The function to apply to the elements of the list is specified as OpenCL kernel, written in OpenCL C (see 2.1.5).

A lot of care has been taken to make the use of the library as friendly as possible to the Erlang programmer. In fact, as we will see in the examples in the next chapter, map and reduce skeletons are nearly perfect drop-in replacement for the functionally equivalent counterparts contained in Erlang's list module.

In addition to the skeleton NIFs, Skel OCL provides a set of low-level functions that make possible to compose the skeletons at will. Using these functions the user can even define new skeletons; in paragraph 3.4 we will show how easily this can be done.

In this chapter we firstly explain Skel OCL Architectural Design (and give an example of usage), secondly we look at the implementation details. Finally we show how to define, using the functions provided by the library, a new data parallel skeleton, namely the `MapReduceLL` skeleton.

3.1 Architectural Design

Skel OCL is implemented as Erlang NIF library (see 2.3 for an intro to NIFs) where the native implementation of the functions is written in C++; to work with GPUs, these functions uses OpenCL. A components diagram for the library is shown in figure 3.1.

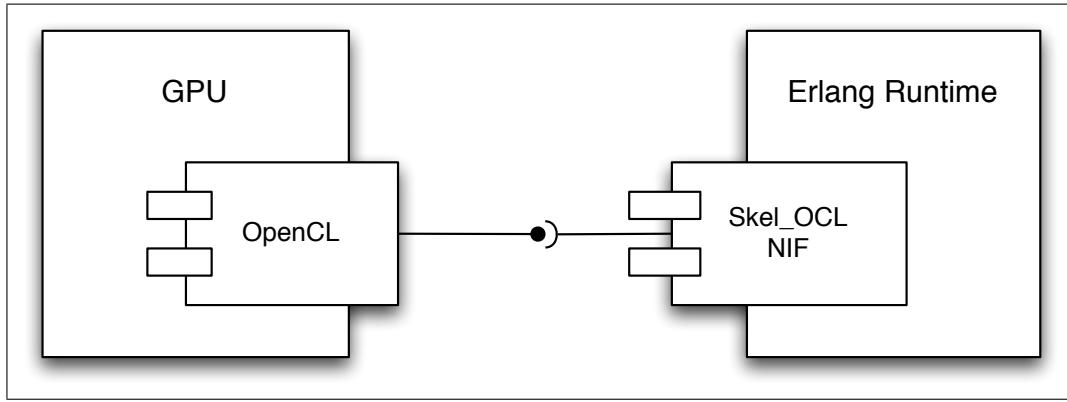


Figure 3.1: Skel OCL components diagram.

A NIF library is a set of natively implemented functions (in C/C++) compiled in a shared library that are loaded at runtime by the Erlang VM. As specified in the documentation [2], each native function must have an Erlang counterpart. These Erlang functions are used as fallback implementation in case the loading of the NIF fails (see also 2.2.3).

In the following sections we details one by one the main components of the library.

3.1.1 The library API: `skel_ocl.erl`

The Erlang interface of the library is defined in the `skel_ocl.erl` file, in which the `skel_ocl` module is contained.

OpenCL bindings

Since OpenCL is used for interacting with GPUs, Skel-OCL provides a set of NIFs to handle OpenCL objects like buffers, programs and kernels (see 2.1 for details on OpenCL). These objects are implemented using the Resource Objects mechanism

(see 2.2.3).

Let's have a look at these NIFs.

- *Program and Kernel NIFs* – In OpenCL the functions are called “kernels”, they run on OpenCL devices and are written in OpenCL C; the program source and executable that implement the kernels are called “program objects”. A program object is like a dynamic library in that it contains a collection of kernel functions. A kernel object instead is like a handle to a function within the dynamic library. Skel OCL provides functions for creating program and kernel objects. Program objects are created reading the code directly from a specified source file. Kernels objects are, instead, created from program objects (same as OpenCL). The specification for these NIFs is shown in figure 3.2

```
%%Program and Kernel Types and functions -----
-type program() :: binary().
-type kernel() :: binary().

%%Build an OpenCL program from the source code in the ProgSrcString iolist()
-spec buildProgramFromString(ProgSrcString) -> {ok, program()} | {error, Why} when
    ProgSrcString :: iolist(),
    Why :: atom().

%%Build an OpenCL program object from the source code in the ProgSrcFile file
-spec buildProgramFromFile(ProgSrcFile) -> {ok, program()} | {error, Why} when
    ProgSrcFile :: nonempty_string(),
    Why :: atom().

%%Create an OpenCL kernel object from a program object.
-spec createKernel(Prog, KerName) -> {ok, kernel()} | {error, Why} when
    Prog :: program(),
    KerName :: nonempty_string(),
    Why :: atom().
```

Figure 3.2: Program and Kernel objects NIFs

- *Buffer NIFs* - OpenCL allows the user to allocate memory buffer either on the host memory or the device memory. This distinction is crucial since the cost of accessing the memory on the device is not negligible, moreover, in order to be used on the host it must be copied first.

Taking that into account, Skel OCL defines different NIFs for allocating host buffers and device buffers. In addition to these functions, NIFs for copying floats

from and to an Erlang list are provided. The functions' specification is shown in figure 3.3

```

%%Buffer Types and functions -----
-type hostBuffer() :: binary().
-type deviceBuffer() :: binary().
-type buffer() :: hostBuffer() | deviceBuffer().

-type rw_flag() :: read | write | read_write.

%%Get the size of Buffer in Bytes
-spec getBufferSize(Buffer) -> {ok, SizeByte} | {error, Why} when
    SizeByte :: non_neg_integer(),
    Buffer :: buffer(),
    Why :: atom().

%%Allocate a buffer on the host of size SizeByte bytes
-spec allocHostBuffer(SizeByte) -> {ok, Buffer} | {error, Why} when
    SizeByte :: non_neg_integer(),
    Buffer :: hostBuffer(),
    Why :: atom().

%%Allocate a buffer on the device of size SizeByte bytes, specifying kernel read/write permissions
-spec allocDeviceBuffer(SizeByte, Flags) -> {ok, Buffer} | {error, Why} when
    SizeByte :: non_neg_integer(),
    Flags :: rw_flag(),
    Buffer :: deviceBuffer(),
    Why :: atom().

%%release a previously allocated buffer
-spec releaseBuffer(Buffer) -> ok when
    Buffer :: buffer().

%%Copy data from a buffer to another having the same size
-spec copyBufferToBuffer(From, To) -> ok | {error, Why} when
    From :: buffer(),
    To :: buffer(),
    Why :: atom().

%%Copy CopySizeByte data from a buffer to another
-spec copyBufferToBuffer(From, To, CopySizeByte) -> ok | {error, Why} when
    From :: buffer(),
    To :: buffer(),
    CopySizeByte :: non_neg_integer(),
    Why :: atom().

%%Copy the content of an Erlang list into a host buffer.
-spec listToBuffer(From, To) -> ok | {error, Why} when
    From :: [float()],
    To :: hostBuffer(),
    Why :: atom().

%%Create a list having as elements the data in the host buffer
-spec bufferToList(From) -> {ok, [float()]} | {error, Why} when
    From :: hostBuffer(),
    Why :: atom().

```

Figure 3.3: Buffer management NIFs

Skeletons NIFs

We now introduce the higher-level NIFs of the library: the *Map and Reduce* Skeleton NIFs. Skeleton NIF functions are named according to the following pattern: to the name of the skeleton two capital letters are appended, the first is the type of the input and the second is the one of the output. The possible types are **List** and **Device buffer**. For example `mapLD` is a map skeleton that takes lists and outputs the result into a device buffer, while `reduceDL` is a reduce skeleton from device buffers to lists.

Skeletons are generic components that can be parameterized with actual computations in the form of user-functions. In Skel OCL such functions are specified in OpenCL C.

NIFs to generate and compile the OpenCL programs are available, one for each type of skeleton. The kernel object returned by these NIFs, is used as functional parameter of the skeletons (see chapter 4 for examples of use).

- *Map* - The map skeleton is functionally equivalent to Erlang's `lists:map` [1] function (see the examples in chapter 4 for a comparison). A user specified function is applied to every element of the input list/buffer and the result of this function is appended to the output list/buffer. In addition to the classical map function, also a binary map (called `map2`) is provided, this skeleton's Erlang counterpart is `lists:zipWith` (see `dotProduct` example in chapter 4). The user-provided binary function is applied to the corresponding elements of the two input lists/buffers and the resulting value is concatenated to the output list/buffer. The specifications for the map NIFs are in figure 3.4.
- *Reduce* – The reduce skeleton is equivalent to the `lists:fold` family of Erlang functions. The fold function, along with map, is the fundamental tool of the functional programmer. The reduce function builds up a return value recombining, through use of a given combining operation, the results of processing its constituent parts in a recursive fashion. The specification for the reduce NIFs are in figure 3.6

```

%%Map functions ----

%%Create a map compatible kernel object from the source in MapSrcFunFile. The kernel name is FunName.
-spec createMapKernel(MapSrcFunFile, FunName) -> {ok, kernel()} | {error, Why} when
    MapSrcFunFile :: nonempty_string(),
    FunName :: nonempty_string(),
    Why :: atom().

%%Create a map compatible kernel object from the source in MapSrcFunFile. The kernel name is FunName, the
%%arity is FunArity
-spec createMapKernel(MapSrcFunFile, FunName, FunArity) -> {ok, kernel()} | {error, Why} when
    MapSrcFunFile :: nonempty_string(),
    FunName :: nonempty_string(),
    FunArity :: pos_integer(),
    Why :: atom().

%%Create a map compatible kernel object from the source in MapSrcFunFile.
%%The kernel name is FunName, the arity is FunArity.
%%Program caching policy must also be specified.
-spec createMapKernel(MapSrcFunFile, FunName, FunArity, CacheKernel) -> {ok, kernel()} | {error, Why} when
    MapSrcFunFile :: nonempty_string(),
    FunName :: nonempty_string(),
    FunArity :: pos_integer(),
    CacheKernel :: cache | no_cache,
    Why :: atom().

%%Map skeleton working on device buffers. User-specified function specified as kernel objects.
-spec mapDD(Kernel, Input, Output) -> ok | {error, Why} when
    Kernel :: kernel(),
    Input :: deviceBuffer(),
    Output :: deviceBuffer(),
    Why :: atom().

%% Map skeleton: input is a list, output too. User-specified function specified as kernel objects.
-spec mapLL(Kernel, Input, InputLength) -> {ok, Result} when
    Kernel :: kernel(),
    Input :: [float()],
    InputLength :: non_neg_integer(),
    Result :: [float()].

%%Map skeleton: input is a list, output is a device buffer. User-specified function specified as kernel objects.
-spec mapLD(Kernel, Input, Output, InputLength) -> ok | {error, Why} when
    Kernel :: kernel(),
    Input :: [float()],
    Output :: deviceBuffer(),
    InputLength :: non_neg_integer(),
    Why :: atom().

%%Binary Map skeleton working on device buffers. User-specified function specified as kernel objects.
%%Similar to lists:zipWith
-spec map2DD(Kernel, Input1, Input2, Output) -> ok | {error, Why} when
    Kernel :: kernel(),
    Input1 :: deviceBuffer(),
    Input2 :: deviceBuffer(),
    Output :: deviceBuffer(),
    Why :: atom().

```

Figure 3.4: Map Skeleton NIFs

```
%%Binary Map skeleton: input is a list, output is a device buffer. User-specified function specified as kernel objects.
-spec map2LD(Kernel, Input1, Input2, Output, InputLength) -> ok | {error, Why} when
    Kernel      :: kernel(),
    Input1     :: [float()],
    Input2     :: [float()],
    Output      :: deviceBuffer(),
    InputLength :: non_neg_integer(),
    Why        :: atom().

%%Binary Map skeleton: input is a list, output too. User-specified function specified as kernel objects.
-spec map2LL(Kernel, Input1, Input2, InputLength) -> {ok, Result} when
    Kernel      :: kernel(),
    Input1     :: [float()],
    Input2     :: [float()],
    InputLength :: non_neg_integer(),
    Result     :: [float()].
```

Figure 3.5: Map Skeleton NIFs (continued)

```
%%Reduce functions ----

%%Create a reduce compatible kernel object from the source in ReduceSrcFunFile.
%%The kernel name is FunName.
-spec createReduceKernel(ReduceSrcFunFile, FunName) -> {ok, kernel()} | {error, Why} when
    ReduceSrcFunFile :: nonempty_string(),
    FunName        :: nonempty_string(),
    Why            :: atom().

%%Create a reduce compatible kernel object from the source in ReduceSrcFunFile.
%%The kernel name is FunName
%%Program caching policy must also be specified.
-spec createReduceKernel(ReduceSrcFunFile, FunName, CacheKernel) -> {ok, kernel()} | {error, Why} when
    ReduceSrcFunFile :: nonempty_string(),
    FunName        :: nonempty_string(),
    CacheKernel    :: cache | no_cache,
    Why            :: atom().

%%Reduce skeleton working on device buffers. User-specified function specified as kernel objects.
-spec reduceDD(Kernel, Input, Output) -> ok | {error, Why} when
    Kernel      :: kernel(),
    Input       :: deviceBuffer(),
    Output      :: deviceBuffer(),
    Why        :: atom().

%%Reduce skeleton: input is a device buffer, output is a list. User-specified function specified as kernel objects.
-spec reduceDL(Kernel, Input) -> {ok, Result} | {error, Why} when
    Kernel      :: kernel(),
    Input       :: deviceBuffer(),
    Result     :: [float()],
    Why        :: atom().

%%Reduce skeleton: both input and output are lists. User function specified as kernel objects.
-spec reduceLL(Kernel, Input, InputLength) -> {ok, Result} | {error, Why} when
    Kernel      :: kernel(),
    Input       :: [float()],
    InputLength :: non_neg_integer(),
    Result     :: [float()],
    Why        :: atom().
```

Figure 3.6: Reduce skeleton NIFs

3.1.2 The Skel OCL shared library: `skel_ocl.so`

The native implementation in C++ of the Skel OCL functions is contained in a compiled shared library as requested by Erlang's NIF mechanism (see 2.2.3). Erlang's VM will load at runtime the library and will delegate to it the execution of NIFs.

Both NIF and OpenCL APIs are actually in C, but being C and C++ compatibles, there isn't any problem in using C++.

The source code file for the Skel OCL is `skel_ocl.cpp`. It includes the code to initialize OpenCL, the implementation of the NIFs and the code for handling the OpenCL objects so to expose them as Resource Objects to Erlang.

3.1.3 OpenCL skeletons

Ultimately, the skeletons are implemented as OpenCL kernel, written in OpenCL. A program is the entity that contains a set of kernel functions. In order to be executed, a program must first be compiled. But, since user-specified functions are just kernels, the code from a program must be generated.

When a user requests the creation of a certain skeleton, Skel OCL generates the source code for the program appending the user defined kernel to the one implementing the skeleton at hand. The generated code is finally compiled using the compiler provided by OpenCL.

Figure 3.7 shows the program code generated for a map skeleton with a user-defined function that computes the square of the input.

```
#pragma OPENCL EXTENSION cl_khr_fp64: enable

double sq(double x) {
    return x*x;
}

_kernel void MapKernel(_global double* input, _global double * output, unsigned int outputOffset,
unsigned int uiNumElements)
{
    _private size_t i = get_global_id(0);
    if(i >= uiNumElements) return;
    output[i + outputOffset] = sq(input[i]);
}
```

Figure 3.7: Generated code for a map skeleton and the user-defined function `sq`

3.2 Example of use

Having introduced Skel OCL and its API, let's now have a look at an example.

In figure 3.8 we use the `mapLL` skeleton and a `sq` kernel to compute the square of each element of a list. We then sum them up thanks to `reduceLL` skeletons and `sum` kernel.

As we can see, using skeleton is quite easy.

Before calling the actual skeletons we need to create the kernel objects representing our user-defined functions (`sq` and `sum` in this case). Kernel objects for a skeleton are created using the `createKernel` function related to that skeleton: (e.g. `createMapKernel` for map and `createReduceKernel` for reduce).

In this example we are using the LL version of the skeletons; that means that each skeleton, to carry out the required computation will:

un-marshal the input list, copy the data on the device, execute the kernel, copy data back from the device and finally marshal the output list.

Knowing this it's clear that, in the cases where we need to use a skeleton's output as input for another one (i.e. composing them), the whole process is extremely inefficient: not only we have to copy intermediate data from and to the device, but also marshal and unmarshal the same data.

Both of these tasks are really expensive, so we must avoid them whenever is possible.

We will show how to do it in the last section, where we use the lower-level functions provided by Skel OCL to efficiently compose skeletons.

```
example(NumVal) ->

    case isPow2(NumVal) of
        false -> erlang:error(input_not_pow2);
        true -> ok
    end,

    MapFunSrcFile = "sq.cl", MapFunName = "sq",
    ReduceFunSrcFile = "sum.cl", ReduceFunName = "sum",

    PWD = element(2,file:get_cwd()) ++ "/",

    MapFunSrcFileAbsPath = PWD ++ MapFunSrcFile,
    ReduceFunSrcFileAbsPath = PWD ++ ReduceFunSrcFile,

    InputList = [ X+0.0 || X <- lists:seq(0, NumVal-1) ], %% [0, ..., NumVal-1]

    %%create Kernels objects
    MapSqKernel = checkResult(createMapKernel(MapFunSrcFileAbsPath, MapFunName)),
    ReduceSumKernel = checkResult(createReduceKernel(ReduceFunSrcFileAbsPath, ReduceFunName)),

    %%Map Skeleton: list -> list
    MapOutputList = checkResult(mapLL(MapSqKernel, InputList, NumVal)),

    %%Reduce skeleton: list -> list
    checkResult(reduceLL(ReduceSumKernel, MapOutputList, NumVal))
.
```

Figure 3.8: Example of Skel OCL usage, mapLL and reduceLL skeletons are used

3.3 Implementation details

In this section, we present some interesting details of Skel OCL's implementation.

Firstly we discuss the choice of using Erlang lists as input data for the skeletons.

Secondly we analyze the implementation of the `mapLL` skeleton because it's interesting as an example of skeleton implementation, and for the techniques adopted

to overlap the handling of Erlang lists with the actual computation.

Finally we examine some of OpenCL's and NIF's shortcomings discovered in developing Skel OCL.

3.3.1 Lists vs. Binaries vs. ROs

One of the firsts design decision was choosing which kind of Erlang data structure to support in our skeletons.

Since one of the goals for the library is to be as friendly as possible to Erlang programmers, it was inevitable to choose Erlang lists as the data structure for input and output data.

Lists are, in fact, the cornerstones of functional programming, and Erlang makes no exception; so overlooking lists is unreasonable.

Keeping in mind that to work with OpenCL we need C arrays, it is inevitable to convert the input list to C array and back.

Let *marshaling* be the process of converting an Erlang list to a C-array, and *un-marshaling* the converse.

So we considered and evaluated several ways to marshal and un-marshal a list.

The most straightforward one is using the NIF list handling functions.

Iterating over the list using the `enif_get_list_cell` function we can also, at each step, extract the element inside each list cell. Once we obtained the term representing the element, we can get the content using the `enif_get_double` function (there are `enif_get_` functions for every Erlang type).

Conversely, to generate a list from a C-array, we firstly create the terms for each element using `enif_make_double`, secondly we pass them as array to `enif_make_list_from_array`, which will generate the list.

See figure 3.9 for an implementation of this technique.

The other technique involves Erlang *binaries*.

A *binary* is a data structure designed for storing large quantities of raw data in a space-efficient manner. They are manipulated using built-in functions (BIFs) or with functions defined in the binary module.

Using binary comprehension it's possible to convert lists of floats into binaries and back, in just a line of code (see figure 3.10).

Binaries are fully supported by the NIF API, and are also implemented as C-arrays so we could use them directly in the native code without any change of representation whatsoever; moreover, converting them to and from lists, is extremely easy.

The last options we examined, was using the Resource Object mechanism (see 2.2.3.1 for an intro to ROs) that we already exploited for exposing OpenCL objects to Erlang. This possibility was ruled out almost immediately owing to ROs totally opaque nature.

ROs are just handles to native data, so they can only be passed back as argument to a NIF. Keeping that in mind, to use the data in Erlang we still need to convert them from and to lists using code similar to the one showed in figure 3.9.

So, to decide between Lists and binaries we set up some tests to measure the performance of the two alternatives.

We measured the time to un-marshall input lists of increasing size and marshaling them back using the two approaches.

The results of the tests (Table in figure 3.11), run on the machine described in section 4.1, show that the technique involving binaries is an order of magnitude slower than the other one.

Moreover, with the list approach, we can work on the list in segments so to overlap the marshaling/un-marshaling with the actual computation on the GPU. In this way, since the costs of converting the list is almost completely hidden by the computation, the conversion process comes essentially for free.

```

ERL_NIF_TERM double_array_to_list(ErlNifEnv *env, double* array, size_t array_size) {

    if(array == NULL)
        return ATOM(error);

    ERL_NIF_TERM* floatTermArray =
        (ERL_NIF_TERM*) enif_alloc(sizeof(ERL_NIF_TERM) * array_size);

    for(uint i = 0; i < array_size; i++)
        floatTermArray[i] = enif_make_double(env, array[i]);

    ERL_NIF_TERM res = enif_make_list_from_array(env, floatTermArray, array_size);

    enif_free(floatTermArray);

    return res;
}

double* double_list_to_array(ErlNifEnv *env, ERL_NIF_TERM list) {
    double *x;
    if(enif_is_list(env, list)) {

        unsigned int len = 0;
        enif_get_list_length(env, list, &len);

        x = (double*) enif_alloc(len*sizeof(double));

        ERL_NIF_TERM curr_list = list;
        for(int i=0; i<len; i++) {

            ERL_NIF_TERM hd, tl;
            enif_get_list_cell(env, curr_list, &hd, &tl);
            if(!enif_get_double(env, hd, &(x[i]))) {
                cerr << "ERROR: trying to read float from something else!" << endl;
            }
            curr_list = tl;
        }
    } else {
        cerr << "ERROR: input is not list" << endl;
    }
    return x;
}

```

Figure 3.9: Marshalling/Unmarshalling using NIF API for lists

```

list2bin (L) ->
    << <<X/native-float>> || X <- L >>.

bin2list (B) ->
    [ X || <<X/native-float>> <= B ].

```

Figure 3.10: Converting binaries to lists and vice-versa using binary and list comprehension.

10[^]	List To List	List To Binary To List
2	3	12
3	38	130
4	575	1200
5	7056	13102
6	71958	245976

Figure 3.11: Running times in microseconds.

We used this streaming technique extensively in implementing map skeletons that works directly on lists, as we are going to see in the next section.

3.3.2 Pipelined Map skeleton: `mapLL`

We now give some details about the implementation of the `mapLL` skeleton. In doing so we will gain some insight into the problems we had to solve during the development of the library.

Following the naming pattern we explained when we earlier introduced skeleton NIFs, `mapLL` is a map skeleton accepting lists as input and returning the result also as lists. This is then the case where marshaling and un-marshaling can potentially take the lion's share of the total running time of a skeleton.

In fact, in order to compute the `mapLL` skeleton the following steps must be completed:

1. *Input un-marshaling* – The elements contained into the input list are extracted and copied into a host buffer
2. *Loading data on the device* – The data is copied from the host buffer to a device buffer on the GPU.
3. *Computation* – The kernel is executed having as input the previous device buffer. The kernel outputs the results in a device buffer.
4. *Unloading data from the device* – The output data is copied from the device back in a host buffer.

5. *Output marshaling* – The output list is created from the data contained into the previous host buffer.

Looking at the figures in Table in figura 3.11 we can immediately understand that, to optimize skeletons' performance, we need to find a way of neutralize (un-)marshaling costs.

To further decrease the running time we exploited the DMA (direct memory access) engine of the GPU. Using DMA is possible to copy data to (or from, but not at the same time, we'd need 2 copy engine for that) while the GPU is busy computing the kernels.

We now recall that, since in a map skeleton there are no data dependencies among the input elements, we can split the input list in an arbitrary number of segments and start computing the user-function on the elements of a segment, independently from the ones belonging to other segments.

Knowing that, we can finally recognize un-marshaling and loading steps (and their inverse unloading and marshaling), as independent from computation, provided that we work on the input in segments.

At this point, it comes almost natural to arrange the three phases we just mentioned, as a three-stage pipeline.

In this way we can overlap computation with both (un-)marshaling and the moving of data to and from the GPU.

As described in 2.1.2, to submit commands for execution on an OpenCL device, command-queues are used. Using in-order queues, to send commands to the device concurrently and independently, more than one command-queue must be associated to a single OpenCL context.

To implement the scheme we just presented, we adopted a dual queue approach. Each segment of the input is split in two: the first half is processed using the first

queue, and the second using the other one. In doing so we can, for instance, issue a copy command for the first half while executing the kernel on the second one and so on.

A representation of the dual-queue pipelining technique is given in Figure 3.12. For simplicity's sake, in the diagram load is comprised of un-marshaling and copying data from host to device; unload is the opposite process, covering copying from device to host and marshaling the output.

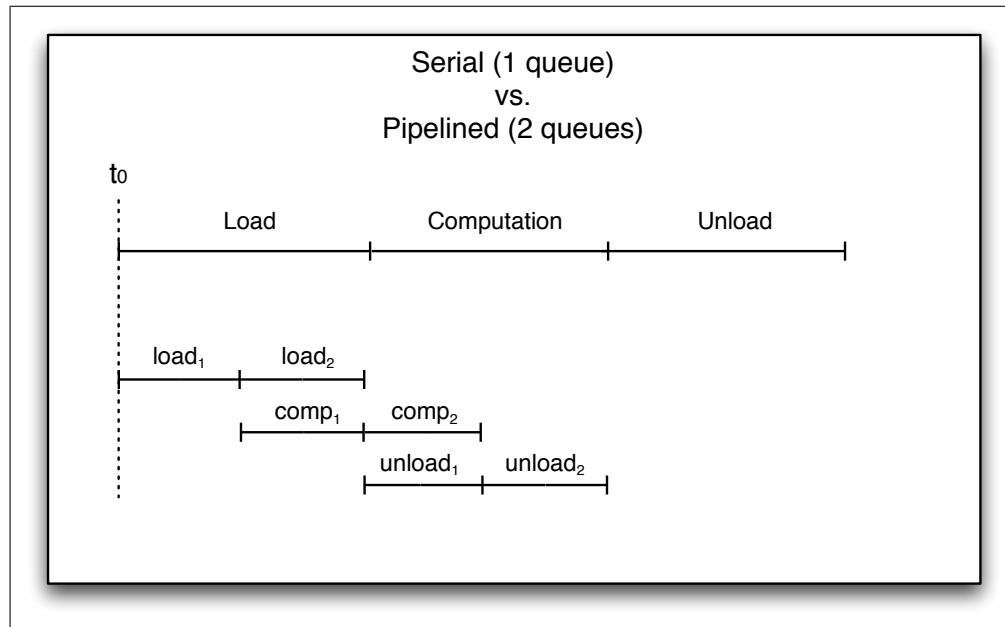


Figure 3.12: Serial vs. pipelined execution using two command queues.

Each of the stage of the pipeline is implemented using the threading facilities provided by the NIF API [2].

Synchronization among the threads is realized exploiting the OpenCL events mechanism: a series of user events (events created calling `clCreateUserEvent`) are used as checkpoints on which a thread must wait until the one responsible for the event fires it (setting them to `CL_COMPLETE`).

Let's now have a look at the code for the `mapLL` NIF in figure 3.13. It uses the same code structure of every NIF in the library; as mandated by the NIF mecha-

nism, the signature of the function is:

```
static ERL_NIF_TERM mapLL(ErlNifEnv*env, int argc, const ERL_NIF_TERM argv[]);
```

A NIF must be a static function returning an Erlang Term and taking 3 parameters:

1. *env* – A pointer to a process bound environment, all function arguments passed to a NIF and its return value belongs to this environment; it's only valid in the thread where it was supplied as argument until the NIF returns.
2. *argc* – The count of terms available in the argv array.
3. *argv* – An array of Erlang terms representing the actual parameters of the NIF.

After some preliminary checks, we read, one by one, the content of the terms in argv using the NIF API.

For example, `mapLL` first parameter is the kernel ROs (created by calling the `createMapKernel` NIF) denoting the user-function to be executed; therefore we need to get back the pointer to the struct that holds the kernel OpenCL object and we do that calling `enif_get_resource`.

Something similar happens to the remaining NIF's parameters: a list and an integer representing the length of that list.

Since we finally got all the information needed to run the map skeleton, the NIF delegates the execution the `mapLL_Impl` function.

3.3.3 NIF and OpenCL issues

In this section we explain some of the issues we had to solve in implementing the library. In particular we discuss some of OpenCL's and NIF's shortcomings we stumbled upon, and how we overcame them.

NIFs

As we have seen in section 2.2.3, one of the main concerns in developing NIFs is implementing them so that they returns quickly (less than 1 ms) otherwise they would interfere with VM's schedulers blocking other running processes; this is the "dirty NIF" problem.

```

static ERL_NIF_TERM mapLL(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
    //mapLL(Kernel::kernel[], InputList::[double[]], InputLength::non_neg_integer[])
    OCL_INIT_CHECK()
    const uint NUM_ARGS = 3;

    NIF_ARITY_CHECK(NUM_ARGS)
    /*get the parameters [Kernel::kernel[], InputList::[double[]], InputLength::non_neg_integer[]****/

    kernel_sync* pKernel_s = NULL;
    if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {
        cerr << "ERROR :: mapLL: 1st parameter is not a kernel_sync" << endl ;
        return enif_make_badarg(env);
    }

    ERL_NIF_TERM inputList = argv[1];
    if (!enif_is_list(env, argv[1])) {
        cerr << "ERROR :: mapLL: 2nd parameter is not a list" << endl ;
        return enif_make_badarg(env);
    }

    uint inputListLength;
    if (!enif_get_uint(env, argv[2], &inputListLength)) {

        cerr << "ERROR :: mapLL: 3rd parameter is not a non_neg_integer()" << endl ;
        return enif_make_badarg(env);
    }

    /****** */
    uint pInputC = 1;
    ERL_NIF_TERM* pInputV[1] = { &inputList };

    return mapLL_Impl(env, pKernel_s, (void**) pInputV, pInputC, NULL, inputListLength);
}

```

Figure 3.13: mapLL NIF.

This is a limitation to the user of the NIF mechanism, so in release R17B, dirty scheduler have been introduced (see 2.2.3). We chose not to use them being a declaredly experimental feature and also because they require to build the Erlang platform with explicit support for them.

Dirty schedulers aside, as a workaround the NIF documentation suggests dispatching the work to another thread, return from the native function, and wait for the result. Then the thread can send the result back to the calling thread using message passing. We investigated the technique but ultimately we found it unfeasible when working with long lists as with `mapLL`.

As explained in 2.2.3, in the NIF API an environment is a container of terms; a term is valid as long as the environment is valid. There are two types of environments: *process bound* and *process independent*.

A process bound environment is passed as the first argument to all NIFs; it contains the function parameters and also the term that the NIF must return. Such kind of environment is only valid in the thread where it was supplied as argument until the NIF returns.

Instead, to store terms between NIF calls and to send them to an Erlang process, a process independent environment must be used. Terms can be copied between environments with `enif_make_copy`, and citing the documentation [2]:

All elements of a list/tuple must belong to the same environment as the list/tuple itself.

This last requirement makes the usage of the suggested workaround, prohibitively expensive. In fact, copying between two environments very long lists including each and every one of their elements is way too costly in a high-performance setting.

So, in the end, we decided to ignore the problem taking also into account Skel OCL’s intended use, namely as a batch system hence multitasking should not be a concern.

OpenCL

With OpenCL, to get a direct pointer into memory buffer, the `clEnqueueMapBuffer` is used; when done working on the buffer, it must be also unmapped calling `clEnqueueUnmapBufferObject`.

We used this function to access the host buffer containing the input and output data for a computation. In our experiments we found that, on our test machine (section 4.1), (un)mapping OpenCL buffers is a serial activity even using different threads each one working on independent command queues. This means, for example, that it is not possible to (un)map a buffer while computing a kernel, not even when the buffer is on the host and it’s unrelated to the kernel: they are executed one after the other. This limitation, forced us to fastidiously design the interactions among `mapLL`’s threads so to carry out as many parallel activities as possible. For instance, still in `mapLL`, we overlap the un-mapping of the host buffer used to store input and output data, with the output list creation (`enif_create_list`).

3.4 Example: The `mapReduceLL` Skeleton

We now show an example of efficient skeleton composition; we are going to define a new skeleton, `mapReduceLL`. While presenting it we will highlight some of `Skel OCL`’s features.

`mapReduceLL` is the combination of a map skeleton with a reduce one (see Figure 3.14); taking a list as input, it applies `MapKernel`, then the result is reduced using `ReduceKernel` and finally it returns the reduced value as a one element list.

In this example we compose `mapLD` with `reduceDL`; `mapLD` takes a list as input and store the result into a specified device buffer, `reduceDL` instead works on a device buffer and outputs a list.

```

%%----- Skeleton Composition example -----

-spec mapReduceLL(MapKernel, ReduceKernel, {InputList, InputListLength}) -> {ok, Result} when
    MapKernel :: kernel(),
    ReduceKernel :: kernel(),
    InputList :: [float()],
    InputListLength :: non_neg_integer(),
    Result :: [float()]
.

mapReduceLL(MapKernel, ReduceKernel, {InputList, InputListLength}) ->
    NumVal = InputListLength,
    SzDouble = 8,
    SzBufferByte = NumVal * SzDouble,
    MapOutputD = checkResult( allocDeviceBuffer(SzBufferByte, read_write) ),
    checkResult( mapLD(MapKernel, InputList, MapOutputD, InputListLength) ),
    Result = reduceDL(ReduceKernel, MapOutputD),
    releaseBuffer(MapOutputD),
    Result
.

```

Figure 3.14: Skeleton composition example: `mapReduceLL`

Knowing that, to combine the two skeletons, we just need to allocate a suitably sized device buffer, specifying it both as `mapLD`'s output buffer and as the `reduceDL`'s input one.

Along these lines, we can go on composing skeletons at will, having just the care of connecting them using intermediate device buffers.

It is now clear why, the distinction between device and host buffers is meaningful: having the power of managing buffers directly on the device we are able to avoid moving intermediate data back and forth from the GPU, as we naively do in the section 3.2.

Chapter 4

Benchmarking Skel OCL

This chapter presents some applications we developed to test and validate Skel OCL performances. These applications are also useful as an example of how to use the skeletons provided by the library to structure parallel computations.

Every sample application is implemented in three different ways:

1. *Erlang/Skel OCL* – Using the library’s facilities to exploit the GPU.
2. *Pure Erlang* – Using Erlang functions from the `lists` module. We compare it against the first one to assess the benefits of using Skel OCL
3. *C++/OpenCL* – This is the reference implementation. Its performance is the target Skel OCL should aim for.

4.1 Test platform and methodology

Let’s have a look at the test platform and methodology used to carry out the experiments.

The machine at hand is `pianosa.di.unipi.it` including:

- *CPU* - 8 core/2-way hyperthreading Intel Xeon E5-2650 running at up to 2.80GHz (Refer to [5] for the complete datasheet).
- *RAM* – 32 GB
- *GPU* – NVIDIA GeForce GTX 660 with 2 GB of on-board memory. It’s based on NVIDIA’s Kepler architecture [21] and is equipped with 5 Streaming Pro-

cessors having 192 CUDA cores each. It is also equipped with one copy engine (see 2.1.6 and [20]).

- *Software* – NVIDIA GeForce driver 304.51 and CUDA 5.5 with OpenCL 1.1 support.

To evaluate Skel OCL performances, we compared the completion time of each implementation of the test applications.

For the ones in Erlang we used the purposely created `test_avg` function; it executes `NumRounds` times a specified function, timing these calls using Erlang’s `timer:tc`. In the C++/OpenCL implementation we used instead the `clock_gettime` function defined in `time.h` getting the wall clock time before and after the calling of the function to measure.

Having the need to simulate various arithmetic intensities in our test applications, we added to user-functions a delay loop that computes `math:sin` a specified number of times. In this way we can easily vary the computation grain and therefore highlight the effect of (un-)marshaling on the total cost of the computation.

Finally, using the running times of the three implementations, we compute the speedup as a measure of the performance improvement achieved using Skel OCL.

For instance, to measure the speedup of Skel OCL versus pure Erlang we use the following formula:

$$\text{speedup} = \frac{T_{erl}}{T_{skel.OCL}}$$

To compare Skel OCL against C++/OpenCL, which we chose as the performance reference, we use:

$$\text{speedup} = \frac{T_{C_{pp}/OCL}}{T_{skel.OCL}}$$

Since we chose the C++/OpenCL implementation as baseline, its running time is always shorter than Skel OCL's resulting in speedups smaller than 1.

4.2 mapLL

The first application is based on the `mapLL` skeleton; it's an example of the library usage and in testing it we can validate the (un-)marshaling techniques discussed in the implementation section (3.2).

4.2.1 Pure Erlang implementation

In Figure 4.1 we see a pure Erlang implementation, it uses the `lists:map` function to apply to each element of the input list, a user-defined function. The user-defined function we chose, simply computes the square of its argument; as explained in 4.1, there is also a delay loop. The list input contains floats from 0 to `NumVal-1` and it's generated by `lists:seq`.

Finally, we measure the running time of each of the 10 runs, using our test facility for Erlang functions, `test_avg`.

```
-define(ERL_DELAY, 10).

sq(X) ->
    dummyLoadLoop(2.0, ?ERL_DELAY),
    X*X
.

erl_mapLL(NumVal) ->

    io:format("~n-----erl_mapLL TEST-----~n"),
    io:format("Number of values: ~w, Bytes: ~w, Input: seq(0, ~w).~n", [NumVal, NumVal*8, NumVal-1]),
    io:format("Map function: fun sq/1 ~n"),
    io:format("Delay: ~w~n~n", [?ERL_DELAY]),

    L1 = [ X+0.0 || X <- lists:seq(0, NumVal-1) ],
    NumRuns = 10,
    Median = test_avg:test_avg(lists, map, [fun sq/1, L1], NumRuns),

    io:format("erl_mapLL_test median total time over ~w executions: ~w usec~n", [NumRuns, Median])
.
```

Figure 4.1: `mapLL` skeleton test in pure Erlang implementation

4.2.2 Erlang / Skel OCL

In Figure 4.2 we see the how to implement the same application described in the previous section, using Skel OCL.

The user-defined function is written in OpenCL C; also in this case there's a delay loop to simulate more complex kernels.

The Erlang code is extremely similar; besides creating the kernel (using `createMapKernel`) we just need to replace `lists:map` with `skel_ocl:mapLL`.

```
//OpenCL C kernel in file sq.cl
double sq(double x) {
    return x*x;
}

%%-----Erlang + Skel OCL-----
ocl_mapLL(NumVal) ->
    case isPow2(NumVal) of
        false -> erlang:error(input_not_pow2);
        true -> ok
    end,
    PWD = element(2,file:get_cwd()) ++ "/",
    FunSrcFile = "sq.cl", FunName = "sq",
    FunSrcFileAbsPath = PWD ++ FunSrcFile,
    io:format("~n-----ocl_mapLL TEST-----~n"),
    io:format("Number of values: ~w, Bytes: ~w, Input: seq([0, ~w], [NumVal, NumVal*8, NumVal-1]),~n",
    io:format("Map function file: ~p~n", [FunSrcFileAbsPath]),
    io:format("Map function name: ~p~n~n", [FunName]),
    L1 = [ X+0.0 || X <- lists:seq(0, NumVal-1) ],
    MapSqKernel =
        checkResult( createMapKernel(FunSrcFileAbsPath, FunName)),
    NumRounds = 10,
    Median = test_avg:test_avg(skel_ocl, mapLL, [MapSqKernel, L1, NumVal], NumRounds),
    io:format("ocl_mapLL_test median total time over ~w executions: ~w usec~n", [NumRounds, Median])
.
```

Figure 4.2: `mapLL` skeleton test in Erlang / Skel OCL

4.2.3 C++/OpenCL

In figure 4.3 we see a C++/OpenCL implementation for the map skeleton. We use it to find a baseline for comparing the performance of the other two implementations.

The map skeleton is implemented in the `map_impl` function, which is a refactored version of the one in Skel OCL.

4.2.4 Tests Results

We tested `mapLL` on input lists of one million (2^{10}) elements, we run the test 10 times for 0, 20, 40 and 80 delay loops so to simulate increasingly complex kernels.

Figure 4.4 shows the speedup of Skel OCL (4.2.2) versus the pure Erlang (4.2.1) implementation of `mapLL` skeleton, while Figure 4.2.4.2 shows the one of Skel OCL relative to the C++/OpenCL version.

As evident from the graphs, to achieve significant speedups the kernel to compute must be complex enough to cover the cost of list (un-)marshaling.

Thanks to the techniques described in 3.2.2 to overlap (un-)marshaling to kernel execution, we neutralize the cost of the former when the cost of the latter is greater.

Turning our attention at figure 4.5, where we see the speedup curve relative to the C++/OpenCL reference implementation, we notice that we get to the breakeven point around the value of 40 delay loops.

As we increase the number of delay loops, the curve flattens after the value 40, showing that the speedup is approaching its limit. That is, what we lose to the reference implementation is due just to Erlang's intrinsic inefficiencies, and not to (un-)marshaling cost.

```

int main(int argc, char** argv) {
    const uint NUM_ELEMS_EXP = 10;

    if(argc == 2) NUM_ELEMS_EXP = (uint) atoi(argv[1]);
    else {
        cout << "\nUsage: map_test NUM_ELEMS_POW2_EXPONENT\n";
        cout << "Using default value: 2^10\n";
    }

    if(!init_OCL()) { cout << "Can't initialize OpenCL"; return 0; }

    const uint TEST_RUNS = 10;

    const string mapFunFile = "sq.cl", mapFunName = "sq";

    const uint NUM_ELEMS = pow2(NUM_ELEMS_EXP);

    cout << "\n-----ocl_map TEST-----\n";
    cout << "Number of values: " << NUM_ELEMS << ", Bytes: " << NUM_ELEMS * 8 << ", Input: seq(0, "
    << NUM_ELEMS-1 << ").\n";
    cout << "Map function file: " << mapFunFile << "\n";
    cout << "Map function name: " << mapFunName << "\n\n";
    cout << TEST_RUNS << " rounds test\nRuntimes in usecs\n\n";

    //init buffer with double values from 1 to arrayLen
    const double* input = seqArray(NUM_ELEMS);

    timespec runs_counters[TEST_RUNS+1];

    for (uint i = 0; i < TEST_RUNS; ++i) {
        GET_TIME(runs_counters[i]);

        mapImpl(input, NUM_ELEMS, mapFunFile, mapFunName);
    }
    GET_TIME(runs_counters[TEST_RUNS]);

    delete[] input;

    long runs_times[TEST_RUNS];

    long total = 0;

    cout << "[ ";
    for (int i = 0; i < TEST_RUNS; ++i) {
        total += runs_times[i] = DIFF(runs_counters[i], runs_counters[i+1]);
        cout << runs_times[i] << " ";
    }
    cout << "] ";

    sort(std::begin(runs_times), std::end(runs_times));

    cout << "\nMedian: " << runs_times[(int) round(TEST_RUNS/2)] << endl;
    cout << "Avg: " << total/TEST_RUNS << endl << endl;

    return 0;
}

```

Figure 4.3: Map skeleton test C++/OpenCL implementation

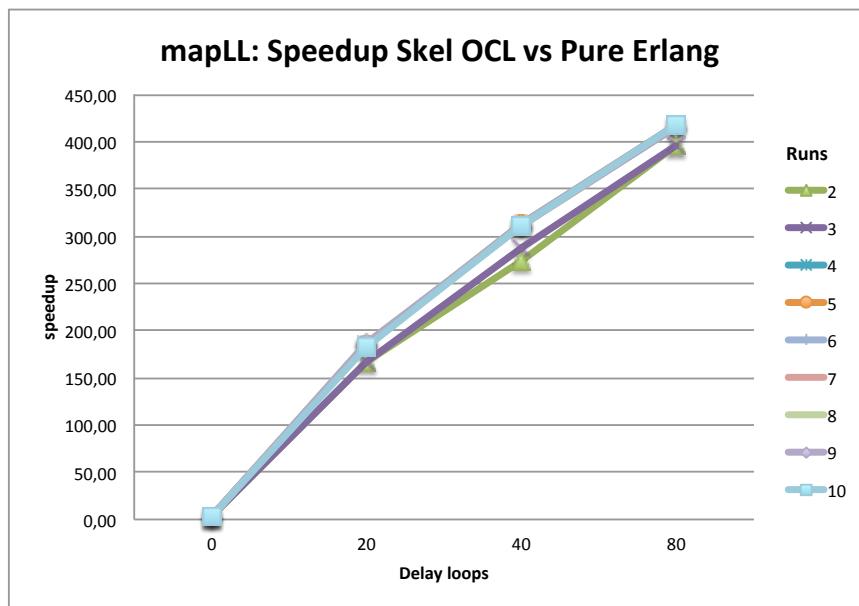


Figure 4.4: Speedup of Skel OCL vs. Pure Erlang implementation of mapLL skeleton

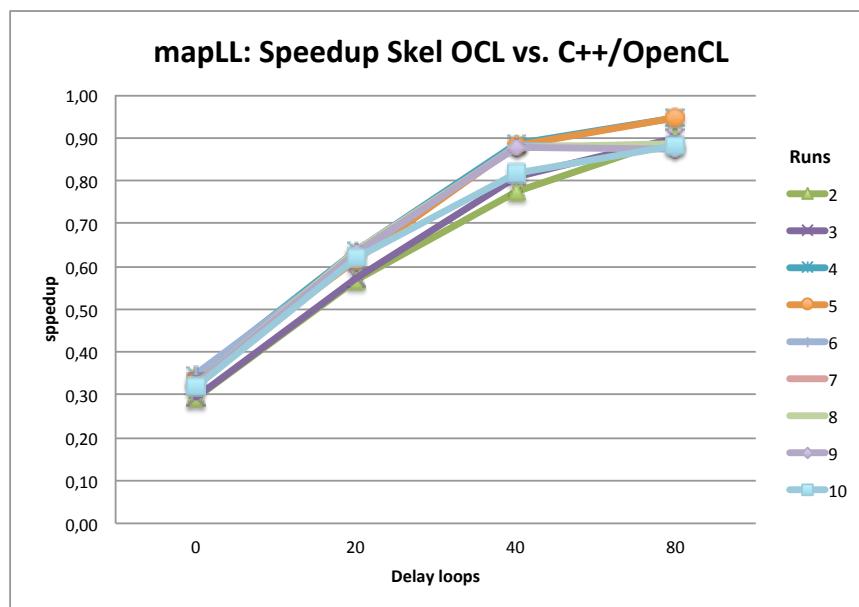


Figure 4.5: Speedup of Skel OCL vs. C++/OpenCL Erlang implementation of mapLL skeleton

4.3 Numerical Integration

The second application is numerical integration using the midpoint method:

$$\int_a^b f(x) dx \approx h \sum_{j=1}^n f(m_j)$$

$$\text{where } h = \frac{(a - b)}{n} \text{ and } m_j = a + (j - \frac{1}{2}) h, 1 \leq j \leq m$$

4.3.1 Pure Erlang implementation

We implement the numerical integration application using `lists:map` and `lists:foldl` combining them in an Erlang implemented `mapReduce` skeleton similar to the one presented in 3.3.

Firstly we generate a list containing the m_j points, secondly we map the integrand function over it and finally we multiply the sum of the function values by h , obtaining the approximation of the integral.

The code for this implementation is shown in figure 4.6.

4.3.2 Erlang / Skel OCL

In figure 4.7 we see the implementation in Erlang / Skel OCL. Here we exploit the `mapReduceLL` skeleton presented in section 3.3 to compute the numerical integral in the same way we described in the previous section.

Like in the `mapLL` example, kernels creation aside, the code is extremely similar to the one in pure Erlang.

4.3.3 C++/OpenCL

In figure 4.8 we see a C++/OpenCL implementation for the numerical integration using skeletons as described earlier. As before, we use it to find a baseline for comparing the performance of the other two implementations.

```

%%Pure erlang implementation
integrandFun(X) ->
    dummyLoadLoop(2.0,?ERL_DELAY),
    X+2
.

erl_integr_test(A,B,N_nodes) ->

    Integrand = fun integrandFun/1,
    Nodes_list = generate_nodes(A, B, N_nodes),

    io:format("~-n-----erl_integr_test-----~n"),
    io:format("Integrand fun: \"fun(X)-> X+2 end\"~n"),
    io:format("Delay: ~w~n", [?ERL_DELAY]),
    io:format("Bounds: [~w,~w]~n",[A,B]),
    io:format("Number of Nodes: ~w~n~n",[N_nodes]),

    Runs = 10,

    Median = test_avg:test_avg(integration, integrate, [erl, Integrand, A, B, {Nodes_list, N_nodes}], Runs),
    io:format("erl_integr_test median total time (~w rounds): ~w usec~n", [Runs, Median])

generate_nodes(A, B, N_nodes) ->

    H = (B-A)/N_nodes,
    MidPoint = fun(I) -> A + ( (I+1.0) - 0.5 ) * H end,
    Points_method = MidPoint,
    for(0, N_nodes-1, Points_method) %% from 0 to n-1
.

%% used also to test Skel OCL
integrate(Type, Integrand, A, B, {Nodes_list, N_nodes}) ->

    H = (B-A)/N_nodes,
    Sum =
        case Type of
            erl -> erl_integr_mapReduce(Integrand, {Nodes_list, N_nodes});
            ocl -> ocl_integr_mapReduce(Integrand, {Nodes_list, N_nodes});
            _ -> erlang:error(wrong_type_error)
        end,
    Sum*H
.

erl_integr_mapReduce(F, {Nodes_List, _N_nodes}) ->

    Sum_fun = fun(X,Y)-> X+Y end,
    Values_list = lists:map(F, Nodes_List),
    Sum = lists:foldl(Sum_fun, 0.0, Values_list),
    Sum
.

```

Figure 4.6: Numerical integration using pure Erlang

```

%%Numerical Integration implemented whith skel_ocl mapReduceLL.

ocl_integr_test(A, B, N_nodes) ->

    Integrand = "f_integrand",%name of the function to be integrated, must be the same of .cl kernel file
    Nodes_list = generate_nodes(A, B, N_nodes),

    io:format("~n-----ocl_integr_test-----~n"),
    io:format("Integrand OCL Kernel file: ~p~n",[Integrand++".cl"]),
    io:format("Bounds: [~w,~w]~n",[A,B]),
    io:format("Numeber of Nodes: ~w~n~n",[N_nodes]),

    Runs = 10,

    Median = test_avg:test_avg(integration, integrate, [ocl, Integrand, A, B, {Nodes_list, N_nodes}], Runs),
    io:format("ocl_integr_test median total time (~w rounds): ~w usec~n", [Runs, Median])

.

ocl_integr_mapReduce(IntegrandFunName, { Nodes_List, N_nodes }) ->

    ReduceFunSrcFile = "sum.cl", ReduceFunName = "sum",
    PWD = element(2,file:get_cwd()) ++ "/",

    %%The integrand's src file has the same name of the integrand function
    Integrand_SrcFileAbsPath = PWD ++ IntegrandFunName ++ ".cl",

    ReduceFunSrcFileAbsPath = PWD ++ ReduceFunSrcFile,
    IntegrandMapKernel = checkResult(createMapKernel([Integrand_SrcFileAbsPath, IntegrandFunName])),
    ReduceSumKernel = checkResult(createReduceKernel([ReduceFunSrcFileAbsPath, ReduceFunName])),

    mapReduceLL(IntegrandMapKernel, ReduceSumKernel, {Nodes_List, N_nodes}),
.

```

Figure 4.7: Numerical Integration test in Erlang / Skel OCL

The mapReduce skeleton is implemented in the `mapReduce_impl` function, which is the composition of refactored versions of Skel OCL's `mapDD` and `reduceDD` skeletons.

4.3.4 Tests Results

We computed an approximation of

$$\int_{10}^{20} (x + 2) dx$$

on 1 million points (2^{10}) using the method described in 4.3.1.

We run the test 10 times for 0, 20, 40 and 80 delay loops so to simulate increasingly heavier computational loads.

Figure 4.9 shows the speedup of Skel OCL (4.3.2) versus the pure Erlang (4.3.1) implementation of the numerical integration test application, while Figure 4.10 shows the one of Skel OCL relative to the C++/OpenCL version.

Also here the kernel load is not heavier enough to achieve a significant speedup so we need the delay loop.

In this application we get higher speedups thanks to the highly optimized reduce kernel based on a NVIDIA reference implementation [18].

Again, we see in Figure 4.3.4.2 how marshaling cost stops being the performance bottleneck after 40 delay loops.

4.4 Dot product

The last test application computes the dot product of two n -dimensional vectors u and v using the formula:

$$u * v = u_1 v_1 + u_2 v_2 + \cdots + u_n v_n = \sum_{j=1}^n f(u_j v_j)$$

```

double mapReduce_impl(double** inputs, uint inputC, uint INPUT_LEN, const string& mapFunFile, const
string& mapFunName, const string& reduceFunFile, const string& reduceFunName);

int main(int argc, char** argv) {
    uint N_NODES_EXP = 10;

    if(argc == 2)
        N_NODES_EXP = (uint) atoi(argv[1]);
    else {
        cout << "\nUsage: integration_test N_NODES_POW2_EXPONENT\n";
        cout << "Using default value: 2^10\n";
    }

    if(!init_OCL()) {
        cout << "Can't initialize OpenCL";
        return 0;
    }

    const double A = 10, B = 20;//integration bounds
    const uint N_NODES = pow2(N_NODES_EXP);

    const uint TEST_RUNS = 10;

    const string mapFunFile = "f_integrand.cl", mapFunName = "f_integrand";
    const string reduceFunFile = "sum.cl", reduceFunName = "sum";

    cout << "\n-----ocl_integr TEST-----\n";
    cout << "Integrand fun: " << mapFunName << "\n";
    cout << "Bounds: [" << A << "," << B << "]\n";
    cout << "Number of Nodes: " << N_NODES << "\n\n";

    cout << TEST_RUNS << " rounds test\nRuntimes in usecs\n\n";

    //generate nodes
    const double H = (B-A)/N_NODES;

    double* nodes = new double[N_NODES];

    for (int i = 0; i < N_NODES; ++i)
        nodes[i] = A + ( (i+1.0) - 0.5 ) * H;

    //test runs
    timespec runs_counters[TEST_RUNS+1];

    for (uint i = 0; i < TEST_RUNS; ++i) {
        GET_TIME(runs_counters[i]);

        mapReduce_impl(&nodes, 1, N_NODES, mapFunFile, mapFunName, reduceFunFile,
                      reduceFunName);

    }
    GET_TIME(runs_counters[TEST_RUNS]);

    .....
    return 0;
}

```

Figure 4.8: Numerical Integration test in C++/OpenCL

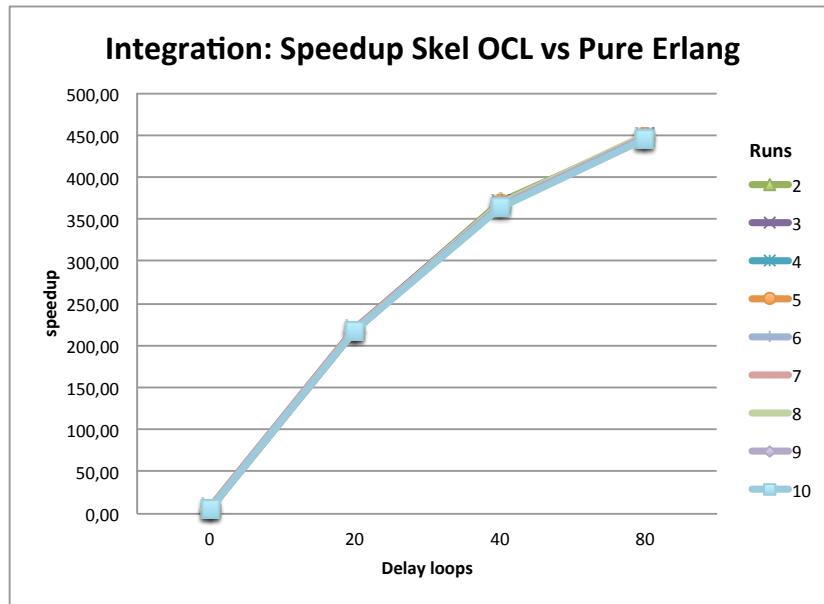


Figure 4.9: Speedup of Skel OCL vs. Pure Erlang implementation of numerical integration

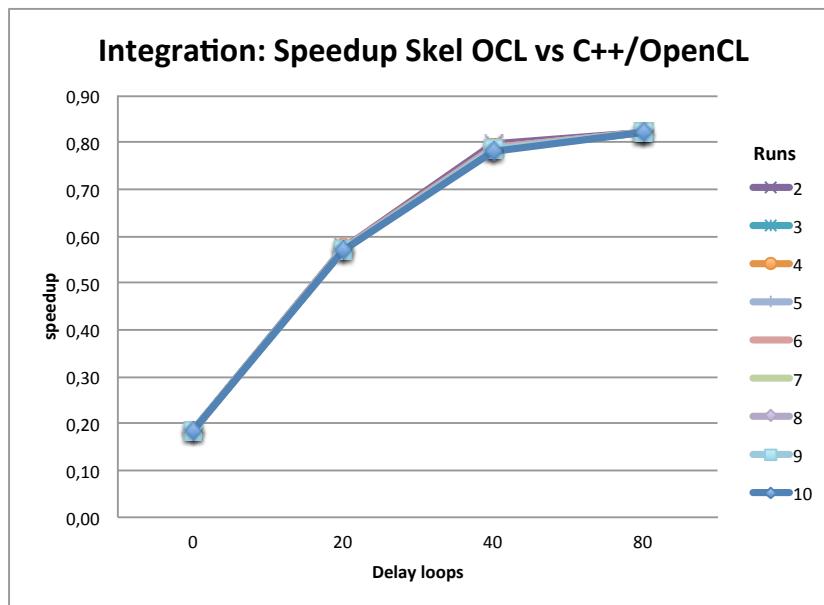


Figure 4.10: Speedup of Skel OCL vs. C++/OpenCL implementation of the numerical integration

We can easily express this computation using the binary map and reduce kernel.

The binary Map skeleton multiplies the corresponding elements of the two vectors and then the Reduce skeleton sums all the intermediate products resulting in the dot product of the two vectors.

4.4.1 Pure Erlang implementation

To implement the computation of the dot product of two vectors represented as lists, we used `lists:zipWith` and `lists:foldl`.

`ZipWith` is the Erlang counterpart of Skel OCL `map2` skeleton; it combines the corresponding elements of two lists of equal length into one list. To each input pair a specified binary function is applied; the result is then appended to the output list.

The code for this implementation is shown in figure 4.11.

4.4.2 Erlang / Skel OCL

Figure 4.12 shows the implementation in Erlang / Skel OCL.

Here we see another instance of skeleton composition: we compose `map2LD` skeleton with a `reduceDL` through a device buffer as we did in section 3.4 where we implemented the `MapReduceLL` skeleton.

The kernel required by a `map2` skeleton must be created using the three-argument version of `createMapKernel` specifying the arity of the function, in this case 2.

4.4.3 C++/OpenCL

In figure 4.13 we see a C++/OpenCL implementation for the dot product application using skeletons as described earlier. As always, we use it to find a baseline for comparing the performance of the other two implementations.

The `mapReduce` skeleton is implemented in the `mapReduce_impl` function, which is the composition of refactored versions of Skel OCL's `mapDD` and `reduceDD` skeletons.

```
%%%-----pure Erlang version-----
sum(X,Y) -> X+Y.

zipFun(X,Y) ->
    dummyLoadLoop(2.0,?ERL_DELAY),
X*Y.

erl_dotProduct(V1,V2) ->
    Sums = lists:zipwith(fun zipFun/2, V1, V2),
    lists:foldl(fun sum/2, 0.0, Sums)
.

erl_dotProduct_test(NumVal) ->
    Val = 2,
    io:format("~n-----erl_dotProduct TEST-----~n"),
    io:format("Vectors' dimensions: ~w~nV1:seq(0,~w)~nV2:duplicate(~w,~w): ~n", [NumVal, NumVal-1,
    NumVal, Val]),
    io:format("Delay: ~w~n~n", [?ERL_DELAY]),
    V1 = [ X+0.0 || X <- lists:seq(0,NumVal-1) ],
    V2 = [ X+0.0 || X <- lists:duplicate(NumVal,Val) ],
    Runs = 10,
    Median = test_avg:test_avg(dotProduct, erl_dotProduct, [V1, V2], Runs),
    io:format("erl_dotProduct_test median total time (~w rounds): ~w usec~n", [Runs, Median])
.
```

Figure 4.11: Erlang implementation of the Dot Product test application

```

ocl_dotProduct(NumVal, V1, V2) ->

    SzDouble = 8,
    SzBufferByte = NumVal * SzDouble,

    MapFunSrcFile = "mul.cl", MapFunName = "mul",
    ReduceFunSrcFile = "sum.cl", ReduceFunName = "sum",

    PWD = element(2,file:get_cwd()) ++ "/",

    MapFunSrcFileAbsPath = PWD ++ MapFunSrcFile,
    ReduceFunSrcFileAbsPath = PWD ++ ReduceFunSrcFile,

    %%Create and compile kernels
    MapMulKernel = checkResult(createMapKernel(MapFunSrcFileAbsPath, MapFunName, 2)),
    ReduceSumKernel = checkResult(createReduceKernel(ReduceFunSrcFileAbsPath, ReduceFunName)),

    %%allocate work buffer on device
    MapOutputD = checkResult(allocDeviceBuffer(SzBufferByte, read_write),

    checkResult( map2LD(MapMulKernel, V1, V2, MapOutputD, NumVal) ),

    Result = checkResult( reduceDL(ReduceSumKernel, MapOutputD) ),

    releaseBuffer(MapOutputD),

    Result

.

ocl_dotProduct_test(NumVal) ->

    %% must be a power of 2
    case isPow2(NumVal) of
        false -> erlang:error(input_not_pow2);
        true -> ok
    end,
    Val = 2,

    io:format("~n-----ocl_dotProduct TEST-----~n"),
    io:format("Vectors' dimension: ~w~nV1 = seq(0,~w)~nV2 = duplicate(~w,~w): ~n~n", [NumVal,
    NumVal-1, NumVal, Val]),

    V1 = [ X+0.0 || X <- lists:seq(0,NumVal-1) ],
    V2 = [ X+0.0 || X <- lists:duplicate(NumVal,Val) ],

    Runs = 10,

    Median = test_avg:test_avg(dotProduct, ocl_dotProduct, [NumVal, V1, V2], Runs),
    io:format("ocl_dotProduct_test median total time (~w rounds): ~w usec~n", [Runs, Median])
.

```

Figure 4.12: Dot Product test application Erlang / Skel OCL implementation

```

double mapReduce_<im>pl(double** inputs, uint inputC, uint INPUT_LEN, const string& mapFunFile, const string& mapFunName, const string& reduceFunFile, const string& reduceFunName);

int main(int argc, char** argv) {
    uint NUM_ELEMS_EXP = 10;

    if(argc == 2)
        NUM_ELEMS_EXP = (uint) atoi(argv[1]);
    else {
        cout << "\nUsage: dotProduct_test DIMENSION_POW2_EXPONENT\n";
        cout << "Using default value: 2^10\n";
    }

    if(init_OCL()) {
        cout << "Can't initialize OpenCL";
        return 0;
    }

    const uint NUM_ELEMS = pow2(NUM_ELEMS_EXP);

    const uint TEST_RUNS = 10;

    const string mapFunFile = "mul.cl", mapFunName = "mul";
    const string reduceFunFile = "sum.cl", reduceFunName = "sum";

    cout << "\n-----ocl_dotProduct TEST-----\n";
    cout << "Vectors' dimensions: " << NUM_ELEMS << " V1 = V2 : seq(0, "<< NUM_ELEMS-1 <<")\n";
    cout << TEST_RUNS << " rounds test\nRuntimes in usecs\n\n";

    double** inputs = new double*[2];
    inputs[0] = new double[NUM_ELEMS];
    inputs[1] = new double[NUM_ELEMS];

    for (int i = 0; i < NUM_ELEMS; ++i) {
        inputs[0][i] = (double) i;
        inputs[1][i] = 2.0;
    }
    //test runs
    timespec runs_counters[TEST_RUNS+1];

    for (uint i = 0; i < TEST_RUNS; ++i) {
        GET_TIME(runs_counters[i]);

        mapReduce_<im>pl(inputs, 2, NUM_ELEMS, mapFunFile, mapFunName, reduceFunFile,
                           reduceFunName);
    }
    GET_TIME(runs_counters[TEST_RUNS]);

    ....
    return 0;
}

```

Figure 4.13: Dot Product test application in C++/OpenCL

4.4.4 Tests Results

To test the Dot Product application we computed the dot product of two 2^{10} dimensional vectors u and v :

$$u * v = \begin{pmatrix} 0 \\ 1 \\ \vdots \\ 2^{10} - 1 \end{pmatrix} * \begin{pmatrix} 2 \\ 2 \\ \vdots \\ 2 \end{pmatrix}$$

As usual, we run the test 10 times for 0, 20, 40 and 80 delay loops so to simulate increasingly heavier computational loads.

Figure 4.14 shows the speedup of Skel OCL (4.4.2) versus the pure Erlang (4.4.1) implementation of the dotProduct test application, while Figure 4.15 shows the one of Skel OCL relative to the C++/OpenCL version.

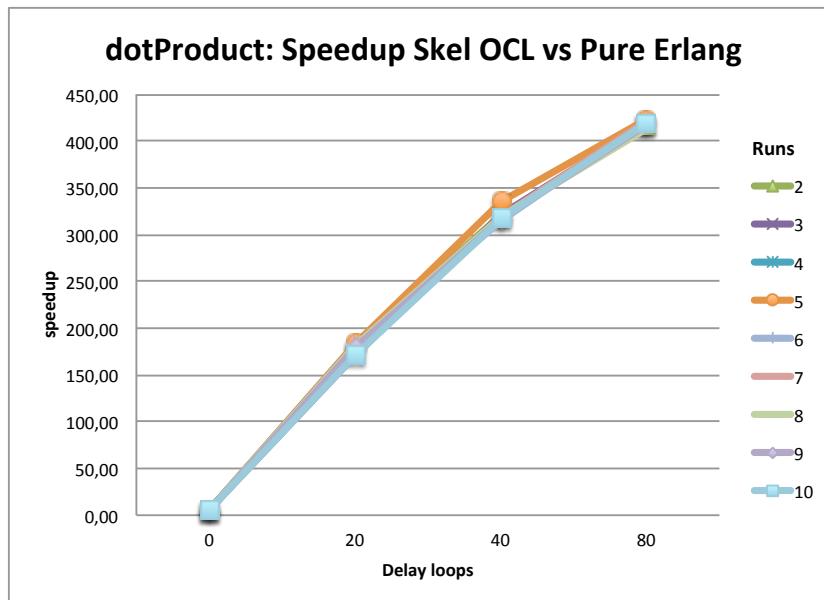


Figure 4.14: Speedup of Skel OCL vs. Pure Erlang implementation of dotProduct

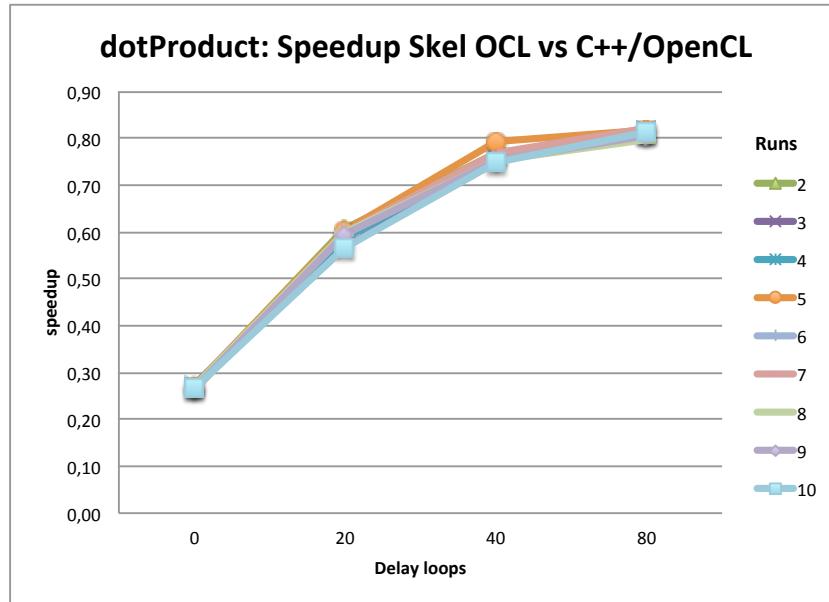


Figure 4.15: Speedup of Skel OCL vs. C++/OpenCL implementation of dotProduct

4.5 Conclusion

In this chapter we presented three applications we developed to test and validate Skel OCL performances:

- *MapLL* – we analyzed the costs of (un-)marshaling, evaluating the techniques discussed in 3.2.
- *Numerical Integration* – An example of usage of the MapReduceLL skeleton defined in 3.3.
- *Dot Product* - An example of skeleton composition using Skel OCL low-level function and of the usage of map2 skeleton.

The outcome of the experiments we presented, shows, on one hand, a significant performance improvement over pure Erlang implementations and, on the other hand, an acceptable performance loss over C++/OpenCL reference implementations.

Chapter 5

Conclusion

In this thesis we presented the Skel OCL skeleton library for the Erlang programming language.

The goal of the project we developed in this thesis was to provide Erlang with an easy to use, high performance, data parallel skeleton library exploiting GPUs' processing power. A prime requirement was vendor neutrality, so OpenCL was the only possible choice.

The skeletons should be easy to use, possibly be a nearly perfect drop-in replacement to standard Erlang functions with similar semantic.

In chapter 1 we discussed how the rise of ubiquitous parallel hardware called for new programming methodologies.

In particular we focused our attention on on GPUs, which, from their origins as special purpose devices, evolved into fully-fledged parallel computing engines.

In chapter 2 we introduced the technologies we used in the project.

To work with GPUs, we chose the OpenCL programming framework thanks to its industry standard status and broad support by the vendors.

As the target programming language we chose Erlang: developed in the telecommunication industry, Erlang has seen a rapid increase in adoption thank to its built-in support for concurrency and distribution.

Erlang is also one of the focuses of the European Paraphrase project. The Paraphrase project is particularly interested in Erlang's support for parallel processors including GPUs.

In chapter 3 we presented Skel OCL, discussing its architecture and implementation. We showed that Skel OCL is an Erlang NIF library providing the map and reduce data parallel skeletons to execute user-defined functions on a GPU. In addition to the built-in skeleton Skel OCL also provides low-level functions to efficiently combine them.

Chapter 4 presented some applications we developed to test and validate Skel OCL performances. In addition to a Skel OCL implementation, we also saw a pure Erlang and a C++/OpenCL one; comparing the completion times of the three we evaluated Skel OCL performances.

Analyzing the tests' results we found out that using Skel OCL we achieved a significant performance improvement over pure Erlang, without losing too much to the C++/OpenCL implementation we chose as baseline.

In summary, we developed Skel OCL, a prototype Erlang NIF library using OpenCL. Skel OCL provides data parallel skeletons to accelerate computations on lists of Erlang floats.

The user code for such computations is specified writing OpenCL kernels using the OpenCL C (see 2.1.5) programming language.

The library has a very user-friendly API being it modeled after Erlang's list module, with which every Erlang programmer is familiar.

Skel OCL also provides lower-level functions for efficient skeleton composition, using these functions the user can easily define new skeletons having the same efficiency as the built-in ones (see 3.3 for an example).

We finally validated the design and implementation of Skel OCL carrying out experiments using three test applications.

The outcome of the tests (chapter 4) shows, on one hand, a significant performance improvement over a pure Erlang implementation and, on the other hand, an acceptable performance loss over a C++/OpenCL reference implementation.

Future work

Owing to its prototypal nature, the directions along which Skel OCL can be further developed are numerous, for example:

- *Broader data types support* – increase the number of data types supported by the skeletons. Skel OCL currently supports computations on lists of Erlang floats. A wider choice of data types would improve Skel OCL suitability, increasing the number of scenarios.
- *More skeletons* – Provide additional optimized built-in skeletons either to be used directly or to be combined employing Skel OCL skeletons composition mechanisms.
- *Kernels in Erlang* – Currently, user-defined functions must be written in OpenCL C so the option of implementing kernels directly in Erlang would be desirable to shorten the learning curve.

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Appendix A

Usage example

In this appendix we are going to see how to get started with Skel OCL.

We show how to setup the library, define an Erlang function that uses a skeleton and the OpenCL C kernel for the computation.

First of all we need to build Skel OCL shared library using `make`. Several target are available, but here we are just interested in the default one.

Using the default target we force Skel OCL to look, during initialization, for a NVIDIA GPU, and to fail when none are found.

So invoking the `make` command from the shell we build the shared library `skel_ocl.so`.

Now we need to write an OpenCL C kernel that implements our computation.

We have to put the functions definitions in a file. In this example we define a simple kernel that performs arithmetic instructions in a file named `userFun.cl` (Figure A.1)

The next step is to write the Erlang code for our computation (Figure A.2). Using the high-level skeletons provided by Skel OCL is really easy.

The skeletons working on Erlang lists (the `LL` type) just need three parameters:

1. The user-defined function as a kernel object
2. The input list (2 in case of `map2`)
3. The length of the input list.

So we firstly to create the kernel object using one of the available NIF. Since we are going to use it in a map skeleton, what we need is `createMapKernel`: As with any `createKernel` NIF we have to pass as arguments the path to the source file and the name of the kernel function.

We now need a floats input list; currently there are limitations on the length of input lists: they must have a number of elements that is a power of two and mustn't be too short (how short depends on the OpenCL device used).

Having all set up, we the `mapLL` skeleton and obtain the list containing the results of our computation.

Lastly, before executing `exampleFun/1` (e.g. invoking it from the Erlang shell) we need to call `skel_ocl:cl_init()` to initialize the library.

```
//foo kernel function
double foo(double x){

    return (x+1)*2;

}
```

Figure A.1: The `foo` kernel function in `exampleKernel.cl`

```

-module(usageExample).

-export([exampleFun/1]).

-import(erl_utils,
      [
       isPow2/1,
       pow2/1
      ]).

-import(skel_ocl,
      [checkResult/1,
       createMapKernel/2,
       mapLL/3
      ]).

exampleFun(NumVal) ->

  case isPow2(NumVal) of
    false -> erlang:error(input_not_pow2);
    true -> ok
  end,

  MapFunSrcFile = "exampleKernel.cl", MapFunName = "foo",
  PWD = element(2,file:get_cwd()) ++ "/",
  MapFunSrcFileAbsPath = PWD ++ MapFunSrcFile,
  InputList = [ X+0.0 || X <- lists:seq(0, NumVal-1) ],
  %%create Kernels objects
  MapSqKernel = checkResult(createMapKernel(MapFunSrcFileAbsPath, MapFunName)),

  %%Map Skeleton: list -> list
  OutputList = checkResult(mapLL(MapSqKernel, InputList, NumVal))
.
```

Figure A.2: The exampleUsage module

Appendix B

Skel OCL: Source Code

This appendix contains the Skel OCL source code. Two are the main sections:

B.1 *skel_ocl.erl* - The Erlang part of the library, containing the API definitions.

B.2 *skel_ocl.so* The native implementation of Skel OCL NIFs, including C++ NIFs and OpenCL C kernel code

B.1 skel_ocl.erl

```
1 -module(skel_ocl).
2 -on_load(init/0).
3
4 -export([
5   %%initializations NIFs
6   init/0,
7   cl_init/0,
8   cl_release/0,
9   skeletonlib/1,
10  checkResult/1,
11
12  %%Buffer NIFs
13  listToBuffer/2,
14  bufferToList/1,
15  bufferToList/2,
16  getBufferSize/1,
17  allocDeviceBuffer/2,
18  allocHostBuffer/1,
19  releaseBuffer/1,
20  copyBufferToBuffer/2,
```

```

21 copyBufferToBuffer/3,
22
23 %%Program and Kernel NIFs
24 buildProgramFromString/1,
25 buildProgramFromFile/1,
26 createKernel/2,
27
28 %%Map NIF
29 createMapKernel/2,
30 createMapKernel/3,
31 createMapKernel/4,
32
33 mapDD/3,
34 mapLL/3,
35 mapLD/4,
36
37 map2DD/4,
38 map2LD/5,
39 map2LL/4,
40
41 %%MapReduce
42 mapReduceLL/3,
43
44 %%Reduce NIF
45 createReduceKernel/2,
46 createReduceKernel/3,
47
48 reduceDD/3,
49 reduceDL/2,
50 reduceLL/3
51
52 ]).
53
54 init()->
55 erlang:load_nif("./skel_ocl",0).
56
57
58
59 %%Buffer Types and functions ——————
60
61 -type hostBuffer() :: binary().
62 -type deviceBuffer() :: binary().
63 -type buffer() :: hostBuffer() | deviceBuffer().
64
65 -type rw_flag() :: read | write | read_write.
66
67 %%Get the size of Buffer in Bytes

```

```

68 -spec getBufferSize(Buffer) -> {ok, SizeByte} | {error, Why} when
69     SizeByte :: non_neg_integer(),
70     Buffer :: buffer(),
71     Why :: atom().
72
73 getBufferSize(_) -> {error, ecl_lib_not_loaded}.
74
75
76 %%Allocate a buffer on the host of size SizeByte bytes
77 -spec allocHostBuffer(SizeByte) -> {ok, Buffer} | {error, Why} when
78     SizeByte :: non_neg_integer(),
79     Buffer :: hostBuffer(),
80     Why :: atom().
81
82 allocHostBuffer(_) -> {error, ecl_lib_not_loaded}.
83
84 %%Allocate a buffer on the device of size SizeByte bytes, specifying kernel read/
85             write permissions
85 -spec allocDeviceBuffer(SizeByte, Flags) -> {ok, Buffer} | {error, Why} when
86     SizeByte :: non_neg_integer(),
87     Flags :: rw_flag(),
88     Buffer :: deviceBuffer(),
89     Why :: atom().
90
91 allocDeviceBuffer(_, _) -> {error, ecl_lib_not_loaded}.
92
93 %%release a previously allocated buffer
94 -spec releaseBuffer(Buffer) -> ok when
95     Buffer :: buffer()
96 .
97
98 releaseBuffer(_) -> {error, ecl_lib_not_loaded}.
99
100 %%Copy data from a buffer to another having the same size
101 -spec copyBufferToBuffer(From, To) -> ok | {error, Why} when
102     From :: buffer(),
103     To :: buffer(),
104     Why :: atom().
105
106 copyBufferToBuffer(_, _) -> {error, ecl_lib_not_loaded}.
107
108 %%Copy CopySizeByte data from a buffer to another
109 -spec copyBufferToBuffer(From, To, CopySizeByte) -> ok | {error, Why} when
110     From :: buffer(),
111     To :: buffer(),
112     CopySizeByte :: non_neg_integer(),
113     Why :: atom().

```

```

114
115 copyBufferToBuffer(., ., .) -> {error, ecl_lib_not_loaded}.
116
117 %%Copy the content of an Erlang list into a host buffer.
118 -spec listToBuffer(From, To) -> ok | {error, Why} when
119     From :: [float()],
120     To :: hostBuffer(),
121     Why :: atom().
122
123 listToBuffer(., .) -> {error, ecl_lib_not_loaded}.
124
125
126 %%Create a list having as elements the data in the host buffer
127 -spec bufferToList(From) -> {ok, [float()]} | {error, Why} when
128     From :: hostBuffer(),
129     Why :: atom().
130
131 bufferToList(.) -> {error, ecl_lib_not_loaded}.
132
133
134 %%Create a list containing the first ListLength elements in the buffer
135 -spec bufferToList(From, ListLength) -> {ok, [float()]} | {error, Why} when
136     From :: hostBuffer(),
137     ListLength :: pos_integer(),
138     Why :: atom().
139
140 bufferToList(., .) -> {error, ecl_lib_not_loaded}.
141
142 %%Program and Kernel Types and functions -----
143 -type program() :: binary().
144 -type kernel() :: binary().
145
146 %%Build an OpenCL program from the source code in the ProgSrcString iolist()
147 -spec buildProgramFromString(ProgSrcString) -> {ok, program()} | {error, Why} when
148     ProgSrcString :: iolist(),
149     Why :: atom().
150
151 buildProgramFromString(.) -> {error, ecl_lib_not_loaded}.
152
153 %%Build an OpenCL program object from the source code in the ProgSrcFile file
154 -spec buildProgramFromFile(ProgSrcFile) -> {ok, program()} | {error, Why} when
155     ProgSrcFile :: nonempty_string(),
156     Why :: atom().
157
158 buildProgramFromFile(.) -> {error, ecl_lib_not_loaded}.
159
160 %%Create an OpenCL kernel object from a program object.

```

```

161 -spec createKernel(Prog, KerName) -> {ok, kernel()} | {error, Why} when
162     Prog      :: program(),
163     KerName   :: nonempty_string(),
164     Why       :: atom().
165
166 createKernel(_, _) -> {error, ecl_lib_not_loaded}.
167
168
169 %%Map functions -----
170
171 %%Create a map compatible kernel object from the source in MapSrcFunFile. The kernel
172 %%name is FunName.
173 -spec createMapKernel(MapSrcFunFile, FunName) -> {ok, kernel()} | {error, Why}
174     when
175         MapSrcFunFile :: nonempty_string(),
176         FunName      :: nonempty_string(),
177         Why          :: atom().
178
179 createMapKernel(MapSrcFunFile, FunName) ->
180     createMapKernel(MapSrcFunFile, FunName, 1, cache).
181
182 %%Create a map compatible kernel object from the source in MapSrcFunFile. The kernel
183 %%name is FunName, the arity is FunArity
184 -spec createMapKernel(MapSrcFunFile, FunName, FunArity) -> {ok, kernel()} | {error
185     , Why} when
186         MapSrcFunFile :: nonempty_string(),
187         FunName      :: nonempty_string(),
188         FunArity     :: pos_integer(),
189         Why          :: atom().
190
191 %%Create a map compatible kernel object from the source in MapSrcFunFile.
192 %%The kernel name is FunName, the arity is FunArity.
193 %%Program caching policy must also be specified.
194 -spec createMapKernel(MapSrcFunFile, FunName, FunArity, CacheKernel) -> {ok,
195     kernel()} | {error, Why} when
196         MapSrcFunFile :: nonempty_string(),
197         FunName      :: nonempty_string(),
198         FunArity     :: pos_integer(),
199         CacheKernel  :: cache | no_cache,
200         Why          :: atom().
201
202 createMapKernel(_, _, _, _) -> {error, ecl_lib_not_loaded}.

```

```

203 %%Map skeleton working on device buffers. User function specified as kernel objects.
204 -spec mapDD(Kernel, Input, Output) -> ok | {error, Why} when
205     Kernel :: kernel(),
206     Input :: deviceBuffer(),
207     Output :: deviceBuffer(),
208     Why :: atom().
209
210 mapDD(_,_,-)-> {error, ecl_lib_not_loaded}.
211
212
213 %%Similar to lists:zipWith(Combine, L1, L2) -> L3 but restricted to double
214 %%Binary Map skeleton working on device buffers. User function specified as kernel
215 %%objects.
216 %%Similar to lists:zipWith
217 -spec map2DD(Kernel, Input1, Input2, Output) -> ok | {error, Why} when
218     Kernel :: kernel(),
219     Input1 :: deviceBuffer(),
220     Input2 :: deviceBuffer(),
221     Output :: deviceBuffer(),
222     Why :: atom().
223
224 map2DD(_,-,-,-)-> {error, ecl_lib_not_loaded}.
225
226 %%Map skeleton: input is a list, output is a device buffer. User function specified
227 %%as kernel objects.
228 -spec mapLD(Kernel, Input, Output, InputLength) -> ok | {error, Why} when
229     Kernel :: kernel(),
230     Input :: [float()],
231     Output :: deviceBuffer(),
232     InputLength :: non_neg_integer(),
233     Why :: atom().
234
235 mapLD(_,-,-,-)-> {error, ecl_lib_not_loaded}.
236
237 %%Binary Map skeleton: input is a list, output is a device buffer. User function
238 %%specified as kernel objects.
239 -spec map2LD(Kernel, Input1, Input2, Output, InputLength) -> ok | {error, Why} when
240     Kernel :: kernel(),
241     Input1 :: [float()],
242     Input2 :: [float()],
243     Output :: deviceBuffer(),
244     InputLength :: non_neg_integer(),
245     Why :: atom().
246
247 map2LD(_,-,-,-,-)-> {error, ecl_lib_not_loaded}.
248

```

```

246 %% Map skeleton: input is a list , output too. User function specified as kernel
247 objects .
248 -spec mapLL(Kernel , Input , InputLength) -> {ok, Result} when
249     Kernel :: kernel() ,
250     Input :: [float()] ,
251     InputLength :: non_neg_integer() ,
252     Result :: [float()] .
253
254
255 %%Binary Map skeleton: input is a list , output too. User function specified as
256 %% kernel objects .
257 -spec map2LL(Kernel , Input1 , Input2 , InputLength) -> {ok, Result} when
258     Kernel :: kernel() ,
259     Input1 :: [float()] ,
260     Input2 :: [float()] ,
261     InputLength :: non_neg_integer() ,
262     Result :: [float()] .
263
264
265
266 %%Reduce functions -----
267
268 %%Create a reduce compatible kernel object from the source in ReduceSrcFunFile .
269 %%The kernel name is FunName .
270 -spec createReduceKernel(ReduceSrcFunFile , FunName ) -> {ok, kernel()} | {error ,
271 Why} when
272     ReduceSrcFunFile :: nonempty_string() ,
273     FunName :: nonempty_string() ,
274     Why :: atom() .
275
276 createReduceKernel(ReduceSrcFunFile , FunName ) ->
277     createReduceKernel(ReduceSrcFunFile , FunName , cache) .
278
279 %%Create a reduce compatible kernel object from the source in ReduceSrcFunFile .
280 %%The kernel name is FunName
281 %%Program caching policy must also be specified .
282 -spec createReduceKernel(ReduceSrcFunFile , FunName , CacheKernel) -> {ok, kernel()} |
283 {error , Why} when
284     ReduceSrcFunFile :: nonempty_string() ,
285     FunName :: nonempty_string() ,
286     CacheKernel :: cache | no_cache ,
287     Why :: atom() .
288 createReduceKernel( , , ) -> {error , ecl_lib_not_loaded} .

```

```

289
290
291 %% -spec reduceHH(Kernel, Input) -> ok | {error, Why} when
292 %%     Kernel    :: kernel(),
293 %%     Input     :: hostBuffer(),
294 %%     Why       :: atom().
295 %%
296 %% reduceHH(_, _) -> {error, ecl_lib_not_loaded}.
297
298 %%Reduce skeleton working on device buffers. User function specified as kernel
299 %% objects.
300 -spec reduceDD(Kernel, Input, Output) -> ok | {error, Why} when
301     Kernel    :: kernel(),
302     Input     :: deviceBuffer(),
303     Output    :: deviceBuffer(),
304     Why       :: atom().
305
306 reduceDD(_, _, _) -> {error, ecl_lib_not_loaded}.
307
308 %%Reduce skeleton: input is a device buffer, output is a list. User function
309 %% specified as kernel objects.
310 -spec reduceDL(Kernel, Input) -> {ok, Result} | {error, Why} when
311     Kernel    :: kernel(),
312     Input     :: deviceBuffer(),
313     Result    :: [float()],
314     Why       :: atom().
315
316 %%Reduce skeleton: both input and output are lists. User function specified as
317 %% kernel objects.
318 -spec reduceLL(Kernel, Input, InputLength) -> {ok, Result} | {error, Why} when
319     Kernel    :: kernel(),
320     Input     :: [float()],
321     InputLength :: non_neg_integer(),
322     Result    :: [float()],
323     Why       :: atom().
324
325 reduceLL(ReduceKernel, InputList, InputListLength) ->
326     SzDouble = 8,
327     SzBufferByte = InputListLength * SzDouble,
328
329 %%allocate host buffer and initialize it using InputList
330 InputH = checkResult( allocHostBuffer(SzBufferByte) ),
331 checkResult(listToBuffer(InputList, InputH)),
332

```

```

333 %%allocate device buffer and copy the input data from the host buffer.
334 InputD = checkResult( allocDeviceBuffer(SzBufferByte , read_write) ) ,
335 copyBufferToBuffer(InputH , InputD) ,
336
337 %%don't need host buffer anymore
338 releaseBuffer (InputH) ,
339
340 %%compute Reduce using the DL version
341 ResultList = reduceDL(ReduceKernel , InputD) ,
342
343 releaseBuffer (InputD) ,
344
345 ResultList
346 .
347
348
349 %%Library info
350 skeletonlib(_)->
351     {error , ecl_lib_not_loaded}.
352
353 %%initialize Skel COL
354 cl_init()->
355     {error , ecl_lib_not_loaded}.
356
357 cl_release()->
358     {error , ecl_lib_not_loaded}.
359
360
361 %%----- Skeleton Composition example -----
362 %%
363 -spec mapReduceLL(MapKernel , ReduceKernel , {InputList , InputListLength}) -> {ok ,
364     Result} when
365     MapKernel :: kernel() ,
366     ReduceKernel:: kernel() ,
367     InputList :: [float()] ,
368     InputListLength :: non_neg_integer() ,
369     Result :: [float()]
370 .
371
372 mapReduceLL(MapKernel , ReduceKernel , {InputList , InputListLength}) ->
373     SzDouble = 8 ,
374     SzBufferByte = InputListLength * SzDouble ,
375
376     MapOutputD = checkResult( allocDeviceBuffer(SzBufferByte , read_write) ) ,
377
378     checkResult( mapLD(MapKernel , InputList , MapOutputD , InputListLength) ) ,

```

```
379
380     Result = reduceDL( ReduceKernel , MapOutputD ) ,
381
382     releaseBuffer (MapOutputD) ,
383
384     Result
385 .
386
387 %%Unwrap the Result tuple returning the second element , if the first is an error
388 %%atom throw an exception
388 -spec checkResult( Result ) -> Result | Why when
389     Result :: {ok, Result} | {error, Why} ,
390     Result :: any() ,
391     Why :: atom() .
392
393 checkResult({ok, Ref}) when is_reference(Ref)->
394     Ref;
395 checkResult({ok, Result}) ->
396     Result ;
397 checkResult({error, Why}) ->
398     erlang:error(Why) ;
399 checkResult(ok) ->
400     ok.
```

B.2 skel_ocl.so

B.2.1 skel_ocl.cpp

```

1 #include "erl_nif.h"
2
3 #include <iostream>
4 #include <fstream>
5 #include <sstream>
6 #include <cstring>
7
8 #include <cerrno>
9
10 #include <vector>
11 #include <unordered_map>
12
13 #include "OCL2.h"
14
15 #include "errors.cpp"
16
17 #include "utils.cpp"
18
19 #include "TimeUtils.h"
20
21 // Kernels ' code
22 #include "map_kernel.h"
23 #include "reduce_kernel.h"
24
25 using namespace std;
26
27
28 #define ONE 1
29 #define TWO 2
30
31 const uint NUM_QUEUES = TWO; //ONE or TWO
32
33 const uint MAX_QUEUES = TWO;
34
35
36 static bool ocl_initialised = false;
37
38 #define OCL_INIT_CHECK() if(! ocl_initialised) return make_error(env, ATOM(
39             skel_ocl_not_initialized));
40
41 #define NIF_ARITY_CHECK(ARITY) if(argc != ARITY) return enif_make_badarg(env);
42

```

```

43 // defined in OCL2.cpp
44 extern uint device_minBaseAddrAlignByte;
45
46
47
48 /***** RESOURCE OBJECTS *****/
49
50 // Define a destructor for OBJTYPE using BODY
51 #define DEFDTOR(OBJTYPE, BODY) \
52     static void dtor_##OBJTYPE(ErlNifEnv* env, void* obj) { \
53         OBJTYPE* _obj = (OBJTYPE*) obj; \
54         if(*_obj) { \
55             BODY \
56         } \
57     }
58
59 // Host and device buffers
60 static ErlNifResourceType* hostBuffer_rt = NULL;
61 static ErlNifResourceType* deviceBuffer_rt = NULL;
62
63 DEFDTOR(cl_mem, (NULL) ;)
64
65 //CL Program
66 static ErlNifResourceType* program_rt = NULL;
67 DEFDTOR(cl_program, clReleaseProgram(*_obj);)
68
69
70 //CL Kernel with synchronized access
71 static ErlNifResourceType* kernel_sync_rt = NULL;
72
73 typedef struct _kernel_sync {
74     cl_kernel kernel;
75     ErlNifMutex *mtx;
76 } kernel_sync;
77
78 static kernel_sync* ctor_kernel_sync(cl_kernel kernel) {
79
80     kernel_sync* pKernel_s = (kernel_sync*) enif_alloc_resource( kernel_sync_rt,
81         sizeof(kernel_sync) );
82
83     pKernel_s->mtx = enif_mutex_create((char*)"setkernelArgMutex");
84
85     pKernel_s->kernel = kernel;
86
87     return pKernel_s;
88 }
```

```

89 static void dtor_kernel_sync(ErlNifEnv* env, void* obj)
90 {
91     kernel_sync* ker_s = (kernel_sync*) obj;
92     if(ker_s) {
93
94 #ifdef DEBUG
95         cerr << "DEBUG :: kernel_sync_dtor" << endl;
96 #endif
97     enif_mutex_destroy(ker_s->mtx);
98
99     clReleaseKernel(ker_s->kernel);
100 }
101 }
102
103 //CL event
104 static ErlNifResourceType* event_rt = NULL;
105 DEFDTOR(cl_event, clReleaseEvent(*_obj));
106
107 /*
108 * Worker thread handler, exploits NIF garbage collection mechanism
109 * so to be sure to join a thread before module unload (as required by NIF spec)
110 */
111 static ErlNifResourceType* thread_handler_rt = NULL;
112 /*thread handler: it just joins a thread*/
113 DEFDTOR(ErlNifTid, enif_thread_join(*_obj, NULL));
114
115
116 //Resource object declaration macro, used in the load function
117 #define DECLRESOURCEOBJ(RTYPEVAR, NAME, DTOR, FLAGS) \
118 { ErlNifResourceType* rt = enif_open_resource_type(env, NULL, NAME, DTOR, FLAGS, \
119     NULL); \
120     if (rt == NULL) return -1; \
121     RTYPEVAR = rt; \
122 }
123 static int load(ErlNifEnv* env, void** priv_data, ERL_NIF_TERM load_info)
124 {
125
126     /*BUFFERS RO */
127     DECLRESOURCEOBJ(hostBuffer_rt, "host_buffer", dtor_cl_mem, ERL_NIF_RT_CREATE)
128     DECLRESOURCEOBJ(deviceBuffer_rt, "device_buffer", dtor_cl_mem, ERL_NIF_RT_CREATE
129     )
130
131     /*Program and Kernel RO*/
132     DECLRESOURCEOBJ(program_rt, "program", dtor_cl_program, ERL_NIF_RT_CREATE)
133     DECLRESOURCEOBJ(kernel_sync_rt, "kernel", dtor_kernel_sync, ERL_NIF_RT_CREATE)

```

```

134  /*Event RO*/
135  DECL_RESOURCE_OBJ(event_rt, "event", dtor_cl_event, ERL_NIF_RT.CREATE)
136
137  /*worker thread RO*/
138  DECL_RESOURCE_OBJ(thread_handler_rt, "thread_handler", dtor_ErlNifTid,
139                      ERL_NIF_RT.CREATE)
140
141 }
142
143
144 /***** BUFFER MANAGEMENT NIFS *****/
145
146 static const char* MEM_READ_FLAG = "read";
147 static const char* MEM_WRITE_FLAG = "write";
148 static const char* MEM_READ_WRITE_FLAG = "read-write";
149
150 static const size_t MEM_FLAG_ATM_MAX_LEN = strlen(MEM_READ_WRITE_FLAG)+1;
151
152 static bool parseMemFlags(const char* flag, cl_mem_flags* parsedFlag) {
153
154     //parse flag
155     if(strcmp(MEM_READ_FLAG, flag) == 0)
156         *parsedFlag |= CL_MEM_READ_ONLY;
157     else
158         if(strcmp(MEM_WRITE_FLAG, flag) == 0)
159             *parsedFlag |= CL_MEM_WRITE_ONLY;
160     else
161         if(strcmp(MEM_READ_WRITE_FLAG, flag) == 0)
162             *parsedFlag |= CL_MEM_READ_WRITE;
163     else
164         return false;
165
166     return true;
167 }
168
169
170 static ERL_NIF_TERM getBufferSize(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[])
171 {
172     OCL_INIT_CHECK()
173
174     /*get the parameter (Buffer::buffer())***** */
175     NIF_ARITY_CHECK(1)
176
177     cl_mem* pBuffer = NULL;

```

```

178  if (!enif_get_resource(env, argv[0], hostBuffer_rt, (void**) &pBuffer)) { //is it
179    a hostBuffer?
180
180  if (!enif_get_resource(env, argv[0], deviceBuffer_rt, (void**) &pBuffer)) { //or
181    a deviceBuffer?
181    cerr << "ERROR :: releaseBuffer: 1st parameter is not a buffer" << endl ;
182    return enif_make_badarg(env);
183  }
184 }
185 if(*pBuffer == NULL)
186   return enif_make_badarg(env);
187
188 /****** */
189
190
191 size_t szBufferByte = 0;
192
193 CHK_SUCCESS(getBufferSizeByte(*pBuffer, &szBufferByte));
194
195
196 return enif_make_ulong(env, szBufferByte);
197 }
198
199
200 static ERL_NIF_TERM allocDeviceBuffer(ErlNifEnv * env, int argc, const ERL_NIF_TERM
201                                     argv[]) {
202
202  // -type rw_flags() :: read | write | read_write
203  // allocBuffer(SizeByte::non_neg_integer(), Flags::rw_flag()) -> {ok, Buffer} | {
203    error, Why}
204
205  OCL_INIT_CHECK()
206
207 /*get the parameter (SizeByte::non_neg_integer(), Flags::rw_flag())***** */
207 */
208  NIF_ARITY_CHECK(2)
209
210 //SizeByte::non_neg_integer()
211 size_t szBufferByte;
212 if (!enif_get_ulong(env, argv[0], &szBufferByte)) {
213
214   cerr << "ERROR :: allocDeviceBuffer: 1st parameter is not a non-neg-integer()"
214     << endl ;
215   return enif_make_badarg(env);
216 }
217
218 //Flags::rw_flag()

```

```

219  char flag [MEM_FLAG_ATM_MAX_LEN];
220  if (!enif_get_atom(env, argv[1], flag, MEM_FLAG_ATM_MAX_LEN, ERL_NIF_LATIN1)) {
221
222      cerr << "ERROR :: allocDeviceBuffer: 2nd parameter is not an rw_flag()" << endl
223      ;
224
225  //*****
226
227
228  //parse flag
229 cl_mem_flags mem_flags = 0;
230
231 if (!parseMemFlags(flag, &mem_flags)) {
232     cerr << "ERROR :: allocDeviceBuffer: 2nd parameter is not an rw_flag()" << endl
233     ;
234
235     return enif_make_badarg(env);
236 }
237
238 /*allocate the memory for the resource (a cl_mem) */
239 cl_mem* pBuffer = NULL;
240
241 //allocate the buffer
242 CHK_SUCCESS_CLEANUP(createBuffer(szBufferByte, mem_flags, pBuffer); ,
243 enif_release_resource(pBuffer));
244
245 /*grant co-ownership to erlang*/
246 ERL_NIF_TERM bufferRO = enif_make_resource(env, pBuffer);
247
248 /*DON'T transfer the ownership to erlang because a releaseBuffer NIF is provided*/
249
250 return enif_make_tuple2(env, ATOM(ok), bufferRO);
251
252 static ERL_NIF_TERM allocHostBuffer(ErlNifEnv * env, int argc, const ERL_NIF_TERM
253 argv[]) {
254
255     //allocHostBuffer(SizeByte::non_neg_integer()) -> {ok, Buffer} | {error, Why}
256
257     OCL_INIT_CHECK()
258
259     /*get the parameter (SizeByte::non_neg_integer())******/
260     NIF_ARITY_CHECK(1)
261
262 }
```

```

261 // SizeByte::non_neg_integer()
262 size_t szBufferByte;
263 if (!enif_get_ulong(env, argv[0], &szBufferByte)) {
264     cerr << "ERROR :: allocHostBuffer: first parameter is not a non_neg_integer()"
265         << endl ;
266     return enif_make_badarg(env);
267 }
268 /*****
269 */
270 /*allocate the memory for the resource (a cl_mem) */
271 cl_mem* pBuffer = NULL;
272
273 pBuffer = (cl_mem*) enif_alloc_resource( hostBuffer_rt, sizeof(cl_mem));
274
275 //request buffer in (pinned) host memory
276 CHK_SUCCESS_CLEANUP(createBuffer(szBufferByte, CL_MEM_ALLOC_HOST_PTR, pBuffer), ,
277     enif_release_resource(pBuffer));
278
279 /*grant co-ownership to erlang*/
280 ERL_NIF_TERM bufferRO = enif_make_resource(env, pBuffer);
281
282 /*DON'T transfer the ownership to erlang because a releaseBuffer NIF is provided*/
283
284 return enif_make_tuple2(env, ATOM(ok), bufferRO);
285
286
287 static ERL_NIF_TERM releaseBuffer(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv
288 [] ) {
289 //releaseBuffer(Buffer::buffer())
290
291 OCL_INIT_CHECK()
292
293 /*get the parameter (Buffer::buffer())***** */
294 NIF_ARITY_CHECK(1)
295
296 cl_mem* pBuffer = NULL;
297 if (!enif_get_resource(env, argv[0], hostBuffer_rt, (void**) &pBuffer)) { //is it
298     a hostBuffer?
299     if (!enif_get_resource(env, argv[0], deviceBuffer_rt, (void**) &pBuffer)) { // or a
300         deviceBuffer?
301         cerr << "ERROR :: releaseBuffer: 1st parameter is not a buffer" << endl ;
302         return enif_make_badarg(env);

```

```

302     }
303 }
304 if(*pBuffer == NULL)
305     return enif_make_badarg(env);
306
307 /******
308 */
309
310 cl_int err =
311     clReleaseMemObject(*pBuffer);
312
313 *pBuffer = NULL;
314 enif_release_resource(pBuffer);
315
316 if(err != CL_SUCCESS)
317     return make_error_cl(env, err);
318
319 return ATOM(ok);
320
321 }
322
323
324 static ERL_NIF_TERM copyBufferToBufferSameSize(ErlNifEnv * env, int argc, const
325     ERL_NIF_TERM argv[]) {
326 //copyBufferToBuffer(From::buffer(), To::buffer()) -> ok
327
328 OCL_INIT_CHECK()
329
330 /*get the parameters (From::buffer(), To::buffer())******/
331 NIF_ARITY_CHECK(2)
332
333 cl_mem* src_buffer = NULL;
334 if (!enif_get_resource(env, argv[0], hostBuffer_rt, (void**) &src_buffer)) { //is
335     it a hostBuffer?
336
337     if (!enif_get_resource(env, argv[0], deviceBuffer_rt, (void**) &src_buffer)) {
338         cerr << "ERROR :: copyBufferToBufferSameSize: first parameter is not a buffer"
339             << endl ;
340
341     return enif_make_badarg(env);
342 }
343 if(*src_buffer == NULL)
344     return enif_make_badarg(env);
345

```

```

344     cl_mem* dst_buffer = NULL;
345     if (!enif_get_resource(env, argv[1], hostBuffer_rt, (void**) &dst_buffer)) { // is
346         a hostBuffer?
347     if (!enif_get_resource(env, argv[1], deviceBuffer_rt, (void**) &dst_buffer)) {
348         cerr << "ERROR :: copyBufferToBufferSameSize: 2nd parameter is not a buffer"
349             << endl ;
350     return enif_make_badarg(env);
351 }
352 if(*dst_buffer == NULL)
353     return enif_make_badarg(env);
354
355 //*****
356 */
357
358 #ifdef TIME
359     timespec fun_start, fun_end;
360     GET_TIME((fun_start));
361 #endif
362
363     size_t szSrcBufferByte = 0, szDestBufferByte = 0;
364
365     CHK_SUCCESS(getBufferSizeByte(*src_buffer, &szSrcBufferByte));
366     CHK_SUCCESS(getBufferSizeByte(*dst_buffer, &szDestBufferByte));
367
368 // Input Buffers must have the same size
369 if(szSrcBufferByte != szDestBufferByte)
370     return make_error(env, ATOM(skel_ocl_buffers_different_size));
371
372     CHK_SUCCESS(copyBuffer(*src_buffer, *dst_buffer, szSrcBufferByte));
373
374
375 #ifdef SEQ
376     clFinish(getCommandQueue(0));
377 #endif
378 #ifdef TIME
379     GET_TIME((fun_end));
380
381     cerr << "copy: " << nsec2usec(diff(&(fun_start), &(fun_end))) << endl;
382 #endif
383
384
385     return ATOM(ok);
386 }
```

```

387
388
389
390 static ERL_NIF_TERM copyBufferSize(ErlNifEnv * env, int argc, const
391 ERL_NIF_TERM argv[]) {
392 //copyBufferToBuffer(From:: buffer(), To:: buffer(), CopySizeByte:: non_neg_integer()
393 //) -> ok
394
395 OCL_INIT_CHECK()
396 /*get the parameters (From:: buffer(), To:: buffer(), CopySizeByte:: non_neg_integer()
397 () *****/
398 NIF_ARITY_CHECK(3)
399
400 cl_mem* src_buffer = NULL;
401 if (!enif_get_resource(env, argv[0], hostBuffer_rt, (void**) &src_buffer)) { //is
402 // it a hostBuffer?
403
404 if (!enif_get_resource(env, argv[0], deviceBuffer_rt, (void**) &src_buffer)) {
405 // or a deviceBuffer?
406 cerr << "ERROR :: copyBufferSize: first parameter is not a buffer" <<
407 endl ;
408 return enif_make_badarg(env);
409 }
410
411 if (*src_buffer == NULL)
412 return enif_make_badarg(env);
413
414 cl_mem* dst_buffer = NULL;
415 if (!enif_get_resource(env, argv[1], hostBuffer_rt, (void**) &dst_buffer)) { //is
416 // a hostBuffer?
417
418 if (!enif_get_resource(env, argv[1], deviceBuffer_rt, (void**) &dst_buffer)) {
419 // or a deviceBuffer?
420 cerr << "ERROR :: copyBufferSize: 2nd parameter is not a buffer" <<
421 endl ;
422 return enif_make_badarg(env);
423
424 //CopySizeByte:: non_neg_integer()
425 size_t szCopyByte;
426 if (!enif_get_ulong(env, argv[2], &szCopyByte)) {
427

```

```

425     cerr << "ERROR :: copyBufferToBuffer: 3rd parameter is not a non_neg_integer()">>
426     << endl ;
427     return enif_make_badarg(env);
428 }
429 /*****************************************************************************
430 */
431 CHK_SUCCESS(copyBuffer(*src_buffer, *dst_buffer, szCopyByte));
432
433 return ATOM(ok);
434 }
435
436
437
438 /** Copies the content of a list into a buffer*/
439 static ERL_NIF_TERM listToBuffer(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[])
440 {
441     //listToBuffer(From::[ float ()], To::hostBuffer() ) -> ok | {error , Why}
442     OCL_INIT_CHECK()
443
444     /*get the parameters (To::list())******/
445     NIF_ARITY_CHECK(2)
446
447     ERL_NIF_TERM list = argv[0];
448     if (!enif_is_list(env, list)) {
449         cerr << "ERROR :: listToBuffer: first parameter is not a list" << endl;
450         return enif_make_badarg(env);
451     }
452
453     cl_mem* pBuffer = NULL;
454     if (!enif_get_resource(env, argv[1], hostBuffer_rt, (void**) &pBuffer) || *pBuffer
455         == NULL) {
456         cerr << "ERROR :: listToBuffer: second parameter is not a buffer" << endl ;
457         return enif_make_badarg(env);
458     }
459     *****
460 #ifdef TIME
461     timespec fun_start, fun_end;
462     GET_TIME((fun_start));
463 #endif
464
465     //copy
466     unsigned int list_len = 0, szListBytes = 0;
467

```

```

468     enif_get_list_length(env, list, &list_len);
469     szListBytes = list_len * sizeof(double);
470
471
472     double* pMappedBuffer = NULL;
473     //obtain a pointer to work on the buffer (out of bounds check done by OpenCL,
474     //returns CL_INVALID_VALUE)
474     CHK_SUCCESS( mapBufferBlocking(*pBuffer, 0, szListBytes, CL_MAP_WRITE, &
475         pMappedBuffer) );
475
476     if(pMappedBuffer == NULL) {
477         cerr << "pMappedBuffer is NULL." << endl;
478         return make_error(env, ATOM(cl_error));
479     }
480
481
482     // copies the content of input list into pBuffer
483     list_to_double_array(env, list, pMappedBuffer);
484
485     //unMap when done copying
486     CHK_SUCCESS( unMapBuffer(*pBuffer, pMappedBuffer) );
487
488
489 #ifdef TIME
490     GET_TIME((fun_end));
491
492     cerr << "Unmarshalling: " << nsec2usec(diff(&(fun_start), &(fun_end))) << endl;
493 #endif
494
495     return ATOM(ok);
496 }
497
498 /** Create a list containing the elements of a buffer*/
499 static ERL_NIF_TERM _bufferToListLength(ErlNifEnv * env, cl_mem* pBuffer, size_t
500     listLen) {
501
502     size_t listLenByte = listLen * sizeof(double);
503
504     double* pDoubleArray = NULL;
505
506     //obtain a pointer to work on the buffer
507     CHK_SUCCESS( mapBufferBlocking(*pBuffer, 0, listLenByte, CL_MAP_READ, &
508         pDoubleArray) );
509
510     ERL_NIF_TERM list = double_array_to_list(env, pDoubleArray, listLen);

```

```

511     CHK_SUCCESS( unMapBuffer(*pBuffer , pDoubleArray); )
512
513     return enif_make_tuple2(env , ATOM(ok) , list );
514
515 }
516
517
518 /** Create a list containing the elements of a buffer*/
519 static ERL_NIF_TERM bufferToListLength(ErlNifEnv * env , int argc , const ERL_NIF_TERM
520     argv []) {
521     //%%bufferToList(From::buffer() , Length::pos_integer()) -> {ok , [float()]}
```

522 OCL_INIT_CHECK()

523 /*get the parameters (From::Buffer , Length::pos_integer())******/

524 NIF_ARITY_CHECK(2)

525

526 cl_mem* pBuffer = NULL;

527 **if** (!enif_get_resource(env , argv[0] , hostBuffer_rt , (**void****) &pBuffer) || *pBuffer
 == NULL) {

528 cerr << "ERROR :: bufferToList: 1st parameter is not a host buffer" << endl ;
 return enif_make_badarg(env);
 }

529 //Length::pos_integer()

530 size_t listLen;

531 **if** (!enif_get_ulong(env , argv[1] , &listLen)) {

532 cerr << "ERROR :: bufferToList: 2nd parameter is not a non_neg_integer()" <<
 endl ;
 return enif_make_badarg(env);
 }

533 // check that buffer contains enough elements

534

535 size_t szBufferByte = 0;

536 CHK_SUCCESS(getBufferSizeByte(*pBuffer , &szBufferByte);)

537 size_t listLenByte = listLen * **sizeof(double)**;

538 **if**(listLenByte > szBufferByte) {

539 cerr << "ERROR :: bufferToList: buffer doesn't contain Length elements" << endl
 ;

```

554     return enif_make_badarg(env);
555 }
556
557 return _bufferToListLength(env, pBuffer, listLen);
558 }
559
560 /** Create a list containing the elements of a buffer*/
561 static ERL_NIF_TERM bufferToList(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv
562 [ ]) {
563     //bufferToList(From::Buffer) -> {ok, []} | {error, Why}
564     OCL_INIT_CHECK()
565     /*get the parameter (From::Buffer)******/
566
567     NIF_ARITY_CHECK(1)
568
569     cl_mem* pBuffer = NULL;
570     if (!enif_get_resource(env, argv[0], hostBuffer_rt, (void**) &pBuffer) || *pBuffer
571         == NULL) {
572
573         cerr << "ERROR :: bufferToList: first parameter is not a host buffer" << endl ;
574         return enif_make_badarg(env);
575     } *****
576
577
578 #ifdef TIME
579     timespec fun_start, fun_end;
580     GET_TIME((fun_start));
581 #endif
582
583     size_t szBufferByte;
584
585     CHK_SUCCESS( getBufferSizeByte(*pBuffer, &szBufferByte); )
586
587     size_t numBufferElements = szBufferByte / sizeof(double);
588
589     ERL_NIF_TERM toReturn = _bufferToListLength(env, pBuffer, numBufferElements);
590
591
592 #ifdef TIME
593     GET_TIME((fun_end));
594
595     cerr << "Marshalling: "<< nsec2usec(diff(&(fun_start), &(fun_end))) << endl;
596 #endif
597
598     return toReturn;

```

```

599 }
600
601
602
603
604 ***** PROGRAM AND KERNEL NIFS *****
605
606 static ERL_NIF_TERM buildProgramFromCharP(ErlNifEnv * env, const char* progSrc){
607
608 /*allocate the memory for the resource (a cl_program) */
609 cl_program* pProgram = NULL;
610
611 pProgram = (cl_program*) enif_alloc_resource(
612         program_rt,
613         sizeof(cl_program)
614 );
615
616 //create and build the program releasing the resource in case of errors
617 CHK_SUCCESS_CLEANUP(createBuildProgramFromString(progSrc, pProgram), ,
618     enif_release_resource(pProgram));
619
620 // enif_free(progSrc); should be freed by the caller
621
622 /*grant co-ownership to erlang*/
623 ERL_NIF_TERM programRO = enif_make_resource(env, pProgram);
624
625 /*transfer the ownership to erlang*/
626 enif_release_resource(pProgram);
627
628 return enif_make_tuple2(env, ATOM(ok), programRO);
629
630
631 //Unmarshal an Erlang list allocating the needed space
632 #define ENIF_GET_STRING(CHAR_P, ARG) \
633 { \
634     unsigned int strLen = 0; /*without \0 */ \
635     if(!enif_get_list_length(env, ARG, &strLen)) \
636         return enif_make_badarg(env); \
637     CHAR_P = (char *) enif_alloc(strLen++); /*including null terminator*/ \
638     enif_get_string(env, ARG, CHAR_P, strLen, ERL_NIF_LATIN1); \
639 }
640
641 #define ENIF_GET_ATOM(CHAR_P, ARG) \
642 { \
643     uint atom_len = 0; \
644     if (!enif_get_atom_length(env, ARG,&atom_len, ERL_NIF_LATIN1)) \

```

```

645     return enif_make_badarg(env); \
646     CHAR_P = (char *) enif_alloc(atom_len++); /*including null terminator*/ \
647     enif_get_atom(env, ARG, CHAR_P, atom_len, ERL_NIF_LATIN1); \
648 }
649
650 static ERL_NIF_TERM buildProgramFromString(ErlNifEnv * env, int argc, const
651   ERL_NIF_TERM argv[]) {
652   //buildProgramFromString(ProgramSrcString::iolist())
653   OCL_INIT_CHECK()
654   /*get the parameter (ProgramSrcString::iolist())***** */
655   /*
656   NIF_ARITY_CHECK(1)
657   //ProgramSrcString may be a binary
658   char *progSrc = NULL;
659   ErlNifBinary bin;
660
661   if(enif_is_binary(env, argv[0])) {
662
663     enif_inspect_iolist_as_binary(env, argv[0], &bin);
664
665     progSrc = (char*) enif_alloc(bin.size+1);
666
667     strncpy(progSrc, (char*) bin.data, bin.size);
668
669   }
670   else //ProgramSrcString must then be a list
671     ENIF_GET_STRING(progSrc, argv[0])
672
673   /****** */
674
675   ERL_NIF_TERM toReturnTerm = buildProgramFromCharP(env, progSrc);
676
677   if (!progSrc)
678     enif_free(progSrc);
679
680   return toReturnTerm;
681
682 }
683
684 static ERL_NIF_TERM buildProgramFromFile(ErlNifEnv * env, int argc, const
685   ERL_NIF_TERM argv[]) {
686   //buildProgram(ProgramSrcPath::nonempty_string())
687   OCL_INIT_CHECK()

```

```

688  /*get the parameter (ProgramSrcPath::nonempty_string())
   *****/
689  NIF_ARITY_CHECK(1)
690
691  char *progPath = NULL;
692
693  ENIF_GET_STRING(progPath, argv[0]) // progPath must be enif_free'd
694
695  /*****
696  */
697 //Read from the file
698 char* progSrc = NULL;
699
700 if(! (progSrc = readFromFile(progPath)) )
701     return enif_make_tuple2(env, ATOM(error), enif_make_atom(env,"opening_file"));
702
703
704 ERL_NIF_TERM toReturn = buildProgramFromCharP(env, progSrc);
705
706 enif_free(progPath);
707 enif_free(progSrc);
708
709 return toReturn;
710
711 }
712
713 static ERL_NIF_TERM _createKernel(ErlNifEnv * env, cl_program* pProgram, const char*
    kerName) {
714
715     cl_kernel kernel = NULL;
716     CHK_SUCCESS(createKernel(*pProgram, kerName, &kernel));
717
718     /*crate a kernel_sync*/
719     kernel_sync* pKernel_s = ctor_kernel_sync(kernel);
720
721     /*grant co-ownership to erlang*/
722     ERL_NIF_TERM kernelRO = enif_make_resource(env, pKernel_s);
723
724     /*transfer the ownership to erlang*/
725     enif_release_resource(pKernel_s);
726
727
728     return enif_make_tuple2(env, ATOM(ok), kernelRO);
729 }
730

```

```

731 static ERL_NIF_TERM createKernel(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv
    []) {
732
733 //%%createKernel(Prog::program() , KerName::nonempty_string())
734 OCL_INIT_CHECK()
735 /*get the parameters (Prog::program() , KerName::nonempty_string())******/
736 NIF_ARITY_CHECK(2)
737
738 // get Program
739
740 cl_program* pProgram = NULL;
741 if (!enif_get_resource(env, argv[0], program_rt, (void**) &pProgram)) {
742     cerr << "ERROR :: createKernel: 1st parameter is not a program" << endl ;
743     return enif_make_badarg(env);
744 }
745
746 char* kerName = NULL;
747 ENIF_GET_STRING(kerName, argv[1])
748
749 /*****
749 */
750
751 ERL_NIF_TERM toReturn = _createKernel(env, pProgram, kerName);
752
753 enif_free(kerName);
754
755 return toReturn;
756
757 }
758
759
760 static ERL_NIF_TERM generateProgramSrcFromFile(ErlNifEnv * env, char* funSrcPath,
761         string fun_name, string type, string mapKernelStr, string& progSrcStr) {
762
763     string funSrcStr = readFromFileStr(funSrcPath);
764
765     if( funSrcStr.compare("NULL") == 0)
766         return enif_make_tuple2(env, ATOM(error), ATOM(opening_file));
767
768     string kernelSrcStr(mapKernelStr);
769
770     replaceTextInString(kernelSrcStr, std::string("TYPE"), type);
771     replaceTextInString(kernelSrcStr, std::string("FUN_NAME"), fun_name);
772
773     //generate the final program from the src from file and kernel definition
774
775     if(type.compare("double") == 0)

```

```

775     progSrcStr.append("#pragma OPENCL EXTENSION cl_khr_fp64 : enable\n");
776
777     progSrcStr.append(funSrcStr);
778     progSrcStr.append(kernelSrcStr);
779
780     return ATOM(ok);
781 }
783
784
785 typedef unordered_map<std::string, cl_program> CLPrograms_Cache;
786
787 //Store built programs so to avoid building them again when used more than once
788 static CLPrograms_Cache* clPrograms_Cache;
789
790
791 static ERL_NIF_TERM createSkeletonKernelFromFile(ErlNifEnv * env, char*
792     skeletonFunSrcFilePath, const string& skeletonFunName, const string&
793     skeletonSrcStr, const char* kernelName, const string& type, bool cache_program)
794 {
795
796     cl_program program;
797
798     string skeletonFunSrcPath_str = string(skeletonFunSrcFilePath);
799
800     if(cache_program && (clPrograms_Cache->count(skeletonFunSrcPath_str) > 0)) { //1
801         found, 0 not found
802
803         #ifdef DEBUG
804             cerr<< "DEBUG - looking for cached program.\n";
805         #endif
806         program = (*clPrograms_Cache)[skeletonFunSrcPath_str]; //use cached program
807
808     } else {
809
810         #ifdef DEBUG
811             cerr<< "DEBUG - generating program.\n";
812
813             ERL_NIF_TERM result_atom =
814                 generateProgramSrcFromFile(env, skeletonFunSrcFilePath, skeletonFunName,
815                                         type, skeletonSrcStr, progSrcStr);
816
817         if(result_atom != ATOM(ok))

```

```

817     return result_atom;
818
819     const char* progSrc = progSrcStr.c_str();
820
821     CHK_SUCCESS(createBuildProgramFromString(progSrc, &program)); // program never
822         released
823
824     if(cache_program) {
825         #ifdef DEBUG
826             cerr<< "DEBUG - adding new program to cache.\n";
827         #endif
828         (*clPrograms_Cache)[skeletonFunSrcPath_str] = program; //save built program
829             into cache
830     }
831
832     return _createKernel(env, &program, kernelName);
833
834
835 }
836
837 static ERL_NIF_TERM createMapKernelFromFile(ErlNifEnv * env, int argc, const
838     ERL_NIF_TERM argv[]) {
839
840     //createMapKernel(MapSrcFunFile::nonempty_string(), FunName::nonempty_string(),
841         FunArity::pos_integer(), ProgCaching::atom())
842     OCL_INIT_CHECK()
843     /*get the parameter (MapFunSrcPath::nonempty_string())***** */
844     NIF_ARITY_CHECK(4)
845
846     char *mapFunSrcPath = NULL;
847     ENIF_GET_STRING(mapFunSrcPath, argv[0]) // mapFunSrcPath must be enif_free'd
848
849     char *mapFunName = NULL;
850     ENIF_GET_STRING(mapFunName, argv[1]) // mapFunName must be enif_free'd
851
852     uint funArity = 0;
853     enif_get_uint(env, argv[2], &funArity);
854
855     if (! enif_is_atom(env, argv[3]))
856         return enif_make_badarg(env);
857
858     ERL_NIF_TERM prog_caching_atom = argv[3];
859

```

```

858  /*************************************************************************/
859  */
860
861 const string* pSkeletonSrcStr;
862 const char* mapKernelName;
863
864 switch (funArity)
865 { case 1 : mapKernelName = "MapKernel";
866     pSkeletonSrcStr = &MapKernelStr;
867     break;
868
869 case 2 : mapKernelName = "MapKernel2";
870     pSkeletonSrcStr = &Map2KernelStr;
871     break;
872
873 default: enif_free(mapFunSrcPath); //FunArity not supported
874     enif_free(mapFunName);
875     return make_error(env, ATOM(skel_ocl_fun_arity_not_supported));
876 }
877
878 const char* type = "double";
879
880 bool cache_program = enif_is_identical(ATOM(cache), prog_caching_atom);
881
882 ERL_NIF_TERM toReturn =
883     createSkeletonKernelFromFile(env, mapFunSrcPath, string(mapFunName), *
884         pSkeletonSrcStr, mapKernelName, type, cache_program);
885
886 enif_free(mapFunSrcPath);
887 enif_free(mapFunName);
888
889 return toReturn;
890 }
891
892 static ERL_NIF_TERM createReduceKernelFromFile(ErlNifEnv * env, int argc, const
893 ERL_NIF_TERM argv[]) {
894     //buildProgram(ReduceFunSrcPath::nonempty_string(), ReduceFunName::nonempty_string
895     ())
896     OCL_INIT_CHECK()
897     /*get the parameter (ReduceFunSrcPath::nonempty_string())***** */
898     */
899     NIF_ARITY_CHECK(3)
900
901     char *reduceFunSrcPath = NULL;

```

```

900
901     ENIF_GET_STRING(reduceFunSrcPath, argv[0]) // reduceFunSrcPath must be enif_free'd
902
903     char *reduceFunName = NULL;
904
905     ENIF_GET_STRING(reduceFunName, argv[1]) // reduceFunName must be enif_free'd
906
907
908     if (!enif_is_atom(env, argv[2]))
909         return enif_make_badarg(env);
910
911     ERL_NIF_TERM prog_caching_atom = argv[2];
912
913     /*****
914     */
915
916     const char* reduceKernelName = "ReduceKernel";
917     const char* type = "double";
918
919     bool cache_program = enif_is_identical(ATOM(cache), prog_caching_atom);
920
921     ERL_NIF_TERM toReturn =
922         createSkeletonKernelFromFile(env, reduceFunSrcPath, string(reduceFunName),
923                                     ReduceKernelStr, reduceKernelName, type, cache_program);
924
925     enif_free(reduceFunSrcPath);
926     enif_free(reduceFunName);
927
928     return toReturn;
929 }
930
931 /***** SKELETONS *****/
932
933
934 #include "map_nifs.cpp"
935
936 #include "reduce_nifs.cpp"
937
938
939
940
941
942
943 /***** MISCs *****/
944

```

```

945 /**
946 * Initialize OpenCL.
947 * Returns
948 * - ok
949 * - {error, "OpenCL already initialized."} when called more than once
950 * - badarg in the case of any error during the initialization
951 */
952 static ERL_NIF_TERM initOCL(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
953
954     if (!ocl_initialised) {
955
956         if (!init())
957             return make_error(env, ATOM(openCL_init_failed));
958         else {
959             ocl_initialised = true;
960
961             //init cl_program cache
962             clPrograms_Cache = new CLPrograms_Cache();
963
964             return ATOM(ok);
965         }
966     }
967     else return
968         make_error(env, ATOM(openCL_already_init));
969 }
970
971 static ERL_NIF_TERM releaseOCL(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[])
972 {
973     if (ocl_initialised) {
974
975 #ifdef DEBUG
976         std::cerr << "calling releaseOCL()" << std::endl;
977 #endif
978         releaseOCL();
979 #ifdef DEBUG
980         std::cerr << "releaseOCL() ok" << std::endl;
981 #endif
982
983         delete clPrograms_Cache;
984
985         ocl_initialised = false;
986
987         return ATOM(ok);
988     }
989     else return enif_make_tuple2(
990         env,

```

```

991     ATOM(error),
992     ATOM(OpenCL_not_initialized)
993 );
994
995 }
996
997 /* NIF interface bindings */
998 static ErlNifFunc nif_funcs[] = {
999
1000 {"skeletonlib", 1, skeletonlib},
1001 {"cl_init", 0, initOCL},
1002 {"cl_release", 0, releaseOCL},
1003
1004 //Buffer
1005 {"listToBuffer", 2, listToBuffer},
1006 {"bufferToList", 1, bufferToList},
1007 {"bufferToListLength", 2, bufferToListLength},
1008 {"getBufferSize", 1, getBufferSize},
1009 {"allocDeviceBuffer", 2, allocDeviceBuffer},
1010 {"allocHostBuffer", 1, allocHostBuffer},
1011 {"releaseBuffer", 1, releaseBuffer},
1012 {"copyBufferToBuffer", 2, copyBufferToBufferSameSize},
1013 {"copyBufferToBuffer", 3, copyBufferToBufferSize},
1014
1015 //Program and kernel
1016 {"buildProgramFromFile", 1, buildProgramFromFile},
1017 {"buildProgramFromString", 1, buildProgramFromString},
1018 {"createKernel", 2, createKernel},
1019
1020 //MAP
1021 {"createMapKernel", 4, createMapKernelFromFile},
1022 {"mapDD", 3, mapDD},
1023 {"mapLD", 4, mapLD},
1024 {"mapLL", 3, mapLL},
1025
1026 {"map2DD", 4, map2DD},
1027 {"map2LD", 5, map2LD},
1028 {"map2LL", 4, map2LL},
1029
1030 //REDUCE
1031 {"createReduceKernel", 3, createReduceKernelFromFile},
1032 {"reduceDD", 3, reduceDD},
1033 {"reduceDL", 2, reduceDL}
1034
1035 };
1036
1037 ERL_NIF_INIT(skel_ocl, nif_funcs, load, NULL, NULL, NULL)

```

B.2.2 Map NIFs: map_nifs.cpp

```

1 /***** MAP *****/
2
3
4 /*Map on device buffers*/
5
6 /*Implementation of n-ary map using device buffers */
7 /* Requirements:
8 * - All buffers in pInputBuffrV must have the same size ,
9 * - pInputBufferC > 0
10 */
11 static ERL_NIF_TERM mapDD_Impl(ErlNifEnv * env, kernel_sync* pKernel_s, void**
12     -pInputBufferV, uint pInputBufferC, void* -pOutputBuffer) {
13
14 #ifdef TIME
15     timespec
16     fun_start,
17     fun_prologue_end,
18     run_end,
19     fun_end;
20
21     GET_TIME((fun_start));
22 #endif
23
24     cl_mem** pInputBufferV = (cl_mem**) -pInputBufferV;
25     cl_mem* pOutputBuffer = (cl_mem*) -pOutputBuffer;
26
27     const uint OFFSET = 0;
28
29     //All buffers have the same size , then check just the first one
30     size_t szInputBufferByte = 0;
31     CHK_SUCCESS(getBufferSizeByte(*(pInputBufferV[0]), &szInputBufferByte));
32
33     //*****set map kernel parameters*****
34     uint uiNumElem = szInputBufferByte / sizeof(double);
35
36     cl_int ciErrNum = 0;
37
38     uint i = 0;
39
40     // clSetKernelArg is not thread-safe
41     enif_mutex_lock(pKernel_s->mtx);
42
43     for(i=0; i < pInputBufferC; i++)
44         ciErrNum |= clSetKernelArg(pKernel_s->kernel, i, sizeof(cl_mem), (void*)
45             pInputBufferV[i]);

```

```

44
45     i--;
46     ciErrNum |= clSetKernelArg(pKernel_s->kernel, i+1, sizeof(cl_mem), (void*)
47                               pOutputBuffer);
47     ciErrNum |= clSetKernelArg(pKernel_s->kernel, i+2, sizeof(unsigned int), (void*)
48                               &OFFSET); // Offset
48     ciErrNum |= clSetKernelArg(pKernel_s->kernel, i+3, sizeof(unsigned int), (void*)
49                               &uiNumElem);
50
50     enif_mutex_unlock(pKernel_s->mtx);
51
51
52     CHK_SUCCESS(ciErrNum);
53
54
55 #ifdef TIME
56     GET_TIME((fun_prologue_end));
57 #endif
58
59 // **** execute the kernel ****
60
61     size_t szGlobalWorkSize = uiNumElem;
62 //     size_t szLocalWorkSize = 1;
63
64 //     szGlobalWorkSize = roundUp((int)szLocalWorkSize, (int)(szGlobalWorkSize));
65
66 // esegui calcolo
67     CHK_SUCCESS(computeKernel(pKernel_s->kernel, &szGlobalWorkSize, NULL/*&
68             szLocalWorkSize*/)); // let OpenCL decide the local work-size
69
69 #ifdef TIME
70     GET_TIME((run_end));
71
72     fun_end = run_end;
73 #endif
74
75 #ifdef TIME
76     cerr <<
77         "prologue + kernel time: " << nsec2usec(diff(&(fun_start), &(fun_end))) <<
78         endl <<
79         "\tFunction prologue time: " << nsec2usec(diff(&(fun_start), &fun_prologue_end))
80             <<
81         endl <<
82         "\tKERNEL computation time: " << nsec2usec(diff(&(fun_prologue_end), &run_end))
83             <<
84         endl
83 ;
84 #endif

```

```

85
86     return ATOM(ok);
87
88 }
89
90 /* unary map adapter (in/out on device)
91 * It just interprets erlang NIF parameters and calls mapDD_Impl
92 */
93 static ERL_NIF_TERM mapDD(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
94
95     //mapDD( Kernel::kernel() , InputBuffer::deviceBuffer() , OutputBuffer::
96     //          deviceBuffer() )
97     OCL_INIT_CHECK()
98     NIF_ARITY_CHECK(3)
99
100    /*get the parameters (Kernel::kernel() , InputBuffer::deviceBuffer() ,
101       OutputBuffer::deviceBuffer())******/
102
103    kernel_sync* pKernel_s = NULL;
104    if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {
105        cerr << "ERROR :: mapDD: 1st parameter is not a kernel_sync" << endl ;
106        return enif_make_badarg(env);
107    }
108
109    cl_mem* pInputBuffer = NULL;
110    if (!enif_get_resource(env, argv[1], deviceBuffer_rt, (void**) &pInputBuffer) ||
111        *pInputBuffer == NULL) {
112        cerr << "ERROR :: mapDD: 2nd parameter is not a device buffer" << endl ;
113        return enif_make_badarg(env);
114    }
115
116    cl_mem* pOutputBuffer = NULL;
117    if (!enif_get_resource(env, argv[2], deviceBuffer_rt, (void**) &pOutputBuffer) ||
118        *pOutputBuffer == NULL) {
119        cerr << "ERROR :: mapDD: 3nd parameter is not a device buffer" << endl ;
120        return enif_make_badarg(env);
121    }
122
123    /****** */
124
125    uint pInputBufferC = 1;
126    cl_mem* pInputBufferV[1] = { pInputBuffer };
127
128    return mapDD_Impl(env, pKernel_s, (void**) pInputBufferV, pInputBufferC,
129                      pOutputBuffer);
130}
131
132

```

```

126
127
128 /*Binary map adapter (in/out on device)
129 *It just interprets erlang NIF parameters and calls mapDD_impl
130 */
131 static ERL_NIF_TERM map2DD(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
132
133     //mapDD( Kernel::kernel() , InputBuffer1::deviceBuffer() , InputBuffer2::
134     //          deviceBuffer() , OutputBuffer::deviceBuffer() )
135
136     OCL_INIT_CHECK()
137     NIF_ARITY_CHECK(4)
138
139     /*get the parameters (Kernel::kernel() , InputBuffer :: deviceBuffer() ,
140      OutputBuffer :: deviceBuffer())******/
141
142     kernel_sync* pKernel_s = NULL;
143     if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {
144         cerr << "ERROR :: map2DD: 1st parameter is not a kernel_sync" << endl ;
145         return enif_make_badarg(env);
146     }
147
148     cl_mem* pInputBuffer1 = NULL;
149     if (!enif_get_resource(env, argv[1], deviceBuffer_rt, (void**) &pInputBuffer1)
150         || *pInputBuffer1 == NULL) {
151         cerr << "ERROR :: map2DD: 2nd parameter is not a device buffer" << endl ;
152         return enif_make_badarg(env);
153     }
154
155     cl_mem* pInputBuffer2 = NULL;
156     if (!enif_get_resource(env, argv[2], deviceBuffer_rt, (void**) &pInputBuffer2)
157         || *pInputBuffer2 == NULL) {
158         cerr << "ERROR :: map2DD: 3nd parameter is not a device buffer" << endl ;
159         return enif_make_badarg(env);
160     }
161
162     /******
163     */
164     size_t szInputBuffer1Byte = 0, szInputBuffer2Byte = 0;
165
166     CHK_SUCCESS(getBufferSizeByte(*pInputBuffer1, &szInputBuffer1Byte);)

```

```

167     CHK_SUCCESS(getBufferSizeByte(*pInputBuffer2, &szInputBuffer2Byte);)
168
169 // Input Buffers must have the same size
170 if(szInputBuffer1Byte != szInputBuffer2Byte)
171     return make_error(env, ATOM(skel_ocl_buffers_different_size));
172
173 uint pInputBufferC = 2; //Map2
174 cl_mem * pInputBufferV[2] = { pInputBuffer1, pInputBuffer2 };
175
176 return mapDD_Impl(env, pKernel_s, (void**) pInputBufferV, pInputBufferC,
177                   pOutputBuffer);
178 }
179
180 //counters struct for map profiling
181 struct Map_fun_counters {
182
183     struct Map_fun_round_counters {
184         uint numInput;
185         timespec
186         start[2],
187         load_start[2],
188         *load_input_start[2],
189         *load_unmarshal_end[2],
190         *load_copyHD_end[2],
191         load_end[2],
192
193         run_start[2],
194         run_end[2],
195
196         unload_start,
197         unload_copy_start[2],
198         unload_copy_end[2],
199         unload_marshal_end[2],
200         unload_end,
201         end[2];
202
203
204     Map_fun_round_counters(uint _numInput) {
205
206         numInput = _numInput;
207
208         for(uint i = 0; i < 2; i++) {
209             load_input_start[i] = new timespec[numInput];
210             load_unmarshal_end[i] = new timespec[numInput];
211             load_copyHD_end[i] = new timespec[numInput];
212         }

```

```

213
214
215     };
216
217     ~Map_fun_round_counters() {
218
219         for(uint i = 0; i< 2; i++) {
220             delete [] load_input_start[i];
221             delete [] load_unmarshal_end[i];
222             delete [] load_copyHD_end[i];
223         }
224
225     };
226
227
228 }; //end Map_fun_round_counters
229
230 timespec
231 fun_start, fun_prologue, fun_end,
232 mapInput_start, *mapInput_end,
233 unmapInput_start, unmapInput_end,
234 mapOutput_start, mapOutput_end,
235 unmapOutput_end,
236
237 releaseInBuffersH_T, releaseOutBufferH_T,
238 releaseOutBufferD_T, releaseInBuffersD_T_start, releaseInBuffersD_T_end
239 ;
240
241 uint numRounds;
242
243 Map_fun_round_counters** round;
244
245 Map_fun_counters(const uint& _numRounds, const uint& numInput) {
246
247     mapInput_end = new timespec[numInput];
248
249     numRounds = _numRounds;
250
251     round = new Map_fun_round_counters*[numRounds];
252
253     for(uint i = 0; i< numRounds; i++)
254         round[i] = new Map_fun_round_counters(numInput);
255 }
256
257 ~Map_fun_counters() {
258
259

```

```

260     for( uint i = 0; i < numRounds; i++)
261         delete round[ i ];
262
263     delete [] round;
264
265     delete [] mapInput_end;
266
267 }
268
269 };
270
271 //Struct containing data shared among mapLL's threads
272 class MapLL_thread_params{
273
274     public:
275
276         bool errorSignal;
277         bool requestedExit;
278
279         ERL_NIF_TERM errorTerm;
280
281         cl_event mapOutput_evt;
282         cl_event unmapInput_evt;
283         cl_event lastMarshallng_evt;
284
285         cl_event** run_start_evt; //NUM_QUEUES for each round
286         cl_event** run_done_evt; //NUM_QUEUES for each round
287
288         ErlNifEnv* env;
289         ErlNifMutex* env_mtx;
290
291         uint pInputC;
292         cl_command_queue cmdQs[NUM_QUEUES];
293         uint cmdQsC;
294         cl_kernel* kernels;
295         cl_mem** inputBuffersD;
296         cl_mem outputBufferD;
297
298         //Shared InputH/outputH. set by load, used by unload in marshalling.
299         //Unload never writes a segment that Load hasn't copied yet on the device;
300         //ensured by stage sync: Load < run < unload
301         double* pMappedOutputBufferH;
302         cl_mem* inputBuffersH;
303         cl_mem outputBufferH;
304
305         ERL_NIF_TERM* currList;
306

```

```
307
308     Map_fun_counters* counters;
309
310     uint numQueues;
311     uint numRounds;
312
313     uint uiNumElem;
314     uint uiNumElemSegment;
315     size_t szSegment;
316
317     uint numSegments;
318
319     ERL_NIF_TERM* outputTerms;
320
321
322     explicit MapLL_thread_params(
323
324         cl_event _mapOutput_evt,
325         cl_event _unmapInput_evt,
326         cl_event _lastMarshallling_evt,
327         cl_event** _run_start_evt,
328         cl_event** _run_done_evt,
329
330         ErlNifEnv* _env,
331         ErlNifMutex* _env_mtx,
332         uint _pInputC,
333         cl_command_queue* _cmdQs,
334         cl_kernel* _kernels,
335         cl_mem** _inputBuffersD,
336         cl_mem _outputBufferD,
337
338         cl_mem* _inputBuffersH,
339         cl_mem _outputBufferH,
340
341         ERL_NIF_TERM* _currList,
342
343         Map_fun_counters* _counters,
344
345         uint _numRounds,
346         uint _numQueues,
347
348         uint _uiNumElem,
349         uint _uiNumElemSegment,
350         size_t _szSegment,
351
352         ERL_NIF_TERM* _outputTerms
353     )
```

```

354     {
355
356         errorSignal = false;
357         requestedExit = false;
358
359         errorTerm = 0;
360
361         mapOutput_evt = _mapOutput_evt;
362         unmapInput_evt = _unmapInput_evt;
363
364         lastMarshalling_evt = _lastMarshalling_evt;
365
366         run_start_evt = _run_start_evt;
367         run_done_evt = _run_done_evt;
368
369         env = _env;
370         env_mtx = _env_mtx;
371
372         pInputC = _pInputC;
373
374         for(int i = 0; i < NUM_QUEUES; i++)
375             cmdQs[ i ] = _cmdQs[ i ];
376
377
378         kernels = _kernels;
379         inputBuffersD = _inputBuffersD;
380         outputBufferD = _outputBufferD;
381
382
383         inputBuffersH = _inputBuffersH;
384         outputBufferH = _outputBufferH;
385
386         currList = _currList;
387
388         counters = _counters;
389
390         numRounds = _numRounds;
391         numQueues = _numQueues;
392
393         uiNumElem = _uiNumElem;
394         uiNumElemSegment = _uiNumElemSegment;
395         szSegment = _szSegment;
396
397         outputTerms = _outputTerms;
398
399         numSegments = uiNumElem/uiNumElemSegment;
400

```

```

401
402      }
403
404
405      void signalError() {
406
407 #ifdef DEBUG
408         cerr << "DEBUG: MapLL - signalError()\n";
409 #endif
410
411         errorSignal = true;
412     }
413
414     void requestExit() {
415         requestedExit = true;
416     }
417
418 };
419
420 void destroyMutexes(ErlNifMutex* mtxs[], uint len) {
421
422     for (uint i = 0; i < len; ++i) {
423         enif_mutex_destroy(mtxs[i]);
424         mtxs[i] = NULL;
425     }
426 }
427
428
429 // struct containing unmarshalling threads' shared data
430 struct MapLL_loadStage_unmarshalling_thread_params {
431
432     Barrier* barrier;
433     MapLL_thread_params* map_params;
434
435     uint* segmentOffset;
436     double** pMappedInputBuffersH;
437
438     uint* idx;
439
440     cl_command_queue curr_cmdQ;
441
442     bool* seqMode;
443
444     cl_event** round_start;
445
446     bool errorSignal;
447     bool* global_errorSignal;

```

```

448     ERL_NIF_TERM errorTerm;
449 };
450
451 static void* mapLL_loadStage_unmarshalling_Thread(void* obj) {
452
453
454 #ifdef DEBUG
455     char msg[128];
456 #endif
457
458     std::pair<uint, MapLL_loadStage_unmarshalling_thread_params*> *id_load_params =
459     (std::pair<uint, MapLL_loadStage_unmarshalling_thread_params*>*) obj;
460
461     uint& i_input = id_load_params->first; //thread's ID
462     MapLL_loadStage_unmarshalling_thread_params*& load_params = id_load_params->
463         second;
464
465 // local variables
466     MapLL_thread_params*& params = load_params->map_params;
467
468     Barrier*& barrier = load_params->barrier;
469
470     uint*& segmentOffset = (load_params->segmentOffset);
471
472     ErlNifEnv*& env = params->env;
473     ErlNifMutex*& env_mtx = params->env_mtx;
474     Map_fun_counters* &counters = params->counters;
475
476     cl_command_queue& curr_cmdQ = load_params->curr_cmdQ;
477     bool*& seqMode = (load_params->seqMode);
478     cl_event**& round_start_evts = (load_params->round_start);
479
480     uint*& idx = (load_params->idx);
481
482
483
484 #ifdef DEBUG
485     sprintf(msg, "DEBUG: Load[%u] - starting\n", i_input); cerr << msg;
486 #endif
487
488 // unmarshalling: L -> H
489 // copy data: H -> D
490     for (uint i_round = 0 ; i_round < params->numRounds; ++i_round) {
491
492
493         for (uint i_queue = 0 ; i_queue < params->numQueues; i_queue++) {

```

```

494
495 #ifdef DEBUG
496         sprintf(msg, "DEBUG: Load[%u] - awaiting signal [%u][%u] \n", i_input,
497         i_round, i_queue); cerr << msg;
498
499 //wait for signal from load master thread
500 clWaitForEvents(1, &(round_start_evt[i_round][i_queue]));
501
502 #ifdef DEBUG
503         sprintf(msg, "DEBUG: Load[%u] - starting [%u][%u] \n", i_input, i_round,
504         i_queue); cerr << msg;
505
506
507 //some other thread (unmarshalling or one of map's threads) has signaled
508 //an error, exit
509 if(load_params->errorSignal || *(load_params->global_errorSignal)) {
510
511 #ifdef DEBUG
512         sprintf(msg, "DEBUG: Load[%u] - Someone signaled an error, exit.\n",
513         i_input); cerr << msg;
514
515
516
517 #ifdef TIME
518     GET_TIME(( counters->round[i_round]->load_input_start[i_queue][i_input])) ;
519
520
521 //***** list -> H *****
522
523     double* pMappedInputBufferSegmentH = load_params->pMappedInputBuffersH[
524     i_input] + (*segmentOffset);
525
526     // copy current input list segment into the corresponding part of
527     // bufferH
528     uint numUnmarshElems=
529     list_to_double_arrayN(
530             params->env,
531             params->currList[i_input],
532             pMappedInputBufferSegmentH, params->
533             uiNumElemSegment,
534             &(params->currList[i_input])
535     );

```

```

533         //currList is updated to the first element of the next segment (of the
534         //list)
535         if (numUnmarshalElems != params->uiNumElemSegment) { //some problem
536             happened during unmarshalling
537             #ifdef DEBUG
538                 sprintf(msg, "DEBUG: Load[%u] - [%u][%u] - Error unmarshalling\n",
539                         i_input, i_round, i_queue); cerr << msg;
540             #endif
541             load_params->errorSignal = true;
542             load_params->errorTerm = sync_make_error(env, params->env_mtx, ATOM(
543                             unmarshalling_error));
544             return NULL;
545         }
546         //print_buffer_debug("inputBuffersH", inputBuffersH[i_input], oneIdx *
547         //params->szSegment, params->szSegment); // DEBUG
548     #ifdef DEBUG
549         sprintf(msg, "DEBUG: Load[%u] - [%u][%u] - L -> H\n", i_input, i_round,
550                 i_queue );
551         cerr << msg;
552     #endif
553     #ifdef TIME
554         GET_TIME((counters->round[i_round]->load_unmarshal_end[i_queue][i_input]))
555     ;
556     #endif
557
558     #ifdef SEQ
559         *seqMode = true;
560     #endif
561
562         cl_event copy_Evt;
563         //***** H -> D *****
564         THREAD.CHK.SUCCESS.CLEANUP.GOTO(
565             clEnqueueWriteBuffer(curr_cmdQ, params->inputBuffersD[i_input][*idx
566                         ], CL_FALSE, 0, params->szSegment,
567             (void*)pMappedInputBufferSegmentH, 0, NULL, *seqMode ? &copy_Evt :
568                         NULL
569             ) ;
570         load_params->errorTerm ,
571     {

```

```

570             sprintf(msg, "DEBUG: Load[%u] - [%u][%u] - ERROR H -> D \n",
571                         i_input, i_round, i_queue); cerr << msg;
572         load_params->errorSignal = true;
573     }
574     clFlush(params->cmdQs[i_queue]);
575
576
577 #ifdef DEBUG
578     sprintf(msg, "DEBUG: Load[%u] - [%u][%u] - H -> D\n", i_input, i_round,
579                         i_queue);
580     cerr << msg;
581 #endif
582 #ifdef SEQ
583     clWaitForEvents(1, &copy_Evt);
584     clReleaseEvent(copy_Evt);
585 #endif
586
587 #ifdef TIME
588     GET_TIME(( counters->round[i_round]->load_copyHD_end[i_queue][i_input]));
589 #endif
590
591 #ifdef DEBUG
592     sprintf(msg, "DEBUG: Load[%u] - WAITING at barrier [%u][%u]\n", i_input,
593                         i_queue, i_round); cerr << msg;
594 //synch with others unmarshalling threads
595 barrier->await();
596
597 #ifdef DEBUG
598     sprintf(msg, "DEBUG: Load[%u] - CROSSED barrier [%u][%u]\n", i_input,
599                         i_queue, i_round); cerr << msg;
600 #endif
601     }
602 }
603
604 #ifdef DEBUG
605     sprintf(msg, "DEBUG: Load[%u] - DONE\n", i_input); cerr << msg;
606 #endif
607
608 cleanup:
609
610 //in case of error must clear the skipped barrier, so to avoid deadlocking other
611 //load threads
612 if(load_params->errorSignal || *(load_params->global_errorSignal)) {

```

```

612 #ifdef DEBUG
613     sprintf(msg, "DEBUG: Load[%u] - CLEARING skipped barrier\n", i_input); cerr
614     << msg;
615     barrier->await();
616
617 }
618
619 return NULL;
620
621 }
622
623
624
625 static void* mapLL_loadStage_Thread(void* obj) {
626
627 #ifdef DEBUG
628     char msg[256];
629 #endif
630
631 #ifdef DEBUG
632     cerr << "DEBUG: Load - Starting.\n";
633 #endif
634
635
636     MapLL_thread_params* params = (MapLL_thread_params *) obj;
637
638 //parametri
639
640 //env, env_mtx needed by CLEANUP macros
641 ErlNifEnv* env = params->env;
642 ErlNifMutex* env_mtx = params->env_mtx;
643
644 Map_fun_counters*& counters = params->counters;
645
646 //local variables
647 size_t szInput_local = (params->szSegment / params->uiNumElemSegment) * params->
648 uiNumElem;
649
650
651 double* pMappedInputBuffersH [params->pInputC]; //pointers mapped into input
652     buffers
653
654 bool inputBufferIsMapped = false;
655
656 bool seqMode = false; //is sequential execution activated?
657

```

```

656     // initialize synchronization barrier for coordinating #pInputC load threads +
657     // master load thread
658     Barrier round_end_barrier(params->pInputC + 1);
659
660     //check if load stage is configured (in mapLL)
661     bool unload_stage_isPresent = params->outputTerms != NULL ? true : false;
662
663
664 #ifdef TIME
665     GET_TIME(counters->mapInput_start);
666 #endif
667
668     //map InputBuffer
669     for(uint i_list = 0; i_list < params->pInputC; ++i_list) {
670
671         pMappedInputBuffersH[ i_list ] = (cl_double*)
672             clEnqueueMapBuffer(params->cmdQs[0], params->inputBuffersH[ i_list ], CLTRUE,
673                             CL_MAP_WRITE, 0, szInput_local, 0, NULL, &ciErrNum);
674
675         THREAD_CHK_SUCCESS_CLEANUP_GOTO(
676             ciErrNum, ,
677             params->errorTerm,
678             {
679                 cerr << "DEBUG: Load - clEnqueueMapBuffer\n";
680                 params->signalError();
681             }
682         )
683 #ifdef TIME
684         GET_TIME(counters->mapInput_end[ i_list ]);
685 #endif
686
687         inputBufferIsMapped = true;
688
689         //set output buffer's pointer to mapped buffer memory, since inputH[0] is
690         //outputH
691         params->pMappedOutputBufferH = pMappedInputBuffersH[0];
692
693         //get OpenCL context needed by clCreateUserEvent
694         cl_context context;
695         clGetCommandQueueInfo(params->cmdQs[0], CL_QUEUE_CONTEXT, sizeof(cl_context), &
696                               context, NULL);
697
698         cl_event* round_start_evts[params->numRounds];
699
700         for (uint i_round = 0 ; i_round < params->numRounds; i_round++) {

```

```

699
700     round_start_evts[i_round] = new cl_event [params->numQueues];
701
702     for( uint i_queue = 0 ; i_queue < params->numQueues; i_queue++)
703         round_start_evts[i_round][i_queue] = clCreateUserEvent(context , NULL);
704 }
705
706
707
708 MapLL_loadStage_unmarshalling_thread_params unmarsh_params;
709
710 // initialize params' loop-independent variables
711 unmarsh_params.barrier = &round_end_barrier;
712 unmarsh_params.map_params = params;
713 unmarsh_params.pMappedInputBuffersH = pMappedInputBuffersH;
714
715 unmarsh_params.seqMode = &seqMode;
716 unmarsh_params.round_start = round_start_evts;
717
718 unmarsh_params.errorSignal = false; //unmarshalling threads signal (local)
719 unmarsh_params.global_errorSignal = &params->errorSignal; //map threads signal (
720     global)
721
722 //create unmarshalling threads
723 ErlNifTid unmarshalling_threads[params->pInputC];
724
725 std :: pair<uint , MapLL_loadStage_unmarshalling_thread_params*> *id_params[params
726     ->pInputC]; // (Thread-ID, thread params)
727
728 for( uint i_input = 0; i_input < params->pInputC; i_input++) {
729
730     id_params[i_input] = new std :: pair<uint ,
731         MapLL_loadStage_unmarshalling_thread_params*>(i_input , &unmarsh_params);
732
733     enif_thread_create((char*)"Unmarshalling" , &unmarshalling_threads[i_input] ,
734         mapLL_loadStage_unmarshalling_Thread , (void*) id_params[i_input] , NULL);
735 }
736
737
738
739 //*****LOAD LOOP*****
740 uint zeroIdx;// first segment index (the one that will be computed by the first
741     queue)
742 uint i_round;
743 for (i_round = 0 ; i_round < params->numRounds; ++i_round) {
744
745     zeroIdx = (params->numQueues * i_round);
746

```

```

741         uint idx;
742         uint segmentOffset;
743
744
745 #ifdef DEBUG
746         sprintf(msg, "DEBUG: Load - Starting Round[%d]\n", i_round);
747         cerr << msg;
748 #endif
749
750         for (uint i_queue = 0 ; i_queue < params->numQueues; i_queue++) {
751
752 #ifdef TIME
753             GET_TIME((params->counters->round[i_round]->start[i_queue]));
754
755             params->counters->round[i_round]->load_start[i_queue] =
756             params->counters->round[i_round]->start[i_queue];
757
758 #endif
759
760             idx = zeroIdx + i_queue;
761
762             segmentOffset = (idx * params->uiNumElemSegment);
763
764             cl_command_queue& curr_cmdQ = params->cmdQs[i_queue];
765
766             cl_event& curr_run_start_evt = params->run_start_evt[i_queue][i_round];
767
768             if(i_round > 0) { //from the first round onwards
769                 uint round_to_wait = i_round - 1;
770                 cl_event& curr_run_done_evt = params->run_done_evt[i_queue][
771                     round_to_wait];
772
773 #ifdef DEBUG
774                 sprintf(msg, "DEBUG: Load - %d - Waiting for unload[%d][%d] to start
775 .\n", i_queue, i_queue, round_to_wait);
776                 cerr << msg;
777
778                 //WAIT for RUN stage, so to synch with unload stage
779                 clWaitForEvents(1, &(curr_run_done_evt));
780             }
781
782             if(params->errorSignal) { //someone has signaled an error, cleanup and
783                 exit
784 #ifdef DEBUG

```

```

785         sprintf(msg, "DEBUG: Load - %d - Someone has signaled an error ,  

786             cleanup and exit\n", i_queue);  

787 #endif  

788         goto cleanup;  

789     }  

790  

791     /*  

792      * list -> H  

793      * H -> D  

794      */  

795  

796     //update unmarshalling threads' loop-dependent variables  

797     unmarshal_params.curr_cmdQ = curr_cmdQ;  

798     unmarshal_params.idx = &idx;  

799     unmarshal_params.segmentOffset = &segmentOffset;  

800  

801  

802 #ifdef DEBUG  

803     sprintf(msg, "DEBUG: Load - signaling unmarshalling threads [%d][%d]\n",  

804             i_round, i_queue); cerr << msg;  

805 #endif  

806     //signal unmarshalling threads  

807     clSetUserEventStatus(round_start_evt[i_round][i_queue], CL_SUCCESS);  

808  

809 #ifdef DEBUG  

810     sprintf(msg, "DEBUG: Load - WAITING at barrier [%d][%d]\n", i_round,  

811             i_queue); cerr << msg;  

812 #endif  

813     //await unmarshalling completion  

814     round_end_barrier.await();  

815 #ifdef DEBUG  

816     sprintf(msg, "DEBUG: Load - barrier CROSSED [%d][%d]\n", i_round, i_queue)  

817 #endif  

818  

819     //check if everything went fine, otherwise set error term (i.e.  

820         propagate the error) and cleanup  

821     if(unmarshal_params.errorSignal){  

822 #ifdef DEBUG  

823         sprintf(msg, "DEBUG: Load - some unmarshalling thread signaled an  

824             error , cleanup.\n"); cerr << msg;  

825 #endif

```

```

826             params->errorTerm = unmarshal_params.errorTerm;
827             params->signalError();
828
829             goto cleanup;
830         }
831
832         if (params->errorTerm) {
833
834             goto cleanup;
835         }
836
837         //*****Segment is now loaded on device*****
838
839 #ifdef TIME
840         GET_TIME(( counters->round[i_round]->load_end[i_queue]));
841 #endif
842
843         //signal run stage, so it can process the segment i just loaded
844         clSetUserEventStatus(curr_run_start_evt, CL_SUCCESS);
845
846
847     } //queue loop end
848
849 } //round loop end
850
851 #ifdef DEBUG
852     sprintf(msg, "DEBUG: Load - WAITING for unload, must UnmapInput\n");
853     cerr << msg;
854
855
856     if (unload_stage_isPresent) {
857         //overlap unmapping with last marshalling (which is done by unload thread)
858         clWaitForEvents(1, &params->unmapInput_evt);
859     }
860
861
862 #ifdef TIME
863     GET_TIME(( counters->unmapInput_start));
864 #endif
865
866     if (inputBufferIsMapped) {
867
868 #ifdef DEBUG
869         sprintf(msg, "DEBUG: Load - UNMAPPING Input\n");
870         cerr << msg;
871

```

```

872         //unmap input buffers
873         for( uint i_list = 1/*0*/; i_list < params->pInputC; ++i_list)
874             // Starts from 1; 0 is released by unload since it's used also as output
875             clEnqueueUnmapMemObject(params->cmdQs[0], params->inputBuffersH[i_list],
876                                     pMappedInputBuffersH[i_list], 0, NULL, NULL);
877
878 #ifdef TIME
879     GET_TIME((counters->unmapInput_end));
880 #endif
881     }
882     pBufferIsMapped = false;
883
884
885 #ifdef DEBUG
886     sprintf(msg, "DEBUG: Load - CLEANING UP\n"); cerr << msg;
887 #endif
888
889
890
891
892 cleanup:
893
894     //in case of error
895     if(params->errorSignal || unmarsh_params.errorSignal) { //release every lock
896         before exiting to avoid deadlocks
897
898 #ifdef DEBUG
899     sprintf(msg, "DEBUG: Load - Error detected, set all events\n"); cerr << msg;
900 #endif
901
902         for (uint i_round = 0 ; i_round < params->numRounds; ++i_round)
903             for (uint i_queue = 0 ; i_queue < params->numQueues; ++i_queue) {
904
905                 clSetUserEventStatus(params->run_start_evt[i_queue][i_round],
906                                     CL_SUCCESS);
907
908                 clSetUserEventStatus(round_start_evts[i_round][i_queue], CL_SUCCESS)
909                     ;
910
911             }
912
913 #ifdef DEBUG
914     sprintf(msg, "DEBUG: Load - joining unmarshalling threads\n"); cerr << msg;
915 #endif
916     //join unmarshalling threads

```



```

957         delete [] round_start_evts[i_round];
958     }
959
960
961
962 #ifdef DEBUG
963     sprintf(msg, "DEBUG: Load - DONE\n");
964     cerr << msg;
965 #endif
966
967     return NULL;
968 }
969
970
971
972
973 static void* mapLL_runStage_Thread(void* obj) {
974 #ifdef DEBUG
975     char msg[256];
976 #endif
977
978     MapLL_thread_params* params = (MapLL_thread_params *) obj;
979
980     //local
981     // bool seqMode = false;
982
983     //needed by CLEANUP macros
984     ErlNifEnv*& env = params->env;
985     ErlNifMutex*& env_mtx = params->env_mtx;
986
987     Map_fun_counters*& counters = params->counters;
988
989     size_t szGlobalWorkSize = params->uiNumElemSegment;
990
991     bool isLastStage = params->outputTerms == NULL ? true : false; //Am I the last
992             stage of the pipeline?
993
994     uint zeroIdx = 0;
995     for (uint i_round = 0 ; i_round < params->numRounds; ++i_round) {
996
997         zeroIdx = (params->numQueues * i_round);
998
999 #ifdef DEBUG
1000     sprintf(msg, "DEBUG: Run - Starting round[%d]\n", i_round);
1001     cerr << msg;
1002 #endif

```

```

1003
1004     uint idx = 0;
1005     for (uint i_queue = 0 ; i_queue < params->numQueues; ++i_queue) {
1006
1007
1008 #ifdef DEBUG
1009         sprintf(msg, "DEBUG: Run - %d - waiting to start\n", i_queue);
1010         cerr << msg;
1011 #endif
1012
1013     idx = zeroIdx + i_queue;
1014
1015     //current loop values
1016     cl_kernel& curr_kernel = params->kernels[i_queue];
1017     cl_command_queue& curr_cmdQ = params->cmdQs[i_queue];
1018
1019     cl_event& curr_run_start_evt = params->run_start_evt[i_queue][i_round];
1020     cl_event& curr_run_done_evt = params->run_done_evt[i_queue][i_round];
1021
1022     uint segmentOffset = idx * params->uiNumElemSegment;
1023
1024
1025     //wait for load stage
1026     clWaitForEvents(1, &(curr_run_start_evt));
1027
1028
1029     if(params->errorSignal) { //Someone signaled an error, exit.
1030 #ifdef DEBUG
1031         sprintf(msg, "DEBUG: Run - %d - Someone signaled an error, exit.\n",
1032                 i_queue);
1033         cerr << msg;
1034     #endif
1035     goto cleanup;
1036
1037
1038     //*****set map kernels parameters*****
1039
1040
1041     uint i_input = 0;
1042
1043     uint numKerParams = params->pInputC + 3;
1044
1045     cl_int ciErrNum [numKerParams];
1046
1047
1048

```

```

1049         for( i_input=0; i_input < params->pInputC; i_input++)
1050             ciErrNum[ i_input ] = clSetKernelArg( curr_kernel , i_input , sizeof(
1051                                         cl_mem) , (void*) &(params->inputBuffersD[ i_input ][ idx ]) );
1052             i_input--;
1053
1054             ciErrNum[ i_input+1 ] = clSetKernelArg( curr_kernel , i_input+1, sizeof(
1055                                         cl_mem) , (void*) &(params->outputBufferD));//round independent
1056             ciErrNum[ i_input+2 ] = clSetKernelArg( curr_kernel , i_input+2, sizeof(
1057                                         unsigned int) , (void*) &(segmentOffset));
1058
1059
1060             for( uint i = 0; i < numKerParams; i++) {
1061                 if( ciErrNum[ i ] != CL_SUCCESS) {
1062 #ifdef DEBUG
1063                     sprintf(msg, "DEBUG: Run - %d - Error SetKernel #%d\n", i_queue ,
1064                                         i); cerr << msg;
1065
1066                     params->signalError();
1067                     params->errorTerm = sync_make_error_cl(env , env_mtx , ciErrNum[ i
1068                                         ]);
1069
1070                     goto cleanup;
1071
1072
1073 #ifdef DEBUG
1074                     sprintf(msg, "DEBUG: Run - %d - Kernel Set\n", i_queue );
1075                     cerr << msg;
1076 #endif
1077
1078
1079 #ifdef TIME
1080             GET_TIME(( counters->round[ i_round ]->run_start[ i_queue ]));
1081 #endif
1082
1083
1084 #ifdef SEQ
1085             seqMode = true;
1086 #endif
1087
1088             //***** Run *****

```

```

1089         cl_event exec_Evt;
1090         THREAD_CHK_SUCCESS_CLEANUP_GOTO(
1091             clEnqueueNDRangeKernel(curr_cmdQ, curr_kernel, 1, NULL, &
1092             szGlobalWorkSize, NULL/*&szLocalWorkSize*/ , 0, NULL, &exec_Evt); ,
1093             params->errorTerm ,
1094             {
1095                 sprintf(msg, "DEBUG: Run - %d - Error NDRangeKernel\n", i_queue); cerr
1096                 << msg;
1097             }
1098             clFlush(curr_cmdQ);
1099
1100
1101 #ifdef DEBUG
1102         sprintf(msg, "DEBUG: Run - %d - run\n", i_queue);
1103         cerr << msg;
1104 #endif
1105
1106
1107         clWaitForEvents(1, &exec_Evt);
1108         clReleaseEvent(exec_Evt);
1109
1110
1111
1112 #ifdef TIME
1113         GET_TIME(( counters->round[i_round]->run_end[i_queue]) );
1114         if(isLastStage)
1115             counters->round[i_round]->end[i_queue] = counters->round[i_round]->
1116             run_end[i_queue];
1117 #endif
1118         //    clWaitForEvents(1, &exec1_Evt);
1119
1120
1121
1122 #ifdef DEBUG
1123         sprintf(msg, "DEBUG: Run - %d - wake up unload\n", i_queue);
1124         cerr << msg;
1125 #endif
1126         //signal unload stage
1127         clSetUserEventStatus(curr_run_done_evt, CL_SUCCESS);
1128
1129     } //end queue
1130
1131 } //end round
1132

```

```

1133 cleanup:
1134
1135     //in case of error
1136     if(params->errorSignal)  { //to avoid deadlocks, release every lock before
1137         exiting
1138 #ifdef DEBUG
1139         sprintf(msg, "DEBUG: Run - Error detected, set all events\n");
1140         cerr << msg;
1141 #endif
1142
1143     for ( uint i_round = 0 ; i_round < params->numRounds; ++i_round)
1144         for ( uint i_queue = 0 ; i_queue < params->numQueues; ++i_queue)
1145             clSetUserEventStatus(params->run_done_evt [ i_queue ] [ i_round ] ,
1146                                 CL_SUCCESS);
1147     }
1148 #ifdef DEBUG
1149     sprintf(msg, "DEBUG: Run - Done\n");
1150     cerr << msg;
1151 #endif
1152
1153     return NULL;
1154 }
1155
1156
1157 static ERL_NIF_TERM mapLD_Impl_3Thread(ErlNifEnv * env, kernel_sync* pKernel_s, void
1158 ** _pInputV, uint pInputC, void* _pOutputBuffer, uint listLength) {
1159 #ifdef DEBUG
1160     cerr << "DEBUG: M - mapLD_Impl_3Thread" << endl;
1161 #endif
1162
1163     ERL_NIF_TERM toReturn;
1164
1165
1166     const uint NUMSEGMENTS = NUMSEGM;
1167
1168     const uint NUMQUEUESLOCAL = NUMQUEUES;
1169
1170     const uint numRounds = NUMSEGMENTS/NUMQUEUESLOCAL;
1171
1172
1173 #if defined(TIME) || defined(TIME_FUN)
1174
1175 #define DIFF_T0(COUNTER) nsec2usec( diff( &(counters->fun_start) , &(COUNTER) ) )
1176

```

```

1177     Map_fun_counters* counters = new Map_fun_counters(numRounds, pInputC);
1178 #endif
1179
1180 #ifdef TIME
1181
1182     timespec createEventsT_start, createEventsT_end,
1183     createInputBuffersH [pInputC],
1184     createInputBuffersD [pInputC],
1185
1186     createOutputBufferD;
1187
1188     timespec joinT [2], //3
1189     //list_createdT_start, list_createdT_end,
1190     releaseEventsT;
1191
1192     // clock_getres(CLOCK_MONOTONIC, &(counters->fun_start));
1193     // cerr << "INFO : Time resolution is " << counters->fun_start.tv_nsec <<
1194     //nsecs" << endl;
1195
1196 #endif
1197 #if defined(TIME) || defined(TIME_FUN)
1198     GET_TIME(counters->fun_start);
1199 #endif
1200
1201
1202     std :: vector<cl_event> events;
1203
1204     //cast input/output to type relevant here
1205     cl_mem outputBufferD = *((cl_mem*) _pOutputBuffer);
1206
1207     //pInputV is an array of ERL_NIF_TERM* representing, each one, an erlang list
1208     ERL_NIF_TERM** inputLists = (ERL_NIF_TERM**) _pInputV;
1209
1210     //input lists must have at least minListLen elements,
1211     //otherwise they can't be split in NUMSEGMENTS segments
1212     uint listLen = listLength;
1213
1214
1215     uint minListLen = device_minBaseAddrAlignByte / sizeof(double);
1216     if(listLen/NUMSEGMENTS < minListLen)
1217         return make_error(env, ATOM(skel_ocl_input_list_too_short));
1218
1219
1220     uint uiNumElem = listLen;
1221
1222     uint uiNumElemSegment = uiNumElem / NUMSEGMENTS;

```

```

1223
1224     size_t szInput = uiNumElem * sizeof(double);
1225
1226     size_t szSegment = szInput / NUMSEGMENTS;
1227
1228
1229 #ifdef DEBUG
1230     cerr << "DEBUG: M - mapLD_impl:" << "#Elem: " << uiNumElem << " #ElemSegment: " <<
1231         uiNumElemSegment << " szInput: " << szInput << " szSegment: " << szSegment <<
1232         endl;
1233 #endif
1234
1235     cl_int ciErrNum = 0;
1236
1237
1238     cl_context context = getContext();
1239
1240     //mutex protecting env from concurrent modifications (enif_make_xxx)
1241     ErlNifMutex* env_mtx = enif_mutex_create((char*)"env_mtx");
1242
1243
1244 #ifdef TIME
1245     GET_TIME(createEventsT_start);
1246 #endif
1247
1248     //*****events setup [NUMQUEUES] [numRounds])
1249     cl_event* run_start_evt [NUMQUEUES_LOCAL];
1250     cl_event* run_done_evt [NUMQUEUES_LOCAL];
1251
1252     for (int i = 0; i < NUMQUEUES_LOCAL; ++i) {
1253         run_start_evt [i] = new cl_event [numRounds];
1254         run_done_evt [i] = new cl_event [numRounds];
1255
1256         for (int i_round = 0; i_round < numRounds; ++i_round) {
1257             run_start_evt [i][i_round] = clCreateUserEvent(context, &ciErrNum);
1258             run_done_evt [i][i_round] = clCreateUserEvent(context, &ciErrNum);
1259
1260             events.push_back(run_start_evt [i][i_round]);
1261             events.push_back(run_done_evt [i][i_round]);
1262         }
1263     }
1264
1265 #ifdef TIME
1266     GET_TIME(createEventsT_end);
1267 #endif

```

```

1268
1269 //*****threads setup
1270 //Load, run. Each stage works on 2 command queues
1271 const uint NUMSTAGES = 2;
1272 ErlNifTid tid [NUMSTAGES];
1273
1274 ErlNifTid& loadStage = tid [0];
1275 ErlNifTid& runStage = tid [1];
1276
1277 //object holding everything needed by the threads
1278 MapLL_thread_params* conf = NULL;
1279
1280
1281 //*****for each input list , create:
1282 // one inputBufferH
1283 // NUMSEGMENTS inputBufferD
1284
1285 cl_mem inputBuffersH [pInputC];
1286
1287 // [pInputC] [NUMSEGMENTS]
1288 cl_mem* inputBuffersD [pInputC];
1289 for (int i = 0; i < pInputC; ++i)
1290     inputBuffersD [i] = new cl_mem [NUMSEGMENTS];
1291
1292 for (uint i_list = 0; i_list < pInputC; ++i_list) {
1293     // create one input host buffer of size szInput , released by Unload
1294     CHK_SUCCESS_CLEANUP_GOTO(
1295         createBuffer (szInput , CL_MEM_READ_WRITE |
1296                     CL_MEM_ALLOC_HOST_PTR, &(inputBuffersH [i_list]))
1297         );
1298 }
1299 #ifdef TIME
1300     GET_TIME(createInputBuffersH [i_list]);
1301 #endif
1302
1303
1304 #ifdef DEBUG
1305     cerr << "DEBUG: M - " << "inputBuffersH [" << i_list << "] CREATED: size: "
1306     << szInput << endl;
1307
1308
1309 //*****create NUMSEGMENTS input device buffers having size szSegment
1310 //*****release by Unload

```

```

1311     for (int i_segm = 0; i_segm < NUMSEGMENTS; ++i_segm) {
1312
1313         CHK_SUCCESS.CLEANUP.GOTO(
1314             createBuffer( szSegment , CL_MEM_READ_ONLY, &(
1315                 inputBuffersD [ i_list ][ i_segm ]) ) , ,
1316             ;
1317             )
1318
1319 #ifdef TIME
1320         GET_TIME( createInputBuffersD [ i_list ] );
1321 #endif
1322
1323 #ifdef DEBUG
1324         cerr << "DEBUG: M - " << "inputBuffersD ["<<i_list <<" ][ " <<i_segm<<" ]
1325             CREATED: size: " << szSegment<< endl;
1326 #endif
1327 }
1328 #ifdef DEBUG
1329     cerr << "DEBUG: M - " << "inputBuffersD CREATED"<< endl;
1330 #endif
1331 }
1332
1333
1334 #ifdef TIME
1335     GET_TIME( createOutputBufferD );
1336 #endif
1337
1338
1339 //Clone kernel , 'cause i need two of them, one per queue
1340 cl_kernel kernels[NUM_QUEUES_LOCAL];
1341
1342 kernels [0] = pKernel_s->kernel;
1343
1344 if (NUM_QUEUES_LOCAL == TWO)
1345     cloneKernel( pKernel_s->kernel , &kernels [1]);
1346
1347 if (NUM_QUEUES_LOCAL > MAX_QUEUES) {
1348     toReturn = make_error( env , ATOM( ocl_num_queue_not_supported));
1349     goto cleanup;
1350 }
1351
1352
1353
1354 //*****Main LOOP*****
1355

```

```

1356
1357     ERL_NIF_TERM currList [pInputC]; //iterator on the list , points to the head of
1358     // the current segment of the list
1359     for (uint i = 0; i < pInputC; ++i)
1360         currList [i] = *(inputLists [i]);
1360
1361
1362     conf = new MapLL_thread_params (
1363             NULL, //mapOutput_evt ,
1364             NULL, //unmapInput_evt ,
1365             NULL, //lastMarshallng_evt ,
1366             run_start_evt ,
1367             run_done_evt ,
1368             env, env_mtx,
1369             pInputC, cmdQs, kernels ,
1370             inputBuffersD , outputBufferD ,
1371             inputBuffersH ,
1372             NULL, //outputBufferH ,
1373             currList ,
1374 #ifdef TIME
1375             counters ,
1376 #else
1377             NULL,
1378 #endif
1379             numRounds ,
1380             NUM_QUEUES_LOCAL,
1381             uiNumElem , uiNumElemSegment , szSegment ,
1382             NULL//outputTerms
1383         );
1384
1385
1386 #ifdef DEBUG
1387     cerr << "DEBUG: M - Starting threads." << endl;
1388 #endif
1389
1390     enif_thread_create ((char*) "mapLL_loadStage" , &loadStage , mapLL_loadStage_Thread
1391             , conf , NULL);
1391
1392     enif_thread_create ((char*) "mapLL_runStage" , &runStage , mapLL_runStage_Thread ,
1393             conf , NULL);
1393
1394
1395 #ifdef TIME
1396     GET_TIME((counters->fun_prologue));
1397 #endif
1398
1399

```

```

1400     if(conf->errorSignal) {
1401         //something's gone wrong, cleanup and return the error term (the last one
1402         generated)
1402 #ifdef DEBUG
1403         cerr << "DEBUG: M - Got an error: cleanup and exit" << endl;
1404 #endif
1405
1406         toReturn = conf->errorTerm;
1407     }
1408
1409     //wait for termination
1410     for(uint i = 0; i < NUMSTAGES; ++i) {
1411
1412         enif_thread_join(tid[i], NULL);
1413
1414 #ifdef TIME
1415         GET_TIME((joinT[i]));
1416 #endif
1417
1418 #ifdef DEBUG
1419         char msg[64];
1420         sprintf(msg, "DEBUG: M - Thread %d joined\n", i);
1421         cerr << msg;
1422 #endif
1423     }
1424
1425     //check again after joining, shouldn't be useful since list creation starts
1426     //after the last marshalling, just before outputBuffer's unmapping,
1427     //at the very end of the computation when nobody can signal an error anymore.
1427     if(conf->errorSignal) {
1428         //something's gone wrong, cleanup and return the error term (the last one
1429         generated)
1429 #ifdef DEBUG
1430         cerr << "DEBUG: M - Got an error: cleanup and exit" << endl;
1431 #endif
1432         toReturn = conf->errorTerm;
1433     }
1434
1435     //*****CLEANUP label*****
1436 cleanup:
1437
1438
1439     //*****EVENTS*****
1440     if(!events.empty())
1441         releaseEvents(events);
1442
1443     for(uint i = 0; i < NUMQUEUES_LOCAL; i++){

```

```

1444         delete [] run_start_evt[i];
1445         delete [] run_done_evt[i];
1446     }
1447 #ifdef TIME
1448     GET_TIME((releaseEventsT));
1449 #endif
1450
1451 // **** KERNELS ****
1452 if (NUM_QUEUES.LOCAL == 2){
1453     if (kernels[1])
1454         clReleaseKernel(kernels[1]); //2 usecs
1455     kernels[1] = NULL;
1456 }
1457
1458 // **** ENV MUTEX ****
1459 if (env_mtx)
1460     destroyMutexes(&env_mtx, 1); //1 usec
1461
1462 // **** CONF ****
1463 if (conf)
1464     delete conf;
1465
1466
1467 #ifdef DEBUG
1468     cerr << "DEBUG: M - Cleanup done. Return" << endl;
1469 #endif
1470
1471
1472 #if defined(TIME) || defined(TIME_FUN)
1473     GET_TIME((counters->fun_end));
1474 #endif
1475
1476
1477
1478 // **** RUN Statistics ****
1479
1480
1481 #if defined(TIME) || defined(TIME_FUN)
1482
1483     cerr << endl <<
1484     "----- map" << pInputC << "LD: (usec)
1485     -----" <<
1486     " Processing " << pInputC << " X " << uiNumElem << " elements in " << NUM_SEGMENTS
1487     << " segments (" << uiNumElemSegment << " per segment)." <<
1488     endl <<
1489     " Total run time: " << DIFF_T0(counters->fun_end) <<

```

```

1489     endl;
1490
1491 #endif
1492
1493 #ifdef TIME
1494
1495     long prologueTime = DIFF_T0(counters-> fun_prologue);
1496
1497
1498     cerr << endl <<
1499     "Function prologue time: " << prologueTime << "\t\t\t\t\tT: " << prologueTime <<
1500     endl;
1501
1502     char str[128];
1503     sprintf(str, "\tcreateEvents: %lu \t\t\t\t\tT: %lu\n", DIFF(createEventsT_start
1504             , createEventsT_end), DIFF_T0( createEventsT_end)); cerr << str;
1505
1506     for (uint i_input = 0; i_input < pInputC; i_input++) {
1507         sprintf(str, "\tcreateInputBuffersH[%d]: \t\t\t\t\tT: %lu\n", i_input, DIFF_T0
1508                 ( createInputBuffersH[i_input])); cerr << str;
1509
1510         sprintf(str,
1511                 "Load - mapInputBuffers[%d]: %lu \t\t\t\t\tT: %lu\n", i_input,
1512                 i_input == 0 ?
1513                 DIFF( counters->mapInput_start , counters->mapInput_end[i_input] ) :
1514                 DIFF( counters->mapInput_end[i_input - 1], counters->mapInput_end[
1515                     i_input ] ),
1516                 DIFF_T0( counters->mapInput_end[i_input] )
1517             );
1518         cerr << str;
1519
1520         uint kernelTotal = 0, inputTotal = 0, outputTotal = 0, marshTotal = 0,
1521             unmarshalTotal = 0;
1522         uint i_round;
1523         //round statistics
1524         for (i_round = 0; i_round < numRounds; i_round++) {
1525             uint tKernel[2] = {0,0};
1526             uint tInput = 0, tOutput = 0;
1527             uint tStartRound[2];
1528             cerr <<

```

```

1529     "\nRound [ " << i_round << " ] total time: " << DIFF(( counters->round[i_round
1530         ]->start[0]), ( counters->round[i_round]->end[NUM_QUEUES_LOCAL-1])) <<
1531         endl;
1532
1533     for (int i_queue = 0; i_queue < NUM_QUEUES_LOCAL; i_queue++) {
1534
1535         inputTotal +=
1536             tInput = DIFF(( counters->round[i_round]->load_start[i_queue]), (
1537                 counters->round[i_round]->load_end[i_queue]));
1538
1539         tStartRound[i_queue] = DIFF_T0(counters->round[i_round]-> start[i_queue
1540             ]);
1541         cerr <<
1542             "\n\t" << i_queue << " - R[" << i_round << "] Start round\t\t\t\t\tT: " <<
1543                 tStartRound[i_queue] <<
1544                     endl <<
1545                     "\t" << i_queue << " - R[" << i_round << "] Load total time: " << tInput <<
1546                     endl;
1547
1548     for (uint i_input = 0; i_input < pInputC; i_input++) {
1549         long unmarshal = 0, t_Unmarshal;
1550         long cpDH_0 = 0, t_cpDH_0;
1551
1552         unmarshalTotal +=
1553             unmarshal = DIFF((counters->round[i_round]->load_input_start[i_queue][
1554                 i_input]), (counters->round[i_round]->load_unmarshal_end[i_queue][
1555                     i_input]));
1556
1557         t_Unmarshal = DIFF_T0(counters->round[i_round]-> load_unmarshal_end[
1558             i_queue][i_input]);
1559
1560         cpDH_0 = DIFF((counters->round[i_round]->load_unmarshal_end[i_queue][
1561             i_input]), (counters->round[i_round]->load_copyHD_end[i_queue][
1562                 i_input]));
1563         t_cpDH_0 = DIFF_T0(counters->round[i_round]-> load_copyHD_end[
1564             i_queue][i_input]);
1565
1566         cerr <<
1567             "\t\t" << i_queue << " - R[" << i_round << "] Input [" << i_input << "
1568                 unmarshal: " << unmarshal << "\t\tT: " << t_Unmarshal <<
1569                     endl <<
1570                     "\t\t" << i_queue << " - R[" << i_round << "] Input [" << i_input << "
1571                         -> D: " << cpDH_0 << "\t\tT: " << t_cpDH_0 <<
1572                             endl;

```

```

1564         }
1565
1566
1567
1568     long t_Kernel = DIFF_T0(counters->round[i_round]-> run_end[i_queue]);
1569     kernelTotal += 
1570     tKernel[i_queue] = DIFF(( counters->round[i_round]->run_start[i_queue]),
1571                               (counters->round[i_round]->run_end[i_queue]));
1572
1573     cerr <<
1574         "\t" << i_queue << " - R[" << i_round << "] KERNEL computation time: " <<
1575             tKernel[i_queue] << "\t\t\tT: " << t_Kernel <<
1576             endl;
1577             endl;
1578
1579         } // end queues
1580 } //end rounds
1581
1582 //FROM last load (lastRound, last_queue) TO unmapInput_end
1583 cerr << "Load - unmapInputBuffers: " << DIFF(counters->unmapInput_start, counters
1584             ->unmapInput_end) <<
1585             "\t\t\t\t\tT: " << DIFF_T0( counters->unmapInput_end)
1586             << endl;
1587
1588 cerr << "Load - releaseInBuffersH: " << DIFF(counters->unmapInput_end, counters
1589             ->releaseInBuffersH_T) << "\t\t\t\t\tT: " << DIFF_T0( counters->
1590             releaseInBuffersH_T) << "\n";
1591
1592
1593 for(int i = 0; i < NUMSTAGES; i++)
1594     cerr << "M - join " << i << "\t\t\t\t\t\t\tT: " << DIFF_T0( joinT[i]) << "\n"
1595             ;
1596
1597 cerr << "M - releaseEvents: " << DIFF(joinT[NUMSTAGES-1], releaseEventsT ) << "\\
1598             t\t\t\t\t\t\tT: " << DIFF_T0( releaseEventsT) << "\n";
1599
1600 cerr << endl <<
1601 "KERNEL total time: " << kernelTotal << ". avg on " << numRounds << " rounds: " <<
1602             kernelTotal/ NUMSEGMENTS <<
1603             endl <<
1604             "Load time: " << inputTotal << ". avg: " << inputTotal / NUMSEGMENTS <<
1605             endl <<
1606             "\tunmarshalTime: " << unmarshalTotal << ". avg: " << unmarshalTotal / (NUMSEGMENTS *
1607                 pInputC) <<
1608             endl <<

```

```

1601     "\tcopyHDTIME: " << inputTotal - unmarshalTotal << ". avg: " << (inputTotal -
1602         unmarshalTotal) / (NUMSEGMENTS * pInputC) <<
1603         endl <<
1604         endl ;
1605
1606 #endif
1607
1608 #if defined(TIME) || defined(TIME_FUN)
1609
1610 #undef DIFF_TO
1611
1612     delete counters;
1613     counters = NULL;
1614 #endif
1615
1616
1617     return ATOM(ok);
1618
1619 }
1620
1621 static ERL_NIF_TERM mapLD(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
1622
1623     //mapLD_S(Kernel::kernel(), InputList ::[ double() ], OutputBuffer :: deviceBuffer())
1624     OCL_INIT_CHECK()
1625     NIF_ARITY_CHECK(4)
1626
1627     /*get the parameters (Kernel::kernel(), InputList ::[ double() ], OutputBuffer::
1628         deviceBuffer(), InputLength :: non_neg_integer())****/
1629
1630     kernel_sync* pKernel_s = NULL;
1631     if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**)&pKernel_s)) {
1632         cerr << "ERROR :: mapLD_S: 1st parameter is not a kernel_sync" << endl ;
1633         return enif_make_badarg(env);
1634     }
1635
1636     ERL_NIF_TERM inputList = argv[1];
1637     if (!enif_is_list(env, argv[1])) {
1638         cerr << "ERROR :: mapLD_S: 2nd parameter is not a list" << endl ;
1639         return enif_make_badarg(env);
1640     }
1641
1642     cl_mem* pOutputBufferD = NULL;
1643     if (!enif_get_resource(env, argv[2], deviceBuffer_rt, (void**)&pOutputBufferD)
1644         || *pOutputBufferD == NULL) {
1645         cerr << "ERROR :: mapLD_S: 3nd parameter is not a device buffer" << endl ;
1646         return enif_make_badarg(env);

```

```

1645     }
1646
1647     uint inputListLength;
1648     if (!enif_get_uint(env, argv[3], &inputListLength)) {
1649
1650         cerr << "ERROR :: mapLL: 4th parameter is not a non-neg-integer()" << endl ;
1651         return enif_make_badarg(env);
1652     }
1653     //*****
1654     /*
1655
1656     uint pInputC = 1;
1657     ERL_NIF_TERM* pInputV[1] = { &inputList };
1658
1659 #ifdef DEBUG
1660     cerr << "DEBUG: mapLD" << endl;
1661 #endif
1662
1663     return mapLD_Impl_3Thread(env, pKernel_s,(void**) pInputV, pInputC,
1664                             pOutputBufferD, inputListLength);
1665 }
1666
1667
1668
1669 static ERL_NIF_TERM map2LD(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
1670
1671     //map2LD_S(Kernel::kernel(), InputList1::[ double() ], InputList2::[ double() ],
1672     //           OutputBuffer::deviceBuffer(), InputLength::non-neg-integer())
1673     OCL_INIT_CHECK()
1674     NIF_ARITY_CHECK(5)
1675
1676     /*get the parameters (Kernel::kernel(), InputList1::[ double() ], InputList2::[
1677     double() ], OutputBuffer::deviceBuffer())******/
1678
1679     kernel_sync* pKernel_s = NULL;
1680     if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {
1681         cerr << "ERROR :: map2LD: 1st parameter is not a kernel_sync" << endl ;
1682         return enif_make_badarg(env);
1683     }
1684
1685     inputLists[2];
1686
1687     inputLists[0] = argv[1];
1688     if (!enif_is_list(env, argv[1])) {
1689         cerr << "ERROR :: map2LD: 2nd parameter is not a list" << endl ;

```

```

1688     return enif_make_badarg(env);
1689 }
1690
1691     inputLists[1] = argv[2];
1692     if (!enif_is_list(env, argv[2])) {
1693         cerr << "ERROR :: map2LD: 2nd parameter is not a list" << endl;
1694         return enif_make_badarg(env);
1695     }
1696
1697     cl_mem* pOutputBuffer = NULL;
1698     if (!enif_get_resource(env, argv[3], deviceBuffer_rt, (void**)&pOutputBuffer)
1699         || *pOutputBuffer == NULL) {
1700         cerr << "ERROR :: map2LD: 4th parameter is not a device buffer" << endl;
1701         return enif_make_badarg(env);
1702     }
1703
1704     uint inputListLength;
1705     if (!enif_get_uint(env, argv[4], &inputListLength)) {
1706
1707         cerr << "ERROR :: mapLL: 5th parameter is not a non_neg_integer()" << endl;
1708         return enif_make_badarg(env);
1709     }
1710     //*****
1711     uint pInputC = 2; //Map2
1712     ERL_NIF_TERM* pInputV[2] = { &(inputLists[0]), &(inputLists[1]) };
1713
1714 #ifdef DEBUG
1715     cerr << "DEBUG: map2LD" << endl;
1716 #endif
1717
1718
1719     return mapLD_Impl_3Thread(env, pKernel_s, (void**) pInputV, pInputC,
1720                             pOutputBuffer, inputListLength); //ok | {error, Why}
1721
1722
1723 static void* mapLL_unloadStage_Thread(void* obj) {
1724
1725 #ifdef DEBUG
1726     char msg[256];
1727 #endif
1728     MapLL_thread_params* params = (MapLL_thread_params *) obj;
1729
1730     ErlNifEnv* env = params->env; //needed by CLEANUP macros
1731     ErlNifMutex*& env_mtx = params->env_mtx;

```

```

1732
1733     Map_fun_counters*& counters = params->counters;
1734
1735     //locals
1736     size_t szInput_local = (params->szSegment / params->uiNumElemSegment) * params->
1737     uiNumElem;
1738
1739     cl_int ciErrNum;
1740
1741
1742     uint zeroIdx;
1743     uint segmentOffset;
1744
1745     for (uint i_round = 0 ; i_round < params->numRounds; ++i_round) {
1746
1747         zeroIdx = (params->numQueues * i_round);
1748
1749 #ifdef DEBUG
1750         sprintf(msg, "DEBUG: Unload - Starting round[%d]\n", i_round);
1751         cerr << msg;
1752 #endif
1753
1754
1755     uint idx = 0;
1756     for (uint i_queue = 0 ; i_queue < params->numQueues; ++i_queue) {
1757
1758
1759 #ifdef DEBUG
1760         sprintf(msg, "DEBUG: Unload - %d - waiting to start round[%d]\n",
1761                 i_queue, i_round);
1762         cerr << msg;
1763
1764     idx = zeroIdx + i_queue;
1765
1766     segmentOffset = idx * params->uiNumElemSegment;
1767
1768     cl_command_queue& curr_cmdQ = params->cmdQs[i_queue];
1769     cl_event& curr_run_done_evt = params->run_done_evt[i_queue][i_round];
1770
1771
1772     //WAIT for RUN stage
1773     clWaitForEvents(1, &(curr_run_done_evt));
1774
1775
1776     if(params->errorSignal) { //Someone signaled an error, exit.

```

```

1777 #ifdef DEBUG
1778         sprintf(msg, "DEBUG: Unload - %d - Someone signaled an error, exit.\\
1779                         n", i_queue);
1780     cerr << msg;
1780 #endif
1781     goto cleanup;
1782 }
1783
1784 #ifdef TIME
1785     GET_TIME(counters->round[i_round]->unload_copy_start[i_queue]);
1786 #endif
1787
1788 //***** D -> H *****
1789 double* pMappedOutputBufferSegmentH = params->pMappedOutputBufferH +
1790     segmentOffset;
1791
1792 //blocking read because i need the data immediately for the marshalling
1792 THREAD_CHK_SUCCESS.CLEANUP.GOTO(
1793     clEnqueueReadBuffer(curr_cmdQ, params->outputBufferD, CL_TRUE,
1794     segmentOffset, params->szSegment, (void*)
1795     pMappedOutputBufferSegmentH, 0, NULL, NULL); ,
1796
1797     params->errorTerm ,
1798     {
1799         sprintf(msg, "DEBUG: Unload - %d - Error D -> H\n", i_queue); cerr
1800             << msg;
1801         params->signalError();
1802     }
1803 )
1804     clFlush(curr_cmdQ);
1805
1806 #ifdef DEBUG
1807     sprintf(msg, "DEBUG: Unload - %d - D -> H\n", i_queue);
1808     cerr << msg;
1806 #endif
1807
1808
1809 #ifdef TIME
1810     GET_TIME(counters->round[i_round]->unload_copy_end[i_queue]);
1811 #endif
1812
1813 //overlap last marshalling with input buffer unmapping
1814 if(i_round == params->numRounds - 1 && i_queue == NUMQUEUES-1)
1815     clSetUserEventStatus(params->unmapInput_evt, CL_SUCCESS);
1816
1817
1818 //***** H -> L *****

```

```

1819
1820         ERL_NIF_TERM* outputTermsSegment = params->outputTerms + segmentOffset;
1821
1822         enif_mutex_lock(env_mtx);
1823
1824         for (uint i = 0; i < params->uiNumElemSegment; i++)
1825             outputTermsSegment[i] = enif_make_double(env,
1826                                         pMappedOutputBufferSegmentH[i]);
1827
1828         enif_mutex_unlock(env_mtx);
1829
1830 #ifdef TIME
1831     GET_TIME(( counters->round[i_round]->unload_marshal_end[i_queue]));
1832 #endif
1833
1834
1835 #ifdef TIME
1836     GET_TIME(counters->round[i_round]->end[i_queue]);
1837 #endif
1838
1839 #ifdef DEBUG
1840     sprintf(msg, "DEBUG: Unload - %d - H -> L\nDEBUG: %d - Round END\n",
1841             i_queue, i_queue);
1842     cerr << msg;
1843 #endif
1844
1845
1846 } //end queue
1847
1848 //signal master to start list creation, overlaps with the following unmap
1849 clSetUserEventStatus(params->lastMarshalling_evt, CL_SUCCESS);
1850
1851 cleanup:
1852
1853     if (params->errorSignal) { //release every lock before exiting to avoid deadlocks
1854
1855 #ifdef DEBUG
1856     sprintf(msg, "DEBUG: Unload - Error detected, set all events\n");
1857     cerr << msg;
1858 #endif
1859
1860     clSetUserEventStatus(params->unmapInput_evt, CL_SUCCESS);
1861     clSetUserEventStatus(params->lastMarshalling_evt, CL_SUCCESS);
1862 }
1863

```

```

1864     //UNMAP using last queue
1865     clEnqueueUnmapMemObject( params->cmdQs[ NUM_QUEUES - 1 ], params->outputBufferH ,
1866                               params->pMappedOutputBufferH , 0 , NULL , NULL ) ;
1867 #ifdef TIME
1868     GET_TIME( counters->unmapOutput_end ) ;
1869 #endif
1870
1871     //*****RELEASE BUFFERS*****
1872     // outputBufferH szInput
1873     clReleaseMemObject( params->outputBufferH ) ;
1874
1875 #ifdef TIME
1876     GET_TIME( counters->releaseOutBufferH_T ) ;
1877 #endif
1878
1879     //outputBufferD szInput
1880     clReleaseMemObject( params->outputBufferD ) ;
1881 #ifdef TIME
1882     GET_TIME( counters->releaseOutBufferD_T ) ;
1883 #endif
1884
1885
1886 #ifdef TIME
1887     GET_TIME( ( counters->releaseInBuffersD_T_start ) );
1888 #endif
1889     //inputBuffersD [ pInputC ] [ NUM_SEGMENTS ] * szSegment = szInput
1890     for ( uint i_input = 0; i_input < params->pInputC; ++i_input ) {
1891
1892         for ( int i_segm = 0; i_segm < params->numSegments; ++i_segm ) {
1893             clReleaseMemObject( params->inputBuffersD [ i_input ] [ i_segm ] ) ;
1894         }
1895
1896         delete [] params->inputBuffersD [ i_input ] ;
1897     }
1898 #ifdef TIME
1899     GET_TIME( ( counters->releaseInBuffersD_T_end ) );
1900 #endif
1901
1902
1903 #ifdef DEBUG
1904     sprintf( msg , "DEBUG: Unload - Done\n" );
1905     cerr << msg ;
1906 #endif
1907
1908     return NULL;
1909 }
```

```

1910
1911
1912 //Number of segments of list splitting in mapLL
1913 #define NUMSEGM 8
1914
1915 static ERL_NIF_TERM mapLL_Impl_3Thread(ErlNifEnv * env, kernel_sync* pKernel_s, void
** _pInputV, uint pInputC, void* _pOutputBuffer, uint listLength) {
1916
1917 #ifdef DEBUG
1918     cerr << "DEBUG: M - mapLL_Impl_3Thread" << endl;
1919 #endif
1920
1921     ERL_NIF_TERM toReturn;
1922
1923     const uint NUMSEGMENTS = NUMSEGM;
1924
1925     const uint NUMQUEUES.LOCAL = NUMQUEUES;
1926
1927     const uint numRounds = NUMSEGMENTS/NUMQUEUES.LOCAL;
1928
1929
1930 #if defined(TIME) || defined(TIME_FUN)
1931
1932 #define DIFF_T0(COUNTER) nsec2usec( diff( &(counters->fun_start), &(COUNTER) ) )
1933
1934     Map_fun_counters* counters = new Map_fun_counters(numRounds, pInputC);
1935 #endif
1936
1937 #ifdef TIME
1938
1939     timespec createEventsT_start, createEventsT_end,
1940     createInputBuffersH[pInputC],
1941     createInputBuffersD[pInputC],
1942
1943     createOutputBufferD;
1944
1945     timespec joinT[3],
1946     list_createdT_start, list_createdT_end,
1947     releaseEventsT;
1948
1949
1950 #endif
1951
1952 #if defined(TIME) || defined(TIME_FUN)
1953     GET_TIME(counters->fun_start);
1954 #endif
1955

```

```

1956
1957     std::vector<cl_event> events;
1958
1959     //cast input/output to type relevant here
1960     //pInputV is an array of ERL_NIF_TERM* representing, each one, an erlang list
1961     ERL_NIF_TERM** inputLists = (ERL_NIF_TERM**) _pInputV;
1962
1963     //input lists must have at least minListLen elements,
1964     //otherwise they can't be split in NUMSEGMENTS segments
1965     uint listLen = listLength;
1966
1967
1968     //check buffer allignment requirements (depends on the device)
1969     uint minListLen = device_minBaseAddrAlignByte / sizeof(double);
1970     if(listLen/NUMSEGMENTS < minListLen)
1971         return make_error(env, ATOM(skel_ocl_input_list_too_short));
1972
1973
1974     uint uiNumElem = listLen;
1975
1976     uint uiNumElemSegment = uiNumElem / NUMSEGMENTS;
1977
1978     size_t szInput = uiNumElem * sizeof(double);
1979
1980     size_t szSegment = szInput / NUMSEGMENTS;
1981
1982
1983 #ifdef DEBUG
1984     cerr << "DEBUG: M - mapLL_impl:" << "#Elem: " << uiNumElem << " #ElemSegment: " <<
1985             uiNumElemSegment << " szInput: " << szInput << " szSegment: " << szSegment <<
1986             endl;
1987 #endif
1988
1989     cl_int ciErrNum = 0;
1990
1991     //working queues
1992     cl_command_queue cmdQs[2] = {getCommandQueue(0), getCommandQueue(1)};
1993
1994     //mutex protecting env from concurrent modifications (enif_make_xxx)
1995     ErlNifMutex* env_mtx = enif_mutex_create((char*)"env_mtx");
1996
1997     //terms array from which to create the output list using
1998         //enif_make_list_from_array, filled by unload stage
1999     ERL_NIF_TERM* outputTerms = NULL;

```

```

2000
2001 #ifdef TIME
2002     GET_TIME(createEventsT_start);
2003 #endif
2004
2005     //*****events setup [NUMQUEUES] [numRounds])
2006     cl_event* run_start_evt [NUMQUEUES_LOCAL];
2007     cl_event* run_done_evt [NUMQUEUES_LOCAL];
2008
2009     for (int i = 0; i < NUMQUEUES_LOCAL; ++i) {
2010         run_start_evt [i] = new cl_event [numRounds];
2011         run_done_evt [i] = new cl_event [numRounds];
2012
2013         for (int i_round = 0; i_round < numRounds; ++i_round) {
2014             run_start_evt [i][i_round] = clCreateUserEvent(context, &ciErrNum);
2015             run_done_evt [i][i_round] = clCreateUserEvent(context, &ciErrNum);
2016
2017             events.push_back(run_start_evt [i][i_round]);
2018             events.push_back(run_done_evt [i][i_round]);
2019
2020     }
2021 }
2022
2023 //to allow overlapping between unmarshalling and mapping the outputBuffer (can't
2024 // map more in parallel)
2025 // fired by load stage after finishing the mapping of inputBuffer, wait by
2026 // unload
2027     cl_event mapOutput_evt = clCreateUserEvent(context, &ciErrNum);
2028     events.push_back(mapOutput_evt);
2029
2030 //to allow overlapping between input buffer unmapping and last marshalling (set
2031 // by unload, wait by load).
2032     cl_event unmapInput_evt = clCreateUserEvent(context, &ciErrNum);
2033     events.push_back(unmapInput_evt);
2034
2035 //to overlap output list creation to unmapOutputBuffer (set by unload, wait by
2036 // master)
2037     cl_event lastMarshalling_evt = clCreateUserEvent(context, &ciErrNum);
2038     events.push_back(lastMarshalling_evt);
2039
2040 //*****threads setup: for now, one thread per stage.
2041 //Load, run, unload. Each stage works on 2 command queues
2042     const uint NUMSTAGES = 3;

```

```

2043     ErlNifTid tid [NUMSTAGES];
2044
2045     ErlNifTid& loadStage = tid [0];
2046     ErlNifTid& runStage = tid [1];
2047     ErlNifTid& unloadStage = tid [2];
2048
2049     //object holding everything needed by the threads
2050     MapLL_thread_params* conf = NULL;
2051
2052
2053     //*****for each input list , create:
2054     // 1 inputBufferH
2055     // NUMSEGMENTS inputBufferD
2056
2057     cl_mem inputBuffersH [pInputC];
2058
2059     // [pInputC] [NUMSEGMENTS]
2060     cl_mem* inputBuffersD [pInputC];
2061     for (int i = 0; i < pInputC; ++i)
2062         inputBuffersD [i] = new cl_mem [NUMSEGMENTS];
2063
2064
2065     //reuse first inputBuffersH as output buffer
2066
2067     cl_mem outputBufferH;
2068
2069     // device output buffer (same size as input)
2070     cl_mem outputBufferD;
2071
2072
2073     for (uint i_list = 0; i_list < pInputC; ++i_list) {
2074         // create one input host buffer of size szInput , released by Unload
2075         CHK_SUCCESS.CLEANUP.GOTO(
2076             createBuffer (szInput , CL_MEM_READ_WRITE |
2077                         CL_MEM_ALLOC_HOST_PTR, &(inputBuffersH [i_list])
2078                         ); ,
2079         )
2080 #ifdef TIME
2081     GET_TIME(createInputBuffersH [i_list]);
2082 #endif
2083
2084
2085 #ifdef DEBUG
2086     cerr << "DEBUG: M - " << "inputBuffersH [" << i_list << "] CREATED: size: "
2087             << szInput << endl;

```

```

2087 #endif
2088
2089
2090     //*****create NUMSEGMENTS input device buffers having size szSegment
2091     *****
2092     //released by Unload
2093     for (int i_segm = 0; i_segm < NUMSEGMENTS; ++i_segm) {
2094
2095         CHK_SUCCESS_CLEANUP_GOTO(
2096             createBuffer(szSegment, CL_MEM_READ_ONLY, &(

2097                 inputBuffersD[i_list][i_segm])) , ,
2098
2099             );
2100
2101     #ifdef TIME
2102         GET_TIME(createInputBuffersD[i_list]);
2103     #endif
2104     #ifdef DEBUG
2105         cerr << "DEBUG: M - " << "inputBuffersD["<<i_list<<" ][ " <<i_segm<<" ]
2106         CREATED: size: " << szSegment<< endl;
2107     #endif
2108     }
2109     #ifdef DEBUG
2110         cerr << "DEBUG: M - " << "inputBuffersD CREATED"<< endl;
2111     #endif
2112 }
2113
2114     outputBufferH = inputBuffersH[0]; //reuse first inputBufferH as outputBufferH
2115     clRetainMemObject(outputBufferH);
2116
2117     //*****Create one OUTPUT buffer on DEVICE*****
2118     //released by Unload
2119     CHK_SUCCESS_CLEANUP_GOTO(
2120         createBuffer(szInput, CL_MEM_WRITE_ONLY, &(

2121             outputBufferD)) , ,
2122
2123     );
2124     #ifdef TIME
2125         GET_TIME(createOutputBufferD);
2126     #endif
2127
2128

```

```

2129 //****Create the array that will store the ERL_NIF_TERM representing the output
2130 //elements*****
2131
2132     outputTerms =
2133         (ERL_NIF_TERM*) enif_alloc(sizeof(ERL_NIF_TERM) * uiNumElem);
2134
2135     if(outputTerms == NULL) {
2136         cerr << "DEBUG: M - " << "enif_alloc returned NULL." << ciErrNum << endl;
2137
2138         toReturn = make_error(env, ATOM(erl_enomem));
2139         goto cleanup;
2140     }
2141
2142
2143 //Clone kernel , 'cause i need two of them, one per queue
2144 cl_kernel kernels[NUM_QUEUES_LOCAL];
2145
2146 kernels[0] = pKernel_s->kernel;
2147
2148 if(NUM_QUEUES.LOCAL == TWO)
2149     cloneKernel(pKernel_s->kernel, &kernels[1]);
2150
2151 if(NUM_QUEUES.LOCAL > MAX_QUEUES) {
2152     toReturn = make_error(env, ATOM(ocl_num_queue_not_supported));
2153     goto cleanup;
2154 }
2155
2156
2157
2158 //*****Main LOOP*****
2159
2160 //iterator on the list , points to the head of the current segment of the list
2161 ERL_NIF_TERM currList[pInputC];
2162 for (uint i = 0; i < pInputC; ++i)
2163     currList[i] = *(inputLists[i]);
2164
2165
2166 conf = new MapLL_thread_params (
2167             mapOutput_evt,
2168             unmapInput_evt,
2169             lastMarshallng_evt,
2170             run_start_evt, run_done_evt,
2171             env, env_mtx,
2172             pInputC, cmdQs, kernels,
2173             inputBuffersD, outputBufferD,
2174             inputBuffersH,
```

```

2175                               outputBufferH ,
2176                               currList ,
2177 #ifdef TIME
2178                               counters ,
2179 #else
2180                               NULL,
2181 #endif
2182                               numRounds ,
2183                               NUM_QUEUES_LOCAL,
2184                               0 ,
2185                               uiNumElem , uiNumElemSegment , szSegment ,
2186                               outputTerms
2187                               ) ;
2188
2189
2190
2191 #ifdef DEBUG
2192     cerr << "DEBUG: M - Starting threads." << endl;
2193 #endif
2194
2195     enif_thread_create((char*) "mapLL loadStage" , &loadStage , mapLL_loadStage_Thread
2196                         , conf , NULL);
2197
2198     enif_thread_create((char*) "mapLL runStage" , &runStage , mapLL_runStage_Thread ,
2199                         conf , NULL);
2200
2201
2202 #ifdef TIME
2203     GET_TIME((counters->fun_prologue));
2204 #endif
2205
2206
2207 //wait last marshalling then start creating the list
2208 clWaitForEvents(1 , &conf->lastMarshalling_evt);
2209
2210
2211     if(conf->errorSignal) {
2212         //something's gone wrong, cleanup and return the error term (the last one
2213         //generated)
2214
2215 #ifdef DEBUG
2216         cerr << "DEBUG: M - Got an error: cleanup and exit" << endl;
2217 #endif
2218
2219     toReturn = conf->errorTerm;

```

```

2218     }
2219     else {
2220         //everything went fine
2221 #ifdef DEBUG
2222         cerr << "DEBUG: M - Creating output list\n";
2223 #endif
2224
2225
2226 #ifdef TIME
2227     GET_TIME(list_createdT_start);
2228 #endif
2229
2230     toReturn = enif_make_tuple2(env,ATOM(ok),enif_make_list_from_array(env,
2231                             outputTerms, uiNumElem));
2232
2233 #ifdef TIME
2234     GET_TIME(list_createdT_end);
2235 #endif
2236
2237
2238     //wait for termination
2239     for(uint i = 0; i < NUMSTAGES; ++i) {
2240
2241         enif_thread_join(tid[i], NULL);
2242
2243 #ifdef TIME
2244         GET_TIME((joinT[i]));
2245 #endif
2246
2247 #ifdef DEBUG
2248         char msg[64];
2249         sprintf(msg, "DEBUG: M - Thread %d joined\n", i);
2250         cerr << msg;
2251 #endif
2252     }
2253
2254     //check again after joining, shouldn't be useful since list creation starts
2255     //after the last marshalling, just before outputBuffer's unmapping,
2256     //at the very end of the computation when nobody can signal an error anymore.
2257     if(conf->errorSignal) {
2258         //something's gone wrong, cleanup and return the error term (the last one
2259         //generated)
2260 #ifdef DEBUG
2261         cerr << "DEBUG: M - Got an error: cleanup and exit" << endl;
2262 #endif
2263     toReturn = conf->errorTerm;

```

```

2262     }
2263
2264     // ****CLEANUP label*****
2265 cleanup :
2266
2267     if (outputTerms) {
2268         enif_free (outputTerms); //2-3 usecs
2269         outputTerms = NULL;
2270     }
2271
2272
2273     // *****EVENTS*****
2274     if (!events.empty())
2275         releaseEvents(events);
2276
2277     for (uint i = 0; i < NUM_QUEUES_LOCAL; i++){
2278         delete [] run_start_evt[i];
2279         delete [] run_done_evt[i];
2280     }
2281 #ifdef TIME
2282     GET_TIME((releaseEventsT));
2283 #endif
2284
2285     // *****KERNELS*****
2286     if (NUM_QUEUES_LOCAL == 2){
2287         if (kernels[1])
2288             clReleaseKernel(kernels[1]); //2 usecs
2289         kernels[1] = NULL;
2290     }
2291
2292     // *****ENV MUTEX*****
2293     if (env_mtx)
2294         destroyMutexes(&env_mtx, 1); //1 usec
2295
2296     // *****CONF*****
2297     if (conf)
2298         delete conf;
2299
2300
2301 #ifdef DEBUG
2302     cerr << "DEBUG: M - Cleanup done. Return" << endl;
2303 #endif
2304
2305
2306 #if defined(TIME) || defined(TIME_FUN)
2307     GET_TIME((counters->fun_end));
2308#endif

```

```

2309
2310
2311
2312 // ****RUN STATISTITCS*****
2313
2314
2315 #if defined(TIME) || defined(TIME_FUN)
2316
2317     cerr << endl <<
2318     "----- map" << pInputC << "LL: (usec)
2319     "-----" <<
2320     endl <<
2321     " Processing " << pInputC << " X " << uiNumElem << " elements in " << NUMSEGMENTS
2322     << " segments (" << uiNumElemSegment << " per segment)" <<
2323     endl <<
2324     " Total run time: " << DIFF_T0(counters->fun_end) <<
2325     endl;
2326
2327
2328 #endif
2329
2330
2331
2332
2333     cerr << endl <<
2334     "Function prologue time: " << prologueTime << "\t\t\t\t\tT: " << prologueTime <<
2335     endl;
2336
2337     char str[128];
2338     sprintf(str, "\tcreateEvents: %lu \t\t\t\t\tT: %lu\n", DIFF(createEventsT_start
2339     , createEventsT_end), DIFF_T0( createEventsT_end)); cerr << str;
2340
2341     for (uint i_input = 0; i_input < pInputC; i_input++) {
2342         sprintf(str, "\tcreateInputBuffersH[%d]: \t\t\t\t\tT: %lu\n", i_input, DIFF_T0(
2343             createInputBuffersH[i_input])); cerr << str;
2344
2345         sprintf(str,
2346                 "Load - mapInputBuffers[%d]: %lu \t\t\t\t\tT: %lu\n", i_input,
2347                 i_input == 0 ?
2348                 DIFF( counters->mapInput_start, counters->mapInput_end[i_input] ) :
2349                 DIFF( counters->mapInput_end[i_input - 1], counters->mapInput_end[
2350                     i_input]),
2351                 DIFF_T0( counters->mapInput_end[i_input])
2352             );
2353
2354         cerr << str;

```

```

2350
2351     sprintf(str, "\tcreateInputBuffersD[%d]: \t\t\t\tT: %lu\n", i_input, DIFF_T0
2352         ( createInputBuffersD[i_input])); cerr << str;
2353 }
2354
2355 sprintf(str, "\tcreateOutputBufferD: \t\t\t\tT: %lu\n", DIFF_T0(
2356     createOutputBufferD)); cerr << str;
2357
2358 uint kernelTotal = 0, inputTotal = 0, outputTotal = 0, marshTotal = 0,
2359     unmarshalTotal = 0;
2360 uint i_round;
2361 //round statistics
2362 for (i_round = 0; i_round < numRounds; i_round++) {
2363     uint tKernel[2] = {0,0};
2364     uint tInput = 0, tOutput = 0;
2365     uint tStartRound[2];
2366
2367     cerr <<
2368     "\nRound["<< i_round << "] total time: " << DIFF(( counters->round[i_round]
2369         ]->start[0]), ( counters->round[i_round]->end[NUM_QUEUES_LOCAL-1]))<<
2370     endl;
2371
2372     for (int i_queue = 0; i_queue < NUM_QUEUES_LOCAL; i_queue++) {
2373
2374         inputTotal +=
2375             tInput = DIFF(( counters->round[i_round]->load_start[i_queue]), (
2376                 counters->round[i_round]->load_end[i_queue]));
2377
2378         tStartRound[i_queue] = DIFF_T0( counters->round[i_round]-> start[i_queue
2379             ]);
2380         cerr <<
2381         "\n\t"<< i_queue << " - R["<< i_round << "] Start round\t\t\t\tT: " <<
2382             tStartRound[i_queue]<<
2383             endl <<
2384             "\t"<< i_queue << " - R["<< i_round << "] Load total time: " << tInput <<
2385             endl;
2386
2387         for (uint i_input = 0; i_input < pInputC; i_input++) {
2388             long unmarshal = 0, t_Unmarshal;
2389             long cpDH_0 = 0, t_cpDH_0;
2390
2391             unmarshalTotal +=

```

```

2389         unmarshal = DIFF((counters->round[i_round]->load_input_start[i_queue][
2390                         i_input]), (counters->round[i_round]->load_unmarshal_end[i_queue][
2391                         i_input]));
2392
2393
2394         cpDH_0 = DIFF((counters->round[i_round]->load_unmarshal_end[i_queue][
2395                         i_input]), (counters->round[i_round]->load_copyHD_end[i_queue][
2396                         i_input]));
2397
2398         t_cpDH_0 = DIFF_T0(counters->round[i_round]-> load_copyHD_end[
2399                         i_queue][i_input]);
2400
2401         cerr <<
2402             "\t\t<< i_queue << " - R["<< i_round << "] Input["<< i_input <<"]
2403                 unmarshal: " << unmarshal << "\t\tT: "<< t_Unmarshal <<
2404                     endl <<
2405             "\t\t<< i_queue << " - R["<< i_round << "] Input["<< i_input <<"] H
2406                 -> D: " << cpDH_0 << "\t\tT: "<< t_cpDH_0 <<
2407                     endl;
2408     }
2409
2410
2411     long t_Kernel = DIFF_T0(counters->round[i_round]-> run_end[i_queue]);
2412     kernelTotal +=
2413         tKernel[i_queue] = DIFF((counters->round[i_round]->run_start[i_queue]),
2414                               (counters->round[i_round]->run_end[i_queue]));
2415
2416     cerr <<
2417         "\t" << i_queue << " - R["<< i_round << "] KERNEL computation time: " <<
2418             tKernel[i_queue] << "\t\tT: "<< t_Kernel <<
2419             endl;
2420
2421
2422     long marsh, cpDH;
2423     long t_marshal, t_cpDH;
2424
2425
2426     cpDH = DIFF((counters->round[i_round]->unload_copy_start[i_queue]), (
2427                     counters->round[i_round]->unload_copy_end[i_queue]));
2428     t_cpDH = DIFF_T0(counters->round[i_round]->unload_copy_end[i_queue]);
2429
2430     marshTotal +=
2431     marsh = DIFF((counters->round[i_round]->unload_copy_end[i_queue]), (
2432                     counters->round[i_round]->unload_marshall_end[i_queue]));

```

```

2424         t_marsh = DIFF_T0(counters->round[i_round]->unload_marshal_end[i_queue]) ;
2425
2426
2427         //sprintf(msg, "\t%d - R[%d] Unload time: %li", i_queue, i_round, cpDH +
2428         //marsh) ;
2429         cerr <<
2430         "\t" << i_queue << " - R[" << i_round << "] Unload time: " << cpDH + marsh
2431         <<
2432         endl <<
2433         "\t\t" << i_queue << " - R[" << i_round << "] D -> H: " << cpDH << "\t\t\
2434         t\T: " << t_cpDH <<
2435         endl <<
2436         "\t\t" << i_queue << " - R[" << i_round << "] marshall: " << marsh << "\t\
2437         t\T: " << t_marsh <<
2438         endl ;
2439
2440     } // end queues
2441 } //end rounds
2442
2443
2444 //FROM last_load (lastRound, last_queue) TO unmapInput_end
2445 cerr << "Load - unmapInputBuffers: " << DIFF(counters->unmapInput_start, counters
2446         ->unmapInput_end) <<
2447         "\t\t\t\t\T: " << DIFF_T0( counters->unmapInput_end)
2448         << endl ;
2449
2450
2451 cerr << "Load - releaseInBuffersH: " << DIFF(counters->unmapInput_end, counters
2452         ->releaseInBuffersH_T) << "\t\t\t\t\T: " << DIFF_T0( counters->
2453         releaseInBuffersH_T) << "\n";
2454
2455
2456 cerr << "Unload - unmapOutputBuffers: " << DIFF(counters->round[numRounds-1]->end
2457         [NUM_QUEUES.LOCAL-1], counters->unmapOutput_end) <<
2458         "\t\t\t\t\T: " << DIFF_T0( counters->unmapOutput_end)
2459         << endl ;
2460
2461
2462 cerr << "Unload - releaseOutBufferH: " << DIFF(counters->unmapOutput_end,
2463         counters->releaseOutBufferH_T) << "\t\t\t\t\T: " << DIFF_T0(counters->
2464         releaseOutBufferH_T) << "\n";
2465
2466
2467 cerr << "Unload - releaseOutBufferD: " << DIFF( counters->releaseOutBufferH_T ,
2468         counters->releaseOutBufferD_T) << "\t\t\t\t\T: " << DIFF_T0(counters->
2469         releaseOutBufferD_T) << "\n";
2470
2471
2472 cerr << "Unload - releaseInBuffersD: " << DIFF(counters->
2473         releaseInBuffersD_T_start , counters->releaseInBuffersD_T_end) << "\t\t\t\t\T:
2474         " << DIFF_T0(counters->releaseInBuffersD_T_end) << "\n";
2475
2476

```

```

2457
2458     cerr << "M - make_list: "<< DIFF(list_createdT_start , list_createdT_end )<< "\t\
2459             t\t\t\t\t\tT: " << DIFF_T0( list_createdT_end ) << "\n";
2460
2461     for(int i = 0; i < NUMSTAGES; i++)
2462         cerr << "M - join " << i << "\t\t\t\t\t\t\t\tT: " << DIFF_T0( joinT [ i ]) << "\n"
2463         ;
2464
2465     cerr << "M - releaseEvents: "<< DIFF(joinT [NUM_STAGES-1], releaseEventsT ) << "\\
2466             t\t\t\t\t\t\tT: " << DIFF_T0( releaseEventsT ) << "\n";
2467
2468     cerr << endl <<
2469     "KERNEL total time: " << kernelTotal << ". avg on " << numRounds << " rounds: " <<
2470             kernelTotal/ NUMSEGMENTS <<
2471             endl <<
2472             "Load time: " << inputTotal << ". avg: " << inputTotal / NUMSEGMENTS <<
2473             endl <<
2474             "\tunmarshalTime: " << unmarshalTotal << ". avg: " << unmarshalTotal / (NUMSEGMENTS *
2475                 pInputC) <<
2476                 endl <<
2477                 "\tcopyHDTIME: " << inputTotal - unmarshalTotal << ". avg: " << (inputTotal -
2478                     unmarshalTotal) / (NUMSEGMENTS * pInputC) <<
2479                     endl <<
2480                     "\tmarshalTime: " << marshalTotal << ". avg: " << (marshalTotal) / NUMSEGMENTS <<
2481                     endl ;
2482
2483 #endif
2484
2485
2486 #if defined(TIME) || defined(TIME_FUN)
2487
2488 #undef DIFF_TO
2489     delete counters;
2490     counters = NULL;
2491
2492 #endif
2493
2494 static ERL_NIF_TERM mapLL(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[])
2495 {
2496     OCL_INIT_CHECK()
2497     const uint NUMARGS = 3;
2498     //mapLL(Kernel::kernel(), InputList::[ double() ], InputLength::non_neg_integer())

```

```

2498
2499     NIF_ARITY_CHECK(NUM_ARGS)
2500
2501     /*get the parameters (Kernel::kernel() , InputList ::[ double() ] , InputLength ::  

2502        non_neg_integer ()****/  

2503
2504     kernel_sync* pKernel_s = NULL;  

2505     if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {  

2506         cerr << "ERROR :: mapLL: 1st parameter is not a kernel_sync" << endl ;  

2507         return enif_make_badarg(env);  

2508     }  

2509
2510     ERL_NIF_TERM inputList = argv[1];  

2511     if (!enif_is_list(env, argv[1])) {  

2512         cerr << "ERROR :: mapLL: 2nd parameter is not a list" << endl ;  

2513         return enif_make_badarg(env);  

2514     }  

2515
2516     uint inputListLength;  

2517     if (!enif_get_uint(env, argv[2], &inputListLength)) {  

2518
2519         cerr << "ERROR :: mapLL: 3nd parameter is not a non-neg-integer()" << endl ;  

2520         return enif_make_badarg(env);  

2521     }  

2522
2523     //*****  

2524
2525
2526     uint pInputC = 1;  

2527     ERL_NIF_TERM* pInputV[1] = { &inputList };  

2528
2529 #ifdef DEBUG  

2530     cerr << "DEBUG: mapLL" << endl;  

2531 #endif  

2532
2533     return mapLL_Impl_3Thread(env, pKernel_s, (void**) pInputV, pInputC, NULL,  

2534         inputListLength);  

2535 }
2536
2537
2538 static ERL_NIF_TERM map2LL(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
2539
2540     //map2LL_S(Kernel::kernel() , InputList1 ::[ double() ] , InputList2 ::[ double() ])
2541     OCL_INIT_CHECK()

```

```

2542     const int ARGS_NUM = 4;
2543
2544     NIF_ARITY_CHECK(ARGS_NUM)
2545
2546     /*get the parameters (Kernel::kernel() , InputList::[ double() ]****/
2547
2548     kernel_sync* pKernel_s = NULL;
2549     if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {
2550         cerr << "ERROR :: map2LL: 1st parameter is not a kernel_sync" << endl ;
2551         return enif_make_badarg(env);
2552     }
2553
2554     ERL_NIF_TERM inputList1 = argv[1];
2555     if (!enif_is_list(env, argv[1])) {
2556         cerr << "ERROR :: map2LL: 2nd parameter is not a list" << endl ;
2557         return enif_make_badarg(env);
2558     }
2559
2560     ERL_NIF_TERM inputList2 = argv[2];
2561     if (!enif_is_list(env, argv[2])) {
2562         cerr << "ERROR :: map2LL: 3rd parameter is not a list" << endl ;
2563         return enif_make_badarg(env);
2564     }
2565
2566     uint inputListLength;
2567     if (!enif_get_uint(env, argv[3], &inputListLength)) {
2568
2569         cerr << "ERROR :: map2LL: 4th parameter is not a non_neg_integer()" << endl
2570         ;
2571         return enif_make_badarg(env);
2572     }
2573     //*****
2574
2575     uint pInputC = 2;
2576     ERL_NIF_TERM* pInputV[2] = { &inputList1, &inputList2 };
2577
2578 #ifdef DEBUG
2579     cerr << "DEBUG: map2LL" << endl;
2580 #endif
2581
2582     return mapLL_Impl_3Thread(env, pKernel_s, (void**) pInputV, pInputC, NULL,
2583     inputListLength);
2584
2585 }
```

B.2.3 Reduce Skeleton NIFs: `reduce_nifs.cpp`

```

1  /*****      REDUCE      *****/
2
3
4 static ERL_NIF_TERM reduceDD_Impl(ErlNifEnv * env, kernel_sync* pKernel_s, cl_mem*
5   pInputBuffer, cl_mem* pOutputBuffer) {
6
7   // %%reduceDD(Kernel::kernel(), InputBuffer::deviceBuffer(), OutputBuffer::
8   // deviceBuffer())
9
10  //Works with modified reduce6 kernel
11
12  //Max values from NVIDIA reduce example
13  const size_t maxThreads = 256;
14  const size_t maxBlocks = 64;
15
16  size_t uiNumBlocks = 0;
17  size_t uiNumThreads = 0;
18
19  size_t globalWorkSize[1];
20  size_t localWorkSize[1];
21
22  //*****Check buffers dimensions: Input must be power of 2
23  size_t szInputBufferByte = 0;
24
25  CHK_SUCCESS(getBufferSizeByte(*pInputBuffer, &szInputBufferByte));
26
27  uint uiN = szInputBufferByte / sizeof(double); //problem size is a power of 2
28
29  getNumBlocksAndThreads(uiN, maxBlocks, maxThreads, uiNumBlocks, uiNumThreads);
30
31  //allocate device output buffer having szOutputBufferByte size
32  size_t szOutputBufferByte = sizeof(double) * uiNumBlocks; //one value per block
33
34  cl_mem outputBufferD = NULL;
35
36  CHK_SUCCESS(createBuffer(szOutputBufferByte, CL_MEM_READ_WRITE, &outputBufferD))
37  ;
38
39  // Decide size of shared memory (one value per thread but kernel needs at least
40  // 64*sizeof(double) bytes)
41  size_t szLocalMemByte = (uiNumThreads <= 32) ? 2 * uiNumThreads * sizeof(double)
42  : uiNumThreads * sizeof(double);
43
44  //*****set reduce kernel parameters (first reduction)*****

```

```

41     cl_int ciErrNum = 0;
42
43 // clSetKernelArg is not thread-safe
44 enif_mutex_lock(pKernel_s->mtx);
45
46 ciErrNum |= clSetKernelArg(pKernel_s->kernel, 0, sizeof(cl_mem), (void*)
47     pInputBuffer);
48 ciErrNum |= clSetKernelArg(pKernel_s->kernel, 1, sizeof(cl_mem), (void*) &
49     outputBufferD);
50 ciErrNum |= clSetKernelArg(pKernel_s->kernel, 2, sizeof(unsigned int), (void*) &
51     uiN);
52 ciErrNum |= clSetKernelArg(pKernel_s->kernel, 3, szLocalMemByte, NULL);
53
54 enif_mutex_unlock(pKernel_s->mtx);
55
56 //****First reduce all elements so that each block produces one element
57 //*****
58 globalWorkSize[0] = uiNumThreads * uiNumBlocks;
59 localWorkSize[0] = uiNumThreads;
60
61
62 cl_command_queue cmdQ_0 = getCommandQueue(0);
63
64 //event which the second round awaits, needed to express the dependency that
65 //exists between the two rounds.
66 //It's mandatory only when using an out-of-order queue.
67 cl_event firstRoundEvt = NULL;
68
69 CHK_SUCCESS_CLEANUP(
70     clEnqueueNDRangeKernel(cmdQ_0, pKernel_s->kernel, 1, NULL,
71         &(globalWorkSize[0]), &(localWorkSize[0]), 0, NULL, &
72         firstRoundEvt),
73     clReleaseMemObject(outputBufferD));
74
75
76 //*****Second reduction*****
77 uiN = uiNumBlocks;
78
79 //since problem size (uiN) has changed, recompute uiNumThreads and
80 //szLocalMemByte
81 getNumBlocksAndThreads(uiN, maxBlocks, maxThreads, uiNumBlocks, uiNumThreads);
82 szLocalMemByte = (uiNumThreads <= 32) ? 2 * uiNumThreads * sizeof(double) :
83     uiNumThreads * sizeof(double);

```

```

79
80     enif_mutex_lock(pKernel_s->mtx);
81
82     //use as input previous reduction's output buffer
83     ciErrNum |= clSetKernelArg(pKernel_s->kernel, 0, sizeof(cl_mem), (void*) &
84         outputBufferD);
85     ciErrNum |= clSetKernelArg(pKernel_s->kernel, 1, sizeof(cl_mem), (void*) &
86         outputBufferD);
87     ciErrNum |= clSetKernelArg(pKernel_s->kernel, 2, sizeof(unsigned int), (void*) &
88         uiN);
89     ciErrNum |= clSetKernelArg(pKernel_s->kernel, 3, szLocalMemByte, NULL);
90
91
92     enif_mutex_unlock(pKernel_s->mtx);
93
94     CHK_SUCCESS_CLEANUP(ciErrNum, , clReleaseMemObject(outputBufferD));
95
96
97     //*****compute the result scalar from numBlock vector*****
98     globalWorkSize[0] = 1 * uiNumThreads; //only one block left
99     localWorkSize[0] = uiNumThreads;
100
101     CHK_SUCCESS_CLEANUP(
102             clEnqueueNDRangeKernel(cmdQ_0, pKernel_s->kernel, 1, NULL,
103             &(globalWorkSize[0]), &(localWorkSize[0]), 1, &
104             firstRoundEvt, NULL), ,
105             clReleaseMemObject(outputBufferD));
106
107
108     clReleaseMemObject(outputBufferD);
109
110     return ATOM(ok);
111 }
112
113
114 static ERL_NIF_TERM reduceDD(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[]) {
115
116     // %reduceDD(Kernel::kernel(), InputBuffer::deviceBuffer(), OutputBuffer::
117     // deviceBuffer())
118     OCL_INIT_CHECK()
119     NIF_ARITY_CHECK(3)

```

```

119  /*get the parameters (Kernel::kernel() , InputBuffer::deviceBuffer() ,
120   OutputBuffer::deviceBuffer())***** */
121
122 kernel_sync* pKernel_s = NULL;
123 if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {
124     cerr << "ERROR :: reduceDD: 1st parameter is not a kernel_sync" << endl ;
125     return enif_make_badarg(env);
126 }
127
128 cl_mem* pInputBuffer = NULL;
129 if (!enif_get_resource(env, argv[1], deviceBuffer_rt, (void**) &pInputBuffer) ||
130     *pInputBuffer == NULL) {
131     cerr << "ERROR :: reduceDD: 2nd parameter is not a device buffer" << endl ;
132     return enif_make_badarg(env);
133 }
134
135 cl_mem* pOutputBuffer = NULL;
136 if (!enif_get_resource(env, argv[2], deviceBuffer_rt, (void**) &pOutputBuffer)
137     || *pOutputBuffer == NULL) {
138     cerr << "ERROR :: reduceDD: 3nd parameter is not a device buffer" << endl ;
139     return enif_make_badarg(env);
140 }
141
142 CHK_SUCCESS(getBufferSizeByte(*pInputBuffer, &szInputBufferByte));
143
144 uint uiN = szInputBufferByte / sizeof(double);
145
146 if (!isPow2(uiN)) { //problem size MUST BE A POWER OF 2
147     cerr << "ERROR :: reduceDD: problem size is not a power of two." << endl ;
148     return enif_make_badarg(env);
149 }
150
151 return reduceDD_Impl(env, pKernel_s, pInputBuffer, pOutputBuffer);
152
153 }
154
155 static ERL_NIF_TERM reduceDL_Impl(ErlNifEnv * env, kernel_sync* pKernel_s, cl_mem*
156 pInputBufferD, cl_mem* pOutputBuffer) {
157 #if defined(TIME) || defined(TIME_FUN)
158     timespec
159     fun_start,
160     fun_end;

```

```

161
162     GET_TIME(( fun_start));
163 #endif
164
165 #ifdef TIME
166     timespec
167     fun_prologue_end ,
168     run_end ,
169     marsh_end ;
170#endif
171
172     (void) pOutputBuffer; //unused
173
174     std :: vector<cl_mem> buffers ;
175
176     size_t szInputBufferDByte = 0;
177
178     CHK_SUCCESS(getBufferSizeByte(*pInputBufferD , &szInputBufferDByte))
179
180     //*****allocate device input and output buffers having the size of the input
181     // and output ones*****
182     cl_mem outputBufferD = NULL;
183
184     size_t szOutputBufferDByte = sizeof(double);
185
186     CHK_SUCCESS_CLEANUP(createBuffer(szOutputBufferDByte , CL_MEM_READ_WRITE, &
187                                     outputBufferD); , releaseBuffers(buffers); )
188     buffers.push_back(outputBufferD);
189
190 #ifdef TIME
191     GET_TIME(( fun_prologue_end));
192 #endif
193
194     ERL_NIF_TERM returnTerm =
195     reduceDD_Impl( env , pKernel_s , pInputBufferD , &outputBufferD);
196
197 #ifdef TIME
198     GET_TIME(( run_end));
199 #endif
200
201     if(returnTerm == ATOM(ok)) { // no errors occurred , return the resulting double
202         double *pBuffer = NULL;
203         CHK_SUCCESS_CLEANUP(mapBufferBlocking(outputBufferD , 0, szOutputBufferDByte ,
204                                             CL_MAP_READ, &pBuffer); , releaseBuffers(buffers); )
205
206         returnTerm = enif_make_list1(env , enif_make_double(env , pBuffer[0]));

```

```

205
206         unMapBuffer( outputBufferD , pBuffer );
207     }
208
209 #ifdef TIME
210     GET_TIME(( marsh_end )) ;
211 #endif
212
213
214     releaseBuffers( buffers ) ;
215
216     ERL_NIF_TERM toReturn = enif_make_tuple2( env , ATOM(ok) , returnTerm ) ;
217
218 #if defined(TIME) || defined(TIME_FUN)
219     GET_TIME(( fun_end )) ;
220 #endif
221
222 #if defined(TIME) || defined(TIME_FUN)
223
224     cerr <<
225     endl <<
226     "----- ReduceDL ( usec )
227             " <<
228     "Total run time: " << DIFF(fun_start , fun_end ) << endl ;
229 #endif
230
231 #ifdef TIME
232     cerr <<
233     "Function prologue time: " << DIFF(fun_start , fun_prologue_end ) <<
234     endl <<
235     "KERNEL computation time: " << DIFF(fun_prologue_end , run_end ) <<
236     endl <<
237     "Marshalling time: " << DIFF(run_end , marsh_end ) <<
238     endl
239     ;
240 #endif
241
242     return toReturn ;
243
244 }
245
246 static ERL_NIF_TERM reduceDL( ErlNifEnv * env , int argc , const ERL_NIF_TERM argv [] ) {
247
248 // %%reduceHH( Kernel::kernel() , InputBuffer::deviceBuffer() )
249 OCL_INIT_CHECK()
250 NIF_ARITY_CHECK(2)

```

```

251  /*get the parameters (Kernel::kernel() , InputBuffer::deviceBuffer())
   *****/
252
253  kernel_sync* pKernel_s = NULL;
254  if (!enif_get_resource(env, argv[0], kernel_sync_rt, (void**) &pKernel_s)) {
255      cerr << "ERROR :: reduceDL: 1st parameter is not a kernel_sync" << endl ;
256      return enif_make_badarg(env);
257  }
258
259  cl_mem* pInputBufferD = NULL;
260  if (!enif_get_resource(env, argv[1], deviceBuffer_rt, (void**) &pInputBufferD)
261      || *pInputBufferD == NULL) {
262      cerr << "ERROR :: reduceDL: 2nd parameter is not a host buffer" << endl ;
263      return enif_make_badarg(env);
264  }
265  //*****
266
267  ERL_NIF_TERM result =
268  reduceDL_Impl(env, pKernel_s, pInputBufferD, NULL);
269
270  if( enif_is_list(env, result))
271      return enif_make_tuple2(env, ATOM(ok), result);
272  else return result;
273
274 }
275
276
277 static ERL_NIF_TERM skeletonlib(ErlNifEnv * env, int argc, const ERL_NIF_TERM argv[])
278 {
279     return enif_make_string(env, "Skel OCL ready", ERL_NIF_LATIN1);
}

```

B.2.4 Utilities: utils.cpp

```

1 //A synchronization barrier implemented using erlang NIF threading facilities
2 class Barrier {
3
4 private:
5     uint N_THREADS;
6
7     uint waiting_threads_counts[2];
8     uint current_counter;
9
10    ErlNifMutex* mtx;
11    ErlNifCond* cond;

```

```

12
13 public:
14     explicit Barrier( uint n_Threads) {
15         N_THREADS = n_Threads;
16         waiting_threads_counts[0] = waiting_threads_counts[1] = 0;
17         current_counter = 0;
18
19         mtx = enif_mutex_create((char*) "map_load_thread_cond_mtx");
20         cond = enif_cond_create((char*) "map_load_thread_cond");
21
22     }
23     ~Barrier() {
24         enif_cond_destroy(cond);
25         enif_mutex_destroy(mtx);
26     }
27
28     void await(){
29
30         enif_mutex_lock(mtx);
31
32         uint local_counter = current_counter;
33
34         waiting_threads_counts[local_counter]++;
35
36
37         if(waiting_threads_counts[local_counter] < N_THREADS) {
38
39             while(waiting_threads_counts[local_counter] < N_THREADS) {
40
41                 enif_cond_wait(cond, mtx);
42             }
43         }
44         else {
45
46             current_counter ^= 1;
47             waiting_threads_counts[current_counter] = 0;
48
49             enif_cond_broadcast(cond);
50         }
51
52         enif_mutex_unlock(mtx);
53     }
54
55 };
56
57 /*
58 * Erlang list unmarshaling function.

```

```

59 * lastListCell is a pointer to the last cell of the list:
60 * valid when arrayLen is less than list length list , NULL when the list is over.
61 */
62 uint list_to_double_arrayN(ErlNifEnv *env, ERL_NIF_TERM list, double* array, uint
   arrayLen, ERL_NIF_TERM* lastListCell) {
63
64     if(array == NULL)
65         return 0;
66
67     uint len = 0;
68
69     if(enif_is_list(env, list)) {
70
71         ERL_NIF_TERM curr_cell = list;
72
73         uint i = 0;
74         for(i = 0; enif_is_list(env, curr_cell) && i < arrayLen; i++) {
75
76             ERL_NIF_TERM hd, tl;
77             if(!enif_get_list_cell(env, curr_cell, &hd, &tl))
78                 break;//list is empty
79
80             if(!(enif_get_double(env, hd, &(array[i])) )) {
81                 std::cerr << "DEBUG: list_to_double_arrayN - Error: attempt to read float
82                 from something else!" << std::endl;
83             }
84             curr_cell = tl;
85         }
86
87         if(lastListCell != NULL) //user requested the term denoting the last cell
88             *lastListCell = curr_cell;
89
90         return i;
91
92     } else { //term "list" is not a list
93
94         #ifdef DEBUG
95             std::cerr << "DEBUG: list_to_double_arrayN - Error: trying to convert a non list
96             as a list"<< std::endl ;
97         #endif
98
99         return 0;
100    }
101
102 //synchronized version

```

```

103 uint sync_list_to_double_arrayN(ErlNifEnv *env, ErlNifMutex *mtx, ERL_NIF_TERM list,
104                                     double* array, uint arrayLen, ERL_NIF_TERM* lastListCell) {
105     enif_mutex_lock(mtx);
106
107     uint result = list_to_double_arrayN(env, list, array, arrayLen, lastListCell);
108
109
110     enif_mutex_unlock(mtx);
111
112
113     return result;
114
115 }
116
117 //Unmarshaling function
118 void list_to_double_array(ErlNifEnv *env, ERL_NIF_TERM list, double* array, uint
119                           listlen) {
120     if(array == NULL)
121         return;
122
123     if(enif_is_list(env, list)) {
124
125         list_to_double_arrayN(env, list, array, listlen, NULL);
126     }
127
128 }
129
130 //Marshaling function
131 ERL_NIF_TERM double_array_to_list(ErlNifEnv *env, double* array, size_t array_size)
132 {
133     if(array == NULL)
134         return ATOM(error);
135
136     ERL_NIF_TERM* floatTermArray = (ERL_NIF_TERM*) enif_alloc(sizeof(ERL_NIF_TERM) *
137                                                               array_size);
138
139     for(uint i = 0; i < array_size; i++)
140         floatTermArray[i] = enif_make_double(env, array[i]);
141
142     ERL_NIF_TERM res = enif_make_list_from_array(env, floatTermArray, array_size);
143
144     enif_free(floatTermArray);
145
146     return res;

```

```

146 }
147
148 /*returns a string containing the content of the file at filePath
149 * otherwise "NULL".*/
150 string readFromFileStr(char* filePath) {
151
152     ifstream file(filePath, ios::in);
153     if (!file.is_open()) {
154         cerr << "Failed to open file for reading: " << filePath << endl;
155         return "NULL";
156     }
157
158     ostringstream oss;
159     oss << file.rdbuf();
160
161     return oss.str();
162 }
163
164 /*returns a dinamically allocated string containing the content of the file at
   filePath
165 * otherwise NULL.
166 * The returned sting must be enif_free'd when not needed.
167 */
168 char* readFromFile(char* filePath) {
169
170     string srcStdStr = readFromFileStr(filePath);
171
172     char* toReturn = (char*) enif_alloc(srcStdStr.length() + 1);
173
174     return strcpy(toReturn, srcStdStr.c_str());
175
176 }
177
178
179 #define MIN(a, b) ((a < b) ? a : b)
180
181 int inline isPow2(unsigned int v) { return v && !(v & (v - 1)); }
182
183 /*
184  * A helper to return a value that is nearest value that is power of 2.
185 */
186 unsigned int nextPow2( unsigned int x )
187 {
188     --x;
189     x |= x >> 1;
190     x |= x >> 2;
191     x |= x >> 4;

```

```

192     x |= x >> 8;
193     x |= x >> 16;
194     return x+1;
195 }
196
197 // Round Up Division function
198 size_t roundUp(uint group_size, uint global_size)
199 {
200     uint r = global_size % group_size;
201     if(r == 0)
202         return global_size;
203     else
204         return global_size + group_size - r;
205 }
206
207 /*
208 * Compute the number of threads and blocks to use for the REDUCTION kernel.
209 * We set threads / block to the minimum of maxThreads and n/2 where n is
210 * problem size. We observe the maximum specified number of blocks , because
211 * each kernel thread can process more than 1 elements.
212 */
213 void getNumBlocksAndThreads(int n, int maxBlocks, int maxThreads, long unsigned int
214     &blocks, long unsigned int &threads)
215 {
216     threads = (n < maxThreads*2) ? nextPow2((n + 1)/ 2) : maxThreads;
217     blocks = (n + (threads * 2 - 1)) / (threads * 2);
218
219     if(maxBlocks > 0)
220         blocks = MIN(maxBlocks, blocks);
221 }
222 /*
223 * It finds all instances of a string in another string and replaces it with
224 * a third string.
225 */
226 void replaceTextInString(std::string& text, std::string find, std::string replace)
227 {
228     std::string::size_type pos=0;
229     while((pos = text.find(find, pos)) != std::string::npos)
230     {
231         text.erase(pos, find.length());
232         text.insert(pos, replace);
233         pos+=replace.length();
234     }
235 }
```

B.2.5 OpenCL Kernels

Map Kernel

```

1 //Map Kernels , FUN_NAME is replaced at runtime with user-defined function's name;
2           TYPE with double
3
4 //Unary Map kernel
5 __kernel void MapKernel(__global TYPE* input, __global TYPE* output, unsigned int
6   outputOffset, unsigned int uiNumElements)
7 {
8     __private size_t i = get_global_id(0);
9
10    if(i >= uiNumElements)
11        return;
12 }
13
14
15 //Binary Map kernel
16 __kernel void MapKernel2(__global TYPE* input1, __global TYPE* input2, __global TYPE
17   * output, unsigned int outputOffset, unsigned int uiNumElements)
18 {
19     __private size_t i = get_global_id(0);
20     if(i >= uiNumElements) return;
21     output[i + outputOffset] = FUN_NAME(input1[i], input2[i]);
22 }
```

Reduce Kernel

```

1 //Reduce kernel based on NVIDIA reduce6 kernel example
2 __kernel void ReduceKernel(__global TYPE* input, __global TYPE* output, unsigned int
3   n, __local TYPE* localData) {
4     __unsigned int blockSize = get_local_size(0);
5     __unsigned int tid = get_local_id(0);
6     __unsigned int i = get_global_id(0);
7     __unsigned int gridSize = blockSize * get_num_groups(0);
8
9     TYPE result = 0;
10    if(i < n) {
11        result = input[i];
12        i += gridSize;
13    }
14    while(i < n) {
15        result = FUN_NAME(result, input[i]);
16        i += gridSize;
17    }
18 }
```

```

18     localData[ tid ] = result ;
19
20     barrier(CLK_LOCAL_MEM_FENCE) ;
21
22 // unrolled loop
23     if( blockSize >= 512 ) {
24         if ( tid < 256 && tid + 256 < n )
25             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 256 ] ) ;
26             barrier(CLK_LOCAL_MEM_FENCE) ;
27     }
28     if( blockSize >= 256 ) {
29         if ( tid < 128 && tid + 128 < n )
30             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 128 ] ) ;
31             barrier(CLK_LOCAL_MEM_FENCE) ;
32     }
33     if( blockSize >= 128 ) {
34         if ( tid < 64 && tid + 64 < n )
35             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 64 ] ) ;
36             barrier(CLK_LOCAL_MEM_FENCE) ; }
37
38     if( blockSize >= 64 ) {
39         if ( tid < 32 && tid + 32 < n )
40             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 32 ] ) ;
41             barrier(CLK_LOCAL_MEM_FENCE) ;
42     }
43     if( blockSize >= 32 ) {
44         if ( tid < 16 && tid + 16 < n )
45             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 16 ] ) ;
46             barrier(CLK_LOCAL_MEM_FENCE) ;
47     }
48     if( blockSize >= 16 ) {
49         if ( tid < 8 && tid + 8 < n )
50             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 8 ] ) ;
51             barrier(CLK_LOCAL_MEM_FENCE) ;
52     }
53     if( blockSize >= 8 ) {
54         if ( tid < 4 && tid + 4 < n )
55             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 4 ] ) ;
56             barrier(CLK_LOCAL_MEM_FENCE) ;
57     }
58     if( blockSize >= 4 ) {
59         if ( tid < 2 && tid + 2 < n )
60             localData[ tid ] = FUN_NAME( localData[ tid ] , localData[ tid + 2 ] ) ;
61             barrier(CLK_LOCAL_MEM_FENCE) ;
62     }
63     if( blockSize >= 2 ) {
64         if ( tid < 1 && tid + 1 < n )

```

```
65         localData [ tid ] = FUN_NAME( localData [ tid ] , localData [ tid + 1 ]) ;
66         barrier (CLK.LOCAL.MEM.FENCE) ;
67     }
68
69     if (tid == 0)  output [ get_group_id (0) ] = localData [ tid ];
70 }
```